# Data Integration

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# Outline

- **➤**Introduction
- Examples of data integration applications
- Schema heterogeneity
- Goal of data integration
- Data integration architectures
- Review of basic database concepts

## Data Integration

- Databases are great?
  - Assuming you've put it all into your schema.
- Data sets are often created independently
  - Only to discover later that they need to combine their data!
  - Data in different systems, different schemata and limited interfaces to data.
- Data integration: tie together different sources, controlled by many people, under a common schema.

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#### DBMS: it's all about abstraction

· Logical vs. Physical; What vs. How.

Students:

SSN	Name	Category
123-45-6789	Charles	undergrad
234-56-7890	Dan	grad

Takes:

SSN	CID
123-45-6789	CSE444
123-45-6789	CSE444
234-56-7890	CSE142

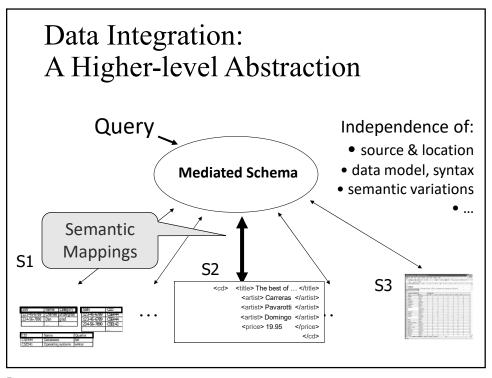
Courses:

CI	D	Name	Quarter
CS	E444	Databases	fall
CS	E541	Operating systems	winter

SELECT C.name

FROM Students S, Takes T, Courses C WHERE S.name="Mary" and

S.ssn = T.ssn and T.cid = C.cid



#### Outline

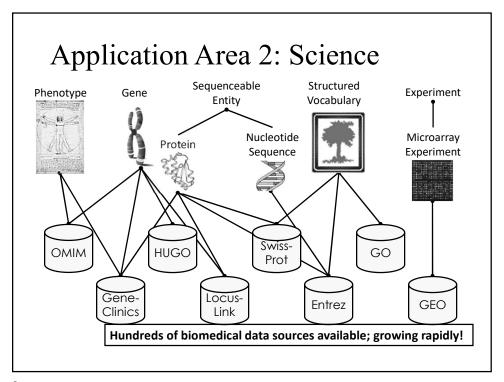
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# Applications of Data Integration

- Business
- Science
- Government
- The Web
- Pretty much everywhere

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# Application Area 1: Business Enterprise Databases EII Apps: CRM ERP Portals ... Legacy Databases Services and Applications 50% of all IT \$\$\$ spent here!

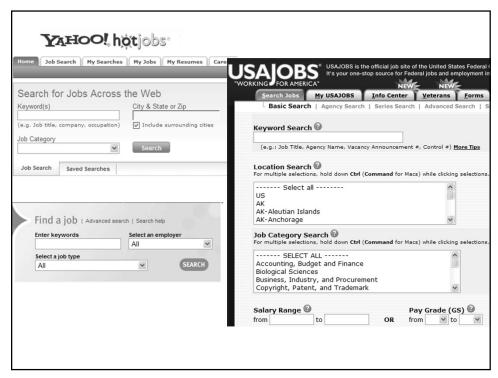




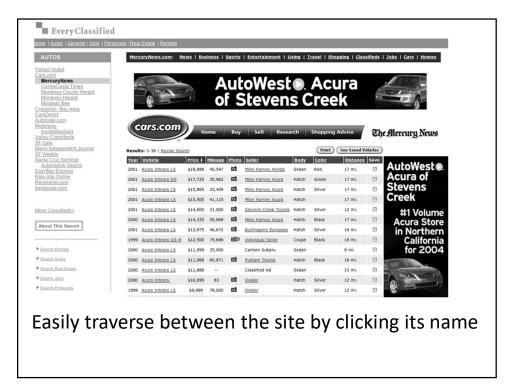


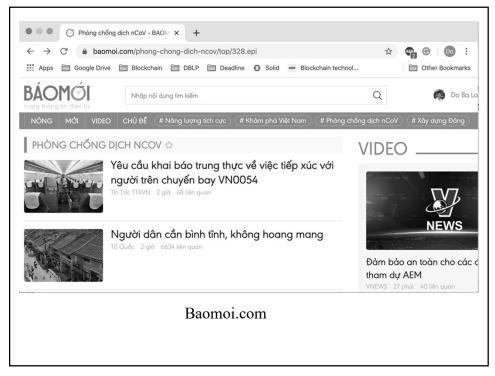
#### The Deep Web

- Millions of high quality HTML forms out there
- Each form has its own special interface
  - Hard to explore data across sites.
- Goal (for some domains):
  - A single interface into a multitude of deep-web sources.









