1 Pseudocode

Algorithm Tree Fields Description

◇ *Shared*

- *Tree* `tree` : A binary tree of `Node`s. `root` is the root node.

◇ *Local*

- *Node* `leaf`: process's leaf in the tree.

◇ *Structures*

► *Node*

- **Node* `left, right, parent` : initialized when creating the tree.
- *BlockList* `blocks` implemented with an array.
- *int* `head`= 1: #`blocks` in `blocks`(-1). `blocks[0]` is a block with all fields equal to zero.
- *int* $\texttt{num}_{\texttt{propagated}}$= 0 : # groups of blocks that have been propagated from the node to its parent. Since it is incremented after propagating, it may be behind by 1.
- *int[]* `super`: `super[i]` stores an approximate index of the superblock of the blocks in `blocks` whose `group` field have value `i`.

► *Block* ▷ For a `block` in a `blocklist` we define *the prefix for the block* to be the blocks in the `BlockList` up to and including the `block`. *put the definitions before the pseudocode*

- *int* `group` : the value read from $\texttt{num}_{\texttt{propagated}}$ when appending this block to the node.

► *LeafBlock* extends *Block*

- *Object* `element` : Each block in a leaf represents a single operation. If the operation is `enqueue(x)` then `element=x`, otherwise `element=null`.
- *int* $\texttt{sum}_{\texttt{enq}}$, $\texttt{sum}_{\texttt{deq}}$: # enqueue, dequeue operations in the prefix for the block

► *InternalBlock* extends *Block*

- *int* $\texttt{end}_{\texttt{left}}$, $\texttt{end}_{\texttt{right}}$: index of the last subblock of the block in the left and right child
- *int* $\texttt{sum}_{\texttt{enq-left}}$: # enqueue operations in the prefix for $\texttt{left.blocks[end}_{\texttt{left}}\texttt{]}$
- *int* $\texttt{sum}_{\texttt{deq-left}}$: # dequeue operations in the prefix for $\texttt{left.blocks[end}_{\texttt{left}}\texttt{]}$
- *int* $\texttt{sum}_{\texttt{enq-right}}$: # enqueue operations in the prefix for $\texttt{right.blocks[end}_{\texttt{right}}\texttt{]}$
- *int* $\texttt{sum}_{\texttt{deq-right}}$: # dequeue operations in the prefix for $\texttt{right.blocks[end}_{\texttt{right}}\texttt{]}$

► *RootBlock* extends *InternalBlock*

- *int* `size` : size of the queue after performing all operations in the prefix for this block

Abbreviations:

- $\texttt{blocks[b].sum}_{\texttt{x}}\texttt{=blocks[b].sum}_{\texttt{x-left}}\texttt{+blocks[b].sum}_{\texttt{x-right}}$ (for b$\geq$0 and x $\in$ {enq, deq})
- $\texttt{blocks[b].sum=blocks[b].sum}_{\texttt{enq}}\texttt{+blocks[b].sum}_{\texttt{deq}}$ (for b$\geq$0)
- $\texttt{blocks[b].num}_{\texttt{x}}\texttt{=blocks[b].sum}_{\texttt{x}}\texttt{-blocks[b-1].sum}_{\texttt{x}}$
  (for b>0 and x $\in$ {$\emptyset$, enq, deq, enq-left, enq-right, deq-left, deq-right}, $\texttt{blocks[0].num}_{\texttt{x}}$=0)

Algorithm *Queue*

```
void Enqueue(Object e)    ▷ Creates a block with element e and appends it to the tree.
    block newBlock= new(LeafBlock)
    newBlock.element= e
    newBlock.sum_enq= leaf.blocks[leaf.head].sum_enq+1
    newBlock.sum_deq= leaf.blocks[leaf.head].sum_deq
    leaf.head+=1
    leaf.Append(newBlock)
end Enqueue

<int, int> FindResponse(int b_d, int i_d)    ▷ If E(root, b_e, i_e) is the response to the D(root, b_d, i_d) returns <b_e,i_e>. Returns <-1,--> if the queue is empty.
    if root.blocks[b_d-1].size + root.blocks[b_d].num_enq - (i + root.blocks[b_d-1].sum_deq) < 0 then
        return <-1,-->
    else
        r_enq= i + root.blocks[b_d-1].sum_deq - (root.blocks[b_d-1].size - root.blocks[b_d-1].sum_enq + root.blocks[b_d-1].sum_deq)
                                                    ▷ size-enqs+deqs=null deqs
        return root.DSearch(r_enq, b_d)
    end if
end FindResponse

Object Dequeue()
    block newBlock= new(LeafBlock)    ▷ Creates a block with null value element, appends it to the tree, computes its order among operations, then computes and returns its response.
    newBlock.element= null
    newBlock.sum_enq= leaf.blocks[leaf.head].sum_enq
    newBlock.sum_deq= leaf.blocks[leaf.head].sum_deq+1
    leaf.Append(newBlock)
    <b_deq, r_deq>= IndexDeq(leaf.head, 1)    ▷ r_deq is the rank among the dequeues of the dequeue of the b_deq th block in the root containing.
    <b_enq, r_enq>= FindResponse(b_deq, r_deq)    ▷ E(root, b_enq, i_enq) is response to the D(root, b_deq, i_deq). If the response is null then
deqRest    r_enq is -1.
    if r_enq==-1 then
        output= null
    else
        output= GetEnq(b_enq, r_enq)    ▷ getting the reponse's element.
    end if
    return output
end Dequeue
```

