COMP170 Discrete Mathematical Tools for Computer Science

Lecture 14 Version 5 Last updated, Nov 17 2005

Discrete Math for Computer Science K. Bogart, C. Stein and R.L. Drysdale Section 5.1, pp. 213-221

Introduction to Probability

Why Study Probability?

- Probability Spaces and Distributions
- Complementary Probabilities
- Probability and Hashing
- The Uniform Probability Distribution

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In Computer Science we often deal with random events. Some involve randomness imposed from the outside, e.g., networking, when requests from computers on the network enter the network at "random" time.

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Studying the performance of computer systems in the presence of these types of randomness, requires understanding randomness, which is the study of probability.

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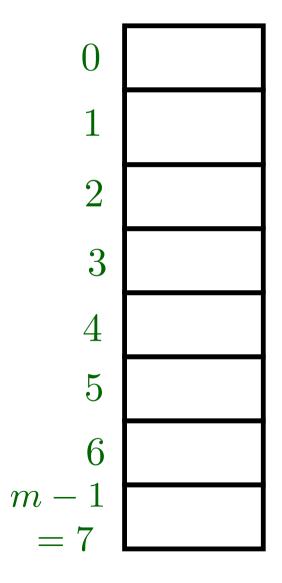
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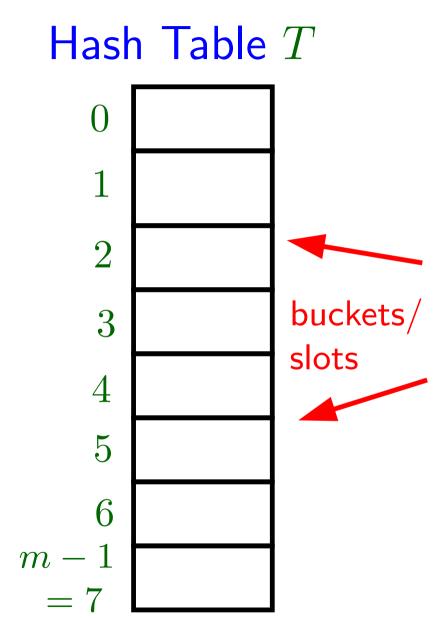
The records are stored in a table. each table location, called a bucket or slot, holds a list of records. We are also given a hash function h(x). A record with key key is stored in the bucket with index h(key).

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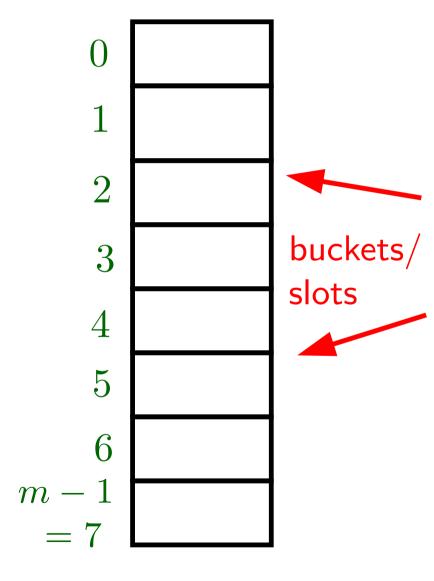


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Our Hash Function:

$$h(x) = x \mod m$$



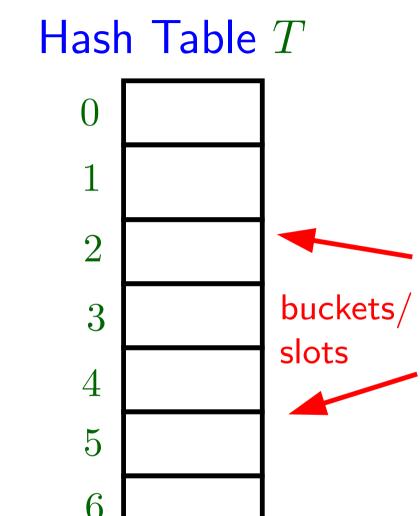


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 $m = 8$

Data (with Keys)



