The Python Language Reference

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This reference manual describes the syntax and "core semantics" of the language. It is terse, but attempts to be exact and complete. The semantics of non-essential built-in object types and of the built-in functions and modules are described in *library-index*. For an informal introduction to the language, see *tutorial-index*. For C or C++ programmers, two additional manuals exist: *extending-index* describes the high-level picture of how to write a Python extension module, and the *c-api-index* describes the interfaces available to C/C++ programmers in detail.

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INTRODUCTION

This reference manual describes the Python programming language. It is not intended as a tutorial.

While I am trying to be as precise as possible, I chose to use English rather than formal specifications for everything except syntax and lexical analysis. This should make the document more understandable to the average reader, but will leave room for ambiguities. Consequently, if you were coming from Mars and tried to re-implement Python from this document alone, you might have to guess things and in fact you would probably end up implementing quite a different language. On the other hand, if you are using Python and wonder what the precise rules about a particular area of the language are, you should definitely be able to find them here. If you would like to see a more formal definition of the language, maybe you could volunteer your time — or invent a cloning machine:-).

It is dangerous to add too many implementation details to a language reference document — the implementation may change, and other implementations of the same language may work differently. On the other hand, CPython is the one Python implementation in widespread use (although alternate implementations continue to gain support), and its particular quirks are sometimes worth being mentioned, especially where the implementation imposes additional limitations. Therefore, you'll find short "implementation notes" sprinkled throughout the text.

Every Python implementation comes with a number of built-in and standard modules. These are documented in *library-index*. A few built-in modules are mentioned when they interact in a significant way with the language definition.

1.1 Alternate Implementations

Though there is one Python implementation which is by far the most popular, there are some alternate implementations which are of particular interest to different audiences.

Known implementations include:

- **CPython** This is the original and most-maintained implementation of Python, written in C. New language features generally appear here first.
- **Jython** Python implemented in Java. This implementation can be used as a scripting language for Java applications, or can be used to create applications using the Java class libraries. It is also often used to create tests for Java libraries. More information can be found at the Jython website.
- **Python for .NET** This implementation actually uses the CPython implementation, but is a managed .NET application and makes .NET libraries available. It was created by Brian Lloyd. For more information, see the Python for .NET home page.
- **IronPython** An alternate Python for .NET. Unlike Python.NET, this is a complete Python implementation that generates IL, and compiles Python code directly to .NET assemblies. It was created by Jim Hugunin, the original creator of Jython. For more information, see the IronPython website.

PyPy An implementation of Python written completely in Python. It supports several advanced features not found in other implementations like stackless support and a Just in Time compiler. One of the goals of the project is to encourage experimentation with the language itself by making it easier to modify the interpreter (since it is written in Python). Additional information is available on the PyPy project's home page.

Each of these implementations varies in some way from the language as documented in this manual, or introduces specific information beyond what's covered in the standard Python documentation. Please refer to the implementation-specific documentation to determine what else you need to know about the specific implementation you're using.

1.2 Notation

The descriptions of lexical analysis and syntax use a modified BNF grammar notation. This uses the following style of definition:

The first line says that a name is an lc_letter followed by a sequence of zero or more lc_letters and underscores. An lc_letter in turn is any of the single characters 'a' through 'z'. (This rule is actually adhered to for the names defined in lexical and grammar rules in this document.)

Each rule begins with a name (which is the name defined by the rule) and ::=. A vertical bar (|) is used to separate alternatives; it is the least binding operator in this notation. A star (*) means zero or more repetitions of the preceding item; likewise, a plus (+) means one or more repetitions, and a phrase enclosed in square brackets ([]) means zero or one occurrences (in other words, the enclosed phrase is optional). The * and + operators bind as tightly as possible; parentheses are used for grouping. Literal strings are enclosed in quotes. White space is only meaningful to separate tokens. Rules are normally contained on a single line; rules with many alternatives may be formatted alternatively with each line after the first beginning with a vertical bar.

In lexical definitions (as the example above), two more conventions are used: Two literal characters separated by three dots mean a choice of any single character in the given (inclusive) range of ASCII characters. A phrase between angular brackets (<...>) gives an informal description of the symbol defined; e.g., this could be used to describe the notion of 'control character' if needed.

Even though the notation used is almost the same, there is a big difference between the meaning of lexical and syntactic definitions: a lexical definition operates on the individual characters of the input source, while a syntax definition operates on the stream of tokens generated by the lexical analysis. All uses of BNF in the next chapter ("Lexical Analysis") are lexical definitions; uses in subsequent chapters are syntactic definitions.

LEXICAL ANALYSIS

A Python program is read by a *parser*. Input to the parser is a stream of *tokens*, generated by the *lexical analyzer*. This chapter describes how the lexical analyzer breaks a file into tokens.

Python reads program text as Unicode code points; the encoding of a source file can be given by an encoding declaration and defaults to UTF-8, see **PEP 3120** for details. If the source file cannot be decoded, a SyntaxError is raised.

2.1 Line structure

A Python program is divided into a number of logical lines.

2.1.1 Logical lines

The end of a logical line is represented by the token NEWLINE. Statements cannot cross logical line boundaries except where NEWLINE is allowed by the syntax (e.g., between statements in compound statements). A logical line is constructed from one or more *physical lines* by following the explicit or implicit *line joining* rules.

2.1.2 Physical lines

A physical line is a sequence of characters terminated by an end-of-line sequence. In source files, any of the standard platform line termination sequences can be used - the Unix form using ASCII LF (linefeed), the Windows form using the ASCII sequence CR LF (return followed by linefeed), or the old Macintosh form using the ASCII CR (return) character. All of these forms can be used equally, regardless of platform.

When embedding Python, source code strings should be passed to Python APIs using the standard C conventions for newline characters (the \n character, representing ASCII LF, is the line terminator).

2.1.3 Comments

A comment starts with a hash character (#) that is not part of a string literal, and ends at the end of the physical line. A comment signifies the end of the logical line unless the implicit line joining rules are invoked. Comments are ignored by the syntax; they are not tokens.

2.1.4 Encoding declarations

If a comment in the first or second line of the Python script matches the regular expression $coding[=:]\star$ ($[-\w.]+$), this comment is processed as an encoding declaration; the first group of this expression names the encoding of the source code file. The recommended forms of this expression are

```
# -*- coding: <encoding-name> -*-
```

which is recognized also by GNU Emacs, and

```
# vim:fileencoding=<encoding-name>
```

which is recognized by Bram Moolenaar's VIM.

If no encoding declaration is found, the default encoding is UTF-8. In addition, if the first bytes of the file are the UTF-8 byte-order mark (b'\xef\xbb\xbf'), the declared file encoding is UTF-8 (this is supported, among others, by Microsoft's **notepad**).

If an encoding is declared, the encoding name must be recognized by Python. The encoding is used for all lexical analysis, including string literals, comments and identifiers. The encoding declaration must appear on a line of its own

2.1.5 Explicit line joining

Two or more physical lines may be joined into logical lines using backslash characters (\), as follows: when a physical line ends in a backslash that is not part of a string literal or comment, it is joined with the following forming a single logical line, deleting the backslash and the following end-of-line character. For example:

```
if 1900 < year < 2100 and 1 <= month <= 12 \
    and 1 <= day <= 31 and 0 <= hour < 24 \
    and 0 <= minute < 60 and 0 <= second < 60:  # Looks like a valid date
    return 1</pre>
```

A line ending in a backslash cannot carry a comment. A backslash does not continue a comment. A backslash does not continue a token except for string literals (i.e., tokens other than string literals cannot be split across physical lines using a backslash). A backslash is illegal elsewhere on a line outside a string literal.

2.1.6 Implicit line joining

Expressions in parentheses, square brackets or curly braces can be split over more than one physical line without using backslashes. For example:

Implicitly continued lines can carry comments. The indentation of the continuation lines is not important. Blank continuation lines are allowed. There is no NEWLINE token between implicit continuation lines. Implicitly continued lines can also occur within triple-quoted strings (see below); in that case they cannot carry comments.

2.1.7 Blank lines

A logical line that contains only spaces, tabs, formfeeds and possibly a comment, is ignored (i.e., no NEWLINE token is generated). During interactive input of statements, handling of a blank line may differ depending on the

implementation of the read-eval-print loop. In the standard interactive interpreter, an entirely blank logical line (i.e. one containing not even whitespace or a comment) terminates a multi-line statement.

2.1.8 Indentation

Leading whitespace (spaces and tabs) at the beginning of a logical line is used to compute the indentation level of the line, which in turn is used to determine the grouping of statements.

Tabs are replaced (from left to right) by one to eight spaces such that the total number of characters up to and including the replacement is a multiple of eight (this is intended to be the same rule as used by Unix). The total number of spaces preceding the first non-blank character then determines the line's indentation. Indentation cannot be split over multiple physical lines using backslashes; the whitespace up to the first backslash determines the indentation.

Indentation is rejected as inconsistent if a source file mixes tabs and spaces in a way that makes the meaning dependent on the worth of a tab in spaces; a TabError is raised in that case.

Cross-platform compatibility note: because of the nature of text editors on non-UNIX platforms, it is unwise to use a mixture of spaces and tabs for the indentation in a single source file. It should also be noted that different platforms may explicitly limit the maximum indentation level.

A formfeed character may be present at the start of the line; it will be ignored for the indentation calculations above. Formfeed characters occurring elsewhere in the leading whitespace have an undefined effect (for instance, they may reset the space count to zero).

The indentation levels of consecutive lines are used to generate INDENT and DEDENT tokens, using a stack, as follows.

Before the first line of the file is read, a single zero is pushed on the stack; this will never be popped off again. The numbers pushed on the stack will always be strictly increasing from bottom to top. At the beginning of each logical line, the line's indentation level is compared to the top of the stack. If it is equal, nothing happens. If it is larger, it is pushed on the stack, and one INDENT token is generated. If it is smaller, it *must* be one of the numbers occurring on the stack; all numbers on the stack that are larger are popped off, and for each number popped off a DEDENT token is generated. At the end of the file, a DEDENT token is generated for each number remaining on the stack that is larger than zero.

Here is an example of a correctly (though confusingly) indented piece of Python code:

```
def perm(1):
          # Compute the list of all permutations of 1
    if len(1) <= 1:
                return [1]
    r = []
    for i in range(len(1)):
                s = 1[:i] + 1[i+1:]
                p = perm(s)
                for x in p:
                     r.append(l[i:i+1] + x)
    return r</pre>
```

The following example shows various indentation errors:

```
def perm(1):  # error: first line indented
for i in range(len(l)):  # error: not indented
    s = l[:i] + l[i+1:]  # error: unexpected indent
    for x in p:
        r.append(l[i:i+1] + x)
    return r  # error: inconsistent dedent
```

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(Actually, the first three errors are detected by the parser; only the last error is found by the lexical analyzer — the indentation of return r does not match a level popped off the stack.)

2.1.9 Whitespace between tokens

Except at the beginning of a logical line or in string literals, the whitespace characters space, tab and formfeed can be used interchangeably to separate tokens. Whitespace is needed between two tokens only if their concatenation could otherwise be interpreted as a different token (e.g., ab is one token, but a b is two tokens).

2.2 Other tokens

Besides NEWLINE, INDENT and DEDENT, the following categories of tokens exist: *identifiers*, *keywords*, *literals*, *operators*, and *delimiters*. Whitespace characters (other than line terminators, discussed earlier) are not tokens, but serve to delimit tokens. Where ambiguity exists, a token comprises the longest possible string that forms a legal token, when read from left to right.

2.3 Identifiers and keywords

Identifiers (also referred to as *names*) are described by the following lexical definitions.

The syntax of identifiers in Python is based on the Unicode standard annex UAX-31, with elaboration and changes as defined below; see also **PEP 3131** for further details.

Within the ASCII range (U+0001..U+007F), the valid characters for identifiers are the same as in Python 2.x: the uppercase and lowercase letters A through \mathbb{Z} , the underscore $\underline{}$ and, except for the first character, the digits \mathbb{Q} through \mathbb{Q}

Python 3.0 introduces additional characters from outside the ASCII range (see

PEP 3131). For these characters, the classification uses the version of the Unicode Character Database as included in the unicodedata module.

Identifiers are unlimited in length. Case is significant.

```
identifier ::= xid_start xid_continue*
id_start ::= <all characters in general categories Lu, Ll, Lt, Lm, Lo, Nl, the u
id_continue ::= <all characters in id_start, plus characters in the categories Mn,
xid_start ::= <all characters in id_start whose NFKC normalization is in "id_start
xid_continue ::= <all characters in id_continue whose NFKC normalization is in "id_continue"</pre>
```

The Unicode category codes mentioned above stand for:

- Lu uppercase letters
- Ll lowercase letters
- Lt titlecase letters
- Lm modifier letters
- Lo other letters
- Nl letter numbers
- Mn nonspacing marks

- Mc spacing combining marks
- Nd decimal numbers
- Pc connector punctuations
- Other_ID_Start explicit list of characters in PropList.txt to support backwards compatibility
- Other ID Continue likewise

All identifiers are converted into the normal form NFKC while parsing; comparison of identifiers is based on NFKC.

A non-normative HTML file listing all valid identifier characters for Unicode 4.1 can be found at http://www.dcl.hpi.uni-potsdam.de/home/loewis/table-3131.html.

2.3.1 Keywords

The following identifiers are used as reserved words, or *keywords* of the language, and cannot be used as ordinary identifiers. They must be spelled exactly as written here:

False	class	finally	is	return
None	continue	for	lambda	try
True	def	from	nonlocal	while
and	del	global	not	with
as	elif	if	or	yield
assert	else	import	pass	
break	except	in	raise	

2.3.2 Reserved classes of identifiers

Certain classes of identifiers (besides keywords) have special meanings. These classes are identified by the patterns of leading and trailing underscore characters:

_* Not imported by from module import *. The special identifier _ is used in the interactive interpreter to store the result of the last evaluation; it is stored in the builtins module. When not in interactive mode, _ has no special meaning and is not defined. See section *The import statement*.

Note: The name _ is often used in conjunction with internationalization; refer to the documentation for the gettext module for more information on this convention.

- __*__ System-defined names. These names are defined by the interpreter and its implementation (including the standard library). Current system names are discussed in the *Special method names* section and elsewhere. More will likely be defined in future versions of Python. *Any* use of ___*__ names, in any context, that does not follow explicitly documented use, is subject to breakage without warning.
- __* Class-private names. Names in this category, when used within the context of a class definition, are re-written to use a mangled form to help avoid name clashes between "private" attributes of base and derived classes. See section *Identifiers (Names)*.

2.4 Literals

Literals are notations for constant values of some built-in types.

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2.4.1 String and Bytes literals

String literals are described by the following lexical definitions:

```
stringliteral
                     [stringprefix] (shortstring | longstring)
                     "r" | "u" | "R" | "U"
stringprefix
                     ""' shortstringitem* ""' | '"' shortstringitem* '"'
shortstring
                ::=
                     "'" longstringitem* "'" | '"" longstringitem* """
longstring
                ::=
                     shortstringchar | stringescapeseq
shortstringitem ::=
longstringitem ::=
                     longstringchar | stringescapeseq
                     <any source character except "\" or newline or the quote>
shortstringchar ::=
                     <any source character except "\">
longstringchar
                ::=
stringescapeseq ::=
                     "\" <any source character>
bytesliteral
                    bytesprefix(shortbytes | longbytes)
               ::=
                    "b" | "B" | "br" | "Br" | "bR" | "BR" | "rb" | "rB" | "Rb" | "RB"
bytesprefix
               ::=
                    ""' shortbytesitem* ""' | '"' shortbytesitem* '"'
shortbytes
               ::=
longbytes
               ::=
                    "'" longbytesitem* "'" | '""" longbytesitem* """"
shortbytesitem ::=
                    shortbyteschar | bytesescapeseq
longbytesitem ::=
                    longbyteschar | bytesescapeseq
                    <any ASCII character except "\" or newline or the quote>
shortbyteschar ::=
                    <any ASCII character except "\">
longbyteschar
               ::=
bytesescapeseq ::=
                    "\" <any ASCII character>
```

One syntactic restriction not indicated by these productions is that whitespace is not allowed between the stringprefix or bytesprefix and the rest of the literal. The source character set is defined by the encoding declaration; it is UTF-8 if no encoding declaration is given in the source file; see section *Encoding declarations*.

In plain English: Both types of literals can be enclosed in matching single quotes (') or double quotes ("). They can also be enclosed in matching groups of three single or double quotes (these are generally referred to as *triple-quoted strings*). The backslash (\) character is used to escape characters that otherwise have a special meaning, such as newline, backslash itself, or the quote character.

Bytes literals are always prefixed with 'b' or 'B'; they produce an instance of the bytes type instead of the str type. They may only contain ASCII characters; bytes with a numeric value of 128 or greater must be expressed with escapes.

As of Python 3.3 it is possible again to prefix unicode strings with a u prefix to simplify maintenance of dual 2.x and 3.x codebases.

Both string and bytes literals may optionally be prefixed with a letter 'r' or 'R'; such strings are called *raw strings* and treat backslashes as literal characters. As a result, in string literals, '\U' and '\u' escapes in raw strings are not treated specially. Given that Python 2.x's raw unicode literals behave differently than Python 3.x's the 'ur' syntax is not supported.

New in version 3.3: The 'rb' prefix of raw bytes literals has been added as a synonym of 'br'. New in version 3.3: Support for the unicode legacy literal (u'value') was reintroduced to simplify the maintenance of dual Python 2.x and 3.x codebases. See PEP 414 for more information.

In triple-quoted strings, unescaped newlines and quotes are allowed (and are retained), except that three unescaped quotes in a row terminate the string. (A "quote" is the character used to open the string, i.e. either ' or ".)

Unless an 'r' or 'R' prefix is present, escape sequences in strings are interpreted according to rules similar to those used by Standard C. The recognized escape sequences are:

Escape Sequence	Meaning	Notes
\newline	Backslash and newline ignored	
\\	Backslash (\)	
\'	Single quote (')	
\"	Double quote (")	
\a	ASCII Bell (BEL)	
\b	ASCII Backspace (BS)	
\f	ASCII Formfeed (FF)	
\n	ASCII Linefeed (LF)	
\r	ASCII Carriage Return (CR)	
\t	ASCII Horizontal Tab (TAB)	
\v	ASCII Vertical Tab (VT)	
\000	Character with octal value ooo	(1,3)
\xhh	Character with hex value hh	(2,3)

Escape sequences only recognized in string literals are:

Escape Sequence	Meaning	Notes
\N{name}	Character named <i>name</i> in the Unicode database	(4)
\uxxxx	Character with 16-bit hex value xxxx	(5)
\Uxxxxxxx	Character with 32-bit hex value xxxxxxxx	(6)

Notes:

- 1. As in Standard C, up to three octal digits are accepted.
- 2. Unlike in Standard C, exactly two hex digits are required.
- 3. In a bytes literal, hexadecimal and octal escapes denote the byte with the given value. In a string literal, these escapes denote a Unicode character with the given value.
- 4. Changed in version 3.3: Support for name aliases ¹ has been added.
- 5. Individual code units which form parts of a surrogate pair can be encoded using this escape sequence. Exactly four hex digits are required.
- 6. Any Unicode character can be encoded this way. Exactly eight hex digits are required.

Unlike Standard C, all unrecognized escape sequences are left in the string unchanged, i.e., the backslash is left in the string. (This behavior is useful when debugging: if an escape sequence is mistyped, the resulting output is more easily recognized as broken.) It is also important to note that the escape sequences only recognized in string literals fall into the category of unrecognized escapes for bytes literals.

Even in a raw string, string quotes can be escaped with a backslash, but the backslash remains in the string; for example, r"\"" is a valid string literal consisting of two characters: a backslash and a double quote; r"\" is not a valid string literal (even a raw string cannot end in an odd number of backslashes). Specifically, a raw string cannot end in a single backslash (since the backslash would escape the following quote character). Note also that a single backslash followed by a newline is interpreted as those two characters as part of the string, not as a line continuation.

2.4.2 String literal concatenation

Multiple adjacent string or bytes literals (delimited by whitespace), possibly using different quoting conventions, are allowed, and their meaning is the same as their concatenation. Thus, "hello" 'world' is equivalent to "helloworld". This feature can be used to reduce the number of backslashes needed, to split long strings conveniently across long lines, or even to add comments to parts of strings, for example:

2.4. Literals

¹ http://www.unicode.org/Public/6.1.0/ucd/NameAliases.txt

Note that this feature is defined at the syntactical level, but implemented at compile time. The '+' operator must be used to concatenate string expressions at run time. Also note that literal concatenation can use different quoting styles for each component (even mixing raw strings and triple quoted strings).

2.4.3 Numeric literals

There are three types of numeric literals: integers, floating point numbers, and imaginary numbers. There are no complex literals (complex numbers can be formed by adding a real number and an imaginary number).

Note that numeric literals do not include a sign; a phrase like -1 is actually an expression composed of the unary operator '-' and the literal 1.

2.4.4 Integer literals

Integer literals are described by the following lexical definitions:

```
integer
               ::=
                    decimalinteger | octinteger | hexinteger | bininteger
decimalinteger ::= nonzerodigit digit* | "0"+
nonzerodigit ::= "1"..."9"
               ::= "0"..."9"
digit
octinteger
hexinteger
                    "0" ("o" | "0") octdigit+
               ::=
                    "0" ("x" | "X") hexdigit+
              ::=
bininteger
              ::=
                    "0" ("b" | "B") bindigit+
octdigit
                    "0"..."7"
               ::=
                    digit | "a"..."f" | "A"..."F"
hexdigit
               ::=
               ::= "0" | "1"
bindigit
```

There is no limit for the length of integer literals apart from what can be stored in available memory.

Note that leading zeros in a non-zero decimal number are not allowed. This is for disambiguation with C-style octal literals, which Python used before version 3.0.

Some examples of integer literals:

```
7 2147483647 00177 0b100110111
3 79228162514264337593543950336 00377 0x10000000
79228162514264337593543950336 0xdeadbeef
```

2.4.5 Floating point literals

Floating point literals are described by the following lexical definitions:

```
floatnumber
pointfloat ::= pointfloat | exponentfloat
pointfloat ::= [intpart] fraction | intpart "."
exponentfloat ::= (intpart | pointfloat) exponent
intpart ::= digit+
fraction ::= "." digit+
exponent ::= ("e" | "E") ["+" | "-"] digit+
```

Note that the integer and exponent parts are always interpreted using radix 10. For example, 077e010 is legal, and denotes the same number as 77e10. The allowed range of floating point literals is implementation-dependent. Some examples of floating point literals:

```
3.14 10. .001 1e100 3.14e-10 0e0
```

Note that numeric literals do not include a sign; a phrase like -1 is actually an expression composed of the unary operator – and the literal 1.

2.4.6 Imaginary literals

Imaginary literals are described by the following lexical definitions:

```
imagnumber ::= (floatnumber | intpart) ("j" | "J")
```

An imaginary literal yields a complex number with a real part of 0.0. Complex numbers are represented as a pair of floating point numbers and have the same restrictions on their range. To create a complex number with a nonzero real part, add a floating point number to it, e.g., (3+4j). Some examples of imaginary literals:

```
3.14j 10.j 10j .001j 1e100j 3.14e-10j
```

2.5 Operators

The following tokens are operators:

```
+ - * ** / // §
<< >> & | ^ ~
< < >> == !=
```

2.6 Delimiters

The following tokens serve as delimiters in the grammar:

```
( ) [ ] { }
, : . ; @ =
+= -= *= /= //= %=
&= |= ^= >>= <<= **=
```

The period can also occur in floating-point and imaginary literals. A sequence of three periods has a special meaning as an ellipsis literal. The second half of the list, the augmented assignment operators, serve lexically as delimiters, but also perform an operation.

The following printing ASCII characters have special meaning as part of other tokens or are otherwise significant to the lexical analyzer:

```
' # \
```

The following printing ASCII characters are not used in Python. Their occurrence outside string literals and comments is an unconditional error:

```
$ ?
```

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CHAPTER

THREE

DATA MODEL

3.1 Objects, values and types

Objects are Python's abstraction for data. All data in a Python program is represented by objects or by relations between objects. (In a sense, and in conformance to Von Neumann's model of a "stored program computer," code is also represented by objects.)

Every object has an identity, a type and a value. An object's *identity* never changes once it has been created; you may think of it as the object's address in memory. The 'is' operator compares the identity of two objects; the id() function returns an integer representing its identity.

CPython implementation detail: For CPython, id(x) is the memory address where x is stored.

An object's type determines the operations that the object supports (e.g., "does it have a length?") and also defines the possible values for objects of that type. The type () function returns an object's type (which is an object itself). Like its identity, an object's *type* is also unchangeable. ¹

The *value* of some objects can change. Objects whose value can change are said to be *mutable*; objects whose value is unchangeable once they are created are called *immutable*. (The value of an immutable container object that contains a reference to a mutable object can change when the latter's value is changed; however the container is still considered immutable, because the collection of objects it contains cannot be changed. So, immutability is not strictly the same as having an unchangeable value, it is more subtle.) An object's mutability is determined by its type; for instance, numbers, strings and tuples are immutable, while dictionaries and lists are mutable.

Objects are never explicitly destroyed; however, when they become unreachable they may be garbage-collected. An implementation is allowed to postpone garbage collection or omit it altogether — it is a matter of implementation quality how garbage collection is implemented, as long as no objects are collected that are still reachable.

CPython implementation detail: CPython currently uses a reference-counting scheme with (optional) delayed detection of cyclically linked garbage, which collects most objects as soon as they become unreachable, but is not guaranteed to collect garbage containing circular references. See the documentation of the gc module for information on controlling the collection of cyclic garbage. Other implementations act differently and CPython may change. Do not depend on immediate finalization of objects when they become unreachable (ex: always close files).

Note that the use of the implementation's tracing or debugging facilities may keep objects alive that would normally be collectable. Also note that catching an exception with a 'try...except' statement may keep objects alive.

Some objects contain references to "external" resources such as open files or windows. It is understood that these resources are freed when the object is garbage-collected, but since garbage collection is not guaranteed to happen, such objects also provide an explicit way to release the external resource, usually a close () method. Programs are strongly recommended to explicitly close such objects. The 'try...finally' statement and the 'with' statement provide convenient ways to do this.

¹ It is possible in some cases to change an object's type, under certain controlled conditions. It generally isn't a good idea though, since it can lead to some very strange behaviour if it is handled incorrectly.

Some objects contain references to other objects; these are called *containers*. Examples of containers are tuples, lists and dictionaries. The references are part of a container's value. In most cases, when we talk about the value of a container, we imply the values, not the identities of the contained objects; however, when we talk about the mutability of a container, only the identities of the immediately contained objects are implied. So, if an immutable container (like a tuple) contains a reference to a mutable object, its value changes if that mutable object is changed.

