

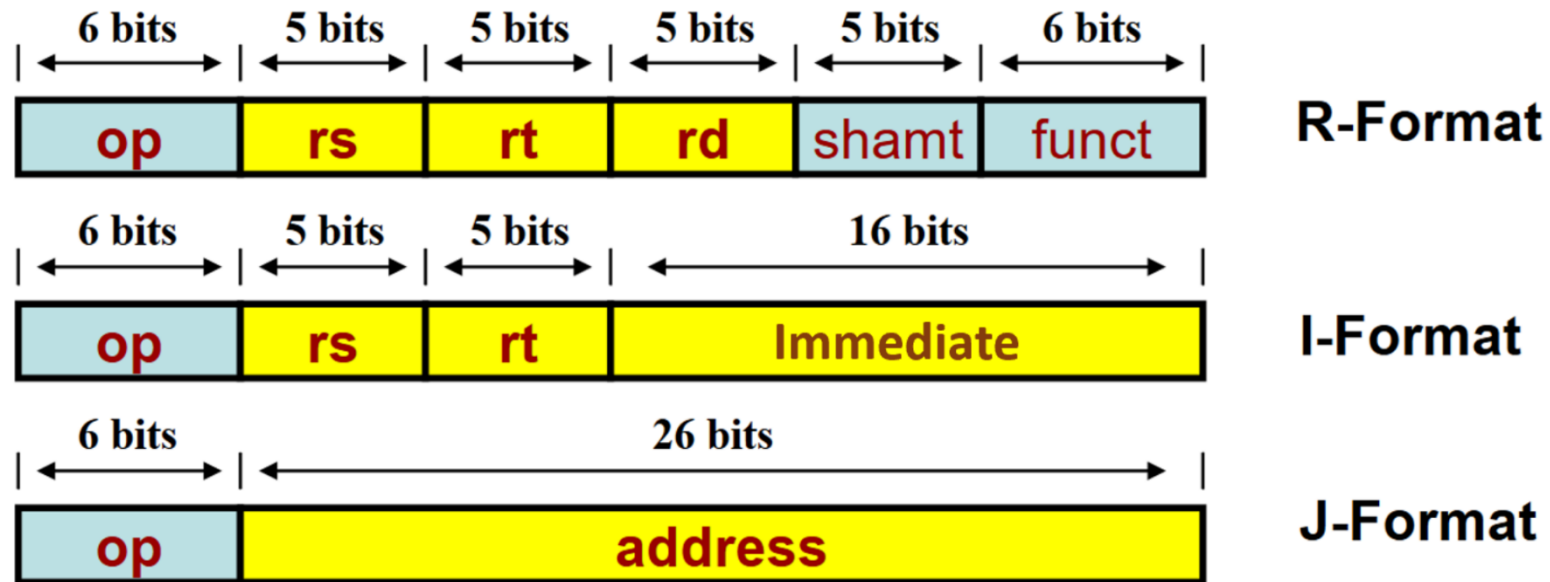
TUTORIAL 7 MIPS MACHINE CODE AND PROCEDURE

Overview

- **You will review the following in this tutorial:**
 - MIPS machine code
 - Sign extension and zero extension
- **You will practice to write simple MIPS procedures**

MIPS Instruction Format

- MIPS has three instruction formats. All uses 32 bits.



Warm up exercise

■ Write down the corresponding MIPS machine code (in binary) of the instructions

a) add \$s0, \$s1, \$t2

a) 000000 10001 01010 10000 00000 100000

b) opcode rs rt rd shamt funct

b) lw \$s0, 16(\$t0)

a) 100011 01000 10000 0000 0000 0001 0000

b) opcode rs rt immediate

MIPS Assembly Examples: Symbolic

```
1  .text # text segment
2  .globl _main
3  _main:
4  # it's not a valid code sequence
5  # R-format
6  add $s2, $s1, $s0
7  sub $s2, $s1, $s0
8  or $s2, $s1, $s0
9  sll $s1, $s0, 2
10 # I-format
11 lw $s1, 16($s0)
12 lb $s1, 16($s0)
13 addi $s2, $s1, -1
14 beq $s1, $s0, exit
15 # J-format
16 j exit
17 exit:
18 jr $ra # is it really J-type?
```

