

# 南京大学 ACM-ICPC 集训队代码模版库



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## 1 General

### 1.1 Code library checksum

```
ab14 #!/usr/bin/python3
c502 import re, sys, hashlib
427e
f7db for line in sys.stdin.read().strip().split("\n") :
ddf5     print(hashlib.md5(re.sub(r'\s|//[.]*', '', line).encode('utf8')).hexdigest()
        [-4:], line)
```

### 1.2 Makefile

```
dab2 .PHONY : run
427e
207e $(t) : $(t).cpp
2d16     g++ --std=c++14 -Wall -D__LOCAL_DEBUG__ -fsanitize=undefined -fsanitize=
        address -ggdb -pipe -o $@ $<
427e
5f25 run : $(t)
bf3e     ./$$(t) < $(t).in
```

### 1.3 .vimrc

```
914c set nocompatible
733d syntax on
6bbc colorscheme slate
7db5 set number
b0e3 set cursorline
061b set shiftwidth=2
8011 set softtabstop=2
a66d set tabstop=2
d23a set expandtab
5245 set magic
740c set smartindent
bee8 set backspace=indent,eol,start
815d set cmdheight=1
0a40 set laststatus=2
1c67 set whichwrap=b,s,<,>,[,]
```

### 1.4 Stack

```
const int STK_SZ = 2000000;
char STK[STK_SZ * sizeof(void)];
void *STK_BAK;

#if defined(__i386__)
#define SP "%esp"
#elif defined(__x86_64__)
#define SP "%rsp"
#endif

int main() {
    asm volatile("movl SP, %0; movl %1, SP: =g(STK_BAK):g(STK+sizeof(STK));");
    ;

    // main program

    asm volatile("movl %0, SP::g(STK_BAK);");
    return 0;
}
```

```
bebe
effc
4e99
427e
7bc9
0894
ac7a
a9ea
1937
427e
3117
3750
427e
427e
427e
6856
7021
95cf
```

### 1.5 Template

```
#include <bits/stdc++.h>
using namespace std;

#ifdef __LOCAL_DEBUG__
# define _debug(fmt, ...) fprintf(stderr, "[%s] " fmt "\n", \
    __func__, __VA_ARGS__)
#else
# define _debug(...) ((void) 0)
#endif

#define rep(i, n) for (int i=0; i<(n); i++)
#define Rep(i, n) for (int i=1; i<=(n); i++)
#define range(x) begin(x), end(x)
typedef long long LL;
typedef unsigned long long ULL;
```

```
302f
421c
427e
426f
3341
611f
a8cb
e6b5
1937
0d6c
cfe3
3505
5cad
b773
```

## 2 Miscellaneous Algorithms

### 2.1 2-SAT

```

0f42 const int MAXN = 100005;
03a9 struct twoSAT{
5c83     int n;
8f72     vector<int> G[MAXN*2];
d060     bool mark[MAXN*2];
b42d     int S[MAXN*2], c;
427e
d34f     void init(int n){
b985         this->n = n;
f9ec         for (int i=0; i<n*2; i++) G[i].clear();
0609         memset(mark, 0, sizeof(mark));
95cf     }
427e
3bd5     bool dfs(int x){
bd70         if (mark[x^1]) return false;
c96a         if (mark[x]) return true;
fd23         mark[x] = true;
4bea         S[c++] = x;
1ce6         for (int i=0; i<G[x].size(); i++)
d942             if (!dfs(G[x][i])) return false;
3361         return true;
95cf     }
427e
5894     void add_clause(int x, bool xval, int y, bool yval){
6afe         x = x * 2 + xval;
e680         y = y * 2 + yval;
81cc         G[x^1].push_back(y);
6835         G[y^1].push_back(x);
95cf     }
427e
d0cb     bool solve() {
7c39         for (int i=0; i<n*2; i+=2){
e63f             if (!mark[i] && !mark[i+1]){
88fb                 c = 0;
f4b9                 if (!dfs(i)){
3f03                     while (c > 0) mark[S[--c]] = false;
86c5                     if (!dfs(i+1)) return false;
95cf                 }
95cf             }

```

```

    }
    return true;
}

inline bool value(unsigned i){return mark[2*i+1];}
};

```

95cf  
3361  
95cf  
427e  
5f0a  
329b

### 2.2 Knuth's optimization

```

int n;
int dp[256][256], dc[256][256];

template <typename T>
void compute(T cost) {
    for (int i = 0; i <= n; i++) {
        dp[i][i] = 0;
        dc[i][i] = i;
    }
    rep (i, n) {
        dp[i][i+1] = 0;
        dc[i][i+1] = i;
    }
    for (int len = 2; len <= n; len++) {
        for (int i = 0; i + len <= n; i++) {
            int j = i + len;
            int lbnd = dc[i][j-1], rbnd = dc[i+1][j];
            dp[i][j] = INT_MAX / 2;
            int c = cost(i, j);
            for (int k = lbnd; k <= rbnd; k++) {
                int res = dp[i][k] + dp[k][j] + c;
                if (res < dp[i][j]) {
                    dp[i][j] = res;
                    dc[i][j] = k;
                }
            }
        }
    }
}
};

```

5c83  
d77c  
427e  
b7ec  
0bc7  
0423  
8f5e  
9488  
95cf  
be8e  
95b5  
aa0f  
95cf  
ec08  
88b8  
d3da  
9824  
a24a  
f933  
90d2  
9bd0  
26b5  
e6af  
9c88  
95cf  
95cf  
95cf  
95cf  
329b

## 2.3 Mo's algorithm

All intervals are closed on both sides. When running functions `enter()` and `leave()`, the global `l` and `r` has not changed yet.

### Usage:

```
add_query(id, l, r)    Add id-th query [l, r].
run()                 Run Mo's algorithm.
init()                TODO. Initialize the range [l, r].
yield(id)              TODO. Yield answer for id-th query.
enter(o)               TODO. Add o-th element.
leave(o)               TODO. Remove o-th element.
```

```
5194 constexpr int BLOCK_SZ = 300;
427e
3ec4 struct query { int l, r, id; };
d26a vector<query> queries;
427e
1e30 void add_query(int id, int l, int r) {
54c9     queries.push_back(query{l, r, id});
95cf }
427e
9f6b int l, r;
427e
427e // ----- functions to implement -----
62b4 inline void init();
50e1 inline void yield(int id);
b20d inline void enter(int o);
13af inline void leave(int o);
427e
37f0 void run() {
ab0b     if (queries.empty()) return;
8508     sort(range(queries), [](query lhs, query rhs) {
c7f8         int lb = lhs.l / BLOCK_SZ, rb = rhs.l / BLOCK_SZ;
03e7         if (lb != rb) return lb < rb;
0780         return lhs.r < rhs.r;
b251     });
6196     l = queries[0].l;
9644     r = queries[0].r;
07e2     init();
5bc9     for (query q : queries) {
7bc7         while (l > q.l) enter(l - 1), l--;
d646         while (r < q.r) enter(r + 1), r++;
13f0         while (l < q.l) leave(l), l++;
e1c6         while (r > q.r) leave(r), r--;
```

```
        yield(q.id);
    }
}
```

```
82f5
95cf
95cf
```

## 3 String

### 3.1 Knuth-Morris-Pratt algorithm

```
const int SIZE = 10005;

struct kmp_matcher {
    char p[SIZE];
    int fail[SIZE];
    int len;

    void construct(const char* needle) {
        len = strlen(p);
        strcpy(p, needle);
        fail[0] = fail[1] = 0;
        for (int i = 1; i < len; i++) {
            int j = fail[i];
            while (j && p[i] != p[j]) j = fail[j];
            fail[i + 1] = p[i] == p[j] ? j + 1 : 0;
        }
    }

    inline void found(int pos) {
        // ! add codes for having found at pos
    }

    void match(const char* haystack) { // must be called after construct
        const char* t = haystack;
        int n = strlen(t);
        int j = 0;
        rep(i, n) {
            while (j && p[j] != t[i]) j = fail[j];
            if (p[j] == t[i]) j++;
            if (j == len) found(i - len + 1);
        }
    }
};
```

```
2836
427e
d02b
2d81
9847
57b7
427e
60cf
aaa1
3a87
3dd4
d8a8
147f
3c79
4643
95cf
95cf
427e
c464
427e
95cf
427e
2daf
700f
8482
8fd0
be8e
4e19
b5d5
f024
95cf
95cf
329b
```

### 3.2 Manacher algorithm

```

81d4 struct Manacher {
cd09     int Len;
9255     vector<int> lc;
b301     string s;
427e
ec07     void work() {
c033         lc[1] = 1;
6bef         int k = 1;
427e
491f         for (int i = 2; i <= Len; i++) {
7957             int p = k + lc[k] - 1;
5e04             if (i <= p) {
24a1                 lc[i] = min(lc[2 * k - i], p - i + 1);
8e2e             } else {
e0e5                 lc[i] = 1;
95cf             }
74ff             while (s[i + lc[i]] == s[i - lc[i]]) lc[i]++;
2b9a             if (i + lc[i] > k + lc[k]) k = i;
95cf         }
95cf     }
427e
bfd5     void init(const char *tt) {
aaaf         int len = strlen(tt);
f701         s.resize(len * 2 + 10);
7045         lc.resize(len * 2 + 10);
8e13         s[0] = '*';
ae54         s[1] = '#';
1321         for (int i = 0; i < len; i++) {
e995             s[i * 2 + 2] = tt[i];
69fd             s[i * 2 + 1] = '#';
95cf         }
43fd         s[len * 2 + 1] = '#';
75d1         s[len * 2 + 2] = '\0';
61f7         Len = len * 2 + 2;
3e7a         work();
95cf     }
427e
b194     pair<int, int> maxpal(int l, int r) {
901a         int center = l + r + 1;
ffb2         int rad = lc[center] / 2;
ab54         int rmid = (l + r + 1) / 2;

```

```

    int r1 = rmid - rad, rr = rmid + rad - 1;
    if ((r ^ 1) & 1) {
    } else rr++;
    return {max(l, r1), min(r, rr)};
}
};

```

```

17e4
3908
69f3
69dc
95cf
329b

```

### 3.3 Aho-corasick automaton

```

struct AC : Trie {
    int fail[MAXN];
    int last[MAXN];

    void construct() {
        queue<int> q;
        fail[0] = 0;
        rep(c, CHARN) {
            if (int u = tr[0][c]) {
                fail[u] = 0;
                q.push(u);
                last[u] = 0;
            }
        }
        while (!q.empty()) {
            int r = q.front();
            q.pop();
            rep(c, CHARN) {
                int u = tr[r][c];
                if (!u) {
                    tr[r][c] = tr[fail[r]][c];
                    continue;
                }
                q.push(u);
                int v = fail[r];
                while (v && !tr[v][c]) v = fail[v];
                fail[u] = tr[v][c];
                last[u] = tag[fail[u]] ? fail[u] : last[fail[u]];
            }
        }
    }

    void found(int pos, int j) {

```

```

a1ad
9143
daca
427e
8690
93d2
a7a6
ce3c
b1c6
a506
3e14
f689
95cf
95cf
cc78
31f0
15dd
ce3c
ab59
0ef5
9d58
b333
95cf
3e14
b3ff
d2ea
c275
654c
95cf
95cf
95cf
427e
7752

