

A* Shortest Path Search

Course 4, Module 3, Lesson 3

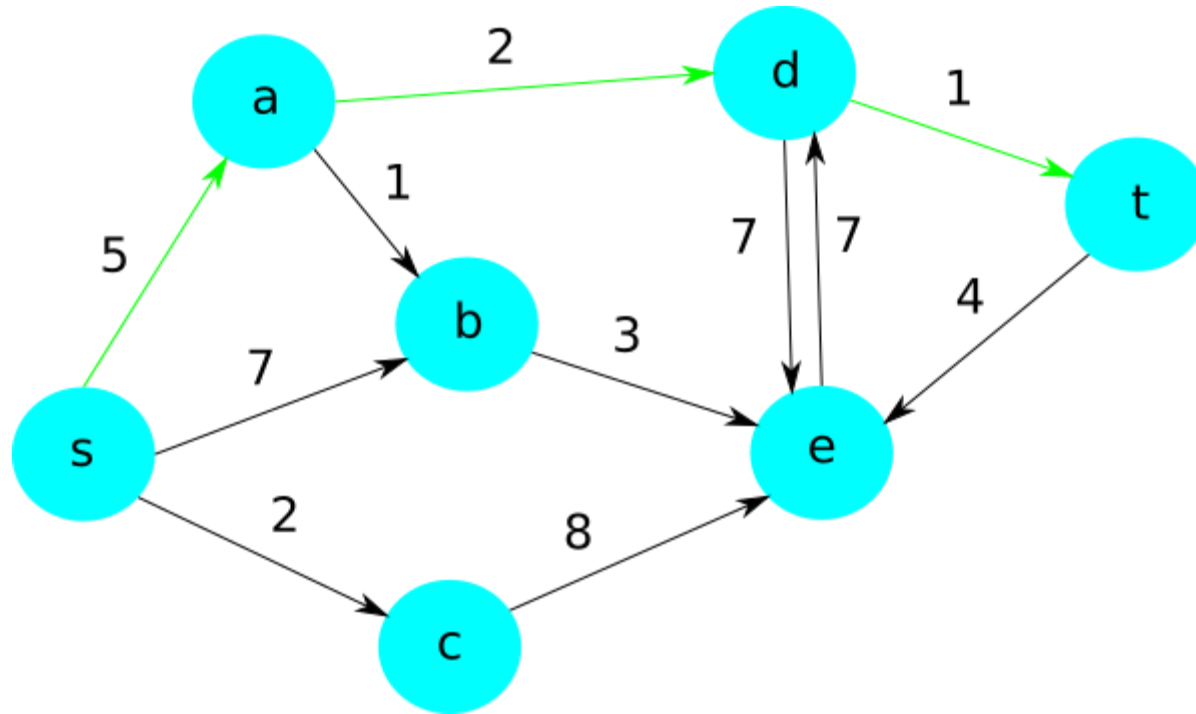


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Learning Objectives

- Understand what admissible heuristics are in the context of graph search
- Understand how to use the Euclidean heuristic to improve our mission planning speed in practice
- Implement the A^* search algorithm, leveraging the Euclidean heuristic
- Understand how to apply A^* search to variants on the mission planning problem involving time instead of distance

