COMP170
Discrete Mathematical Tools for Computer Science

Inference

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Discrete Math for Computer Science
K. Bogart, C. Stein and R.L. Drysdale
Section 3.3, pp. 117-124

3.3 Inference

- What is a Proof?
- Direct Inference (Modus Ponens)
- Rules of Inference for Direct Proofs
- Contrapositive Rule of Inference
- Proof by Contradiction

What is a *Mathematical Proof*, really?

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- In this section we will introduce various techniques used to develop mathematical proofs.

  Some of these techniques will actually be variations on similar ideas (so don't get confused if they look similar to each other).

- We start by examining a simple mathematical proof and its components

Prove that if m is even, then m^2 is even.

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If m is even, then $\exists k$ with $m = 2k$.

Then $\exists k$ such that $m = 2k$.

Prove that if m is even, then m^2 is even.

Let m be an integer.

Suppose that m is even.

If m is even, then $\exists k$ with $m = 2k$.

Then $\exists k$ such that $m = 2k$.

Then $\exists k$ such that $m^2 = 4k^2$.

Prove that if m is even, then m^2 is even.

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If m is even, then $\exists k$ with $m = 2k$.

Then $\exists k$ such that $m = 2k$.

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Then, there is an integer $h = 2k^2$ s.t. $m^2 = 2h$.

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Suppose that m is even.

If m is even, then $\exists k$ with $m = 2k$.

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Then, there is an integer $h = 2k^2$ s.t. $m^2 = 2h$.

Thus, if m is even, then m^2 is even.

3.3 Inference

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Then we can rewrite the three statements as

1) p

2) If p then q $\quad (p \Rightarrow q)$

3) q

Direct Inference (Modus Ponens)

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Principle 3.3 (Direct inference)
From p and $p \Rightarrow q$ we may conclude q.

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Why is this valid?

IMPLIES

p	q	$p \Rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

In our example proof we showed that

If m is even then m^2 is even.

Essentially, we assumed m is even and derived that m^2 is even.

In symbols, we showed that

$(m \text{ is even}) \Rightarrow (m^2 \text{ is even})$.

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Principle 3.4 (Conditional Proof)
If by assuming p we may prove q, then the statement $p \Rightarrow q$ is true

We showed that

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Principle 3.5 (Universal Generalization)
If we can prove a statement $p(x)$ about x by assuming only that x is a member of our universe, then we can conclude that $p(x)$ is true for every member of our universe.

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4. From either $q(x)$ or $\neg p(x)$
   we may conclude $p(x) \Rightarrow q(x)$.

5. From $p(x) \Rightarrow q(x)$ and $q(x) \Rightarrow p(x)$,
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7. From $p(x) \Rightarrow q(x)$ and $q(x) \Rightarrow r(x)$,
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8. If we can derive $q(x)$ from hypothesis that x satisfies $p(x)$,
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9. If we can derive $p(x)$ from the hypothesis that
x is a (generic) member of our universe U,
we may conclude $\forall x \in U\ (p(x))$.

5. From $p(x) \Rightarrow q(x)$ and $q(x) \Rightarrow p(x)$,
   we may conclude $p(x) \Leftrightarrow q(x)$.

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7. From $p(x) \Rightarrow q(x)$ and $q(x) \Rightarrow r(x)$,
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9. If we can derive $p(x)$ from the hypothesis that
   x is a (generic) member of our universe U,
   we may conclude $\forall x \in U\ (p(x))$.

10. From an example of an $x \in U$ satisfying $p(x)$,
    we may conclude $\exists x \in U\ (p(x))$.

Prove that $\forall m \in Z$, if m is even, then m^2 is even.

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Let m be an integer.

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Setup for rule 9

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If m is even, then $\exists k$ with $m = 2k$.

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If m is even, then $\exists k$ with $m = 2k$. Definition

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Then $\exists k$ such that $m = 2k$.

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If m is even, then $\exists k$ with $m = 2k$. Definition

Then $\exists k$ such that $m = 2k$. Rule 6 (m.p.)