```

```

043e     if (j) {
427e         // ! add codes for having found word with tag[j]
4a96         found(pos, last[j]);
95cf     }
95cf }
427e
9785 void find(const char* text) { // must be called after construct()
80a4     int p = 0, c, len = strlen(text);
9c94     rep(i, len) {
b3db         c = id(text[i]);
f119         p = tr[p][c];
f08e         if (tag[p])
389b             found(i, p);
1e67         else if (last[p])
299e             found(i, last[p]);
95cf     }
95cf }
329b };

```

### 3.4 Trie

```

e6f1 const int MAXN = 12000;
dd87 const int CHARN = 26;
427e
8ff5 inline int id(char c) { return c - 'a'; }
427e
a281 struct Trie {
5c83     int n;
f4f5     int tr[MAXN][CHARN]; // Trie tree, 0 denotes fail
35a5     int tag[MAXN];
427e
4fee     Trie() {
3ccc         memset(tr[0], 0, sizeof(tr[0]));
4d52         tag[0] = 0;
46bf         n = 1;
95cf     }
427e
427e     // tag should not be 0
30b0 void add(const char* s, int t) {
d50a     int p = 0, c, len = strlen(s);
9c94     rep(i, len) {
3140         c = id(s[i]);

```

```

if (!tr[p][c]) {
    memset(tr[n], 0, sizeof(tr[n]));
    tag[n] = 0;
    tr[p][c] = n++;
}
p = tr[p][c];
}
tag[p] = t;
}

// returns 0 if not found
// AC automaton does not need this function
int search(const char* s) {
    int p = 0, c, len = strlen(s);
    rep(i, len) {
        c = id(s[i]);
        if (!tr[p][c]) return 0;
        p = tr[p][c];
    }
    return tag[p];
}
};

```

```

d6c8
26dd
2e5c
73bb
95cf
f119
95cf
35ef
95cf
427e
427e
427e
216c
d50a
9c94
3140
f339
f119
95cf
840e
95cf
329b

```

### 3.5 Suffix array

The character immediately after the end of the string **MUST** be set to the **UNIQUE SMALLEST** element.

#### Usage:

<code>s[]</code>	the source string
<code>sa[i]</code>	the index of starting position of $i$ -th suffix
<code>rk[i]</code>	the number of suffixes less than the suffix starting from $i$
<code>h[i]</code>	the longest common prefix between the $i$ -th and $(i-1)$ -th lexicographically smallest suffixes
<code>n</code>	size of source string
<code>m</code>	size of character set

```

void radix_sort(int x[], int y[], int sa[], int n, int m) {
    static int cnt[1000005]; // size > max(n, m)
    fill(cnt, cnt + m, 0);
    rep(i, n) cnt[x[y[i]]]++;
    partial_sum(cnt, cnt + m, cnt);
    for (int i = n - 1; i >= 0; i--) sa[--cnt[x[y[i]]]] = y[i];
}

```

```

de09
ec00
6066
93b7
9154
acac
95cf

```

```

427e void suffix_array(int s[], int sa[], int rk[], int n, int m) {
c939     static int y[1000005]; // size > n
a69a     copy(s, s + n, rk);
7306     iota(y, y + n, 0);
afb6     radix_sort(rk, y, sa, n, m);
7b42     for (int j = 1, p = 0; j <= n; j <= 1, m = p, p = 0) {
c8c2         for (int i = n - j; i < n; i++) y[p++] = i;
8c3a         rep (i, n) if (sa[i] >= j) y[p++] = sa[i] - j;
9323         radix_sort(rk, y, sa, n, m + 1);
9e9d         swap_ranges(rk, rk + n, y);
ae41         rk[sa[0]] = p = 1;
ffd2         for (int i = 1; i < n; i++)
445e             rk[sa[i]] = ((y[sa[i]] == y[sa[i-1]] and y[sa[i]+j] == y[sa[i-1]+j])
f8dc                 ? p : ++p);
02f0         if (p == n) break;
95cf     }
97d9     rep (i, n) rk[sa[i]] = i;
95cf }
427e
1715 void calc_height(int s[], int sa[], int rk[], int h[], int n) {
c41f     int k = 0;
f313     h[0] = 0;
be8e     rep (i, n) {
0883         k = max(k - 1, 0);
527d         if (rk[i]) while (s[i+k] == s[sa[rk[i]-1]+k]) ++k;
56b7         h[rk[i]] = k;
95cf     }
95cf }

```

### 3.6 Rolling hash

**PLEASE** call `init_hash()` in `int main()`!

**Usage:**

`build(str)` Construct the hasher with given string.  
`operator()(l, r)` Get hash value of substring  $[l, r)$ .

```

1e42 const LL mod = 1006658951440146419, g = 967;
9f60 const int MAXN = 200005;
0291 LL pg[MAXN];
427e
dfe7 inline LL mul(LL x, LL y) { return __int128_t(x) * y % mod; }
427e

```

```

void init_hash() { // must be called in `int main()`
    pg[0] = 1;
    for (int i = 1; i < MAXN; i++) pg[i] = mul(pg[i-1], g);
}

struct hasher {
    LL val[MAXN];

    void build(const char *str) { // assume lower-case letter only
        for (int i = 0; str[i]; i++)
            val[i+1] = (mul(val[i], g) + str[i]) % mod;
    }

    LL operator() (int l, int r) { // [l, r)
        return (val[r] - mul(val[l], pg[r-l]) + mod) % mod;
    }
};

```

## 4 Math

### 4.1 Extended Euclidean algorithm and Chinese remainder theorem

```

void exgcd(LL a, LL b, LL &g, LL &x, LL &y) {
    if (!b) g = a, x = 1, y = 0;
    else {
        exgcd(b, a % b, g, y, x);
        y -= x * (a / b);
    }
}

LL crt(LL r[], LL p[], int n) {
    LL q = 1, ret = 0;
    rep (i, n) q *= p[i];
    rep (i, n) {
        LL m = q / p[i];
        LL d, x, y;
        exgcd(p[i], m, d, x, y);
        ret = (ret + y * m * r[i]) % q;
    }
    return (q + ret) % q;
}

```



## 4.2 Matrix powermod

```

44b4 const int MAXN = 105;
92df const LL modular = 1000000007;
5c83 int n; // order of matrices
427e
8864 struct matrix{
3180     LL m[MAXN][MAXN];
427e
43c5     void operator *=(matrix& a){
e735         static LL t[MAXN][MAXN];
34d7         Rep (i, n){
4c11             Rep (j, n){
ee1e                 t[i][j] = 0;
c4a7                 Rep (k, n){
fcaf                     t[i][j] += (m[i][k] * a.m[k][j]) % modular;
199e                     t[i][j] %= modular;
95cf                 }
95cf             }
95cf         }
dad4         memcpy(m, t, sizeof(t));
95cf     }
329b };
427e
63d8 matrix r;
3ec2 void m_powmod(matrix& b, LL e){
83f0     memset(r.m, 0, sizeof(r.m));
a7c3     Rep(i, n)
de64         r.m[i][i] = 1;
3e90     while (e){
5a0e         if (e & 1) r *= b;
35c5         b *= b;
16fc         e >>= 1;
95cf     }
95cf }

```

## 4.3 Linear basis

```

8b44 const int MAXD = 30;
03a6 struct linearbasis {
3558     ULL b[MAXD] = {};
427e

```

```

bool insert(LL v) {
    for (int j = MAXD - 1; j >= 0; j--) {
        if (!(v & (1ll << j))) continue;
        if (b[j]) v ^= b[j]
        else {
            for (int k = 0; k < j; k++)
                if (v & (1ll << k)) v ^= b[k];
            for (int k = j + 1; k < MAXD; k++)
                if (b[k] & (1ll << j)) b[k] ^= v;
            b[j] = v;
            return true;
        }
    }
    return false;
}

```

## 4.4 Gauss elimination over finite field

```

const LL p = 1000000007;

LL powmod(LL b, LL e) {
    LL r = 1;
    while (e) {
        if (e & 1) r = r * b % p;
        b = b * b % p;
        e >>= 1;
    }
    return r;
}

typedef vector<LL> VLL;
typedef vector<VLL> WLL;

```

```

LL gauss(WLL &a, WLL &b) {
    const int n = a.size(), m = b[0].size();
    vector<int> irow(n), icol(n), ipiv(n);
    LL det = 1;

    rep (i, n) {
        int pj = -1, pk = -1;
        rep (j, n) if (!ipiv[j])

```

1566  
9b2b  
de36  
ee78  
037f  
7836  
f0b4  
b0aa  
46c9  
8295  
3361  
95cf  
95cf  
438e  
95cf  
329b

b784  
427e  
2a2c  
95a2  
3e90  
1783  
5549  
16fc  
95cf  
547e  
95cf  
427e  
c130  
42ac  
427e  
2c62  
561b  
a25e  
2976  
427e  
be8e  
d2b5  
6b4a

```

e582     rep (k, n) if (!ipiv[k])
6112         if (pj == -1 || a[j][k] > a[pj][pk]) {
a905             pj = j;
657b             pk = k;
95cf         }
d480     if (a[pj][pk] == 0) return 0;
0305     ipiv[pk]++;
8dad     swap(a[pj], a[pk]);
aad8     swap(b[pj], b[pk]);
be4d     if (pj != pk) det = (p - det) % p;
d080     irow[i] = pj;
f156     icol[i] = pk;
427e
4ecd     LL c = powmod(a[pk][pk], p - 2);
865b     det = det * a[pk][pk] % p;
c36a     a[pk][pk] = 1;
dd36     rep (j, n) a[pk][j] = a[pk][j] * c % p;
1b23     rep (j, m) b[pk][j] = b[pk][j] * c % p;
f8f3     rep (j, n) if (j != pk) {
e97f         c = a[j][pk];
c449         a[j][pk] = 0;
820b         rep (k, n) a[j][k] = (a[j][k] + p - a[pk][k] * c % p) % p;
f039         rep (k, m) b[j][k] = (b[j][k] + p - b[pk][k] * c % p) % p;
95cf     }
95cf }
427e
37e1     for (int j = n - 1; j >= 0; j--) if (irow[j] != icol[j]) {
50dc         for (int k = 0; k < n; k++) swap(a[k][irow[j]], a[k][icol[j]]);
95cf     }
f27f     return det;
95cf }

```

## 4.5 Berlekamp-Massey algorithm

Call `berlekamp()` with input sequence  $(x_0, x_1, \dots, x_{n-1})$ . Return a vector of coefficients  $(c_0 = 1, c_1, \dots, c_{m-1})$  with minimum  $m$ , such that  $\sum_{i=0}^m c_i x_{j-i} = 0$  for all possible  $j$ .

```

6e50     LL mod = 1000000007;
97db     vector<LL> berlekamp(const vector<LL>& a) {
8904         vector<LL> p = {1}, r = {1};
075b         LL dif = 1;
8bc9         rep (i, a.size()) {
1b35             LL u = 0;

```

```

rep (j, p.size()) u = (u + p[j] * a[i-j]) % mod;
if (u == 0) {
    r.insert(r.begin(), 0);
} else {
    auto op = p;
    p.resize(max(p.size(), r.size() + 1));
    LL idif = powmod(dif, mod - 2);
    rep (j, r.size())
        p[j+1] = (p[j+1] - r[j] * idif % mod * u % mod + mod) % mod;
    dif = u; r = op;
}
}
return p;
}

```

```

bd0b
eae9
b14c
8e2e
0c78
02f6
0a2e
9b57
dacc
bcd1
95cf
95cf
e149
95cf

```

## 4.6 Fast Walsh-Hadamard transform

```

void fwt(int* a, int n){
    for (int d = 1; d < n; d <= 1)
        for (int i = 0; i < n; i += d < 1)
            rep (j, d){
                int x = a[i+j], y = a[i+j+d];
                // a[i+j] = x+y, a[i+j+d] = x-y;    // xor
                // a[i+j] = x+y;                    // and
                // a[i+j+d] = x+y;                    // or
            }
}

