

# 南京大学 ACM-ICPC 集训队代码模版库



OS: Linux-4.15.0-45-generic-x86\_64-with-Ubuntu-18.04-bionic  
python: 2.7.15rc1  
2019-02-25 18:27:33.343420, build 0021

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## 1 General

### 1.1 Code library checksum

```
ab14 #!/usr/bin/python3
c502 import re, sys, hashlib
427e
f7db for line in sys.stdin.read().strip().split("\n") :
ddf5     print(hashlib.md5(re.sub(r'\s|//[.]*', '', line).encode('utf8')).hexdigest()
        [-4:], line)
```

### 1.2 Makefile

```
dab2 .PHONY : run
427e
207e $(t) : $(t).cpp
2d16     g++ --std=c++14 -Wall -D__LOCAL_DEBUG__ -fsanitize=undefined -fsanitize=
        address -ggdb -pipe -o $@ $<
427e
5f25 run : $(t)
bf3e     ./$<(t) < $(t).in
```

### 1.3 .vimrc

```
914c set nocompatible
733d syntax on
6bbc colorscheme slate
7db5 set number
b0e3 set cursorline
061b set shiftwidth=2
8011 set softtabstop=2
a66d set tabstop=2
d23a set expandtab
5245 set magic
740c set smartindent
bee8 set backspace=indent,eol,start
815d set cmdheight=1
0a40 set laststatus=2
1c67 set whichwrap=b,s,<,>,[,]
```

### 1.4 Stack

```
const int STK_SZ = 2000000;
char STK[STK_SZ * sizeof(void)];
void *STK_BAK;

#if defined(__i386__)
#define SP "%esp"
#elif defined(__x86_64__)
#define SP "%rsp"
#endif

int main() {
    asm volatile("movl SP, %0; movl %1, SP: "=g"(STK_BAK):"g"(STK+sizeof(STK)):");
    ;

    // main program

    asm volatile("movl %0, SP: "=g"(STK_BAK));
    return 0;
}
```

### 1.5 Template

```
#include <bits/stdc++.h>
using namespace std;

#ifdef __LOCAL_DEBUG__
# define _debug(fmt, ...) fprintf(stderr, "[%s] " fmt "\n", \
    __func__, ##__VA_ARGS__)
#else
# define _debug(...) ((void) 0)
#endif

#define rep(i, n) for (int i=0; i<(n); i++)
#define Rep(i, n) for (int i=1; i<=(n); i++)
#define range(x) begin(x), end(x)
typedef long long LL;
typedef unsigned long long ULL;
```

## 2 Miscellaneous Algorithms

### 2.1 2-SAT

```

0f42 const int MAXN = 100005;
03a9 struct twoSAT{
5c83     int n;
8f72     vector<int> G[MAXN*2];
d060     bool mark[MAXN*2];
b42d     int S[MAXN*2], c;
427e
d34f     void init(int n){
b985         this->n = n;
f9ec         for (int i=0; i<n*2; i++) G[i].clear();
0609         memset(mark, 0, sizeof(mark));
95cf     }
427e
3bd5     bool dfs(int x){
bd70         if (mark[x^1]) return false;
c96a         if (mark[x]) return true;
fd23         mark[x] = true;
4bea         S[c++] = x;
1ce6         for (int i=0; i<G[x].size(); i++)
d942             if (!dfs(G[x][i])) return false;
3361         return true;
95cf     }
427e
5894     void add_clause(int x, bool xval, int y, bool yval){
6afe         x = x * 2 + xval;
e680         y = y * 2 + yval;
81cc         G[x^1].push_back(y);
6835         G[y^1].push_back(x);
95cf     }
427e
d0cb     bool solve() {
7c39         for (int i=0; i<n*2; i+=2){
e63f             if (!mark[i] && !mark[i+1]){
88fb                 c = 0;
f4b9                 if (!dfs(i)){
3f03                     while (c > 0) mark[S[--c]] = false;
86c5                     if (!dfs(i+1)) return false;
95cf                 }
95cf             }

```

```

    }
    return true;
}

inline bool value(unsigned i){return mark[2*i+1];}
};

```

95cf  
3361  
95cf  
427e  
5f0a  
329b

### 2.2 Knuth's optimization

```

int n;
int dp[256][256], dc[256][256];

template <typename T>
void compute(T cost) {
    for (int i = 0; i <= n; i++) {
        dp[i][i] = 0;
        dc[i][i] = i;
    }
    rep (i, n) {
        dp[i][i+1] = 0;
        dc[i][i+1] = i;
    }
    for (int len = 2; len <= n; len++) {
        for (int i = 0; i + len <= n; i++) {
            int j = i + len;
            int lbnd = dc[i][j-1], rbnd = dc[i+1][j];
            dp[i][j] = INT_MAX / 2;
            int c = cost(i, j);
            for (int k = lbnd; k <= rbnd; k++) {
                int res = dp[i][k] + dp[k][j] + c;
                if (res < dp[i][j]) {
                    dp[i][j] = res;
                    dc[i][j] = k;
                }
            }
        }
    }
};

```

5c83  
d77c  
427e  
b7ec  
0bc7  
0423  
8f5e  
9488  
95cf  
be8e  
95b5  
aa0f  
95cf  
ec08  
88b8  
d3da  
9824  
a24a  
f933  
90d2  
9bd0  
26b5  
e6af  
9c88  
95cf  
95cf  
95cf  
95cf  
329b

## 2.3 Mo's algorithm

All intervals are closed on both sides. When running functions `enter()` and `leave()`, the global `l` and `r` has not changed yet.

### Usage:

```
add_query(id, l, r)    Add id-th query [l, r].
run()                 Run Mo's algorithm.
init()                TODO. Initialize the range [l, r].
yield(id)             TODO. Yield answer for id-th query.
enter(o)              TODO. Add o-th element.
leave(o)              TODO. Remove o-th element.
```

```
5194 constexpr int BLOCK_SZ = 300;
427e
3ec4 struct query { int l, r, id; };
d26a vector<query> queries;
427e
1e30 void add_query(int id, int l, int r) {
54c9     queries.push_back(query{l, r, id});
95cf }
427e
9f6b int l, r;
427e
427e // ----- functions to implement -----
62b4 inline void init();
50e1 inline void yield(int id);
b20d inline void enter(int o);
13af inline void leave(int o);
427e
37f0 void run() {
ab0b     if (queries.empty()) return;
8508     sort(range(queries), [](query lhs, query rhs) {
c7f8         int lb = lhs.l / BLOCK_SZ, rb = rhs.l / BLOCK_SZ;
03e7         if (lb != rb) return lb < rb;
0780         return lhs.r < rhs.r;
b251     });
6196     l = queries[0].l;
9644     r = queries[0].r;
07e2     init();
5bc9     for (query q : queries) {
7bc7         while (l > q.l) enter(l - 1), l--;
d646         while (r < q.r) enter(r + 1), r++;
13f0         while (l < q.l) leave(l), l++;
e1c6         while (r > q.r) leave(r), r--;
```

```
        yield(q.id);
    }
}
```

```
82f5
95cf
95cf
```

## 3 String

### 3.1 Knuth-Morris-Pratt algorithm

```
const int SIZE = 10005;

struct kmp_matcher {
    char p[SIZE];
    int fail[SIZE];
    int len;

    void construct(const char* needle) {
        len = strlen(p);
        strcpy(p, needle);
        fail[0] = fail[1] = 0;
        for (int i = 1; i < len; i++) {
            int j = fail[i];
            while (j && p[i] != p[j]) j = fail[j];
            fail[i + 1] = p[i] == p[j] ? j + 1 : 0;
        }
    }

    inline void found(int pos) {
        // ! add codes for having found at pos
    }

    void match(const char* haystack) { // must be called after construct
        const char* t = haystack;
        int n = strlen(t);
        int j = 0;
        rep(i, n) {
            while (j && p[j] != t[i]) j = fail[j];
            if (p[j] == t[i]) j++;
            if (j == len) found(i - len + 1);
        }
    }
};
```

```
2836
427e
d02b
2d81
9847
57b7
427e
60cf
aaa1
3a87
3dd4
d8a8
147f
3c79
4643
95cf
95cf
427e
c464
427e
95cf
427e
2daf
700f
8482
8fd0
be8e
4e19
b5d5
f024
95cf
95cf
329b
```

### 3.2 Manacher algorithm

```

81d4 struct Manacher {
cd09     int Len;
9255     vector<int> lc;
b301     string s;
427e
ec07     void work() {
c033         lc[1] = 1;
6bef         int k = 1;
427e
491f         for (int i = 2; i <= Len; i++) {
7957             int p = k + lc[k] - 1;
5e04             if (i <= p) {
24a1                 lc[i] = min(lc[2 * k - i], p - i + 1);
8e2e             } else {
e0e5                 lc[i] = 1;
95cf             }
74ff             while (s[i + lc[i]] == s[i - lc[i]]) lc[i]++;
2b9a             if (i + lc[i] > k + lc[k]) k = i;
95cf         }
95cf     }
427e
bfd5     void init(const char *tt) {
aaaf         int len = strlen(tt);
f701         s.resize(len * 2 + 10);
7045         lc.resize(len * 2 + 10);
8e13         s[0] = '*';
ae54         s[1] = '#';
1321         for (int i = 0; i < len; i++) {
e995             s[i * 2 + 2] = tt[i];
69fd             s[i * 2 + 1] = '#';
95cf         }
43fd         s[len * 2 + 1] = '#';
75d1         s[len * 2 + 2] = '\0';
61f7         Len = len * 2 + 2;
3e7a         work();
95cf     }
427e
b194     pair<int, int> maxpal(int l, int r) {
901a         int center = l + r + 1;
ffb2         int rad = lc[center] / 2;
ab54         int rmid = (l + r + 1) / 2;

```

```

    int r1 = rmid - rad, rr = rmid + rad - 1;
    if ((r ^ 1) & 1) {
    } else rr++;
    return {max(l, r1), min(r, rr)};
}
};

```

```

17e4
3908
69f3
69dc
95cf
329b

```

### 3.3 Aho-corasick automaton

```

struct AC : Trie {
    int fail[MAXN];
    int last[MAXN];

    void construct() {
        queue<int> q;
        fail[0] = 0;
        rep(c, CHARN) {
            if (int u = tr[0][c]) {
                fail[u] = 0;
                q.push(u);
                last[u] = 0;
            }
        }
        while (!q.empty()) {
            int r = q.front();
            q.pop();
            rep(c, CHARN) {
                int u = tr[r][c];
                if (!u) {
                    tr[r][c] = tr[fail[r]][c];
                    continue;
                }
                q.push(u);
                int v = fail[r];
                while (v && !tr[v][c]) v = fail[v];
                fail[u] = tr[v][c];
                last[u] = tag[fail[u]] ? fail[u] : last[fail[u]];
            }
        }
    }

    void found(int pos, int j) {

```

```

a1ad
9143
daca
427e
8690
93d2
a7a6
ce3c
b1c6
a506
3e14
f689
95cf
95cf
cc78
31f0
15dd
ce3c
ab59
0ef5
9d58
b333
95cf
3e14
b3ff
d2ea
c275
654c
95cf
95cf
95cf
427e
7752

```

```

043e     if (j) {
427e         // ! add codes for having found word with tag[j]
4a96         found(pos, last[j]);
95cf     }
95cf }
427e
9785 void find(const char* text) { // must be called after construct()
80a4     int p = 0, c, len = strlen(text);
9c94     rep(i, len) {
b3db         c = id(text[i]);
f119         p = tr[p][c];
f08e         if (tag[p])
389b             found(i, p);
1e67         else if (last[p])
299e             found(i, last[p]);
95cf     }
95cf }
329b };

```

### 3.4 Suffix array

The character immediately after the end of the string **MUST** be set to the **UNIQUE SMALLEST** element.

