# COMP170 Discrete Mathematical Tools for Computer Science Intro to Probability

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Discrete Math for Computer Science K. Bogart, C. Stein and R.L. Drysdale Section 5.1, pp. 213-221

# Introduction to Probability

Why Study Probability?

- Probability Spaces and Distributions
- Complementary Probabilities
- Probability and Hashing
- The Uniform Probability Distribution

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In Computer Science we often deal with random events. Some involve randomness imposed from the outside, e.g., networking, when requests from computers on the network enter the network at "random" time.

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Studying the performance of computer systems in the presence of these types of randomness, requires understanding randomness, which is the study of probability.

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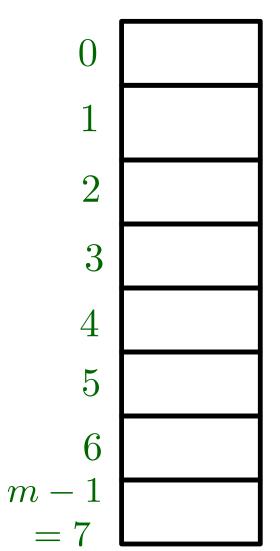
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The records are stored in a table. Each table location, called a bucket or slot, holds a list of records. We are also given a hash function h(x). A record with key key is stored in the bucket with index h(key).

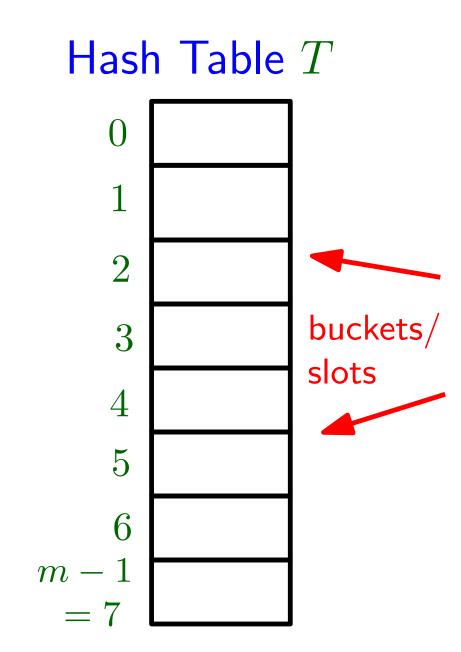
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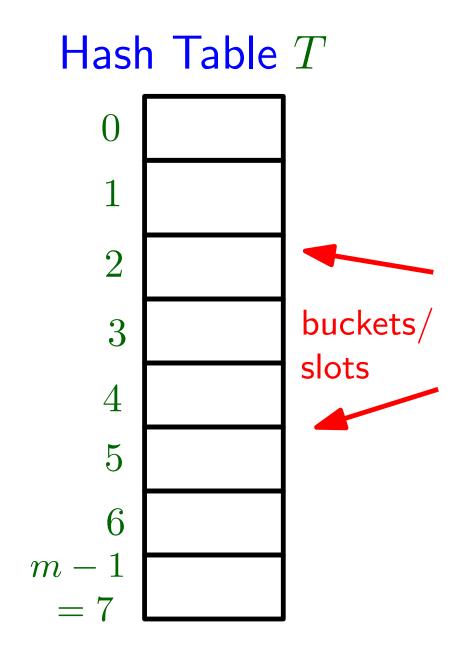
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#### Our Hash Function:

$$h(x) = x \bmod m$$

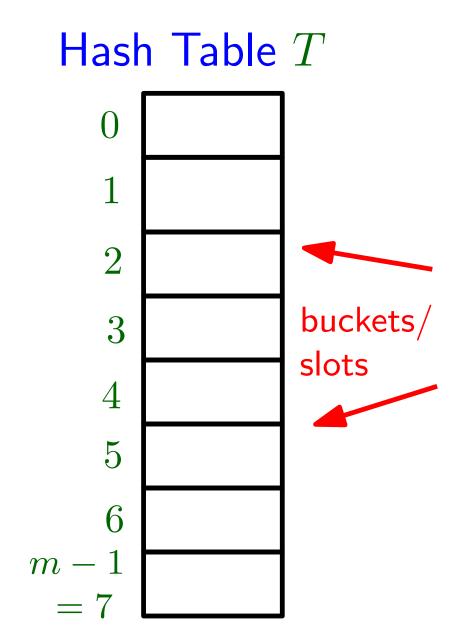


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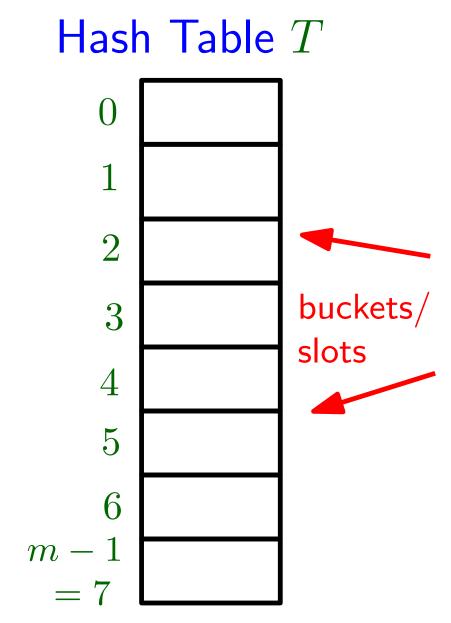


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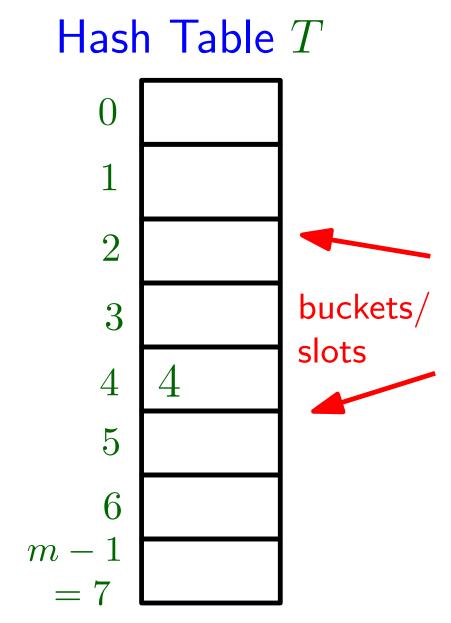


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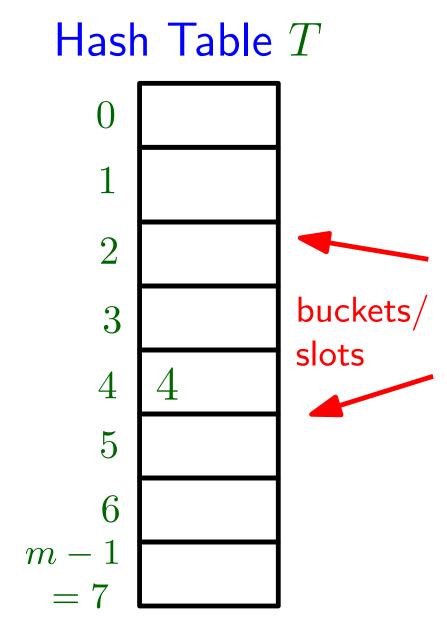


