

南京大学 ACM-ICPC 集训队代码模版库



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1 General

1.1 Code library checksum

```
ab14 #!/usr/bin/python3
c502 import re, sys, hashlib
427e
f7db for line in sys.stdin.read().strip().split("\n") :
ddf5     print(hashlib.md5(re.sub(r'\s|//[.]*', '', line).encode('utf8')).hexdigest()
        [-4:], line)
```

1.2 Makefile

```
dab2 .PHONY : run
427e
207e $(t) : $(t).cpp
2d16     g++ --std=c++14 -Wall -D__LOCAL_DEBUG__ -fsanitize=undefined -fsanitize=
        address -ggdb -pipe -o $@ $<
427e
5f25 run : $(t)
bf3e     ./$$(t) < $(t).in
```

1.3 .vimrc

```
914c set nocompatible
733d syntax on
6bbc colorscheme slate
7db5 set number
b0e3 set cursorline
061b set shiftwidth=2
8011 set softtabstop=2
a66d set tabstop=2
d23a set expandtab
5245 set magic
740c set smartindent
bee8 set backspace=indent,eol,start
815d set cmdheight=1
0a40 set laststatus=2
e458 set statusline=\ %<%F[%1*%M%*%n%R%H]%=\ %y\ %0{&fileformat}\ %&encoding}\ %c
        :%l/%L%\
```

```
set whichwrap=b,s,<,>[,]
```

1c67

1.4 Stack

```
const int STK_SZ = 2000000;
char STK[STK_SZ * sizeof(void*)];
void *STK_BAK;

#if defined(__i386__)
#define SP "%esp"
#elif defined(__x86_64__)
#define SP "%rsp"
#endif

int main() {
    asm volatile("movl SP, %0; movl 1, SP: =g(STK_BAK):g(STK+sizeof(STK)):");
    ;

    // main program

    asm volatile("movl %0, SP: =g(STK_BAK)");
    return 0;
}
```

bebe
effc
4e99
427e
7bc9
0894
ac7a
a9ea
1937
427e
3117
3750
427e
427e
427e
6856
7021
95cf

1.5 Template

```
#include <bits/stdc++.h>
using namespace std;

#ifdef __LOCAL_DEBUG__
# define _debug(fmt, ...) fprintf(stderr, "[%s] " fmt "\n", \
    __func__, ##__VA_ARGS__)
#else
# define _debug(...) ((void) 0)
#endif
#define rep(i, n) for (int i=0; i<(n); i++)
#define Rep(i, n) for (int i=1; i<=(n); i++)
#define range(x) begin(x), end(x)
typedef long long LL;
typedef unsigned long long ULL;
```

302f
421c
427e
426f
3341
611f
a8cb
e6b5
1937
0d6c
cfe3
3505
5cad
b773

2 Miscellaneous Algorithms

2.1 2-SAT

```

0f42 const int MAXN = 100005;
03a9 struct twoSAT{
5c83     int n;
8f72     vector<int> G[MAXN*2];
d060     bool mark[MAXN*2];
b42d     int S[MAXN*2], c;
427e
d34f     void init(int n){
b985         this->n = n;
f9ec         for (int i=0; i<n*2; i++) G[i].clear();
0609         memset(mark, 0, sizeof(mark));
95cf     }
427e
3bd5     bool dfs(int x){
bd70         if (mark[x^1]) return false;
c96a         if (mark[x]) return true;
fd23         mark[x] = true;
4bea         S[c++] = x;
1ce6         for (int i=0; i<G[x].size(); i++)
d942             if (!dfs(G[x][i])) return false;
3361         return true;
95cf     }
427e
5894     void add_clause(int x, bool xval, int y, bool yval){
6afe         x = x * 2 + xval;
e680         y = y * 2 + yval;
81cc         G[x^1].push_back(y);
6835         G[y^1].push_back(x);
95cf     }
427e
d0cb     bool solve() {
7c39         for (int i=0; i<n*2; i+=2){
e63f             if (!mark[i] && !mark[i+1]){
88fb                 c = 0;
f4b9                 if (!dfs(i)){
3f03                     while (c > 0) mark[S[--c]] = false;
86c5                     if (!dfs(i+1)) return false;
95cf                 }
95cf             }

```

```

    }
    return true;
}

inline bool value(unsigned i){return mark[2*i+1];}
};

```

95cf
3361
95cf
427e
5f0a
329b

2.2 Knuth's optimization

```

int n;
int dp[256][256], dc[256][256];

template <typename T>
void compute(T cost) {
    for (int i = 0; i <= n; i++) {
        dp[i][i] = 0;
        dc[i][i] = i;
    }
    rep (i, n) {
        dp[i][i+1] = 0;
        dc[i][i+1] = i;
    }
    for (int len = 2; len <= n; len++) {
        for (int i = 0; i + len <= n; i++) {
            int j = i + len;
            int lbnd = dc[i][j-1], rbnd = dc[i+1][j];
            dp[i][j] = INT_MAX / 2;
            int c = cost(i, j);
            for (int k = lbnd; k <= rbnd; k++) {
                int res = dp[i][k] + dp[k][j] + c;
                if (res < dp[i][j]) {
                    dp[i][j] = res;
                    dc[i][j] = k;
                }
            }
        }
    }
};

```

5c83
d77c
427e
b7ec
0bc7
0423
8f5e
9488
95cf
be8e
95b5
aa0f
95cf
ec08
88b8
d3da
9824
a24a
f933
90d2
9bd0
26b5
e6af
9c88
95cf
95cf
95cf
95cf
329b

2.3 Mo's algorithm

All intervals are closed on both sides. When running functions `enter()` and `leave()`, the global `l` and `r` has not changed yet.

Usage:

```
add_query(id, l, r)    Add id-th query [l, r].
run()                 Run Mo's algorithm.
init()                TODO. Initialize the range [l, r].
yield(id)             TODO. Yield answer for id-th query.
enter(o)              TODO. Add o-th element.
leave(o)              TODO. Remove o-th element.
```

```
5194 constexpr int BLOCK_SZ = 300;
427e
3ec4 struct query { int l, r, id; };
d26a vector<query> queries;
427e
1e30 void add_query(int id, int l, int r) {
54c9     queries.push_back(query{l, r, id});
95cf }
427e
9f6b int l, r;
427e
427e // ----- functions to implement -----
62b4 inline void init();
50e1 inline void yield(int id);
b20d inline void enter(int o);
13af inline void leave(int o);
427e
37f0 void run() {
ab0b     if (queries.empty()) return;
8508     sort(range(queries), [](query lhs, query rhs) {
c7f8         int lb = lhs.l / BLOCK_SZ, rb = rhs.l / BLOCK_SZ;
03e7         if (lb != rb) return lb < rb;
0780         return lhs.r < rhs.r;
b251     });
6196     l = queries[0].l;
9644     r = queries[0].r;
07e2     init();
5bc9     for (query q : queries) {
7bc7         while (l > q.l) enter(l - 1), l--;
d646         while (r < q.r) enter(r + 1), r++;
13f0         while (l < q.l) leave(l), l++;
e1c6         while (r > q.r) leave(r), r--;
```

```
        yield(q.id);
    }
}
```

```
82f5
95cf
95cf
```

3 String

3.1 Knuth-Morris-Pratt algorithm

```
const int SIZE = 10005;

struct kmp_matcher {
    char p[SIZE];
    int fail[SIZE];
    int len;

    void construct(const char* needle) {
        len = strlen(p);
        strcpy(p, needle);
        fail[0] = fail[1] = 0;
        for (int i = 1; i < len; i++) {
            int j = fail[i];
            while (j && p[i] != p[j]) j = fail[j];
            fail[i + 1] = p[i] == p[j] ? j + 1 : 0;
        }
    }

    inline void found(int pos) {
        // ! add codes for having found at pos
    }

    void match(const char* haystack) { // must be called after construct
        const char* t = haystack;
        int n = strlen(t);
        int j = 0;
        rep(i, n) {
            while (j && p[j] != t[i]) j = fail[j];
            if (p[j] == t[i]) j++;
            if (j == len) found(i - len + 1);
        }
    }
};
```

```
2836
427e
d02b
2d81
9847
57b7
427e
60cf
aaa1
3a87
3dd4
d8a8
147f
3c79
4643
95cf
95cf
427e
c464
427e
95cf
427e
2daf
700f
8482
8fd0
be8e
4e19
b5d5
f024
95cf
95cf
329b
```

3.2 Manacher algorithm

```

81d4 struct Manacher {
cd09     int Len;
9255     vector<int> lc;
b301     string s;
427e
ec07     void work() {
c033         lc[1] = 1;
6bef         int k = 1;
427e
491f         for (int i = 2; i <= Len; i++) {
7957             int p = k + lc[k] - 1;
5e04             if (i <= p) {
24a1                 lc[i] = min(lc[2 * k - i], p - i + 1);
8e2e             } else {
e0e5                 lc[i] = 1;
95cf             }
74ff             while (s[i + lc[i]] == s[i - lc[i]]) lc[i]++;
2b9a             if (i + lc[i] > k + lc[k]) k = i;
95cf         }
95cf     }
427e
bfd5     void init(const char *tt) {
aaaf         int len = strlen(tt);
f701         s.resize(len * 2 + 10);
7045         lc.resize(len * 2 + 10);
8e13         s[0] = '*';
ae54         s[1] = '#';
1321         for (int i = 0; i < len; i++) {
e995             s[i * 2 + 2] = tt[i];
69fd             s[i * 2 + 1] = '#';
95cf         }
43fd         s[len * 2 + 1] = '#';
75d1         s[len * 2 + 2] = '\0';
61f7         Len = len * 2 + 2;
3e7a         work();
95cf     }
427e
b194     pair<int, int> maxpal(int l, int r) {
901a         int center = l + r + 1;
ffb2         int rad = lc[center] / 2;
ab54         int rmid = (l + r + 1) / 2;

```

```

int rl = rmid - rad, rr = rmid + rad - 1;
if ((r ^ 1) & 1) {
} else rr++;
return {max(l, rl), min(r, rr)};
}
};

