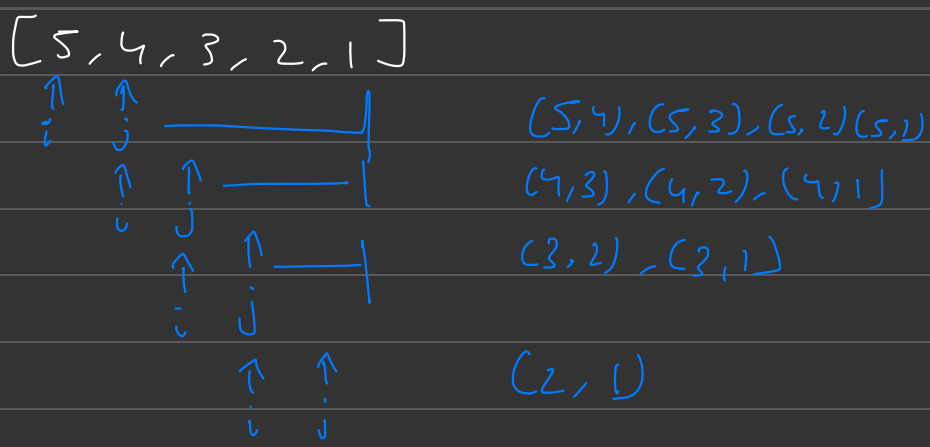
Brute force

An element a[i] can form pair with further elements if (a[i] > a[j])



inversion_ct = 10 //

T.C → O(n²) , SC → O(1)

Optimised approach (using Merge sort)

- In merge sort, while merging two halves, one thing is sure that both the halves are sorted

eg: [2, 3, 5] + [1, 4]

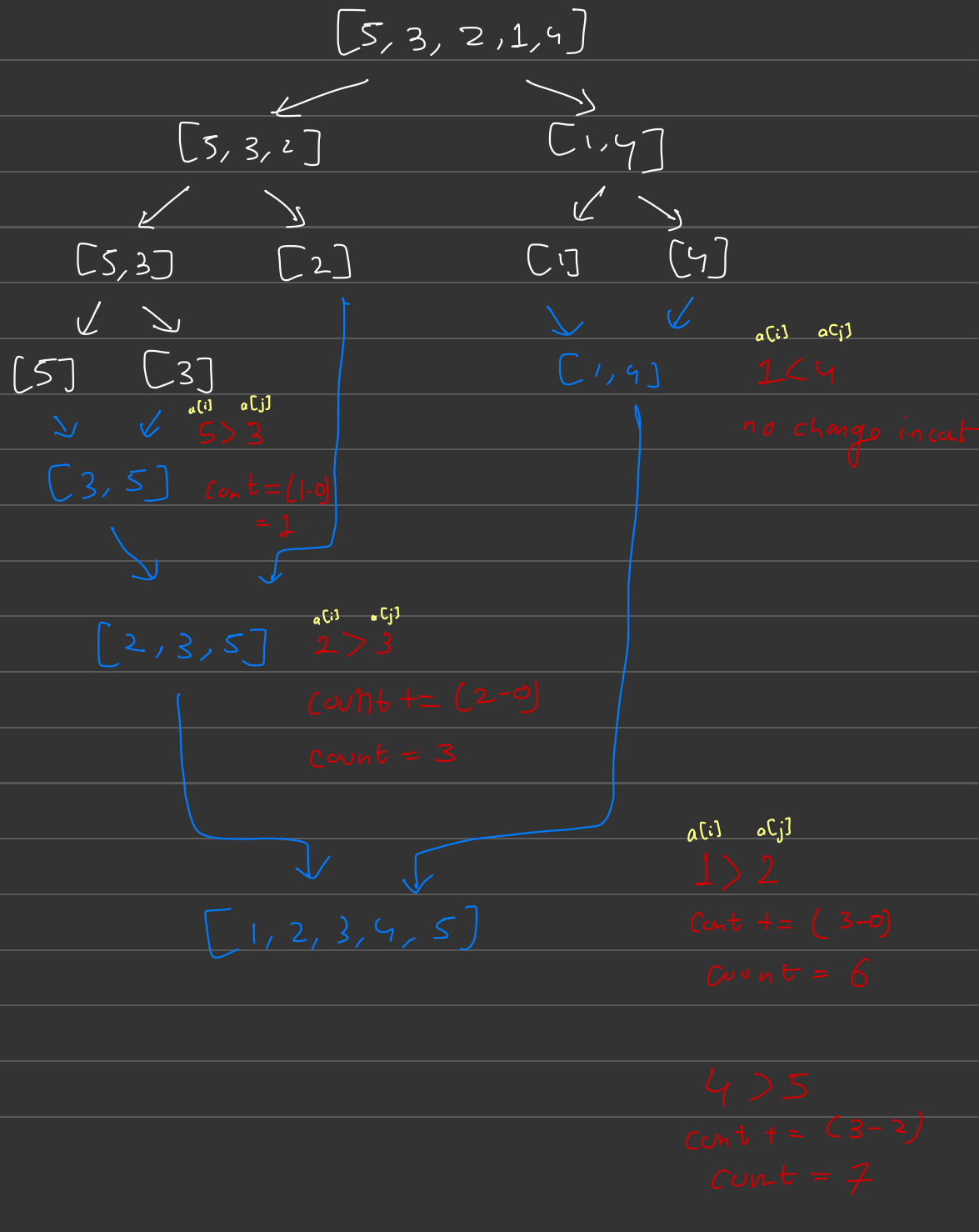
(i at 2, (mid − i) spanning 3, 5; j at 1; mid above)

