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Linux Programmer's Manual

NAMESPACES (7)

NAME

top

namespaces - overview of Linux namespaces

#### DESCRIPTION

top

A namespace wraps a global system resource in an abstraction that makes it appear to the processes within the namespace that they have their own isolated instance of the global resource. Changes to the global resource are visible to other processes that are members of the namespace, but are invisible to other processes. One use of namespaces is to implement containers.

Linux provides the following namespaces:

Namespace	Constant	Isolates
Cgroup	CLONE_NEWCGROUP	Cgroup root directory
IPC	CLONE_NEWIPC	System V IPC, POSIX message queues
Network	CLONE_NEWNET	Network devices, stacks, ports, etc.
Mount	CLONE_NEWNS	Mount points
PID	CLONE_NEWPID	Process IDs
User	CLONE_NEWUSER	User and group IDs
UTS	CLONE_NEWUTS	Hostname and NIS domain name

This page describes the various namespaces and the associated /proc files, and summarizes the APIs for working with namespaces.

#### The namespaces API

As well as various /proc files described below, the namespaces API includes the following system calls:

#### clone(2)

The clone(2) system call creates a new process. If the flags argument of the call specifies one or more of the CLONE\_NEW\* flags listed below, then new namespaces are created for each flag, and the child process is made a member of those namespaces. (This system call also implements a number of features unrelated to namespaces.)

#### setns(2)

The setns(2) system call allows the calling process to join an existing namespace. The namespace to join is specified via a file descriptor that refers to one of the /proc/[pid]/ns files described below.

#### unshare(2)

The unshare(2) system call moves the calling process to a new namespace. If the *flags* argument of the call specifies one or

more of the **CLONE\_NEW\*** flags listed below, then new namespaces are created for each flag, and the calling process is made a member of those namespaces. (This system call also implements a number of features unrelated to namespaces.)

Creation of new namespaces using clone(2) and unshare(2) in most cases requires the CAP\_SYS\_ADMIN capability. User namespaces are the exception: since Linux 3.8, no privilege is required to create a user namespace.

## The /proc/[pid]/ns/ directory

Each process has a /proc/[pid]/ns/ subdirectory containing one entry for each namespace that supports being manipulated by setns(2):

```
$ ls -1 /proc/$$/ns
```

```
total 0
lrwxrwxrwx. 1 mtk mtk 0 Apr 28 12:46 cgroup -> cgroup:[4026531835]
lrwxrwxrwx. 1 mtk mtk 0 Apr 28 12:46 ipc -> ipc:[4026531839]
lrwxrwxrwx. 1 mtk mtk 0 Apr 28 12:46 mnt -> mnt:[4026531840]
lrwxrwxrwx. 1 mtk mtk 0 Apr 28 12:46 net -> net:[4026531969]
lrwxrwxrwx. 1 mtk mtk 0 Apr 28 12:46 pid -> pid:[4026531836]
lrwxrwxrwx. 1 mtk mtk 0 Apr 28 12:46 pid_for_children -> pid:[4026531834]
lrwxrwxrwx. 1 mtk mtk 0 Apr 28 12:46 user -> user:[4026531837]
lrwxrwxrwx. 1 mtk mtk 0 Apr 28 12:46 uts -> uts:[4026531838]
```

Bind mounting (see mount(2)) one of the files in this directory to somewhere else in the filesystem keeps the corresponding namespace of the process specified by pid alive even if all processes currently in the namespace terminate.

Opening one of the files in this directory (or a file that is bind mounted to one of these files) returns a file handle for the corresponding namespace of the process specified by pid. As long as this file descriptor remains open, the namespace will remain alive, even if all processes in the namespace terminate. The file descriptor can be passed to setns(2).

In Linux 3.7 and earlier, these files were visible as hard links. Since Linux 3.8, they appear as symbolic links. If two processes are in the same namespace, then the device IDs and inode numbers of their \( \frac{proc}{[pid]} \frac{lns}{xxx} \) symbolic links will be the same; an application can check this using the \( \frac{stat.st\_dev}{dev} \) and \( \frac{stat.st\_ino}{dev} \) fields returned by \( \frac{stat(2)}{dev} \). The content of this symbolic link is a string containing the namespace type and inode number as in the following example:

```
$ readlink /proc/$$/ns/uts
uts:[4026531838]
```

The symbolic links in this subdirectory are as follows:

```
/proc/[pid]/ns/cgroup (since Linux 4.6)
This file is a handle for the cgroup namespace of the process.
```

```
/proc/[pid]/ns/ipc (since Linux 3.0)
    This file is a handle for the IPC namespace of the process.
```

```
/proc/[pid]/ns/mnt (since Linux 3.8)
```

This file is a handle for the mount namespace of the process.

/proc/[pid]/ns/net (since Linux 3.0)

This file is a handle for the network namespace of the process.

/proc/[pid]/ns/pid (since Linux 3.8)

This file is a handle for the PID namespace of the process. This handle is permanent for the lifetime of the process (i.e., a process's PID namespace membership never changes).

/proc/[pid]/ns/pid\_for\_children (since Linux 4.12)

This file is a handle for the PID namespace of child processes created by this process. This can change as a consequence of calls to unshare(2) and setns(2) (see pid\_namespaces(7)), so the file may differ from /proc/[pid]/ns/pid. The symbolic link gains a value only after the first child process is created in the namespace. (Beforehand, readlink(2) of the symbolic link will return an empty buffer.)

