

# 南京大学 ACM-ICPC 集训队代码模版库



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## 1 General

### 1.1 Code library checksum

```
ab14 #!/usr/bin/python3
c502 import re, sys, hashlib
427e
f7db for line in sys.stdin.read().strip().split("\n") :
ddf5     print(hashlib.md5(re.sub(r'\s|//[.]*', '', line).encode('utf8')).hexdigest()
        [-4:], line)
```

### 1.2 Makefile

```
dab2 .PHONY : run
427e
207e $(t) : $(t).cpp
2d16     g++ --std=c++14 -Wall -D__LOCAL_DEBUG__ -fsanitize=undefined -fsanitize=
        address -ggdb -pipe -o $@ $<
427e
5f25 run : $(t)
bf3e     ./$$(t) < $(t).in
```

### 1.3 .vimrc

```
914c set nocompatible
733d syntax on
6bbc colorscheme slate
7db5 set number
b0e3 set cursorline
061b set shiftwidth=2
8011 set softtabstop=2
a66d set tabstop=2
d23a set expandtab
5245 set magic
740c set smartindent
bee8 set backspace=indent,eol,start
815d set cmdheight=1
0a40 set laststatus=2
1c67 set whichwrap=b,s,<,>[,]
```

### 1.4 Stack

```
const int STK_SZ = 2000000;
char STK[STK_SZ * sizeof(void)];
void *STK_BAK;

#if defined(__i386__)
#define SP "%esp"
#elif defined(__x86_64__)
#define SP "%rsp"
#endif

int main() {
    asm volatile("movl SP, %0; movl %1, SP: =g(STK_BAK):g(STK+sizeof(STK));");
    ;

    // main program

    asm volatile("movl %0, SP::g(STK_BAK);");
    return 0;
}
```

### 1.5 Template

```
#include <bits/stdc++.h>
using namespace std;

#ifdef __LOCAL_DEBUG__
# define _debug(fmt, ...) fprintf(stderr, "[%s] " fmt "\n", \
    __func__, __VA_ARGS__)
#else
# define _debug(...) ((void) 0)
#endif

#define rep(i, n) for (int i=0; i<(n); i++)
#define Rep(i, n) for (int i=1; i<=(n); i++)
#define range(x) begin(x), end(x)
typedef long long LL;
typedef unsigned long long ULL;
```

## 2 Miscellaneous Algorithms

### 2.1 2-SAT

```

0f42 const int MAXN = 100005;
03a9 struct twoSAT{
5c83     int n;
8f72     vector<int> G[MAXN*2];
d060     bool mark[MAXN*2];
b42d     int S[MAXN*2], c;
427e
d34f     void init(int n){
b985         this->n = n;
f9ec         for (int i=0; i<n*2; i++) G[i].clear();
0609         memset(mark, 0, sizeof(mark));
95cf     }
427e
3bd5     bool dfs(int x){
bd70         if (mark[x^1]) return false;
c96a         if (mark[x]) return true;
fd23         mark[x] = true;
4bea         S[c++] = x;
1ce6         for (int i=0; i<G[x].size(); i++)
d942             if (!dfs(G[x][i])) return false;
3361         return true;
95cf     }
427e
5894     void add_clause(int x, bool xval, int y, bool yval){
6afe         x = x * 2 + xval;
e680         y = y * 2 + yval;
81cc         G[x^1].push_back(y);
6835         G[y^1].push_back(x);
95cf     }
427e
d0cb     bool solve() {
7c39         for (int i=0; i<n*2; i+=2){
e63f             if (!mark[i] && !mark[i+1]){
88fb                 c = 0;
f4b9                 if (!dfs(i)){
3f03                     while (c > 0) mark[S[--c]] = false;
86c5                     if (!dfs(i+1)) return false;
95cf                 }
95cf             }

```

```

    }
    return true;
}

inline bool value(unsigned i){return mark[2*i+1];}
};

```

95cf  
3361  
95cf  
427e  
5f0a  
329b

### 2.2 Knuth's optimization

```

int n;
int dp[256][256], dc[256][256];

template <typename T>
void compute(T cost) {
    for (int i = 0; i <= n; i++) {
        dp[i][i] = 0;
        dc[i][i] = i;
    }
    rep (i, n) {
        dp[i][i+1] = 0;
        dc[i][i+1] = i;
    }
    for (int len = 2; len <= n; len++) {
        for (int i = 0; i + len <= n; i++) {
            int j = i + len;
            int lbnd = dc[i][j-1], rbnd = dc[i+1][j];
            dp[i][j] = INT_MAX / 2;
            int c = cost(i, j);
            for (int k = lbnd; k <= rbnd; k++) {
                int res = dp[i][k] + dp[k][j] + c;
                if (res < dp[i][j]) {
                    dp[i][j] = res;
                    dc[i][j] = k;
                }
            }
        }
    }
};

```

5c83  
d77c  
427e  
b7ec  
0bc7  
0423  
8f5e  
9488  
95cf  
be8e  
95b5  
aa0f  
95cf  
ec08  
88b8  
d3da  
9824  
a24a  
f933  
90d2  
9bd0  
26b5  
e6af  
9c88  
95cf  
95cf  
95cf  
95cf  
329b

## 2.3 Mo's algorithm

All intervals are closed on both sides. When running functions `enter()` and `leave()`, the global `l` and `r` has not changed yet.

### Usage:

`add_query(id, l, r)`    Add id-th query  $[l, r]$ .  
`run()`    Run Mo's algorithm.  
`init()`    **TODO.** Initialize the range  $[l, r]$ .  
`yield(id)`    **TODO.** Yield answer for id-th query.  
`enter(o)`    **TODO.** Add o-th element.  
`leave(o)`    **TODO.** Remove o-th element.

```
5194 constexpr int BLOCK_SZ = 300;
427e
3ec4 struct query { int l, r, id; };
d26a vector<query> queries;
427e
1e30 void add_query(int id, int l, int r) {
54c9     queries.push_back(query{l, r, id});
95cf }
427e
9f6b int l, r;
427e
427e // ----- functions to implement -----
62b4 inline void init();
50e1 inline void yield(int id);
b20d inline void enter(int o);
13af inline void leave(int o);
427e
37f0 void run() {
ab0b     if (queries.empty()) return;
8508     sort(range(queries), [](query lhs, query rhs) {
c7f8         int lb = lhs.l / BLOCK_SZ, rb = rhs.l / BLOCK_SZ;
03e7         if (lb != rb) return lb < rb;
0780         return lhs.r < rhs.r;
b251     });
6196     l = queries[0].l;
9644     r = queries[0].r;
07e2     init();
5bc9     for (query q : queries) {
7bc7         while (l > q.l) enter(l - 1), l--;
d646         while (r < q.r) enter(r + 1), r++;
13f0         while (l < q.l) leave(l), l++;
e1c6         while (r > q.r) leave(r), r--;
```

```
        yield(q.id);
    }
}
```

82f5  
95cf  
95cf

## 2.4 Matroid Intersection

Find the maximum cardinality common independent set of two matroids. Matroids are given by independence oracle.

### Usage:

`MatroidOracle`    The independence oracle maintaining an independent set.  
**Note** that the default constructor must properly initialize inner state to an empty set.  
`insert(x)`    Insert element labeled  $x$  to the independent set.  
`test(x)`    Test whether the set is still independent if  $x$  is inserted.  
`MatroidIntersection< MT1, MT2>(n)`    Construct the matroid intersection solver with  $n$  elements labeled from 0 and matroid oracles MT1 and MT2.  
`run()`    Run the algorithm and return the matroid intersection.

```
struct MatroidOracle {
    MatroidOracle() { /* TODO */ }
    void insert(int x) { /* TODO */ }
    bool test(int x) const { /* TODO */ }
};

const int MAXN = 8192;
template <typename MT1, typename MT2>
struct MatroidIntersection {
    int n;
    bool in[MAXN] = {}, t[MAXN], vis[MAXN];
    int pre[MAXN];
    vector<int> adj[MAXN];
    queue<int> q;

    MatroidIntersection(int n) : n(n) { }

    vector<int> getcur() {
        vector<int> ret;
        rep (i, n) if (in[i]) ret.push_back(i);
        return ret;
    }

    void enqueue(int x, int p) {
```

0935  
297b  
53e5  
ff18  
329b  
427e  
a015  
94cc  
3288  
5c83  
5550  
fe84  
0b32  
93d2  
427e  
c152  
427e  
2ed1  
995a  
a585  
ee0f  
95cf  
427e  
ca2b

```

e5da         if (vis[x]) return;
f4a6         vis[x] = true; pre[x] = p; q.push(x);
ff59         if (t[x]) throw x;
329b     };
427e
9081     vector<int> run() {
1026         while (true) {
c40f             vector<int> cur = getcur();
6f47             fill(vis, vis + n, 0);
943b             rep (i, n) adj[i].clear();
0e02             MT2 mt2;
3e54             for (int i : cur) mt2.insert(i);
191d             rep (i, n) t[i] = mt2.test(i);
e167             vector<MT1> mt1s(cur.size());
46d2             vector<MT2> mt2s(cur.size());
660b             rep (i, cur.size()) rep (j, cur.size()) if (i != j) {
3cd7                 mt1s[i].insert(cur[j]);
9680                 mt2s[i].insert(cur[j]);
95cf             }
e8d7             rep (i, n) if (!lin[i]) rep (j, cur.size()) {
3fe9                 if (mt1s[j].test(i)) adj[cur[j]].push_back(i);
645e                 if (mt2s[j].test(i)) adj[i].push_back(cur[j]);
95cf             }
cf76             q = {};
85eb             try {
2f4f                 MT1 mt1;
2f34                 for (int i : cur) mt1.insert(i);
4053                 rep (i, n) if (mt1.test(i)) enqueue(i, -1);
1c7d                 while (q.size()) {
c048                     int u = q.front(); q.pop();
a697                     for (int v : adj[u]) enqueue(v, u);
95cf                 }
5a9a             } catch (int v) {
a8f3                 while (v >= 0) { in[v] ^= 1; v = pre[v]; }
b333                 continue;
95cf             }
6173             break;
329b         };
f2de         return getcur();
95cf     }
329b };

```

## 3 String

### 3.1 Knuth-Morris-Pratt algorithm

```

const int SIZE = 10005;

struct kmp_matcher {
    char p[SIZE];
    int fail[SIZE];
    int len;

    void construct(const char* needle) {
        len = strlen(p);
        strcpy(p, needle);
        fail[0] = fail[1] = 0;
        for (int i = 1; i < len; i++) {
            int j = fail[i];
            while (j && p[i] != p[j]) j = fail[j];
            fail[i + 1] = p[i] == p[j] ? j + 1 : 0;
        }
    }

    inline void found(int pos) {
        // ! add codes for having found at pos
    }

    void match(const char* haystack) { // must be called after construct
        const char* t = haystack;
        int n = strlen(t);
        int j = 0;
        rep(i, n) {
            while (j && p[j] != t[i]) j = fail[j];
            if (p[j] == t[i]) j++;
            if (j == len) found(i - len + 1);
        }
    }
};

