Title: Binary search Tree (BST) operations

problem Statement: Beginning with an empty binary search
tree, construct binary search tree by inserting the values
In the order given lafter constructing a binary face.

I Insert new node. ill find number of nodes in longest
31 minimum data value found in the tree. Path from the root
41 change a tree so that the roles of the right and left
printers are swapped at every node
51 search a value.

Objectives: To construct and understand the concept the Binary search Tree (BST) and its operations

Software Requisement + 9++ 1900 compiler, 64-bit Fedora
eclipse IDE.

Theory: A binary segret tree is a binary tree where the value of each node is greater or equal to the values in its left subtree and less than or equal to the values in the right subtrees. This property allows ratues in the searching and insertion of values in the free.

- when inserting values into a binary search tree,

the values are compared to the value of the root node

The value is less than the root node, it is inserted.

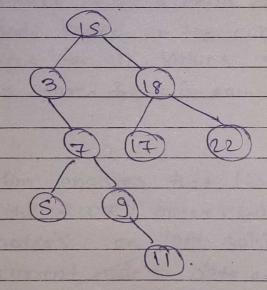
In the left subtree otherwise is inserted in the right

subtree.

This process is repeated recursively until a lead node is seached at which point the new node is inserted as a child of the leaf node.

The order in which values are inserted into the tree tree an affect the height and the balance of the tree. In the worst case scenario, when the values are inserted in ascending or deconding orders, the resulting tree can becomes degenerate tree, where all the nodes have only one child and the tree, where all that nodes have only one child and the tree become essentially a linked list. To avoid this, it is important to insert values in a balanced order, such as randomly or by using a balanced binary search tree algorith.

Example: Keys + 15, 3, 7, 9, 11, 18, 22, 17, 5



Binary search Tree (BST)

Algorithm :

- 1. Start
- 2. Create an empty binary search tree.
- 3. Read the values to be insented in the binary search tree in the order given
- 9. For each value to be inserted do this following
 - a. If the binary search tree is empty, create a new nodes with the value as its data and make it the boot node of the tree.
 - b. If the binary search tree is not empty compare. the value, to be inserted with the data of the corrent node.
 - C. create a new node with the value as its data, and insert it as a child node of the last visited leaf node based on the companisons in step b.
- 5. Ropeat step 3 for all values to be inserted
- 6. The binoxy search frees is now constructed
- 7. stop.

This algorithm ensures that the binary search tree is always in the correct order, with values less than the current nodes's on 'left subtree and values greater than the current mode's data on the right subtree. The time comprexity of this algorithm is o (logn) in. the querage case, where n is the number of values to be inserted and coorst - case time complexity is o(n2)

Conclusion: Binary search frees are useful data structures for algorithm sparehing duta and their efficient construction and traversa