/proc/[pid]/ns/user (since Linux 3.8)

This file is a handle for the user namespace of the process.

/proc/[pid]/ns/uts (since Linux 3.0)

This file is a handle for the UTS namespace of the process.

Permission to dereference or read (readlink(2)) these symbolic links is governed by a ptrace access mode PTRACE\_MODE\_READ\_FSCREDS check; see ptrace(2).

#### The /proc/sys/user directory

The files in the /proc/sys/user directory (which is present since Linux 4.9) expose limits on the number of namespaces of various types that can be created. The files are as follows:

max cgroup namespaces

The value in this file defines a per-user limit on the number of cgroup namespaces that may be created in the user names-pace.

max ipc namespaces

The value in this file defines a per-user limit on the number of ipc namespaces that may be created in the user namespace.

max mnt namespaces

The value in this file defines a per-user limit on the number of mount namespaces that may be created in the user namespace.

max net namespaces

The value in this file defines a per-user limit on the number of network namespaces that may be created in the user names-pace.

max pid namespaces

The value in this file defines a per-user limit on the number of pid namespaces that may be created in the user namespace.

#### max user namespaces

The value in this file defines a per-user limit on the number of user namespaces that may be created in the user namespace.

## max\_uts\_namespaces

The value in this file defines a per-user limit on the number of user namespaces that may be created in the user namespace.

Note the following details about these files:

- \* The values in these files are modifiable by privileged processes.
- \* The values exposed by these files are the limits for the user namespace in which the opening process resides.
- \* The limits are per-user. Each user in the same user namespace can create namespaces up to the defined limit.
- \* The limits apply to all users, including UID 0.
- \* These limits apply in addition to any other per-namespace limits (such as those for PID and user namespaces) that may be enforced.
- \* Upon encountering these limits, clone(2) and unshare(2) fail with the error ENOSPC.
- \* For the initial user namespace, the default value in each of these files is half the limit on the number of threads that may be created (/proc/sys/kernel/threads-max). In all descendant user namespaces, the default value in each file is MAXINT.
- \* When a namespace is created, the object is also accounted against ancestor namespaces. More precisely:
  - + Each user namespace has a creator UID.
  - + When a namespace is created, it is accounted against the creator UIDs in each of the ancestor user namespaces, and the kernel ensures that the corresponding namespace limit for the creator UID in the ancestor namespace is not exceeded.
  - + The aforementioned point ensures that creating a new user namespace cannot be used as a means to escape the limits in force in the current user namespace.

#### Cgroup namespaces (CLONE NEWCGROUP)

See cgroup\_namespaces(7).

#### IPC namespaces (CLONE NEWIPC)

IPC namespaces isolate certain IPC resources, namely, System V IPC objects (see <a href="scient">svipc(7)</a>) and (since Linux 2.6.30) POSIX message queues (see <a href="mailto:mq\_overview(7)">mq\_overview(7)</a>). The common characteristic of these IPC mechanisms is that IPC objects are identified by mechanisms other than filesystem pathnames.

Each IPC namespace has its own set of System V IPC identifiers and its own POSIX message queue filesystem. Objects created in an IPC

namespace are visible to all other processes that are members of that namespace, but are not visible to processes in other IPC namespaces.

The following /proc interfaces are distinct in each IPC namespace:

- \* The POSIX message queue interfaces in /proc/sys/fs/mqueue.
- \* The System V IPC interfaces in /proc/sys/kernel, namely: msgmax, msgmnb, msgmni, sem, shmall, shmmax, shmmni, and shm\_rmid\_forced.
- \* The System V IPC interfaces in /proc/sysvipc.

When an IPC namespace is destroyed (i.e., when the last process that is a member of the namespace terminates), all IPC objects in the namespace are automatically destroyed.

Use of IPC namespaces requires a kernel that is configured with the CONFIG\_IPC\_NS option.

## Network namespaces (CLONE\_NEWNET)

See network namespaces (7).

## Mount namespaces (CLONE\_NEWNS)

See mount\_namespaces(7).

## PID namespaces (CLONE\_NEWPID)

See pid namespaces(7).

## User namespaces (CLONE\_NEWUSER)

See user namespaces (7).

#### UTS namespaces (CLONE NEWUTS)

UTS namespaces provide isolation of two system identifiers: the hostname and the NIS domain name. These identifiers are set using sethostname(2) and setdomainname(2), and can be retrieved using uname(2), gethostname(2), and getdomainname(2).

Use of UTS namespaces requires a kernel that is configured with the CONFIG UTS NS option.

#### EXAMPLE top

See clone(2) and user namespaces(7).

#### SEE ALSO top

nsenter(1), readlink(1), unshare(1), clone(2), ioctl\_ns(2), setns(2),
unshare(2), proc(5), capabilities(7), cgroup\_namespaces(7),
cgroups(7), credentials(7), network\_namespaces(7), pid\_namespaces(7),
user namespaces(7), lsns(8), switch root(8)

# COLOPHON top

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Pages that refer to this page: nsenter(1), systemd-detect-virt(1), unshare(1), clone(2), getdomainname(2), gethostname(2), ioctl\_ns(2), setns(2), uname(2), unshare(2), proc(5), systemd.exec(5), cgroup\_namespaces(7), cgroups(7), credentials(7), mount\_namespaces(7), mq\_overview(7), network\_namespaces(7), pid\_namespaces(7), svipc(7), user\_namespaces(7), lsns(8), mount(8)

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