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If m is even, then $\exists k$ with $m = 2k$. Definition

Then $\exists k$ such that $m = 2k$. Rule 6 (m.p.)

Then $\exists k$ such that $m^2 = 4k^2$.

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Then $\exists k$ such that $m^2 = 4k^2$. Algebra

Then, there is an integer $h = 2k^2$ s.t. $m^2 = 2h$.

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If m is even, then $\exists k$ with $m = 2k$. Definition

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Then, if m is even, then m^2 is even. Rule 8

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p implies q is actually equivalent to $\neg q$ implies $\neg p$.

double truth table

p	q	$p \Rightarrow q$	$\neg p$	$\neg q$	$\neg q \Rightarrow \neg p$
T	T	T	F	F	T
T	F	F	F	T	F
F	T	T	T	F	T
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Principle 3.6 (Proof by Contraposition)
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We Adopt Principle 3.6 as a rule of inference, called the **contrapositive rule of inference**.

11. From $\neg q(x) \Rightarrow \neg p(x)$,
    we may conclude $p(x) \Rightarrow q(x)$.

Example:

If n is a positive integer with $n^2 > 100$, then $n > 10$.

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$p(n)$ $q(n)$

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Proof (by contraposition):

Suppose n is not greater than 10. $\neg q(n)$

Then, because $1 \leq n \leq 10$, we have $n \cdot n \leq n \cdot 10 \leq 10 \cdot 10 = 100$.

(Using: "If $x \leq y$ and $c \geq 0$, then $cx \leq cy$.")

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Thus, n^2 is not greater than 100.

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Thus, n^2 is not greater than 100. $\neg p(n)$

Thus, if $n \not> 10$ then $n^2 \not> 100$

Example:

If n is a positive integer with $\underset{p(n)}{n^2 > 100}$, then $\underset{q(n)}{n > 10}$.

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Thus, if $n \not> 10$ then $n^2 \not> 100$ $\neg q(n) \Rightarrow \neg p(n)$

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Thus, if $n \not> 10$ then $n^2 \not> 100$ $\neg q(n) \Rightarrow \neg p(n)$

By the principle of proof by contraposition,
if $n^2 > 100$, then $n > 10$.

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If n is a positive integer with $\underset{p(n)}{n^2 > 100}$, then $\underset{q(n)}{n > 10}$.

Proof (by contraposition):

Suppose n is not greater than 10. $\quad \neg q(n)$

Then, because $1 \le n \le 10$, we have $n \cdot n \le n \cdot 10 \le 10 \cdot 10 = 100$.

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Thus, if $n \not> 10$ then $n^2 \not> 100$ $\quad \neg q(n) \Rightarrow \neg p(n)$

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if $n^2 > 100$, then $n > 10$. $\quad p(n) \Rightarrow q(n)$

Is p implies q equivalent to q implies p?

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$q(x) \Rightarrow p(x)$: If x has 4 legs then x is a cat

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$q \Rightarrow p$ is called the **converse** of $p \Rightarrow q$.

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- Adopt the principle of proof by contradiction (also called the principle of reduction to absurdity) as last rule of inference

12. If by assuming $p(x)$ and $\neg q(x)$, we can derive
    both $r(x)$ and $\neg r(x)$ for some statement $r(x)$,
    we may conclude $p(x) \Rightarrow q(x)$.

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We will now see 3 different proofs by contradiction that $p \Rightarrow q$ where p is the statement $x^2 + x - 2 = 0$, and q is the statement $x \neq 0$.

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We will now see 3 different proofs by contradiction that $p \Rightarrow q$ where p is the statement $x^2 + x - 2 = 0$, and q is the statement $x \neq 0$.

Each of the three proofs by contradiction work by getting slightly different contradictions.

Prove that if $x^2 + x - 2 = 0$, then $x \neq 0$.

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1. We will assume p is True and q is False;
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Thus, by the principle of proof by contradiction, $\sqrt{5}$ is not rational.