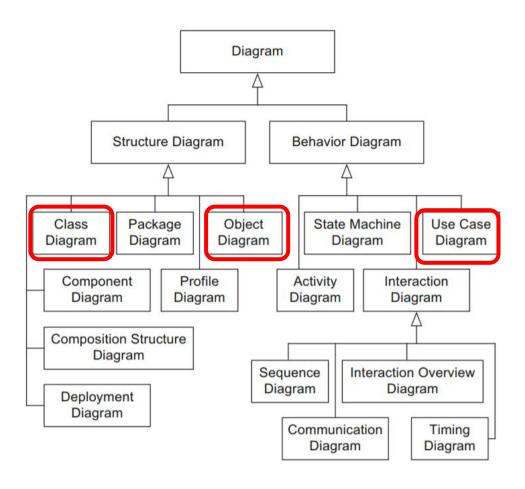


### Lecture 5: UML Part 2





# UML Diagrams





# Summary: Use Case Diagram

			<u> </u>	21 20	1 3
Name	Notation	Description		1 9	
System	System A X	Boundaries between the system and the users of the system	Generalization (actor)	x x	Y inherits from X; Y participates in all use cases in which X participates
Use case	A	Unit of functionality of the system		Y	1
Actor	«actor» or X	Role of the users of the system	Extend relationship	A settends	B extends A: optional incorporation of use case B into use case A
Association	$\begin{vmatrix} \bigcirc \\ X \end{vmatrix}$ $\longrightarrow$ A	X participates in the execution of A		В	
Generalization (use case)	A	B inherits all properties and the entire behavior of A	Include relationship	A sincludes B	A includes B: required incorporation of use case B into use case A



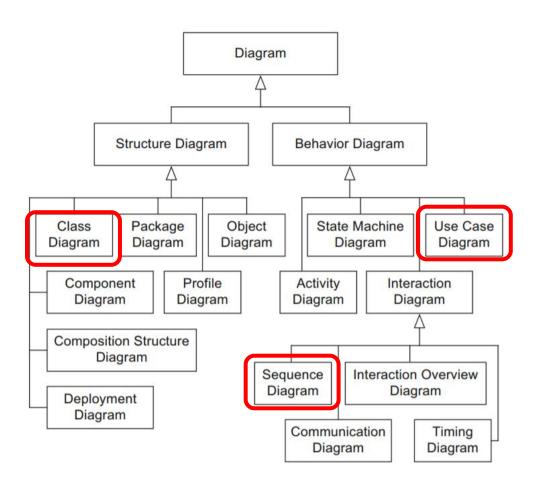
# Summary: Class Diagram

Name	Notation	Description	A	_
	A - a1: T1	Description of the structure and be-	Association class	(
Class	- a2: T2 + o1(): void + o2(): void	havior of a set of objects	xor relationship	{xc
Abstract class	A {abstract}	Class that cannot be instantiated	Strong aggregation = A	
	A B			
Association	(a) B (b)	Relationship between classes: navigability unspecified (a), navigable in both directions (b), not navigable in	Shared aggregation A	
	A × B one direction (c)	Generalization A		
N any association	A -> B	Relationship between $N$ (in this case	Object	0:
N-ary association	С	3) classes	Link 01	_

Association class	A B	More detailed description of an association
xor relationship	B {xor}, C	An object of A is in a relationship with an object of B or with an object of C but not with both
Strong aggregation = composition	A B	Existence-dependent parts-whole relationship (A is part of B; if B is deleted, related instances of A are also deleted)
Shared aggregation	_ A _ → B	Parts-whole relationship (A is part of B; if B is deleted, related instances of A need not be deleted)
Generalization	A	Inheritance relationship (A inherits from B)
Object	<u>o:C</u>	Instance of a class
Link	<u>01</u> <u>02</u>	Relationship between objects



# UML Diagrams





## Sequence Diagram

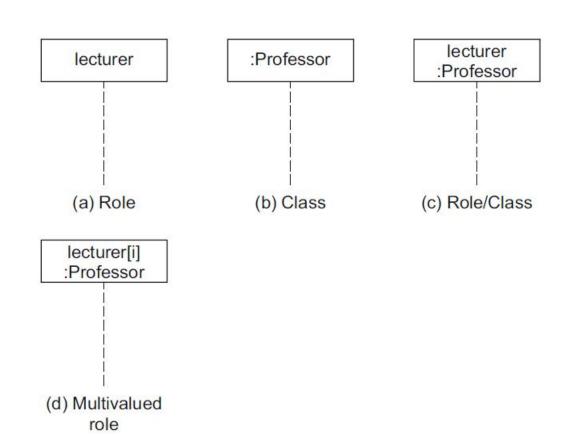
- Message among interaction partners
- Can be constructed with different granularity at different design stages
  - Interaction between the system and its environment
  - Interaction among system parts
  - Interaction between design objects



#### Interaction Partners

- Lifeline
  - r: role
  - C: class
- Use roles instead of objects
  - Each object can play different roles

