

Enhancing Crowdsourcing with the Zero-Determinant Game Theory

Dissertation Proposal

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Committee:

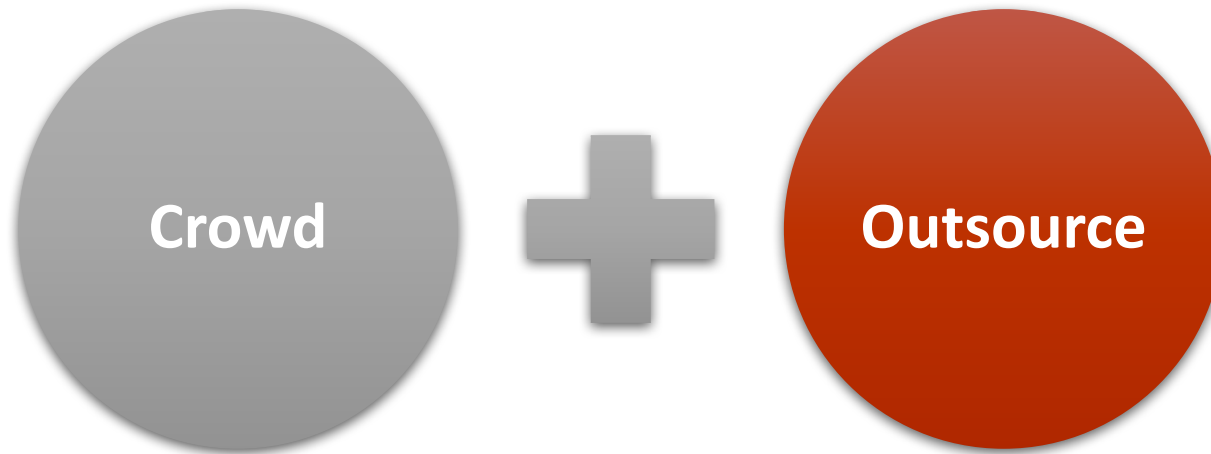
Xiuzhen Cheng

Hyeong-Ah Choi

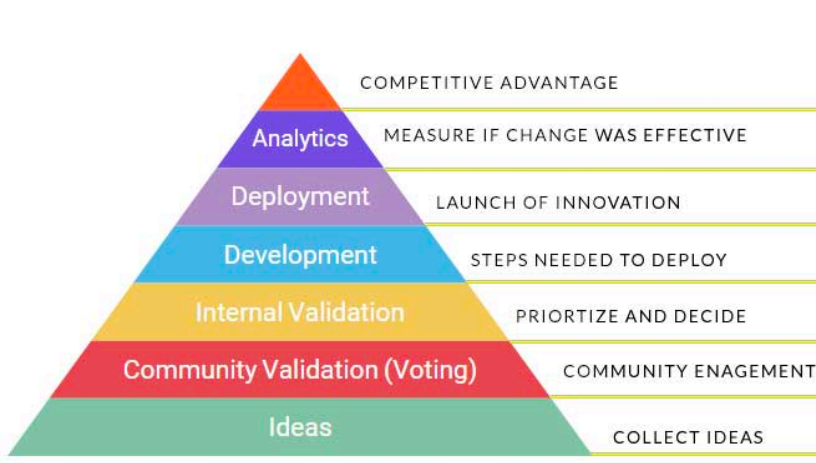
Arkady Yerukhimovich

Xiang Chen

What is crowdsourcing?



Accomplish a complex task via eliciting services from a large group of contributors.



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Challenges



- Skills
- Intents
- Backgrounds
- ...



Overview

- **Eliminating malicious attacks in crowdsourcing**
 - Problem Formulation
 - Strategies for the Worker
 - Intuitive Strategies for the Requestor
 - ZD-Based Algorithm
 - Evaluation
- **Quality control of crowdsourcing**
 - Problem Formulation
 - Extension of the Sequential ZD
 - Sequential ZD based Algorithms and the Evaluation

Problem Formulation



Problem Formulation

Requestor

Success

- Gather distributed human resources

Openness

- Open to all crowds

Vulnerability

- Malicious attacks

- ❑ DARPA Network Challenge, 2009
- ❑ DARPA Shredder Challenge, 2011, UCSD

Worker

Problem Formulation

Requestor

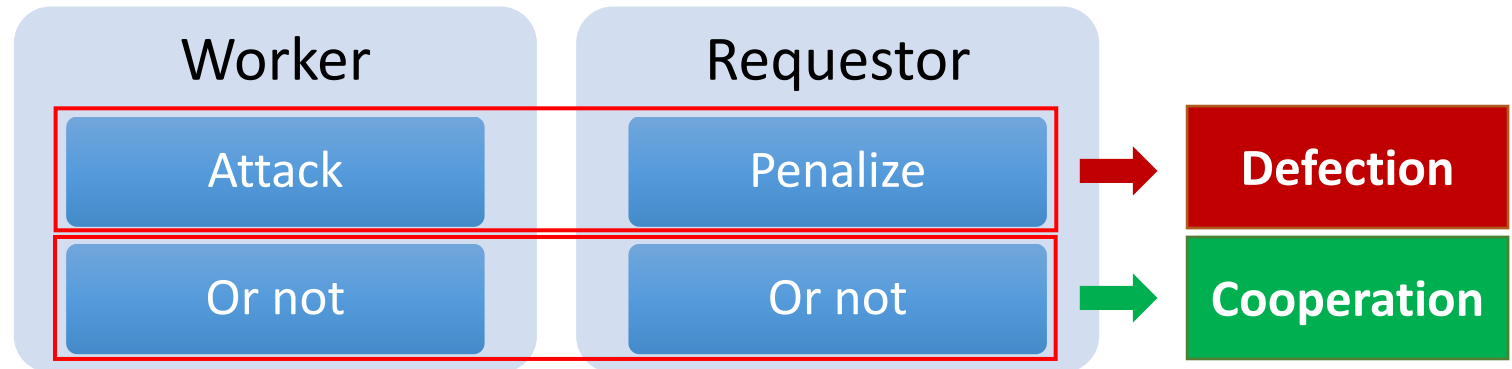


Worker

- Game theory
 - The requestor and any one of the worker
 - Two-player simultaneous game
 - Round by round

