

# HAETAE and SMAUG-T: Korean PQC Standards

Hyeongmin Choe

CryptoLab, Inc.

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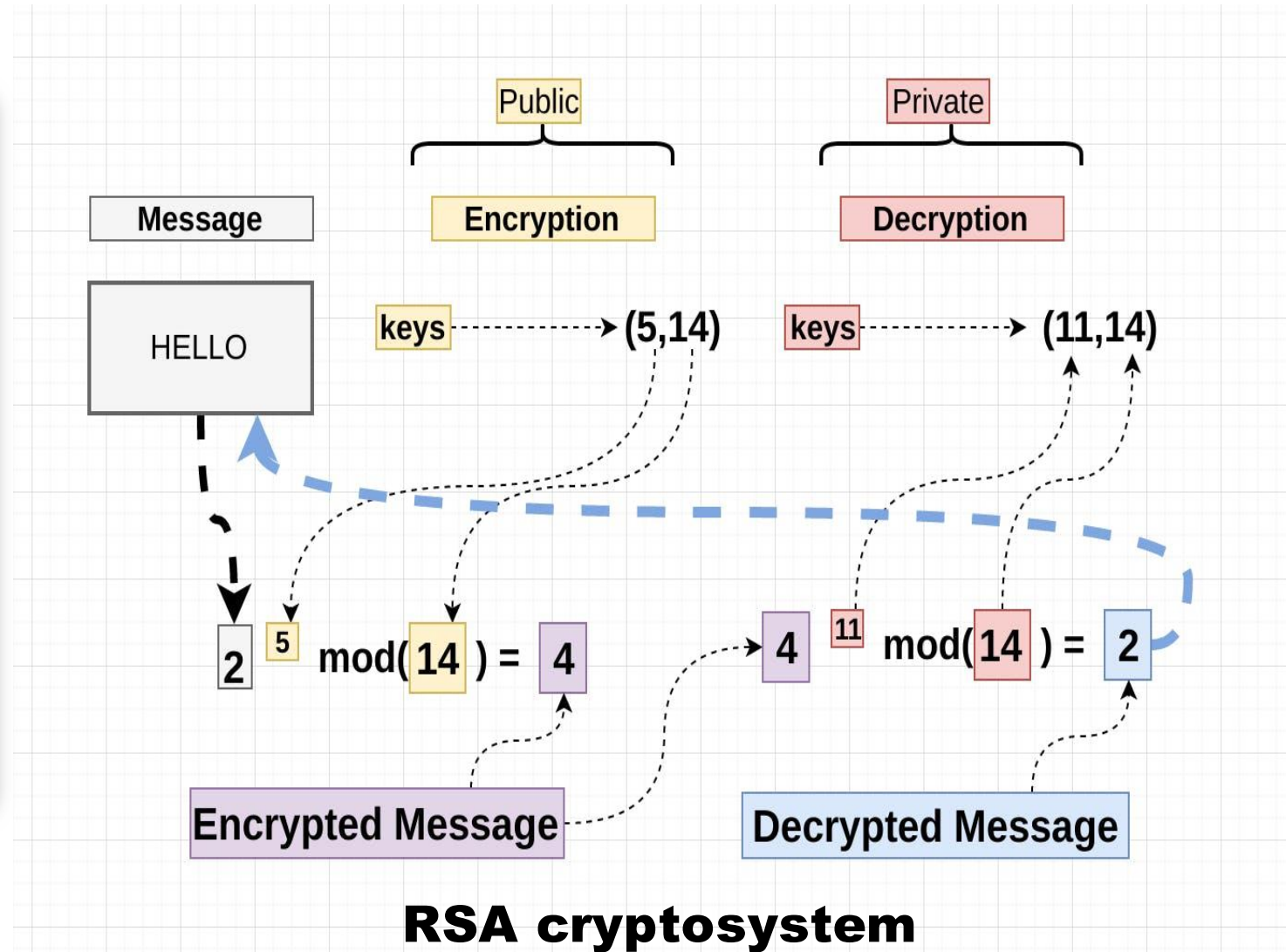
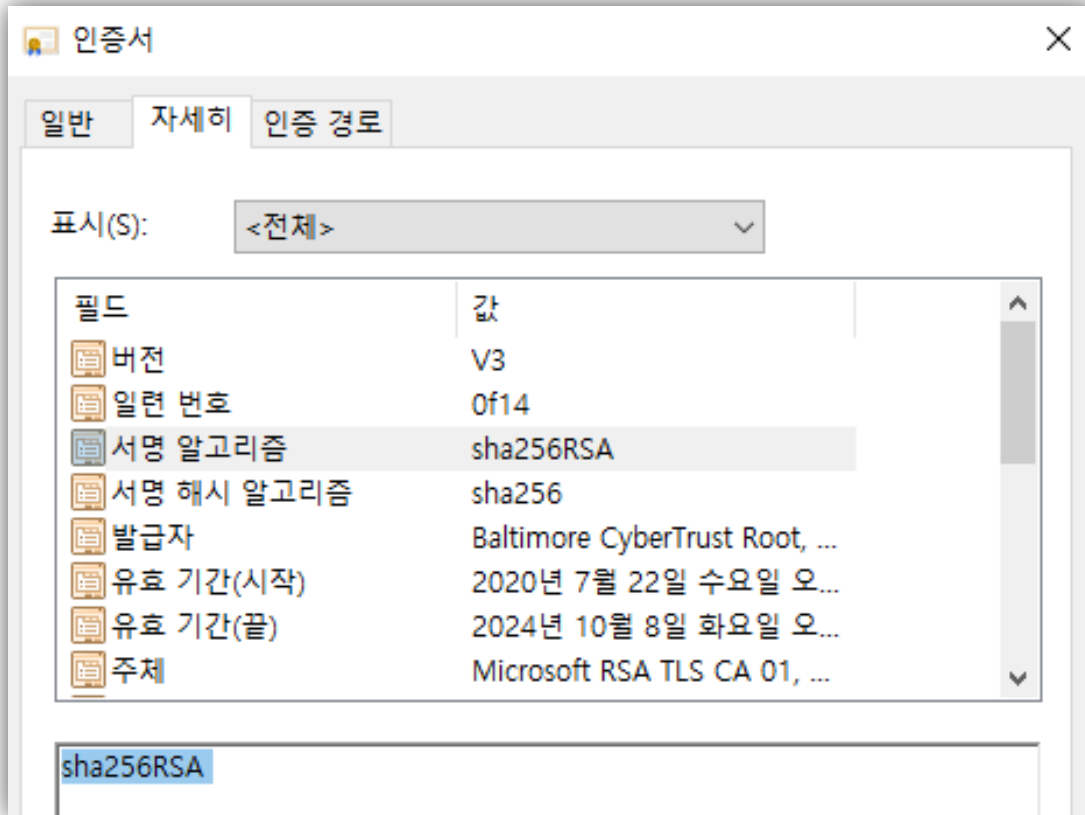
# Why Post-Quantum Cryptography?