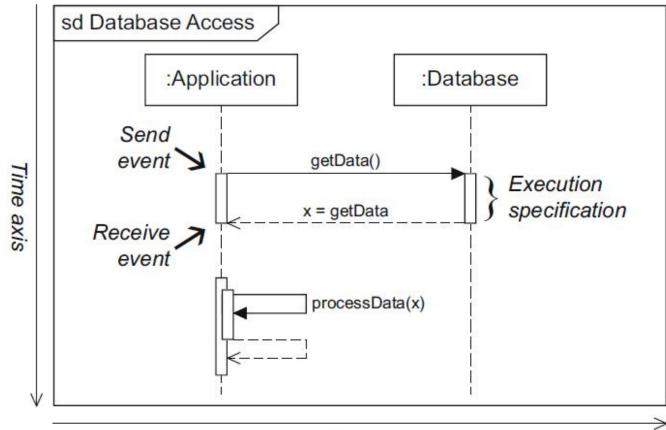
r:C





# Message Exchange

- Two dimensions
  - Time
  - Interaction partners
- Execution specification
  - Self message

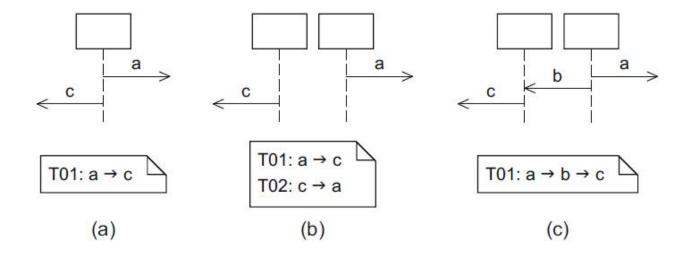


Interaction partners



## Message Exchange: Order

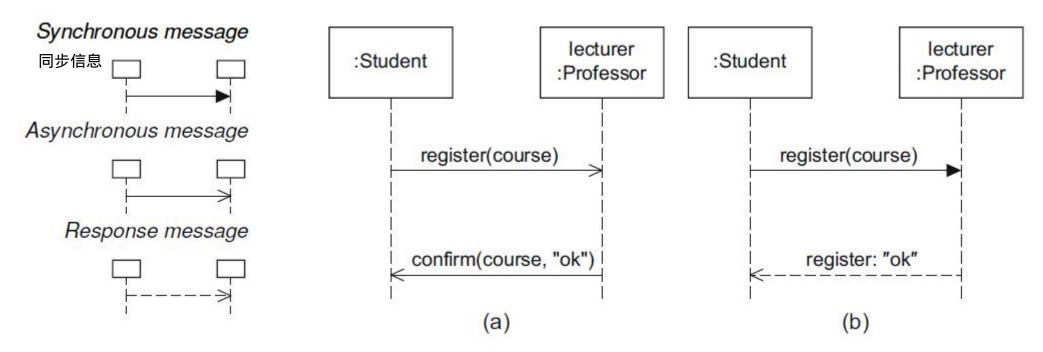
- Message order is chronical if messages on the same lifeline
  - It's a transitive relationship





# Message Exchange: Types

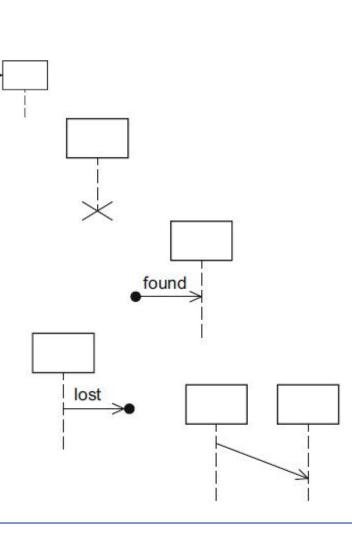
- (a) Register a course via email
- (b) Register a course in person





# Special Messages

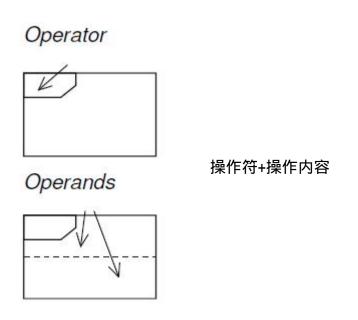
- Create message
  - Creating new object
- Destruction event
  - Destruction of an object
- Found message
  - Unknown/irrelevant sender
- Lost message
  - Unknown/irrelevant receiver
- Time-consuming message





# Combined Fragments

• Each operand has a guard



	Operator	Purpose
	alt	Alternative interaction
Duonahas and lagus	opt	Optional interaction
Branches and loops	loop	Iterative interaction
	break	Exception interaction
	seq	Weak order
Consumoner and order	strict	Strict order
Concurrency and order	par	Concurrent interaction
	critical	Atomic interaction

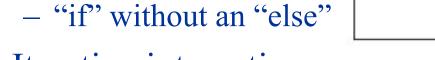


## Branches and loops

- Alternative interactions
  - If-else

上下分别对应if、else,条件写在方括号中

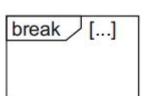
- Optional interactions



- Iterative interactions
  - For loop

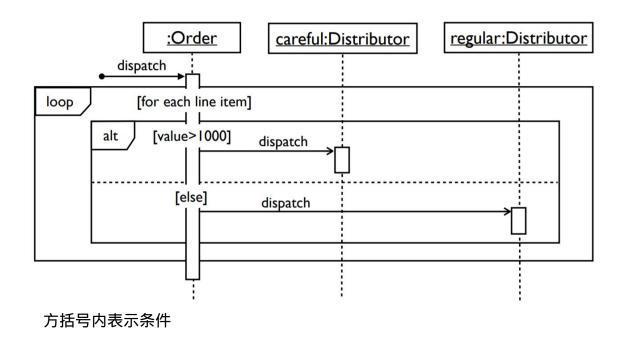


Omit the remaining



loop(...) / [...]

