



UVM Analysis Port Functionality and Using Transaction Copy Commands

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Austin

**Life is too short for bad
or boring training!**



Agenda

Basic queues, mailboxes and TLM FIFOs

← 1st pass

Subscriber satellite TV analogy

Analysis paths & analysis ports, exports, and imps

TLM FIFOs

← More detail

Importance of the `copy()` method

How analysis port connections work - `write()` method

Summary & Conclusions

The paper has more details
and more examples

Important SystemVerilog Features

PAY ATTENTION !!



- Queues

Queues will be used to store **component** handles

Can store class handles -
great for storing connected components

- `push_back()` method to put a handle into the queue
- `foreach()` method to walk through all stored handles
- Does not have blocking `get()` method

Not too useful for scoreboards

- Mailboxes

Mailboxes will be used to store **transaction** handles

Can store class handles -
great for storing transactions

- Has nonblocking `try_put()` method
- Has blocking `get()` method

Important for scoreboards

- Analysis path considerations:

Must include `write()` method

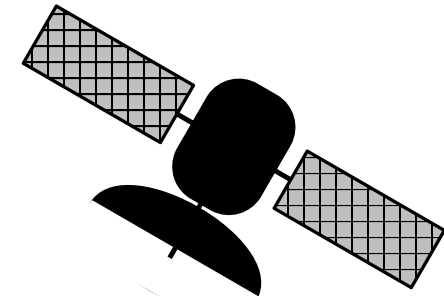
- Must start with `uvm_analysis_port` and end with `uvm_analysis_imp`

- `uvm_tlm_fifo` **cannot** terminate an analysis path
- `uvm_tlm_analysis_fifo` **CAN** terminate an analysis path

Built using mailboxes

Very useful for scoreboards !!

Subscriber Satellite TV Analogy



- Two ways to watch a broadcast satellite TV program
 - Watch the program live
 - Record the program to a DVR to view later