Corresponding MIPS Machine Code

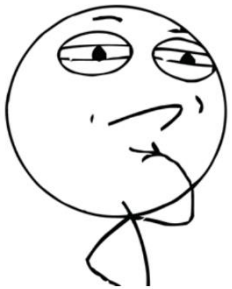
Text Segment				
Bkpt	Address	Code	Basic	Source
<input type="checkbox"/>	0x00400000	0x02309020	add \$18,\$17,\$16	7: add \$s2, \$s1, \$s0
<input type="checkbox"/>	0x00400004	0x02309022	sub \$18,\$17,\$16	8: sub \$s2, \$s1, \$s0
<input type="checkbox"/>	0x00400008	0x02309025	or \$18,\$17,\$16	9: or \$s2, \$s1, \$s0
<input type="checkbox"/>	0x0040000c	0x00108880	sll \$17,\$16,0x00000002	10: sll \$s1, \$s0, 2
<input type="checkbox"/>	0x00400010	0x8e110010	lw \$17,0x00000010(\$16)	12: lw \$s1, 16(\$s0)
<input type="checkbox"/>	0x00400014	0x82110010	lb \$17,0x00000010(\$16)	13: lb \$s1, 16(\$s0)
<input type="checkbox"/>	0x00400018	0x2232ffff	addi \$18,\$17,0xffff...	14: addi \$s2, \$s1, -1
<input type="checkbox"/>	0x0040001c	0x12300001	beq \$17,\$16,0x00000001	15: beq \$s1, \$s0, exit
<input type="checkbox"/>	0x00400020	0x08100009	j 0x00400024	17: j exit
<input type="checkbox"/>	0x00400024	0x03e00008	jr \$31	19: jr \$ra

Zero Extension or Sign Extension

■ Extension is needed when there is a size mis-match

- Source and destination in assignment
- Two operands of a calculation

Type	Example	Extension
Arithmetic Instructions	addi/addiu	Always sign-extend
Load/store Instructions	lb	Sign-extend
	lbu	Zero-extend
Logical instructions	ori, andi	Always zero-extend



Zero Extension

Assembly Code: `ori $t8, $0, 0x4`

Binary: 0011010000011000000000000000100

Hex: 0x34180004

Immediate: -2^{15} to $+2^{15} - 1$

31	26 25	21 20	16 15	0
ORI	zero	t8	immediate	
001101	00000	11000	0000000000000100	
6	5	5	16	

- ❑ **ori** is I format
- ❑ 0x4 is stored in the instruction with 16 bits 0000 0000 0000 0100
- ❑ It is to be bitwise ORed with the thirty-two bits of register zero, 0000 0000 0000 0000 0000 0000 0000 0000
- ❑ This would not ordinarily be possible because the operands are different lengths. However, MIPS zero extends the sixteen-bit operand so the operands are the same length.
- ❑ Sometimes this is also called **padding with zeros**.

0000 0000 0000 0000 0000 0000 0000 0100 -- zero extended immediate operand

0000 0000 0000 0000 0000 0000 0000 0000 -- data in register \$0

0000 0000 0000 0000 0000 0000 0000 0100 -- result, put in register \$t8

Sign Extension for Signed Numbers

- We often need to represent a value given in a certain number of bits by using a larger number of bits.
- You can **repeat the sign bit** of the number as many times as it is needed to the left.
- **Sign extension preserves the original value!!**

- **Examples:**

- Positive Number:

- 0111 → 0000 0000 0000 0111

- Negative Number:

- 1010 → 1111 1111 1111 1010

Example: Sign/Zero Extension

```
Edit Execute
proc_sum procedure_print.asm procedure_cuboid machinecode.asm extension.asm
1 .data
2 bytearray: .byte 0x7f 0xff
3 wordarray: .word 0xffffffff
4 .text # text segment
5 .globl _main
6 _main:
7 # it's not a valid code sequence
8 ori $s0, $0, 0xFF41 # zero extension
9 addi $s1, $0, -1 # sign extension
10
11 la $s2, bytearray
12 lb $s3, 0($s2) # sign extension
13 lbu $s4, 0($s2) # zero extension
14 lb $s5, 1($s2) # sign extension
15 lbu $s6, 1($s2) # zero extension
16
17 la $s7, wordarray
18 li $t0, 0x1234ABCD # pseudo instruction
19 sb $t0, 0($s7) # store the rightmost byte to memory
20
```