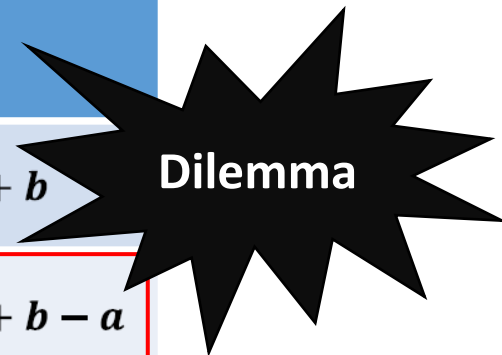
Problem Formulation

- Actions

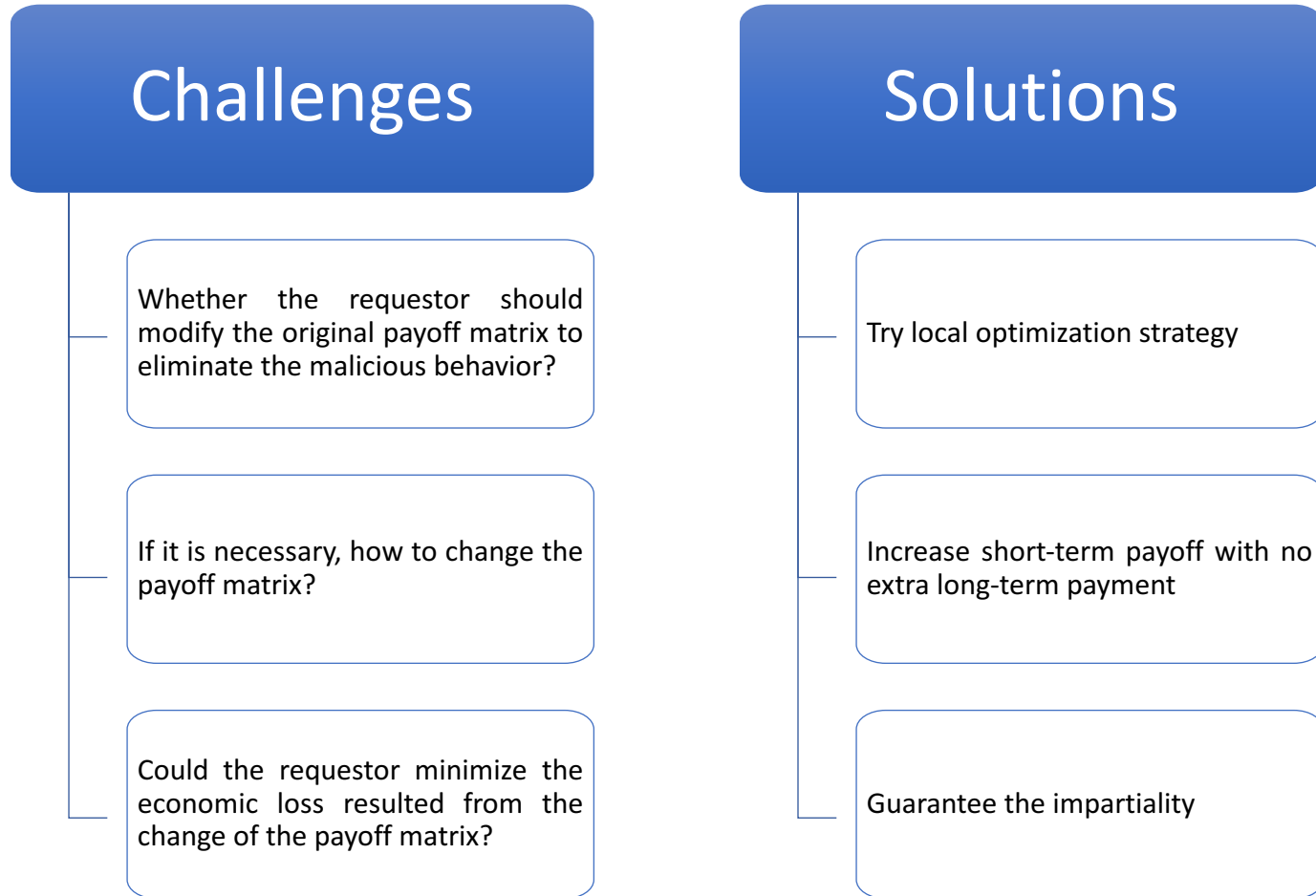


- Payoff Matrix

Requestor \ Worker	Cooperation	Defection
Cooperation	R_r, R_w	$R_r - m, R_w + b$
Defection	$R_r + n, R_w - a$	$R_r - m + n, R_w + b - a$



Problem Formulation



Strategy of the Worker

- Situation

Lack global
information

```
graph TD; A[Lack global information] --> B[Adaptively search best action]; B --> C[Natural selection, survival of the fittest];
```

The diagram illustrates a three-step process for a worker's strategy. It begins with a grey box labeled 'Lack global information'. A white arrow points down to an orange box labeled 'Adaptively search best action'. Another white arrow points down from the orange box to a red box labeled 'Natural selection, survival of the fittest'.

Adaptively search
best action

Natural selection,
survival of the fittest

- Strategy

Evolutionary Player

- Adjust strategy to maximize payoff regardless of the opponent's strategy and/or payoff
- Two examples
 - E Strategy
 - E' Strategy

Strategy of the Worker

- E Strategy

- Cooperation probability $q_w^{t+1} = q_w^t \frac{W_c^t}{E_w^t}$

Expected cooperation payoff $W_c^t = p_r^t E(cc) + (1 - p_r^t) E(cd)$

Total expected payoff $E_w^t = q_w^t W_c^t + (1 - q_w^t) W_d^t$

- E' Strategy

- Cooperation probability $q_w^t = \frac{e^{A_c^t - A_d^t}}{1 + e^{A_c^t - A_d^t}}$

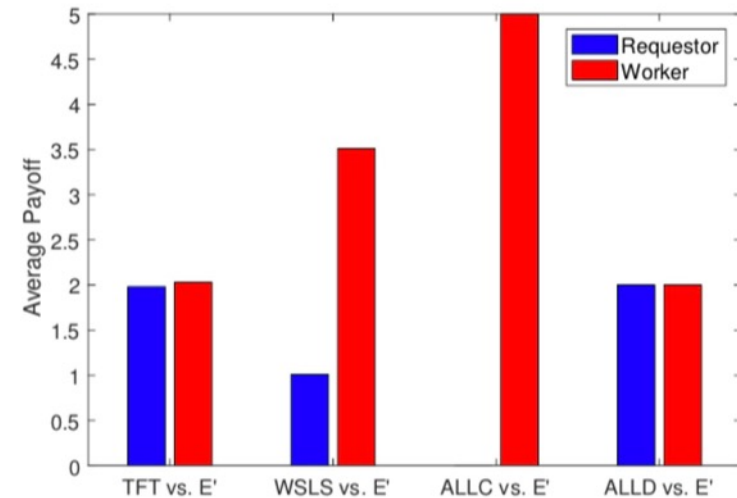
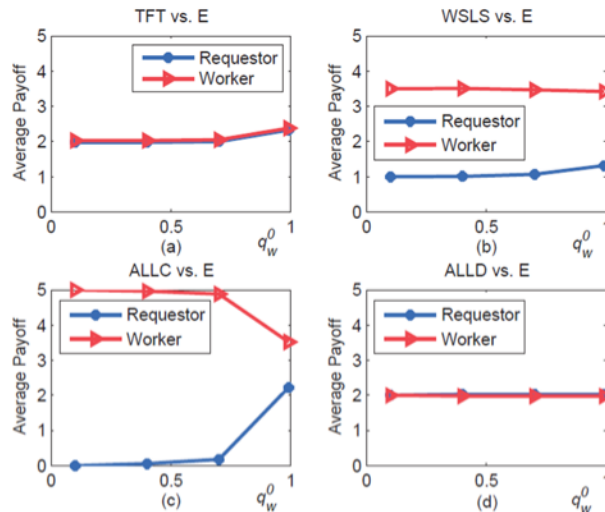
Accumulative expected cooperation payoff $A_c^t = \sum_{\tau=0}^t W_c^\tau$

Accumulative expected defection payoff $A_d^t = \sum_{\tau=0}^t W_d^\tau$

Strategy of the Worker

- Reasonability
 - Play with the classic-strategy requestor
 - TFT: tit-for-tat
 - WSLS: win-stay-lose-shift
 - ALLC: all cooperation
 - ALLD: all defection