- Satellite programs are broadcast as scheduled

There might be 1,000's of viewers

There might be **NO** viewers

- No way to restart a broadcast program

No way to communicate back to the satellite

Other viewers would object to restarting the program

- Subscribers not allowed to change the live program

With the right equipment, you can modify your copy

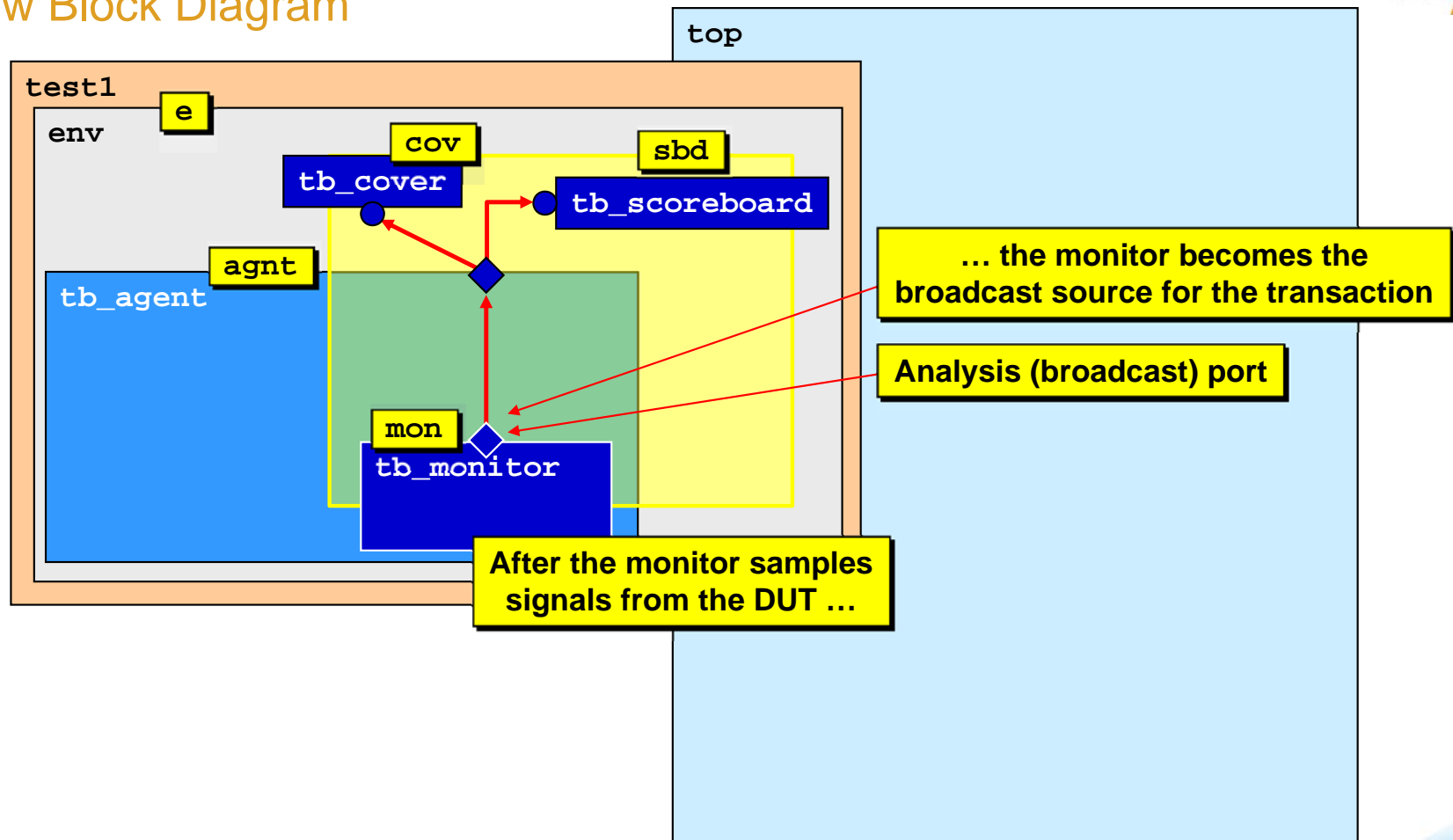


Analysis Port Connections

and TLM FIFOs

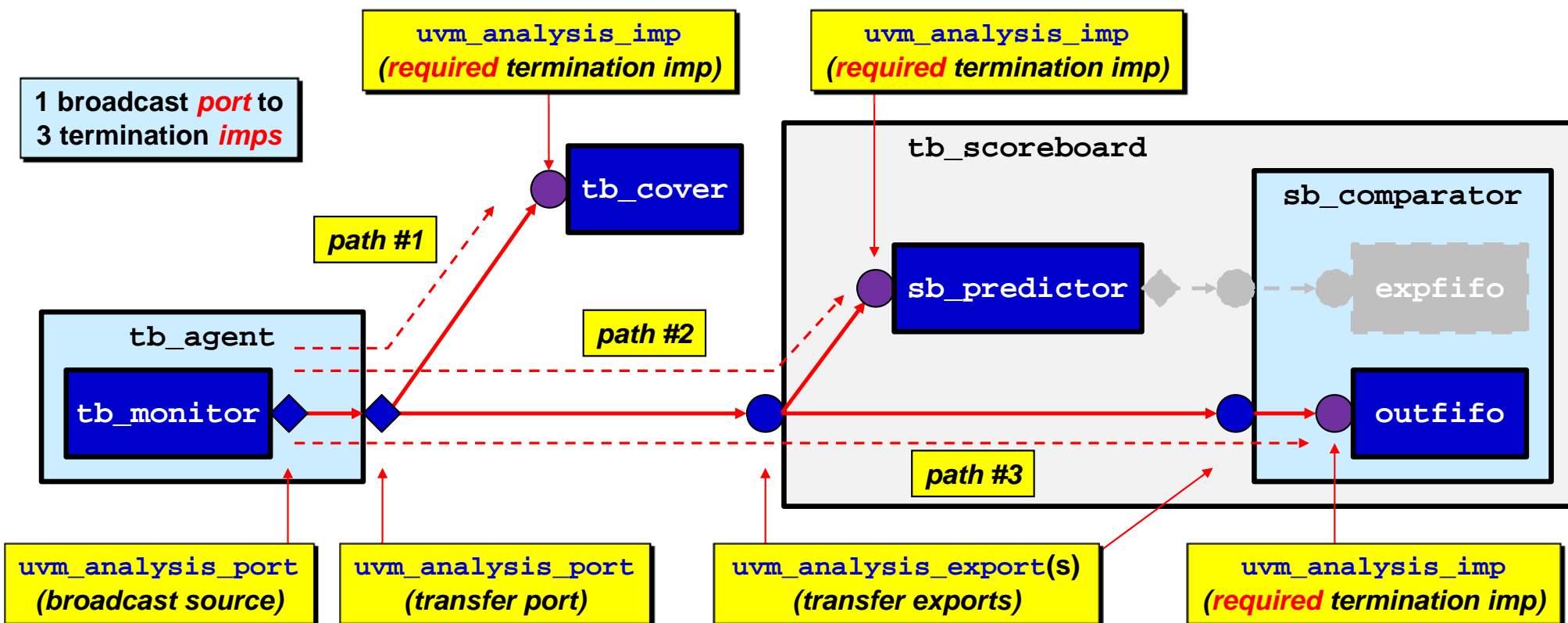


Overview Block Diagram



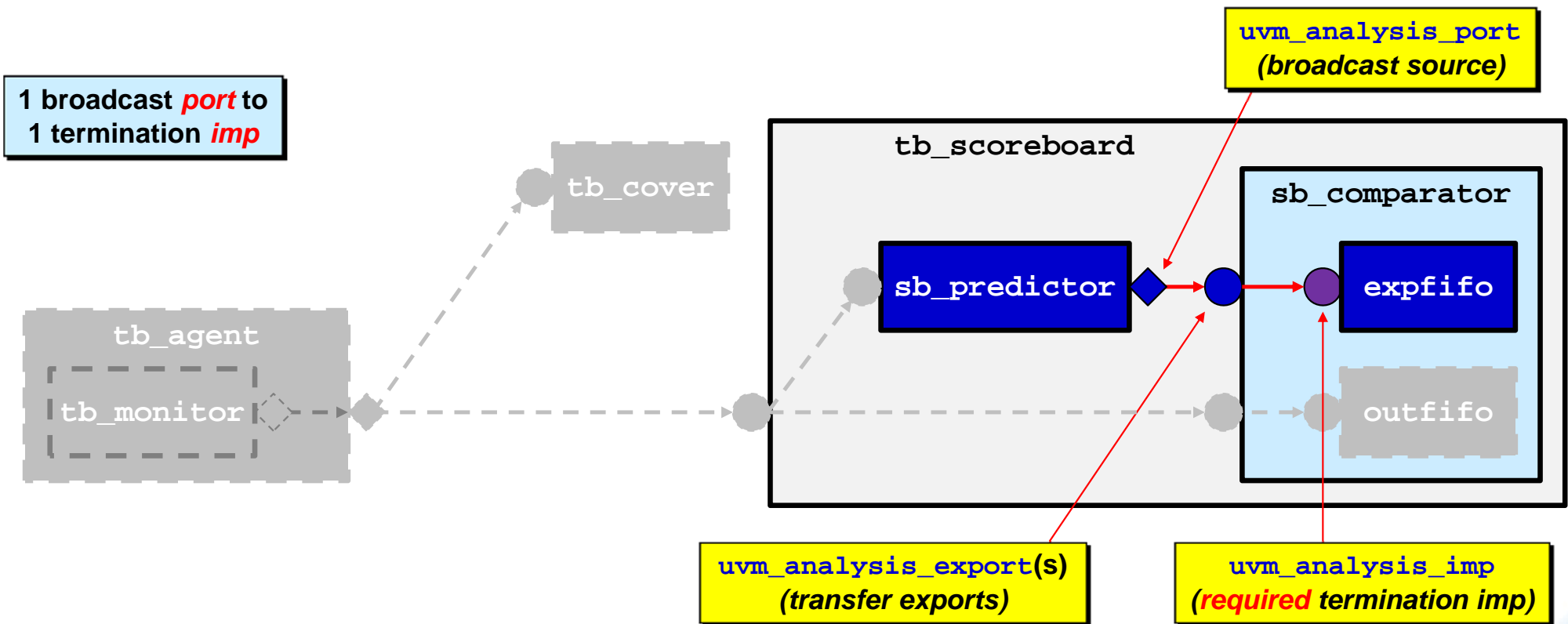
UVM Testbench Analysis Port Paths

Common Paths - Monitor to Multiple Subscribers



UVM Testbench Analysis Port Paths

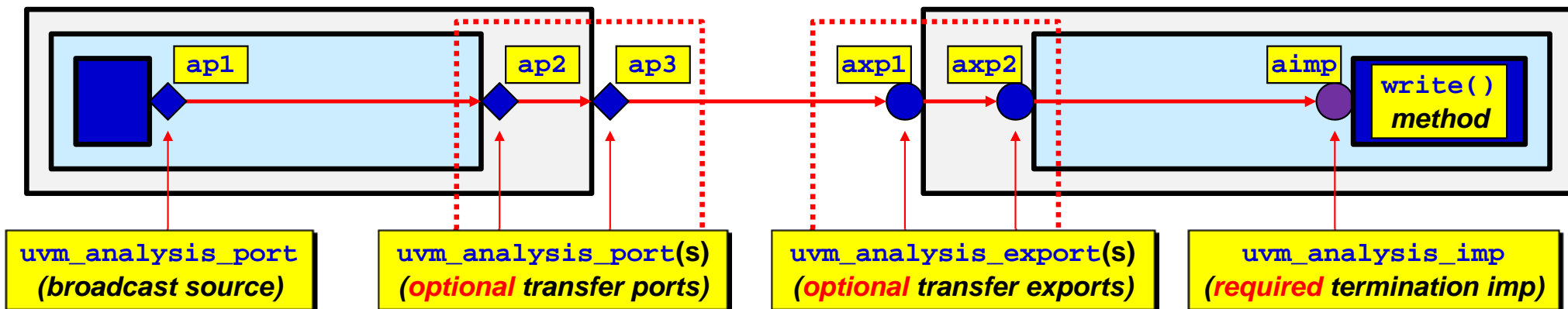
Common Paths - Predictor to Expected Transaction FIFO



