

# ECE2810J

## Data Structures and Algorithms

### **k-d Trees**

#### **Learning Objectives:**

- Know what a k-d tree is and its difference over basic binary search tree
- Know how to implement search, insertion, and removal for a k-d tree

# Multidimensional Search

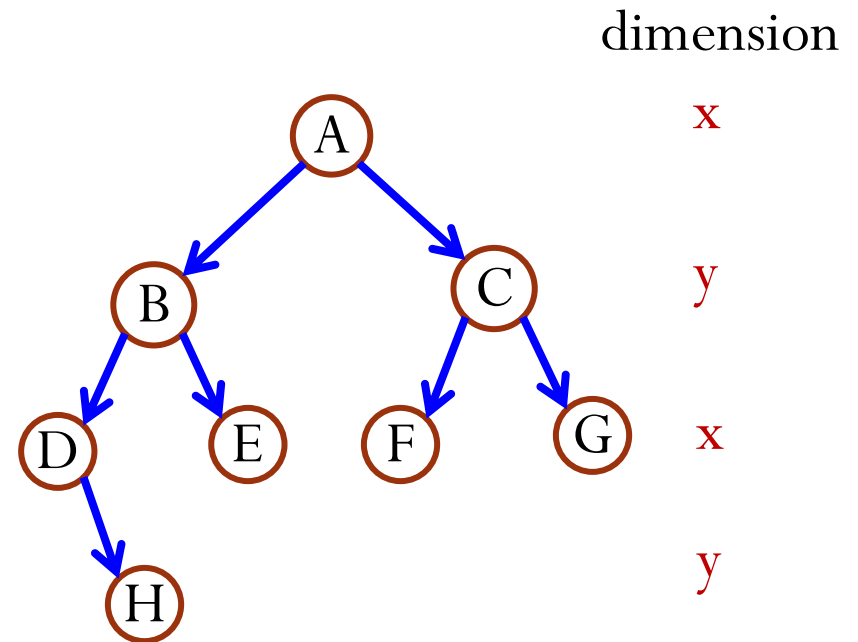
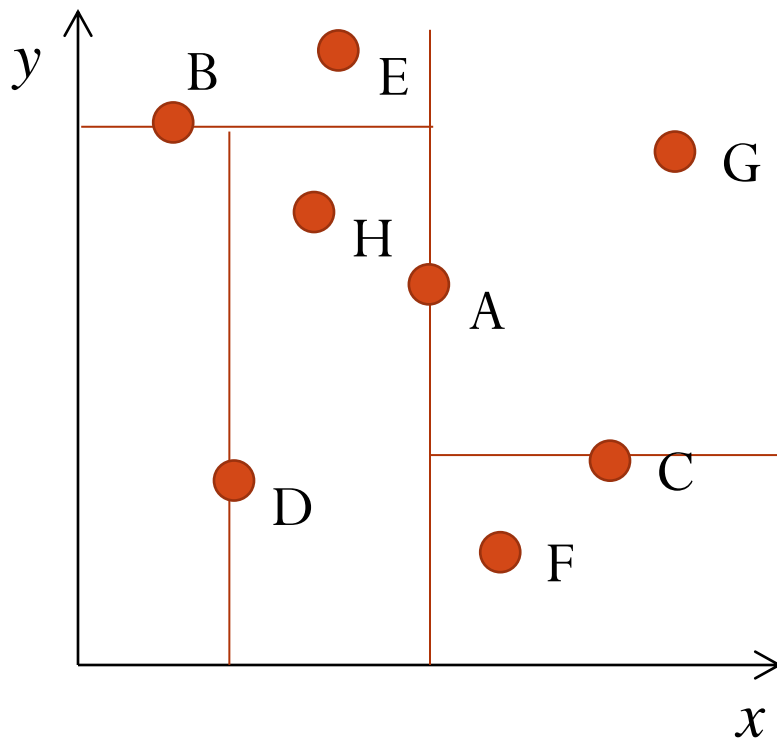
- Example applications:
  - find person by **last name** and **first name** (2D)
  - find location by **latitude** and **longitude** (2D)
  - find book by **author**, **title**, **year published** (3D)
  - find restaurant by **city**, **cuisine**, **popularity**, **sanitation**, **price** (5D)
- Solution:  $k$ -d tree
  - $O(\log n)$  insert and search times

