

Date: 05/05/21

# Black Board

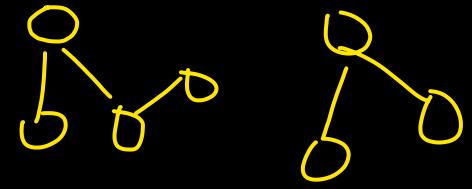
Design and Analysis of Algorithms

Topics:

- Graph Algorithms III
  - Connected Components in Undirected Graphs
  - Strongly Connected Components (SCCs) in directed graphs.

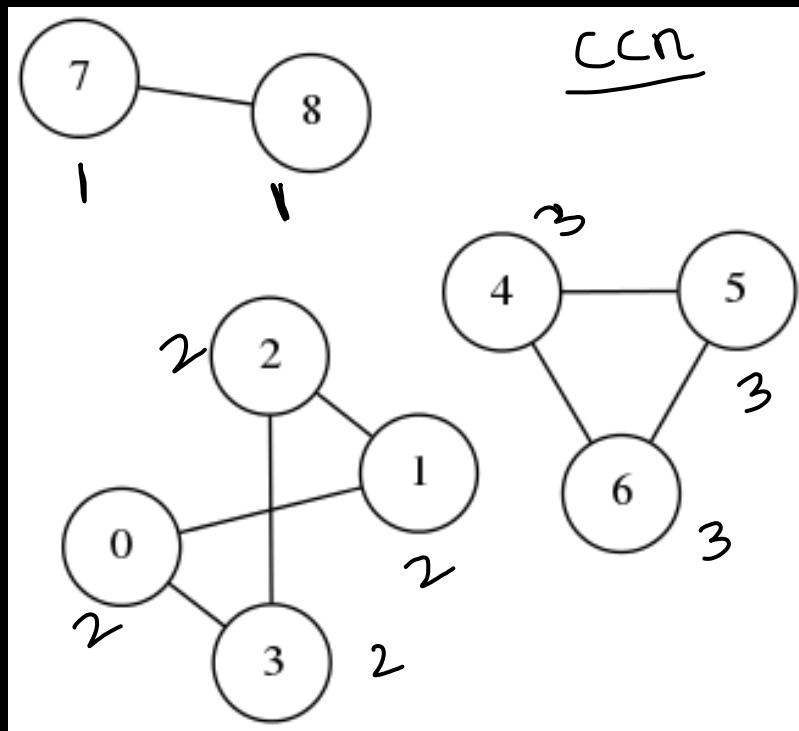
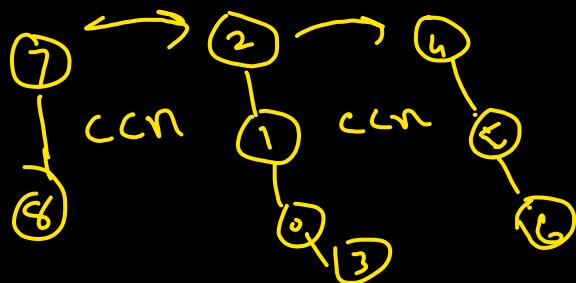
# Connected Components and their applications

- A very important way to analyze graph structure
- Example applications:
  - Internet Documents Clustering
  - Social Media Community Detection
  - Protein-Protein Interaction Analysis
  - Image Segmentation
  - Natural Language Processing: word and character separation, etc.



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# Connected Components in Undirected Graphs



Def:  
Two nodes may  
are in the same  
connected compo  
when there is  
a path any  
in the graph.

# Undirected graphs

**DFS( $G=(V, E)$ )**

For each  $x \in V$ :

$\text{visited}[x] = \text{false}$

$ccn = 1$

For each  $x \in V$ :

    if  $\text{!visited}[x]$

        explore( $G, y$ )

