

1 Set Theory

1.1 Set Axioms

1.1.1 Undefined notions

Set: A, B, C, \dots

1.1.2 Axioms

1. *Extension:* $\forall A \forall B [\forall C (C \in A \Leftrightarrow C \in B) \Rightarrow A = B]$
2. *Regularity:* $\forall A [\exists C (C \in A) \Rightarrow \exists B (B \in A \wedge \neg \exists D (D \in B \wedge D \in A))]$
(Every nonempty set contains a set that is disjoint from it. Also know as "Axiom of Foundation.")
3. *Schema of Specification:* $\forall B \forall X_1 \forall X_2 \dots \forall X_n \exists A \forall C [C \in A \Leftrightarrow (C \in B \wedge \phi)]$
4. *Pairing:* $\forall X_1 \forall X_2 \exists A (X_1 \in A \wedge X_2 \in A)$
5. *Union:* $\forall \mathcal{F}_A \exists U \forall A \forall X [(X \in A \wedge A \in \mathcal{F}_A) \Rightarrow X \in U]$
6. *Schema of Replacement:* $\forall A \forall X_1 \forall X_2 \dots \forall X_n [\forall B (B \in A \Rightarrow \exists! D \phi) \Rightarrow \exists B \forall C (C \in A \Rightarrow \exists D (D \in B \wedge \phi))]$
7. *Infinity:* $\exists \omega_0 [\emptyset \in \omega_0 \wedge \forall X (X \in \omega_0 \Rightarrow X \cup X) \in \omega_0]$
8. *Power Set:* $\forall X \exists \mathcal{P}(X) \forall S [S \subseteq X \Rightarrow S \in \mathcal{P}(X)]$
9. *Empty Set:* $\exists A \forall X (X \notin A)$
10. *Choice:* $\forall X [\emptyset \notin X \Rightarrow \exists (f : X \rightarrow \bigcup X) \forall A \in X (f(A) \in A)]$

Proposition 1.1.1. *The empty set axiom is implied by the other nine axioms.*

Proof. Just choose any formula that is always false such as $\phi(X) = X \in B \wedge X \notin B$ and apply the axiom schema of specification. This will give the empty set. The axiom of extension proves uniqueness vacuously. \square

1.1.3 Universe

A set U is defined with the following properties...

1. $x \in u \in U \Rightarrow x \in U$
2. $u \in U \wedge v \in U \Rightarrow \{u, v\}, \langle u, v \rangle, u \times v \in U$
3. $X \in U \Rightarrow \mathcal{P}(X) \in U \wedge \bigcup X \in U$
4. $\omega_0 \in U$ is the set of finite ordinals
5. if $f : A \rightarrow B$ is a surjective function with $A \in U \wedge B \subset U$, then $B \in U$
(See: Set Constructions.)

In category theory, *small sets* are members of U .

1.2 Set Constructions

1.2.1 Union

- $A \cup B := \{x | x \in A \vee x \in B\}$
- $\bigcup \mathcal{F} := \{x | x \in X \text{ for some } X \in \mathcal{F}\}$

Proposition 1.2.1. *For sets A, B, C , the following hold...*

- Identity: $A \cup \emptyset = A$
- Idempotence: $A \cup A = A$
- Absorption: $A \subseteq B \Leftrightarrow A \cup B = B$
- Commutative: $A \cup B = B \cup A$
- Associative: $A \cup (B \cup C) = (A \cup B) \cup C$

1.2.2 Intersection

- $A \cap B := \{x \in A | x \in B\} = \{x \in B | x \in A\}$
- $\bigcap \mathcal{F} := \{x | x \in X \text{ for all } X \in \mathcal{F}\}$

Proposition 1.2.2. *For sets A, B, C , the following hold...*

- Zero: $A \cap \emptyset = \emptyset$
- Idempotence: $A \cap A = A$
- Absorption: $A \subseteq B \Leftrightarrow A \cap B = A$
- Commutative: $A \cap B = B \cap A$
- Associative: $A \cap (B \cap C) = (A \cap B) \cap C$

1.2.3 Complement

- *Relative Complement:* $A \setminus B := \{x \in A | x \notin B\}$
- *Absolute Complement:* For some universe U and $A \subseteq U$, $A^c := U \setminus A$

