

# 1 Set Theory

## 1.1 Set Axioms

### 1.1.1 Undefined notions

Set:  $A, B, C, \dots$

### 1.1.2 Axioms

1. *Extension:*  $\forall A \forall B [\forall C (C \in A \Leftrightarrow C \in B) \Rightarrow A = B]$
2. *Regularity:*  $\forall A [\exists C (C \in A) \Rightarrow \exists B (B \in A \wedge \neg \exists D (D \in B \wedge D \in A))]$   
(Every nonempty set contains a set that is disjoint from it. Also known as "Axiom of Foundation.")
3. *Schema of Specification:*  $\forall B \forall X_1 \forall X_2 \dots \forall X_n \exists A \forall C [C \in A \Leftrightarrow (C \in B \wedge \phi)]$
4. *Pairing:*  $\forall X_1 \forall X_2 \exists A (X_1 \in A \wedge X_2 \in A)$
5. *Union:*  $\forall \mathcal{F}_A \exists U \forall A \forall X [(X \in A \wedge A \in \mathcal{F}_A) \Rightarrow X \in U]$
6. *Schema of Replacement:*  $\forall A \forall X_1 \forall X_2 \dots \forall X_n [\forall B (B \in A \Rightarrow \exists! D \phi) \Rightarrow \exists B \forall C (C \in A \Rightarrow \exists D (D \in B \wedge \phi))]$
7. *Infinity:*  $\exists \omega [\emptyset \in \omega \wedge \forall X (X \in \omega \Rightarrow X \cup X \in \omega)]$
8. *Power Set:*  $\forall X \exists \mathcal{P}(X) \forall S [S \subseteq X \Rightarrow S \in \mathcal{P}(X)]$
9. *Empty Set:*  $\exists A \forall X (X \notin A)$
10. *Choice:*  $\forall X [\emptyset \notin X \Rightarrow \exists (f : X \rightarrow \bigcup X) \forall A \in X (f(A) \in A)]$

**Proposition 1.1.1.** *The empty set axiom is implied by the other nine axioms.*

*Proof.* Just choose any formula that is always false such as  $\phi(X) = X \in B \wedge X \notin B$  and apply the axiom schema of specification. This will give the empty set. The axiom of extension proves uniqueness vacuously.  $\square$

### 1.1.3 Universe

A set  $U$  is defined with the following properties...

1.  $x \in u \in U \Rightarrow x \in U$
2.  $u \in U \wedge v \in U \Rightarrow \{u, v\}, \langle u, v \rangle, u \times v \in U$
3.  $X \in U \Rightarrow \mathcal{P}(X) \in U \wedge \bigcup X \in U$
4.  $\omega \in U$  is the set of finite ordinals
5. if  $f : A \rightarrow B$  is a surjective function with  $A \in U \wedge B \subset U$ , then  $B \in U$   
(See: Set Constructions.)

In category theory, *small sets* are members of  $U$ .

## 1.2 Set Constructions

### 1.2.1 Union

- $A \cup B := \{x | x \in A \vee x \in B\}$
- $\bigcup \mathcal{F} := \{x | x \in X \text{ for some } X \in \mathcal{F}\}$

**Proposition 1.2.1.** *For sets  $A, B, C$ , the following hold...*

- Identity:  $A \cup \emptyset = A$
- Idempotence:  $A \cup A = A$
- Absorption:  $A \subseteq B \Leftrightarrow A \cup B = B$
- Commutative:  $A \cup B = B \cup A$
- Associative:  $A \cup (B \cup C) = (A \cup B) \cup C$

### 1.2.2 Intersection

- $A \cap B := \{x \in A | x \in B\} = \{x \in B | x \in A\}$
- $\bigcap \mathcal{F} := \{x | x \in X \text{ for all } X \in \mathcal{F}\}$

**Proposition 1.2.2.** *For sets  $A, B, C$ , the following hold...*

- Zero:  $A \cap \emptyset = \emptyset$
- Idempotence:  $A \cap A = A$
- Absorption:  $A \subseteq B \Leftrightarrow A \cap B = A$
- Commutative:  $A \cap B = B \cap A$
- Associative:  $A \cap (B \cap C) = (A \cap B) \cap C$

### 1.2.3 Complement

- *Relative Complement:*  $A \setminus B := \{x \in A | x \notin B\}$
- *Absolute Complement:* For some universe  $U$  and  $A \subseteq U$ ,  $A^c := U \setminus A$

**Proposition 1.2.3.** *For a universe  $U$  and sets  $A, B \subseteq U$ ...*

- $(A^c)^c = A$
- $\emptyset^c = U$
- $U^c = \emptyset$
- $A \cap A^c = \emptyset$

- $A \cup A^c = U$
- $A \subseteq B \Leftrightarrow B^c \subseteq A^c$

**Proposition 1.2.4** (DeMorgan's Laws). *For a universe  $U$  and sets  $A, B \subseteq U$ ...*

- $(A \cup B)^c = A^c \cap B^c$
- $(A \cap B)^c = A^c \cup B^c$

**Proposition 1.2.5.** *For sets  $A, B$ ...*

- $A \setminus B = A \cap B^c$
- $A \subseteq B \Leftrightarrow A \setminus B = \emptyset$
- $A \setminus (A \setminus B) = A \cap B$
- $A \cap (B \setminus C) = (A \cap B) \setminus (A \cap C)$
- $A \cap B \subseteq (A \cap C) \cup (B \cap C^c)$
- $(A \cup C) \cap (B \cup C^c) \subseteq A \cup B$

**Proposition 1.2.6.** *For a family  $\mathcal{F}$ ...*

- $\forall X \in \mathcal{F}, \bigcup_{k \in K} X_k = \bigcup_{j \in J} (\bigcup_{i \in I_j} X_i)$
- $\forall X \in \mathcal{F}, \bigcap_{k \in K} X_k = \bigcap_{j \in J} (\bigcap_{i \in I_j} X_i)$
- $\forall X \in \mathcal{F}, \bigcup_{i \in I} X_i = \bigcup_{j \in J} X_j$
- $\forall X \in \mathcal{F}, \bigcap_{i \in I} X_i = \bigcap_{j \in J} X_j$
- $(\bigcup_{i \in I} A_i) \cap (\bigcup_{j \in J} B_j) = \bigcup_{i,j} (A_i \cap B_j)$
- $(\bigcap_{i \in I} A_i) \cup (\bigcap_{j \in J} B_j) = \bigcap_{i,j} (A_i \cup B_j)$

**Proposition 1.2.7** (Generalized DeMorgan's Laws). *For a universe  $U$  and a family  $\mathcal{F}$ ...*

- $(\bigcup_{X \in \mathcal{F}} X)^c = \bigcap_{X \in \mathcal{F}} X^c$
- $(\bigcap_{X \in \mathcal{F}} X)^c = \bigcup_{X \in \mathcal{F}} X^c$

#### 1.2.4 Symmetric Difference

$$A \triangle B := (A \setminus B) \cup (B \setminus A)$$

### 1.2.5 Power Set

$$\mathcal{P}(X) := \{S \mid S \subseteq X\}$$

**Proposition 1.2.8.** *For sets  $A, B$  and a family  $\mathcal{F} \dots$*

- $\mathcal{P}(A) \cap \mathcal{P}(B) = \mathcal{P}(A \cap B)$
- $\mathcal{P}(A) \cup \mathcal{P}(B) \subseteq \mathcal{P}(A \cup B)$
- $\bigcap_{X \in \mathcal{F}} \mathcal{P}(X) = \mathcal{P}(\bigcap_{X \in \mathcal{F}} X)$
- $\bigcup_{X \in \mathcal{F}} \mathcal{P}(X) \subseteq \mathcal{P}(\bigcup_{X \in \mathcal{F}} X)$

#### 1.2.5.1 Characteristic Function of a subset

For  $A \subseteq X$ ,  $\chi_A : X \rightarrow 2$  where...

$$\chi_A(x) := \begin{cases} 0 & x \in X \setminus A \\ 1 & x \in A \end{cases}$$

### 1.2.6 $n$ -Tuple

- *Ordered pair:*  $\langle a, b \rangle := \{\{a\}, \{a, b\}\}$
- $\langle a_1, a_2, a_3, \dots, a_n \rangle := \langle \langle \langle a_1, a_2 \rangle, a_3 \rangle \dots \rangle, a_n \rangle$

### 1.2.7 Cartesian Product

- $A \times B := \{\langle a, b \rangle \mid \text{for some } a \in A \text{ and for some } b \in B\}$
- $\times \mathcal{F} := \{\langle a_1, a_2, \dots, a_n \rangle \mid \text{for } a_1 \in A_1, a_2 \in A_2, \dots, a_n \in A_n \text{ where } A_1, A_2, \dots, A_n \in \mathcal{F}\}$

**Proposition 1.2.9.** *For sets  $A, B \dots$*

- $(A \cup B) \times X = (A \times X) \cup (B \times X)$
- $(A \cap B) \times (X \cap Y) = (A \times X) \cap (B \times X)$
- $(A \setminus B) \times X = (A \times X) \setminus (B \times X)$

**Proposition 1.2.10.** *For families  $\{A_i\}_{i \in I}, \{B_j\}_{j \in J}, \{X_i\}_{i \in I} \dots$*

- $(\bigcup_{i \in I} A_i) \times (\bigcup_{j \in J} B_j) = \bigcup_{i,j} (A_i \times B_j)$
- $(\bigcap_{i \in I} A_i) \times (\bigcap_{j \in J} B_j) = \bigcap_{i,j} (A_i \times B_j)$
- $\bigcap_i X_i \subseteq X_j \subseteq \bigcup_i X_i$

### 1.2.8 Quotient by Equivalence Relation

$X / \sim := \{[a]_{\sim} \mid a \in X\}$  (See: equivalence relations)

### 1.2.9 Family

Given a set  $X$  and an index set  $I$ , a family is a function  $\mathcal{F} : I \rightarrow X$ . A cleaner way of denoting the concept is...

$$\mathcal{F}(i) := S_i, \{S_i\}_{i \in I}$$

## 1.3 Relations

$\mathcal{R} : \subseteq A \times B$  for some  $A \times B$

### 1.3.1 Equivalence Relations

Relations  $\sim \subseteq A \times A$  such that  $\forall a, b, c \in A \dots$

- *Reflexive:*  $a \sim a$
- *Symmetric:*  $a \sim b \Rightarrow b \sim a$
- *Transitive:*  $a \sim b \wedge b \sim c \Rightarrow a \sim c$

#### 1.3.1.1 Equivalence Class

$$[a]_{\sim} := \{b \in S \mid b \sim a\}$$

#### 1.3.1.2 Set of Equivalence Classes

$$[A] = \{[a]_{\sim} \mid a \in A\}$$

#### 1.3.1.3 Set Partition

A set  $P : \subseteq \mathcal{P}(X)$  such that...

- $\bigcup P = X$
- $\forall S_1, S_2 \in P (S_1 \cap S_2 \neq \emptyset \Rightarrow S_1 = S_2)$

**Proposition 1.3.1.** *Let  $A$  be a set and  $\sim$  an equivalence relation on  $A$ . Then  $[A]$  is a partition of  $A$ .*

**Proposition 1.3.2.** *Let  $A$  be a set and  $P$  be a partition of  $A$ . Define a relation  $x \sim y$  if and only if  $x, y \in C \in P$ . Then  $\sim$  is an equivalence relation.*

#### 1.3.1.4 Congruence Relation

A congruence  $\sim$  of a set  $A$  with a binary operation  $\mu : A \times A \rightarrow A$  is an equivalence relation such that...

$$\bar{\mu}([a], [b]) = [\mu(a, b)]$$

induces a well-defined binary operation on  $[A]$ .

**Proposition 1.3.3.** *An equivalence relation  $\sim$  on  $A$  with  $\mu : A \times A \rightarrow A$  is a congruence relation if for any  $a, a', b, b' \in A$ , whenever  $[a] = [a']$  and  $[b] = [b']$ , we have  $[\mu(a, b)] = [\mu(a', b')]$ .*

### 1.3.2 Functions

A relation  $f : A \rightarrow B$  satisfying  $\forall a \in A \exists! b \in B$  such that  $afb$ , denoted  $f(a) = b$ .

#### 1.3.2.1 Injection

A function  $f : A \hookrightarrow B$  such that  $\forall x, y \in A$  if  $x \neq y$ , then  $f(x) \neq f(y)$ . (See: monomorphism. Injections have right inverses.)

#### 1.3.2.2 Surjection

A function  $f : A \twoheadrightarrow B$  such that  $\forall b \in B \exists a \in A$  such that  $f(a) = b$ . (See: epimorphism, Stirling numbers of the second kind. Surjections have left inverses, called *sections*.)

#### 1.3.2.3 Bijection

A function  $f : A \xrightarrow{\sim} B$  which is an injection and a surjection. (See: isomorphism)

#### 1.3.2.4 Restriction

For  $C \subseteq A$  and  $f : A \rightarrow B$ ,  $f|_C : C \rightarrow B$  where  $\forall c \in C f|_C(c) := f(c)$

#### 1.3.2.5 Image

$$f(A) := \{f(a) | a \in A\}$$

**Proposition 1.3.4.** *For a function  $f : A \rightarrow B$  and a family  $\{X_i\}_{i \in I}$  where  $\forall i \in I X_i \subseteq A$ ...*

- $f(\bigcup_i X_i) = \bigcup_i f(X_i)$
- In general,  $f(\bigcap_i X_i) \neq \bigcap_i f(X_i)$
- In general,  $f(X)^c \neq f(X^c)$

#### 1.3.2.6 Preimage

$$f^{-1}(A) := \{a \in A | f(a) \in B\}$$

**Proposition 1.3.5.** *Given a function  $f : X \rightarrow Y$ ,  $f$  is surjective if and only if  $\forall A \subseteq Y$ , where  $A \neq \emptyset$ ,  $f^{-1}(A) \neq \emptyset$ .*

**Proposition 1.3.6.** *Given a function  $f : X \rightarrow Y$ ,  $f$  is injective if and only if  $\forall A \subseteq \text{ran } f$ , where  $A$  is a singleton,  $f^{-1}(A)$  is a singleton.*

**Proposition 1.3.7.** *Given a function  $f : X \rightarrow Y \dots$*

- *If  $B \subseteq Y$ , then  $f(f^{-1}(B)) \subseteq B$ .*
- *If  $f$  is surjective, then  $f(f^{-1}(B)) = B$ .*
- *If  $A \subseteq X$ , then  $A \subseteq f^{-1}(f(A))$ .*
- *If  $f$  is injective, then  $A = f^{-1}(f(A))$ .*
- *If  $\{B_i\}$  is a family of subset of  $Y$ , then  $f^{-1}(\bigcup_i B_i) = \bigcup_i f^{-1}(B_i)$  and  $f^{-1}(\bigcap_i B_i) = \bigcap_i f^{-1}(B_i)$ .*

### 1.3.2.7 Function Composition

$f : X \rightarrow Y$  and  $g : Y \rightarrow Z \Rightarrow g \circ f : X \rightarrow Z$  where  $\forall x \in X, g \circ f(x) := g(f(x))$

## 1.4 Natural Numbers

### 1.4.1 Successor

For a set  $n$ , its *successor*  $n^+$  is defined by...

$$n^+ = n \cup \{n\}$$

### 1.4.2 Inductive

A set  $N$  is *inductive* if and only if  $\emptyset \in N$  and  $(\forall n \in N) n^+ \in N$ .

