

南京大学 ACM-ICPC 集训队代码模版库



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1 General

1.1 Code library checksum

```
ab14 #!/usr/bin/python3
c502 import re, sys, hashlib
427e
f7db for line in sys.stdin.read().strip().split("\n") :
ddf5     print(hashlib.md5(re.sub(r'\s|//[.]*', '', line).encode('utf8')).hexdigest()
        [-4:], line)
```

1.2 Makefile

```
dab2 .PHONY : run
427e
207e $(t) : $(t).cpp
2d16     g++ --std=c++14 -Wall -D__LOCAL_DEBUG__ -fsanitize=undefined -fsanitize=
        address -ggdb -pipe -o $@ $<
427e
5f25 run : $(t)
bf3e     ./$$(t) < $(t).in
```

1.3 .vimrc

```
914c set nocompatible
733d syntax on
6bbc colorscheme slate
7db5 set number
b0e3 set cursorline
061b set shiftwidth=2
8011 set softtabstop=2
a66d set tabstop=2
d23a set expandtab
5245 set magic
740c set smartindent
bee8 set backspace=indent,eol,start
815d set cmdheight=1
0a40 set laststatus=2
e458 set statusline=\ %<%F[%1*%M%*%n%R%H]%=\ %y\ %0{&fileformat}\ %&encoding}\ %c
        :%l/%L%\
```

```
set whichwrap=b,s,<,>[,]
```

1c67

1.4 Stack

```
const int STK_SZ = 2000000;
char STK[STK_SZ * sizeof(void*)];
void *STK_BAK;

#ifdef __i386__
#define SP "%esp"
#elif defined(__x86_64__)
#define SP "%rsp"
#endif

int main() {
    asm volatile("movl SP, %0; movl 1, SP: =g(STK_BAK):g(STK+sizeof(STK)):");
    ;

    // main program

    asm volatile("movl %0, SP: =g(STK_BAK)");
    return 0;
}
```

bebe
effc
4e99
427e
7bc9
0894
ac7a
a9ea
1937
427e
3117
3750
427e
427e
427e
6856
7021
95cf

1.5 Template

```
#include <bits/stdc++.h>
using namespace std;

#ifdef __LOCAL_DEBUG__
# define _debug(fmt, ...) fprintf(stderr, "[%s] " fmt "\n", \
    __func__, ##__VA_ARGS__)
#else
# define _debug(...) ((void) 0)
#endif
#define rep(i, n) for (int i=0; i<(n); i++)
#define Rep(i, n) for (int i=1; i<=(n); i++)
#define range(x) begin(x), end(x)
typedef long long LL;
typedef unsigned long long ULL;
```

302f
421c
427e
426f
3341
611f
a8cb
e6b5
1937
0d6c
cfe3
3505
5cad
b773

2 Miscellaneous Algorithms

2.1 2-SAT

```

0f42 const int MAXN = 100005;
03a9 struct twoSAT{
5c83     int n;
8f72     vector<int> G[MAXN*2];
d060     bool mark[MAXN*2];
b42d     int S[MAXN*2], c;
427e
d34f     void init(int n){
b985         this->n = n;
f9ec         for (int i=0; i<n*2; i++) G[i].clear();
0609         memset(mark, 0, sizeof(mark));
95cf     }
427e
3bd5     bool dfs(int x){
bd70         if (mark[x^1]) return false;
c96a         if (mark[x]) return true;
fd23         mark[x] = true;
4bea         S[c++] = x;
1ce6         for (int i=0; i<G[x].size(); i++)
d942             if (!dfs(G[x][i])) return false;
3361         return true;
95cf     }
427e
5894     void add_clause(int x, bool xval, int y, bool yval){
6afe         x = x * 2 + xval;
e680         y = y * 2 + yval;
81cc         G[x^1].push_back(y);
6835         G[y^1].push_back(x);
95cf     }
427e
d0cb     bool solve() {
7c39         for (int i=0; i<n*2; i+=2){
e63f             if (!mark[i] && !mark[i+1]){
88fb                 c = 0;
f4b9                 if (!dfs(i)){
3f03                     while (c > 0) mark[S[--c]] = false;
86c5                     if (!dfs(i+1)) return false;
95cf                 }
95cf             }

```

```

    }
    return true;
}

inline bool value(unsigned i){return mark[2*i+1];}
};

```

95cf
3361
95cf
427e
5f0a
329b

2.2 Knuth's optimization

```

int n;
int dp[256][256], dc[256][256];

template <typename T>
void compute(T cost) {
    for (int i = 0; i <= n; i++) {
        dp[i][i] = 0;
        dc[i][i] = i;
    }
    rep (i, n) {
        dp[i][i+1] = 0;
        dc[i][i+1] = i;
    }
    for (int len = 2; len <= n; len++) {
        for (int i = 0; i + len <= n; i++) {
            int j = i + len;
            int lbnd = dc[i][j-1], rbnd = dc[i+1][j];
            dp[i][j] = INT_MAX / 2;
            int c = cost(i, j);
            for (int k = lbnd; k <= rbnd; k++) {
                int res = dp[i][k] + dp[k][j] + c;
                if (res < dp[i][j]) {
                    dp[i][j] = res;
                    dc[i][j] = k;
                }
            }
        }
    }
};

```

5c83
d77c
427e
b7ec
0bc7
0423
8f5e
9488
95cf
be8e
95b5
aa0f
95cf
ec08
88b8
d3da
9824
a24a
f933
90d2
9bd0
26b5
e6af
9c88
95cf
95cf
95cf
95cf
329b

2.3 Mo's algorithm

All intervals are closed on both sides. When running functions `enter()` and `leave()`, the global `l` and `r` has not changed yet.

Usage:

```
add_query(id, l, r)    Add id-th query [l, r].
run()                 Run Mo's algorithm.
init()                TODO. Initialize the range [l, r].
yield(id)             TODO. Yield answer for id-th query.
enter(o)              TODO. Add o-th element.
leave(o)              TODO. Remove o-th element.
```

```
5194 constexpr int BLOCK_SZ = 300;
427e
3ec4 struct query { int l, r, id; };
d26a vector<query> queries;
427e
1e30 void add_query(int id, int l, int r) {
54c9     queries.push_back(query{l, r, id});
95cf }
427e
9f6b int l, r;
427e
427e // ----- functions to implement -----
62b4 inline void init();
50e1 inline void yield(int id);
b20d inline void enter(int o);
13af inline void leave(int o);
427e
37f0 void run() {
ab0b     if (queries.empty()) return;
8508     sort(range(queries), [](query lhs, query rhs) {
c7f8         int lb = lhs.l / BLOCK_SZ, rb = rhs.l / BLOCK_SZ;
03e7         if (lb != rb) return lb < rb;
0780         return lhs.r < rhs.r;
b251     });
6196     l = queries[0].l;
9644     r = queries[0].r;
07e2     init();
5bc9     for (query q : queries) {
7bc7         while (l > q.l) enter(l - 1), l--;
d646         while (r < q.r) enter(r + 1), r++;
13f0         while (l < q.l) leave(l), l++;
e1c6         while (r > q.r) leave(r), r--;
```

```
        yield(q.id);
    }
}
```

```
82f5
95cf
95cf
```

3 String

3.1 Knuth-Morris-Pratt algorithm

```
const int SIZE = 10005;

struct kmp_matcher {
    char p[SIZE];
    int fail[SIZE];
    int len;

    void construct(const char* needle) {
        len = strlen(p);
        strcpy(p, needle);
        fail[0] = fail[1] = 0;
        for (int i = 1; i < len; i++) {
            int j = fail[i];
            while (j && p[i] != p[j]) j = fail[j];
            fail[i + 1] = p[i] == p[j] ? j + 1 : 0;
        }
    }

    inline void found(int pos) {
        // ! add codes for having found at pos
    }

    void match(const char* haystack) { // must be called after construct
        const char* t = haystack;
        int n = strlen(t);
        int j = 0;
        rep(i, n) {
            while (j && p[j] != t[i]) j = fail[j];
            if (p[j] == t[i]) j++;
            if (j == len) found(i - len + 1);
        }
    }
};
```

```
2836
427e
d02b
2d81
9847
57b7
427e
60cf
aaa1
3a87
3dd4
d8a8
147f
3c79
4643
95cf
95cf
427e
c464
427e
95cf
427e
2daf
700f
8482
8fd0
be8e
4e19
b5d5
f024
95cf
95cf
329b
```

3.2 Manacher algorithm

```

81d4 struct Manacher {
cd09     int Len;
9255     vector<int> lc;
b301     string s;
427e
ec07     void work() {
c033         lc[1] = 1;
6bef         int k = 1;
427e
491f         for (int i = 2; i <= Len; i++) {
7957             int p = k + lc[k] - 1;
5e04             if (i <= p) {
24a1                 lc[i] = min(lc[2 * k - i], p - i + 1);
8e2e             } else {
e0e5                 lc[i] = 1;
95cf             }
74ff             while (s[i + lc[i]] == s[i - lc[i]]) lc[i]++;
2b9a             if (i + lc[i] > k + lc[k]) k = i;
95cf         }
95cf     }
427e
bfd5     void init(const char *tt) {
aaaf         int len = strlen(tt);
f701         s.resize(len * 2 + 10);
7045         lc.resize(len * 2 + 10);
8e13         s[0] = '*';
ae54         s[1] = '#';
1321         for (int i = 0; i < len; i++) {
e995             s[i * 2 + 2] = tt[i];
69fd             s[i * 2 + 1] = '#';
95cf         }
43fd         s[len * 2 + 1] = '#';
75d1         s[len * 2 + 2] = '\0';
61f7         Len = len * 2 + 2;
3e7a         work();
95cf     }
427e
b194     pair<int, int> maxpal(int l, int r) {
901a         int center = l + r + 1;
ffb2         int rad = lc[center] / 2;
ab54         int rmid = (l + r + 1) / 2;