# $k$ -d Tree

- A  $k$ -d tree is a **binary search tree**
- At each level, keys from a different search dimension is used as the **discriminator**
- We **cycle** through the dimensions as we go down the tree
  - For example, given keys consisting of x- and y-coordinates, level 0 discriminates by the x-coordinate, level 1 by the y-coordinate, level 2 again by the x-coordinate, etc.
- For any node  $v$ , suppose that at its level, we discriminate on the  $i$ -th dimension of the key
  - Nodes on the left subtree of node  $v$  have values in dimension  $i$  of their keys  $<$  value in dimension  $i$  of the key of node  $v$
  - Nodes on the right subtree have values in dimension  $i$  of their keys  $\geq$  value in dimension  $i$  of the key of node  $v$

# Example

- k-d tree for points in a 2-D plane



# k-d Tree Insert

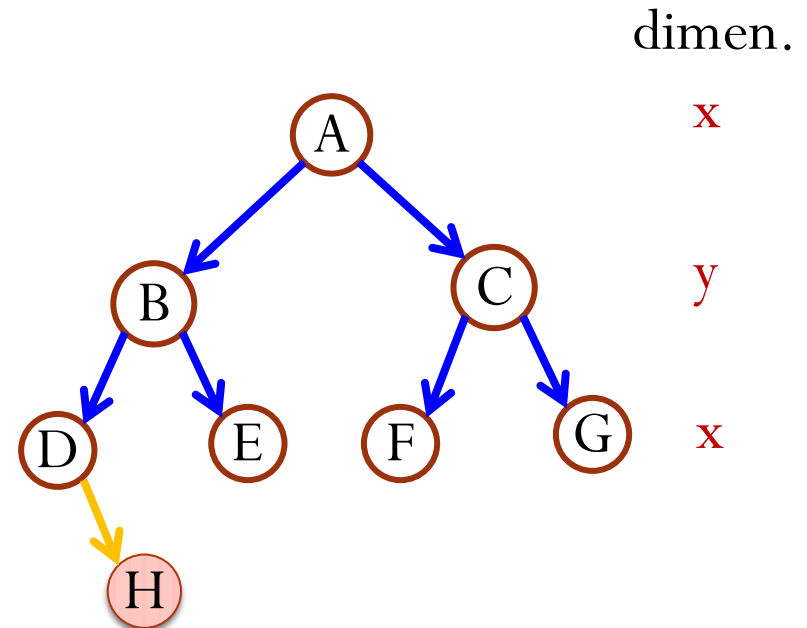
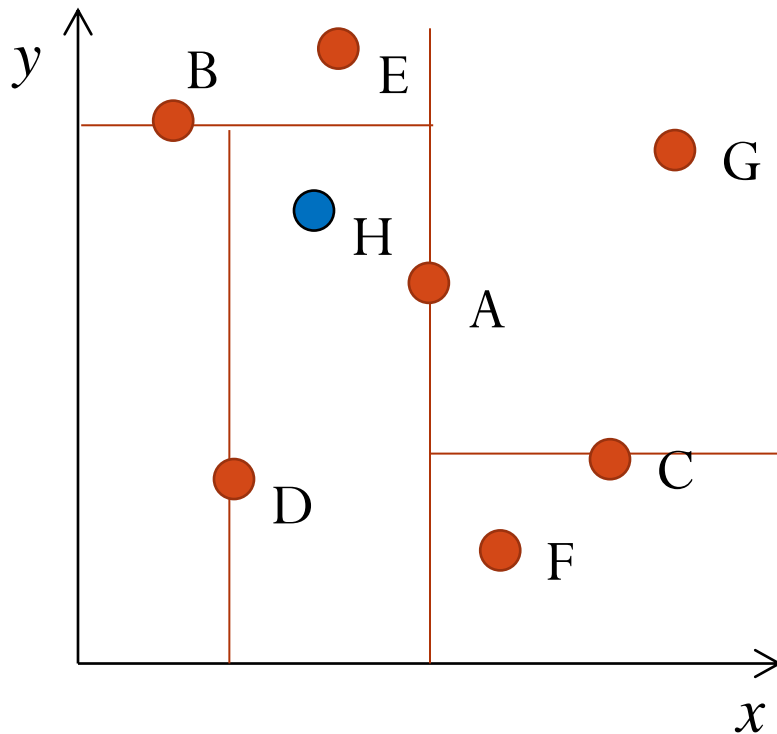
- If new item's key is equal to the root's key, return;
- If new item has a key smaller than that of root along the dimension of the current level, recursive call on left subtree
- Else, recursive call on the right subtree
- In recursive call, cyclically increment the dimension

```
void insert(node *&root, Item item, int dim) {  
    if(root == NULL) {  
        root = new node(item);  
        return;  
    }  
    if(item.key == root->item.key) // equal in all  
        return;                  // dimensions  
    if(item.key[dim] < root->item.key[dim])  
        insert(root->left, item, (dim+1)%numDim);  
    else  
        insert(root->right, item, (dim+1)%numDim);  
}
```