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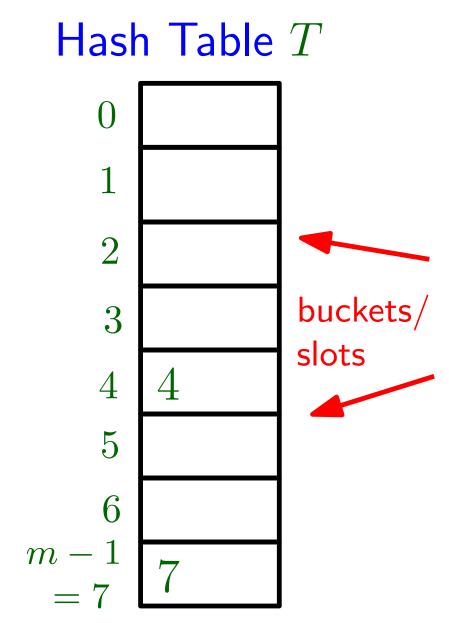


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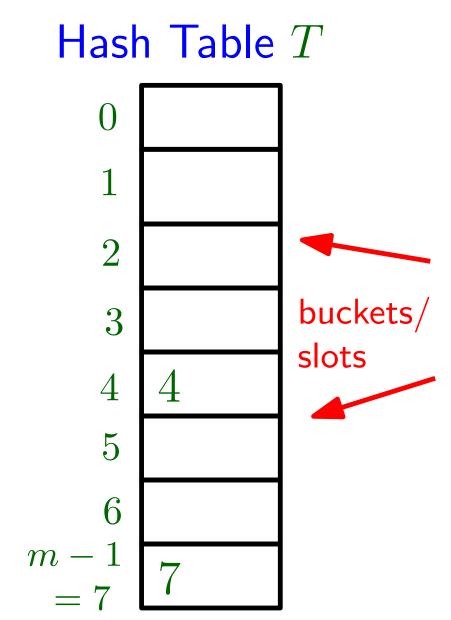
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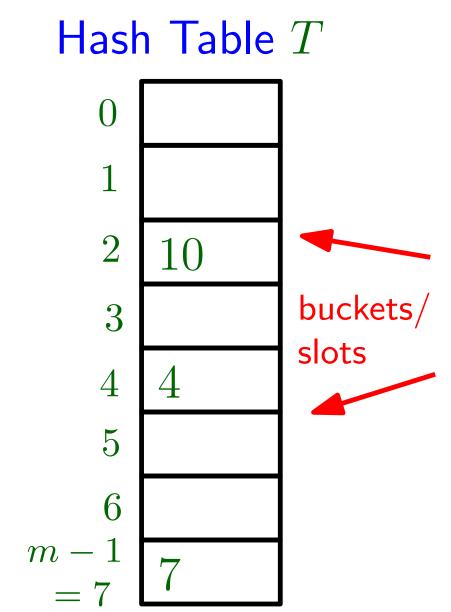
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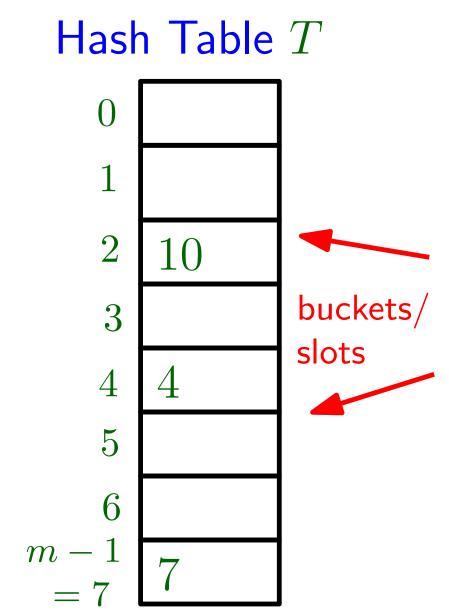
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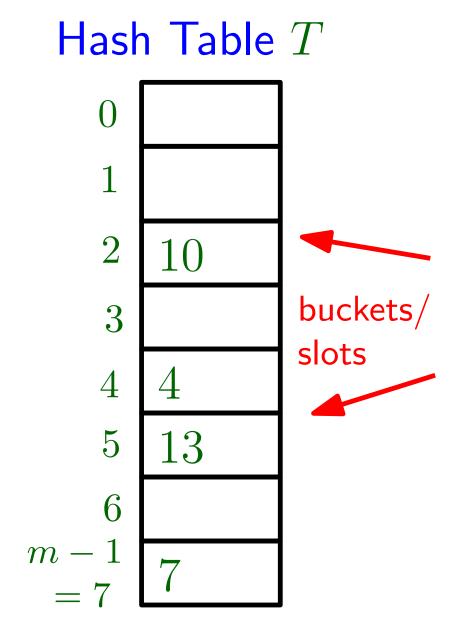
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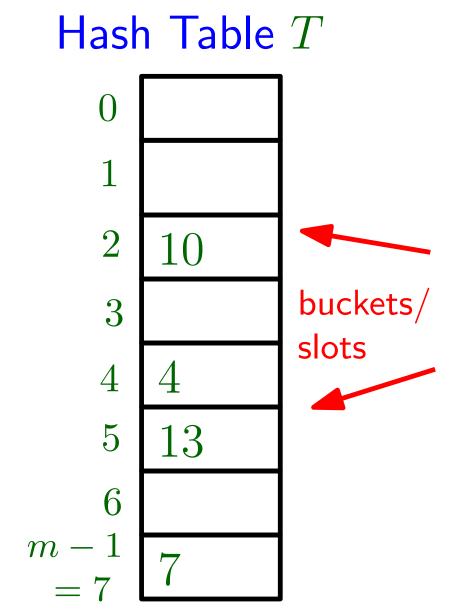