```

```

17e4
3908
69f3
69dc
95cf
329b

```

3.3 Aho-corasick automaton

```

struct AC : Trie {
    int fail[MAXN];
    int last[MAXN];

    void construct() {
        queue<int> q;
        fail[0] = 0;
        rep(c, CHARN) {
            if (int u = tr[0][c]) {
                fail[u] = 0;
                q.push(u);
                last[u] = 0;
            }
        }
        while (!q.empty()) {
            int r = q.front();
            q.pop();
            rep(c, CHARN) {
                int u = tr[r][c];
                if (!u) {
                    tr[r][c] = tr[fail[r]][c];
                    continue;
                }
                q.push(u);
                int v = fail[r];
                while (v && !tr[v][c]) v = fail[v];
                fail[u] = tr[v][c];
                last[u] = tag[fail[u]] ? fail[u] : last[fail[u]];
            }
        }
    }

    void found(int pos, int j) {

```

```

a1ad
9143
daca
427e
8690
93d2
a7a6
ce3c
b1c6
a506
3e14
f689
95cf
95cf
cc78
31f0
15dd
ce3c
ab59
0ef5
9d58
b333
95cf
3e14
b3ff
d2ea
c275
654c
95cf
95cf
95cf
427e
7752

```

```

043e     if (j) {
427e         // ! add codes for having found word with tag[j]
4a96         found(pos, last[j]);
95cf     }
95cf }
427e
9785 void find(const char* text) { // must be called after construct()
80a4     int p = 0, c, len = strlen(text);
9c94     rep(i, len) {
b3db         c = id(text[i]);
f119         p = tr[p][c];
f08e         if (tag[p])
389b             found(i, p);
1e67         else if (last[p])
299e             found(i, last[p]);
95cf     }
95cf }
329b };

```

3.4 Suffix array

The character immediately after the end of the string **MUST** be set to the **UNIQUE SMALLEST** element.

Usage:

s[]	the source string
sa[i]	the index of starting position of i -th suffix
rk[i]	the number of suffixes less than the suffix starting from i
h[i]	the longest common prefix between the i -th and $(i-1)$ -th lexicographically smallest suffixes
n	size of source string
m	size of character set

```

de09 void radix_sort(int x[], int y[], int sa[], int n, int m) {
ec00     static int cnt[1000005]; // size > max(n, m)
6066     fill(cnt, cnt + m, 0);
93b7     rep(i, n) cnt[x[y[i]]]++;
9154     partial_sum(cnt, cnt + m, cnt);
acac     for (int i = n - 1; i >= 0; i--) sa[--cnt[x[y[i]]]] = y[i];
95cf }
427e
c939 void suffix_array(int s[], int sa[], int rk[], int n, int m) {
a69a     static int y[1000005]; // size > n
7306     copy(s, s + n, rk);

```

```

iota(y, y + n, 0);
radix_sort(rk, y, sa, n, m);
for (int j = 1, p = 0; j <= n; j <= 1, m = p, p = 0) {
    for (int i = n - j; i < n; i++) y[p++] = i;
    rep(i, n) if (sa[i] >= j) y[p++] = sa[i] - j;
    radix_sort(rk, y, sa, n, m + 1);
    swap_ranges(rk, rk + n, y);
    rk[sa[0]] = p = 1;
    for (int i = 1; i < n; i++)
        rk[sa[i]] = ((y[sa[i]] == y[sa[i-1]] and y[sa[i]+j] == y[sa[i-1]+j])
            ? p : ++p);
    if (p == n) break;
}
rep(i, n) rk[sa[i]] = i;
}

void calc_height(int s[], int sa[], int rk[], int h[], int n) {
    int k = 0;
    h[0] = 0;
    rep(i, n) {
        k = max(k - 1, 0);
        if (rk[i]) while (s[i+k] == s[sa[rk[i]-1]+k]) ++k;
        h[rk[i]] = k;
    }
}

```

3.5 Trie

```

const int MAXN = 12000;
const int CHARN = 26;

inline int id(char c) { return c - 'a'; }

struct Trie {
    int n;
    int tr[MAXN][CHARN]; // Trie tree, 0 denotes fail
    int tag[MAXN];

    Trie() {
        memset(tr[0], 0, sizeof(tr[0]));
        tag[0] = 0;
        n = 1;
    }
}

```

```

95cf }
427e
427e // tag should not be 0
30b0 void add(const char* s, int t) {
d50a     int p = 0, c, len = strlen(s);
9c94     rep(i, len) {
3140         c = id(s[i]);
d6c8         if (!tr[p][c]) {
26dd             memset(tr[n], 0, sizeof(tr[n]));
2e5c             tag[n] = 0;
73bb             tr[p][c] = n++;
95cf         }
f119         p = tr[p][c];
95cf     }
35ef     tag[p] = t;
95cf }
427e
427e // returns 0 if not found
427e // AC automaton does not need this function
216c int search(const char* s) {
d50a     int p = 0, c, len = strlen(s);
9c94     rep(i, len) {
3140         c = id(s[i]);
f339         if (!tr[p][c]) return 0;
f119         p = tr[p][c];
95cf     }
840e     return tag[p];
95cf }
329b };

```

3.6 Rolling hash

PLEASE call `init_hash()` in `int main()`!

Usage:

`build(str)` Construct the hasher with given string.
`operator()(l, r)` Get hash value of substring $[l, r)$.

```

1e42 const LL mod = 1006658951440146419, g = 967;
9f60 const int MAXN = 200005;
0291 LL pg[MAXN];
427e
6832 inline LL mul(LL x, LL y) {
c919     return __int128_t(x) * y % mod;

```

```

}

void init_hash() { // must be called in `int main()`
    pg[0] = 1;
    for (int i = 1; i < MAXN; i++)
        pg[i] = pg[i - 1] * g % mod;
}

struct hasher {
    LL val[MAXN];

    void build(const char *str) { // assume lower-case letter only
        for (int i = 0; str[i]; i++)
            val[i+1] = (mul(val[i], g) + str[i]) % mod;
    }

    LL operator() (int l, int r) { // [l, r)
        return (val[r] - mul(val[l], pg[r - l]) + mod) % mod;
    }
} ha;

```

4 Math

4.1 Extended Euclidean algorithm and Chinese remainder theorem

```

void exgcd(LL a, LL b, LL &g, LL &x, LL &y) {
    if (!b) g = a, x = 1, y = 0;
    else {
        exgcd(b, a % b, g, y, x);
        y -= x * (a / b);
    }
}

LL crt(LL r[], LL p[], int n) {
    LL q = 1, ret = 0;
    rep(i, n) q *= p[i];
    rep(i, n) {
        LL m = q / p[i];
        LL d, x, y;
        exgcd(p[i], m, d, x, y);
        ret = (ret + y * m * r[i]) % q;
    }
}

```



```

95cf    }
2e47    return (q + ret) % q;
95cf    }

```

4.2 Matrix powermod

```

44b4    const int MAXN = 105;
92df    const LL modular = 1000000007;
5c83    int n; // order of matrices

427e    struct matrix{
3180        LL m[MAXN][MAXN];
427e
43c5        void operator *=(matrix& a){
e735            static LL t[MAXN][MAXN];
34d7            Rep (i, n){
4c11                Rep (j, n){
ee1e                    t[i][j] = 0;
c4a7                    Rep (k, n){
fcaf                        t[i][j] += (m[i][k] * a.m[k][j]) % modular;
199e                        t[i][j] %= modular;
95cf                    }
95cf                }
95cf            }
dad4        memcpy(m, t, sizeof(t));
95cf    }
329b    };
427e
63d8    matrix r;
3ec2    void m_powmod(matrix& b, LL e){
83f0        memset(r.m, 0, sizeof(r.m));
a7c3        Rep(i, n)
de64            r.m[i][i] = 1;
3e90        while (e){
5a0e            if (e & 1) r *= b;
35c5            b *= b;
16fc            e >>= 1;
95cf        }
95cf    }

```

4.3 Linear basis

```

const int MAXD = 30;
struct linearbasis {
    ULL b[MAXD] = {};

    bool insert(LL v) {
        for (int j = MAXD - 1; j >= 0; j--) {
            if (!(v & (1ll << j))) continue;
            if (b[j]) v ^= b[j]
            else {
                for (int k = 0; k < j; k++)
                    if (v & (1ll << k)) v ^= b[k];
                for (int k = j + 1; k < MAXD; k++)
                    if (b[k] & (1ll << j)) b[k] ^= v;
                b[j] = v;
                return true;
            }
        }
        return false;
    }
};

```

4.4 Gauss elimination over finite field

```

const LL p = 1000000007;

LL powmod(LL b, LL e) {
    LL r = 1;
    while (e) {
        if (e & 1) r = r * b % p;
        b = b * b % p;
        e >>= 1;
    }
    return r;
}

typedef vector<LL> VLL;
typedef vector<VLL> VVLL;

LL gauss(VVLL &a, VVLL &b) {
    const int n = a.size(), m = b[0].size();
    vector<int> irow(n), icol(n), ipiv(n);

```

```

2976 LL det = 1;
427e
be8e rep (i, n) {
d2b5     int pj = -1, pk = -1;
6b4a     rep (j, n) if (!ipiv[j])
e582         rep (k, n) if (!ipiv[k])
6112             if (pj == -1 || a[j][k] > a[pj][pk]) {
a905                 pj = j;
657b                 pk = k;
95cf             }
d480     if (a[pj][pk] == 0) return 0;
0305     ipiv[pk]++;
8dad     swap(a[pj], a[pk]);
aad8     swap(b[pj], b[pk]);
be4d     if (pj != pk) det = (p - det) % p;
d080     irow[i] = pj;
f156     icol[i] = pk;
427e
4ecd     LL c = powmod(a[pk][pk], p - 2);
865b     det = det * a[pk][pk] % p;
c36a     a[pk][pk] = 1;
dd36     rep (j, n) a[pk][j] = a[pk][j] * c % p;
1b23     rep (j, m) b[pk][j] = b[pk][j] * c % p;
f8f3     rep (j, n) if (j != pk) {
e97f         c = a[j][pk];
c449         a[j][pk] = 0;
820b         rep (k, n) a[j][k] = (a[j][k] + p - a[pk][k] * c % p) % p;
f039         rep (k, m) b[j][k] = (b[j][k] + p - b[pk][k] * c % p) % p;
95cf     }
95cf }
427e
37e1 for (int j = n - 1; j >= 0; j--) if (irow[j] != icol[j]) {
50dc     for (int k = 0; k < n; k++) swap(a[k][irow[j]], a[k][icol[j]]);
95cf }
f27f return det;
95cf }