Recall: Dijkstra's for Weighted Graph



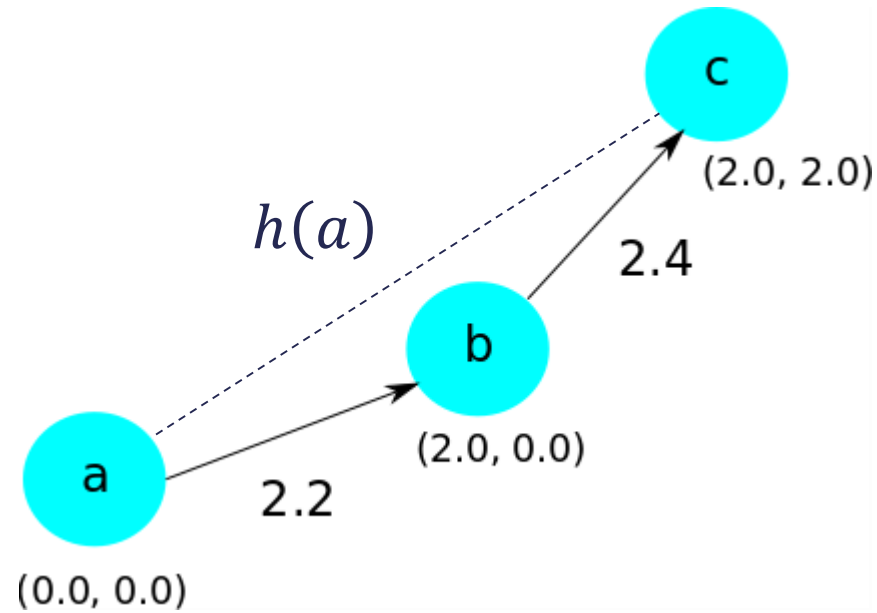
Euclidean Heuristic

- Exploits structure of the problem
- Fast to calculate
- Straight-line distance between two vertices is a useful estimate of true distance along the graph

$$h(v) = \|t - v\|$$

Euclidean Heuristic - Example

$$h(a) = \sqrt{2^2 + 2^2} = 2.828$$



A* Algorithm

A* Algorithm

Algorithm A*(G,s,t)

```
1.  open ← MinHeap()
2.  closed ← Set()
3.  predecessors ← Dict()
4.  open.push(s, 0)
5.  while ! open.isEmpty() do
6.       $u, uCost \leftarrow \text{open.pop}()$ 
7.      if isGoal( $u$ ) then
8.          return extractPath( $u$ , predecessors)
9.      for all  $v \in u.\text{successors}()$ 
10.         if  $v \in \text{closed}$  then
11.             continue
12.          $uvCost \leftarrow \text{edgeCost}(G, u, v)$ 
13.         if  $v \in \text{open}$  then
14.             if  $uCost + uvCost + h(v) < \text{open}[v]$  then
