NESTED QUERIES

A complete SELECT query, called a <u>nested query</u>, can be specified within the WHERE-clause of another query, called the <u>outer query</u>

Retrieve the name of all employees who earn more than the Average Salary

SELECT FNAME, LNAME

FROM EMPLOYEE

WHERE Salary > (SELECT AVG(SALARY)

FROM EMPLOYEE)

DEPARTMENT

| EMPLO' | YEE | DNAME | DN | IUMBER_ | MGRSSN | MGRST | ARTDATE | |
|--------|-----|-------|----|---------|--------|-------|---------|---|
| | | | | | | 1 | | Т |



NESTING OF QUERIES

A complete SELECT query, called a <u>nested query</u>, can be specified within the WHERE-clause of another query

FNAME

Franklin

Ramesh

DNUMBER

Joyce

John

LNAME

Smith

Wong

English

Narayan

Retrieve the name of all employees who work for the

'Research' department.

SELECT FNAME, LNAME

FROM EMPLOYEE

WHERE DNO = (SELECT DNUMBER

FROM DEPARTMENT

WHERE DNAME='Research')

If `=` is used the inner query must return one value

If more than one value is returned then an error msg is generated

| | | DEP | ARTMEN | IT | | | | | |
|-------|--------|-------|--|-------|---------|-----|--------|----------|-----|
| EMI | PLOYEE | DNAM | DNAME <u>DNUMBER</u> MGRSSN MGRSTARTDATE | | | | | | ı |
| FNAME | MINIT | LNAME | SSN | BDATE | ADDRESS | SEX | SALARY | SUPERSSN | DNO |

NESTING OF QUERIES

Retrieve the name of all employees who work for the 'Research' or Administration 'department .

SELECT FNAME, LNAME

FROM EMPLOYEE

WHERE DNO IN (SELECT DNUMBER

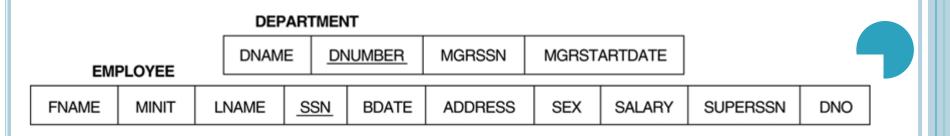
FROM DEPARTMENT

WHERE DNAME='Research' OR

DNAME='Administration')

If `=` is used the inner query must return one value.

If inner query returns more than one value then use IN



NESTING OF QUERIES

Retrieve the name of all employees who do not work for the 'Research' department.

SELECT FNAME, LNAME

FROM EMPLOYEE

WHERE DNO NOT IN (SELECT DNUMBER

FROM DEPARTMENT

WHERE DNAME='Research')

DEPARTMENT

| DNAME | DNUMBER | MGRSSN | MGRSTARTDATE |
|-------|---------|--------|--------------|
|-------|---------|--------|--------------|

EMPLOYEE

| FNAME MINIT LNAME <u>SSN</u> BDATE ADDRESS SEX SALARY SUPERSSN DNC | | FNAME | MINIT | LNAME | SSN | BDATE | ADDRESS | SEX | SALARY | SUPERSSN | DNO |
|--|--|-------|-------|-------|-----|-------|---------|-----|--------|----------|-----|
|--|--|-------|-------|-------|-----|-------|---------|-----|--------|----------|-----|

NESTED QUERIES

You can also use: s > ALL R (means greater than every value)

s > ANY R (means greater than any value)

= ANY is same as IN, <> ALL is same as NOT IN

Find name of employees whose salary is greater than the salary of all employees in department 5

EMPLOYEE

| FNAME | MINIT LNAME | SSN BE | DATE ADDRESS | SEX | SALARY | SUPERSSN | DNO | |
|-------|-------------|--------|--------------|-----|--------|----------|-----|--|
|-------|-------------|--------|--------------|-----|--------|----------|-----|--|

SELECT Fname
FROM Employee
WHERE Salary > ALL (SELECT Salary
FROM Employee
where Dno=5)

| | Fname |
|---|----------|
| 1 | James |
| 2 | Jennifer |

NESTED QUERIES & TSQL

A subquery can be nested inside the **WHERE or HAVING** clause of an outer SELECT, INSERT, UPDATE, or DELETE statement.