```

4.5 Berlekamp-Massey algorithm

```

2b86 const LL MOD = 1000000007;
427e
391d LL inverse(LL b) {

```

```

LL e = MOD - 2, r = 1;
while (e) {
    if (e & 1) r = r * b % MOD;
    b = b * b % MOD;
    e >>= 1;
}
return r;
}

struct Poly {
    vector<int> a;

    Poly() { a.clear(); }

    Poly(vector<int> &a) : a(a) {}

    int length() const { return a.size(); }

    Poly move(int d) {
        vector<int> na(d, 0);
        na.insert(na.end(), a.begin(), a.end());
        return Poly(na);
    }

    int calc(vector<int> &d, int pos) {
        int ret = 0;
        for (int i = 0; i < (int)a.size(); ++i) {
            if ((ret += (long long)d[pos - i] * a[i] % MOD) >= MOD) {
                ret -= MOD;
            }
        }
        return ret;
    }

    Poly operator - (const Poly &b) {
        vector<int> na(max(this->length(), b.length()));
        for (int i = 0; i < (int)na.size(); ++i) {
            int aa = i < this->length() ? this->a[i] : 0;
            int bb = i < b.length() ? b.a[i] : 0;
            na[i] = (aa + MOD - bb) % MOD;
        }
        return Poly(na);
    }
};

```

```

32d3
3e90
9a62
29ea
16fc
95cf
547e
95cf
427e
32a6
afe0
427e
9794
427e
de81
427e
8087
427e
16de
b31d
f915
cecf
95cf
427e
fa1a
5b57
501c
5de5
3041
95cf
95cf
ee0f
95cf
427e
c856
bd55
d1a7
3507
2bee
9526
95cf
cecf
95cf
329b

```

```

427e Poly operator * (const int &c, const Poly &p) {
5473     vector<int> na(p.length());
72de     for (int i = 0; i < (int)na.size(); ++i) {
b1a7         na[i] = (long long)c * p.a[i] % MOD;
b10c     }
95cf     return na;
aaab }
95cf }
427e
a1ff vector<int> solve(vector<int> a) {
9f23     int n = a.size();
58d0     Poly s, b;
4e8f     s.a.push_back(1), b.a.push_back(1);
c2aa     for (int i = 1, j = 0, ld = a[0]; i < n; ++i) {
4158         int d = s.calc(a, i);
d503         if (d) {
c29d             if ((s.length() - 1) * 2 <= i) {
db9d                 Poly ob = b;
6bce                 b = s;
1d0e                 s = s - (long long)d * inverse(ld) % MOD * ob.move(i - j);
0889                 j = i;
64f1                 ld = d;
8e2e             } else {
714e                 s = s - (long long)d * inverse(ld) % MOD * b.move(i - j);
95cf             }
95cf         }
95cf     }
427e     // Caution: s.a might be shorter than expected
e235     return s.a;
95cf }

```

4.6 Fast Walsh-Hadamard transform

```

061e void fwt(int* a, int n){
5595     for (int d = 1; d < n; d <= 1)
05f2         for (int i = 0; i < n; i += d < 1)
b833             rep (j, d){
7796                 int x = a[i+j], y = a[i+j+d];
427e                 // a[i+j] = x+y, a[i+j+d] = x-y; // xor
427e                 // a[i+j] = x+y; // and
427e                 // a[i+j+d] = x+y; // or
95cf             }

```

```

}

void ifwt(int* a, int n){
    for (int d = 1; d < n; d <= 1)
        for (int i = 0; i < n; i += d < 1)
            rep (j, d){
                int x = a[i+j], y = a[i+j+d];
                // a[i+j] = (x+y)/2, a[i+j+d] = (x-y)/2; // xor
                // a[i+j] = x-y; // and
                // a[i+j+d] = y-x; // or
            }
}

void conv(int* a, int* b, int n){
    fwt(a, n);
    fwt(b, n);
    rep(i, n) a[i] *= b[i];
    ifwt(a, n);
}

```

4.7 Fast fourier transform

```

const int NMAX = 1<<20;

typedef complex<double> cplx;

const double PI = 2*acos(0.0);

struct FFT{
    int rev[NMAX];
    cplx omega[NMAX], oinv[NMAX];
    int K, N;

    FFT(int k){
        K = k; N = 1 << k;
        rep (i, N){
            rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
            omega[i] = polar(1.0, 2.0 * PI / N * i);
            oinv[i] = conj(omega[i]);
        }
    }

    void dft(cplx* a, cplx* w){

```

```

a215     rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
ac6e     for (int l = 2; l <= N; l *= 2){
2969         int m = l/2;
b3cf         for (cplx* p = a; p != a + N; p += l)
c24f             rep (k, m){
fe06                 cplx t = w[N/l*k] * p[k+m];
ecbf                 p[k+m] = p[k] - t; p[k] += t;
95cf             }
95cf     }
427e
617b     void fft(cplx* a){dft(a, omega);}
a123     void ifft(cplx* a){
3b2f         dft(a, oinv);
57fc         rep (i, N) a[i] /= N;
95cf     }
427e
bdc0     void conv(cplx* a, cplx* b){
6497         fft(a); fft(b);
12a5         rep (i, N) a[i] *= b[i];
f84e         ifft(a);
95cf     }
329b };

```

4.8 Number theoretic transform

```

4ab9     const int NMAX = 1<<21;
427e
427e     // 998244353 = 7*17*2^23+1, G = 3
fb9a     const int P = 1004535809, G = 3; // = 479*2^21+1
427e
87ab     struct NTT{
c47c         int rev[NMAX];
0eda         LL omega[NMAX], oinv[NMAX];
81af         int g, g_inv; // g:  $g_n = G^{(P-1)/n}$ 
9827         int K, N;
427e
2a2c         LL powmod(LL b, LL e){
95a2             LL r = 1;
3e90             while (e){
6624                 if (e&1) r = r * b % P;
489e                 b = b * b % P;

```

```

         e >>= 1;
     }
     return r;
}

NTT(int k){
    K = k; N = 1 << k;
    g = powmod(G, (P-1)/N);
    g_inv = powmod(g, N-1);
    omega[0] = oinv[0] = 1;
    rep (i, N){
        rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
        if (i){
            omega[i] = omega[i-1] * g % P;
            oinv[i] = oinv[i-1] * g_inv % P;
        }
    }
}

void _ntt(LL* a, LL* w){
    rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
    for (int l = 2; l <= N; l *= 2){
        int m = l/2;
        for (LL* p = a; p != a + N; p += l)
            rep (k, m){
                LL t = w[N/l*k] * p[k+m] % P;
                p[k+m] = (p[k] - t + P) % P;
                p[k] = (p[k] + t) % P;
            }
    }
}

void ntt(LL* a){_ntt(a, omega);}
void intt(LL* a){
    LL inv = powmod(N, P-2);
    _ntt(a, oinv);
    rep (i, N) a[i] = a[i] * inv % P;
}

void conv(LL* a, LL* b){
    ntt(a); ntt(b);
    rep (i, N) a[i] = a[i] * b[i] % P;
    intt(a);
}

```

```

16fc
95cf
547e
95cf
427e
f420
e209
7652
4b3a
e04f
b393
7ba3
ad4f
8d8b
9e14
95cf
95cf
95cf
427e
9668
a215
ac6e
2969
7a1d
c24f
0ad3
6209
fa1b
95cf
95cf
95cf
427e
92ea
5daf
1f2a
9910
a873
95cf
427e
3a5b
ad16
e49e
5748
95cf

```

329b };

4.9 Sieve of Euler

```
cfc3 const int MAXX = 1e7+5;
5861 bool p[MAXX];
73ae int prime[MAXX], sz;
427e
9bc6 void sieve(){
9628     p[0] = p[1] = 1;
1ec8     for (int i = 2; i < MAXX; i++){
bf28         if (!p[i]) prime[sz++] = i;
e82c         for (int j = 0; j < sz && i*prime[j] < MAXX; j++){
b6a9             p[i*prime[j]] = 1;
5f51             if (i % prime[j] == 0) break;
95cf         }
95cf     }
95cf }
```

4.10 Sieve of Euler (General)

```
b62e namespace sieve {
6589     constexpr int MAXN = 10000007;
e982     bool p[MAXN]; // true if not prime
6ae8     int prime[MAXN], sz;
cbf7     int pval[MAXN], pcnt[MAXN];
6030     int f[MAXN];
427e
76f6     void exec(int N = MAXN) {
9628         p[0] = p[1] = 1;
427e
8a8a         pval[1] = 1;
bdda         pcnt[1] = 0;
c6b9         f[1] = 1;
427e
a643         for (int i = 2; i < N; i++) {
01d6             if (!p[i]) {
b2b2                 prime[sz++] = i;
37d9                 for (LL j = i; j < N; j *= i) {
758c                     int b = j / i;
81fd                     pval[j] = i * pval[b];
```

```
pcnt[j] = pcnt[b] + 1;
f[j] = _____; // f[j] = f(i^pcnt[j])
    }
}
for (int j = 0; i * prime[j] < N; j++) {
    int x = i * prime[j]; p[x] = 1;
    if (i % prime[j] == 0) {
        pval[x] = pval[i] * prime[j];
        pcnt[x] = pcnt[i] + 1;
    } else {
        pval[x] = prime[j];
        pcnt[x] = 1;
    }
    if (x != pval[x]) {
        f[x] = f[x / pval[x]] * f[pval[x]]
    }
    if (i % prime[j] == 0) break;
}
}
}
}
```

e0f3
a96c
95cf
95cf
34c0
f87a
20cc
9985
3f93
8e2e
cc91
6322
95cf
6191
d614
95cf
5f51
95cf
95cf
95cf
95cf