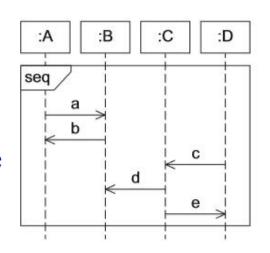
opt [...]



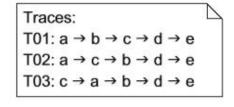


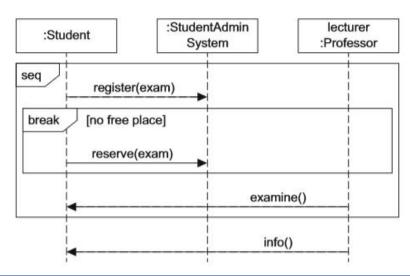
## Concurrency and Order

- Seq fragment
- seq
- Weak order
- The ordering of events within each of the operands is maintained in the result.
- Events on different lifelines from different operands may come in any order.
- Events on the same life line from different operands are ordered such that an event of the first operand comes before that of the second operand.





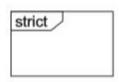






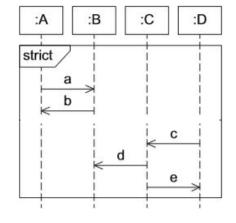
# Concurrency and Order (cont.)

• Strict fragment



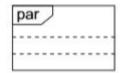
Strong & strict order

顺序严格按照流程图

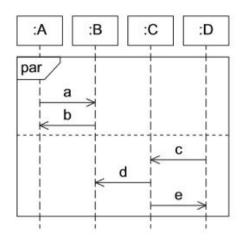


Traces: T01:  $a \rightarrow b \rightarrow c \rightarrow d \rightarrow e$ 

• Par fragment



- Order within the operands are respected
- The order of operands does not matter 并行的 -> 只要不在同一区段,顺序自由

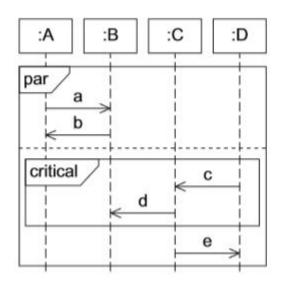


Traces: T01:  $a \rightarrow b \rightarrow c \rightarrow d \rightarrow e$ T02:  $a \rightarrow c \rightarrow b \rightarrow d \rightarrow e$ T03:  $a \rightarrow c \rightarrow d \rightarrow b \rightarrow e$ T04:  $a \rightarrow c \rightarrow d \rightarrow e \rightarrow b$ T05:  $c \rightarrow a \rightarrow b \rightarrow d \rightarrow e$ T06:  $c \rightarrow a \rightarrow d \rightarrow b \rightarrow e$ T07:  $c \rightarrow a \rightarrow d \rightarrow e \rightarrow b$ T08:  $c \rightarrow d \rightarrow a \rightarrow b \rightarrow e$ T09:  $c \rightarrow d \rightarrow a \rightarrow e \rightarrow b$ T10:  $c \rightarrow d \rightarrow e \rightarrow a \rightarrow b$ 



# Concurrency and Order (cont.)

- Critical fragment
- critical
- Atomic interaction
- No other messages can happen during the execution



```
Traces:

T01: a \rightarrow b \rightarrow c \rightarrow d \rightarrow e

T02: a \rightarrow c \rightarrow d \rightarrow b \rightarrow e

T03: a \rightarrow c \rightarrow d \rightarrow e \rightarrow b

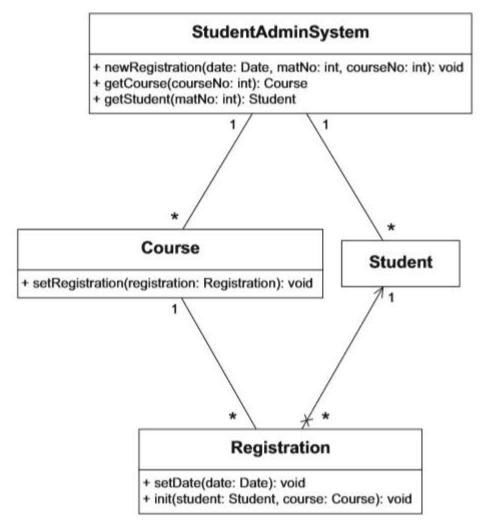
T04: c \rightarrow d \rightarrow a \rightarrow b \rightarrow e

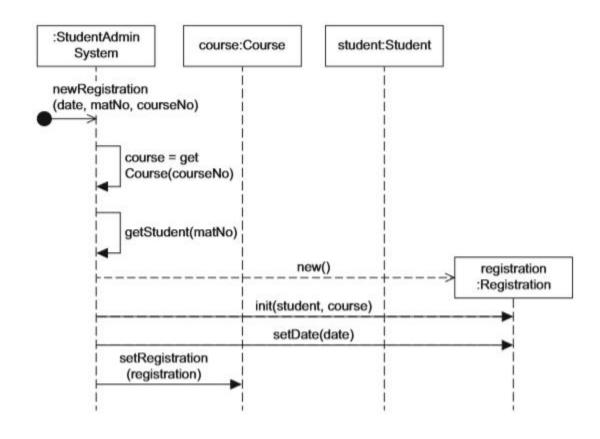
T05: c \rightarrow d \rightarrow a \rightarrow e \rightarrow b

T06: c \rightarrow d \rightarrow e \rightarrow a \rightarrow b
```