The Axiom of Infinity may be restated in terms of "inductiveness," i.e....

*There exists an inductive set  $\omega$ .*

### 1.4.3 Natural Number

A *natural number* is a set that belongs to every inductive set, i.e. the intersection of them all.

The following theorem is a consequence of the definition...

**Theorem 1.4.1** (Induction on  $\omega$ ). *Any inductive subset of  $\omega$  coincides with  $\omega$ .*

**Proposition 1.4.1.** *Every natural number except 0 is the successor of some natural number.*

*Proof.* Let  $T = \{n \in \omega \mid n = 0 \vee (\exists p \in \omega) n = p^+\}$  and use induction. □

#### 1.4.4 Peano's Postulates

##### 1.4.4.1 Peano System

An ordered triple  $\langle N, S, e \rangle$  consisting of a set  $N$ , a function  $S : N \rightarrow N$ , and a member  $e \in N$  such that the following three conditions are met:

1.  $e \notin \text{ran} S$ .
2.  $S$  is injective.
3. Any subset  $A \subseteq N$  that contains  $e$  and is closed under  $S$  equals  $N$  itself.

**Proposition 1.4.2.** *Let  $\sigma = \{\langle n, n^+ \rangle \mid n \in \omega\}$ . Then  $\langle \omega, \sigma, 0 \rangle$  is a Peano system.*

##### 1.4.4.2 Transitive Set

A set  $A$  is said to be a *transitive set* if and only if  $x \in a \in A \Rightarrow x \in A$ .

**Proposition 1.4.3.** *For a transitive set  $a$ ,*

$$\bigcup (a^+) = a.$$

**Proposition 1.4.4.** *Every natural number is a transitive set and  $\omega$  is a transitive set.*

*Proof.* Use induction. □

#### 1.4.5 Recursion

**Theorem 1.4.2** (Recursion Theorem on  $\omega$ ). *Let  $A$  be a set,  $a \in A$ , and  $F : A \rightarrow A$ . Then there exists a unique function  $h : \omega \rightarrow A$  such that...*

$$h(0) = a,$$

and for every  $n \in \omega$ ,

$$h(n^+) = F(h(n)).$$

*Proof.* The idea is to let  $h$  be the union of many approximating functions. For the purposes of this proof, call a function  $v$  *acceptable* if and only if  $\text{dom } v \subseteq \omega$ ,  $\text{ran } v \subseteq A$ , and the following conditions hold:

1. If  $0 \in \text{dom } v$ , then  $v(0) = a$ .
2. If  $n^+ \in \text{dom } v$  (where  $n \in \omega$ ), then also  $n \in \text{dom } v$  and  $v(n^+) = F(v(n))$ .

Let  $\mathcal{H}$  be the collection of all acceptable functions, and let  $h = \bigcup \mathcal{H}$ . Thus...

$$\begin{aligned} (\star) \quad \langle n, y \rangle \in h &\Leftrightarrow \langle n, y \rangle \text{ is a member of some acceptable } v \\ &\Leftrightarrow v(n) = y \text{ for some acceptable } v. \end{aligned}$$



We claim that this  $h$  meets the demands of the theorem. This claim can be broken down into four parts. The four parts involve showing that (I)  $h$  is a function, (II)  $h$  is acceptable, (III)  $\text{dom } h$  is all of  $\omega$ , and (IV)  $h$  is unique.

I. We first claim that  $h$  is a function. Let. . .

$$S = \{n \in \omega \mid \text{for at most one } y, \langle n, y \rangle \in h\}.$$

We must check that  $S$  is inductive. If  $\langle 0, y_1 \rangle \in h$  and  $\langle 0, y_2 \rangle \in h$ , then by  $(\star)$  there exist acceptable  $v_1$  and  $v_2$  such that  $v_1(0) = y_1$  and  $v_2(0) = y_2$ . But by (1) it follows that  $y_1 = a = y_2$ . Thus  $0 \in S$ .

Next suppose that  $k \in S$ . Consider  $\langle k^+, y_1 \rangle \in h$  and  $\langle k^+, y_2 \rangle \in h$ . As before there must exist acceptable  $v_1$  and  $v_2$  such that  $v_1(k^+) = y_1$  and  $v_2(k^+) = y_2$ . By condition (2) it follows that. . .

$$y_1 = v_1(k^+) = F(v_1(k)) \quad \text{and} \quad y_2 = v_2(k^+) = F(v_2(k)).$$

But since  $k \in S$ , we have  $v_1(k) = v_2(k)$ . Therefore. . .

$$y_1 = F(v_1(k)) = F(v_2(k)) = y_2.$$

So  $k^+ \in S$ , proving  $S$  is inductive and coincides with  $\omega$ . Consequently  $h$  is a function.

II. Next we claim that  $h$  itself is acceptable. We have just seen that  $h$  is a function, and it is clear from  $(\star)$  that  $\text{dom } h \subseteq \omega$  and  $\text{ran } h \subseteq A$ .

First examine (1). If  $0 \in \text{dom } h$ , then there must be some acceptable  $v$  with  $v(0) = h(0)$ . Since  $v(0) = a$ , we have  $h(0) = a$ .

Next examine (2). Assume  $n^+ \in \text{dom } h$ . Again there must be some acceptable  $v$  with  $v(n^+) = h(n^+)$ . Since  $v$  is acceptable we have  $n \in \text{dom } v$  (and  $v(n) = h(n)$ ) and

$$h(n^+) = v(n^+) = F(v(n)) = F(h(n)).$$

Thus  $h$  satisfies (2) and so is acceptable.

III. We now claim that  $\text{dom } h = \omega$  (the function is nonempty). It suffices to show that  $\text{dom } h$  is inductive. The function  $\{\langle 0, a \rangle\}$  is acceptable and hence  $0 \in \text{dom } h$ . Suppose the  $k \in \text{dom } h$ . If  $k^+ \notin \text{dom } h$ , then let. . .

$$v = h \cup \{\langle k^+, F(h(k)) \rangle\}.$$

Then  $v$  is a function,  $\text{dom } v \subseteq \omega$ , and  $\text{ran } v \subseteq A$ . We will show that  $v$  is acceptable.

Condition (1) holds since  $v(0) = h(0) = a$ . For condition (2) there are two cases. If  $n^+ \in \text{dom } v$  where  $n^+ \neq k^+$ , then  $n^+ \in \text{dom } h$  and  $v(n^+) = h(n^+) = F(h(n)) = F(v(n))$ . The other case occurs if  $n^+ = k^+$ . Since the successor operation is injective,  $n = k$ . By assumption  $k \in \text{dom } h$ . Thus. . .

$$v(k^+) = F(h(k)) = F(v(k))$$

and (2) holds. Hence  $v$  is acceptable. But then  $v \subseteq h$ , so that  $k^+ \in \text{dom } h$  after all. So  $\text{dom } h$  is inductive and therefore coincides with  $\omega$ .

IV. Finally we claim that  $h$  is unique. For let  $h_1$  and  $h_2$  both satisfy the conclusion of the theorem. Let...

$$S = \{n \in \omega \mid h_1(n) = h_2(n)\}.$$

$S$  is inductive, showing  $h_1 = h_2$ . Thus  $h$  is unique.  $\square$

**Example 1.4.2.1.** *There is no function  $h : \mathbb{Z} \rightarrow \mathbb{Z}$  such that for every  $a \in \mathbb{Z}$ ,*

$$h(a+1) = h(a)^2 + 1.$$

*Proof.* Note  $h(a) > h(a-1) > h(a-2) > \dots > 0$ . Recursion on  $\omega$  relies on there being a starting point 0.  $\mathbb{Z}$  has no analogous starting point.  $\square$

**Theorem 1.4.3.** *Let  $\langle N, S, e \rangle$  be a Peano system. Then  $\langle \omega, \sigma, 0 \rangle$  is isomorphic to  $\langle N, S, e \rangle$ , i.e. there is a function  $h$  mapping  $\omega$  bijectively to  $N$  in a way that preserves the successor operation*

$$h(\sigma(n)) = S(h(n))$$

*and the zero element*

$$h(0) = e.$$

## 1.4.6 Arithmetic

### 1.4.6.1 Addition

*Addition*  $(+)$  is the binary operation on  $\omega$  such that for any  $m$  and  $n \in \omega$ ,

$$m + n = A_m(n),$$

where  $A_m : \omega \rightarrow \omega$  is the unique function given by the recursion theorem for which...

- $A_m(0) = m$
- $A_m(n^+) = A_m(n)^+ \forall n \in \omega$ .

**Proposition 1.4.5.** *For natural numbers  $m$  and  $n$ ,*

- $m + 0 = m$ ,
- $m + n^+ = (m + n)^+$

### 1.4.6.2 Multiplication

*Multiplication*  $(\cdot)$  is the binary operation on  $\omega$  such that for any  $m$  and  $n \in \omega$ ,

$$m \cdot n = M_m(n),$$

where  $M_m : \omega \rightarrow \omega$  is the unique function given by the recursion theorem for which...

- $M_m(0) = 0$
- $M_m(n^+) = M_m(n) + m$ .

**Proposition 1.4.6.** *For natural numbers  $m$  and  $n$ ,*

- $m \cdot 0 = 0$ ,
- $m \cdot n^+ = m \cdot n + m$

### 1.4.6.3 Exponentiation

*Exponentiation* is the binary operation on  $\omega$  such that for any  $m$  and  $n \in \omega$ ,

$$m^n = E_m(n),$$

where  $E_m : \omega \rightarrow \omega$  is the unique function given by the recursion theorem for which...

- $E_m(0) = 1$
- $E_m(n^+) = E_m(n) \cdot m$ .

**Proposition 1.4.7.** *For natural numbers  $m$  and  $n$ ,*

- $m^0 = 1$ ,
- $m^{(n^+)} = m^n \cdot m$ .

### 1.4.7 Ordering on the natural numbers

Define  $m < n$  if and only if  $m \in n$ .

**Lemma 1.4.4.** *For any natural numbers  $m$  and  $n$ ...*

- $m \in n \Leftrightarrow m^+ \in n^+$ .
- $n \notin n$

**Theorem 1.4.5** (Trichotomy Law for  $\omega$ ). *For any natural numbers  $m$  and  $n$ , exactly one of the three conditions...*

- $m \in n$
- $m = n$

- $n \in m$

holds.

**Corollary 1.4.5.1.** *For any natural numbers  $m$  and  $n$ ,*

- $m \in n \Leftrightarrow m \subset n$
- $(m \in n) \vee (m = n) \Leftrightarrow m \subseteq n$

**Proposition 1.4.8.** *For any natural numbers  $m, n$  and  $p, \dots$*

- $m \in n \Leftrightarrow m + p \in n + p$ .
- *If, in addition,  $p \neq 0$ , then  $m \in n \Leftrightarrow m \cdot p \in n \cdot p$ .*

**Corollary 1.4.5.2.** *The following cancellation laws hold for  $m, n, p \in \omega \dots$*

- $m + p \in n + p \Rightarrow m = n$
- *If, in addition,  $p \neq 0$ , then  $m \cdot p \in n \cdot p \Rightarrow m = n$*

**Theorem 1.4.6** (Well Ordering of  $\omega$ ). *Let  $A$  be a nonempty set of  $\omega$ . Then there is some  $m \in A$  such that  $(m \in n) \vee (m = n)$  for all  $n \in A$ .*

*Proof.* Assume that  $A$  is a subset of  $\omega$  without a least element; we will show that  $A = \emptyset$ . We could attempt to do this by showing that the complement  $\omega \setminus A$  is inductive. But in order to show that  $k^+ \in \omega - A$ , it is not enough to know merely that  $k \in \omega \setminus A$ , we must know that all numbers smaller than  $k$  are in  $\omega \setminus A$  as well. Given this additional information, we can argue that  $k^+ \in \omega \setminus A$  lest it be a least element of  $A$ .

To write down what is approximately this argument, let...

$$B = \{m \in \omega \mid \text{no number less than } m \text{ belongs to } A\}.$$

We claim that  $B$  is inductive.  $0 \in B$  vacuously. Suppose that  $k \in B$ . Then if  $n$  is less than  $k^+$ , either  $n$  is less than  $k$  (in which case  $n \notin A$  since  $k \in B$ ) or  $n = k$  (in which case  $n \notin A$  lest, by trichotomy, it be least in  $A$ ). In either case,  $n$  is outside of  $A$ . Hence  $k^+ \in B$  and  $B$  is inductive. It clearly follows that  $A = \emptyset$ .  $\square$

**Corollary 1.4.6.1.** *There is no function  $f : \omega \rightarrow \omega$  such that  $f(n^+) \in f(n)$  for every natural number  $n$ .*

**Theorem 1.4.7** (Strong Induction Principle for  $\omega$ ). *Let  $A$  be a subset of  $\omega$ , and assume that for every  $n \in \omega$ , if every number less than  $n$  is in  $A$ , then  $n \in A$ . Then  $A = \omega$ .*

## 1.5 Constructing Number Systems

For the purposes of this subsection let  $\mathbb{N} := \omega$ .

### 1.5.1 The Integers

Let  $\sim_{\mathbb{Z}}$  be the equivalence relation on  $\mathbb{N} \times \mathbb{N}$  for which...

$$\langle m, n \rangle \Leftrightarrow m + q = p + n.$$

Then the set of *Integers*, denoted  $\mathbb{Z}$ , is the set  $\mathbb{N} \times \mathbb{N} / \sim_{\mathbb{Z}}$ .