Exercise 1

■ Write down the shortest sequence (any one) of MIPS instructions for the following C++ code, assuming variable `a` and `b` are stored in `$s0` and `$s1` respectively. You can use some registers for storing temporary values.

```
b = a + 0x37cf0010;
```

Exercise 2

■ Write down the shortest sequence (any one) of MIPS instructions for the following C++ code, assuming variable `a` and `b` are stored in `$s0` and `$s1` respectively. You can use some registers for storing temporary values.

```
b = a + 0x37cff346;
```

Register Convention for Procedures

■ General purpose registers for procedure calling:

- **\$a0 – \$a3**: arguments (reg's 4 – 7)
- **\$v0, \$v1**: result values (reg's 2 and 3)
- **\$t0 – \$t9**: temporaries
 - Can be overwritten by callee
- **\$s0 – \$s7**: saved
 - Must be saved/restored by callee
- **\$gp**: global pointer for static data (reg 28)
- **\$sp**: stack pointer (reg 29)
- **\$fp**: frame pointer (reg 30)
- **\$ra**: return address (reg 31)

■ Program counter (PC) or instruction address register:

- Register that holds address of the current instruction being executed. It is **updated** after executing the current instruction.
- $PC = PC + 4$ or $PC = \text{branch target address}$

Caller and Callee Coordination

■ The calling program (**caller**)

□ Passing parameters:

- Puts the parameter values in **\$a0 - \$a3**
- Invokes **jal X** to jump to procedure X

■ Procedure X (**callee**)

- Performs the calculations
- To return the results, place the results in **\$v0 - \$v1**
- Returns control to the caller using **jr \$ra**
- Caller picks up the result from **\$v0 - \$v1**

Example 1: Print

```
Edit  Execute
proc_sum  procedure_print.asm  procedure_cuboid  machinecode.asm  extension.asm
1  #procedure print
2  .data #data segment
3  message1: .asciiz "\nHello World!\n"
4  message2: .asciiz "\nPrinted by procedure.\n"
5
6  .text #text segment
7  .globl _main
8  _main: # caller
9      la $a0, message1 # caller prepares arguments in $a
10     jal print # print(message1)
11     la $a0, message2
12     jal print # print(message2)
13     li $v0, 10
14     syscall
15 print: # callee void print(char* str)
16     li $v0, 4
17     syscall
18     jr $ra # return
19
```

Exercise 3A

■ Translate the following C++ function into a MIPS function, using the registers `$a0` and `$a1` for its parameters and the register `$v0` for its return value.

```
int equal (int p1, int p2) {  
    if (p1 == p2)  
        return 1;  
    return 0;  
}
```


Exercise 3B

- Write down the MIPS code segment that make the following call to the C++ function in the previous exercise, assuming the variable **b** is stored in the register **\$s0**.

```
int b = equal (3, 4);
```

Example 2: Integer Comparison

```

Edit  Execute
proc_sum  procedure_print.asm  procedure_cuboid  machinecode.asm  extension.asm  procedure_cmp.asm
1  #procedure comparison
2  .data #data segment
3  equal_msg: .asciiz "\nTwo numbers are equal!\n"
4  not_equal_msg: .asciiz "\nTwo numbers are not equal.\n"
5  .text #text segment
6  .globl _main
7  _main: # caller
8      addi $a0, $0, 11 # caller prepares arguments in $a
9      addi $a1, $0, 11
10     jal cmp # cmp($a0, $a1)
11     li $v0, 10
12     syscall #exit
13  cmp:  # callee void cmp(int a, int b)
14     bne $a0, $a1, not_equal
15     la $a0, equal_msg # no need to preserve $a registers
16     j cmp_exit
17  not_equal:
18     la $a0, not_equal_msg
19  cmp_exit:
20     li $v0, 4 # print msg
21     syscall
22     jr $ra # return
23
```