UVM Analysis Port Paths

LEGAL Paths

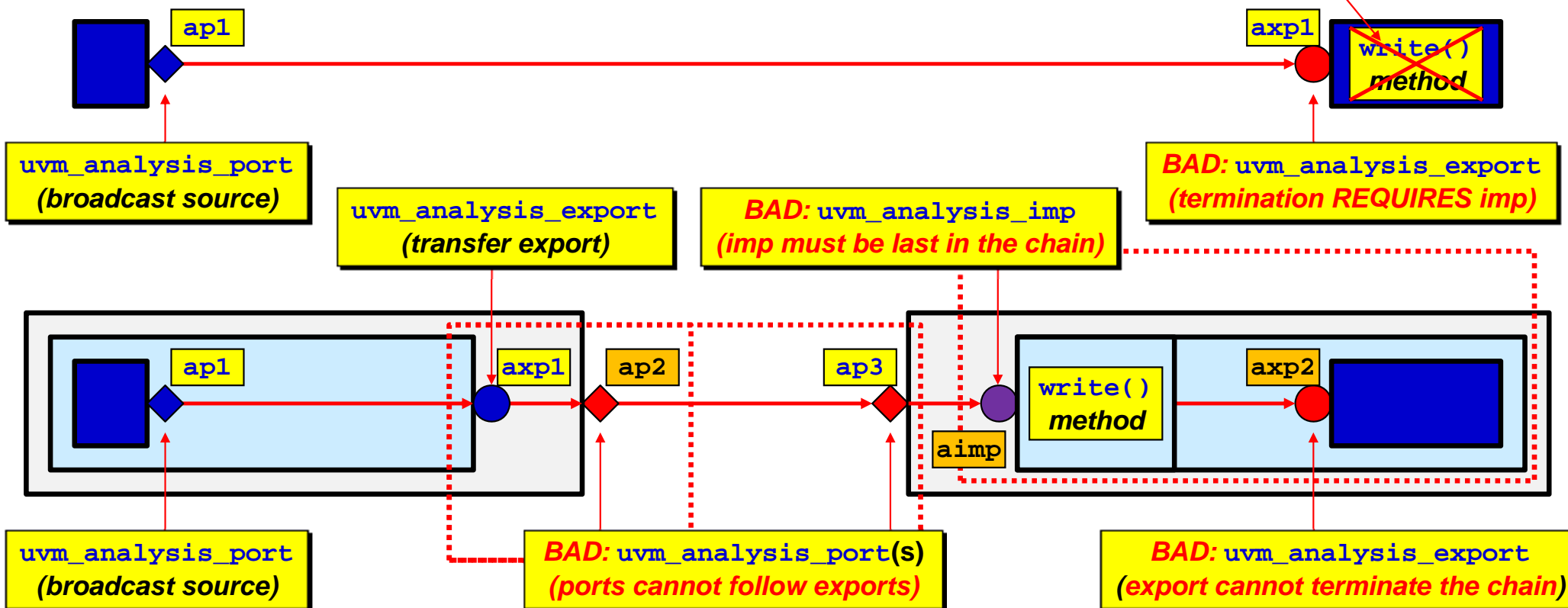
*The most simple analysis path is a
uvm_analysis_port (broadcast source)
with no subscribers*