m-1

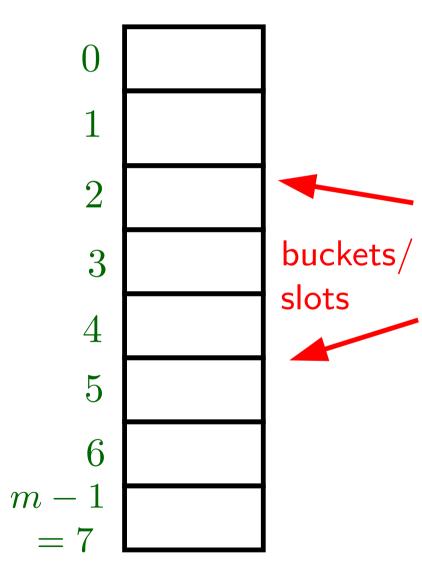
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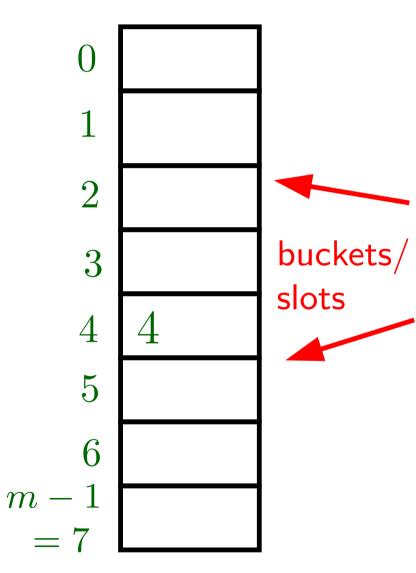
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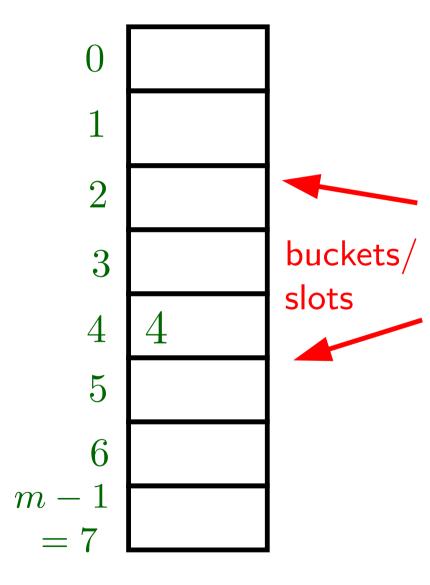
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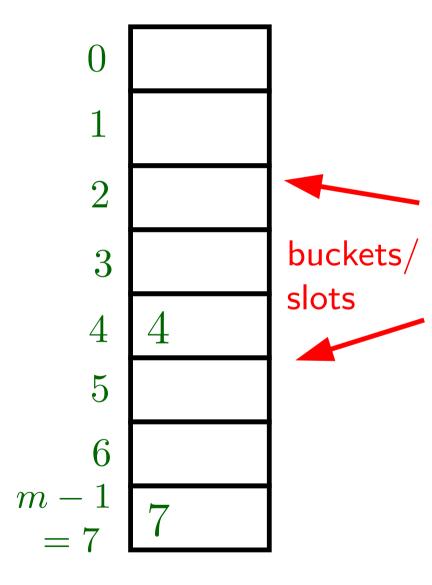
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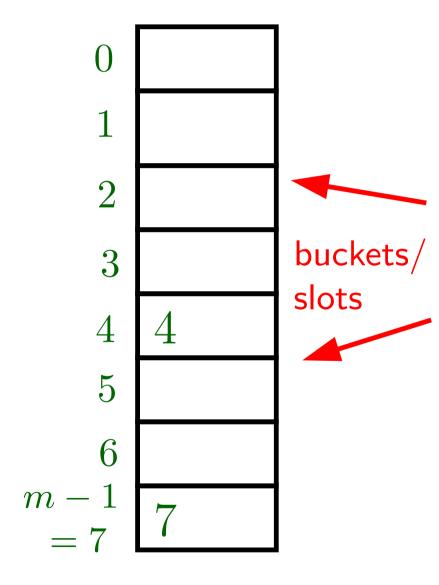
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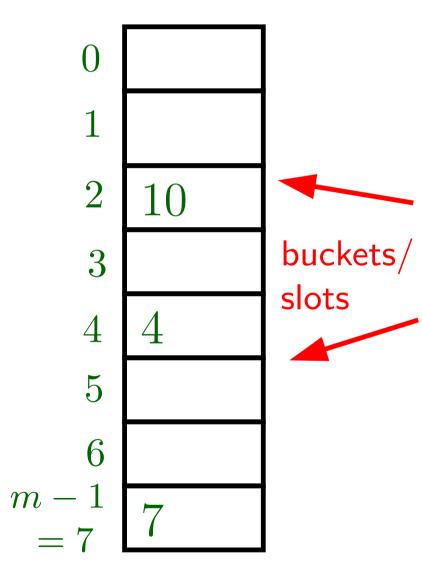
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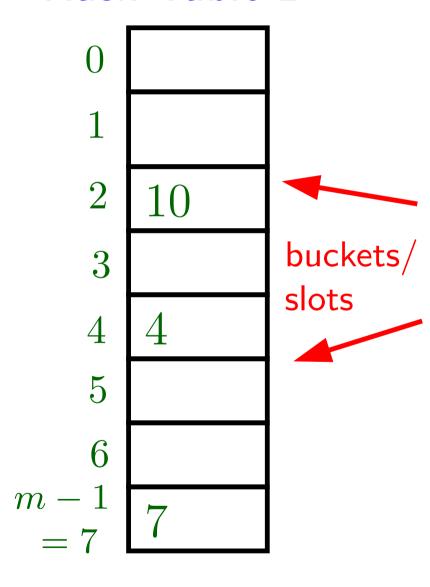