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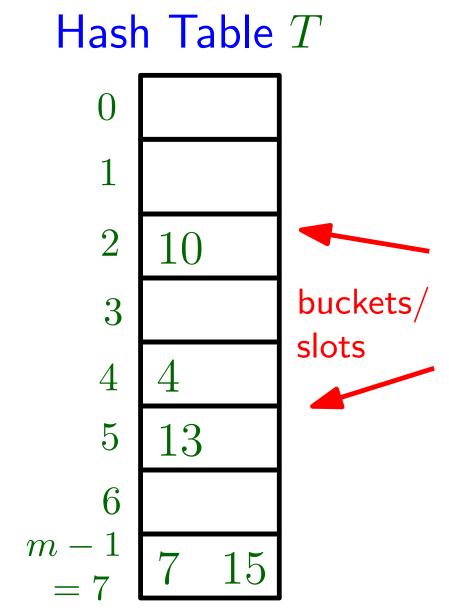
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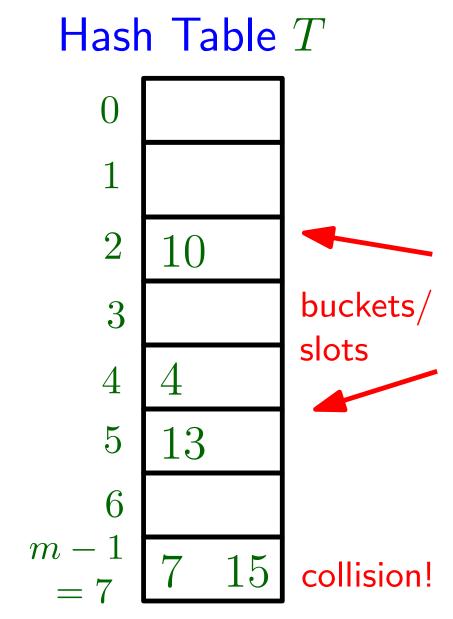
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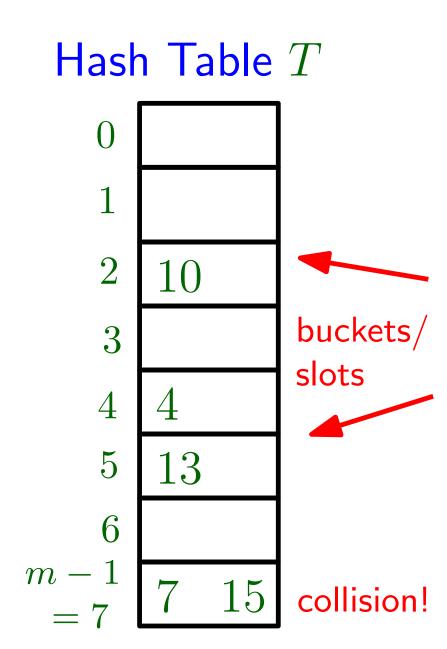
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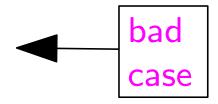
When searching for a record you might have to look at *every* record in the appropriate bucket, so