#### Usage:

s[]	the source string
sa[i]	the index of starting position of $i$ -th suffix
rk[i]	the number of suffixes less than the suffix starting from $i$
h[i]	the longest common prefix between the $i$ -th and $(i-1)$ -th lexicographically smallest suffixes
n	size of source string
m	size of character set

```

de09 void radix_sort(int x[], int y[], int sa[], int n, int m) {
ec00     static int cnt[1000005]; // size > max(n, m)
6066     fill(cnt, cnt + m, 0);
93b7     rep(i, n) cnt[x[y[i]]]++;
9154     partial_sum(cnt, cnt + m, cnt);
acac     for (int i = n - 1; i >= 0; i--) sa[--cnt[x[y[i]]]] = y[i];
95cf }
427e
c939 void suffix_array(int s[], int sa[], int rk[], int n, int m) {
a69a     static int y[1000005]; // size > n
7306     copy(s, s + n, rk);

```

```

iota(y, y + n, 0);
radix_sort(rk, y, sa, n, m);
for (int j = 1, p = 0; j <= n; j <= 1, m = p, p = 0) {
    for (int i = n - j; i < n; i++) y[p++] = i;
    rep(i, n) if (sa[i] >= j) y[p++] = sa[i] - j;
    radix_sort(rk, y, sa, n, m + 1);
    swap_ranges(rk, rk + n, y);
    rk[sa[0]] = p = 1;
    for (int i = 1; i < n; i++)
        rk[sa[i]] = ((y[sa[i]] == y[sa[i-1]] and y[sa[i]+j] == y[sa[i-1]+j])
            ? p : ++p);
    if (p == n) break;
}
rep(i, n) rk[sa[i]] = i;
}

void calc_height(int s[], int sa[], int rk[], int h[], int n) {
    int k = 0;
    h[0] = 0;
    rep(i, n) {
        k = max(k - 1, 0);
        if (rk[i]) while (s[i+k] == s[sa[rk[i]-1]+k]) ++k;
        h[rk[i]] = k;
    }
}

```

### 3.5 Trie

```

const int MAXN = 12000;
const int CHARN = 26;

inline int id(char c) { return c - 'a'; }

struct Trie {
    int n;
    int tr[MAXN][CHARN]; // Trie tree, 0 denotes fail
    int tag[MAXN];

    Trie() {
        memset(tr[0], 0, sizeof(tr[0]));
        tag[0] = 0;
        n = 1;
    }
}

```

```

95cf }
427e
427e // tag should not be 0
30b0 void add(const char* s, int t) {
d50a     int p = 0, c, len = strlen(s);
9c94     rep(i, len) {
3140         c = id(s[i]);
d6c8         if (!tr[p][c]) {
26dd             memset(tr[n], 0, sizeof(tr[n]));
2e5c             tag[n] = 0;
73bb             tr[p][c] = n++;
95cf         }
f119         p = tr[p][c];
95cf     }
35ef     tag[p] = t;
95cf }
427e
427e // returns 0 if not found
427e // AC automaton does not need this function
216c int search(const char* s) {
d50a     int p = 0, c, len = strlen(s);
9c94     rep(i, len) {
3140         c = id(s[i]);
f339         if (!tr[p][c]) return 0;
f119         p = tr[p][c];
95cf     }
840e     return tag[p];
95cf }
329b };

```

### 3.6 Rolling hash

**PLEASE** call `init_hash()` in `int main()`!

**Usage:**

`build(str)` Construct the hasher with given string.  
`operator()(l, r)` Get hash value of substring  $[l, r)$ .

```

1e42 const LL mod = 1006658951440146419, g = 967;
9f60 const int MAXN = 200005;
0291 LL pg[MAXN];
427e
dfe7 inline LL mul(LL x, LL y) { return __int128_t(x) * y % mod; }
427e

```

```

void init_hash() { // must be called in `int main()`
    pg[0] = 1;
    for (int i = 1; i < MAXN; i++) pg[i] = mul(pg[i-1], g);
}

struct hasher {
    LL val[MAXN];

    void build(const char *str) { // assume lower-case letter only
        for (int i = 0; str[i]; i++)
            val[i+1] = (mul(val[i], g) + str[i]) % mod;
    }

    LL operator() (int l, int r) { // [l, r)
        return (val[r] - mul(val[l], pg[r-l]) + mod) % mod;
    }
};

```

599a  
286f  
4af8  
95cf  
427e  
7e62  
534a  
427e  
4554  
f937  
9645  
95cf  
427e  
19f8  
9986  
95cf  
329b

## 4 Math

### 4.1 Extended Euclidean algorithm and Chinese remainder theorem

```

void exgcd(LL a, LL b, LL &g, LL &x, LL &y) {
    if (!b) g = a, x = 1, y = 0;
    else {
        exgcd(b, a % b, g, y, x);
        y -= x * (a / b);
    }
}

LL crt(LL r[], LL p[], int n) {
    LL q = 1, ret = 0;
    rep (i, n) q *= p[i];
    rep (i, n) {
        LL m = q / p[i];
        LL d, x, y;
        exgcd(p[i], m, d, x, y);
        ret = (ret + y * m * r[i]) % q;
    }
    return (q + ret) % q;
}

```

4fba  
7db6  
037f  
ffca  
d798  
95cf  
95cf  
427e  
e491  
84e6  
00d9  
be8e  
98b4  
9f4f  
b082  
3cd3  
95cf  
2e47  
95cf



## 4.2 Matrix powermod

```

44b4 const int MAXN = 105;
92df const LL modular = 1000000007;
5c83 int n; // order of matrices
427e
8864 struct matrix{
3180     LL m[MAXN][MAXN];
427e
43c5     void operator *=(matrix& a){
e735         static LL t[MAXN][MAXN];
34d7         Rep (i, n){
4c11             Rep (j, n){
ee1e                 t[i][j] = 0;
c4a7                 Rep (k, n){
fcac                     t[i][j] += (m[i][k] * a.m[k][j]) % modular;
199e                     t[i][j] %= modular;
95cf                 }
95cf             }
95cf         }
dad4         memcpy(m, t, sizeof(t));
95cf     }
329b };
427e
63d8 matrix r;
3ec2 void m_powmod(matrix& b, LL e){
83f0     memset(r.m, 0, sizeof(r.m));
a7c3     Rep(i, n)
de64         r.m[i][i] = 1;
3e90     while (e){
5a0e         if (e & 1) r *= b;
35c5         b *= b;
16fc         e >>= 1;
95cf     }
95cf }

```

## 4.3 Linear basis

```

8b44 const int MAXD = 30;
03a6 struct linearbasis {
3558     ULL b[MAXD] = {};
427e

```

```

bool insert(LL v) {
    for (int j = MAXD - 1; j >= 0; j--) {
        if (!(v & (1ll << j))) continue;
        if (b[j]) v ^= b[j]
        else {
            for (int k = 0; k < j; k++)
                if (v & (1ll << k)) v ^= b[k];
            for (int k = j + 1; k < MAXD; k++)
                if (b[k] & (1ll << j)) b[k] ^= v;
            b[j] = v;
            return true;
        }
    }
    return false;
}

```

## 4.4 Gauss elimination over finite field

```

const LL p = 1000000007;

LL powmod(LL b, LL e) {
    LL r = 1;
    while (e) {
        if (e & 1) r = r * b % p;
        b = b * b % p;
        e >>= 1;
    }
    return r;
}

typedef vector<LL> VLL;
typedef vector<VLL> WLL;

```

```

LL gauss(WLL &a, WLL &b) {
    const int n = a.size(), m = b[0].size();
    vector<int> irow(n), icol(n), ipiv(n);
    LL det = 1;

    rep (i, n) {
        int pj = -1, pk = -1;
        rep (j, n) if (!ipiv[j])

```

1566  
9b2b  
de36  
ee78  
037f  
7836  
f0b4  
b0aa  
46c9  
8295  
3361  
95cf  
95cf  
438e  
95cf  
329b

b784  
427e  
2a2c  
95a2  
3e90  
1783  
5549  
16fc  
95cf  
547e  
95cf  
427e  
c130  
42ac  
427e  
2c62  
561b  
a25e  
2976  
427e  
be8e  
d2b5  
6b4a

```

e582     rep (k, n) if (!ipiv[k])
6112         if (pj == -1 || a[j][k] > a[pj][pk]) {
a905             pj = j;
657b             pk = k;
95cf         }
d480     if (a[pj][pk] == 0) return 0;
0305     ipiv[pk]++;
8dad     swap(a[pj], a[pk]);
aad8     swap(b[pj], b[pk]);
be4d     if (pj != pk) det = (p - det) % p;
d080     irow[i] = pj;
f156     icol[i] = pk;
427e
4ecd     LL c = powmod(a[pk][pk], p - 2);
865b     det = det * a[pk][pk] % p;
c36a     a[pk][pk] = 1;
dd36     rep (j, n) a[pk][j] = a[pk][j] * c % p;
1b23     rep (j, m) b[pk][j] = b[pk][j] * c % p;
f8f3     rep (j, n) if (j != pk) {
e97f         c = a[j][pk];
c449         a[j][pk] = 0;
820b         rep (k, n) a[j][k] = (a[j][k] + p - a[pk][k] * c % p) % p;
f039         rep (k, m) b[j][k] = (b[j][k] + p - b[pk][k] * c % p) % p;
95cf     }
95cf }
427e
37e1     for (int j = n - 1; j >= 0; j--) if (irow[j] != icol[j]) {
50dc         for (int k = 0; k < n; k++) swap(a[k][irow[j]], a[k][icol[j]]);
95cf     }
f27f     return det;
95cf }

```

## 4.5 Berlekamp-Massey algorithm

```

d790 vector<int> berlekamp(const vector<int>& a) {
4166     vector<int> p = {1}, r = {1};
baed     int dif = 1;
8bc9     rep (i, a.size()) {
3e58         int u = 0;
ac8e         rep (j, p.size())
a488             u = (u + 111 * p[j] * a[i-j]) % mod;
eae9         if (u == 0) {

```

```

        r.insert(r.begin(), 0);
    } else {
        auto op = p;
        p.resize(max(p.size(), r.size() + 1));
        int idif = inv(dif);
        rep (j, r.size())
            p[j+1] =
                (p[j+1] - 111 * r[j] * idif % mod * u % mod + mod) % mod;
        dif = u;
        r = op;
    }
}
return p;
}

```

```

b14c
8e2e
0c78
02f6
786b
9b57
793c
1836
644c
bc58
95cf
95cf
e149
95cf

```

## 4.6 Fast Walsh-Hadamard transform

```

void fwt(int* a, int n){
    for (int d = 1; d < n; d <= 1)
        for (int i = 0; i < n; i += d < 1)
            rep (j, d){
                int x = a[i+j], y = a[i+j+d];
                // a[i+j] = x+y, a[i+j+d] = x-y;    // xor
                // a[i+j] = x+y;                    // and
                // a[i+j+d] = x+y;                    // or
            }
}

```

```

061e
5595
05f2
b833
7796
427e
427e
427e
95cf
95cf
427e
4db1

```

```

void ifwt(int* a, int n){
    for (int d = 1; d < n; d <= 1)
        for (int i = 0; i < n; i += d < 1)
            rep (j, d){
                int x = a[i+j], y = a[i+j+d];
                // a[i+j] = (x+y)/2, a[i+j+d] = (x-y)/2;    // xor
                // a[i+j] = x-y;                                // and
                // a[i+j+d] = y-x;                                // or
            }
}