```

### 3.2 Manacher algorithm

```

struct Manacher {
    int Len;

```

```

9255 vector<int> lc;
b301 string s;
427e
ec07 void work() {
c033     lc[1] = 1;
6bef     int k = 1;
427e
491f     for (int i = 2; i <= Len; i++) {
7957         int p = k + lc[k] - 1;
5e04         if (i <= p) {
24a1             lc[i] = min(lc[2 * k - i], p - i + 1);
8e2e         } else {
e0e5             lc[i] = 1;
95cf         }
74ff         while (s[i + lc[i]] == s[i - lc[i]]) lc[i]++;
2b9a         if (i + lc[i] > k + lc[k]) k = i;
95cf     }
95cf }
427e
bfd5 void init(const char *tt) {
aaaf     int len = strlen(tt);
f701     s.resize(len * 2 + 10);
7045     lc.resize(len * 2 + 10);
8e13     s[0] = '*';
ae54     s[1] = '#';
1321     for (int i = 0; i < len; i++) {
e995         s[i * 2 + 2] = tt[i];
69fd         s[i * 2 + 1] = '#';
95cf     }
43fd     s[len * 2 + 1] = '#';
75d1     s[len * 2 + 2] = '\0';
61f7     Len = len * 2 + 2;
3e7a     work();
95cf }
427e
b194 pair<int, int> maxpal(int l, int r) {
901a     int center = l + r + 1;
ffb2     int rad = lc[center] / 2;
ab54     int rmid = (l + r + 1) / 2;
17e4     int rl = rmid - rad, rr = rmid + rad - 1;
3908     if ((r ^ l) & 1) {
69f3     } else rr++;
69dc     return {max(l, rl), min(r, rr)};
95cf }

```

```
};
```

329b

### 3.3 Aho-corasick automaton

```

struct AC : Trie {
    int fail[MAXN];
    int last[MAXN];

    void construct() {
        queue<int> q;
        fail[0] = 0;
        rep(c, CHARN) {
            if (int u = tr[0][c]) {
                fail[u] = 0;
                q.push(u);
                last[u] = 0;
            }
        }
        while (!q.empty()) {
            int r = q.front();
            q.pop();
            rep(c, CHARN) {
                int u = tr[r][c];
                if (!u) {
                    tr[r][c] = tr[fail[r]][c];
                    continue;
                }
                q.push(u);
                int v = fail[r];
                while (v && !tr[v][c]) v = fail[v];
                fail[u] = tr[v][c];
                last[u] = tag[fail[u]] ? fail[u] : last[fail[u]];
            }
        }
    }

    void found(int pos, int j) {
        if (j) {
            // ! add codes for having found word with tag[j]
            found(pos, last[j]);
        }
    }
}

```

```

a1ad
9143
daca
427e
8690
93d2
a7a6
ce3c
b1c6
a506
3e14
f689
95cf
95cf
cc78
31f0
15dd
ce3c
ab59
0ef5
9d58
b333
95cf
3e14
b3ff
d2ea
c275
654c
95cf
95cf
427e
7752
043e
427e
4a96
95cf
95cf

```

```

427e void find(const char* text) { // must be called after construct()
9785     int p = 0, c, len = strlen(text);
80a4     rep(i, len) {
9c94         c = id(text[i]);
b3db         p = tr[p][c];
f119         if (tag[p])
f08e             found(i, p);
389b         else if (last[p])
299e             found(i, last[p]);
95cf     }
95cf }
329b };

```

### 3.4 Trie

```

e6f1 const int MAXN = 12000;
dd87 const int CHARN = 26;
427e
8ff5 inline int id(char c) { return c - 'a'; }
427e
a281 struct Trie {
5c83     int n;
f4f5     int tr[MAXN][CHARN]; // Trie tree, 0 denotes fail
35a5     int tag[MAXN];
427e
4fee     Trie() {
3ccc         memset(tr[0], 0, sizeof(tr[0]));
4d52         tag[0] = 0;
46bf         n = 1;
95cf     }
427e
427e // tag should not be 0
30b0 void add(const char* s, int t) {
d50a     int p = 0, c, len = strlen(s);
9c94     rep(i, len) {
3140         c = id(s[i]);
d6c8         if (!tr[p][c]) {
26dd             memset(tr[n], 0, sizeof(tr[n]));
2e5c             tag[n] = 0;
73bb             tr[p][c] = n++;
95cf         }

```

```

        p = tr[p][c];
    }
    tag[p] = t;
}

// returns 0 if not found
// AC automaton does not need this function
int search(const char* s) {
    int p = 0, c, len = strlen(s);
    rep(i, len) {
        c = id(s[i]);
        if (!tr[p][c]) return 0;
        p = tr[p][c];
    }
    return tag[p];
}
};

```

### 3.5 Suffix array

The character immediately after the end of the string **MUST** be set to the **UNIQUE SMALLEST** element.

Usage:

s[]	the source string
sa[i]	the index of starting position of $i$ -th suffix
rk[i]	the number of suffixes less than the suffix starting from $i$
h[i]	the longest common prefix between the $i$ -th and $(i-1)$ -th lexicographically smallest suffixes
n	size of source string
m	size of character set

```

void radix_sort(int x[], int y[], int sa[], int n, int m) {
    static int cnt[1000005]; // size > max(n, m)
    fill(cnt, cnt + m, 0);
    rep(i, n) cnt[x[y[i]]]++;
    partial_sum(cnt, cnt + m, cnt);
    for (int i = n - 1; i >= 0; i--) sa[--cnt[x[y[i]]]] = y[i];
}

void suffix_array(int s[], int sa[], int rk[], int n, int m) {
    static int y[1000005]; // size > n
    copy(s, s + n, rk);
    iota(y, y + n, 0);

```



```

7b42 radix_sort(rk, y, sa, n, m);
c8c2 for (int j = 1, p = 0; j <= n; j <= 1, m = p, p = 0) {
8c3a     for (int i = n - j; i < n; i++) y[p++] = i;
9323     rep (i, n) if (sa[i] >= j) y[p++] = sa[i] - j;
9e9d     radix_sort(rk, y, sa, n, m + 1);
ae41     swap_ranges(rk, rk + n, y);
ffd2     rk[sa[0]] = p = 1;
445e     for (int i = 1; i < n; i++)
f8dc         rk[sa[i]] = ((y[sa[i]] == y[sa[i-1]] and y[sa[i]+j] == y[sa[i-1]+j])
                ? p : ++p);
02f0     if (p == n) break;
95cf }
97d9 rep (i, n) rk[sa[i]] = i;
95cf }
427e
1715 void calc_height(int s[], int sa[], int rk[], int h[], int n) {
c41f     int k = 0;
f313     h[0] = 0;
be8e     rep (i, n) {
0883         k = max(k - 1, 0);
527d         if (rk[i]) while (s[i+k] == s[sa[rk[i]-1]+k]) ++k;
56b7         h[rk[i]] = k;
95cf     }
95cf }

```

### 3.6 Rolling hash

**PLEASE** call `init_hash()` in `int main()`!

**Usage:**

`build(str)` Construct the hasher with given string.  
`operator()(l, r)` Get hash value of substring  $[l, r)$ .

```

1e42 const LL mod = 1006658951440146419, g = 967;
9f60 const int MAXN = 200005;
0291 LL pg[MAXN];
427e
dfe7 inline LL mul(LL x, LL y) { return __int128_t(x) * y % mod; }
427e
599a void init_hash() { // must be called in `int main()`
286f     pg[0] = 1;
4af8     for (int i = 1; i < MAXN; i++) pg[i] = mul(pg[i-1], g);
95cf }
427e

```

```

struct hasher {
    LL val[MAXN];

    void build(const char *str) { // assume lower-case letter only
        for (int i = 0; str[i]; i++)
            val[i+1] = (mul(val[i], g) + str[i]) % mod;
    }

    LL operator() (int l, int r) { // [l, r)
        return (val[r] - mul(val[l], pg[r-l]) + mod) % mod;
    }
};

```

7e62  
534a  
427e  
4554  
f937  
9645  
95cf  
427e  
19f8  
9986  
95cf  
329b

## 4 Math

### 4.1 Extended Euclidean algorithm and Chinese remainder theorem

```

void exgcd(LL a, LL b, LL &g, LL &x, LL &y) {
    if (!b) g = a, x = 1, y = 0;
    else {
        exgcd(b, a % b, g, y, x);
        y -= x * (a / b);
    }
}

```

4fba  
7db6  
037f  
ffca  
d798  
95cf  
95cf  
427e  
e491  
84e6  
00d9  
be8e  
98b4  
9f4f  
b082  
3cd3  
95cf  
2e47  
95cf

```

LL crt(LL r[], LL p[], int n) {
    LL q = 1, ret = 0;
    rep (i, n) q *= p[i];
    rep (i, n) {
        LL m = q / p[i];
        LL d, x, y;
        exgcd(p[i], m, d, x, y);
        ret = (ret + y * m * r[i]) % q;
    }
    return (q + ret) % q;
}

```

### 4.2 Linear basis

```
const int MAXD = 30;
```

8b44

```

03a6 struct linearbasis {
3558     ULL b[MAXD] = {};
427e
1566     bool insert(LL v) {
9b2b         for (int j = MAXD - 1; j >= 0; j--) {
de36             if (!(v & (1ll << j))) continue;
ee78             if (b[j] v ^= b[j]
037f                 else {
7836                 for (int k = 0; k < j; k++)
f0b4                     if (v & (1ll << k)) v ^= b[k];
b0aa                 for (int k = j + 1; k < MAXD; k++)
46c9                     if (b[k] & (1ll << j)) b[k] ^= v;
8295                 b[j] = v;
3361                 return true;
95cf             }
95cf         }
438e     return false;
95cf }
329b };

```

### 4.3 Gauss elimination over finite field

```

b784 const LL p = 1000000007;
427e
2a2c LL powmod(LL b, LL e) {
95a2     LL r = 1;
3e90     while (e) {
1783         if (e & 1) r = r * b % p;
5549         b = b * b % p;
16fc         e >>= 1;
95cf     }
547e     return r;
95cf }
427e
c130 typedef vector<LL> VLL;
42ac typedef vector<VLL> VVLL;
427e
2c62 LL gauss(VVLL &a, VVLL &b) {
561b     const int n = a.size(), m = b[0].size();
a25e     vector<int> irow(n), icol(n), ipiv(n);
2976     LL det = 1;
427e

```

```

rep (i, n) {
    int pj = -1, pk = -1;
    rep (j, n) if (!ipiv[j])
        rep (k, n) if (!ipiv[k])
            if (pj == -1 || a[j][k] > a[pj][pk]) {
                pj = j;
                pk = k;
            }
    if (a[pj][pk] == 0) return 0;
    ipiv[pk]++;
    swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
    if (pj != pk) det = (p - det) % p;
    irow[i] = pj;
    icol[i] = pk;

    LL c = powmod(a[pk][pk], p - 2);
    det = det * a[pk][pk] % p;
    a[pk][pk] = 1;
    rep (j, n) a[pk][j] = a[pk][j] * c % p;
    rep (j, m) b[pk][j] = b[pk][j] * c % p;
    rep (j, n) if (j != pk) {
        c = a[j][pk];
        a[j][pk] = 0;
        rep (k, n) a[j][k] = (a[j][k] + p - a[pk][k] * c % p) % p;
        rep (k, m) b[j][k] = (b[j][k] + p - b[pk][k] * c % p) % p;
    }
}

for (int j = n - 1; j >= 0; j--) if (irow[j] != icol[j]) {
    for (int k = 0; k < n; k++) swap(a[k][irow[j]], a[k][icol[j]]);
}
return det;
}