void PROPAGATE()
firstRefresh   **if not** REFRESH() **then**
secondRefresh   REFRESH() ▷ Lemma Double Refresh
   end if
   if `this` **is not** `root` **then**
     `parent`.PROPAGATE()
   end if
end PROPAGATE

boolean REFRESH()
readHead   `h= head`
   `<new,` $\texttt{np}_{\texttt{left}}$`,` $\texttt{np}_{\texttt{right}}$`>=` CREATEBLOCK(`s`) ▷ $\texttt{np}_{\texttt{left}}$, $\texttt{np}_{\texttt{right}}$ are the values read from the children's $\texttt{num}_{\texttt{propagated}}$ field.
addOP   **if** `new.num==0` **then return** `true` ▷ The block contains nothing.
cas   **else if** `blocks.tryAppend(new, h)` **then**
okcas   **for each** `dir` **in** {`left, right`} **do**
setSuper   `CAS(dir.super[`$\texttt{np}_{\texttt{dir}}$`], null, h+1)` ▷ Write would work too.
incNP   `CAS(dir.`$\texttt{num}_{\texttt{propagated}}$`,` $\texttt{np}_{\texttt{dir}}$`,` $\texttt{np}_{\texttt{dir}}$`+1)`
     end for
incrementHead1   `CAS(head, h, h+1)`
     return `true`
   else
     `CAS(head, h, h+1)` ▷ Even if another process wins, help to increase the head. The winner might have fallen sleep before increasing
incrementHead2 head.
     return `false`
   end if
end REFRESH

⇝ Precondition: `blocks[start..end]` contains a block with field `f` $\geq$ `i`
int BSEARCH(*field* `f`, *int* `i`, *int* `start`, *int* `end`) ▷ Does binary search for the value `i` of the given prefix sum `field`. Returns the index of the leftmost block in `blocks[start..end]` whose *field* `f` is $\geq$ `i`.
end BSEARCH

<Block, int, int> CREATEBLOCK(*int* `i`) ▷ Creates a block to be inserted into as ith `block` in `blocks`. Returns the created block as well as values read from each child's $\text{num}_{\text{propagated}}$ field. These values are used for incrementing the children's $\text{num}_{\text{propagated}}$ field if the block was appended to `blocks` successfully.
   `block newBlock=` NEW(*block*)
setGroup   `newBlock.group=` $\texttt{num}_{\texttt{propagated}}$
   for each `dir` **in** {`left, right`} **do**
lastLine $\texttt{index}_{\texttt{last}}$`= dir.head`
prevLine $\texttt{index}_{\texttt{prev}}$`= blocks[i-1].`$\texttt{end}_{\texttt{dir}}$
endDefLine   `newBlock.`$\texttt{end}_{\texttt{dir}}$`=` $\texttt{index}_{\texttt{last}}$
     $\texttt{block}_{\texttt{last}}$`= dir.blocks[`$\texttt{index}_{\texttt{last}}$`]`
     $\texttt{block}_{\texttt{prev}}$`= dir.blocks[`$\texttt{index}_{\texttt{prev}}$`]`
     ▷ `newBlock` includes `dir.blocks[`$\texttt{index}_{\texttt{prev}}$`+1..`$\texttt{index}_{\texttt{last}}$`]`.
setNP $\texttt{np}_{\texttt{dir}}$`= dir.`$\texttt{num}_{\texttt{propagated}}$
     `newBlock.`$\texttt{sum}_{\texttt{enq-dir}}$`= blocks[i-1].`$\texttt{sum}_{\texttt{enq-dir}}$ `+` $\texttt{block}_{\texttt{last}}$`.`$\texttt{sum}_{\texttt{enq}}$ `-` $\texttt{block}_{\texttt{prev}}$`.`$\texttt{sum}_{\texttt{enq}}$
     `newBlock.`$\texttt{sum}_{\texttt{deq-dir}}$`= blocks[i-1].`$\texttt{sum}_{\texttt{deq-dir}}$ `+` $\texttt{block}_{\texttt{last}}$`.`$\texttt{sum}_{\texttt{deq}}$ `-` $\texttt{block}_{\texttt{prev}}$`.`$\texttt{sum}_{\texttt{deq}}$
   end for
   if `this` **is** `root` **then**
computeLength   `newBlock.size= max(root.blocks[i-1].size+ newBlock.`$\texttt{num}_{\texttt{enq}}$ `- newBlock.`$\texttt{num}_{\texttt{deq}}$`, 0)`
   end if
   return `<b,` $\texttt{np}_{\texttt{left}}$`,` $\texttt{np}_{\texttt{right}}$`>`
end CREATEBLOCK

Algorithm Root

```
⇝ Precondition: root.blocks[end].sum_enq ≥ r_enq
<int, int> DSearch(int i, int end)
   ▷ Searches for the ith enqueue of the given prefix sum of bth block in the root. Returns the index of the leftmost block in root.blocks whose sum_enq is ≥ i.
    start= end-1
    while root.blocks[start].sum_enq≥r_enq do
        start= start-(end-start)
    end while
    return root.BSearch(sum_enq, r_enq, start, end)
end DSearch
```