```

```

5595
05f2
b833
7796
427e
427e
427e
95cf
95cf
427e
2ab6
950a
e427

```

```

void conv(int* a, int* b, int n){
    fwt(a, n);
    fwt(b, n);

```

```

8a42     rep(i, n) a[i] *= b[i];
430f     ifwt(a, n);
95cf }

```

## 4.7 Fast fourier transform

```

4e09 const int NMAX = 1<<20;
427e
3fbf typedef complex<double> cplx;
427e
abd1 const double PI = 2*acos(0.0);
12af struct FFT{
c47c     int rev[NMAX];
27d7     cplx omega[NMAX], oinv[NMAX];
9827     int K, N;
427e
1442     FFT(int k){
e209         K = k; N = 1 << k;
b393         rep (i, N){
7ba3             rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
1908             omega[i] = polar(1.0, 2.0 * PI / N * i);
a166             oinv[i] = conj(omega[i]);
95cf         }
95cf     }
427e
b941     void dft(cplx* a, cplx* w){
a215         rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
ac6e         for (int l = 2; l <= N; l *= 2){
2969             int m = l/2;
b3cf             for (cplx* p = a; p != a + N; p += l)
c24f                 rep (k, m){
fe06                     cplx t = w[N/l*k] * p[k+m];
ecbf                     p[k+m] = p[k] - t; p[k] += t;
95cf                 }
95cf             }
95cf         }
427e
617b     void fft(cplx* a){dft(a, omega);}
a123     void ifft(cplx* a){
3b2f         dft(a, oinv);
57fc         rep (i, N) a[i] /= N;
95cf     }

```

```

void conv(cplx* a, cplx* b){
    fft(a); fft(b);
    rep (i, N) a[i] *= b[i];
    ifft(a);
}
};

```

```

427e
bdc0
6497
12a5
f84e
95cf
329b

```

## 4.8 Number theoretic transform

```

const int NMAX = 1<<21;

// 998244353 = 7*17*2^23+1, G = 3
const int P = 1004535809, G = 3; // = 479*2^21+1

struct NTT{
    int rev[NMAX];
    LL omega[NMAX], oinv[NMAX];
    int g, g_inv; // g: g_n = G^((P-1)/n)
    int K, N;

    LL powmod(LL b, LL e){
        LL r = 1;
        while (e){
            if (e&1) r = r * b % P;
            b = b * b % P;
            e >>= 1;
        }
        return r;
    }

    NTT(int k){
        K = k; N = 1 << k;
        g = powmod(G, (P-1)/N);
        g_inv = powmod(g, N-1);
        omega[0] = oinv[0] = 1;
        rep (i, N){
            rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
            if (i){
                omega[i] = omega[i-1] * g % P;
                oinv[i] = oinv[i-1] * g_inv % P;
            }
        }
    }
}

```

```

4ab9
427e
427e
fb9a
427e
87ab
c47c
0eda
81af
9827
427e
2a2c
95a2
3e90
6624
489e
16fc
95cf
547e
95cf
427e
f420
e209
7652
4b3a
e04f
b393
7ba3
ad4f
8d8b
9e14
95cf

```

```

95cf    }
95cf    }
427e
9668    void _ntt(LL* a, LL* w){
a215        rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
ac6e        for (int l = 2; l <= N; l *= 2){
2969            int m = l/2;
7a1d            for (LL* p = a; p != a + N; p += l)
c24f                rep (k, m){
0ad3                    LL t = w[N/l*k] * p[k+m] % P;
6209                    p[k+m] = (p[k] - t + P) % P;
fa1b                    p[k] = (p[k] + t) % P;
95cf                }
95cf            }
95cf        }
427e
92ea    void ntt(LL* a){_ntt(a, omega);}
5daf    void intt(LL* a){
1f2a        LL inv = powmod(N, P-2);
9910        _ntt(a, oinv);
a873        rep (i, N) a[i] = a[i] * inv % P;
95cf    }
427e
3a5b    void conv(LL* a, LL* b){
ad16        ntt(a); ntt(b);
e49e        rep (i, N) a[i] = a[i] * b[i] % P;
5748        intt(a);
95cf    }
329b    };

```

#### 4.9 Sieve of Euler

```

cfc3    const int MAXX = 1e7+5;
5861    bool p[MAXX];
73ae    int prime[MAXX], sz;
427e
9bc6    void sieve(){
9628        p[0] = p[1] = 1;
1ec8        for (int i = 2; i < MAXX; i++){
bf28            if (!p[i]) prime[sz++] = i;
e82c            for (int j = 0; j < sz && i*prime[j] < MAXX; j++){
b6a9                p[i*prime[j]] = 1;

```

```

        if (i % prime[j] == 0) break;
    }
}

```

#### 4.10 Sieve of Euler (General)

```

namespace sieve {
constexpr int MAXN = 10000007;
bool p[MAXN]; // true if not prime
int prime[MAXN], sz;
int pval[MAXN], pcnt[MAXN];
int f[MAXN];

void exec(int N = MAXN) {
    p[0] = p[1] = 1;

    pval[1] = 1;
    pcnt[1] = 0;
    f[1] = 1;

    for (int i = 2; i < N; i++) {
        if (!p[i]) {
            prime[sz++] = i;
            for (LL j = i; j < N; j *= i) {
                int b = j / i;
                pval[j] = i * pval[b];
                pcnt[j] = pcnt[b] + 1;
                f[j] = _____; // f[j] = f(i^pcnt[j])
            }
        }
        for (int j = 0; i * prime[j] < N; j++) {
            int x = i * prime[j]; p[x] = 1;
            if (i % prime[j] == 0) {
                pval[x] = pval[i] * prime[j];
                pcnt[x] = pcnt[i] + 1;
            } else {
                pval[x] = prime[j];
                pcnt[x] = 1;
            }
            if (x != pval[x]) {
                f[x] = f[x / pval[x]] * f[pval[x]]

```

5f51  
95cf  
95cf  
95cf

b62e  
6589  
e982  
6ae8  
cbf7  
6030  
427e  
76f6  
9628  
427e  
8a8a  
bdda  
c6b9  
427e  
a643  
01d6  
b2b2  
37d9  
758c  
81fd  
e0f3  
a96c  
95cf  
95cf  
34c0  
f87a  
20cc  
9985  
3f93  
8e2e  
cc91  
6322  
95cf  
6191  
d614

```

95cf      }
5f51      if (i % prime[j] == 0) break;
95cf    }
95cf  }
95cf  }
95cf  }

```

## 4.11 Miller-Rabin primality test

The array `a[]` (excluding `senitel`, i.e. `LLONG_MAX`) should be

```

{2}                when  $n < 2,047$ .
{2, 7, 61}         when  $n < 4,759,123,141 (2^{32})$ .
{2, 3, 5, 7, 11}   when  $n < 2.1 \times 10^{12}$ .
{2, 325, 9375, 28178, 450775, 9780504, 1795265022} when  $n < 2^{64}$ .

```

```

f16f bool test(LL n){
59f2   if (n < 3) return n==2;
427e   // ! The array a[] should be modified if the range of x changes.
3f11   const LL a[] = {2LL, 7LL, 61LL, LLONG_MAX};
c320   LL r = 0, d = n-1, x;
f410   while (~d & 1) d >>= 1, r++;
2975   for (int i=0; a[i] < n; i++){
ece1     x = powmod(a[i], d, n); // ! powmod must use for 64bit mulmod
7f99     if (x == 1 || x == n-1) goto next;
e257     rep (i, r) {
d7ff       x = mulmod(x, x, n);
8d2e       if (x == n-1) goto next;
95cf     }
438e     return false;
d490 next:;
95cf   }
3361   return true;
95cf }

```

## 4.12 Pollard's rho algorithm

```

2e6b ULL gcd(ULL a, ULL b) {return b ? gcd(b, a % b) : a;}
427e
54a5 ULL PollardRho(ULL n){
45eb   ULL c, x, y, d = n;

```

```

if (~n&1) return 2;
while (d == n){
  x = y = 2;
  d = 1;
  c = rand() % (n - 1) + 1;
  while (d == 1){
    x = (mulmod(x, x, n) + c) % n;
    y = (mulmod(y, y, n) + c) % n;
    y = (mulmod(y, y, n) + c) % n;
    d = gcd(x>y ? x-y : y-x, n);
  }
}
return d;
}

```

```

d3e5
3c69
0964
4753
5952
9e5b
33d5
e1bf
e1bf
a313
95cf
95cf
5d89
95cf

```

# 5 Graph Theory

## 5.1 Strongly connected component

```

const int MAXV = 100005;

struct graph{
  vector<int> adj[MAXV];
  stack<int> s;
  int V; // number of vertices
  int pre[MAXV], lnk[MAXV], scc[MAXV];
  int time, sccn;

  void add_edge(int u, int v){
    adj[u].push_back(v);
  }

  void dfs(int u){
    pre[u] = lnk[u] = ++time;
    s.push(u);
    for (int v : adj[u]){
      if (!pre[v]){
        dfs(v);
        lnk[u] = min(lnk[u], lnk[v]);
      } else if (!scc[v]){
        lnk[u] = min(lnk[u], pre[v]);
      }
    }
  }
}

```

```

837c
427e
2ea0
88e3
9cad
3d02
8b6c
27ee
427e
bfab
c71a
95cf
427e
d714
7e41
80f6
18f6
173e
5f3c
002c
6068
d5df

```

```

95cf      }
95cf      }
8de2      if (lnk[u] == pre[u]){
660f          sccn++;
3c9e          int x;
a69f          do {
3834              x = s.top(); s.pop();
b0e9              scc[x] = sccn;
6757          } while (x != u);
95cf      }
95cf      }
427e
4c88      void find_scc(){
f4a2          time = sccn = 0;
8de7          memset(scc, 0, sizeof scc);
8c2f          memset(pre, 0, sizeof pre);
6901          Rep (i, V){
56d1              if (!pre[i]) dfs(i);
95cf          }
95cf      }
427e
27ce      vector<int> adjc[MAXV];
364d      void contract(){
1a1e          Rep (i, V)
21a2              rep (j, adj[i].size()){
b730                  if (scc[i] != scc[adj[i][j]])
b46e                      adjc[scc[i]].push_back(scc[adj[i][j]]);
95cf              }
95cf          }
329b      };

```

## 5.2 Vertex biconnected component

```

0f42      const int MAXN = 100005;
2ea0      struct graph {
33ae          int pre[MAXN], iscut[MAXN], bccno[MAXN], dfs_clock, bcc_cnt;
848f          vector<int> adj[MAXN], bcc[MAXN];
6b06          set<pair<int, int>> bcce[MAXN];
427e
76f7          stack<pair<int, int>> s;
427e
bfab          void add_edge(int u, int v) {

```

```

adj[u].push_back(v);
adj[v].push_back(u);
}

int dfs(int u, int fa) {
    int lowu = pre[u] = ++dfs_clock;
    int child = 0;
    for (int v : adj[u]) {
        if (!pre[v]) {
            s.push({u, v});
            child++;
            int lowv = dfs(v, u);
            lowu = min(lowu, lowv);
            if (lowv >= pre[u]) {
                iscut[u] = 1;
                bcc[bcc_cnt].clear();
                bcce[bcc_cnt].clear();
                while (1) {
                    int xu, xv;
                    tie(xu, xv) = s.top(); s.pop();
                    bcce[bcc_cnt].insert({min(xu, xv), max(xu, xv)});
                    if (bccno[xu] != bcc_cnt) {
                        bcc[bcc_cnt].push_back(xu);
                        bccno[xu] = bcc_cnt;
                    }
                    if (bccno[xv] != bcc_cnt) {
                        bcc[bcc_cnt].push_back(xv);
                        bccno[xv] = bcc_cnt;
                    }
                    if (xu == u && xv == v) break;
                }
                bcc_cnt++;
            }
        } else if (pre[v] < pre[u] && v != fa) {
            s.push({u, v});
            lowu = min(lowu, pre[v]);
        }
    }
    if (fa < 0 && child == 1) iscut[u] = 0;
    return lowu;
}

void find_bcc(int n) {
    memset(pre, 0, sizeof pre);

```

```

c71a
a717
95cf
427e
7d3c
9fe6
ec14
18f6
173e
e7f8
fdcf
f851
189c
b687
6323
57eb
90b8
a147
a6a3
a0c3
0ef5
3db2
e0db
d27f
95cf
f357
752b
57c9
95cf
7096
95cf
03f5
95cf
7470
e7f8
f115
95cf
95cf
e104
1160
95cf
427e
17be
8c2f

```

```
e2d2     memset(iscut, 0, sizeof iscut);
40d3     memset(bccno, -1, sizeof bccno);
fae2     dfs_clock = bcc_cnt = 0;
5c63     rep (i, n) if (!pre[i]) dfs(i, -1);
95cf     }
329b     };
```

### 5.3 Cut vertices

If the graph is unconnected, the algorithm should be run on each component.