```

### 4.4 Berlekamp-Massey algorithm

Call `berlekamp()` with input sequence  $(x_0, x_1, \dots, x_{n-1})$ . Return a vector of coefficients  $(c_0 = 1, c_1, \dots, c_{m-1})$  with minimum  $m$ , such that  $\sum_{i=0}^m c_i x_{j-i} = 0$  for all possible  $j$ .

```

LL mod = 1000000007;
vector<LL> berlekamp(const vector<LL>& a) {
    vector<LL> p = {1}, r = {1};

```

```

be8e
d2b5
6b4a
e582
6112
a905
657b
95cf
d480
0305
8dad
aad8
be4d
d080
f156
427e
4ecd
865b
c36a
dd36
1b23
f8f3
e97f
c449
820b
f039
95cf
95cf
427e
37e1
50dc
95cf
f27f
95cf

```

```

6e50
97db
8904

```

```

075b     LL dif = 1;
8bc9     rep (i, a.size()) {
1b35         LL u = 0;
bd0b         rep (j, p.size()) u = (u + p[j] * a[i-j]) % mod;
eae9         if (u == 0) {
b14c             r.insert(r.begin(), 0);
8e2e         } else {
0c78             auto op = p;
02f6             p.resize(max(p.size(), r.size() + 1));
0a2e             LL idif = powmod(dif, mod - 2);
9b57             rep (j, r.size())
dacc                 p[j+1] = (p[j+1] - r[j] * idif % mod * u % mod + mod) % mod;
bcd1             dif = u; r = op;
95cf         }
95cf     }
e149     return p;
95cf }

```

## 4.5 Fast Walsh-Hadamard transform

```

061e void fwt(int* a, int n){
5595     for (int d = 1; d < n; d <= 1)
05f2         for (int i = 0; i < n; i += d << 1)
b833             rep (j, d){
7796                 int x = a[i+j], y = a[i+j+d];
427e                 // a[i+j] = x+y, a[i+j+d] = x-y; // xor
427e                 // a[i+j] = x+y; // and
427e                 // a[i+j+d] = x+y; // or
95cf             }
95cf }
427e
4db1 void ifwt(int* a, int n){
5595     for (int d = 1; d < n; d <= 1)
05f2         for (int i = 0; i < n; i += d << 1)
b833             rep (j, d){
7796                 int x = a[i+j], y = a[i+j+d];
427e                 // a[i+j] = (x+y)/2, a[i+j+d] = (x-y)/2; // xor
427e                 // a[i+j] = x-y; // and
427e                 // a[i+j+d] = y-x; // or
95cf             }
95cf }
427e

```

```

void conv(int* a, int* b, int n){
    fwt(a, n);
    fwt(b, n);
    rep(i, n) a[i] *= b[i];
    ifwt(a, n);
}

```

## 4.6 Fast fourier transform

```

const int NMAX = 1<<20;

typedef complex<double> cplx;

const double PI = 2*acos(0.0);

struct FFT{
    int rev[NMAX];
    cplx omega[NMAX], oinv[NMAX];
    int K, N;

    FFT(int k){
        K = k; N = 1 << k;
        rep (i, N){
            rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
            omega[i] = polar(1.0, 2.0 * PI / N * i);
            oinv[i] = conj(omega[i]);
        }
    }

    void dft(cplx* a, cplx* w){
        rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
        for (int l = 2; l <= N; l *= 2){
            int m = l/2;
            for (cplx* p = a; p != a + N; p += l)
                rep (k, m){
                    cplx t = w[N/l*k] * p[k+m];
                    p[k+m] = p[k] - t; p[k] += t;
                }
        }
    }

    void fft(cplx* a){dft(a, omega);}
    void ifft(cplx* a){

```

2ab6  
950a  
e427  
8a42  
430f  
95cf

4e09  
427e  
3fbf  
427e  
abd1  
12af  
c47c  
27d7  
9827  
427e  
1442  
e209  
b393  
7ba3  
1908  
a166  
95cf  
95cf  
427e  
b941  
a215  
ac6e  
2969  
b3cf  
c24f  
fe06  
ecbf  
95cf  
95cf  
95cf  
427e  
617b  
a123

```

3b2f     dft(a, oinv);
57fc     rep (i, N) a[i] /= N;
95cf     }
427e
bdc0     void conv(cplx* a, cplx* b){
6497         fft(a); fft(b);
12a5         rep (i, N) a[i] *= b[i];
f84e         fft(a);
95cf     }
329b };

```

## 4.7 Number theoretic transform

```

4ab9     const int NMAX = 1<<21;
427e
427e     // 998244353 = 7*17*2^23+1, G = 3
fb9a     const int P = 1004535809, G = 3; // = 479*2^21+1
427e
87ab     struct NTT{
c47c         int rev[NMAX];
0eda         LL omega[NMAX], oinv[NMAX];
81af         int g, g_inv; // g: g_n = G^((P-1)/n)
9827         int K, N;
427e
2a2c         LL powmod(LL b, LL e){
95a2             LL r = 1;
3e90             while (e){
6624                 if (e&1) r = r * b % P;
489e                 b = b * b % P;
16fc                 e >>= 1;
95cf             }
547e             return r;
95cf         }
427e
f420         NTT(int k){
e209             K = k; N = 1 << k;
7652             g = powmod(G, (P-1)/N);
4b3a             g_inv = powmod(g, N-1);
e04f             omega[0] = oinv[0] = 1;
b393             rep (i, N){
7ba3                 rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
ad4f                 if (i){

```

```

         omega[i] = omega[i-1] * g % P;
         oinv[i] = oinv[i-1] * g_inv % P;
     }
}

void _ntt(LL* a, LL* w){
    rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
    for (int l = 2; l <= N; l *= 2){
        int m = l/2;
        for (LL* p = a; p != a + N; p += l)
            rep (k, m){
                LL t = w[N/l*k] * p[k+m] % P;
                p[k+m] = (p[k] - t + P) % P;
                p[k] = (p[k] + t) % P;
            }
    }
}

void ntt(LL* a){_ntt(a, omega);}
void intt(LL* a){
    LL inv = powmod(N, P-2);
    _ntt(a, oinv);
    rep (i, N) a[i] = a[i] * inv % P;
}

void conv(LL* a, LL* b){
    ntt(a); ntt(b);
    rep (i, N) a[i] = a[i] * b[i] % P;
    intt(a);
}
};

```

## 4.8 Sieve of Euler

```

const int MAXX = 1e7+5;
bool p[MAXX];
int prime[MAXX], sz;

void sieve(){
    p[0] = p[1] = 1;
    for (int i = 2; i < MAXX; i++){

```

```

bf28         if (!p[i]) prime[sz++] = i;
e82c         for (int j = 0; j < sz && i*prime[j] < MAXX; j++){
b6a9             p[i*prime[j]] = 1;
5f51             if (i % prime[j] == 0) break;
95cf         }
95cf     }
95cf }

```

## 4.9 Sieve of Euler (General)

```

b62e namespace sieve {
6589     constexpr int MAXN = 10000007;
e982     bool p[MAXN]; // true if not prime
6ae8     int prime[MAXN], sz;
cbf7     int pval[MAXN], pcnt[MAXN];
6030     int f[MAXN];

427e     void exec(int N = MAXN) {
9628         p[0] = p[1] = 1;
427e
8a8a         pval[1] = 1;
bdda         pcnt[1] = 0;
c6b9         f[1] = 1;
427e
a643         for (int i = 2; i < N; i++) {
01d6             if (!p[i]) {
b2b2                 prime[sz++] = i;
37d9                 for (LL j = i; j < N; j *= i) {
758c                     int b = j / i;
81fd                     pval[j] = i * pval[b];
e0f3                     pcnt[j] = pcnt[b] + 1;
a96c                     f[j] = _____; // f[j] = f(i^pcnt[j])
95cf                 }
95cf             }
34c0             for (int j = 0; i * prime[j] < N; j++) {
f87a                 int x = i * prime[j]; p[x] = 1;
20cc                 if (i % prime[j] == 0) {
9985                     pval[x] = pval[i] * prime[j];
3f93                     pcnt[x] = pcnt[i] + 1;
8e2e                 } else {
cc91                     pval[x] = prime[j];
6322                     pcnt[x] = 1;

```

```

    }
    if (x != pval[x]) {
        f[x] = f[x / pval[x]] * f[pval[x]]
    }
    if (i % prime[j] == 0) break;
    }
    }
}

```

```

95cf
6191
d614
95cf
5f51
95cf
95cf
95cf
95cf
95cf

```

## 4.10 Miller-Rabin primality test

The array `a[]` (excluding sentinel, i.e. `LLONG_MAX`) should be

{2}	when $n < 2,047$ .
{2, 7, 61}	when $n < 4,759,123,141 (2^{32})$ .
{2, 3, 5, 7, 11}	when $n < 2.1 \times 10^{12}$ .
{2, 325, 9375, 28178, 450775, 9780504, 1795265022}	when $n < 2^{64}$ .

```

bool test(LL n){
    if (n < 3) return n==2;
    // ! The array a[] should be modified if the range of x changes.
    const LL a[] = {2LL, 7LL, 61LL, LLONG_MAX};
    LL r = 0, d = n-1, x;
    while (~d & 1) d >>= 1, r++;
    for (int i=0; a[i] < n; i++){
        x = powmod(a[i], d, n); // ! powmod must use for 64bit mulmod
        if (x == 1 || x == n-1) goto next;
        rep (i, r) {
            x = mulmod(x, x, n);
            if (x == n-1) goto next;
        }
        return false;
    }
next:;
    }
    return true;
}

```

```

f16f
59f2
427e
3f11
c320
f410
2975
ece1
7f99
e257
d7ff
8d2e
95cf
438e
d490
95cf
3361
95cf

```

## 4.11 Integer factorization (Pollard's rho)

```

ULL gcd(ULL a, ULL b) {return b ? gcd(b, a % b) : a;}

```

```

2e6b

```

```

427e ULL PollardRho(ULL n){
54a5     ULL c, x, y, d = n;
45eb     if (~n&1) return 2;
d3e5     while (d == n){
3c69         x = y = 2;
0964         d = 1;
4753         c = rand() % (n - 1) + 1;
5952         while (d == 1){
9e5b             x = (mulmod(x, x, n) + c) % n;
33d5             y = (mulmod(y, y, n) + c) % n;
e1bf             y = (mulmod(y, y, n) + c) % n;
e1bf             d = gcd(x>y ? x-y : y-x, n);
a313         }
95cf     }
95cf     return d;
5d89 }
95cf

```

## 5 Graph Theory

### 5.1 Strongly connected component

```

837c const int MAXV = 100005;
427e
2ea0 struct graph{
88e3     vector<int> adj[MAXV];
9cad     stack<int> s;
3d02     int V; // number of vertices
8b6c     int pre[MAXV], lnk[MAXV], scc[MAXV];
27ee     int time, sccn;
427e
bfab     void add_edge(int u, int v){
c71a         adj[u].push_back(v);
95cf     }
427e
d714     void dfs(int u){
7e41         pre[u] = lnk[u] = ++time;
80f6         s.push(u);
18f6         for (int v : adj[u]){
173e             if (!pre[v]){
5f3c                 dfs(v);