```

```

061e
5595
05f2
b833
7796
427e
427e
427e
95cf
95cf
427e
4db1
5595
05f2
b833
7796
427e
427e
427e
95cf
95cf
427e

```

```

void ifwt(int* a, int n){
    for (int d = 1; d < n; d <= 1)
        for (int i = 0; i < n; i += d < 1)
            rep (j, d){
                int x = a[i+j], y = a[i+j+d];
                // a[i+j] = (x+y)/2, a[i+j+d] = (x-y)/2;    // xor
                // a[i+j] = x-y;                                // and
                // a[i+j+d] = y-x;                                // or
            }
}

```

```

427e
4db1
5595
05f2
b833
7796
427e
427e
427e
95cf
95cf
427e
2ab6
950a
e427

```

```

void conv(int* a, int* b, int n){
    fwt(a, n);
    fwt(b, n);

```

```

8a42     rep(i, n) a[i] *= b[i];
430f     ifwt(a, n);
95cf }

```

## 4.7 Fast fourier transform

```

4e09 const int NMAX = 1<<20;
427e
3fbf typedef complex<double> cplx;
427e
abd1 const double PI = 2*acos(0.0);
12af struct FFT{
c47c     int rev[NMAX];
27d7     cplx omega[NMAX], oinv[NMAX];
9827     int K, N;
427e
1442     FFT(int k){
e209         K = k; N = 1 << k;
b393         rep (i, N){
7ba3             rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
1908             omega[i] = polar(1.0, 2.0 * PI / N * i);
a166             oinv[i] = conj(omega[i]);
95cf         }
95cf     }
427e
b941     void dft(cplx* a, cplx* w){
a215         rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
ac6e         for (int l = 2; l <= N; l *= 2){
2969             int m = l/2;
b3cf             for (cplx* p = a; p != a + N; p += l)
c24f                 rep (k, m){
fe06                     cplx t = w[N/l*k] * p[k+m];
ecbf                     p[k+m] = p[k] - t; p[k] += t;
95cf                 }
95cf             }
95cf         }
427e
617b     void fft(cplx* a){dft(a, omega);}
a123     void ifft(cplx* a){
3b2f         dft(a, oinv);
57fc         rep (i, N) a[i] /= N;
95cf     }

```

```

void conv(cplx* a, cplx* b){
    fft(a); fft(b);
    rep (i, N) a[i] *= b[i];
    ifft(a);
}
};

```

```

427e
bdc0
6497
12a5
f84e
95cf
329b

```

## 4.8 Number theoretic transform

```

const int NMAX = 1<<21;
4ab9
427e
// 998244353 = 7*17*2^23+1, G = 3
427e
const int P = 1004535809, G = 3; // = 479*2^21+1
fb9a
427e
struct NTT{
87ab     int rev[NMAX];
c47c     LL omega[NMAX], oinv[NMAX];
0eda     int g, g_inv; // g:  $g_n = G^{(P-1)/n}$ 
81af     int K, N;
9827
427e     LL powmod(LL b, LL e){
2a2c         LL r = 1;
95a2         while (e){
3e90             if (e&1) r = r * b % P;
6624             b = b * b % P;
489e             e >>= 1;
16fc         }
95cf         return r;
547e     }
95cf
427e     NTT(int k){
f420         K = k; N = 1 << k;
e209         g = powmod(G, (P-1)/N);
7652         g_inv = powmod(g, N-1);
4b3a         omega[0] = oinv[0] = 1;
e04f         rep (i, N){
b393             rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
7ba3             if (i){
ad4f                 omega[i] = omega[i-1] * g % P;
8d8b                 oinv[i] = oinv[i-1] * g_inv % P;
9e14             }
95cf         }

```

```

95cf    }
95cf    }
427e
9668    void _ntt(LL* a, LL* w){
a215        rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
ac6e        for (int l = 2; l <= N; l *= 2){
2969            int m = l/2;
7a1d            for (LL* p = a; p != a + N; p += l)
c24f                rep (k, m){
0ad3                    LL t = w[N/l*k] * p[k+m] % P;
6209                    p[k+m] = (p[k] - t + P) % P;
fa1b                    p[k] = (p[k] + t) % P;
95cf                }
95cf            }
95cf        }
427e
92ea    void ntt(LL* a){_ntt(a, omega);}
5daf    void intt(LL* a){
1f2a        LL inv = powmod(N, P-2);
9910        _ntt(a, oinv);
a873        rep (i, N) a[i] = a[i] * inv % P;
95cf    }
427e
3a5b    void conv(LL* a, LL* b){
ad16        ntt(a); ntt(b);
e49e        rep (i, N) a[i] = a[i] * b[i] % P;
5748        intt(a);
95cf    }
329b    };

```

#### 4.9 Sieve of Euler

```

cfc3    const int MAXX = 1e7+5;
5861    bool p[MAXX];
73ae    int prime[MAXX], sz;
427e
9bc6    void sieve(){
9628        p[0] = p[1] = 1;
1ec8        for (int i = 2; i < MAXX; i++){
bf28            if (!p[i]) prime[sz++] = i;
e82c            for (int j = 0; j < sz && i*prime[j] < MAXX; j++){
b6a9                p[i*prime[j]] = 1;

```

```

        if (i % prime[j] == 0) break;
    }
}

```

#### 4.10 Sieve of Euler (General)

```

namespace sieve {
constexpr int MAXN = 10000007;
bool p[MAXN]; // true if not prime
int prime[MAXN], sz;
int pval[MAXN], pcnt[MAXN];
int f[MAXN];

void exec(int N = MAXN) {
    p[0] = p[1] = 1;

    pval[1] = 1;
    pcnt[1] = 0;
    f[1] = 1;

    for (int i = 2; i < N; i++) {
        if (!p[i]) {
            prime[sz++] = i;
            for (LL j = i; j < N; j *= i) {
                int b = j / i;
                pval[j] = i * pval[b];
                pcnt[j] = pcnt[b] + 1;
                f[j] = _____; // f[j] = f(i^pcnt[j])
            }
        }
        for (int j = 0; i * prime[j] < N; j++) {
            int x = i * prime[j]; p[x] = 1;
            if (i % prime[j] == 0) {
                pval[x] = pval[i] * prime[j];
                pcnt[x] = pcnt[i] + 1;
            } else {
                pval[x] = prime[j];
                pcnt[x] = 1;
            }
            if (x != pval[x]) {
                f[x] = f[x / pval[x]] * f[pval[x]]

```

5f51  
95cf  
95cf  
95cf

b62e  
6589  
e982  
6ae8  
cbf7  
6030  
427e  
76f6  
9628  
427e  
8a8a  
bdda  
c6b9  
427e  
a643  
01d6  
b2b2  
37d9  
758c  
81fd  
e0f3  
a96c  
95cf  
95cf  
34c0  
f87a  
20cc  
9985  
3f93  
8e2e  
cc91  
6322  
95cf  
6191  
d614

```

95cf      }
5f51      if (i % prime[j] == 0) break;
95cf    }
95cf  }
95cf  }
95cf  }

```

## 4.11 Miller-Rabin primality test

The array `a[]` (excluding `senitel`, i.e. `LLONG_MAX`) should be

```

{2}                when  $n < 2,047$ .
{2, 7, 61}          when  $n < 4,759,123,141 (2^{32})$ .
{2, 3, 5, 7, 11}    when  $n < 2.1 \times 10^{12}$ .
{2, 325, 9375, 28178, 450775, 9780504, 1795265022} when  $n < 2^{64}$ .

```

```

f16f bool test(LL n){
59f2   if (n < 3) return n==2;
427e   // ! The array a[] should be modified if the range of x changes.
3f11   const LL a[] = {2LL, 7LL, 61LL, LLONG_MAX};
c320   LL r = 0, d = n-1, x;
f410   while (~d & 1) d >>= 1, r++;
2975   for (int i=0; a[i] < n; i++){
ece1     x = powmod(a[i], d, n); // ! powmod must use for 64bit mulmod
7f99     if (x == 1 || x == n-1) goto next;
e257     rep (i, r) {
d7ff       x = mulmod(x, x, n);
8d2e       if (x == n-1) goto next;
95cf     }
438e     return false;
d490 next:;
95cf   }
3361   return true;
95cf }

```

## 4.12 Pollard's rho algorithm

```

2e6b ULL gcd(ULL a, ULL b) {return b ? gcd(b, a % b) : a;}
427e
54a5 ULL PollardRho(ULL n){
45eb   ULL c, x, y, d = n;

```

```

if (~n&1) return 2;
while (d == n){
  x = y = 2;
  d = 1;
  c = rand() % (n - 1) + 1;
  while (d == 1){
    x = (mulmod(x, x, n) + c) % n;
    y = (mulmod(y, y, n) + c) % n;
    y = (mulmod(y, y, n) + c) % n;
    d = gcd(x>y ? x-y : y-x, n);
  }
}
return d;
}

```

```

d3e5
3c69
0964
4753
5952
9e5b
33d5
e1bf
e1bf
a313
95cf
95cf
5d89
95cf

```

# 5 Graph Theory

## 5.1 Strongly connected component

```

const int MAXV = 100005;

struct graph{
  vector<int> adj[MAXV];
  stack<int> s;
  int V; // number of vertices
  int pre[MAXV], lnk[MAXV], scc[MAXV];
  int time, sccn;

  void add_edge(int u, int v){
    adj[u].push_back(v);
  }

  void dfs(int u){
    pre[u] = lnk[u] = ++time;
    s.push(u);
    for (int v : adj[u]){
      if (!pre[v]){
        dfs(v);
        lnk[u] = min(lnk[u], lnk[v]);
      } else if (!scc[v]){
        lnk[u] = min(lnk[u], pre[v]);
      }
    }
  }
}

```

```

837c
427e
2ea0
88e3
9cad
3d02
8b6c
27ee
427e
bfab
c71a
95cf
427e
d714
7e41
80f6
18f6
173e
5f3c
002c
6068
d5df

```

```

95cf      }
95cf      }
8de2      if (lnk[u] == pre[u]){
660f          sccn++;
3c9e          int x;
a69f          do {
3834              x = s.top(); s.pop();
b0e9              scc[x] = sccn;
6757          } while (x != u);
95cf      }
95cf      }
427e
4c88      void find_scc(){
f4a2          time = sccn = 0;
8de7          memset(scc, 0, sizeof scc);
8c2f          memset(pre, 0, sizeof pre);
6901          Rep (i, V){
56d1              if (!pre[i]) dfs(i);
95cf          }
95cf      }
427e
27ce      vector<int> adjc[MAXV];
364d      void contract(){
1a1e          Rep (i, V)
21a2              rep (j, adj[i].size()){
b730                  if (scc[i] != scc[adj[i][j]])
b46e                      adjc[scc[i]].push_back(scc[adj[i][j]]);
95cf              }
95cf          }
329b      };

```

## 5.2 Vertex biconnected component

```

0f42      const int MAXN = 100005;
2ea0      struct graph {
33ae          int pre[MAXN], iscut[MAXN], bccno[MAXN], dfs_clock, bcc_cnt;
848f          vector<int> adj[MAXN], bcc[MAXN];
6b06          set<pair<int, int>> bcce[MAXN];
427e
76f7          stack<pair<int, int>> s;
427e
bfab          void add_edge(int u, int v) {