```

```

    int r1 = rmid - rad, rr = rmid + rad - 1;
    if ((r ^ 1) & 1) {
    } else rr++;
    return {max(l, r1), min(r, rr)};
}
};

```

```

17e4
3908
69f3
69dc
95cf
329b

```

3.3 Aho-corasick automaton

```

struct AC : Trie {
    int fail[MAXN];
    int last[MAXN];

    void construct() {
        queue<int> q;
        fail[0] = 0;
        rep(c, CHARN) {
            if (int u = tr[0][c]) {
                fail[u] = 0;
                q.push(u);
                last[u] = 0;
            }
        }
        while (!q.empty()) {
            int r = q.front();
            q.pop();
            rep(c, CHARN) {
                int u = tr[r][c];
                if (!u) {
                    tr[r][c] = tr[fail[r]][c];
                    continue;
                }
                q.push(u);
                int v = fail[r];
                while (v && !tr[v][c]) v = fail[v];
                fail[u] = tr[v][c];
                last[u] = tag[fail[u]] ? fail[u] : last[fail[u]];
            }
        }
    }

    void found(int pos, int j) {

```

```

a1ad
9143
daca
427e
8690
93d2
a7a6
ce3c
b1c6
a506
3e14
f689
95cf
95cf
cc78
31f0
15dd
ce3c
ab59
0ef5
9d58
b333
95cf
3e14
b3ff
d2ea
c275
654c
95cf
95cf
95cf
427e
7752

```

```

043e     if (j) {
427e         // ! add codes for having found word with tag[j]
4a96         found(pos, last[j]);
95cf     }
95cf }
427e
9785 void find(const char* text) { // must be called after construct()
80a4     int p = 0, c, len = strlen(text);
9c94     rep(i, len) {
b3db         c = id(text[i]);
f119         p = tr[p][c];
f08e         if (tag[p])
389b             found(i, p);
1e67         else if (last[p])
299e             found(i, last[p]);
95cf     }
95cf }
329b };

```

3.4 Suffix array

Usage:

s[]	the source string
sa[i]	the index of starting position of i -th suffix
rk[i]	the number of suffixes less than the suffix starting from i
n	size of source string
m	size of character set

```

b5a7 void suffix_array(const int s[], int sa[], int rk[], int n, int m) {
427e     // size of x, y must be at least n, size of cnt must be at least max(n, m)
0a1b     static int x[1000005], y[1000005], cnt[1000005];
7306     copy(s, s + n, rk);
afbb     iota(y, y + n, 0);
1e09     for (int k = 0; ; k ? (k <= 1) : (k = 1)) {
abc8         if (k) {
2f70             int t = 0;
8b5e             for (int i = n - k; i < n; i++) y[t++] = i;
1c33             rep(i, n) if (sa[i] >= k) y[t++] = sa[i] - k;
95cf         }
6066         fill(cnt, cnt + m, 0);
6c75         rep(i, n) cnt[rk[i]]++;
9154         partial_sum(cnt, cnt + m, cnt);
3c00         rep(i, n) x[i] = rk[y[i]];

```

```

        for (int i = n - 1; i >= 0; i--)
            sa[--cnt[x[i]]] = y[i];
        swap_ranges(rk, rk + n, y);
        rk[sa[0]] = m = 0;
        for (int i = 1; i < n; i++) {
            int t1 = sa[i], t2 = sa[i-1];
            rk[t1] = (y[t1] == y[t2] and y[t1 + k] == y[t2 + k]) ? m : ++m;
        }
        if (++m == n) break;
    }
}

```

```

66a7
28cb
ae41
4f97
324a
ed69
c4f2
95cf
7912
95cf
95cf

```

3.5 Trie

```

const int MAXN = 12000;
const int CHARN = 26;

```

```

inline int id(char c) { return c - 'a'; }

```

```

struct Trie {
    int n;
    int tr[MAXN][CHARN]; // Trie tree, 0 denotes fail
    int tag[MAXN];

```

```

    Trie() {
        memset(tr[0], 0, sizeof(tr[0]));
        tag[0] = 0;
        n = 1;
    }

```

// tag should not be 0

```

void add(const char* s, int t) {
    int p = 0, c, len = strlen(s);
    rep(i, len) {
        c = id(s[i]);
        if (!tr[p][c]) {
            memset(tr[n], 0, sizeof(tr[n]));
            tag[n] = 0;
            tr[p][c] = n++;
        }
        p = tr[p][c];
    }
}

```

```

e6f1
dd87
427e
8ff5
427e
a281
5c83
f4f5
35a5
427e
4fee
3ccc
4d52
46bf
95cf
427e
427e
30b0
d50a
9c94
3140
d6c8
26dd
2e5c
73bb
95cf
f119
95cf

```

```

35ef     tag[p] = t;
95cf }
427e
427e // returns 0 if not found
427e // AC automaton does not need this function
216c int search(const char* s) {
d50a     int p = 0, c, len = strlen(s);
9c94     rep(i, len) {
3140         c = id(s[i]);
f339         if (!tr[p][c]) return 0;
f119         p = tr[p][c];
95cf     }
840e     return tag[p];
95cf }
329b };

```

3.6 Rolling hash

PLEASE call `init_hash()` in `int main()`!

Usage:

`build(str)` Construct the hasher with given string.
`operator()(l, r)` Get hash value of substring $[l, r)$.

```

1e42 const LL mod = 1006658951440146419, g = 967;
9f60 const int MAXN = 200005;
0291 LL pg[MAXN];
427e
6832 inline LL mul(LL x, LL y) {
c919     return __int128_t(x) * y % mod;
95cf }
427e
599a void init_hash() { // must be called in `int main()`
286f     pg[0] = 1;
d00f     for (int i = 1; i < MAXN; i++)
4aa9         pg[i] = pg[i - 1] * g % mod;
95cf }
427e
7e62 struct hasher {
534a     LL val[MAXN];
427e
4554     void build(const char *str) { // assume lower-case letter only
f937         for (int i = 0; str[i]; i++)
9645             val[i+1] = (mul(val[i], g) + str[i]) % mod;

```

```

}

LL operator() (int l, int r) { // [l, r)
    return (val[r] - mul(val[l], pg[r - l]) + mod) % mod;
}
} ha;

```

4 Math

4.1 Matrix powermod

```

const int MAXN = 105;
const LL modular = 1000000007;
int n; // order of matrices

struct matrix{
    LL m[MAXN][MAXN];

    void operator *=(matrix& a){
        static LL t[MAXN][MAXN];
        Rep (i, n){
            Rep (j, n){
                t[i][j] = 0;
                Rep (k, n){
                    t[i][j] += (m[i][k] * a.m[k][j]) % modular;
                    t[i][j] %= modular;
                }
            }
            memcpy(m, t, sizeof(t));
        }
    }
};

matrix r;
void m_powmod(matrix& b, LL e){
    memset(r.m, 0, sizeof(r.m));
    Rep(i, n)
        r.m[i][i] = 1;
    while (e){
        if (e & 1) r *= b;
        b *= b;
    }
}

```



```

16fc         e >>= 1;
95cf     }
95cf }

```

4.2 Linear basis

```

8b44 const int MAXD = 30;
03a6 struct linearbasis {
3558     ULL b[MAXD] = {};
427e
842f     bool insert(1l v) {
9b2b         for (int j = MAXD - 1; j >= 0; j--) {
de36             if (!(v & (1ll << j))) continue;
ee78             if (b[j]) v ^= b[j]
037f             else {
7836                 for (int k = 0; k < j; k++)
f0b4                     if (v & (1ll << k)) v ^= b[k];
b0aa                 for (int k = j + 1; k < MAXD; k++)
46c9                     if (b[k] & (1ll << j)) b[k] ^= v;
8295                 b[j] = v;
3361                 return true;
95cf             }
95cf         }
438e         return false;
95cf     }
329b };

```

4.3 Gauss elimination over finite field

```

b784 const LL p = 1000000007;
427e
2a2c LL powmod(LL b, LL e) {
95a2     LL r = 1;
3e90     while (e) {
1783         if (e & 1) r = r * b % p;
5549         b = b * b % p;
16fc         e >>= 1;
95cf     }
547e     return r;
95cf }
427e

```

```

typedef vector<LL> VLL;
typedef vector<VLL> VVLL;

```

```

LL gauss(VVLL &a, VVLL &b) {
    const int n = a.size(), m = b[0].size();
    vector<int> irow(n), icol(n), ipiv(n);
    LL det = 1;

```

```

    rep (i, n) {
        int pj = -1, pk = -1;
        rep (j, n) if (!ipiv[j])
            rep (k, n) if (!ipiv[k])
                if (pj == -1 || a[j][k] > a[pj][pk]) {
                    pj = j;
                    pk = k;
                }
        if (a[pj][pk] == 0) return 0;
        ipiv[pk]++;
        swap(a[pj], a[pk]);
        swap(b[pj], b[pk]);
        if (pj != pk) det = (p - det) % p;
        irow[i] = pj;
        icol[i] = pk;