**Usage:**

tarjan(u, fa)                      Run Tarjan's algorithm on tree rooted at fa. Please call with identical u and fa.

```
9f60     const int MAXN = 200005;
0b32     vector<int> adj[MAXN];
18e4     int dfn[MAXN], low[MAXN], idx;
d39d     bool cut[MAXN];
427e
bfa8     void add_edge(int u, int v) {
c71a         adj[u].push_back(v);
a717         adj[v].push_back(u);
95cf     }
427e
50aa     void tarjan(int u, int fa) {
9891         dfn[u] = low[u] = ++idx;
ec14         int child = 0;
18f6         for (int v : adj[u]) {
3c64             if (!dfn[v]) {
9636                 tarjan(v, fa); low[u] = min(low[u], low[v]);
f368                 if (low[v] >= dfn[u] && u != fa) cut[u] = true;
7923                 child += u == fa;
95cf             }
769a             low[u] = min(low[u], dfn[v]);
95cf         }
7927         if (u == fa && child > 1) cut[u] = true;
95cf     }
```

### 5.4 Minimum spanning arborescence (Chu-Liu)

All vertices are 1-based.

**Usage:**

getans(n, root, edges)                      Compute the total weight of MSA rooted at root.

**Time Complexity:**  $O(|V||E|)$

```
struct edge {
    int u, v;
    LL w;
};

const int MAXN = 10005;
LL in[MAXN];
int pre[MAXN], vis[MAXN], id[MAXN];

LL getans(int n, int rt, vector<edge>& edges) {
    LL ans = 0;
    int cnt = 0;
    while (1) {
        Rep (i, n) in[i] = LLONG_MAX, id[i] = vis[i] = 0;
        for (auto e : edges) {
            if (e.u != e.v and e.w < in[e.v]) {
                pre[e.v] = e.u;
                in[e.v] = e.w;
            }
        }
        in[rt] = 0;
        Rep (i, n) {
            if (in[i] == LLONG_MAX) return -1;
            ans += in[i];
            int u;
            for (u = i; u != rt && vis[u] != i && !id[u]; u = pre[u])
                vis[u] = i;
            if (u != rt && !id[u]) {
                id[u] = ++cnt;
                for (int v = pre[u]; v != u; v = pre[v])
                    id[v] = cnt;
            }
        }
        if (!cnt) return ans;
        Rep (i, n) if (!id[i]) id[i] = ++cnt;
        for (auto& e : edges) {
            LL laz = in[e.v];
            e.u = id[e.u];
            e.v = id[e.v];
            if (e.u != e.v) e.w -= laz;
        }
    }
}
```

```

95cf     }
6cc4     n = cnt; rt = id[rt]; cnt = 0;
95cf     }
95cf }

```

## 5.5 Minimum spanning arborescence, faster

All vertices are 1-based. Clear the fields when reuse the struct.

### Usage:

`add_edge(u, v, w)` Add an edge from  $u$  to  $v$  with weight  $w$ .  
`run(n, rt)` Compute the total weight of MSA rooted at  $rt$ . If not exist, return `LLONG_MIN`.

**Time Complexity:**  $O(|E| \log^2 |V|)$

```

5ece const int MAXN = 300005;
2fef typedef pair<LL, int> pii;
1495 struct MDST {
01b2     priority_queue<pii, vector<pii>, greater<pii>> heap[MAXN];
321d     LL shift[MAXN];
fc06     int fa[MAXN], vis[MAXN];
427e
38dd     int find(int x) { return fa[x] == x ? x : fa[x] = find(fa[x]); }
427e
29b0     void unite(int x, int y) {
0c14         x = find(x); y = find(y); fa[y] = x; if (x == y) return;
6fa0         if (heap[x].size() < heap[y].size()) {
9c26             swap(heap[x], heap[y]);
2ffc             swap(shift[x], shift[y]);
95cf         }
9959         while (heap[y].size()) {
175b             auto p = heap[y].top(); heap[y].pop();
c353             heap[x].emplace(p.first + shift[y] - shift[x], p.second);
95cf         }
95cf     }
427e
0bbd     void add_edge(int u, int v, LL w) { heap[v].emplace(w, u); }
427e
a526     LL run(int n, int rt) {
f7ff         LL ans = 0;
81f2         iota(fa, fa + n + 1, 0);
19b3         Rep (i, n) if (find(i) != find(rt)) {
a7b1             int u = find(i);

```

```

stack<int, vector<int>> s;
while (find(u) != find(rt)) {
    if (vis[u]) while (s.top() != u) {
        vis[s.top()] = 0; unite(u, s.top()); s.pop();
    } else { vis[u] = 1; s.push(u); }
    while (heap[u].size()) {
        ans += heap[u].top().first + shift[u];
        shift[u] = -heap[u].top().first;
        if (find(heap[u].top().second) != u) break;
        heap[u].pop();
    }
    if (heap[u].empty()) return LLONG_MIN;
    u = find(heap[u].top().second);
}
while (s.size()) { vis[s.top()] = 0; unite(rt, s.top()); s.pop(); }
}
return ans;
};

```

## 5.6 Maximum flow (Dinic)

### Usage:

`add_edge(u, v, c)` Add an edge from  $u$  to  $v$  with capacity  $c$ .  
`max_flow(s, t)` Compute maximum flow from  $s$  to  $t$ .

**Time Complexity:** For general graph,  $O(V^2 E)$ ; for network with unit capacity,  $O(\min\{V^{2/3}, \sqrt{E}\} E)$ ; for bipartite network,  $O(\sqrt{V} E)$ .

```

struct edge{
    int from, to;
    LL cap, flow;
};

const int MAXN = 1005;
struct Dinic {
    int n, m, s, t;
    vector<edge> edges;
    vector<int> G[MAXN];
    bool vis[MAXN];
    int d[MAXN];
    int cur[MAXN];

    void add_edge(int from, int to, LL cap) {

```



```

7b55     edges.push_back(edge{from, to, cap, 0});
1db7     edges.push_back(edge{to, from, 0, 0});
fe77     m = edges.size();
dff5     G[from].push_back(m-2);
8f2d     G[to].push_back(m-1);
95cf }
427e
1836 bool bfs() {
3b73     memset(vis, 0, sizeof(vis));
93d2     queue<int> q;
5d13     q.push(s);
2cd2     vis[s] = 1;
721d     d[s] = 0;
cc78     while (!q.empty()) {
66ba         int x = q.front(); q.pop();
3b61         for (int i = 0; i < G[x].size(); i++) {
b510             edge& e = edges[G[x][i]];
bba9             if (!vis[e.to] && e.cap > e.flow) {
cd72                 vis[e.to] = 1;
cf26                 d[e.to] = d[x] + 1;
ca93                 q.push(e.to);
95cf             }
95cf         }
95cf     }
b23b     return vis[t];
95cf }
427e
9252 LL dfs(int x, LL a) {
6904     if (x == t || a == 0) return a;
8bf9     LL flow = 0, f;
f515     for (int& i = cur[x]; i < G[x].size(); i++) {
b510         edge& e = edges[G[x][i]];
2374         if(d[x] + 1 == d[e.to] && (f = dfs(e.to, min(a, e.cap-e.flow))) > 0)
        {
1cce             e.flow += f;
e16d             edges[G[x][i]^1].flow -= f;
a74d             flow += f;
23e5             a -= f;
97ed             if(a == 0) break;
95cf         }
95cf     }
84fb     return flow;
95cf }
427e

```

```

LL max_flow(int s, int t) {
    this->s = s; this->t = t;
    LL flow = 0;
    while (bfs()) {
        memset(cur, 0, sizeof(cur));
        flow += dfs(s, LLONG_MAX);
    }
    return flow;
}

vector<int> min_cut() { // call this after maxflow
    vector<int> ans;
    for (int i = 0; i < edges.size(); i++) {
        edge& e = edges[i];
        if(vis[e.from] && !vis[e.to] && e.cap > 0) ans.push_back(i);
    }
    return ans;
}
};

```

## 5.7 Maximum cardinality bipartite matching (Hungarian)

```

#include <bits/stdc++.h>
using namespace std;

#define rep(i, n) for (int i = 0; i < (n); i++)
#define Rep(i, n) for (int i = 1; i <= (n); i++)
#define range(x) (x).begin(), (x).end()
typedef long long LL;

struct Hungarian{
    int nx, ny;
    vector<int> mx, my;
    vector<vector<int>> > e;
    vector<bool> mark;

    void init(int nx, int ny){
        this->nx = nx;
        this->ny = ny;
        mx.resize(nx); my.resize(ny);
        e.clear(); e.resize(nx);
        mark.resize(nx);
    }
};

```

```

95cf    }
427e
4589    inline void add(int a, int b){
486c        e[a].push_back(b);
95cf    }
427e
0c2b    bool augment(int i){
207c        if (!mark[i]) {
dae4            mark[i] = true;
6a1e            for (int j : e[i]){
0892                if (my[j] == -1 || augment(my[j])){
9ca3                    mx[i] = j; my[j] = i;
3361                    return true;
95cf                }
95cf            }
95cf        }
438e        return false;
95cf    }
427e
3fac    int match(){
5b57        int ret = 0;
b0f1        fill(range(mx), -1);
b957        fill(range(my), -1);
4ed1        rep (i, nx){
13a5            fill(range(mark), false);
cc89            if (augment(i)) ret++;
95cf        }
ee0f        return ret;
95cf    }
329b    };

```

## 5.8 Maximum matching of general graph (Edmond's blossom)

### Usage:

init(n)	Initialize the template with $n$ vertices, numbered from 1.
add_edge(u, v)	Add an undirected edge $uv$ .
solve()	Find the maximum matching. Return the number of matched edges.
mate[]	The mate of a matched vertex. If it is not matched, then the value is 0.