1.2.4 Symmetric Difference

$$A \triangle B := (A \setminus B) \cup (B \setminus A)$$

1.2.5 Power Set

$$\mathcal{P}(X) := \{S | S \subseteq X\}$$

1.2.6 n -Tuple

- *Ordered pair:* $\langle a, b \rangle := \{\{a\}, \{a, b\}\}$
- $\langle a_1, a_2, a_3, \dots, a_n \rangle := \langle \langle \langle a_1, a_2 \rangle, a_3 \rangle \dots \rangle, a_n \rangle$

1.2.7 Cartesian Product

- $A \times B := \{\langle a, b \rangle \mid \text{for some } a \in A \text{ and for some } b \in B\}$
- $\times \mathcal{F} := \{\langle a_1, a_2, \dots, a_n \rangle \mid \text{for } a_1 \in A_1, a_2 \in A_2, \dots, a_n \in A_n \text{ where } A_1, A_2, \dots, A_n \in \mathcal{F}\}$

1.2.8 Quotient by Equivalence Relation

$X / \sim := \{[a]_{\sim} \mid a \in X\}$ (See: equivalence relations)

1.2.9 Family

Given a set X and an index set I , a family is a function $\mathcal{F} : I \rightarrow X$. A cleaner way of denoting the concept is...

$$\mathcal{F}(i) := S_i, \{S_i\}_{i \in I}$$

1.3 Relations

$\mathcal{R} \subseteq A \times B$ for some $A \times B$

1.3.1 Equivalence Relations

Relations $\sim \subseteq A \times A$ such that $\forall a, b, c \in A \dots$

- $a \sim a$
- $a \sim b \Rightarrow b \sim a$
- $a \sim b \wedge b \sim c \Rightarrow a \sim c$

Equivalence Class: $[a]_{\sim} := \{b \in S \mid b \sim a\}$

Set Partition: A set $P \subseteq \mathcal{P}(X)$ such that...

- $\bigcup P = X$
- $\forall S_1, S_2 \in P (S_1 \cap S_2 \neq \emptyset \Rightarrow S_1 = S_2)$

1.3.2 Functions

A relation $f : A \rightarrow B$ satisfying $\forall a \in A \exists! b \in B$ such that afb , denoted $f(a) = b$.

Injection: A function $f : A \hookrightarrow B$ such that $\forall x, y \in A$ if $x \neq y$, then $f(x) \neq f(y)$.
(See: monomorphism. Injections have right inverses.)

Surjection: A function $f : A \twoheadrightarrow B$ such that $\forall b \in B \exists a \in A$ such that $f(a) = b$.
(See: epimorphism. Surjections have left inverses, called *sections*.)

Bijection: A function $f : A \xrightarrow{\sim} B$ which is an injection and a surjection. (See: isomorphism)

Restriction: For $C \subseteq A$ and $f : A \rightarrow B$, $f|_C : C \rightarrow B$ where $\forall c \in C f|_C(c) := f(c)$

Image: $f(A) := \{f(a) | a \in A\}$

Preimage: $f^{-1}(A) := \{a \in A | f(a) \in B\}$

Inclusion Map: For $A \subseteq X$, $\iota : A \hookrightarrow X$ where $\forall a \in A \iota(a) := a \in X$

Function Composition: $f : X \rightarrow Y$ and $g : Y \rightarrow Z \Rightarrow g \circ f : X \rightarrow Z$ where \dots
 $\forall x \in X, g \circ f(x) := g(f(x))$

Characteristic Function of a subset: For $A \subseteq X$, $\chi_A : X \rightarrow 2$ where \dots

$$\chi_A(x) := \begin{cases} 0 & x \in X \setminus A \\ 1 & x \in A \end{cases}$$

2 Category Theory

2.1 Metacategories

2.1.1 Undefined notions

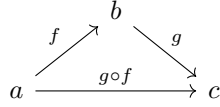
- *Objects:* $a, b, c \dots$
- *Arrows:* $f, g, h \dots$

2.1.2 Operations

Given $f : a \rightarrow b \dots$

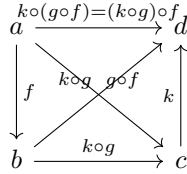
- *Domain:* **dom:** arrows \rightarrow objects, $f \mapsto a$
- *Codomain:* **cod:** arrows \rightarrow objects, $f \mapsto b$

- *Identity*: **id**: objects \rightarrow arrows, $a \mapsto \text{id}_a = 1_a$
- *Composition*: **comp**: arrows \times : arrows \rightarrow arrows, $\langle g, f \rangle \mapsto g \circ f$,
 $g \circ f : \text{dom} f \rightarrow \text{cod} g$

