

# CIS 313, Loop Invariant Examples

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- Example 1

Our goal is to prove the correctness of the ArrayFind algorithm<sup>1</sup>:

Algorithm arrayFind(x,A):

Input: An element  $x$  and an  $n$ -element array,  $A$ .

Output: The index  $i$  such that  $x = A[i]$  or  $-1$  if no element of  $A$  is equal to  $x$ .

$i = 0$

while  $i < n$  do

if  $x == A[i]$  then

return  $i$

else

$i = i + 1$

return  $-1$

This example proof is a little more verbose than necessary, but we will include extra detail to make it as clear as possible.

We will prove that, if the loop terminates (without returning early), then  $x$  is not in  $A$ .<sup>2</sup> The loop condition,  $\gamma$ , is: “ $i < n$ ”. We define the loop invariant,  $\alpha$ , as: “ $0 \leq i \leq n$  and  $x$  is not stored in  $A[0..(i-1)]$ ”, where  $A[0..(i-1)]$  refers to the values stored in array  $A$  from indices 0 through  $i-1$ , inclusive. We now use this loop invariant to prove the correctness of the loop by showing that it satisfies all three properties:

- (i) Initialization: From the assignment statement before the start of the loop,  $i = 0$ . Clearly,  $0 \leq 0 \leq n$ . Because  $i = 0$ ,  $A[0..i-1]$  is an array of zero elements and the second half of the invariant is trivially true.
- (ii) Maintenance: Let  $i$  refer to the value of variable  $i$  at the beginning of the loop and  $i'$  to its value at the end of the loop. From the execution of the loop body,  $i' = i + 1$ . Since  $0 \leq i$  (from  $\alpha$ ) and  $i < n$  (from  $\gamma$ ),  $0 \leq i + 1 \leq n$ , so  $0 \leq i' \leq n$ . Furthermore, from  $\alpha$  we know that  $x$  is not contained in  $A[0..(i-1)]$ . From the loop body, if the loop continues then  $x$  is not in  $A[i]$  either. Therefore,  $x$  is not contained in  $A[0..i]$ . From the definition of  $i'$ , we can conclude that  $x$  is not contained in  $A[0..(i'-1)]$  and the invariant remains true with the new value of  $i$ .
- (iii) Termination: Since  $0 \leq i \leq n$  (from  $\alpha$ ) and  $i \not< n$  (from  $\neg\gamma$ ),  $i = n$ . Substituting  $i = n$  back into the loop invariant, we can conclude that  $x$  is not stored in  $A[0..(n-1)]$ , which is the entire array.

- Example 2

**Question:** Use a loop invariant to prove the correctness of the following algorithm for finding the maximum element on an array:

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<sup>1</sup>This example is taken from Chapter 1 of the data structures textbook by Goodrich and Tamassia, which used to be the standard textbook for CIS 313.

<sup>2</sup>If the algorithm *does* return early, then from the preceding if statement we know that it must have found  $x$  and is returning the correct index. Therefore, whenever the algorithm returns early it returns the correct result.

Algorithm arrayMax(A, n)  
Input array A of n integers  
Output maximum element A

```
currentMax = A[0]
i = 0
while i < n-1 do
    if A[i+1] > currentMax
        currentMax = A[i+1]
    i = i + 1
```

**Answer:** We use the loop invariant: “ $0 \leq i \leq n - 1$  and currentMax is the maximum value stored in  $A[0..i]$ ”, where  $A[0..i]$  refers to the values stored in array  $A$  from indices 0 through  $i$ , inclusive. We now use this loop invariant to prove the correctness of the loop by showing that it satisfies all three properties:

- (i) From the initialization statements,  $i = 0$  and  $\text{currentMax} = A[0]$ . Because  $i = 0$ ,  $A[0..i] = A[0]$ , so the maximum value stored in  $A[0..i]$  is simply  $A[0]$ . From the precondition, currentMax already equals this value, satisfying the loop invariant. Also, since  $i = 0$ , clearly  $0 \leq i \leq n - 1$ .
- (ii) Let the primed variables  $i'$  and  $\text{currentMax}'$  refer to the values of  $i$  and currentMax at the end of the loop, and the unprimed variables  $i$  and currentMax to their values at the beginning of the loop. From the execution of the loop body,  $i' = i + 1$ . Since  $0 \leq i$  (from  $\alpha$ ) and  $i < n - 1$  (from  $\gamma$ ),  $0 \leq i + 1 \leq n$ , so  $0 \leq i' \leq n$ . Furthermore,

$\text{currentMax}' = \max(\text{currentMax}, A[i + 1])$	Execution of the loop
$= \max(\max(A[0..i]), A[i + 1])$	Follows from $\alpha$
$= \max(A[0..i + 1])$	max is associative
$= \max(A[0..i'])$	Substitution

Since the new values  $\text{currentMax}'$  and  $i'$  satisfy the loop invariant,  $\alpha$  remains true at the end of the loop.

- (iii) Since  $0 \leq i \leq n - 1$  (from  $\alpha$ ) and  $i \not\leq n - 1$  (from  $\neg\gamma$ ),  $i = n - 1$ . Therefore, currentMax is the maximum value stored in  $A[0..n - 1]$ , which is the maximum value in  $A$ .