4.11 Miller-Rabin primality test

The array `a[]` (excluding sentinel, i.e. `LLONG_MAX`) should be

{2}	when $n < 2,047$.
{2, 7, 61}	when $n < 4,759,123,141 (2^{32})$.
{2, 3, 5, 7, 11}	when $n < 2.1 \times 10^{12}$.
{2, 325, 9375, 28178, 450775, 9780504, 1795265022}	when $n < 2^{64}$.

```
bool test(LL n){
    if (n < 3) return n==2;
    // ! The array a[] should be modified if the range of x changes.
    const LL a[] = {2LL, 7LL, 61LL, LLONG_MAX};
    LL r = 0, d = n-1, x;
    while (~d & 1) d >>= 1, r++;
    for (int i=0; a[i] < n; i++){
        x = powmod(a[i], d, n); // ! powmod must use for 64bit mulmod
        if (x == 1 || x == n-1) goto next;
        rep (i, r) {
            x = mulmod(x, x, n);
            if (x == n-1) goto next;
```

f16f
59f2
427e
3f11
c320
f410
2975
ece1
7f99
e257
d7ff
8d2e

```

95cf    }
438e    return false;
d490    next++;
95cf    }
3361    return true;
95cf    }

```

4.12 Pollard's rho algorithm

```

2e6b    ULL gcd(ULL a, ULL b) {return b ? gcd(b, a % b) : a;}
427e
54a5    ULL PollardRho(ULL n){
45eb        ULL c, x, y, d = n;
d3e5        if (~n&1) return 2;
3c69        while (d == n){
0964            x = y = 2;
4753            d = 1;
5952            c = rand() % (n - 1) + 1;
9e5b            while (d == 1){
33d5                x = (mulmod(x, x, n) + c) % n;
e1bf                y = (mulmod(y, y, n) + c) % n;
e1bf                y = (mulmod(y, y, n) + c) % n;
a313                d = gcd(x>y ? x-y : y-x, n);
95cf            }
95cf        }
5d89        return d;
95cf    }

```

5 Graph Theory

5.1 Strongly connected component

```

837c    const int MAXV = 100005;
427e
2ea0    struct graph{
88e3        vector<int> adj[MAXV];
9cad        stack<int> s;
3d02        int V; // number of vertices
8b6c        int pre[MAXV], lnk[MAXV], scc[MAXV];
27ee        int time, sccn;

```

```

void add_edge(int u, int v){
    adj[u].push_back(v);
}

void dfs(int u){
    pre[u] = lnk[u] = ++time;
    s.push(u);
    for (int v : adj[u]){
        if (!pre[v]){
            dfs(v);
            lnk[u] = min(lnk[u], lnk[v]);
        } else if (!scc[v]){
            lnk[u] = min(lnk[u], pre[v]);
        }
    }
    if (lnk[u] == pre[u]){
        sccn++;
        int x;
        do {
            x = s.top(); s.pop();
            scc[x] = sccn;
        } while (x != u);
    }
}

```

```

void find_scc(){
    time = sccn = 0;
    memset(scc, 0, sizeof scc);
    memset(pre, 0, sizeof pre);
    Rep (i, V){
        if (!pre[i]) dfs(i);
    }
}

```

```

vector<int> adjc[MAXV];
void contract(){
    Rep (i, V)
        rep (j, adj[i].size()){
            if (scc[i] != scc[adj[i][j]])
                adjc[scc[i]].push_back(scc[adj[i][j]]);
        }
}
};

```

```

427e
bfab
c71a
95cf
427e
d714
7e41
80f6
18f6
173e
5f3c
002c
6068
d5df
95cf
95cf
8de2
660f
3c9e
a69f
3834
b0e9
6757
95cf
95cf
427e
4c88
f4a2
8de7
8c2f
6901
56d1
95cf
95cf
427e
27ce
364d
1a1e
21a2
b730
b46e
95cf
95cf
329b

```

5.2 Vertex biconnected component

```

0f42 const int MAXN = 100005;
2ea0 struct graph {
33ae     int pre[MAXN], iscut[MAXN], bccno[MAXN], dfs_clock, bcc_cnt;
848f     vector<int> adj[MAXN], bcc[MAXN];
6b06     set<pair<int, int>> bcce[MAXN];
427e
76f7     stack<pair<int, int>> s;
427e
bfab     void add_edge(int u, int v) {
c71a         adj[u].push_back(v);
a717         adj[v].push_back(u);
95cf     }
427e
7d3c     int dfs(int u, int fa) {
9fe6         int lowu = pre[u] = ++dfs_clock;
ec14         int child = 0;
18f6         for (int v : adj[u]) {
173e             if (!pre[v]) {
e7f8                 s.push({u, v});
fdcf                 child++;
f851                 int lowv = dfs(v, u);
189c                 lowu = min(lowu, lowv);
b687                 if (lowv >= pre[u]) {
6323                     iscut[u] = 1;
57eb                     bcc[bcc_cnt].clear();
90b8                     bcce[bcc_cnt].clear();
a147                     while (1) {
a6a3                         int xu, xv;
a0c3                         tie(xu, xv) = s.top(); s.pop();
0ef5                         bcce[bcc_cnt].insert({min(xu, xv), max(xu, xv)});
3db2                         if (bccno[xu] != bcc_cnt) {
e0db                             bcc[bcc_cnt].push_back(xu);
d27f                             bccno[xu] = bcc_cnt;
95cf                         }
f357                         if (bccno[xv] != bcc_cnt) {
752b                             bcc[bcc_cnt].push_back(xv);
57c9                             bccno[xv] = bcc_cnt;
95cf                         }

```

```

        if (xu == u && xv == v) break;
    }
    bcc_cnt++;
}
} else if (pre[v] < pre[u] && v != fa) {
    s.push({u, v});
    lowu = min(lowu, pre[v]);
}
}
if (fa < 0 && child == 1) iscut[u] = 0;
return lowu;
}

void find_bcc(int n) {
    memset(pre, 0, sizeof pre);
    memset(iscut, 0, sizeof iscut);
    memset(bccno, -1, sizeof bccno);
    dfs_clock = bcc_cnt = 0;
    rep(i, n) if (!pre[i]) dfs(i, -1);
}
};

```

7096
95cf
03f5
95cf
7470
e7f8
f115
95cf
95cf
e104
1160
95cf
427e
17be
8c2f
e2d2
40d3
fae2
5c63
95cf
329b

5.3 Minimum spanning arborescence (Chu-Liu)

All vertices are 1-based.

Usage:

getans(n, root, edges) Compute the total size of MSA rooted at root.

Time Complexity: $O(|V||E|)$

```

struct edge {
    int u, v;
    LL w;
};

const int MAXN = 10005;
LL in[MAXN];
int pre[MAXN], vis[MAXN], id[MAXN];

LL getans(int n, int rt, vector<edge>& edges) {
    LL ans = 0;
    int cnt = 0;
    while (1) {

```

bcf8
54f1
309c
329b
427e
f5a4
7124
1c1d
427e
5a43
f7ff
8abb
a147

```

641a     Rep (i, n) in[i] = LLONG_MAX, id[i] = vis[i] = 0;
0705     for (auto e : edges) {
073a         if (e.u != e.v and e.w < in[e.v]) {
c1df             pre[e.v] = e.u;
5fbc             in[e.v] = e.w;
95cf         }
95cf     }
3fdb     in[rt] = 0;
34d7     Rep (i, n) {
3c97         if (in[i] == LLONG_MAX) return -1;
cf57         ans += in[i];
a763         int u;
4b0e         for (u = i; u != rt && vis[u] != i && !id[u]; u = pre[u])
88a2             vis[u] = i;
4b22         if (u != rt && !id[u]) {
b66e             id[u] = ++cnt;
0443             for (int v = pre[u]; v != u; v = pre[v])
5c22                 id[v] = cnt;
95cf         }
95cf     }
91e9     if (!cnt) return ans;
5e22     Rep (i, n) if (!id[i]) id[i] = ++cnt;
7400     for (auto& e : edges) {
7750         LL laz = in[e.v];
97ae         e.u = id[e.u];
fae6         e.v = id[e.v];
bdd2         if (e.u != e.v) e.w -= laz;
95cf     }
6cc4     n = cnt; rt = id[rt]; cnt = 0;
95cf }
95cf

```

5.4 Maximum flow (Dinic)

Usage:

add_edge(u, v, c) Add an edge from *u* to *v* with capacity *c*.

max_flow(s, t) Compute maximum flow from *s* to *t*.

Time Complexity: For general graph, $O(V^2E)$; for network with unit capacity, $O(\min\{V^{2/3}, \sqrt{E}\}E)$; for bipartite network, $O(\sqrt{VE})$.

```

bcf8 struct edge{
60e2     int from, to;
5e6d     LL cap, flow;

```

```

};

const int MAXN = 1005;
struct Dinic {
    int n, m, s, t;
    vector<edge> edges;
    vector<int> G[MAXN];
    bool vis[MAXN];
    int d[MAXN];
    int cur[MAXN];

    void add_edge(int from, int to, LL cap) {
        edges.push_back(edge{from, to, cap, 0});
        edges.push_back(edge{to, from, 0, 0});
        m = edges.size();
        G[from].push_back(m-2);
        G[to].push_back(m-1);
    }

    bool bfs() {
        memset(vis, 0, sizeof(vis));
        queue<int> q;
        q.push(s);
        vis[s] = 1;
        d[s] = 0;
        while (!q.empty()) {
            int x = q.front(); q.pop();
            for (int i = 0; i < G[x].size(); i++) {
                edge& e = edges[G[x][i]];
                if (!vis[e.to] && e.cap > e.flow) {
                    vis[e.to] = 1;
                    d[e.to] = d[x] + 1;
                    q.push(e.to);
                }
            }
        }
        return vis[t];
    }

    LL dfs(int x, LL a) {
        if (x == t || a == 0) return a;
        LL flow = 0, f;
        for (int& i = cur[x]; i < G[x].size(); i++) {
            edge& e = edges[G[x][i]];

```

```

329b
427e
e2cd
9062
4dbf
9f0c
b891
bbb6
b40a
ddec
427e
5973
7b55
1db7
fe77
dff5
8f2d
95cf
427e
1836
3b73
93d2
5d13
2cd2
721d
cc78
66ba
3b61
b510
bba9
cd72
cf26
ca93
95cf
95cf
95cf
b23b
95cf
427e
9252
6904
8bf9
f515
b510