Good hash function spreads keys evenly among buckets.



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→ Study of Probability

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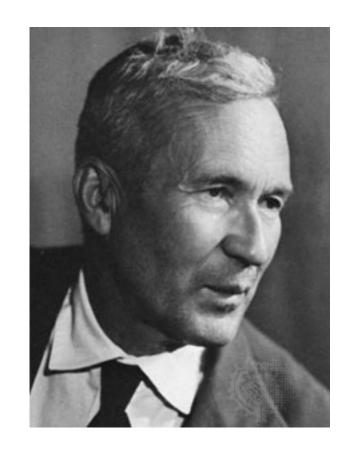
- The underlying Sample Space and Elements (Outcomes) in the sample space
- An event in the Sample Space
- The Weight of an element in the sample space Gives a Probability Distribution (Measure)

#### Andrei Nikolaevich Kolmogorov

Russian Mathematician

b. 1903. d. 1987

The birth of probability theory is often dated to 1654, when Pascal and Fermat, trying to solve a gambling problem, developed the fundamentals.



It wasn't until the work of Kolmogorav in 1933, though, that we had a "definition of probability that is precise enough for use in mathematics, yet comprehensive enough to be applicable to a wide range of phenomena".

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#### Example 1:

Professor starts each class
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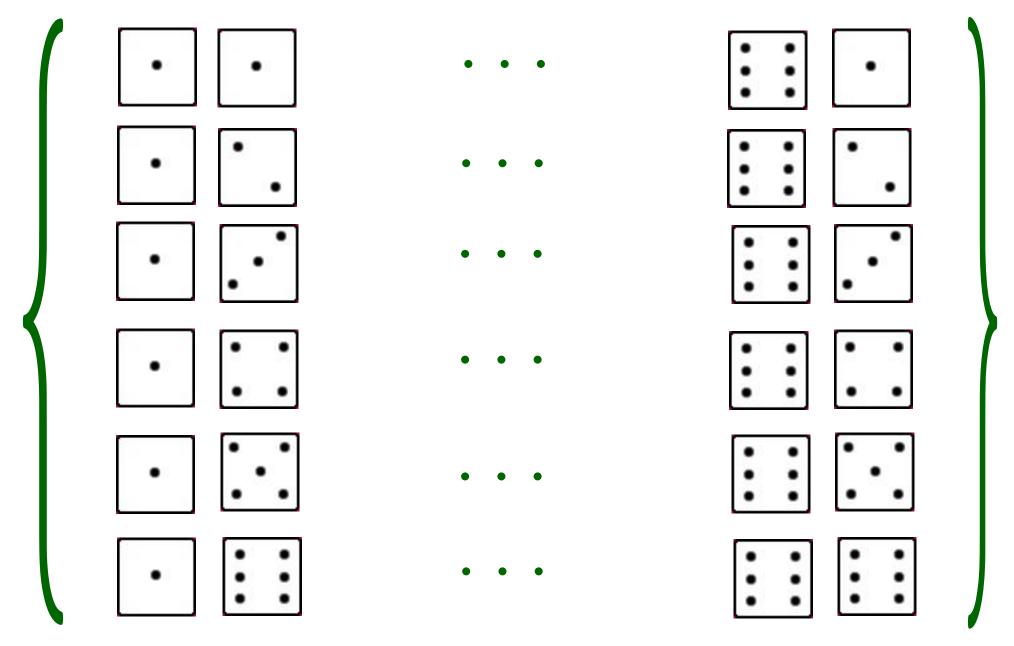
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Note: TTT corresponds to all answers being true, etc..