15.                  $\text{open}[v] \leftarrow uCost + uvCost + h(v)$ 
16.                  $\text{costs}[v] \leftarrow uCost + uvCost$ 
17.                  $\text{predecessors}[v] \leftarrow u$ 
18.         else
19.              $\text{open.push}(v, uCost + uvCost)$ 
20.              $\text{costs}[v] \leftarrow uCost + uvCost$ 
21.              $\text{predecessors}[v] \leftarrow u$ 
22.     closed.add( $u$ )
```

A* Algorithm

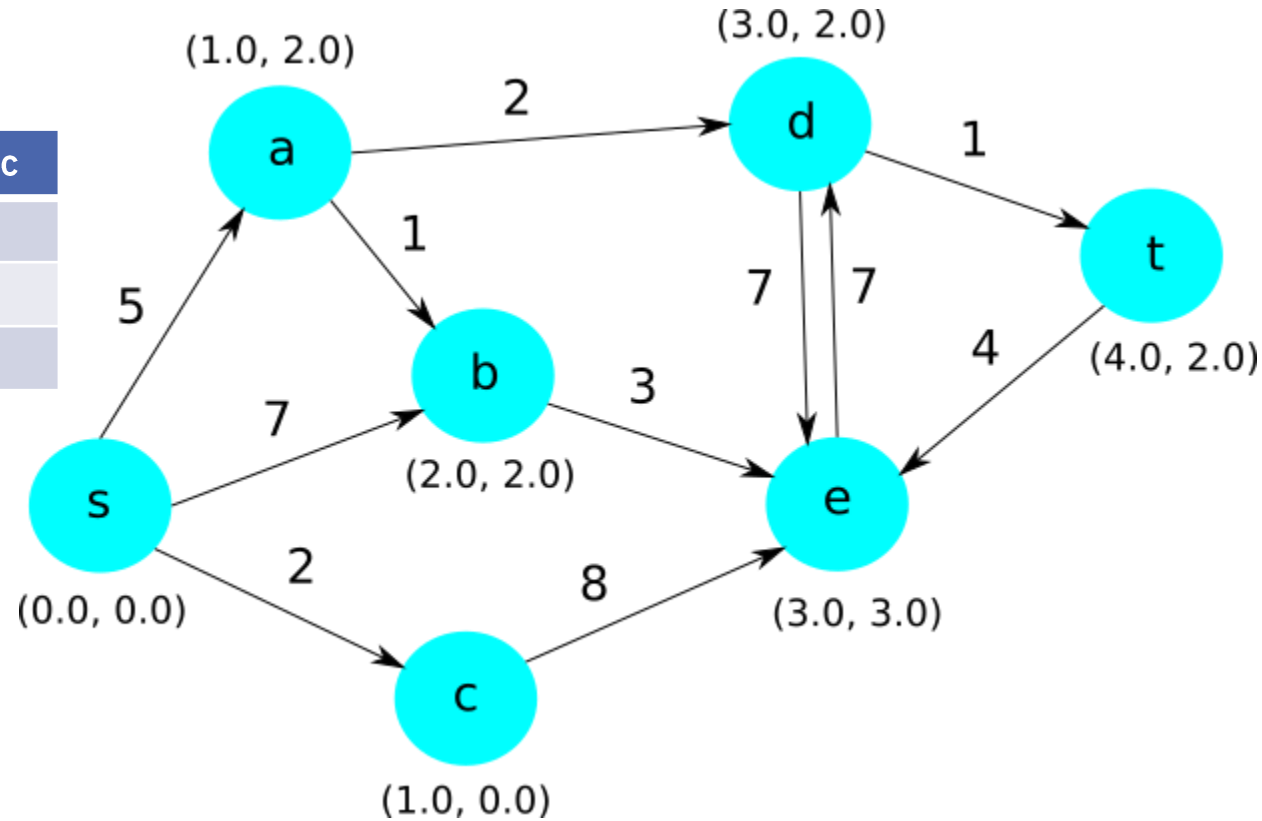
```
1.  if  $v \in \text{open}$  then
2.      if  $u\text{Cost} + uv\text{Cost} + h(v) < \text{open}[v]$  then
3.           $\text{open}[v] \leftarrow u\text{Cost} + uv\text{Cost} + h(v)$ 
4.           $\text{costs}[v] \leftarrow u\text{Cost} + uv\text{Cost}$ 
5.           $\text{predecessors}[v] \leftarrow u$ 
6.      else
7.           $\text{open.push}(v, u\text{Cost} + uv\text{Cost})$ 
8.           $\text{costs}[v] \leftarrow u\text{Cost} + uv\text{Cost}$ 
9.           $\text{predecessors}[v] \leftarrow u$ 
```