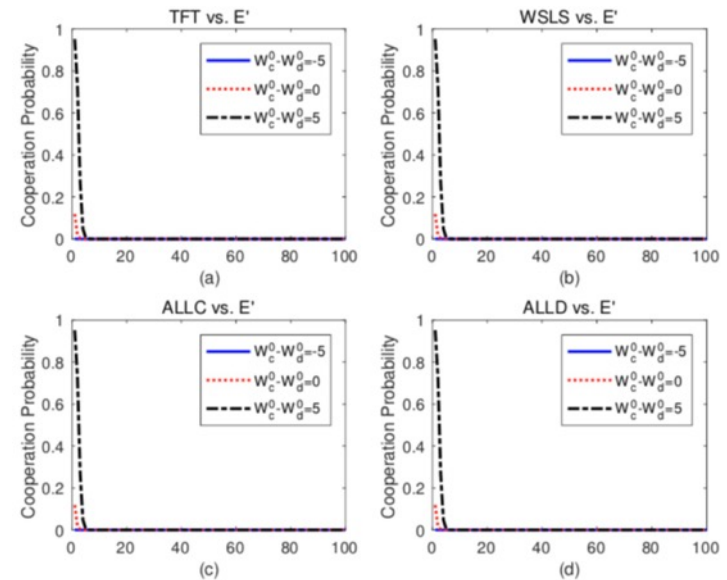
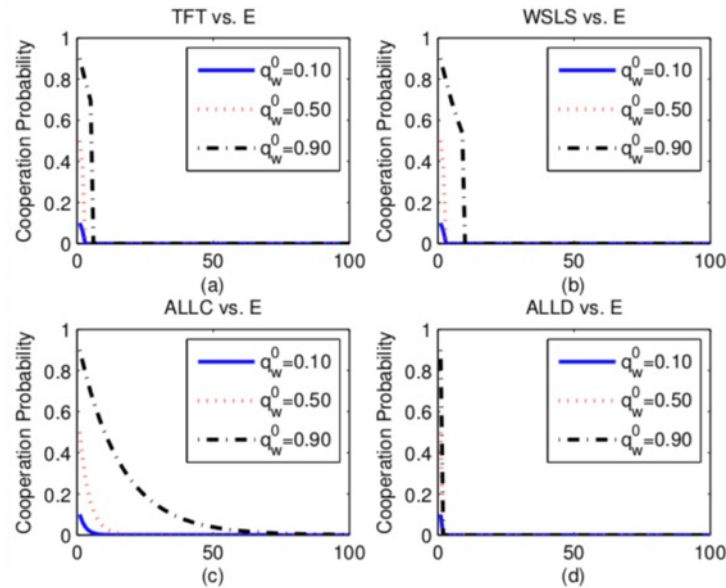
Theoretical analysis also proved these results.



Intuitive Strategies for the Requestor

- What should the requestor do when faced with such a strong opponent?

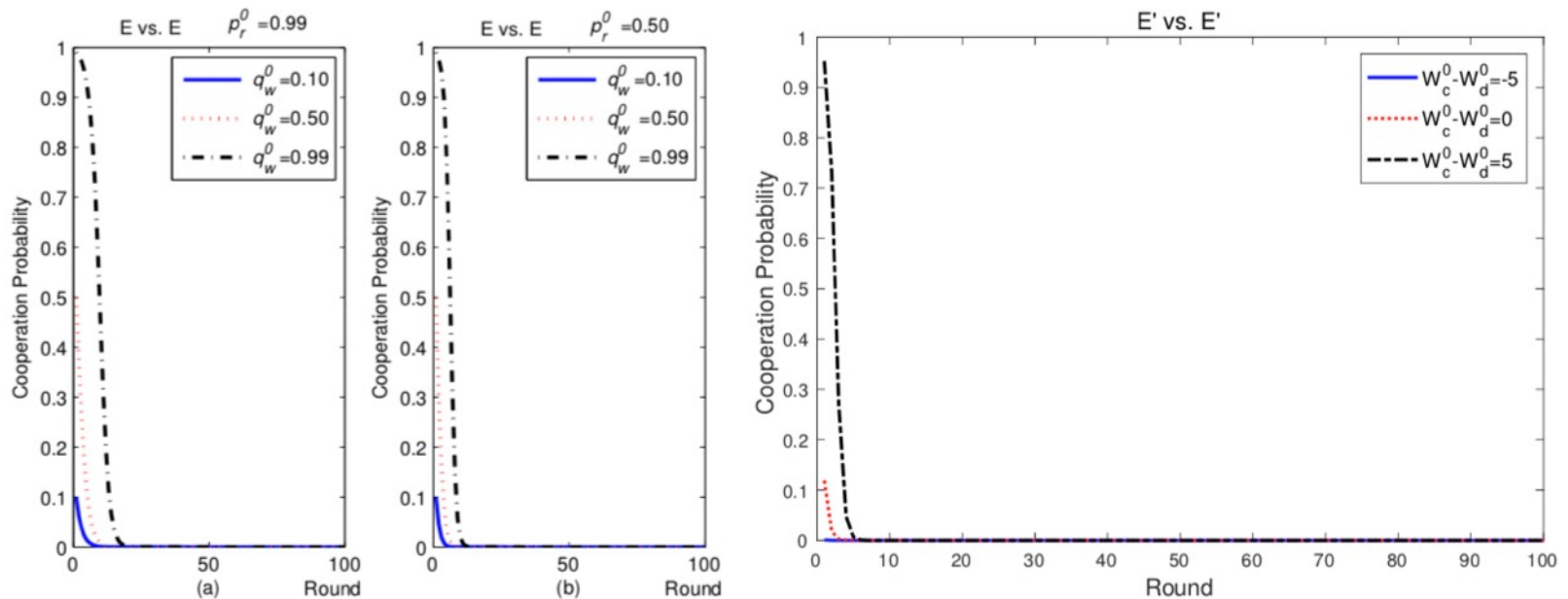
Classic Strategies



Intuitive Strategies for the Requestor

- What should the requestor do when faced with such a strong opponent?

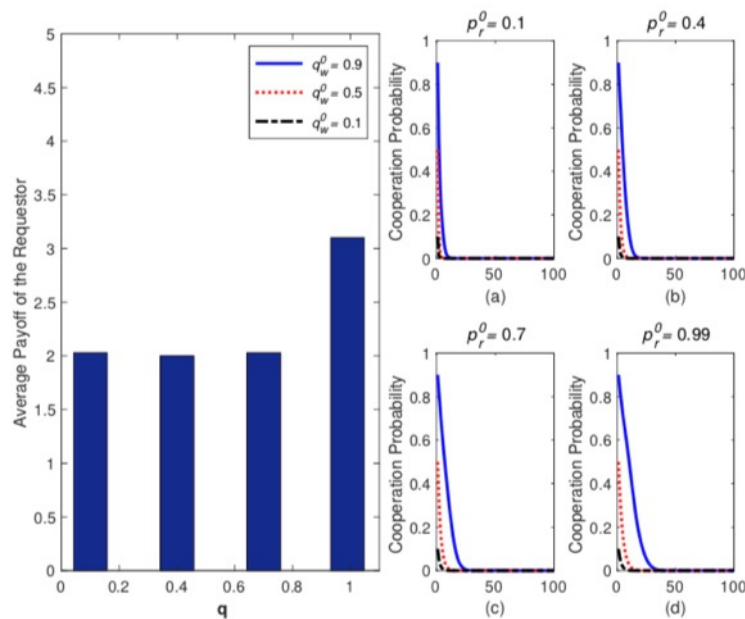
Evolutionary
Strategies



Intuitive Strategies for the Requestor

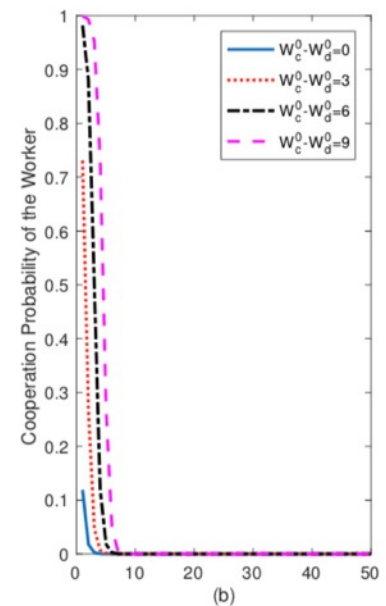
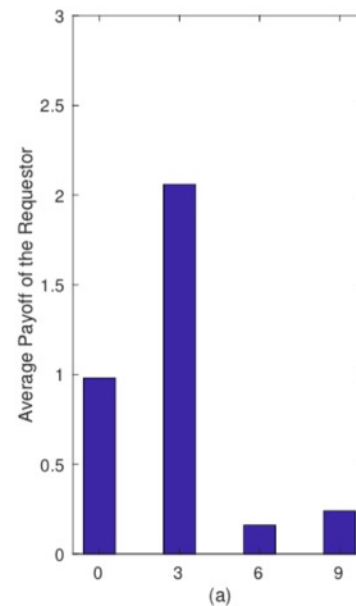
- What should the requestor do when faced with such a strong opponent?