```



```

2374         if(d[x] + 1 == d[e.to] && (f = dfs(e.to, min(a, e.cap-e.flow))) > 0)
        {
1cce             e.flow += f;
e16d             edges[G[x][i]^1].flow -= f;
a74d             flow += f;
23e5             a -= f;
97ed             if(a == 0) break;
95cf         }
95cf     }
84fb     return flow;
95cf }
427e
5bf2 LL max_flow(int s, int t) {
590d     this->s = s; this->t = t;
62e2     LL flow = 0;
ed58     while (bfs()) {
f326         memset(cur, 0, sizeof(cur));
fb3a         flow += dfs(s, LLONG_MAX);
95cf     }
84fb     return flow;
95cf }
427e
c72e vector<int> min_cut() { // call this after maxflow
1df9     vector<int> ans;
df9a     for (int i = 0; i < edges.size(); i++) {
56d8         edge& e = edges[i];
46a2         if(vis[e.from] && !vis[e.to] && e.cap > 0) ans.push_back(i);
95cf     }
4206     return ans;
95cf }
329b };

```

5.5 Maximum cardinality bipartite matching (Hungarian)

```

302f #include <bits/stdc++.h>
421c using namespace std;
427e
0d6c #define rep(i, n) for (int i = 0; i < (n); i++)
cfe3 #define Rep(i, n) for (int i = 1; i <= (n); i++)
8843 #define range(x) (x).begin(), (x).end()
5cad typedef long long LL;
427e

```

```

struct Hungarian{
    int nx, ny;
    vector<int> mx, my;
    vector<vector<int> > e;
    vector<bool> mark;

    void init(int nx, int ny){
        this->nx = nx;
        this->ny = ny;
        mx.resize(nx); my.resize(ny);
        e.clear(); e.resize(nx);
        mark.resize(nx);
    }

    inline void add(int a, int b){
        e[a].push_back(b);
    }

    bool augment(int i){
        if (!mark[i]) {
            mark[i] = true;
            for (int j : e[i]){
                if (my[j] == -1 || augment(my[j])){
                    mx[i] = j; my[j] = i;
                    return true;
                }
            }
        }
        return false;
    }

    int match(){
        int ret = 0;
        fill(range(mx), -1);
        fill(range(my), -1);
        rep (i, nx){
            fill(range(mark), false);
            if (augment(i)) ret++;
        }
        return ret;
    }
};

```

84ee
fbf6
9ec6
9d4c
edec
427e
8324
c1d1
f9c1
ac92
3f11
1023
95cf
427e
4589
486c
95cf
427e
0c2b
207c
dae4
6a1e
0892
9ca3
3361
95cf
95cf
95cf
438e
95cf
427e
3fac
5b57
b0f1
b957
4ed1
13a5
cc89
95cf
ee0f
95cf
329b

5.6 Maximum matching of general graph (Edmond's blossom)

Usage:

`init(n)` Initialize the template with n vertices, numbered from 1.
`add_edge(u, v)` Add an undirected edge uv .
`solve()` Find the maximum matching. Return the number of matched edges.
`mate[]` The mate of a matched vertex. If it is not matched, then the value is 0.

Time Complexity: $O(|V|^3)$, but extremely fast in practice.

```
c041 const int MAXN = 1024;
6ab1 struct Blossom {
0b32     vector<int> adj[MAXN];
93d2     queue<int> q;
5c83     int n; // set n to number of vertices before use
0de2     int label[MAXN], mate[MAXN], save[MAXN], used[MAXN];

427e
427e
2186     void init(int nv) {
6646         n = nv;
e962         Rep (i, n) adj[i].clear();
f7e2         memset(label, 0, sizeof label);
5f6a         memset(mate, 0, sizeof mate);
c4b9         memset(save, 0, sizeof save);
ee13         memset(used, 0, sizeof used);
95cf     }

427e
bfab     void add_edge(int u, int v) {
c71a         adj[u].push_back(v);
a717         adj[v].push_back(u);
95cf     }

427e
2a48     void rematch(int x, int y){
8af8         int m = mate[x]; mate[x] = y;
1aa4         if (mate[m] == x) {
f4ba             if (label[x] <= n) {
e7ce                 mate[m] = label[x];
bec9                 rematch(label[x], m);
8e2e             } else {
3341                 int a = 1 + (label[x] - n - 1) / n;
2885                 int b = 1 + (label[x] - n - 1) % n;
ef33                 rematch(a, b); rematch(b, a);
95cf             }
}
```

```
    }
}

void traverse(int x) {
    Rep (i, n) save[i] = mate[i];
    rematch(x, x);
    Rep (i, n) {
        if (mate[i] != save[i]) used[i]++;
        mate[i] = save[i];
    }
}

void relabel(int x, int y) {
    Rep (i, n) used[i] = 0;
    traverse(x); traverse(y);
    Rep (i, n) {
        if (used[i] == 1 and label[i] < 0) {
            label[i] = n + x + (y - 1) * n;
            q.push(i);
        }
    }
}

int solve() {
    Rep (i, n) {
        if (mate[i]) continue;
        Rep (j, n) label[j] = -1;
        label[i] = 0; q = queue<int>(); q.push(i);
        while (q.size()) {
            int x = q.front(); q.pop();
            for (int y : adj[x]) {
                if (mate[y] == 0 and i != y) {
                    mate[y] = x;
                    rematch(x, y);
                    q = queue<int>();
                    break;
                }
            }
            if (label[y] >= 0) {
                relabel(x, y);
                continue;
            }
        }
        if (label[mate[y]] < 0) {
            label[mate[y]] = x;
            q.push(mate[y]);
        }
    }
}
```

95cf
95cf
427e
8a50
43c0
2ef7
34d7
62c5
97ef
95cf
95cf
427e
8bf8
d101
c4ea
34d7
dee9
1c22
eb31
95cf
95cf
95cf
427e
a0ce
34d7
a073
1fc0
7676
1c7d
66ba
b98c
c07f
0593
2b14
8ea8
6173
95cf
9079
a72e
b333
95cf
58ec
9773
086d

```

95cf      }
95cf      }
95cf      }
95cf      }
8abb      int cnt = 0;
c816      Rep (i, n) if (mate[i] > i) cnt++;
6808      return cnt;
95cf      }
329b };
    
```

5.7 Minimum cost maximum flow

```

bcf8 struct edge{
60e2     int from, to;
d698     int cap, flow;
32cc     LL cost;
329b };
427e
cc3e const LL INF = LLONG_MAX / 2;
2aa8 const int MAXN = 5005;
c6cb struct MCMF {
9ceb     int s, t, n, m;
9f0c     vector<edge> edges;
b891     vector<int> G[MAXN];
f74f     bool inq[MAXN]; // queue
8f67     LL d[MAXN];     // distance
9524     int p[MAXN];     // previous
b330     int a[MAXN];    // improvement
427e
f7f2 void add_edge(int from, int to, int cap, LL cost) {
24f0     edges.push_back(edge{from, to, cap, 0, cost});
95f0     edges.push_back(edge{to, from, 0, 0, -cost});
fe77     m = edges.size();
dff5     G[from].push_back(m-2);
8f2d     G[to].push_back(m-1);
95cf }
427e
3c52 bool spfa(){
93d2     queue<int> q;
8494     fill(d, d + MAXN, INF); d[s] = 0;
fd48     memset(inq, 0, sizeof(inq));
5e7c     q.push(s); inq[s] = true;
    
```

```

p[s] = 0; a[s] = INT_MAX;
while (!q.empty()){
    int u = q.front(); q.pop(); inq[u] = false;
    for (int i : G[u]) {
        edge& e = edges[i];
        if (e.cap > e.flow && d[e.to] > d[u] + e.cost){
            d[e.to] = d[u] + e.cost;
            p[e.to] = G[u][i];
            a[e.to] = min(a[u], e.cap - e.flow);
            if (!inq[e.to]) q.push(e.to), inq[e.to] = true;
        }
    }
}
return d[t] != INF;
}

void augment(){
    int u = t;
    while (u != s){
        edges[p[u]].flow += a[t];
        edges[p[u]^1].flow -= a[t];
        u = edges[p[u]].from;
    }
}

#ifdef GIVEN_FLOW
bool min_cost(int s, int t, int f, LL& cost) {
    this->s = s; this->t = t;
    int flow = 0;
    cost = 0;
    while (spfa()) {
        augment();
        if (flow + a[t] >= f){
            cost += (f - flow) * d[t]; flow = f;
            return true;
        } else {
            flow += a[t]; cost += a[t] * d[t];
        }
    }
    return false;
}
#else
int min_cost(int s, int t, LL& cost) {
    this->s = s; this->t = t;
    
```

2dae
 cc78
 b0aa
 3bba
 56d8
 3601
 55bc
 0bea
 8249
 e5d3
 95cf
 95cf
 95cf
 6d7c
 95cf
 427e
 71a4
 06f1
 b19d
 db09
 25a9
 e6c9
 95cf
 95cf
 427e
 6e20
 5972
 590d
 21d4
 23cb
 22dc
 bcdb
 a671
 b14d
 3361
 8e2e
 2a83
 95cf
 95cf
 438e
 95cf
 a8cb
 f9a9
 590d

```

21d4     int flow = 0;
23cb     cost = 0;
22dc     while (spfa()) {
bcd8         augment();
2a83         flow += a[t]; cost += a[t] * d[t];
95cf     }
84fb     return flow;
95cf     }
1937 #endif
329b };
    
```

```

        bestc = c;
        bestw = wt[last];
    }
    } else {
        rep (j, n) wt[j] += w[last][j];
        added[last] = true;
    }
}
return {bestw, bestc};
}
    
```

bab6
 372e
 95cf
 8e2e
 caeb
 8b92
 95cf
 95cf
 95cf
 038c
 95cf

5.8 Global minimum cut (Stoer-Wagner)

Usage:

stoer(w) Compute the global minimum cut of the graph specified by the **symmetric** adjacent matrix w (0-based). Return the capacity of the cut and the indices of one part of the cut.