```

```

        lnk[u] = min(lnk[u], lnk[v]);
    } else if (!scc[v]){
        lnk[u] = min(lnk[u], pre[v]);
    }
}
if (lnk[u] == pre[u]){
    sccn++;
    int x;
    do {
        x = s.top(); s.pop();
        scc[x] = sccn;
    } while (x != u);
}
}

void find_scc(){
    time = sccn = 0;
    memset(scc, 0, sizeof scc);
    memset(pre, 0, sizeof pre);
    Rep (i, V){
        if (!pre[i]) dfs(i);
    }
}

vector<int> adjc[MAXV];
void contract(){
    Rep (i, V)
        rep (j, adj[i].size()){
            if (scc[i] != scc[adj[i][j]])
                adjc[scc[i]].push_back(scc[adj[i][j]]);
        }
}
};

```

002c  
6068  
d5df  
95cf  
95cf  
8de2  
660f  
3c9e  
a69f  
3834  
b0e9  
6757  
95cf  
95cf  
427e  
4c88  
f4a2  
8de7  
8c2f  
6901  
56d1  
95cf  
95cf  
427e  
27ce  
364d  
1a1e  
21a2  
b730  
b46e  
95cf  
95cf  
329b

### 5.2 Vertex biconnected component

```

const int MAXN = 100005;
struct graph {
    int pre[MAXN], iscut[MAXN], bccno[MAXN], dfs_clock, bcc_cnt;
    vector<int> adj[MAXN], bcc[MAXN];
    set<pair<int, int>> bcce[MAXN];

```

0f42  
2ea0  
33ae  
848f  
6b06  
427e

```

76f7     stack<pair<int, int>> s;
427e
bfab     void add_edge(int u, int v) {
c71a         adj[u].push_back(v);
a717         adj[v].push_back(u);
95cf     }
427e
7d3c     int dfs(int u, int fa) {
9fe6         int lowu = pre[u] = ++dfs_clock;
ec14         int child = 0;
18f6         for (int v : adj[u]) {
173e             if (!pre[v]) {
e7f8                 s.push({u, v});
fdcf                 child++;
f851                 int lowv = dfs(v, u);
189c                 lowu = min(lowu, lowv);
b687                 if (lowv >= pre[u]) {
6323                     iscut[u] = 1;
57eb                     bcc[bcc_cnt].clear();
90b8                     bcce[bcc_cnt].clear();
a147                     while (1) {
a6a3                         int xu, xv;
a0c3                         tie(xu, xv) = s.top(); s.pop();
0ef5                         bcce[bcc_cnt].insert({min(xu, xv), max(xu, xv)});
3db2                         if (bccno[xu] != bcc_cnt) {
e0db                             bcc[bcc_cnt].push_back(xu);
d27f                             bccno[xu] = bcc_cnt;
95cf                         }
f357                         if (bccno[xv] != bcc_cnt) {
752b                             bcc[bcc_cnt].push_back(xv);
57c9                             bccno[xv] = bcc_cnt;
95cf                         }
7096                         if (xu == u && xv == v) break;
95cf                     }
03f5                     bcc_cnt++;
95cf                 }
7470             } else if (pre[v] < pre[u] && v != fa) {
e7f8                 s.push({u, v});
f115                 lowu = min(lowu, pre[v]);
95cf             }
95cf         }
e104         if (fa < 0 && child == 1) iscut[u] = 0;
1160         return lowu;
95cf     }

```

```

void find_bcc(int n) {
    memset(pre, 0, sizeof pre);
    memset(iscut, 0, sizeof iscut);
    memset(bccno, -1, sizeof bccno);
    dfs_clock = bcc_cnt = 0;
    rep (i, n) if (!pre[i]) dfs(i, -1);
}
};

```

427e  
17be  
8c2f  
e2d2  
40d3  
fae2  
5c63  
95cf  
329b

### 5.3 Cut vertices

If the graph is unconnected, the algorithm should be run on each component. One may run Rep (i, n) if (!dfn[i]) tarjan(i, i) for unconnected graph.

#### Usage:

add_edge(u, v)	Add an undirected edge ( $u, v$ ).
tarjan(u, fa)	Run Tarjan's algorithm on tree rooted at fa. Please call with identical u and fa.
cut[v]	Whether $v$ is a cut vertex.

```

const int MAXN = 200005;
vector<int> adj[MAXN];
int dfn[MAXN], low[MAXN], idx;
bool cut[MAXN];

void add_edge(int u, int v) {
    adj[u].push_back(v);
    adj[v].push_back(u);
}

void tarjan(int u, int fa) {
    dfn[u] = low[u] = ++idx;
    int child = 0;
    for (int v : adj[u]) {
        if (!dfn[v]) {
            tarjan(v, fa); low[u] = min(low[u], low[v]);
            if (low[v] >= dfn[u] && u != fa) cut[u] = true;
            child += u == fa;
        }
        low[u] = min(low[u], dfn[v]);
    }
    if (u == fa && child > 1) cut[u] = true;
}

```

9f60  
0b32  
18e4  
d39d  
427e  
bfab  
c71a  
a717  
95cf  
427e  
50aa  
9891  
ec14  
18f6  
3c64  
9636  
f368  
7923  
95cf  
769a  
95cf  
7927  
95cf

## 5.4 Minimum spanning arborescence, faster

All vertices are 1-based. Clear the fields when reuse the struct.

### Usage:

`add_edge(u, v, w)` Add an edge from  $u$  to  $v$  with weight  $w$ .  
`run(n, rt)` Compute the total weight of MSA rooted at  $rt$ . If not exist, return `LLONG_MIN`.

**Time Complexity:**  $O((|E| + |V| \log |V|) \log |V|)$

```
5ece const int MAXN = 300005;
2fef typedef pair<LL, int> pii;
1495 struct MDST {
01b2     priority_queue<pii, vector<pii>, greater<pii>> heap[MAXN];
321d     LL shift[MAXN];
fc06     int fa[MAXN], vis[MAXN];

427e
38dd     int find(int x) { return fa[x] == x ? x : fa[x] = find(fa[x]); }

427e
29b0     void unite(int x, int y) {
0c14         x = find(x); y = find(y); fa[y] = x; if (x == y) return;
6fa0         if (heap[x].size() < heap[y].size()) {
9c26             swap(heap[x], heap[y]);
2ffc             swap(shift[x], shift[y]);
95cf         }
9959         while (heap[y].size()) {
175b             auto p = heap[y].top(); heap[y].pop();
c0c5             heap[x].emplace(p.first - shift[y] + shift[x], p.second);
95cf         }
95cf     }

427e
0bbd     void add_edge(int u, int v, LL w) { heap[v].emplace(w, u); }

427e
a526     LL run(int n, int rt) {
f7ff         LL ans = 0;
81f2         iota(fa, fa + n + 1, 0);
19b3         Rep (i, n) if (find(i) != find(rt)) {
a7b1             int u = find(i);
010e             stack<int, vector<int>> s;
eff5             while (find(u) != find(rt)) {
0dda                 if (vis[u]) while (s.top() != u) {
c593                     vis[s.top()] = 0; unite(u, s.top()); s.pop();
83c4                 } else { vis[u] = 1; s.push(u); }
c76e                 while (heap[u].size()) {
b385                     ans += heap[u].top().first - shift[u];
```

```

shift[u] = heap[u].top().first;
if (find(heap[u].top().second) != u) break;
heap[u].pop();
}
if (heap[u].empty()) return LLONG_MIN;
u = find(heap[u].top().second);
}
while (s.size()) { vis[s.top()] = 0; unite(rt, s.top()); s.pop(); }
}
return ans;
}
};
```

dde2  
da47  
9fbb  
95cf  
6961  
87e6  
95cf  
2d46  
95cf  
4206  
95cf  
329b

## 5.5 Maximum flow (Dinic)

### Usage:

`add_edge(u, v, c)` Add an edge from  $u$  to  $v$  with capacity  $c$ .  
`max_flow(s, t)` Compute maximum flow from  $s$  to  $t$ .

**Time Complexity:** For general graph,  $O(V^2E)$ ; for network with unit capacity,  $O(\min\{V^{2/3}, \sqrt{E}\}E)$ ; for bipartite network,  $O(\sqrt{V}E)$ .

```
struct edge{
    int from, to;
    LL cap, flow;
};

const int MAXN = 1005;
struct Dinic {
    int n, m, s, t;
    vector<edge> edges;
    vector<int> G[MAXN];
    bool vis[MAXN];
    int d[MAXN];
    int cur[MAXN];

    void add_edge(int from, int to, LL cap) {
        edges.push_back(edge{from, to, cap, 0});
        edges.push_back(edge{to, from, 0, 0});
        m = edges.size();
        G[from].push_back(m-2);
        G[to].push_back(m-1);
    }
```

bcf8  
60e2  
5e6d  
329b  
427e  
e2cd  
9062  
4dbf  
9f0c  
b891  
bbb6  
b40a  
ddcc  
427e  
5973  
7b55  
1db7  
fe77  
dff5  
8f2d  
95cf  
427e



```

1836 bool bfs() {
3b73     memset(vis, 0, sizeof(vis));
93d2     queue<int> q;
5d13     q.push(s);
2cd2     vis[s] = 1;
721d     d[s] = 0;
cc78     while (!q.empty()) {
66ba         int x = q.front(); q.pop();
3b61         for (int i = 0; i < G[x].size(); i++) {
b510             edge& e = edges[G[x][i]];
bba9             if (!vis[e.to] && e.cap > e.flow) {
cd72                 vis[e.to] = 1;
cf26                 d[e.to] = d[x] + 1;
ca93                 q.push(e.to);
95cf             }
95cf         }
95cf     }
b23b     return vis[t];
95cf }
427e
9252 LL dfs(int x, LL a) {
6904     if (x == t || a == 0) return a;
8bf9     LL flow = 0, f;
f515     for (int& i = cur[x]; i < G[x].size(); i++) {
b510         edge& e = edges[G[x][i]];
2374         if(d[x] + 1 == d[e.to] && (f = dfs(e.to, min(a, e.cap-e.flow))) > 0)
        {
1cce             e.flow += f;
e16d             edges[G[x][i]^1].flow -= f;
a74d             flow += f;
23e5             a -= f;
97ed             if(a == 0) break;
95cf         }
95cf     }
84fb     return flow;
95cf }
427e
5bf2 LL max_flow(int s, int t) {
590d     this->s = s; this->t = t;
62e2     LL flow = 0;
ed58     while (bfs()) {
f326         memset(cur, 0, sizeof(cur));
fb3a         flow += dfs(s, LLONG_MAX);
95cf     }

```

```

        return flow;
    }

    vector<int> min_cut() { // call this after maxflow
        vector<int> ans;
        for (int i = 0; i < edges.size(); i++) {
            edge& e = edges[i];
            if(vis[e.from] && !vis[e.to] && e.cap > 0) ans.push_back(i);
        }
        return ans;
    }
};

```

## 5.6 Maximum cardinality bipartite matching (Hungarian)

```

#include <bits/stdc++.h>
using namespace std;

#define rep(i, n) for (int i = 0; i < (n); i++)
#define Rep(i, n) for (int i = 1; i <= (n); i++)
#define range(x) (x).begin(), (x).end()
typedef long long LL;

struct Hungarian{
    int nx, ny;
    vector<int> mx, my;
    vector<vector<int>> > e;
    vector<bool> mark;

    void init(int nx, int ny){
        this->nx = nx;
        this->ny = ny;
        mx.resize(nx); my.resize(ny);
        e.clear(); e.resize(nx);
        mark.resize(nx);
    }

    inline void add(int a, int b){
        e[a].push_back(b);
    }

    bool augment(int i){

```

```

207c     if (!mark[i]) {
dae4         mark[i] = true;
6a1e         for (int j : e[i]){
0892             if (my[j] == -1 || augment(my[j])){
9ca3                 mx[i] = j; my[j] = i;
3361                 return true;
95cf             }
95cf         }
95cf     }
438e     return false;
95cf }
427e
3fac int match(){
5b57     int ret = 0;
b0f1     fill(range(mx), -1);
b957     fill(range(my), -1);
4ed1     rep (i, nx){
13a5         fill(range(mark), false);
cc89         if (augment(i)) ret++;
95cf     }
ee0f     return ret;
95cf }
329b };

```

## 5.7 Maximum matching of general graph (Edmond's blossom)

### Usage:

init(n)	Initialize the template with $n$ vertices, numbered from 1.
add_edge(u, v)	Add an undirected edge $uv$ .
solve()	Find the maximum matching. Return the number of matched edges.
mate[]	The mate of a matched vertex. If it is not matched, then the value is 0.