$ccn = ccn + 1$

**Explore( $G=(V, E), x$ )**

$\text{visited}[x] = \text{true}$

$comp[x] = ccn$

    For each  $(x, y) \in E$ :

        if  $\text{!visited}[y]$

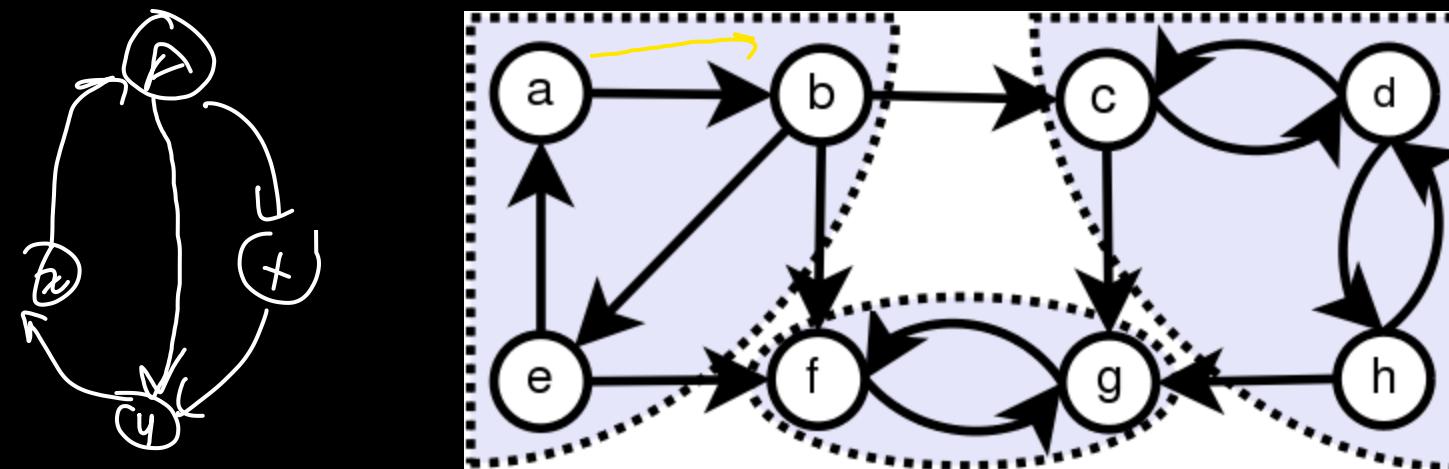
            explore( $G, y$ )

# Connected Components in Undirected Graphs

Can be found in  $O(|V| + |E|)$   
with small additions | the  
regular DFS code.

# Connected Components in Directed Graphs

- Strongly Connected Components (SCCs)



SCC:

In any pair of vertices  $x, y$  in a SCC of  $G$ ,  
there is a ~~path~~ from  $x$  to  $y$  & a path  
from  $y$  to  $x$ .

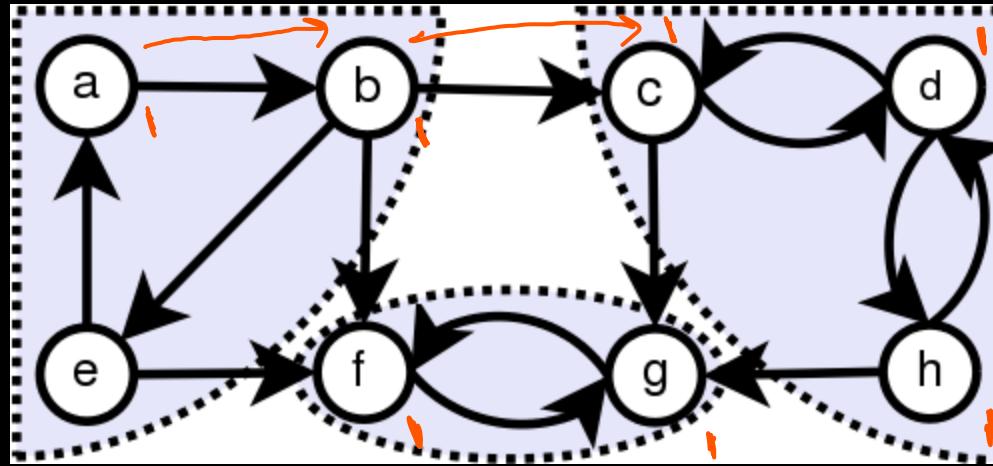
# The Challenge of finding SCCs in Linear Time

CCN=1

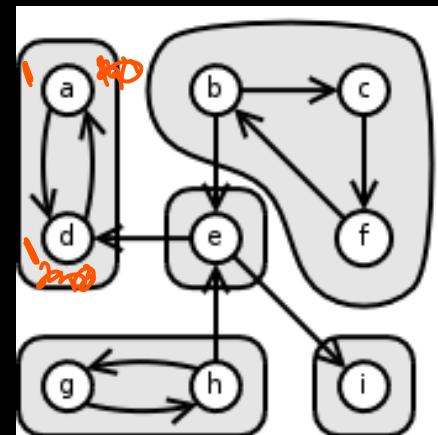
(return to DFS)

flow to  
stop the  
chain

exploses after  
labeling every component  
separately.