# Why Post-Quantum Cryptography?

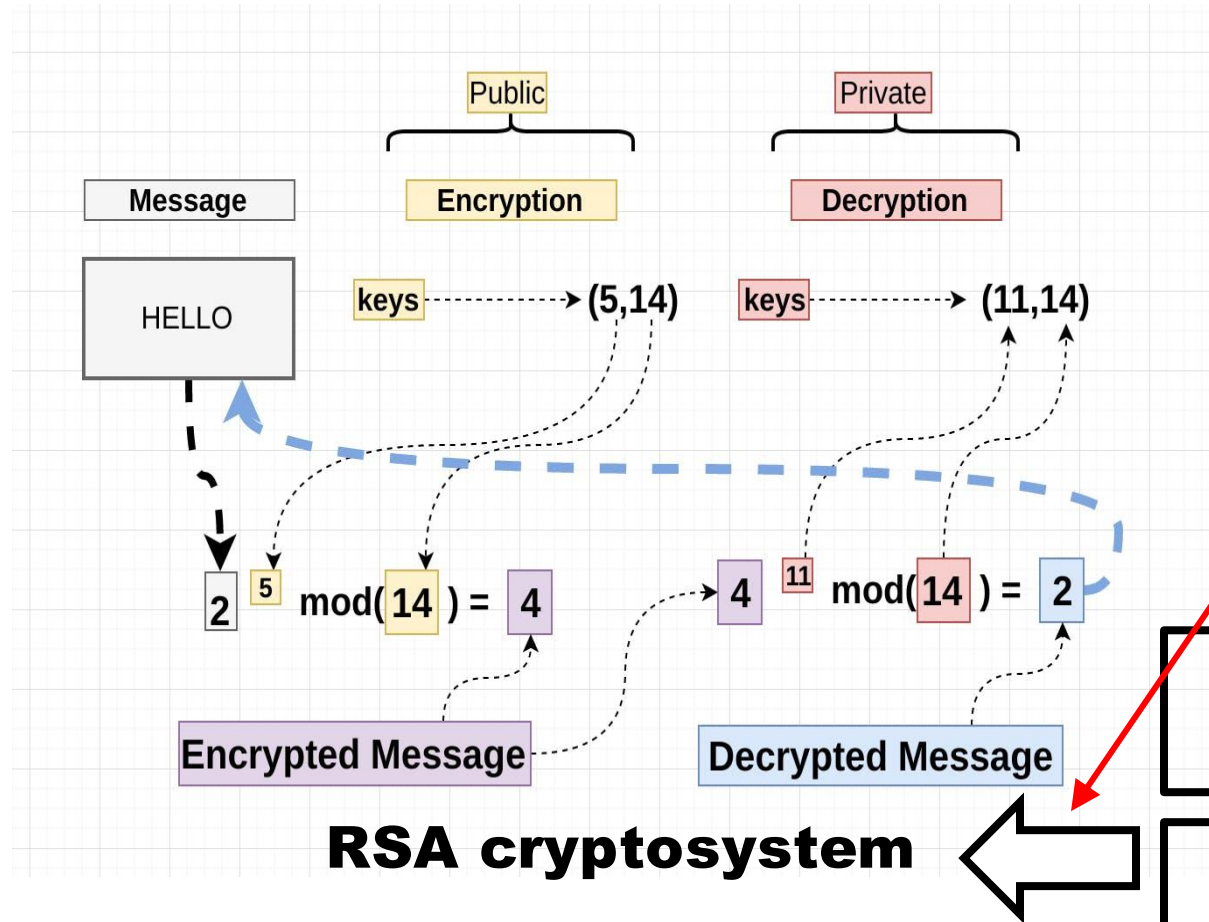
## “Classical” Cryptography





# Why Post-Quantum Cryptography?

## "Classical" Cryptography



**Reduction**

**Integer Factorization problem**

**RSA problem: *for given***  
 $a, e, n \in \mathbb{Z}$ , ***find***  $b \in \mathbb{Z}$   
***satisfying***  $a \equiv b^e \bmod n$ .

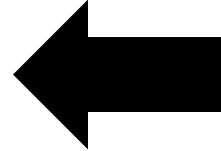


# Why Post-Quantum Cryptography?

“Classical” Cryptography

**Cryptosystem**

**Reduction**



**Hard problems**

- RSA Encryption/Signature
- Diffie-Hellman Key Exchange
- ECDSA/ECDH

- Integer Factorization
- Discrete Logarithm (DL)
- DL in Elliptic Curve  $\mathbb{F}_p$  (ECDL)