# Example



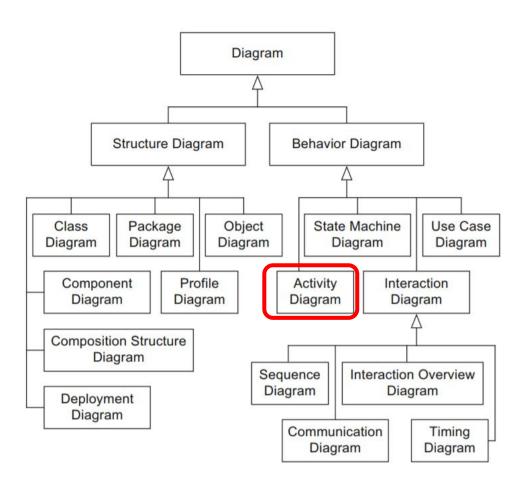




Name	Notation	Description
Lifeline	r:C	Interaction partners involved in the communication
Destruction event	×	Time at which an interaction partner ceases to exist
Combined fragment	[]	Control constructs
Synchronous message		Sender waits for a response message
Response message	ļ	Response to a synchronous message
Asynchronous mes- sage	- <del>k</del> -	Sender continues its own work after sending the asynchronous message
Lost message	lost	Message to an unknown receiver
Found message	found	Message from an unknown sender



# UML Diagrams





# Activity Diagram

• What are the procedures of a system?

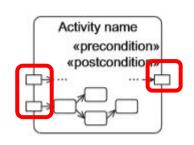
• At conceptual level: How to implement use case?

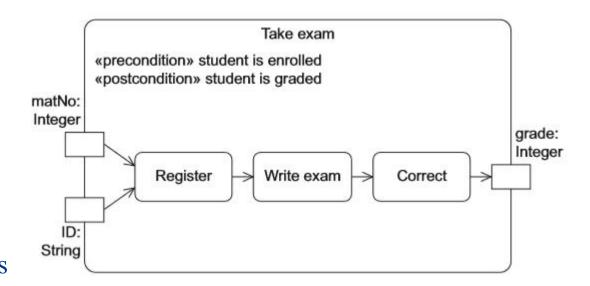
- At implementation level: How to implement an operation
- A flow-oriented language



# Activity Diagram: Syntax

- Activity
  - Parameters
  - Precondition
  - Postcondition
  - Actions Action
    - No language restrictions
    - Atomic: may be further broken down in other contexts
  - − EdgesAB
    - Control flow edge: order between actions
    - Object flow edge: can exchange data

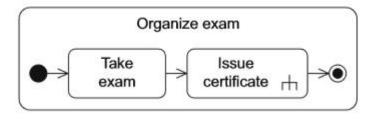


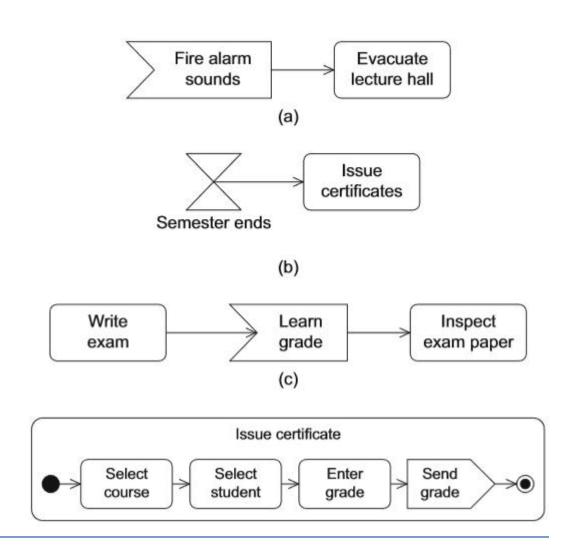




#### Predefined Actions

- Event-based actions
  - Accept event action
    - Wait for a specific event E
  - Accept time event action
    - For time-based events
  - Send signal action
- Call behavior actions
  - Calling other activities



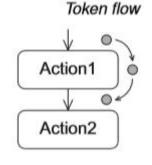


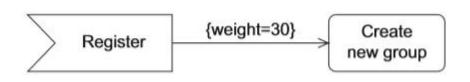


#### Control Flow

- Semantics for executing activities
- Token
  - Multiple incoming edges: token must be present for all incoming edges
  - Multiple outgoing edges: tokens are given to all edges
- Edge A B
  - Guard
    - Similar to the guard definition in the State Machine diagram
  - Weight of an edge

