

Distance Recoloring

*in collaboration with
Niranka Banerjee (Kyoto University, Japan)
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1 Introduction to Reconfiguration

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3 Distance Recoloring

4 Concluding Remarks

Introduction to Reconfiguration

Reconfiguration Setting

- A description of what *states* (\equiv *configurations*) are
- One or more *allowed moves* between states (\equiv *reconfiguration rule(s)*)

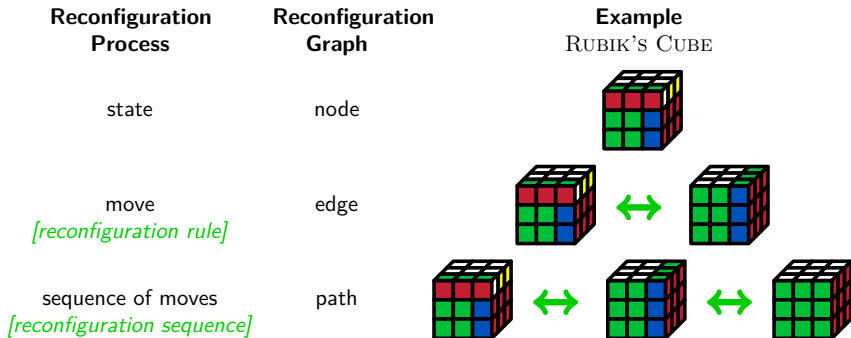


Figure: Reconfiguration Setting

Introduction to Reconfiguration

Reconfiguration vs. Solution Space

For a computational problem \mathcal{P} (e.g., INDEPENDENT SET, DOMINATING SET, VERTEX-COLORING, etc.)

Reconfiguration	Solution Space
States/Configurations	Feasible solutions of \mathcal{P}
Allowed Moves	Slight modifications of a solution <i>without changing its feasibility</i>

Algorithmic Questions

- **REACHABILITY:** Given two states S and T , is there a sequence of moves that *transforms S into T* ?
- **SHORTEST TRANSFORMATION:** Given two states S and T and a positive integer ℓ , is there a sequence of moves that *transforms S into T using at most ℓ moves*?
- **CONNECTIVITY:** Is there a sequence of moves *between any pair of states*?
- and so on

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- See [this page](#).
- <http://www.ecei.tohoku.ac.jp/alg/coresurvey/>

- <https://lists.uwaterloo.ca/mailman/listinfo/reconf>

- <http://www.ecei.tohoku.ac.jp/alg/core/>.

- See [this page](#).
- [YouTube「組合せ遷移」channel](#) (in Japanese, available since October 2021).

- See [this page](#).

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Figure: Online Reconfiguration Wiki Page (<https://reconf.wikidot.com/>)

Distance Recoloring — Duc A. Hoang



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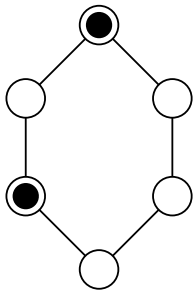
Some Examples

TOKEN RECONFIGURATION

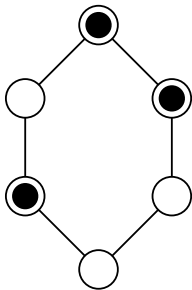
Each vertex has at most one token



Each state is a set of tokens satisfying certain property



Independent Set



Dominating Set

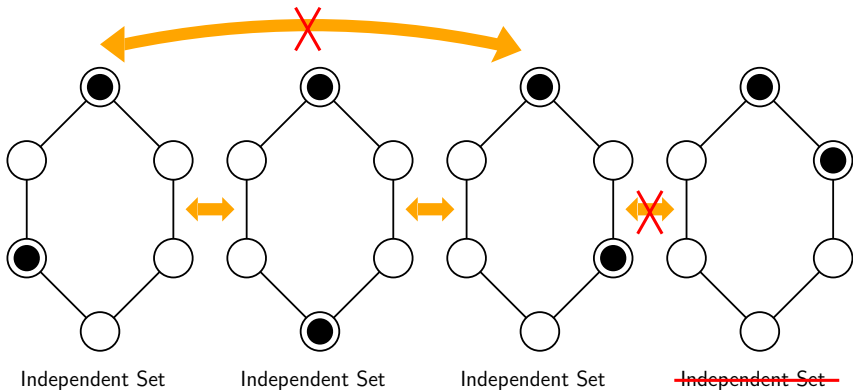
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Some Examples