**NP!**

**Shor's algorithm  
in Quantum Computer**



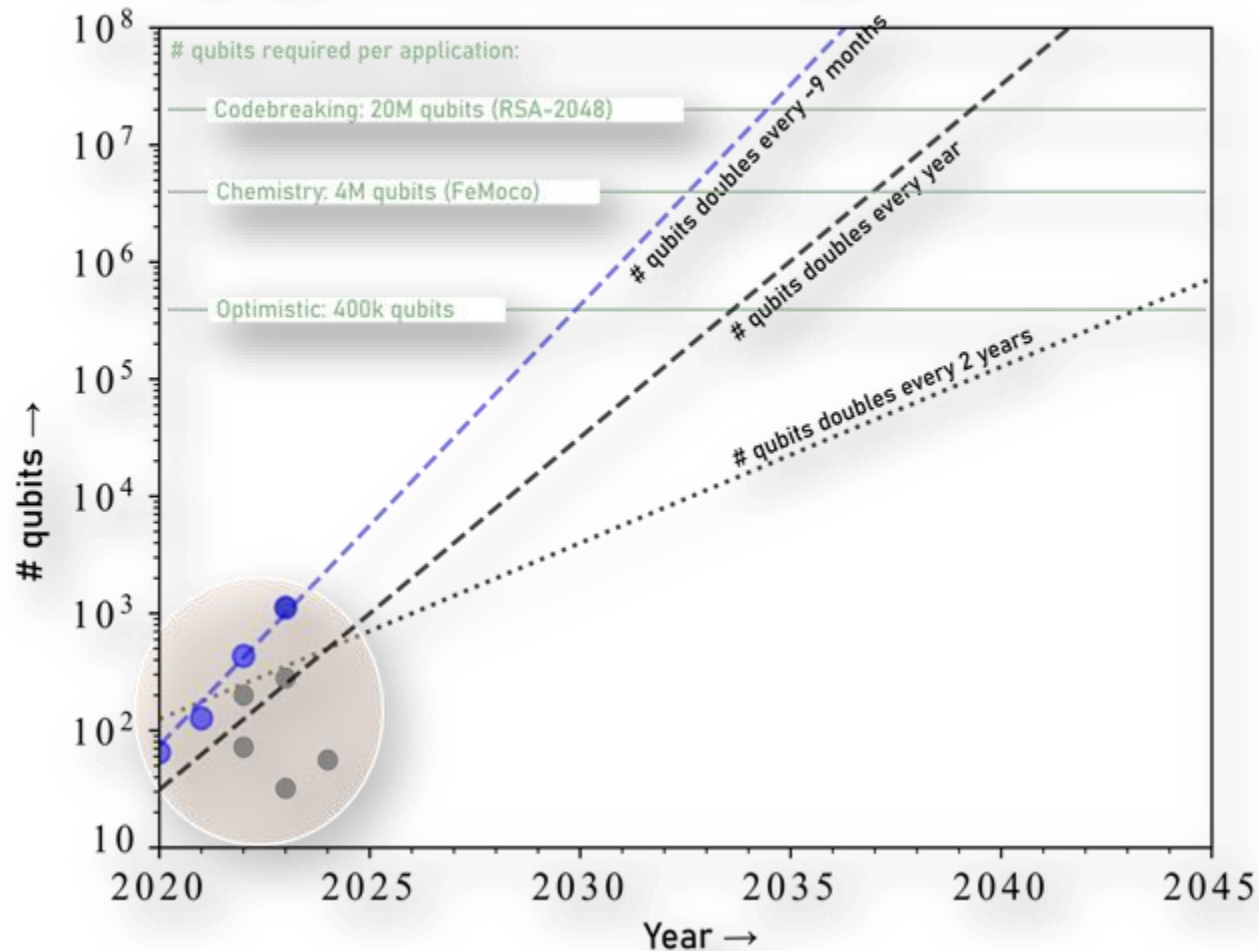
# Why Post-Quantum Cryptography?

“Classical” Cryptography

Crypto

- RSA Encryption
- Diffie-Hellman
- ECDSA/ECDSA

Qubit growth estimates, according to Moore's Law



ms

NP!  
DL)  
over  $\mathbb{F}_p$  (ECDL)

RSA2048: billions of years vs. several seconds in  
4,000-logical-cubit quantum computer.

(expected, when quantum computers are commercialized... but when?)



