## **Mathematical Proofs**

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## Some Inference Rules in Propositional Logic

- Since p is true and  $p \rightarrow q$ , we conclude that q is also true.
  - called, modus ponens.
- Since q is false and  $p \rightarrow q$ , we conclude that p is also false
  - called, modus tollens.
- Since  $p \rightarrow q$  and  $q \rightarrow r$ , we conclude that  $p \rightarrow r$ .
  - called, hypothetical syllogism.
- Since  $p \lor q$  is true and p is false, we conclude that q must be true.
  - called, disjunctive syllogism.

### Remark

In predicate logic, we have similar inference rules. In this case, whenever we say that P(x) is true or false, we view x as a specific element in its domain.

# **Arguments**

### **Definition 1**

- An <u>argument</u> in propositional logic is a sequence of propositions.
- All but the final proposition in the argument are called <u>premises</u>, and the final proposition is called the <u>conclusion</u>.
- An argument is <u>valid</u> if the truth of all its premises implies that the conclusion is true.

## Example 2

The following is an argument for the conclusion that  $5 \le \sum_{i=1}^{5} i \le 25$ .

Proposition 1:  $\sum_{i=1}^{5} i \ge \sum_{i=1}^{5} 1 = 5$  (Premise 1).

Proposition 2:  $\sum_{i=1}^{5} i \le \sum_{i=1}^{5} 5 = 25$  (Premise 2).

Proposition 3:  $5 \le \sum_{i=1}^{5} i \le 25$  (Conclusion).

Combining Propositions 1 and 2 yields the desired conclusion.

## **Proofs and Proof Methods**

### **Definition 3**

- A <u>proof</u> is a valid and clear argument that demonstrates the truth of a theorem (statement).
- A proof is based on premises/axioms/definitions (i.e., statements already assumed/known to be true), and inference rules.
- A <u>proof method</u> usually has a form that can be justified by an **inference** rule.

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# Proving that " $P(x) \Rightarrow Q(x)$ " Is True or False

## Clarification of terminology

Recall that for any predicate P(x) with domain D, P(x) is either true or false for any specific  $x \in D$ . So whenever we say P(x) is true or false, we mean the specific proposition P(x) for a specific x in the domain.

By ' $P(x) \Rightarrow Q(x)$ ' being true or false, we mean the same.

### Remarks

- Recall P(x) implies Q(x) means " $\forall x \in D, P(x) \to Q(x)$ ".
- Recall P(x) implies Q(x) is true if, whenever P(x) is true, Q(x) is also true.
- Recall P(x) implies Q(x) is false if there is any counterexample x = a where P(a) is true and Q(a) is false.

# Counter Example Proof for " $P(x) \Rightarrow Q(x)$ " Being False

### **Problem**

Prove that " $P(x) \Rightarrow Q(x)$ " is false.

### How?

Find an element a in the common domain such that P(a) is true and Q(a) is false.

## Example 4

Let P(n) be "n is a multiple of 4, and Q(n) be "n is a multiple of 8 with common domain  $\mathbb{N}$ . Prove that " $P(x) \Rightarrow Q(x)$ " is false.

### Proof.

Note that P(4) is true, but Q(4) is false. Hence, n=4 is an counterexample.

# Direct Proof for " $P(x) \Rightarrow Q(x)$ "

#### **Problem**

Prove that " $P(x) \Rightarrow Q(x)$ ".

### How?

Assume P(x) is true. Derive a chain of implications, which ends with Q(x).

## Example 5

Prove x < 0 implies x < 1.

### Proof.

Assume that x < 0. We want to prove that x < 1.

- **1** By assumption, x < 0.
- ② We know that 0 < 1.
- **3** Combining the two implications above yields x < 0 < 1.

# Proof by Contraposition for " $P(x) \Rightarrow Q(x)$ "

### **Problem**

Prove that " $P(x) \Rightarrow Q(x)$ " (equivalent to its contrapositive " $\sim Q(x) \Rightarrow \sim P(x)$ ")

### How?

Assume Q(x) is false. Prove that P(x) is also false (it is an indirect proof).