UVM Analysis Port Paths

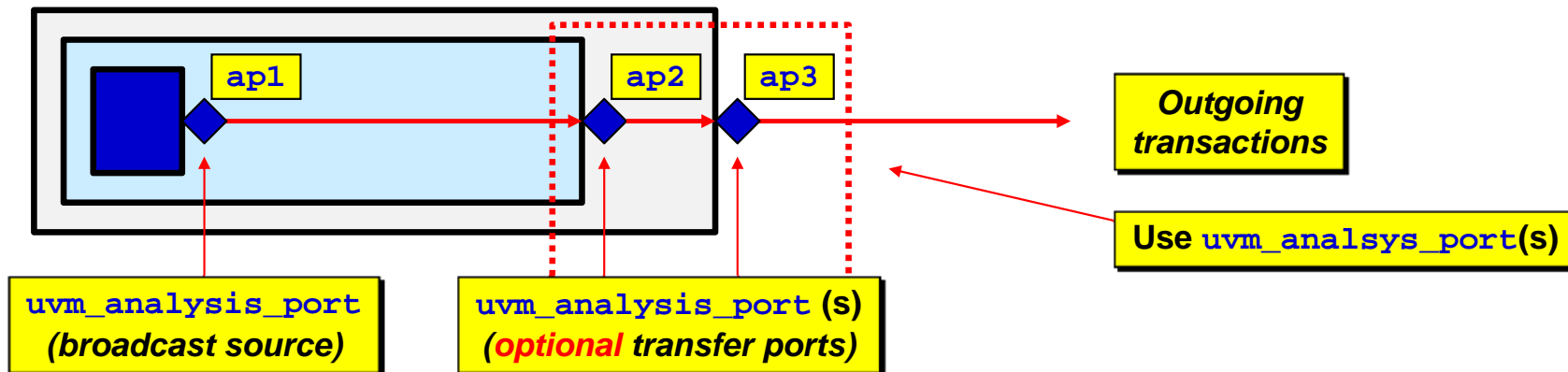
ILLEGAL Paths

Not automatically called from
a `uvm_analysis_export`



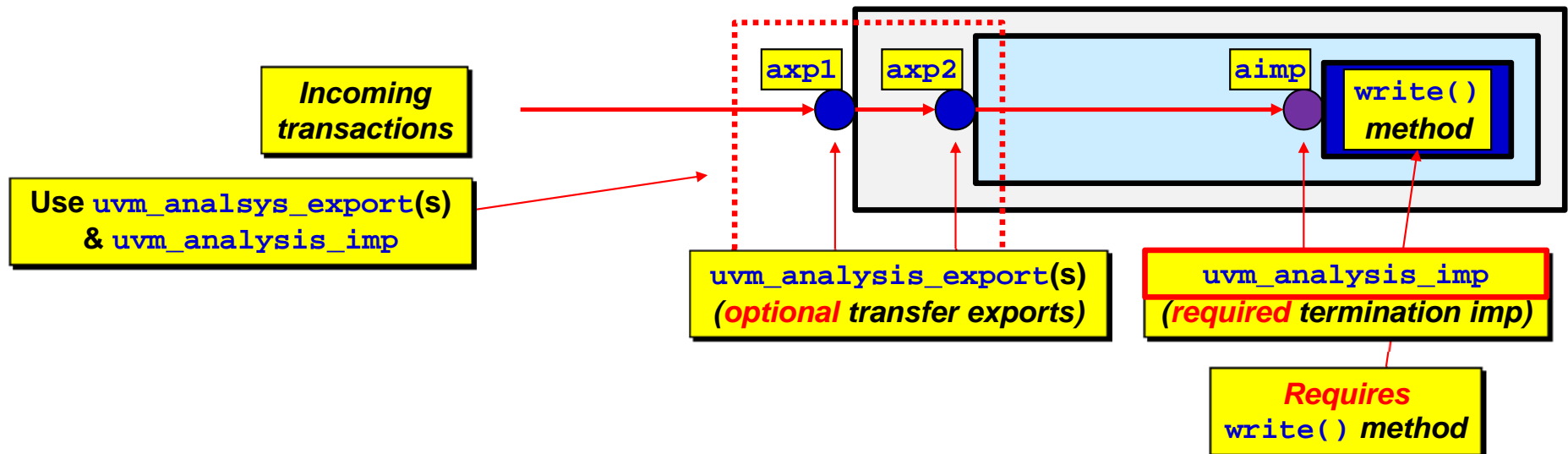
UVM Analysis Ports

Recommended Usage



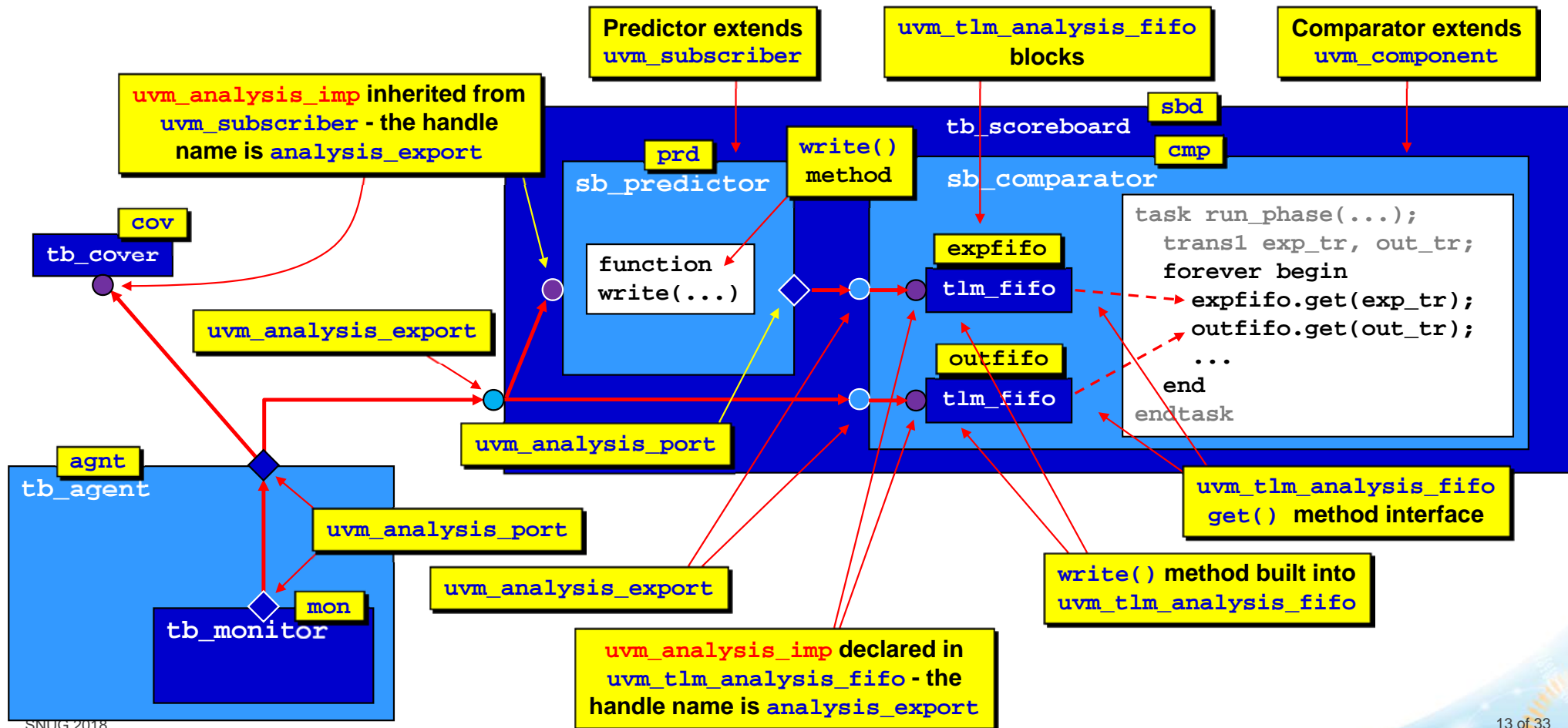
UVM Analysis Exports & Imps

Recommended Usage



Common Analysis Port Connections

Recommended Connections



TLM FIFOs - Definitions & Usage

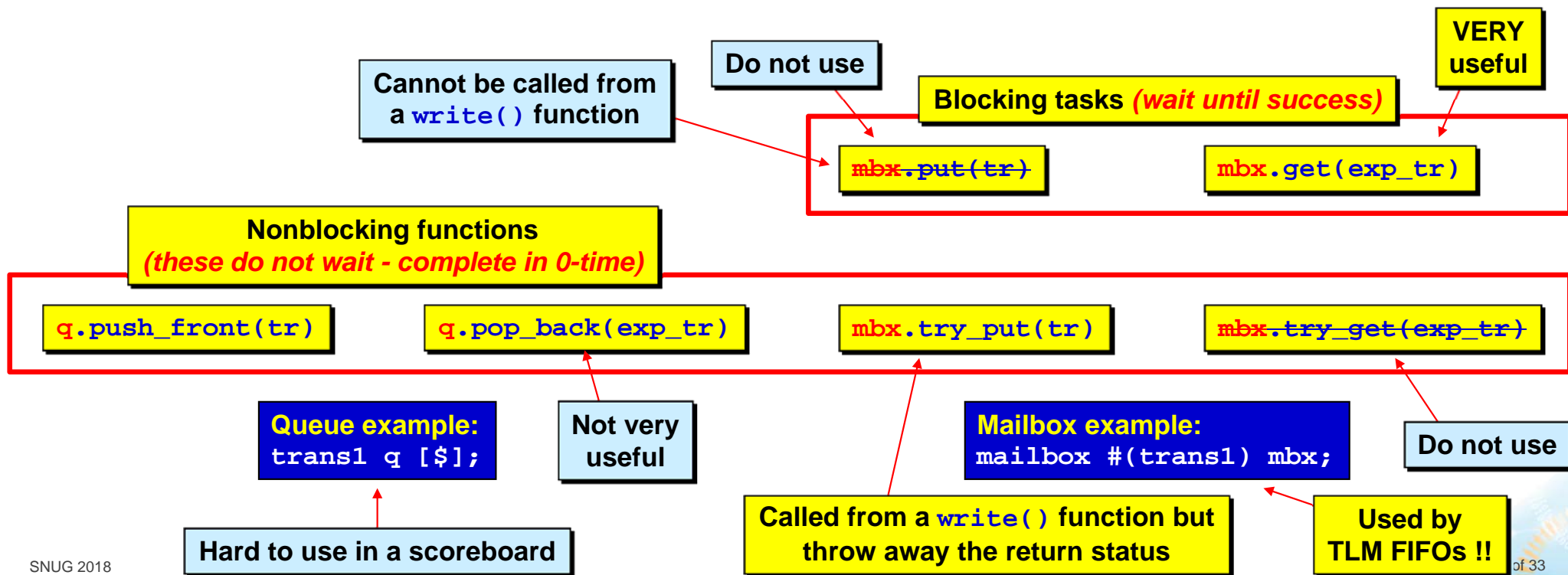


TLM FIFOs & Scoreboards

SystemVerilog Queues & Mailboxes

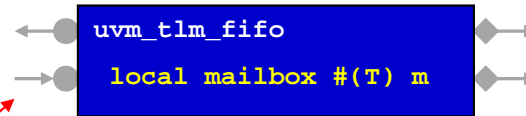


- Scoreboards typically store *expected* and *actual* transactions
- SystemVerilog has *queues* and *mailboxes* ← Which should be used?



uvm_tlm_fifo

Most Common Usage



Although there are 2 **non-analysis** `imp` ports and 2 **analysis** ports on the `uvm_tlm_fifo`, they typically are not used

The `uvm_tlm_fifo`, will be constructed to be unbounded. Example:
`exp_fifo=new("exp_fifo", this, 0);`

0 means unbounded

`try_put()` method stores into the mailbox

`uvm_tlm_fifo`
`local mailbox #(T) m`

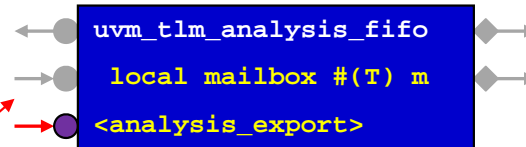
`get()` method retrieves from the mailbox

The `uvm_analysis_imp` `write(tr)` method will call `void'(try_put(tr))`

`void`-cast to throw away the `try_put()` return-status
(`try_put()` always succeeds on unbounded fifo's mailbox)