    A (weight=X)
    - Tokens consumed on the edge

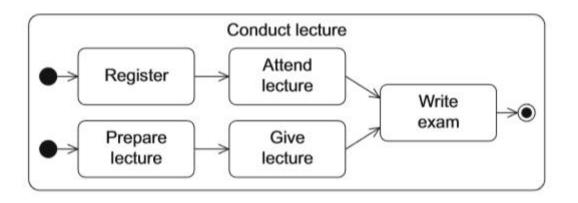


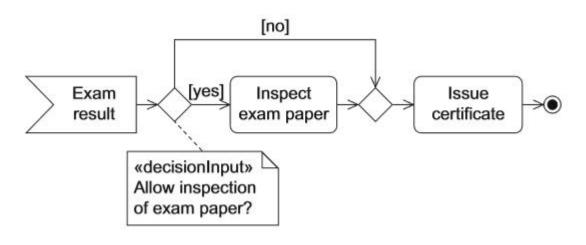




# Control Flow (Cont.)

- Connector
- $A \rightarrow y y \Rightarrow B$
- Just to make the diagram clearer
- Control nodes
  - Initial node
  - Activity final node →
  - Decision node
  - Decision behavior
    - Save space & provide clarity
  - Merge node





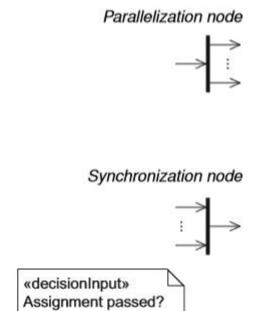
«decisionInput»

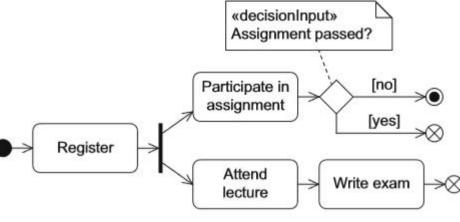
decision behavior



# Control flow (cont.)

- Parallelization & Synchronization node
  - Decision node can take only one edge
- Activity final node
  - Multiple final node: first reached final node terminates the activity
- Flow final node ⊗
  - For concurrent activities only
  - Only terminate one concurrent path

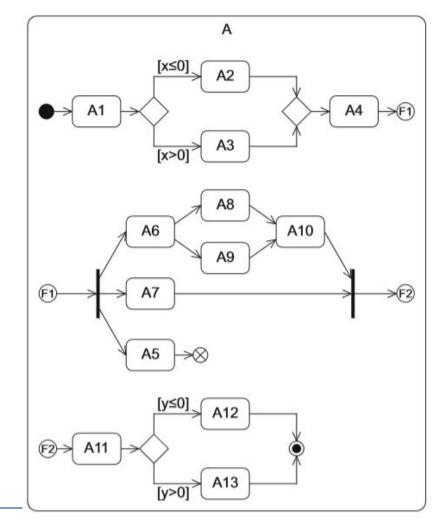






## Example: Control flow

• If A5 is still executing when the activity final node is reached, A5 is interrupted

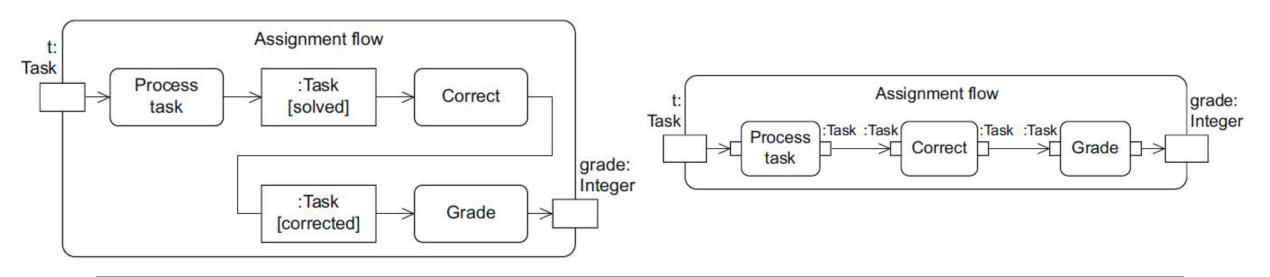




# Object Flow: Syntax

- Exchange of data among actions
- ObjectObject
- Pin notation

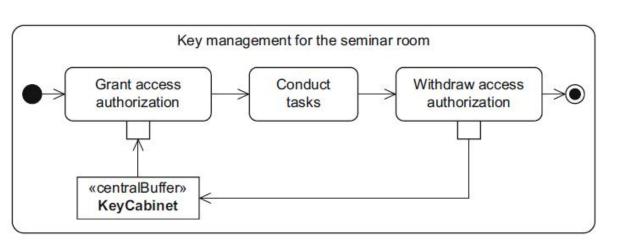
- Action
- Parameters are also objects

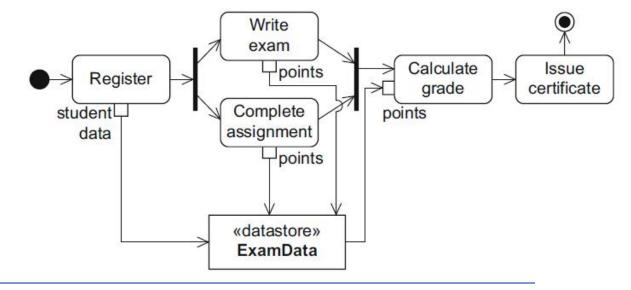




# Object Flow: Syntax (cont.)

- Central buffer
- «centralBuffer» CB
- Data exited the central buffer is no longer in there
- Data store datastore»
  - Data exited the data store still has a copy in there

