

# Lecture 7a: Trees!

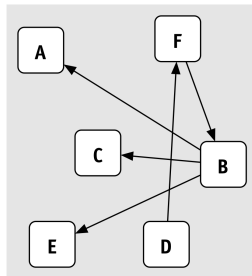
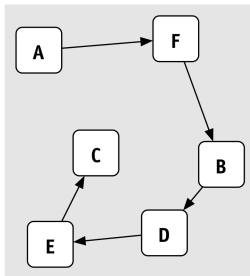
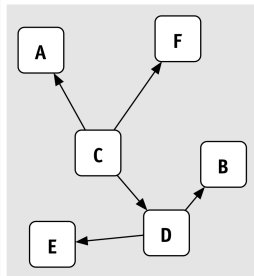
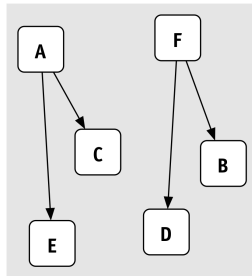
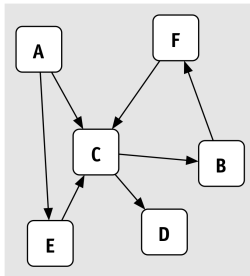
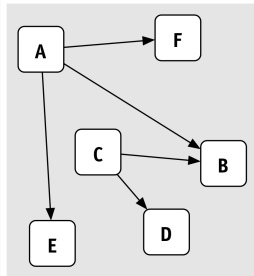
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CS 70: Data Structures and Program Development

Tuesday, March 3, 2020

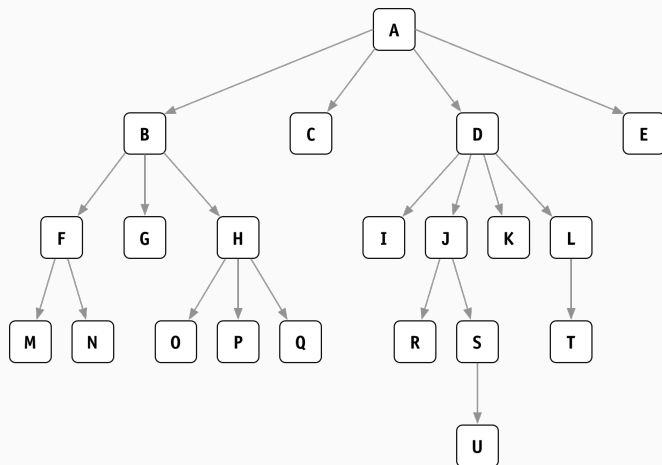
# What is a tree?