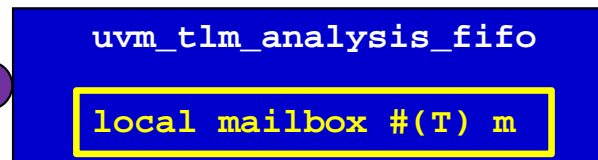
The scoreboard comparator will call the blocking `get(tr)` method and wait to retrieve a `uvm_tlm_fifo` transaction

Most Common Usage



The `uvm_tlm_analysis_fifo` is unbounded by default

Termination of an analysis-path



get() method retrieves
from the mailbox

`uvm_analysis_imp` with handle name `analysis_export` is almost always used

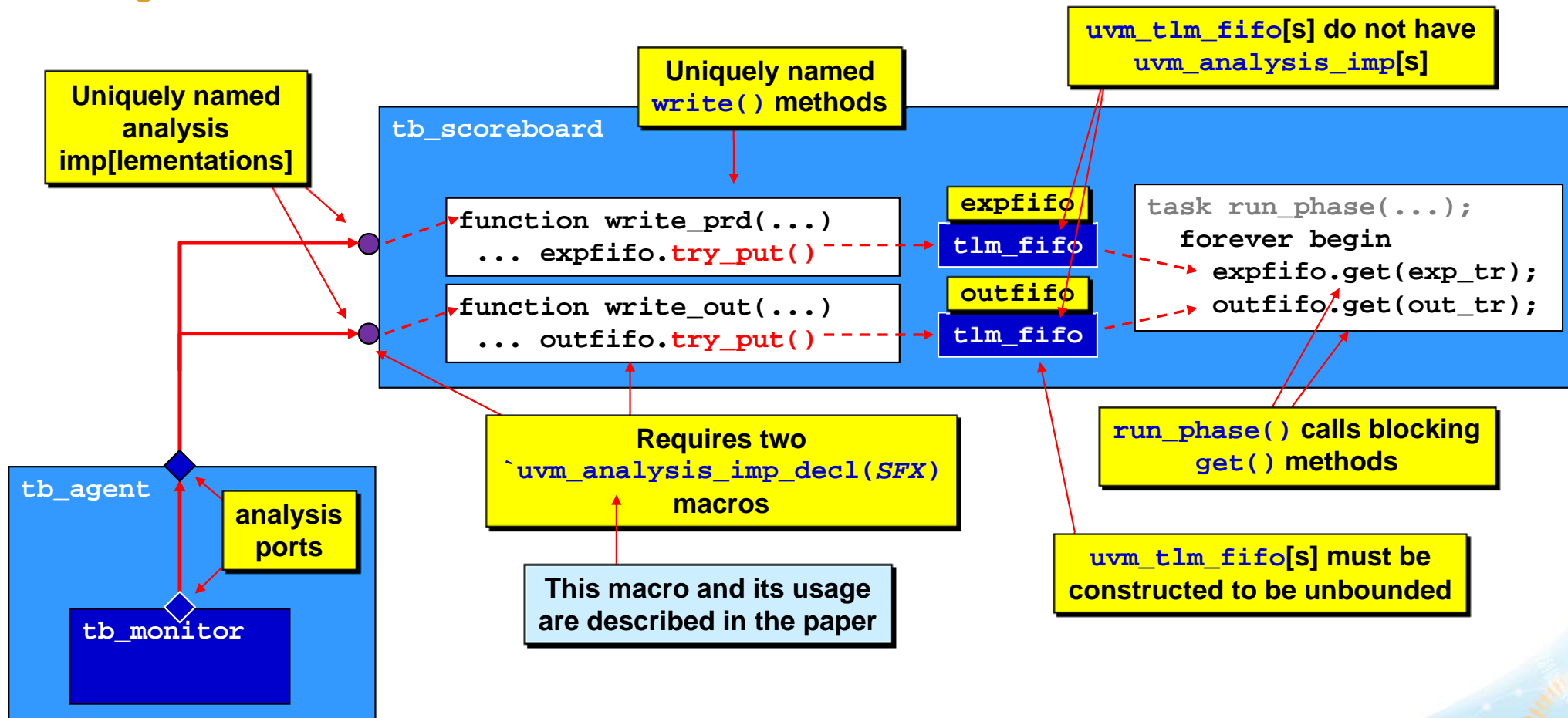
`imp` already has `write(t)` method built-in

Internally executes:
`void'(this.try_put(t));`

The scoreboard comparator still calls blocking `get(tr)` method and waits to retrieve a `uvm_tlm_analysis_fifo` transaction