# Why Post-Quantum Cryptography?

## Post-Quantum Cryptography

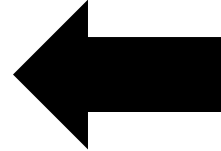
**Post-Quantum  
Cryptography**

- **Lattice-based cryptography**

- **Code-based cryptography**

•  
•  
•

**Reduction**



**Hard problems  
(even) against  
Quantum  
Algorithms**

- **Shortest/Closest Vector Problem (SVP/CVP)**

≈ **Learning With Errors (LWE)**

≈ **Learning With Rounding (LWR)**

≈ **Short Integer Solution (SIS)**

- **Syndrome Decoding Problem (SDP)**

•  
•  
•

**NP-hard!\***

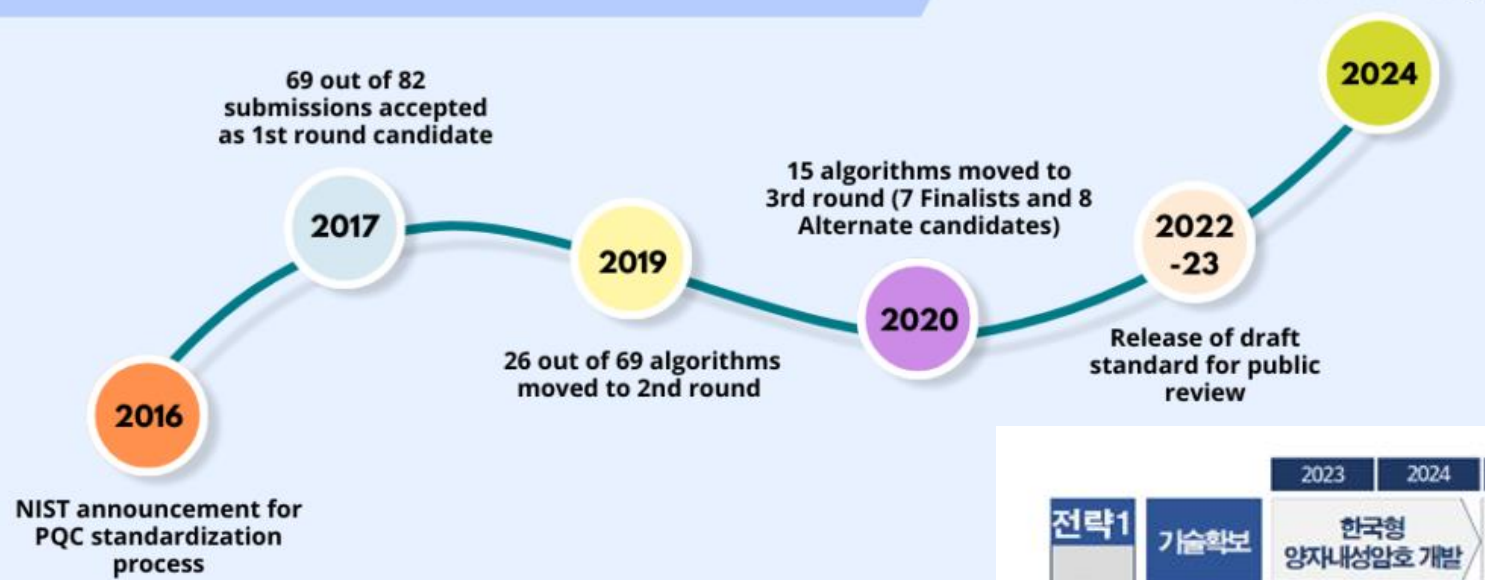




# Why Post-Quantum Cryptography?

## Global Movement

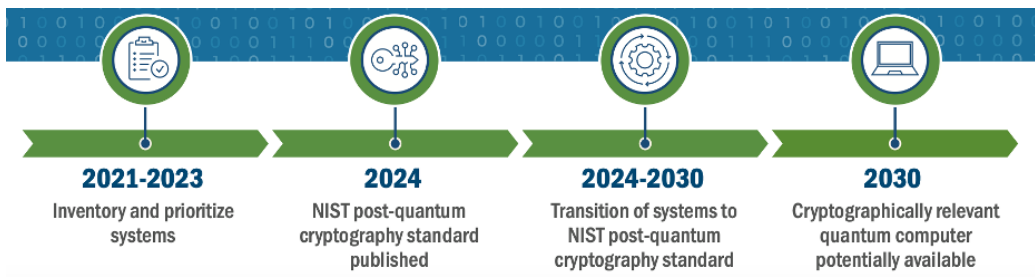
### NIST PQC Standardization Milestones



ETSI White Paper No. 8

### Quantum Safe Cryptography and Security

An introduction, benefits, enablers and challenges





- 2021년, 국립보안기술연구소에서 한국형 양자내성암호 확보를 목표로 공모
  - 공개 키 암호(PKE)와 키 캡슐화(KEM) 분야: 7개 알고리즘
  - 전자서명(Digital Signature) 분야: 9개 알고리즘
- 2025년, 최종 4개 알고리즘 선정

KpqC 2025 Selected	PKE/KEM	Signatures	Overall
Lattice-based	SMAUG-T NTRU+	HAETAE	3
Symmetric-based		AlMer	1



# SMAUG-T

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Defense Counterintelligence  
Command

## SMAUG-T



서울대학교  
SEOUL NATIONAL UNIVERSITY



HEAΛN  
CRYPTO LAB

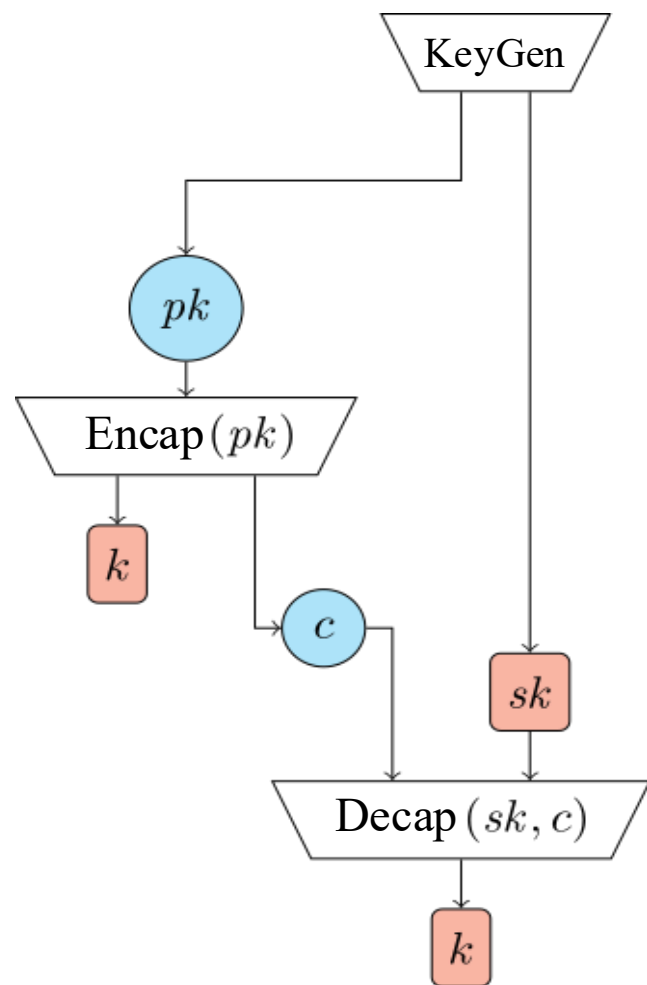


국군방첩사령부  
Defense Counterintelligence Command

Joint work between **Seoul National University** (서울대학교), **CryptoLab Inc.** ((주)크립토랩),  
and **Defense Counterintelligence Command** (국군방첩사령부).



- Key Encapsulation Mechanism (KEM):

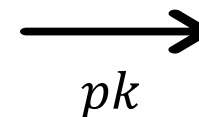


0.  $(sk, pk) \leftarrow \text{KeyGen}$

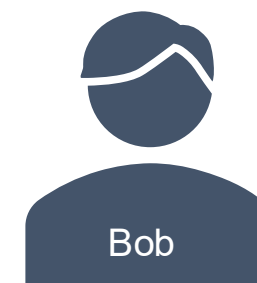
(knows  $sk$ )



Alice

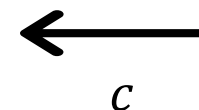


(knows  $pk$ )