Statements with subquery usually take one of these formats:

- WHERE expression [NOT] IN (subquery)
- •WHERE expression comparison_operator [ANY | ALL] (subquery)
- WHERE [NOT] EXISTS (subquery)

- Up to 32 levels of nesting is possible,
 - This limit depends on available memory and the complexity of other expressions in the query.

Why NESTED QUERIES?

- Many Transact-SQL statements that include subqueries can be alternatively formulated as joins.
- Other questions can be posed only with subqueries.
- An aggregate may not appear in the WHERE clause
 - unless it is in a subquery contained in a HAVING clause or a select list, and the column being aggregated is an outer reference



CORRELATED NESTED QUERIES

- If a condition in the nested query references an attribute of a relation declared in the outer query =>
 - Then two queries are said to be correlated

Retrieve the name of each employee who has a dependent with the same first name as the employee.

EMPLOYEE

| Fname | Minit | Lname | Ssn | Bdate | Address | Sex | Salary | Super_ssn | Dno |
|----------|-------|---------|-----------|------------|--------------------------|-----|--------|-----------|-----|
| John | В | Smith | 123456789 | 1965-01-09 | 731 Fondren, Houston, TX | М | 30000 | 333445555 | 5 |
| Franklin | Т | Wong | 333445555 | 1955-12-08 | 638 Voss, Houston, TX | М | 40000 | 888665555 | 5 |
| Alicia | J | Zelaya | 999887777 | 1968-01-19 | 3321 Castle, Spring, TX | F | 25000 | 987654321 | 4 |
| Jennifer | S | Wallace | 987654321 | 1941-06-20 | 291 Berry, Bellaire, TX | F | 43000 | 888665555 | 4 |
| Ramesh | К | Narayan | 666884444 | 1962-09-15 | 975 Fire Oak, Humble, TX | М | 38000 | 333445555 | 5 |
| Joyce | Α | English | 453453453 | 1972-07-31 | 5631 Rice, Houston, TX | F | 25000 | 333445555 | 5 |
| Ahmad | ٧ | Jabbar | 987987987 | 1969-03-29 | 980 Dallas, Houston, TX | М | 25000 | 987654321 | 4 |
| Januar | | Dava | 000005555 | 10071110 | AEO Chana Hauston TV | | 55000 | NII II | |

Nested Correlated query is evaluated once for each tuple in outer query

CORRELATED NESTED QUERIES

A **correlated subquery** (also called repeating subquery) depends on the outer query for its values.

 This means that the nested subquery is executed repeatedly, once for each row that might be selected by the outer query.

Retrieve the name of each employee who has a dependent with the same first name as the employee.

SELECT E.FNAME, E.LNAME

FROM EMPLOYEE AS E

WHERE E.SSN IN (SELECT ESSN FROM DEPENDENT

WHERE SSN = ESSN AND FNAME=DEPENDENT_NAME)

EMPLOYEE

| -[| Fname | Minit | Lname | Ssn | Bdate | Address | Sex | Salary | Super_ssn | Dno |
|----|----------|-------|---------|-----------|------------|--------------------------|-----|--------|-----------|-----|
| | John | В | Smith | 123456789 | 1965-01-09 | 731 Fondren, Houston, TX | М | 30000 | 333445555 | 5 |
| | Franklin | Т | Wong | 333445555 | 1955-12-08 | 638 Voss, Houston, TX | М | 40000 | 888665555 | 5 |
| | Alicia | J | Zelaya | 999887777 | 1968-01-19 | 3321 Castle, Spring, TX | F | 25000 | 987654321 | 4 |
| | Jennifer | S | Wallace | 987654321 | 1941-06-20 | 291 Berry, Bellaire, TX | F | 43000 | 888665555 | 4 |
| | Ramesh | K | Narayan | 666884444 | 1962-09-15 | 975 Fire Oak, Humble, TX | М | 38000 | 333445555 | 5 |
| | Joyce | Α | English | 453453453 | 1972-07-31 | 5631 Rice, Houston, TX | F | 25000 | 333445555 | 5 |
| | Ahmad | ٧ | Jabbar | 987987987 | 1969-03-29 | 980 Dallas, Houston, TX | М | 25000 | 987654321 | 4 |
| | James | E | Borg | 888665555 | 1937-11-10 | 450 Stone, Houston, TX | М | 55000 | NULL | 1 |

NESTED QUERIES

 A query written with nested SELECT... FROM... WHERE... blocks and using the = or IN comparison operators can *always* be expressed as a single block query.