Example 3: Multiplication

```
Edit  Execute
procedure_mult.asm  procedure_sum.asm  procedure_print.asm  procedure_cuboid.asm  procedure_cmp.asm

1  #procedure multiplication
2  .text #text segment
3  .globl _main
4  _main: # caller
5      addi $a0, $0, 2 # calculate 2x6
6      addi $a1, $0, 6 # caller prepares arguments in $a
7      jal mult_with_add # mult_with_add($a0, $a1)
8      add $s0, $0, $v0 # grab return value in $v
9      li $v0, 10 # exit main
10     syscall
11 mult_with_add: # callee int mult_with_add(int a, int b)
12     add $v0, $0, $0 # init $v0=0
13 loop:
14     beq $a1, $0, loop_done
15     add $v0, $v0, $a0 # add $a0 for $a1 times
16     addi $a1, $a1, -1 # $a1--
17     j loop
18 loop_done:
19     jr $ra # $v0 = $a0x$a1, return back to caller
```



Example 4: Integer Cuboid

```

Edit  Execute
proc_sum  procedure_print.asm  procedure_cuboid*  procedure_cmp.asm  procedure_mult.asm
1  # procedure cuboid
2  # use procedure mult_with_add
3  .data
4  sides: .word 2 3 4 # sides of the cuboid
5  result: .word 0 # volume of the cuboid
6  .text #text segment
7  .globl _main
8  _main: # caller
9      la $s0, sides
10     la $s1, result
11     lw $a0, 0($s0) # side a
12     lw $a1, 4($s0) # side b
13     jal mult_with_add # mult_with_add($a0, $a1)
14     add $a0, $v0, $0 # $a0 = a x b
15     lw $a1, 8($s0)
16     jal mult_with_add # $v0 = a x b x c
17     sw $v0, 0($s1)
18     li $v0, 10 # exit main
19     syscall

```

Example 5: Array Sum

```

Edit  Execute
proc_sum  procedure_print.asm  procedure_cuboid  procedure_cmp.asm  procedure_mult.asm
2  .data
3  arr: .word 1 -2 3 -4 5 -6 7 -8
4  msg1: .asciiz "\n The sum of positive values: "
5  msg2: .asciiz "\n The sum of negative values: "
6  .text #text segment
7  .globl _main
8  _main: # caller
9      la $a0, arr # starting addr of array
10     addi $a1, $0, 8 # array size
11     jal sum # call sum(arr, size), it returns two values in $v
12     add $s0, $v0, $0 # $s0 positive sum
13     add $s1, $v1, $0 # $s1 negative sum
14     la $a0, msg1
15     li $v0, 4
16     syscall
17     add $a0, $s0, $0 # print positive sum
18     li $v0, 1
19     syscall
20     la $a0, msg2
21     li $v0, 4
22     syscall
23     add $a0, $s1, $0 # print negative sum
24     li $v0, 1
25     syscall
26     li $v0, 10 # exit
27     syscall
```

Edit	Execute			
proc_sum	procedure_print.asm	procedure_cuboid	procedure_cmp.asm	procedure_mult.asm

```
27  sum: # callee
28      add $v0, $0, $0 # positive sum
29      add $v1, $0, $0 # negative sum
30      add $t0, $0, $0 # loop iterator
31  loop:
32      slt $t1, $t0, $a1
33      beq $t1, $0, loop_done # while (i < size)
34      sll $t2, $t0, 2 # $t2 = 4 x i
35      add $t2, $t2, $a0 # $t2 = addr of arr[i]
36      lw $t3, 0($t2) # $t3 = arr[i]
37      bltz $t3, negative # arr[i] <= 0?
38      add $v0, $v0, $t3 # $v0 += positive element
39      j increment
40  negative:
41      add $v1, $v1, $t3 # $v1 += negative element
42  increment:
43      addi $t0, $t0, 1 # i++
44      j loop
45  loop_done:
46      jr $ra
```

Extra Exercise

■ The following C++ function takes as inputs the base address of an int array **A** and returns the minimum value in **A**. Use the registers **\$a0** and **\$a1** as arguments to the function, **\$v0** as return value, **\$ra** as function return address. Translate the C++ function into a MIPS function.

```
int minArray (int A[], int arraySize){  
    int min = A[0];  
    int i = 1;  
    while (i < arraySize) {  
        if (min > A[i])  
            min = A[i];  
        i++;  
    }  
    return min;  
}
```