

南京大学 ACM-ICPC 集训队代码模版库



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1 General

1.1 Code library checksum

```
ab14 #!/usr/bin/python3
c502 import re, sys, hashlib
427e
f7db for line in sys.stdin.read().strip().split("\n") :
ddf5     print(hashlib.md5(re.sub(r'\s|//[.]*', '', line).encode('utf8')).hexdigest()
        [-4:], line)
```

1.2 Makefile

```
dab2 .PHONY : run
427e
207e $(t) : $(t).cpp
2d16     g++ --std=c++14 -Wall -D__LOCAL_DEBUG__ -fsanitize=undefined -fsanitize=
        address -ggdb -pipe -o $@ $<
427e
5f25 run : $(t)
bf3e     ./$<(t) < $(t).in
```

1.3 .vimrc

```
914c set nocompatible
733d syntax on
6bbc colorscheme slate
7db5 set number
b0e3 set cursorline
061b set shiftwidth=2
8011 set softtabstop=2
a66d set tabstop=2
d23a set expandtab
5245 set magic
740c set smartindent
bee8 set backspace=indent,eol,start
815d set cmdheight=1
0a40 set laststatus=2
1c67 set whichwrap=b,s,<,>,[,]
```

1.4 Stack

```
const int STK_SZ = 2000000;
char STK[STK_SZ * sizeof(void)];
void *STK_BAK;

#if defined(__i386__)
#define SP "%esp"
#elif defined(__x86_64__)
#define SP "%rsp"
#endif

int main() {
    asm volatile("movl SP, %0; movl %1, SP: =g(STK_BAK):g(STK+sizeof(STK));");
    ;

    // main program

    asm volatile("movl %0, SP::g(STK_BAK);");
    return 0;
}
```

```
bebe
effc
4e99
427e
7bc9
0894
ac7a
a9ea
1937
427e
3117
3750
427e
427e
427e
6856
7021
95cf
```

1.5 Template

```
#include <bits/stdc++.h>
using namespace std;

#ifdef __LOCAL_DEBUG__
# define _debug(fmt, ...) fprintf(stderr, "[%s] " fmt "\n", \
    __func__, __VA_ARGS__)
#else
# define _debug(...) ((void) 0)
#endif

#define rep(i, n) for (int i=0; i<(n); i++)
#define Rep(i, n) for (int i=1; i<=(n); i++)
#define range(x) begin(x), end(x)
typedef long long LL;
typedef unsigned long long ULL;
```

```
302f
421c
427e
426f
3341
611f
a8cb
e6b5
1937
0d6c
cfe3
3505
5cad
b773
```

2 Miscellaneous Algorithms

2.1 2-SAT

```

0f42 const int MAXN = 100005;
03a9 struct twoSAT{
5c83     int n;
8f72     vector<int> G[MAXN*2];
d060     bool mark[MAXN*2];
b42d     int S[MAXN*2], c;
427e
d34f     void init(int n){
b985         this->n = n;
f9ec         for (int i=0; i<n*2; i++) G[i].clear();
0609         memset(mark, 0, sizeof(mark));
95cf     }
427e
3bd5     bool dfs(int x){
bd70         if (mark[x^1]) return false;
c96a         if (mark[x]) return true;
fd23         mark[x] = true;
4bea         S[c++] = x;
1ce6         for (int i=0; i<G[x].size(); i++)
d942             if (!dfs(G[x][i])) return false;
3361         return true;
95cf     }
427e
5894     void add_clause(int x, bool xval, int y, bool yval){
6afe         x = x * 2 + xval;
e680         y = y * 2 + yval;
81cc         G[x^1].push_back(y);
6835         G[y^1].push_back(x);
95cf     }
427e
d0cb     bool solve() {
7c39         for (int i=0; i<n*2; i+=2){
e63f             if (!mark[i] && !mark[i+1]){
88fb                 c = 0;
f4b9                 if (!dfs(i)){
3f03                     while (c > 0) mark[S[--c]] = false;
86c5                     if (!dfs(i+1)) return false;
95cf                 }
95cf             }