[Hearn and Demaine 2005]

TOKEN RECONFIGURATION (Example: INDEPENDENT SET RECONFIGURATION)

Token Sliding (TS)



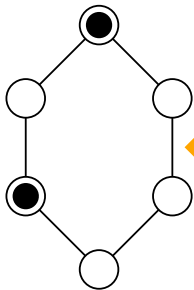
Some Examples

[Hearn and Demaine 2005]

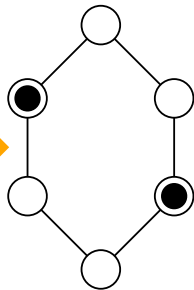
TOKEN RECONFIGURATION (Example: INDEPENDENT SET RECONFIGURATION)

Token Sliding (TS)

PSPACE-complete on general graphs



Independent Set



Independent Set

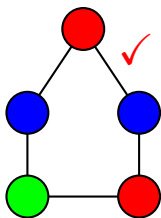
Some Examples

[Cereceda, van den Heuvel, and Johnson 2008]

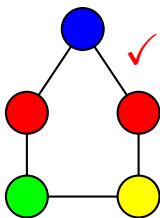
VERTEX-COLORING RECONFIGURATION

Each vertex is colored by one of the k given colors

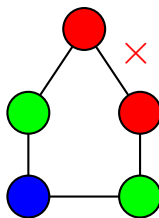
Each state is a k -coloring of all vertices such that no two adjacent vertices share the same color



3-coloring



4-coloring



~~3-coloring~~

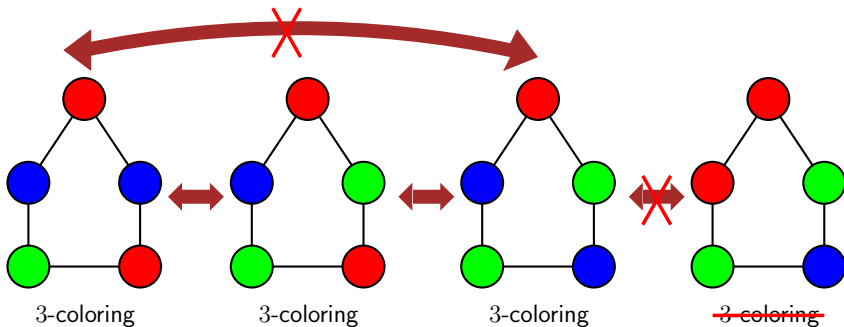
Some Examples

[Cereceda, van den Heuvel, and Johnson 2008]

VERTEX-COLORING RECONFIGURATION

Recoloring using $\leq k$ colors

Example: $k = 3$



Some Examples

[Cereceda, van den Heuvel, and Johnson 2008]

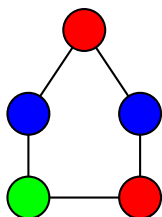
VERTEX-COLORING RECONFIGURATION

Recoloring using $\leq k$ colors

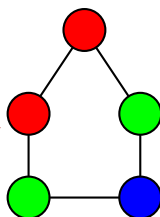
PSPACE-complete on general graphs for $k \geq 4$

Example: $k = 3$

P on general graphs for $k \leq 3$



3-coloring



3-coloring

Some Examples

[Hearn and Demaine 2005]

NONDETERMINISTIC CONSTRAINT LOGIC (NCL)

Each state is a weighted, oriented graph

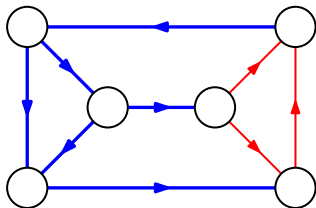


weight 2



weight 1

Total incoming weight at each vertex ≥ 2



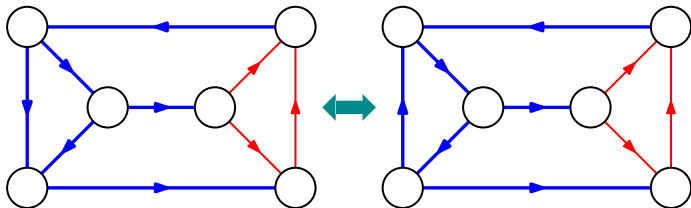
Some Examples

[Hearn and Demaine 2005]