Example - Origin Node

Open Min Heap:

Node	Cost + Heuristic
s	4.472

Closed Set:

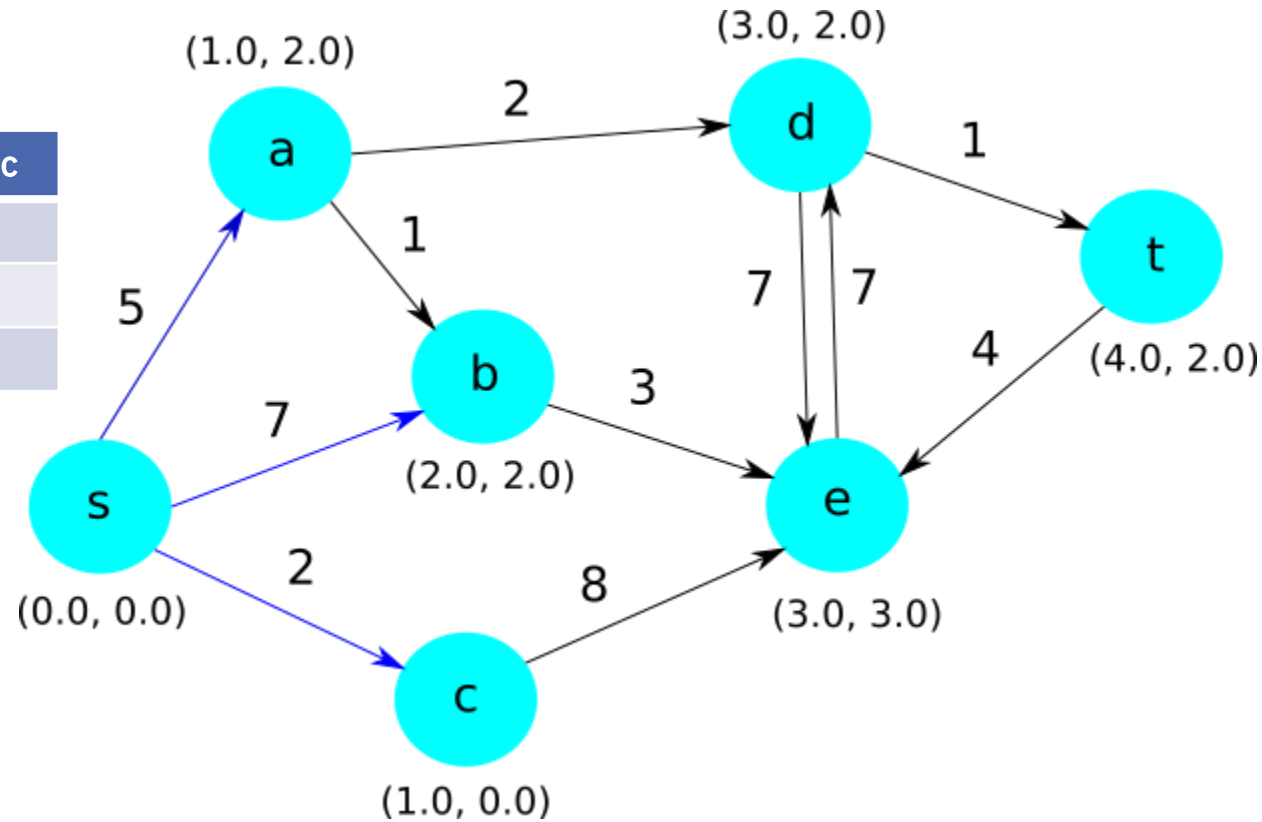


Example - Processing s

Open Min Heap:

Node	Cost + Heuristic
c	5.606
a	8
b	9

Closed Set: s

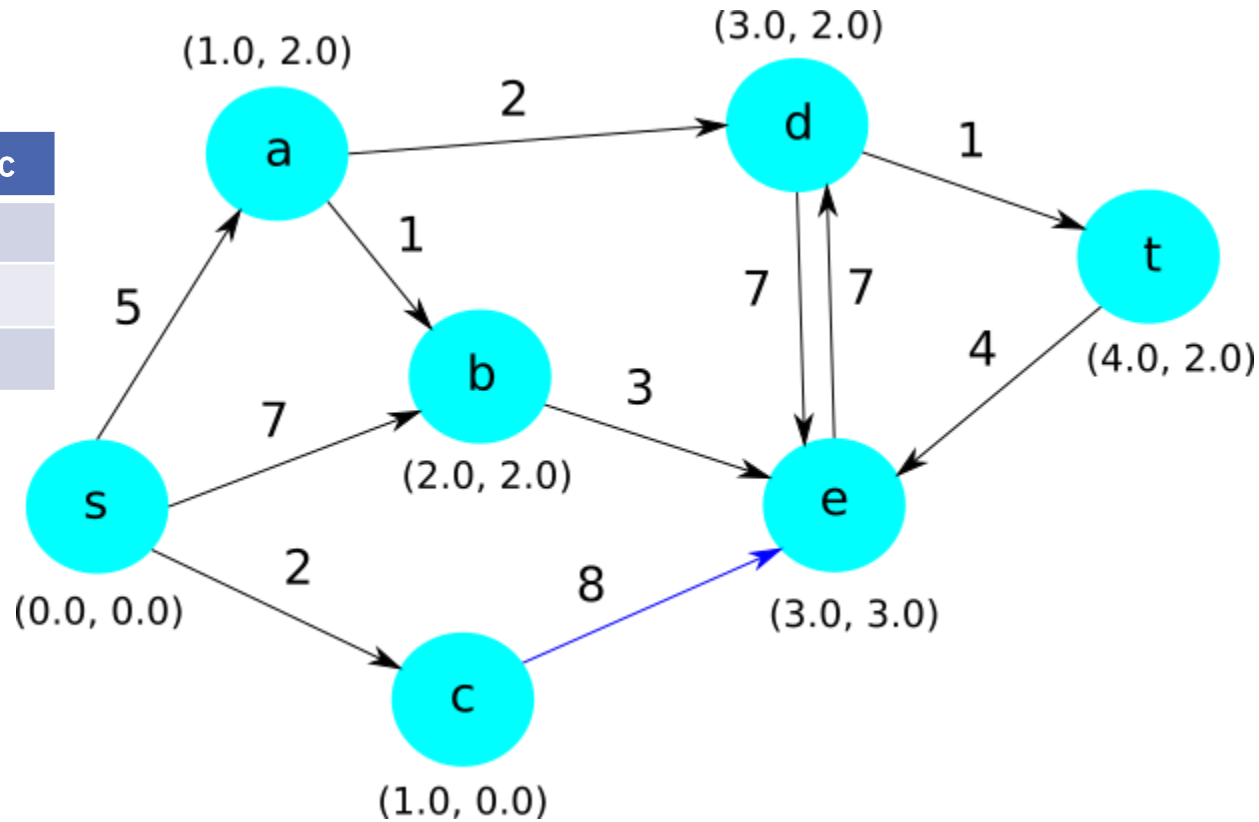


Example - Processing c

Open Min Heap:

