Name:	 	 	
Today's Date:	 	 	

Today's Goals

- Explain how Splay and AVL trees maintain balance
- Describe the tradeoffs between Splay and AVL trees
- Explain how AVL trees have amortized performance.

Today's Question(s)

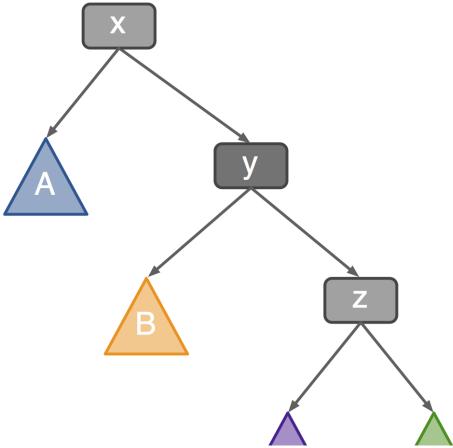
Of the trees we've seen before, which give:

• guaranteed constant time insert and lookup

• expected constant time insert and lookup

Lingering Questions

How can we bring z to the root?

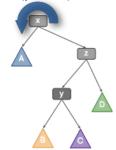


Splay rotations

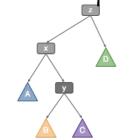
left(parent)



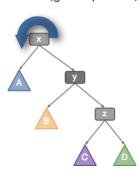
left(parent)



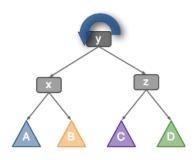
"normal" left-left



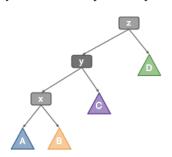
left(grandparent)



left(parent)



"splay" left-left



Splay Trees: Insert and Lookup

Always insert at the root, but use different rotations

- ▶ Never perform a regular left-left rotation use . . .
- ▶ Never perform a regular right-right rotation use . . .

For lookup, bring found item to root using ...

Applet demo

(Applet is great for gaining intuition, but off-limits for the final.)

Splay Trees: Properties

Splay trees are a **self-adjusting** data structure

- Recently added or accessed items migrate to the top of the tree
- Lookups balance the tree, where necessary

Advantages

Disadvantages

Well now my son is my grandson and my Daughter is my ma and my Granddaughter's my sister and my Dad is my son and now my Brother is my grandson and my Mom is my daughter and my Grandson is my brother and my Son is my new pa

I'm a node in a tree
This is not about a person
It's a splay rotation on a BST

– A musical interlude from Dani Demas (CS 70, Spring 2014)
Listen Here!

AVL Trees

(named after their inventors, G.M. Adelson-Velsky and E.M. Landis)

Balance conditions

- Each subtree is balanced
- Height of each node's subtrees differs by at most 1

Advantages

- Fast find keeps tree height minimized
- ► No risk of ending up with a "bad" tree

Disadvantages

- Rigid balance requirements slow down insert and delete
- Need to keep track of subtree heights