Algorithm Node

```
⇝ Precondition: blocks[b].num_enq≥i
element GetEnq(int b, int i)
    if this is leaf then
        return blocks[b].element
    else if i ≤ blocks[b].num_enq-left then                                    ▷ i exists in the left child of this node
        subBlock= left.BSearch(sum_enq, i, blocks[b-1].end_left+1, blocks[b].end_left)    ▷ Search range of left child's subblocks of blocks[b].
        return left.Get(i-left.blocks[subBlock-1].sum_enq, subBlock)
    else
        i= i-blocks[b].num_enq-left
tChildGet        subBlock= right.BSearch(sum_enq, i, blocks[b-1].end_right+1, blocks[b].end_right)    ▷ Search range of right child's subblocks of blocks[b].
        return right.Get(i-right.blocks[subBlock-1].sum_enq, subBlock)
    end if
end GetEnq

⇝ Precondition: bth block of the node has propagated up to the root and blocks[b].num_enq≥i.
<int, int> IndexDeq(int b, int i)          ▷ Returns the rank of ith dequeue in the bth block of the node, among the dequeues in the root.
    if this is root then
xBaseCase        return <b, i>
    else
        dir= (parent.left==n)?  left:  right                                   ▷ check if a left or a right child
nputeSuper        superBlock= parent.BSearch(sum_deq-dir, i, super[blocks[b].group]-p, super[blocks[b].group]+p)
                                        ▷ superblock's group has at most p difference with the value stored in super[].
        if dir is right then
            i+= blocks[superBlock].sum_deq-left                                ▷ consider the dequeues from the right child
        end if
        return this.parent.IndexDeq(superBlock, i)
    end if
end Index
```

Algorithm Leaf

```
void Append(block blk)                                       ▷ Append is only called by the owner of the leaf.
appendEnd    head+=1
pendStart    blk.group= head
    blocks[head]= blk
    parent.Propagate()
end Append
```

Algorithm BlockList

▷ : Supports two operations `blocks.tryAppend(Block b)`, `blocks[i]`. Initially empty, when `blocks.tryAppend(b, n)` returns true `b` is appended to `blocks[n]` and `blocks[i]` returns ith block in the `blocks`. If some instance of `blocks.tryAppend(b, n)` returns `false` there is a concurrent instance of `blocks.tryAppend(b′, n)` which has returned `true`.`blocks[0]` contains an empty block with all fields equal to 0 and $\texttt{end}_{\texttt{left}}$, $\texttt{end}_{\texttt{right}}$ pointers to the first block of the corresponding children.

```
blocks[]:   array of blocks
boolean TryAppend(block blk, int n)
    return CAS(blocks[n], null, blk)
end TryAppend
```

2 Proof of Linearizability

$\boxed{\text{TEST}}$ Fix the logical order of definitions (cyclic refrences).

$\boxed{\text{TEST}}$ Is it better to show $\texttt{ops(EST}_{\texttt{n, t}}\texttt{)}$ with $\texttt{EST}_{\texttt{n, t}}$?

$\boxed{\text{Question}}$ A good notation for *the index of the* `b`?

$\boxed{\text{Question}}$ How to remove the notion of time? To say pre(n,i) contains n.blocks[0..i] instead of EST(n,t) which head=i at time t. Is it good? Furthermore, can we remove the notion of established blocks?

Definition 1 (Block)**.** A block is an object storing some statistics, as described in Algorithm Queue. It implicitly represents a set of operations. If `n.blocks[i]==b` we call `i` the *index* of block `b`. Block `b` is before block $\texttt{b}'$ in node `n` if and only if the index of the `b` is smaller than the index of the $\texttt{b}'$'s.

:subblock **Definition 2** (Subblock)**.** Block `b` is a *direct subblock* of `n.blocks[i]` if it is $\in \texttt{n.left.blocks[n.blocks[i-1].end}_{\texttt{left}}\texttt{+1..n.blocks[i].end}_{\texttt{left}}\texttt{]}$ $\cup\ \texttt{n.right.blocks[n.blocks[i-1].end}_{\texttt{right}}\texttt{+1..n.blocks[i].end}_{\texttt{right}}\texttt{]}$ (See line endDefLine 333 for the defined range). Block `b` is a subblock of a `n.blocks[i]` if it is a direct subblock of it or subblock of a direct subblock of it.

Definition 3 (Superblock)**.** Block b is direct superblock of block c if c is a direct subblock of b. Block b is superblock of block c if c is a subblock of b.

$\boxed{\texttt{def::ops}}$ **Definition 4** (Operations of a block)**.** A block lb in a leaf represents one operation which if it is `enqueue(x)` then `lb.element=x`, otherwise `element=null`. The set of operations of block *b* are the operations in the subblocks of *b*. We show the set of operations of block b by `ops(b)`.

For simplicity we say block `b` is propagated to node `n` or to a set of blocks `S` if `b` is in `n.blocks` or `S` or is a subblock of a block in `n.blocks` or `S`. We also say *b* contains *op* if $op \in ops(b)$.

Definition 5. A block `b` in `n.blocks` is *established* at time *t* if the last value written into `n.head` before *t* is greater than the index of `b` in `n.blocks` at time *t*. $EST_{n,\ t}$ is the set of established blocks at time *t* of node *n*.

$\boxed{\texttt{head}}$ **Observation 6.** *Once a block* `b` *is written in* `n.blocks[i]` *then* `n.blocks[i]` *never changes.*

dProgress **Lemma 7** (headProgress)**.** `n.head` *is non-decreasing over time and* $\texttt{n.blocks[i].end}_{\texttt{left}} \geq \texttt{n.blocks[i-1].end}_{\texttt{left}}$,$\texttt{n.blocks[i].end}_{\texttt{right}}$ $\geq \texttt{n.blocks[i-1].end}_{\texttt{right}}$.