Types affect almost all aspects of object behavior. Even the importance of object identity is affected in some sense: for immutable types, operations that compute new values may actually return a reference to any existing object with the same type and value, while for mutable objects this is not allowed. E.g., after a = 1; b = 1, a and b may or may not refer to the same object with the value one, depending on the implementation, but after c = []; d = [], c and d are guaranteed to refer to two different, unique, newly created empty lists. (Note that c = d = [] assigns the same object to both c and d.)

3.2 The standard type hierarchy

Below is a list of the types that are built into Python. Extension modules (written in C, Java, or other languages, depending on the implementation) can define additional types. Future versions of Python may add types to the type hierarchy (e.g., rational numbers, efficiently stored arrays of integers, etc.), although such additions will often be provided via the standard library instead.

Some of the type descriptions below contain a paragraph listing 'special attributes.' These are attributes that provide access to the implementation and are not intended for general use. Their definition may change in the future.

None This type has a single value. There is a single object with this value. This object is accessed through the built-in name None. It is used to signify the absence of a value in many situations, e.g., it is returned from functions that don't explicitly return anything. Its truth value is false.

NotImplemented This type has a single value. There is a single object with this value. This object is accessed through the built-in name NotImplemented. Numeric methods and rich comparison methods may return this value if they do not implement the operation for the operands provided. (The interpreter will then try the reflected operation, or some other fallback, depending on the operator.) Its truth value is true.

Ellipsis This type has a single value. There is a single object with this value. This object is accessed through the literal . . . or the built-in name Ellipsis. Its truth value is true.

numbers.Number These are created by numeric literals and returned as results by arithmetic operators and arithmetic built-in functions. Numeric objects are immutable; once created their value never changes. Python numbers are of course strongly related to mathematical numbers, but subject to the limitations of numerical representation in computers.

Python distinguishes between integers, floating point numbers, and complex numbers:

numbers.Integral These represent elements from the mathematical set of integers (positive and negative).

There are two types of integers:

Integers (int)

These represent numbers in an unlimited range, subject to available (virtual) memory only. For the purpose of shift and mask operations, a binary representation is assumed, and negative numbers are represented in a variant of 2's complement which gives the illusion of an infinite string of sign bits extending to the left.

Booleans (bool) These represent the truth values False and True. The two objects representing the values False and True are the only Boolean objects. The Boolean type is a subtype of the integer type, and Boolean values behave like the values 0 and 1, respectively, in almost all contexts, the exception being that when converted to a string, the strings "False" or "True" are returned, respectively.

The rules for integer representation are intended to give the most meaningful interpretation of shift and mask operations involving negative integers.

- numbers.Real (float) These represent machine-level double precision floating point numbers. You are at the mercy of the underlying machine architecture (and C or Java implementation) for the accepted range and handling of overflow. Python does not support single-precision floating point numbers; the savings in processor and memory usage that are usually the reason for using these is dwarfed by the overhead of using objects in Python, so there is no reason to complicate the language with two kinds of floating point numbers.
- **numbers.Complex** (complex) These represent complex numbers as a pair of machine-level double precision floating point numbers. The same caveats apply as for floating point numbers. The real and imaginary parts of a complex number z can be retrieved through the read-only attributes z.real and z.imag.
- **Sequences** These represent finite ordered sets indexed by non-negative numbers. The built-in function len() returns the number of items of a sequence. When the length of a sequence is n, the index set contains the numbers 0, 1, ..., n-1. Item i of sequence a is selected by a[i].

Sequences also support slicing: a[i:j] selects all items with index k such that $i \le k \le j$. When used as an expression, a slice is a sequence of the same type. This implies that the index set is renumbered so that it starts at 0.

Some sequences also support "extended slicing" with a third "step" parameter: a[i:j:k] selects all items of a with index x where x = i + n*k, n >= 0 and i <= x < j.

Sequences are distinguished according to their mutability:

Immutable sequences An object of an immutable sequence type cannot change once it is created. (If the object contains references to other objects, these other objects may be mutable and may be changed; however, the collection of objects directly referenced by an immutable object cannot change.)

The following types are immutable sequences:

- Strings A string is a sequence of values that represent Unicode codepoints. All the codepoints in range U+0000 U+10FFFF can be represented in a string. Python doesn't have a chr type, and every character in the string is represented as a string object with length 1. The built-in function ord() converts a character to its codepoint (as an integer); chr() converts an integer in range 0 10FFFF to the corresponding character. str.encode() can be used to convert a str to bytes using the given encoding, and bytes.decode() can be used to achieve the opposite.
- **Tuples** The items of a tuple are arbitrary Python objects. Tuples of two or more items are formed by comma-separated lists of expressions. A tuple of one item (a 'singleton') can be formed by affixing a comma to an expression (an expression by itself does not create a tuple, since parentheses must be usable for grouping of expressions). An empty tuple can be formed by an empty pair of parentheses.
- Bytes A bytes object is an immutable array. The items are 8-bit bytes, represented by integers in the range $0 \le x \le 256$. Bytes literals (like b'abc') and the built-in function bytes () can be used to construct bytes objects. Also, bytes objects can be decoded to strings via the decode () method.
- **Mutable sequences** Mutable sequences can be changed after they are created. The subscription and slicing notations can be used as the target of assignment and del (delete) statements.

There are currently two intrinsic mutable sequence types:

- **Lists** The items of a list are arbitrary Python objects. Lists are formed by placing a comma-separated list of expressions in square brackets. (Note that there are no special cases needed to form lists of length 0 or 1.)
- **Byte Arrays** A bytearray object is a mutable array. They are created by the built-in bytearray () constructor. Aside from being mutable (and hence unhashable), byte arrays otherwise provide the same interface and functionality as immutable bytes objects.

The extension module array provides an additional example of a mutable sequence type, as does the collections module.

Set types These represent unordered, finite sets of unique, immutable objects. As such, they cannot be indexed by any subscript. However, they can be iterated over, and the built-in function len() returns the number of items in a set. Common uses for sets are fast membership testing, removing duplicates from a sequence, and computing mathematical operations such as intersection, union, difference, and symmetric difference.

For set elements, the same immutability rules apply as for dictionary keys. Note that numeric types obey the normal rules for numeric comparison: if two numbers compare equal (e.g., 1 and 1.0), only one of them can be contained in a set.

There are currently two intrinsic set types:

Sets These represent a mutable set. They are created by the built-in set () constructor and can be modified afterwards by several methods, such as add().

Frozen sets These represent an immutable set. They are created by the built-in frozenset () constructor. As a frozenset is immutable and *hashable*, it can be used again as an element of another set, or as a dictionary key.

Mappings These represent finite sets of objects indexed by arbitrary index sets. The subscript notation a [k] selects the item indexed by k from the mapping a; this can be used in expressions and as the target of assignments or del statements. The built-in function len() returns the number of items in a mapping.

There is currently a single intrinsic mapping type:

Dictionaries These represent finite sets of objects indexed by nearly arbitrary values. The only types of values not acceptable as keys are values containing lists or dictionaries or other mutable types that are compared by value rather than by object identity, the reason being that the efficient implementation of dictionaries requires a key's hash value to remain constant. Numeric types used for keys obey the normal rules for numeric comparison: if two numbers compare equal (e.g., 1 and 1.0) then they can be used interchangeably to index the same dictionary entry.

Dictionaries are mutable; they can be created by the { . . . } notation (see section *Dictionary displays*).

The extension modules dbm.ndbm and dbm.gnu provide additional examples of mapping types, as does the collections module.

Callable types These are the types to which the function call operation (see section *Calls*) can be applied:

User-defined functions A user-defined function object is created by a function definition (see section *Function definitions*). It should be called with an argument list containing the same number of items as the function's formal parameter list.

Special attributes:

Attribute	Meaning	
doc	The function's documentation string, or None if unavailable	Writable
name	The function's name	Writable
qualnam	e <u>Th</u> e function's <i>qualified name</i> New in version 3.3.	Writable
module_	The name of the module the function was defined in, or None if unavailable.	Writable
default	s A tuple containing default argument values for those arguments that have	Writable
	defaults, or None if no arguments have a default value	
code	The code object representing the compiled function body.	Writable
globals	A reference to the dictionary that holds the function's global variables — the	Read-
	global namespace of the module in which the function was defined.	only
dict	The namespace supporting arbitrary function attributes.	Writable
closure	None or a tuple of cells that contain bindings for the function's free variables.	Read-
		only
annotat	Andict containing annotations of parameters. The keys of the dict are the	Writable
	parameter names, or 'return' for the return annotation, if provided.	
kwdefau	tAs <u>dict</u> containing defaults for keyword-only parameters.	Writable

Most of the attributes labelled "Writable" check the type of the assigned value.

Function objects also support getting and setting arbitrary attributes, which can be used, for example, to attach metadata to functions. Regular attribute dot-notation is used to get and set such attributes. *Note that the current implementation only supports function attributes on user-defined functions. Function attributes on built-in functions may be supported in the future.*

Additional information about a function's definition can be retrieved from its code object; see the description of internal types below.

Instance methods An instance method object combines a class, a class instance and any callable object (normally a user-defined function).

Special read-only attributes: __self__ is the class instance object, __func__ is the function object; __doc__ is the method's documentation (same as __func__.__doc__); __name__ is the method name (same as __func__.__name__); __module__ is the name of the module the method was defined in, or None if unavailable.

Methods also support accessing (but not setting) the arbitrary function attributes on the underlying function object.

User-defined method objects may be created when getting an attribute of a class (perhaps via an instance of that class), if that attribute is a user-defined function object or a class method object.

When an instance method object is created by retrieving a user-defined function object from a class via one of its instances, its __self__ attribute is the instance, and the method object is said to be bound. The new method's __func__ attribute is the original function object.

When a user-defined method object is created by retrieving another method object from a class or instance, the behaviour is the same as for a function object, except that the ___func__ attribute of the new instance is not the original method object but its ___func__ attribute.

When an instance method object is created by retrieving a class method object from a class or instance, its __self__ attribute is the class itself, and its __func__ attribute is the function object underlying the class method.

When an instance method object is called, the underlying function ($__func_$) is called, inserting the class instance ($__self_$) in front of the argument list. For instance, when C is a class which contains a definition for a function f(), and x is an instance of C, calling x.f(1) is equivalent to calling C.f(x, 1).

When an instance method object is derived from a class method object, the "class instance" stored in __self__ will actually be the class itself, so that calling either x.f(1) or C.f(1) is equivalent to

calling f(C, 1) where f is the underlying function.

Note that the transformation from function object to instance method object happens each time the attribute is retrieved from the instance. In some cases, a fruitful optimization is to assign the attribute to a local variable and call that local variable. Also notice that this transformation only happens for user-defined functions; other callable objects (and all non-callable objects) are retrieved without transformation. It is also important to note that user-defined functions which are attributes of a class instance are not converted to bound methods; this *only* happens when the function is an attribute of the class.

- Generator functions A function or method which uses the <code>yield</code> statement (see section *The yield statement*) is called a *generator function*. Such a function, when called, always returns an iterator object which can be used to execute the body of the function: calling the iterator's <code>iterator__next__</code>() method will cause the function to execute until it provides a value using the <code>yield</code> statement. When the function executes a <code>return</code> statement or falls off the end, a <code>StopIteration</code> exception is raised and the iterator will have reached the end of the set of values to be returned.
- Built-in functions A built-in function object is a wrapper around a C function. Examples of built-in functions are len() and math.sin() (math is a standard built-in module). The number and type of the arguments are determined by the C function. Special read-only attributes: __doc__ is the function's documentation string, or None if unavailable; __name__ is the function's name; __self__ is set to None (but see the next item); __module__ is the name of the module the function was defined in or None if unavailable.
- **Built-in methods** This is really a different disguise of a built-in function, this time containing an object passed to the C function as an implicit extra argument. An example of a built-in method is alist.append(), assuming *alist* is a list object. In this case, the special read-only attribute __self__ is set to the object denoted by *alist*.
- Classes Classes are callable. These objects normally act as factories for new instances of themselves, but variations are possible for class types that override __new__(). The arguments of the call are passed to __new__() and, in the typical case, to __init__() to initialize the new instance.
- Class Instances Instances of arbitrary classes can be made callable by defining a __call__() method in their class.
- Modules Modules are a basic organizational unit of Python code, and are created by the *import system* as invoked either by the import statement (see import), or by calling functions such as importlib.import_module() and built-in __import__(). A module object has a namespace implemented by a dictionary object (this is the dictionary referenced by the __globals__ attribute of functions defined in the module). Attribute references are translated to lookups in this dictionary, e.g., m.x is equivalent to m.__dict__["x"]. A module object does not contain the code object used to initialize the module (since it isn't needed once the initialization is done).

Attribute assignment updates the module's namespace dictionary, e.g., m.x = 1 is equivalent to $m._dict_["x"] = 1$.

Special read-only attribute: ___dict__ is the module's namespace as a dictionary object.

CPython implementation detail: Because of the way CPython clears module dictionaries, the module dictionary will be cleared when the module falls out of scope even if the dictionary still has live references. To avoid this, copy the dictionary or keep the module around while using its dictionary directly.

Predefined (writable) attributes: __name__ is the module's name; __doc__ is the module's documentation string, or None if unavailable; __file__ is the pathname of the file from which the module was loaded, if it was loaded from a file. The __file__ attribute may be missing for certain types of modules, such as C modules that are statically linked into the interpreter; for extension modules loaded dynamically from a shared library, it is the pathname of the shared library file.

Custom classes Custom class types are typically created by class definitions (see section *Class definitions*). A class has a namespace implemented by a dictionary object. Class attribute references are translated to lookups in this

dictionary, e.g., C.x is translated to C.__dict__["x"] (although there are a number of hooks which allow for other means of locating attributes). When the attribute name is not found there, the attribute search continues in the base classes. This search of the base classes uses the C3 method resolution order which behaves correctly even in the presence of 'diamond' inheritance structures where there are multiple inheritance paths leading back to a common ancestor. Additional details on the C3 MRO used by Python can be found in the documentation accompanying the 2.3 release at http://www.python.org/download/releases/2.3/mro/.

When a class attribute reference (for class C, say) would yield a class method object, it is transformed into an instance method object whose __self__ attributes is C. When it would yield a static method object, it is transformed into the object wrapped by the static method object. See section *Implementing Descriptors* for another way in which attributes retrieved from a class may differ from those actually contained in its __dict__.

Class attribute assignments update the class's dictionary, never the dictionary of a base class.

A class object can be called (see above) to yield a class instance (see below).

Special attributes: __name__ is the class name; __module__ is the module name in which the class was defined; __dict__ is the dictionary containing the class's namespace; __bases__ is a tuple (possibly empty or a singleton) containing the base classes, in the order of their occurrence in the base class list; __doc__ is the class's documentation string, or None if undefined.

Class instances A class instance is created by calling a class object (see above). A class instance has a namespace implemented as a dictionary which is the first place in which attribute references are searched. When an attribute is not found there, and the instance's class has an attribute by that name, the search continues with the class attributes. If a class attribute is found that is a user-defined function object, it is transformed into an instance method object whose __self__ attribute is the instance. Static method and class method objects are also transformed; see above under "Classes". See section *Implementing Descriptors* for another way in which attributes of a class retrieved via its instances may differ from the objects actually stored in the class's __dict__. If no class attribute is found, and the object's class has a __getattr__ () method, that is called to satisfy the lookup.

Attribute assignments and deletions update the instance's dictionary, never a class's dictionary. If the class has a __setattr__() or __delattr__() method, this is called instead of updating the instance dictionary directly.

Class instances can pretend to be numbers, sequences, or mappings if they have methods with certain special names. See section *Special method names*.

Special attributes: __dict__ is the attribute dictionary; __class__ is the instance's class.

I/O objects (also known as file objects) A *file object* represents an open file. Various shortcuts are available to create file objects: the open() built-in function, and also os.popen(), os.fdopen(), and the makefile() method of socket objects (and perhaps by other functions or methods provided by extension modules).

The objects sys.stdin, sys.stdout and sys.stderr are initialized to file objects corresponding to the interpreter's standard input, output and error streams; they are all open in text mode and therefore follow the interface defined by the io.TextIOBase abstract class.

Internal types A few types used internally by the interpreter are exposed to the user. Their definitions may change with future versions of the interpreter, but they are mentioned here for completeness.

Code objects Code objects represent *byte-compiled* executable Python code, or *bytecode*. The difference between a code object and a function object is that the function object contains an explicit reference to the function's globals (the module in which it was defined), while a code object contains no context; also the default argument values are stored in the function object, not in the code object (because they represent values calculated at run-time). Unlike function objects, code objects are immutable and contain no references (directly or indirectly) to mutable objects.

Special read-only attributes: co_name gives the function name; co_argcount is the number of positional arguments (including arguments with default values); co_nlocals is the number of local variables

used by the function (including arguments); co_varnames is a tuple containing the names of the local variables (starting with the argument names); co_cellvars is a tuple containing the names of local variables that are referenced by nested functions; co_freevars is a tuple containing the names of free variables; co_code is a string representing the sequence of bytecode instructions; co_consts is a tuple containing the literals used by the bytecode; co_names is a tuple containing the names used by the bytecode; co_filename is the filename from which the code was compiled; co_firstlineno is the first line number of the function; co_lnotab is a string encoding the mapping from bytecode offsets to line numbers (for details see the source code of the interpreter); co_stacksize is the required stack size (including local variables); co_flags is an integer encoding a number of flags for the interpreter.

The following flag bits are defined for co_flags: bit 0x04 is set if the function uses the *arguments syntax to accept an arbitrary number of positional arguments; bit 0x08 is set if the function uses the **keywords syntax to accept arbitrary keyword arguments; bit 0x20 is set if the function is a generator.

Future feature declarations (from __future__ import division) also use bits in co_flags to indicate whether a code object was compiled with a particular feature enabled: bit 0×2000 is set if the function was compiled with future division enabled; bits 0×10 and 0×1000 were used in earlier versions of Python.

Other bits in co flags are reserved for internal use.

If a code object represents a function, the first item in co_consts is the documentation string of the function, or None if undefined.

Frame objects Frame objects represent execution frames. They may occur in traceback objects (see below).

Special read-only attributes: f_back is to the previous stack frame (towards the caller), or None if this is the bottom stack frame; f_code is the code object being executed in this frame; f_locals is the dictionary used to look up local variables; f_globals is used for global variables; f_builtins is used for built-in (intrinsic) names; f_lasti gives the precise instruction (this is an index into the bytecode string of the code object).

Special writable attributes: f_trace, if not None, is a function called at the start of each source code line (this is used by the debugger); f_lineno is the current line number of the frame — writing to this from within a trace function jumps to the given line (only for the bottom-most frame). A debugger can implement a Jump command (aka Set Next Statement) by writing to f_lineno.

Traceback objects Traceback objects represent a stack trace of an exception. A traceback object is created when an exception occurs. When the search for an exception handler unwinds the execution stack, at each unwound level a traceback object is inserted in front of the current traceback. When an exception handler is entered, the stack trace is made available to the program. (See section *The try statement*.) It is accessible as the third item of the tuple returned by <code>sys.exc_info()</code>. When the program contains no suitable handler, the stack trace is written (nicely formatted) to the standard error stream; if the interpreter is interactive, it is also made available to the user as <code>sys.last_traceback</code>.

Special read-only attributes: tb_next is the next level in the stack trace (towards the frame where the exception occurred), or None if there is no next level; tb_frame points to the execution frame of the current level; tb_lineno gives the line number where the exception occurred; tb_lasti indicates the precise instruction. The line number and last instruction in the traceback may differ from the line number of its frame object if the exception occurred in a try statement with no matching except clause or with a finally clause.

Slice objects Slice objects are used to represent slices for __getitem__() methods. They are also created by the built-in slice() function.

Special read-only attributes: start is the lower bound; stop is the upper bound; step is the step value; each is None if omitted. These attributes can have any type.

Slice objects support one method:

```
slice.indices (self, length)
```

This method takes a single integer argument *length* and computes information about the slice that the slice object would describe if applied to a sequence of *length* items. It returns a tuple of three integers; respectively these are the *start* and *stop* indices and the *step* or stride length of the slice. Missing or out-of-bounds indices are handled in a manner consistent with regular slices.

Static method objects Static method objects provide a way of defeating the transformation of function objects to method objects described above. A static method object is a wrapper around any other object, usually a user-defined method object. When a static method object is retrieved from a class or a class instance, the object actually returned is the wrapped object, which is not subject to any further transformation. Static method objects are not themselves callable, although the objects they wrap usually are. Static method objects are created by the built-in staticmethod() constructor.

Class method objects A class method object, like a static method object, is a wrapper around another object that alters the way in which that object is retrieved from classes and class instances. The behaviour of class method objects upon such retrieval is described above, under "User-defined methods". Class method objects are created by the built-in classmethod() constructor.

3.3 Special method names

A class can implement certain operations that are invoked by special syntax (such as arithmetic operations or subscripting and slicing) by defining methods with special names. This is Python's approach to *operator overloading*, allowing classes to define their own behavior with respect to language operators. For instance, if a class defines a method named __getitem__(), and x is an instance of this class, then x[i] is roughly equivalent to type(x).__getitem__(x, i). Except where mentioned, attempts to execute an operation raise an exception when no appropriate method is defined (typically AttributeError or TypeError).

When implementing a class that emulates any built-in type, it is important that the emulation only be implemented to the degree that it makes sense for the object being modelled. For example, some sequences may work well with retrieval of individual elements, but extracting a slice may not make sense. (One example of this is the NodeList interface in the W3C's Document Object Model.)

3.3.1 Basic customization

```
object.__new__(cls[,...])
```

Called to create a new instance of class *cls*. __new__() is a static method (special-cased so you need not declare it as such) that takes the class of which an instance was requested as its first argument. The remaining arguments are those passed to the object constructor expression (the call to the class). The return value of __new__() should be the new object instance (usually an instance of *cls*).

Typical implementations create a new instance of the class by invoking the superclass's __new__() method using super(currentclass, cls).__new__(cls[, ...]) with appropriate arguments and then modifying the newly-created instance as necessary before returning it.

If $_new__()$ returns an instance of cls, then the new instance's $_init__()$ method will be invoked like $_init__(self[, ...])$, where self is the new instance and the remaining arguments are the same as were passed to $_new__()$.

If __new__ () does not return an instance of *cls*, then the new instance's __init__ () method will not be invoked.

__new__() is intended mainly to allow subclasses of immutable types (like int, str, or tuple) to customize instance creation. It is also commonly overridden in custom metaclasses in order to customize class creation.

object.__init__(self[,...])

Called when the instance is created. The arguments are those passed to the class constructor expression. If a base class has an __init__() method, the derived class's __init__() method, if any, must explicitly call it to ensure proper initialization of the base class part of the instance; for example: BaseClass.__init__(self, [args...]). As a special constraint on constructors, no value may be returned; doing so will cause a TypeError to be raised at runtime.

object.__**del**__(*self*)

Called when the instance is about to be destroyed. This is also called a destructor. If a base class has a __del__() method, the derived class's __del__() method, if any, must explicitly call it to ensure proper deletion of the base class part of the instance. Note that it is possible (though not recommended!) for the __del__() method to postpone destruction of the instance by creating a new reference to it. It may then be called at a later time when this new reference is deleted. It is not guaranteed that __del__() methods are called for objects that still exist when the interpreter exits.

Note: del x doesn't directly call x.__del__() — the former decrements the reference count for x by one, and the latter is only called when x's reference count reaches zero. Some common situations that may prevent the reference count of an object from going to zero include: circular references between objects (e.g., a doubly-linked list or a tree data structure with parent and child pointers); a reference to the object on the stack frame of a function that caught an exception (the traceback stored in sys.exc_info() [2] keeps the stack frame alive); or a reference to the object on the stack frame that raised an unhandled exception in interactive mode (the traceback stored in sys.last_traceback keeps the stack frame alive). The first situation can only be remedied by explicitly breaking the cycles; the latter two situations can be resolved by storing None in sys.last_traceback. Circular references which are garbage are detected when the option cycle detector is enabled (it's on by default), but can only be cleaned up if there are no Python-level __del__() methods involved. Refer to the documentation for the gc module for more information about how __del__() methods are handled by the cycle detector, particularly the description of the garbage value.

Warning: Due to the precarious circumstances under which __del__() methods are invoked, exceptions that occur during their execution are ignored, and a warning is printed to sys.stderr instead. Also, when __del__() is invoked in response to a module being deleted (e.g., when execution of the program is done), other globals referenced by the __del__() method may already have been deleted or in the process of being torn down (e.g. the import machinery shutting down). For this reason, __del__() methods should do the absolute minimum needed to maintain external invariants. Starting with version 1.5, Python guarantees that globals whose name begins with a single underscore are deleted from their module before other globals are deleted; if no other references to such globals exist, this may help in assuring that imported modules are still available at the time when the __del__() method is called.

object.__repr__(self)

Called by the repr() built-in function to compute the "official" string representation of an object. If at all possible, this should look like a valid Python expression that could be used to recreate an object with the same value (given an appropriate environment). If this is not possible, a string of the form <...some useful description...> should be returned. The return value must be a string object. If a class defines __repr__() but not __str__(), then __repr__() is also used when an "informal" string representation of instances of that class is required.

This is typically used for debugging, so it is important that the representation is information-rich and unambiguous.

object.__**str**__(*self*)

Called by the str() built-in function and by the print() function to compute the "informal" string representation of an object. This differs from __repr__() in that it does not have to be a valid Python expression: a more convenient or concise representation may be used instead. The return value must be a string object.

object.__bytes__(self)

Called by bytes () to compute a byte-string representation of an object. This should return a bytes object.

```
object.__format__(self, format_spec)
```

Called by the format () built-in function (and by extension, the str.format() method of class str) to produce a "formatted" string representation of an object. The format_spec argument is a string that contains a description of the formatting options desired. The interpretation of the format_spec argument is up to the type implementing __format__(), however most classes will either delegate formatting to one of the built-in types, or use a similar formatting option syntax.

See *formatspec* for a description of the standard formatting syntax.

The return value must be a string object.

```
object.__lt__ (self, other)
object.__le__ (self, other)
object.__eq__ (self, other)
object.__ne__ (self, other)
object.__gt__ (self, other)
object.__ge__ (self, other)
```

These are the so-called "rich comparison" methods. The correspondence between operator symbols and method names is as follows: x < y calls $x . __lt __(y)$, x <= y calls $x . __lt __(y)$, x == y calls $x . __lt __(y)$, x == y calls $x . __lt __(y)$, x > y calls $x . __lt __(y)$, and x >= y calls $x . __lt __(y)$.