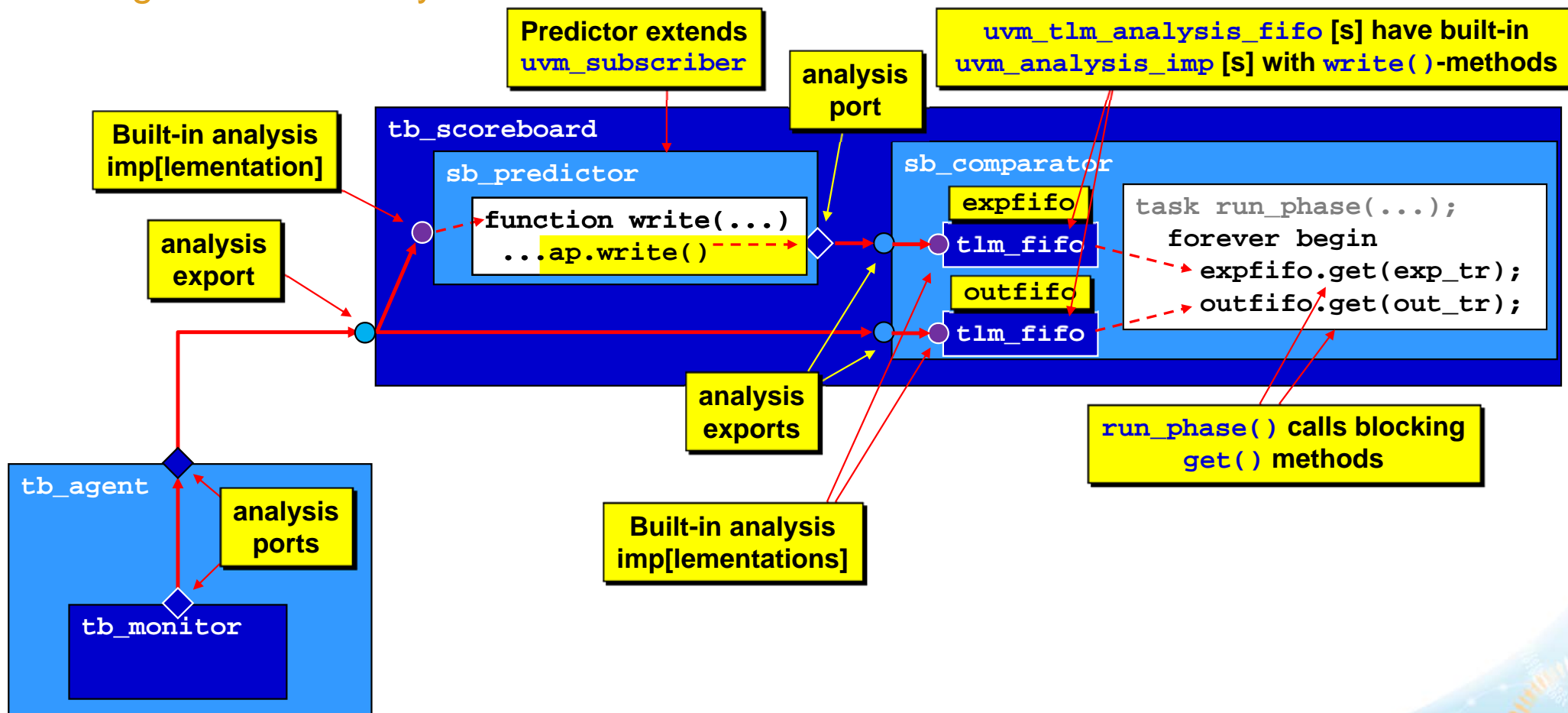
Typical Scoreboard

Using uvm_tlm_fifos



Typical Scoreboard

Using `uvm_tlm_analysis_fifos`



Creating & Copying Transactions



Predictor `write()` method creates new `etr` and copies `trans`

Pointers to `etr` #1

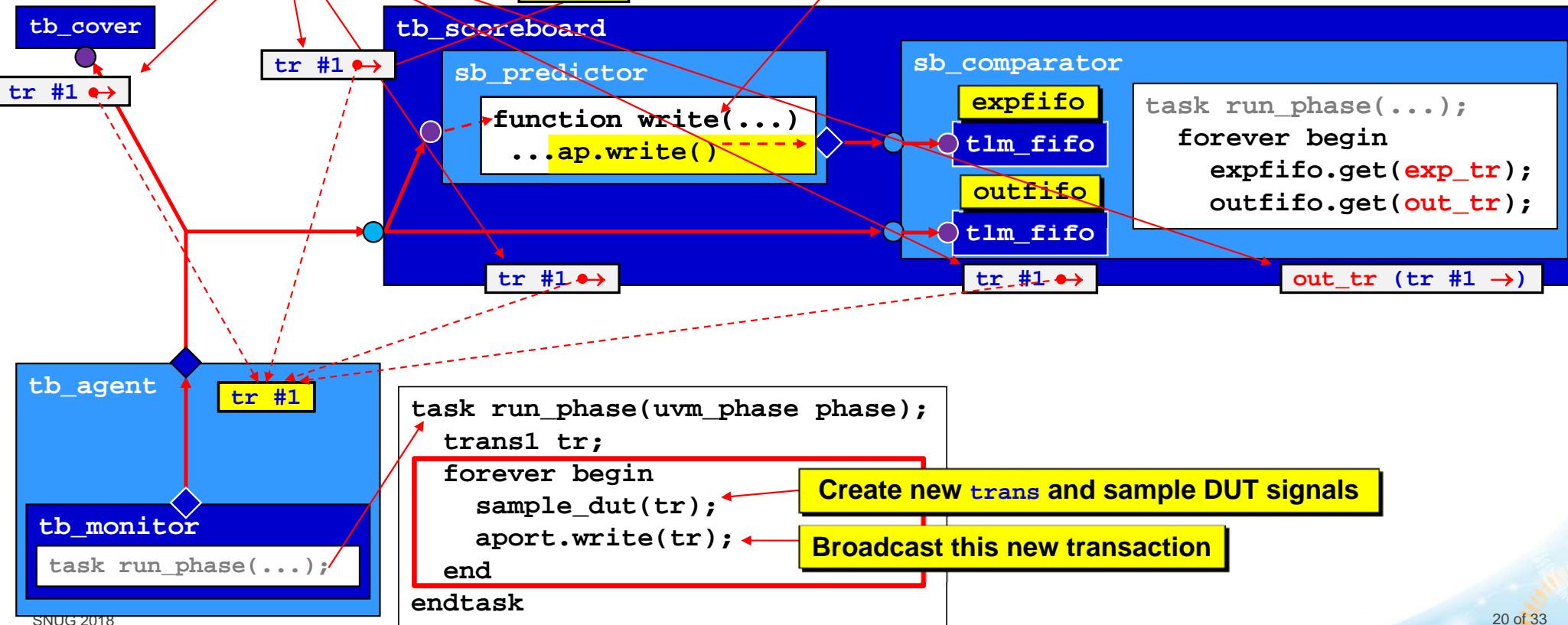
These are just pointers to `tr` #1

`copy()`

`etr` #1

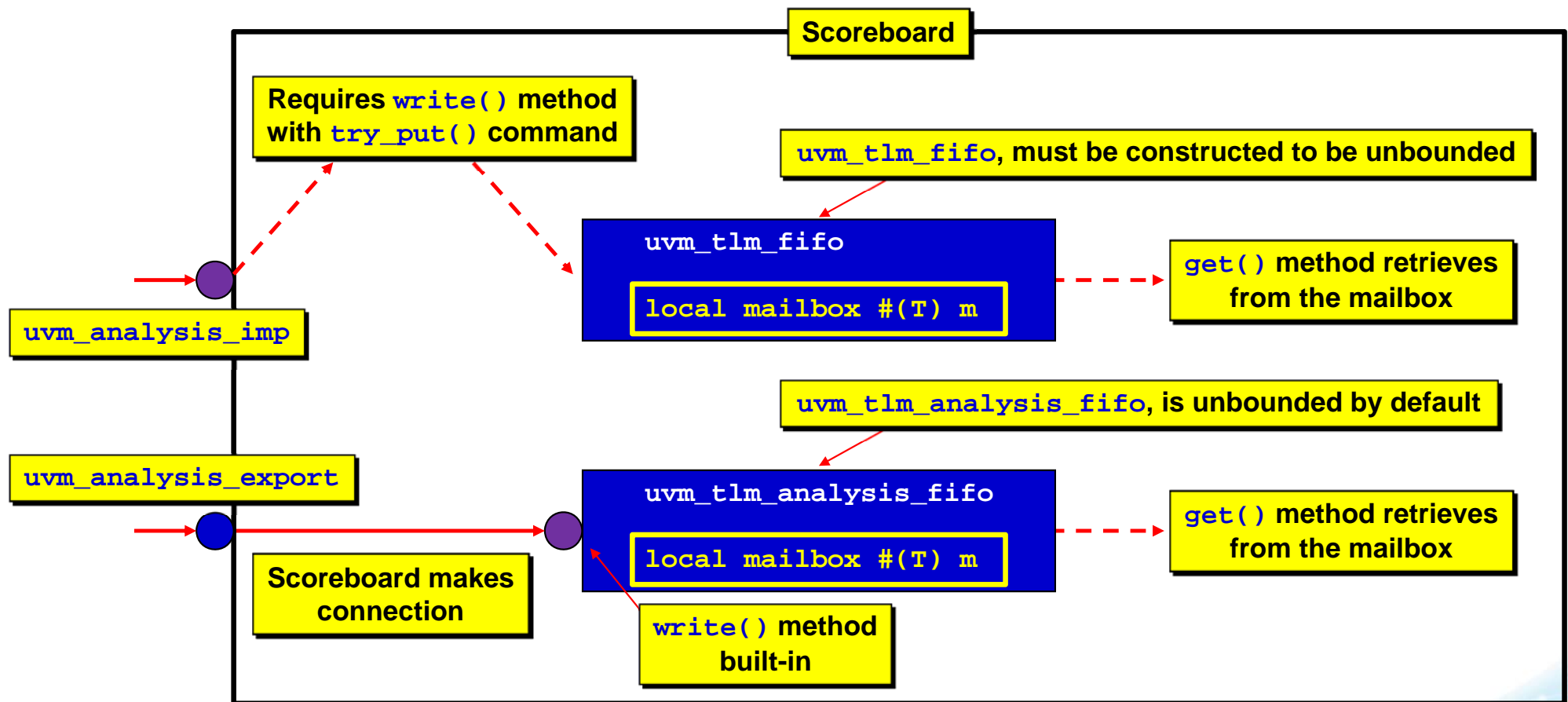
`etr` #1

`exp_tr (etr #1 →)`



Comparing TLM FIFOs

uvm_tlm_fifo -vs- uvm_tlm_analysis_fifo



Analysis Path Basics

How do analysis port-paths work?