## Example 6

Prove x < 0 implies x < 1.

### Proof.

Assume that  $x \ge 1$ . We want to prove that x > 0.

- By assumption,  $x \ge 1$ .
- ② We know that 1 > 0.
- **3** Combining the two implications above yields  $x \ge 1 > 0$ .

# Proof by Contradiction for " $P(x) \Rightarrow Q(x)$ "

### **Problem**

Prove that " $P(x) \Rightarrow Q(x)$ " (it is an indirect proof).

## How?

Assume that P(x) is true but Q(x) is false. Then show a contradiction.

## Example 7

Prove xy = 0 implies  $x = 0 \lor y = 0$ .

### Proof.

Assume that xy = 0 and  $x \neq 0 \land y \neq 0$ . We want to derive a contradiction.

- By assumption,  $x \neq 0$  and  $y \neq 0$ .
- ② It then follows that  $xy \neq 0$ , which is contrary to the assumption that xy = 0.

# Proof of " $\exists x, P(x)$ "

### How?

Find a value of x such that P(x) is true.

## Example 8

Prove that there exists an  $x \in \mathbb{N}$  such that  $x^2 - 3x + 2 = 0$ .

### Proof.

We have  $x^2-3x+2=(x-1)(x-2)$ . Hence  $x=2\in\mathbb{N}$  is a solution of the equation  $x^2-3x+2=0$ .

# Proof of " $\forall x, P(x)$ "

## **Direct proof**

Show that P(x) is true for all values of x in the domain.

## Example 9

Prove that  $\lfloor (n+1)/2 \rfloor \geq n/2$  for all  $n \in \mathbb{N}$ .

### Proof.

When *n* is even, |(n+1)/2| = n/2.

When *n* is odd,  $\lfloor (n+1)/2 \rfloor = (n+1)/2 > n/2$ .

Combining the conclusions in the two cases completes the proof.

# Proof of " $\forall x, P(x)$ "

## Proof by contradiction

Assume P(x) is false for some value of x in the domain. We want to derive a contradiction.

## Example 10

Prove that  $x < x^2 + 1$  for all  $x \in \mathbb{R}$ , the set of real numbers.

## Proof.

Assume that  $x \ge x^2 + 1$  for some  $x \in \mathbb{R}$ . We want to derive a contradiction.

- $x \ge x^2 + 1$  and  $x^2 + 1 \ge 1$  implies that  $x \ge 1$ .
- $x \ge x^2 + 1$  and  $x^2 + 1 > x^2$  implies that  $x > x^2$ .
- x > 1 and  $x > x^2$  implies that x < 1.

 $x \ge 1$  and x < 1 form a contradiction.

## Indirect Proofs: Two Classical Theorems

### **Definition 11**

Rational numbers are those of the form  $\frac{m}{n}$ , where  $n \in \mathbb{N}$  and  $m \in \mathbb{Z}$ .

### Theorem 12

 $\sqrt{2}$  is irrational (not rational).

### Proof.

Suppose  $\sqrt{2}$  is rational. Then there are two integers m and n such that  $\gcd(m,n)=1$  and  $\sqrt{2}=m/n$ . We want to derive a contradiction. We have then  $m^2=2n^2$ . It then follows that m is even. Let m=2k for some integer k. We obtain then

$$n^2 = 2k^2$$
.

Hence, n is also even. Consequently, gcd(m, n) has the factor 2. This is contrary to our assumption that gcd(m, n) = 1.



## Indirect Proofs: Two Classical Theorems

### Theorem 13

There are infinitely many prime numbers.

### Proof.

Suppose there are only a finite number of primes. Then some prime number p is the largest of all the prime numbers, and hence we can list the prime numbers in ascending order:

$$2,3,5,7,11,\ldots,p.$$

Let

$$n = (2 \times 3 \times 5 \times 7 \times 11 \times \cdots \times p) + 1.$$

Then n > 1, and n cannot be divided by any prime number in the list above. Therefore, n is also a prime. Clearly, n is larger than all the primes in the list. This is contrary to the assumption that all primes are in the list above.