A rich comparison method may return the singleton NotImplemented if it does not implement the operation for a given pair of arguments. By convention, False and True are returned for a successful comparison. However, these methods can return any value, so if the comparison operator is used in a Boolean context (e.g., in the condition of an if statement), Python will call bool () on the value to determine if the result is true or false.

There are no implied relationships among the comparison operators. The truth of x==y does not imply that x!=y is false. Accordingly, when defining $__eq__()$, one should also define $__ne__()$ so that the operators will behave as expected. See the paragraph on $__hash__()$ for some important notes on creating *hashable* objects which support custom comparison operations and are usable as dictionary keys.

There are no swapped-argument versions of these methods (to be used when the left argument does not support the operation but the right argument does); rather, $__lt_{__}()$ and $__gt_{__}()$ are each other's reflection, $__le_{__}()$ and $__ge_{__}()$ are each other's reflection, and $__eq_{__}()$ and $__ne_{__}()$ are their own reflection.

Arguments to rich comparison methods are never coerced.

To automatically generate ordering operations from a single root operation, see functools.total_ordering().

```
object.__hash__(self)
```

Called by built-in function hash() and for operations on members of hashed collections including set, frozenset, and dict. __hash__() should return an integer. The only required property is that objects which compare equal have the same hash value; it is advised to somehow mix together (e.g. using exclusive or) the hash values for the components of the object that also play a part in comparison of objects.

If a class does not define an __eq__() method it should not define a __hash__() operation either; if it defines __eq__() but not __hash__(), its instances will not be usable as items in hashable collections. If a class defines mutable objects and implements an __eq__() method, it should not implement __hash__(), since the implementation of hashable collections requires that a key's hash value is immutable (if the object's hash value changes, it will be in the wrong hash bucket).

User-defined classes have $_eq_()$ and $_hash_()$ methods by default; with them, all objects compare unequal (except with themselves) and $x._hash_()$ returns an appropriate value such that x == y implies both that x is y and hash(x) == hash(y).

A class that overrides __eq_ () and does not define __hash__ () will have its __hash__ () implicitly set to None. When the __hash__ () method of a class is None, instances of the class will raise an appropriate TypeError when a program attempts to retrieve their hash value, and will also be correctly identified as unhashable when checking isinstance (obj, collections.Hashable).

If a class that overrides __eq_ () needs to retain the implementation of __hash__() from a parent class, the interpreter must be told this explicitly by setting __hash__ = <ParentClass>.__hash__.

If a class that does not override __eq_ () wishes to suppress hash support, it should include __hash__ = None in the class definition. A class which defines its own __hash__ () that explicitly raises a TypeError would be incorrectly identified as hashable by an isinstance (obj, collections.Hashable) call.

Note: By default, the __hash__() values of str, bytes and datetime objects are "salted" with an unpredictable random value. Although they remain constant within an individual Python process, they are not predictable between repeated invocations of Python.

This is intended to provide protection against a denial-of-service caused by carefully-chosen inputs that exploit the worst case performance of a dict insertion, O(n^2) complexity. See http://www.ocert.org/advisories/ocert-2011-003.html for details.

Changing hash values affects the iteration order of dicts, sets and other mappings. Python has never made guarantees about this ordering (and it typically varies between 32-bit and 64-bit builds).

See also PYTHONHASHSEED.

Changed in version 3.3: Hash randomization is enabled by default.

```
object.__bool__(self)
```

Called to implement truth value testing and the built-in operation bool(); should return False or True. When this method is not defined, __len__() is called, if it is defined, and the object is considered true if its result is nonzero. If a class defines neither __len__() nor __bool__(), all its instances are considered true.

3.3.2 Customizing attribute access

The following methods can be defined to customize the meaning of attribute access (use of, assignment to, or deletion of x.name) for class instances.

```
object.__getattr__(self, name)
```

Called when an attribute lookup has not found the attribute in the usual places (i.e. it is not an instance attribute nor is it found in the class tree for self). name is the attribute name. This method should return the (computed) attribute value or raise an AttributeError exception.

Note that if the attribute is found through the normal mechanism, __getattr__() is not called. (This is an intentional asymmetry between __getattr__() and __setattr__().) This is done both for efficiency reasons and because otherwise __getattr__() would have no way to access other attributes of the instance. Note that at least for instance variables, you can fake total control by not inserting any values in the instance attribute dictionary (but instead inserting them in another object). See the __getattribute__() method below for a way to actually get total control over attribute access.

```
object.__getattribute__(self, name)
```

Called unconditionally to implement attribute accesses for instances of the class. If the class also defines __getattr__(), the latter will not be called unless __getattribute__() either calls it explicitly or raises an AttributeError. This method should return the (computed) attribute value or raise an AttributeError exception. In order to avoid infinite recursion in this method, its implementation should always call the base class method with the same name to access any attributes it needs, for example, object.__getattribute__(self, name).

Note: This method may still be bypassed when looking up special methods as the result of implicit invocation via language syntax or built-in functions. See *Special method lookup*.

```
object. setattr (self, name, value)
```

Called when an attribute assignment is attempted. This is called instead of the normal mechanism (i.e. store the value in the instance dictionary). *name* is the attribute name, *value* is the value to be assigned to it.

```
If __setattr__() wants to assign to an instance attribute, it should call the base class method with the same name, for example, object.__setattr__(self, name, value).
```

```
object.__delattr__(self, name)
```

Like __setattr__() but for attribute deletion instead of assignment. This should only be implemented if del obj.name is meaningful for the object.

```
object.__dir__(self)
```

Called when dir() is called on the object. A sequence must be returned. dir() converts the returned sequence to a list and sorts it.

Implementing Descriptors

The following methods only apply when an instance of the class containing the method (a so-called *descriptor* class) appears in an *owner* class (the descriptor must be in either the owner's class dictionary or in the class dictionary for one of its parents). In the examples below, "the attribute" refers to the attribute whose name is the key of the property in the owner class' ___dict___.

```
object.__get__ (self, instance, owner)
```

Called to get the attribute of the owner class (class attribute access) or of an instance of that class (instance attribute access). *owner* is always the owner class, while *instance* is the instance that the attribute was accessed through, or None when the attribute is accessed through the *owner*. This method should return the (computed) attribute value or raise an AttributeError exception.

```
object.__set__(self, instance, value)
```

Called to set the attribute on an instance *instance* of the owner class to a new value, value.

```
object.__delete__(self, instance)
```

Called to delete the attribute on an instance *instance* of the owner class.

Invoking Descriptors

In general, a descriptor is an object attribute with "binding behavior", one whose attribute access has been overridden by methods in the descriptor protocol: __get__(), __set__(), and __delete__(). If any of those methods are defined for an object, it is said to be a descriptor.

The default behavior for attribute access is to get, set, or delete the attribute from an object's dictionary. For instance, a.x has a lookup chain starting with a.__dict__['x'], then type(a).__dict__['x'], and continuing through the base classes of type(a) excluding metaclasses.

However, if the looked-up value is an object defining one of the descriptor methods, then Python may override the default behavior and invoke the descriptor method instead. Where this occurs in the precedence chain depends on which descriptor methods were defined and how they were called.

The starting point for descriptor invocation is a binding, a.x. How the arguments are assembled depends on a:

Direct Call The simplest and least common call is when user code directly invokes a descriptor method: $x \cdot get(a)$.

```
Instance Binding If binding to an object instance, a.x is transformed into the call: type(a).__dict__['x'].__get__(a, type(a)).
```

Class Binding If binding to a class, A.x is transformed into the call: A.__dict__['x'].__get__(None, A).

Super Binding If a is an instance of super, then the binding super(B, obj).m() searches obj.__class__._mro__ for the base class A immediately preceding B and then invokes the descriptor with the call: A.__dict__['m'].__get__(obj, obj.__class__).

For instance bindings, the precedence of descriptor invocation depends on the which descriptor methods are defined. A descriptor can define any combination of __get__(), __set__() and __delete__(). If it does not define __get__(), then accessing the attribute will return the descriptor object itself unless there is a value in the object's instance dictionary. If the descriptor defines __set__() and/or __delete__(), it is a data descriptor; if it defines neither, it is a non-data descriptor. Normally, data descriptors define both __get__() and __set__(), while non-data descriptors have just the __get__() method. Data descriptors with __set__() and __get__() defined always override a redefinition in an instance dictionary. In contrast, non-data descriptors can be overridden by instances.

Python methods (including staticmethod() and classmethod()) are implemented as non-data descriptors. Accordingly, instances can redefine and override methods. This allows individual instances to acquire behaviors that differ from other instances of the same class.

The property () function is implemented as a data descriptor. Accordingly, instances cannot override the behavior of a property.

slots

By default, instances of classes have a dictionary for attribute storage. This wastes space for objects having very few instance variables. The space consumption can become acute when creating large numbers of instances.

The default can be overridden by defining <u>__slots__</u> in a class definition. The <u>__slots__</u> declaration takes a sequence of instance variables and reserves just enough space in each instance to hold a value for each variable. Space is saved because <u>__dict__</u> is not created for each instance.

```
object.__slots_
```

This class variable can be assigned a string, iterable, or sequence of strings with variable names used by instances. If defined in a class, __slots__ reserves space for the declared variables and prevents the automatic creation of __dict__ and __weakref__ for each instance.

Notes on using slots

- When inheriting from a class without <u>__slots__</u>, the <u>__dict__</u> attribute of that class will always be accessible, so a <u>__slots__</u> definition in the subclass is meaningless.
- Without a __dict__ variable, instances cannot be assigned new variables not listed in the __slots__ definition. Attempts to assign to an unlisted variable name raises AttributeError. If dynamic assignment of new variables is desired, then add '__dict__' to the sequence of strings in the __slots__ declaration.
- Without a __weakref__ variable for each instance, classes defining __slots__ do not support weak references to
 its instances. If weak reference support is needed, then add ' __weakref__' to the sequence of strings in the
 __slots__ declaration.
- __slots__ are implemented at the class level by creating descriptors (*Implementing Descriptors*) for each variable name. As a result, class attributes cannot be used to set default values for instance variables defined by __slots__; otherwise, the class attribute would overwrite the descriptor assignment.

- The action of a __slots__ declaration is limited to the class where it is defined. As a result, subclasses will have a __dict__ unless they also define __slots__ (which must only contain names of any additional slots).
- If a class defines a slot also defined in a base class, the instance variable defined by the base class slot is inaccessible (except by retrieving its descriptor directly from the base class). This renders the meaning of the program undefined. In the future, a check may be added to prevent this.
- Nonempty __slots__ does not work for classes derived from "variable-length" built-in types such as int, str and tuple.
- Any non-string iterable may be assigned to <u>__slots__</u>. Mappings may also be used; however, in the future, special meaning may be assigned to the values corresponding to each key.
- __class__ assignment works only if both classes have the same __slots__.

3.3.3 Customizing class creation

By default, classes are constructed using type (). The class body is executed in a new namespace and the class name is bound locally to the result of type (name, bases, namespace).

The class creation process can be customised by passing the metaclass keyword argument in the class definition line, or by inheriting from an existing class that included such an argument. In the following example, both MyClass and MySubclass are instances of Meta:

```
class Meta(type):
    pass

class MyClass(metaclass=Meta):
    pass

class MySubclass(MyClass):
    pass
```

Any other keyword arguments that are specified in the class definition are passed through to all metaclass operations described below.

When a class definition is executed, the following steps occur:

- the appropriate metaclass is determined
- the class namespace is prepared
- the class body is executed
- the class object is created

Determining the appropriate metaclass

The appropriate metaclass for a class definition is determined as follows:

- if no bases and no explicit metaclass are given, then type () is used
- if an explicit metaclass is given and it is *not* an instance of type (), then it is used directly as the metaclass
- if an instance of type () is given as the explicit metaclass, or bases are defined, then the most derived metaclass is used

The most derived metaclass is selected from the explicitly specified metaclass (if any) and the metaclasses (i.e. type(cls)) of all specified base classes. The most derived metaclass is one which is a subtype of *all* of these

candidate metaclasses. If none of the candidate metaclasses meets that criterion, then the class definition will fail with TypeError.

Preparing the class namespace

Once the appropriate metaclass has been identified, then the class namespace is prepared. If the metaclass has a __prepare__ attribute, it is called as namespace = metaclass.__prepare__ (name, bases, **kwds) (where the additional keyword arguments, if any, come from the class definition).

If the metaclass has no __prepare__ attribute, then the class namespace is initialised as an empty dict () instance.

See Also:

PEP 3115 - Metaclasses in Python 3000 Introduced the __prepare__ namespace hook

Executing the class body

The class body is executed (approximately) as exec (body, globals(), namespace). The key difference from a normal call to exec() is that lexical scoping allows the class body (including any methods) to reference names from the current and outer scopes when the class definition occurs inside a function.

However, even when the class definition occurs inside the function, methods defined inside the class still cannot see names defined at the class scope. Class variables must be accessed through the first parameter of instance or class methods, and cannot be accessed at all from static methods.

Creating the class object

Once the class namespace has been populated by executing the class body, the class object is created by calling metaclass (name, bases, namespace, **kwds) (the additional keywords passed here are the same as those passed to __prepare__).

This class object is the one that will be referenced by the zero-argument form of <code>super()</code>. __class__ is an implicit closure reference created by the compiler if any methods in a class body refer to either __class__ or <code>super</code>. This allows the zero argument form of <code>super()</code> to correctly identify the class being defined based on lexical scoping, while the class or instance that was used to make the current call is identified based on the first argument passed to the method.

After the class object is created, it is passed to the class decorators included in the class definition (if any) and the resulting object is bound in the local namespace as the defined class.

See Also:

PEP 3135 - New super Describes the implicit __class__ closure reference

Metaclass example

The potential uses for metaclasses are boundless. Some ideas that have been explored include logging, interface checking, automatic delegation, automatic property creation, proxies, frameworks, and automatic resource locking/synchronization.

Here is an example of a metaclass that uses an collections. OrderedDict to remember the order that class members were defined:

```
class OrderedClass(type):
```

@classmethod

```
def __prepare__ (metacls, name, bases, **kwds):
    return collections.OrderedDict()

def __new__ (cls, name, bases, namespace, **kwds):
    result = type.__new__ (cls, name, bases, dict(namespace))
    result.members = tuple(namespace)
    return result

class A(metaclass=OrderedClass):
    def one(self): pass
    def two(self): pass
    def three(self): pass
    def four(self): pass

//_ module__', 'one', 'two', 'three', 'four')
```

When the class definition for A gets executed, the process begins with calling the metaclass's $_prepare_()$ method which returns an empty collections.OrderedDict. That mapping records the methods and attributes of A as they are defined within the body of the class statement. Once those definitions are executed, the ordered dictionary is fully populated and the metaclass's $_new_()$ method gets invoked. That method builds the new type and it saves the ordered dictionary keys in an attribute called members.

3.3.4 Customizing instance and subclass checks

The following methods are used to override the default behavior of the isinstance() and issubclass() built-in functions.

In particular, the metaclass abc. ABCMeta implements these methods in order to allow the addition of Abstract Base Classes (ABCs) as "virtual base classes" to any class or type (including built-in types), including other ABCs.

```
class.__instancecheck__(self, instance)
```

Return true if *instance* should be considered a (direct or indirect) instance of *class*. If defined, called to implement isinstance (instance, class).

```
class.__subclasscheck__(self, subclass)
```

Return true if *subclass* should be considered a (direct or indirect) subclass of *class*. If defined, called to implement issubclass (subclass, class).

Note that these methods are looked up on the type (metaclass) of a class. They cannot be defined as class methods in the actual class. This is consistent with the lookup of special methods that are called on instances, only in this case the instance is itself a class.

See Also:

PEP 3119 - Introducing Abstract Base Classes Includes the specification for customizing isinstance() and issubclass() behavior through __instancecheck__() and __subclasscheck__(), with motivation for this functionality in the context of adding Abstract Base Classes (see the abc module) to the language.

3.3.5 Emulating callable objects

```
object.__call__(self[, args...])

Called when the instance is "called" as a function; if this method is defined, x(arg1, arg2, ...) is a shorthand for x.__call__(arg1, arg2, ...).
```

3.3.6 Emulating container types

The following methods can be defined to implement container objects. Containers usually are sequences (such as lists or tuples) or mappings (like dictionaries), but can represent other containers as well. The first set of methods is used either to emulate a sequence or to emulate a mapping; the difference is that for a sequence, the allowable keys should be the integers k for which $0 \le k \le N$ where N is the length of the sequence, or slice objects, which define a range of items. It is also recommended that mappings provide the methods keys (), values (), items (), get(), clear(), setdefault(), pop(), popitem(), copy(), and update() behaving similar to those for Python's standard dictionary objects. The collections module provides a MutableMapping abstract base class to help create those methods from a base set of __qetitem__(), __setitem__(), __delitem__(), and keys(). Mutable sequences should provide methods append(), count(), index(), extend(), insert(), pop(), remove(), reverse() and sort(), like Python standard list objects. Finally, sequence types should implement addition (meaning concatenation) and multiplication (meaning repetition) by defining the methods __add__(), __radd__(), __iadd__(), __mul__(), __rmul__() and __imul__() described below; they should not define other numerical operators. It is recommended that both mappings and sequences implement the contains () method to allow efficient use of the in operator; for mappings, in should search the mapping's keys; for sequences, it should search through the values. It is further recommended that both mappings and sequences implement the iter () method to allow efficient iteration through the container; for mappings, iter () should be the same as keys (); for sequences, it should iterate through the values.

```
object.__len__(self)
```

Called to implement the built-in function len(). Should return the length of the object, an integer >= 0. Also, an object that doesn't define a __bool__() method and whose __len__() method returns zero is considered to be false in a Boolean context.

```
object.__length_hint__(self)
```

Called to implement operator.length_hint(). Should return an estimated length for the object (which may be greater or less than the actual length). The length must be an integer \geq 0. This method is purely an optimization and is never required for correctness. New in version 3.4.

Note: Slicing is done exclusively with the following three methods. A call like

```
a[1:2] = b
```

is translated to

```
a[slice(1, 2, None)] = b
```

and so forth. Missing slice items are always filled in with None.

```
object.__getitem__(self, key)
```

Called to implement evaluation of self[key]. For sequence types, the accepted keys should be integers and slice objects. Note that the special interpretation of negative indexes (if the class wishes to emulate a sequence type) is up to the __getitem__() method. If key is of an inappropriate type, TypeError may be raised; if of a value outside the set of indexes for the sequence (after any special interpretation of negative values), IndexError should be raised. For mapping types, if key is missing (not in the container), KeyError should be raised.

Note: for loops expect that an IndexError will be raised for illegal indexes to allow proper detection of the end of the sequence.

```
object.__setitem__(self, key, value)
```

Called to implement assignment to self[key]. Same note as for $_getitem_$ (). This should only be implemented for mappings if the objects support changes to the values for keys, or if new keys can be added, or for sequences if elements can be replaced. The same exceptions should be raised for improper key values as for the $_getitem_$ () method.

```
object. delitem (self, key)
```

Called to implement deletion of <code>self[key]</code>. Same note as for <code>__getitem__()</code>. This should only be implemented for mappings if the objects support removal of keys, or for sequences if elements can be removed from the sequence. The same exceptions should be raised for improper <code>key</code> values as for the <code>__getitem__()</code> method.

```
object. iter (self)
```

This method is called when an iterator is required for a container. This method should return a new iterator object that can iterate over all the objects in the container. For mappings, it should iterate over the keys of the container, and should also be made available as the method keys ().

Iterator objects also need to implement this method; they are required to return themselves. For more information on iterator objects, see *typeiter*.

```
object.__reversed__(self)
```

Called (if present) by the reversed () built-in to implement reverse iteration. It should return a new iterator object that iterates over all the objects in the container in reverse order.

```
If the __reversed__() method is not provided, the reversed() built-in will fall back to using the sequence protocol (__len__() and __getitem__()). Objects that support the sequence protocol should only provide __reversed__() if they can provide an implementation that is more efficient than the one provided by reversed().
```

The membership test operators (in and not in) are normally implemented as an iteration through a sequence. However, container objects can supply the following special method with a more efficient implementation, which also does not require the object be a sequence.

```
object.__contains__(self, item)
```

Called to implement membership test operators. Should return true if *item* is in *self*, false otherwise. For mapping objects, this should consider the keys of the mapping rather than the values or the key-item pairs.

For objects that don't define __contains__(), the membership test first tries iteration via __iter__(), then the old sequence iteration protocol via __getitem__(), see this section in the language reference.

3.3.7 Emulating numeric types

The following methods can be defined to emulate numeric objects. Methods corresponding to operations that are not supported by the particular kind of number implemented (e.g., bitwise operations for non-integral numbers) should be left undefined.

```
object.__add__ (self, other)
object.__sub__ (self, other)
object.__mul__ (self, other)
object.__truediv__ (self, other)
object.__floordiv__ (self, other)
object.__mod__ (self, other)
object.__divmod__ (self, other)
object.__pow__ (self, other[, modulo])
object.__lshift__ (self, other)
object.__rshift__ (self, other)
object.__and__ (self, other)
object.__xor__ (self, other)
object.__or__ (self, other)
```

These methods are called to implement the binary arithmetic operations $(+, -, *, /, //, %, divmod(), pow(), **, <<, >>, &, ^, |)$. For instance, to evaluate the expression x + y, where x is an instance of a class that has an __add__() method, x.__add__(y) is called. The __divmod__() method should be the equivalent to using __floordiv__() and __mod__(); it should not be related to __truediv__(). Note

that __pow__() should be defined to accept an optional third argument if the ternary version of the built-in pow() function is to be supported.

If one of those methods does not support the operation with the supplied arguments, it should return NotImplemented.

```
object.__radd__ (self, other)
object.__rsub__ (self, other)
object.__rmul__ (self, other)
object.__rtruediv__ (self, other)
object.__rfloordiv__ (self, other)
object.__rmod__ (self, other)
object.__rdivmod__ (self, other)
object.__rlshift__ (self, other)
object.__rshift__ (self, other)
object.__rand__ (self, other)
object.__rand__ (self, other)
object.__rxor__ (self, other)
object.__ror__ (self, other)
```

These methods are called to implement the binary arithmetic operations $(+, -, *, /, //, %, divmod(), pow(), **, <<, >>, &, ^, |)$ with reflected (swapped) operands. These functions are only called if the left operand does not support the corresponding operation and the operands are of different types. ² For instance, to evaluate the expression x - y, where y is an instance of a class that has an $__rsub__()$ method, y. $__rsub__(x)$ is called if x. $__sub__(y)$ returns *NotImplemented*.

Note that ternary pow () will not try calling __rpow__() (the coercion rules would become too complicated).

Note: If the right operand's type is a subclass of the left operand's type and that subclass provides the reflected method for the operation, this method will be called before the left operand's non-reflected method. This behavior allows subclasses to override their ancestors' operations.

```
object.__iadd__ (self, other)
object.__isub__ (self, other)
object.__imul__ (self, other)
object.__itruediv__ (self, other)
object.__ifloordiv__ (self, other)
object.__imod__ (self, other)
object.__ipow__ (self, other[, modulo])
object.__ilshift__ (self, other)
object.__irshift__ (self, other)
object.__iand__ (self, other)
object.__ixor__ (self, other)
object.__ior__ (self, other)
```

These methods are called to implement the augmented arithmetic assignments (+=, -=, *=, /=, //=, %=, **=, <<=, >>=, &=, ^=, |=). These methods should attempt to do the operation in-place (modifying self) and return the result (which could be, but does not have to be, self). If a specific method is not defined, the augmented assignment falls back to the normal methods. For instance, to execute the statement x += y, where x is an instance of a class that has an $_iadd__()$ method, $x._iadd__(y)$ is called. If x is an instance of a class that does not define a $_iadd__()$ method, $x._iadd__(y)$ and $y._iadd__(x)$ are considered, as with the evaluation of x + y.

```
object.__neg__(self)
object.__pos__(self)
```

 $^{^2}$ For operands of the same type, it is assumed that if the non-reflected method (such as $__add__()$) fails the operation is not supported, which is why the reflected method is not called.

```
object.__abs__(self)
object.__invert__(self)
     Called to implement the unary arithmetic operations (-, +, abs()) and \sim).
object.__complex__(self)
object.__int__(self)
object. float (self)
object. round (self, n)
     Called to implement the built-in functions complex(), int(), float() and round(). Should return a
     value of the appropriate type.
```

```
object.__index__(self)
```

Called to implement operator.index(). Also called whenever Python needs an integer object (such as in slicing, or in the built-in bin (), hex () and oct () functions). Must return an integer.

3.3.8 With Statement Context Managers

A context manager is an object that defines the runtime context to be established when executing a with statement. The context manager handles the entry into, and the exit from, the desired runtime context for the execution of the block of code. Context managers are normally invoked using the with statement (described in section The with statement), but can also be used by directly invoking their methods.

Typical uses of context managers include saving and restoring various kinds of global state, locking and unlocking resources, closing opened files, etc.

For more information on context managers, see *typecontextmanager*.

```
object.__enter__(self)
```

Enter the runtime context related to this object. The with statement will bind this method's return value to the target(s) specified in the as clause of the statement, if any.

```
object.__exit__(self, exc_type, exc_value, traceback)
```

Exit the runtime context related to this object. The parameters describe the exception that caused the context to be exited. If the context was exited without an exception, all three arguments will be None.

If an exception is supplied, and the method wishes to suppress the exception (i.e., prevent it from being propagated), it should return a true value. Otherwise, the exception will be processed normally upon exit from this method.

Note that exit () methods should not reraise the passed-in exception; this is the caller's responsibility.

See Also:

PEP 0343 - The "with" statement The specification, background, and examples for the Python with statement.