NONDETERMINISTIC CONSTRAINT LOGIC (NCL)

Reversing edge direction

Total incoming weight at each vertex ≥ 2



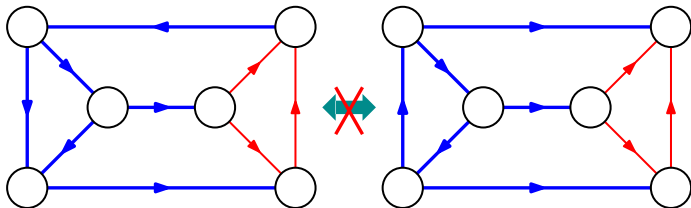
Some Examples

[Hearn and Demaine 2005]

NONDETERMINISTIC CONSTRAINT LOGIC (NCL)

Reversing edge direction

Total incoming weight at each vertex ≥ 2



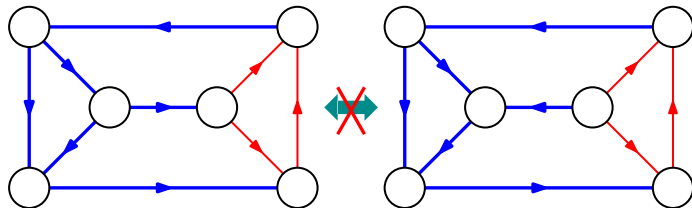
Some Examples

[Hearn and Demaine 2005]

NONDETERMINISTIC CONSTRAINT LOGIC (NCL)

Reversing edge direction

Total incoming weight at each vertex ≥ 2

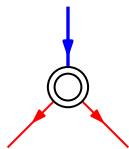


Some Examples

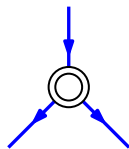
[van der Zanden 2015]

NONDETERMINISTIC CONSTRAINT LOGIC (NCL)

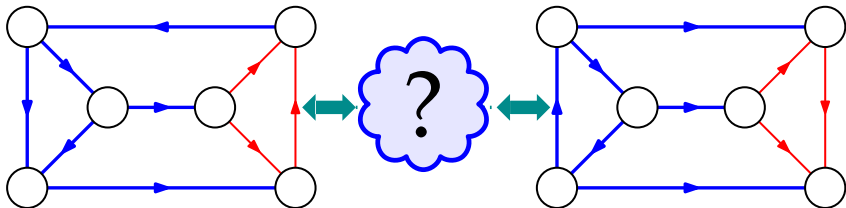
PSPACE-complete on
planar graphs
having only two types of vertices
max degree 3
bounded bandwidth



AND vertex



OR vertex

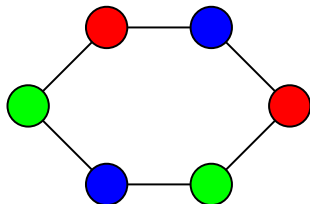


Distance Recoloring

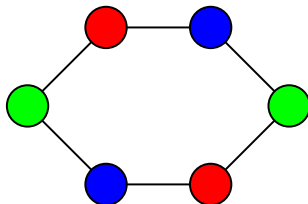
Distance constraints for *vertex colorings*

	k -Coloring	(d, k) -Coloring
Distance between two vertices having the same color	≥ 2	$\geq d + 1$

- (d, k) -coloring was first studied in [F. Kramer and H. Kramer 1969]
- Has applications in *frequency assignment problem* (or radio channel assignment) [F. Kramer and H. Kramer 2008]