**dim** refers to the dimension of the root

# Insert Example

- Insert H
- Initial function call: `insert(A, H, 0)` // 0 indicates dimension x



# k-d Tree Search

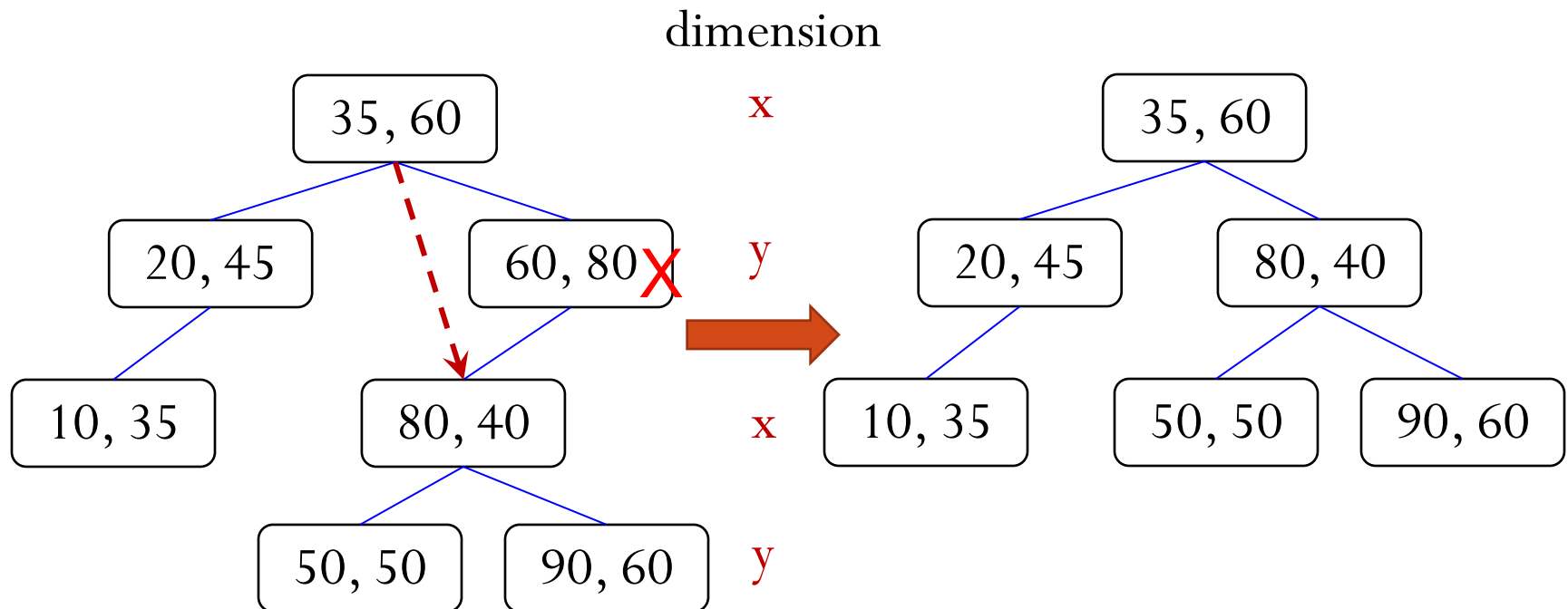
- Search works similarly to insert
  - In recursive call, cyclically increment the dimension

```
node *search(node *root, Key k, int dim) {  
    if(root == NULL) return NULL;  
    if(k == root->item.key)  
        return root;  
    if(k[dim] < root->item.key[dim])  
        return search(root->left, k, (dim+1)%numDim);  
    else  
        return search(root->right, k, (dim+1)%numDim);  
}
```

Time complexities of insert and search are all  $O(\log n)$

# k-d Tree Remove

- If the node is a leaf, simply remove it (e.g., remove (50,50))
- If the node has only one child, can we do the same thing as BST (i.e., connect the node's parent to the node's child)?
  - Consider remove (60, 80) Answer: No!