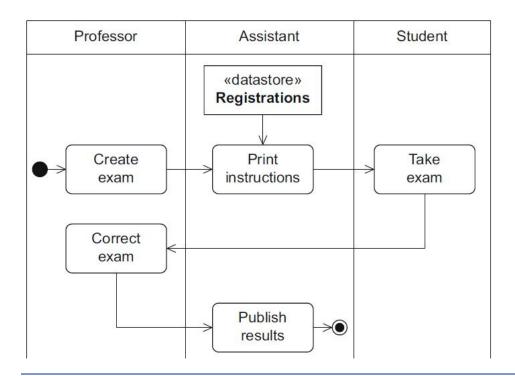






#### Swimlane/Partition

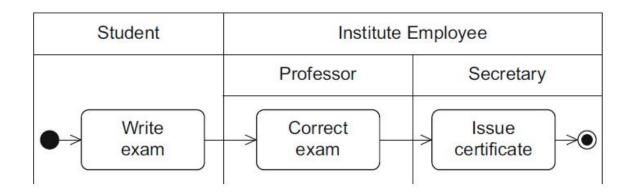
- Group actions in terms of who's performing them
- A much clearer view of the activity diagram

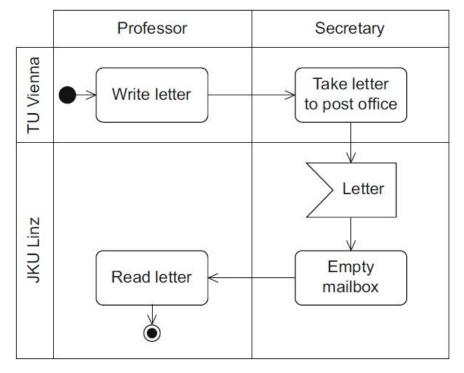




## Swimlane (cont.)

- Swim lanes can have sub-partitions
- Swim lanes can also have multiple dimensions



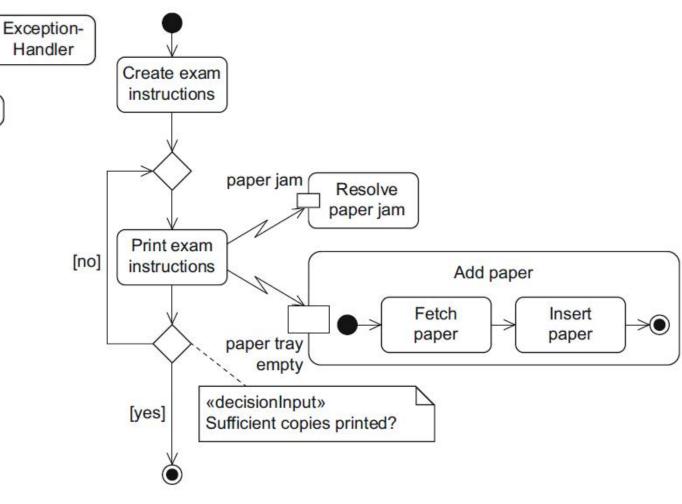




# **Exception Handling**

Action

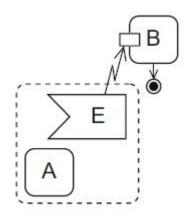
- Exception handler node
  - Error situation e
  - Execute exception handler when e happens
  - Then the sequence continues as if the action ended normally
  - One action can have multiple exception handlers

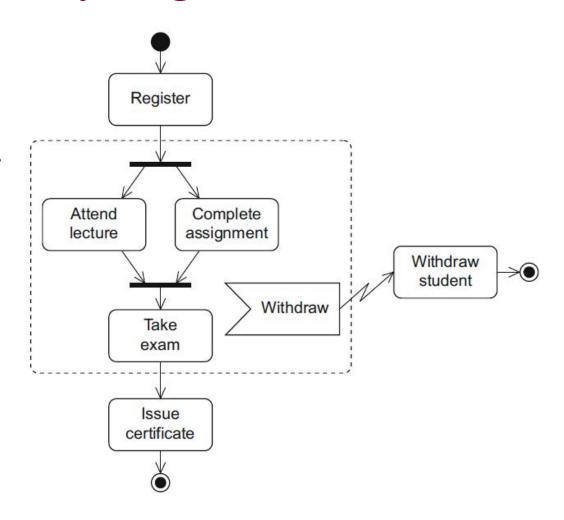




## Interruptible activity region

- Interruptible activity region
- Activities within the dashed region terminate immediately when E occur