**Time Complexity:**  $O(|V|^3)$ , but extremely fast in practice.

```

const int MAXN = 1024;
struct Blossom {
    vector<int> adj[MAXN];
    queue<int> q;
    int n;
    int label[MAXN], mate[MAXN], save[MAXN], used[MAXN];

    void init(int nv) {
        n = nv; for (auto& v : adj) v.clear();
        fill(range(label), 0); fill(range(mate), 0);
        fill(range(save), 0); fill(range(used), 0);
    }

    void add_edge(int u, int v) { adj[u].push_back(v); adj[v].push_back(u); }

    void rematch(int x, int y) {
        int m = mate[x]; mate[x] = y;
        if (mate[m] == x) {
            if (label[x] <= n) {
                mate[m] = label[x]; rematch(label[x], m);
            } else {
                int a = 1 + (label[x] - n - 1) / n;
                int b = 1 + (label[x] - n - 1) % n;
                rematch(a, b); rematch(b, a);
            }
        }
    }

    void traverse(int x) {
        Rep (i, n) save[i] = mate[i];
        rematch(x, x);
        Rep (i, n) {
            if (mate[i] != save[i]) used[i] ++;
            mate[i] = save[i];
        }
    }

    void relabel(int x, int y) {
        Rep (i, n) used[i] = 0;
        traverse(x); traverse(y);
        Rep (i, n) {
            if (used[i] == 1 and label[i] < 0) {
                label[i] = n + x + (y - 1) * n;
                q.push(i);
            }
        }
    }
};

```

c041  
6ab1  
0b32  
93d2  
5c83  
0de2  
427e  
2186  
3728  
477d  
bb35  
95cf  
427e  
c2dd  
427e  
2a48  
8af8  
1aa4  
f4ba  
740a  
8e2e  
3341  
2885  
ef33  
95cf  
95cf  
95cf  
427e  
8a50  
43c0  
2ef7  
34d7  
62c5  
97ef  
95cf  
95cf  
427e  
8bf8  
d101  
c4ea  
34d7  
dee9  
1c22  
eb31

```

95cf    }
95cf    }
95cf    }
427e
a0ce    int solve() {
34d7        Rep (i, n) {
a073            if (mate[i]) continue;
1fc0            Rep (j, n) label[j] = -1;
7676            label[i] = 0; q = queue<int>(); q.push(i);
1c7d            while (q.size()) {
66ba                int x = q.front(); q.pop();
b98c                for (int y : adj[x]) {
c07f                    if (mate[y] == 0 and i != y) {
7f36                        mate[y] = x; rematch(x, y); q = queue<int>(); break;
95cf                    }
d315                    if (label[y] >= 0) { relabel(x, y); continue; }
58ec                    if (label[mate[y]] < 0) {
c9c4                        label[mate[y]] = x; q.push(mate[y]);
95cf                    }
95cf                }
95cf            }
95cf        }
8abb        int cnt = 0;
b52f        Rep (i, n) cnt += (mate[i] > i);
6808        return cnt;
95cf    }
329b    };

```

## 5.9 Minimum cost maximum flow

```

bcf8    struct edge{
60e2        int from, to;
d698        int cap, flow;
32cc        LL cost;
329b    };
427e
cc3e    const LL INF = LLONG_MAX / 2;
2aa8    const int MAXN = 5005;
c6cb    struct MCMF {
9ceb        int s, t, n, m;
9f0c        vector<edge> edges;
b891        vector<int> G[MAXN];

```

```

bool inq[MAXN]; // queue
LL d[MAXN];     // distance
int p[MAXN];    // previous
int a[MAXN];    // improvement

void add_edge(int from, int to, int cap, LL cost) {
    edges.push_back(edge{from, to, cap, 0, cost});
    edges.push_back(edge{to, from, 0, 0, -cost});
    m = edges.size();
    G[from].push_back(m-2);
    G[to].push_back(m-1);
}

bool spfa(){
    queue<int> q;
    fill(d, d + MAXN, INF); d[s] = 0;
    memset(inq, 0, sizeof(inq));
    q.push(s); inq[s] = true;
    p[s] = 0; a[s] = INT_MAX;
    while (!q.empty()){
        int u = q.front(); q.pop(); inq[u] = false;
        for (int i : G[u]) {
            edge& e = edges[i];
            if (e.cap > e.flow && d[e.to] > d[u] + e.cost){
                d[e.to] = d[u] + e.cost;
                p[e.to] = G[u][i];
                a[e.to] = min(a[u], e.cap - e.flow);
                if (!inq[e.to]) q.push(e.to), inq[e.to] = true;
            }
        }
    }
    return d[t] != INF;
}

void augment(){
    int u = t;
    while (u != s){
        edges[p[u]].flow += a[t];
        edges[p[u]^1].flow -= a[t];
        u = edges[p[u]].from;
    }
}

#ifdef GIVEN_FLOW

```

f74f  
8f67  
9524  
b330  
427e  
f7f2  
24f0  
95f0  
fe77  
dff5  
8f2d  
95cf  
427e  
3c52  
93d2  
8494  
fd48  
5e7c  
2dae  
cc78  
b0aa  
3bba  
56d8  
3601  
55bc  
0bea  
8249  
e5d3  
95cf  
95cf  
95cf  
6d7c  
95cf  
427e  
71a4  
06f1  
b19d  
db09  
25a9  
e6c9  
95cf  
95cf  
427e  
6e20

```

5972 bool min_cost(int s, int t, int f, LL& cost) {
590d     this->s = s; this->t = t;
21d4     int flow = 0;
23cb     cost = 0;
22dc     while (spfa()) {
bcd8         augment();
a671         if (flow + a[t] >= f){
b14d             cost += (f - flow) * d[t]; flow = f;
3361             return true;
8e2e         } else {
2a83             flow += a[t]; cost += a[t] * d[t];
95cf         }
95cf     }
438e     return false;
95cf }
a8cb #else
f9a9     int min_cost(int s, int t, LL& cost) {
590d         this->s = s; this->t = t;
21d4         int flow = 0;
23cb         cost = 0;
22dc         while (spfa()) {
bcd8             augment();
2a83             flow += a[t]; cost += a[t] * d[t];
95cf         }
84fb         return flow;
95cf     }
1937 #endif
329b };

```

## 5.10 Global minimum cut (Stoer-Wagner)

### Usage:

stoer(w) Compute the global minimum cut of the graph specified by the **symmetric** adjacency matrix w (0-based). Return the capacity of the cut and the indices of one part of the cut.

**Time Complexity:**  $O(|V|^3)$

```

f9d7 typedef vector<LL> VI;
045e typedef vector<VI> VWI;
427e
f012 pair<LL, VI> stoer(WWI &w) {
66f7     int n = w.size();

```

```

VI used(n), c, bestc;
LL bestw = -1;

for (int ph = n - 1; ph >= 0; ph--) {
    VI wt = w[0], added = used;
    int prev, last = 0;
    rep (i, ph) {
        prev = last;
        last = -1;
        for (int j = 1; j < n; j++)
            if (!added[j] && (last == -1 || wt[j] > wt[last]))
                last = j;
        if (i == ph - 1) {
            rep (j, n) w[prev][j] += w[last][j];
            rep (j, n) w[j][prev] = w[prev][j];
            used[last] = true;
            c.push_back(last);
            if (bestw == -1 || wt[last] < bestw) {
                bestc = c;
                bestw = wt[last];
            }
        } else {
            rep (j, n) wt[j] += w[last][j];
            added[last] = true;
        }
    }
}

return {bestw, bestc};
}

```

## 5.11 Fast LCA

All indices of the tree are 1-based.

### Usage:

preprocess(root) Initialize with tree rooted at root.  
lca(u, v) Query the lowest common ancestor of  $u$  and  $v$ .

```

const int MAXN = 500005;
vector<int> adj[MAXN];
int id[MAXN], nid;
pair<int, int> st[MAXN << 1][33 - __builtin_clz(MAXN)];

void dfs(int u, int p, int d) {

```

```

0df2     st[id[u] = nid++][0] = {d, u};
18f6     for (int v : adj[u]) {
bd87         if (v == p) continue;
f58c         dfs(v, u, d + 1);
08ad         st[nid++][0] = {d, u};
95cf     }
95cf }
427e
3d1b void preprocess(int root) {
3269     nid = 0;
91e1     dfs(root, 0, 1);
5e98     int l = 31 - __builtin_clz(nid);
213b     rep (j, l) rep (i, 1+nid-(1<<j))
1131         st[i][j+1] = min(st[i][j], st[i+(1<<j)][j]);
95cf }
427e
0f0b int lca(int u, int v) {
cfc4     tie(u, v) = minmax(id[u], id[v]);
be9b     int k = 31 - __builtin_clz(v-u+1);
8ebc     return min(st[u][k], st[v-(1<<k)+1][k]).second;
95cf }

```

## 5.12 Heavy-light decomposition

**Time Complexity:** The decomposition itself takes linear time. Each query takes  $O(\log n)$  operations.

```

0f42 const int MAXN = 100005;
0b32 vector<int> adj[MAXN];
42f2 int sz[MAXN], top[MAXN], fa[MAXN], son[MAXN], depth[MAXN], id[MAXN];
427e
be5c void dfs1(int x, int dep, int par){
7489     depth[x] = dep;
2ee7     sz[x] = 1;
adb4     fa[x] = par;
b79d     int maxn = 0, s = 0;
c861     for (int c: adj[x]){
fe45         if (c == par) continue;
fd2f         dfs1(c, dep + 1, x);
b790         sz[x] += sz[c];
f0f1         if (sz[c] > maxn){
c749             maxn = sz[c];
fe19             s = c;

```

```

    }
    }
    son[x] = s;
}

int cid = 0;
void dfs2(int x, int t){
    top[x] = t;
    id[x] = ++cid;
    if (son[x]) dfs2(son[x], t);
    for (int c: adj[x]){
        if (c == fa[x]) continue;
        if (c == son[x]) continue;
        else dfs2(c, c);
    }
}

void decomp(int root){
    dfs1(root, 1, 0);
    dfs2(root, root);
}

void query(int u, int v){
    while (top[u] != top[v]){
        if (depth[top[u]] < depth[top[v]]) swap(u, v);
        // id[top[u]] to id[u]
        u = fa[top[u]];
    }
    if (depth[u] > depth[v]) swap(u, v);
    // id[u] to id[v]
}

```

95cf  
95cf  
0e08  
95cf  
427e  
ba54  
3644  
8d96  
d314  
c4a1  
c861  
9881  
5518  
13f9  
95cf  
95cf  
427e  
0f04  
9fa4  
1c88  
95cf  
427e  
2c98  
03a1  
45ec  
427e  
005b  
95cf  
6083  
427e  
95cf

## 5.13 Centroid decomposition

Note that the centroid here is not the exact centroid of the graph. It only guarantees that the size of each subtree does not exceed half of that of the original tree. This is enough to guarantee the correct time complexity. All vertices are numbered from 1. Call `decomp(root)` to use.

**Usage:**

`decomp(u, p)`      Decompose the tree rooted at  $u$  with parent  $p$ .

**Time Complexity:** The decomposition itself takes  $O(n \log n)$  time.

```

1fb6 vector<int> adj[100005];
88e0 int sz[100005], sum;
427e
f93d void getsz(int u, int p) {
5b36     sz[u] = 1; sum++;
18f6     for (int v : adj[u]) {
bd87         if (v == p) continue;
e3cb         getsz(v, u);
8449         sz[u] += sz[v];
95cf     }
95cf }
427e
67f9 int getcent(int u, int p) {
d51f     for (int v : adj[u])
76e4         if (v != p and sz[v] > sum / 2)
18e3             return getcent(v, u);
81b0     return u;
95cf }
427e
4662 void decompose(int u) {
618e     sum = 0; getsz(u, 0);
303c     u = getcent(u, 0); // update u to the centroid
427e
18f6     for (int v : adj[u]) {
427e         // get answer for subtree v
95cf     }
427e     // get answer for the whole tree
427e     // don't forget to count the centroid itself
427e
18f6     for (int v : adj[u]) { // divide and conquer
c375         adj[v].erase(find(range(adj[v]), u));
fa6b         decompose(v);
a717         adj[v].push_back(u); // restore deleted edge
95cf     }
95cf }

```

## 5.14 DSU on tree

This implementation avoids parallel existence of multiple data structures but requires that the data structure is invertible. To use this template, implement merge, enter, leave as needed; first call decomp(root, 0), then call work(root, 0, false). Labels of vertices start from 1.

