# Loops

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What's next...



### LC3 overview

Instruction Set

| Op   | Format   | Description  | Example  |
|------|--|--|--|
| ADD  | ADD DR, SR1, SR2<br>ADD DR, SR1, imm5  | Adds the values in SR1 and SR2/imm5 and sets DR to that value.   | ADD R1, R2, #5<br>The value 5 is added to the value in<br>R2 and stored in R1.   |
| AND  | AND DR, SR1, SR2<br>AND DR, SR1, imm5  | Performs a bitwise and on the values in SR1 and SR2/imm5 and sets DR to the result.  | AND R0, R1, R2 A bitwise and is preformed on the values in R1 and R2 and the result stored in R0.  |
| BR   | BR(n/z/p) LABEL<br>Note: (n/z/p) means<br>any combination of<br>those letters can appear<br>there, but must be in<br>that order. | Branch to the code section indicated by LABEL, if the bit indicated by (n/z/p) has been set by a previous instruction. n: negative bit, z: zero bit, p: positive bit. Note that some instructions do not set condition codes bits. | BRZ LPBODY Branch to LPBODY if the last instruction that modified the condition codes resulted in zero. BRnp ALT1 Branch to ALT1 if last instruction that modified the condition codes resulted in a positive or negative (non-zero) number. |
| JMP  | JMP SR1  | Unconditionally jump to the instruction based upon the address in SR1.   | JMP R1<br>Jump to the code indicated by the<br>address in R1.  |
| JSR  | JSR LABEL  | Put the address of the next<br>instruction after the JSR<br>instruction into R7 and jump to<br>the subroutine indicated by<br>LABEL.   | JSR POP Store the address of the next instruction into R7 and jump to the subroutine POP.  |
| JSRR | JSSR SR1   | Similar to JSR except the address stored in SR1 is used instead of using a LABEL.  | JSSR R3 Store the address of the next instruction into R7 and jump to the subroutine indicated by R3's value.  |

Covered already

Covered today

Covered later

Not covered

| LD  | LD DR, LABEL         | Load the value indicated by LABEL into the DR register.                                      | LD R2, VAR1<br>Load the value at VAR1 into R2.   |
|-----|----------------------|--|--|
| LDI | LDI DR, LABEL        | Load the value indicated by the address at LABEL's memory location into the DR register.     | LDI R3, ADDR1 Suppose ADDR1 points to a memory location with the value x3100. Suppose also that memory location x3100 has the value 8. 8 then would be loaded into R3. |
| LDR | LDR DR, SR1, offset6 | Load the value from the memory location found by adding the value of SR1 to offset6 into DR. | LDR R3, R4, #-2<br>Load the value found at the address<br>(R4 -2) into R3.   |
| LEA | LEA DR, LABEL        | Load the address of LABEL into DR.   | LEA R1, DATA1<br>Load the address of DATA1 into<br>R1.   |
| NOT | NOT DR, SR1          | Performs a bitwise not on SR1 and stores the result in DR.                                   | NOT R0, R1 A bitwise not is preformed on R1 and the result is stored in R0.  |
| RET | RET                  | Return from a subroutine using the value in R7 as the base address.                          | RET Equivalent to JMP R7.  |

| RTI                   |                                      | RTI  |
|-----------------------|--------------------------------------|--|
|                       | 1                                    |  |
|                       | address to return to is obtained     | Note: RTI can only be used if the  |
|                       | by popping it off the supervisor     | processor is in supervisor mode.   |
|                       | stack, which is automatically        |  |
|                       | done by RTI.                         |  |
| ST SR1, LABEL         | Store the value in SR1 into the      | ST R1, VAR3  |
|                       | memory location indicated by         | Store R1's value into the memory   |
|                       | LABEL.                               | location of VAR3.  |
| STI SR1, LABEL        | Store the value in SR1 into the      | STI R2, ADDR2  |
|                       | memory location indicated by         | Suppose ADDR2's memory   |
|                       | the value that LABEL's memory        | location contains the value x3101.   |
|                       | location contains.                   | R2's value would then be stored  |
|                       |                                      | into memory location x3101.  |
| STR SR1, SR2, offset6 | The value in SR1 is stored in the    | STR R2, R1, #4   |
|                       | memory location found by             | The value of R2 is stored in   |
|                       | adding SR2 and offest6 together.     | memory location (R1 + 4).  |
| TRAP trapvector8      | Performs the trap service            | TRAP x25   |
|                       | specified by trapvector8. Each       | Calls a trap service to end the  |
|                       | trapvector8 service has its own      | program. The assembly instruction  |
|                       | assembly instruction that can        | HALT can also be used to replace   |
|                       | replace the trap instruction.        | TRAP x25.  |
|                       | STI SR1, LABEL STR SR1, SR2, offset6 | code that was interrupted. The address to return to is obtained by popping it off the supervisor stack, which is automatically done by RTI.  ST SR1, LABEL  Store the value in SR1 into the memory location indicated by LABEL.  Store the value in SR1 into the memory location indicated by the value that LABEL's memory location contains.  STR SR1, SR2, offset6  The value in SR1 is stored in the memory location found by adding SR2 and offest6 together.  Performs the trap service specified by trapvector8. Each trapvector8 service has its own assembly instruction that can |

## What will the following program do?

```
.ORIG x3000
LD R0, MY NUM
SOME LABEL
    ADD, R0, R0, #-1
    BRp SOME_LABEL
HALT
MY NUM .FILL #5
. END
```

## Example 1

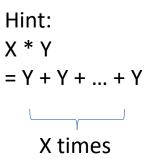
Write an LC-3 assembly program that prints out "Hello, World!" *n* times, where N is defined via .FILL.

## Example 2

Write an LC-3 assembly program that calculates the sum of all integers from 1 to N, where N is stored at address x3100.

## Multiplication

The integers X and Y are stored at locations x3100 and x3101, respectively. Write an LC-3 assembler program that calculates X \* Y and stores the result at x3102.



## Some tougher loop practice

### LAB 4

### Fibonacci Numbers

### 4.1 Problem Statement

- 1. Write a program in LC-3 assembly language that computes  $F_n$ , the n-th Fibonacci number.
- 2. Find the largest  $F_n$  such that no overflow occurs, i.e. find n = N such that  $F_N$  is the largest Fibonacci number to be correctly represented with 16 bits in two's complement format.

#### **4.1.1** Inputs

The integer n is in memory location **x3100**:

x3100 n

#### 4.1.2 Outputs

| x3101 | $F_n$ |
|-------|-------|
| x3102 | N     |
| x3103 | $F_N$ |

### 4.2 Example

| x3100 | 6     |
|-------|-------|
| x3101 | 8     |
| x3102 | N     |
| x3103 | $F_N$ |

Starting with 6 in location **x3100** means that we intend to compute  $F_6$  and place that result in location **x3101**. Indeed,  $F_6 = 8$ . (See below.) The actual values of N and  $F_N$  should be found by your program, and be placed in their corresponding locations.