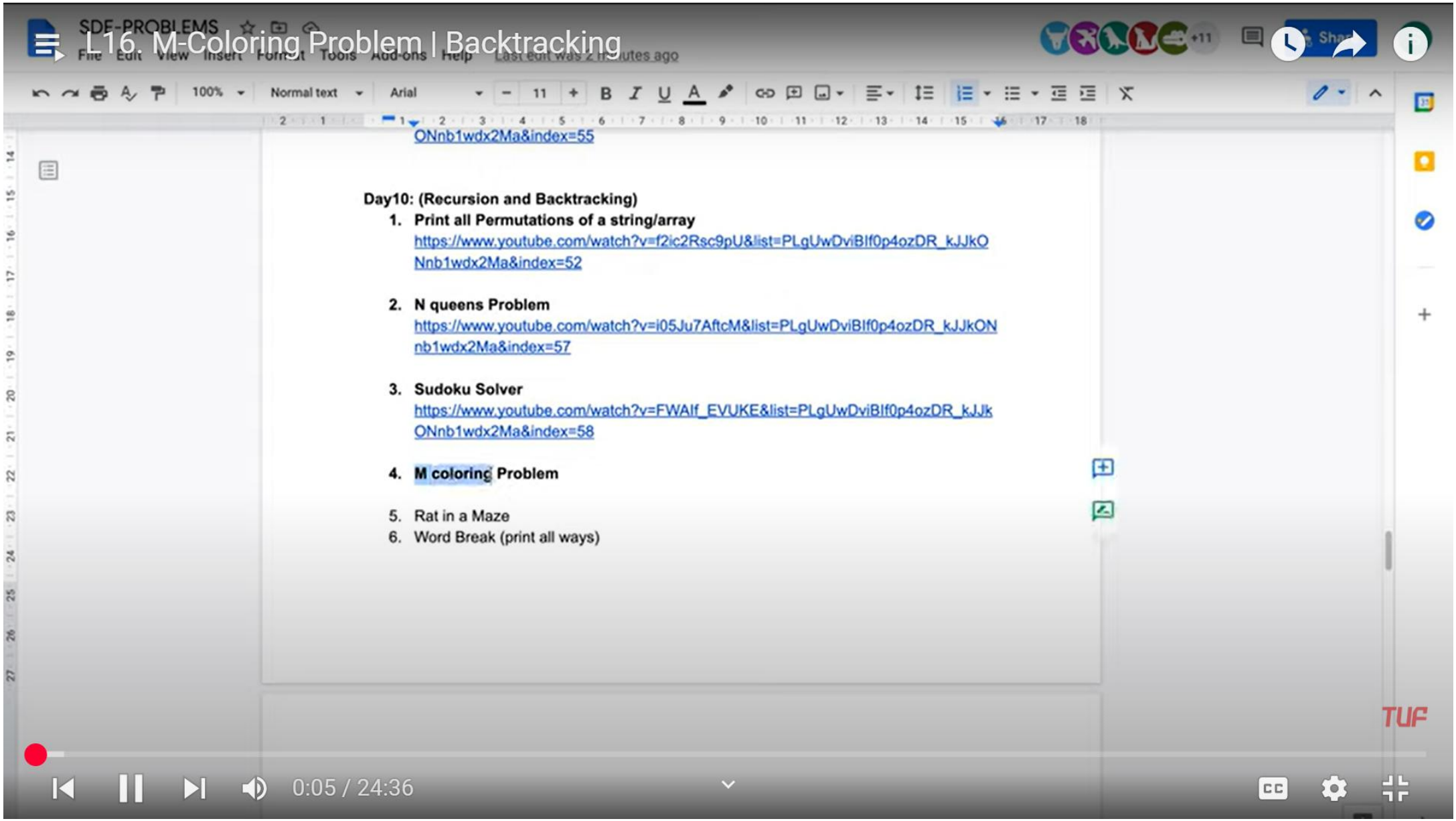

SDE-PROBLEMS
L16. M-Coloring Problem | Backtracking
ONnb1wdx2Ma&index=55
Day10: (Recursion and Backtracking)
1. Print all Permutations of a string/array
https://www.youtube.com/watch?v=f2ic2Rsc9pU&list=PLgUwDviBIf0p4ozDR_kJJkONnb1wdx2Ma&index=52
2. N queens Problem
https://www.youtube.com/watch?v=i05Ju7AftcM&list=PLgUwDviBIf0p4ozDR_kJJkONnb1wdx2Ma&index=57
3. Sudoku Solver
https://www.youtube.com/watch?v=FWAlf_EVUKE&list=PLgUwDviBIf0p4ozDR_kJJkONnb1wdx2Ma&index=58
4. M coloring Problem
5. Rat in a Maze
6. Word Break (print all ways)
TUF
0:05 / 24:36