Local
Optimization
Strategy



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Intuitive Strategies for the Requestor

- Failure analysis

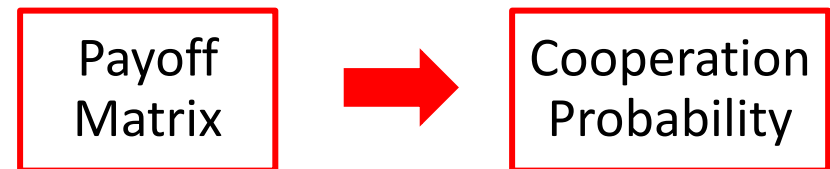
- E strategy

$$W_c^t \leq E_w^t \text{ for } p_r^t, q_w^t \in [0, 1]$$
$$\lim_{t \rightarrow \infty} q_w^{t+1} = q_w^t \frac{W_c^t}{E_w^t} = 0$$

- E' strategy

$$A_c^t - A_d^t \downarrow$$
$$\lim_{t \rightarrow \infty} q_w^t = \frac{e^{A_c^t - A_d^t}}{1 + e^{A_c^t - A_d^t}} = 0$$

- Key factor



ZD-Based Algorithm

- Aim
 - Elicits the worker's cooperation by increasing its short-term payoff but no any extra long-term payment
- Zero-determinant strategy
 - Proposed by Press and Dyson (2012, PNAS)
 - Set strategy satisfying required condition, we have

$$\forall t, \alpha E_r^t + \beta E_w^t + \gamma = 0$$

- Pinning strategy $\alpha = 0, E_w^t = -\frac{\gamma}{\beta}$
- Extortion strategy $\alpha = \phi, \beta = -\phi\chi, \gamma = \phi(\chi(R_w + b - a) - (R_r - m + n)),$

$$E_r^t - (R_r - m + n) = \chi(E_w^t - (R_w + b - a)), \chi \geq 1$$

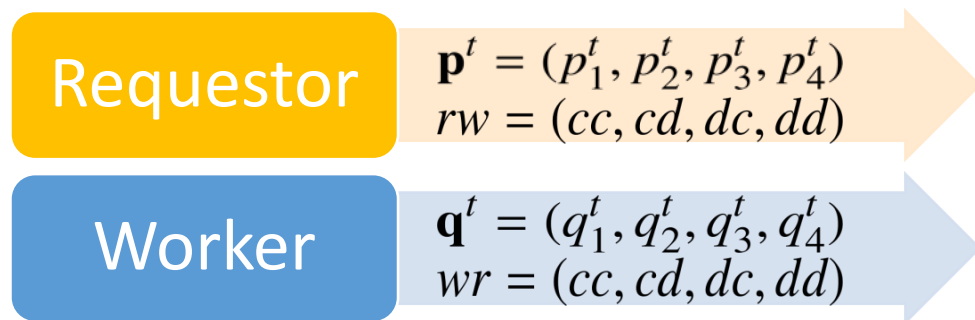


ZD player's
strategy

Opponent's
expected
payoff

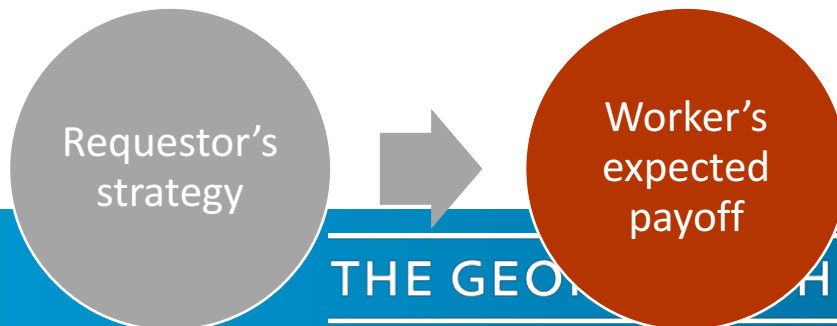
ZD-Based Algorithm

- Mixed strategy



- Expected payoff $E_r^t = \frac{D(\mathbf{p}^t, \mathbf{q}^t, \mathbf{S}_r)}{D(\mathbf{p}^t, \mathbf{q}^t, \mathbf{1})}, E_w^t = \frac{D(\mathbf{p}^t, \mathbf{q}^t, \mathbf{S}_w)}{D(\mathbf{p}^t, \mathbf{q}^t, \mathbf{1})}$
- ZD strategy $\tilde{\mathbf{p}}^t = (-1 + p_1^t, -1 + p_2^t, p_3^t, p_4^t) = \beta \mathbf{S}_w + \gamma \mathbf{1}$

or $\phi[(\mathbf{S}_r - (R_r - m + n)) - \chi(\mathbf{S}_w - (R_w + b - a))]$



ZD-Based Algorithm

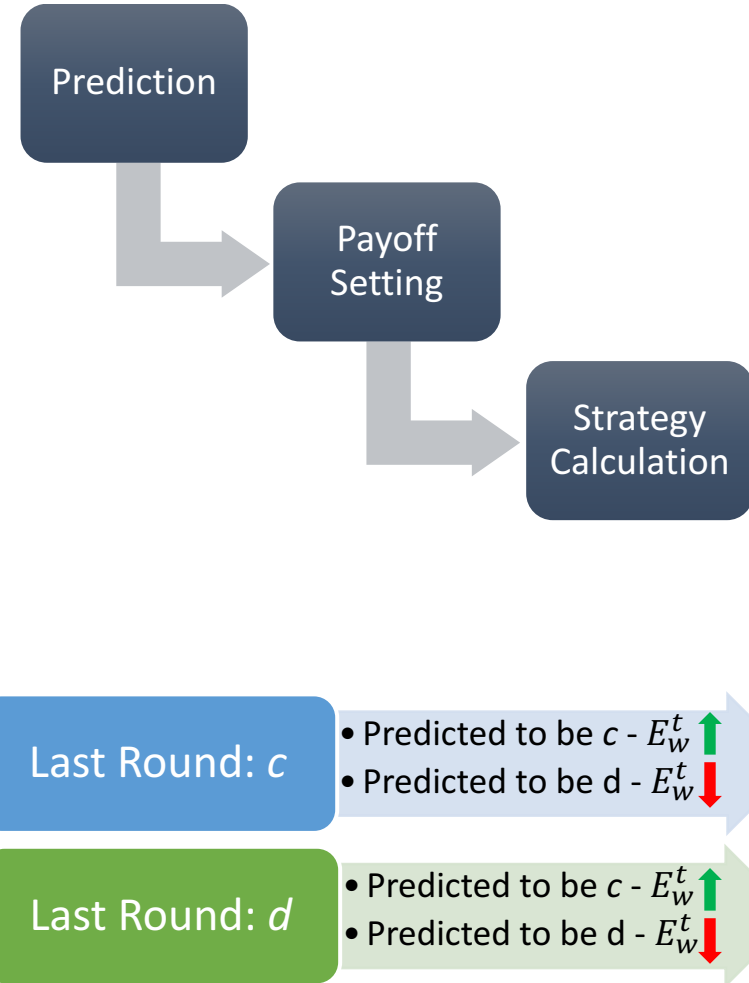
Algorithm 2 Reward-Penalty Expected Payoff Algorithm based on ZD Strategies

Require: p^t : the requestor's strategy at round t and its initial value is the one used in the last round of the preparatory stage; $\mathbf{P}_s = (P_{cc}, P_{cd}, P_{dc}, P_{dd})$: the state transition probability of the worker, and its initial value is statistically calculated by the data collected in the preparatory stage; N : the total number of rounds.

