

Building a physics engine - part 2: narrow phase of collision detection

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Narrow phase

Narrow phase

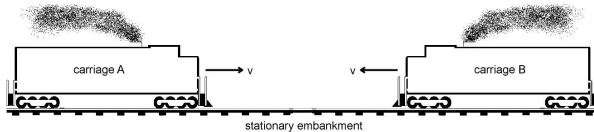
- Find intersections for each pair of rigid bodies
 - Whether they intersect
 - When they intersect
 - Where they intersect
- Acceptable precision
- Very high performance

Narrow phase

Response

- At all points of contact the *relative velocity* projected along the normal must be non-negative
- No pair of bodies in contact is getting closer
 - This would mean that they are penetrating each other

Relative Velocity



Narrow phase

Response

- *Relatively* easy to handle between pairs of bodies
- Multiple bodies in contact make it much harder

Relative Velocity

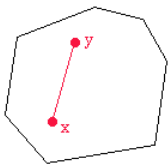


Convex polyhedra

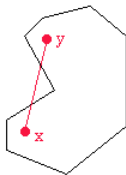
Convex polyhedra

- Also known as *polytopes*
- $\forall P, Q \in C. \forall \alpha \in [0..1]. \text{lerp}(\alpha, P, Q) \in C$
- Complex bodies can be treated as multiple polytopes with distance constraints

Convex polygon

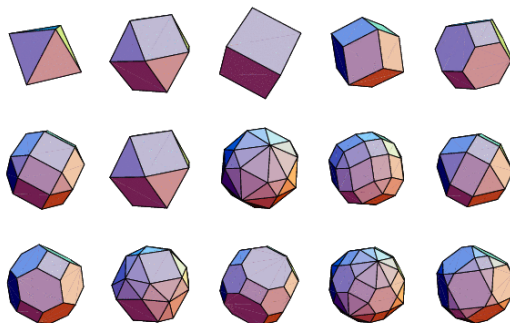


A convex polygon



A non-convex polygon

Convex polyhedra

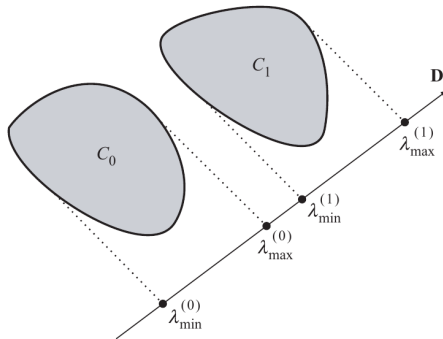


Separating axis

Method of separating axis

- We use a method that works in 2D and 3D
 - We look for an axis/plane that separates the bodies

Separating axis

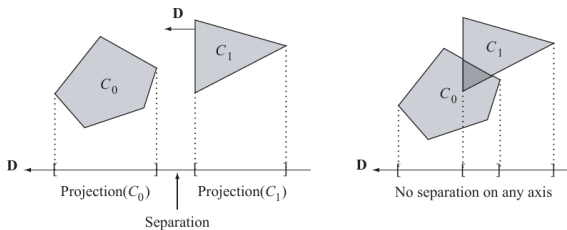


Separating axis

Method of separating axis

- Given a (not necessarily unit-length) direction D and two polytopes C_0 and C_1
- We project them over the direction D : $I_i = [\lambda_{min}^i, \lambda_{max}^i] = [\min_{x \in C_i} \{D \cdot (X - O)\}, \max_{x \in C_i} \{D \cdot (X - O)\}]$
- There is no collision if $\exists D. \lambda_{min}^0 > \lambda_{max}^1 \vee \lambda_{min}^1 > \lambda_{max}^0$

Separating axis

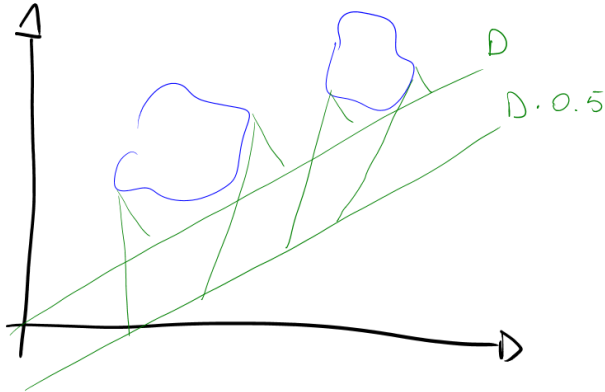


Separating axis

Method of separating axis

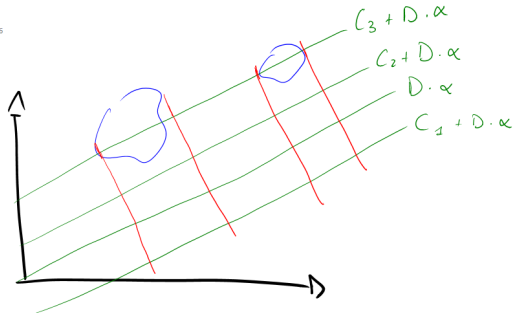
- D may be multiplied by a (non-zero) constant
- The projection is scaled uniformly, but no contact still translates into a separated projection
- The translation of a separating axis remains a separating axis, so we only deal with lines that go through the origin