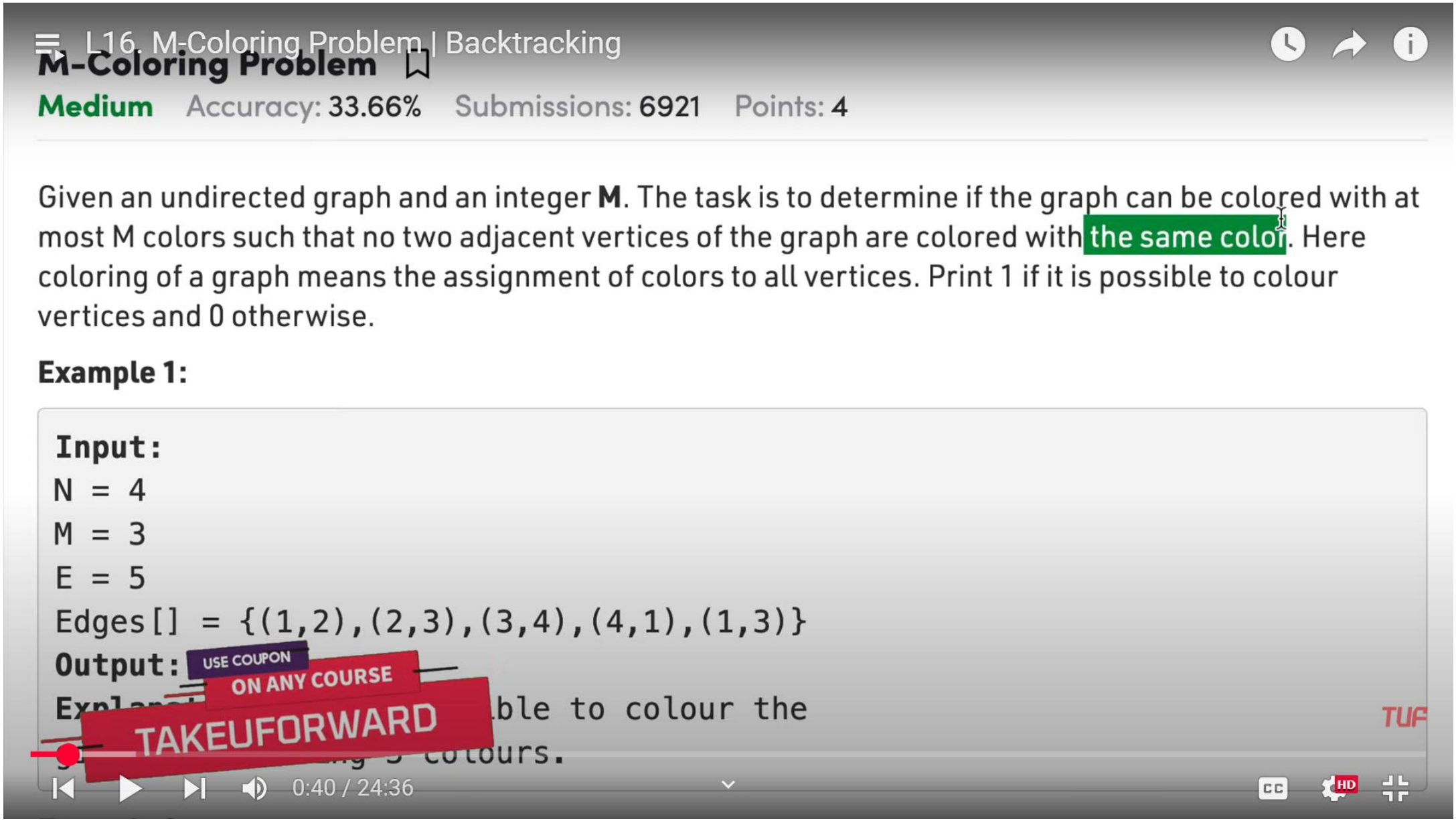
M-Coloring Problem

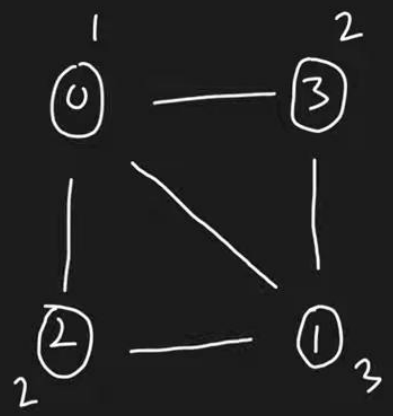
Medium Accuracy: 33.66% Submissions: 6921 Points: 4

Given an undirected graph and an integer **M**. The task is to determine if the graph can be colored with at most M colors such that no two adjacent vertices of the graph are colored with the same color. Here coloring of a graph means the assignment of colors to all vertices. Print 1 if it is possible to colour vertices and 0 otherwise.