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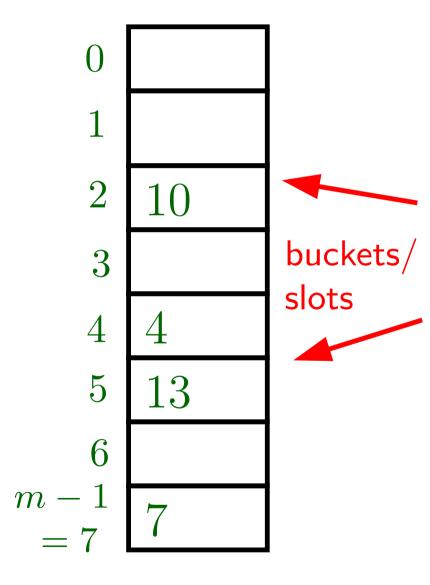
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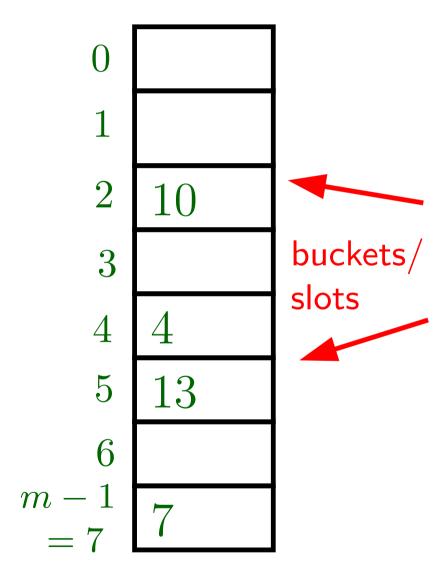
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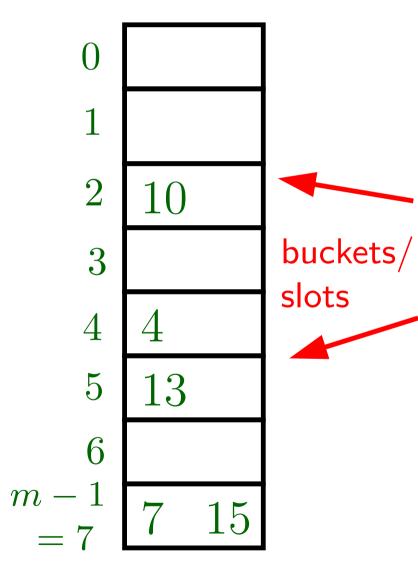
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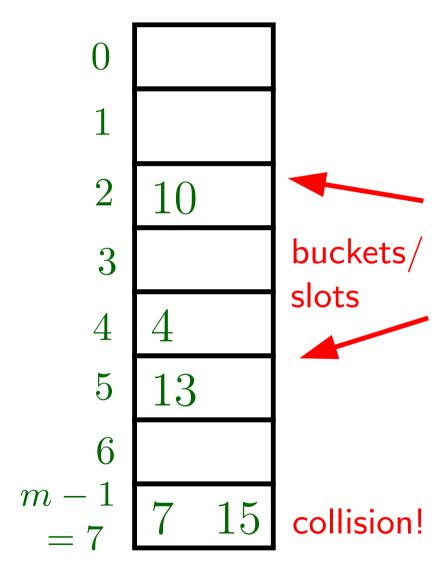
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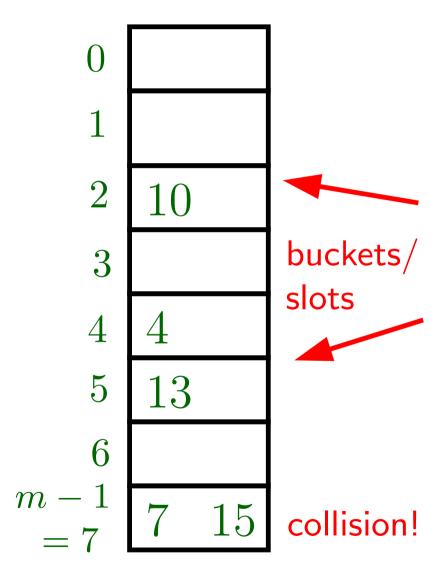
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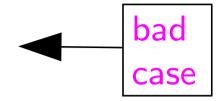
When searching for a record you might have to look at *every* record in the appropriate bucket, so

