

Lecture 12: February 12th , 2020

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12.1 Context-sensitive analysis continued

12.1.1 Name Resolution

- In Java : A2 and A3
 1. Build global environment
 2. Resolve type names
 3. Build/check class hierarchy (methods/files)
 4. Disambiguate ambiguous namespace
 - In Java you can't simply determine the namespace based on location of usage
 5. Resolve "expressions" (Variables, static fields)
 6. Type checking
 7. Resolve methods instance (non-static) fields
- (2) Resolving Name Types Continued : More Rules

$$\frac{\begin{array}{llll} S \in \text{super}(T) & m \in \text{contain}(S) & \text{abstract} \notin \text{mods}(m) & \text{nodecl}(T, m) \\ S' \in \text{super}(T) & m' \in \text{contain}(S') & \text{abstract} \notin \text{mods}(m') & \text{sig}(m) = \text{sig}(m') \end{array}}{(m, m') \in \text{replace}}$$

$$\text{nodecl}(T, m) = \forall m' \in \text{declare}(T) : \text{sig}(m) \neq \text{sig}(m')$$

- An interface without super interfaces implicitly declares an abstract version of every public method in `java.lang.Object` (Refer to JLS 9.2)
- (3) Checking Hierarchy
 1. not $\exists T : T < T$ (Basically there should be no cycles)
 2. $\forall m, m' \in \text{declare}(T) : m \neq m' \implies \text{sig}(m) \neq \text{sig}(m')$
 3. $\forall m, m' \in \text{contains}(T) : \text{sig}(m) = \text{sig}(m') \implies \text{type}(m) = \text{type}(m')$
 4. $\forall m \in \text{contains}(T) : \text{abstract} \in \text{mods}(m) \implies \text{abstract} \in \text{mods}(T)$
 5. $\forall (m, m') \in \text{replace} : \text{static} \in \text{mods}(m') \iff \text{static} \in \text{mods}(m)$
 6. $\forall (m, m') \in \text{replace} : \text{type}(m) = \text{type}(m')$
 7. $\forall (m, m') \in \text{public} \in \text{mods}(m') \implies \text{public} \in \text{mods}(m)$
 8. No 8, because the course is attempting to be consistent with numbering schemes in other offerings. 8 would normally involve exception handling

- 9. $\forall(m, m') \in \text{replace} : \text{final} \notin \text{mods}(m')$
- 10. $\forall f, f' \in \text{declare}(T) : f' \neq f \implies \text{name}(f) \neq \text{name}(f')$
- (4) Disambiguating ambiguous name spaces
 - Notation
 - * P = package
 - * T = Type
 - * E = Expression
 - * M = Method
 - (Refer to JLS 6.5.2) $(a_1 \cdot a_2 \cdot a_3 \cdot a_4 \cdot a_k)$
 1. If local variables/parameters named a_1 exists , resolve to it
 2. If field name $a_1 \in \text{contain}(\text{enclosing class})$, resolve to it
 3. For each k in increasing order
 - * if $a_1 \cdot \dots \cdot a_k$ resolves to a type, resolve to it
 - a_{k+1} is a static field
 - $a_{k+2} \dots$ are non static fields
 4. Resolve instance fields, methods after type declaration
 5. General search outwards from inner most scope