3.3.9 Special method lookup

For custom classes, implicit invocations of special methods are only guaranteed to work correctly if defined on an object's type, not in the object's instance dictionary. That behaviour is the reason why the following code raises an exception:

```
>>> class C:
        pass
>>> C = C()
>>> c.__len__ = lambda: 5
>>> len(c)
Traceback (most recent call last):
```

```
File "<stdin>", line 1, in <module>
TypeError: object of type 'C' has no len()
```

The rationale behind this behaviour lies with a number of special methods such as __hash__() and __repr__() that are implemented by all objects, including type objects. If the implicit lookup of these methods used the conventional lookup process, they would fail when invoked on the type object itself:

```
>>> 1 .__hash__() == hash(1)
True
>>> int.__hash__() == hash(int)
Traceback (most recent call last):
  File "<stdin>", line 1, in <module>
TypeError: descriptor '__hash__' of 'int' object needs an argument
```

Incorrectly attempting to invoke an unbound method of a class in this way is sometimes referred to as 'metaclass confusion', and is avoided by bypassing the instance when looking up special methods:

```
>>> type(1).__hash__(1) == hash(1)
True
>>> type(int).__hash__(int) == hash(int)
True
```

In addition to bypassing any instance attributes in the interest of correctness, implicit special method lookup generally also bypasses the __getattribute__() method even of the object's metaclass:

```
>>> class Meta(type):
       def __getattribute__(*args):
          print("Metaclass getattribute invoked")
. . .
          return type.__getattribute__(*args)
>>> class C(object, metaclass=Meta):
        def __len__(self):
            return 10
        def __getattribute__(*args):
            print("Class getattribute invoked")
            return object.__getattribute__(*args)
. . .
>>> C = C()
                                 # Explicit lookup via instance
>>> c.__len__()
Class getattribute invoked
                                 # Explicit lookup via type
>>> type(c).__len__(c)
Metaclass getattribute invoked
>>> len(c)
                                 # Implicit lookup
```

Bypassing the <u>__getattribute___()</u> machinery in this fashion provides significant scope for speed optimisations within the interpreter, at the cost of some flexibility in the handling of special methods (the special method *must* be set on the class object itself in order to be consistently invoked by the interpreter).

EXECUTION MODEL

4.1 Naming and binding

Names refer to objects. Names are introduced by name binding operations. Each occurrence of a name in the program text refers to the *binding* of that name established in the innermost function block containing the use.

A *block* is a piece of Python program text that is executed as a unit. The following are blocks: a module, a function body, and a class definition. Each command typed interactively is a block. A script file (a file given as standard input to the interpreter or specified on the interpreter command line the first argument) is a code block. A script command (a command specified on the interpreter command line with the '-c' option) is a code block. The string argument passed to the built-in functions eval() and exec() is a code block.

A code block is executed in an *execution frame*. A frame contains some administrative information (used for debugging) and determines where and how execution continues after the code block's execution has completed.

A *scope* defines the visibility of a name within a block. If a local variable is defined in a block, its scope includes that block. If the definition occurs in a function block, the scope extends to any blocks contained within the defining one, unless a contained block introduces a different binding for the name. The scope of names defined in a class block is limited to the class block; it does not extend to the code blocks of methods – this includes comprehensions and generator expressions since they are implemented using a function scope. This means that the following will fail:

```
class A:
    a = 42
    b = list(a + i for i in range(10))
```

When a name is used in a code block, it is resolved using the nearest enclosing scope. The set of all such scopes visible to a code block is called the block's *environment*.

If a name is bound in a block, it is a local variable of that block, unless declared as nonlocal. If a name is bound at the module level, it is a global variable. (The variables of the module code block are local and global.) If a variable is used in a code block but not defined there, it is a *free variable*.

When a name is not found at all, a NameError exception is raised. If the name refers to a local variable that has not been bound, a UnboundLocalError exception is raised. UnboundLocalError is a subclass of NameError.

The following constructs bind names: formal parameters to functions, import statements, class and function definitions (these bind the class or function name in the defining block), and targets that are identifiers if occurring in an assignment, for loop header, or after as in a with statement or except clause. The import statement of the form from ... import * binds all names defined in the imported module, except those beginning with an underscore. This form may only be used at the module level.

A target occurring in a del statement is also considered bound for this purpose (though the actual semantics are to unbind the name).

Each assignment or import statement occurs within a block defined by a class or function definition or at the module level (the top-level code block).

If a name binding operation occurs anywhere within a code block, all uses of the name within the block are treated as references to the current block. This can lead to errors when a name is used within a block before it is bound. This rule is subtle. Python lacks declarations and allows name binding operations to occur anywhere within a code block. The local variables of a code block can be determined by scanning the entire text of the block for name binding operations.

If the global statement occurs within a block, all uses of the name specified in the statement refer to the binding of that name in the top-level namespace. Names are resolved in the top-level namespace by searching the global namespace, i.e. the namespace of the module containing the code block, and the builtins namespace, the namespace of the module builtins. The global namespace is searched first. If the name is not found there, the builtins namespace is searched. The global statement must precede all uses of the name.

The builtins namespace associated with the execution of a code block is actually found by looking up the name __builtins__ in its global namespace; this should be a dictionary or a module (in the latter case the module's dictionary is used). By default, when in the __main__ module, __builtins__ is the built-in module builtins; when in any other module, __builtins__ is an alias for the dictionary of the builtins module itself. __builtins__ can be set to a user-created dictionary to create a weak form of restricted execution.

CPython implementation detail: Users should not touch __builtins__; it is strictly an implementation detail. Users wanting to override values in the builtins namespace should import the builtins module and modify its attributes appropriately.

The namespace for a module is automatically created the first time a module is imported. The main module for a script is always called __main__.

The global statement has the same scope as a name binding operation in the same block. If the nearest enclosing scope for a free variable contains a global statement, the free variable is treated as a global.

A class definition is an executable statement that may use and define names. These references follow the normal rules for name resolution. The namespace of the class definition becomes the attribute dictionary of the class. Names defined at the class scope are not visible in methods.

4.1.1 Interaction with dynamic features

There are several cases where Python statements are illegal when used in conjunction with nested scopes that contain free variables.

If a variable is referenced in an enclosing scope, it is illegal to delete the name. An error will be reported at compile time.

If the wild card form of import — import * — is used in a function and the function contains or is a nested block with free variables, the compiler will raise a SyntaxError.

The eval() and exec() functions do not have access to the full environment for resolving names. Names may be resolved in the local and global namespaces of the caller. Free variables are not resolved in the nearest enclosing namespace, but in the global namespace. 1 The exec() and eval() functions have optional arguments to override the global and local namespace. If only one namespace is specified, it is used for both.

4.2 Exceptions

Exceptions are a means of breaking out of the normal flow of control of a code block in order to handle errors or other exceptional conditions. An exception is *raised* at the point where the error is detected; it may be *handled* by

¹ This limitation occurs because the code that is executed by these operations is not available at the time the module is compiled.

the surrounding code block or by any code block that directly or indirectly invoked the code block where the error occurred.

The Python interpreter raises an exception when it detects a run-time error (such as division by zero). A Python program can also explicitly raise an exception with the raise statement. Exception handlers are specified with the try ... except statement. The finally clause of such a statement can be used to specify cleanup code which does not handle the exception, but is executed whether an exception occurred or not in the preceding code.

Python uses the "termination" model of error handling: an exception handler can find out what happened and continue execution at an outer level, but it cannot repair the cause of the error and retry the failing operation (except by reentering the offending piece of code from the top).

When an exception is not handled at all, the interpreter terminates execution of the program, or returns to its interactive main loop. In either case, it prints a stack backtrace, except when the exception is SystemExit.

Exceptions are identified by class instances. The except clause is selected depending on the class of the instance: it must reference the class of the instance or a base class thereof. The instance can be received by the handler and can carry additional information about the exceptional condition.

Note: Exception messages are not part of the Python API. Their contents may change from one version of Python to the next without warning and should not be relied on by code which will run under multiple versions of the interpreter.

See also the description of the try statement in section *The try statement* and raise statement in section *The raise statement*.

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CHAPTER

FIVE

THE IMPORT SYSTEM

Python code in one *module* gains access to the code in another module by the process of *importing* it. The import statement is the most common way of invoking the import machinery, but it is not the only way. Functions such as importlib.import_module() and built-in __import__() can also be used to invoke the import machinery.

The import statement combines two operations; it searches for the named module, then it binds the results of that search to a name in the local scope. The search operation of the import statement is defined as a call to the __import___() function, with the appropriate arguments. The return value of __import___() is used to perform the name binding operation of the import statement. See the import statement for the exact details of that name binding operation.

A direct call to __import__() performs only the module search and, if found, the module creation operation. While certain side-effects may occur, such as the importing of parent packages, and the updating of various caches (including sys.modules), only the import statement performs a name binding operation.

When calling __import__() as part of an import statement, the import system first checks the module global namespace for a function by that name. If it is not found, then the standard builtin __import__() is called. Other mechanisms for invoking the import system (such as importlib.import_module()) do not perform this check and will always use the standard import system.

When a module is first imported, Python searches for the module and if found, it creates a module object ¹, initializing it. If the named module cannot be found, an ImportError is raised. Python implements various strategies to search for the named module when the import machinery is invoked. These strategies can be modified and extended by using various hooks described in the sections below. Changed in version 3.3: The import system has been updated to fully implement the second phase of PEP 302. There is no longer any implicit import machinery - the full import system is exposed through sys.meta_path. In addition, native namespace package support has been implemented (see PEP 420).

5.1 importlib

The importlib module provides a rich API for interacting with the import system. For example importlib.import_module() provides a recommended, simpler API than built-in __import__() for invoking the import machinery. Refer to the importlib library documentation for additional detail.

5.2 Packages

Python has only one type of module object, and all modules are of this type, regardless of whether the module is implemented in Python, C, or something else. To help organize modules and provide a naming hierarchy, Python has

¹ See types.ModuleType.

a concept of packages.

You can think of packages as the directories on a file system and modules as files within directories, but don't take this analogy too literally since packages and modules need not originate from the file system. For the purposes of this documentation, we'll use this convenient analogy of directories and files. Like file system directories, packages are organized hierarchically, and packages may themselves contain subpackages, as well as regular modules.

It's important to keep in mind that all packages are modules, but not all modules are packages. Or put another way, packages are just a special kind of module. Specifically, any module that contains a __path__ attribute is considered a package.

All modules have a name. Subpackage names are separated from their parent package name by dots, akin to Python's standard attribute access syntax. Thus you might have a module called sys and a package called email, which in turn has a subpackage called email.mime and a module within that subpackage called email.mime.text.

5.2.1 Regular packages

Python defines two types of packages, *regular packages* and *namespace packages*. Regular packages are traditional packages as they existed in Python 3.2 and earlier. A regular package is typically implemented as a directory containing an __init__.py file. When a regular package is imported, this __init__.py file is implicitly executed, and the objects it defines are bound to names in the package's namespace. The __init__.py file can contain the same Python code that any other module can contain, and Python will add some additional attributes to the module when it is imported.

For example, the following file system layout defines a top level parent package with three subpackages:

```
parent/
    __init__.py
    one/
    __init__.py
    two/
    __init__.py
    three/
        init .py
```

Importing parent.one will implicitly execute parent/__init__.py and parent/one/__init__.py. Subsequent imports of parent.two or parent.three will execute parent/two/__init__.py and parent/three/__init__.py respectively.

5.2.2 Namespace packages

A namespace package is a composite of various *portions*, where each portion contributes a subpackage to the parent package. Portions may reside in different locations on the file system. Portions may also be found in zip files, on the network, or anywhere else that Python searches during import. Namespace packages may or may not correspond directly to objects on the file system; they may be virtual modules that have no concrete representation.

Namespace packages do not use an ordinary list for their __path__ attribute. They instead use a custom iterable type which will automatically perform a new search for package portions on the next import attempt within that package if the path of their parent package (or sys.path for a top level package) changes.

With namespace packages, there is no parent/__init__.py file. In fact, there may be multiple parent directories found during import search, where each one is provided by a different portion. Thus parent/one may not be physically located next to parent/two. In this case, Python will create a namespace package for the top-level parent package whenever it or one of its subpackages is imported.

See also PEP 420 for the namespace package specification.

5.3 Searching

To begin the search, Python needs the *fully qualified* name of the module (or package, but for the purposes of this discussion, the difference is immaterial) being imported. This name may come from various arguments to the import statement, or from the parameters to the importlib.import_module() or __import__() functions.

This name will be used in various phases of the import search, and it may be the dotted path to a submodule, e.g. foo.bar.baz. In this case, Python first tries to import foo, then foo.bar, and finally foo.bar.baz. If any of the intermediate imports fail, an ImportError is raised.

5.3.1 The module cache

The first place checked during import search is sys.modules. This mapping serves as a cache of all modules that have been previously imported, including the intermediate paths. So if foo.bar.baz was previously imported, sys.modules will contain entries for foo, foo.bar, and foo.bar.baz. Each key will have as its value the corresponding module object.

During import, the module name is looked up in sys.modules and if present, the associated value is the module satisfying the import, and the process completes. However, if the value is None, then an ImportError is raised. If the module name is missing, Python will continue searching for the module.

sys.modules is writable. Deleting a key may not destroy the associated module (as other modules may hold references to it), but it will invalidate the cache entry for the named module, causing Python to search anew for the named module upon its next import. The key can also be assigned to None, forcing the next import of the module to result in an ImportError.

Beware though, as if you keep a reference to the module object, invalidate its cache entry in sys.modules, and then re-import the named module, the two module objects will *not* be the same. By contrast, imp.reload() will reuse the *same* module object, and simply reinitialise the module contents by rerunning the module's code.

5.3.2 Finders and loaders

If the named module is not found in sys.modules, then Python's import protocol is invoked to find and load the module. This protocol consists of two conceptual objects, *finders* and *loaders*. A finder's job is to determine whether it can find the named module using whatever strategy it knows about. Objects that implement both of these interfaces are referred to as *importers* - they return themselves when they find that they can load the requested module.

Python includes a number of default finders and importers. One knows how to locate frozen modules, and another knows how to locate built-in modules. A third default finder searches an *import path* for modules. The *import path* is a list of locations that may name file system paths or zip files. It can also be extended to search for any locatable resource, such as those identified by URLs.

The import machinery is extensible, so new finders can be added to extend the range and scope of module searching.

Finders do not actually load modules. If they can find the named module, they return a *loader*, which the import machinery then invokes to load the module and create the corresponding module object.

The following sections describe the protocol for finders and loaders in more detail, including how you can create and register new ones to extend the import machinery.

5.3.3 Import hooks

The import machinery is designed to be extensible; the primary mechanism for this are the *import hooks*. There are two types of import hooks: *meta hooks* and *import path hooks*.

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Meta hooks are called at the start of import processing, before any other import processing has occurred, other than sys.modules cache look up. This allows meta hooks to override sys.path processing, frozen modules, or even built-in modules. Meta hooks are registered by adding new finder objects to sys.meta path, as described below.

Import path hooks are called as part of sys.path (or package.__path__) processing, at the point where their associated path item is encountered. Import path hooks are registered by adding new callables to sys.path_hooks as described below.

5.3.4 The meta path

When the named module is not found in sys.modules, Python next searches sys.meta_path, which contains a list of meta path finder objects. These finders are queried in order to see if they know how to handle the named module. Meta path finders must implement a method called find_module() which takes two arguments, a name and an import path. The meta path finder can use any strategy it wants to determine whether it can handle the named module or not.

If the meta path finder knows how to handle the named module, it returns a loader object. If it cannot handle the named module, it returns None. If sys.meta_path processing reaches the end of its list without returning a loader, then an ImportError is raised. Any other exceptions raised are simply propagated up, aborting the import process.

The find_module () method of meta path finders is called with two arguments. The first is the fully qualified name of the module being imported, for example foo.bar.baz. The second argument is the path entries to use for the module search. For top-level modules, the second argument is None, but for submodules or subpackages, the second argument is the value of the parent package's __path__ attribute. If the appropriate __path__ attribute cannot be accessed, an ImportError is raised.

The meta path may be traversed multiple times for a single import request. For example, assuming none of the modules involved has already been cached, importing foo.bar.baz will first perform a top level import, calling mpf.find_module("foo", None) on each meta path finder (mpf). After foo has been imported, foo.bar will be imported by traversing the meta path a second time, calling mpf.find_module("foo.bar", foo.__path___). Once foo.bar has been imported, the final traversal will call mpf.find_module("foo.bar.baz", foo.bar.__path___).

Some meta path finders only support top level imports. These importers will always return None when anything other than None is passed as the second argument.

Python's default sys.meta_path has three meta path finders, one that knows how to import built-in modules, one that knows how to import frozen modules, and one that knows how to import modules from an *import path* (i.e. the *path based finder*).

5.4 Loaders

If and when a module loader is found its <code>load_module()</code> method is called, with a single argument, the fully qualified name of the module being imported. This method has several responsibilities, and should return the module object it has loaded ². If it cannot load the module, it should raise an <code>ImportError</code>, although any other exception raised during <code>load_module()</code> will be propagated.

In many cases, the finder and loader can be the same object; in such cases the finder.find_module() would just return self.

Loaders must satisfy the following requirements:

² The importlib implementation avoids using the return value directly. Instead, it gets the module object by looking the module name up in sys.modules. The indirect effect of this is that an imported module may replace itself in sys.modules. This is implementation-specific behavior that is not guaranteed to work in other Python implementations.

• If there is an existing module object with the given name in sys.modules, the loader must use that existing module. (Otherwise, imp.reload() will not work correctly.) If the named module does not exist in sys.modules, the loader must create a new module object and add it to sys.modules.

Note that the module *must* exist in sys.modules before the loader executes the module code. This is crucial because the module code may (directly or indirectly) import itself; adding it to sys.modules beforehand prevents unbounded recursion in the worst case and multiple loading in the best.

If loading fails, the loader must remove any modules it has inserted into sys.modules, but it must remove **only** the failing module, and only if the loader itself has loaded it explicitly. Any module already in the sys.modules cache, and any module that was successfully loaded as a side-effect, must remain in the cache.

- The loader may set the __file__ attribute of the module. If set, this attribute's value must be a string. The loader may opt to leave __file__ unset if it has no semantic meaning (e.g. a module loaded from a database).
- The loader may set the __name__ attribute of the module. While not required, setting this attribute is highly recommended so that the repr() of the module is more informative.
- If the module is a package (either regular or namespace), the loader must set the module object's __path__ attribute. The value must be iterable, but may be empty if __path__ has no further significance to the loader. If __path__ is not empty, it must produce strings when iterated over. More details on the semantics of __path__ are given below.
- The __loader__ attribute must be set to the loader object that loaded the module. This is mostly for introspection and reloading, but can be used for additional loader-specific functionality, for example getting data associated with a loader.
- The module's __package__ attribute should be set. Its value must be a string, but it can be the same value as its __name__. If the attribute is set to None or is missing, the import system will fill it in with a more appropriate value. When the module is a package, its __package__ value should be set to its __name__. When the module is not a package, __package__ should be set to the empty string for top-level modules, or for submodules, to the parent package's name. See PEP 366 for further details.

This attribute is used instead of __name__ to calculate explicit relative imports for main modules, as defined in PEP 366.

• If the module is a Python module (as opposed to a built-in module or a dynamically loaded extension), the loader should execute the module's code in the module's global name space (module.__dict__).

5.4.1 Module reprs

By default, all modules have a usable repr, however depending on the attributes set above, and hooks in the loader, you can more explicitly control the repr of module objects.

Loaders may implement a module_repr() method which takes a single argument, the module object. When repr(module) is called for a module with a loader supporting this protocol, whatever is returned from module.__loader__.module_repr(module) is returned as the module's repr without further processing. This return value must be a string.

If the module has no __loader__ attribute, or the loader has no module_repr() method, then the module object implementation itself will craft a default repr using whatever information is available. It will try to use the module.__name__, module.__file__, and module.__loader__ as input into the repr, with defaults for whatever information is missing.

Here are the exact rules used:

• If the module has a __loader__ and that loader has a module_repr() method, call it with a single argument, which is the module object. The value returned is used as the module's repr.

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- If an exception occurs in module_repr(), the exception is caught and discarded, and the calculation of the module's repr continues as if module_repr() did not exist.
- If the module has a ___file__ attribute, this is used as part of the module's repr.
- If the module has no __file__ but does have a __loader__, then the loader's repr is used as part of the module's repr.
- Otherwise, just use the module's ___name__ in the repr.

This example, from PEP 420 shows how a loader can craft its own module repr:

class NamespaceLoader:

```
@classmethod
def module_repr(cls, module):
    return "<module ' { } ' (namespace) > ".format(module.__name__)
```

5.4.2 module.__path__

By definition, if a module has an __path__ attribute, it is a package, regardless of its value.

A package's __path__ attribute is used during imports of its subpackages. Within the import machinery, it functions much the same as sys.path, i.e. providing a list of locations to search for modules during import. However, __path__ is typically much more constrained than sys.path.

__path__ must be an iterable of strings, but it may be empty. The same rules used for sys.path also apply to a package's __path__, and sys.path_hooks (described below) are consulted when traversing a package's __path__.

A package's __init__.py file may set or alter the package's __path__ attribute, and this was typically the way namespace packages were implemented prior to **PEP 420**. With the adoption of **PEP 420**, namespace packages no longer need to supply __init__.py files containing only __path__ manipulation code; the namespace loader automatically sets __path__ correctly for the namespace package.

5.5 The Path Based Finder

As mentioned previously, Python comes with several default meta path finders. One of these, called the *path based finder*, searches an *import path*, which contains a list of *path entries*. Each path entry names a location to search for modules.

The path based finder itself doesn't know how to import anything. Instead, it traverses the individual path entries, associating each of them with a path entry finder that knows how to handle that particular kind of path.

The default set of path entry finders implement all the semantics for finding modules on the file system, handling special file types such as Python source code (.py files), Python byte code (.pyc and .pyo files) and shared libraries (e.g. .so files). When supported by the zipimport module in the standard library, the default path entry finders also handle loading all of these file types (other than shared libraries) from zipfiles.

Path entries need not be limited to file system locations. They can refer to URLs, database queries, or any other location that can be specified as a string.

The path based finder provides additional hooks and protocols so that you can extend and customize the types of searchable path entries. For example, if you wanted to support path entries as network URLs, you could write a hook that implements HTTP semantics to find modules on the web. This hook (a callable) would return a *path entry finder* supporting the protocol described below, which was then used to get a loader for the module from the web.

A word of warning: this section and the previous both use the term *finder*, distinguishing between them by using the terms *meta path finder* and *path entry finder*. These two types of finders are very similar, support similar protocols,

and function in similar ways during the import process, but it's important to keep in mind that they are subtly different. In particular, meta path finders operate at the beginning of the import process, as keyed off the sys.meta_path traversal.

By contrast, path entry finders are in a sense an implementation detail of the path based finder, and in fact, if the path based finder were to be removed from sys.meta_path, none of the path entry finder semantics would be invoked.

5.5.1 Path entry finders

The *path based finder* is responsible for finding and loading Python modules and packages whose location is specified with a string *path entry*. Most path entries name locations in the file system, but they need not be limited to this.

As a meta path finder, the *path based finder* implements the find_module() protocol previously described, however it exposes additional hooks that can be used to customize how modules are found and loaded from the *import path*.

Three variables are used by the *path based finder*, sys.path, sys.path_hooks and sys.path_importer_cache. The __path__ attributes on package objects are also used. These provide additional ways that the import machinery can be customized.

sys.path contains a list of strings providing search locations for modules and packages. It is initialized from the PYTHONPATH environment variable and various other installation- and implementation-specific defaults. Entries in sys.path can name directories on the file system, zip files, and potentially other "locations" (see the site module) that should be searched for modules, such as URLs, or database queries.

The path based finder is a meta path finder, so the import machinery begins the import path search by calling the path based finder's find_module() method as described previously. When the path argument to find_module() is given, it will be a list of string paths to traverse - typically a package's __path__ attribute for an import within that package. If the path argument is None, this indicates a top level import and sys.path is used.

The path based finder iterates over every entry in the search path, and for each of these, looks for an appropriate *path entry finder* for the path entry. Because this can be an expensive operation (e.g. there may be *stat()* call overheads for this search), the path based finder maintains a cache mapping path entries to path entry finders. This cache is maintained in sys.path_importer_cache (despite the name, this cache actually stores finder objects rather than being limited to *importer* objects). In this way, the expensive search for a particular *path entry* location's *path entry finder* need only be done once. User code is free to remove cache entries from sys.path_importer_cache forcing the path based finder to perform the path entry search again ³.

If the path entry is not present in the cache, the path based finder iterates over every callable in sys.path_hooks. Each of the *path entry hooks* in this list is called with a single argument, the path entry to be searched. This callable may either return a *path entry finder* that can handle the path entry, or it may raise ImportError. An ImportError is used by the path based finder to signal that the hook cannot find a *path entry finder* for that *path entry*. The exception is ignored and *import path* iteration continues.

If sys.path_hooks iteration ends with no path entry finder being returned, then the path based finder's find_module() method will store None in sys.path_importer_cache (to indicate that there is no finder for this path entry) and return None, indicating that this meta path finder could not find the module.

If a path entry finder is returned by one of the path entry hook callables on sys.path_hooks, then the following protocol is used to ask the finder for a module loader, which is then used to load the module.

5.5.2 Path entry finder protocol

In order to support imports of modules and initialized packages and also to contribute portions to namespace packages, path entry finders must implement the find loader () method.

³ In legacy code, it is possible to find instances of imp.NullImporter in the sys.path_importer_cache. It is recommended that code be changed to use None instead. See *portingpythoncode* for more details.

find_loader() takes one argument, the fully qualified name of the module being imported. find_loader() returns a 2-tuple where the first item is the loader and the second item is a namespace *portion*. When the first item (i.e. the loader) is None, this means that while the path entry finder does not have a loader for the named module, it knows that the path entry contributes to a namespace portion for the named module. This will almost always be the case where Python is asked to import a namespace package that has no physical presence on the file system. When a path entry finder returns None for the loader, the second item of the 2-tuple return value must be a sequence, although it can be empty.

If find_loader() returns a non-None loader value, the portion is ignored and the loader is returned from the path based finder, terminating the search through the path entries.

For backwards compatibility with other implementations of the import protocol, many path entry finders also support the same, traditional find_module() method that meta path finders support. However path entry finder find_module() methods are never called with a path argument (they are expected to record the appropriate path information from the initial call to the path hook).

The find_module() method on path entry finders is deprecated, as it does not allow the path entry finder to contribute portions to namespace packages. Instead path entry finders should implement the find_loader() method as described above. If it exists on the path entry finder, the import system will always call find_loader() in preference to find module().

5.6 Replacing the standard import system

The most reliable mechanism for replacing the entire import system is to delete the default contents of sys.meta_path, replacing them entirely with a custom meta path hook.

If it is acceptable to only alter the behaviour of import statements without affecting other APIs that access the import system, then replacing the builtin __import__ () function may be sufficient. This technique may also be employed at the module level to only alter the behaviour of import statements within that module.

To selectively prevent import of some modules from a hook early on the meta path (rather than disabling the standard import system entirely), it is sufficient to raise ImportError directly from find_module() instead of returning None. The latter indicates that the meta path search should continue. while raising an exception terminates it immediately.

5.7 Open issues

XXX It would be really nice to have a diagram.