```

```

        LL c = powmod(a[pk][pk], p - 2);
        det = det * a[pk][pk] % p;
        a[pk][pk] = 1;
        rep (j, n) a[pk][j] = a[pk][j] * c % p;
        rep (j, m) b[pk][j] = b[pk][j] * c % p;
        rep (j, n) if (j != pk) {
            c = a[j][pk];
            a[j][pk] = 0;
            rep (k, n) a[j][k] = (a[j][k] + p - a[pk][k] * c % p) % p;
            rep (k, m) b[j][k] = (b[j][k] + p - b[pk][k] * c % p) % p;
        }
    }

```

```

    for (int j = n - 1; j >= 0; j--) if (irow[j] != icol[j]) {
        for (int k = 0; k < n; k++) swap(a[k][irow[j]], a[k][icol[j]]);
    }
    return det;
}

```

c130
42ac
427e
2c62
561b
a25e
2976
427e
be8e
d2b5
6b4a
e582
6112
a905
657b
95cf
d480
0305
8dad
aad8
be4d
d080
f156
427e
4ecd
865b
c36a
dd36
1b23
f8f3
e97f
c449
820b
f039
95cf
95cf
427e
37e1
50dc
95cf
f27f
95cf

4.4 Berlekamp-Massey algorithm

```

2b86 const LL MOD = 1000000007;
427e
391d LL inverse(LL b) {
32d3     LL e = MOD - 2, r = 1;
3e90     while (e) {
9a62         if (e & 1) r = r * b % MOD;
29ea         b = b * b % MOD;
16fc         e >>= 1;
95cf     }
547e     return r;
95cf }

32a6 struct Poly {
afe0     vector<int> a;
427e
9794     Poly() { a.clear(); }
427e
de81     Poly(vector<int> &a) : a(a) {}
427e
8087     int length() const { return a.size(); }
427e
16de     Poly move(int d) {
b31d         vector<int> na(d, 0);
f915         na.insert(na.end(), a.begin(), a.end());
cecf         return Poly(na);
95cf     }
427e
fa1a     int calc(vector<int> &d, int pos) {
5b57         int ret = 0;
501c         for (int i = 0; i < (int)a.size(); ++i) {
5de5             if ((ret += (long long)d[pos - i] * a[i] % MOD) >= MOD) {
3041                 ret -= MOD;
95cf             }
95cf         }
ee0f         return ret;
95cf     }
427e
c856     Poly operator - (const Poly &b) {
bd55         vector<int> na(max(this->length(), b.length()));
d1a7         for (int i = 0; i < (int)na.size(); ++i) {
3507             int aa = i < this->length() ? this->a[i] : 0,

```

```

        bb = i < b.length() ? b.a[i] : 0;
        na[i] = (aa + MOD - bb) % MOD;
    }
    return Poly(na);
}

Poly operator * (const int &c, const Poly &p) {
    vector<int> na(p.length());
    for (int i = 0; i < (int)na.size(); ++i) {
        na[i] = (long long)c * p.a[i] % MOD;
    }
    return na;
}

vector<int> solve(vector<int> a) {
    int n = a.size();
    Poly s, b;
    s.a.push_back(1), b.a.push_back(1);
    for (int i = 1, j = 0, ld = a[0]; i < n; ++i) {
        int d = s.calc(a, i);
        if (d) {
            if ((s.length() - 1) * 2 <= i) {
                Poly ob = b;
                b = s;
                s = s - (long long)d * inverse(ld) % MOD * ob.move(i - j);
                j = i;
                ld = d;
            } else {
                s = s - (long long)d * inverse(ld) % MOD * b.move(i - j);
            }
        }
    }
    // Caution: s.a might be shorter than expected
    return s.a;
}

```

2bee
9526
95cf
cecf
95cf
329b
427e
5473
72de
d1a7
bf0c
95cf
aaab
95cf
427e
afff
9f23
58d0
4e8f
c2aa
4158
d503
c29d
db9d
6bce
1d0e
0889
64f1
8e2e
714e
95cf
95cf
95cf
427e
e235
95cf

4.5 Fast Walsh-Hadamard transform

```

void fwt(int* a, int n){
    for (int d = 1; d < n; d <= 1)
        for (int i = 0; i < n; i += d < 1)

```

061e
5595
05f2

```

b833     rep (j, d){
7796         int x = a[i+j], y = a[i+j+d];
427e         // a[i+j] = x+y, a[i+j+d] = x-y;    // xor
427e         // a[i+j] = x+y;                    // and
427e         // a[i+j+d] = x+y;                    // or
95cf     }
95cf }
427e
4db1 void ifwt(int* a, int n){
5595     for (int d = 1; d < n; d <= 1)
05f2         for (int i = 0; i < n; i += d < 1)
b833             rep (j, d){
7796                 int x = a[i+j], y = a[i+j+d];
427e                 // a[i+j] = (x+y)/2, a[i+j+d] = (x-y)/2;    // xor
427e                 // a[i+j] = x-y;                            // and
427e                 // a[i+j+d] = y-x;                            // or
95cf             }
95cf }
427e
2ab6 void conv(int* a, int* b, int n){
950a     fwt(a, n);
e427     fwt(b, n);
8a42     rep(i, n) a[i] *= b[i];
430f     ifwt(a, n);
95cf }

```

4.6 Fast fourier transform

```

4e09 const int NMAX = 1<<20;
427e
3fbf typedef complex<double> cplx;
427e
abd1 const double PI = 2*acos(0.0);
12af struct FFT{
c47c     int rev[NMAX];
27d7     cplx omega[NMAX], oinv[NMAX];
9827     int K, N;
427e
1442     FFT(int k){
e209         K = k; N = 1 << k;
b393         rep (i, N){
7ba3             rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));

```

```

        omega[i] = polar(1.0, 2.0 * PI / N * i);
        oinv[i] = conj(omega[i]);
    }
}

void dft(cplx* a, cplx* w){
    rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
    for (int l = 2; l <= N; l *= 2){
        int m = l/2;
        for (cplx* p = a; p != a + N; p += l)
            rep (k, m){
                cplx t = w[N/l*k] * p[k+m];
                p[k+m] = p[k] - t; p[k] += t;
            }
    }
}

void fft(cplx* a){dft(a, omega);}
void ifft(cplx* a){
    dft(a, oinv);
    rep (i, N) a[i] /= N;
}

void conv(cplx* a, cplx* b){
    fft(a); fft(b);
    rep (i, N) a[i] *= b[i];
    ifft(a);
}
};

```

4.7 Number theoretic transform

```

const int NMAX = 1<<21;

// 998244353 = 7*17*2^23+1, G = 3
const int P = 1004535809, G = 3; // = 479*2^21+1

struct NTT{
    int rev[NMAX];
    LL omega[NMAX], oinv[NMAX];
    int g, g_inv; // g: g_n = G^((P-1)/n)
    int K, N;

```

```

427e  LL powmod(LL b, LL e){
2a2c     LL r = 1;
95a2     while (e){
3e90         if (e&1) r = r * b % P;
6624         b = b * b % P;
489e         e >>= 1;
16fc     }
95cf     return r;
547e }
95cf
427e NTT(int k){
f420     K = k; N = 1 << k;
e209     g = powmod(G, (P-1)/N);
7652     g_inv = powmod(g, N-1);
4b3a     omega[0] = oinv[0] = 1;
e04f     rep (i, N){
b393         rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
7ba3         if (i){
ad4f             omega[i] = omega[i-1] * g % P;
8d8b             oinv[i] = oinv[i-1] * g_inv % P;
9e14         }
95cf     }
95cf }
427e
9668 void _ntt(LL* a, LL* w){
a215     rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
ac6e     for (int l = 2; l <= N; l *= 2){
2969         int m = l/2;
7a1d         for (LL* p = a; p != a + N; p += l)
c24f             rep (k, m){
0ad3                 LL t = w[N/l*k] * p[k+m] % P;
6209                 p[k+m] = (p[k] - t + P) % P;
fa1b                 p[k] = (p[k] + t) % P;
95cf             }
95cf         }
95cf     }
427e
92ea void ntt(LL* a){_ntt(a, omega);}
5daf void intt(LL* a){
1f2a     LL inv = powmod(N, P-2);
9910     _ntt(a, oinv);
a873     rep (i, N) a[i] = a[i] * inv % P;
95cf }

```

```

void conv(LL* a, LL* b){
    ntt(a); ntt(b);
    rep (i, N) a[i] = a[i] * b[i] % P;
    intt(a);
}
};

```

```

427e
3a5b
ad16
e49e
5748
95cf
329b

```

4.8 Sieve of Euler

```

const int MAXX = 1e7+5;
bool p[MAXX];
int prime[MAXX], sz;

void sieve(){
    p[0] = p[1] = 1;
    for (int i = 2; i < MAXX; i++){
        if (!p[i]) prime[sz++] = i;
        for (int j = 0; j < sz && i*prime[j] < MAXX; j++){
            p[i*prime[j]] = 1;
            if (i % prime[j] == 0) break;
        }
    }
}

```

```

cfc3
5861
73ae
427e
9bc6
9628
1ec8
bf28
e82c
b6a9
5f51
95cf
95cf
95cf

```

4.9 Sieve of Euler (General)

```

namespace sieve {
    constexpr int MAXN = 10000007;
    bool p[MAXN]; // true if not prime
    int prime[MAXN], sz;
    int pval[MAXN], pcnt[MAXN];
    int f[MAXN];

    void exec(int N = MAXN) {
        p[0] = p[1] = 1;

        pval[1] = 1;
        pcnt[1] = 0;
        f[1] = 1;
    }
}