Proof. The first claim follows trivially from the pseudocode since `n.head` is only incremented. Also when `n.blocks[i]` is created its $\texttt{end}_{\texttt{left}}$, $\texttt{end}_{\texttt{right}}$ are greater than or equal to the values in `n.blocks[i-1]`. Since $\texttt{blocks[i-1].end}_{\texttt{dir}}\texttt{<dir.head= blocks[i].end}_{\texttt{dir}}$ (Lines lastLine, prevLine ??). □

Lemma 8. *Every block has most one direct superblock.*

Proof. To show this we are going to refer to the way `n.blocks[]` is partitioned while propagating blocks up to `n.parent`. `n.CreateBlock(i)` merges the blocks in $\texttt{n.left.blocks[n.blocks[i-1].end}_{\texttt{left}}\texttt{..n.blocks[i].end}_{\texttt{left}}\texttt{]}$ and $\texttt{n.right.blocks[n.blocks[i-1].end}_{\texttt{right}}\texttt{..n.blocks[i]}$ (Lines lastLine, pr ??). Since $\texttt{end}_{\texttt{left}}\texttt{,end}_{\texttt{right}}$ are non-decreasing, so the range of the subblocks of `n.blocks[i]` which is $\texttt{(n.blocks[i-1].end}_{\texttt{dir}}\texttt{+1..n.blocks[i]}$ does not overlap with the range of the subblocks of `n.blocks[i-1]`. □

$\boxed{\texttt{append}}$ **Corollary 9** (No Duplicates)**.** *If* `op` *is in* `n.blocks[i]` *then there is no* $\texttt{j}\neq\texttt{i}$ *such that* $\texttt{op}\in\texttt{ops(n.blocks[j])}$.

dPosition **Invariant 10** (headPosition)**.** If the value of `n.head` is `h` then, `n.blocks[i]=null` for `i>h` and $\texttt{n.blocks[i]}\neq\texttt{null}$ for `i<h`.

Proof. The invariant is true initially since 1 is assigned to `n.head` and `n.blocks[x]` is null for every `x`. The truth of the invariant may be affected by writing into `n.blocks` or incrementing `n.head`.

Some value is written into `n.blocks[head]` only in Line 313. It is obvious that writing into `n.blocks[head]` presereves the invariant. The value of `n.head` is modified only in lines incrementHead1 incrementHead2 318], 321. Depending on wether the `TryAppend()`in Line cas 313 succeeded or not we show that the claim holds after the increment lines of `n.head` in either case. If `head` is incremented to h it is sufficient to show `n.blocks[h]`$\neq$`null` to prove the invariant still holds. In the first case the process applied a successful `TryAppend(new,h)` in line okcas 314, which means `n.blocks[h]` is not null anymore. Note that wether incrementHead1 318 returns true or false after Line `n.head` we know has been incremented from Line readHead 310. The failure case is also the same since it means some value is written into `n.blocks[head]` by some process. □

Explain More

shedOrder **Lemma 11** (establishedOrder)**.** *If time $t <$ time t', then* $\texttt{ops}(\text{EST}_{\texttt{n, t}}) \subseteq \texttt{ops}(\text{EST}_{\texttt{n, t}'})$.

Proof. Blocks are only appended (not modified) with CAS to `n.blocks[n.head]` and `n.head` is non-decreasing, so the set of operations in established blocks of a node can only grows. □

`CreateBlock()` aggregates the blocks in the children that are not already established in the parent into one block. If a `Refresh()` procedure returns true it means it has appended the block created by `CreateBlock()` into the parent node's sequence. So suppose two `Refresh`es fail. Since the first `Refresh()` was not successful, it means another CAS operation by a `Refresh`, concurrent to the first `Refresh()`, was successful before the second `Refresh()`. So it means the second failed `Refresh` is concurrent with another successful `Refresh()` that assuredly has read block before the mentioned line 35. After all it means if any of the `Refresh()` attempts were successful the claim is true, and also if both fail the mentioned claim still holds.

To review

:::headInc **Lemma 12** (head Increment)**.** *If an n.Refresh instance reaches Line 313 instance and reads head=h (310) after it terminates head is greater than h.*

Proof. If Line 318 or 318 succeeded the claim holds, otherwise another process has incremented the head. □

ueRefresh **Lemma 13** (trueRefresh)**.** *Let t_i be the time an instance of* `n.Refresh()` *is invoked and t_t be the time it terminates. Suppose the* `TryAppend(new, s)` *of the* `n.Refresh()` *returns* `true`*, then* $\texttt{ops(EST}_{\texttt{n.left, t}_i}\texttt{)} \cup \texttt{ops(EST}_{\texttt{n.right, t}_i}\texttt{)} \subseteq \texttt{ops(EST}_{\texttt{n, t}_t}\texttt{)}$.

Proof. From Lemma 11 we know that $\texttt{ops(EST}_{\texttt{n, t}_i}\texttt{)} \subseteq \texttt{ops(EST}_{\texttt{n, t}_t}\texttt{)}$. So it remains to show the operations of $\texttt{ops(EST}_{\texttt{n.left, t}_i}\texttt{)} \cup \texttt{ops(EST}_{\texttt{n.right, t}_i}\texttt{) - ops(EST}_{\texttt{n, t}_i}\texttt{)}$, which we call *new operations*, are all in $\texttt{ops(EST}_{\texttt{n, t}_t}\texttt{)}$. If `TryAppend` returns true a block `new` is written into `n.blocks[h]` (Line 313). We show $\texttt{ops(EST}_{\texttt{n.left, t}_i}\texttt{)} \subseteq \texttt{ops(EST}_{\texttt{n, t}_t}\texttt{)}$. The proof for the right child's claim is the same. Let `n.left.head` at $\texttt{t}_\texttt{i}$ be `hli`. Let `n.Refresh()` read `head` equal to h(Line 310). By the lines 332,331 the `new` block in `n.blocks[h]` contains $\texttt{n.left.blocks[n.blocks[h-1].end}_{\texttt{left}}\texttt{+1..left.head]}$. Since `left.head` is read after $\texttt{t}_\texttt{i}$ then $\texttt{ops(EST}_{\texttt{n.left, t}_i}\texttt{)} \subseteq \texttt{ops(n.left.blocks [0..left.head])}$. By Lemma 11 $\texttt{ops(n.left.blocks[0..n.blocks[h-1].end}_{\texttt{left}}\texttt{])} \subseteq \texttt{ops(EST}_{\texttt{n, t}_i}\texttt{)} \subseteq \texttt{ops(EST}_{\texttt{n, t}_t}\texttt{)}$. Since after line 321 we are sure that the `head` is incremented (Lemma 12) and `n.head=h+1` at t_t so the new block is established at t_t and the new block contains the new operations which is what we wanted to show.