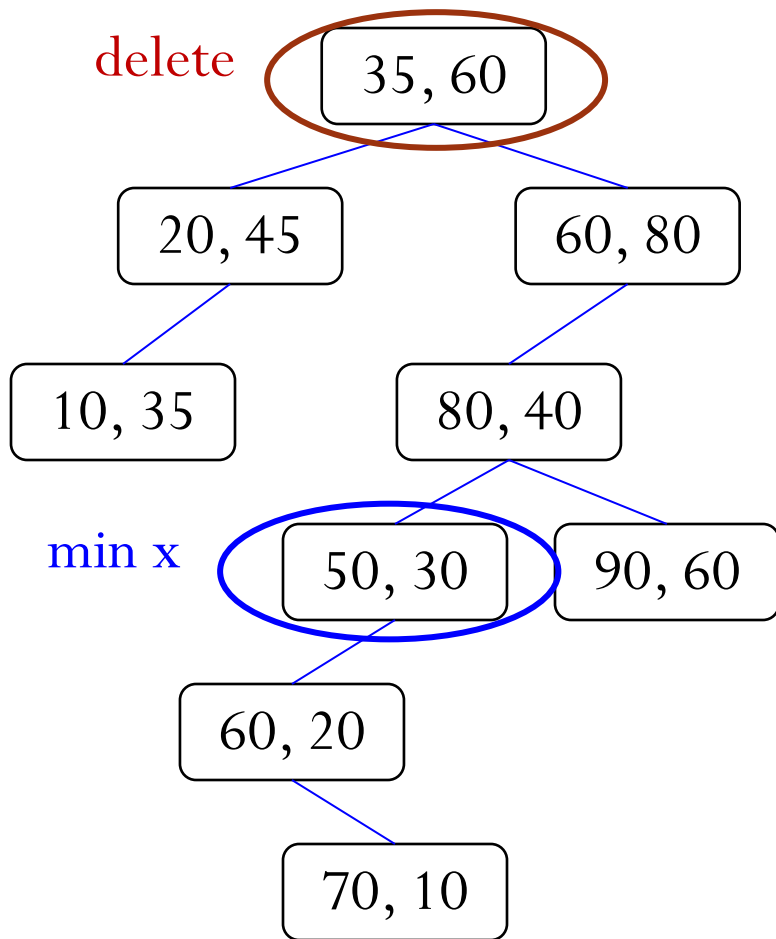




# $k$ -d Tree Removal of Non-leaf Node

- If the node  $R$  to be removed has right subtree, find the node  $M$  in right subtree with the **minimum** value of the current dimension
  - Replace the (key, value) pair of  $R$  with the (key, value) pair of  $M$
  - Recurse on  $M$  until a leaf is reached. Then remove the leaf
- Else, find the node  $M$  in left subtree with the **maximum** value of the current dimension. Then replace and recurse