Time Complexity: $O(|V|^3)$

```

f9d7 typedef vector<LL> VI;
045e typedef vector<VI> WI;
427e
f012 pair<LL, VI> stoer(WI &w) {
66f7     int n = w.size();
4d98     VI used(n), c, bestc;
329d     LL bestw = -1;
427e
cd21     for (int ph = n - 1; ph >= 0; ph--) {
ec6e         VI wt = w[0], added = used;
f20e         int prev, last = 0;
4b32         rep (i, ph) {
8bfc             prev = last;
0706             last = -1;
4942             for (int j = 1; j < n; j++)
c4b9                 if (!added[j] && (last == -1 || wt[j] > wt[last]))
887d                     last = j;
71bc             if (i == ph - 1) {
9cfa                 rep (j, n) w[prev][j] += w[last][j];
1f25                 rep (j, n) w[j][prev] = w[prev][j];
5613                 used[last] = true;
8e11                 c.push_back(last);
bb8e                 if (bestw == -1 || wt[last] < bestw) {
    
```

5.9 Fast LCA

All indices of the tree are 1-based.

Usage:

preprocess(root) Initialize with tree rooted at root.
lca(u, v) Query the lowest common ancestor of u and v.

```

const int MAXN = 500005;
vector<int> adj[MAXN];
int id[MAXN], nid;
pair<int, int> st[MAXN << 1][33 - __builtin_clz(MAXN)];

void dfs(int u, int p, int d) {
    st[id[u] = nid++][0] = {d, u};
    for (int v : adj[u]) {
        if (v == p) continue;
        dfs(v, u, d + 1);
        st[nid++][0] = {d, u};
    }
}

void preprocess(int root) {
    nid = 0;
    dfs(root, 0, 1);
    int l = 31 - __builtin_clz(nid);
    rep (j, l) rep (i, 1+nid-(1<<j))
        st[i][j+1] = min(st[i][j], st[i+(1<<j)][j]);
}

int lca(int u, int v) {
    tie(u, v) = minmax(id[u], id[v]);
    
```

0e34
 0b32
 fccb
 1356
 427e
 e16d
 0df2
 18f6
 bd87
 f58c
 08ad
 95cf
 95cf
 427e
 3d1b
 3269
 91e1
 5e98
 213b
 1131
 95cf
 427e
 0f0b
 cfc4

```

be9b     int k = 31 - __builtin_clz(v-u+1);
8ebc     return min(st[u][k], st[v-(1<<k)+1][k]).second;
95cf }

```

5.10 Heavy-light decomposition

Time Complexity: The decomposition itself takes linear time. Each query takes $O(\log n)$ operations.

```

0f42 const int MAXN = 100005;
0b32 vector<int> adj[MAXN];
42f2 int sz[MAXN], top[MAXN], fa[MAXN], son[MAXN], depth[MAXN], id[MAXN];
427e
be5c void dfs1(int x, int dep, int par){
7489     depth[x] = dep;
2ee7     sz[x] = 1;
adb4     fa[x] = par;
b79d     int maxn = 0, s = 0;
c861     for (int c: adj[x]){
fe45         if (c == par) continue;
fd2f         dfs1(c, dep + 1, x);
b790         sz[x] += sz[c];
f0f1         if (sz[c] > maxn){
c749             maxn = sz[c];
fe19             s = c;
95cf         }
95cf     }
0e08     son[x] = s;
95cf }
427e
ba54 int cid = 0;
3644 void dfs2(int x, int t){
8d96     top[x] = t;
d314     id[x] = ++cid;
c4a1     if (son[x]) dfs2(son[x], t);
c861     for (int c: adj[x]){
9881         if (c == fa[x]) continue;
5518         if (c == son[x]) continue;
13f9         else dfs2(c, c);
95cf     }
95cf }
427e
0f04 void decomp(int root){

```

```

dfs1(root, 1, 0);
dfs2(root, root);
}

void query(int u, int v){
    while (top[u] != top[v]){
        if (depth[top[u]] < depth[top[v]]) swap(u, v);
        // id[top[u]] to id[u]
        u = fa[top[u]];
    }
    if (depth[u] > depth[v]) swap(u, v);
    // id[u] to id[v]
}

```

5.11 Centroid decomposition

Note that the centroid here is not the exact centroid of the graph. It only guarantees that the size of each subtree does not exceed half of that of the original tree. This is enough to guarantee the correct time complexity. All vertices are numbered from 1. Call `decomp(root)` to use.

Usage:

`decomp(u, p)` Decompose the tree rooted at u with parent p .

Time Complexity: The decomposition itself takes $O(n \log n)$ time.

```

vector<int> adj[100005];
int sz[100005], sum;

void getsz(int u, int p) {
    sz[u] = 1; sum++;
    for (int v : adj[u]) {
        if (v == p) continue;
        getsz(v, u);
        sz[u] += sz[v];
    }
}

int getcent(int u, int p) {
    for (int v : adj[u])
        if (v != p and sz[v] > sum / 2)
            return getcent(v, u);
    return u;
}

```

```

4662 void decompose(int u) {
618e     sum = 0; getsz(u, 0);
303c     u = getcent(u, 0); // update u to the centroid
427e
18f6     for (int v : adj[u]) {
427e         // get answer for subtree v
95cf     }
427e     // get answer for the whole tree
427e     // don't forget to count the centroid itself
427e
18f6     for (int v : adj[u]) { // divide and conquer
c375         adj[v].erase(find(range(adj[v]), u));
fa6b         decompose(v);
a717         adj[v].push_back(u); // restore deleted edge
95cf     }
95cf }

```

5.12 DSU on tree

This implementation avoids parallel existence of multiple data structures but requires that the data structure is invertible. To use this template, implement merge, enter, leave as needed; first call decomp(root, 0), then call work(root, 0, false). Labels of vertices start from 1.

Usage:

decomp(u, p) Decompose the tree *u*.
work(u, p, keep) Work for subtree *u*. When keep is set, information is not cleared.

Time Complexity: $O(n \log n)$ times the complexity for merge, enter, leave.

```

1fb6 vector<int> adj[100005];
901d int sz[100005], son[100005];
427e
5559 void decomp(int u, int p) {
50c0     sz[u] = 1;
18f6     for (int v : adj[u]) {
bd87         if (v == p) continue;
a851         decomp(v, u);
8449         sz[u] += sz[v];
d28c         if (sz[v] > sz[son[u]]) son[u] = v;
95cf     }
95cf }
427e

```

```

template <typename T>
void trav(T fn, int u, int p) {
    fn(u);
    for (int v : adj[u]) if (v != p) trav(fn, v, u);
}

#define for_light(v) for (int v : adj[u]) if (v != p and v != son[u])
void work(int u, int p, bool keep) {
    for_light(v) work(v, u, 0); // process light children

    // process heavy child
    // current data structure contains info of heavy child
    if (son[u]) work(son[u], u, 1);

    auto merge = [u] (int c) { /* count contribution of c */ };
    auto enter = [] (int c) { /* add vertex c */ };
    auto leave = [] (int c) { /* remove vertex c */ };

    for_light(v) {
        trav(merge, v, u);
        trav(enter, v, u);
    }

    // count answer for root and add it
    // Warning: special check may apply to root!
    merge(u);
    enter(u);

    // Leave current tree
    if (!keep) trav(leave, u, p);
}

```

b7ec
62f5
4412
30b3
95cf
427e
7467
33ff
72a2
427e
427e
427e
9866
427e
18a9
1ab0
f241
427e
3d3b
74c6
c13d
95cf
427e
427e
427e
c54f
9dec
427e
427e
4e3e
95cf

6 Data Structures

6.1 Fenwick tree (point update range query)

```

struct bit_purq { // point update, range query
    int N;
    vector<LL> tr;

    void init(int n) { // fill the array with 0

```

9976
d7af
99ff
427e
d34f

```

1010     tr.resize(N = n + 5);
95cf }
427e
63d0 LL sum(int n) {
f7ff     LL ans = 0;
e290     while (n) {
0715         ans += tr[n];
c0d4         n &= n - 1;
95cf     }
4206     return ans;
95cf }
427e
f4bd void add(int n, LL x){
ad20     while (n < N) {
6c81         tr[n] += x;
0af5         n += n & -n;
95cf     }
95cf }
329b };

```

6.2 Fenwick tree (range update point query)

```

3d03 struct bit_rupq{ // range update, point query
d7af     int N;
99ff     vector<LL> tr;
427e
d34f void init(int n) { // fill the array with 0
1010     tr.resize(N = n + 5);
95cf }
427e
38d4 LL query(int n) {
f7ff     LL ans = 0;
ad20     while (n < N) {
0715         ans += tr[n];
0af5         n += n & -n;
95cf     }
4206     return ans;
95cf }
427e
f4bd void add(int n, LL x) {
e290     while (n){
6c81         tr[n] += x;

```

```

        n &= n - 1;
    }
};

```

```

c0d4
95cf
95cf
329b

```

6.3 Segment tree

```

LL p;
const int MAXN = 4 * 100006;
struct segtree {
    int l[MAXN], m[MAXN], r[MAXN];
    LL val[MAXN], tadd[MAXN], tmul[MAXN];

#define lson (o<<1)
#define rson (o<<1|1)

    void pull(int o) {
        val[o] = (val[lson] + val[rson]) % p;
    }

    void push_add(int o, LL x) {
        val[o] = (val[o] + x * (r[o] - l[o])) % p;
        tadd[o] = (tadd[o] + x) % p;
    }

    void push_mul(int o, LL x) {
        val[o] = val[o] * x % p;
        tadd[o] = tadd[o] * x % p;
        tmul[o] = tmul[o] * x % p;
    }

    void push(int o) {
        if (l[o] == m[o]) return;
        if (tmul[o] != 1) {
            push_mul(lson, tmul[o]);
            push_mul(rson, tmul[o]);
            tmul[o] = 1;
        }
        if (tadd[o]) {
            push_add(lson, tadd[o]);
            push_add(rson, tadd[o]);
            tadd[o] = 0;