### Usage:

decomp(u, p)                      Decompose the tree *u*.  
work(u, p, keep)                  Work for subtree *u*. When keep is set, information is not cleared.

**Time Complexity:**  $O(n \log n)$  times the complexity for merge, enter, leave.

```

vector<int> adj[100005];
int sz[100005], son[100005];

void decomp(int u, int p) {
    sz[u] = 1;
    for (int v : adj[u]) {
        if (v == p) continue;
        decomp(v, u);
        sz[u] += sz[v];
        if (sz[v] > sz[son[u]]) son[u] = v;
    }
}

template <typename T>
void trav(T fn, int u, int p) {
    fn(u);
    for (int v : adj[u]) if (v != p) trav(fn, v, u);
}

#define for_light(v) for (int v : adj[u]) if (v != p and v != son[u])
void work(int u, int p, bool keep) {
    for_light(v) work(v, u, 0); // process light children

    // process heavy child
    // current data structure contains info of heavy child
    if (son[u]) work(son[u], u, 1);

    auto merge = [u] (int c) { /* count contribution of c */ };
    auto enter = [] (int c) { /* add vertex c */ };
    auto leave = [] (int c) { /* remove vertex c */ };

    for_light(v) {
        trav(merge, v, u);
        trav(enter, v, u);
    }

    // count answer for root and add it
    // Warning: special check may apply to root!

```

```

c54f     merge(u);
9dec     enter(u);
427e
427e     // Leave current tree
4e3e     if (!keep) trav(leave, u, p);
95cf     }

```

## 6 Data Structures

### 6.1 Fenwick tree (point update range query)

```

9976 struct bit_purq { // point update, range query
d7af     int N;
99ff     vector<LL> tr;
427e
d34f     void init(int n) { // fill the array with 0
1010         tr.resize(N = n + 5);
95cf     }
427e
63d0     LL sum(int n) {
f7ff         LL ans = 0;
e290         while (n) {
0715             ans += tr[n];
c0d4             n &= n - 1;
95cf         }
4206         return ans;
95cf     }
427e
f4bd     void add(int n, LL x){
ad20         while (n < N) {
6c81             tr[n] += x;
0af5             n += n & -n;
95cf         }
95cf     }
329b };

```

### 6.2 Fenwick tree (range update point query)

```

3d03 struct bit_rupq{ // range update, point query
d7af     int N;

```

```

vector<LL> tr;

void init(int n) { // fill the array with 0
    tr.resize(N = n + 5);
}

LL query(int n) {
    LL ans = 0;
    while (n < N) {
        ans += tr[n];
        n += n & -n;
    }
    return ans;
}

void add(int n, LL x) {
    while (n){
        tr[n] += x;
        n &= n - 1;
    }
}
};

```

### 6.3 Segment tree

```

LL p;
const int MAXN = 4 * 100006;
struct segtree {
    int l[MAXN], m[MAXN], r[MAXN];
    LL val[MAXN], tadd[MAXN], tmul[MAXN];

#define lson (o<<1)
#define rson (o<<1|1)

    void pull(int o) {
        val[o] = (val[lson] + val[rson]) % p;
    }

    void push_add(int o, LL x) {
        val[o] = (val[o] + x * (r[o] - l[o])) % p;
        tadd[o] = (tadd[o] + x) % p;
    }
}

```

```

427e void push_mul(int o, LL x) {
d658     val[o] = val[o] * x % p;
b82c     tadd[o] = tadd[o] * x % p;
aa86     tmul[o] = tmul[o] * x % p;
649f }
95cf
427e void push(int o) {
b149     if (l[o] == m[o]) return;
3159     if (tmul[o] != 1) {
0a90         push_mul(lson, tmul[o]);
0f4a         push_mul(rson, tmul[o]);
045e         tmul[o] = 1;
ac0a     }
95cf
1b82     if (tadd[o]) {
9547         push_add(lson, tadd[o]);
0e73         push_add(rson, tadd[o]);
6234         tadd[o] = 0;
95cf     }
95cf }
427e
471c void build(int o, int ll, int rr) {
0e87     int mm = (ll + rr) / 2;
9d27     l[o] = ll; r[o] = rr; m[o] = mm;
ac0a     tmul[o] = 1;
5c92     if (ll == mm) {
001f         scanf("%lld", val + o);
e5b6         val[o] %= p;
8e2e     } else {
7293         build(lson, ll, mm);
5e67         build(rson, mm, rr);
ba26         pull(o);
95cf     }
95cf }
427e
4406 void add(int o, int ll, int rr, LL x) {
3c16     if (ll <= l[o] && r[o] <= rr) {
db32         push_add(o, x);
8e2e     } else {
c4b0         push(o);
4305         if (m[o] > ll) add(lson, ll, rr, x);
d5a6         if (m[o] < rr) add(rson, ll, rr, x);
ba26         pull(o);
95cf     }

```

```

}
void mul(int o, int ll, int rr, LL x) {
    if (ll <= l[o] && r[o] <= rr) {
        push_mul(o, x);
    } else {
        push(o);
        if (ll < m[o]) mul(lson, ll, rr, x);
        if (m[o] < rr) mul(rson, ll, rr, x);
        pull(o);
    }
}

LL query(int o, int ll, int rr) {
    if (ll <= l[o] && r[o] <= rr) {
        return val[o];
    } else {
        push(o);
        if (rr <= m[o]) return query(lson, ll, rr);
        if (ll >= m[o]) return query(rson, ll, rr);
        return query(lson, ll, rr) + query(rson, ll, rr);
    }
}
} seg;

```

95cf  
427e  
48cd  
3c16  
e7d0  
8e2e  
c4b0  
d1ba  
67f3  
ba26  
95cf  
95cf  
427e  
0f62  
3c16  
6dfe  
8e2e  
c4b0  
462a  
5cca  
bbf9  
95cf  
95cf  
4d99

## 6.4 Treap

Self-balanced binary search tree which supports split and merge.

### Usage:

push(x)	Push lazy tags to children.
pull(x)	Update statistics of node $x$ .
Init(x, v)	Initialize node $x$ with value $v$ .
Add(x, v)	Apply addition to subtree $x$ .
Reverse(x)	Apply reversion to subtree $x$ .
Merge(x, y)	Merge trees rooted at $x$ and $y$ . Return the root of new tree.
Split(t, k, x, y)	Split out the left $k$ elements of tree $t$ . The roots of left part and right part are stored in $x$ and $y$ , respectively.
init(n)	Initialize the treap with array of size $n$ .
work(op, l, r)	Range operation over $[l, r)$ .

**Time Complexity:** Expected  $O(\log n)$  per operation.

```
const int MAXN = 200005;
```

9f60



```

a7c5 mt19937 gen(time(NULL));
9542 struct Treap {
6d61     int ch[MAXN][2];
3948     int sz[MAXN], key[MAXN], val[MAXN];
5d9a     int add[MAXN], rev[MAXN];
2b1b     LL sum[MAXN] = {0};
a773     int maxv[MAXN] = {INT_MIN}, minv[MAXN] = {INT_MAX};
427e
a629     void Init(int x, int v) {
5a00         ch[x][0] = ch[x][1] = 0;
d8cd         key[x] = gen(); val[x] = v; pull(x);
95cf     }
427e
3bf9     void pull(int x) {
e1c3         sz[x] = 1 + sz[ch[x][0]] + sz[ch[x][1]];
99f8         sum[x] = val[x] + sum[ch[x][0]] + sum[ch[x][1]];
94e9         maxv[x] = max({val[x], maxv[ch[x][0]], maxv[ch[x][1]]});
6bb9         minv[x] = min({val[x], minv[ch[x][0]], minv[ch[x][1]]});
95cf     }
427e
8c8e     void Add(int x, int a) {
a7b1         val[x] += a; add[x] += a;
832a         sum[x] += LL(sz[x]) * a; maxv[x] += a; minv[x] += a;
95cf     }
427e
aaf6     void Reverse(int x) {
52c6         rev[x] ^= 1;
7850         swap(ch[x][0], ch[x][1]);
95cf     }
427e
1a53     void push(int x) {
5fe5         for (int c : ch[x]) if (c) {
fd76             Add(c, add[x]);
7a53             if (rev[x]) Reverse(c);
95cf         }
49ee         add[x] = 0; rev[x] = 0;
95cf     }
427e
9d2c     int Merge(int x, int y) {
1b09         if (!x || !y) return x | y;
cd7e         push(x); push(y);
bfffa        if (key[x] > key[y]) {
a3df             ch[x][1] = Merge(ch[x][1], y); pull(x); return x;
8e2e         } else {

```

```

        ch[y][0] = Merge(x, ch[y][0]); pull(y); return y;
    }
}

void Split(int t, int k, int &x, int &y) {
    if (t == 0) { x = y = 0; return; }
    push(t);
    if (sz[ch[t][0]] < k) {
        x = t; Split(ch[t][1], k - sz[ch[t][0]] - 1, ch[t][1], y);
    } else {
        y = t; Split(ch[t][0], k, x, ch[t][0]);
    }
    if (x) pull(x); if (y) pull(y);
}

} treap;

int root;

void init(int n) {
    Rep(i, n) {
        int x; scanf("%d", &x);
        treap.Init(i, x);
        root = (i == 1) ? 1 : treap.Merge(root, i);
    }
}

void work(int op, int l, int r) {
    int tl, tm, tr;
    treap.Split(root, l, tl, tm);
    treap.Split(tm, r - l, tm, tr);
    if (op == 1) {
        int x; scanf("%d", &x); treap.Add(tm, x);
    } else if (op == 2) {
        treap.Reverse(tm);
    } else if (op == 3) {
        printf("%lld %d %d\n",
            treap.sum[tm], treap.minv[tm], treap.maxv[tm]);
    }
    root = treap.Merge(treap.Merge(tl, tm), tr);
}

```

```

bf9e
95cf
95cf
427e
dc7e
6303
f26b
3465
ffdc
8e2e
8a23
95cf
89e3
95cf
b1f4
427e
24b6
427e
d34f
34d7
7681
0ed8
bcc8
95cf
95cf
427e
d030
6639
b6c4
8de3
3658
c039
1dcb
ae78
581d
e092
867f
95cf
6188
95cf

```

## 6.5 Link/cut tree

Dynamic connectivity of undirected acyclic graph. Support single-vertex update, path aggregation and relative LCA query. Vertices are numbered from 1. Zero initialization is enough except for the statistic information.

### Usage:

<code>pull(x)</code>	Update statistics of node $x$ .
<code>Root(u)</code>	Get the root of tree where vertex $u$ is in.
<code>Link(u, v)</code>	Link two unconnected trees.
<code>Cut(u, v)</code>	Cut an existent edge.
<code>Query(u, v)</code>	Path aggregation.
<code>Update(u, x)</code>	Single point modification.
<code>LCA(u, v, root)</code>	Get the lowest common ancestor of $u$ and $v$ in tree rooted at root.