# Which of these are trees?

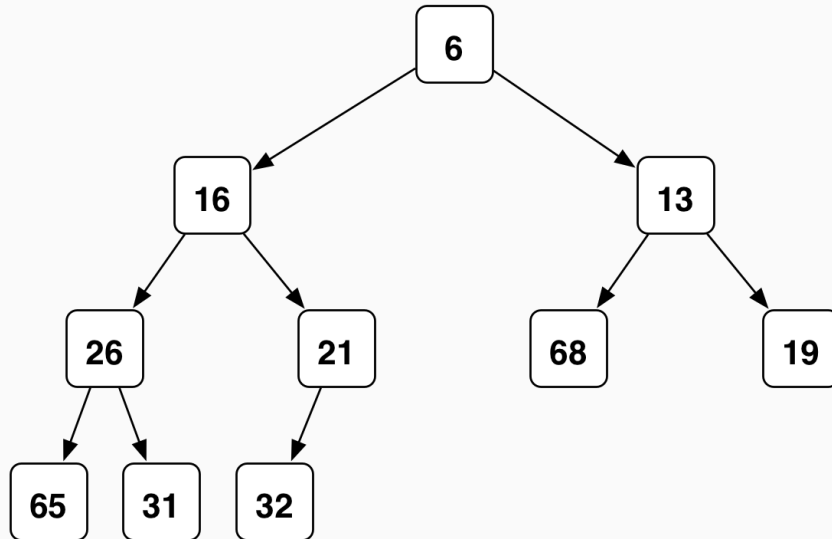


# Terminology

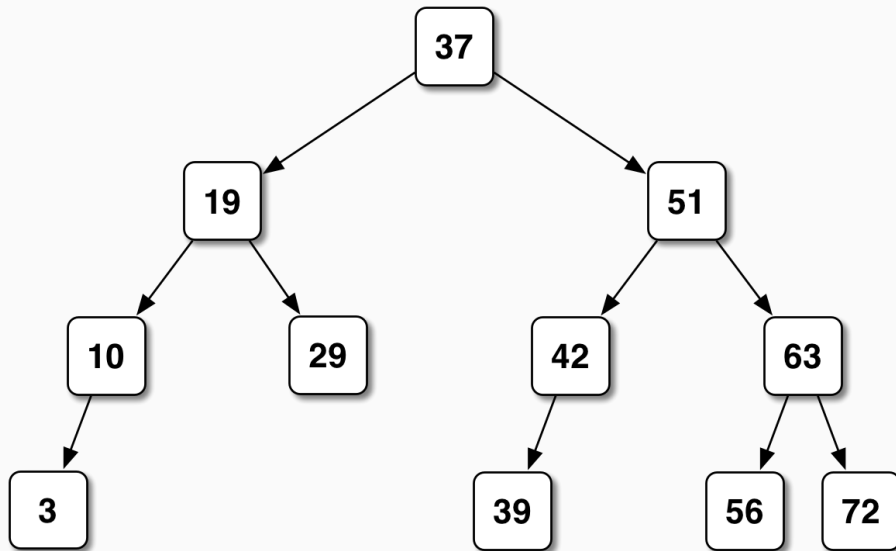
node, edge; root, leaf, tree, subtree; parent, child, ancestor; height, balance; binary