```

```

3942
1ebb
451a
27be
4510
427e
ac35
1294
427e
1344
bbe9
95cf
427e
e4bc
5dd6
6eff
95cf
427e
d658
b82c
aa86
649f
95cf
427e
b149
3159
0a90
0f4a
045e
ac0a
95cf
1b82
9547
0e73
6234

```

```

95cf     }
95cf     }
427e
471c void build(int o, int ll, int rr) {
0e87     int mm = (ll + rr) / 2;
9d27     l[o] = ll; r[o] = rr; m[o] = mm;
ac0a     tmul[o] = 1;
5c92     if (ll == mm) {
001f         scanf("%lld", val + o);
e5b6         val[o] %= p;
8e2e     } else {
7293         build(lson, ll, mm);
5e67         build(rson, mm, rr);
ba26         pull(o);
95cf     }
95cf }
427e
4406 void add(int o, int ll, int rr, LL x) {
3c16     if (ll <= l[o] && r[o] <= rr) {
db32         push_add(o, x);
8e2e     } else {
c4b0         push(o);
4305         if (m[o] > ll) add(lson, ll, rr, x);
d5a6         if (m[o] < rr) add(rson, ll, rr, x);
ba26         pull(o);
95cf     }
95cf }
427e
48cd void mul(int o, int ll, int rr, LL x) {
3c16     if (ll <= l[o] && r[o] <= rr) {
e7d0         push_mul(o, x);
8e2e     } else {
c4b0         push(o);
d1ba         if (ll < m[o]) mul(lson, ll, rr, x);
67f3         if (m[o] < rr) mul(rson, ll, rr, x);
ba26         pull(o);
95cf     }
95cf }
427e
0f62 LL query(int o, int ll, int rr) {
3c16     if (ll <= l[o] && r[o] <= rr) {
6dfe         return val[o];
8e2e     } else {
c4b0         push(o);

```

```

        if (rr <= m[o]) return query(lson, ll, rr);
        if (ll >= m[o]) return query(rson, ll, rr);
        return query(lson, ll, rr) + query(rson, ll, rr);
    }
}
} seg;

```

462a
5cca
bbf9
95cf
95cf
4d99

6.4 Link/cut tree

Usage:

pull(x)	Collect information of subtrees.
Link(u, v)	Link two unconnected trees.
Cut(u, v)	Cut an existent edge.
Query(u, v)	Path aggregation.
Update(u, x)	Single point modification.

// about 0.13s per 100k ops @Luogu.org

```

namespace LCT {
    const int MAXN = 300005;
    int fa[MAXN], ch[MAXN][2], val[MAXN], sum[MAXN];
    bool rev[MAXN];

    bool isroot(int x) {
        return ch[fa[x]][0] == x || ch[fa[x]][1] == x;
    }

    void pull(int x) {
        sum[x] = val[x] ^ sum[ch[x][0]] ^ sum[ch[x][1]];
    }

    void reverse(int x) {
        swap(ch[x][0], ch[x][1]);
        rev[x] ^= 1;
    }

    void push(int x) {
        if (rev[x]) {
            if (ch[x][0]) reverse(ch[x][0]);
            if (ch[x][1]) reverse(ch[x][1]);
            rev[x] = 0;
        }
    }
}

```

427e
427e
ed4d
5ece
6a6d
c6e1
427e
7839
45a9
95cf
427e
3bf9
6664
95cf
427e
3698
7850
52c6
95cf
427e
1a53
8f1f
ebf3
6eb0
8fc1
95cf
95cf


```

427e void rotate(int x) {
425f     int y = fa[x], z = fa[y], k = ch[y][1] == x, w = ch[x][!k];
51af     if (isroot(y)) ch[z][ch[z][1] == y] = x;
e1fe     ch[x][!k] = y; ch[y][k] = w;
af46     if (w) fa[w] = y;
fa6f     fa[y] = x; fa[x] = z;
3540     pull(y);
72ef }
95cf
427e void pushall(int x) {
bc1b     if (isroot(x)) pushall(fa[x]);
a316     push(x);
a97b }
95cf
427e void splay(int x) {
f69c     int y = x, z = 0;
d095     pushall(y);
8ab3     while (isroot(x)) {
f244         y = fa[x]; z = fa[y];
ceef         if (isroot(y)) rotate((ch[y][0] == x) ^ (ch[z][0] == y) ? x : y);
4449         rotate(x);
cf90     }
95cf     pull(x);
78a0 }
95cf
427e void access(int x) {
6229     int z = x;
1548     for (int y = 0; x; x = fa[y = x]) {
ba78         splay(x);
8fec         ch[x][1] = y;
b05d         pull(x);
78a0     }
95cf     splay(z);
7afd }
95cf
427e void chroot(int x) {
502e     access(x);
766a     reverse(x);
cb0d }
95cf
427e void split(int x, int y) {
471a     chroot(x);
3015     access(y);
29b5

```

```

}

int Root(int x) {
    access(x);
    while (ch[x][0]) {
        push(x);
        x = ch[x][0];
    }
    splay(x);
    return x;
}

void Link(int u, int v) { // assume unconnected before
    chroot(u);
    fa[u] = v;
}

void Cut(int u, int v) { // assume connected before
    split(u, v);
    fa[u] = ch[v][0] = 0;
    pull(v);
}

int Query(int u, int v) {
    split(u, v);
    return sum[v];
}

void Update(int u, int x) {
    splay(u);
    val[u] = x;
}
};

```

```

95cf
427e
d87a
766a
874d
a97b
b83a
95cf
8fec
d074
95cf
427e
70d3
b8a5
2448
95cf
427e
c2f4
e8ce
fd95
743b
95cf
427e
6ca2
e8ce
a5ba
95cf
427e
eaba
46ce
1d62
95cf
329b

```

6.5 Balanced binary search tree from pb_ds

```

#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;

tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update>
    rkt;
// null_tree_node_update

```

```

0475
332d
427e
43a7
427e

```

```

427e // SAMPLE USAGE
190e rkt.insert(x);      // insert element
05d4 rkt.erase(x);      // erase element
add5 rkt.order_of_key(x); // obtain the number of elements less than x
b064 rkt.find_by_order(i); // iterator to i-th (numbered from 0) smallest element
c103 rkt.lower_bound(x);
4ff4 rkt.upper_bound(x);
b19b rkt.join(rkt2);    // merge tree (only if their ranges do not intersect)
cb47 rkt.split(x, rkt2); // split all elements greater than x to rkt2

```

6.6 Persistent segment tree, range k-th query

```

f1a7 struct node {
2ff6     static int n, pos;
427e
7cec     int value;
70e2     node *left, *right;
427e
20b0     void* operator new(size_t size);
427e
3dc0     static node* Build(int l, int r) {
b6c5         node* a = new node;
ce96         if (r > l + 1) {
181e             int mid = (l + r) / 2;
3ba2             a->left = Build(l, mid);
8aaf             a->right = Build(mid, r);
8e2e         } else {
bfc4             a->value = 0;
95cf         }
5ffd         return a;
95cf     }
427e
5a45     static node* init(int size) {
2c46         n = size;
7ee3         pos = 0;
be52         return Build(0, n);
95cf     }
427e
93c0     static int Query(node* lt, node *rt, int l, int r, int k) {
d30c         if (r == l + 1) return l;
181e         int mid = (l + r) / 2;

```

```

if (rt->left->value - lt->left->value < k) {
    k -= rt->left->value - lt->left->value;
    return Query(lt->right, rt->right, mid, r, k);
} else {
    return Query(lt->left, rt->left, l, mid, k);
}
}

```

```

static int query(node* lt, node *rt, int k) {
    return Query(lt, rt, 0, n, k);
}

```

```

node *Inc(int l, int r, int pos) const {
    node* a = new node(*this);
    if (r > l + 1) {
        int mid = (l + r) / 2;
        if (pos < mid)
            a->left = left->Inc(l, mid, pos);
        else
            a->right = right->Inc(mid, r, pos);
    }
    a->value++;
    return a;
}

```

```

node *inc(int index) {
    return Inc(0, n, index);
}
nodes[8000000];

```

```

int node::n, node::pos;
inline void* node::operator new(size_t size) {
    return nodes + (pos++);
}

```

6.7 Persistent block list

Block list that supports persistence. All indices are 0-based. `std::shared_ptr` is used to ease memory management. One should modify the constructor of `block` to maintain extra information. Here we use this policy that the size of each block does not exceed `BLOCK`, while the sum of sizes of two adjacent blocks does not less than `BLOCK`.

When some operation that breaks block list property, please call `maintain` in time to restore

cb5a
8edb
2412
8e2e
0119
95cf
95cf
427e
c9ad
9e27
95cf
427e
b19c
5794
ce96
181e
203d
f44a
649a
1024
95cf
2b3e
5ffd
95cf
427e
e80f
c246
95cf
865a
427e
99ce
1987
bb3c
95cf

the property.

Usage:

`maintain()` Maintain the block list property.
`split(pos)` Split the block list at position `pos`. Returns an iterator to a block starting at `pos`.
`sum(l, r)` An example function of list traversal between $[l, r)$.