```

```

adj[u].push_back(v);
adj[v].push_back(u);
}

int dfs(int u, int fa) {
    int lowu = pre[u] = ++dfs_clock;
    int child = 0;
    for (int v : adj[u]) {
        if (!pre[v]) {
            s.push({u, v});
            child++;
            int lowv = dfs(v, u);
            lowu = min(lowu, lowv);
            if (lowv >= pre[u]) {
                iscut[u] = 1;
                bcc[bcc_cnt].clear();
                bcce[bcc_cnt].clear();
                while (1) {
                    int xu, xv;
                    tie(xu, xv) = s.top(); s.pop();
                    bcce[bcc_cnt].insert({min(xu, xv), max(xu, xv)});
                    if (bccno[xu] != bcc_cnt) {
                        bcc[bcc_cnt].push_back(xu);
                        bccno[xu] = bcc_cnt;
                    }
                    if (bccno[xv] != bcc_cnt) {
                        bcc[bcc_cnt].push_back(xv);
                        bccno[xv] = bcc_cnt;
                    }
                    if (xu == u && xv == v) break;
                }
                bcc_cnt++;
            }
        } else if (pre[v] < pre[u] && v != fa) {
            s.push({u, v});
            lowu = min(lowu, pre[v]);
        }
    }
    if (fa < 0 && child == 1) iscut[u] = 0;
    return lowu;
}

void find_bcc(int n) {
    memset(pre, 0, sizeof pre);

```

```

c71a
a717
95cf
427e
7d3c
9fe6
ec14
18f6
173e
e7f8
fdcf
f851
189c
b687
6323
57eb
90b8
a147
a6a3
a0c3
0ef5
3db2
e0db
d27f
95cf
f357
752b
57c9
95cf
7096
95cf
03f5
95cf
7470
e7f8
f115
95cf
95cf
e104
1160
95cf
427e
17be
8c2f

```

```
e2d2     memset(iscut, 0, sizeof iscut);
40d3     memset(bccno, -1, sizeof bccno);
fae2     dfs_clock = bcc_cnt = 0;
5c63     rep (i, n) if (!pre[i]) dfs(i, -1);
95cf     }
329b     };
```

### 5.3 Cut vertices

If the graph is unconnected, the algorithm should be run on each component. One may run `Rep (i, n) if (!dfn[i]) tarjan(i, i)` for unconnected graph.

#### Usage:

`add_edge(u, v)` Add an undirected edge  $(u, v)$ .  
`tarjan(u, fa)` Run Tarjan's algorithm on tree rooted at `fa`. Please call with identical `u` and `fa`.  
`cut[v]` Whether  $v$  is a cut vertex.

```
9f60 const int MAXN = 200005;
0b32 vector<int> adj[MAXN];
18e4 int dfn[MAXN], low[MAXN], idx;
d39d bool cut[MAXN];
427e
bfa8 void add_edge(int u, int v) {
c71a     adj[u].push_back(v);
a717     adj[v].push_back(u);
95cf }
427e
50aa void tarjan(int u, int fa) {
9891     dfn[u] = low[u] = ++idx;
ec14     int child = 0;
18f6     for (int v : adj[u]) {
3c64         if (!dfn[v]) {
9636             tarjan(v, fa); low[u] = min(low[u], low[v]);
f368             if (low[v] >= dfn[u] && u != fa) cut[u] = true;
7923             child += u == fa;
95cf         }
769a         low[u] = min(low[u], dfn[v]);
95cf     }
7927     if (u == fa && child > 1) cut[u] = true;
95cf }
```

### 5.4 Minimum spanning arborescence, faster

All vertices are 1-based. Clear the fields when reuse the struct.

#### Usage:

`add_edge(u, v, w)` Add an edge from  $u$  to  $v$  with weight  $w$ .  
`run(n, rt)` Compute the total weight of MSA rooted at `rt`. If not exist, return `LLONG_MIN`.

**Time Complexity:**  $O((|E| + |V| \log |V|) \log |V|)$

```
const int MAXN = 300005;
typedef pair<LL, int> pii;
struct MDST {
    priority_queue<pii, vector<pii>, greater<pii>> heap[MAXN];
    LL shift[MAXN];
    int fa[MAXN], vis[MAXN];

    int find(int x) { return fa[x] == x ? x : fa[x] = find(fa[x]); }

    void unite(int x, int y) {
        x = find(x); y = find(y); fa[y] = x; if (x == y) return;
        if (heap[x].size() < heap[y].size()) {
            swap(heap[x], heap[y]);
            swap(shift[x], shift[y]);
        }
        while (heap[y].size()) {
            auto p = heap[y].top(); heap[y].pop();
            heap[x].emplace(p.first - shift[y] + shift[x], p.second);
        }
    }

    void add_edge(int u, int v, LL w) { heap[v].emplace(w, u); }

    LL run(int n, int rt) {
        LL ans = 0;
        iota(fa, fa + n + 1, 0);
        Rep (i, n) if (find(i) != find(rt)) {
            int u = find(i);
            stack<int, vector<int>> s;
            while (find(u) != find(rt)) {
                if (vis[u]) while (s.top() != u) {
                    vis[s.top()] = 0; unite(u, s.top()); s.pop();
                } else { vis[u] = 1; s.push(u); }
                while (heap[u].size()) {
                    ans += heap[u].top().first - shift[u];
```

```
5ece
2fef
1495
01b2
321d
fc06
427e
38dd
427e
29b0
0c14
6fa0
9c26
2ffc
95cf
9959
175b
c0c5
95cf
95cf
427e
0bbd
427e
a526
f7ff
81f2
19b3
a7b1
010e
eff5
0dda
c593
83c4
c76e
b385
```

```

dde2         shift[u] = heap[u].top().first;
da47         if (find(heap[u].top().second) != u) break;
9fbb         heap[u].pop();
95cf     }
6961         if (heap[u].empty()) return LLONG_MIN;
87e6         u = find(heap[u].top().second);
95cf     }
2d46     while (s.size()) { vis[s.top()] = 0; unite(rt, s.top()); s.pop(); }
95cf }
4206     return ans;
95cf }
329b };

```

## 5.5 Maximum flow (Dinic)

### Usage:

add\_edge(u, v, c)      Add an edge from  $u$  to  $v$  with capacity  $c$ .

max\_flow(s, t)      Compute maximum flow from  $s$  to  $t$ .

**Time Complexity:** For general graph,  $O(V^2E)$ ; for network with unit capacity,  $O(\min\{V^{2/3}, \sqrt{E}\}E)$ ; for bipartite network,  $O(\sqrt{VE})$ .

```

bcf8 struct edge{
60e2     int from, to;
5e6d     LL cap, flow;
329b };
427e
e2cd const int MAXN = 1005;
9062 struct Dinic {
4dbf     int n, m, s, t;
9f0c     vector<edge> edges;
b891     vector<int> G[MAXN];
bbb6     bool vis[MAXN];
b40a     int d[MAXN];
ddec     int cur[MAXN];
427e
5973     void add_edge(int from, int to, LL cap) {
7b55         edges.push_back(edge{from, to, cap, 0});
1db7         edges.push_back(edge{to, from, 0, 0});
fe77         m = edges.size();
dff5         G[from].push_back(m-2);
8f2d         G[to].push_back(m-1);
95cf     }
427e

```

```

bool bfs() {
    memset(vis, 0, sizeof(vis));
    queue<int> q;
    q.push(s);
    vis[s] = 1;
    d[s] = 0;
    while (!q.empty()) {
        int x = q.front(); q.pop();
        for (int i = 0; i < G[x].size(); i++) {
            edge& e = edges[G[x][i]];
            if (!vis[e.to] && e.cap > e.flow) {
                vis[e.to] = 1;
                d[e.to] = d[x] + 1;
                q.push(e.to);
            }
        }
    }
    return vis[t];
}

LL dfs(int x, LL a) {
    if (x == t || a == 0) return a;
    LL flow = 0, f;
    for (int& i = cur[x]; i < G[x].size(); i++) {
        edge& e = edges[G[x][i]];
        if (d[x] + 1 == d[e.to] && (f = dfs(e.to, min(a, e.cap-e.flow))) > 0)
        {
            e.flow += f;
            edges[G[x][i]^1].flow -= f;
            flow += f;
            a -= f;
            if (a == 0) break;
        }
    }
    return flow;
}

LL max_flow(int s, int t) {
    this->s = s; this->t = t;
    LL flow = 0;
    while (bfs()) {
        memset(cur, 0, sizeof(cur));
        flow += dfs(s, LLONG_MAX);
    }
}

```



```

84fb     return flow;
95cf }
427e
c72e vector<int> min_cut() { // call this after maxflow
1df9     vector<int> ans;
df9a     for (int i = 0; i < edges.size(); i++) {
56d8         edge& e = edges[i];
46a2         if(vis[e.from] && !vis[e.to] && e.cap > 0) ans.push_back(i);
95cf     }
4206     return ans;
95cf }
329b };

```

## 5.6 Maximum cardinality bipartite matching (Hungarian)

```

302f #include <bits/stdc++.h>
421c using namespace std;
427e
0d6c #define rep(i, n) for (int i = 0; i < (n); i++)
cfe3 #define Rep(i, n) for (int i = 1; i <= (n); i++)
8843 #define range(x) (x).begin(), (x).end()
5cad typedef long long LL;
427e
84ee struct Hungarian{
fbf6     int nx, ny;
9ec6     vector<int> mx, my;
9d4c     vector<vector<int>> > e;
edec     vector<bool> mark;
427e
8324     void init(int nx, int ny){
c1d1         this->nx = nx;
f9c1         this->ny = ny;
ac92         mx.resize(nx); my.resize(ny);
3f11         e.clear(); e.resize(nx);
1023         mark.resize(nx);
95cf     }
427e
4589     inline void add(int a, int b){
486c         e[a].push_back(b);
95cf     }
427e
0c2b     bool augment(int i){

```

```

        if (!mark[i]) {
            mark[i] = true;
            for (int j : e[i]){
                if (my[j] == -1 || augment(my[j])){
                    mx[i] = j; my[j] = i;
                    return true;
                }
            }
        }
        return false;
    }

    int match(){
        int ret = 0;
        fill(range(mx), -1);
        fill(range(my), -1);
        rep (i, nx){
            fill(range(mark), false);
            if (augment(i)) ret++;
        }
        return ret;
    }
};

```

207c  
dae4  
6a1e  
0892  
9ca3  
3361  
95cf  
95cf  
95cf  
438e  
95cf  
427e  
3fac  
5b57  
b0f1  
b957  
4ed1  
13a5  
cc89  
95cf  
ee0f  
95cf  
329b

## 5.7 Maximum matching of general graph (Edmond's blossom)

### Usage:

init(n)	Initialize the template with $n$ vertices, numbered from 1.
add_edge(u, v)	Add an undirected edge $uv$ .
solve()	Find the maximum matching. Return the number of matched edges.
mate[]	The mate of a matched vertex. If it is not matched, then the value is 0.