```

```

b62e
6589
e982
6ae8
cbf7
6030
427e
76f6
9628
427e
8a8a
bdda
c6b9
427e

```

```

a643   for (int i = 2; i < N; i++) {
01d6       if (!p[i]) {
b2b2           prime[sz++] = i;
37d9           for (LL j = i; j < N; j *= i) {
758c               int b = j / i;
81fd               pval[j] = i * pval[b];
e0f3               pcnt[j] = pcnt[b] + 1;
a96c               f[j] = _____; // f[j] = f(i^pcnt[j])
95cf           }
95cf       }
34c0       for (int j = 0; i * prime[j] < N; j++) {
f87a           int x = i * prime[j]; p[x] = 1;
20cc           if (i % prime[j] == 0) {
9985               pval[x] = pval[i] * prime[j];
3f93               pcnt[x] = pcnt[i] + 1;
8e2e           } else {
cc91               pval[x] = prime[j];
6322               pcnt[x] = 1;
95cf           }
6191           if (x != pval[x]) {
d614               f[x] = f[x / pval[x]] * f[pval[x]]
95cf           }
5f51           if (i % prime[j] == 0) break;
95cf       }
95cf   }
95cf }

```

4.10 Miller-Rabin primality test

The array `a[]` (excluding sentinel, i.e. `LLONG_MAX`) should be

{2}	when $n < 2,047$.
{2, 7, 61}	when $n < 4,759,123,141$ (2^{32}).
{2, 3, 5, 7, 11}	when $n < 2.1 \times 10^{12}$.
{2, 325, 9375, 28178, 450775, 9780504, 1795265022}	when $n < 2^{64}$.

```

f16f   bool test(LL n){
59f2       if (n < 3) return n==2;
427e       // ! The array a[] should be modified if the range of x changes.
3f11       const LL a[] = {2LL, 7LL, 61LL, LLONG_MAX};
c320       LL r = 0, d = n-1, x;
f410       while (~d & 1) d >>= 1, r++;

```

```

for (int i=0; a[i] < n; i++){
    x = powmod(a[i], d, n); // ! powmod must use for 64bit mulmod
    if (x == 1 || x == n-1) goto next;
    rep (i, r) {
        x = mulmod(x, x, n);
        if (x == n-1) goto next;
    }
    return false;
next;;
}
return true;
}

```

4.11 Pollard's rho algorithm

```

ULL gcd(ULL a, ULL b) {return b ? gcd(b, a % b) : a;}

ULL PollardRho(ULL n){
    ULL c, x, y, d = n;
    if (~n&1) return 2;
    while (d == n){
        x = y = 2;
        d = 1;
        c = rand() % (n - 1) + 1;
        while (d == 1){
            x = (mulmod(x, x, n) + c) % n;
            y = (mulmod(y, y, n) + c) % n;
            y = (mulmod(y, y, n) + c) % n;
            d = gcd(x>y ? x-y : y-x, n);
        }
    }
    return d;
}

```

4.12 Qusai-polynomial sum

Must call `init()` before use!

```

namespace polysum {
#define rep(i, a, n) for (int i = a; i < n; i++)
#define per(i, a, n) for (int i = n - 1; i >= a; i--)

```

```

3946 const int D = 2010;
c076 ll a[D], f[D], g[D], p[D], p1[D], p2[D], b[D], h[D][2], C[D];
c4cb ll powmod(ll a, ll b) {
e4b7     ll res = 1;
af5c     a %= mod;
6e39     assert(b >= 0);
b1fa     for (; b >= 1) {
0684         if (b & 1) res = res * a % mod;
05a8         a = a * a % mod;
95cf     }
244d     return res;
95cf }
e88b ll calcn(int d, ll *a, ll n) { // a[0].. a[d] a[n]
b4aa     if (n <= d) return a[n];
d6be     p1[0] = p2[0] = 1;
3245     rep(i, 0, d + 1) {
ffec         ll t = (n - i + mod) % mod;
532d         p1[i + 1] = p1[i] * t % mod;
95cf     }
3245     rep(i, 0, d + 1) {
9800         ll t = (n - d + i + mod) % mod;
9f60         p2[i + 1] = p2[i] * t % mod;
95cf     }
19f3     ll ans = 0;
3245     rep(i, 0, d + 1) {
860e         ll t = g[i] * g[d - i] % mod * p1[i] % mod * p2[d - i] % mod * a[i] % mod;
752a         if ((d - i) & 1)
a69f             ans = (ans - t + mod) % mod;
649a         else
29fe             ans = (ans + t) % mod;
95cf     }
4206     return ans;
95cf }
1901 void init(int M) {
6323     f[0] = f[1] = g[0] = g[1] = 1;
fe69     rep(i, 2, M + 5) f[i] = f[i - 1] * i % mod;
b375     g[M + 4] = powmod(f[M + 4], mod - 2);
7e87     per(i, 1, M + 4) g[i] = g[i + 1] * (i + 1) % mod;
95cf }
5f6d ll polysum(ll m, ll *a, ll n) { // a[0].. a[m] \sum_{i=0}^{n-1} a[i]
2f0c     ll b[D];
a950     for (int i = 0; i <= m; i++) b[i] = a[i];
96b8     b[m + 1] = calcn(m, b, m + 1);
7785     rep(i, 1, m + 2) b[i] = (b[i - 1] + b[i]) % mod;

```

```

return calcn(m + 1, b, n - 1);
}
ll qpolysum(ll R, ll n, ll *a, ll m) { // a[0].. a[m] \sum_{i=0}^{n-1} a[i]*R^i
    if (R == 1) return polysum(n, a, m);
    a[m + 1] = calcn(m, a, m + 1);
    ll r = powmod(R, mod - 2), p3 = 0, p4 = 0, c, ans;
    h[0][0] = 0;
    h[0][1] = 1;
    rep(i, 1, m + 2) {
        h[i][0] = (h[i - 1][0] + a[i - 1]) * r % mod;
        h[i][1] = h[i - 1][1] * r % mod;
    }
    rep(i, 0, m + 2) {
        ll t = g[i] * g[m + 1 - i] % mod;
        if (i & 1)
            p3 = ((p3 - h[i][0] * t) % mod + mod) % mod,
            p4 = ((p4 - h[i][1] * t) % mod + mod) % mod;
        else
            p3 = (p3 + h[i][0] * t) % mod, p4 = (p4 + h[i][1] * t) % mod;
    }
    c = powmod(p4, mod - 2) * (mod - p3) % mod;
    rep(i, 0, m + 2) h[i][0] = (h[i][0] + h[i][1] * c) % mod;
    rep(i, 0, m + 2) C[i] = h[i][0];
    ans = (calcn(m, C, n) * powmod(R, n) - c) % mod;
    if (ans < 0) ans += mod;
    return ans;
}
} // namespace polysum

```

cc07
95cf
c704
356d
ee67
2f7b
c222
c576
4d99
dcbdb
3f1a
95cf
dc94
2d72
59aa
60b1
19f7
649a
b9ee
95cf
6eed
a893
9267
8a10
2dc8
4206
95cf
95cf

5 Graph Theory

5.1 Strongly connected component

```

const int MAXV = 100005;

struct graph{
    vector<int> adj[MAXV];
    stack<int> s;
    int V; // number of vertices
    int pre[MAXV], lnk[MAXV], scc[MAXV];
    int time, sccn;

```

837c
427e
2ea0
88e3
9cad
3d02
8b6c
27ee

```

427e void add_edge(int u, int v){
bfab     adj[u].push_back(v);
c71a }
95cf
427e void dfs(int u){
d714     pre[u] = lnk[u] = ++time;
7e41     s.push(u);
80f6     for (int v : adj[u]){
18f6         if (!pre[v]){
173e             dfs(v);
5f3c             lnk[u] = min(lnk[u], lnk[v]);
002c         } else if (!scc[v]){
6068             lnk[u] = min(lnk[u], pre[v]);
d5df         }
95cf     }
95cf     if (lnk[u] == pre[u]){
8de2         sccn++;
660f         int x;
3c9e         do {
a69f             x = s.top(); s.pop();
3834             scc[x] = sccn;
b0e9         } while (x != u);
6757     }
95cf }
95cf
427e void find_scc(){
4c88     time = sccn = 0;
f4a2     memset(scc, 0, sizeof scc);
8de7     memset(pre, 0, sizeof pre);
8c2f     Rep (i, V){
6901         if (!pre[i]) dfs(i);
56d1     }
95cf }
95cf
427e vector<int> adjc[MAXV];
27ce void contract(){
364d     Rep (i, V)
1a1e         rep (j, adj[i].size()){
21a2             if (scc[i] != scc[adj[i][j]])
b730                 adjc[scc[i]].push_back(scc[adj[i][j]]);
b46e         }
95cf     }
95cf }
329b };