# k-d Tree Removal Example



x

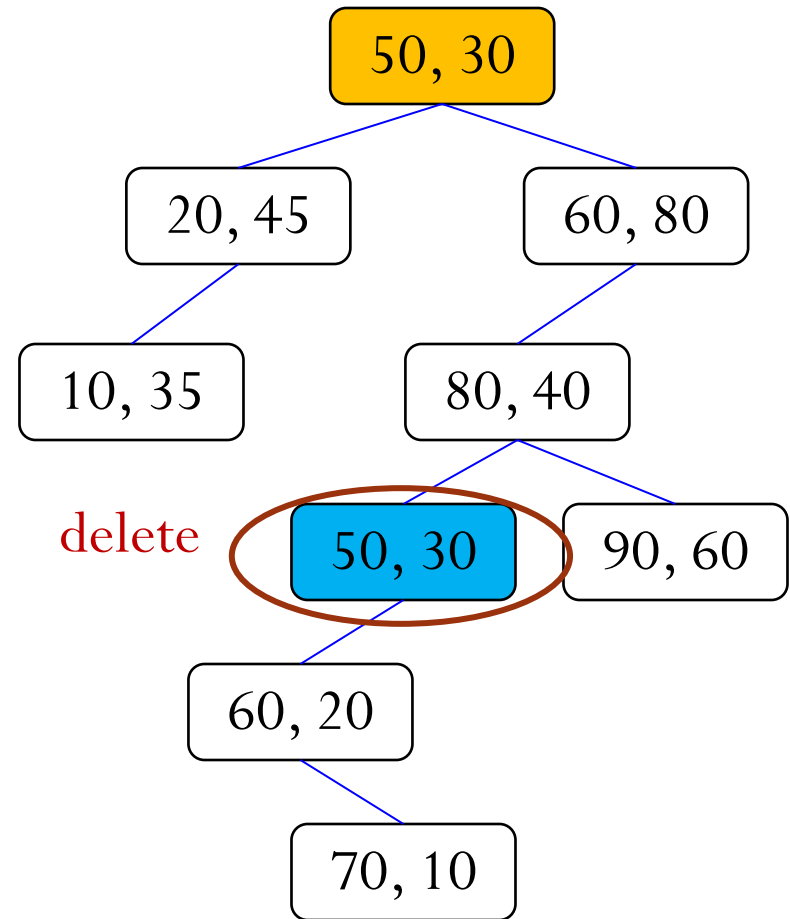
y

x