Good hash function spreads keys evenly among buckets.

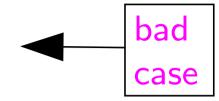


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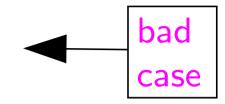


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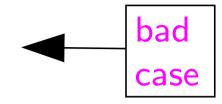
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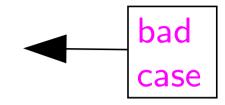


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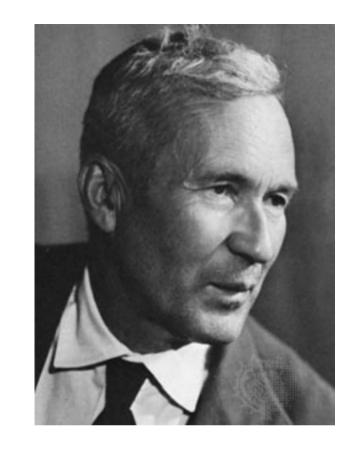
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- The Weight of an element in the sample space Gives a Probability Distribution (Measure)

Andrei Nikolaevich Kolmogorov

Russian Mathematician

b. 1903. d. 1987

The birth of probability theory is often dated to 1654, when Pascal and Fermat, trying to solve a gambling problem, developed the fundamentals.



It wasn't until the work of Kolmogorav in 1933, though, that we had a "definition of probability that is precise enough for use in mathematics, yet comprehensive enough to be applicable to a wide range of phenomena".

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Professor starts each class
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The sample space of all possible patterns
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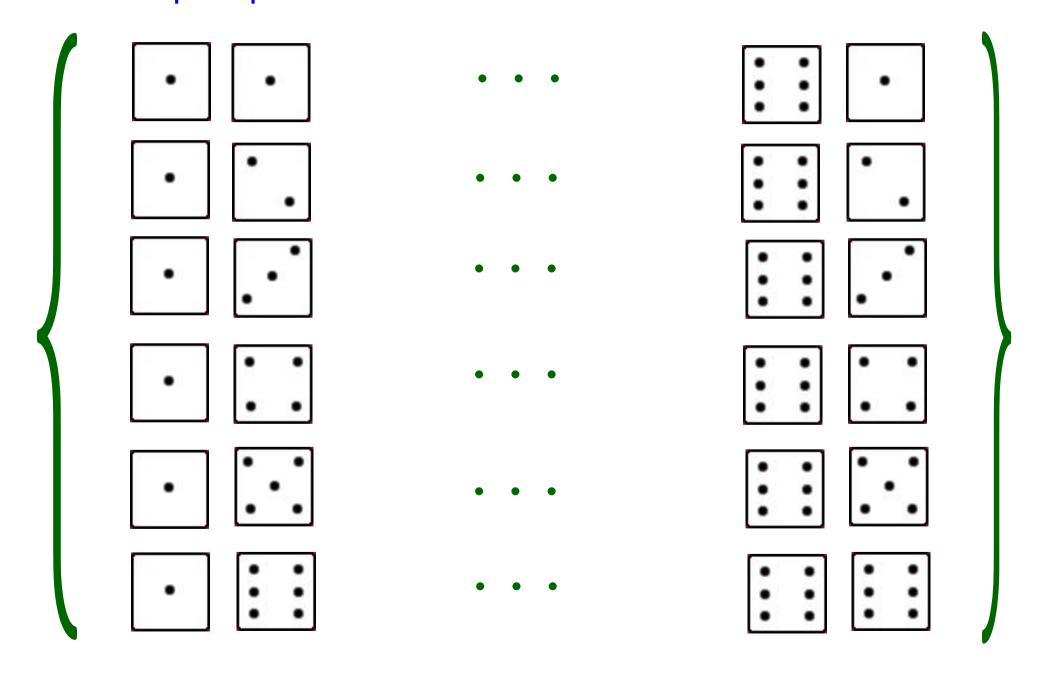
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Note: TTT corresponds to all answers being true, etc...