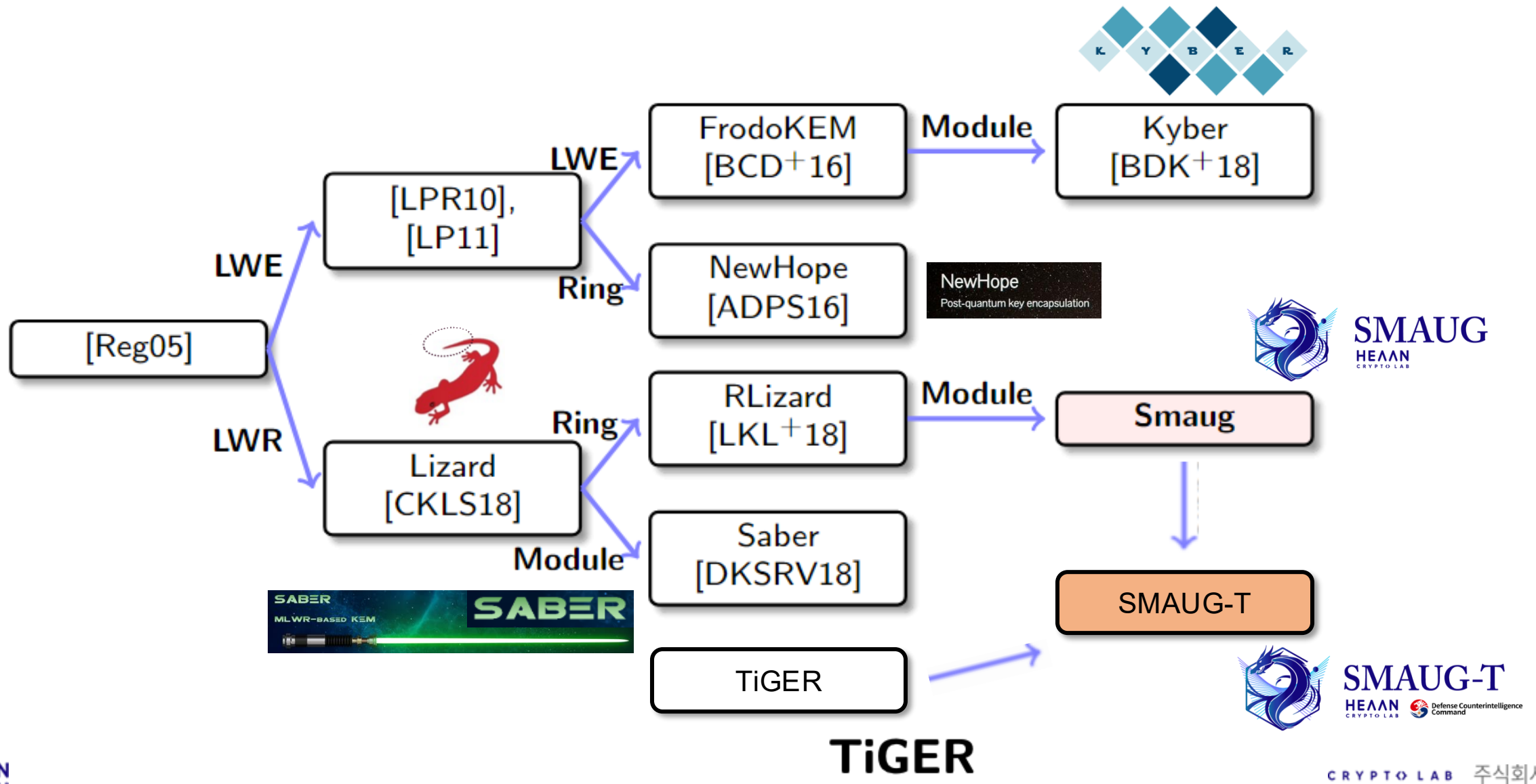
Bob

1. key  $k$  and ciphertext  $c$   
 $(k, c) \leftarrow \text{Encap}(pk)$



2. Shared key  
 $k \leftarrow \text{Decap}(sk, c)$

→ Both have the same key  $k$





- **Learning With Errors (LWE)**

- Gaussian noise is added:

$$\left( \begin{matrix} n \\ m \end{matrix} \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{A}, \mathbf{b} = \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{A} \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{s} + \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{e} \right) \in \mathbb{Z}_q^{m \times n} \times \mathbb{Z}_q^m$$

- **Learning With Rounding (LWR)**

- Noise comes from rounding:

$$\left( \begin{matrix} n \\ m \end{matrix} \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{A}, \mathbf{b} = \left\lfloor \frac{p}{q} \cdot \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{A} \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{s} \right\rfloor \right) \in \mathbb{Z}_q^{m \times n} \times \mathbb{Z}_p^m$$

- **Module LWE (MLWE)**

- Gaussian noise is added:

$$\left( \begin{matrix} n \\ m \end{matrix} \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{A}, \mathbf{b} = \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{A} \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{s} + \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{e} \right) \in R_q^{m \times n} \times R_q^m$$

$a_{0,0}(x) \in R_q$  (indicated by a red box on the first element of  $\mathbf{A}$ )

- **Module LWR (MLWR)**

- Noise comes from rounding:

$$\left( \begin{matrix} n \\ m \end{matrix} \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{A}, \mathbf{b} = \left\lfloor \frac{p}{q} \cdot \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{A} \begin{bmatrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{bmatrix} \mathbf{s} \right\rfloor \right) \in R_q^{m \times n} \times R_p^m$$

$a_{0,0}(x) \in R_q$  (indicated by a red box on the first element of  $\mathbf{A}$ )

$n$ : power of two integer  
 $R = \mathbb{Z}[x]/(x^n + 1)$



- **SMAUG-T.PKE**

$$n = 256$$

$$R = \mathbb{Z}[x]/(x^n + 1)$$

- MLWE-based Public Key

pk:  $\left( \begin{matrix} \overset{k}{\underbrace{\begin{matrix} k \\ \mathbf{A} \end{matrix}}} , \mathbf{b} = \begin{matrix} \mathbf{A} \end{matrix} \begin{matrix} \mathbf{s} \end{matrix} + \begin{matrix} \mathbf{e} \end{matrix} \end{matrix} \right), \text{ sk: } \begin{matrix} \mathbf{s} \end{matrix}$

Each element. is polynomial in  $R_q$

$$\mathbf{A} \in R_q^{k \times k}, \mathbf{b} \in R_q^k \text{ for } k = 2, 3, 4$$

\* secret (s): sparse ternary (HWT)  
 \* noise (e): discrete Gaussian (dG)

- MLWR-based Encryption

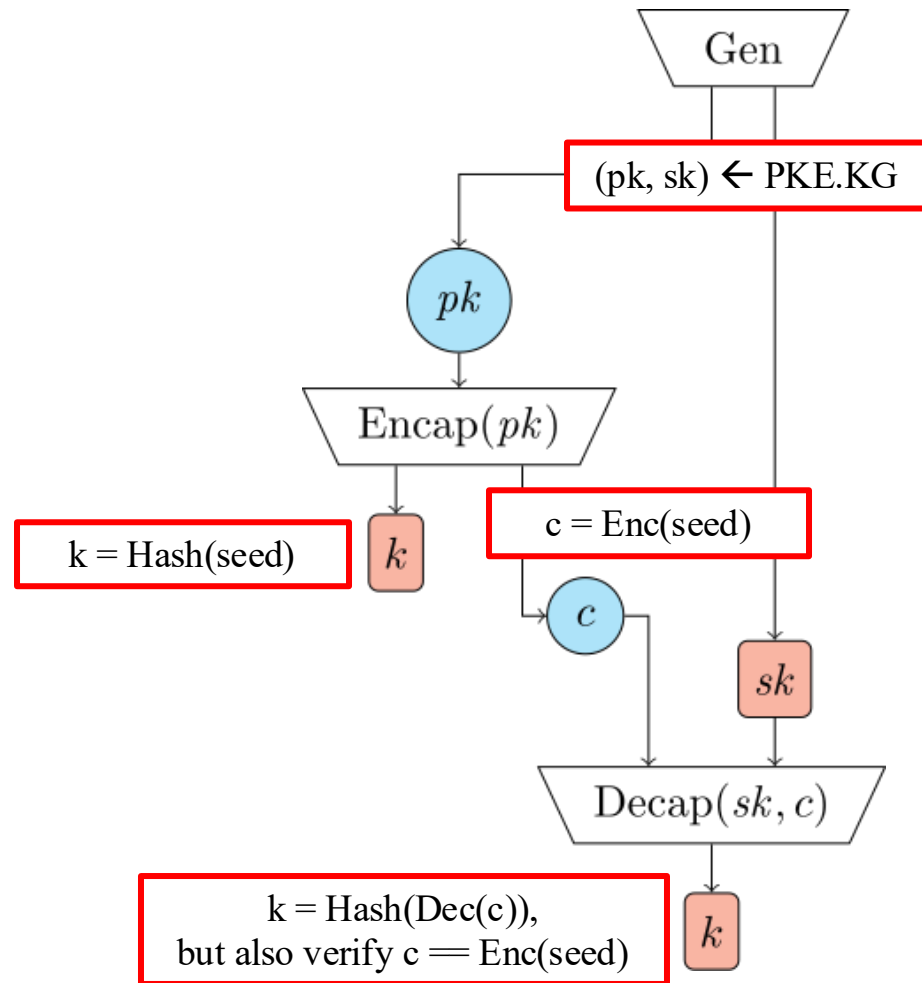
$$\left( \left[ \frac{p}{q} \cdot \begin{matrix} \mathbf{r} \end{matrix} \begin{matrix} \mathbf{A} \end{matrix} \right], \left[ \frac{p'}{q} \cdot \begin{matrix} \mathbf{r} \end{matrix} \begin{matrix} \mathbf{b} \end{matrix} + \frac{p'}{q} \cdot \begin{matrix} \Delta M \end{matrix} \right] \right)$$