□

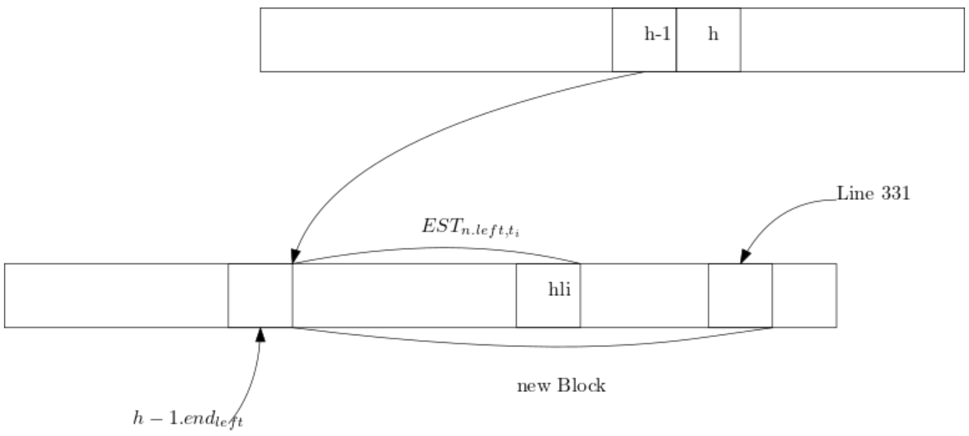

Figure 1: New established operations of the left child are in the new block.

ueRefresh **Lemma 14** (Precise True Refresh)**.** *Let t_i be the time an instance of* `n.Refresh()` *read the head (Line 310) and t_t be the time its* `TryAppend(new, s)` *terminates with and returns* `true` *(Line 313). We have* $\texttt{ops(EST}_{\texttt{n.left, t}_i}\texttt{)} \cup \texttt{ops(EST}_{\texttt{n.right, t}_i}\texttt{)} \subseteq \texttt{ops(n.blocks)}$.

Lemma 15 (Double Refresh)**.** *Consider two consecutive failed instances R_1, R_2 of* `n.Refresh()` *by some process. Let t_1 be the time R_1 is invoked and t_2 be the time R_2 terminated. We have* $\texttt{ops}(\texttt{EST}_{\texttt{n.left},\ t_1}) \cup \texttt{ops}(\texttt{EST}_{\texttt{n.right},\ t_1}) \subseteq \texttt{ops}(\texttt{EST}_{\texttt{n},\ t_2})$.

Proof.

If Line 313 of R_1 or R_2 returns `true`, then the claim is held by Lemma 13. Let R_1 read i and R_2 read $i+1$ from Line 310. If R_2 reads some value greater than $i+1$ in Line 310 it means a successful instance of `Refresh()` started after Line 310 of R_1 and finished its Line 318 or 321 before 310 of R_2, from Lemma 13 by the end of this instance $\texttt{ops}(\texttt{EST}_{\texttt{n.left},\ t_1}) \cup \texttt{ops}(\texttt{EST}_{\texttt{n.right},\ t_1})$ has been propagated.

Since R_2's `TryAppend()` returns `false` then there is another successful instance R_2' of `n.Refresh()` that has done `TryAppend()` successfully into `n.blocks[i+1]` before R_2 tries to append. In Figure 1 we see why the block R_2' is appending contains established block in the n's children at t_1, since it create a block reading the head after t_1. By Lemma 14 after R_2''s `CAS` we have $\texttt{ops}(\texttt{EST}_{\texttt{n.left},\ t_1}) \cup \texttt{ops}(\texttt{EST}_{\texttt{n.right},\ t_1}) \subseteq \texttt{ops}(\texttt{n.blocks})$. Also by Lemma 12 of R_2 head is more than $i+1$ after R_2ś 321 line, so the block appended by R_2' to `n` is established by then. To summarized t_1 is before R_2''s read head and R_2''s `CAS` is before R_2's termination. So $\texttt{ops}(\texttt{EST}_{\texttt{n.left},\ t_1}) \cup \texttt{ops}(\texttt{EST}_{\texttt{n.right},\ t_1}) \subseteq \texttt{ops}(\texttt{EST}_{\texttt{n},\ t_2})$. □

last sentence need more detail and should be earlier. define i and tell why R2prime exists