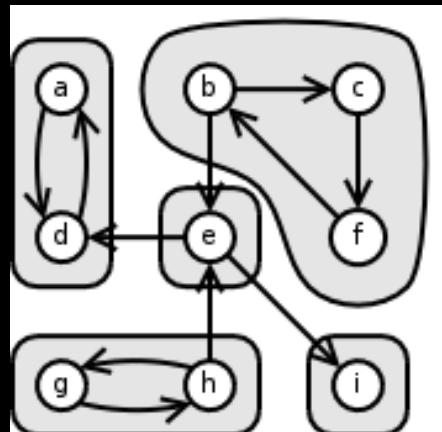
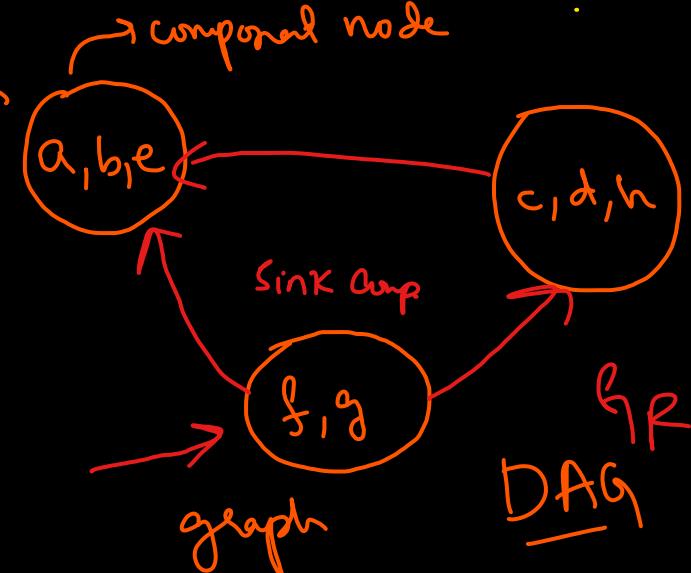
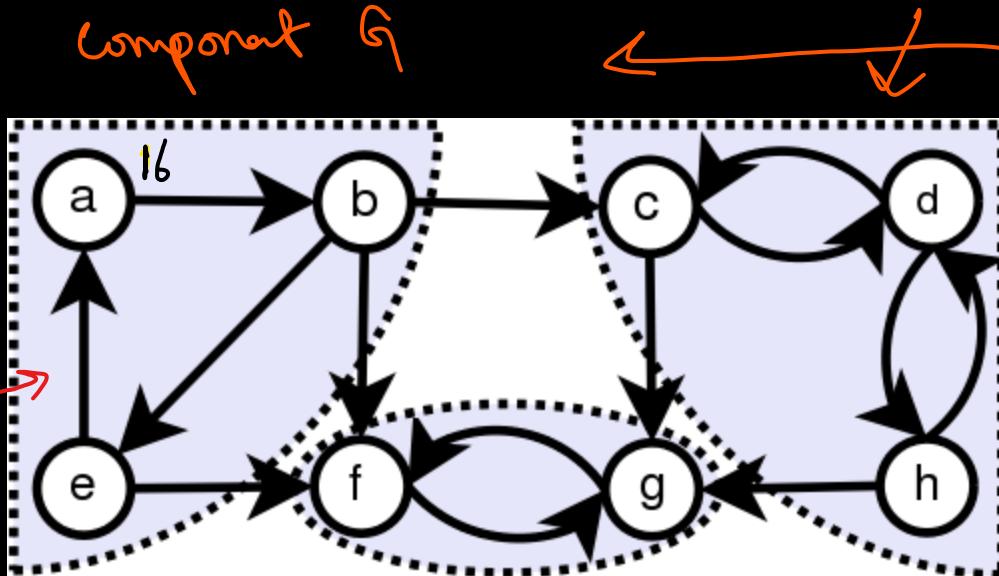


Example 1



Example 2

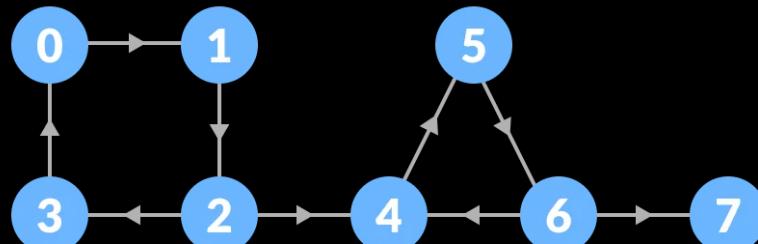
# Relationship of SCCs with DAGs



$G \uparrow = (V, E)$

$G_R = (V, E_R)$

What is the Component DAG of the following graph?