**Time Complexity:**  $O(|V|^3)$ , but extremely fast in practice.

```

const int MAXN = 1024;
struct Blossom {
    vector<int> adj[MAXN];
    queue<int> q;
    int n;
    int label[MAXN], mate[MAXN], save[MAXN], used[MAXN];

```

c041  
6ab1  
0b32  
93d2  
5c83  
0de2  
427e

```

2186 void init(int nv) {
3728     n = nv; for (auto& v : adj) v.clear();
477d     fill(range(label), 0); fill(range(mate), 0);
bb35     fill(range(save), 0); fill(range(used), 0);
95cf }
427e
c2dd void add_edge(int u, int v) { adj[u].push_back(v); adj[v].push_back(u); }
427e
2a48 void rematch(int x, int y) {
8af8     int m = mate[x]; mate[x] = y;
1aa4     if (mate[m] == x) {
f4ba         if (label[x] <= n) {
740a             mate[m] = label[x]; rematch(label[x], m);
8e2e         } else {
3341             int a = 1 + (label[x] - n - 1) / n;
2885             int b = 1 + (label[x] - n - 1) % n;
ef33             rematch(a, b); rematch(b, a);
95cf         }
95cf     }
95cf }
427e
8a50 void traverse(int x) {
43c0     Rep (i, n) save[i] = mate[i];
2ef7     rematch(x, x);
34d7     Rep (i, n) {
62c5         if (mate[i] != save[i]) used[i] ++;
97ef         mate[i] = save[i];
95cf     }
95cf }
427e
8bf8 void relabel(int x, int y) {
d101     Rep (i, n) used[i] = 0;
c4ea     traverse(x); traverse(y);
34d7     Rep (i, n) {
dee9         if (used[i] == 1 and label[i] < 0) {
1c22             label[i] = n + x + (y - 1) * n;
eb31             q.push(i);
95cf         }
95cf     }
95cf }
427e
a0ce int solve() {
34d7     Rep (i, n) {
a073         if (mate[i]) continue;

```

```

Rep (j, n) label[j] = -1;
label[i] = 0; q = queue<int>(); q.push(i);
while (q.size()) {
    int x = q.front(); q.pop();
    for (int y : adj[x]) {
        if (mate[y] == 0 and i != y) {
            mate[y] = x; rematch(x, y); q = queue<int>(); break;
        }
        if (label[y] >= 0) { relabel(x, y); continue; }
        if (label[mate[y]] < 0) {
            label[mate[y]] = x; q.push(mate[y]);
        }
    }
}
int cnt = 0;
Rep (i, n) cnt += (mate[i] > i);
return cnt;
}
};

```

## 5.8 Minimum cost maximum flow

```

struct edge{
    int from, to;
    int cap, flow;
    LL cost;
};

const LL INF = LLONG_MAX / 2;
const int MAXN = 5005;
struct MCMF {
    int s, t, n, m;
    vector<edge> edges;
    vector<int> G[MAXN];
    bool inq[MAXN]; // queue
    LL d[MAXN]; // distance
    int p[MAXN]; // previous
    int a[MAXN]; // improvement

    void add_edge(int from, int to, int cap, LL cost) {
        edges.push_back(edge{from, to, cap, 0, cost});
    }
};

```

```

95f0     edges.push_back(edge{to, from, 0, 0, -cost});
fe77     m = edges.size();
dff5     G[from].push_back(m-2);
8f2d     G[to].push_back(m-1);
95cf }
427e
3c52     bool spfa(){
93d2         queue<int> q;
8494         fill(d, d + MAXN, INF); d[s] = 0;
fd48         memset(inq, 0, sizeof(inq));
5e7c         q.push(s); inq[s] = true;
2dae         p[s] = 0; a[s] = INT_MAX;
cc78         while (!q.empty()){
b0aa             int u = q.front(); q.pop(); inq[u] = false;
3bba             for (int i : G[u]) {
56d8                 edge& e = edges[i];
3601                 if (e.cap > e.flow && d[e.to] > d[u] + e.cost){
55bc                     d[e.to] = d[u] + e.cost;
0bea                     p[e.to] = G[u][i];
8249                     a[e.to] = min(a[u], e.cap - e.flow);
e5d3                     if (!inq[e.to]) q.push(e.to), inq[e.to] = true;
95cf                 }
95cf             }
95cf         }
6d7c         return d[t] != INF;
95cf     }
427e
71a4     void augment(){
06f1         int u = t;
b19d         while (u != s){
db09             edges[p[u]].flow += a[t];
25a9             edges[p[u]^1].flow -= a[t];
e6c9             u = edges[p[u]].from;
95cf         }
95cf     }
427e
6e20     #ifndef GIVEN_FLOW
5972         bool min_cost(int s, int t, int f, LL& cost) {
590d             this->s = s; this->t = t;
21d4             int flow = 0;
23cb             cost = 0;
22dc             while (spfa()) {
bcd8                 augment();
a671                 if (flow + a[t] >= f){

```

```

cost += (f - flow) * d[t]; flow = f;
return true;
} else {
    flow += a[t]; cost += a[t] * d[t];
}
}
return false;
}
#else
int min_cost(int s, int t, LL& cost) {
    this->s = s; this->t = t;
    int flow = 0;
    cost = 0;
    while (spfa()) {
        augment();
        flow += a[t]; cost += a[t] * d[t];
    }
    return flow;
}
#endif
};

```

b14d  
3361  
8e2e  
2a83  
95cf  
95cf  
438e  
95cf  
a8cb  
f9a9  
590d  
21d4  
23cb  
22dc  
bcd8  
2a83  
95cf  
84fb  
95cf  
1937  
329b

## 5.9 Global minimum cut (Stoer-Wagner)

### Usage:

stoer(w)

Compute the global minimum cut of the graph specified by the **symmetric** adjacent matrix w (0-based). Return the capacity of the cut and the indices of one part of the cut.

**Time Complexity:**  $O(|V|^3)$

```

typedef vector<LL> VI;
typedef vector<VI> VVI;

```

```

pair<LL, VI> stoer(VI &w) {
    int n = w.size();
    VI used(n), c, bestc;
    LL bestw = -1;

    for (int ph = n - 1; ph >= 0; ph--) {
        VI wt = w[0], added = used;
        int prev, last = 0;
        rep (i, ph) {

```

f9d7  
045e  
427e  
f012  
66f7  
4d98  
329d  
427e  
cd21  
ec6e  
f20e  
4b32

```

8bfc     prev = last;
0706     last = -1;
4942     for (int j = 1; j < n; j++)
c4b9         if (!added[j] && (last == -1 || wt[j] > wt[last]))
887d             last = j;
71bc     if (i == ph - 1) {
9cfa         rep (j, n) w[prev][j] += w[last][j];
1f25         rep (j, n) w[j][prev] = w[prev][j];
5613         used[last] = true;
8e11         c.push_back(last);
bb8e         if (bestw == -1 || wt[last] < bestw) {
bab6             bestc = c;
372e             bestw = wt[last];
95cf         }
8e2e     } else {
caeb         rep (j, n) wt[j] += w[last][j];
8b92         added[last] = true;
95cf     }
95cf }
95cf }
038c     return {bestw, bestc};
95cf }

```

## 5.10 Fast LCA

All indices of the tree are 1-based.

### Usage:

preprocess(root)      Initialize with tree rooted at root.  
lca(u, v)              Query the lowest common ancestor of  $u$  and  $v$ .

```

0e34     const int MAXN = 500005;
0b32     vector<int> adj[MAXN];
fccb     int id[MAXN], nid;
1356     pair<int, int> st[MAXN << 1][33 - __builtin_clz(MAXN)];
427e
e16d     void dfs(int u, int p, int d) {
0df2         st[id[u] = nid++][0] = {d, u};
18f6         for (int v : adj[u]) {
bd87             if (v == p) continue;
f58c             dfs(v, u, d + 1);
08ad             st[nid++][0] = {d, u};
95cf         }
95cf     }

```

```

void preprocess(int root) {
    nid = 0;
    dfs(root, 0, 1);
    int l = 31 - __builtin_clz(nid);
    rep (j, l) rep (i, 1+nid-(1<<j))
        st[i][j+1] = min(st[i][j], st[i+(1<<j)][j]);
}

int lca(int u, int v) {
    tie(u, v) = minmax(id[u], id[v]);
    int k = 31 - __builtin_clz(v-u+1);
    return min(st[u][k], st[v-(1<<k)+1][k]).second;
}

```

## 5.11 Heavy-light decomposition

**Time Complexity:** The decomposition itself takes linear time. Each query takes  $O(\log n)$  operations.

```

const int MAXN = 100005;
vector<int> adj[MAXN];
int sz[MAXN], top[MAXN], fa[MAXN], son[MAXN], depth[MAXN], id[MAXN];

void dfs1(int x, int dep, int par){
    depth[x] = dep;
    sz[x] = 1;
    fa[x] = par;
    int maxn = 0, s = 0;
    for (int c: adj[x]){
        if (c == par) continue;
        dfs1(c, dep + 1, x);
        sz[x] += sz[c];
        if (sz[c] > maxn){
            maxn = sz[c];
            s = c;
        }
    }
    son[x] = s;
}

int cid = 0;
void dfs2(int x, int t){

```

```

8d96     top[x] = t;
d314     id[x] = ++cid;
c4a1     if (son[x]) dfs2(son[x], t);
c861     for (int c: adj[x]){
9881         if (c == fa[x]) continue;
5518         if (c == son[x]) continue;
13f9         else dfs2(c, c);
95cf     }
95cf }
427e
0f04 void decomp(int root){
9fa4     dfs1(root, 1, 0);
1c88     dfs2(root, root);
95cf }
427e
2c98 void query(int u, int v){
03a1     while (top[u] != top[v]){
45ec         if (depth[top[u]] < depth[top[v]]) swap(u, v);
427e         // id[top[u]] to id[u]
005b         u = fa[top[u]];
95cf     }
6083     if (depth[u] > depth[v]) swap(u, v);
427e     // id[u] to id[v]
95cf }

```

## 5.12 Centroid decomposition

Note that the centroid here is not the exact centroid of the graph. It only guarantees that the size of each subtree does not exceed half of that of the original tree. This is enough to guarantee the correct time complexity. All vertices are numbered from 1. Call `decomp(root)` to use.

### Usage:

`decomp(u, p)` Decompose the tree rooted at  $u$  with parent  $p$ .

**Time Complexity:** The decomposition itself takes  $O(n \log n)$  time.

```

1fb6 vector<int> adj[100005];
88e0 int sz[100005], sum;
427e
f93d void getsz(int u, int p) {
5b36     sz[u] = 1; sum++;
18f6     for (int v : adj[u]) {
bd87         if (v == p) continue;
e3cb         getsz(v, u);

```

```

        sz[u] += sz[v];
    }
}

int getcent(int u, int p) {
    for (int v : adj[u])
        if (v != p and sz[v] > sum / 2)
            return getcent(v, u);
    return u;
}

void decompose(int u) {
    sum = 0; getsz(u, 0);
    u = getcent(u, 0); // update u to the centroid

    for (int v : adj[u]) {
        // get answer for subtree v
    }
    // get answer for the whole tree
    // don't forget to count the centroid itself

    for (int v : adj[u]) { // divide and conquer
        adj[v].erase(find(range(adj[v]), u));
        decompose(v);
        adj[v].push_back(u); // restore deleted edge
    }
}

```

## 5.13 DSU on tree

This implementation avoids parallel existence of multiple data structures but requires that the data structure is invertible. To use this template, implement `merge`, `enter`, `leave` as needed; first call `decomp(root, 0)`, then call `work(root, 0, false)`. Labels of vertices start from 1.

### Usage:

`decomp(u, p)` Decompose the tree  $u$ .  
`work(u, p, keep)` Work for subtree  $u$ . When `keep` is set, information is not cleared.