#### 1.5.1.1 Addition

Addition of integers  $a = \langle m, n \rangle$  and  $b = \langle p, q \rangle$  is defined as...

$$a +_{\mathbb{Z}} b = [\langle m + p, n + q \rangle]$$

**Lemma 1.5.1.** *Addition of integers ( $+_{\mathbb{Z}}$ ) is well defined, i.e. if  $\langle m, n \rangle \sim_{\mathbb{Z}} \langle m', n' \rangle$  and  $\langle p, q \rangle \sim_{\mathbb{Z}} \langle p', q' \rangle$ , then...*

$$\langle m + p, n + q \rangle \sim_{\mathbb{Z}} \langle m' + p', n' + q' \rangle$$

The integers under addition form an abelian group.

#### 1.5.1.2 Multiplication

Multiplication of integers  $a = \langle m, n \rangle$  and  $b = \langle p, q \rangle$  is defined as...

$$a \cdot_{\mathbb{Z}} b = [\langle mp + nq, mq + np \rangle]$$

**Lemma 1.5.2.** *Multiplication of integers ( $\cdot_{\mathbb{Z}}$ ) is well defined, i.e. if  $\langle m, n \rangle \sim_{\mathbb{Z}} \langle m', n' \rangle$  and  $\langle p, q \rangle \sim_{\mathbb{Z}} \langle p', q' \rangle$ , then...*

$$\langle mp + nq, mq + np \rangle \sim_{\mathbb{Z}} \langle m'p' + n'q', m'q' + n'p' \rangle$$

The integers under multiplication form an abelian group.

#### 1.5.1.3 Order

Order of integers  $a = \langle m, n \rangle$  and  $b = \langle p, q \rangle$  is defined as...

$$a <_{\mathbb{Z}} b \Leftrightarrow m + q \in p + n$$

**Lemma 1.5.3.** *Order of integers ( $<_{\mathbb{Z}}$ ) is well defined, i.e. if  $\langle m, n \rangle \sim_{\mathbb{Z}} \langle m', n' \rangle$  and  $\langle p, q \rangle \sim_{\mathbb{Z}} \langle p', q' \rangle$ , then...*

$$m + q \in p + n \Leftrightarrow m' + q' \in p' + n'$$

The order relation so defined linearly orders the integers.

### 1.5.2 The Rational Numbers

Let  $\sim_{\mathbb{Q}}$  be the equivalence relation on  $\mathbb{Z} \times (\mathbb{Z} \setminus \{0_{\mathbb{Z}}\})$  for which...

$$\langle a, b \rangle \sim \langle c, d \rangle \Leftrightarrow a \cdot_{\mathbb{Z}} d = c \cdot_{\mathbb{Z}} b.$$

Then the set of *Rational Numbers*, denoted  $\mathbb{Q}$ , is the set  $\mathbb{Z} \times (\mathbb{Z} \setminus \{0_{\mathbb{Z}}\}) / \sim_{\mathbb{Q}}$ .

### 1.5.2.1 Addition

Addition of rational numbers  $p = \langle a, b \rangle$  and  $q = \langle c, d \rangle$  is defined as...

$$p +_{\mathbb{Q}} q = [\langle ad + cb, bd \rangle]$$

**Lemma 1.5.4.** *Addition of rational numbers is well defined.*

The rational numbers under addition form an abelian group.

### 1.5.2.2 Multiplication

Multiplication of rational numbers  $p = \langle a, b \rangle$  and  $q = \langle c, d \rangle$  is defined as...

$$p \cdot_{\mathbb{Q}} q = [\langle ac, bd \rangle]$$

**Lemma 1.5.5.** *Multiplication of rational numbers is well defined.*

The rational numbers under addition and multiplication form a field.

### 1.5.2.3 Order

Order of rational numbers  $p = \langle a, b \rangle$  and  $q = \langle c, d \rangle$  is defined as...

$$p <_{\mathbb{Q}} q \Leftrightarrow ad < cb.$$

**Lemma 1.5.6.** *The order of rational numbers is well-defined.*

The order relation so defined linearly orders the rational numbers.

## 1.5.3 The Real Numbers with Cauchy Sequences

Define a *Cauchy sequence* to be a function  $s : \omega \rightarrow \mathbb{Q}$  such that...

$$(\forall \varepsilon > 0)(\exists k \in \omega)(\forall m > k)(\forall n > k)|s_m - s_n| < \varepsilon.$$

Let  $C$  be the set of all Cauchy sequences. For  $r, s \in C$ , define  $r \sim_{\mathbb{R}} s$  if and only if  $|r_n - s_n|$  is arbitrarily small for large  $n$ .

With more work we can define  $\mathbb{R} := C / \sim$ .

## 1.5.4 The Real Numbers with Dedekind Cuts

A *Dedekind cut* is a subset  $x$  of  $\mathbb{Q}$  such that:

1.  $\emptyset \neq x \neq \mathbb{Q}$
2.  $x$  is "closed downward," i.e.,

$$q \in x \wedge r < q \Rightarrow r \in x.$$

3.  $x$  has no largest member

$\mathbb{R}$  is the set of Dedekind cuts.

### 1.5.4.1 Order

Define an ordering on  $\mathbb{R}$  as...

$$x <_{\mathbb{R}} y \Leftrightarrow x \subset y$$

**Proposition 1.5.1.**  $<_{\mathbb{R}}$  is a linear ordering.

*Proof.*  $<_{\mathbb{R}}$  is clearly transitive; so it suffices to show that  $<_{\mathbb{R}}$  satisfies trichotomy on  $\mathbb{R}$ . So consider  $x, y \in \mathbb{R}$ . Obviously *at most* one of the alternatives,

$$x \subset y, \quad x = y, \quad y \subset x,$$

can hold, but we must prove that at least one holds. Without loss of generality, suppose that the first two fail, i.e., that  $x \not\subseteq y$ .

Since  $x \not\subseteq y$  there is some rational  $r$  in the relative complement  $x \setminus y$ . Consider any  $q \in y$ . If  $r \subseteq q$ , then since  $y$  is closed downward, we would have  $r \in y$ . But  $r \notin y$ , so we must have  $q < r$ . Since  $x$  is closed downward, it follows that  $q \in x$ . Since  $q$  was arbitrary (and  $x \neq y$ ), we have  $y \subset x$ .  $\square$

**Theorem 1.5.7** (Least Upper Bound Property). *Any bounded nonempty subset of  $\mathbb{R}$  has a least upper bound in  $\mathbb{R}$ .*

*Proof.* Let  $A$  be a set of real numbers. The least upper bound is just  $\bigcup A$ .  $\square$

### 1.5.4.2 Addition

Addition of real number  $x, y$  is defined as...

$$x +_{\mathbb{R}} y = q + r \mid q \in x \wedge r \in y$$

### 1.5.4.3 Multiplication

The *absolute value* of a real number  $x$  is defined as...

$$|x| = x \cup -x$$

Multiplication of real number  $x, y$  is defined as follows...

- If  $x$  and  $y$  are nonnegative real numbers, then...

$$x \cdot_{\mathbb{R}} y = 0_{\mathbb{R}} \cup \{rs \mid 0 \leq r \in x \wedge 0 \leq s \in y\}.$$

- If  $x$  and  $y$  are both negative real numbers, then...

$$x \cdot_{\mathbb{R}} y = |x| \cdot_{\mathbb{R}} |y|.$$

- If one of the real numbers  $x$  and  $y$  is negative and one is nonnegative, then...

$$x \cdot_{\mathbb{R}} y = -(|x| \cdot_{\mathbb{R}} |y|).$$

Real numbers under addition, multiplication, and their order relation form an ordered field.

## 1.6 Cardinality

### 1.6.1 Equinumerosity

Two sets  $A$  and  $B$  are *equinumerous*, denoted  $A \approx B$ , if and only if there is a bijection  $f : A \rightarrow B$ .

**Proposition 1.6.1.** *Equinumerosity is an equivalence relation. (See: isomorphism)*

**Theorem 1.6.1** (Diagonalization). *The set  $\omega$  is not equinumerous to the set  $\mathbb{R}$  of real numbers.*

*Proof.* Suppose for the sake of contradiction that there is a bijection  $f : \omega \rightarrow \mathbb{R}$ . Thus we can imagine a list of successive values...

$$f(0) = 236.001\dots$$

$$f(1) = -7.777\dots$$

$$f(2) = 3.1415\dots$$

$$\vdots$$

Then consider the real number  $0.a_1a_2a_3\dots$  where:

$$a_n = \begin{cases} 7 & \text{if the } n\text{th decimal of } f(n) \neq 7 \\ 6 & \text{otherwise.} \end{cases}$$

This number cannot be in the range of  $f$ , so it is not a bijection. ✚

**Theorem 1.6.2** (Diagonalization). *No set is equinumerous to its power set.*

*Proof.* Let  $g : A \rightarrow \mathcal{P}(A)$ . Consider...

$$B = \{x \in A \mid x \notin g(x)\}.$$

Then  $B \subseteq A$ , but for each  $x \in A$ ,

$$x \in B \Leftrightarrow x \notin g(x).$$

Hence  $B \notin \text{ran } g$  and  $g$  is not a bijection. □

### 1.6.2 Finite/Infinite

A set is *finite* if and only if it is equinumerous to some natural number. Otherwise it is *infinite*.

**Theorem 1.6.3** (Pigeonhole Principle). *No natural number is equinumerous to a proper subset of itself.*



*Proof.* Suppose  $f : N \rightarrow N$  is a bijection from a finite set to itself. We will show that  $\text{ran } f$  is all of the set  $n$ . This suffices to prove the theorem.

We use the induction on  $n$ . Define:

$$T = \{n \in \omega \mid \text{every injection from } n \text{ into } n \text{ has range } n\}$$

We have that  $0 \in T$ ; the only function from the set 0 into the set 0 is the empty function, which has range 0. Now suppose that  $k \in T$  and that  $f$  is an injection from  $k^+$  into  $k^+$ . Note that the restriction  $f|_k$  maps  $k$  injectively into  $k^+$ . There are two cases...

Case I: The set  $k$  is closed under  $f$ . Then  $f|_k$  maps the set  $k$  into the set  $k$ . Then because  $k \in T$  we may conclude that  $\text{ran } (f|_k) = k$ . Since  $f$  is injective, the only possible value for  $f(k)$  is the number  $k$ . Hence  $\text{ran } f$  is  $k \cup \{k\}$ , which is the set  $k^+$ .

Case II: Otherwise  $f(p) = k$  for some number  $p$  less than  $k$ . In this case we interchange two values of the function. Define  $\hat{f}$  by...

$$\hat{f}(p) = f(k),$$

$$\hat{f}(k) = f(p) = k,$$

$$\hat{f}(x) = f(x) \text{ for other } x \in k^+.$$

The  $\hat{f}$  maps the set  $k^+$  injectively into the set  $k^+$ , and the set  $k$  is closed under  $\hat{f}$ . So we can apply Case I.

Thus  $\text{ran } f = k^+$ . □

**Corollary 1.6.3.1.** *No finite set is equinumerous to a proper subset of itself.*

**Corollary 1.6.3.2.** *Any set equinumerous to a proper subset of itself is infinite.*

**Corollary 1.6.3.3.** *The set  $\omega$  is infinite.*

**Corollary 1.6.3.4.** *Any finite set is equinumerous to a unique natural number.*

**Lemma 1.6.4.** *If  $C$  is a proper subset of a natural number  $n$ , the  $C \approx m$  for some  $m$  less than  $n$ .*

**Corollary 1.6.4.1.** *Any subset of a finite set is finite.*

### 1.6.3 Cardinal Numbers

For any set  $A$ , the cardinal number of  $A$ , denoted  $\text{card } A$ , is a set...

1. For any sets  $A, B$ ...

$$\text{card } A = \text{card } B \Leftrightarrow A \approx B.$$

2. For a finite set  $A$ ,  $\text{card } A$  is the natural number  $n$  for which  $A \approx n$ .

(See: cardinal number definition using ordinals)

### 1.6.3.1 Cardinal Arithmetic

Let  $\kappa$  and  $\lambda$  be any cardinal numbers.

- $\kappa + \lambda = \text{card}(K \cup L)$ , where  $K$  and  $L$  are any disjoint sets of cardinality  $\kappa$  and  $\lambda$ , respectively.
- $\kappa \cdot \lambda = \text{card}(K \times L)$ , where  $K$  and  $L$  are any sets of cardinality  $\kappa$  and  $\lambda$ , respectively.
- $\kappa^\lambda = \text{card}^L K$ , where  $K$  and  $L$  are any sets of cardinality  $\kappa$  and  $\lambda$ , respectively.

**Proposition 1.6.2.** *Assume that  $K_1 \approx K_2$  and  $L_1 \approx L_2$ .*

1. *If  $K_1 \cap L_1 = K_2 \cap L_2 = \emptyset$ , then  $K_1 \cup L_1 \approx K_2 \cup L_2$ .*
2.  *$K_1 \times L_1 \approx K_2 \times L_2$ .*
3.  *${}^{L_1}K_1 \approx {}^{L_2}K_2$ .*

**Proposition 1.6.3.** *For any cardinal numbers  $\kappa, \lambda$ , and  $\mu \dots$*

- $\kappa + \lambda = \lambda + \kappa$  and  $\kappa \cdot \lambda = \lambda \cdot \kappa$ .
- $\kappa + (\lambda + \mu) = (\kappa + \lambda) + \mu$  and  $\kappa \cdot (\lambda \cdot \mu) = (\kappa \cdot \lambda) \cdot \mu$ .
- $\kappa \cdot (\lambda + \mu) = \kappa \cdot \lambda + \kappa \cdot \mu$ .
- $\kappa^{\lambda+\mu} = \kappa^\lambda \cdot \kappa^\mu$ .
- $(\kappa \cdot \lambda)^\mu = \kappa^\mu \cdot \lambda^\mu$ .
- $(\kappa^\lambda)^\mu = \kappa^{\lambda \cdot \mu}$

**Proposition 1.6.4.** *Let  $m$  and  $n$  be finite cardinals. Then...*

- $m + n = m +_\omega n$
- $m \cdot n = m \cdot_\omega n$
- $m^n = m^n$

(See: natural number arithmetic.)