# Binary Tree

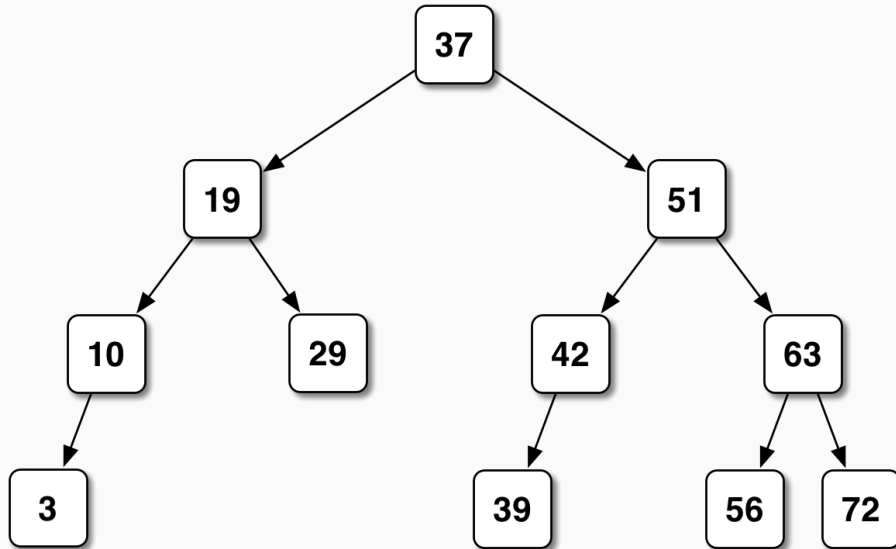


# Binary Search Tree (a.k.a. Ordered Binary Tree)



# Basic Algorithms

To do: find 56; find 35; insert 47

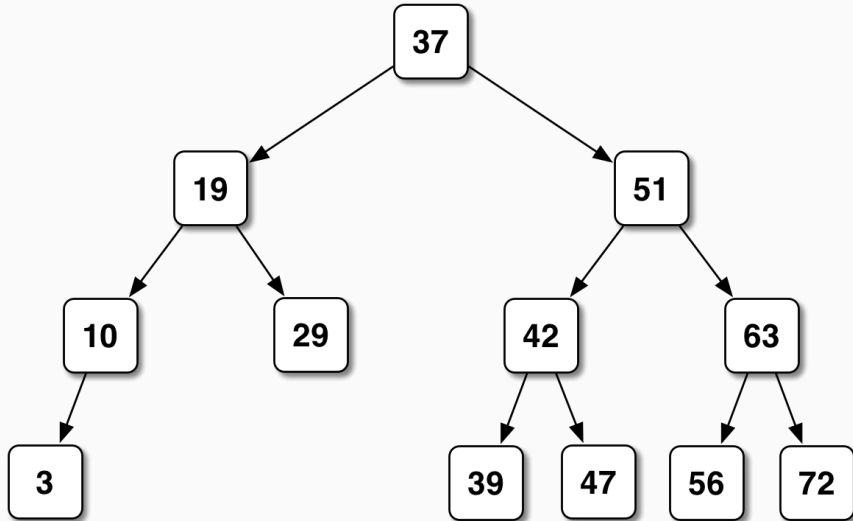




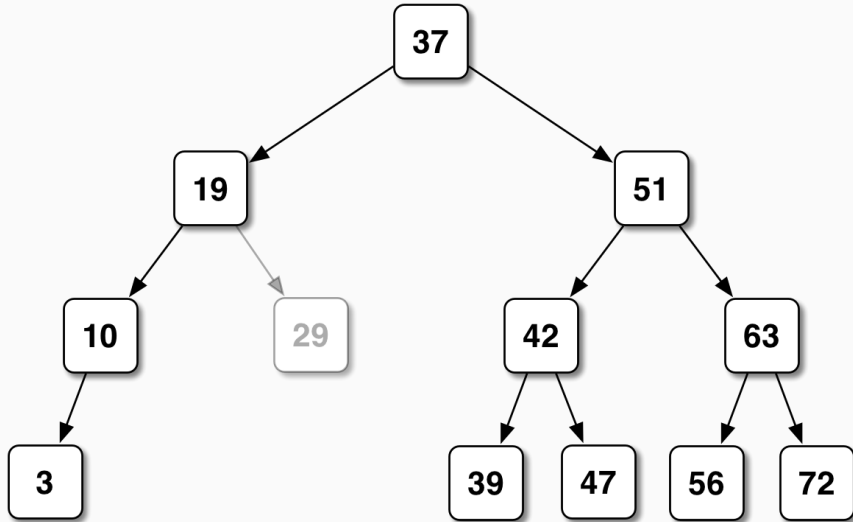
## insert pseudocode

```
insert(tree, x):  
    if tree is empty:  
        make x its new root.  
  
    else if  $x < \text{tree's root}$ :  
        insert(left subtree, x)  
  
    else if  $\text{tree's root} < x$ :  
        insert(right subtree, x)
```