XXX * (import_machinery.rst) how about a section devoted just to the attributes of modules and packages, perhaps expanding upon or supplanting the related entries in the data model reference page?

XXX runpy, pkgutil, et al in the library manual should all get "See Also" links at the top pointing to the new import system section.

5.8 References

The import machinery has evolved considerably since Python's early days. The original specification for packages is still available to read, although some details have changed since the writing of that document.

The original specification for sys.meta_path was PEP 302, with subsequent extension in PEP 420.

PEP 420 introduced *namespace packages* for Python 3.3. **PEP 420** also introduced the find_loader() protocol as an alternative to find_module().

PEP 366 describes the addition of the __package__ attribute for explicit relative imports in main modules.

PEP 328 introduced absolute and explicit relative imports and initially proposed __name__ for semantics PEP 366 would eventually specify for __package__.

PEP 338 defines executing modules as scripts.

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CHAPTER

SIX

EXPRESSIONS

This chapter explains the meaning of the elements of expressions in Python.

Syntax Notes: In this and the following chapters, extended BNF notation will be used to describe syntax, not lexical analysis. When (one alternative of) a syntax rule has the form

```
name ::= othername
```

and no semantics are given, the semantics of this form of name are the same as for othername.

6.1 Arithmetic conversions

When a description of an arithmetic operator below uses the phrase "the numeric arguments are converted to a common type," this means that the operator implementation for built-in types works that way:

- If either argument is a complex number, the other is converted to complex;
- otherwise, if either argument is a floating point number, the other is converted to floating point;
- otherwise, both must be integers and no conversion is necessary.

Some additional rules apply for certain operators (e.g., a string left argument to the '%' operator). Extensions must define their own conversion behavior.

6.2 Atoms

Atoms are the most basic elements of expressions. The simplest atoms are identifiers or literals. Forms enclosed in parentheses, brackets or braces are also categorized syntactically as atoms. The syntax for atoms is:

6.2.1 Identifiers (Names)

An identifier occurring as an atom is a name. See section *Identifiers and keywords* for lexical definition and section *Naming and binding* for documentation of naming and binding.

When the name is bound to an object, evaluation of the atom yields that object. When a name is not bound, an attempt to evaluate it raises a NameError exception.

Private name mangling: When an identifier that textually occurs in a class definition begins with two or more underscore characters and does not end in two or more underscores, it is considered a *private name* of that class. Private names are transformed to a longer form before code is generated for them. The transformation inserts the class name in front of the name, with leading underscores removed, and a single underscore inserted in front of the class name. For example, the identifier __spam occurring in a class named Ham will be transformed to _Ham__spam. This transformation is independent of the syntactical context in which the identifier is used. If the transformed name is extremely long (longer than 255 characters), implementation defined truncation may happen. If the class name consists only of underscores, no transformation is done.

6.2.2 Literals

Python supports string and bytes literals and various numeric literals:

Evaluation of a literal yields an object of the given type (string, bytes, integer, floating point number, complex number) with the given value. The value may be approximated in the case of floating point and imaginary (complex) literals. See section *Literals* for details.

All literals correspond to immutable data types, and hence the object's identity is less important than its value. Multiple evaluations of literals with the same value (either the same occurrence in the program text or a different occurrence) may obtain the same object or a different object with the same value.

6.2.3 Parenthesized forms

A parenthesized form is an optional expression list enclosed in parentheses:

```
parenth_form ::= "(" [expression_list] ")"
```

A parenthesized expression list yields whatever that expression list yields: if the list contains at least one comma, it yields a tuple; otherwise, it yields the single expression that makes up the expression list.

An empty pair of parentheses yields an empty tuple object. Since tuples are immutable, the rules for literals apply (i.e., two occurrences of the empty tuple may or may not yield the same object).

Note that tuples are not formed by the parentheses, but rather by use of the comma operator. The exception is the empty tuple, for which parentheses *are* required — allowing unparenthesized "nothing" in expressions would cause ambiguities and allow common typos to pass uncaught.

6.2.4 Displays for lists, sets and dictionaries

For constructing a list, a set or a dictionary Python provides special syntax called "displays", each of them in two flavors:

- either the container contents are listed explicitly, or
- they are computed via a set of looping and filtering instructions, called a comprehension.

Common syntax elements for comprehensions are:

```
comprehension ::= expression comp_for
comp_for ::= "for" target_list "in" or_test [comp_iter]
comp_iter ::= comp_for | comp_if
comp_if ::= "if" expression_nocond [comp_iter]
```

The comprehension consists of a single expression followed by at least one for clause and zero or more for or if clauses. In this case, the elements of the new container are those that would be produced by considering each of the for or if clauses a block, nesting from left to right, and evaluating the expression to produce an element each time the innermost block is reached.

Note that the comprehension is executed in a separate scope, so names assigned to in the target list don't "leak" in the enclosing scope.

6.2.5 List displays

A list display is a possibly empty series of expressions enclosed in square brackets:

```
list_display ::= "[" [expression_list | comprehension] "]"
```

A list display yields a new list object, the contents being specified by either a list of expressions or a comprehension. When a comma-separated list of expressions is supplied, its elements are evaluated from left to right and placed into the list object in that order. When a comprehension is supplied, the list is constructed from the elements resulting from the comprehension.

6.2.6 Set displays

A set display is denoted by curly braces and distinguishable from dictionary displays by the lack of colons separating keys and values:

```
set_display ::= "{" (expression_list | comprehension) "}"
```

A set display yields a new mutable set object, the contents being specified by either a sequence of expressions or a comprehension. When a comma-separated list of expressions is supplied, its elements are evaluated from left to right and added to the set object. When a comprehension is supplied, the set is constructed from the elements resulting from the comprehension.

An empty set cannot be constructed with { }; this literal constructs an empty dictionary.

6.2.7 Dictionary displays

A dictionary display is a possibly empty series of key/datum pairs enclosed in curly braces:

A dictionary display yields a new dictionary object.

If a comma-separated sequence of key/datum pairs is given, they are evaluated from left to right to define the entries of the dictionary: each key object is used as a key into the dictionary to store the corresponding datum. This means that you can specify the same key multiple times in the key/datum list, and the final dictionary's value for that key will be the last one given.

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A dict comprehension, in contrast to list and set comprehensions, needs two expressions separated with a colon followed by the usual "for" and "if" clauses. When the comprehension is run, the resulting key and value elements are inserted in the new dictionary in the order they are produced.

Restrictions on the types of the key values are listed earlier in section *The standard type hierarchy*. (To summarize, the key type should be *hashable*, which excludes all mutable objects.) Clashes between duplicate keys are not detected; the last datum (textually rightmost in the display) stored for a given key value prevails.

6.2.8 Generator expressions

A generator expression is a compact generator notation in parentheses:

```
generator_expression ::= "(" expression comp_for ")"
```

A generator expression yields a new generator object. Its syntax is the same as for comprehensions, except that it is enclosed in parentheses instead of brackets or curly braces.

Variables used in the generator expression are evaluated lazily when the $__next__()$ method is called for generator object (in the same fashion as normal generators). However, the leftmost for clause is immediately evaluated, so that an error produced by it can be seen before any other possible error in the code that handles the generator expression. Subsequent for clauses cannot be evaluated immediately since they may depend on the previous for loop. For example: (x*y for x in range(10) for y in bar(x)).

The parentheses can be omitted on calls with only one argument. See section *Calls* for the detail.

6.2.9 Yield expressions

The yield expression is only used when defining a generator function, and can only be used in the body of a function definition. Using a yield expression in a function definition is sufficient to cause that definition to create a generator function instead of a normal function.

When a generator function is called, it returns an iterator known as a generator. That generator then controls the execution of a generator function. The execution starts when one of the generator's methods is called. At that time, the execution proceeds to the first yield expression, where it is suspended again, returning the value of $expression_list$ to generator's caller. By suspended we mean that all local state is retained, including the current bindings of local variables, the instruction pointer, and the internal evaluation stack. When the execution is resumed by calling one of the generator's methods, the function can proceed exactly as if the yield expression was just another external call. The value of the yield expression after resuming depends on the method which resumed the execution. If $_next_()$ is used (typically via either a for or the next() builtin) then the result is None, otherwise, if send() is used, then the result will be the value passed in to that method.

All of this makes generator functions quite similar to coroutines; they yield multiple times, they have more than one entry point and their execution can be suspended. The only difference is that a generator function cannot control where should the execution continue after it yields; the control is always transferred to the generator's caller.

yield expressions are allowed in the try clause of a try ... finally construct. If the generator is not resumed before it is finalized (by reaching a zero reference count or by being garbage collected), the generator-iterator's close() method will be called, allowing any pending finally clauses to execute.

When yield from <expr> is used, it treats the supplied expression as a subiterator. All values produced by that subiterator are passed directly to the caller of the current generator's methods. Any values passed in with send() and any exceptions passed in with throw() are passed to the underlying iterator if it has the appropriate methods. If

this is not the case, then send() will raise AttributeError or TypeError, while throw() will just raise the passed in exception immediately.

When the underlying iterator is complete, the value attribute of the raised StopIteration instance becomes the value of the yield expression. It can be either set explicitly when raising StopIteration, or automatically when the sub-iterator is a generator (by returning a value from the sub-generator).

Changed in version 3.3: Added yield from <expr> to delegate control flow to a subiterator

The parentheses can be omitted when the yield expression is the sole expression on the right hand side of an assignment statement.

Generator-iterator methods

This subsection describes the methods of a generator iterator. They can be used to control the execution of a generator function.

Note that calling any of the generator methods below when the generator is already executing raises a ValueError exception.

```
generator.__next__()
```

Starts the execution of a generator function or resumes it at the last executed yield expression. When a generator function is resumed with a __next__() method, the current yield expression always evaluates to None. The execution then continues to the next yield expression, where the generator is suspended again, and the value of the expression_list is returned to next() 's caller. If the generator exits without yielding another value, a StopIteration exception is raised.

This method is normally called implicitly, e.g. by a for loop, or by the built-in next () function.

```
generator.send(value)
```

Resumes the execution and "sends" a value into the generator function. The value argument becomes the result of the current <code>yield</code> expression. The <code>send()</code> method returns the next value yielded by the generator, or raises <code>StopIteration</code> if the generator exits without yielding another value. When <code>send()</code> is called to start the generator, it must be called with <code>None</code> as the argument, because there is no <code>yield</code> expression that could receive the value.

```
generator.throw(type[, value[, traceback]])
```

Raises an exception of type type at the point where generator was paused, and returns the next value yielded by the generator function. If the generator exits without yielding another value, a StopIteration exception is raised. If the generator function does not catch the passed-in exception, or raises a different exception, then that exception propagates to the caller.

```
generator.close()
```

Raises a GeneratorExit at the point where the generator function was paused. If the generator function then raises StopIteration (by exiting normally, or due to already being closed) or GeneratorExit (by not catching the exception), close returns to its caller. If the generator yields a value, a RuntimeError is raised. If the generator raises any other exception, it is propagated to the caller. close() does nothing if the generator has already exited due to an exception or normal exit.

Here is a simple example that demonstrates the behavior of generators and generator functions:

```
>>> def echo(value=None):
...     print("Execution starts when 'next()' is called for the first time.")
...     try:
...     while True:
...         try:
...         value = (yield value)
...         except Exception as e:
...         value = e
```

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See Also:

- **PEP 0255 Simple Generators** The proposal for adding generators and the yield statement to Python.
- **PEP 0342 Coroutines via Enhanced Generators** The proposal to enhance the API and syntax of generators, making them usable as simple coroutines.
- **PEP 0380 Syntax for Delegating to a Subgenerator** The proposal to introduce the yield_from syntax, making delegation to sub-generators easy.

6.3 Primaries

Primaries represent the most tightly bound operations of the language. Their syntax is:

```
primary ::= atom | attributeref | subscription | slicing | call
```

6.3.1 Attribute references

An attribute reference is a primary followed by a period and a name:

```
attributeref ::= primary "." identifier
```

The primary must evaluate to an object of a type that supports attribute references, which most objects do. This object is then asked to produce the attribute whose name is the identifier (which can be customized by overriding the __getattr__() method). If this attribute is not available, the exception AttributeError is raised. Otherwise, the type and value of the object produced is determined by the object. Multiple evaluations of the same attribute reference may yield different objects.

6.3.2 Subscriptions

A subscription selects an item of a sequence (string, tuple or list) or mapping (dictionary) object:

```
subscription ::= primary "[" expression_list "]"
```

The primary must evaluate to an object that supports subscription, e.g. a list or dictionary. User-defined objects can support subscription by defining a __getitem__() method.

For built-in objects, there are two types of objects that support subscription:

If the primary is a mapping, the expression list must evaluate to an object whose value is one of the keys of the mapping, and the subscription selects the value in the mapping that corresponds to that key. (The expression list is a tuple except if it has exactly one item.)

If the primary is a sequence, the expression (list) must evaluate to an integer or a slice (as discussed in the following section).

The formal syntax makes no special provision for negative indices in sequences; however, built-in sequences all provide a $_getitem_$ () method that interprets negative indices by adding the length of the sequence to the index (so that x[-1] selects the last item of x). The resulting value must be a nonnegative integer less than the number of items in the sequence, and the subscription selects the item whose index is that value (counting from zero). Since the support for negative indices and slicing occurs in the object's $__getitem_$ () method, subclasses overriding this method will need to explicitly add that support.

A string's items are characters. A character is not a separate data type but a string of exactly one character.

6.3.3 Slicings

A slicing selects a range of items in a sequence object (e.g., a string, tuple or list). Slicings may be used as expressions or as targets in assignment or del statements. The syntax for a slicing:

```
primary "[" slice list "]"
slicing
                   slice_item ("," slice_item)* [","]
slice_list
              ::=
slice item
                   expression | proper_slice
              ::=
                   [lower bound] ":" [upper bound] [ ":" [stride] ]
proper_slice
             ::=
lower bound
                   expression
              ::=
upper_bound
              ::=
                   expression
stride
                   expression
```

There is ambiguity in the formal syntax here: anything that looks like an expression list also looks like a slice list, so any subscription can be interpreted as a slicing. Rather than further complicating the syntax, this is disambiguated by defining that in this case the interpretation as a subscription takes priority over the interpretation as a slicing (this is the case if the slice list contains no proper slice).

The semantics for a slicing are as follows. The primary must evaluate to a mapping object, and it is indexed (using the same __getitem__() method as normal subscription) with a key that is constructed from the slice list, as follows. If the slice list contains at least one comma, the key is a tuple containing the conversion of the slice items; otherwise, the conversion of the lone slice item is the key. The conversion of a slice item that is an expression is that expression. The conversion of a proper slice is a slice object (see section *The standard type hierarchy*) whose start, stop and step attributes are the values of the expressions given as lower bound, upper bound and stride, respectively, substituting None for missing expressions.

6.3.4 Calls

A call calls a callable object (e.g., a function) with a possibly empty series of arguments:

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```
primary "(" [argument list [","] | comprehension] ")"
call
                           positional_arguments ["," keyword_arguments]
argument list
                      ::=
                            ["," "*" expression] ["," keyword_arguments]
                            ["," "**" expression]
                            | keyword_arguments ["," "*" expression]
                            ["," keyword_arguments] ["," "**" expression]
                            "*" expression ["," keyword_arguments] ["," "**" expressi
                            | "∗∗" expression
                           expression ("," expression) *
positional_arguments
                      ::=
                           keyword_item ("," keyword_item) *
keyword_arguments
                      ::=
                           identifier "=" expression
keyword_item
                      ::=
```

A trailing comma may be present after the positional and keyword arguments but does not affect the semantics.

The primary must evaluate to a callable object (user-defined functions, built-in functions, methods of built-in objects, class objects, methods of class instances, and all objects having a __call__() method are callable). All argument expressions are evaluated before the call is attempted. Please refer to section *Function definitions* for the syntax of formal parameter lists.

If keyword arguments are present, they are first converted to positional arguments, as follows. First, a list of unfilled slots is created for the formal parameters. If there are N positional arguments, they are placed in the first N slots. Next, for each keyword argument, the identifier is used to determine the corresponding slot (if the identifier is the same as the first formal parameter name, the first slot is used, and so on). If the slot is already filled, a TypeError exception is raised. Otherwise, the value of the argument is placed in the slot, filling it (even if the expression is None, it fills the slot). When all arguments have been processed, the slots that are still unfilled are filled with the corresponding default value from the function definition. (Default values are calculated, once, when the function is defined; thus, a mutable object such as a list or dictionary used as default value will be shared by all calls that don't specify an argument value for the corresponding slot; this should usually be avoided.) If there are any unfilled slots for which no default value is specified, a TypeError exception is raised. Otherwise, the list of filled slots is used as the argument list for the call.

CPython implementation detail: An implementation may provide built-in functions whose positional parameters do not have names, even if they are 'named' for the purpose of documentation, and which therefore cannot be supplied by keyword. In CPython, this is the case for functions implemented in C that use PyArg_ParseTuple() to parse their arguments.

If there are more positional arguments than there are formal parameter slots, a TypeError exception is raised, unless a formal parameter using the syntax *identifier is present; in this case, that formal parameter receives a tuple containing the excess positional arguments (or an empty tuple if there were no excess positional arguments).

If any keyword argument does not correspond to a formal parameter name, a TypeError exception is raised, unless a formal parameter using the syntax **identifier is present; in this case, that formal parameter receives a dictionary containing the excess keyword arguments (using the keywords as keys and the argument values as corresponding values), or a (new) empty dictionary if there were no excess keyword arguments.

If the syntax *expression appears in the function call, expression must evaluate to an iterable. Elements from this iterable are treated as if they were additional positional arguments; if there are positional arguments xI, ..., xN, and expression evaluates to a sequence yI, ..., yM, this is equivalent to a call with M+N positional arguments xI, ..., xN, yI, ..., yM.

A consequence of this is that although the *expression syntax may appear after some keyword arguments, it is processed before the keyword arguments (and the **expression argument, if any – see below). So:

```
>>> def f(a, b):
... print(a, b)
...
>>> f(b=1, *(2,))
2 1
```

```
>>> f(a=1, *(2,))
Traceback (most recent call last):
   File "<stdin>", line 1, in ?
TypeError: f() got multiple values for keyword argument 'a'
>>> f(1, *(2,))
1 2
```

It is unusual for both keyword arguments and the *expression syntax to be used in the same call, so in practice this confusion does not arise.

If the syntax **expression appears in the function call, expression must evaluate to a mapping, the contents of which are treated as additional keyword arguments. In the case of a keyword appearing in both expression and as an explicit keyword argument, a TypeError exception is raised.

Formal parameters using the syntax *identifier or **identifier cannot be used as positional argument slots or as keyword argument names.

A call always returns some value, possibly None, unless it raises an exception. How this value is computed depends on the type of the callable object.

If it is-

- a user-defined function: The code block for the function is executed, passing it the argument list. The first thing the code block will do is bind the formal parameters to the arguments; this is described in section *Function definitions*. When the code block executes a return statement, this specifies the return value of the function call.
- **a built-in function or method:** The result is up to the interpreter; see *built-in-funcs* for the descriptions of built-in functions and methods.
- a class object: A new instance of that class is returned.
- **a class instance method:** The corresponding user-defined function is called, with an argument list that is one longer than the argument list of the call: the instance becomes the first argument.
- a class instance: The class must define a __call__() method; the effect is then the same as if that method was called.

6.4 The power operator

The power operator binds more tightly than unary operators on its left; it binds less tightly than unary operators on its right. The syntax is:

```
power ::= primary ["**" u expr]
```

Thus, in an unparenthesized sequence of power and unary operators, the operators are evaluated from right to left (this does not constrain the evaluation order for the operands): -1 **2 results in -1.

The power operator has the same semantics as the built-in pow () function, when called with two arguments: it yields its left argument raised to the power of its right argument. The numeric arguments are first converted to a common type, and the result is of that type.

For int operands, the result has the same type as the operands unless the second argument is negative; in that case, all arguments are converted to float and a float result is delivered. For example, 10**2 returns 100, but 10**-2 returns 0.01.

Raising 0.0 to a negative power results in a ZeroDivisionError. Raising a negative number to a fractional power results in a complex number. (In earlier versions it raised a ValueError.)

6.5 Unary arithmetic and bitwise operations

All unary arithmetic and bitwise operations have the same priority:

```
u_expr ::= power | "-" u_expr | "+" u_expr | "~" u_expr
```

The unary – (minus) operator yields the negation of its numeric argument.

The unary + (plus) operator yields its numeric argument unchanged.

The unary \sim (invert) operator yields the bitwise inversion of its integer argument. The bitwise inversion of \times is defined as -(x+1). It only applies to integral numbers.

In all three cases, if the argument does not have the proper type, a TypeError exception is raised.

6.6 Binary arithmetic operations

The binary arithmetic operations have the conventional priority levels. Note that some of these operations also apply to certain non-numeric types. Apart from the power operator, there are only two levels, one for multiplicative operators and one for additive operators:

The * (multiplication) operator yields the product of its arguments. The arguments must either both be numbers, or one argument must be an integer and the other must be a sequence. In the former case, the numbers are converted to a common type and then multiplied together. In the latter case, sequence repetition is performed; a negative repetition factor yields an empty sequence.

The / (division) and // (floor division) operators yield the quotient of their arguments. The numeric arguments are first converted to a common type. Integer division yields a float, while floor division of integers results in an integer; the result is that of mathematical division with the 'floor' function applied to the result. Division by zero raises the ZeroDivisionError exception.

The % (modulo) operator yields the remainder from the division of the first argument by the second. The numeric arguments are first converted to a common type. A zero right argument raises the ZeroDivisionError exception. The arguments may be floating point numbers, e.g., 3.14%0.7 equals 0.34 (since 3.14 equals 4*0.7 + 0.34.) The modulo operator always yields a result with the same sign as its second operand (or zero); the absolute value of the result is strictly smaller than the absolute value of the second operand 1.

The floor division and modulo operators are connected by the following identity: x == (x//y) * y + (x * y). Floor division and modulo are also connected with the built-in function divmod(): divmod(x, y) == (x//y), x * y.

In addition to performing the modulo operation on numbers, the % operator is also overloaded by string objects to perform old-style string formatting (also known as interpolation). The syntax for string formatting is described in the Python Library Reference, section *old-string-formatting*.

 $^{^1}$ While abs (x%y) < abs (y) is true mathematically, for floats it may not be true numerically due to roundoff. For example, and assuming a platform on which a Python float is an IEEE 754 double-precision number, in order that -1e-100 % 1e100 have the same sign as 1e100, the computed result is -1e-100 + 1e100, which is numerically exactly equal to 1e100. The function math.fmod() returns a result whose sign matches the sign of the first argument instead, and so returns -1e-100 in this case. Which approach is more appropriate depends on the application.

² If x is very close to an exact integer multiple of y, it's possible for x//y to be one larger than (x-x*y)//y due to rounding. In such cases, Python returns the latter result, in order to preserve that divmod (x,y)[0] * y + x * y be very close to x.

The floor division operator, the modulo operator, and the divmod() function are not defined for complex numbers. Instead, convert to a floating point number using the abs() function if appropriate.

The + (addition) operator yields the sum of its arguments. The arguments must either both be numbers or both sequences of the same type. In the former case, the numbers are converted to a common type and then added together. In the latter case, the sequences are concatenated.

The – (subtraction) operator yields the difference of its arguments. The numeric arguments are first converted to a common type.

6.7 Shifting operations

The shifting operations have lower priority than the arithmetic operations:

```
shift_expr ::= a_expr | shift_expr ( "<<" | ">>" ) a_expr
```

These operators accept integers as arguments. They shift the first argument to the left or right by the number of bits given by the second argument.

A right shift by n bits is defined as division by pow (2, n). A left shift by n bits is defined as multiplication with pow (2, n).

Note: In the current implementation, the right-hand operand is required to be at most sys.maxsize. If the right-hand operand is larger than sys.maxsize an OverflowError exception is raised.

6.8 Binary bitwise operations

Each of the three bitwise operations has a different priority level:

```
and_expr ::= shift_expr | and_expr "%" shift_expr
xor_expr ::= and_expr | xor_expr "^" and_expr
or_expr ::= xor_expr | or_expr "|" xor_expr
```

The & operator yields the bitwise AND of its arguments, which must be integers.

The ^ operator yields the bitwise XOR (exclusive OR) of its arguments, which must be integers.

The | operator yields the bitwise (inclusive) OR of its arguments, which must be integers.

6.9 Comparisons

Unlike C, all comparison operations in Python have the same priority, which is lower than that of any arithmetic, shifting or bitwise operation. Also unlike C, expressions like a < b < c have the interpretation that is conventional in mathematics:

Comparisons yield boolean values: True or False.

Comparisons can be chained arbitrarily, e.g., x < y <= z is equivalent to x < y and y <= z, except that y is evaluated only once (but in both cases z is not evaluated at all when x < y is found to be false).

Formally, if a, b, c, ..., y, z are expressions and op1, op2, ..., opN are comparison operators, then a op1 b op2 c ... y opN z is equivalent to a op1 b and b op2 c and ... y opN z, except that each expression is evaluated at most once.

Note that a op1 b op2 c doesn't imply any kind of comparison between a and c, so that, e.g., x < y > z is perfectly legal (though perhaps not pretty).

The operators <, >, ==, >=, <=, and != compare the values of two objects. The objects need not have the same type. If both are numbers, they are converted to a common type. Otherwise, the == and != operators always consider objects of different types to be unequal, while the <, >, >= and <= operators raise a TypeError when comparing objects of different types that do not implement these operators for the given pair of types. You can control comparison behavior of objects of non-built-in types by defining rich comparison methods like $__gt$ __(), described in section $Basic\ customization$.

Comparison of objects of the same type depends on the type:

- Numbers are compared arithmetically.
- The values float ('NaN') and Decimal ('NaN') are special. The are identical to themselves, x is x but are not equal to themselves, x != x. Additionally, comparing any value to a not-a-number value will return False. For example, both 3 < float ('NaN') and float ('NaN') < 3 will return False.
- Bytes objects are compared lexicographically using the numeric values of their elements.
- Strings are compared lexicographically using the numeric equivalents (the result of the built-in function ord()) of their characters. ³ String and bytes object can't be compared!
- Tuples and lists are compared lexicographically using comparison of corresponding elements. This means that to compare equal, each element must compare equal and the two sequences must be of the same type and have the same length.

If not equal, the sequences are ordered the same as their first differing elements. For example, [1, 2, x] < [1, 2, y] has the same value as x < y. If the corresponding element does not exist, the shorter sequence is ordered first (for example, [1, 2] < [1, 2, 3]).

- Mappings (dictionaries) compare equal if and only if they have the same (key, value) pairs. Order comparisons ('<', '<=', '>=', '>') raise TypeError.
- Sets and frozensets define comparison operators to mean subset and superset tests. Those relations do not define total orderings (the two sets {1,2} and {2,3} are not equal, nor subsets of one another, nor supersets of one another). Accordingly, sets are not appropriate arguments for functions which depend on total ordering. For example, min(), max(), and sorted() produce undefined results given a list of sets as inputs.
- Most other objects of built-in types compare unequal unless they are the same object; the choice whether one
 object is considered smaller or larger than another one is made arbitrarily but consistently within one execution
 of a program.

Comparison of objects of the differing types depends on whether either of the types provide explicit support for the comparison. Most numeric types can be compared with one another. When cross-type comparison is not supported, the comparison method returns NotImplemented. The operators in and not in test for membership. x in s evaluates to true if x is a member of s, and false otherwise. x not in s returns the negation of x in s. All built-in sequences and set types support this as well as dictionary, for which in tests whether a the dictionary has a given key. For container types such as list, tuple, set, frozenset, dict, or collections deque, the expression x in y is equivalent to any (x is y or y == y for y in y or y == y for y for y == y for y for y for y == y for y for y for y == y for y for

 $^{^3}$ While comparisons between strings make sense at the byte level, they may be counter-intuitive to users. For example, the strings "\u00C7" and "\u00327\u0043" compare differently, even though they both represent the same unicode character (LATIN CAPITAL LETTER C WITH CEDILLA). To compare strings in a human recognizable way, compare using unicodedata.normalize().

For the string and bytes types, x in y is true if and only if x is a substring of y. An equivalent test is y. find (x) != -1. Empty strings are always considered to be a substring of any other string, so "" in "abc" will return True.

For user-defined classes which define the $__contains__()$ method, x in y is true if and only if y. $__contains__(x)$ is true.

For user-defined classes which do not define $__{contains}_{_}$ () but do define $__{iter}_{_}$ (), x in y is true if some value z with x == z is produced while iterating over y. If an exception is raised during the iteration, it is as if in raised that exception.

Lastly, the old-style iteration protocol is tried: if a class defines $__{getitem}$ (), x in y is true if and only if there is a non-negative integer index i such that x == y[i], and all lower integer indices do not raise IndexError exception. (If any other exception is raised, it is as if in raised that exception).

The operator not in is defined to have the inverse true value of in.

The operators is and is not test for object identity: x is y is true if and only if x and y are the same object. x is not y yields the inverse truth value.

6.10 Boolean operations