In the software world, this is known
as the "*Observer Pattern*"



How Does UVM Work?

On previous slides

- We have learned about analysis ports & TLM FIFOs
- You do not have to know how UVM works
- The best engineers want to have *some* understanding on how UVM works
- The remaining slides show how UVM makes subscribers work

You now know how to use:

`uvm_analysis_port`
`uvm_analysis_export`
`uvm_analysis_imp`
`uvm_tlm_fifo`
`uvm_tlm_analysis_fifo`

You now have enough knowledge to use analysis components

This is ***NOT*** UVM code !!

This is a *basic* version of what UVM does internally

These slides show how UVM uses *queues* and *foreach* loops to call each subscriber's *write()* method

This is a high-level tutorial on how monitors and subscribers work

This is not exactly how UVM works, *but it is close*

Monitor with Multiple Subscribers



Goal

- Create a *Monitor* that can connect to any number of subscribers and can call a `write()` method from each subscriber *without modifying the Monitor code*

– Version #1 ← top module must know subscriber handle names in the Monitor ✗

The monitor ...

Must declare each subscriber handle ✗

Has no `connect()` method
Must copy handles by name ✗

Must call `write()` method for each subscriber ✗

– Version #2 ← Monitor w/ generic `connect()` method to hide subscriber handle names ✓

The monitor ...

Has queue of subscriber handles ✓

Defines common `connect()` method for all subscribers ✓

Uses `foreach` loop to call `write()` methods using queued subscriber handles ✓

UVM
Like!

Monitor & Subscribers

Version 1 - No connect() method

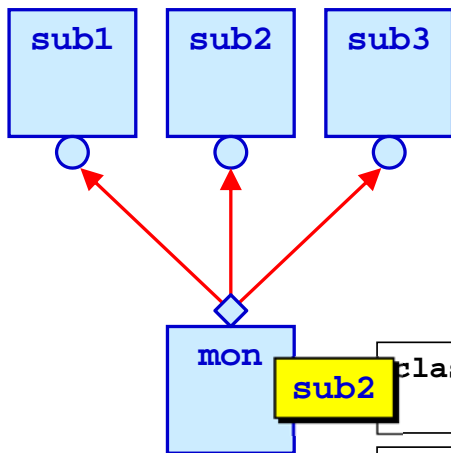
Extended classes must implement `write()` method

```
virtual class analysis_if;
  pure virtual task write(transl t);
endclass
```

```
class subscriber1 extends analysis_if;
  virtual task write(transl t);
    $display("subscriber1: ",
      "received ...", ...);
  endtask
endclass
```

sub1

top



Each subscriber handle is copied to the `ap1-3` handles in `monitor1`

```
...
mon.ap1 = sub1;
mon.ap2 = sub2;
mon.ap3 = sub3;
...
```

In top module

```
class subscriber2 extends analysis_if;
  ... virtual task write(...) ...
```

```
class subscriber3 extends analysis_if;
  ... virtual task write(...) ...
```

sub2

sub3

virtual
analysis_if
base class

```
class monitor1;
```

```
  analysis_if ap1;
  analysis_if ap2;
  analysis_if ap3;
```

subscriber1 sub1 to ap1

subscriber2 sub2 to ap2

subscriber3 sub3 to ap3

```
task run();
  transl t = new();
  repeat(5) begin
    void'(t.randomize());
    $display("monitor:  ",
      "***BROADCAST** ...", ...);
    ap1.write(t);
    ap2.write(t);
    ap3.write(t);
  end
endtask
endclass
```

Monitor & Subscribers

Version 1 - No connect() method

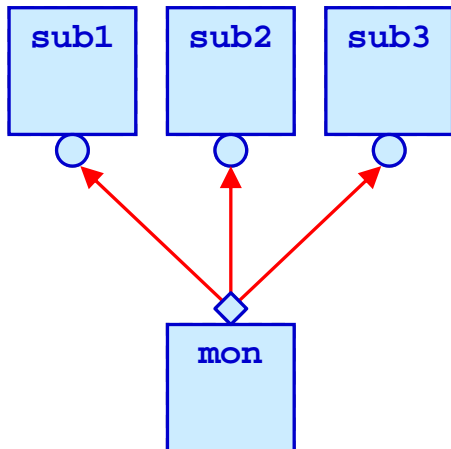
Declare `monitor1` and `subscriber1-3` handles

`new()` - construct `mon` and `sub1-3`

Copy `sub1-3` handles to `ap1-3` handles in `monitor1`

Call the `mon.run()` task

`top`



```
module top;
import tb_pkg::*;

monitor1    mon;
subscriber1 sub1;
subscriber2 sub2;
subscriber3 sub3;

initial begin
    mon = new();
    sub1 = new();
    sub2 = new();
    sub3 = new();
    mon.ap1 = sub1;
    mon.ap2 = sub2;
    mon.ap3 = sub3;
    mon.run();
end
endmodule
```

```
class monitor1;
analysis_if ap1;
analysis_if ap2;
analysis_if ap3;
```

Monitor must declare each `analysis_if`

With no `connect()` method in `monitor1`, the `top` module must reference names declared in `monitor1`

```
task run();
    trans1 t = new();
    repeat(5) begin
        void'(t.randomize());
        $display("monitor:  ", ...);
        "***BROADCAST** ...", ...);
        ap1.write(t);
        ap2.write(t);
        ap3.write(t);
```

Repeat 5 times

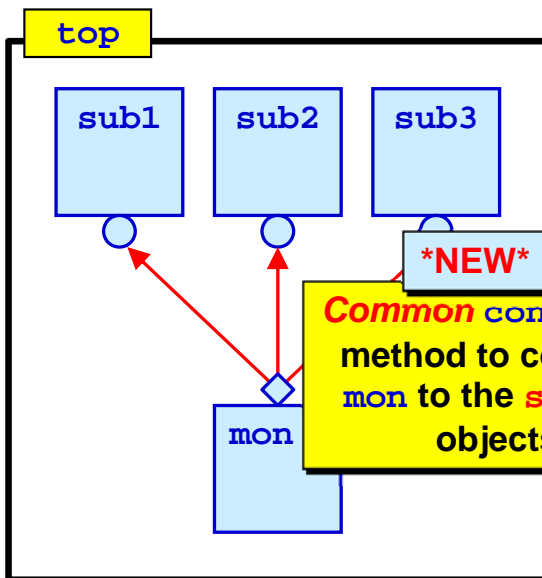
`randomize()` transaction

Separately call each `ap[#].write()` method

```
    end
endtask
endclass
```

Monitor & Subscribers

Version 2 - Adds analysis_if queue



NEW
Common connect() method to connect mon to the sub1-3 objects

No change from Version 1

```

module top;
    import tb_pkg::*;

    monitor2    mon;
    subscriber1 sub1;
    subscriber2 sub2;
    subscriber3 sub3;

    initial begin
        mon = new();
        sub1 = new();
        sub2 = new();
        sub3 = new();

        mon.connect(sub1);
        mon.connect(sub2);
        mon.connect(sub3);
        mon.run();
    end
endmodule
    
```

```

class monitor2;
    analysis_if ap[$];
    
```

NEW
Monitor declares queue of analysis_if ports

NEW

Each call to connect() method will push_back another analysis_if onto the ap-queue

```

function void connect (analysis_if port);
    ap.push_back(port);
endfunction
    
```