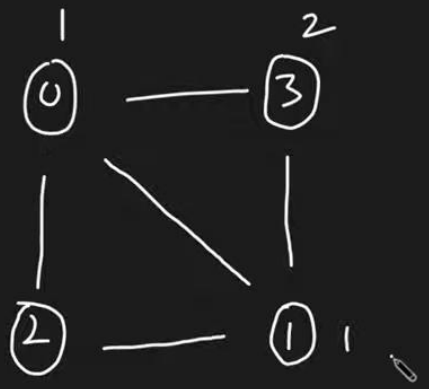
Example 1:

```
Input:
N = 4
M = 3
E = 5
Edges[] = {(1,2),(2,3),(3,4),(4,1),(1,3)}
Output: 1
Explanation: It is possible to colour the
given graph using 3 colours.
```

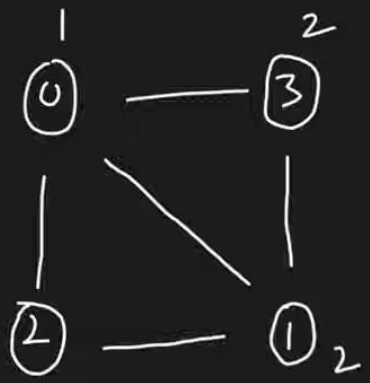

L16. M-Coloring Problem | Backtracking
M-Coloring Problem
Medium Accuracy: 33.66% Submissions: 6921 Points: 4
Given an undirected graph and an integer M. The task is to determine if the graph can be colored with at most M colors such that no two adjacent vertices of the graph are colored with the same color. Here coloring of a graph means the assignment of colors to all vertices. Print 1 if it is possible to colour vertices and 0 otherwise.
Example 1:
Input:
N = 4
M = 3
E = 5
Edges[] = {(1,2),(2,3),(3,4),(4,1),(1,3)}
Output:
USE COUPON
ON ANY COURSE
TAKEUFORWARD
TUF
0:40 / 24:36
CC
HD