**Time Complexity:**  $O(n \log n)$  times the complexity for `merge`, `enter`, `leave`.

```

vector<int> adj[100005];
int sz[100005], son[100005];

```

```

8449
95cf
95cf
427e
67f9
d51f
76e4
18e3
81b0
95cf
427e
4662
618e
303c
427e
18f6
427e
95cf
427e
427e
427e
18f6
c375
fa6b
a717
95cf
95cf

```

```

1fb6
901d

```

```

427e void decomp(int u, int p) {
5559     sz[u] = 1;
50c0     for (int v : adj[u]) {
18f6         if (v == p) continue;
bd87         decomp(v, u);
a851         sz[u] += sz[v];
8449         if (sz[v] > sz[son[u]]) son[u] = v;
d28c     }
95cf }
95cf
427e template <typename T>
b7ec void trav(T fn, int u, int p) {
62f5     fn(u);
4412     for (int v : adj[u]) if (v != p) trav(fn, v, u);
30b3 }
95cf
427e #define for_light(v) for (int v : adj[u]) if (v != p and v != son[u])
7467 void work(int u, int p, bool keep) {
33ff     for_light(v) work(v, u, 0); // process light children
72a2
427e     // process heavy child
427e     // current data structure contains info of heavy child
9866     if (son[u]) work(son[u], u, 1);
427e
18a9     auto merge = [u] (int c) { /* count contribution of c */ };
1ab0     auto enter = [] (int c) { /* add vertex c */ };
f241     auto leave = [] (int c) { /* remove vertex c */ };
427e
3d3b     for_light(v) {
74c6         trav(merge, v, u);
c13d         trav(enter, v, u);
95cf     }
427e
427e     // count answer for root and add it
427e     // Warning: special check may apply to root!
c54f     merge(u);
9dec     enter(u);
427e
427e     // Leave current tree
4e3e     if (!keep) trav(leave, u, p);
95cf }

```

## 6 Data Structures

### 6.1 Fenwick tree (point update range query)

```

struct bit_purq { // point update, range query
    int N;
    vector<LL> tr;

    void init(int n) { tr.resize(N = n + 5); }

    LL sum(int n) {
        LL ans = 0;
        while (n) { ans += tr[n]; n &= n - 1; }
        return ans;
    }

    void add(int n, LL x){
        while (n < N) { tr[n] += x; n += n & -n; }
    }
};

```

9976  
d7af  
99ff  
427e  
456d  
427e  
63d0  
f7ff  
6770  
4206  
95cf  
427e  
f4bd  
968e  
95cf  
329b

### 6.2 Fenwick tree (range update point query)

```

struct bit_rupq{ // range update, point query
    int N;
    vector<LL> tr;

    void init(int n) { tr.resize(N = n + 5);}

    LL query(int n) {
        LL ans = 0;
        while (n < N) { ans += tr[n]; n += n & -n; }
        return ans;
    }

    void add(int n, LL x) {
        while (n) { tr[n] += x; n &= n - 1; }
    }
};

```

3d03  
d7af  
99ff  
427e  
456d  
427e  
38d4  
f7ff  
3667  
4206  
95cf  
427e  
f4bd  
0a2b  
95cf  
329b

### 6.3 Segment tree

```

3942 LL p;
1ebb const int MAXN = 4 * 100006;
451a struct segtree {
27be     int l[MAXN], m[MAXN], r[MAXN];
4510     LL val[MAXN], tadd[MAXN], tmul[MAXN];
427e
ac35 #define lson (o<<1)
1294 #define rson (o<<1|1)
427e
1344 void pull(int o) {
bbe9     val[o] = (val[lson] + val[rson]) % p;
95cf }
427e
e4bc void push_add(int o, LL x) {
5dd6     val[o] = (val[o] + x * (r[o] - l[o])) % p;
6eff     tadd[o] = (tadd[o] + x) % p;
95cf }
427e
d658 void push_mul(int o, LL x) {
b82c     val[o] = val[o] * x % p;
aa86     tadd[o] = tadd[o] * x % p;
649f     tmul[o] = tmul[o] * x % p;
95cf }
427e
b149 void push(int o) {
3159     if (l[o] == m[o]) return;
0a90     if (tmul[o] != 1) {
0f4a         push_mul(lson, tmul[o]);
045e         push_mul(rson, tmul[o]);
ac0a         tmul[o] = 1;
95cf     }
1b82     if (tadd[o]) {
9547         push_add(lson, tadd[o]);
0e73         push_add(rson, tadd[o]);
6234         tadd[o] = 0;
95cf     }
95cf }
427e
471c void build(int o, int ll, int rr) {
0e87     int mm = (ll + rr) / 2;
9d27     l[o] = ll; r[o] = rr; m[o] = mm;

```

```

tmul[o] = 1;
if (ll == mm) {
    scanf("%lld", val + o);
    val[o] %= p;
} else {
    build(lson, ll, mm);
    build(rson, mm, rr);
    pull(o);
}
}

void add(int o, int ll, int rr, LL x) {
    if (ll <= l[o] && r[o] <= rr) {
        push_add(o, x);
    } else {
        push(o);
        if (m[o] > ll) add(lson, ll, rr, x);
        if (m[o] < rr) add(rson, ll, rr, x);
        pull(o);
    }
}

void mul(int o, int ll, int rr, LL x) {
    if (ll <= l[o] && r[o] <= rr) {
        push_mul(o, x);
    } else {
        push(o);
        if (ll < m[o]) mul(lson, ll, rr, x);
        if (m[o] < rr) mul(rson, ll, rr, x);
        pull(o);
    }
}

LL query(int o, int ll, int rr) {
    if (ll <= l[o] && r[o] <= rr) {
        return val[o];
    } else {
        push(o);
        if (rr <= m[o]) return query(lson, ll, rr);
        if (ll >= m[o]) return query(rson, ll, rr);
        return query(lson, ll, rr) + query(rson, ll, rr);
    }
}
} seg;

```

```

ac0a
5c92
001f
e5b6
8e2e
7293
5e67
ba26
95cf
95cf
427e
4406
3c16
db32
8e2e
c4b0
4305
d5a6
ba26
95cf
95cf
427e
48cd
3c16
e7d0
8e2e
c4b0
d1ba
67f3
ba26
95cf
95cf
427e
0f62
3c16
6dfe
8e2e
c4b0
462a
5cca
bbf9
95cf
95cf
4d99

```

## 6.4 Treap

Self-balanced binary search tree which supports split and merge.

### Usage:

push(x)	Push lazy tags to children.
pull(x)	Update statistics of node $x$ .
Init(x, v)	Initialize node $x$ with value $v$ .
Add(x, v)	Apply addition to subtree $x$ .
Reverse(x)	Apply reversion to subtree $x$ .
Merge(x, y)	Merge trees rooted at $x$ and $y$ . Return the root of new tree.
Split(t, k, x, y)	Split out the left $k$ elements of tree $t$ . The roots of left part and right part are stored in $x$ and $y$ , respectively.
init(n)	Initialize the treap with array of size $n$ .
work(op, l, r)	Range operation over $[l, r)$ .

**Time Complexity:** Expected  $O(\log n)$  per operation.

```

9f60 const int MAXN = 200005;
a7c5 mt19937 gen(time(NULL));
9542 struct Treap {
6d61     int ch[MAXN][2];
3948     int sz[MAXN], key[MAXN], val[MAXN];
5d9a     int add[MAXN], rev[MAXN];
2b1b     LL sum[MAXN] = {0};
a773     int maxv[MAXN] = {INT_MIN}, minv[MAXN] = {INT_MAX};
427e
a629     void Init(int x, int v) {
5a00         ch[x][0] = ch[x][1] = 0;
d8cd         key[x] = gen(); val[x] = v; pull(x);
95cf     }
427e
3bf9     void pull(int x) {
e1c3         sz[x] = 1 + sz[ch[x][0]] + sz[ch[x][1]];
99f8         sum[x] = val[x] + sum[ch[x][0]] + sum[ch[x][1]];
94e9         maxv[x] = max({val[x], maxv[ch[x][0]], maxv[ch[x][1]]});
6bb9         minv[x] = min({val[x], minv[ch[x][0]], minv[ch[x][1]]});
95cf     }
427e
8c8e     void Add(int x, int a) {
a7b1         val[x] += a; add[x] += a;
832a         sum[x] += LL(sz[x]) * a; maxv[x] += a; minv[x] += a;

```

```

}

void Reverse(int x) {
    rev[x] ^= 1;
    swap(ch[x][0], ch[x][1]);
}

void push(int x) {
    for (int c : ch[x]) if (c) {
        Add(c, add[x]);
        if (rev[x]) Reverse(c);
    }
    add[x] = 0; rev[x] = 0;
}

int Merge(int x, int y) {
    if (!x || !y) return x | y;
    push(x); push(y);
    if (key[x] > key[y]) {
        ch[x][1] = Merge(ch[x][1], y); pull(x); return x;
    } else {
        ch[y][0] = Merge(x, ch[y][0]); pull(y); return y;
    }
}

void Split(int t, int k, int &x, int &y) {
    if (t == 0) { x = y = 0; return; }
    push(t);
    if (sz[ch[t][0]] < k) {
        x = t; Split(ch[t][1], k - sz[ch[t][0]] - 1, ch[t][1], y);
    } else {
        y = t; Split(ch[t][0], k, x, ch[t][0]);
    }
    if (x) pull(x); if (y) pull(y);
}

} treap;

int root;

void init(int n) {
    Rep(i, n) {
        int x; scanf("%d", &x);
        treap.Init(i, x);
        root = (i == 1) ? 1 : treap.Merge(root, i);
    }
}

```

95cf  
427e  
aaf6  
52c6  
7850  
95cf  
427e  
1a53  
5fe5  
fd76  
7a53  
95cf  
49ee  
95cf  
427e  
9d2c  
1b09  
cd7e  
bfffa  
a3df  
8e2e  
bf9e  
95cf  
95cf  
427e  
dc7e  
6303  
f26b  
3465  
ffd8  
8e2e  
8a23  
95cf  
89e3  
95cf  
b1f4  
427e  
24b6  
427e  
d34f  
34d7  
7681  
0ed8  
bcc8



```

95cf     }
95cf }
427e
d030 void work(int op, int l, int r) {
6639     int tl, tm, tr;
b6c4     treap.Split(root, l, tl, tm);
8de3     treap.Split(tm, r - 1, tm, tr);
3658     if (op == 1) {
c039         int x; scanf("%d", &x); treap.Add(tm, x);
1dcb     } else if (op == 2) {
ae78         treap.Reverse(tm);
581d     } else if (op == 3) {
e092         printf("%lld_%d_%d\n",
867f             treap.sum[tm], treap.minv[tm], treap.maxv[tm]);
95cf     }
6188     root = treap.Merge(treap.Merge(tl, tm), tr);
95cf }

```

## 6.5 Link/cut tree

Dynamic connectivity of undirected acyclic graph. Support single-vertex update, path aggregation and relative LCA query. Vertices are numbered from 1. Zero initialization is enough except for the statistic information.

### Usage:

<code>pull(x)</code>	Update statistics of node $x$ .
<code>Root(u)</code>	Get the root of tree where vertex $u$ is in.
<code>Link(u, v)</code>	Link two unconnected trees.
<code>Cut(u, v)</code>	Cut an existent edge.
<code>Query(u, v)</code>	Path aggregation.
<code>Update(u, x)</code>	Single point modification.
<code>LCA(u, v, root)</code>	Get the lowest common ancestor of $u$ and $v$ in tree rooted at root.