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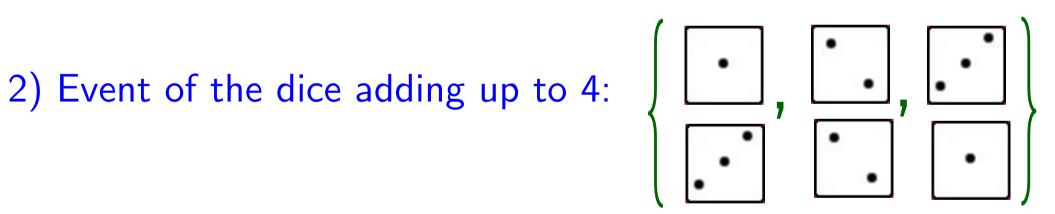
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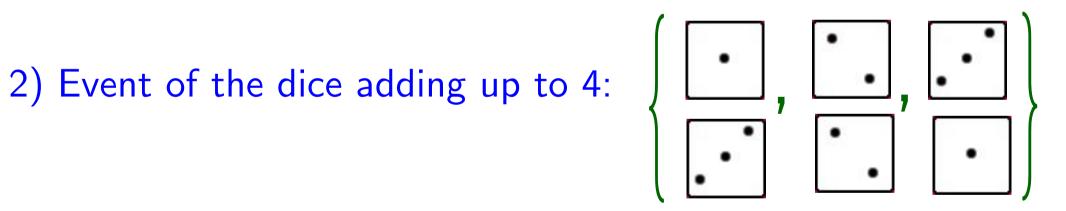
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3) Event of a Head occurring in the first 3 flips: {H, TH, TTH}.

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$$P(E) = \sum_{x:x \in E} P(x)$$

read: "The probability of event E is the sum, over all x such that x is in E, of P(x)."

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A function P satisfying these rules is called a probability distribution or a probability measure.

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Suppose that each sequence of T and F is equally likely. We then want to assign each outcome the same (uniform) probability weight. Since the sum of the weights must add up to 1, we assign each of the 8 outcomes a weight of $\frac{1}{8}$.

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These add up to 1, so this is a legal probability distribution

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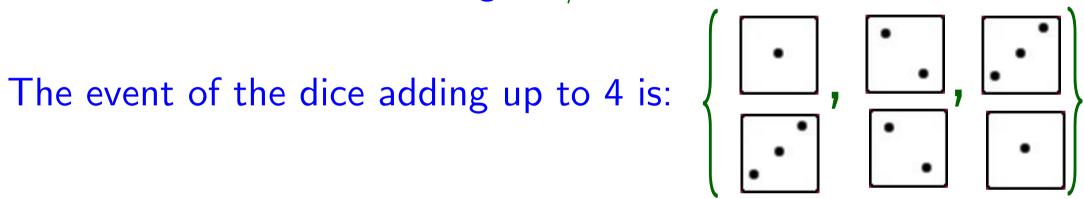
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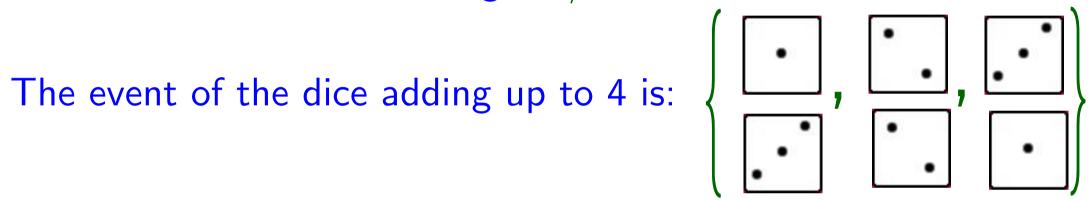
$$\frac{8}{27} + \frac{4}{27} + \frac{4}{27} + \frac{2}{27} = \frac{2}{3}$$

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There are three outcomes in this event, so its probability is

$$\frac{1}{36} + \frac{1}{36} + \frac{1}{36} = \frac{1}{12}$$

When throwing a coin until the first H is seen, the sample space is

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Note that this is a legal probability distribution, since

$$\sum_{i=1}^{\infty} \frac{1}{2^i} = 1$$

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Note that $S = A \cup B$ (or B = S - A) and that P(A) + P(B) = 1.

