# Learned Indexes

Part II Project - Franz Nowak

## What's a database index again?

- data structure that improves speed of data retrieval operations
- cost of additional writes
- cost of storage space to maintain index
- e.g. b-tree on sorted data: lookup O(log(n)) -> optimal for range lookups

## What's a database index again?

**ROOT NODE** 

INTERMEDIATE NODES

F' 'L' 'N'

LEAF
NODES

('F',rid)

b-tree index

source:https://www.ibm.com/support/knowledgecenter

('G',rid)

('l',rid)

('K',rid)

('M',rid)

('N',rid)

. . .

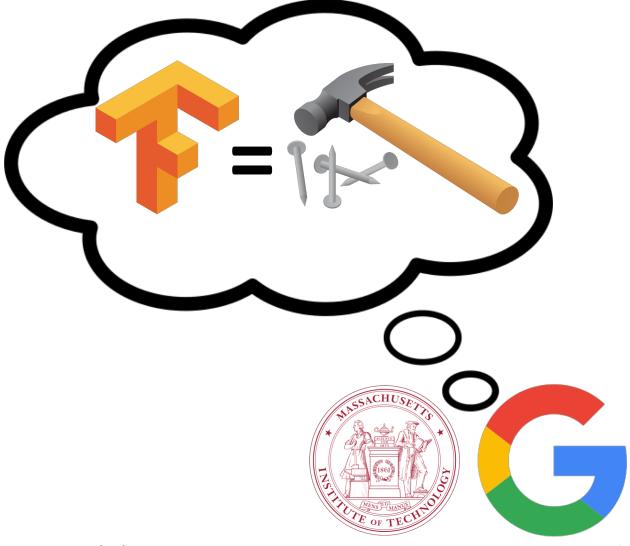
## Exploit inherent structure of data

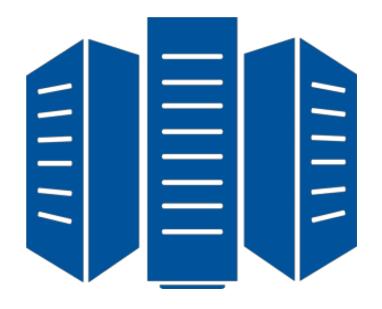
What if data is all integers from 0 to 1,000,000?

data\_array[lookup\_key]

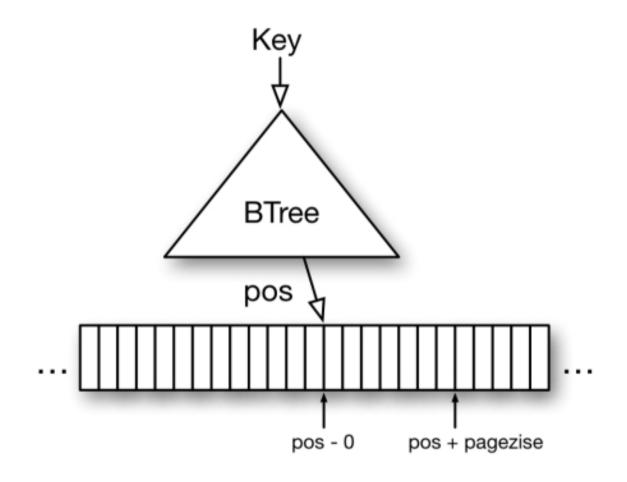
• No need for b-tree, lookup is O(1) time and space!