**Time Complexity:**  $O(|V|^3)$ , but extremely fast in practice.

```

c041 const int MAXN = 1024;
6ab1 struct Blossom {
0b32     vector<int> adj[MAXN];
93d2     queue<int> q;
5c83     int n;
0de2     int label[MAXN], mate[MAXN], save[MAXN], used[MAXN];
427e

```

```

void init(int nv) {
    n = nv; for (auto& v : adj) v.clear();
    fill(range(label), 0); fill(range(mate), 0);
    fill(range(save), 0); fill(range(used), 0);
}

void add_edge(int u, int v) { adj[u].push_back(v); adj[v].push_back(u); }

void rematch(int x, int y) {
    int m = mate[x]; mate[x] = y;
    if (mate[m] == x) {
        if (label[x] <= n) {
            mate[m] = label[x]; rematch(label[x], m);
        } else {
            int a = 1 + (label[x] - n - 1) / n;
            int b = 1 + (label[x] - n - 1) % n;
            rematch(a, b); rematch(b, a);
        }
    }
}

void traverse(int x) {
    Rep (i, n) save[i] = mate[i];
    rematch(x, x);
    Rep (i, n) {
        if (mate[i] != save[i]) used[i] ++;
        mate[i] = save[i];
    }
}

void relabel(int x, int y) {
    Rep (i, n) used[i] = 0;
    traverse(x); traverse(y);
    Rep (i, n) {
        if (used[i] == 1 and label[i] < 0) {
            label[i] = n + x + (y - 1) * n;
            q.push(i);
        }
    }
}

int solve() {
    Rep (i, n) {
        if (mate[i]) continue;

```

```

2186
3728
477d
bb35
95cf
427e
c2dd
427e
2a48
8af8
1aa4
f4ba
740a
8e2e
3341
2885
ef33
95cf
95cf
95cf
427e
8a50
43c0
2ef7
34d7
62c5
97ef
95cf
95cf
427e
8bf8
d101
c4ea
34d7
dee9
1c22
eb31
95cf
95cf
95cf
427e
a0ce
34d7
a073

```

```

1fc0     Rep (j, n) label[j] = -1;
7676     label[i] = 0; q = queue<int>(); q.push(i);
1c7d     while (q.size()) {
66ba         int x = q.front(); q.pop();
b98c         for (int y : adj[x]) {
c07f             if (mate[y] == 0 and i != y) {
7f36                 mate[y] = x; rematch(x, y); q = queue<int>(); break;
95cf             }
d315             if (label[y] >= 0) { relabel(x, y); continue; }
58ec             if (label[mate[y]] < 0) {
c9c4                 label[mate[y]] = x; q.push(mate[y]);
95cf             }
95cf         }
95cf     }
8abb     int cnt = 0;
b52f     Rep (i, n) cnt += (mate[i] > i);
6808     return cnt;
95cf }
329b };

```

## 5.8 Minimum cost maximum flow

```

bcf8 struct edge{
60e2     int from, to;
d698     int cap, flow;
32cc     LL cost;
329b };
427e
cc3e const LL INF = LLONG_MAX / 2;
2aa8 const int MAXN = 5005;
c6cb struct MCMF {
9ceb     int s, t, n, m;
9f0c     vector<edge> edges;
b891     vector<int> G[MAXN];
f74f     bool inq[MAXN]; // queue
8f67     LL d[MAXN];    // distance
9524     int p[MAXN];    // previous
b330     int a[MAXN];    // improvement
427e
f7f2     void add_edge(int from, int to, int cap, LL cost) {
24f0         edges.push_back(edge{from, to, cap, 0, cost});

```

```

edges.push_back(edge{to, from, 0, 0, -cost});
m = edges.size();
G[from].push_back(m-2);
G[to].push_back(m-1);
}

bool spfa(){
    queue<int> q;
    fill(d, d + MAXN, INF); d[s] = 0;
    memset(inq, 0, sizeof(inq));
    q.push(s); inq[s] = true;
    p[s] = 0; a[s] = INT_MAX;
    while (!q.empty()){
        int u = q.front(); q.pop(); inq[u] = false;
        for (int i : G[u]) {
            edge& e = edges[i];
            if (e.cap > e.flow && d[e.to] > d[u] + e.cost){
                d[e.to] = d[u] + e.cost;
                p[e.to] = G[u][i];
                a[e.to] = min(a[u], e.cap - e.flow);
                if (!inq[e.to]) q.push(e.to), inq[e.to] = true;
            }
        }
    }
    return d[t] != INF;
}

void augment(){
    int u = t;
    while (u != s){
        edges[p[u]].flow += a[t];
        edges[p[u]^1].flow -= a[t];
        u = edges[p[u]].from;
    }
}

#ifdef GIVEN_FLOW
bool min_cost(int s, int t, int f, LL& cost) {
    this->s = s; this->t = t;
    int flow = 0;
    cost = 0;
    while (spfa()) {
        augment();
        if (flow + a[t] >= f){

```

```

95f0
fe77
dff5
8f2d
95cf
427e
3c52
93d2
8494
fd48
5e7c
2dae
cc78
b0aa
3bba
56d8
3601
55bc
0bea
8249
e5d3
95cf
95cf
95cf
6d7c
95cf
427e
71a4
06f1
b19d
db09
25a9
e6c9
95cf
95cf
427e
6e20
5972
590d
21d4
23cb
22dc
bcd b
a671

```

```

b14d         cost += (f - flow) * d[t]; flow = f;
3361         return true;
8e2e     } else {
2a83         flow += a[t]; cost += a[t] * d[t];
95cf     }
95cf     }
438e     return false;
95cf }
a8cb #else
f9a9     int min_cost(int s, int t, LL& cost) {
590d         this->s = s; this->t = t;
21d4         int flow = 0;
23cb         cost = 0;
22dc         while (spfa()) {
bcd8             augment();
2a83             flow += a[t]; cost += a[t] * d[t];
95cf         }
84fb         return flow;
95cf     }
1937 #endif
329b };

```

## 5.9 Fast LCA

All indices of the tree are 1-based.

### Usage:

preprocess(root)      Initialize with tree rooted at root.  
lca(u, v)      Query the lowest common ancestor of  $u$  and  $v$ .

```

0e34 const int MAXN = 500005;
0b32 vector<int> adj[MAXN];
fccb int id[MAXN], nid;
1356 pair<int, int> st[MAXN << 1][33 - __builtin_clz(MAXN)];
427e
e16d void dfs(int u, int p, int d) {
0df2     st[id[u] = nid++][0] = {d, u};
18f6     for (int v : adj[u]) {
bd87         if (v == p) continue;
f58c         dfs(v, u, d + 1);
08ad         st[nid++][0] = {d, u};
95cf     }
95cf }
427e

```

```

void preprocess(int root) {
    nid = 0;
    dfs(root, 0, 1);
    int l = 31 - __builtin_clz(nid);
    rep(j, l) rep(i, 1+nid-(1<<j))
        st[i][j+1] = min(st[i][j], st[i+(1<<j)][j]);
}

int lca(int u, int v) {
    tie(u, v) = minmax(id[u], id[v]);
    int k = 31 - __builtin_clz(v-u+1);
    return min(st[u][k], st[v-(1<<k)+1][k]).second;
}

```

## 5.10 Heavy-light decomposition

**Time Complexity:** The decomposition itself takes linear time. Each query takes  $O(\log n)$  operations.

```

const int MAXN = 100005;
vector<int> adj[MAXN];
int sz[MAXN], top[MAXN], fa[MAXN], son[MAXN], depth[MAXN], id[MAXN];

void dfs1(int x, int dep, int par){
    depth[x] = dep;
    sz[x] = 1;
    fa[x] = par;
    int maxn = 0, s = 0;
    for (int c: adj[x]){
        if (c == par) continue;
        dfs1(c, dep + 1, x);
        sz[x] += sz[c];
        if (sz[c] > maxn){
            maxn = sz[c];
            s = c;
        }
    }
    son[x] = s;
}

int cid = 0;
void dfs2(int x, int t){
    top[x] = t;
}

```

```

d314     id[x] = ++cid;
c4a1     if (son[x]) dfs2(son[x], t);
c861     for (int c: adj[x]){
9881         if (c == fa[x]) continue;
5518         if (c == son[x]) continue;
13f9         else dfs2(c, c);
95cf     }
95cf }
427e
0f04 void decomp(int root){
9fa4     dfs1(root, 1, 0);
1c88     dfs2(root, root);
95cf }
427e
2c98 void query(int u, int v){
03a1     while (top[u] != top[v]){
45ec         if (depth[top[u]] < depth[top[v]]) swap(u, v);
427e         // id[top[u]] to id[u]
005b         u = fa[top[u]];
95cf     }
6083     if (depth[u] > depth[v]) swap(u, v);
427e     // id[u] to id[v]
95cf }

```

## 5.11 Centroid decomposition

Note that the centroid here is not the exact centroid of the graph. It only guarantees that the size of each subtree does not exceed half of that of the original tree. This is enough to guarantee the correct time complexity. All vertices are numbered from 1. Call `decomp(root)` to use.

### Usage:

`decomp(u, p)` Decompose the tree rooted at  $u$  with parent  $p$ .

**Time Complexity:** The decomposition itself takes  $O(n \log n)$  time.

```

1fb6 vector<int> adj[100005];
88e0 int sz[100005], sum;
427e
f93d void getsz(int u, int p) {
5b36     sz[u] = 1; sum++;
18f6     for (int v : adj[u]) {
bd87         if (v == p) continue;
e3cb         getsz(v, u);
8449         sz[u] += sz[v];

```

```

    }
}

int getcent(int u, int p) {
    for (int v : adj[u])
        if (v != p and sz[v] > sum / 2)
            return getcent(v, u);
    return u;
}

void decompose(int u) {
    sum = 0; getsz(u, 0);
    u = getcent(u, 0); // update u to the centroid

    for (int v : adj[u]) {
        // get answer for subtree v
    }
    // get answer for the whole tree
    // don't forget to count the centroid itself

    for (int v : adj[u]) { // divide and conquer
        adj[v].erase(find(range(adj[v]), u));
        decompose(v);
        adj[v].push_back(u); // restore deleted edge
    }
}

```

## 5.12 DSU on tree

This implementation avoids parallel existence of multiple data structures but requires that the data structure is invertible. To use this template, implement `merge`, `enter`, `leave` as needed; first call `decomp(root, 0)`, then call `work(root, 0, false)`. Labels of vertices start from 1.

### Usage:

`decomp(u, p)` Decompose the tree  $u$ .  
`work(u, p, keep)` Work for subtree  $u$ . When `keep` is set, information is not cleared.