Retrieve the name of each employee who has a dependent with the same first name as the employee.

SELECT E.FNAME, E.LNAME

FROM EMPLOYEE E, DEPENDENT D

WHERE E.SSN=D.ESSN AND

E.FNAME=D.DEPENDENT_NAME



EXISTS Function checks whether the result of a nested query is empty or not

 Retrieve the name of each employee who has a dependent with the same first name as the employee.

SELECT FNAME, LNAME

FROM EMPLOYEE

WHERE EXISTS (SELECT *

FROM DEPENDENT

WHERE SSN=ESSN AND

FNAME=DEPENDENT_NAME)



Retrieve the names of employees who have no dependents.

SELECT FNAME, LNAME

FROM EMPLOYEE

WHERE NOT EXISTS (SELECT *

FROM DEPENDENT WHERE SSN=ESSN)

EXISTS is necessary for the expressive power of SQL

EMPLOYEE

| Fname | Minit | Lname | <u>Ssn</u> | Bdate | Address | Sex | Salary | Super_ssn | Dno |
|----------|-------|---------|------------|------------|--------------------------|-----|--------|-----------|-----|
| John | В | Smith | 123456789 | 1965-01-09 | 731 Fondren, Houston, TX | М | 30000 | 333445555 | 5 |
| Franklin | Т | Wong | 333445555 | 1955-12-08 | 638 Voss, Houston, TX | М | 40000 | 888665555 | 5 |
| Alicia | J | Zelaya | 999887777 | 1968-01-19 | 3321 Castle, Spring, TX | F | 25000 | 987654321 | 4 |
| Jennifer | S | Wallace | 987654321 | 1941-06-20 | 291 Berry, Bellaire, TX | F | 43000 | 888665555 | 4 |
| Ramesh | K | Narayan | 666884444 | 1962-09-15 | 975 Fire Oak, Humble, TX | М | 38000 | 333445555 | 5 |
| Joyce | Α | English | 453453453 | 1972-07-31 | 5631 Rice, Houston, TX | F | 25000 | 333445555 | 5 |
| Ahmad | ٧ | Jabbar | 987987987 | 1969-03-29 | 980 Dallas, Houston, TX | М | 25000 | 987654321 | 4 |

James

The above correlated nested query retrieves all DEPENDENT tuples related to an EMPLOYEE tuple.

If none exist, the EMPLOYEE tuple is selected

Find the names of managers who have at least one dependents.

DEPARTMENT

| DNAME | DNUM | BER M | MGRSSN | MGRSTA | ARTDATE | | | | | | | |
|-------|--------|-------|--------|---------|---------|-------|-----------|-----|------|----------|----------|-------|
| | | | [| DEPENDE | NT | | | | | | | |
| EMF | PLOYEE | | | Essn | _ | Deper | ndent_nar | me_ | Sex | Bdate | Relation | nship |
| FNAME | MINIT | LNAME | SSN | BDATE | ADDRE | SS | SEX | SAL | _ARY | SUPERSSN | DNO | |

Retrieve the name of each employee who works on all the projects controlled by department number 4.

Set theory: S1 contains S2 if (S2 - S1 = 0)

S1 = set of projects of each employee

S2 = set of Dept 4 projects

PROJECT

| Pname | Pnumber | Plocation | Dnum |
|-----------------|---------|-----------|------|
| ProductX | 1 | Bellaire | 5 |
| ProductY | 2 | Sugarland | 5 |
| ProductZ | 3 | Houston | 5 |
| Computerization | 10 | Stafford | 4 |
| Reorganization | 20 | Houston | 1 |
| Newbenefits | 30 | Stafford | 4 |

WORKS_ON

| Essn | Pno | Hours |
|-----------|-----|-------|
| 123456789 | 1 | 32.5 |
| 123456789 | 2 | 7.5 |
| 666884444 | 3 | 40.0 |
| 453453453 | 1 | 20.0 |
| 453453453 | 2 | 20.0 |
| 333445555 | 2 | 10.0 |
| 333445555 | 3 | 10.0 |
| 333445555 | 10 | 10.0 |
| 333445555 | 20 | 10.0 |
| 999887777 | 30 | 30.0 |
| 999887777 | 10 | 10.0 |
| 987987987 | 10 | 35.0 |
| 987987987 | 30 | 5.0 |
| 987654321 | 30 | 20.0 |
| 987654321 | 20 | 15.0 |
| 888665555 | 20 | NULL |

Retrieve the name of each employee who works on all the projects controlled by department number 4.