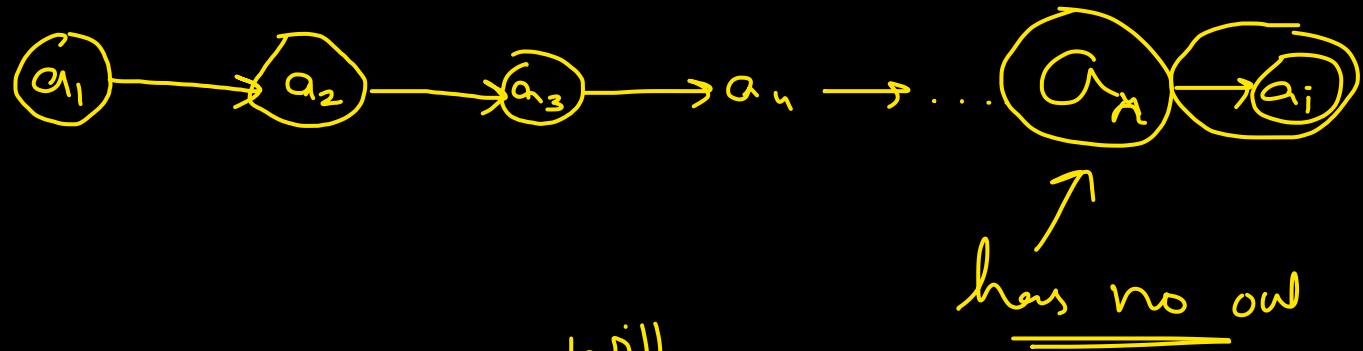
# Sinks, Sources, DAGs

- Where should be start each **explore** call?
  - From a sink SCC.

**FACT:** Every DAG has at least one source and at least one sink

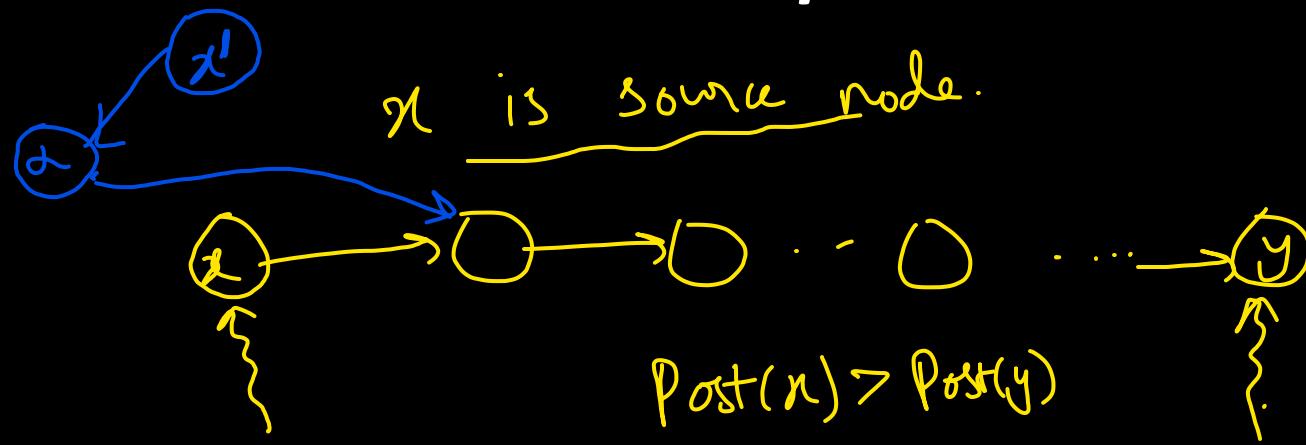
- **Proof:**

→ Think of a DAG  $G = (V, E)$   $V = \{a_1, a_2, \dots, a_n\}$



Our Component DAG <sup>will</sup> ~~must~~ always  
contain a Sink Component & a Source Component.