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# **Enterprise Data Integration:**

FullServe Corporation

Employees Resumes

FullTimeEmp Interview Hire CV

TempEmployees

Training Services
Courses Services
Enrollments Customers

Contracts

Sales HelpLine Products Calls

Sales

# **EuroCard Corporation**

Employees Resumes

Employees Interview

Hire

Credit Cards HelpLine

Customer Calls

CustDetail

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# **Examples of Heterogeneity**

FullServe EuroCard

FullTimeEmp Employees

ssn, empld, firstName
ID, firstNameMiddleInitial,
middleName, lastName
lastName

Hire Hire

empld, hireDate, recruiter ID, hireDate, recruiter

**TempEmployees** 

ssn, hireStart, hireEnd

Find all employees (making over \$100K)

#### Customer Call Center

Agents should have a full view of customer when they call in.

#### Sales

Products Sales

#### **Credit Cards**

Customer CustDetail

#### **Services**

Services Customers Contracts

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## Other Reasons to Integrate Data

- Create a (useful) web site for tracking services
- Collaborate with third parties
  - E.g., create branded services
- Comply with government regulations
  - Find "risky" employees
- Business intelligence
  - What's really wrong with our products?

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# Goal of Data Integration

- Uniform query access to a set of data sources
- Handle:
  - Scale of sources: from tens to millions
  - Heterogeneity
  - Semi-structure

## Why is it Hard?

- Systems-level reasons:
  - Managing different platforms
  - SQL across multiple systems is not so simple
  - Distributed query processing
- Logical reasons:
  - Schema (and data) heterogeneity
- 'Social' reasons:
  - Locating and capturing relevant data in the enterprise.
  - Convincing people to share (data fiefdoms)
    - Security, privacy and performance implications.

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# Data Integration Smorgasbord

#### Something for everyone:

- Theory of modeling data sources
- Systems aspects of data integration
- Architectural issues: e.g., P2P data sharing
- AI @ work: automated schema matching
- Web: latest on data integration & web
- Commercial products: BEA, IBM
- Semantic Web: what does it have to offer?
- New trends in DBMS: uncertainty, dataspaces

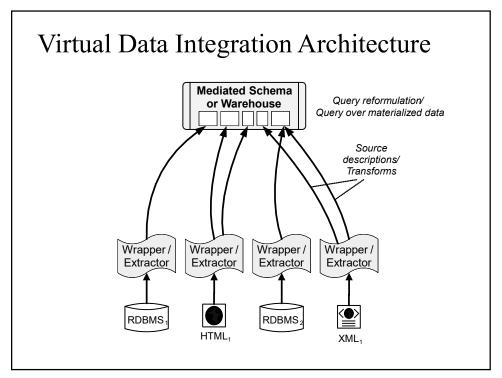
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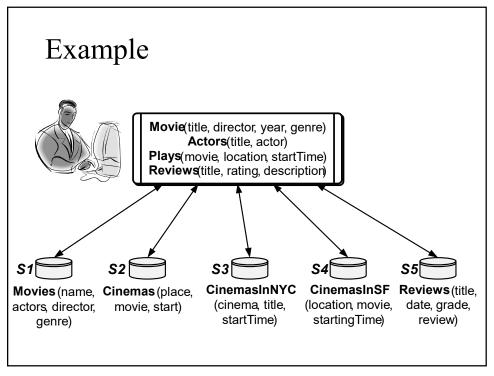
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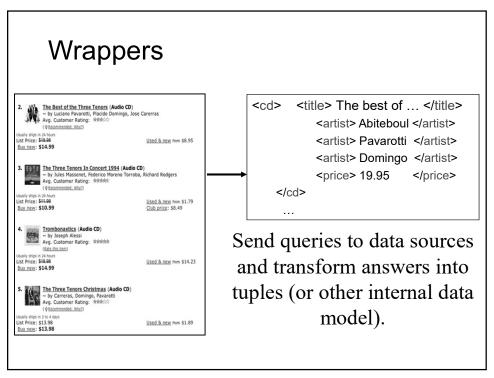
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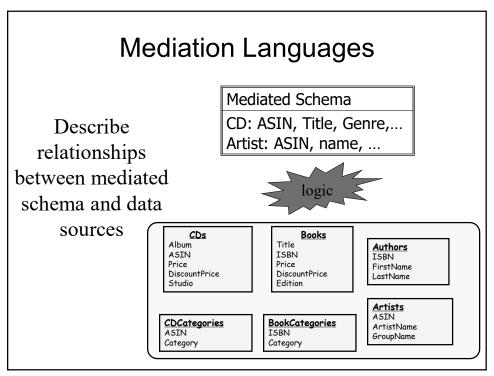
#### Virtual, Warehousing and in Between

- Virtual data integration: leave the data at the sources and access it at query time.
- Data warehousing: integrate by bringing the data into a single physical warehouse
- ❖ semantic heterogeneity
- ❖ Numerous intermediate architectures.