Scale of separating axis



Translation of separating axis

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Separating axis

Method of separating axis

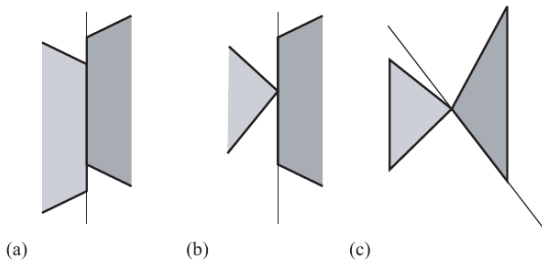
- We now consider C_j polytopes, each with
 - P_i^j vertices, *ordered counter-clockwise* and with wrapping indices
 - $E_i^j = P_{i+1}^j - P_i^j$ edges, *ordered counter-clockwise*
 - N_i^j normals, perpendicular to the edges
- Similarly we store the triangles in 3D

Separating axis

Method of separating axis

- There are infinite candidate directions D
- Fortunately we need only consider a finite set
- In 2D, the normals of each polygon

Separating axis



Separating axis algorithms

Algorithms

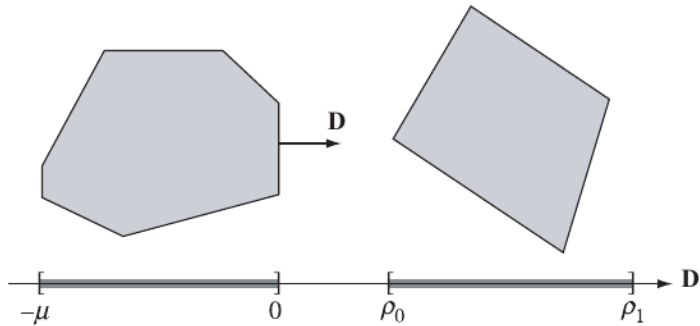
- Naïve algorithm
- For each normal:
 - Project all vertices of both polytopes
 - Compute the intervals
 - Check for intersection

Separating axis algorithms

Algorithms

- Smarter algorithm
 - For each body C_j
 - For each normal P_i^j, N_i^j
 - Check if the closest vertex P_k^l of the other body C_l is too close to the edge
 - $(P_k^l - P_i^j) \cdot N_i^j > \epsilon$

Separation against face



Pseudo-code

```
bool TestIntersection (ConvexPolygon C0, ConvexPolygon  
    C1) {  
    for (i0 = C0.GetN()-1, i1 = 0; i1 < C0.GetN(); i0 =  
        i1++) {  
        P = C0.GetVertex(i1);  
        D = C0.GetNormal(i0);  
        if (WhichSide(C1,P,D) > 0) {  
            return false;  
        }  
    }  
}
```


Pseudo-code

```
for (i0 = C1.GetN()-1, i1 = 0; i1 < C1.GetN(); i0 =  
    i1++) {  
    P = C1.GetVertex(i1);  
    D = C1.GetNormal(i0);  
    if (WhichSide(C0,P,D) > 0) {  
        return false;  
    }  
}  
return true;  
}
```

Pseudo-code

```
int WhichSide (ConvexPolygon C, Point P, Vector D) {  
    posCount = 0;  
    negCount = 0;  
    zeroCount = 0;  
    for (i = 0; i < C.GetN(); ++i) {  
        t = Dot(D, C.GetVertex(i) - P);  
        if (t > 0) { posCount++; }  
        else if (t < 0) { negCount++; }  
        else { zeroCount++; }  
        if ((posCount > 0 and negCount > 0) or zeroCount >  
            0) { return 0; }  
    }  
    return posCount ? 1 : -1;  
}
```