**Time Complexity:**  $O(\log n)$  per operation

```

2e73 const int MAXN = 1000005;
ca06 struct LCT {
6a6d     int fa[MAXN], ch[MAXN][2], val[MAXN], sum[MAXN];
c6e1     bool rev[MAXN];
427e
eba3     bool isroot(int x) { return ch[fa[x]][0] == x || ch[fa[x]][1] == x; }
f19f     void pull(int x) { sum[x] = val[x] ^ sum[ch[x][0]] ^ sum[ch[x][1]]; }
1c4d     void reverse(int x) { swap(ch[x][0], ch[x][1]); rev[x] ^= 1; }

```

```

void push(int x) {
    if (rev[x]) rep (i, 2) if (ch[x][i]) reverse(ch[x][i]); rev[x] = 0;
}
void rotate(int x) {
    int y = fa[x], z = fa[y], k = ch[y][1] == x, w = ch[x][!k];
    if (isroot(y)) ch[z][ch[z][1] == y] = x;
    ch[x][!k] = y; ch[y][k] = w; if (w) fa[w] = y;
    fa[y] = x; fa[x] = z; pull(y);
}
void pushall(int x) { if (isroot(x)) pushall(fa[x]); push(x); }
void splay(int x) {
    int y = x, z = 0;
    for (pushall(y); isroot(x); rotate(x)) {
        y = fa[x]; z = fa[y];
        if (isroot(y)) rotate((ch[y][0] == x) ^ (ch[z][0] == y) ? x : y);
    }
    pull(x);
}
void access(int x) {
    int z = x;
    for (int y = 0; x; x = fa[y = x]) { splay(x); ch[x][1] = y; pull(x); }
    splay(z);
}
void chroot(int x) { access(x); reverse(x); }
void split(int x, int y) { chroot(x); access(y); }

int Root(int x) {
    for (access(x); ch[x][0]; x = ch[x][0]) push(x);
    splay(x); return x;
}
void Link(int u, int v) { chroot(u); fa[u] = v; }
void Cut(int u, int v) { split(u, v); fa[u] = ch[v][0] = 0; pull(v); }
int Query(int u, int v) { split(u, v); return sum[v]; }
void Update(int u, int x) { splay(u); val[u] = x; }
int LCA(int x, int y, int root) {
    chroot(root); access(x); splay(y);
    while (fa[y]) splay(y = fa[y]);
    return y;
}
};

```

1a53  
89a0  
95cf  
425f  
51af  
e1fe  
1e6f  
6d09  
95cf  
52c6  
f69c  
d095  
c494  
ceef  
4449  
95cf  
78a0  
95cf  
6229  
1548  
8854  
7afd  
95cf  
a067  
126d  
427e  
d87a  
f4f1  
0d77  
95cf  
9e46  
7c10  
0691  
a999  
1f42  
6cb2  
02e5  
c218  
95cf  
329b

## 6.6 Balanced binary search tree from pb\_ds

```

0475 #include <ext/pb_ds/assoc_container.hpp>
332d using namespace __gnu_pbds;
427e
43a7 tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update>
    rkt;
427e // null_tree_node_update
427e
427e // SAMPLE USAGE
190e rkt.insert(x); // insert element
05d4 rkt.erase(x); // erase element
add5 rkt.order_of_key(x); // obtain the number of elements less than x
b064 rkt.find_by_order(i); // iterator to i-th (numbered from 0) smallest element
c103 rkt.lower_bound(x);
4ff4 rkt.upper_bound(x);
b19b rkt.join(rkt2); // merge tree (only if their ranges do not intersect)
cb47 rkt.split(x, rkt2); // split all elements greater than x to rkt2

```

## 6.7 Persistent segment tree, range k-th query

```

f1a7 struct node {
2ff6     static int n, pos;
427e
7cec     int value;
70e2     node *left, *right;
427e
20b0     void* operator new(size_t size);
427e
3dc0     static node* Build(int l, int r) {
b6c5         node* a = new node;
ce96         if (r > l + 1) {
181e             int mid = (l + r) / 2;
3ba2             a->left = Build(l, mid);
8aaf             a->right = Build(mid, r);
8e2e         } else {
bfc4             a->value = 0;
95cf         }
5ffd         return a;
95cf     }
427e
5a45     static node* init(int size) {
2c46         n = size;

```

```

pos = 0;
return Build(0, n);
}

```

```

static int Query(node* lt, node *rt, int l, int r, int k) {
    if (r == l + 1) return l;
    int mid = (l + r) / 2;
    if (rt->left->value - lt->left->value < k) {
        k -= rt->left->value - lt->left->value;
        return Query(lt->right, rt->right, mid, r, k);
    } else {
        return Query(lt->left, rt->left, l, mid, k);
    }
}

```

```

static int query(node* lt, node *rt, int k) {
    return Query(lt, rt, 0, n, k);
}

```

```

node *Inc(int l, int r, int pos) const {
    node* a = new node(*this);
    if (r > l + 1) {
        int mid = (l + r) / 2;
        if (pos < mid)
            a->left = left->Inc(l, mid, pos);
        else
            a->right = right->Inc(mid, r, pos);
    }
    a->value++;
    return a;
}

```

```

node *inc(int index) {
    return Inc(0, n, index);
}
nodes[8000000];

```

```

int node::n, node::pos;
inline void* node::operator new(size_t size) {
    return nodes + (pos++);
}

```

```

7ee3
be52
95cf
427e
93c0
d30c
181e
cb5a
8edb
2412
8e2e
0119
95cf
95cf
427e
c9ad
9e27
95cf
427e
b19c
5794
ce96
181e
203d
f44a
649a
1024
95cf
2b3e
5ffd
95cf
427e
e80f
c246
95cf
865a
427e
99ce
1987
bb3c
95cf

```

## 6.8 Block list

All indices are 0-based. All ranges are left-closed right-open.

### Usage:

<code>block::fix()</code>	Apply tags to the current block.
<code>Init(l, r)</code>	Range initializer.
<code>Reverse(l, r)</code>	Reverse the range.
<code>Add(l, r, x)</code>	Add $x$ to the range.
<code>Query(l, r)</code>	Range aggregation.

```
fd9e const int BLOCK = 800;
76b3 typedef vector<int> vi;
427e
a771 struct block {
8fbc     vi data;
e3b5     LL sum; int minv, maxv;
41db     int add; bool rev;
427e
d7eb     block(vi&& vec) : data(move(vec)),
1f0c         sum(accumulate(range(data), 0ll)),
8216         minv(*min_element(range(data))),
527d         maxv(*max_element(range(data))),
6437         add(0), rev(0) { }
427e
b919     void fix() {
0694         if (rev) reverse(range(data));         rev = 0;
0527         if (add) for (int& x : data) x += add;   add = 0;
95cf     }
427e
8bc4     void merge(block& another) {
b895         fix(); another.fix();
f516         vi temp(move(data));
d02c         temp.insert(temp.end(), range(another.data));
88ea         *this = block(move(temp));
95cf     }
427e
42e8     block split(int pos) {
3e79         fix();
ccab         block result(vi(data.begin() + pos, data.end()));
861a         data.resize(pos); *this = block(move(data));
56b0         return result;
95cf     }
329b };
427e
```

```
typedef list<block>::iterator lit;

struct blocklist {
    list<block> blk;

    void maintain() {
        lit it = blk.begin();
        while (it != blk.end() && next(it) != blk.end()) {
            lit it2 = it;
            while (next(it2) != blk.end() &&
                it2->data.size() + next(it2)->data.size() <= BLOCK) {
                it2->merge(*next(it2));
                blk.erase(next(it2));
            }
            ++it;
        }
    }

    lit split(int pos) {
        for (lit it = blk.begin(); ; it++) {
            if (pos == 0) return it;
            while (it->data.size() > pos)
                blk.insert(next(it), it->split(pos));
            pos -= it->data.size();
        }
    }

    void Init(int *l, int *r) {
        for (int *cur = l; cur < r; cur += BLOCK)
            blk.emplace_back(vi(cur, min(cur + BLOCK, r)));
    }

    void Reverse(int l, int r) {
        lit it = split(l), it2 = split(r);
        reverse(it, it2);
        while (it != it2) {
            it->rev ^= 1;
            it++;
        }
        maintain();
    }

    void Add(int l, int r, int x) {
```

```
2a18
427e
ce14
5540
427e
7b8e
3131
4628
852d
188c
3600
93e1
e1fa
95cf
5771
95cf
95cf
427e
b7b3
2273
5502
8e85
2099
a5a1
427e
95cf
95cf
427e
1c7b
9919
8950
95cf
427e
a22f
997b
dfd0
8f89
6a06
5283
95cf
b204
95cf
427e
3cce
```

```

997b     lit it = split(l), it2 = split(r);
8f89     while (it != it2) {
e927         it->sum += LL(x) * it->data.size();
03d3         it->minv += x; it->maxv += x;
4511         it->add += x; it++;
95cf     }
b204     maintain();
95cf }
427e
3ad3     void Query(int l, int r) {
997b         lit it = split(l), it2 = split(r);
c33d         LL sum = 0; int minv = INT_MAX, maxv = INT_MIN;
8f89         while (it != it2) {
e472             sum += it->sum;
72c4             minv = min(minv, it->minv);
e1c4             maxv = max(maxv, it->maxv);
5283             it++;
95cf         }
b204         maintain();
8792         printf("%lld_%d_%d\n", sum, minv, maxv);
95cf     }
958e } lst;

```

## 6.9 Persistent block list

Block list that supports persistence. All indices are 0-based. All ranges are left-closed right-open. `std::shared_ptr` is used to ease memory management. One should modify the constructor of `block` to maintain extra information. Here we use this policy that the size of each block does not exceed `BLOCK`, while the sum of sizes of two adjacent blocks does not less than `BLOCK`.

When some operation that breaks block list property, please call `maintain` in time to restore the property.

### Usage:

<code>maintain()</code>	Maintain the block list property.
<code>split(pos)</code>	Split the block list at position <code>pos</code> . Returns an iterator to a block starting at <code>pos</code> .
<code>sum(l, r)</code>	An example function of list traversal between $[l, r)$ .

