Proof

	First	Follow
javaclass	{C, D}	{\$}
classname	{C, D}	${X} U {B} U FIRST(varref) = {X, B, J, K, ()}$
varlist	{I, S, C, D}	{;} U {) } = { ; ,) }
vardef	{I, S, C, D}	{ , } U FOLLOW(varlist} = {',', ; ,) }
type	{I, S}	$FIRST(varname) \ U \ FIRST(methodname) = \{Y, Z, M, N\}$
varname	$\{Y,Z\}$	FOLLOW(vardef) U {=} U FOLLOW(oprnd) = { ',',;,), =, *,+,<,>,!}
letter	$\{Y,Z\}$	FIRST(char) U FOLLOW(char) U FOLLOW(varname) = {Y, Z, 0, 1, 2, 3, ',', ;,), =, *, +, <, >, !}
char	{Y, Z, 0, 1, 2, 3}	FOLLOW(varname) U FIRST(char) = {Y, Z, 0, 1, 2, 3, ',', ;, }, =, *, +, <, >, !}
digit	{0, 1, 2, 3}	FOLLOW(char) U FIRST(digit) U FOLLOW(integer) = {Y, Z, 0, 1, 2, 3, ',', ; ,), = , *, +, <, >, !}
integer	{0, 1, 2, 3}	FOLLOW(oprnd) = { *, +, ; ,), <, =, >, ! }
varref	{J, K}	FOLLOW(vardef) U {=} U {.} = { = , , ',' , ; ,) }
method	{P, V}	$\{E\}\ U\ FIRST(method) = \{E, P, V\}$
accessor	{P, V}	$FIRST(type) = \{I, S\}$
methodname	{M, N}	{(}
statemt	$\{F, Y, Z, J, K, W\}$	FIRST(returnstatemt) U FIRST(statemt) U FIRST(ifstatemt) U FIRST(assignstatemt) U FIRST(whilestatemt) U FIRST(methodcall) U {E} = {F, Y, Z, J, K, W, R, E}
ifstatemt	{F}	$FOLLOW(statemt) = \{F, Y, Z, J, K, W, R, E\}$
assignstatemt	$\{Y, Z, J, K\}$	{;}
mathexpr	{0,1, 2, 3, Y, Z, (, J, K}	FOLLOW(assignstatemt) U {)} = {;,)}
factor	{0,1, 2, 3, Y, Z, (, J,	{+} U FOLLOW(mathexpr) = {+, ; ,)}

	K}	
oprnd	{0,1, 2, 3, Y, Z, (, J, K)	{*} U FOLLOW(factor) U FIRST(operator) U {) } = { *, +, ; ,), <, =, >, ! }
getvarref	$\{O, J, K\}$	FOLLOW(assignstatemt) = {;}
whilestatemt	{W}	FOLLOW(statemt) = {F, Y, Z, J, K, W, R, E}
cond	{(}	{ T }
operator	{<,=,>,!}	$FIRST(oprnd) = \{0,1,2,3,Y,Z,(,J,K)\}$
returnstatemt	{R}	{E}
methodcall	{J, K}	FOLLOW(statemt) U FOLLOW(oprnd) U FOLLOW(getvarref) = {F, Y, Z, J, K, W, R, E, *, +, ; ,), <, =, >, !}

Proof:<statemt> ::= <ifstatemt> | <assignstatemt>;|<whilestatemt>|<methodcall>

FIRST(ifstatemt), FIRST(assignstatemt), FIRST(whilestatemt), FIRST(methodcall)

 \Rightarrow We have FIRST(assignstatemt) \cap FIRST(methodcall) = $\{Y, Z, J, K\} \cap \{J, K\}$ are not pairwise disjoint. Therefore, we cannot use a Recursive Descent Parser with this grammar