3-coloring
~~(2, 3)-coloring~~

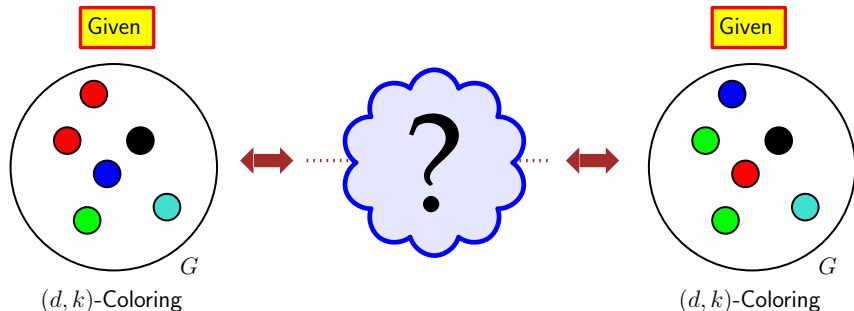


3-coloring
(2, 3)-coloring

Distance Recoloring

(d, k) -COLORING RECONFIGURATION $((d, k)$ -CR)

Reconfiguration Rule: Recoloring a vertex



- The case $d = 1$ (k -COLORING RECONFIGURATION (k -CR)) has been well-studied [Mynhardt and Nasserar 2019]; [Heuvel 2013]
- We focus on the case $d \geq 2$

Distance Recoloring

Graph	k -CR ($= (1, k)$ -CR)	(d, k) -CR ($d \geq 2$)
general	PSPACE-C ($k \geq 3$) [Cereceda, Heuvel, and Johnson 2011]	PSPACE-C ($k = \Omega(d^2)$)
planar	PSPACE-C ($4 \leq k \leq 6$) P ($k \geq 7$) [Bonsma and Cereceda 2009]	
bipartite	PSPACE-C ($k \geq 4$) [Bonsma and Cereceda 2009]	
planar \cap bipartite	PSPACE-C ($k = 4$) P ($k \geq 5$) [Bonsma and Cereceda 2009]	
2-degenerate	P [Hatanaka, Ito, and Zhou 2019]	
planar \cap bipartite \cap 2-degenerate	P (\subseteq 2-degenerate)	
path	P (\subseteq planar \cap bipartite \cap 2-degenerate)	P ($k \geq d + 1$)
split	P [Hatanaka, Ito, and Zhou 2019]	PSPACE-C ($d = 2$, large k)
		P ($d \geq 3$)

Table: Our Results for $d \geq 2$. We provide the status for $d = 1$ for comparison. Here PSPACE-C stands for PSPACE-complete [Banerjee, Engels, and **Hoang** 2024]

Distance Recoloring

Theorem (Banerjee, Engels, and Hoang 2024)

(d, k) -CR is PSPACE-complete for $d \geq 2$ and $k = \Omega(d^2)$ on *graphs which are planar, bipartite, and 2-degenerate*.

Proof Sketch.

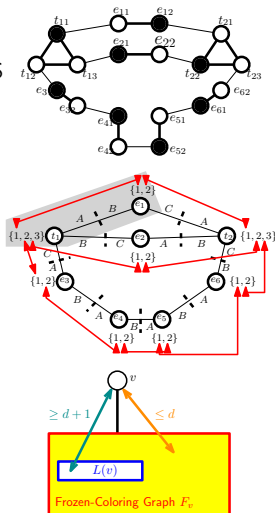
Restricted PSPACE-Complete Variant of ISR under TS
(RESTRICTED SLIDING TOKENS)



LIST $(d, \Omega(d))$ -COLORING RECONFIGURATION



$(d, \Omega(d^2))$ -COLORING RECONFIGURATION



Distance Recoloring

First Phase

RESTRICTED SLIDING TOKENS \Rightarrow LIST $(d, \Omega(d))$ -COLORING RECONFIGURATION

