

# 1 Set Theory

## 1.1 Set Axioms

### 1.1.1 Undefined notions

*Set:*  $A, B, C, \dots$

### 1.1.2 Axioms

1. *Extension:*  $\forall A \forall B [\forall C (C \in A \Leftrightarrow C \in B) \Rightarrow A = B]$
2. *Regularity:*  $\forall A [\exists C (C \in A) \Rightarrow \exists B (B \in A \wedge \neg \exists D (D \in B \wedge D \in A))]$   
(Every nonempty set contains a set that is disjoint from it. Also know as "Axiom of Foundation.")
3. *Schema of Specification:*  $\forall B \forall X_1 \forall X_2 \dots \forall X_n \exists A \forall C [C \in A \Leftrightarrow (C \in B \wedge \phi)]$
4. *Pairing:*  $\forall X_1 \forall X_2 \exists A (X_1 \in A \wedge X_2 \in A)$
5. *Union:*  $\forall \mathcal{F}_A \exists U \forall A \forall X [(X \in A \wedge A \in \mathcal{F}_A) \Rightarrow X \in U]$
6. *Schema of Replacement:*  $\forall A \forall X_1 \forall X_2 \dots \forall X_n [\forall B (B \in A \Rightarrow \exists! D \phi) \Rightarrow \exists B \forall C (C \in A \Rightarrow \exists D (D \in B \wedge \phi))]$
7. *Infinity:*  $\exists \omega_0 [\emptyset \in \omega_0 \wedge \forall X (X \in \omega_0 \Rightarrow X \cup X) \in \omega_0]$
8. *Power Set:*  $\forall X \exists \mathcal{P}(X) \forall S [S \subseteq X \Rightarrow S \in \mathcal{P}(X)]$
9. *Empty Set:*  $\exists A \forall X (X \notin A)$
10. *Choice:*  $\forall X [\emptyset \notin X \Rightarrow \exists (f : X \rightarrow \bigcup X) \forall A \in X (f(A) \in A)]$

**Proposition 1.1.** *The empty set axiom is implied by the other nine axioms.*

*Proof.* Just choose any formula that is always false such as  $\phi(X) = X \in B \wedge X \notin B$  and apply the axiom schema of specification. This will give the empty set. The axiom of extension proves uniqueness vacuously.  $\square$

### 1.1.3 Universe

A set  $U$  is defined with the following properties. . .

1.  $x \in u \in U \Rightarrow x \in U$
2.  $u \in U \wedge v \in U \Rightarrow \{u, v\}, \langle u, v \rangle, u \times v \in U$
3.  $X \in U \Rightarrow \mathcal{P}(X) \in U \wedge \bigcup X \in U$
4.  $\omega_0 \in U$  is the set of finite ordinals
5. if  $f : A \rightarrow B$  is a surjective function with  $A \in U \wedge B \subset U$ , then  $B \in U$   
(See: Set Constructions.)

In category theory, *small sets* are members of  $U$ .

## 1.2 Set Constructions

### 1.2.1 Union

- $A \cup B := \{x | x \in A \vee x \in B\}$
- $\bigcup \mathcal{F} := \{x | x \in X \text{ for some } X \in \mathcal{F}\}$

**Proposition 1.2.** *For sets  $A, B, C \dots$*

- Property Name?:  $A \cup \emptyset = A$
- Idempotence:  $A \cup A = A$
- Property Name?:  $A \subseteq B \Leftrightarrow A \cup B = B$
- Commutative:  $A \cup B = B \cup A$
- Associative:  $A \cup (B \cup C) = (A \cup B) \cup C$

### 1.2.2 Intersection

- $A \cap B := \{x \in A | x \in B\} = \{x \in B | x \in A\}$
- $\bigcap \mathcal{F} := \{x | x \in X \text{ for all } X \in \mathcal{F}\}$

### 1.2.3 Complement

- *Relative Complement:*  $A \setminus B := \{x \in A | x \notin B\}$
- *Absolute Complement:* For some universe  $U$  and  $A \subseteq U$ ,  $A^c := U \setminus A$

### 1.2.4 Symmetric Difference

$$A \triangle B := (A \setminus B) \cup (B \setminus A)$$

### 1.2.5 Power Set

$$\mathcal{P}(X) := \{S | S \subseteq X\}$$

### 1.2.6 $n$ -Tuple

- *Ordered pair:*  $\langle a, b \rangle := \{\{a\}, \{a, b\}\}$
- $\langle a_1, a_2, a_3, \dots, a_n \rangle := \langle \langle \langle a_1, a_2 \rangle, a_3 \rangle \dots \rangle, a_n \rangle$