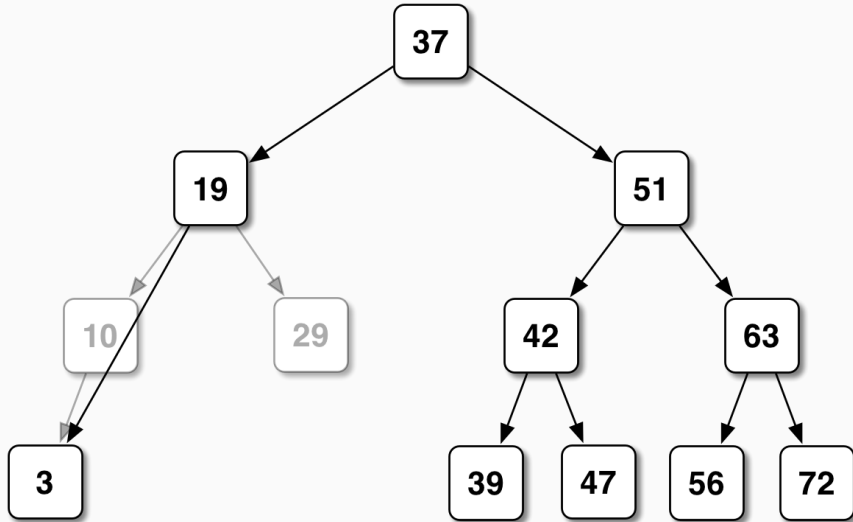
To do: delete 29



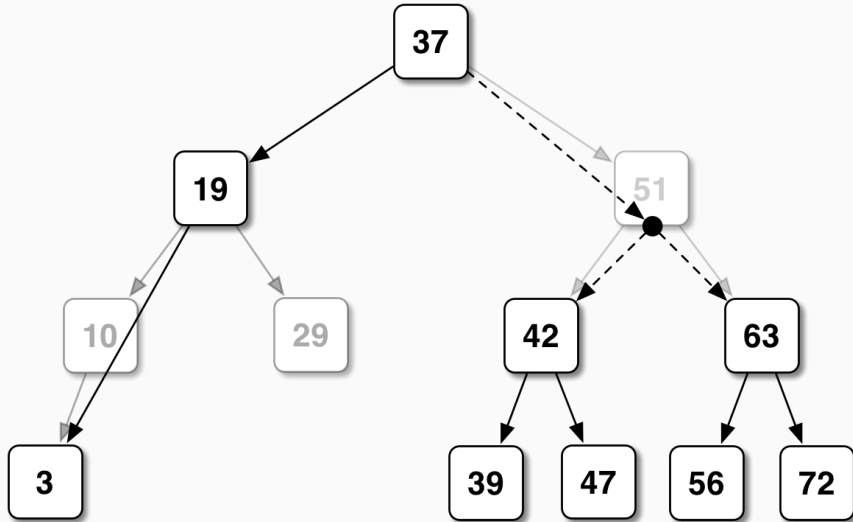
To do: delete 10



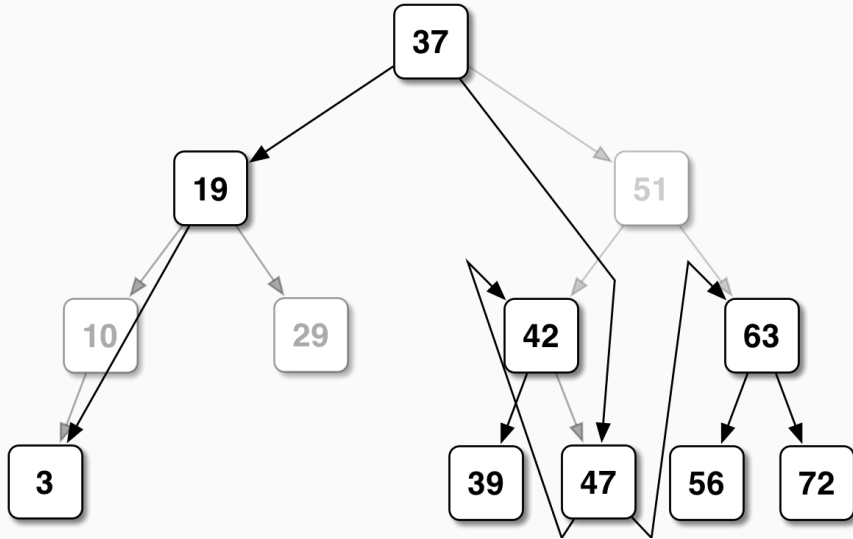
To do: delete 51



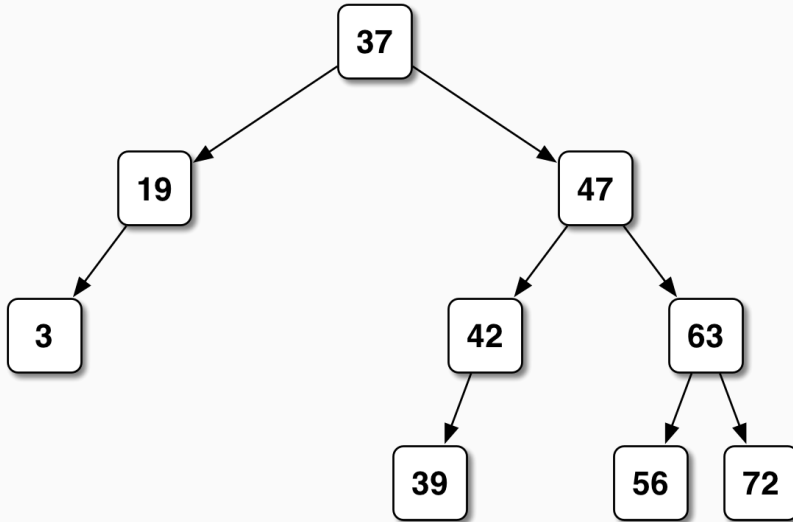
# Oops



# Better (why?)



## Redrawing the same tree...



# Exercise

## Quick Test: Doing Insert

Choose one of these sequences, and insert it into an empty tree.

D, B, A, F, E, C, G	D, B, F, E, A, C, G	D, B, F, C, E, G, A	D, F, B, C, A, E, G
D, F, E, G, B, A, C	D, F, G, E, B, C, A	D, B, F, C, E, A, G	D, B, C, F, E, A, G
D, B, C, A, F, G, E	D, B, F, A, E, C, G	D, F, G, B, C, E, A	D, B, F, C, G, E, A
D, F, B, G, A, E, C	D, B, C, F, G, A, E	D, B, F, G, E, C, A	D, F, G, B, A, E, C
D, B, F, C, A, G, E	D, F, B, A, E, C, G	D, B, C, F, G, E, A	D, B, F, A, G, C, E
D, F, B, G, C, A, E	D, B, F, E, A, G, C	D, F, B, E, A, G, C	D, B, F, G, A, C, E
D, F, G, B, A, C, E	D, F, B, C, G, E, A	D, F, B, C, G, A, E	D, F, G, E, B, A, C