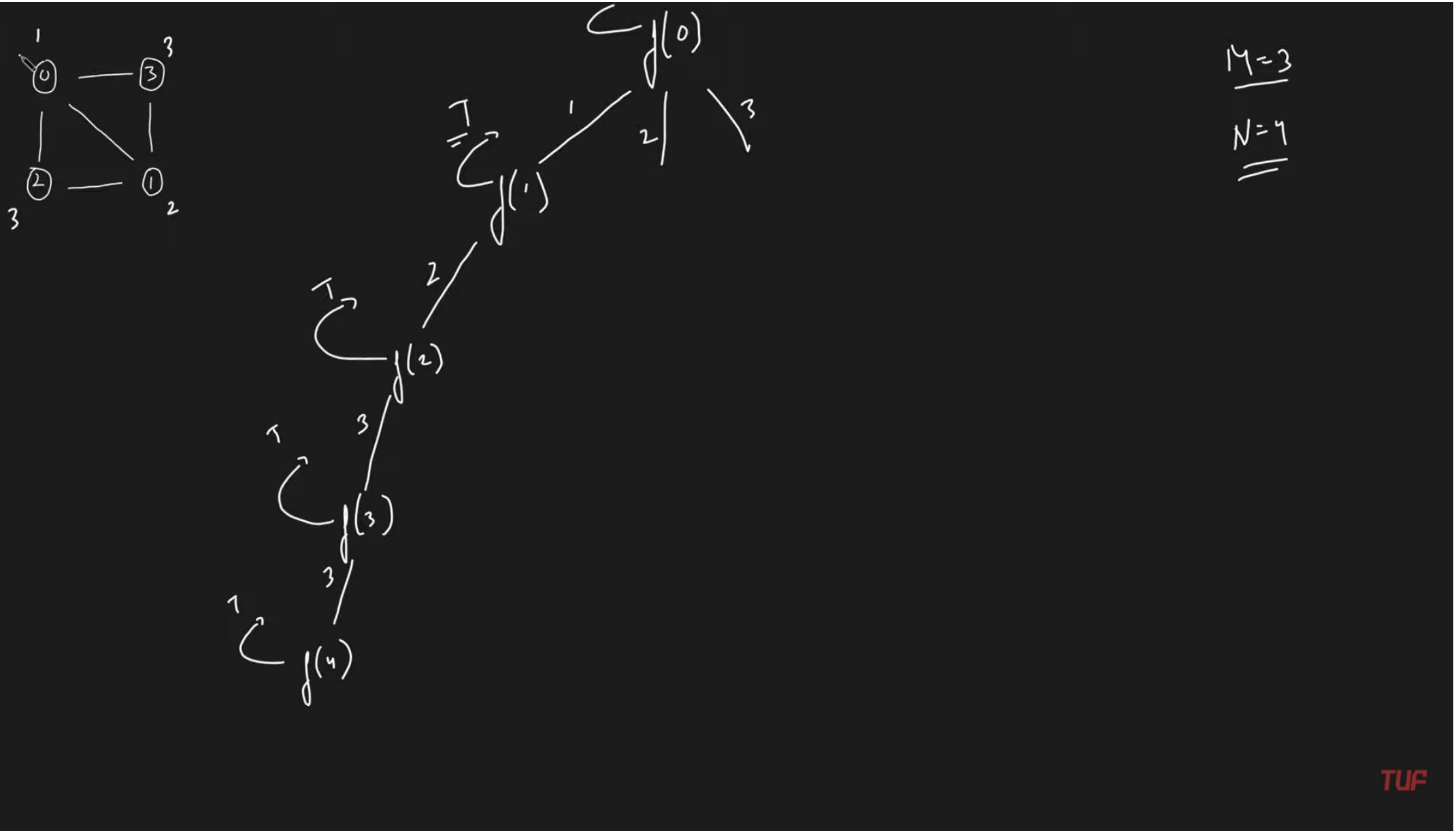
M = 3

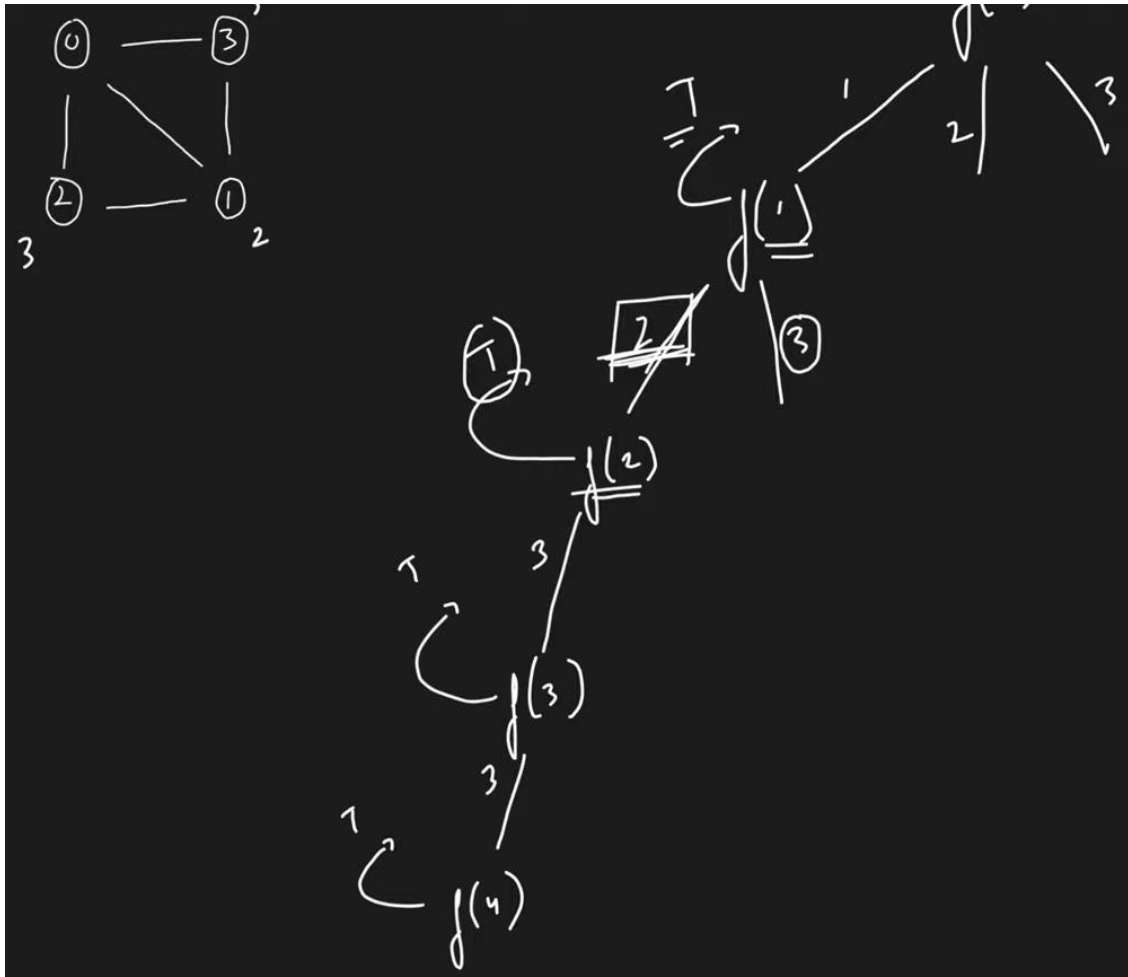


M = 2



M = 2


M=3
N=4
TUF

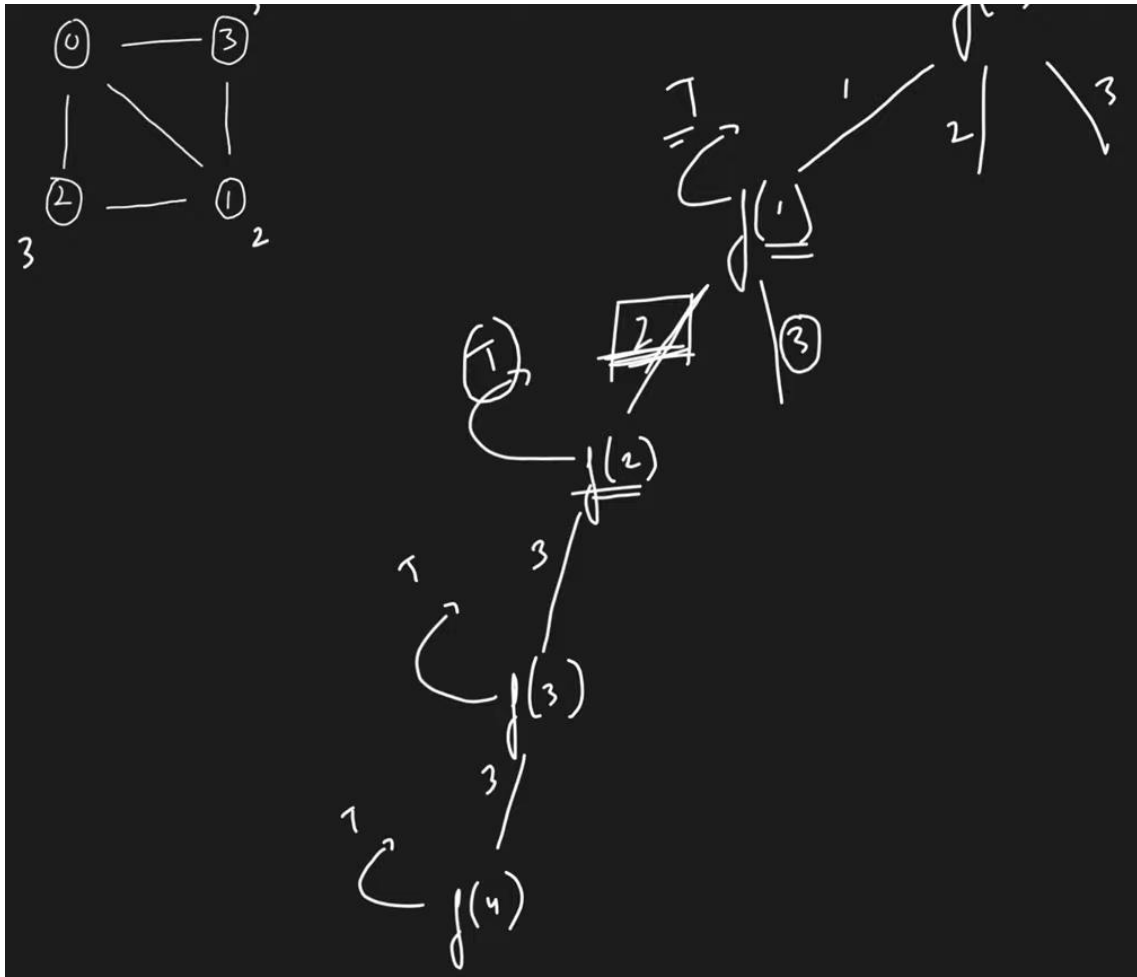


m=3

N=4

```
f(node)
{
  if(node == N) return T;

  for(col = 1 → m)
  {
    if(possible → ✓)
    {
      color[node] = col;
      if(f(node+1) == T)
        return T;
      color[node] = 0;
    }
  }
```