$$\mathbf{r} \in R_q^k$$

\* random (r): sparse CBD (spCBD)  
 \* noise: rounding error



- Fujisaki-Okamoto Transform (FO):
  - SMAUG-T.PKE  $\rightarrow$  SMAUG-T.KEM







HAETAETAE  
HEΛΛN  
CRYPTO LAB

# HAETAETAE

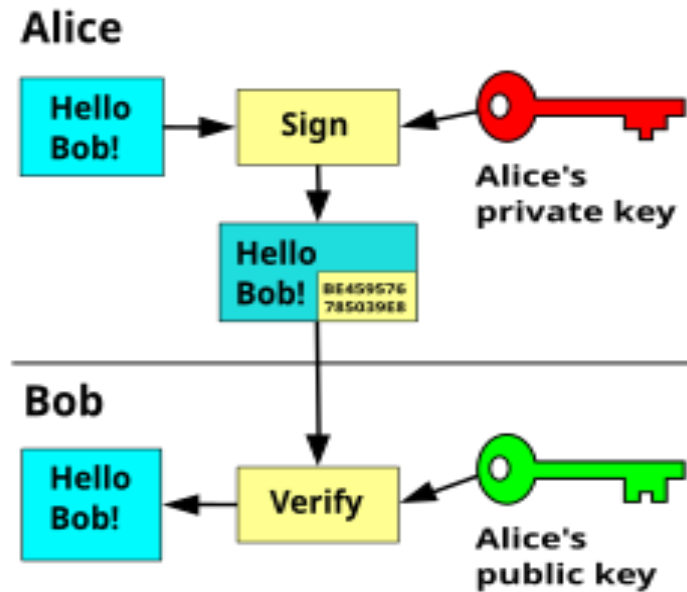
서울대학교  
SEOUL NATIONAL UNIVERSITY



Joint work between **Seoul National University (서울대학교)**, **CryptoLab, Inc. ((주)크립토랩)**,  
**École Normale Supérieure de Lyon (ENS de Lyon)**, **Ruhr-Universität Bochum (RUB)**,  
and **Deutsches Forschungszentrum für Künstliche Intelligenz (DFKI)**.



- Digital signatures



$(sk, vk) \leftarrow \text{KeyGen}$  and broadcast  $vk$

(knows  $sk$ )



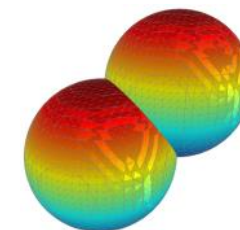
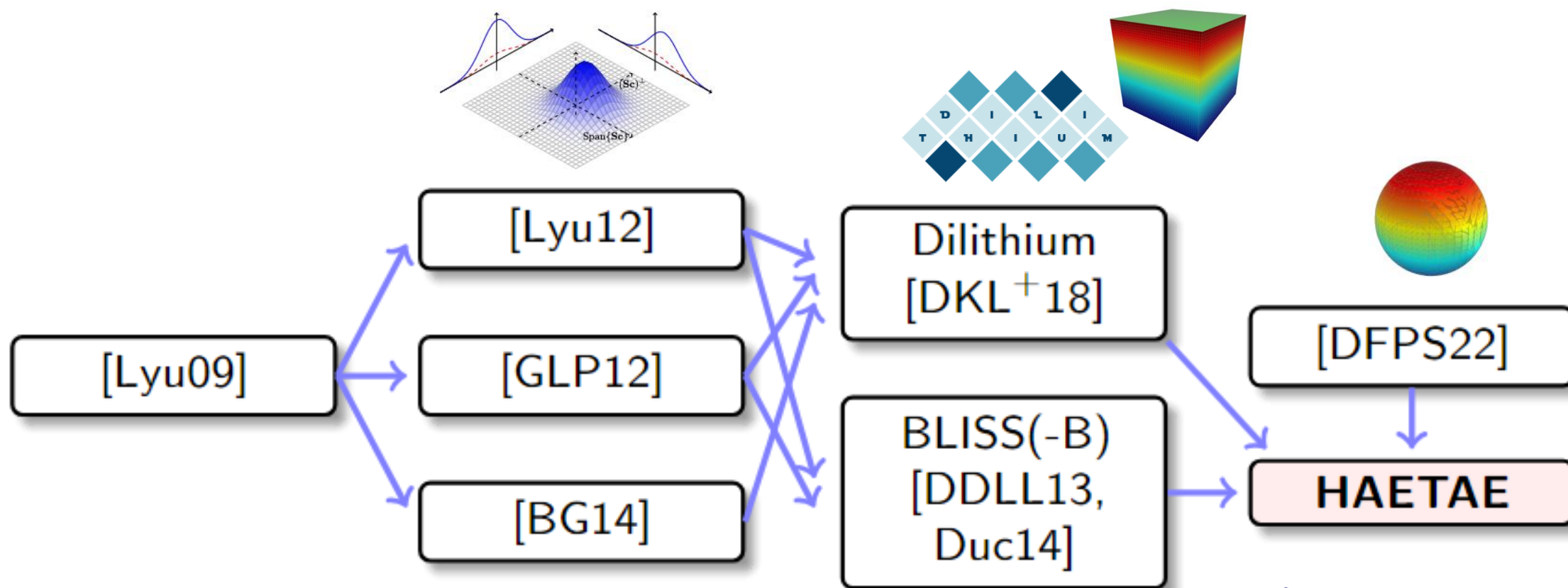
signature  $\sigma \leftarrow \text{Sign}(sk, m)$

(knows  $vk$ )



$(m, \sigma)$

$\text{Verify}(vk, m, \sigma) = \text{accept (or reject)}$





- **Short Integer Solution (SIS)**

- Can we find a short solution  $x \in \mathbb{Z}_q^n \setminus \{0\}$  of  $Ax \equiv 0 \pmod{q}$  for given  $A \in \mathbb{Z}_q^{m \times n}$ ?

