Outline of Lecture

- Division Algorithms
- Division Hardware

Division

• The final arithmetic operation to be included in the ALU is *division*.

• Paper-and-pencil example:

```
1001 Quotient
Divisor 1000 1001010 Dividend
-1000
10
101
1010
-1000
-1000
10 Remainder
```

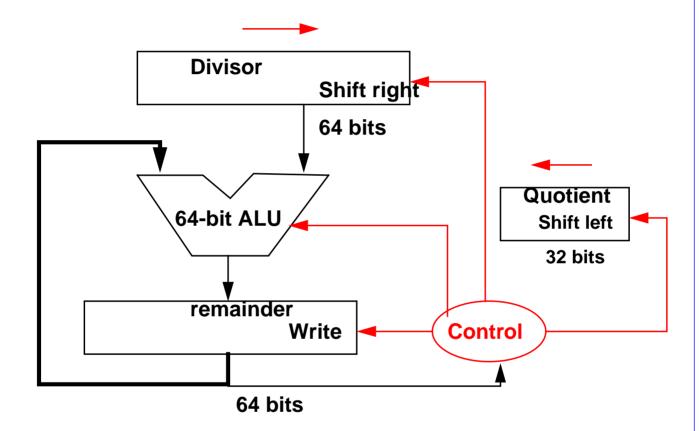
Dividend = Quotient x Divisor + Remainder

- See how big a number can be subtracted, creating quotient bit on each step:
- Binary \Rightarrow 1 x divisor or 0 x divisor

First Version Start 1. Subtract the divisor reg. from the remainder reg. and place the result in the remainder register Remainder > 0 Test Remainder < 0 Remainder 2b. Restore original value by adding the divisor to the 2a. Shift quotient reg. left setting remainder reg. and place the new rightmost bit to 1. the sum in the remainder reg. Also shift the quotient reg. left, setting the new least significant bit to 0. 3. Shift Divisor reg. right 1 bit No 33rd repetition-Yes Done

Division Hardware

• The hardware needed to implement the paperand-pencil algorithm is as follows:



Example

Using 4-bit numbers, divide 0000 0111_2 by 0010_2 .

Solution

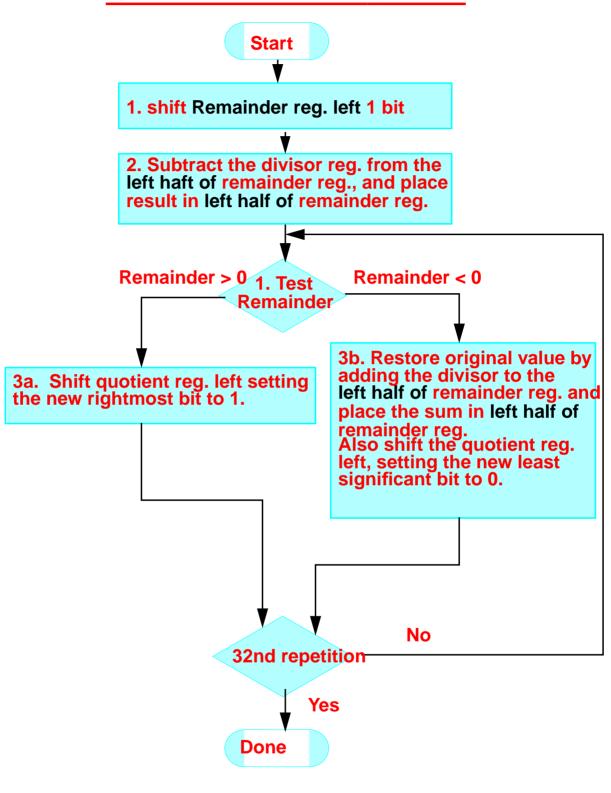
Iteration	Step	Quotient	Divisor	Remainder
0	Initial Values	0000	0010 0000	0000 0111
1	1a: Rem = Rem - Div	0000	0010 0000	1110 0111
	2b: Rem $<0 \Rightarrow +$ Div, sll Q, Q0=0	0000	0010 0000	0000 0111
	3: Shift Div right	0000	0001 0000	0000 0111
2	1a: Rem = Rem - Div	0000	0001 0000	1111 0111
	2b: Rem $<0 \Rightarrow +$ Div, sll Q, Q0=0	0000	0001 0000	0000 0111
	3: Shift Div right	0000	0000 1000	0000 0111
3	1a: Rem = Rem - Div	0000	0000 1000	1111 1111
	2b: Rem $<0 \Rightarrow +$ Div, sll Q, Q0=0	0000	0000 1000	0000 0111
	3: Shift Div right	0000	0000 0100	0000 0111
4	1a: Rem = Rem - Div	0000	0000 0100	0000 0011
	2b: Rem>0 ⇒ sll Q, Q0=1	0001	0000 0100	0000 0011
	3: Shift Div right	0001	0000 0010	0000 0011
5	1a: Rem = Rem - Div	0001	0000 0010	0000 0001
	2b: Rem>0 ⇒ sll Q, Q0=1	0011	0000 0010	0000 0001
	3: Shift Div right	0011	0000 0001	0000 0001