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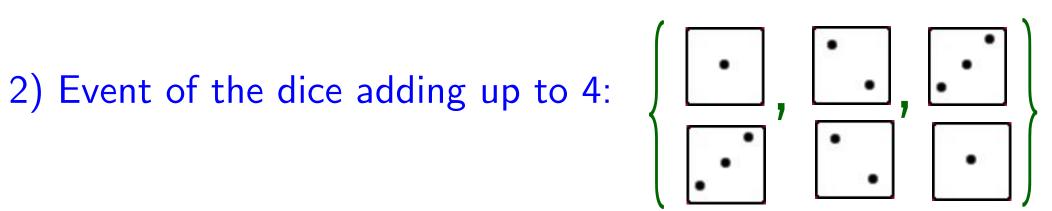
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3) Event of a Head occurring in the first 3 flips: {H, TH, TTH}.

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$$P(E) = \sum_{x:x \in E} P(x)$$

read: "The probability of event E is the sum, over all x such that x is in E, of P(x)."

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A function P satisfying these rules is called a probability distribution or a probability measure.

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Suppose that each sequence of T and F is equally likely. We then want to assign each outcome the same (uniform) probability weight. Since the sum of the weights must add up to 1, we assign each of the 8 outcomes a weight of  $\frac{1}{8}$ .

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These add up to 1, so this is a legal probability distribution

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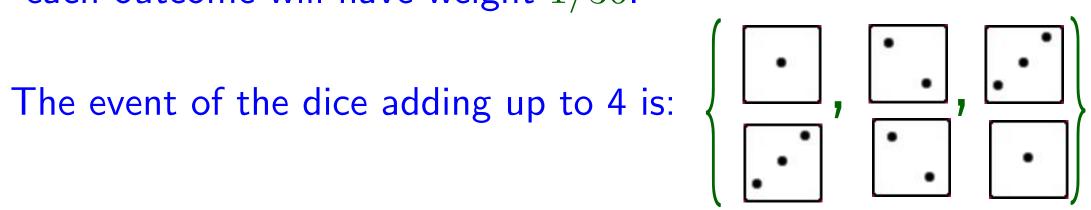
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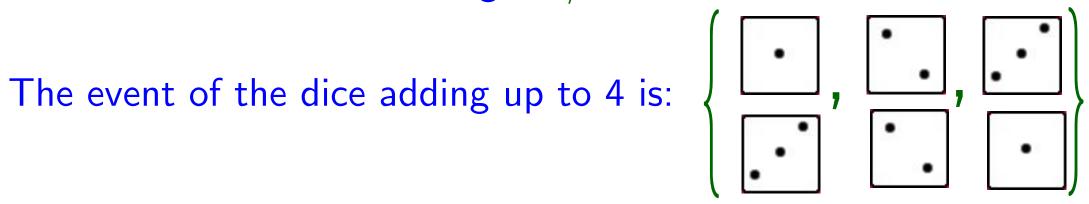
$$\frac{8}{27} + \frac{4}{27} + \frac{4}{27} + \frac{2}{27} = \frac{2}{3}$$

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There are three outcomes in this event, so its probability is

$$\frac{1}{36} + \frac{1}{36} + \frac{1}{36} = \frac{1}{12}$$

When throwing a coin until the first H is seen, the sample space is

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Note that this is a legal probability distribution, since

$$\sum_{i=1}^{\infty} \frac{1}{2^i} = 1$$

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Note that 
$$S = A \cup B$$
 (or  $B = S - A$ ) and that  $P(A) + P(B) = 1$ .

