CS 444 - Compiler Construction

Winter 2020

Lecture 6: January 22nd, 2020

Lecturer: Ondřej Lhoták Notes By: Harsh Mistry

6.1 Analysis Continued

6.1.1 Parsing Continued

6.1.1.1 LR(0) Parsing

Example 6.1 -

- $S \rightarrow E$ \$
- $E \rightarrow E + a$
- \bullet $E \rightarrow a$

You can not parse this grammar with LL because the praser does not how many times to apply the recursive rule.

6.1.1.2 LR(0) Parsing Algorithm

- for each a in ℓ \$ (ℓ is input)
 - While ($Reduce(stack) = \{A \rightarrow \gamma\}$)
 - * Pop $|\gamma|$ times
 - * Push A
 - if (Reject(stack + a)) then throw ERROR
 - push a
- $Reject(\alpha) = \alpha$ is not a viable prefix
- $Reduce(\alpha) = \{A \to \gamma | \exists P, \alpha = \beta \gamma \text{ and } \beta A \text{ is a viable prefix } \}$

Definition 6.2 α is a <u>viable prefix</u> if it is a prefix of a some sentential form which is any sequence that can be derived from the start <u>symbol</u>.

Only some viable prefixes ever occur as stacks and these subsets form a regular language.