Time Complexity: When `BLOCK` is properly selected, the time complexity is $O(\sqrt{n})$ per operation.

```
a19e constexpr int BLOCK = 800;
76b3 typedef vector<int> vi;
0563 typedef shared_ptr<vi> pvi;
013b typedef shared_ptr<const vi> pcvi;
427e
a771 struct block {
2989     pcvi data;
8fd0     LL sum;
427e
427e     // add information to maintain
a613     block(pcvi ptr) :
24b5         data(ptr),
0cf0         sum(accumulate(ptr->begin(), ptr->end(), 0ll))
e93b     { }
427e
5c0f     void merge(const block& another) {
0b18         pvi temp = make_shared<vi>(data->begin(), data->end());
ac21         temp->insert(temp->end(), another.data->begin(), another.data->end());
6467         *this = block(temp);
95cf     }
427e
42e8     block split(int pos) {
dac1         block result(make_shared<vi>(data->begin() + pos, data->end()));
01db         *this = block(make_shared<vi>(data->begin(), data->begin() + pos));
56b0         return result;
95cf     }
329b };
427e
2a18 typedef list<block>::iterator lit;
427e
ce14 struct blocklist {
5540     list<block> blk;
427e
7b8e     void maintain() {
```

```
lit it = blk.begin();
while (it != blk.end() and next(it) != blk.end()) {
    lit it2 = it;
    while (next(it2) != blk.end() and
           it2->data->size() + next(it2)->data->size() <= BLOCK) {
        it2->merge(*next(it2));
        blk.erase(next(it2));
    }
    ++it;
}

lit split(int pos) {
    for (lit it = blk.begin(); ; it++) {
        if (pos == 0) return it;
        while (it->data->size() > pos) {
            blk.insert(next(it), it->split(pos));
        }
        pos -= it->data->size();
    }
}

LL sum(int l, int r) { // traverse
    lit it1 = split(l), it2 = split(r);
    LL res = 0;
    while (it1 != it2) {
        res += it1->sum;
        it1++;
    }
    maintain();
    return res;
}

};
```

3131
5e44
852d
0b03
029f
93e1
e1fa
95cf
5771
95cf
95cf
427e
b7b3
2273
5502
d480
2099
95cf
a1c8
95cf
95cf
427e
fd38
48b4
ac09
9f1d
8284
61fd
95cf
b204
244d
95cf
329b

6.8 Sparse table, range extremum query

The array is 0-based and the range is closed.

```
const int MAXN = 100007;
int a[MAXN];
int st[MAXN][32 - __builtin_clz(MAXN)];

inline int ext(int x, int y){return x>y?x:y;} // ! max
```

db63
b330
69ae
427e
8041

```

427e void init(int n){
d34f     int l = 31 - __builtin_clz(n);
ce01     rep (i, n) st[i][0] = a[i];
cf75     rep (j, l)
b811         rep (i, 1+n-(1<<j))
6937             st[i][j+1] = ext(st[i][j], st[i+(1<<j)][j]);
082a }
95cf
427e
c863 int rmq(int l, int r){
92f5     int k = 31 - __builtin_clz(r-l+1);
baa2     return ext(st[l][k], st[r-(1<<k)+1][k]);
95cf }

```

7 Geometrics

7.1 2D geometric template

```

302f #include <bits/stdc++.h>
421c using namespace std;
427e
4553 typedef int T;
c0ae typedef struct pt {
7a9d     T x, y;
ffa9     T operator , (pt a) { return x*a.x + y*a.y; } // inner product
3ec7     T operator * (pt a) { return x*a.y - y*a.x; } // outer product
221a     pt operator + (pt a) { return {x+a.x, y+a.y}; }
8b34     pt operator - (pt a) { return {x-a.x, y-a.y}; }
427e
368b     pt operator * (T k) { return {x*k, y*k}; }
90f4     pt operator - () { return {-x, -y}; }
ba8c } vec;
427e
0ea6 typedef pair<pt, pt> seg;
427e
8d6e bool ptOnSeg(pt& p, seg& s){
ce77     vec v1 = s.first - p, v2 = s.second - p;
de97     return (v1, v2) <= 0 && v1 * v2 == 0;
95cf }
427e
427e // 0 not on segment

```

```

// 1 on segment except vertices
// 2 on vertices
int ptOnSeg2(pt& p, seg& s){
    vec v1 = s.first - p, v2 = s.second - p;
    T ip = (v1, v2);
    if (v1 * v2 != 0 || ip > 0) return 0;
    return (v1, v2) ? 1 : 2;
}

// if two orthogonal rectangles do not touch, return true
inline bool nIntRectRect(seg a, seg b){
    return min(a.first.x, a.second.x) > max(b.first.x, b.second.x) ||
           min(a.first.y, a.second.y) > max(b.first.y, b.second.y) ||
           min(b.first.x, b.second.x) > max(a.first.x, a.second.x) ||
           min(b.first.y, b.second.y) > max(a.first.y, a.second.y);
}

// >0 in order
// <0 out of order
// =0 not standard
inline double rotOrder(vec a, vec b, vec c){return double(a*b)*(b*c);}

inline bool intersect(seg a, seg b){
    // ! if (nIntRectRect(a, b)) return false; // if commented, assume that a
    // and b are non-collinear
    return rotOrder(b.first-a.first, a.second-a.first, b.second-a.first) >= 0 &&
           rotOrder(a.first-b.first, b.second-b.first, a.second-b.first) >= 0;
}

// 0 not intersect
// 1 standard intersection
// 2 vertex-line intersection
// 3 vertex-vertex intersection
// 4 collinear and have common point(s)
int intersect2(seg& a, seg& b){
    if (nIntRectRect(a, b)) return 0;
    vec va = a.second - a.first, vb = b.second - b.first;
    double j1 = rotOrder(b.first-a.first, va, b.second-a.first),
           j2 = rotOrder(a.first-b.first, vb, a.second-b.first);
    if (j1 < 0 || j2 < 0) return 0;
    if (j1 != 0 && j2 != 0) return 1;
    if (j1 == 0 && j2 == 0){
        if (va * vb == 0) return 4; else return 3;
    } else return 2;
}

```

```

95cf }
427e
2c68 template <typename Tp = T>
5894 inline pt getIntersection(pt P, vec v, pt Q, vec w){
6850     static_assert(is_same<Tp, double>::value, "must_be_double!");
7c9a     return P + v * (w*(P-Q)/(v*w));
95cf }
427e
427e // -1 outside the polygon
427e // 0 on the border of the polygon
427e // 1 inside the polygon
cbdd int ptOnPoly(pt p, pt* poly, int n){
5fb4     int wn = 0;
1294     for (int i = 0; i < n; i++) {
427e         T k, d1 = poly[i].y - p.y, d2 = poly[(i+1)%n].y - p.y;
b957         if (k = (poly[(i+1)%n] - poly[i])*(p - poly[i])){
8c40             if (k > 0 && d1 <= 0 && d2 > 0) wn++;
3c4d             if (k < 0 && d2 <= 0 && d1 > 0) wn--;
aad3         } else return 0;
95cf     }
0a5f     return wn ? 1 : -1;
95cf }
427e
d4a3 istream& operator >> (istream& lhs, pt& rhs){
fa86     lhs >> rhs.x >> rhs.y;
331a     return lhs;
95cf }
427e
07ae istream& operator >> (istream& lhs, seg& rhs){
5cab     lhs >> rhs.first >> rhs.second;
331a     return lhs;
95cf }

```

8 Appendices

8.1 Primes

8.1.1 First primes

p	$g(p)$	p	$g(p)$	p	$g(p)$	p	$g(p)$	p	$g(p)$
2	1	3	2	5	2	7	3	11	2
13	2	17	3	19	2	23	5	29	2
31	3	37	2	41	6	43	3	47	5
53	2	59	2	61	2	67	2	71	7
73	5	79	3	83	2	89	3	97	5
101	2	103	5	107	2	109	6	113	3
127	3	131	2	137	3	139	2	149	2
151	6	157	5	163	2	167	5	173	2
179	2	181	2	191	19	193	5	197	2
199	3	211	2	223	3	227	2	229	6

8.1.2 Arbitrary length primes

$\lg p$	p	$g(p)$	p	$g(p)$
3	967	5	1031	14
4	9859	2	10273	10
5	96331	10	102931	3
6	958543	6	1031137	5
7	9594539	2	10169651	2
8	96243449	3	103211039	7
9	980483981	2	1042484357	2
10	9858935453	2	10261276009	7
11	95748666809	3	101759940101	2
12	950781833849	3	1012797784423	5
13	9739822952371	7	10037217092377	7
14	96181051140397	5	104974966380359	11
15	981030138360889	13	1029038416465403	2
16	9655206098080843	3	10116299875820773	2
17	97687777921994419	3	101506415998163437	2

8.1.3 $\sim 1 \times 10^9$

p	$g(p)$	p	$g(p)$	p	$g(p)$
954854573	3	967607731	2	973215833	3
975831713	3	978949117	2	980766497	3
983879921	3	985918807	3	986608921	29
991136977	5	991752599	13	997137961	11
1003911991	3	1009775293	2	1012423549	6
1021000537	5	1023976897	7	1024153643	2
1037027287	3	1038812881	11	1044754639	3
1045125617	3	1047411427	3	1047753349	6

8.1.4 $\sim 1 \times 10^{18}$

p	$g(p)$	p	$g(p)$
951970612352230049	3	963284339889659609	3
967495386904694119	3	969751761517096213	2
983238274281901499	2	984647442475101409	23
989286107138674069	11	1002507954383424641	3
1006658951440146419	2	1020152326159075903	3
1034876265966119449	7	1042753851435034019	2
1043609016597371563	2	1045571042176595707	2
1048364250160580293	2	1049495624119026949	2

8.2 Pell's equation

$x^2 - ny^2 = 1$, where n is a positive nonsquare integer.

Let (x_0, y_0) be the smallest positive solution of the equation, then the k -th solution is:

$$\begin{pmatrix} x_k \\ y_k \end{pmatrix} = \begin{pmatrix} x_0 & ny_0 \\ y_0 & x_0 \end{pmatrix}^k \begin{pmatrix} x_0 \\ y_0 \end{pmatrix}$$

Some smallest solutions to Pell's equation:

n	2	3	5	6	7	8	10	11	12	13	14	15	17	18	19	20
x	3	2	9	5	8	3	19	10	7	649	15	4	33	17	170	9
y	2	1	4	2	3	1	6	3	2	180	4	1	8	4	39	2

8.3 Burnside's lemma and Polya's enumeration theorem

The Burnside's lemma says that

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

where G is a group acting on X , X^g is the set of elements in X that are fixed by g , i.e. $X^g = \{x \in X : gx = x\}$.

The unweighted version of Pólya enumeration theorem says that

$$|Y^X/G| = \frac{1}{|G|} \sum_{g \in G} m^{c_g}$$

where $m = |X|$ is the number of colors, c_g is the number of the cycles of permutation g .

8.4 Lagrange's interpolation

For sample points $(x_0, y_0), \dots, (x_k, y_k)$, define

$$l_j(x) = \prod_{0 \leq m \leq k, m \neq j} \frac{x - x_m}{x_j - x_m}$$

then the Lagrange polynomial is

$$L(x) = \sum_{j=0}^k y_j l_j(x).$$