```

5.2 Vertex biconnected component

```

const int MAXN = 100005;
struct graph {
    int pre[MAXN], iscut[MAXN], bccno[MAXN], dfs_clock, bcc_cnt;
    vector<int> adj[MAXN], bcc[MAXN];
    set<pair<int, int>> bcce[MAXN];

    stack<pair<int, int>> s;

    void add_edge(int u, int v) {
        adj[u].push_back(v);
        adj[v].push_back(u);
    }

    int dfs(int u, int fa) {
        int lowu = pre[u] = ++dfs_clock;
        int child = 0;
        for (int v : adj[u]) {
            if (!pre[v]) {
                s.push({u, v});
                child++;
                int lowv = dfs(v, u);
                lowu = min(lowu, lowv);
                if (lowv >= pre[u]) {
                    iscut[u] = 1;
                    bcc[bcc_cnt].clear();
                    bcce[bcc_cnt].clear();
                    while (1) {
                        int xu, xv;
                        tie(xu, xv) = s.top(); s.pop();
                        bcce[bcc_cnt].insert({min(xu, xv), max(xu, xv)});
                        if (bccno[xu] != bcc_cnt) {
                            bcc[bcc_cnt].push_back(xu);
                            bccno[xu] = bcc_cnt;
                        }
                        if (bccno[xv] != bcc_cnt) {
                            bcc[bcc_cnt].push_back(xv);
                            bccno[xv] = bcc_cnt;
                        }
                    }
                }
            }
        }
    }
};

```

```

7096         if (xu == u && xv == v) break;
95cf     }
03f5     bcc_cnt++;
95cf     }
7470     } else if (pre[v] < pre[u] && v != fa) {
e7f8         s.push({u, v});
f115         lowu = min(lowu, pre[v]);
95cf     }
95cf     }
e104     if (fa < 0 && child == 1) iscut[u] = 0;
1160     return lowu;
95cf }
427e
17be void find_bcc(int n) {
8c2f     memset(pre, 0, sizeof pre);
e2d2     memset(iscut, 0, sizeof iscut);
40d3     memset(bccno, -1, sizeof bccno);
fae2     dfs_clock = bcc_cnt = 0;
5c63     rep (i, n) if (!pre[i]) dfs(i, -1);
95cf }
329b };

```

5.3 Minimum spanning arborescence (Chu-Liu)

All vertices are 1-based.

Usage:

getans(n, root, edges) Compute the total size of MSA rooted at root.

Time Complexity: $O(|V||E|)$

```

bcf8 struct edge {
54f1     int u, v;
309c     LL w;
329b };
427e
f5a4 const int MAXN = 10005;
7124 LL in[MAXN];
1c1d int pre[MAXN], vis[MAXN], id[MAXN];
427e
5a43 LL getans(int n, int rt, vector<edge>& edges) {
f7ff     LL ans = 0;
8abb     int cnt = 0;
a147     while (1) {

```

```

Rep (i, n) in[i] = LLONG_MAX, id[i] = vis[i] = 0;
for (auto e : edges) {
    if (e.u != e.v and e.w < in[e.v]) {
        pre[e.v] = e.u;
        in[e.v] = e.w;
    }
}
in[rt] = 0;
Rep (i, n) {
    if (in[i] == LLONG_MAX) return -1;
    ans += in[i];
    int u;
    for (u = i; u != rt && vis[u] != i && !id[u]; u = pre[u])
        vis[u] = i;
    if (u != rt && !id[u]) {
        id[u] = ++cnt;
        for (int v = pre[u]; v != u; v = pre[v])
            id[v] = cnt;
    }
}
if (!cnt) return ans;
Rep (i, n) if (!id[i]) id[i] = ++cnt;
for (auto& e : edges) {
    LL laz = in[e.v];
    e.u = id[e.u];
    e.v = id[e.v];
    if (e.u != e.v) e.w -= laz;
}
n = cnt; rt = id[rt]; cnt = 0;
}
}

```

5.4 Maximum flow (Dinic)

Usage:

add_edge(u, v, c) Add an edge from u to v with capacity c .

max_flow(s, t) Compute maximum flow from s to t .

Time Complexity: For general graph, $O(V^2E)$; for network with unit capacity, $O(\min\{V^{2/3}, \sqrt{E}\}E)$; for bipartite network, $O(\sqrt{VE})$.

```

struct edge{
    int from, to;
    LL cap, flow;

```

641a
0705
073a
c1df
5fbc
95cf
95cf
3fdb
34d7
3c97
cf57
a763
4b0e
88a2
4b22
b66e
0443
5c22
95cf
95cf
91e9
5e22
7400
7750
97ae
fae6
bdd2
95cf
6cc4
95cf
95cf

bcf8
60e2
5e6d


```

329b };
427e
e2cd const int MAXN = 1005;
9062 struct Dinic {
4dbf     int n, m, s, t;
9f0c     vector<edge> edges;
b891     vector<int> G[MAXN];
bbb6     bool vis[MAXN];
b40a     int d[MAXN];
ddec     int cur[MAXN];
427e
5973     void add_edge(int from, int to, LL cap) {
7b55         edges.push_back(edge{from, to, cap, 0});
1db7         edges.push_back(edge{to, from, 0, 0});
fe77         m = edges.size();
dff5         G[from].push_back(m-2);
8f2d         G[to].push_back(m-1);
95cf     }
427e
1836     bool bfs() {
3b73         memset(vis, 0, sizeof(vis));
93d2         queue<int> q;
5d13         q.push(s);
2cd2         vis[s] = 1;
721d         d[s] = 0;
cc78         while (!q.empty()) {
66ba             int x = q.front(); q.pop();
3b61             for (int i = 0; i < G[x].size(); i++) {
b510                 edge& e = edges[G[x][i]];
bba9                 if (!vis[e.to] && e.cap > e.flow) {
cd72                     vis[e.to] = 1;
cf26                     d[e.to] = d[x] + 1;
ca93                     q.push(e.to);
95cf                 }
95cf             }
95cf         }
b23b         return vis[t];
95cf     }
427e
9252     LL dfs(int x, LL a) {
6904         if (x == t || a == 0) return a;
8bf9         LL flow = 0, f;
f515         for (int& i = cur[x]; i < G[x].size(); i++) {
b510             edge& e = edges[G[x][i]];

```

```

         if(d[x] + 1 == d[e.to] && (f = dfs(e.to, min(a, e.cap-e.flow))) > 0)
         {
             e.flow += f;
             edges[G[x][i]^1].flow -= f;
             flow += f;
             a -= f;
             if(a == 0) break;
         }
     }
    return flow;
}

LL max_flow(int s, int t) {
    this->s = s; this->t = t;
    LL flow = 0;
    while (bfs()) {
        memset(cur, 0, sizeof(cur));
        flow += dfs(s, LLONG_MAX);
    }
    return flow;
}

vector<int> min_cut() { // call this after maxflow
    vector<int> ans;
    for (int i = 0; i < edges.size(); i++) {
        edge& e = edges[i];
        if(vis[e.from] && !vis[e.to] && e.cap > 0) ans.push_back(i);
    }
    return ans;
}
};

```

5.5 Maximum cardinality bipartite matching (Hungarian)

```

#include <bits/stdc++.h>
using namespace std;

#define rep(i, n) for (int i = 0; i < (n); i++)
#define Rep(i, n) for (int i = 1; i <= (n); i++)
#define range(x) (x).begin(), (x).end()
typedef long long LL;

```

```

84ee struct Hungarian{
fbf6     int nx, ny;
9ec6     vector<int> mx, my;
9d4c     vector<vector<int>> > e;
edec     vector<bool> mark;
427e
8324     void init(int nx, int ny){
c1d1         this->nx = nx;
f9c1         this->ny = ny;
ac92         mx.resize(nx); my.resize(ny);
3f11         e.clear(); e.resize(nx);
1023         mark.resize(nx);
95cf     }
427e
4589     inline void add(int a, int b){
486c         e[a].push_back(b);
95cf     }
427e
0c2b     bool augment(int i){
207c         if (!mark[i]) {
dae4             mark[i] = true;
6a1e             for (int j : e[i]){
0892                 if (my[j] == -1 || augment(my[j])){
9ca3                     mx[i] = j; my[j] = i;
3361                     return true;
95cf                 }
95cf             }
95cf         }
438e         return false;
95cf     }
427e
3fac     int match(){
5b57         int ret = 0;
b0f1         fill(range(mx), -1);
b957         fill(range(my), -1);
4ed1         rep (i, nx){
13a5             fill(range(mark), false);
cc89             if (augment(i)) ret++;
95cf         }
ee0f         return ret;
95cf     }
329b };

```

5.6 Minimum cost maximum flow

```

struct edge{
    int from, to;
    int cap, flow;
    LL cost;
};

const LL INF = LLONG_MAX / 2;
const int MAXN = 5005;
struct MCMF {
    int s, t, n, m;
    vector<edge> edges;
    vector<int> G[MAXN];
    bool inq[MAXN]; // queue
    LL d[MAXN];     // distance
    int p[MAXN];    // previous
    int a[MAXN];    // improvement

    void add_edge(int from, int to, int cap, LL cost) {
        edges.push_back(edge{from, to, cap, 0, cost});
        edges.push_back(edge{to, from, 0, 0, -cost});
        m = edges.size();
        G[from].push_back(m-2);
        G[to].push_back(m-1);
    }

    bool spfa(){
        queue<int> q;
        fill(d, d + MAXN, INF); d[s] = 0;
        memset(inq, 0, sizeof(inq));
        q.push(s); inq[s] = true;
        p[s] = 0; a[s] = INT_MAX;
        while (!q.empty()){
            int u = q.front(); q.pop(); inq[u] = false;
            for (int i : G[u]) {
                edge& e = edges[i];
                if (e.cap > e.flow && d[e.to] > d[u] + e.cost){
                    d[e.to] = d[u] + e.cost;
                    p[e.to] = G[u][i];
                    a[e.to] = min(a[u], e.cap - e.flow);
                    if (!inq[e.to]) q.push(e.to), inq[e.to] = true;
                }
            }
        }
    }
};