### 1.2.7 Cartesian Product

- $A \times B := \{\langle a, b \rangle | \text{for some } a \in A \text{ and for some } b \in B\}$
- $\times \mathcal{F} := \{\langle a_1, a_2, \dots, a_n \rangle | \text{for } a_1 \in A_1, a_2 \in A_2, \dots, a_n \in A_n \text{ where } A_1, A_2, \dots, A_n \in \mathcal{F}\}$

### 1.2.8 Quotient by Equivalence Relation

$X/\sim := \{[a]_\sim | a \in X\}$  (See: equivalence relations)

### 1.2.9 Family

Given a set  $X$  and an index set  $I$ , a family is a function  $\mathcal{F} : I \rightarrow X$ . A cleaner way of denoting the concept is...

$$\mathcal{F}(i) := S_i, \{S_i\}_{i \in I}$$

## 1.3 Relations

$\mathcal{R} : \subseteq A \times B$  for some  $A \times B$

### 1.3.1 Equivalence Relations

Relations  $\sim \subseteq A \times A$  such that  $\forall a, b, c \in A \dots$

- $a \sim a$
- $a \sim b \Rightarrow b \sim a$
- $a \sim b \wedge b \sim c \Rightarrow a \sim c$

*Equivalence Class:*  $[a]_\sim := \{b \in S | b \sim a\}$

*Set Partition:* A set  $P : \subseteq \mathcal{P}(X)$  such that...

- $\bigcup P = X$
- $\forall S_1, S_2 \in P (S_1 \cap S_2 \neq \emptyset \Rightarrow S_1 = S_2)$

### 1.3.2 Functions

A relation  $f : A \rightarrow B$  satisfying  $\forall a \in A \exists! b \in B$  such that  $afb$ , denoted  $f(a) = b$ .

*Injection:* A function  $f : A \hookrightarrow B$  such that  $\forall x, y \in A$  if  $x \neq y$ , then  $f(x) \neq f(y)$ . (See: monomorphism. Injections have right inverses.)

*Surjection:* A function  $f : A \twoheadrightarrow B$  such that  $\forall b \in B \exists a \in A$  such that  $f(a) = b$ . (See: epimorphism. Surjections have left inverses, called *sections*.)

*Bijection:* A function  $f : A \xrightarrow{\sim} B$  which is an injection and a surjection. (See: isomorphism)

*Restriction:* For  $C \subseteq A$  and  $f : A \rightarrow B$ ,  $f|_C : C \rightarrow B$  where  $\forall c \in C f|_C(c) := f(c)$

*Image:*  $f(A) := \{f(a) | a \in A\}$

*Preimage:*  $f^{-1}(A) := \{a \in A \mid f(a) \in B\}$

*Inclusion Map:* For  $A \subseteq X$ ,  $\iota : A \hookrightarrow X$  where  $\forall a \in A \iota(a) := a \in X$

*Function Composition:*  $f : X \rightarrow Y$  and  $g : Y \rightarrow Z \Rightarrow g \circ f : X \rightarrow Z$  where  $\dots$   
 $\forall x \in X, g \circ f(x) := g(f(x))$

*Characteristic Function of a subset:* For  $A \subseteq X$ ,  $\chi_A : X \rightarrow 2$  where  $\dots$

$$\chi_A(x) := \begin{cases} 0 & x \in X \setminus A \\ 1 & x \in A \end{cases}$$

## 2 Category Theory

### 2.1 Metacategories

#### 2.1.1 Undefined notions

- *Objects:*  $a, b, c \dots$
- *Arrows:*  $f, g, h \dots$

#### 2.1.2 Operations

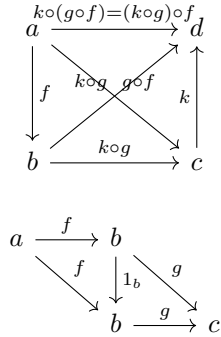
Given  $f : a \rightarrow b \dots$

- *Domain:* **dom:** arrows  $\rightarrow$  objects,  $f \mapsto a$
- *Codomain:* **cod:** arrows  $\rightarrow$  objects,  $f \mapsto b$
- *Identity:* **id:** objects  $\rightarrow$  arrows,  $a \mapsto \text{id}_a = 1_a$
- *Composition:* **comp:** arrows  $\times$  : arrows  $\rightarrow$  arrows,  $\langle g, f \rangle \mapsto g \circ f$ ,  
 $g \circ f : \text{dom} f \rightarrow \text{cod} g$

$$\begin{array}{ccc} & b & \\ f \nearrow & & \searrow g \\ a & \xrightarrow{g \circ f} & c \end{array}$$