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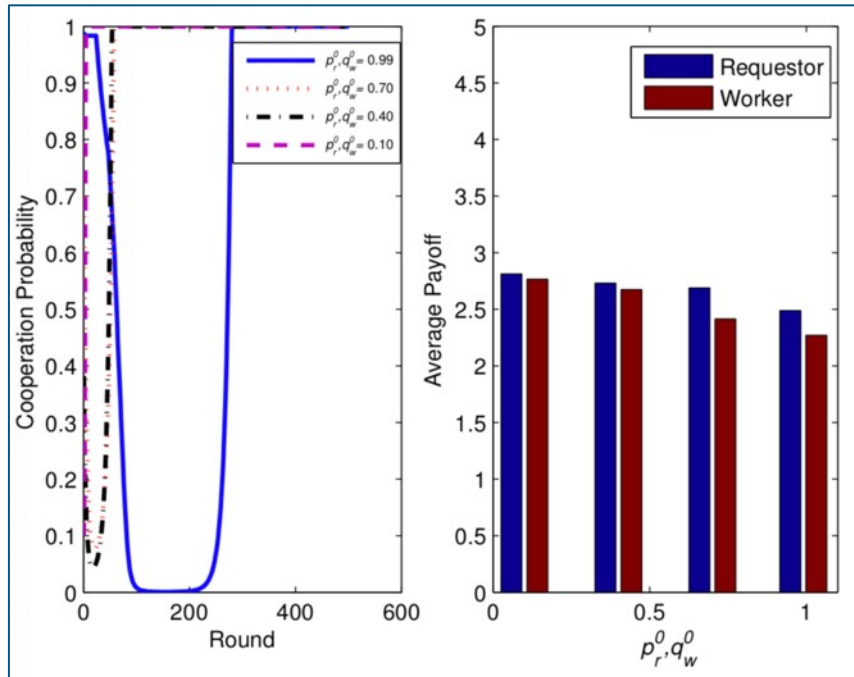
1: Initialize( $E_w^0$ )
2: for  $t = 1$  to  $N$  do
3:   if The worker's previous move is  $c$  then
4:     if  $P_{cc} > P_{cd}$  then
5:       Calculate  $p^t$  which makes
6:        $E_w^t \leftarrow E_w^{t-1} + (max - E_w^{t-1})/2$ 
7:     else
8:       Calculate  $p^t$  which makes
9:        $E_w^t \leftarrow E_w^{t-1} - (E_w^{t-1} - min)/2$ 
10:    end if
11:  else
12:    if  $P_{dc} > P_{dd}$  then
13:      Calculate  $p^t$  which makes
14:       $E_w^t \leftarrow E_w^{t-1} + (max - E_w^{t-1})/2$ 
15:    else
16:      Calculate  $p^t$  which makes
17:       $E_w^t \leftarrow E_w^{t-1} - (E_w^{t-1} - min)/2$ 
18:    end if
19:  end if
20:  if The current round terminates then
21:    Update  $\mathbf{P}_s$ 
22:  end if
23: end for
    