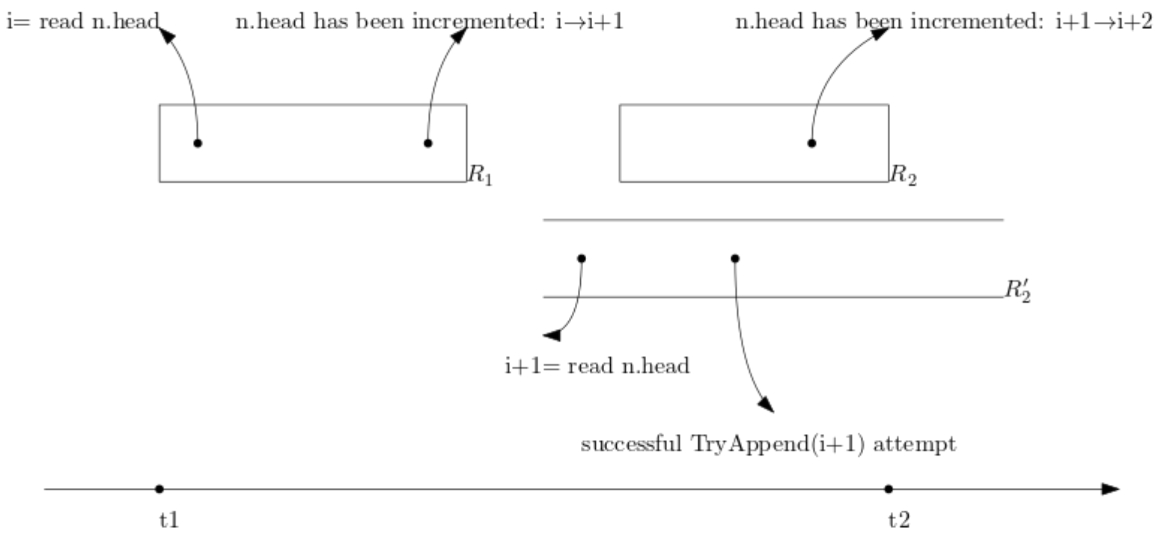

Figure 2: $t1 < r_1$ reading head < incrementing `n.head` from i to $i+1 < R_2'$ reading head < `TryAppend(i+1)` < incrementing `n.head` from $i+1$ to $i+2 < t2$

this chain with more depth should be in the proof

Corollary 16 (Propagate Step)**.** *All operations in* `n`*'s children's established blocks before line 302 are guaranteed to be in* `n`*'s established blocks after line 303.*

Proof. Lines 302 and 303 satisfy the preconditions of Lemma 15. □

Corollary 17. *After* `Append(blk)` *finishes* `ops(blk)`$\subseteq$`ops(root.blocks[x])` *for some* `x` *and only one* `x`*.*

Proof. Follows from Lemma 15, 9. □

Lemma 18 (Block Size Upper Bound)**.** *Each block contains at most one operation from each processs.*

Proof. By proof of contradiction, assume there are more than one operation from process p in block b in node n. A process cannot invoke more than one operations concurrently. From p 's operations in b, let op_1 be the first operation invoked and op_2 be the second one. Note that it is terminated before op_2 started. So before appending op_2 to the tree op_1 exists in every node from the path of p's leaf to the root. So there is some block b' before b in n containing op_1. op_1 existing in b an b' contradicts with 9. □

Lemma 19 (Subblocks Upperbound)**.** *Each block has at most p direct subblocks.*

Proof. It follows directly from Lemma 18 and the observation that each block contains at least one operation, induced from Line 312. □

ordering **Definition 20** (Ordering of operations inside the nodes)**.** ▶ Note that processes are numbered from 1 to p, left to right in the leaves of the tree and from Lemma blockSize 18 we know there is at most one operation from each process in a given block.

- We call operations strictly before op in the sequence of operations S, prefix of the op.
- $E(n, b)$ is the sequence of enqueue operations $\in$ `ops(n.blocks[b])` ordered by their process id.
- $E_{n,b,i}$ is the ith enqueue in $E(n, b)$.
- $D(n, b)$ is the sequence of dequeue operations $\in$ `ops(n.blocks[b])` ordered by their process id.
- $D_{n,b,i}$ is the ith enqueue in $D(n, b)$.
- Order of the enqueue operations in n: $E(n) = E(n, 1).E(n, 2).E(n, 3)...$
- $E_{n,i}$ is the ith enqueue in $E(n)$.
- Order of the dequeue operations in n: $D(n) = D(n, 1).D(n, 2).D(n, 3)...$
- $D_{n,i}$ is the ith dequeue in $D(n)$.
- Linearization: $L = E(root, 1).D(root, 1).E(root, 2).D(root, 2).E(root, 3).D(root, 3)...$

Note that in the non-root nodes we only order enqueues and dequeues among the operations of their own type. Since `GetENQ()` *only searches among enqueues and* `IndexDEQ()` *works on dequeues.*

Preconditions of all invocation of `BSearch` are satisfied.

Lemma 21 (Get correctness). *If* $\texttt{n.blocks[b].num}_{\texttt{enq}} \geq \texttt{i}$ *then* `n.GetENQ(b,i)` *returns* $E_{n,b,i}$.

Proof. We are going to prove this lemma by induction on the height of the tree. The base case for the leaves of the tree is pretty straight forward. Since leaf blocks contain exactly one operation then only `GetENQ(b,1)` can be called on leaves. `leaf.GetENQ(b,1)` returns the operation stored in the bth block of leaf l. For non leaf nodes in Line 404 it is decided that the `i`th enqueue in block `b` of internal node `n`resides in the left child or the right child of `n`. From Definition 20 we know operations in a block are ordered by their process id. Furthermore $\texttt{b.sum}_{\texttt{enq-lef}}$ stores the number of `enqueue()` operations from the `b`'s subblocks of the left child of `n`. So if `i` is greater than $\texttt{b.sum}_{\texttt{enq-lef}}$ it means ith operation is propagated from the right child, otherwise we should search for the ith enqueue i the left child subblocks. By definition 4 and 2 we need to search in subblocks of b which their range is $\texttt{n.left.blocks[n.blocks[i-1].end}_{\texttt{left}}\texttt{+1..n.blocks[i].end}_{\texttt{left}}\texttt{]}$ $\cup$ $\texttt{n.right.blocks[n.blocks[i-1].end}_{\texttt{right}}\texttt{+1..n.blocks[i].end}_{\texttt{right}}\texttt{]}$. If the enqueue we're looking for was in the right child as there are $\texttt{b.sum}_{\texttt{enq-left}}$ enqueues before it we need to search for $\texttt{i-b.sum}_{\texttt{enq-left}}$ (Line 409). of By definition of $E(n,b)$ operations from the left child come before the operations of the right child. Having $\texttt{sum}_{\texttt{enq}}$, the prefix sum of the number of enqueues we can compute the direct subblock containing the enqueue we are finding for with binary search. Then `n.child.GetENQ( block containing, order in the block)` is invoked which returns the correct operation by the hypothesis of the induction. ☐