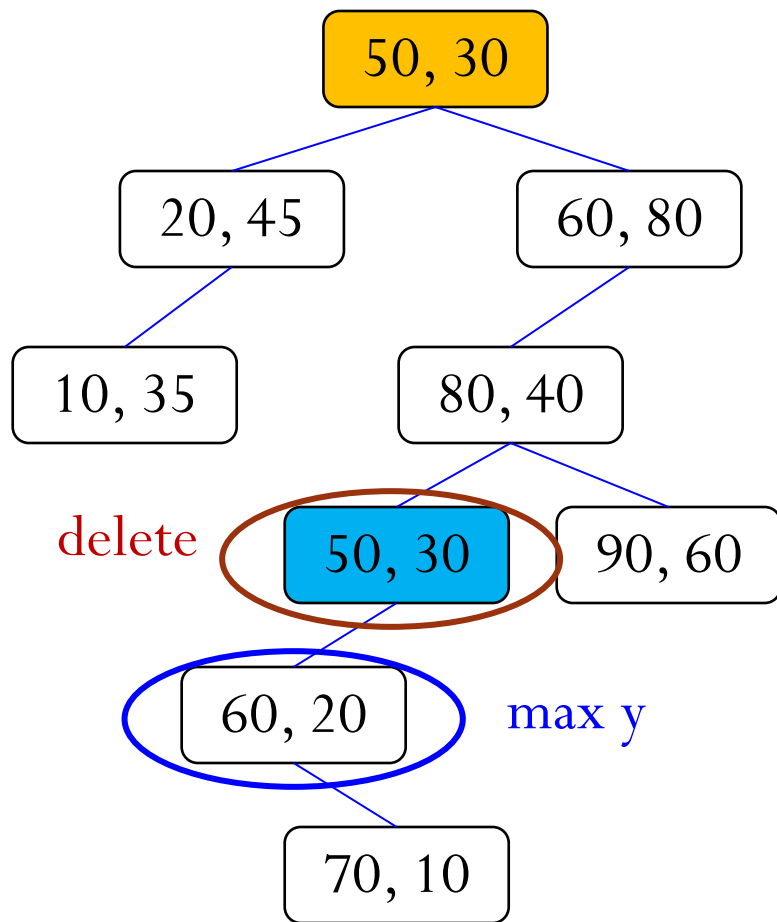
y

x

y



# k-d Tree Removal Example



x

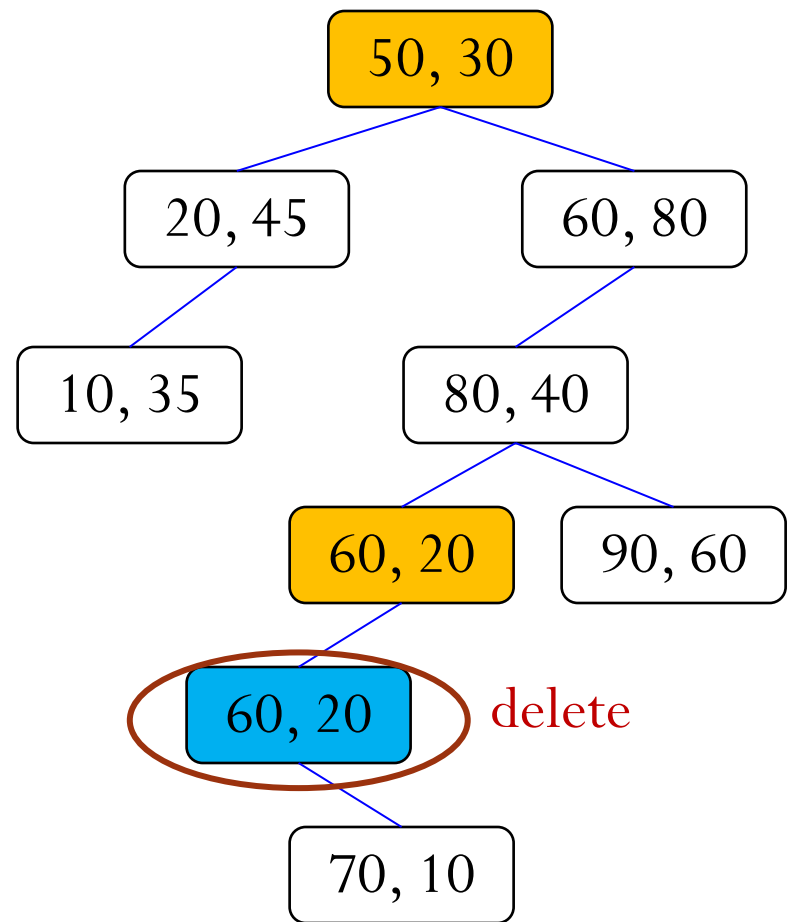