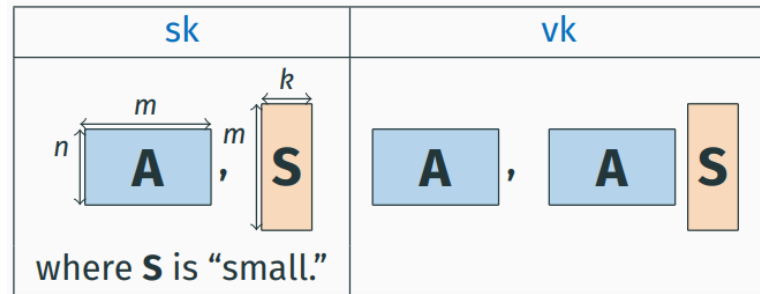
- **Module Variant (MSIS)**

- Can we find a short solution  $x \in R_q^n \setminus \{0\}$  of  $Ax \equiv 0$  in  $R_q$  for given  $A \in R_q^{m \times n}$ ?

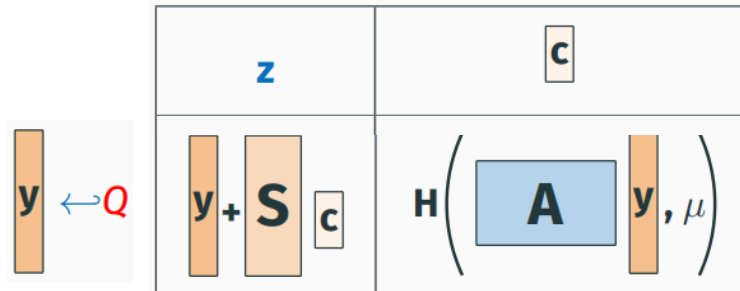


## “Fiat-Shamir with Abort” Paradigm

- KeyGen



- Sign



But with **rejection sampling**, since

$$z = y + S \parallel c \text{ depends on } S.$$

- Verify

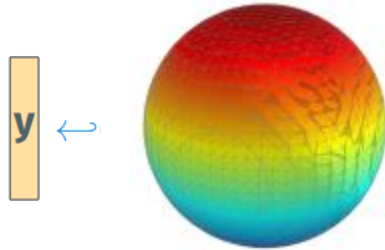
- Is  $z$  short?
- Is  $c = H(Az - Tc)$ ?

$$\begin{aligned}
 A \parallel z - T \parallel c &= A \parallel (y + S \parallel c) - A \parallel S \parallel c \\
 &= A \parallel y
 \end{aligned}$$

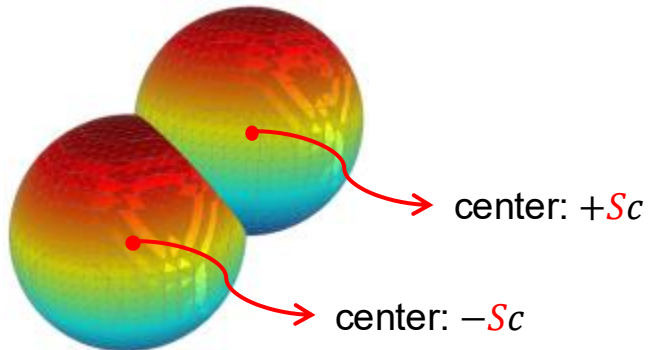


- HAETAE

- Randomness  $y$

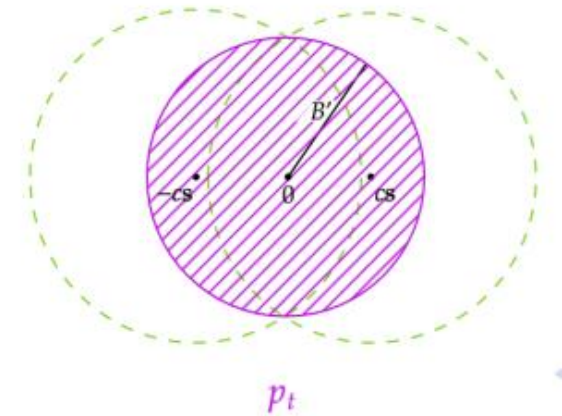
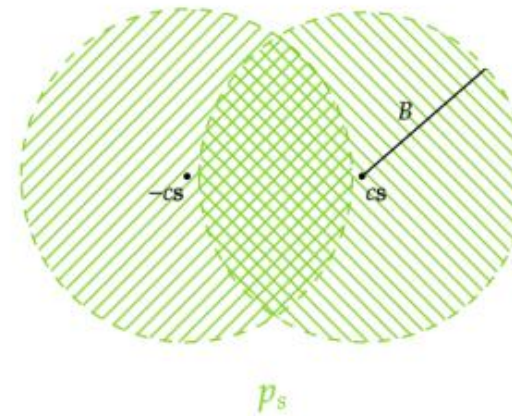