(if a[i] > a[j], a[j] forms inversion pairs with all the elements on the right of i to end of first arr i.e. mid.)

∴ count_inv += (mid − i)

In the above example, when being copied to temp array inv pairs → (1, 2), (1, 3), (1, 5)
(4, 5)

T.C → O(n log n)   S.C → O(n)


[5, 3, 2, 1, 4]
[5, 3, 2]
[1, 4]
[5, 3]
[2]
[1]
[4]
[5]
[3]
[1, 4]
a[i] a[j]
1 < 4
no change in cnt
a[i] a[j]
5 > 3
[3, 5]
cnt = (1-0)
= 1
[2, 3, 5]
a[i] a[j]
2 > 3
count += (2-0)
count = 3
a[i] a[j]
1 > 2
Cnt += (3-0)
Count = 6
[1, 2, 3, 4, 5]
4 > 5
cnt += (3-2)
Cnt = 7

```
int mergeSort(int arr[], int temp[], int l, int r)
    int mid, inv_count = 0;

    while (l < r) {
     mid = (l+r)/2;
     inv_count += mergeSort(arr, temp, l, mid);
     inv_count += mergeSort(arr, temp, mid+1, r);
     inv_count += merge(arr, temp, l, mid+1, r);
    }
   return inv_count;
}
```

```
int merge (int arr[], int temp[],
           int left, int mid, int right)
{
  int i, j, k;
  int inv_count = 0;
  i = left ; j = mid ; k = left;
  // arr 1 -> l to mid-1     // arr 2 -> mid to r
  while ((i <= mid-1) && (j <= right)) {
    if (arr[i] <= arr[j]) {  // No swapping
      temp[k++] = arr[i++];  // copy to temp & i++
    } else {  // Swapping & inv_count calc
      temp[k++] = arr[j++];  // copy to temp & j++
      inv_count += (mid - i);  // inv_count calc
    }
  }

// If elems in arr[l] to arr[mid-1] are left, copy them
while (i <= mid-1)   temp[k++] = arr[i++]
// If elems in arr[mid] to arr[r] are left, copy them
while (j <= right)   temp[k++] = arr[j++]
for (i = left; i <= right; i++)   arr[i] = temp[i]

return inv_count;
}
```

copy temp arr elems to arr.