```

```

bcf8
60e2
d698
32cc
329b
427e
cc3e
2aa8
c6cb
9ceb
9f0c
b891
f74f
8f67
9524
b330
427e
f7f2
24f0
95f0
fe77
dff5
8f2d
95cf
427e
3c52
93d2
8494
fd48
5e7c
2dae
cc78
b0aa
3bba
56d8
3601
55bc
0bea
8249
e5d3
95cf

```

```

95cf    }
95cf    }
6d7c    return d[t] != INF;
95cf    }
427e
71a4    void augment(){
06f1        int u = t;
b19d        while (u != s){
db09            edges[p[u]].flow += a[t];
25a9            edges[p[u]^1].flow -= a[t];
e6c9            u = edges[p[u]].from;
95cf        }
95cf    }
427e
6e20    #ifndef GIVEN_FLOW
5972        bool min_cost(int s, int t, int f, LL& cost) {
590d            this->s = s; this->t = t;
21d4            int flow = 0;
23cb            cost = 0;
22dc            while (spfa()) {
bcbd                augment();
a671                if (flow + a[t] >= f){
b14d                    cost += (f - flow) * d[t]; flow = f;
3361                    return true;
8e2e                } else {
2a83                    flow += a[t]; cost += a[t] * d[t];
95cf                }
95cf            }
438e            return false;
95cf        }
a8cb    #else
f9a9        int min_cost(int s, int t, LL& cost) {
590d            this->s = s; this->t = t;
21d4            int flow = 0;
23cb            cost = 0;
22dc            while (spfa()) {
bcbd                augment();
2a83                flow += a[t]; cost += a[t] * d[t];
95cf            }
84fb            return flow;
95cf        }
1937    #endif
329b    };

```

5.7 Global minimum cut (Stoer-Wagner)

```

typedef vector<LL> VI;
typedef vector<VI> VVI;

pair<LL, VI> stoer(VVI &w) {
    int n = w.size();
    VI used(n), c, bestc;
    LL bestw = -1;

    for (int ph = n - 1; ph >= 0; ph--) {
        VI wt = w[0], added = used;
        int prev, last = 0;
        rep (i, ph) {
            prev = last;
            last = -1;
            for (int j = 1; j < n; j++)
                if (!added[j] && (last == -1 || wt[j] > wt[last]))
                    last = j;
            if (i == ph - 1) {
                rep (j, n) w[prev][j] += w[last][j];
                rep (j, n) w[j][prev] = w[prev][j];
                used[last] = true;
                c.push_back(last);
                if (bestw == -1 || wt[last] < bestw) {
                    bestc = c;
                    bestw = wt[last];
                }
            } else {
                rep (j, n) wt[j] += w[last][j];
                added[last] = true;
            }
        }
    }
    return {bestw, bestc};
}

```

5.8 Heavy-light decomposition

Time Complexity: The decomposition itself takes linear time. Each query takes $O(\log n)$ operations.

```
const int MAXN = 100005;
```

```

0b32 vector<int> adj[MAXN];
42f2 int sz[MAXN], top[MAXN], fa[MAXN], son[MAXN], depth[MAXN], id[MAXN];
427e
be5c void dfs1(int x, int dep, int par){
7489     depth[x] = dep;
2ee7     sz[x] = 1;
adb4     fa[x] = par;
b79d     int maxn = 0, s = 0;
c861     for (int c: adj[x]){
fe45         if (c == par) continue;
fd2f         dfs1(c, dep + 1, x);
b790         sz[x] += sz[c];
f0f1         if (sz[c] > maxn){
c749             maxn = sz[c];
fe19             s = c;
95cf         }
95cf     }
0e08     son[x] = s;
95cf }
427e
ba54 int cid = 0;
3644 void dfs2(int x, int t){
8d96     top[x] = t;
d314     id[x] = ++cid;
c4a1     if (son[x]) dfs2(son[x], t);
c861     for (int c: adj[x]){
9881         if (c == fa[x]) continue;
5518         if (c == son[x]) continue;
13f9         else dfs2(c, c);
95cf     }
95cf }
427e
0f04 void decomp(int root){
9fa4     dfs1(root, 1, 0);
1c88     dfs2(root, root);
95cf }
427e
2c98 void query(int u, int v){
03a1     while (top[u] != top[v]){
45ec         if (depth[top[u]] < depth[top[v]]) swap(u, v);
427e         // id[top[u]] to id[u]
005b         u = fa[top[u]];
95cf     }
6083     if (depth[u] > depth[v]) swap(u, v);

```

```

// id[u] to id[v]
}

```

427e
95cf

5.9 Centroid decomposition

Note that the centroid here is not the exact centroid of the graph. It only guarantees that the size of each subtree does not exceed half of that of the original tree. This is enough to guarantee the correct time complexity. All vertices are numbered from 1. Call `decomp(root)` to use.

Usage:

`decomp(u, p)` Decompose the tree rooted at u with parent p .

Time Complexity: The decomposition itself takes $O(n \log n)$ time.

```

vector<int> adj[100005];
int sz[100005], sum;

void getsz(int u, int p) {
    sz[u] = 1; sum++;
    for (int v : adj[u]) {
        if (v == p) continue;
        getsz(v, u);
        sz[u] += sz[v];
    }
}

int getcent(int u, int p) {
    for (int v : adj[u])
        if (v != p and sz[v] > sum / 2)
            return getcent(v, u);
    return u;
}

```

1fb6
88e0
427e
f93d
5b36
18f6
bd87
e3cb
8449
95cf
95cf
427e
67f9
d51f
76e4
18e3
81b0
95cf
427e
4662
618e
303c
427e
18f6
427e
95cf
427e
427e
427e
18f6

```

void decompose(int u) {
    sum = 0; getsz(u, 0);
    u = getcent(u, 0); // update u to the centroid

    for (int v : adj[u]) {
        // get answer for subtree v
    }
    // get answer for the whole tree
    // don't forget to count the centroid itself

    for (int v : adj[u]) { // divide and conquer

```

```

c375     adj[v].erase(find(range(adj[v]), u));
fa6b     decompose(v);
a717     adj[v].push_back(u); // restore deleted edge
95cf     }
95cf     }

```

5.10 DSU on tree

This implementation avoids parallel existence of multiple data structures but requires that the data structure is invertible. To use this template, implement merge, enter, leave as needed; first call decomp(root, 0), then call work(root, 0, false). Labels of vertices start from 1.

Usage:

decomp(u, p) Decompose the tree *u*.
work(u, p, keep) Work for subtree *u*. When keep is set, information is not cleared.

Time Complexity: $O(n \log n)$ times the complexity for merge, enter, leave.

```

1fb6     vector<int> adj[100005];
901d     int sz[100005], son[100005];
427e
5559     void decomp(int u, int p) {
50c0         sz[u] = 1;
18f6         for (int v : adj[u]) {
bd87             if (v == p) continue;
a851             decomp(v, u);
8449             sz[u] += sz[v];
d28c             if (sz[v] > sz[son[u]]) son[u] = v;
95cf         }
95cf     }
427e
b7ec     template <typename T>
62f5     void trav(T fn, int u, int p) {
4412         fn(u);
30b3         for (int v : adj[u]) if (v != p) trav(fn, v, u);
95cf     }
427e
7467     #define for_light(v) for (int v : adj[u]) if (v != p and v != son[u])
33ff     void work(int u, int p, bool keep) {
72a2         for_light(v) work(v, u, 0); // process light children
427e
427e         // process heavy child

```

```

// current data structure contains info of heavy child
if (son[u]) work(son[u], u, 1);

auto merge = [u] (int c) { /* count contribution of c */ };
auto enter = [] (int c) { /* add vertex c */ };
auto leave = [] (int c) { /* remove vertex c */ };

for_light(v) {
    trav(merge, v, u);
    trav(enter, v, u);
}

// count answer for root and add it
// Warning: special check may apply to root!
merge(u);
enter(u);

// Leave current tree
if (!keep) trav(leave, u, p);
}

```

427e
9866
427e
18a9
1ab0
f241
427e
3d3b
74c6
c13d
95cf
427e
427e
427e
c54f
9dec
427e
427e
4e3e
95cf

6 Data Structures

6.1 Fenwick tree (point update range query)

```

struct bit_purq { // point update, range query
    int N;
    vector<LL> tr;

    void init(int n) { // fill the array with 0
        tr.resize(N = n + 5);
    }

    LL sum(int n) {
        LL ans = 0;
        while (n) {
            ans += tr[n];
            n &= n - 1;
        }
        return ans;
    }
}

```

9976
d7af
99ff
427e
d34f
1010
95cf
427e
63d0
f7ff
e290
0715
c0d4
95cf
4206
95cf

```

427e void add(int n, LL x){
f4bd     while (n < N) {
ad20         tr[n] += x;
6c81         n += n & -n;
0af5     }
95cf }
95cf }
329b };