Separating axis algorithms

Optimizations

- Bisection on the sorted vertices
- Find the vertex closest to an edge faster

Separating axis algorithms

Optimizations

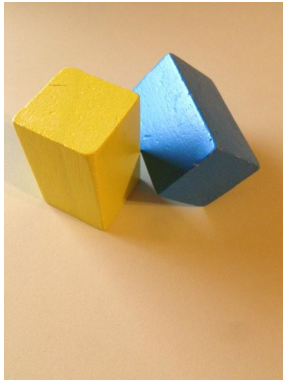
- 2D and 3D: BSP, Gauss map, hash table of the vertices sorted w.r.t. their direction
- Find the vertex closest to an edge much faster
- We will *maybe* see this algorithm, depending on how fast the course goes

Separating axis in 3D

Candidate axes

- In 2D the candidate axes are just the edge normals
- In 3D the face normals are not enough
- Edge-to-edge collisions are not covered

Separating axis

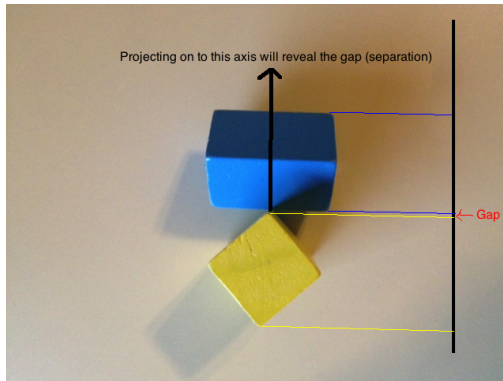


Separating axis in 3D

Candidate axes

- Also consider edge-to-edge cross products
- Exactly the same algorithm, but with more potential axes

Separating axis



Pseudo-code

```
bool TestIntersection (ConvexPolyhedron C0,  
    ConvexPolyhedron C1) {  
    for (i = 0; i < C0.GetFCount(); ++i) {  
        D = C0.GetNormal(i);  
        ComputeInterval(C0,D,min0,max0);  
        ComputeInterval(C1,D,min1,max1);  
        if (max1 < min0 || max0 < min1) { return false; }  
    }  
}
```

Pseudo-code

```
for (j = 0; j < C1.GetFCount(); ++j) {  
    D = C1.GetNormal(j);  
    ComputeInterval(C0,D,min0,max0);  
    ComputeInterval(C1,D,min1,max1);  
    if (max1 < min0 || max0 < min1) { return false; }  
}
```

Pseudo-code

```
for (i = 0; i < C0.GetECount(); ++i) {  
    for (j = 0; j < C1.GetECount(); ++j) {  
        D = Cross(C0.GetEdge(i), C1.Edge(j));  
        ComputeInterval(C0, D, min0, max0);  
        ComputeInterval(C1, D, min1, max1);  
        if (max1 < min0 || max0 < min1){ return false; }  
    }  
}  
return true;  
}
```

Pseudo-code

```
void ComputeInterval (ConvexPolyhedron C, Vector D,  
    double& min, double& max) {  
    min = Dot(D,C.GetVertex(0));  
    max = min;  
    for (i = 1; i < C.GetVCount(); ++i) {  
        value = Dot(D,C.GetVertex(i));  
        if (value < min) {  
            min = value;  
        } else {  
            max = value;  
        }  
    }  
}
```

Moving objects

Moving objects

- Objects are usually moving in a game :)
- We must check for future intersections
- With some simplifying assumptions
 - Rotations can be ignored (phew!)
 - Interpenetration may happen a bit

Moving objects

Moving objects

- Use the *moving projection* on an axis
- Consider the *relative velocity* $V = V_2 - V_1$
- Speed of projection along D is $\sigma = V \cdot D$ when $|D| = 1$
- The distance of the minimum point must be bigger than σ
- $\Delta x = \sigma \Delta t$ $\Delta t \frac{\Delta x}{d} = \text{time of collision}$
- When the time of collision is outside the time of the current frame, then there is no collision

Moving objects

```
bool TestIntersection (ConvexPolygon C0, Vector V0,  
    ConvexPolygon C1, Vector V1, double tmax, double&  
    tfirst, double& tlast) {  
    V = V1 - V0;  
    tfirst = 0;  
    tlast = INFINITY;
```

Pseudo-code

```
for (i0 = C0.GetN() - 1, i1 = 0; i1 < C0.GetN(); i0  
    = i1++) {  
    D = C0.GetNormal(i0);  
    ComputeInterval(C0,D,min0,max0);  
    ComputeInterval(C1,D,min1,max1);  
    speed = Dot(D,V);  
    if (NoIntersect(tmax,speed,min0,max0,min1,max1,  
        tfirst, tlast)) { return false; }  
}
```