## How can we improve on this?

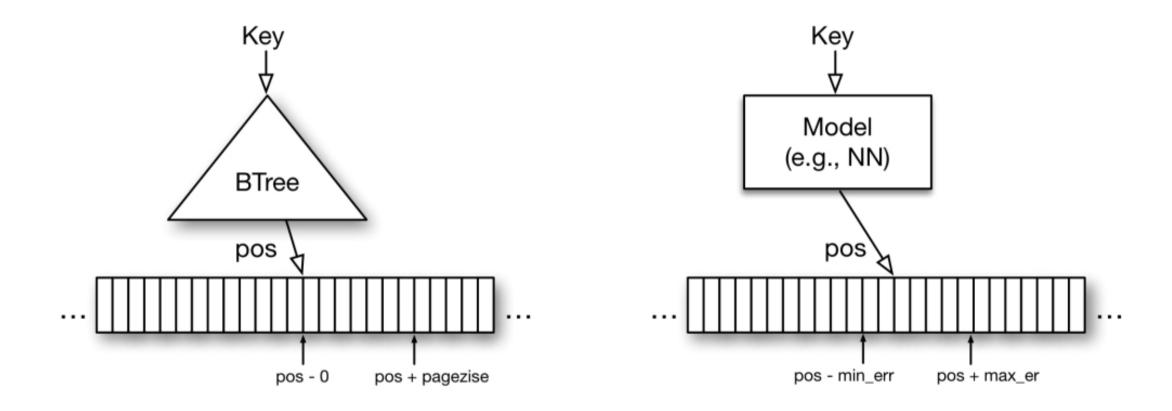




### Core idea: B-Trees are models



## Key idea: B-Trees are models



## To Do: Regression Tree

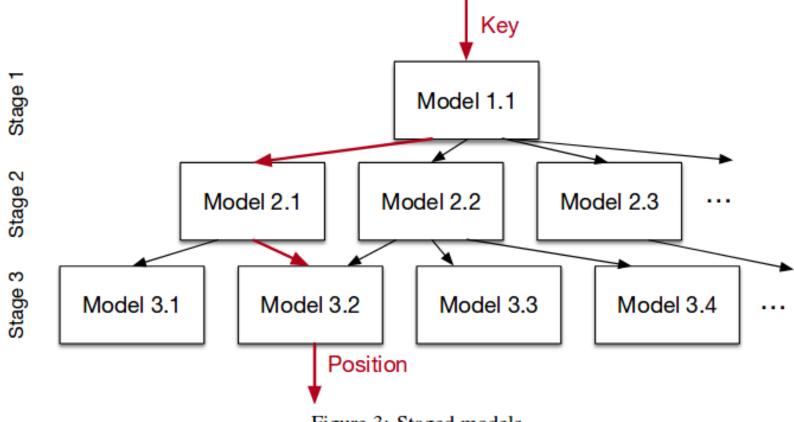


Figure 3: Staged models

## Position in sorted array: CDF

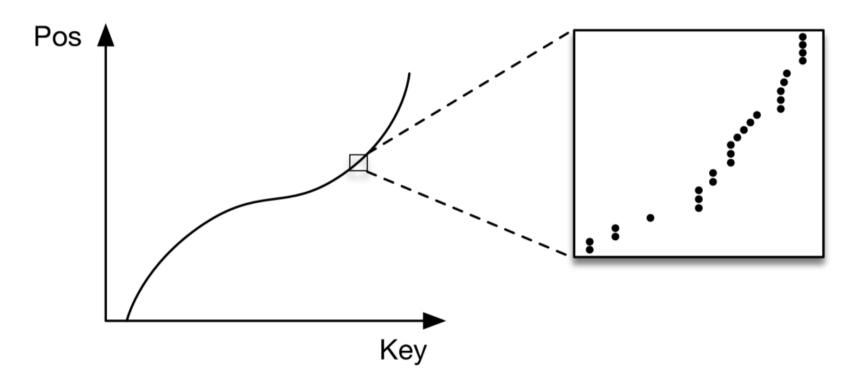
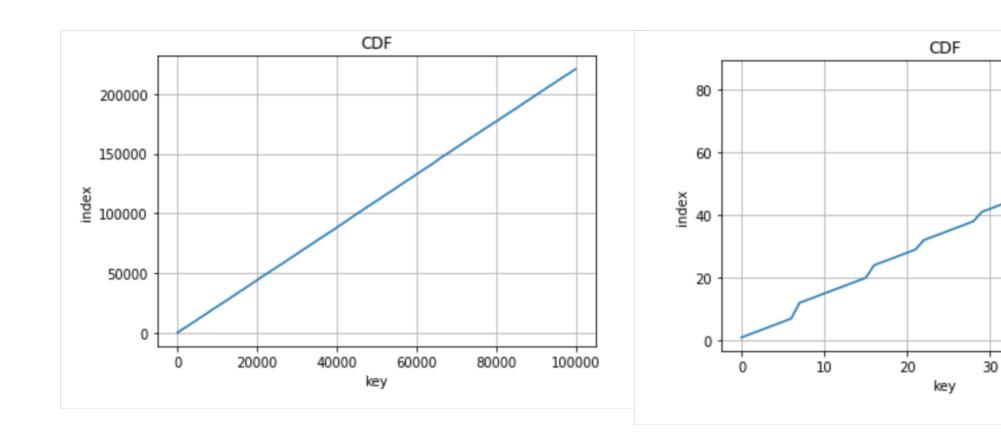