Observations

 The division needs a 64-bit ALU. However, only half of the divisor bits contain useful information.

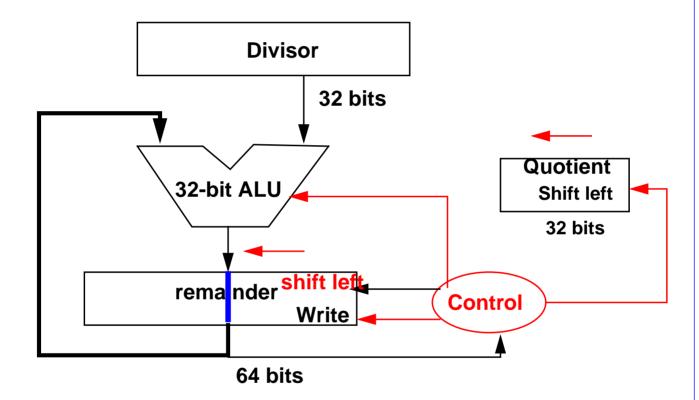
- On the other hand, shifting the remainder to the left instead of shifting the divisor to the right produces the same alignment and accomplishes the same goal.
- Thus, we could limit the size of the divisor to the 32 bits, and in this case, we will need just a 32-bit ALU to accomplish division - which can be done by our 32-bit MIPS ALU.

Second Version



Hardware

• The hardware needed to implement this new version of the division algorithm is as follows:



Example

Using 4-bit numbers, divide 0000 0111_2 by 0010_2 .

Solution

Iteration	Step	Quotient	Divisor	Remainder
0	Initial Values	0000	0010	0000 0111
1	1: Shift Remainder left	0000	0010	0000 1110
	2: Rem = Rem - Div	0000	0010	1 110 1110
	3b: Rem $< 0 \Rightarrow$ + Div, sll Q, Q0 = 0	0000	0010	0000 1110
2	1: Shift Remainder left	0000	0010	0001 1100
	2: Rem = Rem - Div	0000	0010	1 111 1100
	3b: Rem $< 0 \Rightarrow$ + Div, sll Q, Q0 = 0	0000	0010	0001 1100
3	1: Shift Remainder left	0000	0010	0011 1000
	2: Rem = Rem - Div	0000	0010	0001 1000
	3a: Rem > $0 \Rightarrow sll Q, Q0 = 1$	0001	0010	0001 1000
4	1: Shift Remainder left	0001	0010	0011 0000
	2: Rem = Rem - Div	0001	0010	0001 0000
	3a: Rem > $0 \Rightarrow sll Q, Q0 = 1$	0011	0010	0001 0000