**Time Complexity:**  $O(n \log n)$  times the complexity for `merge`, `enter`, `leave`.

```

vector<int> adj[100005];
int sz[100005], son[100005];

```

```

5559 void decomp(int u, int p) {
50c0     sz[u] = 1;
18f6     for (int v : adj[u]) {
bd87         if (v == p) continue;
a851         decomp(v, u);
8449         sz[u] += sz[v];
d28c         if (sz[v] > sz[son[u]]) son[u] = v;
95cf     }
95cf }
427e
b7ec template <typename T>
62f5 void trav(T fn, int u, int p) {
4412     fn(u);
30b3     for (int v : adj[u]) if (v != p) trav(fn, v, u);
95cf }
427e
7467 #define for_light(v) for (int v : adj[u]) if (v != p and v != son[u])
33ff void work(int u, int p, bool keep) {
72a2     for_light(v) work(v, u, 0); // process light children
427e
427e     // process heavy child
427e     // current data structure contains info of heavy child
9866     if (son[u]) work(son[u], u, 1);
427e
18a9     auto merge = [u] (int c) { /* count contribution of c */ };
1ab0     auto enter = [] (int c) { /* add vertex c */ };
f241     auto leave = [] (int c) { /* remove vertex c */ };
427e
3d3b     for_light(v) {
74c6         trav(merge, v, u);
c13d         trav(enter, v, u);
95cf     }
427e
427e     // count answer for root and add it
427e     // Warning: special check may apply to root!
c54f     merge(u);
9dec     enter(u);
427e
427e     // Leave current tree
4e3e     if (!keep) trav(leave, u, p);
95cf }

```

## 6 Data Structures

### 6.1 Fenwick tree (point update range query)

```

struct bit_purq { // point update, range query
    int N;
    vector<LL> tr;

    void init(int n) { tr.resize(N = n + 5); }

    LL sum(int n) {
        LL ans = 0;
        while (n) { ans += tr[n]; n &= n - 1; }
        return ans;
    }

    void add(int n, LL x){
        while (n < N) { tr[n] += x; n += n & -n; }
    }
};

```

9976  
d7af  
99ff  
427e  
456d  
427e  
63d0  
f7ff  
6770  
4206  
95cf  
427e  
f4bd  
968e  
95cf  
329b

### 6.2 Fenwick tree (range update point query)

```

struct bit_rupq{ // range update, point query
    int N;
    vector<LL> tr;

    void init(int n) { tr.resize(N = n + 5);}

    LL query(int n) {
        LL ans = 0;
        while (n < N) { ans += tr[n]; n += n & -n; }
        return ans;
    }

    void add(int n, LL x) {
        while (n) { tr[n] += x; n &= n - 1; }
    }
};

```

3d03  
d7af  
99ff  
427e  
456d  
427e  
38d4  
f7ff  
3667  
4206  
95cf  
427e  
f4bd  
0a2b  
95cf  
329b

### 6.3 Segment tree

```

3942 LL p;
1ebb const int MAXN = 4 * 100006;
451a struct segtree {
27be     int l[MAXN], m[MAXN], r[MAXN];
4510     LL val[MAXN], tadd[MAXN], tmul[MAXN];
427e
ac35 #define lson (o<<1)
1294 #define rson (o<<1|1)
427e
1344     void pull(int o) {
bbe9         val[o] = (val[lson] + val[rson]) % p;
95cf     }
427e
e4bc     void push_add(int o, LL x) {
5dd6         val[o] = (val[o] + x * (r[o] - l[o])) % p;
6eff         tadd[o] = (tadd[o] + x) % p;
95cf     }
427e
d658     void push_mul(int o, LL x) {
b82c         val[o] = val[o] * x % p;
aa86         tadd[o] = tadd[o] * x % p;
649f         tmul[o] = tmul[o] * x % p;
95cf     }
427e
b149     void push(int o) {
3159         if (l[o] == m[o]) return;
0a90         if (tmul[o] != 1) {
0f4a             push_mul(lson, tmul[o]);
045e             push_mul(rson, tmul[o]);
ac0a             tmul[o] = 1;
95cf         }
1b82         if (tadd[o]) {
9547             push_add(lson, tadd[o]);
0e73             push_add(rson, tadd[o]);
6234             tadd[o] = 0;
95cf         }
95cf     }
427e
471c     void build(int o, int ll, int rr) {
0e87         int mm = (ll + rr) / 2;
9d27         l[o] = ll; r[o] = rr; m[o] = mm;

```

```

tmul[o] = 1;
if (ll == mm) {
    scanf("%lld", val + o);
    val[o] %= p;
} else {
    build(lson, ll, mm);
    build(rson, mm, rr);
    pull(o);
}
}

void add(int o, int ll, int rr, LL x) {
    if (ll <= l[o] && r[o] <= rr) {
        push_add(o, x);
    } else {
        push(o);
        if (m[o] > ll) add(lson, ll, rr, x);
        if (m[o] < rr) add(rson, ll, rr, x);
        pull(o);
    }
}

void mul(int o, int ll, int rr, LL x) {
    if (ll <= l[o] && r[o] <= rr) {
        push_mul(o, x);
    } else {
        push(o);
        if (ll < m[o]) mul(lson, ll, rr, x);
        if (m[o] < rr) mul(rson, ll, rr, x);
        pull(o);
    }
}

LL query(int o, int ll, int rr) {
    if (ll <= l[o] && r[o] <= rr) {
        return val[o];
    } else {
        push(o);
        if (rr <= m[o]) return query(lson, ll, rr);
        if (ll >= m[o]) return query(rson, ll, rr);
        return query(lson, ll, rr) + query(rson, ll, rr);
    }
}
} seg;

```

```

ac0a
5c92
001f
e5b6
8e2e
7293
5e67
ba26
95cf
95cf
427e
4406
3c16
db32
8e2e
c4b0
4305
d5a6
ba26
95cf
95cf
427e
48cd
3c16
e7d0
8e2e
c4b0
d1ba
67f3
ba26
95cf
95cf
427e
0f62
3c16
6dfe
8e2e
c4b0
462a
5cca
bbf9
95cf
95cf
4d99

```

## 6.4 Treap

Self-balanced binary search tree which supports split and merge.

### Usage:

push(x)	Push lazy tags to children.
pull(x)	Update statistics of node $x$ .
Init(x, v)	Initialize node $x$ with value $v$ .
Add(x, v)	Apply addition to subtree $x$ .
Reverse(x)	Apply reversion to subtree $x$ .
Merge(x, y)	Merge trees rooted at $x$ and $y$ . Return the root of new tree.
Split(t, k, x, y)	Split out the left $k$ elements of tree $t$ . The roots of left part and right part are stored in $x$ and $y$ , respectively.
init(n)	Initialize the treap with array of size $n$ .
work(op, l, r)	Range operation over $[l, r)$ .

**Time Complexity:** Expected  $O(\log n)$  per operation.

```

9f60 const int MAXN = 200005;
a7c5 mt19937 gen(time(NULL));
9542 struct Treap {
6d61     int ch[MAXN][2];
3948     int sz[MAXN], key[MAXN], val[MAXN];
5d9a     int add[MAXN], rev[MAXN];
2b1b     LL sum[MAXN] = {0};
a773     int maxv[MAXN] = {INT_MIN}, minv[MAXN] = {INT_MAX};
427e
a629     void Init(int x, int v) {
5a00         ch[x][0] = ch[x][1] = 0;
d8cd         key[x] = gen(); val[x] = v; pull(x);
95cf     }
427e
3bf9     void pull(int x) {
e1c3         sz[x] = 1 + sz[ch[x][0]] + sz[ch[x][1]];
99f8         sum[x] = val[x] + sum[ch[x][0]] + sum[ch[x][1]];
94e9         maxv[x] = max({val[x], maxv[ch[x][0]], maxv[ch[x][1]]});
6bb9         minv[x] = min({val[x], minv[ch[x][0]], minv[ch[x][1]]});
95cf     }
427e
8c8e     void Add(int x, int a) {
a7b1         val[x] += a; add[x] += a;
832a         sum[x] += LL(sz[x]) * a; maxv[x] += a; minv[x] += a;

```

```

}

void Reverse(int x) {
    rev[x] ^= 1;
    swap(ch[x][0], ch[x][1]);
}

void push(int x) {
    for (int c : ch[x]) if (c) {
        Add(c, add[x]);
        if (rev[x]) Reverse(c);
    }
    add[x] = 0; rev[x] = 0;
}

int Merge(int x, int y) {
    if (!x || !y) return x | y;
    push(x); push(y);
    if (key[x] > key[y]) {
        ch[x][1] = Merge(ch[x][1], y); pull(x); return x;
    } else {
        ch[y][0] = Merge(x, ch[y][0]); pull(y); return y;
    }
}

void Split(int t, int k, int &x, int &y) {
    if (t == 0) { x = y = 0; return; }
    push(t);
    if (sz[ch[t][0]] < k) {
        x = t; Split(ch[t][1], k - sz[ch[t][0]] - 1, ch[t][1], y);
    } else {
        y = t; Split(ch[t][0], k, x, ch[t][0]);
    }
    if (x) pull(x); if (y) pull(y);
}

} treap;

int root;

void init(int n) {
    Rep(i, n) {
        int x; scanf("%d", &x);
        treap.Init(i, x);
        root = (i == 1) ? 1 : treap.Merge(root, i);

```

95cf  
427e  
aaf6  
52c6  
7850  
95cf  
427e  
1a53  
5fe5  
fd76  
7a53  
95cf  
49ee  
95cf  
427e  
9d2c  
1b09  
cd7e  
bfffa  
a3df  
8e2e  
bf9e  
95cf  
95cf  
427e  
dc7e  
6303  
f26b  
3465  
ffd8  
8e2e  
8a23  
95cf  
89e3  
95cf  
b1f4  
427e  
24b6  
427e  
d34f  
34d7  
7681  
0ed8  
bcc8



```

95cf     }
95cf }
427e
d030 void work(int op, int l, int r) {
6639     int tl, tm, tr;
b6c4     treap.Split(root, l, tl, tm);
8de3     treap.Split(tm, r - 1, tm, tr);
3658     if (op == 1) {
c039         int x; scanf("%d", &x); treap.Add(tm, x);
1dcb     } else if (op == 2) {
ae78         treap.Reverse(tm);
581d     } else if (op == 3) {
e092         printf("%lld_%d_%d\n",
867f             treap.sum[tm], treap.minv[tm], treap.maxv[tm]);
95cf     }
6188     root = treap.Merge(treap.Merge(tl, tm), tr);
95cf }

```

## 6.5 Link/cut tree

Dynamic connectivity of undirected acyclic graph. Support single-vertex update, path aggregation and relative LCA query. Vertices are numbered from 1. Zero initialization is enough except for the statistic information.

### Usage:

<code>pull(x)</code>	Update statistics of node $x$ .
<code>Root(u)</code>	Get the root of tree where vertex $u$ is in.
<code>Link(u, v)</code>	Link two unconnected trees.
<code>Cut(u, v)</code>	Cut an existent edge.
<code>Query(u, v)</code>	Path aggregation.
<code>Update(u, x)</code>	Single point modification.
<code>LCA(u, v, root)</code>	Get the lowest common ancestor of $u$ and $v$ in tree rooted at root.