**Time Complexity:**  $O(\log n)$  per operation

```

2e73 const int MAXN = 1000005;
ca06 struct LCT {
6a6d     int fa[MAXN], ch[MAXN][2], val[MAXN], sum[MAXN];
c6e1     bool rev[MAXN];

eba3     bool isroot(int x) { return ch[fa[x]][0] == x || ch[fa[x]][1] == x; }
f19f     void pull(int x) { sum[x] = val[x] ^ sum[ch[x][0]] ^ sum[ch[x][1]]; }
1c4d     void reverse(int x) { swap(ch[x][0], ch[x][1]); rev[x] ^= 1; }
1a53     void push(int x) {
89a0         if (rev[x]) rep (i, 2) if (ch[x][i]) reverse(ch[x][i]); rev[x] = 0;
95cf     }
425f     void rotate(int x) {
51af         int y = fa[x], z = fa[y], k = ch[y][1] == x, w = ch[x][!k];
e1fe         if (isroot(y)) ch[z][ch[z][1] == y] = x;
1e6f         ch[x][!k] = y; ch[y][k] = w; if (w) fa[w] = y;
6d09         fa[y] = x; fa[x] = z; pull(y);
95cf     }
52c6     void pushall(int x) { if (isroot(x)) pushall(fa[x]); push(x); }
f69c     void splay(int x) {
d095         int y = x, z = 0;
c494         for (pushall(y); isroot(x); rotate(x)) {
ceef             y = fa[x]; z = fa[y];
4449             if (isroot(y)) rotate((ch[y][0] == x) ^ (ch[z][0] == y) ? x : y);
95cf         }
78a0         pull(x);
95cf     }
6229     void access(int x) {

```

```

        int z = x;
        for (int y = 0; x; x = fa[y = x]) { splay(x); ch[x][1] = y; pull(x); }
        splay(z);
    }
    void chroot(int x) { access(x); reverse(x); }
    void split(int x, int y) { chroot(x); access(y); }

    int Root(int x) {
        for (access(x); ch[x][0]; x = ch[x][0]) push(x);
        splay(x); return x;
    }
    void Link(int u, int v) { chroot(u); fa[u] = v; }
    void Cut(int u, int v) { split(u, v); fa[u] = ch[v][0] = 0; pull(v); }
    int Query(int u, int v) { split(u, v); return sum[v]; }
    void Update(int u, int x) { splay(u); val[u] = x; }
    int LCA(int x, int y, int root) {
        chroot(root); access(x); splay(y);
        while (fa[y]) splay(y = fa[y]);
        return y;
    }
};

```

1548  
8854  
7afd  
95cf  
a067  
126d  
427e  
d87a  
f4f1  
0d77  
95cf  
9e46  
7c10  
0691  
a999  
1f42  
6cb2  
02e5  
c218  
95cf  
329b

## 6.6 Balanced binary search tree from pb\_ds

```

#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;

tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update>
    rkt;
// null_tree_node_update

// SAMPLE USAGE
rkt.insert(x);           // insert element
rkt.erase(x);           // erase element
rkt.order_of_key(x);     // obtain the number of elements less than x
rkt.find_by_order(i);    // iterator to i-th (numbered from 0) smallest element
rkt.lower_bound(x);
rkt.upper_bound(x);
rkt.join(rkt2);          // merge tree (only if their ranges do not intersect)
rkt.split(x, rkt2);      // split all elements greater than x to rkt2

```

0475  
332d  
427e  
43a7  
427e  
427e  
427e  
190e  
05d4  
add5  
b064  
c103  
4ff4  
b19b  
cb47

## 6.7 Persistent segment tree, range k-th query

```

f1a7 struct node {
2ff6     static int n, pos;
427e
7cec     int value;
70e2     node *left, *right;
427e
20b0     void* operator new(size_t size);
427e
3dc0     static node* Build(int l, int r) {
b6c5         node* a = new node;
ce96         if (r > l + 1) {
181e             int mid = (l + r) / 2;
3ba2             a->left = Build(l, mid);
8aaf             a->right = Build(mid, r);
8e2e         } else {
bfc4             a->value = 0;
95cf         }
5ffd         return a;
95cf     }
427e
5a45     static node* init(int size) {
2c46         n = size;
7ee3         pos = 0;
be52         return Build(0, n);
95cf     }
427e
93c0     static int Query(node* lt, node *rt, int l, int r, int k) {
d30c         if (r == l + 1) return l;
181e         int mid = (l + r) / 2;
cb5a         if (rt->left->value - lt->left->value < k) {
8edb             k -= rt->left->value - lt->left->value;
2412             return Query(lt->right, rt->right, mid, r, k);
8e2e         } else {
0119             return Query(lt->left, rt->left, l, mid, k);
95cf         }
95cf     }
427e
c9ad     static int query(node* lt, node *rt, int k) {
9e27         return Query(lt, rt, 0, n, k);
95cf     }
427e

```

```

node *Inc(int l, int r, int pos) const {
    node* a = new node(*this);
    if (r > l + 1) {
        int mid = (l + r) / 2;
        if (pos < mid)
            a->left = left->Inc(l, mid, pos);
        else
            a->right = right->Inc(mid, r, pos);
    }
    a->value++;
    return a;
}

node *inc(int index) {
    return Inc(0, n, index);
}
} nodes[8000000];

int node::n, node::pos;
inline void* node::operator new(size_t size) {
    return nodes + (pos++);
}

```

b19c  
5794  
ce96  
181e  
203d  
f44a  
649a  
1024  
95cf  
2b3e  
5ffd  
95cf  
427e  
e80f  
c246  
95cf  
865a  
427e  
99ce  
1987  
bb3c  
95cf

## 6.8 Block list

All indices are 0-based. All ranges are left-closed right-open.

### Usage:

block::fix()	Apply tags to the current block.
Init(l, r)	Range initializer.
Reverse(l, r)	Reverse the range.
Add(l, r, x)	Add $x$ to the range.
Query(l, r)	Range aggregation.

```

const int BLOCK = 800;
typedef vector<int> vi;

struct block {
    vi data;
    LL sum; int minv, maxv;
    int add; bool rev;

    block(vi&& vec) : data(move(vec)),
        sum(accumulate(range(data), 0ll)),

```

fd9e  
76b3  
427e  
a771  
8fbc  
e3b5  
41db  
427e  
d7eb  
1f0c

```

8216     minv(*min_element(range(data))),
527d     maxv(*max_element(range(data))),
6437     add(0), rev(0) { }
427e
b919 void fix() {
0694     if (rev) reverse(range(data));         rev = 0;
0527     if (add) for (int& x : data) x += add;  add = 0;
95cf }
427e
8bc4 void merge(block& another) {
b895     fix(); another.fix();
f516     vi temp(move(data));
d02c     temp.insert(temp.end(), range(another.data));
88ea     *this = block(move(temp));
95cf }
427e
42e8 block split(int pos) {
3e79     fix();
ccab     block result(vi(data.begin() + pos, data.end()));
861a     data.resize(pos); *this = block(move(data));
56b0     return result;
95cf }
329b };
427e
2a18 typedef list<block>::iterator lit;
427e
ce14 struct blocklist {
5540     list<block> blk;
427e
7b8e void maintain() {
3131     lit it = blk.begin();
4628     while (it != blk.end() && next(it) != blk.end()) {
852d         lit it2 = it;
188c         while (next(it2) != blk.end() &&
3600             it2->data.size() + next(it2)->data.size() <= BLOCK) {
93e1             it2->merge(*next(it2));
e1fa             blk.erase(next(it2));
95cf         }
5771         ++it;
95cf     }
95cf }
427e
b7b3 lit split(int pos) {
2273     for (lit it = blk.begin(); ; it++) {

```

```

        if (pos == 0) return it;
        while (it->data.size() > pos)
            blk.insert(next(it), it->split(pos));
        pos -= it->data.size();
    }
}

void Init(int *l, int *r) {
    for (int *cur = l; cur < r; cur += BLOCK)
        blk.emplace_back(vi(cur, min(cur + BLOCK, r)));
}

void Reverse(int l, int r) {
    lit it = split(l), it2 = split(r);
    reverse(it, it2);
    while (it != it2) {
        it->rev ^= 1;
        it++;
    }
    maintain();
}

void Add(int l, int r, int x) {
    lit it = split(l), it2 = split(r);
    while (it != it2) {
        it->sum += LL(x) * it->data.size();
        it->minv += x; it->maxv += x;
        it->add += x; it++;
    }
    maintain();
}

void Query(int l, int r) {
    lit it = split(l), it2 = split(r);
    LL sum = 0; int minv = INT_MAX, maxv = INT_MIN;
    while (it != it2) {
        sum += it->sum;
        minv = min(minv, it->minv);
        maxv = max(maxv, it->maxv);
        it++;
    }
    maintain();
    printf("%lld_%d_%d\n", sum, minv, maxv);

```

```

5502
8e85
2099
a5a1
427e
95cf
95cf
427e
1c7b
9919
8950
95cf
427e
a22f
997b
dfd0
8f89
6a06
5283
95cf
b204
95cf
427e
3cce
997b
8f89
e927
03d3
4511
95cf
b204
95cf
427e
3ad3
997b
c33d
8f89
e472
72c4
e1c4
5283
95cf
b204
8792

```

```
95cf     }
958e } lst;
```

## 6.9 Persistent block list

Block list that supports persistence. All indices are 0-based. All ranges are left-closed right-open. `std::shared_ptr` is used to ease memory management. One should modify the constructor of `block` to maintain extra information. Here we use this policy that the size of each block does not exceed `BLOCK`, while the sum of sizes of two adjacent blocks does not less than `BLOCK`.

When some operation that breaks block list property, please call `maintain` in time to restore the property.

### Usage:

<code>maintain()</code>	Maintain the block list property.
<code>split(pos)</code>	Split the block list at position <code>pos</code> . Returns an iterator to a block starting at <code>pos</code> .
<code>sum(l, r)</code>	An example function of list traversal between $[l, r)$ .