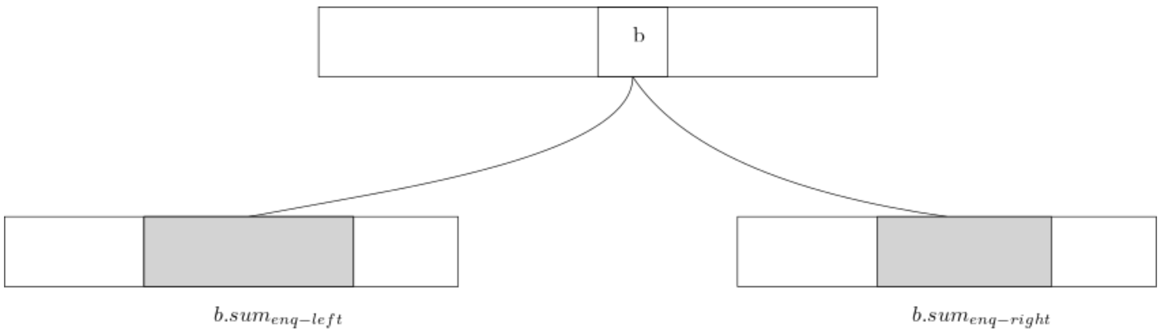

Figure 3: The number of enqueues from the left and the right child

I'm not sure it is going to be long and boring to talk about the parameters, since the reader can find out them.

dsearch **Lemma 22** (DSearch correctness). *If* `root.blocks[end].num`$_{\texttt{enq}}$$\geq$`i` *and* $E_{root,i}$ *is the response to some* `Dequeue()` *in* `root.blocks[end]` *then* `DSearch(i, end)` *returns* `b` *such that* `root.blocks[b]` *contains* $E_{root,b,i}$ *in* Θ(log*(*`root.blocks[b].size+root.blocks[end].size`*)* *steps.*

Proof. First we show `end-b`$\leq$`root.blocks[b].size+root.blocks[end].size`. We know each block size is greater than 0. So every block in `root.blocks[b..end]` contains at least one `Enqueue()` or one `Dequeue()`. There cannot be more than `root.blocks[b].size` `Dequeue()`s in `root.blocks[b+1..end-1]`, since the queue would become empty after bth block end before end and $E(n,i)$ could not be the response to to some `DEQ` in `end`. And since the lentgh of the queue would become `root.blocks[end].size` in the end so there cannot be more than `root.blocks[end].size` Dequeus in `root.blocks[b..end]`. Cause if there was more then the end's lenght would become more than `root.blocks[end].size`.

Now that we know there are at most `root.blocks[b].size+root.blocks[end].size` distance between end and b then with doubling search in log`root.blocks[b].size+root.blocks[end].size` steps we reach a block `c` that the `c.sum`$_{\texttt{enq}}$ is less than i and the distance between c and end is not more than 2$\times$`root.blocks[b].size+root.blocks[end].size`. So the binary search takes Θ(log`root.blocks[b].size+root.blocks[end].size`)) steps. ☐

Preconditions of all invocation of `BSearch` are satisfied.

Lemma 23 (Computing SuperBlock)**.** *After computing line 418 of* `n.IndexDEQ(b,i)`, `n.parent.blocks[superblock]` *contains* $D(n, b, i)$.

Proof. 1. Value read for `super[b.group]` in line 418 is not null.

▶ Values $\texttt{np}_{\texttt{dir}}$ read in lines 337, `super` are set before incrementing in lines 315,316. So before incrementing $\texttt{num}_{\texttt{propagated}}$, $\texttt{super[num}_{\texttt{propagated}}\texttt{]}$ is set so it cannot be null while reading.

2. `super[]` preserves order from child to parent; i.e. if in node `n` block `b` is before `c` then `b.group` $\leq$ `c.group` ▶ Line 329. Since $\texttt{num}_{\texttt{propagated}}$ is increasing.

3. Let `b, c` be in node `n`, if `b.group` $\leq$ `c.group` then `super[b.group]` $\leq$ `super[c.group]`

▶ Line 315.

4. `super[i+1]-super[i]`$\leq p$

▶ In a Refresh with successful CAS in line 46, `super` and `counter` are set for each child in lines 48,49. Assume the current value of the counter in node `n` is `i+1` and still `super[i+1]` is not set. If an instance of successful `Refresh(n)` finishes `super[i+1]` is set a new value and a block is added after `n.parent[sup[i]]`. There could be at most p successful unfinished concurrent instances of `Refresh()` that have not reached line 49. So the distance between `super[i+1]` and `super[i]` is less than p.