```
or_test ::= and_test | or_test "or" and_test
and_test ::= not_test | and_test "and" not_test
not test ::= comparison | "not" not test
```

In the context of Boolean operations, and also when expressions are used by control flow statements, the following values are interpreted as false: False, None, numeric zero of all types, and empty strings and containers (including strings, tuples, lists, dictionaries, sets and frozensets). All other values are interpreted as true. User-defined objects can customize their truth value by providing a __bool__() method.

The operator not yields True if its argument is false, False otherwise.

The expression x and y first evaluates x; if x is false, its value is returned; otherwise, y is evaluated and the resulting value is returned.

The expression $x \circ y$ first evaluates x; if x is true, its value is returned; otherwise, y is evaluated and the resulting value is returned.

(Note that neither and nor or restrict the value and type they return to False and True, but rather return the last evaluated argument. This is sometimes useful, e.g., if s is a string that should be replaced by a default value if it is empty, the expression s or 'foo' yields the desired value. Because not has to invent a value anyway, it does not bother to return a value of the same type as its argument, so e.g., not 'foo' yields False, not".)

6.11 Conditional expressions

```
conditional_expression ::= or_test ["if" or_test "else" expression]
expression ::= conditional_expression | lambda_form
expression_nocond ::= or_test | lambda_form_nocond
```

Conditional expressions (sometimes called a "ternary operator") have the lowest priority of all Python operations.

⁴ Due to automatic garbage-collection, free lists, and the dynamic nature of descriptors, you may notice seemingly unusual behaviour in certain uses of the is operator, like those involving comparisons between instance methods, or constants. Check their documentation for more info.

The expression x if C else y first evaluates the condition, C (not x); if C is true, x is evaluated and its value is returned; otherwise, y is evaluated and its value is returned.

See PEP 308 for more details about conditional expressions.

6.12 Lambdas

```
lambda_form ::= "lambda" [parameter_list]: expression
lambda_form_nocond ::= "lambda" [parameter_list]: expression_nocond
```

Lambda forms (lambda expressions) have the same syntactic position as expressions. They are a shorthand to create anonymous functions; the expression lambda arguments: expression yields a function object. The unnamed object behaves like a function object defined with

```
def <lambda>(arguments):
    return expression
```

See section *Function definitions* for the syntax of parameter lists. Note that functions created with lambda forms cannot contain statements or annotations.

6.13 Expression lists

```
expression_list ::= expression ( "," expression ) * [","]
```

An expression list containing at least one comma yields a tuple. The length of the tuple is the number of expressions in the list. The expressions are evaluated from left to right.

The trailing comma is required only to create a single tuple (a.k.a. a *singleton*); it is optional in all other cases. A single expression without a trailing comma doesn't create a tuple, but rather yields the value of that expression. (To create an empty tuple, use an empty pair of parentheses: ().)

6.14 Evaluation order

Python evaluates expressions from left to right. Notice that while evaluating an assignment, the right-hand side is evaluated before the left-hand side.

In the following lines, expressions will be evaluated in the arithmetic order of their suffixes:

```
expr1, expr2, expr3, expr4
(expr1, expr2, expr3, expr4)
{expr1: expr2, expr3: expr4}
expr1 + expr2 * (expr3 - expr4)
expr1(expr2, expr3, *expr4, **expr5)
expr3, expr4 = expr1, expr2
```

6.15 Summary

The following table summarizes the operator precedences in Python, from lowest precedence (least binding) to highest precedence (most binding). Operators in the same box have the same precedence. Unless the syntax is explicitly given,

operators are binary. Operators in the same box group left to right (except for comparisons, including tests, which all have the same precedence and chain from left to right — see section *Comparisons* — and exponentiation, which groups from right to left).

Operator	Description
lambda	Lambda expression
if-else	Conditional expression
or	Boolean OR
and	Boolean AND
not x	Boolean NOT
in, not in, is, is not, <, <=, >, >=, !=,	Comparisons, including membership tests and identity tests,
==	
	Bitwise OR
^	Bitwise XOR
&	Bitwise AND
<<,>>	Shifts
+, -	Addition and subtraction
*, /, //, %	Multiplication, division, remainder ⁵
+x, -x, ~x	Positive, negative, bitwise NOT
**	Exponentiation ⁶
x[index],x[index:index],	Subscription, slicing, call, attribute reference
x(arguments), x.attribute	
(expressions),	Binding or tuple display, list display, dictionary display, set
[expressions],	display
{key:datum},	
{expressions}	

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⁵The % operator is also used for string formatting; the same precedence applies.

⁶The power operator ** binds less tightly than an arithmetic or bitwise unary operator on its right, that is, 2**-1 is 0.5.

SIMPLE STATEMENTS

Simple statements are comprised within a single logical line. Several simple statements may occur on a single line separated by semicolons. The syntax for simple statements is:

7.1 Expression statements

Expression statements are used (mostly interactively) to compute and write a value, or (usually) to call a procedure (a function that returns no meaningful result; in Python, procedures return the value <code>None</code>). Other uses of expression statements are allowed and occasionally useful. The syntax for an expression statement is:

```
expression_stmt ::= expression_list
```

An expression statement evaluates the expression list (which may be a single expression).

In interactive mode, if the value is not None, it is converted to a string using the built-in repr () function and the resulting string is written to standard output on a line by itself (except if the result is None, so that procedure calls do not cause any output.)

7.2 Assignment statements

Assignment statements are used to (re)bind names to values and to modify attributes or items of mutable objects:

(See section *Primaries* for the syntax definitions for the last three symbols.)

An assignment statement evaluates the expression list (remember that this can be a single expression or a commaseparated list, the latter yielding a tuple) and assigns the single resulting object to each of the target lists, from left to right.

Assignment is defined recursively depending on the form of the target (list). When a target is part of a mutable object (an attribute reference, subscription or slicing), the mutable object must ultimately perform the assignment and decide about its validity, and may raise an exception if the assignment is unacceptable. The rules observed by various types and the exceptions raised are given with the definition of the object types (see section *The standard type hierarchy*).

Assignment of an object to a target list, optionally enclosed in parentheses or square brackets, is recursively defined as follows.

- If the target list is a single target: The object is assigned to that target.
- If the target list is a comma-separated list of targets: The object must be an iterable with the same number of items as there are targets in the target list, and the items are assigned, from left to right, to the corresponding targets.
 - If the target list contains one target prefixed with an asterisk, called a "starred" target: The object must be a sequence with at least as many items as there are targets in the target list, minus one. The first items of the sequence are assigned, from left to right, to the targets before the starred target. The final items of the sequence are assigned to the targets after the starred target. A list of the remaining items in the sequence is then assigned to the starred target (the list can be empty).
 - Else: The object must be a sequence with the same number of items as there are targets in the target list, and the items are assigned, from left to right, to the corresponding targets.

Assignment of an object to a single target is recursively defined as follows.

- If the target is an identifier (name):
 - If the name does not occur in a global or nonlocal statement in the current code block: the name is bound to the object in the current local namespace.
 - Otherwise: the name is bound to the object in the global namespace or the outer namespace determined by nonlocal, respectively.

The name is rebound if it was already bound. This may cause the reference count for the object previously bound to the name to reach zero, causing the object to be deallocated and its destructor (if it has one) to be called.

- If the target is a target list enclosed in parentheses or in square brackets: The object must be an iterable with the same number of items as there are targets in the target list, and its items are assigned, from left to right, to the corresponding targets.
- If the target is an attribute reference: The primary expression in the reference is evaluated. It should yield an object with assignable attributes; if this is not the case, TypeError is raised. That object is then asked to assign the assigned object to the given attribute; if it cannot perform the assignment, it raises an exception (usually but

not necessarily AttributeError). Note: If the object is a class instance and the attribute reference occurs on both sides of the assignment operator, the RHS expression, a.x can access either an instance attribute or (if no instance attribute exists) a class attribute. The LHS target a.x is always set as an instance attribute, creating it if necessary. Thus, the two occurrences of a.x do not necessarily refer to the same attribute: if the RHS expression refers to a class attribute, the LHS creates a new instance attribute as the target of the assignment:

This description does not necessarily apply to descriptor attributes, such as properties created with property().

• If the target is a subscription: The primary expression in the reference is evaluated. It should yield either a mutable sequence object (such as a list) or a mapping object (such as a dictionary). Next, the subscript expression is evaluated.

If the primary is a mutable sequence object (such as a list), the subscript must yield an integer. If it is negative, the sequence's length is added to it. The resulting value must be a nonnegative integer less than the sequence's length, and the sequence is asked to assign the assigned object to its item with that index. If the index is out of range, IndexError is raised (assignment to a subscripted sequence cannot add new items to a list).

If the primary is a mapping object (such as a dictionary), the subscript must have a type compatible with the mapping's key type, and the mapping is then asked to create a key/datum pair which maps the subscript to the assigned object. This can either replace an existing key/value pair with the same key value, or insert a new key/value pair (if no key with the same value existed).

For user-defined objects, the <u>__setitem__()</u> method is called with appropriate arguments.

• If the target is a slicing: The primary expression in the reference is evaluated. It should yield a mutable sequence object (such as a list). The assigned object should be a sequence object of the same type. Next, the lower and upper bound expressions are evaluated, insofar they are present; defaults are zero and the sequence's length. The bounds should evaluate to integers. If either bound is negative, the sequence's length is added to it. The resulting bounds are clipped to lie between zero and the sequence's length, inclusive. Finally, the sequence object is asked to replace the slice with the items of the assigned sequence. The length of the slice may be different from the length of the assigned sequence, thus changing the length of the target sequence, if the object allows it.

CPython implementation detail: In the current implementation, the syntax for targets is taken to be the same as for expressions, and invalid syntax is rejected during the code generation phase, causing less detailed error messages.

WARNING: Although the definition of assignment implies that overlaps between the left-hand side and the right-hand side are 'safe' (for example a, b = b, a swaps two variables), overlaps within the collection of assigned-to variables are not safe! For instance, the following program prints [0, 2]:

```
x = [0, 1]

i = 0

i, x[i] = 1, 2

print(x)
```

See Also:

PEP 3132 - Extended Iterable Unpacking The specification for the *target feature.

7.2.1 Augmented assignment statements

Augmented assignment is the combination, in a single statement, of a binary operation and an assignment statement:

```
augmented_assignment_stmt ::= augtarget augop (expression_list | yield_expression)
augtarget ::= identifier | attributeref | subscription | slicing
augop ::= "+=" | "-=" | "*=" | "/=" | "/=" | "%=" | "**="
| ">>=" | "<<=" | "&=" | "^=" | "^=" | "|="</pre>
```

(See section *Primaries* for the syntax definitions for the last three symbols.)

An augmented assignment evaluates the target (which, unlike normal assignment statements, cannot be an unpacking) and the expression list, performs the binary operation specific to the type of assignment on the two operands, and assigns the result to the original target. The target is only evaluated once.

An augmented assignment expression like x += 1 can be rewritten as x = x + 1 to achieve a similar, but not exactly equal effect. In the augmented version, x is only evaluated once. Also, when possible, the actual operation is performed *in-place*, meaning that rather than creating a new object and assigning that to the target, the old object is modified instead.

With the exception of assigning to tuples and multiple targets in a single statement, the assignment done by augmented assignment statements is handled the same way as normal assignments. Similarly, with the exception of the possible *inplace* behavior, the binary operation performed by augmented assignment is the same as the normal binary operations.

For targets which are attribute references, the same *caveat about class and instance attributes* applies as for regular assignments.

7.3 The assert statement

Assert statements are a convenient way to insert debugging assertions into a program:

```
assert_stmt ::= "assert" expression ["," expression]
The simple form, assert expression, is equivalent to

if __debug__:
    if not expression: raise AssertionError
The extended form, assert expression1, expression2, is equivalent to

if __debug__:
    if not expression1: raise AssertionError(expression2)
```

These equivalences assume that __debug__ and AssertionError refer to the built-in variables with those names. In the current implementation, the built-in variable __debug__ is True under normal circumstances, False when optimization is requested (command line option -O). The current code generator emits no code for an assert statement when optimization is requested at compile time. Note that it is unnecessary to include the source code for the expression that failed in the error message; it will be displayed as part of the stack trace.

Assignments to ___debuq__ are illegal. The value for the built-in variable is determined when the interpreter starts.

7.4 The pass statement

```
pass stmt ::= "pass"
```

pass is a null operation — when it is executed, nothing happens. It is useful as a placeholder when a statement is required syntactically, but no code needs to be executed, for example:

```
def f(arg): pass  # a function that does nothing (yet)
class C: pass  # a class with no methods (yet)
```

7.5 The del statement

```
del_stmt ::= "del" target_list
```

Deletion is recursively defined very similar to the way assignment is defined. Rather than spelling it out in full details, here are some hints.

Deletion of a target list recursively deletes each target, from left to right.

Deletion of a name removes the binding of that name from the local or global namespace, depending on whether the name occurs in a global statement in the same code block. If the name is unbound, a NameError exception will be raised.

Deletion of attribute references, subscriptions and slicings is passed to the primary object involved; deletion of a slicing is in general equivalent to assignment of an empty slice of the right type (but even this is determined by the sliced object). Changed in version 3.2: Previously it was illegal to delete a name from the local namespace if it occurs as a free variable in a nested block.

7.6 The return statement

```
return stmt ::= "return" [expression list]
```

return may only occur syntactically nested in a function definition, not within a nested class definition.

If an expression list is present, it is evaluated, else None is substituted.

return leaves the current function call with the expression list (or None) as return value.

When return passes control out of a try statement with a finally clause, that finally clause is executed before really leaving the function.

In a generator function, the return statement indicates that the generator is done and will cause StopIteration to be raised. The returned value (if any) is used as an argument to construct StopIteration and becomes the StopIteration.value attribute.

7.7 The yield statement

```
yield_stmt ::= yield_expression
```

The yield statement is only used when defining a generator function, and is only used in the body of the generator function. Using a yield statement in a function definition is sufficient to cause that definition to create a generator function instead of a normal function.

When a generator function is called, it returns an iterator known as a generator iterator, or more commonly, a generator. The body of the generator function is executed by calling the next () function on the generator repeatedly until it raises an exception.

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When a yield statement is executed, the state of the generator is frozen and the value of expression_list is returned to next() 's caller. By "frozen" we mean that all local state is retained, including the current bindings of local variables, the instruction pointer, and the internal evaluation stack: enough information is saved so that the next time next() is invoked, the function can proceed exactly as if the yield statement were just another external call.

The yield statement is allowed in the try clause of a try ... finally construct. If the generator is not resumed before it is finalized (by reaching a zero reference count or by being garbage collected), the generator-iterator's close() method will be called, allowing any pending finally clauses to execute.

When yield from <expr> is used, it treats the supplied expression as a subiterator, producing values from it until the underlying iterator is exhausted.

Changed in version 3.3: Added yield from <expr> to delegate control flow to a subiterator

For full details of yield semantics, refer to the Yield expressions section.

See Also:

- **PEP 0255 Simple Generators** The proposal for adding generators and the yield statement to Python.
- **PEP 0342 Coroutines via Enhanced Generators** The proposal to enhance the API and syntax of generators, making them usable as simple coroutines.
- **PEP 0380 Syntax for Delegating to a Subgenerator** The proposal to introduce the yield_from syntax, making delegation to sub-generators easy.

7.8 The raise statement

```
raise_stmt ::= "raise" [expression ["from" expression]]
```

If no expressions are present, raise re-raises the last exception that was active in the current scope. If no exception is active in the current scope, a RuntimeError exception is raised indicating that this is an error.

Otherwise, raise evaluates the first expression as the exception object. It must be either a subclass or an instance of BaseException. If it is a class, the exception instance will be obtained when needed by instantiating the class with no arguments.

The type of the exception is the exception instance's class, the value is the instance itself.

A traceback object is normally created automatically when an exception is raised and attached to it as the __traceback__ attribute, which is writable. You can create an exception and set your own traceback in one step using the with_traceback() exception method (which returns the same exception instance, with its traceback set to its argument), like so:

```
raise Exception("foo occurred").with_traceback(tracebackobj)
```

The from clause is used for exception chaining: if given, the second *expression* must be another exception class or instance, which will then be attached to the raised exception as the __cause__ attribute (which is writable). If the raised exception is not handled, both exceptions will be printed:

```
>>> try:
... print(1 / 0)
... except Exception as exc:
... raise RuntimeError("Something bad happened") from exc
...
Traceback (most recent call last):
  File "<stdin>", line 2, in <module>
ZeroDivisionError: int division or modulo by zero
```

The above exception was the direct cause of the following exception:

```
Traceback (most recent call last):
   File "<stdin>", line 4, in <module>
RuntimeError: Something bad happened
```

A similar mechanism works implicitly if an exception is raised inside an exception handler: the previous exception is then attached as the new exception's __context__ attribute:

```
>>> try:
...     print(1 / 0)
... except:
...     raise RuntimeError("Something bad happened")
...
Traceback (most recent call last):
    File "<stdin>", line 2, in <module>
ZeroDivisionError: int division or modulo by zero

During handling of the above exception, another exception occurred:
Traceback (most recent call last):
    File "<stdin>", line 4, in <module>
RuntimeError: Something bad happened
```

Additional information on exceptions can be found in section *Exceptions*, and information about handling exceptions is in section *The try statement*.

7.9 The break statement

```
break_stmt ::= "break"
```

break may only occur syntactically nested in a for or while loop, but not nested in a function or class definition within that loop.

It terminates the nearest enclosing loop, skipping the optional else clause if the loop has one.

If a for loop is terminated by break, the loop control target keeps its current value.

When break passes control out of a try statement with a finally clause, that finally clause is executed before really leaving the loop.

7.10 The continue statement

```
continue_stmt ::= "continue"
```

continue may only occur syntactically nested in a for or while loop, but not nested in a function or class definition or finally clause within that loop. It continues with the next cycle of the nearest enclosing loop.

When continue passes control out of a try statement with a finally clause, that finally clause is executed before really starting the next loop cycle.

7.11 The import statement

```
"import" module ["as" name] ("," module ["as" name] )*
import stmt
                 ..=
                      | "from" relative_module "import" identifier ["as" name]
                       ("," identifier ["as" name] ) *
                      "from" relative_module "import" "(" identifier ["as" name]
                       ( "," identifier ["as" name] ) * [","] ")"
                      | "from" module "import" "*"
                      (identifier ".") ★ identifier
module
                 ::=
                      "."* module | "."+
relative_module ::=
                      identifier
name
                 ::=
```

The basic import statement (no from clause) is executed in two steps:

- 1. find a module, loading and initializing it if necessary
- 2. define a name or names in the local namespace for the scope where the import statement occurs.

When the statement contains multiple clauses (separated by commas) the two steps are carried out separately for each clause, just as though the clauses had been separated out into individiual import statements.

The details of the first step, finding and loading modules is described in greater detail in the section on the *import system*, which also describes the various types of packages and modules that can be imported, as well as all the hooks that can be used to customize the import system. Note that failures in this step may indicate either that the module could not be located, *or* that an error occurred while initializing the module, which includes execution of the module's code.

If the requested module is retrieved successfully, it will be made available in the local namespace in one of three ways:

- If the module name is followed by as, then the name following as is bound directly to the imported module.
- If no other name is specified, and the module being imported is a top level module, the module's name is bound in the local namespace as a reference to the imported module
- If the module being imported is *not* a top level module, then the name of the top level package that contains the module is bound in the local namespace as a reference to the top level package. The imported module must be accessed using its full qualified name rather than directly

The from form uses a slightly more complex process:

- 1. find the module specified in the from clause loading and initializing it if necessary;
- 2. for each of the identifiers specified in the import clauses:
 - (a) check if the imported module has an attribute by that name
 - (b) if not, attempt to import a submodule with that name and then check the imported module again for that attribute
 - (c) if the attribute is not found, ImportError is raised.
 - (d) otherwise, a reference to that value is bound in the local namespace, using the name in the as clause if it is present, otherwise using the attribute name

Examples:

```
import foo  # foo imported and bound locally
import foo.bar.baz  # foo.bar.baz imported, foo bound locally
import foo.bar.baz as fbb  # foo.bar.baz imported and bound as fbb
from foo.bar import baz  # foo.bar.baz imported and bound as baz
from foo import attr  # foo imported and foo.attr bound as attr
```

If the list of identifiers is replaced by a star (' *'), all public names defined in the module are bound in the local namespace for the scope where the import statement occurs.

The *public names* defined by a module are determined by checking the module's namespace for a variable named __all__; if defined, it must be a sequence of strings which are names defined or imported by that module. The names given in __all__ are all considered public and are required to exist. If __all__ is not defined, the set of public names includes all names found in the module's namespace which do not begin with an underscore character ('_'). __all__ should contain the entire public API. It is intended to avoid accidentally exporting items that are not part of the API (such as library modules which were imported and used within the module).

The from form with * may only occur in a module scope. Attempting to use it in class or function definitions will raise a SyntaxError.

The *public names* defined by a module are determined by checking the module's namespace for a variable named __all__; if defined, it must be a sequence of strings which are names defined or imported by that module. The names given in __all__ are all considered public and are required to exist. If __all__ is not defined, the set of public names includes all names found in the module's namespace which do not begin with an underscore character ('_'). __all__ should contain the entire public API. It is intended to avoid accidentally exporting items that are not part of the API (such as library modules which were imported and used within the module).

The from form with * may only occur in a module scope. The wild card form of import — import * — is only allowed at the module level. Attempting to use it in class or function definitions will raise a SyntaxError.

When specifying what module to import you do not have to specify the absolute name of the module. When a module or package is contained within another package it is possible to make a relative import within the same top package without having to mention the package name. By using leading dots in the specified module or package after from you can specify how high to traverse up the current package hierarchy without specifying exact names. One leading dot means the current package where the module making the import exists. Two dots means up one package level. Three dots is up two levels, etc. So if you execute from . import mod from a module in the pkg package then you will end up importing pkg.mod. If you execute from ..subpkg2 import mod from within pkg.subpkg1 you will import pkg.subpkg2.mod. The specification for relative imports is contained within PEP 328.

importlib.import_module() is provided to support applications that determine which modules need to be loaded dynamically.

7.11.1 Future statements

A *future statement* is a directive to the compiler that a particular module should be compiled using syntax or semantics that will be available in a specified future release of Python. The future statement is intended to ease migration to future versions of Python that introduce incompatible changes to the language. It allows use of the new features on a per-module basis before the release in which the feature becomes standard.

A future statement must appear near the top of the module. The only lines that can appear before a future statement are:

- the module docstring (if any),
- · comments,
- · blank lines, and

• other future statements.

The features recognized by Python 3.0 are absolute_import, division, generators, unicode_literals, print_function, nested_scopes and with_statement. They are all redundant because they are always enabled, and only kept for backwards compatibility.

A future statement is recognized and treated specially at compile time: Changes to the semantics of core constructs are often implemented by generating different code. It may even be the case that a new feature introduces new incompatible syntax (such as a new reserved word), in which case the compiler may need to parse the module differently. Such decisions cannot be pushed off until runtime.

For any given release, the compiler knows which feature names have been defined, and raises a compile-time error if a future statement contains a feature not known to it.

The direct runtime semantics are the same as for any import statement: there is a standard module __future__, described later, and it will be imported in the usual way at the time the future statement is executed.

The interesting runtime semantics depend on the specific feature enabled by the future statement.

Note that there is nothing special about the statement:

```
import __future__ [as name]
```

That is not a future statement; it's an ordinary import statement with no special semantics or syntax restrictions.

Code compiled by calls to the built-in functions <code>exec()</code> and <code>compile()</code> that occur in a module M containing a future statement will, by default, use the new syntax or semantics associated with the future statement. This can be controlled by optional arguments to <code>compile()</code> — see the documentation of that function for details.

A future statement typed at an interactive interpreter prompt will take effect for the rest of the interpreter session. If an interpreter is started with the -i option, is passed a script name to execute, and the script includes a future statement, it will be in effect in the interactive session started after the script is executed.

See Also:

PEP 236 - Back to the __future__ The original proposal for the __future__ mechanism.

7.12 The global statement