D, F, B, A, G, E, C	D, B, F, A, G, E, C	D, F, B, E, G, C, A	D, F, G, B, C, A, E
D, F, E, G, B, C, A	D, B, A, F, G, E, C	D, B, F, A, C, G, E	D, B, A, F, G, C, E
D, F, B, C, E, G, A	D, B, A, F, C, G, E	D, F, B, G, C, E, A	D, B, F, A, C, E, G
D, F, B, E, A, C, G	D, F, E, B, G, A, C	D, F, B, E, C, G, A	D, F, E, B, A, C, G
D, B, F, A, E, G, C	D, F, E, B, C, A, G	D, F, E, B, A, G, C	D, F, E, B, G, C, A
D, B, F, G, E, A, C	D, F, B, E, C, A, G	D, F, B, A, C, G, E	D, B, F, E, G, A, C
D, B, C, F, A, E, G	D, B, A, F, C, E, G	D, F, B, C, A, G, E	D, F, B, A, E, G, C
D, F, B, G, E, A, C	D, B, F, E, C, G, A	D, F, E, B, C, G, A	D, B, F, G, C, E, A
D, F, B, A, C, E, G	D, B, F, C, G, A, E	D, B, F, G, C, A, E	D, F, B, A, G, C, E
D, B, A, C, F, G, E	D, B, C, F, E, G, A	D, F, B, G, E, C, A	D, F, G, B, E, C, A



## Exercise (continued)

What tree results from the following sequences of inserts?