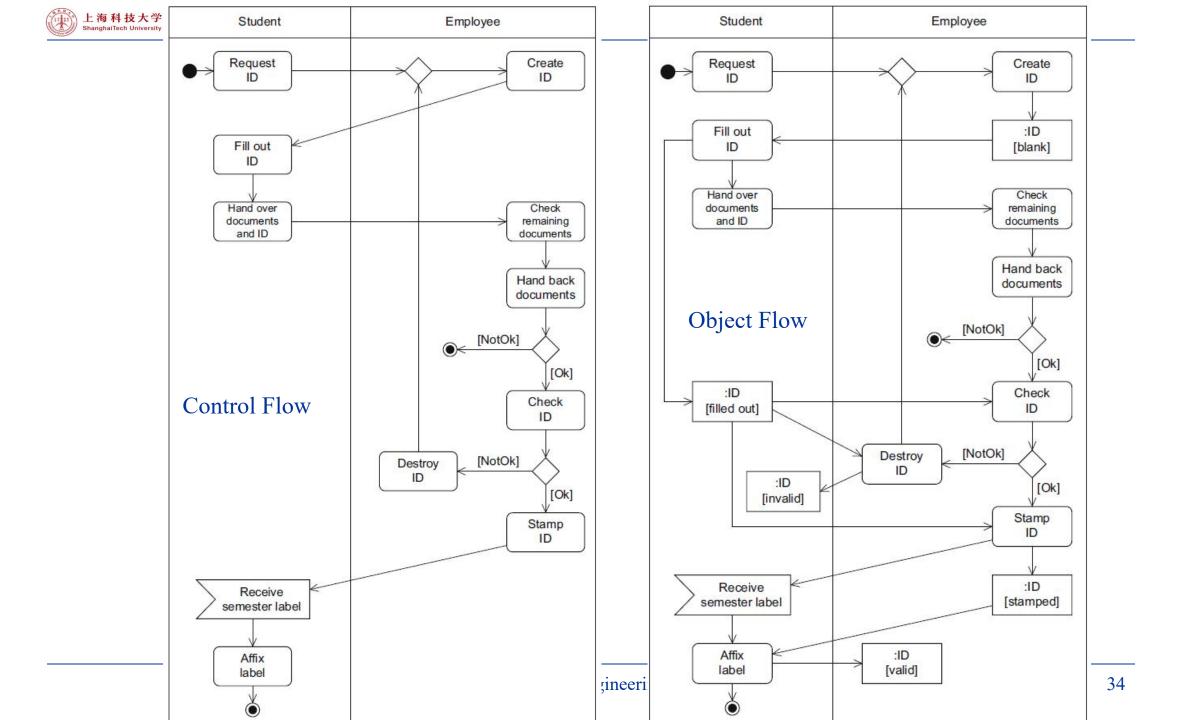






## Example: Student ID

- To obtain a student ID, the student must request this ID from an employee of the student office.
- The employee hands the student the forms that the student has to fill out to register at the university.
- These forms include the student ID itself, which is a small, old-style cardboard card.
- The student has to enter personal data on this card and the employee confirms it with a stamp after checking it against certain documents.
- Once the student has filled out the forms, the student returns them to the employee in the student office and hands over documents such as photo identification, school-leaving certificate, and birth certificate.
- The employee checks the documents. If the documents are incomplete or the student is not authorized to receive a student ID for the university, the process is terminated immediately.
- If the documents are all in order, the employee checks whether the student has filled out the student ID correctly.
- If there are any errors, this ID is destroyed and the student has to fill out another one. Otherwise the ID is stamped.
- However, the student ID is not valid until it bears the semester label sent to the student by post.



Name	Notation	Description
Action node	Action	Actions are atomic, i.e., they cannot be broken down further
Activity node	Activity	Activities can be broken down further
Initial node	•	Start of the execution of an activity
Activity final node	•	End of ALL execution paths of an activity
Flow final node	$\otimes$	End of ONE execution path of an activity
Decision node	<b>→</b>	Splitting of one execution path into two or more alternative execution paths
Merge node	<b>⋙</b>	Merging of two or more alternative execution paths into one execution path
Parallelization node	→ :: ->	Splitting of one execution path into two or more concurrent execution paths
Synchronization node	→ : →	Merging of two or more concurrent execution paths into one execution path
Edge		Connection between the nodes of an activity
Call behavior action	A H	Action A refers to an activity of the same name
Object node	Object	Contains data and objects that are created, changed, and read
Parameters for activi- ties	Activity 🗦	Contain data and objects as input and output parameters
Parameters for actions (pins)	Action	Contain data and objects as input and output parameters

Name	Notation	Description
Partition	A B A	Grouping of nodes and edges within an activity
Send signal action	s	Transmission of a signal to a receiver
Asynchronous accept (time) event action	E or T	Wait for an event E or a time event T
Exception handler	e Exception- Handler  Action	Exception handler is executed instead of the action in the event of an error e
Interruptible activity region	B E A	Flow continues on a different path if event E is detected



## Why do we need models?

#### Prediction

- We know the low-level mechanisms but we want to understand how they affect higher-level behaviors
- Use simulation instead of testing on the real system