```

$\triangleright P_{cc} \leq P_{cd}.$
 \triangleright The worker's previous move is d .
 $\triangleright P_{dc} \leq P_{dd}.$

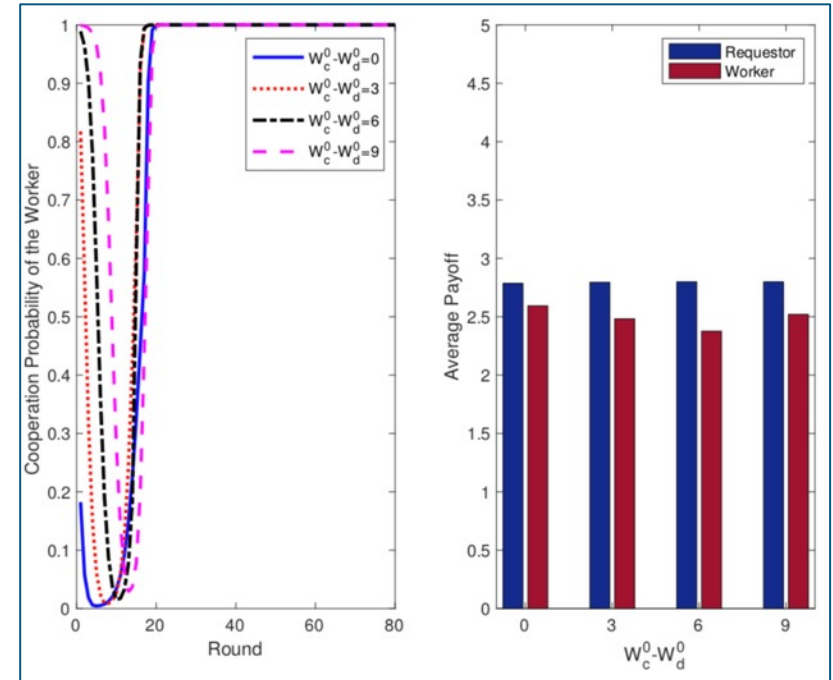


Evaluation

Theoretically proved



Effectiveness



Fairness

Liveliness

Overview

- **Eliminating malicious attacks in crowdsourcing**

- Problem Formulation
- Strategies for the Worker
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- ZD-Based Algorithm
- Evaluation

- **Quality control of crowdsourcing**

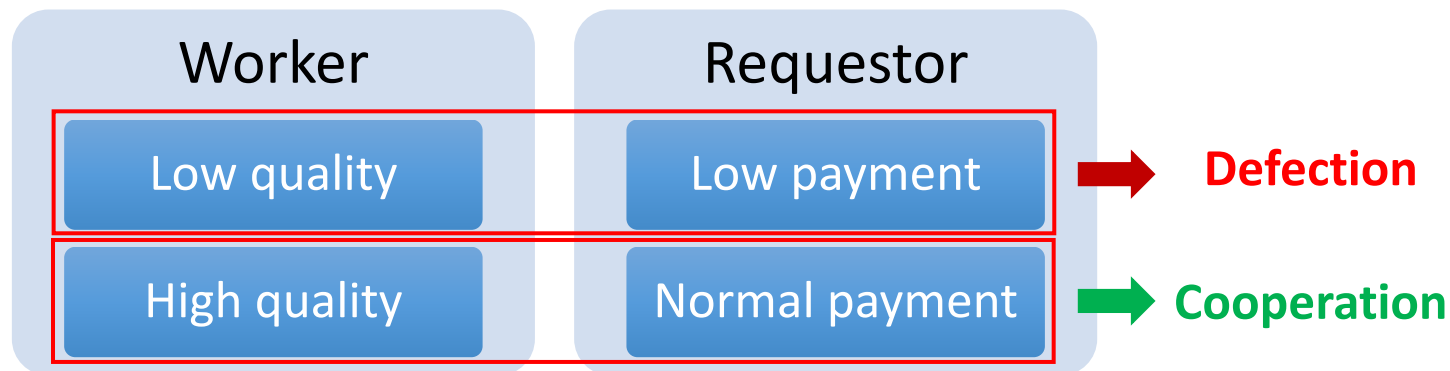
- Problem Formulation
- Extension of the Sequential ZD
- Sequential ZD based Algorithms
- Evaluation

Problem Formulation

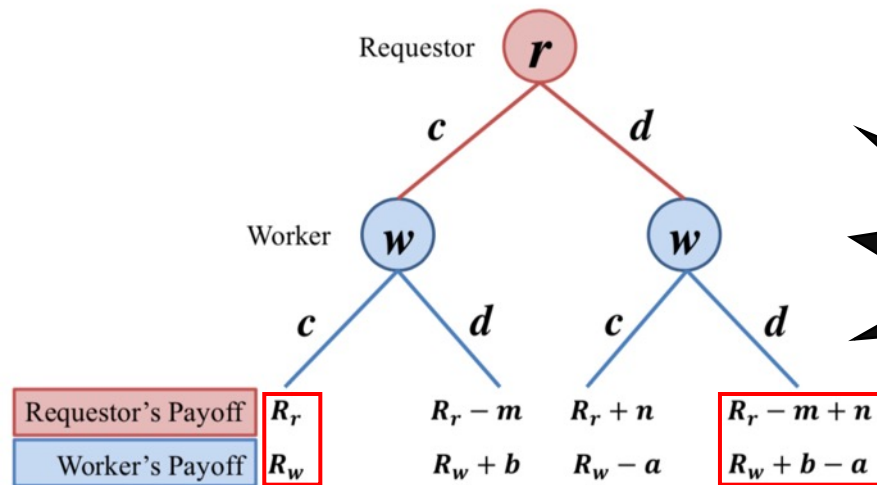


Problem Formulation

- Actions



- Game Tree



Problem Formulation

- Actions



Utility Functions

$$w_r(x, y) = A_r \phi(y) - B_r x$$
$$w_w(x, y) = A_w x - B_w \psi(y)$$



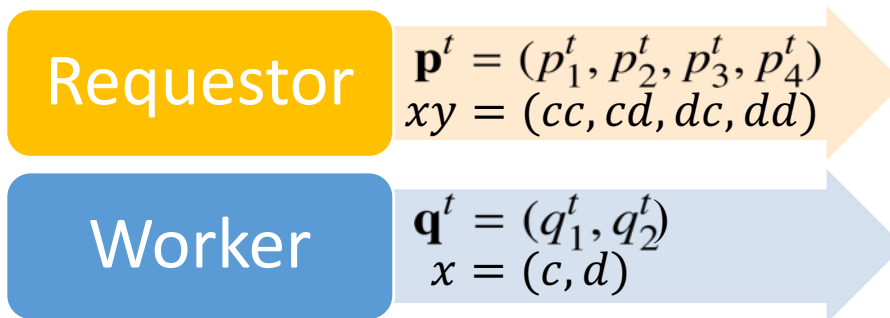
Nash Equilibrium

$$(x^*, y^*) = (l_r, l_w)$$

Sequential dilemma still exists!

Extension of the Sequential ZD

- Mixed strategy



- Only the requestor can adopt the ZD strategy

$$\begin{bmatrix} p_1^t q_1^t - 1 & p_1^t - 1 & q_1^t - 1 & f_1 \\ p_2^t q_3^t & p_2^t - 1 & q_2^t & f_2 \\ p_3^t q_3^t & p_3^t & q_3^t - 1 & f_3 \\ p_4^t q_4^t & p_4^t & q_4^t & f_4 \end{bmatrix} \rightarrow \begin{bmatrix} p_1^t q_1^t - 1 & p_1^t - 1 & (1 - p_1^t)q_2^t + p_1^t q_1^t - 1 & f_1 \\ p_2^t q_1^t & p_2^t - 1 & (1 - p_2^t)q_2^t + p_2^t q_1^t & f_2 \\ p_3^t q_1^t & p_3^t & (1 - p_3^t)q_2^t + p_3^t q_1^t - 1 & f_3 \\ p_4^t q_1^t & p_4^t & (1 - p_4^t)q_2^t + p_4^t q_1^t & f_4 \end{bmatrix}$$

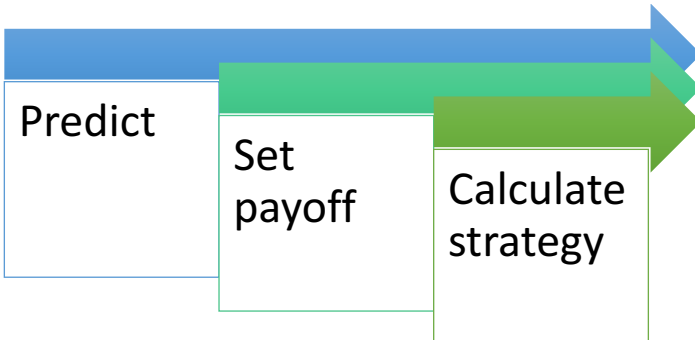
The first mover therefore obtain the advantage.

Sequential ZD based Algorithms

- Binary model

- Evolutionary worker

$$q_w^{t+1} = q_w^t \frac{W_c^t}{\tilde{E}_w^t}$$



Algorithm 3 Sequential ZD based Incentive Algorithm for the Binary Model

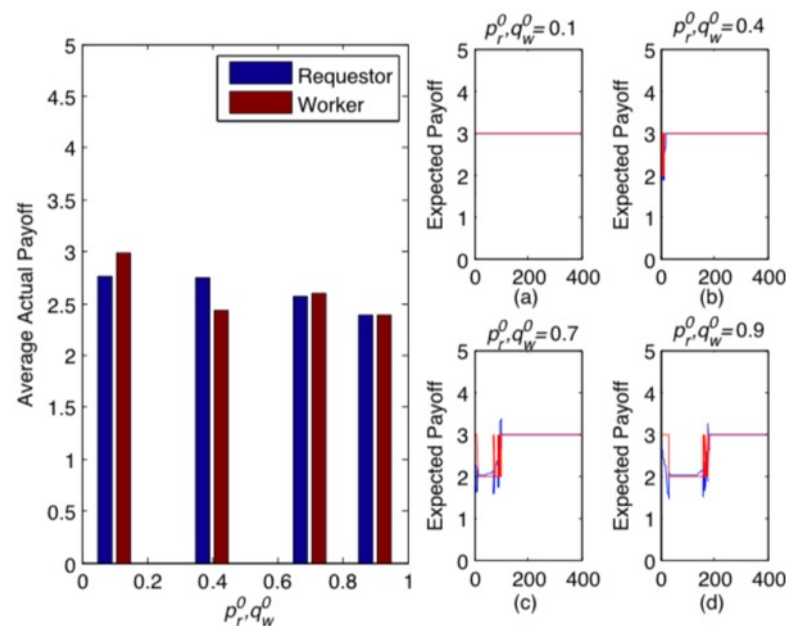
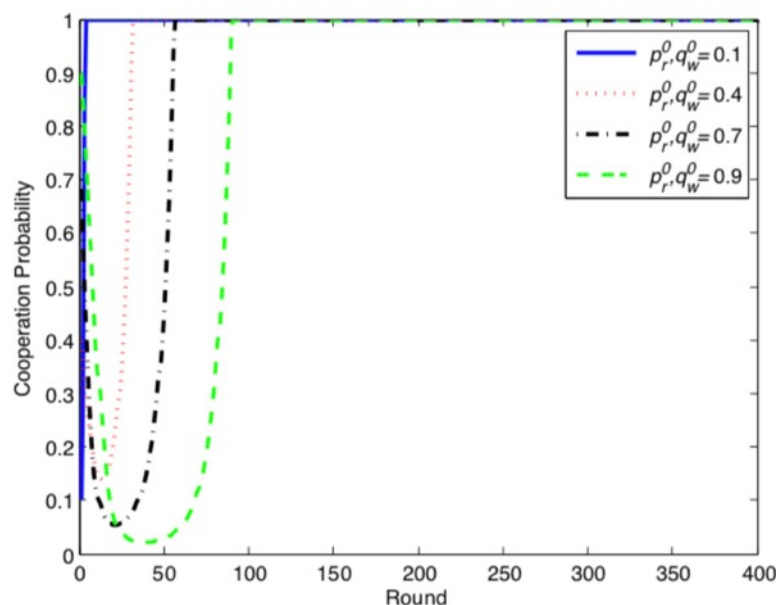
Require: $\mathbf{p}^i = \{p_1^i, p_2^i, p_3^i, p_4^i\}$: the requestor's strategy at round i and its initial values are the ones used in the N_0^{th} round; $\mathbf{P}_s = (P_{cc}, P_{cd}, P_{dc}, P_{dd})$: the state transition probabilities of the worker, and their initial values are calculated statistically through the preparatory N_0 rounds; N : the total number of rounds.