#### 2.1.3 Axioms

- *Associativity:*  $a \xrightarrow{f} b \xrightarrow{g} c \xrightarrow{k} d, k \circ (g \circ f) = (k \circ g) \circ f$
- *Unit Law:*  $1_a \circ f = f$  and  $g \circ 1_b = g$



## 2.2 Categories

### 2.2.1 Directed Graph

- $A$  - a set of arrows
- $O$  - a set of objects
- $\mathbf{dom} : A \rightarrow O$ ,  $\mathbf{cod} : A \rightarrow O$

Set of composable pairs of arrows:

$$A \times_O A = \{\langle g, f \rangle | g, f \in A \text{ and } \mathbf{dom}(g) = \mathbf{cod}(f)\}$$

### 2.2.2 Categories

Add the following structure to a directed graph...

- $O \xrightarrow{id} A$ ,  $c \mapsto id_C$
- $A \times_O A \xrightarrow{\circ} A$ ,  $\langle g, f \rangle \mapsto g \circ f$

which satisfy  $\forall a \in O$  and  $\forall \langle g, f \rangle \in A \times_O A$ ...

- $\mathbf{dom}(\mathbf{id}(a)) = a = \mathbf{cod}(\mathbf{id}(a))$
- $\mathbf{dom}(g \circ f) = \mathbf{dom}(f)$
- $\mathbf{cod}(g \circ f) = \mathbf{cod}(g)$
- metacategorical axioms

Small categories use small sets for their objects.

### 2.2.3 Hom Sets

$$\mathbf{hom}(b, c) = \{f | f \in C, \mathbf{dom}(f) = b, \mathbf{cod}(f) = c\}$$

### 2.2.4 Groupoids

A category in which every arrow is an isomorphism.

## 2.3 Morphisms

Arrows in categories.

### 2.3.1 Isomorphisms

A morphism  $f \in \text{hom}(b, c)$  that has a two-sided inverse  $g \in \text{hom}(c, b)$  under composition such that

$$gf = 1_b, \quad fg = 1_c.$$

**Proposition 2.1.** *The inverse of an isomorphism is unique.*

*Proof.* For inverses  $g_1, g_2$  of  $f$  observe...

$$g_1 = g_1 1_c = g_1 (fg_2) = (g_1 f) g_2 = 1_b g_2 = g_2$$

□

**Proposition 2.2.** *Supposing  $f^{-1}$  is the inverse of  $f$ ...*

- *Each identity  $1_c$  is an isomorphism and is its own inverse.*
- *If  $f$  is an isomorphism, then  $f^{-1}$  is an isomorphism and further  $(f^{-1})^{-1} = f$ .*
- *If  $f \in \text{hom}(a, b)$ ,  $g \in \text{hom}(b, c)$  are isomorphisms, then the composition  $gf$  is an isomorphism and  $(gf)^{-1} = f^{-1}g^{-1}$ .*

### 2.3.2 Automorphisms

An isomorphism of an object to itself. Denoted:

$$\text{hom}(c, c) = \text{aut}(c)$$

Observe  $\text{aut}(c)$  is a group.

### 2.3.3 Monomorphisms

A morphism  $f \in \text{hom}(b, c)$  such that  $\forall z \in C$  and  $\forall \alpha', \alpha'' \in \text{hom}(z, b)$ :

$$f \circ \alpha' = f \circ \alpha'' \Rightarrow \alpha' = \alpha''$$

### 2.3.4 Epimorphisms

A morphism  $f \in \text{hom}(b, c)$  such that  $\forall z \in C$  and  $\forall \beta', \beta'' \in \text{hom}(b, z)$ :

$$\beta' \circ f = \beta'' \circ f \Rightarrow \beta' = \beta''$$

## 2.4 Functors

Morphisms  $T : C \rightarrow B$  with domain and codomain both categories. It consists of two suitably related functions

- object function  $T, c \mapsto Tc$
- arrow function  $T, f : c \rightarrow c' \mapsto Tf : Tc \rightarrow Tc'$

which satisfy...

- $T(1_c) = 1_{Tc}$
- $T(g \circ f) = Tg \circ Tf$

### 2.4.1 Full

$\forall c, c' \in C$  and  $g : Tc \rightarrow Tc' \in B, \exists f : c \rightarrow c' \in C$  s.t.  $g \in Tf$

### 2.4.2 Faithful

$\forall c, c' \in C$  and  $f_1, f_2 : c \rightarrow c', Tf_1 = Tf_2 \Rightarrow f_1 = f_2$

## 3 Group Theory