#### • Explain the data

 Make assumptions and use our knowledge to explain mechanisms that we don't understand

#### Classification

- i.e. definitions, machine learning algorithms



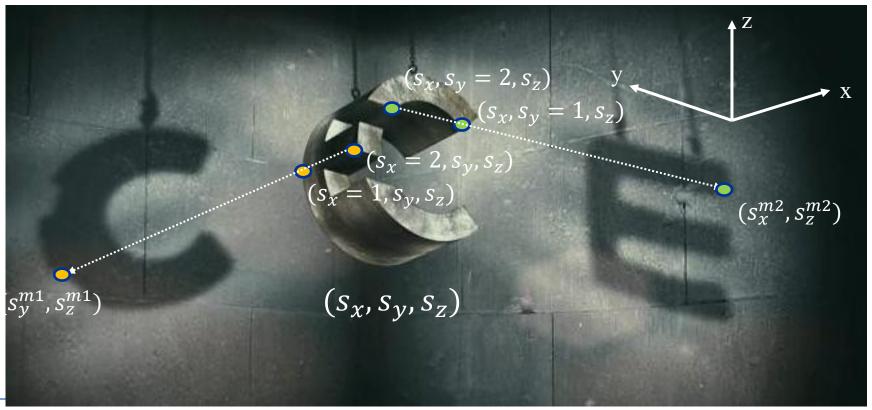
#### What are models?

- A system: (*S*, *I*, *T*, *O*)
  - S: States  $s_1$ ,  $s_2$ ... $s_n$
  - *I*: Inputs (could be ∅)
  - T: Transitions  $S \times I \times S$
  - *O*: Observations  $f(S_o)$ ,  $S_o \subseteq S$

- Model of the system  $(S^m, I^m, T^m)$ 
  - S<sup>m</sup>: Abstraction/interpolation of S
    - Much fewer state variables
  - $I^{\rm m}$ : abstraction of I (could be  $\emptyset$ )
  - $T^m$ : Transitions  $S^m \times I^m \times S^m$

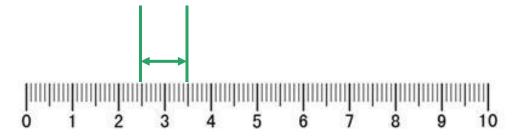
### Abstraction – removal of state variables

- States  $(s_x, s_y, s_z)$  are abstracted to  $(s_y^{m1}, s_z^{m1})$ 
  - $\left( S_{\chi}, S_{\gamma}, S_{Z} \right) \rightarrow \left( S_{\gamma}^{m1}, S_{Z}^{m1} \right)$
- Loss of information



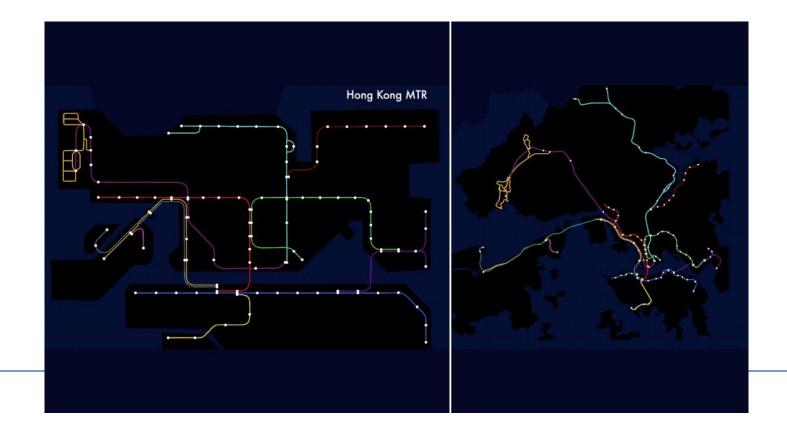
### Abstraction: Approximation of state variable values

- Irrational numbers
  - $-\pi \approx 3.1415$
  - $-\sqrt{2}\approx 1.414$
- Approximation is another way of abstraction



### Interpolation: extracting interpretable information

- Locational information -> topological information
- $S^m = f(S_p), S_p \subseteq S$





# More Interpolation: London MTR





# What is considered as a "good" model?

- Accuracy
  - All models are wrong!
  - Error accumulates over time
  - Initial condition of the model cannot be determined due to limited observability
- Generality
  - The capability to explain not only training data, but also testing data
- Identifiability
  - Model parameters can be identified from data
- Interpretability
  - $-S^m$  are meaningful and interpretable by human

#### Newton vs. Einstein

Newtonian physics is suitable for macro level objects at low speed

$$\bullet \ L = L_0 \sqrt{1 - \frac{v^2}{c^2}}$$

• A model can only be "good" within the context of its designated application

• The definition of "goodness" is changing over time



# Modeling methodologies

- Bottom-up modeling
  - "White-box" model
  - Using first principles
  - Pros:
    - Interpretable
    - Convincing
  - Cons:
    - State space explosion
    - Difficult to be general
    - Low identifiability



- Data driven models (i.e. Neural networks)
  - "Black-box" model
  - From observable data
  - Pros:
    - No need to know domain knowledge
  - Cons:
    - Large and uninterpretable  $S^m$
    - Depends highly on the quality and quantity of data