```

1: Initialize ( $\tilde{E}_w^0$ )
2: for  $i = 1$  to  $N$  do
3:   if The worker's last move is  $c$  then
4:     if  $P_{cc} > P_{cd}$  then
5:       Set  $\{p_1^i, p_2^i, p_3^i, p_4^i\} \leftarrow \frac{(1-p_1^i)(R_w+b-a)+p_4^i R_w}{1-p_1^i+p_4^i} = R_w \wedge 0 \leq p_1^i, p_2^i, p_3^i, p_4^i \leq 1\}$ 
6:        $\tilde{E}_w^i \leftarrow R_w$ 
7:     else  $\triangleright P_{cc} \leq P_{cd}$ .
8:       Set  $\{p_1^i, p_2^i, p_3^i, p_4^i\} \leftarrow \frac{(1-p_1^i)(R_w+b-a)+p_4^i R_w}{1-p_1^i+p_4^i} = R_w+b-a \wedge 0 \leq p_1^i, p_2^i, p_3^i, p_4^i \leq 1\}$ 
9:        $\tilde{E}_w^i \leftarrow R_w + b - a$ 
10:    end if
11:  else  $\triangleright$  The worker's last move is  $d$ .
12:    if  $P_{dc} > P_{dd}$  then
13:      Set  $\{p_1^i, p_2^i, p_3^i, p_4^i\} \leftarrow \frac{(1-p_1^i)(R_w+b-a)+p_4^i R_w}{1-p_1^i+p_4^i} = R_w \wedge 0 \leq p_1^i, p_2^i, p_3^i, p_4^i \leq 1\}$ 
14:       $\tilde{E}_w^i \leftarrow R_w$ 
15:    else  $\triangleright P_{dc} \leq P_{dd}$ .
16:      Set  $\{p_1^i, p_2^i, p_3^i, p_4^i\} \leftarrow \frac{(1-p_1^i)(R_w+b-a)+p_4^i R_w}{1-p_1^i+p_4^i} = R_w+b-a \wedge 0 \leq p_1^i, p_2^i, p_3^i, p_4^i \leq 1\}$ 
17:       $\tilde{E}_w^i \leftarrow R_w + b - a$ 
18:    end if
19:  end if
20:  if The current round ends then
21:    Update  $\mathbf{P}_s$ 
22:  end if
23: end for
  
```

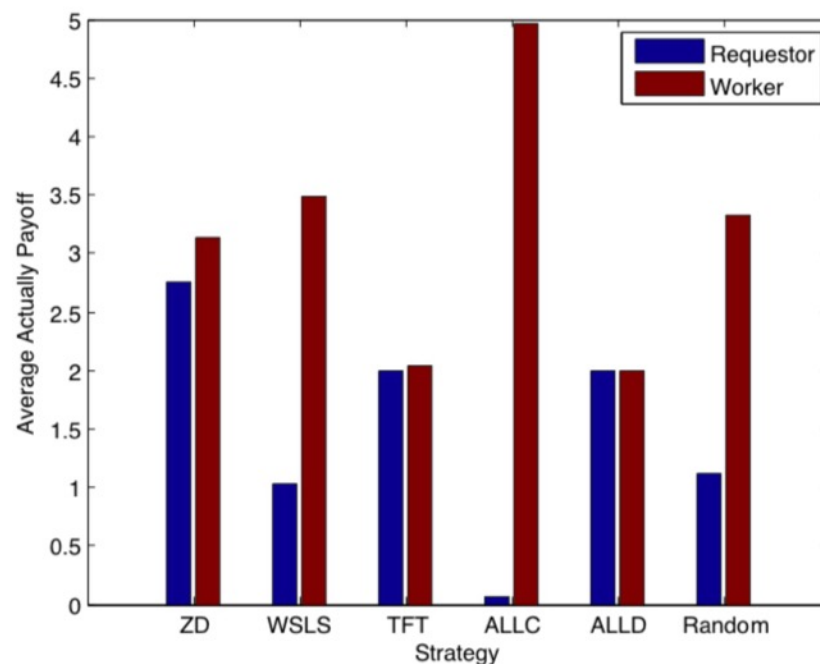
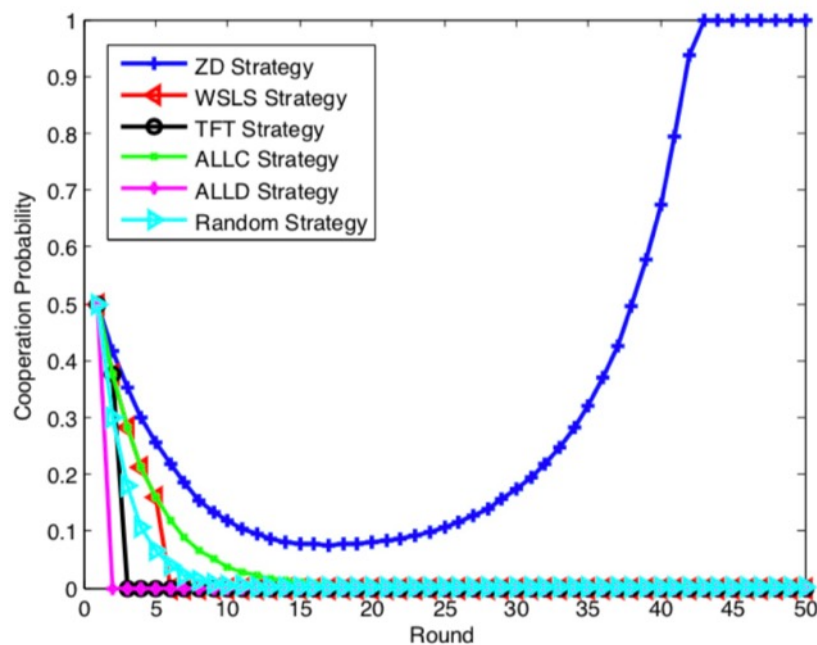
Sequential ZD based Algorithms

- Binary model



Sequential ZD based Algorithms

- Binary model



Sequential ZD based Algorithms

- Continuous model

- Evolutionary worker

$$f_h^{t+1} = f_h^t \frac{W_h^t}{E_w^t}$$

Algorithm 4 Sequential ZD Strategy based Incentive Algorithm for the Continuous Model

Require: $p^t(x|x_{-1}, y_{-1})$: the requestor's strategy and its initial value is the one used in the N_0^{th} round; $\mathbf{P}_s^c = \{P_{ij}^c\}_{\eta \times \eta}$: the state transition probability of the worker, and its initial value is calculated statistically through the preparatory N_0 rounds; N : the total number of rounds for algorithm termination.

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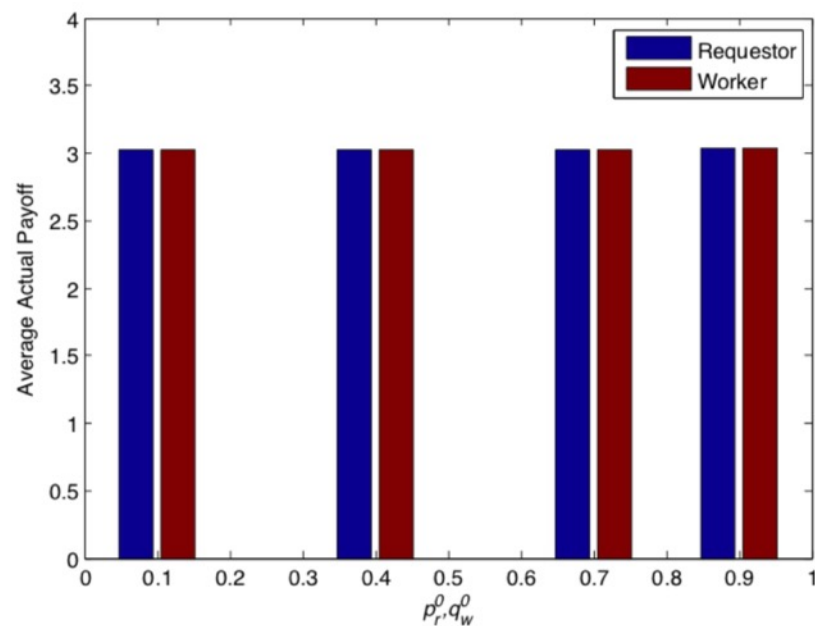
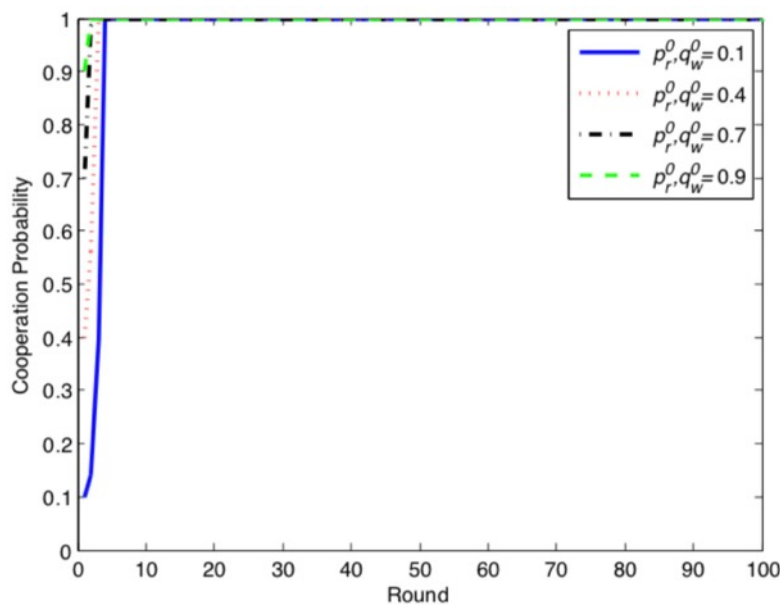
1: Initialize( $E_w^0$ )
2: for  $i = 1$  to  $N$  do
3:   if  $y^{i-1} \in [l_w + (\kappa - 1)\delta, l_w + \kappa\delta]$  then
4:     if  $P_{\kappa-\eta}^c \geq P_{\kappa-j}^c, \forall j \in \{1, w, \dots, \eta\}$  then
5:       Set  $p^i(x|x_{-1}, y_{-1})$  to let
6:          $E_w^i \leftarrow \max(E_w^t)$ 
7:     else
8:       Set  $p^i(x|x_{-1}, y_{-1})$  to let
9:          $E_w^i \leftarrow \min(E_w^t)$ 
10:    end if
11:  end if
12:  if The current round ends then
13:    Update  $\mathbf{P}_s^c$ 
14:  end if
15: end for