Figure 2: Indexes as CDFs

#### Best case vs worst case



40

50

### Dataset

• 100,000 keys in sorted array

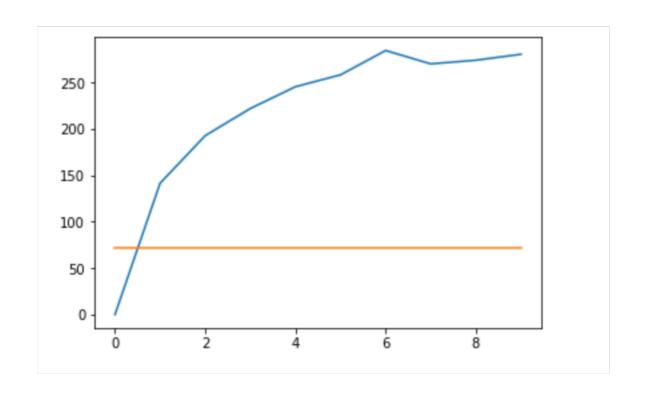
• 10 levels of entropy

• averaged over 100 runs

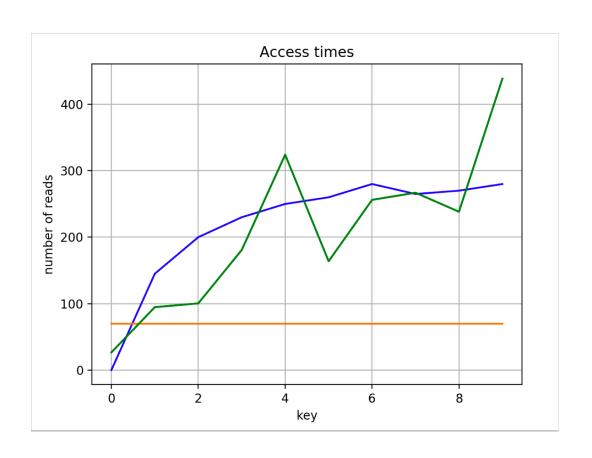
## Implementation-independent evaluation

- What we really want to know: wall-clock time
- different depending on OS, hardware (SIMD), implementation of b-tree, programming language, ...
- Idea: use reads as proxy
- read\_time = f(n) + reads
- f(n) is the base complexity of the data structure (e.g. invocation overhead of NN of size proportional to dataset size)

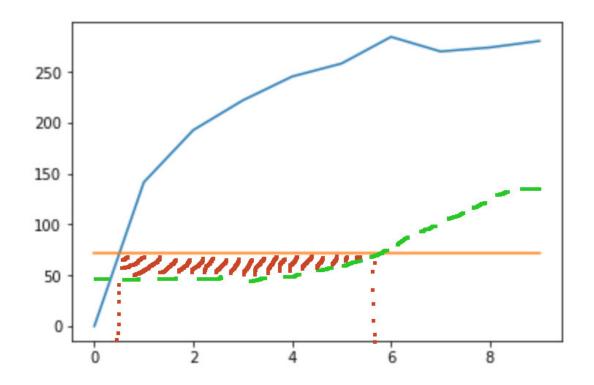
# Preliminary results



# Preliminary results



# Hypothesis



### Main references

• The paper: <a href="https://arxiv.org/abs/1712.01208">https://arxiv.org/abs/1712.01208</a>

• Stanford EE380: Computer Systems Colloquium Seminar

https://www.youtube.com/watch?v=NaqJO7rrXy0