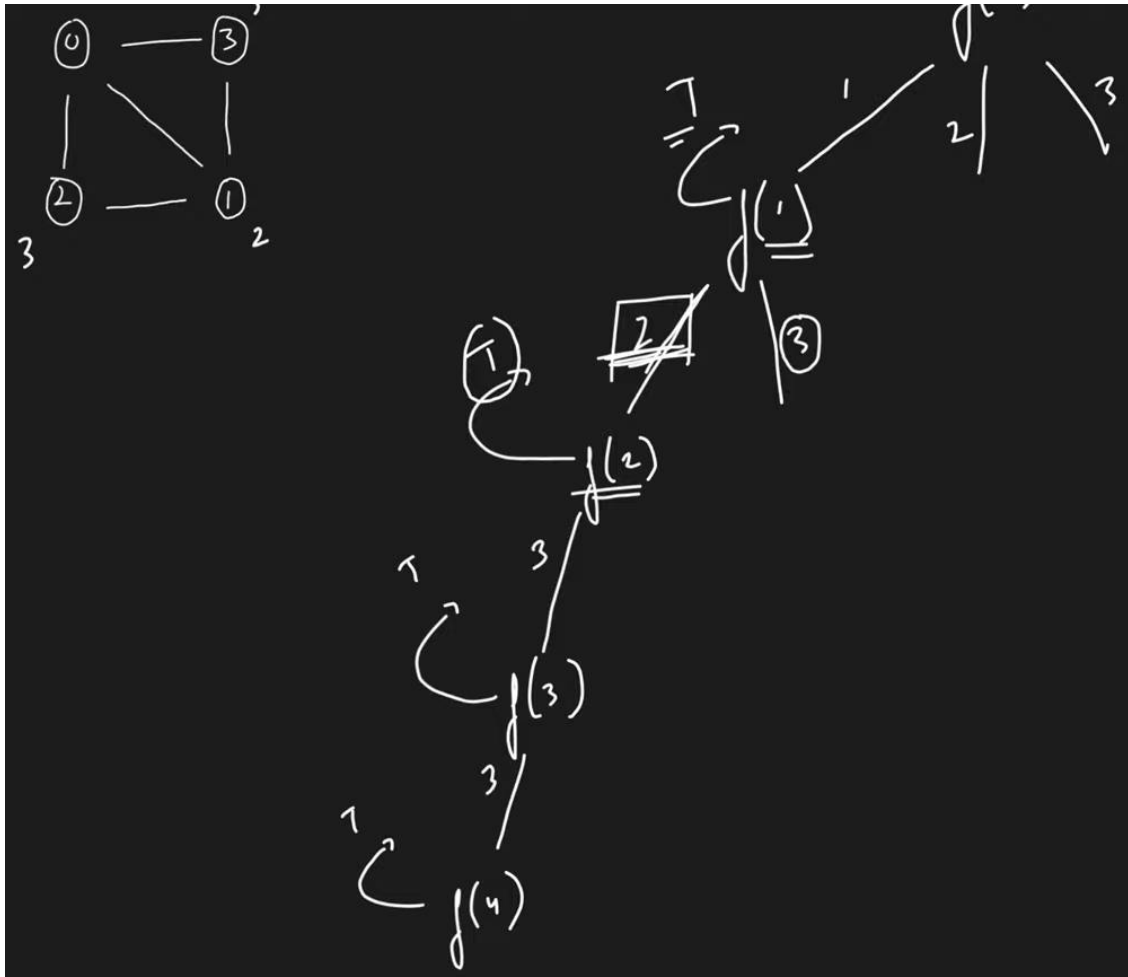
TUF



M=3

N=4

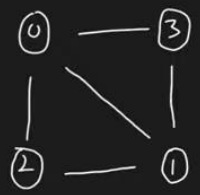
```
f(node)
{
  if(node == N) return T;

  for(col = (1 → m))
  {
    if(possible → v)
    {
      color[node] = col;
      if(f(node+1) == T)
        return T;
      color[node] = 0;
    }
  }
```



M=3

N=4

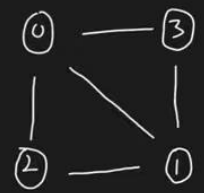
```
f(node)
{
  if(node == N) return T;

  for(col = (1 → m))
  {
    if(possible → v)
    {
      color[node] = col;
      if(f(node+1) == T)
        return T;
      color[node] = 0;
    }
  }
  return F;
}
```



M=3 2

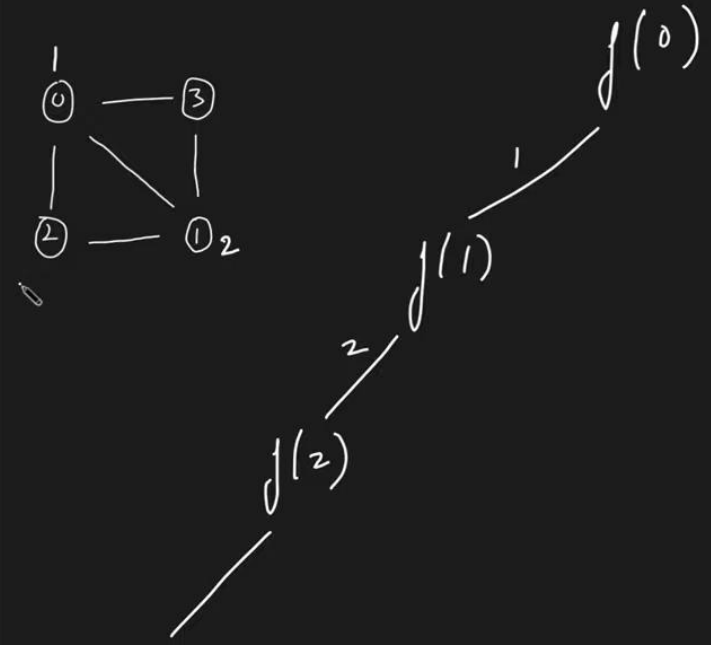
N=4

```
f(node)
{
  if(node == N) return T;

  for(col = (1 → m))
  {
    if(possible → v)
    {
      color[node] = col;
      if(f(node+1) == T)
        return T;
      color[node] = 0;
    }
  }
  return F;
}
```