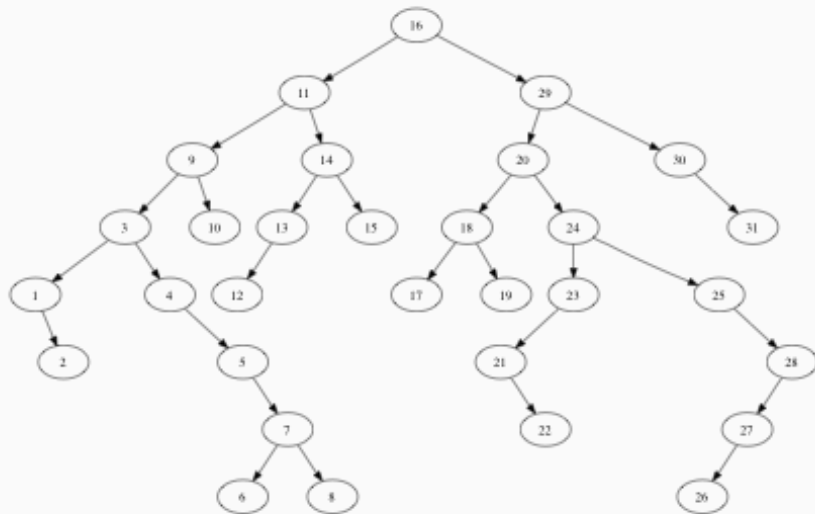
- A, B, C, D, E, F, G
- D, C, A, B, E, F, G

# Suppose we have a BST with $n$ nodes.

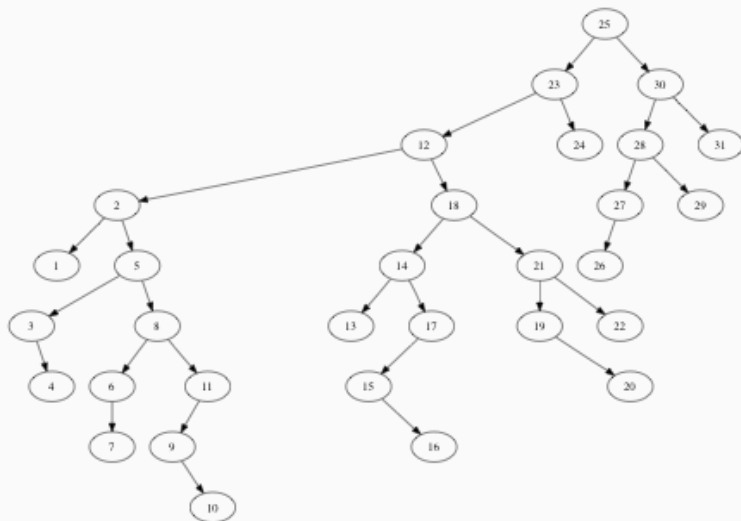
What is the worst-case running time for `find` (and `insert`)

- if we have a really terrible tree?
- if we have a really nice tree?
- if we have a “random” tree?

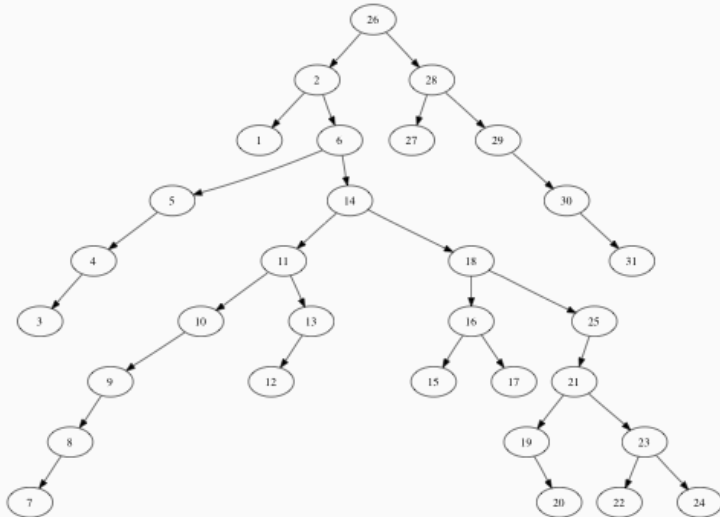
# Random trees average 39% worse than perfect