Common connect() method

```

task run();
    trans1 t = new();
    repeat(5) begin
        void'(t.randomize());
        $display("monitor: ",
            "BROADCAST** ...", ...);
    end
endtask
endclass
    
```

NEW

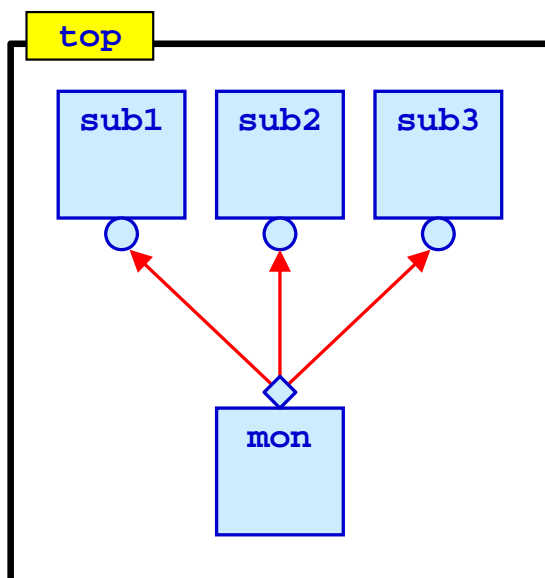
```
foreach(ap[i]) ap[i].write(t);
```

Each subscriber's write() method is called from the ap-queue

More subscribers could be added to top module without modifying monitor2 code

Monitor & Subscribers

Simulation Output



```
Randomized trans1 values addr=f9 data=50
monitor:  **BROADCAST** addr=f9 data=50
subscriber1: received    addr=f9 data=50
subscriber2: received    addr=f9 data=50
subscriber3: received    addr=f9 data=50
```

```
Randomized trans1 values addr=e9 data=27
monitor:  **BROADCAST** addr=e9 data=27
subscriber1: received    addr=e9 data=27
subscriber2: received    addr=e9 data=27
subscriber3: received    addr=e9 data=27
```

...

Subscriber2 BUG

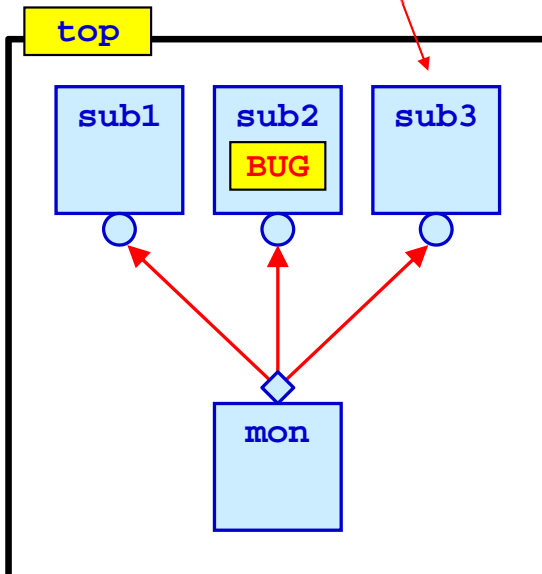
Version 4 - modifies transaction values



subscriber1 has the original transaction
addr & data values

```
class subscriber1 extends analysis_if;
  virtual task write(trans1 t);
    $display("subscriber1: ", "received addr=%2h  data=%2h", t.addr, t.data);
  endtask
endclass
```

sub3 sees corrupted
transaction



```
class subscriber2 extends analysis_if;
  virtual task write(trans1 t);
    $display("subscriber2: ", "received addr=%2h  data=%2h", t.addr, t.data);
```

BUG: subscriber2 modifies the addr &
data of the broadcast transaction

```
`ifdef BUG
  t.addr = 8'hFF;
  t.data = 8'h00;
  $display("subscriber2: ", "set      addr=%2h  data=%2h", t.addr, t.data);
`endif
```

NEVER modify the broadcast transaction !!

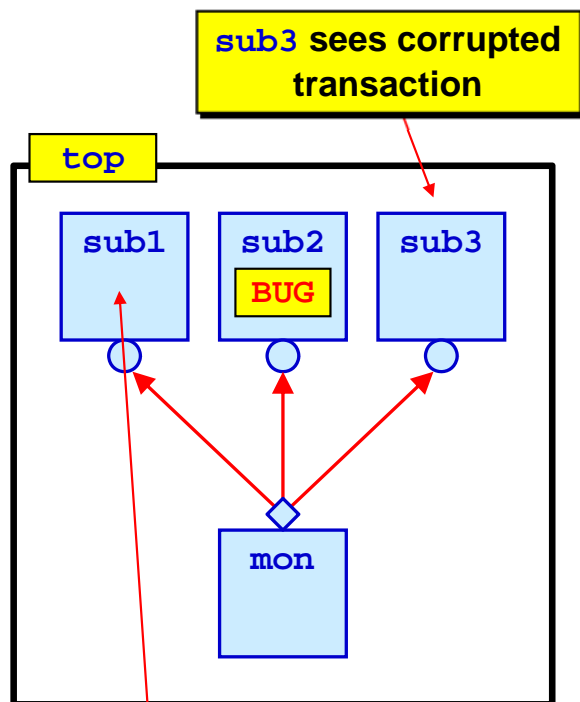
```
endtask
endclass
```

subscriber3 now sees the modified
transaction addr & data values

```
class subscriber3 extends analysis_if;
  virtual task write(trans1 t);
    $display("subscriber3: ", "received addr=%2h  data=%2h", t.addr, t.data);
  endtask
endclass
```

Monitor & Subscribers

BUG: Simulation Output



```
Randomized trans1 values addr=f9 data=50
monitor: **BROADCAST** addr=f9 data=50
subscriber1: received addr=f9 data=50
subscriber2: received addr=f9 data=50
subscriber2: set addr=ff data=00
subscriber3: received addr=ff data=00
```

```
Randomized trans1 values addr=e9 data=27
monitor: **BROADCAST** addr=e9 data=27
subscriber1: received addr=e9 data=27
subscriber2: received addr=e9 data=27
subscriber2: set addr=ff data=00
subscriber3: received addr=ff data=00
```

...

Depending on how the subscribers are pushed onto the `ap` - queue, `sub1` might also see the bug

Transaction Copy() Method



- All subscribers receive a handle to the *same* broadcast transaction
- A subscriber should **NEVER** modify contents of the received transaction
- Any subscriber that modifies transaction contents ***MUST take a copy before making modifications***

Summary & Conclusions



- Analysis ports are ports that broadcast transactions to 0 or more destinations
- Each subscriber chain terminates with a `uvm_analysis_imp` and corresponding `write()` method
- Subscribers should **NEVER** modify the broadcast transaction
- Subscribers need to use the transaction in 0-time

-OR-

- Subscribers need to take a local copy
- If a component has multiple `imp`-inputs, use the macro:

```
`uvm_analysis_imp_decl(SFX)
```

This is described in the paper

- The `uvm_tlm_analysis_fifo` has a built-in `uvm_analysis_imp` port
- Prove that the scoreboard analysis paths are working

Great feature for terminating an analysis path in a scoreboard

DO NOT ASSUME that the analysis paths are working correctly !!

Acknowledgements

Thanks!



- We are grateful to our colleagues Jeff Vance, Kelly Larson and Don Mills for their reviews and suggested improvements to our paper
- Thanks also to David Lee, Don Mills, Jeff Vance, Kelly Larson and Dan Chaplin for helping to improve the presentation content and flow

*If you liked the presentation,
these colleagues deserve recognition*

*If you **DID NOT LIKE** the presentation,
these colleagues **deserve all the blame !!***





Thank You





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**Life is too short for bad
or boring training!**