Pseudo-code

```
for (i0 = C1.N - 1, i1 = 0; i1 < C1.N; i0 = i1++) {  
    D = C1.GetNormal(i0);  
    ComputeInterval(C0,D,min0,max0);  
    ComputeInterval(C1,D,min1,max1);  
    speed = Dot(D,V);  
    if (NoIntersect(tmax,speed,min0,max0,min1,max1,  
        tfirst, tlast)) { return false; }  
}  
return true;  
}
```

Moving objects

Moving objects

- if the polygons intersect at a first time t_{first} , then there is a separating axis $\forall t. t < t_{\text{first}}$
- if the polygons intersect at a last time t_{last} , then there is a separating axis $\forall t. t > t_{\text{last}}$
- if for all direction $t_{\text{first}} < t_{\text{last}}$, then the polygons intersect at time t_{first} (check against Δt)
- if for all direction $t_{\text{first}} > t_{\text{last}}$, then the polygons do not intersect (all axis must intersect at the same time!)

Pseudo-code

```
bool NoIntersect (double tmax, double speed, double
    min0, double max0, double min1, double max1,
    double& tfirst, double& tlast) {
    if (max1 < min0) {
        if (speed <= 0) { return true; }
        t = (min0 - max1)/speed;
        if (t > tfirst) { tfirst = t; }
        if (tfirst > tmax) { return true; }
        t = (max0 - min1)/speed;
        if (t < tlast) { tlast = t; }
        if (tfirst > tlast) { return true; }
```

Pseudo-code

```
} else if ( max0 < min1 ) {  
    if (speed >= 0) { return true; }  
    t = (max0 - min1)/speed;  
    if (t > tfirst) { tfirst = t; }  
    if (tfirst > tmax) { return true; }  
    t = (min0 - max1)/speed;  
    if (t < tlast) { tlast = t; }  
    if (tfirst > tlast) { return true; }
```

Pseudo-code

```
} else {  
    if (speed > 0) {  
        t = (max0 - min1)/speed;  
        if (t < tlast) { tlast = t; }  
        if (tfirst > tlast) { return true; }  
    } else if (speed < 0) {  
        t = (min0 - max1)/speed;  
        if (t < tlast) { tlast = t; }  
        if (tfirst > tlast) { return true; }  
    }  
}  
return false;  
}
```

Finding intersections

Contact manifold

- The collision response system needs the intersections
- So far we have built `TestIntersection`
- Huge numbers of possible algorithms for this (**GJK** is the most diffuse)
- We now *sketch* a description of one intuitive algorithm

Finding intersections

Finding intersections

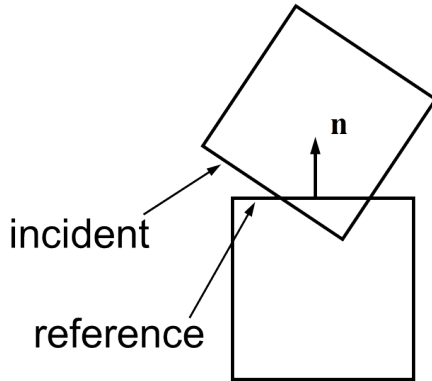
- When the intersection is found at time T , before returning
- We move the bodies forward with the respective velocities
- We compute and return (an approximation) of the contact set
- May be face-to-vertex or edge-to-edge

Finding intersections

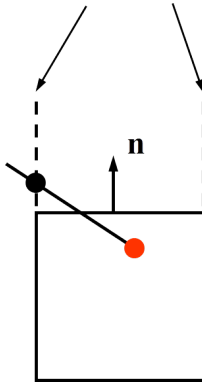
Finding intersections for SAT failures

- Identify *reference face* ($D = N$) and *incident face* (most anti-parallel to D)
- Clip incident face against edges of reference face
- Keep all resulting vertices *below or on* the reference face

Reference face



Reference face sideways



Finding intersections

Edge-to-edge SAT failures

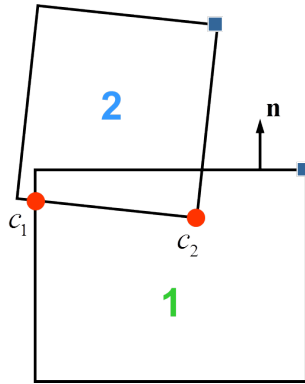
- Just intersect the edges

Finding intersections

Contact manifold

- Two vertices are enough in 2D
- Three vertices are enough in 3D
- Too many more vertices do not improve stability (a few might)
- We always choose the ones penetrating the most
- Deterministic as much as possible (contact caching)