Set theory: S1 contains S2 if (S2 - S1 = 0)

SELECT FNAME, LNAME

FROM EMPLOYEE

WHERE NOT EXISTS (

(SELECT PNUMBER

FROM PROJECT

WHERE DNUM=4)

EXCEPT

(SELECT PNO FROM WORKS_ON WHERE SSN=ESSN) S1 = set of projects of each employee

S2 = set of dept 4 projects

PROJECT

| Pname | Pnumber | Plocation | Dnum |
|-----------------|---------|-----------|------|
| ProductX | 1 | Bellaire | 5 |
| ProductY | 2 | Sugarland | 5 |
| ProductZ | 3 | Houston | 5 |
| Computerization | 10 | Stafford | 4 |
| Reorganization | 20 | Houston | 1 |
| Newbenefits | 30 | Stafford | 4 |

WORKS_ON

| Essn | <u>Pno</u> | Hours |
|-----------|------------|-------|
| 123456789 | 1 | 32.5 |
| 123456789 | 2 | 7.5 |
| 666884444 | 3 | 40.0 |
| 453453453 | 1 | 20.0 |
| 453453453 | 2 | 20.0 |
| 333445555 | 2 | 10.0 |
| 333445555 | 3 | 10.0 |
| 333445555 | 10 | 10.0 |
| 333445555 | 20 | 10.0 |
| 999887777 | 30 | 30.0 |
| 999887777 | 10 | 10.0 |
| 987987987 | 10 | 35.0 |
| 987987987 | 30 | 5.0 |
| 987654321 | 30 | 20.0 |
| 987654321 | 20 | 15.0 |
| 888665555 | 20 | NULL |

HOW to do this in Relational Algebra?



Yet another Example

Find SSN of employees who work on all the projects of Dnum= 4

PROJECT

| Pname | Pnumber | Plocation | Dnum |
|-----------------|---------|-----------|------|
| ProductX | 1 | Bellaire | 5 |
| ProductY | 2 | Sugarland | 5 |
| ProductZ | 3 | Houston | 5 |
| Computerization | 10 | Stafford | 4 |
| Reorganization | 20 | Houston | 1 |
| Newbenefits | 30 | Stafford | 4 |

- PD4(Pno) $\leftarrow \pi_{\text{Pnumber}}$ ($\sigma_{\text{DNUM=4}}$ Project)
- Ssn_Pnos $\leftarrow \pi_{Essn.Pno}$ (Works_on)
- SSNS(ssn) ← Ssn_Pnos ??? PD4

WORKS_ON

| Essn | <u>Pno</u> | Hours |
|-----------|------------|-------|
| 123456789 | 1 | 32.5 |
| 123456789 | 2 | 7.5 |
| 666884444 | 3 | 40.0 |
| 453453453 | 1 | 20.0 |
| 453453453 | 2 | 20.0 |
| 333445555 | 2 | 10.0 |
| 333445555 | 3 | 10.0 |
| 333445555 | 10 | 10.0 |
| 333445555 | 20 | 10.0 |
| 999887777 | 30 | 30.0 |
| 999887777 | 10 | 10.0 |
| 987987987 | 10 | 35.0 |
| 987987987 | 30 | 5.0 |
| 987654321 | 30 | 20.0 |
| 987654321 | 20 | 15.0 |
| 888665555 | 20 | NULL |

DIVISION

Yet an other Example

Find SSN of employees who work on all the projects of Dnum= 4

PROJECT

| Pname | Pnumber | Plocation | Dnum |
|-----------------|---------|-----------|------|
| ProductX | 1 | Bellaire | 5 |
| ProductY | 2 | Sugarland | 5 |
| ProductZ | 3 | Houston | 5 |
| Computerization | 10 | Stafford | 4 |
| Reorganization | 20 | Houston | 1 |
| Newbenefits | 30 | Stafford | 4 |