**Corollary 1.6.4.2.** *If  $A$  and  $B$  are finite, then  $A \cup B$ ,  $A \times B$ , and  ${}^B A$  are also finite.*

### 1.6.3.2 Ordering Cardinal Numbers

A set  $A$  is *dominated* by a set  $B$  (written  $A \preceq B$ ) if and only if there is an injective function from  $A$  into  $B$ .

**Theorem 1.6.5** (Schröder-Bernstein Theorem). *If  $A \preceq B$  and  $B \preceq A$ , then  $A \approx B$ .*

*Proof.* The proof is accomplished with mirrors. Given injections  $f : A \rightarrow B$  and  $g : B \rightarrow A$ . Define  $C_n$  by recursion, using the formulas

$$C_0 = A \setminus \text{ran } g \quad \text{and} \quad C_{n+} = g[f[C_n]].$$

Thus  $C_0$  is the troublesome part that keeps  $g$  from being a bijection. We bounce it back and forth, obtaining  $C_1, C_2, \dots$ . This function showing that  $A \approx B$  is the function  $h : A \rightarrow B$  defined by...

$$h(x) = \begin{cases} f(x) & \text{if } x \in C_n \text{ for some } n, \\ g^{-1}(x) & \text{otherwise.} \end{cases}$$

Note that in the second case ( $x \in A$  but  $x \notin C_n$  for any  $n$ ) it follows that  $x \notin C_0$  and hence  $x \in \text{ran } g$ . So  $g^{-1}(x)$  makes sense in this case. We verify that  $h$  is indeed a bijection. Define  $D_n = f[C_n]$ , so that  $C_{n+} = g[D_n]$ . Consider distinct  $x, y \in A$ . Since both  $f$  and  $g^{-1}$  are injective, the only possible problem arises when, say,  $x \in C_m$  and  $y \in \bigcup_{n \in \omega} C_n$ . In this case,

$$h(x) = f(x) \in D_m,$$

whereas,

$$h(y) = g^{-1}(y) \notin D_m,$$

lest  $y \in C_{m+}$ . So  $h(x) \neq h(y)$ , showing  $h$  is injective.

Finally, we show  $h$  is surjective. Certainly each  $D_n \subseteq \text{ran } h$ , because  $D_n = h[C_n]$ . Consider then a point  $y$  in  $B \setminus \bigcup_{n \in \omega} D_n$ . Where is  $g(y)$ ? Certainly  $g(y) \notin C_0$ . Also  $g(y) \notin C_{n+}$ , because  $C_{n+} = g[D_n]$ ,  $y \notin D_n$ , and  $g$  is injective. So  $g(y) \notin C_n$  for any  $n$ . Therefore  $h(g(y)) = g^{-1}(g(y)) = y$ . This shows that  $y \in \text{ran } h$ , thereby proving part (a).  $\square$

**Theorem 1.6.6** (Restated Schröder-Bernstein Theorem). *For cardinal numbers  $\kappa$  and  $\lambda$ , if  $\kappa \leq \lambda$  and  $\lambda \leq \kappa$ , then  $\kappa = \lambda$ .*

**Proposition 1.6.5.** *Let  $\kappa, \lambda$  and  $\mu$  be cardinal numbers.*

- $\kappa \leq \lambda \Rightarrow \kappa + \mu \leq \lambda + \mu$
- $\kappa \leq \lambda \Rightarrow \kappa \cdot \mu \leq \lambda \cdot \mu$
- $\kappa \leq \lambda \Rightarrow \kappa^\mu \leq \lambda^\mu$
- $\kappa \leq \lambda \Rightarrow \mu^\kappa \leq \mu^\lambda$ ; if not both  $\kappa$  and  $\mu$  equal zero.

### 1.6.3.3 Infinite Cardinal Arithmetic

**Lemma 1.6.7.** *For any infinite cardinal  $\kappa$ , we have  $\kappa \cdot \kappa = \kappa$ .*

**Theorem 1.6.8** (Absorption Law of Cardinal Arithmetic). *Let  $\kappa$  and  $\lambda$  be cardinal numbers, the larger of which is infinite and the smaller of which is nonzero. Then...*

$$\kappa + \lambda = \kappa \cdot \lambda = \max(\kappa, \lambda).$$

## 1.7 Countable Sets

A set  $A$  is *countable* if and only if  $A \preceq \omega$ , i.e. if and only if  $\text{card } A \leq \aleph_0$ .

**Theorem 1.7.1.** *A countable union of countable sets is countable.*

*Proof.* We may suppose that  $\emptyset \notin \mathcal{A}$ , for otherwise we could simply remove it without affecting  $\bigcup \mathcal{A}$ . We may further suppose that  $\mathcal{A} \neq \emptyset$ , since  $\bigcup \emptyset$  is certainly countable. Thus  $\mathcal{A}$  is a countable (but nonempty) function from  $\omega \times \omega$  onto  $\bigcup \mathcal{A}$ . It is easy to find a function from  $\omega$  onto  $\omega \times \omega$ , and the composition will map  $\omega$  onto  $\bigcup \mathcal{A}$ , thereby showing that  $\bigcup \mathcal{A}$  is countable. Since  $\mathcal{A}$  is countable but nonempty, there is a function  $G$  from  $\omega$  onto  $\mathcal{A}$ . We are given that each set  $G(m)$  is countable and nonempty. Hence for each  $m$  there is a function from  $\omega$  onto  $G(m)$ . We must then use the axiom of choice to select such a function for each  $m$ . Let  $H : \omega \rightarrow^\omega (\bigcup \mathcal{A})$  be defined by...

$$H(m) = \{g \mid g \text{ is a function from } \omega \text{ onto } G(m)\}.$$

We know that  $H(m)$  is nonempty for each  $m$ . Hence there is function  $F$  with domain  $\omega$  such that for each  $m$ ,  $F(m)$  is a function from  $\omega$  onto  $G(m)$ . To conclude the proof we have only to let  $f(m, n) = F(m)(n)$ . Then  $f$  is a function from  $\omega \times \omega$  onto  $\bigcup \mathcal{A}$ .  $\square$

## 1.8 Axiom of Choice

(See: set axioms)

**Theorem 1.8.1** (Axiom of Choice). *The following statements are equivalent.*

1. *For any relation  $R$ , there is a function  $F \subseteq R$  with  $\text{dom } F = \text{dom } R$ .*
2. *The Cartesian product of nonempty sets is always nonempty. That is, if  $H$  is a function with domain  $I$  and if  $(\forall i \in I) H(i) \neq \emptyset$ , then there is a function  $f$  with domain  $I$  such that  $(\forall i \in I) f(i) \in H(i)$ .*
3. *For any set  $A$  there is a function  $F$  (a "choice function" for  $A$ ) such that  $F(B) \in B$  for every nonempty  $B \subseteq A$ .*
4. *Let  $\mathcal{A}$  be a set such that (a) each member of  $\mathcal{A}$  is a nonempty set, and (b) any two distinct members of  $\mathcal{A}$  are disjoint. Then there exists a set  $C$  containing exactly one element from each member of  $\mathcal{A}$  (i.e., for each  $B \in \mathcal{A}$  the set  $C \cap B$  is a singleton  $\{x\}$  for some  $x$ ).*

There are other theorems that are equivalent to the axiom of choice.

**Theorem 1.8.2** (Cardinal Comparability). *For any sets  $C$  and  $D$ , either  $C \preceq D$  or  $D \preceq C$ . For any two cardinal numbers  $\kappa$  and  $\lambda$ , either  $\kappa \leq \lambda$  or  $\lambda \leq \kappa$ .*

**Theorem 1.8.3** (Zorn's Lemma). *Let  $\mathcal{A}$  be a set such that for every chain  $\mathcal{B} \subseteq \mathcal{A}$ , we have  $\bigcup \mathcal{B} \in \mathcal{A}$ . ( $\mathcal{B}$  is called a chain if and only if for any  $C$  and  $D$  in  $\mathcal{B}$ , either  $C \subseteq D$  or  $D \subseteq C$ .) Then  $\mathcal{A}$  contains an element  $M$  (a "maximal" element) such that  $M$  is not a subset of any other set in  $\mathcal{A}$ .*

## 1.9 Continuum Hypothesis

**Proposition 1.9.1.** *For any infinite set  $A$ , we have  $\omega \preceq A$ .*

**Proposition 1.9.2.**  $\aleph_0 \leq \kappa$  for any infinite cardinal  $\kappa$ .

**Corollary 1.9.0.1.** *A set is infinite if and only if it is equinumerous to a proper subset of itself.*

The *continuum hypothesis* is:

There is no set  $\mathcal{S}$  such that  $\aleph_0 \prec \text{card } \mathcal{S} \prec 2^{\aleph_0}$ .

## 1.10 Ordinal Numbers

### 1.10.1 Partial Orderings

A *partial ordering* is a relation  $R$  such that...

1.  $R$  is transitive
2.  $R$  is irreflexive, that is for all  $x$  we have  $x \not R x$

**Proposition 1.10.1.** *Assume that  $<$  is a partial ordering. Then for  $x, y$ , and  $z$ :*

1. At most one of the alternatives,

$$x < y, \quad x = y, \quad y < x,$$

*can hold.*

2.  $x \leq y \leq x \Rightarrow x = y$ .

### 1.10.2 Linear Orderings

A *linear ordering* is a partial ordering  $R$  that satisfies trichotomy.

### 1.10.3 Well Orderings

A *well ordering* is a linear ordering  $R$  on  $A$  such that every nonempty subset of  $A$  has a least element.

**Theorem 1.10.1.** *Let  $<$  be a linear ordering on  $A$ . Then  $<$  is a well ordering if and only if there does not exist any function  $f : \omega \rightarrow A$  with  $f(n^+) < f(n)$  for every  $n \in \omega$ .*

**Theorem 1.10.2** (Transfinite Induction Principle). *Assume that  $<$  is a well ordering on  $A$ . Assume that  $B$  is a subset of  $A$  with the special property that for every  $t \in A$ ,*

$$\text{seg } t \subseteq B \Rightarrow t \in B.$$

*Then  $B$  coincides with  $A$ .*

*Proof.* If  $B \subset A$ , then  $A \setminus B$  has a least element  $m$ . But the leastness,  $y \in B$  for any  $y < m$ . But this is to say that  $\text{seg } m \subseteq B$ , so by assumption  $m \in B$  after all.  $\square$

**Proposition 1.10.2.** *Assume that  $<$  is a linear ordering on  $A$ . Further assume that the only subset of  $A$  such that  $\forall t \in A, \text{seg } t \subseteq B \Rightarrow t \in B$  is  $A$  itself. Then  $<$  is a well ordering on  $A$ .*

### 1.10.4 Transfinite Recursion

**Theorem 1.10.3** (Transfinite Recursion Theorem Schema). *For any formula  $\gamma(x, y)$  the following is a theorem:*

*Assume that  $<$  is a well ordering on a set  $A$ . Assume that for any  $f$  there is a unique  $y$  such that  $\gamma(f, y)$ . Then there exists a unique function  $F$  with domain  $A$  such that...*

$$\gamma(F \upharpoonright \text{seg } t, F(t))$$

*for all  $t \in A$ .*

The following axiom is used to prove the transfinite recursion theorem schema.

For any formula  $\varphi(x, y)$  not containing the letter  $B$ , the following is an axiom:

$$\begin{aligned} & \forall[(\forall x \in A) \forall y_1 \forall y_2 (\varphi(x, y_1) \wedge \varphi(x, y_2) \Rightarrow y_1 = y_2) \\ & \Rightarrow \exists B \forall y (y \in B \Leftrightarrow (\exists x \in A) \varphi(x, y))]. \end{aligned}$$

### 1.10.5 Epsilon Images

Let  $<$  be a well ordering on  $A$  and let  $\gamma(x, y)$  be the formula  $y = \text{ran } x$ . Then the transfinite recursion theorem gives an unique function  $E$  with domain  $A$  such that  $\forall t \in A$ :

$$\begin{aligned} E(t) &= \text{ran } (E \upharpoonright \text{seg } t) \\ &= E[\text{seg } t] \\ &= \{E(x) | x < t\}. \end{aligned}$$

The  $\epsilon$ -image of  $\langle A, < \rangle$  is the range of  $E$ .

**Proposition 1.10.3.** *Let  $<$  be a well ordering on  $A$  and let  $E$  be as above and  $\alpha$  its epsilon image.*

1.  $E(t) \notin E(t)$  for any  $t \in A$ .
2.  $E$  maps  $A$  bijectively to  $\alpha$ .
3. For any  $s$  and  $t$  in  $A$ ,

$$s < t \text{ if and only if } E(s) \in E(t)$$

4.  $\alpha$  is a transitive set.

#### 1.10.6 Ordinal Numbers

**Proposition 1.10.4.** *Two well-ordered structures are isomorphic if and only if they have the same  $\epsilon$ -image. That is, if  $<_1$  and  $<_2$  are well orderings on  $A_1$  and  $A_2$ , respectively, then  $\langle A_1, <_1 \rangle \cong \langle A_2, <_2 \rangle$  if and only if the  $\epsilon$ -image of  $\langle A_1, <_1 \rangle$  is the same as the  $\epsilon$ -image of  $\langle A_2, <_2 \rangle$ .*

The *ordinal number* of  $\langle A, < \rangle$  is its  $\epsilon$ -image. An *ordinal number* is a set that is the ordinal number of some well-ordered structure.

#### 1.10.7 Cardinal Numbers

**Theorem 1.10.4** (Numeration Theorem). *Any set is equinumerous to some ordinal number.*

For any set  $A$ , define the cardinal number of  $A$  ( $\text{card } A$ ) to be the least ordinal equinumerous to  $A$ .

## 2 Combinatorics

### 2.1 Basic Methods

Use Cardinality to derive the most basic results.

#### 2.1.1 Addition

**Theorem 2.1.1** (Addition principle). *If  $A$  and  $B$  are two disjoint finite sets, then...*

$$|A \cup B| = |A| + |B|.$$

**Theorem 2.1.2** (Generalized addition principle). *Let  $A_1, A_2, \dots, A_n$  be finite sets that are pairwise disjoint. Then...*

$$|\bigcup_{i=1}^n A_i| = \sum_{i=1}^n |A_i|$$

#### 2.1.2 Subtraction

**Theorem 2.1.3** (Subtraction principle). *Let  $A$  be a finite set, and let  $B \subseteq A$ . Then  $|A \setminus B| = |A| - |B|$ .*

*Proof.* Observe  $|A \setminus B| + |B| = |A|$  by the addition principle.  $\square$

#### 2.1.3 Multiplication

**Theorem 2.1.4** (Product principle). *Let  $X$  and  $Y$  be two finite sets. Then  $|X \times Y| = |X| \times |Y|$ .*

**Theorem 2.1.5** (Generalized product principle). *Let  $X_1, X_2, \dots, X_n$  be finite sets. Then  $|\times_{i \in I}^n X_i| = \prod_{i \in I} |X_i|$ .*

#### 2.1.4 Division

**Theorem 2.1.6.** *Let  $S$  and  $T$  be finite sets so that a  $d$ -to-one function  $f : T \rightarrow S$  exists. Then*

$$|S| = \frac{|T|}{d}.$$

#### 2.1.5 Binomial Coefficients

See permutations.