**Time Complexity:** When `BLOCK` is properly selected, the time complexity is  $O(\sqrt{n})$  per operation.

```

a19e constexpr int BLOCK = 800;
76b3 typedef vector<int> vi;

```

```

typedef shared_ptr<vi> pvi;
typedef shared_ptr<const vi> pcvi;

struct block {
    pcvi data;
    LL sum;

    // add information to maintain
    block(pcvi ptr) :
        data(ptr),
        sum(accumulate(ptr->begin(), ptr->end(), 0ll))
    { }

    void merge(const block& another) {
        pvi temp = make_shared<vi>(data->begin(), data->end());
        temp->insert(temp->end(), another.data->begin(), another.data->end());
        *this = block(temp);
    }

    block split(int pos) {
        block result(make_shared<vi>(data->begin() + pos, data->end()));
        *this = block(make_shared<vi>(data->begin(), data->begin() + pos));
        return result;
    }
};

typedef list<block>::iterator lit;

struct blocklist {
    list<block> blk;

    void maintain() {
        lit it = blk.begin();
        while (it != blk.end() and next(it) != blk.end()) {
            lit it2 = it;
            while (next(it2) != blk.end() and
                it2->data->size() + next(it2)->data->size() <= BLOCK) {
                it2->merge(*next(it2));
                blk.erase(next(it2));
            }
            ++it;
        }
    }
};

```

0563  
013b  
427e  
a771  
2989  
8fd0  
427e  
427e  
a613  
24b5  
0cf0  
e93b  
427e  
5c0f  
0b18  
ac21  
6467  
95cf  
427e  
42e8  
dac1  
01db  
56b0  
95cf  
329b  
427e  
2a18  
427e  
ce14  
5540  
427e  
7b8e  
3131  
5e44  
852d  
0b03  
029f  
93e1  
e1fa  
95cf  
5771  
95cf  
427e

```

b7b3     lit split(int pos) {
2273         for (lit it = blk.begin(); ; it++) {
5502             if (pos == 0) return it;
d480             while (it->data->size() > pos) {
2099                 blk.insert(next(it), it->split(pos));
95cf             }
a1c8             pos -= it->data->size();
95cf         }
95cf     }
427e
fd38     LL sum(int l, int r) { // traverse
48b4         lit it1 = split(l), it2 = split(r);
ac09         LL res = 0;
9f1d         while (it1 != it2) {
8284             res += it1->sum;
61fd             it1++;
95cf         }
b204         maintain();
244d         return res;
95cf     }
329b };

```

## 6.10 Sparse table, range minimum query

The array is 0-based and the range is left-closed right-open.

```

db63     const int MAXN = 100007;
cefd     int a[MAXN], st[MAXN][30];
427e
d34f     void init(int n){
c73d         int l = log2(n);
cf75         rep (i, n) st[i][0] = a[i];
426b         rep (j, l) rep (i, 1+n-(1<<j))
1131             st[i][j+1] = min(st[i][j], st[i+(1<<j)][j]);
95cf     }
427e
c863     int rmq(int l, int r){
f089         int k = log2(r - l);
6117         return min(st[l][k], st[r-(1<<k)][k]);
95cf     }

```

## 7 Geometrics

### 7.1 2D geometric template

```

#include <bits/stdc++.h>
using namespace std;

typedef int T;
typedef struct pt {
    T x, y;
    T operator , (pt a) { return x*a.x + y*a.y; } // inner product
    T operator * (pt a) { return x*a.y - y*a.x; } // outer product
    pt operator + (pt a) { return {x+a.x, y+a.y}; }
    pt operator - (pt a) { return {x-a.x, y-a.y}; }

    pt operator * (T k) { return {x*k, y*k}; }
    pt operator - () { return {-x, -y}; }
} vec;

typedef pair<pt, pt> seg;

bool ptOnSeg(pt& p, seg& s){
    vec v1 = s.first - p, v2 = s.second - p;
    return (v1, v2) <= 0 && v1 * v2 == 0;
}

// 0 not on segment
// 1 on segment except vertices
// 2 on vertices
int ptOnSeg2(pt& p, seg& s){
    vec v1 = s.first - p, v2 = s.second - p;
    T ip = (v1, v2);
    if (v1 * v2 != 0 || ip > 0) return 0;
    return (v1, v2) ? 1 : 2;
}

// if two orthogonal rectangles do not touch, return true
inline bool nIntRectRect(seg a, seg b){
    return min(a.first.x, a.second.x) > max(b.first.x, b.second.x) ||
           min(a.first.y, a.second.y) > max(b.first.y, b.second.y) ||
           min(b.first.x, b.second.x) > max(a.first.x, a.second.x) ||
           min(b.first.y, b.second.y) > max(a.first.y, a.second.y);
}

```

302f  
421c  
427e  
4553  
c0ae  
7a9d  
ffaa  
3ec7  
221a  
8b34  
427e  
368b  
90f4  
ba8c  
427e  
0ea6  
427e  
8d6e  
ce77  
de97  
95cf  
427e  
427e  
427e  
427e  
8421  
ce77  
70ca  
8b14  
0847  
95cf  
427e  
427e  
72bb  
f9ac  
f486  
39ce  
80c7  
95cf

```

427e // >0 in order
427e // <0 out of order
427e // =0 not standard
7538 inline double rotOrder(vec a, vec b, vec c){return double(a*b)*(b*c);}
427e
31ed inline bool intersect(seg a, seg b){
427e     // ! if (nIntRectRect(a, b)) return false; // if commented, assume that a
        and b are non-collinear
cb52     return rotOrder(b.first-a.first, a.second-a.first, b.second-a.first) >= 0 &&
059e         rotOrder(a.first-b.first, b.second-b.first, a.second-b.first) >= 0;
95cf }
427e
427e // 0 not intersect
427e // 1 standard intersection
427e // 2 vertex-line intersection
427e // 3 vertex-vertex intersection
427e // 4 collinear and have common point(s)
4d19 int intersect2(seg& a, seg& b){
5dc4     if (nIntRectRect(a, b)) return 0;
42c0     vec va = a.second - a.first, vb = b.second - b.first;
2096     double j1 = rotOrder(b.first-a.first, va, b.second-a.first),
72fe         j2 = rotOrder(a.first-b.first, vb, a.second-b.first);
5ac6     if (j1 < 0 || j2 < 0) return 0;
9400     if (j1 != 0 && j2 != 0) return 1;
83db     if (j1 == 0 && j2 == 0){
6b0c         if (va * vb == 0) return 4; else return 3;
fb17     } else return 2;
95cf }
427e
2c68 template <typename Tp = T>
5894 inline pt getIntersection(pt P, vec v, pt Q, vec w){
6850     static_assert(is_same<Tp, double>::value, "must_be_double!");
7c9a     return P + v * (w*(P-Q)/(v*w));
95cf }
427e
427e // -1 outside the polygon
427e // 0 on the border of the polygon
427e // 1 inside the polygon
cbdd int ptOnPoly(pt p, pt* poly, int n){
5fb4     int wn = 0;
1294     for (int i = 0; i < n; i++) {
427e
3cae         T k, d1 = poly[i].y - p.y, d2 = poly[(i+1)%n].y - p.y;
```

```

        if (k = (poly[(i+1)%n] - poly[i])*(p - poly[i])){
            if (k > 0 && d1 <= 0 && d2 > 0) wn++;
            if (k < 0 && d2 <= 0 && d1 > 0) wn--;
        } else return 0;
    }
    return wn ? 1 : -1;
}

istream& operator >> (istream& lhs, pt& rhs){
    lhs >> rhs.x >> rhs.y;
    return lhs;
}

istream& operator >> (istream& lhs, seg& rhs){
    lhs >> rhs.first >> rhs.second;
    return lhs;
}
}
```

b957  
8c40  
3c4d  
aad3  
95cf  
0a5f  
95cf  
427e  
d4a3  
fa86  
331a  
95cf  
427e  
07ae  
5cab  
331a  
95cf

## 8 Appendices

### 8.1 Primes

#### 8.1.1 First primes

$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$
2	1	3	2	5	2	7	3	11	2
13	2	17	3	19	2	23	5	29	2
31	3	37	2	41	6	43	3	47	5
53	2	59	2	61	2	67	2	71	7
73	5	79	3	83	2	89	3	97	5
101	2	103	5	107	2	109	6	113	3
127	3	131	2	137	3	139	2	149	2
151	6	157	5	163	2	167	5	173	2
179	2	181	2	191	19	193	5	197	2
199	3	211	2	223	3	227	2	229	6

#### 8.1.2 Arbitrary length primes

$\lg p$	$p$	$g(p)$	$p$	$g(p)$
3	967	5	1031	14
4	9859	2	10273	10
5	96331	10	102931	3
6	958543	6	1031137	5
7	9594539	2	10169651	2
8	96243449	3	103211039	7
9	980483981	2	1042484357	2
10	9858935453	2	10261276009	7
11	95748666809	3	101759940101	2
12	950781833849	3	1012797784423	5
13	9739822952371	7	10037217092377	7
14	96181051140397	5	104974966380359	11
15	981030138360889	13	1029038416465403	2
16	9655206098080843	3	10116299875820773	2
17	97687777921994419	3	101506415998163437	2

#### 8.1.3 $\sim 1 \times 10^9$

$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$
954854573	3	967607731	2	973215833	3
975831713	3	978949117	2	980766497	3
983879921	3	985918807	3	986608921	29
991136977	5	991752599	13	997137961	11
1003911991	3	1009775293	2	1012423549	6
1021000537	5	1023976897	7	1024153643	2
1037027287	3	1038812881	11	1044754639	3
1045125617	3	1047411427	3	1047753349	6

#### 8.1.4 $\sim 1 \times 10^{18}$

$p$	$g(p)$	$p$	$g(p)$
951970612352230049	3	963284339889659609	3
967495386904694119	3	969751761517096213	2
983238274281901499	2	984647442475101409	23
989286107138674069	11	1002507954383424641	3
1006658951440146419	2	1020152326159075903	3
1034876265966119449	7	1042753851435034019	2
1043609016597371563	2	1045571042176595707	2
1048364250160580293	2	1049495624119026949	2

### 8.2 Pell's equation

$x^2 - ny^2 = 1$ , where  $n$  is a positive nonsquare integer.

Let  $(x_0, y_0)$  be the smallest positive solution of the equation, then the  $k$ -th solution is:

$$\begin{pmatrix} x_k \\ y_k \end{pmatrix} = \begin{pmatrix} x_0 & ny_0 \\ y_0 & x_0 \end{pmatrix}^k \begin{pmatrix} x_0 \\ y_0 \end{pmatrix}$$

Some smallest solutions to Pell's equation:

$n$	2	3	5	6	7	8	10	11	12	13	14	15	17	18	19	20
$x$	3	2	9	5	8	3	19	10	7	649	15	4	33	17	170	9
$y$	2	1	4	2	3	1	6	3	2	180	4	1	8	4	39	2

### 8.3 Burnside's lemma and Polya's enumeration theorem

The Burnside's lemma says that

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

where  $G$  is a group acting on  $X$ ,  $X^g$  is the set of elements in  $X$  that are fixed by  $g$ , i.e.  $X^g = \{x \in X : gx = x\}$ .

The unweighted version of Pólya enumeration theorem says that

$$|Y^X/G| = \frac{1}{|G|} \sum_{g \in G} m^{c_g}$$

where  $m = |X|$  is the number of colors,  $c_g$  is the number of the cycles of permutation  $g$ .

### 8.4 Lagrange's interpolation

For sample points  $(x_0, y_0), \dots, (x_k, y_k)$ , define

$$l_j(x) = \prod_{0 \leq m \leq k, m \neq j} \frac{x - x_m}{x_j - x_m}$$

then the Lagrange polynomial is

$$L(x) = \sum_{j=0}^k y_j l_j(x).$$

To use the script below, type two lines

```
x0 x1 x2 ... xn
y0 y1 y2 ... yn
```

the script will print the fractional coefficient of the polynomial in ascending exponent order.

```
#!/usr/bin/python2
from fractions import *

def polymul(a, b) :
    p = [0] * (len(a)+len(b)-1)
    for e1, c1 in enumerate(a) :
        for e2, c2 in enumerate(b) :
            p[e1+e2] += c1*c2
    return p

x, y = [map(Fraction, raw_input().split()) for _ in 0,0]
n = len(x)
lj = [reduce(polymul, [[-x[m]/(x[j]-x[m]), 1/(x[j]-x[m])]
    for m in range(n) if m != j]]) for j in range(n)]
print ' '.join(map(str, map(sum, zip(*map(
    lambda a, b : [x*a for x in b], y, lj)))))
```

6dc9  
4b2b  
427e  
796b  
83e4  
f697  
156c  
dfce  
5849  
427e  
f06d  
e80a  
a649  
9dfa  
3cae  
7c0d