**Time Complexity:**  $O(\log n)$  per operation

```

2e73 const int MAXN = 1000005;
ca06 struct LCT {
6a6d     int fa[MAXN], ch[MAXN][2], val[MAXN], sum[MAXN];
c6e1     bool rev[MAXN];
427e
eba3     bool isroot(int x) { return ch[fa[x]][0] == x || ch[fa[x]][1] == x; }
f19f     void pull(int x) { sum[x] = val[x] ^ sum[ch[x][0]] ^ sum[ch[x][1]]; }
1c4d     void reverse(int x) { swap(ch[x][0], ch[x][1]); rev[x] ^= 1; }

```

```

void push(int x) {
    if (rev[x]) rep (i, 2) if (ch[x][i]) reverse(ch[x][i]); rev[x] = 0;
}
void rotate(int x) {
    int y = fa[x], z = fa[y], k = ch[y][1] == x, w = ch[x][!k];
    if (isroot(y)) ch[z][ch[z][1] == y] = x;
    ch[x][!k] = y; ch[y][k] = w; if (w) fa[w] = y;
    fa[y] = x; fa[x] = z; pull(y);
}
void pushall(int x) { if (isroot(x)) pushall(fa[x]); push(x); }
void splay(int x) {
    int y = x, z = 0;
    for (pushall(y); isroot(x); rotate(x)) {
        y = fa[x]; z = fa[y];
        if (isroot(y)) rotate((ch[y][0] == x) ^ (ch[z][0] == y) ? x : y);
    }
    pull(x);
}
void access(int x) {
    int z = x;
    for (int y = 0; x; x = fa[y = x]) { splay(x); ch[x][1] = y; pull(x); }
    splay(z);
}
void chroot(int x) { access(x); reverse(x); }
void split(int x, int y) { chroot(x); access(y); }

int Root(int x) {
    for (access(x); ch[x][0]; x = ch[x][0]) push(x);
    splay(x); return x;
}
void Link(int u, int v) { chroot(u); fa[u] = v; }
void Cut(int u, int v) { split(u, v); fa[u] = ch[v][0] = 0; pull(v); }
int Query(int u, int v) { split(u, v); return sum[v]; }
void Update(int u, int x) { splay(u); val[u] = x; }
int LCA(int x, int y, int root) {
    chroot(root); access(x); splay(y);
    while (fa[y]) splay(y = fa[y]);
    return y;
}
};

```

1a53  
89a0  
95cf  
425f  
51af  
e1fe  
1e6f  
6d09  
95cf  
52c6  
f69c  
d095  
c494  
ceef  
4449  
95cf  
78a0  
95cf  
6229  
1548  
8854  
7afd  
95cf  
a067  
126d  
427e  
d87a  
f4f1  
0d77  
95cf  
9e46  
7c10  
0691  
a999  
1f42  
6cb2  
02e5  
c218  
95cf  
329b

## 6.6 Balanced binary search tree from pb\_ds

```

0475 #include <ext/pb_ds/assoc_container.hpp>
332d using namespace __gnu_pbds;
427e
43a7 tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update>
    rkt;
427e // null_tree_node_update
427e
427e // SAMPLE USAGE
190e rkt.insert(x); // insert element
05d4 rkt.erase(x); // erase element
add5 rkt.order_of_key(x); // obtain the number of elements less than x
b064 rkt.find_by_order(i); // iterator to i-th (numbered from 0) smallest element
c103 rkt.lower_bound(x);
4ff4 rkt.upper_bound(x);
b19b rkt.join(rkt2); // merge tree (only if their ranges do not intersect)
cb47 rkt.split(x, rkt2); // split all elements greater than x to rkt2

```

## 6.7 Persistent segment tree, range k-th query

```

f1a7 struct node {
2ff6     static int n, pos;
427e
7cec     int value;
70e2     node *left, *right;
427e
20b0     void* operator new(size_t size);
427e
3dc0     static node* Build(int l, int r) {
b6c5         node* a = new node;
ce96         if (r > l + 1) {
181e             int mid = (l + r) / 2;
3ba2             a->left = Build(l, mid);
8aaf             a->right = Build(mid, r);
8e2e         } else {
bfc4             a->value = 0;
95cf         }
5ffd         return a;
95cf     }
427e
5a45     static node* init(int size) {
2c46         n = size;

```

```

pos = 0;
return Build(0, n);
}

static int Query(node* lt, node *rt, int l, int r, int k) {
    if (r == l + 1) return l;
    int mid = (l + r) / 2;
    if (rt->left->value - lt->left->value < k) {
        k -= rt->left->value - lt->left->value;
        return Query(lt->right, rt->right, mid, r, k);
    } else {
        return Query(lt->left, rt->left, l, mid, k);
    }
}

static int query(node* lt, node *rt, int k) {
    return Query(lt, rt, 0, n, k);
}

node *Inc(int l, int r, int pos) const {
    node* a = new node(*this);
    if (r > l + 1) {
        int mid = (l + r) / 2;
        if (pos < mid)
            a->left = left->Inc(l, mid, pos);
        else
            a->right = right->Inc(mid, r, pos);
    }
    a->value++;
    return a;
}

node *inc(int index) {
    return Inc(0, n, index);
}
} nodes[8000000];

int node::n, node::pos;
inline void* node::operator new(size_t size) {
    return nodes + (pos++);
}

```

7ee3  
be52  
95cf  
427e  
93c0  
d30c  
181e  
cb5a  
8edb  
2412  
8e2e  
0119  
95cf  
95cf  
427e  
c9ad  
9e27  
95cf  
427e  
b19c  
5794  
ce96  
181e  
203d  
f44a  
649a  
1024  
95cf  
2b3e  
5ffd  
95cf  
427e  
e80f  
c246  
95cf  
865a  
427e  
99ce  
1987  
bb3c  
95cf

## 6.8 Block list

All indices are 0-based. All ranges are left-closed right-open.

### Usage:

<code>block::fix()</code>	Apply tags to the current block.
<code>Init(l, r)</code>	Range initializer.
<code>Reverse(l, r)</code>	Reverse the range.
<code>Add(l, r, x)</code>	Add $x$ to the range.
<code>Query(l, r)</code>	Range aggregation.

```
fd9e const int BLOCK = 800;
76b3 typedef vector<int> vi;
427e
a771 struct block {
8fbc     vi data;
e3b5     LL sum; int minv, maxv;
41db     int add; bool rev;
427e
d7eb     block(vi&& vec) : data(move(vec)),
1f0c         sum(accumulate(range(data), 0ll)),
8216         minv(*min_element(range(data))),
527d         maxv(*max_element(range(data))),
6437         add(0), rev(0) { }
427e
b919     void fix() {
0694         if (rev) reverse(range(data));         rev = 0;
0527         if (add) for (int& x : data) x += add;   add = 0;
95cf     }
427e
8bc4     void merge(block& another) {
b895         fix(); another.fix();
f516         vi temp(move(data));
d02c         temp.insert(temp.end(), range(another.data));
88ea         *this = block(move(temp));
95cf     }
427e
42e8     block split(int pos) {
3e79         fix();
ccab         block result(vi(data.begin() + pos, data.end()));
861a         data.resize(pos); *this = block(move(data));
56b0         return result;
95cf     }
329b };
427e
```

```
typedef list<block>::iterator lit;

struct blocklist {
    list<block> blk;

    void maintain() {
        lit it = blk.begin();
        while (it != blk.end() && next(it) != blk.end()) {
            lit it2 = it;
            while (next(it2) != blk.end() &&
                    it2->data.size() + next(it2)->data.size() <= BLOCK) {
                it2->merge(*next(it2));
                blk.erase(next(it2));
            }
            ++it;
        }
    }

    lit split(int pos) {
        for (lit it = blk.begin(); ; it++) {
            if (pos == 0) return it;
            while (it->data.size() > pos)
                blk.insert(next(it), it->split(pos));
            pos -= it->data.size();
        }
    }

    void Init(int *l, int *r) {
        for (int *cur = l; cur < r; cur += BLOCK)
            blk.emplace_back(vi(cur, min(cur + BLOCK, r)));
    }

    void Reverse(int l, int r) {
        lit it = split(l), it2 = split(r);
        reverse(it, it2);
        while (it != it2) {
            it->rev ^= 1;
            it++;
        }
        maintain();
    }

    void Add(int l, int r, int x) {
```

```
2a18
427e
ce14
5540
427e
7b8e
3131
4628
852d
188c
3600
93e1
e1fa
95cf
5771
95cf
95cf
427e
b7b3
2273
5502
8e85
2099
a5a1
427e
95cf
95cf
427e
1c7b
9919
8950
95cf
427e
a22f
997b
dfd0
8f89
6a06
5283
95cf
b204
95cf
427e
3cce
```

```

997b     lit it = split(l), it2 = split(r);
8f89     while (it != it2) {
e927         it->sum += LL(x) * it->data.size();
03d3         it->minv += x; it->maxv += x;
4511         it->add += x; it++;
95cf     }
b204     maintain();
95cf }
427e
3ad3 void Query(int l, int r) {
997b     lit it = split(l), it2 = split(r);
c33d     LL sum = 0; int minv = INT_MAX, maxv = INT_MIN;
8f89     while (it != it2) {
e472         sum += it->sum;
72c4         minv = min(minv, it->minv);
e1c4         maxv = max(maxv, it->maxv);
5283         it++;
95cf     }
b204     maintain();
8792     printf("%lld_%d_%d\n", sum, minv, maxv);
95cf }
958e } lst;

```

## 6.9 Persistent block list

Block list that supports persistence. All indices are 0-based. All ranges are left-closed right-open. `std::shared_ptr` is used to ease memory management. One should modify the constructor of `block` to maintain extra information. Here we use this policy that the size of each block does not exceed `BLOCK`, while the sum of sizes of two adjacent blocks does not less than `BLOCK`.

When some operation that breaks block list property, please call `maintain` in time to restore the property.

### Usage:

<code>maintain()</code>	Maintain the block list property.
<code>split(pos)</code>	Split the block list at position <code>pos</code> . Returns an iterator to a block starting at <code>pos</code> .
<code>sum(l, r)</code>	An example function of list traversal between $[l, r)$ .