```

6.2 Fenwick tree (range update point query)

```

3d03 struct bit_rupq{ // range update, point query
d7af     int N;
99ff     vector<LL> tr;
427e
d34f void init(int n) { // fill the array with 0
1010     tr.resize(N = n + 5);
95cf }
427e
38d4 LL query(int n) {
f7ff     LL ans = 0;
ad20     while (n < N) {
0715         ans += tr[n];
0af5         n += n & -n;
95cf     }
4206     return ans;
95cf }
427e
f4bd void add(int n, LL x) {
e290     while (n){
6c81         tr[n] += x;
c0d4         n &= n - 1;
95cf     }
95cf }
329b };

```

6.3 Segment tree

```

3942 LL p;
1ebb const int MAXN = 4 * 100006;
451a struct segtree {

```

```

int l[MAXN], m[MAXN], r[MAXN];
LL val[MAXN], tadd[MAXN], tmul[MAXN];

#define lson (o<<1)
#define rson (o<<1|1)

void pull(int o) {
    val[o] = (val[lson] + val[rson]) % p;
}

void push_add(int o, LL x) {
    val[o] = (val[o] + x * (r[o] - l[o])) % p;
    tadd[o] = (tadd[o] + x) % p;
}

void push_mul(int o, LL x) {
    val[o] = val[o] * x % p;
    tadd[o] = tadd[o] * x % p;
    tmul[o] = tmul[o] * x % p;
}

void push(int o) {
    if (l[o] == m[o]) return;
    if (tmul[o] != 1) {
        push_mul(lson, tmul[o]);
        push_mul(rson, tmul[o]);
        tmul[o] = 1;
    }
    if (tadd[o]) {
        push_add(lson, tadd[o]);
        push_add(rson, tadd[o]);
        tadd[o] = 0;
    }
}

void build(int o, int ll, int rr) {
    int mm = (ll + rr) / 2;
    l[o] = ll; r[o] = rr; m[o] = mm;
    tmul[o] = 1;
    if (ll == mm) {
        scanf("%lld", val + o);
        val[o] %= p;
    } else {
        build(lson, ll, mm);

```

```

27be
4510
427e
ac35
1294
427e
1344
bbe9
95cf
427e
e4bc
5dd6
6eff
95cf
427e
d658
b82c
aa86
649f
95cf
427e
b149
3159
0a90
0f4a
045e
ac0a
95cf
1b82
9547
0e73
6234
95cf
95cf
427e
471c
0e87
9d27
ac0a
5c92
001f
e5b6
8e2e
7293

```

```

5e67     build(rson, mm, rr);
ba26     pull(o);
95cf     }
95cf     }
427e
4406 void add(int o, int ll, int rr, LL x) {
3c16     if (ll <= l[o] && r[o] <= rr) {
db32         push_add(o, x);
8e2e     } else {
c4b0         push(o);
4305         if (m[o] > ll) add(lson, ll, rr, x);
d5a6         if (m[o] < rr) add(rson, ll, rr, x);
ba26         pull(o);
95cf     }
95cf     }
427e
48cd void mul(int o, int ll, int rr, LL x) {
3c16     if (ll <= l[o] && r[o] <= rr) {
e7d0         push_mul(o, x);
8e2e     } else {
c4b0         push(o);
d1ba         if (ll < m[o]) mul(lson, ll, rr, x);
67f3         if (m[o] < rr) mul(rson, ll, rr, x);
ba26         pull(o);
95cf     }
95cf     }
427e
0f62 LL query(int o, int ll, int rr) {
3c16     if (ll <= l[o] && r[o] <= rr) {
6dfe         return val[o];
8e2e     } else {
c4b0         push(o);
462a         if (rr <= m[o]) return query(lson, ll, rr);
5cca         if (ll >= m[o]) return query(rson, ll, rr);
bbf9         return query(lson, ll, rr) + query(rson, ll, rr);
95cf     }
95cf     }
4d99 } seg;

```

6.4 Link/cut tree

Usage:

pull(x)	Collect information of subtrees.
Link(u, v)	Link two unconnected trees.
Cut(u, v)	Cut an existent edge.
Query(u, v)	Path aggregation.
Update(u, x)	Single point modification.

// about 0.13s per 100k ops @Luogu.org

```

namespace LCT {
const int MAXN = 300005;
int fa[MAXN], ch[MAXN][2], val[MAXN], sum[MAXN];
bool rev[MAXN];

bool isroot(int x) {
return ch[fa[x]][0] == x || ch[fa[x]][1] == x;
}

void pull(int x) {
sum[x] = val[x] ^ sum[ch[x][0]] ^ sum[ch[x][1]];
}

void reverse(int x) {
swap(ch[x][0], ch[x][1]);
rev[x] ^= 1;
}

void push(int x) {
if (rev[x]) {
if (ch[x][0]) reverse(ch[x][0]);
if (ch[x][1]) reverse(ch[x][1]);
rev[x] = 0;
}
}

void rotate(int x) {
int y = fa[x], z = fa[y], k = ch[y][1] == x, w = ch[x][!k];
if (isroot(y)) ch[z][ch[z][1] == y] = x;
ch[x][!k] = y; ch[y][k] = w;
if (w) fa[w] = y;
fa[y] = x; fa[x] = z;
pull(y);
}

void pushall(int x) {

```

```

a316     if (isroot(x)) pushall(fa[x]);
a97b     push(x);
95cf     }
427e
f69c     void splay(int x) {
d095         int y = x, z = 0;
8ab3         pushall(y);
f244         while (isroot(x)) {
ceef             y = fa[x]; z = fa[y];
4449             if (isroot(y)) rotate((ch[y][0] == x) ^ (ch[z][0] == y) ? x : y);
cf90             rotate(x);
95cf         }
78a0         pull(x);
95cf     }
427e
6229     void access(int x) {
1548         int z = x;
ba78         for (int y = 0; x; x = fa[y = x]) {
8fec             splay(x);
b05d             ch[x][1] = y;
78a0             pull(x);
95cf         }
7afd         splay(z);
95cf     }
427e
502e     void chroot(int x) {
766a         access(x);
cb0d         reverse(x);
95cf     }
427e
471a     void split(int x, int y) {
3015         chroot(x);
29b5         access(y);
95cf     }
427e
d87a     int Root(int x) {
766a         access(x);
874d         while (ch[x][0]) {
a97b             push(x);
b83a             x = ch[x][0];
95cf         }
8fec         splay(x);
d074         return x;
95cf     }

```

```

void Link(int u, int v) { // assume unconnected before
    chroot(u);
    fa[u] = v;
}

void Cut(int u, int v) { // assume connected before
    split(u, v);
    fa[u] = ch[v][0] = 0;
    pull(v);
}

int Query(int u, int v) {
    split(u, v);
    return sum[v];
}

void Update(int u, int x) {
    splay(u);
    val[u] = x;
}
};

```

427e
70d3
b8a5
2448
95cf
427e
c2f4
e8ce
fd95
743b
95cf
427e
6ca2
e8ce
a5ba
95cf
427e
eaba
46ce
1d62
95cf
329b

6.5 Balanced binary search tree from pb_ds

```

#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;

tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update>
rkt;
// null_tree_node_update

// SAMPLE USAGE
rkt.insert(x); // insert element
rkt.erase(x); // erase element
rkt.order_of_key(x); // obtain the number of elements less than x
rkt.find_by_order(i); // iterator to i-th (numbered from 0) smallest element
rkt.lower_bound(x);
rkt.upper_bound(x);
rkt.join(rkt2); // merge tree (only if their ranges do not intersect)
rkt.split(x, rkt2); // split all elements greater than x to rkt2

```

0475
332d
427e
43a7
427e
427e
427e
190e
05d4
add5
b064
c103
4ff4
b19b
cb47

6.6 Persistent segment tree, range k-th query

```

f1a7 struct node {
2ff6     static int n, pos;
427e
7cec     int value;
70e2     node *left, *right;
427e
20b0     void* operator new(size_t size);
427e
3dc0     static node* Build(int l, int r) {
b6c5         node* a = new node;
ce96         if (r > l + 1) {
181e             int mid = (l + r) / 2;
3ba2             a->left = Build(l, mid);
8aaf             a->right = Build(mid, r);
8e2e         } else {
bfc4             a->value = 0;
95cf         }
5ffd         return a;
95cf     }
427e
5a45     static node* init(int size) {
2c46         n = size;
7ee3         pos = 0;
be52         return Build(0, n);
95cf     }
427e
93c0     static int Query(node* lt, node *rt, int l, int r, int k) {
d30c         if (r == l + 1) return l;
181e         int mid = (l + r) / 2;
cb5a         if (rt->left->value - lt->left->value < k) {
8edb             k -= rt->left->value - lt->left->value;
2412             return Query(lt->right, rt->right, mid, r, k);
8e2e         } else {
0119             return Query(lt->left, rt->left, l, mid, k);
95cf         }
95cf     }
427e
c9ad     static int query(node* lt, node *rt, int k) {
9e27         return Query(lt, rt, 0, n, k);
95cf     }
427e

```

```

node *Inc(int l, int r, int pos) const {
    node* a = new node(*this);
    if (r > l + 1) {
        int mid = (l + r) / 2;
        if (pos < mid)
            a->left = left->Inc(l, mid, pos);
        else
            a->right = right->Inc(mid, r, pos);
    }
    a->value++;
    return a;
}

node *inc(int index) {
    return Inc(0, n, index);
}
nodes[8000000];

int node::n, node::pos;
inline void* node::operator new(size_t size) {
    return nodes + (pos++);
}

```

b19c
5794
ce96
181e
203d
f44a
649a
1024
95cf
2b3e
5ffd
95cf
427e
e80f
c246
95cf
865a
427e
99ce
1987
bb3c
95cf