M = 3
N = 4

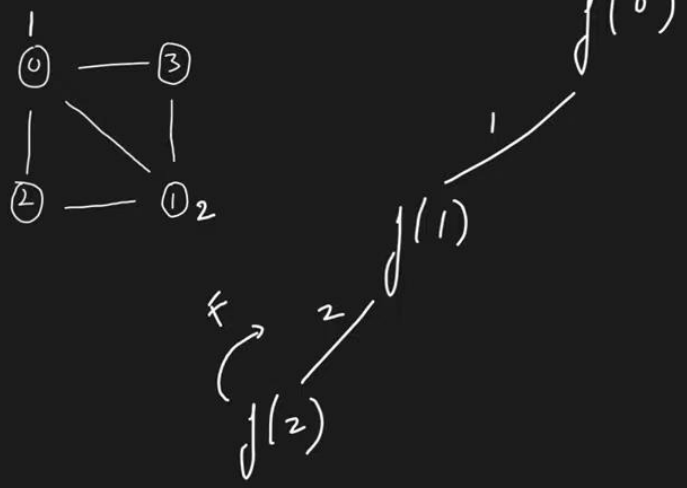
```
f(node)
{
  if(node == N) return T;

  for(col = (1 → m))
  {
    if(possible → v)
    {
      color[node] = col;
      if(f(node+1) == T)
        return T;
      color[node] = 0;
    }
  }
  return F;
}
```



M = 3

N = 4

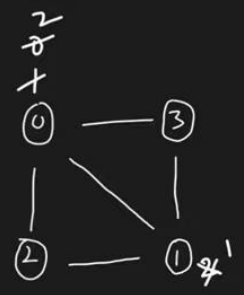
```
f(node)
{
  if(node == N) return T;

  for(col = (1 → m))
  {
    if(possible → v)
    {
      color[node] = col;
      if(f(node+1) == T)
        return T;
      color[node] = 0;
    }
  }
  return F;
}
```

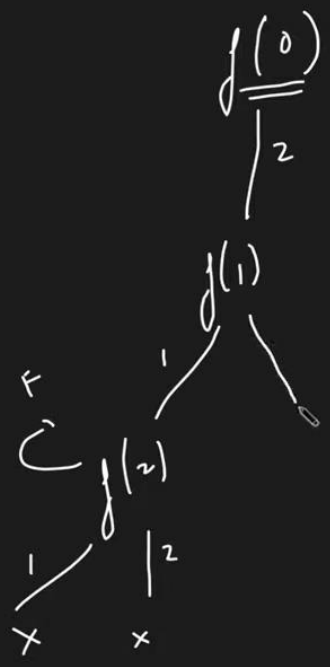


M = 2

N = 4

```
f(node)
{
  if(node == N) return T;

  for(col = (1 → m))
  {
    if(possible → v)
    {
      color[node] = col;
      if(f(node+1) == T)
        return T;
      color[node] = 0;
    }
  }
  return F;
}
```





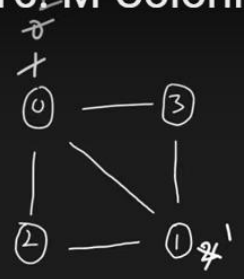
N=4

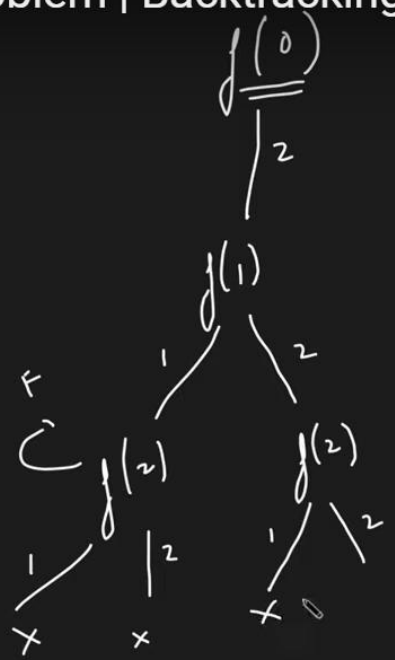
```
f(node)
{
  if(node == N) return T;

  for( col = (1 → m))
  {
    if ( possible → ✓)
    {
      color[node] = col;
      if ( f(node+1) == T)
        return T;
      color[node] = 0;
    }
  }
  return F;
}
```

f(2) is possible ?