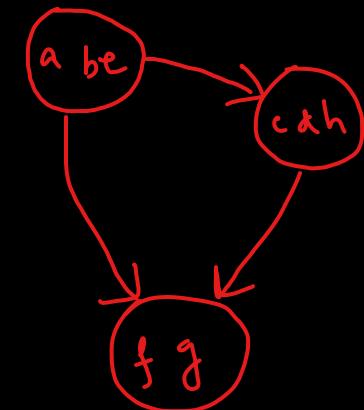
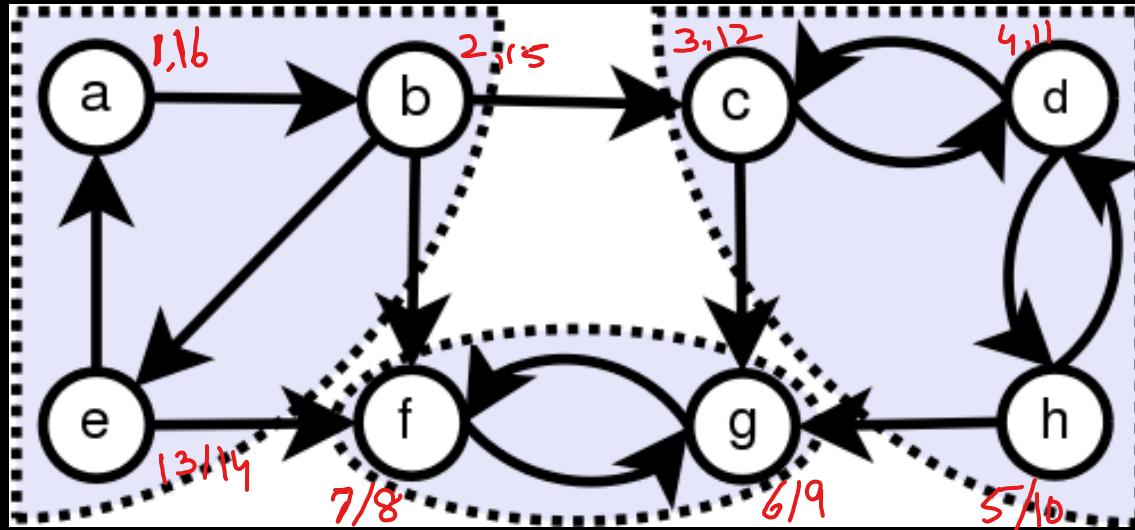
- So the component DAG has at least one sink component node and at least one source component node.
- ***How to detect a sink component node?***