```

$\triangleright P_{\kappa-\eta}^c < P_{\kappa-j}^c$.

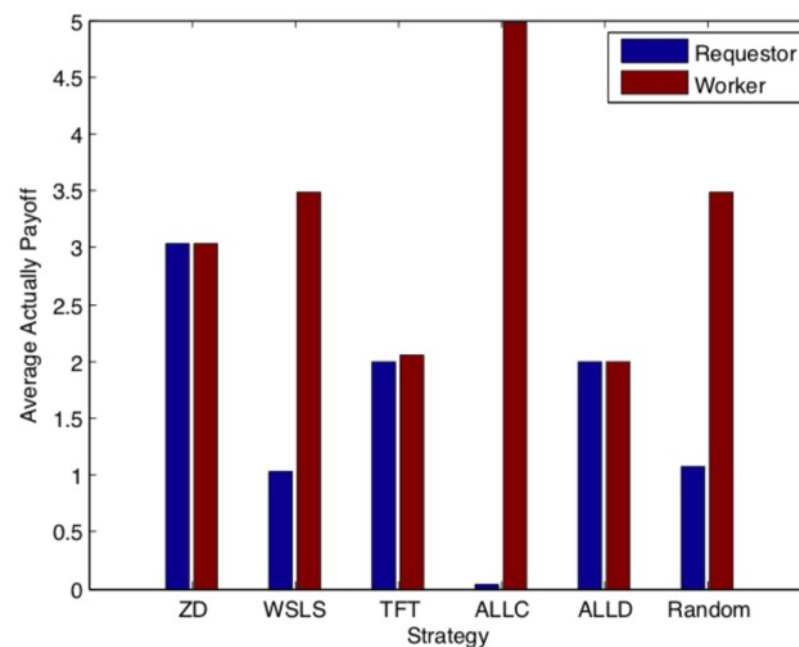
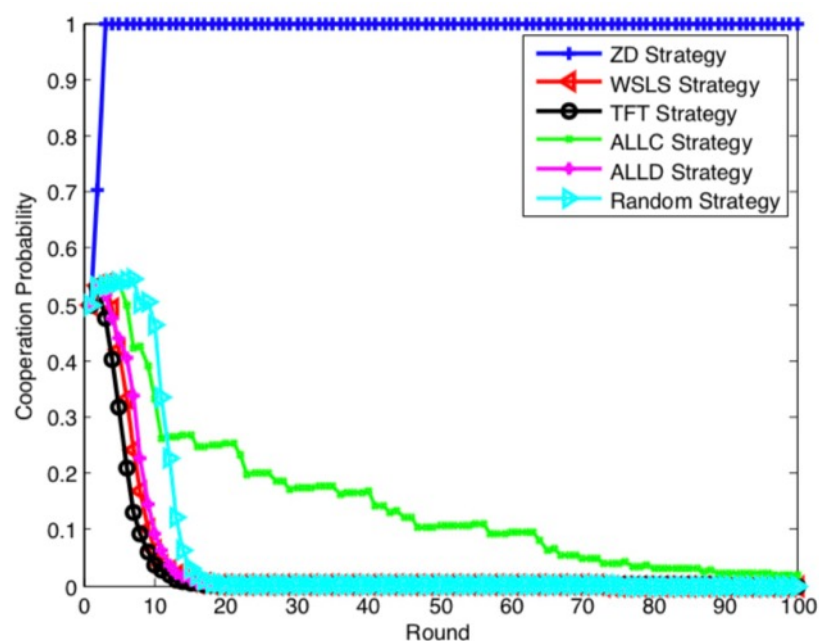
Sequential ZD based Algorithms

- Continuous model



Sequential ZD based Algorithms

- Continuous model



Related Papers

- **Qin Hu**, Shengling Wang*, Peizi Ma, Xiuzhen Cheng, Weifeng Lv, Rongfang Bie. Quality Control in Crowdsourcing Using Sequential Zero-Determinant Strategies, *IEEE Transactions on Knowledge and Data Engineering*, submitted.
- **Qin Hu**, Shengling Wang*, Xiuzhen Cheng, Liran Ma, Rongfang Bie. Solving the Crowdsourcing Dilemma Using the Zero-Determinant Strategies, *IEEE Transactions on Information Forensics and Security*, submitted.
- **Qin Hu**, Shengling Wang*, Chunqiang Hu, Jianhui Huang, Wei Li, Xiuzhen Cheng. Messages in a Concealed Bottle: Achieving Query Content Privacy with Accurate Location-Based Services, *IEEE Transactions on Vehicular Technology*, 2018, 67 (8), 7698-7711.
- **Qin Hu**, Shengling Wang*, Rongfang Bie, Xiuzhen Cheng. Social Welfare Control in Mobile Crowdsensing Using Zero-Determinant Strategy, *Sensors*, 2017, 17(5): 1012.
- **Qin Hu**, Shengling Wang*, Liran Ma, Rongfang Bie, Xiuzhen Cheng. Anti-Malicious Crowdsourcing Using the Zero-Determinant Strategy, *Distributed Computing Systems (ICDCS)*, 2017 IEEE 37th International Conference on. IEEE, 2017: 1137-1146.

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