6.7 Sparse table, range extremum query

The array is 0-based and the range is closed.

```

const int MAXN = 100007;
int a[MAXN];
int st[MAXN][32 - __builtin_clz(MAXN)];

inline int ext(int x, int y){return x>y?x:y;} // ! max

void init(int n){
    int l = 31 - __builtin_clz(n);
    rep (i, n) st[i][0] = a[i];
    rep (j, l)
        rep (i, 1+n-(1<<j))
            st[i][j+1] = ext(st[i][j], st[i+(1<<j)][j]);
}

int rmq(int l, int r){
    int k = 31 - __builtin_clz(r-l+1);

```

db63
b330
69ae
427e
8041
427e
d34f
ce01
cf75
b811
6937
082a
95cf
427e
c863
92f5

```

baa2     return ext(st[l][k], st[r-(1<<k)+1][k]);
95cf }

```

7 Geometrics

7.1 2D geometric template

```

302f #include <bits/stdc++.h>
421c using namespace std;
427e
4553 typedef int T;
c0ae typedef struct pt {
7a9d     T x, y;
ffaa     T operator , (pt a) { return x*a.x + y*a.y; } // inner product
3ec7     T operator * (pt a) { return x*a.y - y*a.x; } // outer product
221a     pt operator + (pt a) { return {x+a.x, y+a.y}; }
8b34     pt operator - (pt a) { return {x-a.x, y-a.y}; }
427e
368b     pt operator * (T k) { return {x*k, y*k}; }
90f4     pt operator - () { return {-x, -y}; }
ba8c } vec;
427e
0ea6 typedef pair<pt, pt> seg;
427e
8d6e bool ptOnSeg(pt& p, seg& s){
ce77     vec v1 = s.first - p, v2 = s.second - p;
de97     return (v1, v2) <= 0 && v1 * v2 == 0;
95cf }
427e
427e // 0 not on segment
427e // 1 on segment except vertices
427e // 2 on vertices
8421 int ptOnSeg2(pt& p, seg& s){
ce77     vec v1 = s.first - p, v2 = s.second - p;
70ca     T ip = (v1, v2);
8b14     if (v1 * v2 != 0 || ip > 0) return 0;
0847     return (v1, v2) ? 1 : 2;
95cf }
427e
427e // if two orthogonal rectangles do not touch, return true
72bb inline bool nIntRectRect(seg a, seg b){

```

```

return min(a.first.x, a.second.x) > max(b.first.x, b.second.x) ||
min(a.first.y, a.second.y) > max(b.first.y, b.second.y) ||
min(b.first.x, b.second.x) > max(a.first.x, a.second.x) ||
min(b.first.y, b.second.y) > max(a.first.y, a.second.y);
}

// >0 in order
// <0 out of order
// =0 not standard
inline double rotOrder(vec a, vec b, vec c){return double(a*b)*(b*c);}

inline bool intersect(seg a, seg b){
    // ! if (nIntRectRect(a, b)) return false; // if commented, assume that a
    // and b are non-collinear
    return rotOrder(b.first-a.first, a.second-a.first, b.second-a.first) >= 0 &&
        rotOrder(a.first-b.first, b.second-b.first, a.second-b.first) >= 0;
}

// 0 not intersect
// 1 standard intersection
// 2 vertex-line intersection
// 3 vertex-vertex intersection
// 4 collinear and have common point(s)
int intersect2(seg& a, seg& b){
    if (nIntRectRect(a, b)) return 0;
    vec va = a.second - a.first, vb = b.second - b.first;
    double j1 = rotOrder(b.first-a.first, va, b.second-a.first),
        j2 = rotOrder(a.first-b.first, vb, a.second-b.first);
    if (j1 < 0 || j2 < 0) return 0;
    if (j1 != 0 && j2 != 0) return 1;
    if (j1 == 0 && j2 == 0){
        if (va * vb == 0) return 4; else return 3;
    } else return 2;
}

template <typename Tp = T>
inline pt getIntersection(pt P, vec v, pt Q, vec w){
    static_assert(is_same<Tp, double>::value, "must_be_double!");
    return P + v * (w*(P-Q)/(v*w));
}

// -1 outside the polygon
// 0 on the border of the polygon
// 1 inside the polygon

```

```

f9ac
f486
39ce
80c7
95cf
427e
427e
427e
427e
7538
427e
31ed
427e
cb52
059e
95cf
427e
427e
427e
427e
4d19
5dc4
42c0
2096
72fe
5ac6
9400
83db
6b0c
fb17
95cf
427e
2c68
5894
6850
7c9a
95cf
427e
427e
427e
427e

```

```

cbdd int ptOnPoly(pt p, pt* poly, int n){
5fb4     int wn = 0;
1294     for (int i = 0; i < n; i++) {
427e         T k, d1 = poly[i].y - p.y, d2 = poly[(i+1)%n].y - p.y;
3cae         if (k = (poly[(i+1)%n] - poly[i])*(p - poly[i])){
b957             if (k > 0 && d1 <= 0 && d2 > 0) wn++;
8c40             if (k < 0 && d2 <= 0 && d1 > 0) wn--;
3c4d         } else return 0;
aad3     }
95cf     return wn ? 1 : -1;
0a5f }
95cf }
427e
d4a3 istream& operator >> (istream& lhs, pt& rhs){
fa86     lhs >> rhs.x >> rhs.y;
331a     return lhs;
95cf }
427e
07ae istream& operator >> (istream& lhs, seg& rhs){
5cab     lhs >> rhs.first >> rhs.second;
331a     return lhs;
95cf }

```

8.1.2 Arbitrary length primes

$\lg p$	p	$g(p)$	p	$g(p)$
3	967	5	1031	14
4	9859	2	10273	10
5	96331	10	102931	3
6	958543	6	1031137	5
7	9594539	2	10169651	2
8	96243449	3	103211039	7
9	980483981	2	1042484357	2
10	9858935453	2	10261276009	7
11	95748666809	3	101759940101	2
12	950781833849	3	1012797784423	5
13	9739822952371	7	10037217092377	7
14	96181051140397	5	104974966380359	11
15	981030138360889	13	1029038416465403	2
16	9655206098080843	3	10116299875820773	2
17	97687777921994419	3	101506415998163437	2

8 Appendices

8.1 Primes

8.1.1 First primes

p	$g(p)$	p	$g(p)$	p	$g(p)$	p	$g(p)$	p	$g(p)$
2	1	3	2	5	2	7	3	11	2
13	2	17	3	19	2	23	5	29	2
31	3	37	2	41	6	43	3	47	5
53	2	59	2	61	2	67	2	71	7
73	5	79	3	83	2	89	3	97	5
101	2	103	5	107	2	109	6	113	3
127	3	131	2	137	3	139	2	149	2
151	6	157	5	163	2	167	5	173	2
179	2	181	2	191	19	193	5	197	2
199	3	211	2	223	3	227	2	229	6

8.1.3 $\sim 1 \times 10^9$

p	$g(p)$	p	$g(p)$	p	$g(p)$
954854573	3	967607731	2	973215833	3
975831713	3	978949117	2	980766497	3
983879921	3	985918807	3	986608921	29
991136977	5	991752599	13	997137961	11
1003911991	3	1009775293	2	1012423549	6
1021000537	5	1023976897	7	1024153643	2
1037027287	3	1038812881	11	1044754639	3
1045125617	3	1047411427	3	1047753349	6

8.1.4 $\sim 1 \times 10^{18}$

p	$g(p)$	p	$g(p)$
951970612352230049	3	963284339889659609	3
967495386904694119	3	969751761517096213	2
983238274281901499	2	984647442475101409	23
989286107138674069	11	1002507954383424641	3
1006658951440146419	2	1020152326159075903	3
1034876265966119449	7	1042753851435034019	2
1043609016597371563	2	1045571042176595707	2
1048364250160580293	2	1049495624119026949	2

8.2 Pell's equation

$x^2 - ny^2 = 1$, where n is a positive nonsquare integer.

Let (x_0, y_0) be the smallest positive solution of the equation, then the k -th solution is:

$$\begin{pmatrix} x_k \\ y_k \end{pmatrix} = \begin{pmatrix} x_0 & ny_0 \\ y_0 & x_0 \end{pmatrix}^k \begin{pmatrix} x_0 \\ y_0 \end{pmatrix}$$

Some smallest solutions to Pell's equation:

n	2	3	5	6	7	8	10	11	12	13	14	15	17	18	19	20
x	3	2	9	5	8	3	19	10	7	649	15	4	33	17	170	9
y	2	1	4	2	3	1	6	3	2	180	4	1	8	4	39	2

8.3 Burnside's lemma and Polya's enumeration theorem

The Burnside's lemma says that

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

where G is a group acting on X , X^g is the set of elements in X that are fixed by g , i.e. $X^g = \{x \in X : gx = x\}$.

The unweighted version of Pólya enumeration theorem says that

$$|Y^X/G| = \frac{1}{|G|} \sum_{g \in G} m^{c_g}$$

where $m = |X|$ is the number of colors, c_g is the number of the cycles of permutation g .

8.4 Lagrange's interpolation

For sample points $(x_0, y_0), \dots, (x_k, y_k)$, define

$$l_j(x) = \prod_{0 \leq m \leq k, m \neq j} \frac{x - x_m}{x_j - x_m}$$

then the Lagrange polynomial is

$$L(x) = \sum_{j=0}^k y_j l_j(x).$$