PD4

Pno

10

30

SSN_PNOS

| Essn | Pno |
|-----------|-----|
| 123456789 | 1 |
| 123456789 | 2 |
| 666884444 | 3 |
| 453453453 | 1 |
| 453453453 | 2 |
| 333445555 | 2 |
| 333445555 | 3 |
| 333445555 | 10 |
| 333445555 | 20 |
| 999887777 | 30 |
| 999887777 | 10 |
| 987987987 | 10 |
| 987987987 | 30 |
| 987654321 | 30 |
| 987654321 | 20 |
| 888665555 | 20 |

- PD4(Pno) $\leftarrow \pi_{\text{Pnumber}}$ ($\sigma_{\text{DNUM=4}}$ Project)
- Ssn_Pnos $\leftarrow \pi_{Essn,Pno}$ (Works_on)
- SSNS(ssn) ← Ssn_Pnos ÷ PD4

DIVISION

DIVISION (Binary Operation)

Division operation is applied to two relations R1 and R2

R1(Attributes_R1) ÷ R2(Attributes_R2)

Let Result = $R1 \div R2$

Attr_Res = Attributes_R1 - Attributes_R2

 Attr_Res is a set of attributes of R1 that are not the attributes of R2. R1

A B
a1 b1
a2 b1
a3 b1
a4 b1
a1 b2
a3 b2
a3 b2
a3 b2
a4 b3
a4 b3
a4 b3
a4 b3
a1 b4

b4

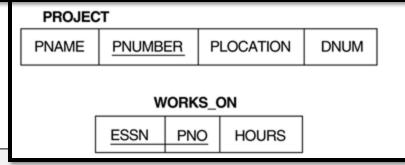
For a **tuple t** to appear in the result of the DIVISION, the values in t must appear in R1 in combination with *every* tuple in R2.

Example of DIVISION

Find **SSN** of employees who work on all the projects that *John Smith* works on



| LIVII | LOILL | | | | | | | | |
|-------|-------|-------|-----|-------|---------|-----|--------|----------|-----|
| FNAME | MINIT | LNAME | SSN | BDATE | ADDRESS | SEX | SALARY | SUPERSSN | DNO |



- Smith $\leftarrow \sigma_{\text{fname='John'} \text{ and } \text{lname='Smith'}}$ (Employee)
- Smith_Pnos $\leftarrow \pi_{Pno}$ (Works_on essn=ssn Smith)
- $Ssn_Pnos \leftarrow \pi_{Essn_Pno}$ (Works_on)
- SSNS(ssn) ← Ssn_Pnos ÷ Smith_Pnos

SSN PNOS

Essn 123456789 123456789

666884444

987987987

987987987

987654321

987654321

888665555

Pno

10

30

30

20

20

| | 453453453 | 1 |
|------------|-----------|----|
| | 453453453 | 2 |
| | 333445555 | 2 |
| SMITH_PNOS | 333445555 | 3 |
| Pno | 333445555 | 10 |
| 1 | 333445555 | 20 |
| 2 | 999887777 | 30 |
| | 999887777 | 10 |
| SSNS | 987987987 | 10 |

Ssn

123456789

453453453

Examples of Queries in RA

Find the names of employees who work on *all* the projects controlled by department number 5.

T1(Pno)
$$\leftarrow \pi_{\text{Pnumber}} (\sigma_{\text{Dnum}=5} (\text{Project}))$$

T2 $\leftarrow \pi_{\text{Essn, Pno}} (\text{Work_On})$
T3 $\leftarrow (\text{T2} \div \text{T1})$
R $\leftarrow \pi_{\text{LNAME, FNAME}} (\text{T3} * \text{Employee})$

PROJECT

| Pname | Pnumber | Plocation | Dnum |
|-----------------|---------|-----------|------|
| ProductX | 1 | Bellaire | 5 |
| ProductY | 2 | Sugarland | 5 |
| ProductZ | 3 | Houston | 5 |
| Computerization | 10 | Stafford | 4 |
| Reorganization | 20 | Houston | 1 |
| Newbenefits | 30 | Stafford | 4 |

| Essn | Pno |
|-----------|-----|
| 123456789 | 1 |
| 123456789 | 2 |
| 666884444 | 3 |
| 453453453 | 1 |
| 453453453 | 2 |
| 333445555 | 2 |
| 333445555 | 3 |
| 333445555 | 10 |
| 333445555 | 20 |
| 999887777 | 30 |
| 999887777 | 10 |
| 987987987 | 10 |
| 987987987 | 30 |
| 987654321 | 30 |
| 987654321 | 20 |
| 888665555 | 20 |
| | |