#### Woody Allen Comedies in NY

#### Mediated schema:

Movie: Title, director, year, genre

Actors: title, actor

**Plays**: movie, location, startTime **Reviews**: title, rating, description

select title, startTime
from Movie, Plays
where Movie.title=Plays.movie AND
location="New York" AND
director="Woody Allen"

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Movie: Title, director, year, genre

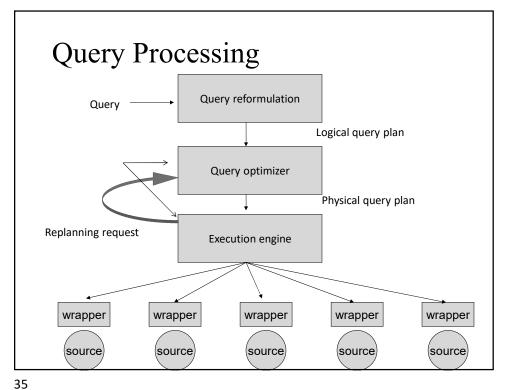
Actors: title, actor

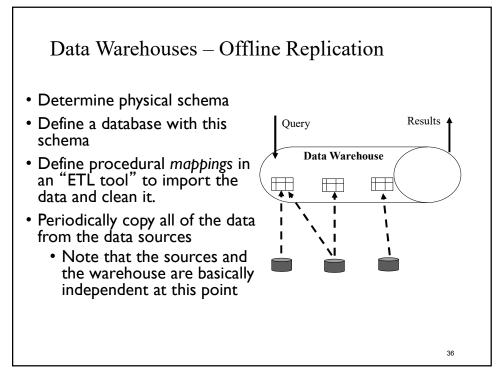
**Plays**: movie, location, startTime **Reviews**: title, rating, description

select title, startTime
from Movie, Plays
where Movie.title=Plays.movie AND
location="New York" AND
director="Woody Allen"

Sources S1 and S3 are relevant, sources S4 and S5 are irrelevant, and source S2 is relevant but possibly redundant.

S1 S2 S3 S5 Movies: Cinemas in NYC: Cinemas in SF: Cinemas: Reviews: name, actors, place, movie, cinema, title, location, movie, title, date director, genre start startTime startingTime grade, review





#### Pros and Cons of Data Warehouses

- \*Need to spend time to design the physical database layout, as well as logical
  - \* This actually takes a lot of effort!
- \*Data is generally not up-to-date (lazy or offline refresh)
- ✓ Queries over the warehouse don't disrupt the data sources
- ✓ Can run very heavy-duty computations, including data mining and cleaning

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# Basic Database Concepts

- Relational data model
- Integrity constraints
- Queries and answers
- Conjunctive queries
- Datalog

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# Relational Terminology

#### Relational schemas

• Tables, attributes

#### Relation instances

• Sets (or multi-sets) of tuples

#### Integrity constraints

• Keys, foreign keys, inclusion dependencies

Prodi	Table/relation name  Product  Attribute names					
	PName	Price	Category	Manufacturer		
	Gizmo	\$19.99	Gadgets	GizmoWorks		
Po	owergizmo	\$29.99	Gadgets	GizmoWorks		
Si	ngleTouch	\$149.99	Photography	Canon		
N	lultiTouch	\$203.99	Household	Hitachi		
Tuple	Tuples or rows					

# SQL (very basic)

#### Interview:

candidate, date, recruiter, hireDecision, grade

#### EmployeePerf:

empID, name, reviewQuarter, grade, reviewer

select recruiter, candidate from Interview, EmployeePerf where recruiter=name AND grade < 2.5

# **Query Answers**

- *Q(D)*: the set (or multi-set) of rows resulting from applying the query *Q* on the database *D*.
- Unless otherwise stated, we will consider sets rather than multi-sets.