Node	Cost + Heuristic
a	8
b	9
e	11.414

Closed Set: s
c

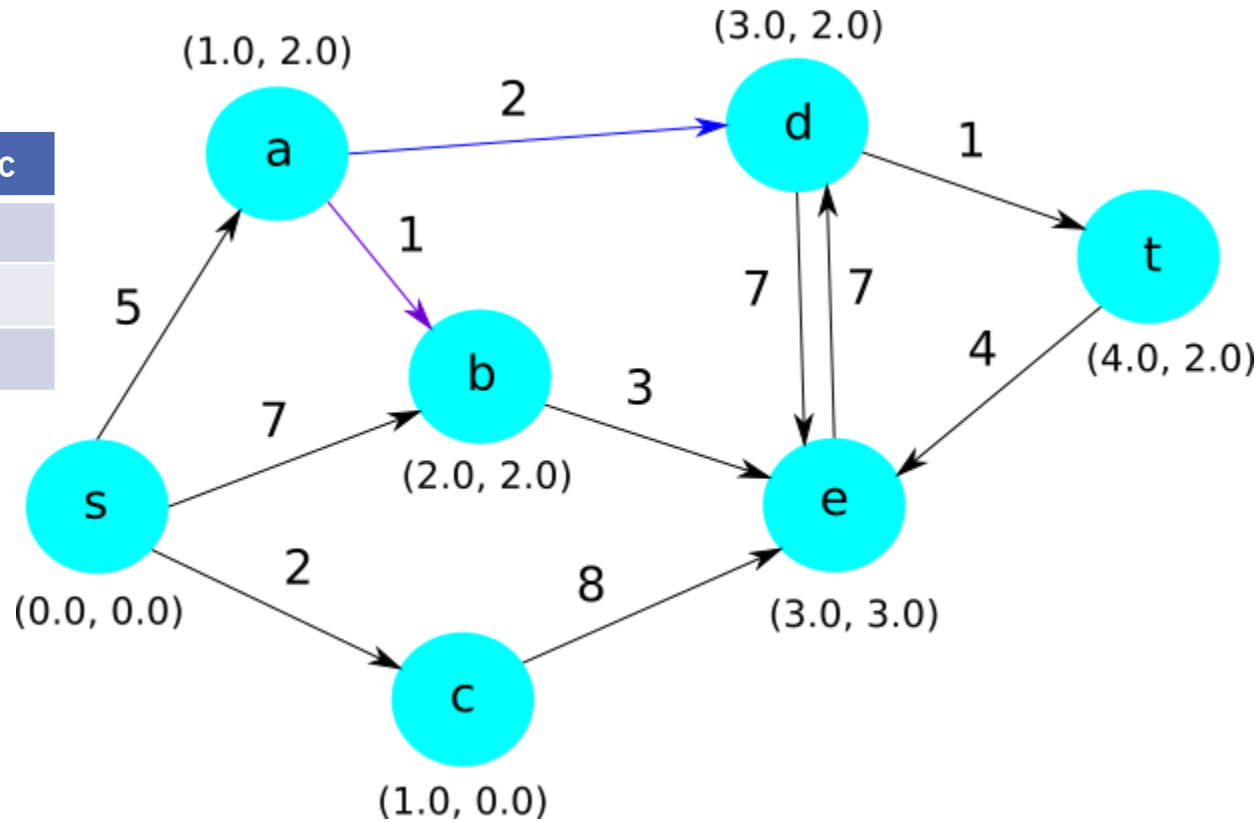


Example - Processing a

Open Min Heap:

Node	Cost + Heuristic
d	7
b	8
e	11.414

Closed Set: s
c
a

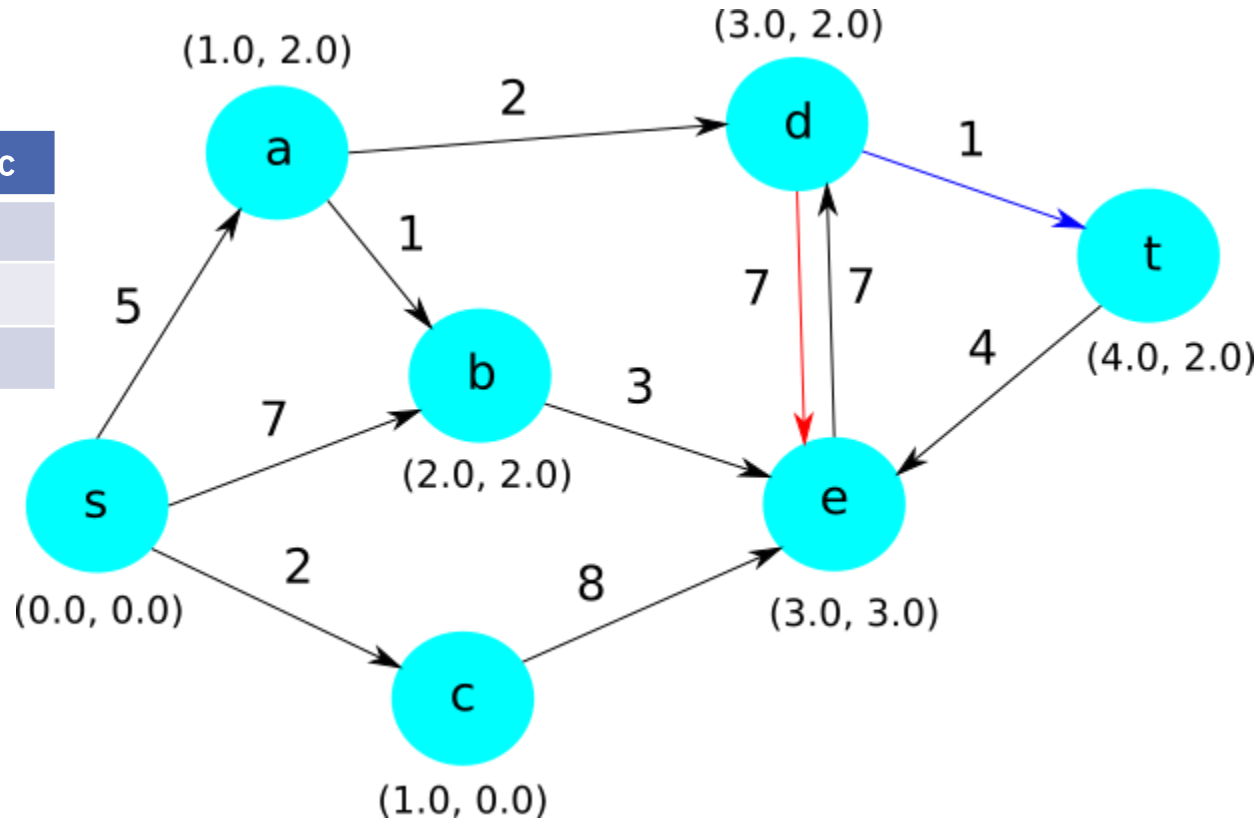


Example - Processing d

Open Min Heap:

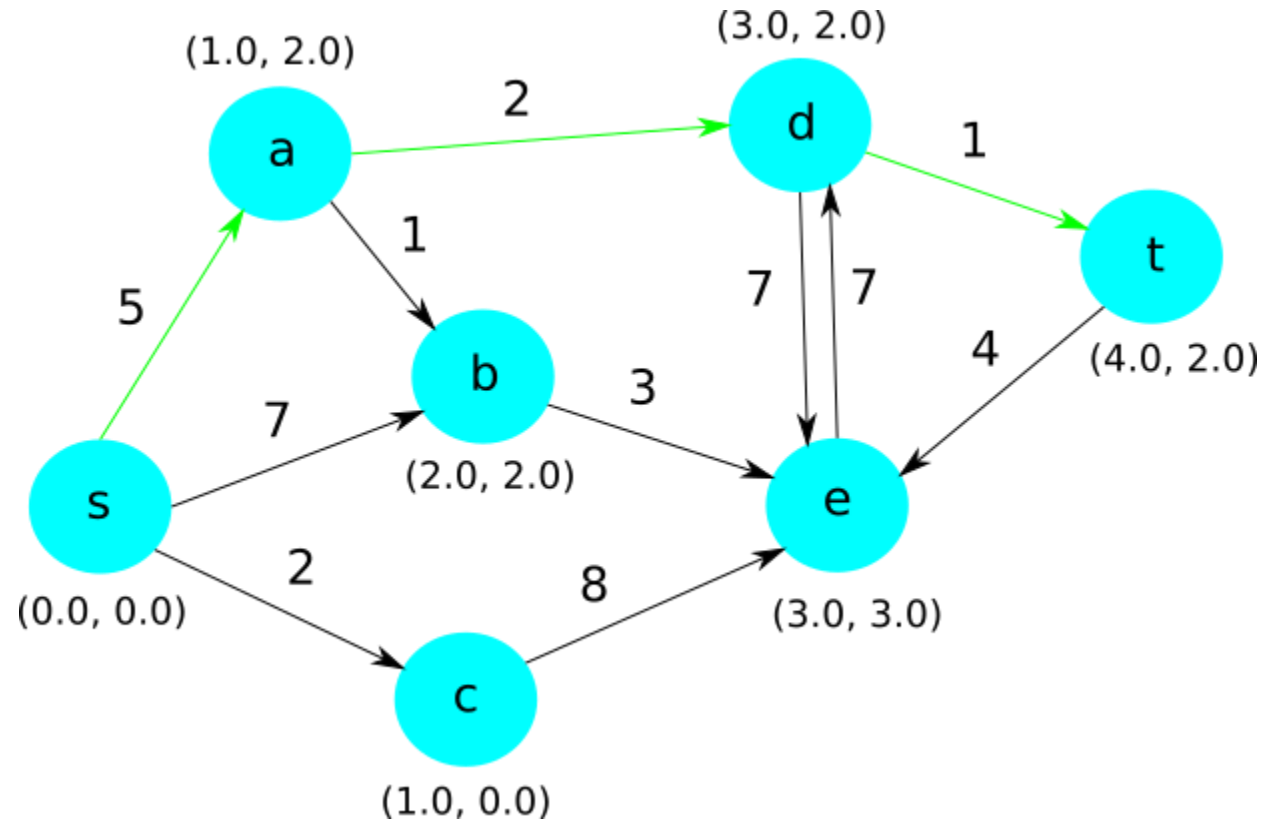
Node	Cost + Heuristic
t	7
b	8
e	11.414

Closed Set: s
c
a
d



Example - Final Path

Final Path: s
a
d
t

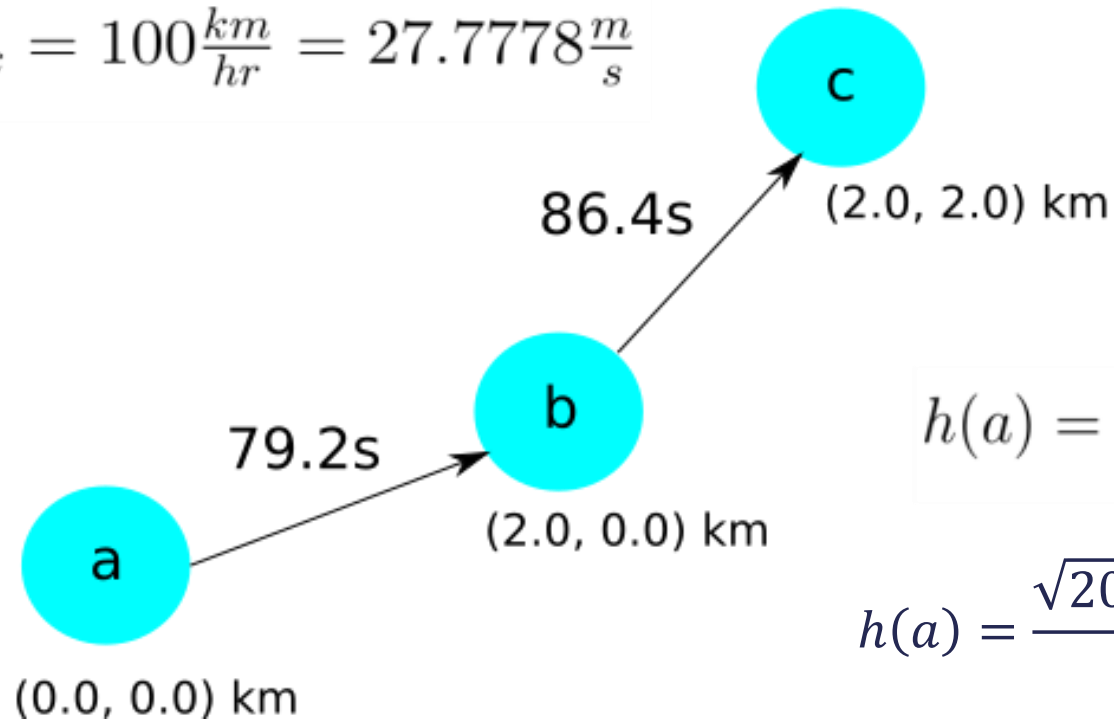


Extensions to Other Factors

- Traffic, speed limits, and weather affect mission planning
- Time rather than distance is better at capturing these factors
- Replace distance edge weights with time estimates

Example

$$v_{max} = 100 \frac{km}{hr} = 27.7778 \frac{m}{s}$$

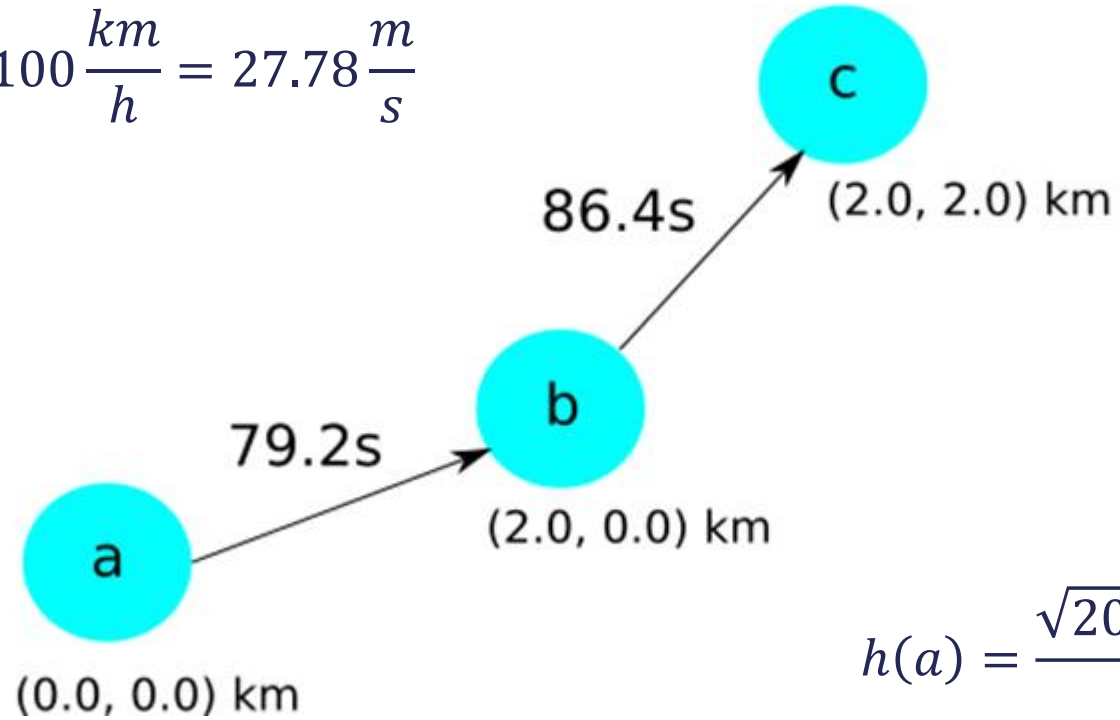


$$h(a) = \frac{\sqrt{2^2 + 2^2}}{v_{max}} = 101.82s$$

$$h(a) = \frac{\sqrt{2000^2 + 2000^2}}{v_{max}} = 101.82s$$

Example

$$v_{max} = 100 \frac{km}{h} = 27.78 \frac{m}{s}$$



$$h(a) = \frac{\sqrt{2000^2 + 2000^2}}{v_{max}} = 101.82s$$

Summary

- Introduced Euclidean heuristic, showed it was admissible to our mission planning problem
- Walked through the A^* search algorithm
- Discussed how to modify the heuristic to handle travel time rather than distance in our search



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