**Time Complexity:** When `BLOCK` is properly selected, the time complexity is  $O(\sqrt{n})$  per operation.

```

a19e constexpr int BLOCK = 800;
76b3 typedef vector<int> vi;

```

```

typedef shared_ptr<vi> pvi;
typedef shared_ptr<const vi> pcvi;

struct block {
    pcvi data;
    LL sum;

    // add information to maintain
    block(pcvi ptr) :
        data(ptr),
        sum(accumulate(ptr->begin(), ptr->end(), 0ll))
    { }

    void merge(const block& another) {
        pvi temp = make_shared<vi>(data->begin(), data->end());
        temp->insert(temp->end(), another.data->begin(), another.data->end());
        *this = block(temp);
    }

    block split(int pos) {
        block result(make_shared<vi>(data->begin() + pos, data->end()));
        *this = block(make_shared<vi>(data->begin(), data->begin() + pos));
        return result;
    }
};

typedef list<block>::iterator lit;

struct blocklist {
    list<block> blk;

    void maintain() {
        lit it = blk.begin();
        while (it != blk.end() and next(it) != blk.end()) {
            lit it2 = it;
            while (next(it2) != blk.end() and
                it2->data->size() + next(it2)->data->size() <= BLOCK) {
                it2->merge(*next(it2));
                blk.erase(next(it2));
            }
            ++it;
        }
    }
};

```

0563  
013b  
427e  
a771  
2989  
8fd0  
427e  
427e  
a613  
24b5  
0cf0  
e93b  
427e  
5c0f  
0b18  
ac21  
6467  
95cf  
427e  
42e8  
dac1  
01db  
56b0  
95cf  
329b  
427e  
2a18  
427e  
ce14  
5540  
427e  
7b8e  
3131  
5e44  
852d  
0b03  
029f  
93e1  
e1fa  
95cf  
5771  
95cf  
427e

```

b7b3     lit split(int pos) {
2273         for (lit it = blk.begin(); ; it++) {
5502             if (pos == 0) return it;
d480             while (it->data->size() > pos) {
2099                 blk.insert(next(it), it->split(pos));
95cf             }
a1c8             pos -= it->data->size();
95cf         }
95cf     }
427e
fd38     LL sum(int l, int r) { // traverse
48b4         lit it1 = split(l), it2 = split(r);
ac09         LL res = 0;
9f1d         while (it1 != it2) {
8284             res += it1->sum;
61fd             it1++;
95cf         }
b204         maintain();
244d         return res;
95cf     }
329b };

```

## 6.10 Sparse table, range minimum query

The array is 0-based and the range is left-closed right-open.

```

db63     const int MAXN = 100007;
cefd     int a[MAXN], st[MAXN][30];
427e
d34f     void init(int n){
c73d         int l = log2(n);
cf75         rep (i, n) st[i][0] = a[i];
426b         rep (j, l) rep (i, 1+n-(1<<j))
1131             st[i][j+1] = min(st[i][j], st[i+(1<<j)][j]);
95cf     }
427e
c863     int rmq(int l, int r){
f089         int k = log2(r - l);
6117         return min(st[l][k], st[r-(1<<k)][k]);
95cf     }

```

## 7 Geometrics

### 7.1 2D geometric template

```

#include <bits/stdc++.h>
using namespace std;

typedef int T;
typedef struct pt {
    T x, y;
    T operator , (pt a) { return x*a.x + y*a.y; } // inner product
    T operator * (pt a) { return x*a.y - y*a.x; } // outer product
    pt operator + (pt a) { return {x+a.x, y+a.y}; }
    pt operator - (pt a) { return {x-a.x, y-a.y}; }

    pt operator * (T k) { return {x*k, y*k}; }
    pt operator - () { return {-x, -y}; }
} vec;

typedef pair<pt, pt> seg;

bool ptOnSeg(pt& p, seg& s){
    vec v1 = s.first - p, v2 = s.second - p;
    return (v1, v2) <= 0 && v1 * v2 == 0;
}

// 0 not on segment
// 1 on segment except vertices
// 2 on vertices
int ptOnSeg2(pt& p, seg& s){
    vec v1 = s.first - p, v2 = s.second - p;
    T ip = (v1, v2);
    if (v1 * v2 != 0 || ip > 0) return 0;
    return (v1, v2) ? 1 : 2;
}

// if two orthogonal rectangles do not touch, return true
inline bool nIntRectRect(seg a, seg b){
    return min(a.first.x, a.second.x) > max(b.first.x, b.second.x) ||
           min(a.first.y, a.second.y) > max(b.first.y, b.second.y) ||
           min(b.first.x, b.second.x) > max(a.first.x, a.second.x) ||
           min(b.first.y, b.second.y) > max(a.first.y, a.second.y);
}

```

302f  
421c  
427e  
4553  
c0ae  
7a9d  
ffaa  
3ec7  
221a  
8b34  
427e  
368b  
90f4  
ba8c  
427e  
0ea6  
427e  
8d6e  
ce77  
de97  
95cf  
427e  
427e  
427e  
427e  
8421  
ce77  
70ca  
8b14  
0847  
95cf  
427e  
427e  
72bb  
f9ac  
f486  
39ce  
80c7  
95cf

```

427e // >0 in order
427e // <0 out of order
427e // =0 not standard
7538 inline double rotOrder(vec a, vec b, vec c){return double(a*b)*(b*c);}
427e
31ed inline bool intersect(seg a, seg b){
427e     // ! if (nIntRectRect(a, b)) return false; // if commented, assume that a
        and b are non-collinear
cb52     return rotOrder(b.first-a.first, a.second-a.first, b.second-a.first) >= 0 &&
059e         rotOrder(a.first-b.first, b.second-b.first, a.second-b.first) >= 0;
95cf }
427e
427e // 0 not intersect
427e // 1 standard intersection
427e // 2 vertex-line intersection
427e // 3 vertex-vertex intersection
427e // 4 collinear and have common point(s)
4d19 int intersect2(seg& a, seg& b){
5dc4     if (nIntRectRect(a, b)) return 0;
42c0     vec va = a.second - a.first, vb = b.second - b.first;
2096     double j1 = rotOrder(b.first-a.first, va, b.second-a.first),
72fe         j2 = rotOrder(a.first-b.first, vb, a.second-b.first);
5ac6     if (j1 < 0 || j2 < 0) return 0;
9400     if (j1 != 0 && j2 != 0) return 1;
83db     if (j1 == 0 && j2 == 0){
6b0c         if (va * vb == 0) return 4; else return 3;
fb17     } else return 2;
95cf }
427e
2c68 template <typename Tp = T>
5894 inline pt getIntersection(pt P, vec v, pt Q, vec w){
6850     static_assert(is_same<Tp, double>::value, "must_be_double!");
7c9a     return P + v * (w*(P-Q)/(v*w));
95cf }
427e
427e // -1 outside the polygon
427e // 0 on the border of the polygon
427e // 1 inside the polygon
cbdd int ptOnPoly(pt p, pt* poly, int n){
5fb4     int wn = 0;
1294     for (int i = 0; i < n; i++) {
427e
3cae         T k, d1 = poly[i].y - p.y, d2 = poly[(i+1)%n].y - p.y;

```

```

        if (k = (poly[(i+1)%n] - poly[i])*(p - poly[i])){
            if (k > 0 && d1 <= 0 && d2 > 0) wn++;
            if (k < 0 && d2 <= 0 && d1 > 0) wn--;
        } else return 0;
    }
    return wn ? 1 : -1;
}

istream& operator >> (istream& lhs, pt& rhs){
    lhs >> rhs.x >> rhs.y;
    return lhs;
}

istream& operator >> (istream& lhs, seg& rhs){
    lhs >> rhs.first >> rhs.second;
    return lhs;
}

```

```

b957
8c40
3c4d
aad3
95cf
0a5f
95cf
427e
d4a3
fa86
331a
95cf
427e
07ae
5cab
331a
95cf

```

## 8 Appendices

### 8.1 Primes

#### 8.1.1 First primes

$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$
2	1	3	2	5	2	7	3	11	2
13	2	17	3	19	2	23	5	29	2
31	3	37	2	41	6	43	3	47	5
53	2	59	2	61	2	67	2	71	7
73	5	79	3	83	2	89	3	97	5
101	2	103	5	107	2	109	6	113	3
127	3	131	2	137	3	139	2	149	2
151	6	157	5	163	2	167	5	173	2
179	2	181	2	191	19	193	5	197	2
199	3	211	2	223	3	227	2	229	6

#### 8.1.2 Arbitrary length primes

$\lg p$	$p$	$g(p)$	$p$	$g(p)$
3	967	5	1031	14
4	9859	2	10273	10
5	96331	10	102931	3
6	958543	6	1031137	5
7	9594539	2	10169651	2
8	96243449	3	103211039	7
9	980483981	2	1042484357	2
10	9858935453	2	10261276009	7
11	95748666809	3	101759940101	2
12	950781833849	3	1012797784423	5
13	9739822952371	7	10037217092377	7
14	96181051140397	5	104974966380359	11
15	981030138360889	13	1029038416465403	2
16	9655206098080843	3	10116299875820773	2
17	97687777921994419	3	101506415998163437	2

#### 8.1.3 $\sim 1 \times 10^9$

$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$
954854573	3	967607731	2	973215833	3
975831713	3	978949117	2	980766497	3
983879921	3	985918807	3	986608921	29
991136977	5	991752599	13	997137961	11
1003911991	3	1009775293	2	1012423549	6
1021000537	5	1023976897	7	1024153643	2
1037027287	3	1038812881	11	1044754639	3
1045125617	3	1047411427	3	1047753349	6

#### 8.1.4 $\sim 1 \times 10^{18}$

$p$	$g(p)$	$p$	$g(p)$
951970612352230049	3	963284339889659609	3
967495386904694119	3	969751761517096213	2
983238274281901499	2	984647442475101409	23
989286107138674069	11	1002507954383424641	3
1006658951440146419	2	1020152326159075903	3
1034876265966119449	7	1042753851435034019	2
1043609016597371563	2	1045571042176595707	2
1048364250160580293	2	1049495624119026949	2

### 8.2 Pell's equation

$x^2 - ny^2 = 1$ , where  $n$  is a positive nonsquare integer.

Let  $(x_0, y_0)$  be the smallest positive solution of the equation, then the  $k$ -th solution is:

$$\begin{pmatrix} x_k \\ y_k \end{pmatrix} = \begin{pmatrix} x_0 & ny_0 \\ y_0 & x_0 \end{pmatrix}^k \begin{pmatrix} x_0 \\ y_0 \end{pmatrix}$$

Some smallest solutions to Pell's equation:

$n$	2	3	5	6	7	8	10	11	12	13	14	15	17	18	19	20
$x$	3	2	9	5	8	3	19	10	7	649	15	4	33	17	170	9
$y$	2	1	4	2	3	1	6	3	2	180	4	1	8	4	39	2

### 8.3 Burnside's lemma and Polya's enumeration theorem

The Burnside's lemma says that

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

where  $G$  is a group acting on  $X$ ,  $X^g$  is the set of elements in  $X$  that are fixed by  $g$ , i.e.  $X^g = \{x \in X : gx = x\}$ .

The unweighted version of Pólya enumeration theorem says that

$$|Y^X/G| = \frac{1}{|G|} \sum_{g \in G} m^{c_g}$$

where  $m = |X|$  is the number of colors,  $c_g$  is the number of the cycles of permutation  $g$ .

### 8.4 Lagrange's interpolation

For sample points  $(x_0, y_0), \dots, (x_k, y_k)$ , define

$$l_j(x) = \prod_{0 \leq m \leq k, m \neq j} \frac{x - x_m}{x_j - x_m}$$

then the Lagrange polynomial is

$$L(x) = \sum_{j=0}^k y_j l_j(x).$$

To use the script below, type two lines

```
x0 x1 x2 ... xn
y0 y1 y2 ... yn
```

the script will print the fractional coefficient of the polynomial in ascending exponent order.

```
#!/usr/bin/python2
from fractions import *

def polymul(a, b) :
    p = [0] * (len(a)+len(b)-1)
    for e1, c1 in enumerate(a) :
        for e2, c2 in enumerate(b) :
            p[e1+e2] += c1*c2
    return p

x, y = [map(Fraction, raw_input().split()) for _ in 0,0]
n = len(x)
lj = [reduce(polymul, [[-x[m]/(x[j]-x[m]), 1/(x[j]-x[m])]
    for m in range(n) if m != j]]) for j in range(n)]
print '\n'.join(map(str, map(sum, zip(*map(
    lambda a, b : [x*a for x in b], y, lj)))))
```

6dc9  
4b2b  
427e  
796b  
83e4  
f697  
156c  
dfce  
5849  
427e  
f06d  
e80a  
a649  
9dfa  
3cae  
7c0d