y

x

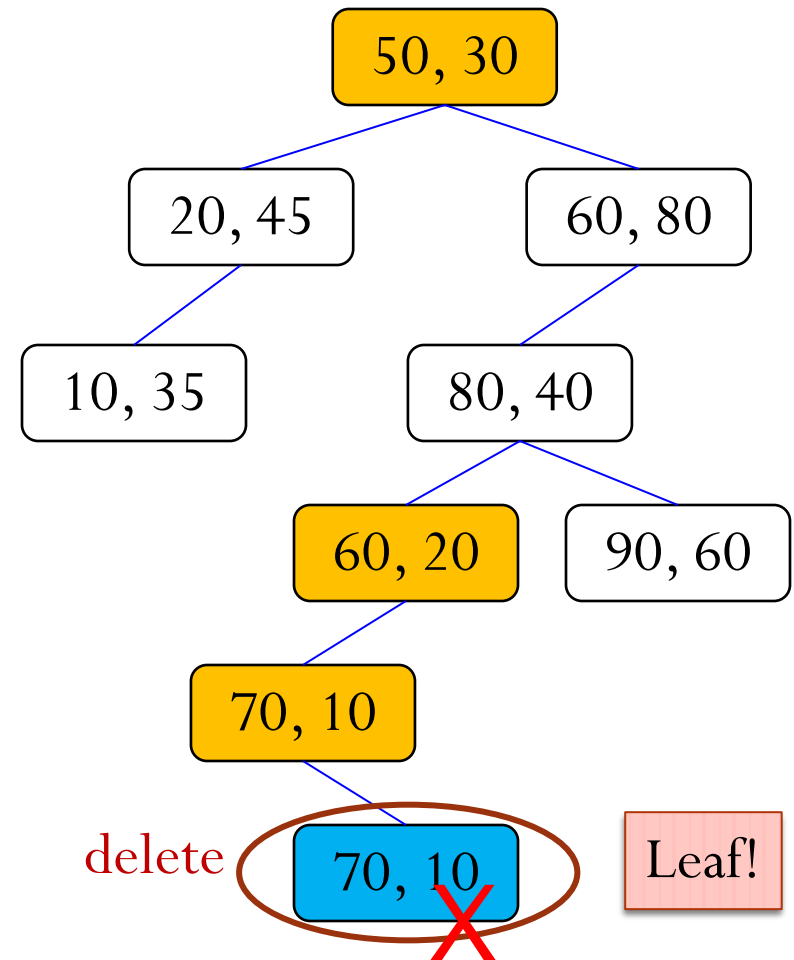
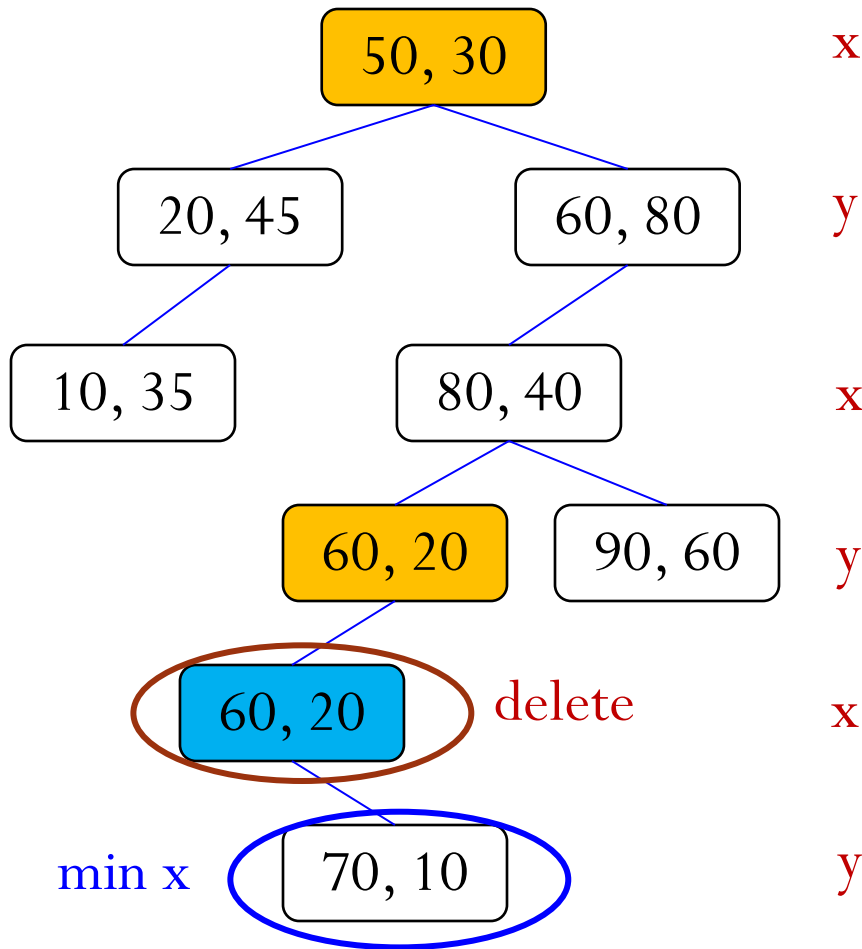
y