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# SQL (w/aggregation)

#### EmployeePerf:

empID, name, reviewQuarter, grade, reviewer

select reviewer, Avg(grade) from **EmployeePerf** where reviewQuarter="1/2007"

## Integrity Constraints (Keys)

- A key is a set of columns that uniquely determine a row in the database:
  - There do not exist two tuples,  $t_1$  and  $t_2$  such that  $t_1 \neq t_2$  and  $t_1$  and  $t_2$  have the same values for the key columns.
  - (EmpID, reviewQuarter) is a key for EmployeePerf

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# Integrity Constraints (Functional Dependencies)

- A set of attribute A functionally determines a set of attributes B if: whenever, t<sub>1</sub> and t<sub>2</sub> agree on the values of A, they must also agree on the values of B.
- For example, (EmpID, reviewQuarter) functionally determine (grade).
- Note: a key dependency is a functional dependency where the key determines all the other columns.

#### **Integrity Constraints (Foreign Keys)**

- Given table **T** with key *B* and table **S** with key *A*: *A* is a foreign key of *B* in **T** if whenever a **S** has a row where the value of *A* is **v**, then **T** must have a row where the value of *B* is **v**.
- Example: the empID attribute of **EmployeePerf** is a foreign key for attribute emp of **Employee**.

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### General Integrity Constraints

Tuple generating dependencies (TGD's)

$$(\forall \overline{X}) s_1(\overline{X}_1),...,s_m(\overline{X}_m) \rightarrow (\exists \overline{Y}) \ t_1(\overline{Y}_1),...,t_1(\overline{Y}_1)$$

Equality generating dependencies (EGD's): right hand side contains only equalities.

$$(\forall \overline{X})s_1(\overline{X}_1), ..., s_m(\overline{X}_m) \rightarrow Y_1^1 = Y_2^1, ..., Y_1^k = Y_2^k$$

Exercise: express the previous constraints using general integrity constraints.

# Conjunctive Queries

#### Q(X,T):

Interview(X,D,Y,H,F), EmployeePerf(E,Y,T,W,Z), W < 2.5.

Joins are expressed with multiple occurrences of the same variable

select recruiter, candidate from **Interview**, **EmployeePerf** where recruiter=name AND grade < 2.5

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# Conjunctive Queries (interpreted predicates)

#### Q(X,T):

Interview(X,D,Y,H,F), EmployeePerf(E,Y,T,W,Z), **W < 2.5**.

Interpreted (or comparison) predicates. Variables must also appear in regular atoms.

select recruiter, candidate from **Interview**, **EmployeePerf** where recruiter=name AND grade < 2.5

# Conjunctive Queries (negated subgoals)

#### Q(X,T):

Interview(X,D,Y,H,F), EmployeePerf(E,Y,T,W,Z), →OfferMade(X, date).

Safety: every head variable must appear in a positive subgoal.

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# Unions of Conjunctive Queries

Multiple rules with the same head predicate express a union

#### Q(R,C):-

Interview(X,D,Y,H,F), EmployeePerf(E,Y,T,W,Z), W < 2.5.

#### Q(R,C):-

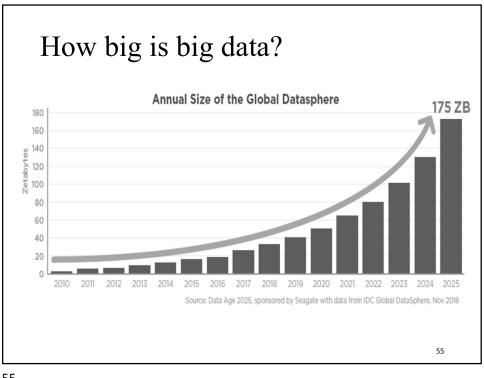
Interview(X,D,Y,H,F), EmployeePerf(E,Y,T,W,Z), Manager(y), W > 3.9.

#### Summary

- Data integration: abstract away the fact that data comes from multiple sources in varying schemata.
- Problem occurs everywhere: it's key to business, science, Web and government.
- Goal: reduce the effort involved in integrating.
- Regardless of the architecture, heterogeneity is a key issue.
- Architectures range from warehousing to virtual integration.

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Big Data .... Integration



#### Big Data Growth Statistics

- An internet user generates ~ 1.7 megabytes (MB) of data / second.
- 2022: ~ 97 zettabytes the estimated volume of data created worldwide
- 2023:
  - $\sim 2/3$  world population be online
  - internet users generate nearly 3 times the volume of data generated in 2019.
- 2025:
  - people will create more than 181 ZB of data. That's 181, followed by 21 zeros.
  - there will be 55.7 billion connected IoT devices. These IoT devices alone will generate almost 80 ZB by 2025.
- An internet user would need more than 180 million years to download all the data from the web.
- Nearly 80% of companies estimate that 50%-90% of their data is unstructured. Think text, video, audio, web server logs, or social media activities.

https://www.brimco.io/analytics/big-data-analytics-statistics/