**Theorem 2.1.7.** *Let  $n$  be a positive integer, and let  $k \leq n$  be a nonnegative integer. Then the number of all  $k$ -element subsets of  $[n]$  is*

$$\binom{n}{k} = \frac{(n)_k}{k!} = \frac{n!}{k!(n-k)!}.$$



Note:  $\binom{n}{k} = \binom{n}{n-k}$  exhibits duality.

**Theorem 2.1.8** (Binomial theorem). *If  $n$  is a positive integer, then...*

$$(x + y)^n = \sum_{k=0}^n \binom{n}{k} x^k y^{n-k}.$$

*Proof.* The left-hand side of the equation contains the factor  $(x + y)$   $n$  times. To compute the product we choose an  $x$  or  $y$  term from each factor and multiply those  $n$  terms together, then do this in all  $2^n$  possible ways, adding all the resulting products. It suffices to show that there are exactly  $\binom{n}{k}$  products of the form  $x^k y^{n-k}$ , which is immediately obvious from the way we compute the product.  $\square$

**Theorem 2.1.9.** *Let  $n$  and  $k$  be nonnegative integers so that  $k < n$ . Then...*

$$\binom{n}{k} + \binom{n}{k+1} = \binom{n+1}{k+1}$$

**Theorem 2.1.10.** *For all positive integers  $n$ ,*

$$\binom{2n}{n} = \sum_{k=0}^n \binom{n}{k}^2.$$

### 2.1.6 Pigeonhole Principle

**Theorem 2.1.11** (Pigeonhole Principle). *Let  $A_1, A_2, \dots, A_k$  be finite sets that are pairwise disjoint. Let us assume that*

$$|A_1 \cup A_2 \cup \dots \cup A_k| > kr.$$

*Then there exists at least one index  $i$  so that  $|A_i| > r$ . (See: Pigeonhole Principle in Set Theory)*

**Example 2.1.11.1.** *Consider the sequence  $1, 3, 7, 15, 31, \dots$ , in other words, the sequence whose  $i$ th element is  $a_i = 2^i - 1$ . Let  $q$  be any odd integer. Then our sequence contains an element that is divisible by  $q$ .*

*Proof.* Consider the first  $q$  elements of our sequence. If one of them is divisible by  $q$ , then we are done. If not, then consider their remainders modulo  $q$ . That is, let us write...

$$a_i = d_i q + r_i$$

where  $0 < r_i < q$ , and  $d_i = \lfloor a_i/q \rfloor$ . As the integers  $r_1, r_2, \dots, r_q$  all come from the open interval  $(0, q)$ , there are  $q - 1$  possibilities for their values. On the other hand, their number is  $q$ , so, by the pigeonhole principle, there have to be two of them that are equal. Say these are  $r_n$  and  $r_m$ , with  $n > m$ . Then  $a_n = d_n q + r_n$  and  $a_m = d_m q + r_m$ , so...

$$a_n - a_m = (d_n - d_m)q$$

or, after rearranging,

$$\begin{aligned}
(d_n - d_m)q &= a_n - a_m \\
&= (2^n - 1) - (2^m - 1) \\
&= 2^m(2^{n-m} - 1) \\
&= 2^m a_{n-m}
\end{aligned}$$

As the first expression of our chain of equations is divisible by  $q$ , so too must be the last expression. Note that  $2^{n-m}$  is relatively prime to any odd number  $q$ , that is, the largest common divisor of  $2^{n-m}$  and  $q$  is 1. Therefore, the equality  $(d_n - d_m)q = 2^{n-m}a_{n-m}$  implies that  $a_{n-m}$  is divisible by  $q$ .  $\square$

## 2.2 Applications of Basic Methods

### 2.2.1 Inclusion-Exclusion

**Theorem 2.2.1** (Inclusion-exclusion principle). *Let  $A_1, A_2, \dots, A_n$  be finite sets. Then...*

$$|A_1 \cup A_2 \cdots \cup A_n| = \sum_{j=1}^n (-1)^{j-1} \sum_{i_1, i_2, \dots, i_j} |A_{i_1} \cap A_{i_2} \cap \cdots \cap A_{i_j}|,$$

where  $(i_1, i_2, \dots, i_j)$  ranges all  $j$ -element subsets of  $[n]$ .

*Proof.* We prove the two following claims:

1. If  $x$  is contained in the set represented on the left side of the equation, then the right side counts it exactly once.
2. If  $x$  is not contained in any  $A_i$ , then the right-hand side counts  $x$  zero times.

(1) Assume that  $x$  is contained in exactly  $k$  of the  $n$   $A_i$ -sets, with  $k > 0$ . Certainly,  $x$  is not in any  $j$ -fold intersection where  $j > k$ . On the other hand  $j \leq k$ , then  $x$  is contained in exactly  $\binom{k}{j}$  different  $j$ -fold intersections. If we take the signs into account, this means that the right side counts  $x$  exactly...

$$m = \sum_{j=1}^k (-1)^{j-1} \binom{k}{j}$$

times. Now we show that  $m = 1$  necessarily. Observe...

$$1 - m = \sum_{j=0}^k (-1)^j \binom{k}{j} = (1 - 1)^k = 0,$$

since  $k$  is positive.

(2) We repeat the above argument with  $k = 0$ . Then the binomial theorem technique we use above gives us  $(1 - 1)^0 = 1$ , implying  $m = 0$ .

Thus the left-hand side and the right-hand side count the same objects.  $\square$

### 2.2.2 Multisets

Given a set  $A$ , a *multiset* is defined via a function  $m : A \rightarrow \mathbb{N} \cup \{0\}$ . It is a set containing  $a \in A$   $m(a)$  many times.

#### 2.2.2.1 Multinomial Coefficients

**Theorem 2.2.2.** *Given a multiset  $A$  of  $n$  elements over a  $k$  element sets. The number of ways to linearly order the elements of  $A$  is...*

$$\binom{n}{a_1, a_2, \dots, a_k} = \frac{n!}{a_1! a_2! \dots a_k!}.$$

### 2.2.3 Weak Compositions

Let  $a_1, a_2, \dots, a_k$  be nonnegative integers satisfying

$$\sum_{i=1}^k a_i = n.$$

Then the ordered  $k$ -tuple  $(a_1, a_2, \dots, a_k)$  is called a *weak composition* of  $n$  into  $k$  parts.

**Theorem 2.2.3.** *The number of weak compositions of  $n$  into  $k$  parts is...*

$$\binom{n+k-1}{n} = \binom{n+k-1}{k-1}.$$

**Corollary 2.2.3.1.** *The number of  $n$ -element multisets over a  $k$ -element set is...*

$$\binom{n+k-1}{n} = \binom{n+k-1}{k-1}.$$

### 2.2.4 Compositions

Let  $a_1, a_2, \dots, a_k$  be positive integers satisfying

$$\sum_{i=1}^k a_i = n.$$

Then the ordered  $k$ -tuple  $(a_1, a_2, \dots, a_k)$  is called a *composition* of  $n$  into  $k$  parts.

**Corollary 2.2.3.2.** *The number of compositions of  $n$  into  $k$  parts is...*

$$\binom{n-1}{k-1}.$$

### 2.2.5 Stirling numbers of the second kind

Given a finite set  $A$ ,  $|A| = n$ , the number of set partitions of  $A$  into  $0 < k \leq n$  classes is denoted  $S(n, k)$ , the *Stirling number of the second kind*.

**Theorem 2.2.4.** *For all positive integers  $n$  and  $k$  satisfying  $n \leq k$ , the equality...*

$$S(n, k) = S(n-1, k-1) + kS(n-1, k)$$

**Theorem 2.2.5.** *For all positive integers  $n$  and  $k$  satisfying  $n \geq k$ .*

$$S(n+1, k) = \sum_{i=0}^n \binom{n}{i} S(n-i, k-1)$$

**Theorem 2.2.6.** *The number of surjections from  $[n]$  to  $[k]$  is equal to*

$$\sum_{j=0}^k (-1)^j \binom{k}{j} (k-j)^n.$$

**Corollary 2.2.6.1.** *For all positive integers  $k \leq n$ ,*

$$S(n, k) = \frac{1}{k!} \sum_{j=0}^k (-1)^j \binom{k}{j} (k-j)^n.$$

#### 2.2.5.1 Bell numbers

The number of all partitions of a finite set  $A$ , where  $|A| = n$ , is denoted  $B(n)$  and is called a *Bell number*.

**Theorem 2.2.7.** *Set  $B(0) = 1$ . Then, for all positive integers  $n$ ,*

$$B(n+1) = \sum_{k=0}^n B(k) \binom{n}{k}.$$

### 2.2.6 Partitions of integers

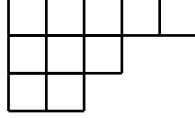
A *partition of an integer  $n$*  is a finite sequence  $(a_1, a_2, \dots, a_k)$  of positive integers satisfying  $a_1 \geq a_2 \geq \dots \geq a_k$  and  $a_1 + a_2 + \dots + a_k = n$ .

**Theorem 2.2.8.** *As  $n \rightarrow \infty$ , the function  $p(n)$  satisfies...*

$$p(n) \sim \frac{1}{4\sqrt{3}} \exp \left( \pi \sqrt{\frac{2n}{3}} \right)$$

### 2.2.7 Ferrers shapes

The *Ferrers shape* of the partition  $(a_1, a_2, \dots, a_k)$  is a row diagram of squares, with non-increasing amounts of squares in lower rows. For example the Ferrers shape for  $(5, 3, 2)$  is...



**Proposition 2.2.1.** *For all positive integers  $k \leq n$ , the number of partitions of  $n$  that have at least  $k$  parts is equal to the number of partitions of  $n$  in which the largest part is at least  $k$ .*

**Proposition 2.2.2.** *For every positive integer  $n$ , the number of partitions of  $n$  in which the first two parts are equal is equal to the number of partitions of  $n$  in which each part is at least 2.*

**Lemma 2.2.9.** *Let  $m > k \geq 1$ . Let  $S$  be the set of partitions of  $n$  into  $m$  parts, the smallest of which is equal to  $k$ , and let  $T$  be the set of partitions of  $n$  into  $m - 1$  parts, in which the  $k$ th part is larger than the  $(k + 1)$ st part and the smallest part is at least  $k$ . Then  $|S| = |T|$ .*

### 2.2.8 Euler's totient function

For any positive integer  $n$ , let  $\phi(n)$  denote the number of positive integers  $k \leq n$  that are relatively prime to  $n$ .

**Proposition 2.2.3.** *Let  $n = pq$ , where  $p$  and  $q$  are distinct primes. Then  $\phi(n) = (p - 1)(q - 1)$ .*

*Proof.* Use the inclusion-exclusion principle on  $[pq]$ , followed by the subtraction principle.  $\square$

*Proof.* Let  $n = p_1 p_2 \dots p_t$ , where the  $p_i$  are pairwise distinct primes. Then...

$$\phi(n) = \prod_{i=1}^t (p_i - 1).$$

$\square$

**Lemma 2.2.10.** *Let  $a$  and  $b$  be two positive integers whose greatest common divisor is 1, and let  $n = ab$ . Then  $\phi(n) = \phi(a)\phi(b)$ .*

**Proposition 2.2.4.** *For any prime  $p$ , and any positive integer  $d$ ,*

$$\phi(p^d) = (p - 1)p^{d-1}.$$

**Proposition 2.2.5.** *Let  $n = p_1^{d_1} p_2^{d_2} \dots p_t^{d_t}$ , where the  $p_i$  are distinct primes. Then...*

$$\phi(n) = \prod_{i=1}^t p_i^{d_i-1} (p_i - 1)$$

## 2.3 Permutations

Given a set  $A$ , a *permutation* of  $A$  is a bijection  $f : A \rightarrow A$ .

**Proposition 2.3.1.** *Given a finite set  $A$ , if  $n = |A|$  the number of permutations of  $A$  is  $n!$ .*

Intuitively permutations represent the reordering of an ordered list. Looking at the idea of "sub-orderings" of lists we come up with the following proposition...

**Proposition 2.3.2** (*k*-lists). *Let  $n$  and  $k$  be positive integers so that  $n \geq k$ . Then the number of injections  $f : [k] \rightarrow [n]$  is...*

$$(n)_k := n(n-1)(n-2) \cdots (n-k+1).$$

## 2.4 Twelfold Way

There are 12 fundamental counting problems. Sometimes they are formulated in terms of putting *balls* into *baskets*.

Let  $N$  and  $K$  be finite sets and  $n$  and  $k$  be their cardinality respectively...

### 2.4.1 Functions from $K$ to $N$

Count with sequences of  $k$  elements in  $N$ ,  $|^K N|$ .

### 2.4.2 Injections from $K$ to $N$

Count with  $k$ -lists,  $(n)_k$ .

### 2.4.3 Surjections from $K$ to $N$

Count with the number of surjections from  $[k]$  to  $[n]$ ,  $\sum_{j=0}^n (-1)^j \binom{n}{j} (n-j)^k$ .

### 2.4.4 Injections from $K$ to $N$ , up to a permutation of $K$

Count subsets,  $k$ -lists without order,  $\binom{n}{k}$ .

### 2.4.5 Functions from $K$ to $N$ , up to a permutation of $K$

Count multisets with  $k$  elements from  $N$ ,  $\binom{n+k-1}{k}$ .

### 2.4.6 Surjections from $K$ to $N$ , up to a permutation of $K$

Count compositions of  $k$  into  $n$  parts,  $\binom{k-1}{n-1}$ .

#### 2.4.7 Injections from $K$ to $N$ , up to a permutation of $N$

Provided  $k \leq n$ , there is only 1 of these.

#### 2.4.8 Surjections from $K$ to $N$ , up to a permutation of $N$

Count partitions of  $K$  into  $n$  non-empty subsets,  $S(k, n)$ .