At f(1), we cannot use color 2, since adjacent nodes cannot have same color, I in a hurry just drew the tree of f(2), ignore that

L16. M-Coloring Problem | Backtracking






```
f(node)
{
  if(node == N) return T;

  for( col = (1 → m))
  {
    if ( possible → V)
    {
      color[node] = col;
      if ( f(node+1) == T)
```

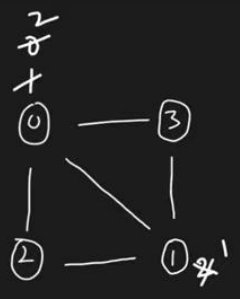
N=4

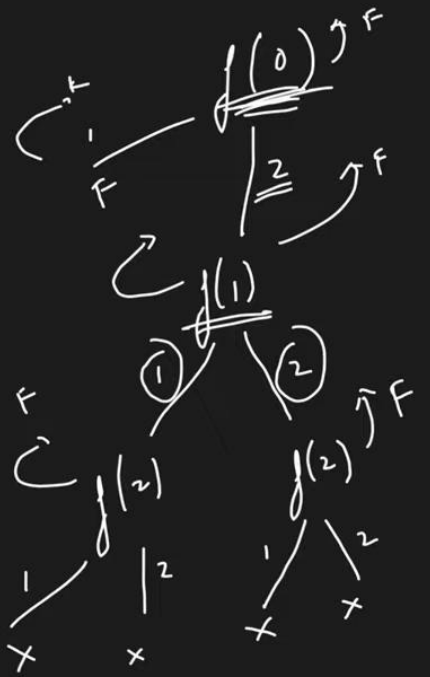
f(2) is possible ?

At f(1), we cannot use color 2, since adjacent nodes cannot have same color, I in a hurry just drew the tree of f(2), ignore that

TUF

16:08 / 24:36





TC. → (N^m)

SC → O(N) + O(N)

N=4

```
f(node)
{
    if(node == N) return T;

    for(col = (1 → m))
    {
        if(possible → ✓)
        {
            color[node] = col;
            if(f(node+1) == T)
                return T;
            color[node] = 0;
        }
    }
    return F;
}
```

```java
class solve
{
    private static boolean isSafe(int node, List<Integer>[] G, int[] color, int n, int col) {
        for(int it: G[node]) {
            if(color[it] == col) return false;
        }
        return true;
    }
    private static boolean solve(int node, List<Integer>[] G, int[] color, int n, int m) {
        if(node == n) return true;

        for(int i = 1;i<=m;i++) {
            if(isSafe(node, G, color, n, i)) {
                color[node] = i;
                if(solve(node+1, G, color, n, m) == true) return true;
                color[node] = 0;
            }
        }
        return false;
    }
    public static boolean graphColoring(List<Integer>[] G, int[] color, int i, int m)
    {
        int n = G.length;
        if(solve(i, G, color, n, m) == true) return true;
        return false;
        // Your code here
    }
}
```

```java
class solve
{
    private static boolean isSafe(int node, List<Integer>[] G, int[] color, int n, int col) {
        for(int it: G[node]) {
            if(color[it] == col) return false;
        }
        return true;
    }
    private static boolean solve(int node, List<Integer>[] G, int[] color, int n, int m) {
        if(node == n) return true;

        for(int i = 1;i<=m;i++) {
            if(isSafe(node, G, color, n, i)) {
                color[node] = i;
                if(solve(node+1, G, color, n, m) == true) return true;
                color[node] = 0;
            }
        }
        return false;
    }
    public static boolean graphColoring(List<Integer>[] G, int[] color, int i, int m)
    {
        int n = G.length;
        if(solve(i, G, color, n, m) == true) return true;
        return false;
        // Your code here
    }
}
```

```
10
11  bool isSafe(int node, int color[], bool graph[101][101], int n, int col) {
12      for(int k = 0;k<n;k++) {
13          if(k != node && graph[k][node] == 1 && color[k] == col) {
14              return false;
15          }
16      }
17      return true;
18  }
19  bool solve(int node, int color[], int m, int N, bool graph[101][101]) {
20      if(node == N) {
21          return true;
22      }
23
24      for(int i = 1;i<=m;i++) {
25          if(isSafe(node, color, graph, N, i)) {
26              color[node] = i;
27              if(solve(node+1, color, m, N, graph)) return true;
28              color[node] = 0;
29          }
30
31      }
32      return false;
33  }
34
35  //Function to determine if graph can be coloured with at most M colours such
36  //that no two adjacent vertices of graph are coloured with same colour.
37  bool graphColoring(bool graph[101][101], int m, int N)
38  {
39      int color[N] = {0};
40      if(solve(0,color,m,N,graph)) return true;
41      return false;
42  }
43  // } Driver Code Ends
```