RESTRICTED SLIDING TOKENS

PSPACE-complete on **very restricted instances**

Three types of gadgets

Token triangle

Token edge

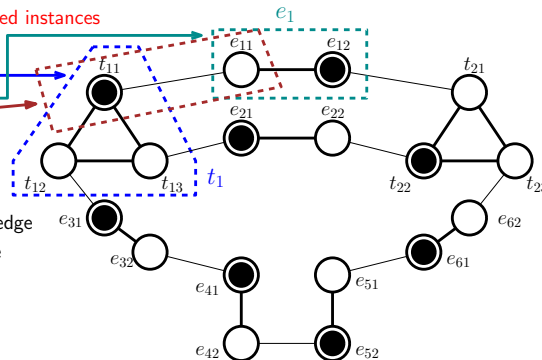
Link edge

Planar

Max degree 3, min degree 2

No two token triangles
are directly joined by a link edge

Each token triangle/token edge
has exactly one token



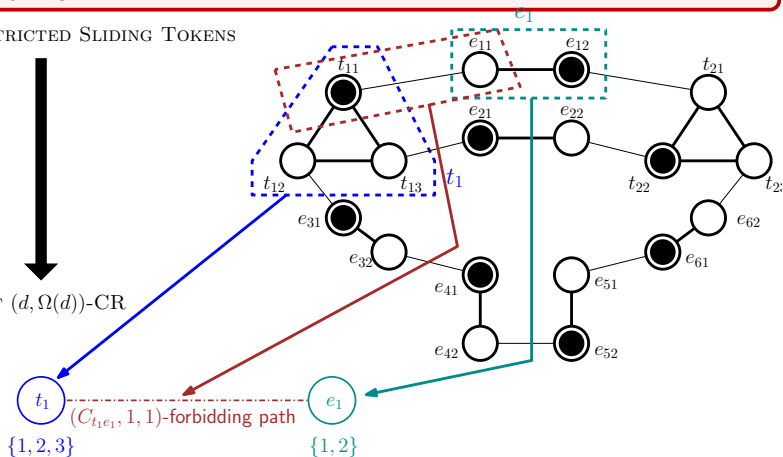
Distance Recoloring

First Phase

RESTRICTED SLIDING TOKENS \Rightarrow LIST $(d, \Omega(d))$ -COLORING RECONFIGURATION

RESTRICTED SLIDING TOKENS

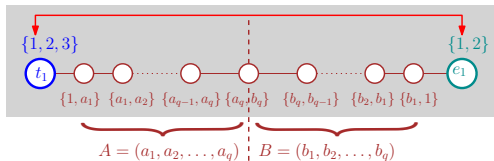
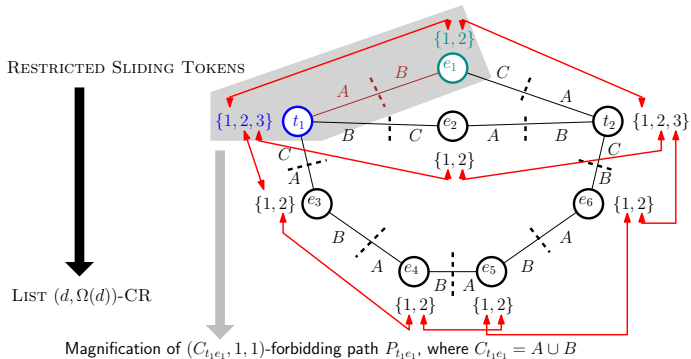
LIST $(d, \Omega(d))$ -CR



Distance Recoloring

First Phase

RESTRICTED SLIDING TOKENS \Rightarrow LIST $(d, \Omega(d))$ -COLORING RECONFIGURATION



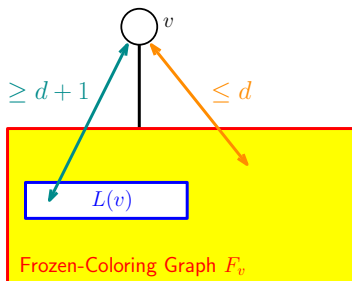
Distance Recoloring

Second Phase

LIST $(d, \Omega(d))$ -COLORING RECONFIGURATION $\Rightarrow (d, \Omega(d^2))$ -COLORING RECONFIGURATION

Key Ideas

- 1 List Coloring \equiv Coloring with constraints on which colors can be used for each vertex
- 2 *Frozen-Coloring Graphs*: Pre-colored graphs where *no vertex can be recolored*



Vertices are pre-colored

Containing all possible colors

No vertex can be recolored

Distance Recoloring

Second Phase

LIST $(d, \Omega(d))$ -COLORING RECONFIGURATION $\Rightarrow (d, \Omega(d^2))$ -COLORING RECONFIGURATION