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# Introduction to Probability

Why Study Probability?

- Probability Spaces and Distributions
- Complementary Probabilities
- Probability and Hashing
- The Uniform Probability Distribution

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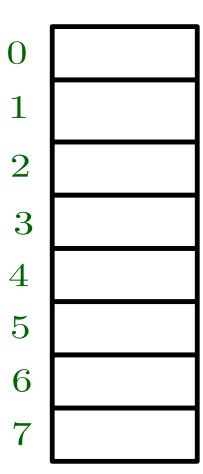
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Example:



# Probability and Hashing

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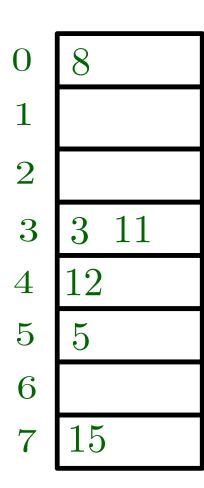
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Example: 3, 12, 15, 8, 11, 5

is outcome (3, 4, 7, 0, 3, 5)



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Weight function: Assuming that hash function is "random" then every n-tuple is equally likely. So, every n-tuple should have (the same) weight  $1/m^n$ .

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$$P(B) = 1 - P(A) = 1 - .855 = .145$$

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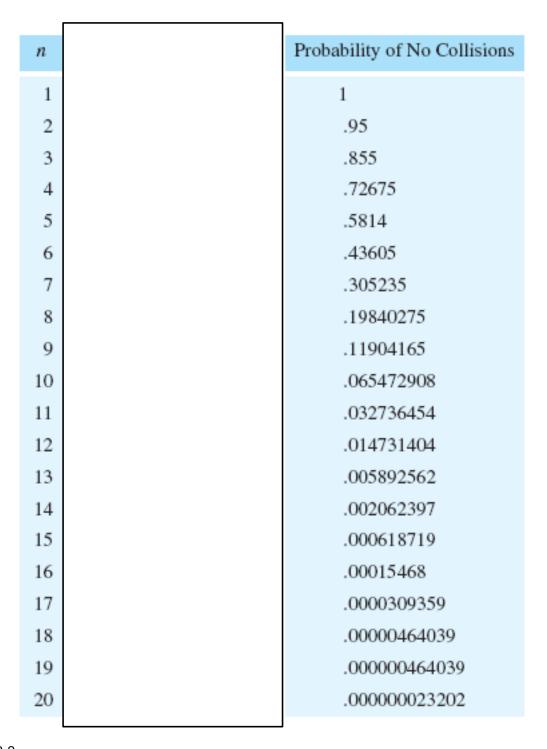
Since events A and B are complementary

$$P(B) = 1 - P(A) = 1 - \frac{20 \text{ }^{\underline{\mathbf{n}}}}{20^n}$$

Probability of No Collisions n1 1 .95 .855 .72675 4 .5814 5 .43605 6 .305235 .19840275 9 .11904165 10 .065472908 11 .032736454 12 .014731404 13 .005892562 14 .002062397 .000618719 15 16 .00015468 17 .0000309359 18 .00000464039 19 .000000464039 .000000023202 20

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 $p_n$  decreases as n increases

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Note: We have implicitly used this theorem many times already

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$$\Rightarrow p = \frac{1}{|S|}$$

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## Combining these equations gives

$$P(E) = |E|p = |E|(1/|S|) = |E|/|S|$$
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#### Use Theorem 5.2

 $\Rightarrow$  probability is 4/8 = 1/2 by Theorem 5.2.

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Theorem 5.2 doesn't apply to this distribution. For example, let E be the event that the outcome is not positive.

Then 
$$E = \{0\}$$
 but

$$P(E) = \frac{1}{8} \neq \frac{1}{4} = \frac{|E|}{|S|}$$

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which is exactly the non-uniform distribution we just saw.