```
global_stmt ::= "global" identifier ("," identifier) *
```

The global statement is a declaration which holds for the entire current code block. It means that the listed identifiers are to be interpreted as globals. It would be impossible to assign to a global variable without global, although free variables may refer to globals without being declared global.

Names listed in a global statement must not be used in the same code block textually preceding that global statement.

Names listed in a global statement must not be defined as formal parameters or in a for loop control target, class definition, function definition, or import statement.

CPython implementation detail: The current implementation does not enforce the latter two restrictions, but programs should not abuse this freedom, as future implementations may enforce them or silently change the meaning of the program.

Programmer's note: the global is a directive to the parser. It applies only to code parsed at the same time as the global statement. In particular, a global statement contained in a string or code object supplied to the built-in exec() function does not affect the code block *containing* the function call, and code contained in such a string is unaffected by global statements in the code containing the function call. The same applies to the eval() and compile() functions.

7.13 The nonlocal statement

```
nonlocal_stmt ::= "nonlocal" identifier ("," identifier) *
```

The nonlocal statement causes the listed identifiers to refer to previously bound variables in the nearest enclosing scope. This is important because the default behavior for binding is to search the local namespace first. The statement allows encapsulated code to rebind variables outside of the local scope besides the global (module) scope.

Names listed in a nonlocal statement, unlike to those listed in a global statement, must refer to pre-existing bindings in an enclosing scope (the scope in which a new binding should be created cannot be determined unambiguously).

Names listed in a nonlocal statement must not collide with pre-existing bindings in the local scope.

See Also:

PEP 3104 - Access to Names in Outer Scopes The specification for the nonlocal statement.

COMPOUND STATEMENTS

Compound statements contain (groups of) other statements; they affect or control the execution of those other statements in some way. In general, compound statements span multiple lines, although in simple incarnations a whole compound statement may be contained in one line.

The if, while and for statements implement traditional control flow constructs. try specifies exception handlers and/or cleanup code for a group of statements, while the with statement allows the execution of initialization and finalization code around a block of code. Function and class definitions are also syntactically compound statements.

Compound statements consist of one or more 'clauses.' A clause consists of a header and a 'suite.' The clause headers of a particular compound statement are all at the same indentation level. Each clause header begins with a uniquely identifying keyword and ends with a colon. A suite is a group of statements controlled by a clause. A suite can be one or more semicolon-separated simple statements on the same line as the header, following the header's colon, or it can be one or more indented statements on subsequent lines. Only the latter form of suite can contain nested compound statements; the following is illegal, mostly because it wouldn't be clear to which if clause a following else clause would belong:

```
if test1: if test2: print(x)
```

Also note that the semicolon binds tighter than the colon in this context, so that in the following example, either all or none of the print () calls are executed:

```
if x < y < z: print(x); print(y); print(z)
Summarizing:</pre>
```

```
if stmt
compound stmt ::=
                     | while_stmt
                     | for_stmt
                     | try_stmt
                     | with_stmt
                     | funcdef
                     | classdef
                    stmt list NEWLINE | NEWLINE INDENT statement+ DEDENT
suite
               ::=
                     stmt_list NEWLINE | compound_stmt
statement
               ::=
                    simple_stmt (";" simple_stmt)* [";"]
stmt_list
               ::=
```

Note that statements always end in a NEWLINE possibly followed by a DEDENT. Also note that optional continuation clauses always begin with a keyword that cannot start a statement, thus there are no ambiguities (the 'dangling else' problem is solved in Python by requiring nested if statements to be indented).

The formatting of the grammar rules in the following sections places each clause on a separate line for clarity.

8.1 The if statement

The if statement is used for conditional execution:

It selects exactly one of the suites by evaluating the expressions one by one until one is found to be true (see section *Boolean operations* for the definition of true and false); then that suite is executed (and no other part of the if statement is executed or evaluated). If all expressions are false, the suite of the else clause, if present, is executed.

8.2 The while statement

The while statement is used for repeated execution as long as an expression is true:

```
while_stmt ::= "while" expression ":" suite
["else" ":" suite]
```

This repeatedly tests the expression and, if it is true, executes the first suite; if the expression is false (which may be the first time it is tested) the suite of the else clause, if present, is executed and the loop terminates.

A break statement executed in the first suite terminates the loop without executing the else clause's suite. A continue statement executed in the first suite skips the rest of the suite and goes back to testing the expression.

8.3 The for statement

The for statement is used to iterate over the elements of a sequence (such as a string, tuple or list) or other iterable object:

```
for_stmt ::= "for" target_list "in" expression_list ":" suite
["else" ":" suite]
```

The expression list is evaluated once; it should yield an iterable object. An iterator is created for the result of the expression_list. The suite is then executed once for each item provided by the iterator, in the order of ascending indices. Each item in turn is assigned to the target list using the standard rules for assignments (see *Assignment statements*), and then the suite is executed. When the items are exhausted (which is immediately when the sequence is empty or an iterator raises a StopIteration exception), the suite in the else clause, if present, is executed, and the loop terminates.

A break statement executed in the first suite terminates the loop without executing the else clause's suite. A continue statement executed in the first suite skips the rest of the suite and continues with the next item, or with the else clause if there was no next item.

The suite may assign to the variable(s) in the target list; this does not affect the next item assigned to it.

Names in the target list are not deleted when the loop is finished, but if the sequence is empty, it will not have been assigned to at all by the loop. Hint: the built-in function range() returns an iterator of integers suitable to emulate the effect of Pascal's for i := a to b do; e.g., list (range(3)) returns the list [0, 1, 2].

Note: There is a subtlety when the sequence is being modified by the loop (this can only occur for mutable sequences, i.e. lists). An internal counter is used to keep track of which item is used next, and this is incremented on each iteration.

When this counter has reached the length of the sequence the loop terminates. This means that if the suite deletes the current (or a previous) item from the sequence, the next item will be skipped (since it gets the index of the current item which has already been treated). Likewise, if the suite inserts an item in the sequence before the current item, the current item will be treated again the next time through the loop. This can lead to nasty bugs that can be avoided by making a temporary copy using a slice of the whole sequence, e.g.,

```
for x in a[:]:
    if x < 0: a.remove(x)</pre>
```

8.4 The try statement

The try statement specifies exception handlers and/or cleanup code for a group of statements:

The except clause(s) specify one or more exception handlers. When no exception occurs in the try clause, no exception handler is executed. When an exception occurs in the try suite, a search for an exception handler is started. This search inspects the except clauses in turn until one is found that matches the exception. An expression-less except clause, if present, must be last; it matches any exception. For an except clause with an expression, that expression is evaluated, and the clause matches the exception if the resulting object is "compatible" with the exception. An object is compatible with an exception if it is the class or a base class of the exception object or a tuple containing an item compatible with the exception.

If no except clause matches the exception, the search for an exception handler continues in the surrounding code and on the invocation stack. ¹

If the evaluation of an expression in the header of an except clause raises an exception, the original search for a handler is canceled and a search starts for the new exception in the surrounding code and on the call stack (it is treated as if the entire try statement raised the exception).

When a matching except clause is found, the exception is assigned to the target specified after the as keyword in that except clause, if present, and the except clause's suite is executed. All except clauses must have an executable block. When the end of this block is reached, execution continues normally after the entire try statement. (This means that if two nested handlers exist for the same exception, and the exception occurs in the try clause of the inner handler, the outer handler will not handle the exception.)

When an exception has been assigned using as target, it is cleared at the end of the except clause. This is as if

```
except E as N:
foo

was translated to

except E as N:
try:
foo
```

¹ The exception is propagated to the invocation stack unless there is a finally clause which happens to raise another exception. That new exception causes the old one to be lost.

```
finally:
del N
```

This means the exception must be assigned to a different name to be able to refer to it after the except clause. Exceptions are cleared because with the traceback attached to them, they form a reference cycle with the stack frame, keeping all locals in that frame alive until the next garbage collection occurs.

Before an except clause's suite is executed, details about the exception are stored in the sys module and can be access via sys.exc_info().sys.exc_info() returns a 3-tuple consisting of the exception class, the exception instance and a traceback object (see section *The standard type hierarchy*) identifying the point in the program where the exception occurred. sys.exc_info() values are restored to their previous values (before the call) when returning from a function that handled an exception.

The optional else clause is executed if and when control flows off the end of the try clause. 2 Exceptions in the else clause are not handled by the preceding except clauses.

If finally is present, it specifies a 'cleanup' handler. The try clause is executed, including any except and else clauses. If an exception occurs in any of the clauses and is not handled, the exception is temporarily saved. The finally clause is executed. If there is a saved exception it is re-raised at the end of the finally clause. If the finally clause raises another exception, the saved exception is set as the context of the new exception. If the finally clause executes a return or break statement, the saved exception is discarded:

The exception information is not available to the program during execution of the finally clause.

When a return, break or continue statement is executed in the try suite of a try...finally statement, the finally clause is also executed 'on the way out.' A continue statement is illegal in the finally clause. (The reason is a problem with the current implementation — this restriction may be lifted in the future).

Additional information on exceptions can be found in section *Exceptions*, and information on using the raise statement to generate exceptions may be found in section *The raise statement*.

8.5 The with statement

The with statement is used to wrap the execution of a block with methods defined by a context manager (see section *With Statement Context Managers*). This allows common try...except...finally usage patterns to be encapsulated for convenient reuse.

```
with_stmt ::= "with" with_item ("," with_item) * ":" suite
with_item ::= expression ["as" target]
```

The execution of the with statement with one "item" proceeds as follows:

- 1. The context expression (the expression given in the with item) is evaluated to obtain a context manager.
- 2. The context manager's __exit__() is loaded for later use.
- 3. The context manager's __enter__() method is invoked.

² Currently, control "flows off the end" except in the case of an exception or the execution of a return, continue, or break statement.

4. If a target was included in the with statement, the return value from __enter__() is assigned to it.

Note: The with statement guarantees that if the __enter__() method returns without an error, then __exit__() will always be called. Thus, if an error occurs during the assignment to the target list, it will be treated the same as an error occurring within the suite would be. See step 6 below.

- 5. The suite is executed.
- 6. The context manager's __exit__() method is invoked. If an exception caused the suite to be exited, its type, value, and traceback are passed as arguments to __exit__(). Otherwise, three None arguments are supplied.

If the suite was exited due to an exception, and the return value from the $__{exit}_{}()$ method was false, the exception is reraised. If the return value was true, the exception is suppressed, and execution continues with the statement following the with statement.

If the suite was exited for any reason other than an exception, the return value from __exit__() is ignored, and execution proceeds at the normal location for the kind of exit that was taken.

With more than one item, the context managers are processed as if multiple with statements were nested:

```
with A() as a, B() as b:
    suite

is equivalent to

with A() as a:
    with B() as b:
    suite
```

Changed in version 3.1: Support for multiple context expressions.

See Also:

PEP 0343 - The "with" statement The specification, background, and examples for the Python with statement.

8.6 Function definitions

A function definition defines a user-defined function object (see section *The standard type hierarchy*):

```
[decorators] "def" funcname "(" [parameter_list] ")" ["->" expres
funcdef
                     decorator+
decorators
                ::=
                     "@" dotted_name ["(" [parameter_list [","]] ")"] NEWLINE
decorator
                ::=
dotted name
                ::=
                     identifier ("."
                                      identifier) *
                     (defparameter ",") *
parameter_list ::=
                     ( "*" [parameter] ("," defparameter)* ["," "**" parameter]
                     | "**" parameter
                     | defparameter [","] )
parameter
                ::=
                     identifier [":" expression]
defparameter
                     parameter ["=" expression]
                ::=
                     identifier
funchame
                ::=
```

A function definition is an executable statement. Its execution binds the function name in the current local namespace to a function object (a wrapper around the executable code for the function). This function object contains a reference to the current global namespace as the global namespace to be used when the function is called.

The function definition does not execute the function body; this gets executed only when the function is called. ³

A function definition may be wrapped by one or more *decorator* expressions. Decorator expressions are evaluated when the function is defined, in the scope that contains the function definition. The result must be a callable, which is invoked with the function object as the only argument. The returned value is bound to the function name instead of the function object. Multiple decorators are applied in nested fashion. For example, the following code

```
@f1(arg)
@f2
def func(): pass
is equivalent to
def func(): pass
func = f1(arg)(f2(func))
```

When one or more parameters have the form *parameter* = *expression*, the function is said to have "default parameter values." For a parameter with a default value, the corresponding argument may be omitted from a call, in which case the parameter's default value is substituted. If a parameter has a default value, all following parameters up until the "*" must also have a default value — this is a syntactic restriction that is not expressed by the grammar.

Default parameter values are evaluated when the function definition is executed. This means that the expression is evaluated once, when the function is defined, and that the same "pre-computed" value is used for each call. This is especially important to understand when a default parameter is a mutable object, such as a list or a dictionary: if the function modifies the object (e.g. by appending an item to a list), the default value is in effect modified. This is generally not what was intended. A way around this is to use None as the default, and explicitly test for it in the body of the function, e.g.:

```
def whats_on_the_telly(penguin=None):
    if penguin is None:
        penguin = []
    penguin.append("property of the zoo")
    return penguin
```

Function call semantics are described in more detail in section *Calls*. A function call always assigns values to all parameters mentioned in the parameter list, either from position arguments, from keyword arguments, or from default values. If the form "*identifier" is present, it is initialized to a tuple receiving any excess positional parameters, defaulting to the empty tuple. If the form "**identifier" is present, it is initialized to a new dictionary receiving any excess keyword arguments, defaulting to a new empty dictionary. Parameters after "*" or "*identifier" are keyword-only parameters and may only be passed used keyword arguments.

Parameters may have annotations of the form ": expression" following the parameter name. Any parameter may have an annotation even those of the form *identifier or **identifier. Functions may have "return" annotation of the form "-> expression" after the parameter list. These annotations can be any valid Python expression and are evaluated when the function definition is executed. Annotations may be evaluated in a different order than they appear in the source code. The presence of annotations does not change the semantics of a function. The annotation values are available as values of a dictionary keyed by the parameters' names in the __annotations__ attribute of the function object.

It is also possible to create anonymous functions (functions not bound to a name), for immediate use in expressions. This uses lambda forms, described in section *Lambdas*. Note that the lambda form is merely a shorthand for a simplified function definition; a function defined in a "def" statement can be passed around or assigned to another name just like a function defined by a lambda form. The "def" form is actually more powerful since it allows the execution of multiple statements and annotations.

Programmer's note: Functions are first-class objects. A "def" form executed inside a function definition defines a local function that can be returned or passed around. Free variables used in the nested function can access the local

³ A string literal appearing as the first statement in the function body is transformed into the function's __doc__ attribute and therefore the function's *docstring*.

variables of the function containing the def. See section *Naming and binding* for details.

See Also:

PEP 3107 - Function Annotations The original specification for function annotations.

8.7 Class definitions

A class definition defines a class object (see section *The standard type hierarchy*):

```
classdef ::= [decorators] "class" classname [inheritance] ":" suite
inheritance ::= "(" [parameter_list] ")"
classname ::= identifier
```

A class definition is an executable statement. The inheritance list usually gives a list of base classes (see *Customizing class creation* for more advanced uses), so each item in the list should evaluate to a class object which allows subclassing. Classes without an inheritance list inherit, by default, from the base class object; hence,

```
class Foo:
    pass
is equivalent to
class Foo(object):
    pass
```

The class's suite is then executed in a new execution frame (see *Naming and binding*), using a newly created local namespace and the original global namespace. (Usually, the suite contains mostly function definitions.) When the class's suite finishes execution, its execution frame is discarded but its local namespace is saved. ⁴ A class object is then created using the inheritance list for the base classes and the saved local namespace for the attribute dictionary. The class name is bound to this class object in the original local namespace.

Class creation can be customized heavily using metaclasses.

Classes can also be decorated: just like when decorating functions,

```
@f1(arg)
@f2
class Foo: pass
is equivalent to
class Foo: pass
Foo = f1(arg)(f2(Foo))
```

The evaluation rules for the decorator expressions are the same as for function decorators. The result must be a class object, which is then bound to the class name.

Programmer's note: Variables defined in the class definition are class attributes; they are shared by instances. Instance attributes can be set in a method with self.name = value. Both class and instance attributes are accessible through the notation "self.name", and an instance attribute hides a class attribute with the same name when accessed in this way. Class attributes can be used as defaults for instance attributes, but using mutable values there can lead to unexpected results. *Descriptors* can be used to create instance variables with different implementation details.

See Also:

```
PEP 3115 - Metaclasses in Python 3
```

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⁴ A string literal appearing as the first statement in the class body is transformed into the namespace's __doc__ item and therefore the class's *docstring*.

PEP 3129 - Class Decorators

CHAPTER

NINE

TOP-LEVEL COMPONENTS

The Python interpreter can get its input from a number of sources: from a script passed to it as standard input or as program argument, typed in interactively, from a module source file, etc. This chapter gives the syntax used in these cases.

9.1 Complete Python programs

While a language specification need not prescribe how the language interpreter is invoked, it is useful to have a notion of a complete Python program. A complete Python program is executed in a minimally initialized environment: all built-in and standard modules are available, but none have been initialized, except for sys (various system services), built-ins (built-in functions, exceptions and None) and __main__. The latter is used to provide the local and global namespace for execution of the complete program.

The syntax for a complete Python program is that for file input, described in the next section.

The interpreter may also be invoked in interactive mode; in this case, it does not read and execute a complete program but reads and executes one statement (possibly compound) at a time. The initial environment is identical to that of a complete program; each statement is executed in the namespace of __main__.

Under Unix, a complete program can be passed to the interpreter in three forms: with the -c string command line option, as a file passed as the first command line argument, or as standard input. If the file or standard input is a tty device, the interpreter enters interactive mode; otherwise, it executes the file as a complete program.

9.2 File input

All input read from non-interactive files has the same form:

```
file_input ::= (NEWLINE | statement) *
```

This syntax is used in the following situations:

- when parsing a complete Python program (from a file or from a string);
- when parsing a module;
- when parsing a string passed to the exec () function;

9.3 Interactive input

Input in interactive mode is parsed using the following grammar:

```
interactive_input ::= [stmt_list] NEWLINE | compound_stmt NEWLINE
```

Note that a (top-level) compound statement must be followed by a blank line in interactive mode; this is needed to help the parser detect the end of the input.

9.4 Expression input

There are two forms of expression input. Both ignore leading whitespace. The string argument to eval () must have the following form:

```
eval_input ::= expression_list NEWLINE*
```

Note: to read 'raw' input line without interpretation, you can use the readline () method of file objects, including sys.stdin.

FULL GRAMMAR SPECIFICATION

This is the full Python grammar, as it is read by the parser generator and used to parse Python source files:

```
# Grammar for Python
# Note: Changing the grammar specified in this file will most likely
         require corresponding changes in the parser module
         (../Modules/parsermodule.c). If you can't make the changes to
         that module yourself, please co-ordinate the required changes
         with someone who can; ask around on python-dev for help. Fred
         Drake <fdrake@acm.org> will probably be listening there.
# NOTE WELL: You should also follow all the steps listed in PEP 306,
# "How to Change Python's Grammar"
# Start symbols for the grammar:
       single_input is a single interactive statement;
       file input is a module or sequence of commands read from an input file;
       eval_input is the input for the eval() functions.
# NB: compound_stmt in single_input is followed by extra NEWLINE!
single_input: NEWLINE | simple_stmt | compound_stmt NEWLINE
file_input: (NEWLINE | stmt) * ENDMARKER
eval_input: testlist NEWLINE* ENDMARKER
decorator: '@' dotted name [ '(' [arglist] ')' ] NEWLINE
decorators: decorator+
decorated: decorators (classdef | funcdef)
funcdef: 'def' NAME parameters ['->' test] ':' suite
parameters: '(' [typedargslist] ')'
typedargslist: (tfpdef ['=' test] (',' tfpdef ['=' test])* [','
       ['*' [tfpdef] (',' tfpdef ['=' test]) * [',' '**' tfpdef] | '**' tfpdef]]
     | '*' [tfpdef] (',' tfpdef ['=' test])* [',' '**' tfpdef] | '**' tfpdef)
tfpdef: NAME [':' test]
varargslist: (vfpdef ['=' test] (',' vfpdef ['=' test])* [','
       ['*' [vfpdef] (',' vfpdef ['=' test])* [',' '**' vfpdef] | '**' vfpdef]]
     | '*' [vfpdef] (',' vfpdef ['=' test]) * [',' '**' vfpdef] | '**' vfpdef)
vfpdef: NAME
stmt: simple_stmt | compound_stmt
simple_stmt: small_stmt (';' small_stmt) * [';'] NEWLINE
small_stmt: (expr_stmt | del_stmt | pass_stmt | flow_stmt |
```

```
import_stmt | global_stmt | nonlocal_stmt | assert_stmt)
expr_stmt: testlist_star_expr (augassign (yield_expr|testlist) |
                     ('=' (yield expr|testlist star expr))*)
testlist_star_expr: (test|star_expr) (',' (test|star_expr))* [',']
augassign: ('+=' | '-=' | '*=' | '/=' | '%=' | '&=' | '\=' |
            '<<=' | '>>=' | ' **=' | ' //=')
# For normal assignments, additional restrictions enforced by the interpreter
del stmt: 'del' exprlist
pass_stmt: 'pass'
flow_stmt: break_stmt | continue_stmt | return_stmt | raise_stmt | yield_stmt
break_stmt: 'break'
continue_stmt: 'continue'
return_stmt: 'return' [testlist]
yield_stmt: yield_expr
raise_stmt: 'raise' [test ['from' test]]
import_stmt: import_name | import_from
import_name: 'import' dotted_as_names
# note below: the ('.' | '...') is necessary because '...' is tokenized as ELLIPSIS
import_from: ('from' (('.' | '...') * dotted_name | ('.' | '...') +)
              'import' ('*' | '(' import_as_names ')' | import_as_names))
import_as_name: NAME ['as' NAME]
dotted_as_name: dotted_name ['as' NAME]
import_as_names: import_as_name (',' import_as_name) * [',']
dotted_as_names: dotted_as_name (',' dotted_as_name) *
dotted name: NAME ('.' NAME) *
global stmt: 'global' NAME (',' NAME) *
nonlocal_stmt: 'nonlocal' NAME (',' NAME) *
assert_stmt: 'assert' test [',' test]
compound_stmt: if_stmt | while_stmt | for_stmt | try_stmt | with_stmt | funcdef | classdef
if_stmt: 'if' test ':' suite ('elif' test ':' suite)* ['else' ':' suite]
while_stmt: 'while' test ':' suite ['else' ':' suite]
for_stmt: 'for' exprlist 'in' testlist ':' suite ['else' ':' suite]
try_stmt: ('try' ':' suite
           ((except_clause ':' suite)+
            ['else' ':' suite]
            ['finally' ':' suite] |
           'finally' ':' suite))
with_stmt: 'with' with_item (',' with_item) * ':' suite
with_item: test ['as' expr]
# NB compile.c makes sure that the default except clause is last
except_clause: 'except' [test ['as' NAME]]
suite: simple_stmt | NEWLINE INDENT stmt+ DEDENT
test: or_test ['if' or_test 'else' test] | lambdef
test_nocond: or_test | lambdef_nocond
lambdef: 'lambda' [varargslist] ':' test
lambdef_nocond: 'lambda' [vararqslist] ':' test_nocond
or_test: and_test ('or' and_test) *
and_test: not_test ('and' not_test) *
not_test: 'not' not_test | comparison
comparison: expr (comp op expr) *
# <> isn't actually a valid comparison operator in Python. It's here for the
# sake of a ___future__ import described in PEP 401
```