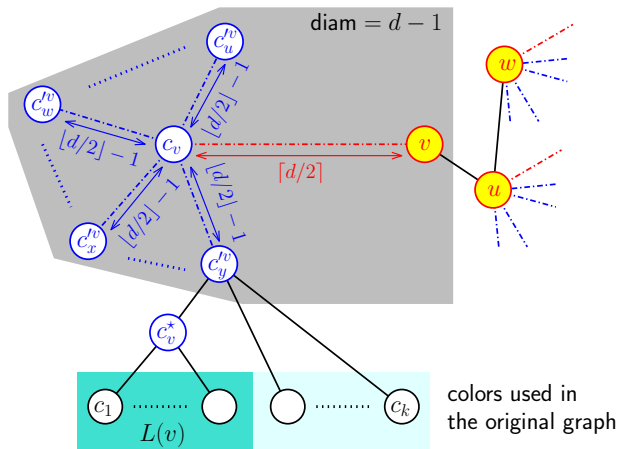


Figure: Construction of the frozen-coloring subgraph F_v for a vertex v . The colors used in this figure are just for illustration of paths.

Theorem (Banerjee, Engels, and Hoang 2024)

On *split graphs*,

	(d, k) -Coloring Reconfiguration
$d = 1$	P for any k [Hatanaka, Ito, and Zhou 2019]
$d = 2$	PSPACE-complete for large k
$d \geq 3$	P for any k

Proof Sketch.

- $d \geq 3$ Trivial. (Any connected split graph has diameter $\leq 3 \Rightarrow$ Reconfiguration is easy!)
- $d = 2$ Reduction from the problem for $d = 1$ and $k \geq 4$ on general graphs (which is known to be PSPACE-complete [Bonsma and Cereceda 2009])



Concluding Remarks

Take-Home Messages

- 1 *Reconfiguration* studies the “solution space” of a problem
 - » Moving from one solution to another *without violating feasibility*
- 2 Under certain *distance constraints*,
 - » *Reconfiguration problems* can be *hard* for very restricted graph classes
 - » *Problems on graphs whose diameters are bounded by some constant c* (e.g., split graphs) are interesting when *restricted to distances close to c*
- 3 *Nondeterministic Constraint Logic* is a powerful tool for *hardness reductions*

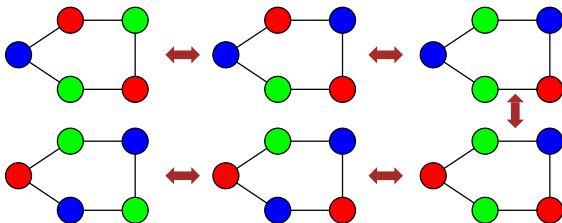
Open Problems

The complexities of the following problem remain *open* for *trees*:

- » (d, k) -CR ($d \geq 2$)



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**Came to my talk, you did.
Thank you, I must!**



Banerjee, Niranka, Christian Engels, and **Duc A. Hoang** (2024). “Distance Recoloring”. In: *arXiv preprint*. arXiv: 2402.12705.



Hatanaka, Tatsuhiko, Takehiro Ito, and Xiao Zhou (2019). “The Coloring Reconfiguration Problem on Specific Graph Classes”. In: *IEICE Transactions on Information and Systems* E102.D.3, pp. 423–429. DOI: 10.1587/transinf.2018FCP0005.



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Bonsma, Paul S. and Luis Cereceda (2009). “Finding Paths Between Graph Colourings: PSPACE-Completeness and Superpolynomial Distances”. In: *Theoretical Computer Science* 410.50, pp. 5215–5226. DOI: 10.1016/j.tcs.2009.08.023.



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Kramer, Florica and Horst Kramer (2008). “A Survey on the Distance-Colouring of Graphs”. In: *Discrete mathematics* 308.2-3, pp. 422–426. DOI: 10.1016/j.disc.2006.11.059.



Hearn, Robert A. and Erik D. Demaine (2005).

“PSPACE-completeness of sliding-block puzzles and other problems through the nondeterministic constraint logic model of computation”.

In: *Theoretical Computer Science* 343.1-2, pp. 72–96. DOI: 10.1016/j.tcs.2005.05.008.



Kramer, Florica and Horst Kramer (1969). “Ein Färbungsproblem der Knotenpunkte eines Graphen bezüglich der Distanz p ”. In: *Rev. Roumaine Math. Pures Appl* 14.2, pp. 1031–1038.