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Introduction to Probability

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- Probability Spaces and Distributions
- Complementary Probabilities
- Probability and Hashing
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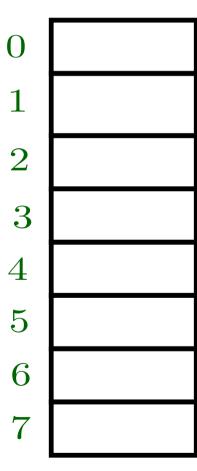
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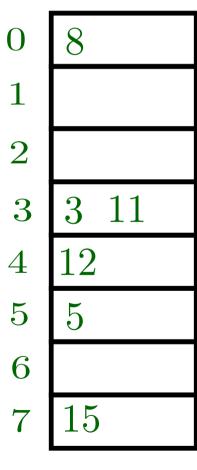
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Example: 3, 12, 15, 8, 11, 5

is outcome (3, 4, 7, 0, 3, 5)



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The size of the sample space is m^n (why?).

Weight function: Assuming that hash function is "random" then every n-tuple is equally likely. So, every n-tuple should have (the same) weight $1/m^n$.

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$$P(B) = 1 - P(A) = 1 - .855 = .145$$

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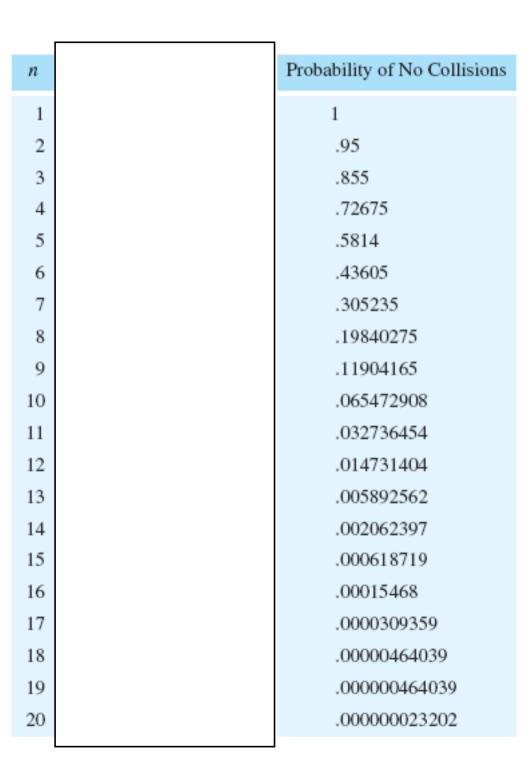
Since events A and B are complementary

$$P(B) = 1 - P(A) = 1 - \frac{20}{20^n}$$

Probability of No Collisions n1 .95 .855 .72675 4 .5814 .43605 6 .305235 .19840275 9 .11904165 10 .065472908 11 .032736454 12 .014731404 13 .005892562 14 .002062397 15 .000618719 16 .00015468 17 .0000309359 18 .00000464039 19 .000000464039 20 .000000023202

Probabilities that all elements of set hash to different entries of hash table of size 20 is

$$p_n = \frac{20 \text{ }^{\underline{\text{II}}}}{20^n}.$$



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Since

$$p_{n+1} = p_n \frac{20 - n}{20} < p_n,$$

 p_n decreases as n increases

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Suppose P is uniform probability measure defined on sample space S. Then for any event E,

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Note: We have implicitly used this theorem many times already

Let $S = \{x_1, x_2, \dots, x_{|S|}\}.$

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$$\Rightarrow p = \frac{1}{|S|}$$

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Combining these equations gives

$$P(E) = |E|p = |E|(1/|S|) = |E|/|S|.$$

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Use Theorem 5.2

 \Rightarrow probability is 4/8 = 1/2 by Theorem 5.2.

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Theorem 5.2 doesn't apply to this distribution. For example, let E be the event that the outcome is not positive.

Then $E = \{0\}$ but

$$P(E) = \frac{1}{8} \neq \frac{1}{4} = \frac{|E|}{|S|}$$

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which is exactly the non-uniform distribution we just saw.