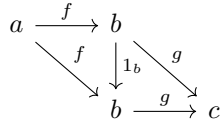


2.1.3 Axioms

- *Associativity*: $a \xrightarrow{f} b \xrightarrow{g} c \xrightarrow{k} d$, $k \circ (g \circ f) = (k \circ g) \circ f$



- *Unit Law*: $1_a \circ f = f$ and $g \circ 1_b = g$



2.2 Categories

2.2.1 Directed Graph

- A - a set of arrows
- O - a set of objects
- **dom** : $A \rightarrow O$, **cod** : $A \rightarrow O$

Set of composable pairs of arrows:

$$A \times_O A = \{\langle g, f \rangle | g, f \in A \text{ and } \mathbf{dom}(g) = \mathbf{cod}(f)\}$$

2.2.2 Categories

Add the following structure to a directed graph. . .

- $O \xrightarrow{id} A, c \mapsto id_C$
- $A \times_O A \xrightarrow{\circ} A, \langle g, f \rangle \mapsto g \circ f$

which satisfy $\forall a \in O$ and $\forall \langle g, f \rangle \in A \times_O A$. . .

- $\mathbf{dom}(\mathbf{id}(a)) = a = \mathbf{cod}(\mathbf{id}(a))$
- $\mathbf{dom}(g \circ f) = \mathbf{dom}(f)$
- $\mathbf{cod}(g \circ f) = \mathbf{cod}(g)$
- metacategorical axioms

Small categories use small sets for their objects.

2.2.3 Hom Sets

$$\mathbf{hom}(b, c) = \{f | f \in C, \mathbf{dom}(f) = b, \mathbf{cod}(f) = c\}$$

2.2.4 Groupoids

A category in which every arrow is an isomorphism.

2.3 Morphisms

Arrows in categories.

2.3.1 Isomorphisms

A morphism $f \in \mathbf{hom}(b, c)$ that has a two-sided inverse $g \in \mathbf{hom}(c, b)$ under composition such that

$$gf = 1_b, fg = 1_c.$$

Proposition 2.3.1. *The inverse of an isomorphism is unique.*

Proof. For inverses g_1, g_2 of f observe. . .

$$g_1 = g_1 1_c = g_1 (f g_2) = (g_1 f) g_2 = 1_b g_2 = g_2$$

□

Proposition 2.3.2. *Supposing f^{-1} is the inverse of f . . .*

- Each identity 1_c is an isomorphism and is its own inverse.
- If f is an isomorphism, then f^{-1} is an isomorphism and further $(f^{-1})^{-1} = f$.
- If $f \in \mathbf{hom}(a, b)$, $g \in \mathbf{hom}(b, c)$ are isomorphisms, then the composition gf is an isomorphism and $(gf)^{-1} = f^{-1}g^{-1}$.

2.3.2 Automorphisms

An isomorphism of an object to itself. Denoted:

$$\text{hom}(c, c) = \text{aut}(c)$$

Observe $\text{aut}(c)$ is a group.

2.3.3 Monomorphisms

A morphism $f \in \text{hom}(b, c)$ such that $\forall z \in C$ and $\forall \alpha', \alpha'' \in \text{hom}(z, b)$:

$$f \circ \alpha' = f \circ \alpha'' \Rightarrow \alpha' = \alpha''$$

2.3.4 Epimorphisms

A morphism $f \in \text{hom}(b, c)$ such that $\forall z \in C$ and $\forall \beta', \beta'' \in \text{hom}(b, z)$:

$$\beta' \circ f = \beta'' \circ f \Rightarrow \beta' = \beta''$$

2.4 Functors

Morphisms $T : C \rightarrow B$ with domain and codomain both categories. It consists of two suitably related functions

- object function $T, c \mapsto Tc$
- arrow function $T, f : c \rightarrow c' \mapsto Tf : Tc \rightarrow Tc'$

which satisfy...

- $T(1_c) = 1_{Tc}$
- $T(g \circ f) = Tg \circ Tf$

2.4.1 Full

$\forall c, c' \in C$ and $g : Tc \rightarrow Tc' \in B, \exists f : c \rightarrow c' \in C$ s.t. $g \in Tf$

2.4.2 Faithful

$\forall c, c' \in C$ and $f_1, f_2 : c \rightarrow c', Tf_1 = Tf_2 \Rightarrow f_1 = f_2$

3 Group Theory