#### 2.4.9 Functions from $K$ to $N$ , up to a permutation of $N$

Count all the partitions of  $K$  up to  $n$  classes,  $\sum_{i=0}^n \binom{k}{i}$ . If  $k \leq n$ ,  $B(k)$ .

#### 2.4.10 Functions from $K$ to $N$ , up to a permutation of $K$ and $N$

Count partitions of  $k$  into  $\leq n$  non-empty subsets,  $\sum_{i=0}^n p_i(k)$ .

#### 2.4.11 Injections from $K$ to $N$ , up to a permutation of $K$ and $N$

Provided  $k \leq n$ , there is only 1 of these.

#### 2.4.12 Surjections from $K$ to $N$ , up to a permutation of $K$ and $N$

Count partitions of  $k$  into  $n$  non-empty subsets,  $p_n(k)$ .

### 2.5 Graphs

A *graph* is an ordered pair  $G = \langle V, E \rangle$  comprising a set  $V$  of *nodes* and  $E$  of *edges*, which are 2-element subsets of  $V$ .

**Proposition 2.5.1.** *Let  $d_1, d_2, \dots, d_n$  be the degrees of the vertices of a graph  $G$  on  $n$  vertices that has  $e$  edges. Then we have...*

$$d_1 + d_2 + \dots + d_n = 2e.$$

#### 2.5.1 Simple Graph

A *simple graph* is a graph that contains no loops and no multiple edges.

##### 2.5.1.1 Walk

A *walk* is a series  $e_1 e_2 \dots e_k$  of edges that lead from a vertex to another one.

##### 2.5.1.2 Cycle

A *cycle* is a walk whose starting point.

#### 2.5.2 Graph Isomorphisms

An isomorphism  $f$  of two graphs  $G, H$  is a bijection from  $V(G)$  to  $V(H)$  such that if  $\{a, b\} \in E(G)$ , then  $\{f(a), f(b)\} \in E(H)$ .

### 2.5.2.1 Group of Automorphisms

Define the group of automorphisms of a graph  $G$ ,  $\text{Aut}(G)$ , as normal.

Let  $J$  be a graph on  $n$  unlabeled vertices. Then define  $\ell(J)$  as the number of possible ways to bijectively label  $J$  so that the resulting graphs are non-isomorphic.

**Proposition 2.5.2.** *For any graph  $H$  on vertex set  $[n]$ ,*

$$|\text{Aut}(H)| \cdot \ell(H) = n!$$

### 2.5.3 Trees

A *tree* is a simple graph that is minimally connected.

#### 2.5.3.1 Minimally Connected Graph

A *minimally connected* graph is contains the least number of edges in order to be connected.

**Lemma 2.5.1.** *Let  $G$  be a connected simple graph on  $n$  vertices. Then the following are equivalent.*

1. *The graph  $G$  is minimally connected.*
2. *There are no cycles in  $G$ .*
3. *The graph  $G$  has exactly  $n - 1$  edges.*

*Proof.* (1)  $\Rightarrow$  (2) Assume there is a cycle  $C$  in  $G$ . Then  $G$  cannot be minimally connected since any one edge  $e$  of  $C$  can be omitted, and the obtained graph  $G'$  is still connected. Indeed, if a path  $uv$  used the edge  $e$ , then there would be a walk from  $u$  to  $v$  in which the edge  $e$  is replaced by the set edges of  $C$  that are different from  $e$ .

(2)  $\Rightarrow$  (3) Pick any vertex  $x \in G$  and start walking in some direction, never revisiting a vertex. As there is no cycle in  $G$ , eventually we will get stuck, meaning that we will hit a vertex of degree 1. This means that a connected simple graph with no cycles contains a vertex of degree 1. Removing such a vertex (and the only edge adjacent to it) from  $G$ , we get a graph  $G^*$  with one less vertex and one less edge, and the statement is proved by induction on  $n$ .

(3)  $\Rightarrow$  (1) Suppose for the sake of contradiction that a graph on  $n$  vertices and  $n-2$  edges cannot be connected. Let  $H$  be such a graph with a minimum number of vertices. Then  $H$  must have more than 3 vertices. As  $H$  has  $n-2$  edges, there has to be a vertex  $y$  of degree 1 in  $H$ , otherwise  $H$  would need to have at least  $n$  edges. Removing  $y$  from  $H$ , we get an even smaller counterexample for our statement, which is a contradiction.  $\square$

**Theorem 2.5.2** (Cayley's formula). *For all positive integers  $n$ , the number of all trees on vertex set  $[n]$  is  $n^{n-2}$ .*



*Proof.* We need to prove that  $T_n = n^{n-2}$ , which is the number of all functions from  $[n-2]$  to  $[n]$ . This is certainly equivalent to proving the identity...

$$n^2 T_n = n^n.$$

Here the right-hand side is the number of all functions from  $[n]$  into  $[n]$ . The left-hand side, on the other hand, is equal to the number of all trees on  $[n]$  in which we select two vertices, called Start and End (which may be identical). Let us call these trees *doubly rooted trees*.

We construct a bijection  $G$  to prove the above formula. Let  $f \in \text{End}_{\text{Set}}([n])$  and draw its *short diagram*, that is, represent  $x \in [n]$  as a vertex in a graph, where there is an arrow  $\langle x, y \rangle$  if and only if  $f(x) = y$ . This creates two kinds of vertices, namely, those that are in a directed cycle and those that are not. Let  $C$  and  $N$ , respectively, denote these two subsets of  $[n]$ .

Now we start creating the doubly rooted tree  $G(f)$ . First, note that  $f$  acts as a permutation on  $C$ . If  $C = \{c_1, c_2, \dots, c_k\}$  so that  $c_1 < c_2 < \dots < c_k$ , call  $f(c_1)$  Start and  $f(c_k)$  End, and create a path with vertices  $f(c_1), f(c_2), \dots, f(c_k)$ . Note that so far we have defined a graph with  $k$  vertices and  $k-1$  edges.

If  $x \in N$ , then simply connect  $x$  to  $f(x)$ , just as in the short diagram of  $f$ . This will define  $n-k$  more edges. Therefore, we now have a graph on  $[n]$  that has  $n-1$  edges and has two vertices (called Start and End, respectively). This is the graph that we want to call  $G(f)$ . In order to justify that name, we must prove that  $G(f)$  is connected. This is true, however, since, in the short diagram of  $f$ , each directed path starting at any  $x \in N$  must reach a vertex of  $C$  at some point (there is no other way it could end). So indeed,  $G(f)$  is a doubly rooted tree for all  $f \in \text{End}_{\text{Set}}([n])$ .

In order to show that  $G$  is a bijection, we prove it has an inverse. Let  $t$  be a doubly rooted tree. Then there is a unique path  $p$  from Start to End in  $t$ . To find  $f = G^{-1}(t)$ , just put the vertices along  $p$  into  $C$ , and put all the other vertices to  $N$ . If  $x \in N$ , then define  $f(x)$  as the unique neighbor of  $x \in t$  that is closer to  $p$  than  $x$ . For the vertices  $x \in C$ , we define  $f$  so that the  $i$ th vertex of the Start-End path is the image of the  $i$ th smallest element of  $C$ . It is a direct consequence of the definition of  $G$  that this way we will get an  $f \in \text{End}_{\text{Set}}([n])$  satisfying  $G(f) = t$ , and that this  $f$  is the only preimage of  $t$  under  $G$ . Therefore,  $G$  is a bijection.  $\square$

## 3 Category Theory

### 3.1 Metacategories

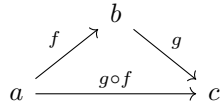
#### 3.1.1 Undefined notions

- *Objects:*  $a, b, c \dots$
- *Arrows:*  $f, g, h \dots$

#### 3.1.2 Operations

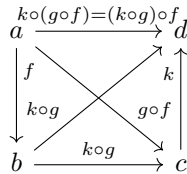
Given  $f : a \rightarrow b \dots$

- *Domain:* **dom:** arrows  $\rightarrow$  objects,  $f \mapsto a$
- *Codomain:* **cod:** arrows  $\rightarrow$  objects,  $f \mapsto b$
- *Identity:* **id:** objects  $\rightarrow$  arrows,  $a \mapsto \text{id}_a = 1_a$
- *Composition:* **comp:** arrows  $\times$  : arrows  $\rightarrow$  arrows,  $\langle g, f \rangle \mapsto g \circ f$ ,  
 $g \circ f : \text{dom} f \rightarrow \text{cod} g$

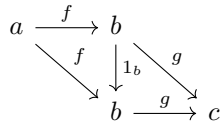


#### 3.1.3 Axioms

- *Associativity:*  $a \xrightarrow{f} b \xrightarrow{g} c \xrightarrow{k} d$ ,  $k \circ (g \circ f) = (k \circ g) \circ f$



- *Unit Law:*  $1_a \circ f = f$  and  $g \circ 1_b = g$



## 3.2 Categories

### 3.2.1 Directed Graph

- $A$  - a set of arrows
- $O$  - a set of objects
- $\mathbf{dom} : A \rightarrow O, \mathbf{cod} : A \rightarrow O$

#### 3.2.1.1 Set of composable pairs of arrows

$$A \times_O A = \{\langle g, f \rangle \mid g, f \in A \text{ and } \mathbf{dom}(g) = \mathbf{cod}(f)\}$$

### 3.2.2 Categories

Add the following structure to a directed graph. . .

- $O \xrightarrow{id} A, c \mapsto id_C$
- $A \times_O A \xrightarrow{\circ} A, \langle g, f \rangle \mapsto g \circ f$

which satisfy  $\forall a \in O$  and  $\forall \langle g, f \rangle \in A \times_O A$  . .

- $\mathbf{dom}(\mathbf{id}(a)) = a = \mathbf{cod}(\mathbf{id}(a))$
- $\mathbf{dom}(g \circ f) = \mathbf{dom}(f)$
- $\mathbf{cod}(g \circ f) = \mathbf{cod}(g)$
- metacategorical axioms

### 3.2.3 Small categories

*Small categories* use small sets for their objects.

### 3.2.4 Hom Sets

$$\mathbf{hom}(b, c) = \{f \mid f \in C, \mathbf{dom}(f) = b, \mathbf{cod}(f) = c\}$$

#### 3.2.4.1 Alternate Definition of Categories

Small categories may be defined with hom-sets as follows. . .

1. A set of objects  $a, b, c, \dots$
2. A function which assigns to each ordered pair  $\langle a, b \rangle$  of objects a set  $\mathbf{hom}(a, b)$

3. For each ordered triple  $\langle a, b, c \rangle$  of objects a function

$$\text{hom}(b, c) \times \text{hom}(a, b) \rightarrow \text{hom}(a, c)$$

called composition, and written  $\langle g, f \rangle \rightarrow g \circ f$  for  $g \in \text{hom}(b, c)$ ,  $f \in \text{hom}(a, b)$

4. For each object  $b$ , an element  $1_b \in \text{hom}(b, b)$ , called the identity of  $b$ .  
 5. If  $\langle a, b \rangle \neq \langle a', b' \rangle$ , then  $\text{hom}(a, b) \cap \text{hom}(a', b') = \emptyset$

The above satisfy the meta-categorical axioms.

Functors in terms of hom-sets are the object function with a collection of functions

$$T_{c,c'} : \text{hom}_C(c, c') \rightarrow \text{hom}_B(Tc, Tc')$$

such that each  $T_{c,c'} 1_c = 1_{Tc}$  and every diagram...

$$\begin{array}{ccc} \text{hom}_C(c', c'') \times \text{hom}_C(c, c') & \xrightarrow{\circ} & \text{hom}_C(c, c'') \\ \downarrow T_{c',c''} \times T_{c,c'} & & \downarrow T_{c',c''} \\ \text{hom}_B(Tc', Tc'') \times \text{hom}_B(Tc, Tc') & \xrightarrow{\circ} & \text{hom}_B(Tc, Tc'') \end{array}$$

is commutative.

### 3.2.5 Groupoids

A category in which every arrow is an isomorphism.

## 3.3 Morphisms

Arrows in categories.

### 3.3.1 Isomorphisms

A morphism  $f \in \text{hom}(b, c)$  that has a two-sided inverse  $g \in \text{hom}(c, b)$  under composition such that

$$gf = 1_b, \quad fg = 1_c.$$

**Proposition 3.3.1.** *The inverse of an isomorphism is unique.*

*Proof.* For inverses  $g_1, g_2$  of  $f$  observe...

$$g_1 = g_1 1_c = g_1 (fg_2) = (g_1 f) g_2 = 1_b g_2 = g_2$$

□

**Proposition 3.3.2.** *Supposing  $f^{-1}$  is the inverse of  $f$ ...*

- Each identity  $1_c$  is an isomorphism and is its own inverse.
- If  $f$  is an isomorphism, then  $f^{-1}$  is an isomorphism and further  $(f^{-1})^{-1} = f$ .
- If  $f \in \text{hom}(a, b)$ ,  $g \in \text{hom}(b, c)$  are isomorphisms, then the composition  $gf$  is an isomorphism and  $(gf)^{-1} = f^{-1}g^{-1}$ .

### 3.3.2 Automorphisms

An isomorphism of an object to itself. Denoted:

$$\text{hom}(c, c) = \text{aut}(c)$$

Observe  $\text{aut}(c)$  is a group.

### 3.3.3 Monomorphisms

A morphism  $f \in \text{hom}(b, c)$  such that  $\forall z \in C$  and  $\forall \alpha', \alpha'' \in \text{hom}(z, b)$ :

$$f \circ \alpha' = f \circ \alpha'' \Rightarrow \alpha' = \alpha''$$

### 3.3.4 Epimorphisms

A morphism  $f \in \text{hom}(b, c)$  such that  $\forall z \in C$  and  $\forall \beta', \beta'' \in \text{hom}(b, z)$ :

$$\beta' \circ f = \beta'' \circ f \Rightarrow \beta' = \beta''$$

### 3.3.5 Split Morphism

A morphism  $f : b \rightarrow b$  such that  $f^2 = f$  and there exist morphisms  $g : b \rightarrow c$ ,  $h : c \rightarrow b$  satisfying...

$$f = hg \quad \wedge \quad gh = 1_c$$

## 3.4 Some Objects in Categories

### 3.4.1 Initial Objects

We say that an object  $i$  of a category  $C$  is *initial* in  $C$  if for every object  $a$  of  $C$  there exists exactly one morphism  $i \rightarrow a$  in  $C$ :

$$\forall a \in \text{Obj}(C) : \quad \text{Hom}_C(i, a) \text{ is a singleton.}$$

### 3.4.2 Final Objects

We say that an object  $f$  of a category  $C$  is *final* in  $C$  if for every object  $a$  of  $C$  there exists exactly one morphism  $a \rightarrow f$  in  $C$ :

$$\forall a \in \text{Obj}(C) : \text{Hom}_C(a, f) \text{ is a singleton.}$$

**Proposition 3.4.1.** *Let  $C$  be a category.*

- If  $i_1, i_2$  are both initial objects in  $C$ , then  $i_1 \cong i_2$ .
- If  $f_1, f_2$  are both final objects in  $C$ , then  $f_1 \cong f_2$ .