**Time Complexity:** When `BLOCK` is properly selected, the time complexity is  $O(\sqrt{n})$  per operation.

```
a19e constexpr int BLOCK = 800;
76b3 typedef vector<int> vi;
0563 typedef shared_ptr<vi> pvi;
013b typedef shared_ptr<const vi> pcvi;
427e
a771 struct block {
2989     pcvi data;
8fd0     LL sum;
427e
427e     // add information to maintain
a613     block(pcvi ptr) :
24b5         data(ptr),
0cf0         sum(accumulate(ptr->begin(), ptr->end(), 0ll))
e93b     { }
427e
5c0f     void merge(const block& another) {
0b18         pvi temp = make_shared<vi>(data->begin(), data->end());
ac21         temp->insert(temp->end(), another.data->begin(), another.data->end());
6467         *this = block(temp);
95cf     }
427e
42e8     block split(int pos) {
```

```
        block result(make_shared<vi>(data->begin() + pos, data->end()));
        *this = block(make_shared<vi>(data->begin(), data->begin() + pos));
        return result;
    }
};

typedef list<block>::iterator lit;

struct blocklist {
    list<block> blk;

    void maintain() {
        lit it = blk.begin();
        while (it != blk.end() and next(it) != blk.end()) {
            lit it2 = it;
            while (next(it2) != blk.end() and
                    it2->data->size() + next(it2)->data->size() <= BLOCK) {
                it2->merge(*next(it2));
                blk.erase(next(it2));
            }
            ++it;
        }
    }

    lit split(int pos) {
        for (lit it = blk.begin(); ; it++) {
            if (pos == 0) return it;
            while (it->data->size() > pos) {
                blk.insert(next(it), it->split(pos));
            }
            pos -= it->data->size();
        }
    }

    LL sum(int l, int r) { // traverse
        lit it1 = split(l), it2 = split(r);
        LL res = 0;
        while (it1 != it2) {
            res += it1->sum;
            it1++;
        }
        maintain();
        return res;
    }
}
```

```
dac1
01db
56b0
95cf
329b
427e
2a18
427e
ce14
5540
427e
7b8e
3131
5e44
852d
0b03
029f
93e1
e1fa
95cf
5771
95cf
95cf
427e
b7b3
2273
5502
d480
2099
95cf
a1c8
95cf
95cf
427e
fd38
48b4
ac09
9f1d
8284
61fd
95cf
b204
244d
95cf
```

329b };

## 6.10 Sparse table, range extremum query

The array is 0-based and the range is closed.

```
db63 const int MAXN = 100007;
b330 int a[MAXN];
69ae int st[MAXN][32 - __builtin_clz(MAXN)];
427e
8041 inline int ext(int x, int y){return x>y?x:y;} // ! max
427e
d34f void init(int n){
ce01     int l = 31 - __builtin_clz(n);
cf75     rep (i, n) st[i][0] = a[i];
b811     rep (j, l)
6937         rep (i, 1+n-(1<<j))
082a             st[i][j+1] = ext(st[i][j], st[i+(1<<j)][j]);
95cf }
427e
c863 int rmq(int l, int r){
92f5     int k = 31 - __builtin_clz(r-l+1);
baa2     return ext(st[l][k], st[r-(1<<k)+1][k]);
95cf }
```

## 7 Geometrics

### 7.1 2D geometric template

```
302f #include <bits/stdc++.h>
421c using namespace std;
427e
4553 typedef int T;
c0ae typedef struct pt {
7a9d     T x, y;
ffaa     T operator , (pt a) { return x*a.x + y*a.y; } // inner product
3ec7     T operator * (pt a) { return x*a.y - y*a.x; } // outer product
221a     pt operator + (pt a) { return {x+a.x, y+a.y}; }
8b34     pt operator - (pt a) { return {x-a.x, y-a.y}; }
427e
368b     pt operator * (T k) { return {x*k, y*k}; }
```

```
pt operator - () { return {-x, -y};}
} vec;
```

```
typedef pair<pt, pt> seg;
```

```
bool ptOnSeg(pt& p, seg& s){
    vec v1 = s.first - p, v2 = s.second - p;
    return (v1, v2) <= 0 && v1 * v2 == 0;
}
```

```
// 0 not on segment
// 1 on segment except vertices
// 2 on vertices
```

```
int ptOnSeg2(pt& p, seg& s){
    vec v1 = s.first - p, v2 = s.second - p;
    T ip = (v1, v2);
    if (v1 * v2 != 0 || ip > 0) return 0;
    return (v1, v2) ? 1 : 2;
}
```

```
// if two orthogonal rectangles do not touch, return true
```

```
inline bool nIntRectRect(seg a, seg b){
    return min(a.first.x, a.second.x) > max(b.first.x, b.second.x) ||
           min(a.first.y, a.second.y) > max(b.first.y, b.second.y) ||
           min(b.first.x, b.second.x) > max(a.first.x, a.second.x) ||
           min(b.first.y, b.second.y) > max(a.first.y, a.second.y);
}
```

```
// >0 in order
```

```
// <0 out of order
```

```
// =0 not standard
```

```
inline double rotOrder(vec a, vec b, vec c){return double(a*b)*(b*c);}
```

```
inline bool intersect(seg a, seg b){
    // ! if (nIntRectRect(a, b)) return false; // if commented, assume that a
    // and b are non-collinear
    return rotOrder(b.first-a.first, a.second-a.first, b.second-a.first) >= 0 &&
           rotOrder(a.first-b.first, b.second-b.first, a.second-b.first) >= 0;
}
```

```
// 0 not intersect
```

```
// 1 standard intersection
```

```
// 2 vertex-Line intersection
```

```
// 3 vertex-vertex intersection
```

90f4  
ba8c  
427e  
0ea6  
427e  
8d6e  
ce77  
de97  
95cf  
427e  
427e  
427e  
427e  
8421  
ce77  
70ca  
8b14  
0847  
95cf  
427e  
427e  
72bb  
f9ac  
f486  
39ce  
80c7  
95cf  
427e  
427e  
427e  
427e  
7538  
427e  
31ed  
427e  
cb52  
059e  
95cf  
427e  
427e  
427e  
427e

```

427e // 4 collinear and have common point(s)
4d19 int intersect2(seg& a, seg& b){
5dc4     if (nIntRectRect(a, b)) return 0;
42c0     vec va = a.second - a.first, vb = b.second - b.first;
2096     double j1 = rotOrder(b.first-a.first, va, b.second-a.first),
72fe         j2 = rotOrder(a.first-b.first, vb, a.second-b.first);
5ac6     if (j1 < 0 || j2 < 0) return 0;
9400     if (j1 != 0 && j2 != 0) return 1;
83db     if (j1 == 0 && j2 == 0){
6b0c         if (va * vb == 0) return 4; else return 3;
fb17     } else return 2;
95cf }
427e
2c68 template <typename Tp = T>
5894 inline pt getIntersection(pt P, vec v, pt Q, vec w){
6850     static_assert(is_same<Tp, double>::value, "must_be_double!");
7c9a     return P + v * (w*(P-Q)/(v*w));
95cf }
427e
427e // -1 outside the polygon
427e // 0 on the border of the polygon
427e // 1 inside the polygon
cbdd int ptOnPoly(pt p, pt* poly, int n){
5fb4     int wn = 0;
1294     for (int i = 0; i < n; i++) {
427e
3cae         T k, d1 = poly[i].y - p.y, d2 = poly[(i+1)%n].y - p.y;
b957         if (k = (poly[(i+1)%n] - poly[i])*(p - poly[i])){
8c40             if (k > 0 && d1 <= 0 && d2 > 0) wn++;
3c4d             if (k < 0 && d2 <= 0 && d1 > 0) wn--;
aad3         } else return 0;
95cf     }
0a5f     return wn ? 1 : -1;
95cf }
427e
d4a3 istream& operator >> (istream& lhs, pt& rhs){
fa86     lhs >> rhs.x >> rhs.y;
331a     return lhs;
95cf }
427e
07ae istream& operator >> (istream& lhs, seg& rhs){
5cab     lhs >> rhs.first >> rhs.second;
331a     return lhs;
95cf }

```

## 8 Appendices

### 8.1 Primes

#### 8.1.1 First primes

$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$
2	1	3	2	5	2	7	3	11	2
13	2	17	3	19	2	23	5	29	2
31	3	37	2	41	6	43	3	47	5
53	2	59	2	61	2	67	2	71	7
73	5	79	3	83	2	89	3	97	5
101	2	103	5	107	2	109	6	113	3
127	3	131	2	137	3	139	2	149	2
151	6	157	5	163	2	167	5	173	2
179	2	181	2	191	19	193	5	197	2
199	3	211	2	223	3	227	2	229	6

#### 8.1.2 Arbitrary length primes

$\lg p$	$p$	$g(p)$	$p$	$g(p)$
3	967	5	1031	14
4	9859	2	10273	10
5	96331	10	102931	3
6	958543	6	1031137	5
7	9594539	2	10169651	2
8	96243449	3	103211039	7
9	980483981	2	1042484357	2
10	9858935453	2	10261276009	7
11	95748666809	3	101759940101	2
12	950781833849	3	1012797784423	5
13	9739822952371	7	10037217092377	7
14	96181051140397	5	104974966380359	11
15	981030138360889	13	1029038416465403	2
16	9655206098080843	3	10116299875820773	2
17	97687777921994419	3	101506415998163437	2

**8.1.3**  $\sim 1 \times 10^9$ 

$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$
954854573	3	967607731	2	973215833	3
975831713	3	978949117	2	980766497	3
983879921	3	985918807	3	986608921	29
991136977	5	991752599	13	997137961	11
1003911991	3	1009775293	2	1012423549	6
1021000537	5	1023976897	7	1024153643	2
1037027287	3	1038812881	11	1044754639	3
1045125617	3	1047411427	3	1047753349	6

**8.1.4**  $\sim 1 \times 10^{18}$ 

$p$	$g(p)$	$p$	$g(p)$
951970612352230049	3	963284339889659609	3
967495386904694119	3	969751761517096213	2
983238274281901499	2	984647442475101409	23
989286107138674069	11	1002507954383424641	3
1006658951440146419	2	1020152326159075903	3
1034876265966119449	7	1042753851435034019	2
1043609016597371563	2	1045571042176595707	2
1048364250160580293	2	1049495624119026949	2

**8.2 Pell's equation**

$x^2 - ny^2 = 1$ , where  $n$  is a positive nonsquare integer.

Let  $(x_0, y_0)$  be the smallest positive solution of the equation, then the  $k$ -th solution is:

$$\begin{pmatrix} x_k \\ y_k \end{pmatrix} = \begin{pmatrix} x_0 & ny_0 \\ y_0 & x_0 \end{pmatrix}^k \begin{pmatrix} x_0 \\ y_0 \end{pmatrix}$$

Some smallest solutions to Pell's equation:

$n$	2	3	5	6	7	8	10	11	12	13	14	15	17	18	19	20
$x$	3	2	9	5	8	3	19	10	7	649	15	4	33	17	170	9
$y$	2	1	4	2	3	1	6	3	2	180	4	1	8	4	39	2

**8.3 Burnside's lemma and Polya's enumeration theorem**

The Burnside's lemma says that

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

where  $G$  is a group acting on  $X$ ,  $X^g$  is the set of elements in  $X$  that are fixed by  $g$ , i.e.  $X^g = \{x \in X : gx = x\}$ .

The unweighted version of Pólya enumeration theorem says that

$$|Y^X/G| = \frac{1}{|G|} \sum_{g \in G} m^{c_g}$$

where  $m = |X|$  is the number of colors,  $c_g$  is the number of the cycles of permutation  $g$ .

**8.4 Lagrange's interpolation**

For sample points  $(x_0, y_0), \dots, (x_k, y_k)$ , define

$$l_j(x) = \prod_{0 \leq m \leq k, m \neq j} \frac{x - x_m}{x_j - x_m}$$

then the Lagrange polynomial is

$$L(x) = \sum_{j=0}^k y_j l_j(x).$$

To use the script below, type two lines

```
x0 x1 x2 ... xn
y0 y1 y2 ... yn
```

the script will print the fractional coefficient of the polynomial in ascending exponent order.

```
#!/usr/bin/python2
from fractions import *

def polyadd(a, b) : return map(lambda x, y : (x or 0) + (y or 0), a, b)

def polymul(a, b) :
    p = [0] * (len(a)+len(b)-1)
    for e1, c1 in enumerate(a) :
```

```
6dc9
4b2b
427e
bbbe
427e
796b
83e4
f697
```



```

156c         for e2, c2 in enumerate(b) :
dfce             p[e1 + e2] += c1 * c2
5849         return p
427e
f06d x, y = [map(Fraction, raw_input().split()) for _ in 0,0]
e80a n = len(x)

```

```

lj = [reduce(polymul, [[-x[m]/(x[j]-x[m]), 1/(x[j]-x[m])]
        for m in range(n) if m != j]) for j in range(n)]
print ' '.join(map(str, reduce(polyadd,
        map(lambda a, b : [x * a for x in b], y, lj))))

```

```

a649
9dfa
46f9
d754

```