# The highest post # must belong to a source node.

# Using post numbers to find sink component nodes

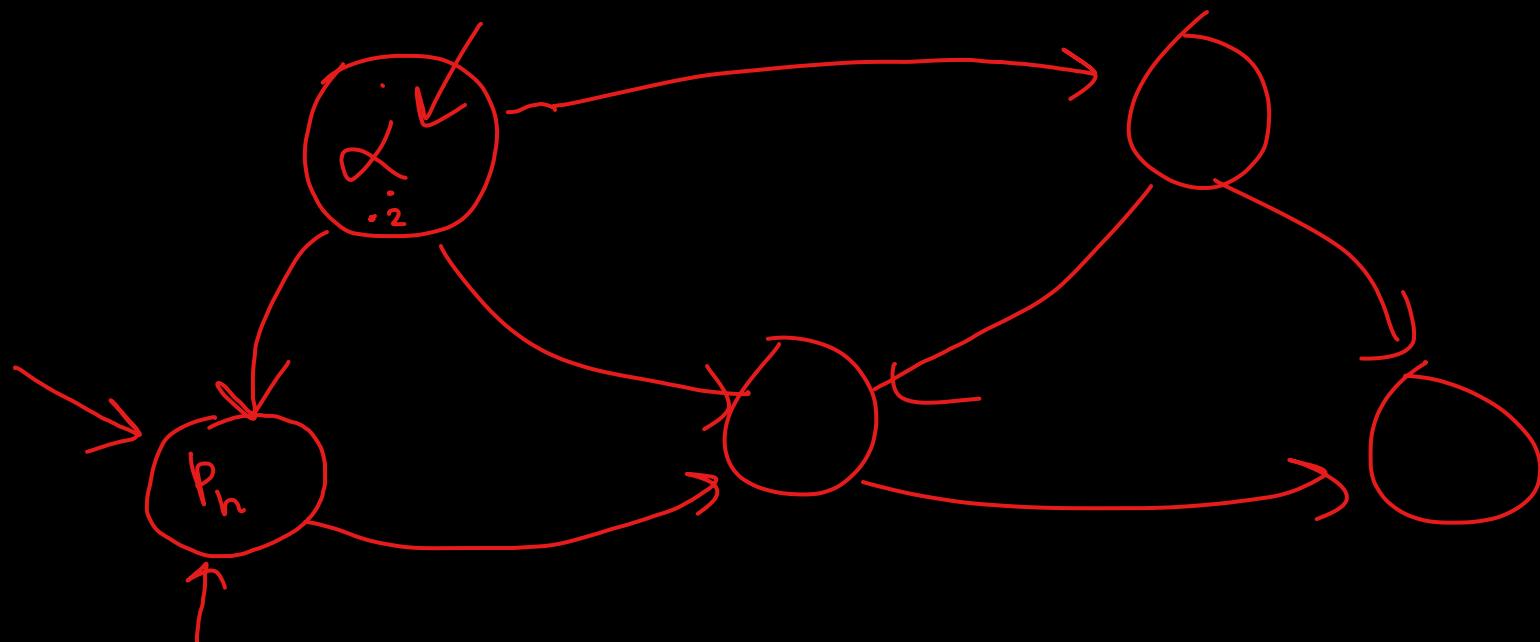
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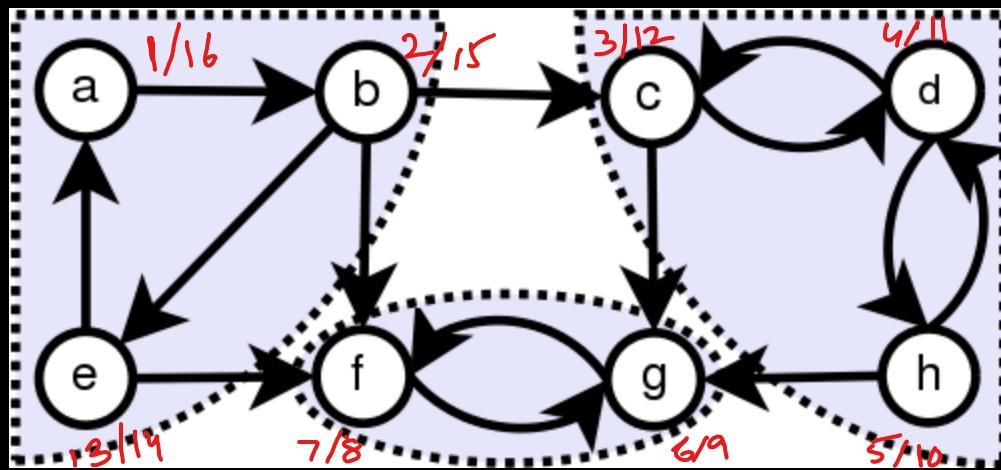
Claim: The node with the highest post # must be in a source component?

# Post Numbers and Source Components

- FACT: the highest post number must be in a source SCC.
- Proof:

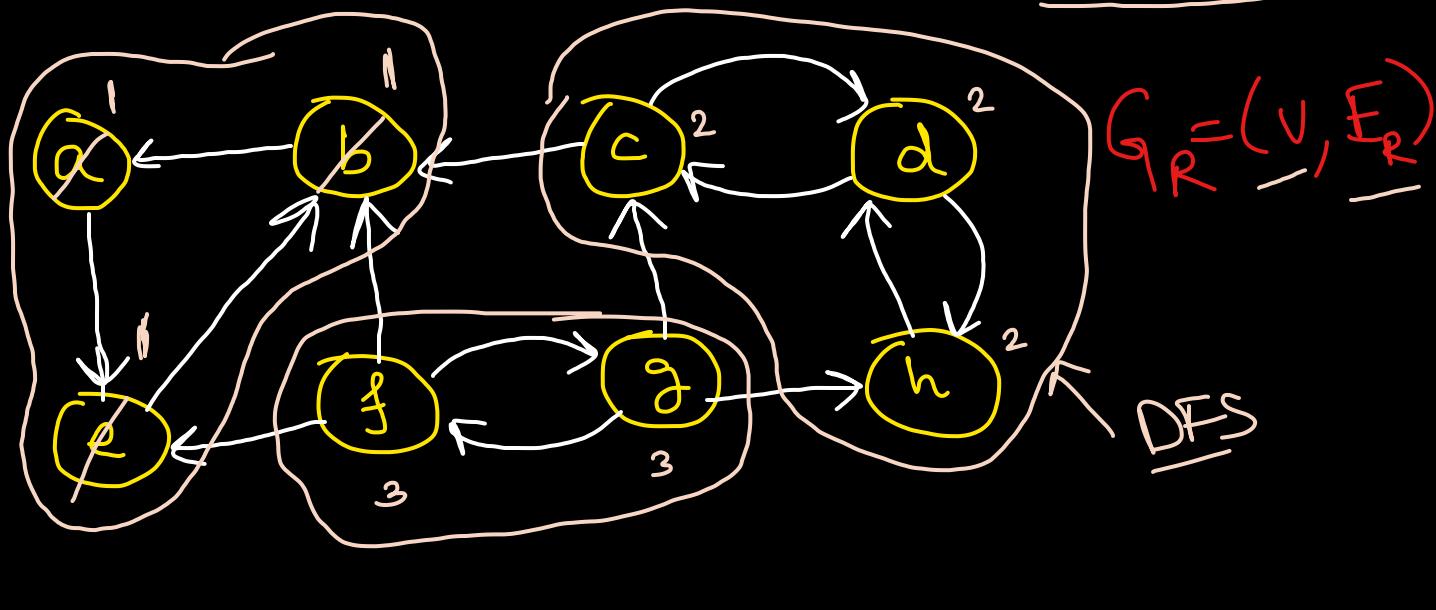


# The Graph Reversal Trick



$$G = (V, E)$$

$$\text{ccnum}[x] = cc_n$$

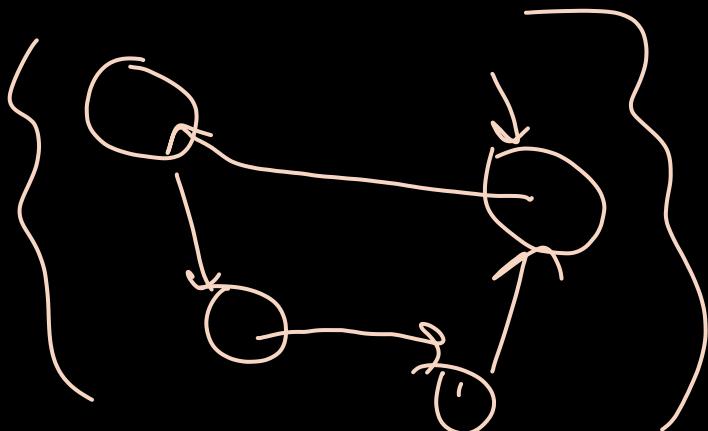
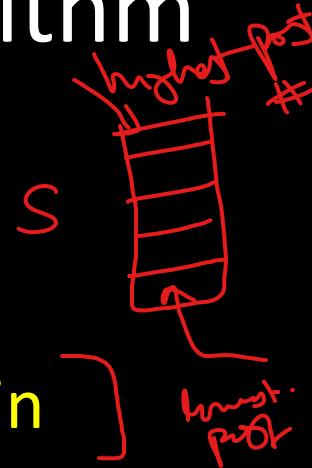


$f_{1,0}$

# The Linear time SCC finding Algorithm (Kosaraju's Algorithm)

## First Two Steps:

- Perform regular DFS on  $G$ , and get vertices in decreasing order of post numbers in a stack  $S$
- Create reverse Graph  $G_R = (V, E_R)$  in  $O(\underline{|V|+|E|})$
- Call ModifiedDFS( $G_R = (V, E_R)$ ,  $S$ )



$$T(G = (V, E)) = O(|V| + |E|)$$

# Modified DFS and Explore Methods

**ModifiedDFS( $G=(V, E)$ ,  $S$ )**

For each  $x \in V$ :

    visited[x] = false

$ccn = 1$

WHILE  $S.\text{empty}()$ :

$x \leftarrow S.\text{top}()$

    if !visited[x]

        Explore( $G$ ,  $x$ )

$ccn = ccn + 1$

$S.\text{pop}()$

**Explore( $G=(V, E)$ ,  $x$ )**

    visited[x] = true

$\text{comp}[x] = ccn$

    For each  $(x, y) \in E$ :

        if !visited[y]

            explore( $G$ ,  $y$ )