- $$z = y + (-1)^b \cdot Sc$$



- Rejection sampling on  $z$

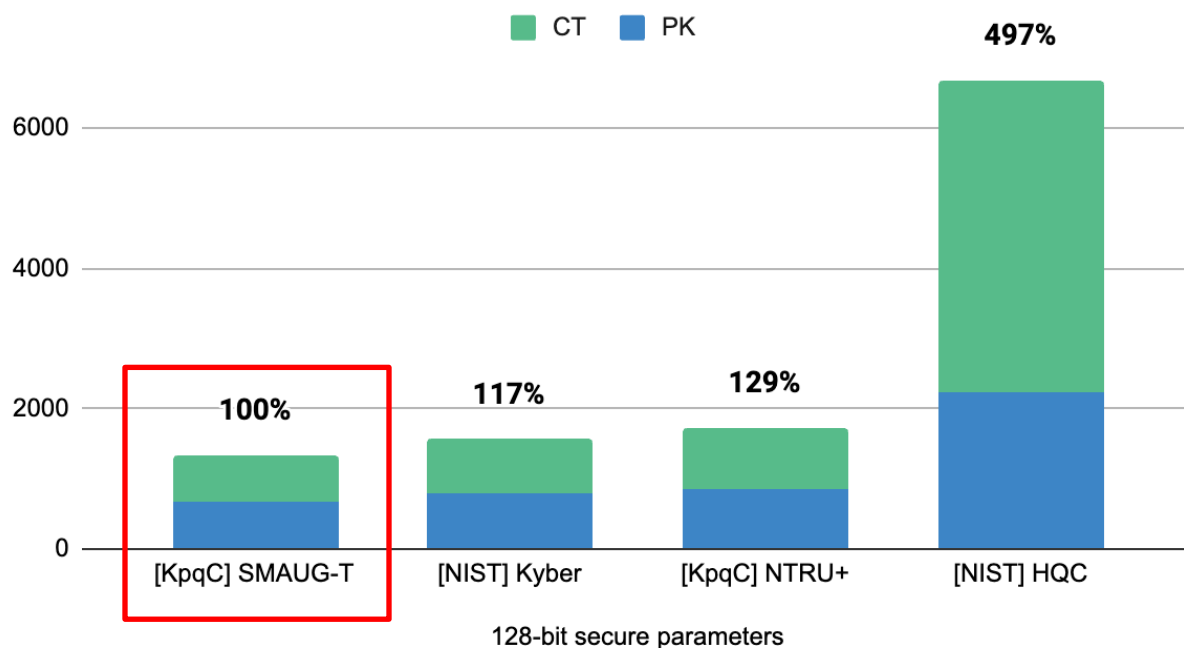
- From: **hyperball** centered at  $\pm Sc$
- To: **smaller hyperball** centered at 0





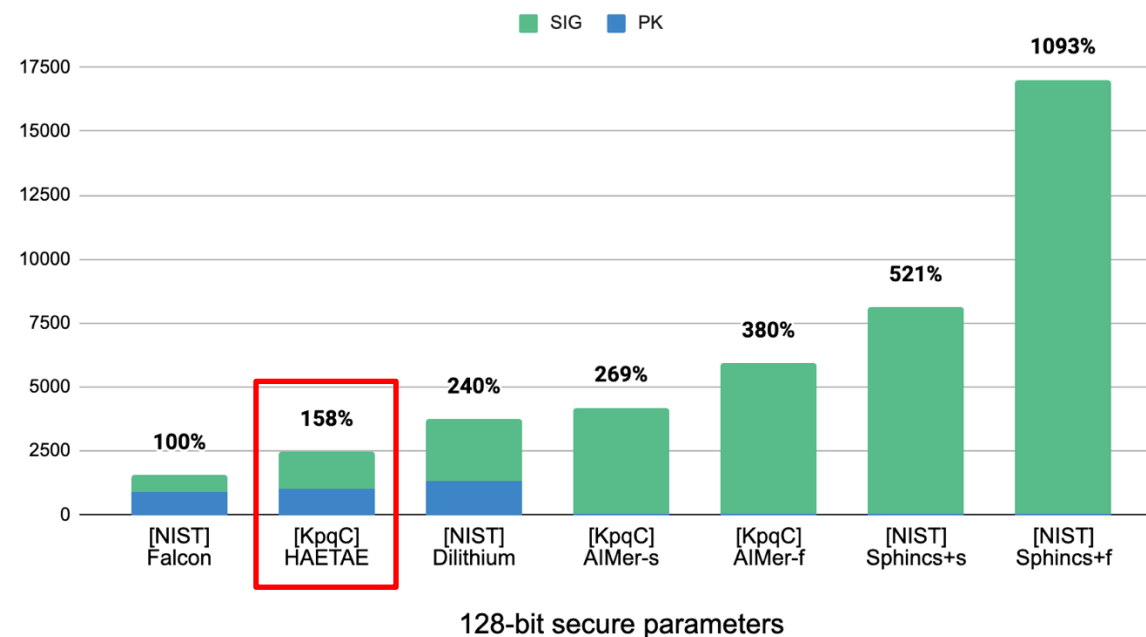
### KEM

Size in bytes



### Digital Signatures

Size in bytes



- Concretely-proven security
- Small sizes with good performance

NIST

Information Technology Laboratory

COMPUTER SECURITY RESOURCE CENTER

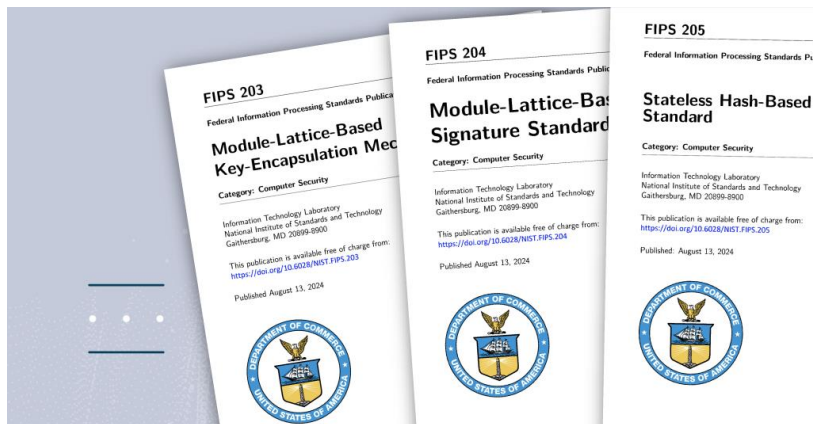
PROJECTS

Post-Quantum Cryptography PQC



양자내성암호연구단

# From Competition to Standard



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산업표준심의회  
2023년 4월 3일 개정





- **Migration to PQC**

- When? ASAP, Harvest Now, Decrypt Later (HNDL) attacks!



x: time that products and data must remain **secure**

y: time it takes to **migrate** to post-quantum cryptography

z: time it takes until cryptographically-relevant **quantum computers will be available**

- See HEaaN PQC Alliance Program at <https://heaanpqc.com> !!!



- **Korean Standards**
  - Ongoing efforts for governmental/industrial usages!
- **Hybrid with..**
  - Classical algorithms?
  - NIST-selected US standards?
- **Some Improvements..**

# Thank You