5. Superblock of `b` is within range $\pm 2p$ of the `super[b.time]`.

▶ `super[i]` is the index of the superblock of a block containing block b, followed by Lemma 26. It is trivial to see that `n.super` and `n.b.counter` are increasing. `super(b)` is the real superblock of b. `super(t)` is the index of the superblock of the last block with time `t`. If `b.time` is `t` we have:

$$super[t] - p \leq super[t-1] \leq super(t-1] \leq super(b) \leq super(t+1) \leq super(t+1] \leq super[t] + p$$

□

We call the dequeues that return some value *non-null dequeues*. rth non-null dequeue returns the element of th rth enqueue. We can compute # non-null dequeues in the prefix for a block this way: #non-null dequeues= size- #enqueues. Note that the ith dequeue in the given block is not a non-null dequeue.

Lemma 24 (Index correctness). `n.IndexDEQ(b,i)` *returns the rank in* $D(root)$ *of* $D_{n,b,i}$.

Proof. We will prove this by induction on the distance of `n` from the `root`. We can see the base case `root.IndexDEQ(b,i)` is trivial (Line indexBaseCase 415). In the non-root nodes `n.IndexDEQ(b,i)` computes the superblock of the ith Dequeue in the bth block of `n` in `n.parent` by Lemma superBlock 23 (Line computeSuper 418). After that the order in $D(n.parent, superblock)$ is computed and `index()` is called on `n.parent` recursively. Then if the operation was propagated from the right child the number of dequeues from the left child are added to it, because the left child operations come before the right child operations (ordering 20). ☐

Do I need to talk about the computation of the order in the parent which is based on the definition of ordering of dequeues in a block?

mputeHead **Lemma 25** (Computing Queue's Head block). *Let* S *be the state of an empty queue if the operations in prefix in* L *of ith dequeue in* $D(root, b)$ *are applied on it.* `FindResponse()` *returns* `(b, i)` *which* $E(root, b, i)$ *is the the head of the queue in* S. *If the queue is empty in* S *then it returns* `<-1,-->`.

Proof. The sizeof the queue if the operations in the prefix for the bth block in the root are applied with the order of L is stored in the `root.blocks[b].size`. It is computed while creating the block in Line computeLength 342. If the sizeof a queue is greater than 0 then a `Dequeue()` would decrease the sizeof the queue, otherwise the sizeof the queue remains 0. Having sizeof the queue after the previous block and number of enqueues and dequeues in the block, Line computeLength 342 computes wether the queue becomes empty or the sizeof it. ☐

HOW? How to prove mathematicaly that $\texttt{ax(root.blocks[i-1].size+ b.num}_{\texttt{enq}}\texttt{ - b.num}_{\texttt{deq}}\texttt{, 0)}$ is the sizeof the queue after the block. I can only explain it here.

erCounter **Lemma 26** (super property). *If* `super[i]` $\neq$ `null`, *then* `super[i]` *in node* `n` *is the index of the superblock of a block with* `time=i`$\pm$`p`.

To write

mputeHead

Lemma 27 (Computing Queue's Head block)**.** *Let S be the state of an empty queue if the operations in prefix in L of ith dequeue in $D(root, b)$ are applied on it.* `FindResponse()` *returns* `(b, i)` *which $E(root, b, i)$ is the the head of the queue in S. If the queue is empty in S then it returns* `<-1,-->`.

Proof. The sizeof the queue if the operations in the prefix for the bth block in the root are applied with the order of L is stored in the `root.blocks[b].size`. It is computed while creating the block in Line `computeLength` 342. If the sizeof a queue is greater than 0 then a `Dequeue()` would decrease the sizeof the queue, otherwise the sizeof the queue remains 0. Having sizeof the queue after the previous block and number of enqueues and dequeues in the block, Line `computeLength` 342 computes wether the queue becomes empty or the sizeof it. ☐

HOW? How to prove mathematicaly that `ax(root.blocks[i-1].size+ b.num`$_{enq}$ `- b.num`$_{deq}$`, 0)` is the sizeof the queue after the block. I can only explain it here.

To review

Theorem 28 (Main)**.** *The queue implementation is linearizable.*

Proof. We choose L in Definition 20 to be linearization ordering of operations and prove if we linearize operations as L the queue works consistently. □

Lemma 29. *Operations in a block have a time point in common (There is a time t all the operations are running).*

Lemma 30 (satisfiability)**.** *L can be a linearization ordering.*

Proof. Once some operations are aggregated in one block they will be propagated together up to the root and we can linearize them in any order among themselves (previous lemma). Furthermore in L we arbitrary choose the order to be by process id, since it makes computations in the blocks faster . □

Lemma 31 (correctness)**.** *If operations are applied as L on a sequential queue, the sequence of the responses would be the same as our algorithm.*

Proof. Old parts to review We show that the ordering L stored in the root, satisfies the properties of a linearizable ordering.

1. If op_1 ends before op_2 begins in E, then op_1 comes before op_2 in T.

   ► This is followed by Lemma 9. The time op_1 ends it is in root, before op_2, by Definition 20 op_1 is before op_2.

2. Responses to operations in E are same as they would be if done sequentially in order of L.

   ► Enqueue operations do not have any response so it does no matter how they are ordered. It remains to prove Dequeue d returns the correct response according to the linearization order. By Lemma 27 it is deduced that the head of the queue at time of the linearization of d is computed properly. If the Queue is not empty by Lemma 21 we know that the returning response is the computed index element.

□

Lemma 32 (Amortized time analysis)**.** `Enqueue()` *and* `Dequeue` *take* $O(\log 62p + q)$ *steps (amortized anlysis), which* p *is the number of processes and* q *is the size of the queue at the time of invocation.*

To write