```
comp op: '<'|'>'|'=='|'>='|'<='|'<>'|'!='|'in'|'not' 'in'|'is'|'is' 'not'
star_expr: '*' expr
expr: xor expr (' | ' xor expr) *
xor_expr: and_expr ('^' and_expr) *
and_expr: shift_expr ('&' shift_expr) *
shift expr: arith expr (('<<'|'>>') arith expr) *
arith expr: term (('+')'-') term) *
term: factor (('*'|'/'|'%'|'//') factor)*
factor: ('+'|'-'|'\sim') factor | power
power: atom trailer* ['**' factor]
atom: ('(' [yield_expr|testlist_comp] ')' |
       '[' [testlist_comp] ']' |
       '{' [dictorsetmaker] '}' |
      NAME | NUMBER | STRING+ | '...' | 'None' | 'True' | 'False')
testlist_comp: (test|star_expr) ( comp_for | (',' (test|star_expr))* [','] )
trailer: '(' [arglist] ')' | '[' subscriptlist ']' | '.' NAME
subscriptlist: subscript (',' subscript)* [',']
subscript: test | [test] ':' [test] [sliceop]
sliceop: ':' [test]
exprlist: (expr|star_expr) (',' (expr|star_expr))* [',']
testlist: test (',' test)* [',']
dictorsetmaker: ( (test ':' test (comp_for | (',' test ':' test) * [','])) |
                  (test (comp_for | (',' test) * [','])) )
classdef: 'class' NAME ['(' [arglist] ')'] ':' suite
arglist: (argument ',')* (argument [',']
                         |'*' test (',' argument)* [',' '**' test]
                         |'**' test)
# The reason that keywords are test nodes instead of NAME is that using NAME
# results in an ambiguity. ast.c makes sure it's a NAME.
argument: test [comp_for] | test '=' test # Really [keyword '='] test
comp_iter: comp_for | comp_if
comp_for: 'for' exprlist 'in' or_test [comp_iter]
comp_if: 'if' test_nocond [comp_iter]
# not used in grammar, but may appear in "node" passed from Parser to Compiler
encoding_decl: NAME
yield_expr: 'yield' [yield_arg]
vield arg: 'from' test | testlist
```

GLOSSARY

- >>> The default Python prompt of the interactive shell. Often seen for code examples which can be executed interactively in the interpreter.
- ... The default Python prompt of the interactive shell when entering code for an indented code block or within a pair of matching left and right delimiters (parentheses, square brackets or curly braces).
- **2to3** A tool that tries to convert Python 2.x code to Python 3.x code by handling most of the incompatibilities which can be detected by parsing the source and traversing the parse tree.
 - 2to3 is available in the standard library as lib2to3; a standalone entry point is provided as Tools/scripts/2to3. See 2to3-reference.
- abstract base class Abstract base classes complement duck-typing by providing a way to define interfaces when other techniques like hasattr() would be clumsy or subtly wrong (for example with magic methods). ABCs introduce virtual subclasses, which are classes that don't inherit from a class but are still recognized by isinstance() and issubclass(); see the abc module documentation. Python comes with many built-in ABCs for data structures (in the collections.abc module), numbers (in the numbers module), streams (in the io module), import finders and loaders (in the importlib.abc module). You can create your own ABCs with the abc module.
- **argument** A value passed to a function or method, assigned to a named local variable in the function body. A function or method may have both positional arguments and keyword arguments in its definition. Positional and keyword arguments may be variable-length: * accepts or passes (if in the function definition or call) several positional arguments in a list, while ** does the same for keyword arguments in a dictionary.
 - Any expression may be used within the argument list, and the evaluated value is passed to the local variable.
- **attribute** A value associated with an object which is referenced by name using dotted expressions. For example, if an object o has an attribute a it would be referenced as o.a.
- **BDFL** Benevolent Dictator For Life, a.k.a. Guido van Rossum, Python's creator.
- bytecode Python source code is compiled into bytecode, the internal representation of a Python program in the CPython interpreter. The bytecode is also cached in .pyc and .pyc files so that executing the same file is faster the second time (recompilation from source to bytecode can be avoided). This "intermediate language" is said to run on a *virtual machine* that executes the machine code corresponding to each bytecode. Do note that bytecodes are not expected to work between different Python virtual machines, nor to be stable between Python releases.
 - A list of bytecode instructions can be found in the documentation for the dis module.
- **class** A template for creating user-defined objects. Class definitions normally contain method definitions which operate on instances of the class.
- **coercion** The implicit conversion of an instance of one type to another during an operation which involves two arguments of the same type. For example, int (3.15) converts the floating point number to the integer 3, but

in 3+4.5, each argument is of a different type (one int, one float), and both must be converted to the same type before they can be added or it will raise a TypeError. Without coercion, all arguments of even compatible types would have to be normalized to the same value by the programmer, e.g., float (3) +4.5 rather than just 3+4.5.

- complex number An extension of the familiar real number system in which all numbers are expressed as a sum of a real part and an imaginary part. Imaginary numbers are real multiples of the imaginary unit (the square root of -1), often written i in mathematics or j in engineering. Python has built-in support for complex numbers, which are written with this latter notation; the imaginary part is written with a j suffix, e.g., 3+1j. To get access to complex equivalents of the math module, use cmath. Use of complex numbers is a fairly advanced mathematical feature. If you're not aware of a need for them, it's almost certain you can safely ignore them.
- context manager An object which controls the environment seen in a with statement by defining __enter__() and __exit__() methods. See PEP 343.
- **CPython** The canonical implementation of the Python programming language, as distributed on python.org. The term "CPython" is used when necessary to distinguish this implementation from others such as Jython or IronPython.
- **decorator** A function returning another function, usually applied as a function transformation using the @wrapper syntax. Common examples for decorators are classmethod() and staticmethod().

The decorator syntax is merely syntactic sugar, the following two function definitions are semantically equivalent:

The same concept exists for classes, but is less commonly used there. See the documentation for *function definitions* and *class definitions* for more about decorators.

descriptor Any object which defines the methods __get__(), __set__(), or __delete__(). When a class attribute is a descriptor, its special binding behavior is triggered upon attribute lookup. Normally, using a.b to get, set or delete an attribute looks up the object named b in the class dictionary for a, but if b is a descriptor, the respective descriptor method gets called. Understanding descriptors is a key to a deep understanding of Python because they are the basis for many features including functions, methods, properties, class methods, static methods, and reference to super classes.

For more information about descriptors' methods, see *Implementing Descriptors*.

- dictionary An associative array, where arbitrary keys are mapped to values. The keys can be any object with __hash__() and __eq__() methods. Called a hash in Perl.
- **docstring** A string literal which appears as the first expression in a class, function or module. While ignored when the suite is executed, it is recognized by the compiler and put into the __doc__ attribute of the enclosing class, function or module. Since it is available via introspection, it is the canonical place for documentation of the object.
- duck-typing A programming style which does not look at an object's type to determine if it has the right interface; instead, the method or attribute is simply called or used ("If it looks like a duck and quacks like a duck, it must be a duck.") By emphasizing interfaces rather than specific types, well-designed code improves its flexibility by allowing polymorphic substitution. Duck-typing avoids tests using type() or isinstance(). (Note, however, that duck-typing can be complemented with abstract base classes.) Instead, it typically employs hasattr() tests or EAFP programming.

- **EAFP** Easier to ask for forgiveness than permission. This common Python coding style assumes the existence of valid keys or attributes and catches exceptions if the assumption proves false. This clean and fast style is characterized by the presence of many try and except statements. The technique contrasts with the *LBYL* style common to many other languages such as C.
- **expression** A piece of syntax which can be evaluated to some value. In other words, an expression is an accumulation of expression elements like literals, names, attribute access, operators or function calls which all return a value. In contrast to many other languages, not all language constructs are expressions. There are also *statements* which cannot be used as expressions, such as if. Assignments are also statements, not expressions.
- **extension module** A module written in C or C++, using Python's C API to interact with the core and with user code.
- **file object** An object exposing a file-oriented API (with methods such as read() or write()) to an underlying resource. Depending on the way it was created, a file object can mediate access to a real on-disk file or to another type of storage or communication device (for example standard input/output, in-memory buffers, sockets, pipes, etc.). File objects are also called *file-like objects* or *streams*.

There are actually three categories of file objects: raw binary files, buffered binary files and text files. Their interfaces are defined in the io module. The canonical way to create a file object is by using the open () function.

file-like object A synonym for *file object*.

- **finder** An object that tries to find the *loader* for a module. It must implement either a method named find_loader() or a method named find_module(). See PEP 302 and PEP 420 for details and importlib.abc.Finder for an abstract base class.
- **floor division** Mathematical division that rounds down to nearest integer. The floor division operator is //. For example, the expression 11 // 4 evaluates to 2 in contrast to the 2.75 returned by float true division. Note that (-11) // 4 is -3 because that is -2.75 rounded *downward*. See PEP 238.
- **function** A series of statements which returns some value to a caller. It can also be passed zero or more arguments which may be used in the execution of the body. See also *argument* and *method*.
- **__future**__ A pseudo-module which programmers can use to enable new language features which are not compatible with the current interpreter.

By importing the __future__ module and evaluating its variables, you can see when a new feature was first added to the language and when it becomes the default:

```
>>> import __future__
>>> __future__.division
_Feature((2, 2, 0, 'alpha', 2), (3, 0, 0, 'alpha', 0), 8192)
```

- **garbage collection** The process of freeing memory when it is not used anymore. Python performs garbage collection via reference counting and a cyclic garbage collector that is able to detect and break reference cycles.
- generator A function which returns an iterator. It looks like a normal function except that it contains yield statements for producing a series a values usable in a for-loop or that can be retrieved one at a time with the next() function. Each yield temporarily suspends processing, remembering the location execution state (including local variables and pending try-statements). When the generator resumes, it picks-up where it left-off (in contrast to functions which start fresh on every invocation).
- **generator expression** An expression that returns an iterator. It looks like a normal expression followed by a for expression defining a loop variable, range, and an optional if expression. The combined expression generates values for an enclosing function:

```
>>> sum(i*i for i in range(10))  # sum of squares 0, 1, 4, ... 81 285
```

GIL See *global interpreter lock*.

global interpreter lock The mechanism used by the *CPython* interpreter to assure that only one thread executes Python *bytecode* at a time. This simplifies the CPython implementation by making the object model (including critical built-in types such as dict) implicitly safe against concurrent access. Locking the entire interpreter makes it easier for the interpreter to be multi-threaded, at the expense of much of the parallelism afforded by multi-processor machines.

However, some extension modules, either standard or third-party, are designed so as to release the GIL when doing computationally-intensive tasks such as compression or hashing. Also, the GIL is always released when doing I/O.

Past efforts to create a "free-threaded" interpreter (one which locks shared data at a much finer granularity) have not been successful because performance suffered in the common single-processor case. It is believed that overcoming this performance issue would make the implementation much more complicated and therefore costlier to maintain.

hashable An object is *hashable* if it has a hash value which never changes during its lifetime (it needs a <u>hash</u> () method), and can be compared to other objects (it needs an <u>eq</u> () method). Hashable objects which compare equal must have the same hash value.

Hashability makes an object usable as a dictionary key and a set member, because these data structures use the hash value internally.

All of Python's immutable built-in objects are hashable, while no mutable containers (such as lists or dictionaries) are. Objects which are instances of user-defined classes are hashable by default; they all compare unequal, and their hash value is their id().

- **IDLE** An Integrated Development Environment for Python. IDLE is a basic editor and interpreter environment which ships with the standard distribution of Python.
- **immutable** An object with a fixed value. Immutable objects include numbers, strings and tuples. Such an object cannot be altered. A new object has to be created if a different value has to be stored. They play an important role in places where a constant hash value is needed, for example as a key in a dictionary.
- import path A list of locations (or path entries) that are searched by the path based finder for modules to import.
 During import, this list of locations usually comes from sys.path, but for subpackages it may also come from the parent package's __path__ attribute.

importing The process by which Python code in one module is made available to Python code in another module.

importer An object that both finds and loads a module; both a *finder* and *loader* object.

- interactive Python has an interactive interpreter which means you can enter statements and expressions at the interpreter prompt, immediately execute them and see their results. Just launch python with no arguments (possibly by selecting it from your computer's main menu). It is a very powerful way to test out new ideas or inspect modules and packages (remember help(x)).
- interpreted Python is an interpreted language, as opposed to a compiled one, though the distinction can be blurry because of the presence of the bytecode compiler. This means that source files can be run directly without explicitly creating an executable which is then run. Interpreted languages typically have a shorter development/debug cycle than compiled ones, though their programs generally also run more slowly. See also interactive.
- iterable An object capable of returning its members one at a time. Examples of iterables include all sequence types (such as list, str, and tuple) and some non-sequence types like dict and file and objects of any classes you define with an __iter__() or __getitem__() method. Iterables can be used in a for loop and in many other places where a sequence is needed (zip(), map(), ...). When an iterable object is passed as an argument to the built-in function iter(), it returns an iterator for the object. This iterator is good for one pass over the set of values. When using iterables, it is usually not necessary to call iter() or deal with iterator objects yourself. The for statement does that automatically for you, creating a temporary unnamed variable to hold the iterator for the duration of the loop. See also iterator, sequence, and generator.

iterator An object representing a stream of data. Repeated calls to the iterator's __next___() method (or passing it to the built-in function next()) return successive items in the stream. When no more data are available a StopIteration exception is raised instead. At this point, the iterator object is exhausted and any further calls to its __next__() method just raise StopIteration again. Iterators are required to have an __iter__() method that returns the iterator object itself so every iterator is also iterable and may be used in most places where other iterables are accepted. One notable exception is code which attempts multiple iteration passes. A container object (such as a list) produces a fresh new iterator each time you pass it to the iter() function or use it in a for loop. Attempting this with an iterator will just return the same exhausted iterator object used in the previous iteration pass, making it appear like an empty container.

More information can be found in typeiter.

key function A key function or collation function is a callable that returns a value used for sorting or ordering. For example, locale.strxfrm() is used to produce a sort key that is aware of locale specific sort conventions.

A number of tools in Python accept key functions to control how elements are ordered or grouped. They include min(), max(), sorted(), list.sort(), heapq.nsmallest(), heapq.nlargest(), and itertools.groupby().

There are several ways to create a key function. For example, the str.lower() method can serve as a key function for case insensitive sorts. Alternatively, an ad-hoc key function can be built from a lambda expression such as lambda r: (r[0], r[2]). Also, the operator module provides three key function constructors: attrgetter(), itemgetter(), and methodcaller(). See the *Sorting HOW TO* for examples of how to create and use key functions.

- **keyword argument** Arguments which are preceded with a variable_name= in the call. The variable name designates the local name in the function to which the value is assigned. ** is used to accept or pass a dictionary of keyword arguments. See *argument*.
- **lambda** An anonymous inline function consisting of a single *expression* which is evaluated when the function is called. The syntax to create a lambda function is lambda [arguments]: expression
- **LBYL** Look before you leap. This coding style explicitly tests for pre-conditions before making calls or lookups. This style contrasts with the *EAFP* approach and is characterized by the presence of many if statements.

In a multi-threaded environment, the LBYL approach can risk introducing a race condition between "the looking" and "the leaping". For example, the code, if key in mapping: return mapping [key] can fail if another thread removes *key* from *mapping* after the test, but before the lookup. This issue can be solved with locks or by using the EAFP approach.

- **list** A built-in Python *sequence*. Despite its name it is more akin to an array in other languages than to a linked list since access to elements are O(1).
- **list comprehension** A compact way to process all or part of the elements in a sequence and return a list with the results. result = $['\{:\#04x\}']$. format(x) for x in range(256) if x % 2 == 0] generates a list of strings containing even hex numbers (0x...) in the range from 0 to 255. The if clause is optional. If omitted, all elements in range(256) are processed.
- **loader** An object that loads a module. It must define a method named load_module(). A loader is typically returned by a *finder*. See PEP 302 for details and importlib.abc.Loader for an *abstract base class*.
- mapping A container object that supports arbitrary key lookups and implements the methods specified in the Mapping or MutableMapping abstract base classes. Examples include dict, collections.defaultdict, collections.OrderedDict and collections.Counter.
- **meta path finder** A finder returned by a search of sys.meta_path. Meta path finders are related to, but different from *path entry finders*.
- **metaclass** The class of a class. Class definitions create a class name, a class dictionary, and a list of base classes. The metaclass is responsible for taking those three arguments and creating the class. Most object oriented programming languages provide a default implementation. What makes Python special is that it is possible to

create custom metaclasses. Most users never need this tool, but when the need arises, metaclasses can provide powerful, elegant solutions. They have been used for logging attribute access, adding thread-safety, tracking object creation, implementing singletons, and many other tasks.

More information can be found in *Customizing class creation*.

- **method** A function which is defined inside a class body. If called as an attribute of an instance of that class, the method will get the instance object as its first *argument* (which is usually called self). See *function* and *nested scope*.
- **method resolution order** Method Resolution Order is the order in which base classes are searched for a member during lookup. See The Python 2.3 Method Resolution Order.
- **module** An object that serves as an organizational unit of Python code. Modules have a namespace containing arbitrary Python objects. Modules are loaded into Python by the process of *importing*.

MRO See method resolution order.

- mutable Mutable objects can change their value but keep their id(). See also immutable.
- **named tuple** Any tuple-like class whose indexable elements are also accessible using named attributes (for example, time.localtime() returns a tuple-like object where the *year* is accessible either with an index such as t[0] or with a named attribute like t.tm_year).
 - A named tuple can be a built-in type such as time.struct_time, or it can be created with a regular class definition. A full featured named tuple can also be created with the factory function collections.namedtuple(). The latter approach automatically provides extra features such as a self-documenting representation like Employee(name='jones', title='programmer').
- namespace The place where a variable is stored. Namespaces are implemented as dictionaries. There are the local, global and built-in namespaces as well as nested namespaces in objects (in methods). Namespaces support modularity by preventing naming conflicts. For instance, the functions builtins.open() and os.open() are distinguished by their namespaces. Namespaces also aid readability and maintainability by making it clear which module implements a function. For instance, writing random.seed() or itertools.islice() makes it clear that those functions are implemented by the random and itertools modules, respectively.
- **namespace package** A **PEP 420** *package* which serves only as a container for subpackages. Namespace packages may have no physical representation, and specifically are not like a *regular package* because they have no __init__.py file.
- **nested scope** The ability to refer to a variable in an enclosing definition. For instance, a function defined inside another function can refer to variables in the outer function. Note that nested scopes by default work only for reference and not for assignment. Local variables both read and write in the innermost scope. Likewise, global variables read and write to the global namespace. The nonlocal allows writing to outer scopes.
- **new-style class** Old name for the flavor of classes now used for all class objects. In earlier Python versions, only new-style classes could use Python's newer, versatile features like __slots__, descriptors, properties, __getattribute__(), class methods, and static methods.
- **object** Any data with state (attributes or value) and defined behavior (methods). Also the ultimate base class of any *new-style class*.
- **package** A Python module which can contain submodules or recursively, subpackages. Technically, a package is a Python module with an __path__ attribute.
- path entry A single location on the *import path* which the *path based finder* consults to find modules for importing.
- path entry finder A finder returned by a callable on sys.path_hooks (i.e. a path entry hook) which knows how to locate modules given a path entry.
- **path entry hook** A callable on the sys.path_hook list which returns a *path entry finder* if it knows how to find modules on a specific *path entry*.

path based finder One of the default meta path finders which searches an import path for modules.

portion A set of files in a single directory (possibly stored in a zip file) that contribute to a namespace package, as defined in **PEP 420**.

positional argument The arguments assigned to local names inside a function or method, determined by the order in which they were given in the call. * is used to either accept multiple positional arguments (when in the definition), or pass several arguments as a list to a function. See *argument*.

provisional package A provisional package is one which has been deliberately excluded from the standard library's backwards compatibility guarantees. While major changes to such packages are not expected, as long as they are marked provisional, backwards incompatible changes (up to and including removal of the package) may occur if deemed necessary by core developers. Such changes will not be made gratuitously – they will occur only if serious flaws are uncovered that were missed prior to the inclusion of the package.

This process allows the standard library to continue to evolve over time, without locking in problematic design errors for extended periods of time. See PEP 411 for more details.

Python 3000 Nickname for the Python 3.x release line (coined long ago when the release of version 3 was something in the distant future.) This is also abbreviated "Py3k".

Pythonic An idea or piece of code which closely follows the most common idioms of the Python language, rather than implementing code using concepts common to other languages. For example, a common idiom in Python is to loop over all elements of an iterable using a for statement. Many other languages don't have this type of construct, so people unfamiliar with Python sometimes use a numerical counter instead:

```
for i in range(len(food)):
    print(food[i])
```

As opposed to the cleaner, Pythonic method:

```
for piece in food:
    print(piece)
```

qualified name A dotted name showing the "path" from a module's global scope to a class, function or method defined in that module, as defined in

PEP 3155. For top-level functions and classes, the qualified name is the same as the object's name:

```
>>> class C:
... class D:
... def meth(self):
... pass
...
>>> C.__qualname___
'C'
>>> C.D.__qualname__
'C.D'
>>> C.D.meth.__qualname__
'C.D.meth'
```

When used to refer to modules, the *fully qualified name* means the entire dotted path to the module, including any parent packages, e.g. email.mime.text:

```
>>> import email.mime.text
>>> email.mime.text.__name__
'email.mime.text'
```

reference count The number of references to an object. When the reference count of an object drops to zero, it is deallocated. Reference counting is generally not visible to Python code, but it is a key element of the *CPython*

- implementation. The sys module defines a getrefcount () function that programmers can call to return the reference count for a particular object.
- **regular package** A traditional *package*, such as a directory containing an __init__.py file.
- **__slots**__ A declaration inside a class that saves memory by pre-declaring space for instance attributes and eliminating instance dictionaries. Though popular, the technique is somewhat tricky to get right and is best reserved for rare cases where there are large numbers of instances in a memory-critical application.
- sequence An iterable which supports efficient element access using integer indices via the __getitem__() special
 method and defines a __len__() method that returns the length of the sequence. Some built-in sequence types
 are list, str, tuple, and bytes. Note that dict also supports __getitem__() and __len__(), but
 is considered a mapping rather than a sequence because the lookups use arbitrary immutable keys rather than
 integers.
- slice An object usually containing a portion of a *sequence*. A slice is created using the subscript notation, [] with colons between numbers when several are given, such as in variable_name[1:3:5]. The bracket (subscript) notation uses slice objects internally.
- **special method** A method that is called implicitly by Python to execute a certain operation on a type, such as addition. Such methods have names starting and ending with double underscores. Special methods are documented in *Special method names*.
- **statement** A statement is part of a suite (a "block" of code). A statement is either an *expression* or a one of several constructs with a keyword, such as if, while or for.
- struct sequence A tuple with named elements. Struct sequences expose an interface similar to named tuple in that elements can either be accessed either by index or as an attribute. However, they do not have any of the named tuple methods like _make() or _asdict(). Examples of struct sequences include sys.float_info and the return value of os.stat().
- **triple-quoted string** A string which is bound by three instances of either a quotation mark (") or an apostrophe ('). While they don't provide any functionality not available with single-quoted strings, they are useful for a number of reasons. They allow you to include unescaped single and double quotes within a string and they can span multiple lines without the use of the continuation character, making them especially useful when writing docstrings.
- type The type of a Python object determines what kind of object it is; every object has a type. An object's type is accessible as its __class__ attribute or can be retrieved with type (obj).
- universal newlines A manner of interpreting text streams in which all of the following are recognized as ending a line: the Unix end-of-line convention '\n', the Windows convention '\r\n', and the old Macintosh convention '\r'. See PEP 278 and PEP 3116, as well as str.splitlines() for an additional use.
- **view** The objects returned from dict.keys(), dict.values(), and dict.items() are called dictionary views. They are lazy sequences that will see changes in the underlying dictionary. To force the dictionary view to become a full list use list (dictview). See *dict-views*.
- **virtual machine** A computer defined entirely in software. Python's virtual machine executes the *bytecode* emitted by the bytecode compiler.
- **Zen of Python** Listing of Python design principles and philosophies that are helpful in understanding and using the language. The listing can be found by typing "import this" at the interactive prompt.

В

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These documents are generated from reStructuredText sources by Sphinx, a document processor specifically written for the Python documentation.

Development of the documentation and its toolchain takes place on the docs@python.org mailing list. We're always looking for volunteers wanting to help with the docs, so feel free to send a mail there!

Many thanks go to:

- Fred L. Drake, Jr., the creator of the original Python documentation toolset and writer of much of the content;
- the Docutils project for creating reStructuredText and the Docutils suite;
- Fredrik Lundh for his Alternative Python Reference project from which Sphinx got many good ideas.

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It is only with the input and contributions of the Python community that Python has such wonderful documentation – Thank You!

HISTORY AND LICENSE

C.1 History of the software

Python was created in the early 1990s by Guido van Rossum at Stichting Mathematisch Centrum (CWI, see http://www.cwi.nl/) in the Netherlands as a successor of a language called ABC. Guido remains Python's principal author, although it includes many contributions from others.

In 1995, Guido continued his work on Python at the Corporation for National Research Initiatives (CNRI, see http://www.cnri.reston.va.us/) in Reston, Virginia where he released several versions of the software.

In May 2000, Guido and the Python core development team moved to BeOpen.com to form the BeOpen PythonLabs team. In October of the same year, the PythonLabs team moved to Digital Creations (now Zope Corporation; see http://www.zope.com/). In 2001, the Python Software Foundation (PSF, see http://www.python.org/psf/) was formed, a non-profit organization created specifically to own Python-related Intellectual Property. Zope Corporation is a sponsoring member of the PSF.

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Release	Derived from	Year	Owner	GPL compatible?		
0.9.0 thru 1.2	n/a	1991-1995	CWI	yes		
1.3 thru 1.5.2	1.2	1995-1999	CNRI	yes		
1.6	1.5.2	2000	CNRI	no		
2.0	1.6	2000	BeOpen.com	no		
1.6.1	1.6	2001	CNRI	no		
2.1	2.0+1.6.1	2001	PSF	no		
2.0.1	2.0+1.6.1	2001	PSF	yes		
2.1.1	2.1+2.0.1	2001	PSF	yes		
2.2	2.1.1	2001	PSF	yes		
2.1.2	2.1.1	2002	PSF	yes		
2.1.3	2.1.2	2002	PSF	yes		
2.2.1	2.2	2002	PSF	yes		
2.2.2	2.2.1	2002	PSF	yes		
2.2.3	2.2.2	2002-2003	PSF	yes		
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2.3.1	2.3	2002-2003	PSF	yes		
2.3.2	2.3.1	2003	PSF	yes		
2.3.3	2.3.2	2003	PSF	yes		
2.3.4	2.3.3	2004	PSF	yes		
2.3.5	2.3.4	2005	PSF	yes		
2.4	2.3	2004	PSF	yes		
Continued on next page						

Table C.1 – continued from previous page						
2.4.1	2.4	2005	PSF	yes		
2.4.2	2.4.1	2005	PSF	yes		
2.4.3	2.4.2	2006	PSF	yes		
2.4.4	2.4.3	2006	PSF	yes		
2.5	2.4	2006	PSF	yes		
2.5.1	2.5	2007	PSF	yes		
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2.6.1	2.6	2008	PSF	yes		
2.6.2	2.6.1	2009	PSF	yes		
2.6.3	2.6.2	2009	PSF	yes		
2.6.4	2.6.3	2009	PSF	yes		
3.0	2.6	2008	PSF	yes		
3.0.1	3.0	2009	PSF	yes		
3.1	3.0.1	2009	PSF	yes		
3.1.1	3.1	2009	PSF	yes		
3.1.2	3.1.1	2010	PSF	yes		
3.1.3	3.1.2	2010	PSF	yes		
3.1.4	3.1.3	2011	PSF	yes		
3.2	3.1	2011	PSF	yes		
3.2.1	3.2	2011	PSF	yes		
3.2.2	3.2.1	2011	PSF	yes		
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C.3.1 Mersenne Twister

The _random module includes code based on a download from http://www.math.sci.hiroshima-u.ac.jp/~m-mat/MT/MT2002/emt19937ar.html. The following are the verbatim comments from the original code:

A C-program for MT19937, with initialization improved 2002/1/26. Coded by Takuji Nishimura and Makoto Matsumoto.

Before using, initialize the state by using init_genrand(seed) or init_by_array(init_key, key_length).

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```
http://www.math.sci.hiroshima-u.ac.jp/~m-mat/MT/emt.html email: m-mat @ math.sci.hiroshima-u.ac.jp (remove space)
```

C.3.2 Sockets

The socket module uses the functions, getaddrinfo(), and getnameinfo(), which are coded in separate source files from the WIDE Project, http://www.wide.ad.jp/.

```
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```

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Modified by Jack Jansen, CWI, July 1995:

- Use binascii module to do the actual line-by-line conversion between ascii and binary. This results in a 1000-fold speedup. The C version is still 5 times faster, though.
- Arguments more compliant with Python standard

C.3.8 XML Remote Procedure Calls

The xmlrpc.client module contains the following notice:

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C.3.11 SHA-3

The module _sha3 and hashlib are using the reference implementation of Keccak. The files at Modules/_sha3/keccak/ contain the following note:

The Keccak sponge function, designed by Guido Bertoni, Joan Daemen, Michaël Peeters and Gilles Van Assche. For more information, feedback or questions, please refer to our website: http://keccak.noekeon.org/

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C.3.12 strtod and dtoa

The file Python/dtoa.c, which supplies C functions dtoa and strtod for conversion of C doubles to and from strings, is derived from the file of the same name by David M. Gay, currently available from http://www.netlib.org/fp/. The original file, as retrieved on March 16, 2009, contains the following copyright and licensing notice:

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APPENDIX

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