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# Building better trees: Off-line algorithm

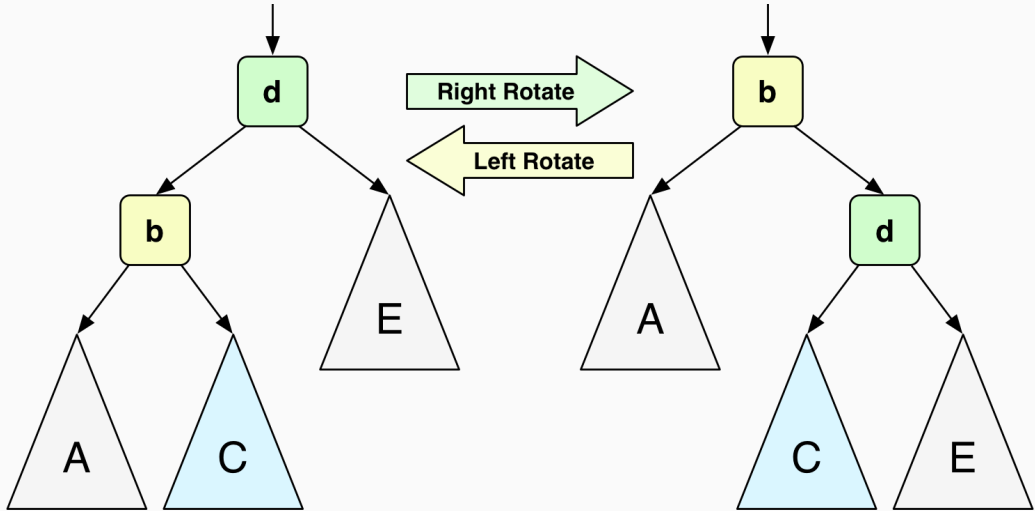
1. Take the inputs we want to put in the tree.
2. Randomly shuffle them.
3. Build tree by inserting in *shuffled* order.

# Building better trees: Off-line algorithm

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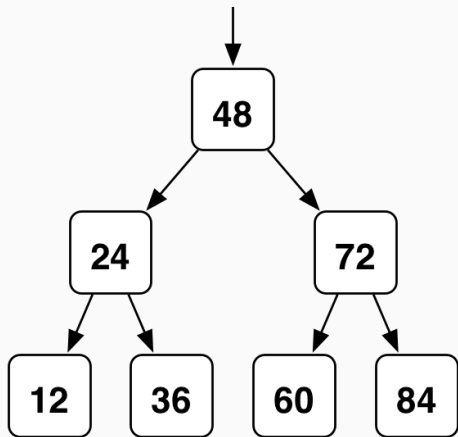
What if we want to keep the tree balanced as new data comes in?

# Tree Rotations





**Insert 40. Rotate left at 36, left at 24, right at 48.**



# insertAtRoot pseudocode

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```
insertAtRoot(tree, x):  
    if tree is empty:  
        make x its new root.  
  
    else if  $x < \text{tree's root}$ :  
        insertAtRoot(left subtree, x)  
        do right rotation at tree's root.  
  
    else if  $\text{tree's root} < x$ :  
        insertAtRoot(right subtree, x)  
        do left rotation at tree's root.
```

# Building better trees: Randomized Binary Trees

Idea: insert each new key “randomly” into the tree-so-far

- Maybe it should become the new root
- Maybe put it somewhere below the existing root

But how often to do each?

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But how often to do each?

Answer: If the tree has  $n$  nodes **before** the insert,

- do insert-at-root with probability  $1/(n + 1)$
- otherwise, insert randomly into the appropriate child.

# Learning Targets

1. Given a tree, I can tell whether it's a valid BST.
2. I can simulate BST lookup, insert, and delete (on paper)
3. I can simulate left and right rotations (on paper)
4. I can simulate `insertAtRoot` (on paper)
5. I can simulate Randomized Binary Tree insertion.