Contact manifold



Finding intersections

Clipping and intersections

- Sutherland-Hogman algorithm for clipping
- Edge-to-edge intersection

Finding intersections

Clipping faces

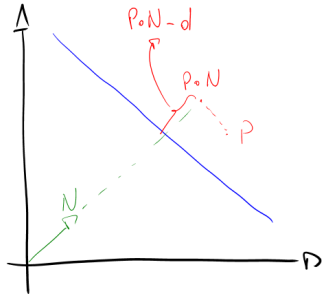
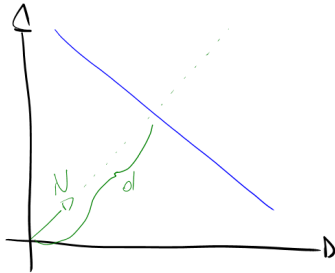
- Project planes from the edges of the reference face
- Vertices of the incident face *inside* all planes are contact points
- Reference plane with incident edge intersections are more contact points
- Keep a subset of the points: (most inside the reference face)

Finding intersections

Clipping a vertex

- Given a plane $\langle N, d \rangle$ where N is the normal of the plane and d is the minimum distance of the plane from the origin
- We determine the side of the plane on which a vertex V lies with the sign of $P \cdot N - d$

Clipping



Finding intersections

Clipping a vertex

- All signs of plane-vertex distance must be the same
- The planes must all be facing inward (or outward)

Finding intersections

Clipping an edge

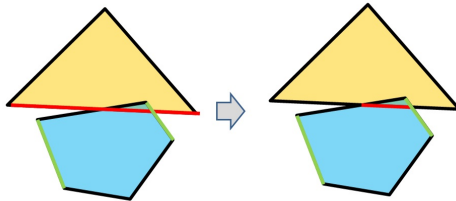
- Given a plane $\langle V \cdot N = d \rangle$ and an edge $V = A + \alpha B$ ($0 \leq \alpha \leq 1$)
- We put these equations together: $(A + \alpha B) \cdot N = d$ and solve for $\alpha = \frac{d - A \cdot N}{B \cdot N}$

Finding intersections

Clipping an edge

- $\alpha = \frac{d - A \cdot N}{B \cdot N}$
- Watch for lack of solutions
 - If $|B \cdot N| \leq \epsilon$ then edge and plane are parallel
 - No intersection possible

Clipping

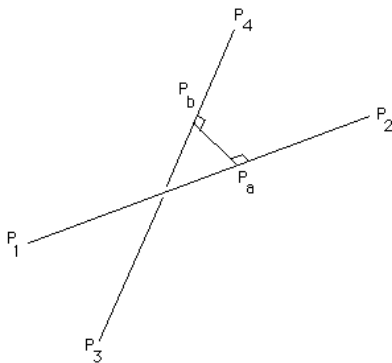


Finding intersections

Edge-to-edge intersections

- We consider two edges: $P_1 + \alpha D_1$ and $P_2 + \beta D_2$
- Intersection in 3D is rather rare
- We look for the shortest connecting axis: $P_a = P_1 + \mu_a D_1$
and $P_b = P_2 + \mu_b D_2$

Clipping



Finding intersections

Edge-to-edge intersections

- The connecting axis is perpendicular to the both edges, so
$$\begin{cases} (P_a - P_b) \cdot D_1 = 0 \\ (P_a - P_b) \cdot D_2 = 0 \end{cases}$$

Finding intersections

Edge-to-edge intersections

- The connecting axis is perpendicular to the both edges, so

$$\begin{cases} (P_a - P_b) \cdot D_1 = 0 \\ (P_a - P_b) \cdot D_2 = 0 \end{cases}$$

- Expanding P_a and P_b results in

$$\begin{cases} (P_1 + \mu_a D_1 - P_2 - \mu_b D_2) \cdot D_1 = 0 \\ (P_1 + \mu_a D_1 - P_2 - \mu_b D_2) \cdot D_2 = 0 \end{cases}$$

Finding intersections

Edge-to-edge intersections

- This is just a system with two unknowns and two equations; the dot products turn everything into numbers:

$$\begin{cases} (P_1 + \mu_a D_1 - P_2 - \mu_b D_2) \cdot D_1 = 0 \\ (P_1 + \mu_a D_1 - P_2 - \mu_b D_2) \cdot D_2 = 0 \end{cases}$$

- The full (tedious) derivation of the solutions is shown in <http://paulbourke.net/geometry/pointlineplane/>

Assignment

Assignment

- Before the end of next week
- Group-work archive/video on Natschool or uploaded somewhere else and linked in your report
- Individual report by each of you on Natschool
- Build a narrow phase collision detector that can compute the contact manifold

Narrow phase collision detection and response
Convex polyhedra
Separating axis
Moving objects
Finding the contact manifold
Assignment

That's it

Thank you!