x

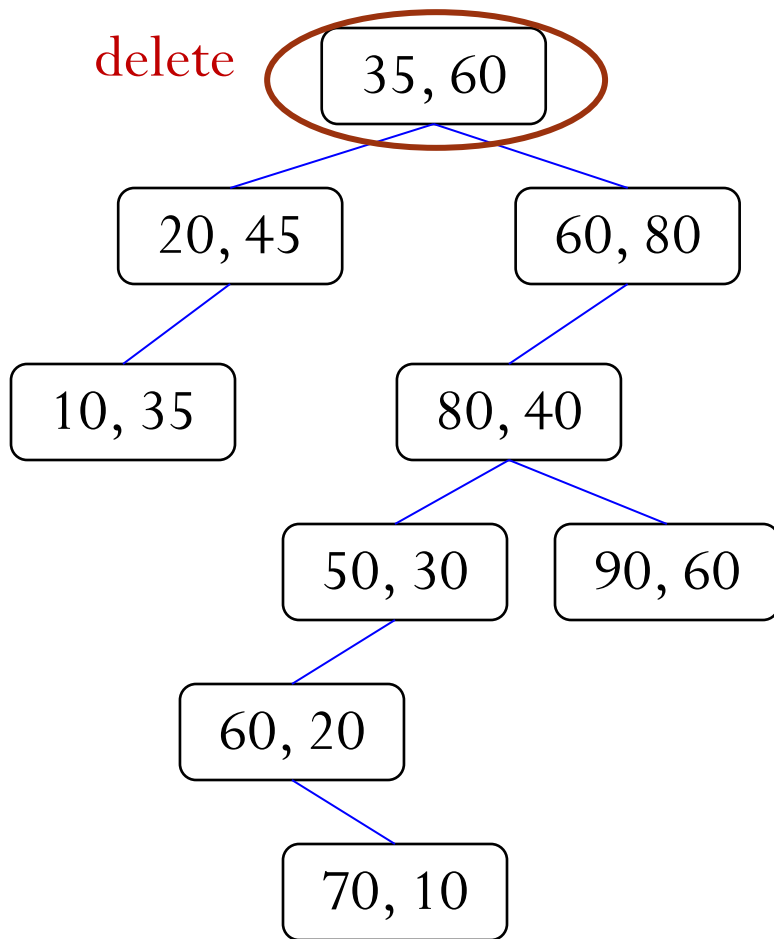
y



# k-d Tree Removal Example



# k-d Tree Removal Example: Summary



x

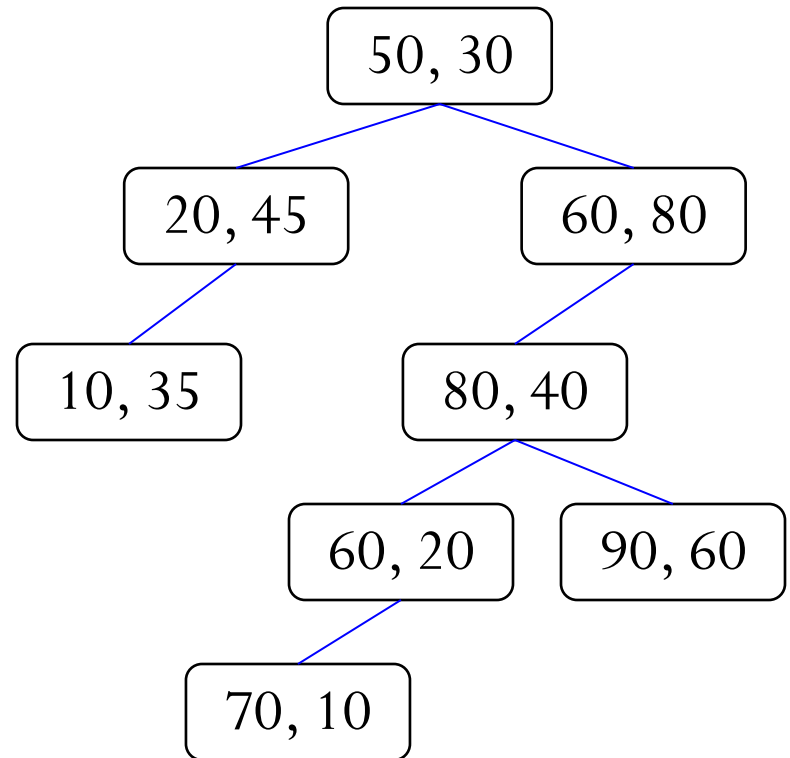
y

x

y

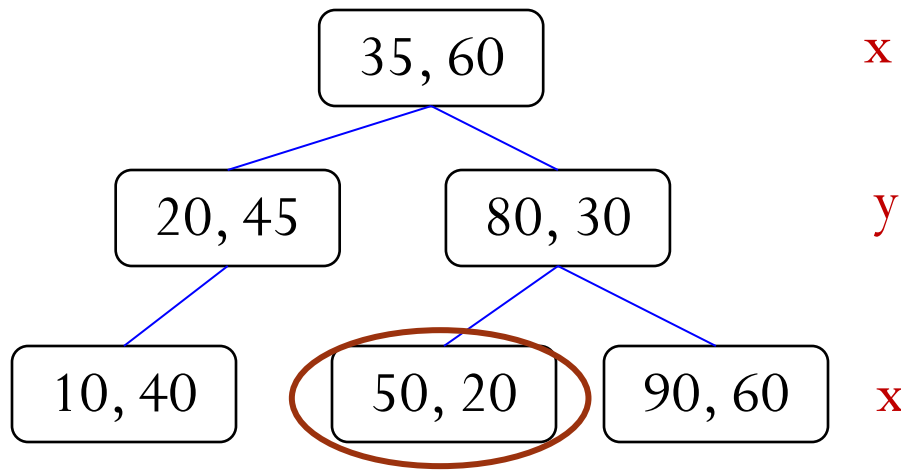
x

y



# Find Minimum Value in a Dimension

- Different from the basic BST, because it may not be the left-most descendent.



Find the node with minimum value in dimension y

# Find Minimum Value in a Dimension

```
node *findMin(node *root, int dimCmp, int dim) {  
    // dimCmp: dimension for comparison  
    if(!root) return NULL;  
    node *min =  
        findMin(root->left, dimCmp, (dim+1)%numDim);  
    if(dimCmp != dim) {  
        rightMin =  
            findMin(root->right, dimCmp, (dim+1)%numDim);  
        min = minNode(min, rightMin, dimCmp);  
    }  
    return minNode(min, root, dimCmp);  
}
```

- **minNode** takes two nodes and a dimension as input, and returns the node with the smaller value in that dimension

# Multidimensional Range Search

- Example
  - Buy ticket for travel between certain dates and certain times
  - Look for apartments within certain price range, certain districts, and number of bedrooms
  - Find all restaurants near you
- k-d tree supports efficient range search, which is similar to that of basic BST



# k-d Tree Range Search

```
void rangeSearch(node *root, int dim,  
    Key searchRange[], Key treeRange[],  
    List results)
```

- Cycle through the dimensions as we go down the level
- **searchRange[]** holds two values (min, max) per dimension
  - Define a hyper-cube
  - min of dimension **j** at **searchRange[2\*j]**, max at **searchRange[2\*j+1]**
- **treeRange[]** holds lower bound and upper bound per dimension for the tree rooted at **root**.
  - Need to be updated as we go down the levels
  - Need to check if a search range overlaps a subtree range