### 3.4.3 Null Objects

An object that is both initial and terminal.

## 3.5 Functors

Morphisms  $T : C \rightarrow B$  with domain and codomain both categories. It consists of two suitably related functions

- object function  $T, c \mapsto Tc$
- arrow function  $T, f : c \rightarrow c' \mapsto Tf : Tc \rightarrow Tc'$

which satisfy...

- $T(1_c) = 1_{Tc}$
- $T(g \circ f) = Tg \circ Tf$

### 3.5.1 Full

$$\forall c, c' \in C \text{ and } g : Tc \rightarrow Tc' \in B, \exists f : c \rightarrow c' \in C \text{ s.t. } g \in Tf$$

### 3.5.2 Faithful

$$\forall c, c' \in C \text{ and } f_1, f_2 : c \rightarrow c', Tf_1 = Tf_2 \Rightarrow f_1 = f_2$$

### 3.5.3 Forgetful

A functor that drops some of the structure of its input. For example, the forgetful functor  $U : \text{Cat} \rightarrow \text{Graph} \dots$

- $C \mapsto UC$  where  $UC$  is comprised of the underlying objects and arrows of the category
- $F : C \rightarrow C' \mapsto UF : UC \rightarrow UC'$  where  $UF$  is a morphism between corresponding graphs

### 3.6 Natural Transformations

Given two functors  $S, T : C \rightarrow B$  a *natural transformation*  $\tau : S \rightarrow T$  is a function which assigns to each object  $c \in C$  an arrow

$$\tau_c = \tau c : Sc \rightarrow Tc$$

of  $B$  in such a way that every arrow  $f : c \rightarrow c'$  in  $C$  yields a diagram...

$$\begin{array}{ccccc} c & & Sc & \xrightarrow{\tau c} & Tc \\ \downarrow f & & \downarrow Sf & & \downarrow Tf \\ c' & & Sc' & \xrightarrow{\tau c'} & Tc' \end{array}$$

which is commutative.

In the following diagram  $\tau a, \tau b, \tau c$  are the components of the natural transformation.

$$\begin{array}{ccccc} a & & Sa & \xrightarrow{\tau a} & Ta \\ & \searrow f & \downarrow Sf & & \downarrow Tf \\ & & Sc & \xrightarrow{\tau c} & Tc \\ \downarrow h & & \downarrow Sh & & \downarrow Th \\ b & \nearrow g & Sb & \xrightarrow{\tau b} & Tb \\ & & \uparrow Sg & & \uparrow Tg \end{array}$$

### 3.7 Duality

Statement $\Sigma$	Dual Statement $\Sigma^*$
$f : a \rightarrow b$	$f : b \rightarrow a$
$a = \text{dom} f$	$a = \text{cod} f$
$i = 1_a$	$i = 1_a$
$h = g \circ f$	$h = f \circ g$
$f$ is a monomorphism	$f$ is an epimorphism
$u$ is a right inverse of $h$	$u$ is a left inverse of $h$
$f$ is invertible	$f$ is invertible
$f$ is a terminal object	$f$ is an initial object

### 3.8 Contravariance and Opposites

#### 3.8.1 Contravariant Functor

Given a functor  $S : C^{op} \rightarrow B$  the *contravariant functor*  $\bar{S} : C \rightarrow B$  satisfies...

- $\overline{S}f = Sf^{op}$ ,
- $c \mapsto \overline{S}c$ ,
- $f : a \rightarrow b \mapsto \overline{S}f : \overline{S}b \rightarrow \overline{S}a$ ,
- $\overline{S}(1_c) = 1_{\overline{S}c}$ ,
- $\overline{S}(fg) = (\overline{S}g)(\overline{S}f)$ .

### 3.8.1.1 Covariant Hom-Functor

A hom-functor  $C(a, -) = \text{hom}(a, -) : C \rightarrow \text{Set}$  satisfying...

- $b \mapsto \text{hom}(a, b)$
- $k : b \rightarrow b' \mapsto \text{hom}(a, k) : \text{hom}(a, b) \rightarrow \text{hom}(a, b')$ ; the right side maps  $f \mapsto k \circ f$  and is denoted  $k^*$

### 3.8.1.2 Contravariant Hom-Functor

A hom-functor  $C(-, b) = \text{hom}(-, b) : C^{op} \rightarrow \text{Set}$  satisfying...

- $a \mapsto \text{hom}(a, b)$
- $g : a \rightarrow a' \mapsto \text{hom}(g, a) : \text{hom}(a', b) \rightarrow \text{hom}(a, b)$ ; the right side maps  $f \mapsto f \circ g$  and is denoted  $g^*$

The functions  $g^*, k^*$  defined above satisfy the following commutative diagram.

$$\begin{array}{ccc}
 \text{hom}(a', b) & \xrightarrow{g^*} & \text{hom}(a, b) \\
 \downarrow k^* & & \downarrow k^* \\
 \text{hom}(a', b') & \xrightarrow{g^*} & \text{hom}(a, b')
 \end{array}$$

## 3.9 Category Constructions

### 3.9.1 Products

Given categories  $B$  and  $C$  we construct the product category  $B \times C$ ...

- Objects: pairs of objects  $\langle b, c \rangle$  ( $b \in B$  and  $c \in C$ )
- Arrows:  $\langle b, c \rangle \rightarrow \langle b', c' \rangle$  are a pair  $\langle f, g \rangle$  of arrows ( $f \in B$  and  $g \in C$ )
- Composition:  $\langle f', g' \rangle \circ \langle f, g \rangle = \langle f' \circ f, g' \circ g \rangle$



The corresponding universal property is: for any functors  $R$  and  $T$ , there is a unique functor  $F$  making the digram commute...

$$\begin{array}{ccccc} & & D & & \\ & \swarrow R & \vdots F & \searrow T & \\ B & \xleftarrow{P} & B \times C & \xrightarrow{Q} & C \end{array}$$

Note:  $P\langle f, g \rangle = f$  and  $Q\langle f, g \rangle = g$  are called the *projections* of the product.

### 3.9.1.1 Products of Functors

Given functors  $U$  and  $V$ , the functor product  $U \times V$  satisfies...

- $(U \times V)\langle b, c \rangle = \langle Ub, Uc \rangle$  for objects
- $(U \times V)\langle f, g \rangle = \langle Uf, Ug \rangle$  for arrows

$$\begin{array}{ccccc} B & \xleftarrow{P} & B \times C & \xrightarrow{Q} & C \\ \downarrow U & & \vdots U \times V & & \downarrow V \\ B & \xleftarrow{P'} & B \times C & \xrightarrow{Q'} & B \end{array}$$

### 3.9.1.2 Bifunctors

A functor  $S : B \times C \rightarrow D$ . Intuitively, "a functor of two variables."

Determined by the functors that result when any one object of exactly one of the categories is fixed. This is recorded more explicitly in the following proposition...

**Proposition 3.9.1.** *Let  $B, C$ , and  $D$  be categories. For all objects  $c \in C$  and  $b \in B$ , let*

$$L_c : B \rightarrow D, \quad M_b : C \rightarrow D$$

*be functors such that  $M_b(c) = L_c(b)$  for all  $b$  and  $c$ . Then there exists a bifunctor  $S : B \times C \rightarrow D$  with  $S(-, c) = L_c$  for all  $c$  and  $S(b, -) = M_b$  for all  $b$  if and only if for every pair of arrows  $f : b \rightarrow b'$  and  $g : c \rightarrow c'$  one has*

$$M_{b'}g \circ L_cf = L_{c'}f \circ M_bg.$$

*These equal arrows in  $D$  are then the value  $S(f, g)$  of the arrow function of  $S$  at  $f$  and  $g$ .*

*Proof.* Observe...

$$\langle b', g \rangle \circ \langle f, c \rangle = \langle b'f, gc \rangle = \langle f, g \rangle = \langle fb, c'g \rangle = \langle f, c' \rangle \circ \langle b, g \rangle$$

(where  $b, b', c, c'$  are identity arrows).

This implies...

$$S(b', g)S(f, c) = S(f, c')S(b, g).$$

Which further implies...

$$\begin{array}{ccc} S(b, c) & \xrightarrow{S(b, g)} & S(b, c') \\ \downarrow S(f, c) & & \downarrow S(f, c') \\ S(b', c) & \xrightarrow{S(b', g)} & S(b', c') \end{array}$$

□

### 3.9.1.3 Natural transformations between bifunctors

Given  $S, S' : B \times C \rightarrow D$ . Consider  $\alpha(b, c) : S(b, c) \rightarrow S'(b, c)$ . We say  $\alpha$  is *natural in  $b$*  if  $\forall c \in C$  the components  $\alpha(b, c)$  for all  $b$  define  $\alpha(-, c) : S(-, c) \rightarrow S'(-, c)$ , a natural transformation of functors  $B \rightarrow D$ .

**Proposition 3.9.2.** *For bifunctors  $S, S'$ , the function  $\alpha$  displayed above is a natural transformation  $\alpha : S \rightarrow S'$  (i.e., of bifunctors) if and only if  $\alpha(b, c)$  is natural in  $b$  for each  $c \in C$  and natural in  $c$  for each  $b \in B$ .*

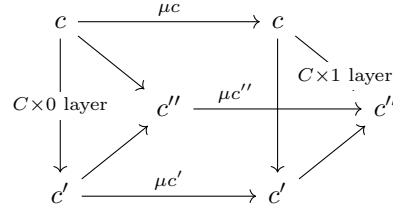
$$\begin{array}{ccc} S(b, c) & \xrightarrow{\alpha(f, g)} & S(b, c) \\ \downarrow S(f, g) & & \downarrow S'(f, g) \\ S(b', c') & \xrightarrow{\alpha(b', c')} & S(b', c') \end{array}$$

### 3.9.1.4 The Universal Natural Transformation

Given any natural transformation  $\tau : S \rightarrow T$  between  $S, T : C \rightarrow B$  there is a unique functor  $F : C \times 2 \rightarrow B$  with  $F\mu c = \tau c$  for any object  $c$ .

- $F\langle f, 0 \rangle = Sf$
- $F\langle f, 1 \rangle = Tf$
- $F\langle f, \downarrow \rangle = Tf \circ \tau c = \tau c' \circ Sf$  (where  $\downarrow : 0 \rightarrow 1$ )

Observe  $C \times 2$  below...

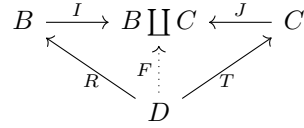


where  $\mu c = \langle c, \downarrow \rangle$

### 3.9.2 Coproducts

Given categories  $B$  and  $C$  the dual of the product category is coproduct category  $B \amalg C$ .

The corresponding universal property is: for any functors  $R$  and  $T$ , there is a unique functor  $F$  making the digram commute...



### 3.9.3 Quotients

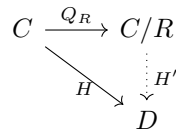
The *quotient category* is specified in the following proposition.

**Proposition 3.9.3.** *For a given category  $C$ , let  $R$  be a function which assigns to each pair of objects  $a, b$  of  $C$  a binary relation  $R_{a,b}$  on the hom-set  $C(a, b)$ . Then there exist a category  $C/R$  and a functor  $Q = Q_R : C \rightarrow C/R$  such that...*

1. *If  $f R_{a,b} f'$  in  $C$ , then  $Qf = Qf'$ .*
2. *If  $H : C \rightarrow D$  is any functor from  $C$  for which  $f R_{a,b} f'$  implies  $Hf = Hf'$  for all  $f$  and  $f'$ , then there is a unique functor  $H' : C/R \rightarrow D$  with  $H' \circ Q_R = H$ .*

Moreover, the functor  $Q_R$  is a bijection on objects.

The corresponding universal property is represented in the following diagram...



### 3.9.3.1 Congruence

A *congruence* is a relation  $R$  on a category  $C$  such that...

- $\forall a, b \in \text{Obj}(C)$ ,  $R_{a,b}$  is an equivalence relation
- if  $f, f' : a \rightarrow b$  have  $f R_{a,b} f'$ , then for all  $g : a' \rightarrow a$  and all  $h : b \rightarrow b'$  one has  $(hfg) R_{a',b'} (hf'g)$ .

### 3.9.4 Free Categories

#### 3.9.4.1 O-graph

The *O-graph* is a graph on a fixed set  $O$  of objects such that  $O$ -graph morphisms  $D_O$  are always equal to  $id_O$ .

We define the *product over O* as a set of composable pairs of arrows...

$$A \times_O B = \{\langle g, f \rangle \mid \delta_0 g = \delta_1 f, g \in A, f \in B\}$$

where  $\delta_0, \delta_1$ , resp., are functions representing the **dom**, **cod**, resp., operations.

A category with objects  $O$  is an  $O$ -graph equipped with two morphisms  $c : A \times_O A \rightarrow A$  and  $i : O \rightarrow A$  of  $O$ -graphs making the following diagrams commutative.

$$\begin{array}{ccc} (A \times_O A) \times_O A & \xrightarrow{\cong} & A \times_O (A \times_O A) \xrightarrow{1 \times c} A \times_O A \\ \downarrow c \times 1 & & \downarrow c \\ A \times_O A & \xrightarrow{\quad c \quad} & A \end{array} \quad \begin{array}{ccc} O \times_O A & \xrightarrow{i \times 1} & A \times_O A \xleftarrow{1 \times i} A \times_O O \\ \downarrow \cong & & \downarrow c \\ A & \xrightarrow{=} & A \xleftarrow{=} A \end{array}$$

#### 3.9.4.2 Free Category

Let  $C(G)$  be the *free category* generated by graph  $G$ , specified in the subsequent theorem...

**Theorem 3.9.1.** *Let  $G = \{A \rightrightarrows O\}$  be a small graph. There is a small category  $C(G)$  with  $O$  as its set of objects and a morphism  $P : G \rightarrow UC$  of graphs from  $G$  to the underlying graph  $UC$  of  $C$  with the following property. Given any category  $B$  and any morphism  $D : G \rightarrow UB$  of graphs, there is a unique functor  $D' : C \rightarrow B$  with  $(UD') \circ P = D$ , as in the commutative diagram*

$$\begin{array}{ccc} C & & G \xrightarrow{P} UC \\ \downarrow D' & & \searrow D \downarrow UD' \\ B & & UB \end{array}$$

In particular, if  $B$  had  $O$  as set of objects and  $D$  is a morphism of  $O$ -graphs, then  $D'$  is the identity on objects.

**Corollary 3.9.1.1.** *To any set  $X$  there is a monoid  $M$  and a function  $p : X \rightarrow UM$ , where  $UM$  is the underlying set of  $M$ , with the following universal property: for any monoid  $L$  and any function  $h : X \rightarrow UL$  there is a unique morphism  $h' : M \rightarrow L$  of monoids with  $h : Uh' \circ p$ .*

$$\text{Hom}_{\text{Cat}}(C(G), B) \cong \text{Grph}(G, UB), \quad D' \mapsto D = UD' \circ P$$

### 3.9.5 Comma Categories

#### 3.9.5.1 Category of objects unber $\mathbf{b}$ ( $b \downarrow C$ )

Objects  $\langle f, c \rangle$ :

$$\begin{array}{c} b \\ \downarrow f \\ c \end{array}$$

Arrows  $\langle f, c \rangle \xrightarrow{h} \langle f', c' \rangle$ :

$$\begin{array}{ccc} & b & \\ f' \swarrow & & \searrow f \\ c & \xrightarrow{h} & c' \end{array}$$

#### 3.9.5.2 Category of objects over $\mathbf{a}$ ( $C \downarrow a$ )

Objects  $\langle f, c \rangle$ :

$$\begin{array}{c} c \\ \downarrow f \\ a \end{array}$$

Arrows  $\langle f, c \rangle \xrightarrow{h} \langle f', c' \rangle$ :

$$\begin{array}{ccc} c & \xrightarrow{h} & c' \\ f \searrow & & \swarrow f' \\ & a & \end{array}$$

#### 3.9.5.3 Category of objects $S$ -unber $\mathbf{b}$ ( $b \downarrow S$ )

Given a functor  $S : D \rightarrow C$ .

Objects  $\langle f, Sd \rangle$ :

$$\begin{array}{c} b \\ \downarrow f \\ Sd \end{array}$$

Arrows  $\langle f, Sd \rangle \xrightarrow{Sh} \langle f', Sd' \rangle$ :

$$\begin{array}{ccc} & b & \\ f' \swarrow & & \searrow f \\ Sd & \xrightarrow{Sh} & Sd' \end{array}$$

#### 3.9.5.4 Category of objects $T$ -over $\mathbf{a}$ ( $T \downarrow \mathbf{a}$ )

Given a functor  $T : E \rightarrow C$ .

Objects  $\langle f, Te \rangle$ :

$$\begin{array}{c} Te \\ \downarrow f \\ \mathbf{a} \end{array}$$

Arrows  $\langle f, Te \rangle \xrightarrow{Th} \langle f', Te' \rangle$ :

$$\begin{array}{ccc} Te & \xrightarrow{Th} & Te' \\ f \searrow & & \swarrow f' \\ & \mathbf{a} & \end{array}$$

#### 3.9.5.5 Comma Category ( $T \downarrow S$ )

Given functors  $S : D \rightarrow C$  and  $T : E \rightarrow C$ .

Objects  $\langle e, d, f \rangle$ :

$$\begin{array}{c} Te \\ \downarrow f \\ Sd \end{array}$$

where  $d \in \text{Obj}(D)$ ,  $e \in \text{Obj}(E)$ ,  $f : Te \rightarrow Sd$ .

Arrows  $\langle e, d, f \rangle \xrightarrow{\langle k, h \rangle} \langle e', d', f' \rangle$ :

$$\begin{array}{ccc}
Te & \xrightarrow{Tk} & Te' \\
\downarrow f & & \downarrow f' \\
Sd & \xrightarrow{Sh} & Sd'
\end{array}$$

where  $k : e \rightarrow e'$ ,  $h : d \rightarrow d'$  such that  $f' \circ Tk = Sh \circ f$ .

Composition  $\langle k', h' \rangle \circ \langle k, h \rangle = \langle k' \circ k, h' \circ h \rangle$  when defined.

$$\begin{array}{ccccccc}
& & T \downarrow S & & & & \\
& P \swarrow & \downarrow Q & \searrow R & & & \\
E & \xleftarrow{T} & C & \xleftarrow{C^{d_0}} & C^2 & \xrightarrow{C^{d_1}} & C \xleftarrow{S} D
\end{array}$$

$P$  and  $Q$  are the *projections* of the comma category.  $C^{d_0}, C^{d_1}$ , resp., send arrows to domain, codomain, resp.

$$\begin{array}{ccccc}
& \langle e, d, f : Te \rightarrow Sd \rangle & & & \\
& \swarrow & \downarrow & \searrow & \\
e & \hookrightarrow Te & \xleftarrow{\quad} (f : Te \rightarrow Sd) & \xrightarrow{\quad} Sd & \hookleftarrow d
\end{array}$$

### 3.10 Higher Level Categories

#### 3.10.1 Functor Categories

A *functor category* is a category whose objects are functors and whose arrows are natural transformations. Since compositions of natural transformations are natural transformations, composition can be defined as in the following diagram...

$$\begin{array}{ccc}
Rc & \xrightarrow{Rf} & Rc' \\
\downarrow \sigma_c & & \downarrow \sigma_{c'} \\
(\tau \circ \sigma)_c \left( Sc \right. & \xrightarrow{Sf} & \left. Sc' \right) (\tau \circ \sigma)_{c'} \\
\downarrow \tau_c & & \downarrow \tau_{c'} \\
Tc & \xrightarrow{Tf} & Tc'
\end{array}$$

### 3.10.2 2-Categories

#### 3.10.2.1 Vertical Composition

For natural transformations  $\tau$  and  $\sigma$ , we have "vertical" composition  $\tau \circ \sigma$ , as in the following diagram...

$$\begin{array}{ccc}
 C & \xrightarrow{\quad} & B \\
 \downarrow \sigma & & \downarrow \sigma \\
 C & \xrightarrow{\quad} & B \\
 \downarrow \tau & & \downarrow \tau \\
 C & \xrightarrow{\quad} & B
 \end{array}
 \begin{array}{c}
 \tau \sigma \\
 \tau \sigma
 \end{array}$$

#### 3.10.2.2 Horizontal Composition

We can also define "horizontal" composition for natural transformations  $\tau$  and  $\tau'$ ,  $\tau' \circ \tau$ , as in the following commutative diagrams...

$$\begin{array}{ccccc}
 C & \xrightarrow{S} & B & \xrightarrow{S'} & A \\
 \downarrow \tau & & \downarrow \tau & \downarrow \tau' & \downarrow \tau' \\
 C & \xrightarrow{T} & B & \xrightarrow{T'} & A
 \end{array}$$

$$\begin{array}{ccc}
 S'Sc & \xrightarrow{\tau'Sc} & T'Sc \\
 \downarrow S'\tau c & \searrow (\tau' \circ \tau)c & \downarrow T'\tau c \\
 S'Tc & \xrightarrow{\tau'Tc} & T'Tc
 \end{array}$$

The next diagram shows  $\tau' \circ \tau : S'S \rightarrow T'T$  is natural.

$$\begin{array}{ccccccc}
 c & S'Sc & \xrightarrow{S'\tau c} & S'Tc & \xrightarrow{\tau'Tc} & T'Tc & \\
 \downarrow f & \downarrow S'Sf & & \downarrow S'Tf & & \downarrow T'Tf & \\
 b & S'Sb & \xrightarrow{T'\tau b} & S'Tb & \xrightarrow{\tau'Tb} & T'Tb &
 \end{array}$$



So  $\tau' \circ \tau = (T' \circ \tau) \cdot (\tau' \circ S) = (\tau' \circ T) \cdot (S' \circ \tau)$ , which leads into our next concept.

### 3.10.2.3 Interchange Law

For natural transformations  $\sigma, \sigma', \tau, \tau'$  satisfying...

$$\begin{array}{ccccc}
 C & \longrightarrow & B & \longrightarrow & A \\
 \downarrow \sigma & & \sigma \left( \downarrow \right) \sigma' & & \downarrow \sigma' \\
 C & \longrightarrow & B & \longrightarrow & A \\
 \downarrow \tau & & \tau \left( \downarrow \right) \tau' & & \downarrow \tau' \\
 C & \longrightarrow & B & \longrightarrow & A
 \end{array}$$

the *interchange law* is  $(\tau' \cdot \sigma') \circ (\tau \cdot \sigma) = (\tau' \circ \tau) \cdot (\sigma' \circ \sigma)$ .

The proof of the interchange law derives from the following diagram. Intuitively, the interchange law occurs along the dotted diagonal lines.

$$\begin{array}{ccccc}
 S'Sc & \xrightarrow{\sigma'S} & T'Sc & \xrightarrow{\tau'S} & R'Sc \\
 \downarrow S'\sigma & \nearrow \text{dotted} & \downarrow T'\sigma & \nearrow \text{dotted} & \downarrow R'\sigma' \\
 S'Tc & \xrightarrow{\sigma'T} & T'Tc & \xrightarrow{\tau'T} & R'Tc \\
 \downarrow S'\tau & \nearrow \text{dotted} & \downarrow T'\tau & \nearrow \text{dotted} & \downarrow R'\tau' \\
 S'Rc & \xrightarrow{\sigma'R} & T'Rc & \xrightarrow{\tau'R} & R'Rc
 \end{array}$$

**Theorem 3.10.1.** *The collection of natural transformations in the set of arrows of two different categories under two different operations of composition,  $\cdot$  and  $\circ$ , which satisfy the interchange law. Moreover, any arrow (transformation) which is an identity for the composition  $\circ$  is also an identity for the composition  $\cdot$ .*

### 3.10.2.4 Double Category

The set of arrows for two different compositions with two different compositions which together satisfy the interchange law.

### 3.10.2.5 2-Category

A double category in which every identity arrow for the first composition is also an identity for the second composition.

### 3.11 Universal Properties

## 4 Group Theory

## 5 Abelian Group Theory

## 6 Ring Theory

## 7 Modules

## 8 Topology

### 8.1 Metric Spaces

A *metric space* is a set  $X$  together with a *metric*  $d : X \times X \rightarrow \mathbb{R}$  satisfying...

1.  $d(x, y) \geq 0$  for all  $x, y \in X$  and  $d(x, y) = 0$  if and only if  $x = y$ .
2.  $d(x, y) = d(y, x)$  for all  $x, y \in X$ .
3. (*The Triangle Inequality*):  $d(x, y) + d(y, z) \geq d(x, z)$  for all  $x, y, z \in X$ .

#### 8.1.1 Open Ball

The *open ball* of radius  $\varepsilon > 0$  centered at a point  $x$  in a metric space  $(X, d)$  is given by...

$$B_\varepsilon(x) = \{y \in X \mid d(x, y) < \varepsilon\}.$$

#### 8.1.2 Continuity

Suppose  $(X, d_X)$  and  $(Y, d_Y)$  are two metric spaces and  $f : X \rightarrow Y$  is a function. Then  $f$  is *continuous at*  $x \in X$  if for any  $\epsilon > 0$ , there is a  $\delta > 0$  so that  $B_\delta(x) \subset f^{-1}(B_\epsilon(f(x)))$ .

The function  $f$  is *continuous* if it is continuous at  $x$  for all  $x \in X$ .

#### 8.1.3 Open Set

The *open set*  $U$  of a metric space  $(X, d)$  is *open* if for any  $u \in U$ , there is  $\varepsilon > 0$  so that  $B_\varepsilon(u) \subseteq U$ .

**Theorem 8.1.1.** *A function  $f : X \rightarrow Y$  between metric spaces  $(X, d)$  and  $(Y, d)$  is continuous if and only if for any open subset  $V$  of  $Y$ , the subset  $f^{-1}(V)$  is open in  $X$ .*

#### 8.1.4 Examples

### 8.2 Topological Spaces

## 9 Homotopy



## 10 Homology

## 11 Dimension

**Theorem 11.0.1.** *There is a one-to-one correspondence  $\mathbb{R} \rightarrow \mathbb{R} \times \mathbb{R}$ .*

*Proof.* Since the mapping  $f : \mathbb{R} \rightarrow (0, 1)$  given by  $r \mapsto \frac{1}{\pi}(\arctan(r) + \frac{\pi}{2})$  is a bijection, it is sufficient to find a bijection  $(0, 1) \rightarrow (0, 1) \times (0, 1)$ . For this we use the Schröder-Bernstein Theorem.

Observe that  $g : (0, 1) \rightarrow (0, 1) \times (0, 1)$  given by  $t \mapsto (t, t)$  is an injection. So the only real work is in constructing an injection  $(0, 1) \times (0, 1) \rightarrow (0, 1)$ .

To start things off, introduce the following injection  $I : (0, 1) \rightarrow (0, 1) \cap (\mathbb{R} \setminus \mathbb{Q}) \dots$

$$I(r) = \begin{cases} [0; a_1 + 2, a_2 + 2, \dots, a_n + 2, 2, 2, \dots] & \text{if } r = [0; a_1, a_2, \dots, a_n], \\ [0; a_1 + 2, a_2 + 2, a_3 + 2, \dots] & \text{if } r = [0; a_1, a_2, a_3, \dots]. \end{cases}$$

Composed with another injection,  $t : (0, 1) \cap (\mathbb{R} \setminus \mathbb{Q}) \times (0, 1) \cap (\mathbb{R} \setminus \mathbb{Q}) \rightarrow (0, 1)$  given by  $([0; a_1, a_2, \dots], [0; a_1, a_2, \dots]) \mapsto [0; a_1, b_1, a_2, b_2, \dots]$ , we get our desired injection  $t \circ (I \times I)$ .  $\square$

**Corollary 11.0.1.1.** *There is a bijection  $f : \mathbb{R}^m \rightarrow \mathbb{R}^n$  for all positive integers  $m$  and  $n$ .*