```

```

    }
    return true;
}

inline bool value(unsigned i){return mark[2*i+1];}
};

```

95cf
3361
95cf
427e
5f0a
329b

2.2 Knuth's optimization

```

int n;
int dp[256][256], dc[256][256];

template <typename T>
void compute(T cost) {
    for (int i = 0; i <= n; i++) {
        dp[i][i] = 0;
        dc[i][i] = i;
    }
    rep (i, n) {
        dp[i][i+1] = 0;
        dc[i][i+1] = i;
    }
    for (int len = 2; len <= n; len++) {
        for (int i = 0; i + len <= n; i++) {
            int j = i + len;
            int lbnd = dc[i][j-1], rbnd = dc[i+1][j];
            dp[i][j] = INT_MAX / 2;
            int c = cost(i, j);
            for (int k = lbnd; k <= rbnd; k++) {
                int res = dp[i][k] + dp[k][j] + c;
                if (res < dp[i][j]) {
                    dp[i][j] = res;
                    dc[i][j] = k;
                }
            }
        }
    }
};

```

5c83
d77c
427e
b7ec
0bc7
0423
8f5e
9488
95cf
be8e
95b5
aa0f
95cf
ec08
88b8
d3da
9824
a24a
f933
90d2
9bd0
26b5
e6af
9c88
95cf
95cf
95cf
95cf
329b

2.3 Mo's algorithm

All intervals are closed on both sides. When running functions `enter()` and `leave()`, the global `l` and `r` has not changed yet.

Usage:

```
add_query(id, l, r)    Add id-th query [l, r].
run()                 Run Mo's algorithm.
init()                TODO. Initialize the range [l, r].
yield(id)             TODO. Yield answer for id-th query.
enter(o)              TODO. Add o-th element.
leave(o)              TODO. Remove o-th element.
```

```
5194 constexpr int BLOCK_SZ = 300;
427e
3ec4 struct query { int l, r, id; };
d26a vector<query> queries;
427e
1e30 void add_query(int id, int l, int r) {
54c9     queries.push_back(query{l, r, id});
95cf }
427e
9f6b int l, r;
427e
427e // ----- functions to implement -----
62b4 inline void init();
50e1 inline void yield(int id);
b20d inline void enter(int o);
13af inline void leave(int o);
427e
37f0 void run() {
ab0b     if (queries.empty()) return;
8508     sort(range(queries), [](query lhs, query rhs) {
c7f8         int lb = lhs.l / BLOCK_SZ, rb = rhs.l / BLOCK_SZ;
03e7         if (lb != rb) return lb < rb;
0780         return lhs.r < rhs.r;
b251     });
6196     l = queries[0].l;
9644     r = queries[0].r;
07e2     init();
5bc9     for (query q : queries) {
7bc7         while (l > q.l) enter(l - 1), l--;
d646         while (r < q.r) enter(r + 1), r++;
13f0         while (l < q.l) leave(l), l++;
e1c6         while (r > q.r) leave(r), r--;
```

```
        yield(q.id);
    }
}
```

```
82f5
95cf
95cf
```

3 String

3.1 Knuth-Morris-Pratt algorithm

```
const int SIZE = 10005;

struct kmp_matcher {
    char p[SIZE];
    int fail[SIZE];
    int len;

    void construct(const char* needle) {
        len = strlen(p);
        strcpy(p, needle);
        fail[0] = fail[1] = 0;
        for (int i = 1; i < len; i++) {
            int j = fail[i];
            while (j && p[i] != p[j]) j = fail[j];
            fail[i + 1] = p[i] == p[j] ? j + 1 : 0;
        }
    }

    inline void found(int pos) {
        // ! add codes for having found at pos
    }

    void match(const char* haystack) { // must be called after construct
        const char* t = haystack;
        int n = strlen(t);
        int j = 0;
        rep(i, n) {
            while (j && p[j] != t[i]) j = fail[j];
            if (p[j] == t[i]) j++;
            if (j == len) found(i - len + 1);
        }
    }
};
```

```
2836
427e
d02b
2d81
9847
57b7
427e
60cf
aaa1
3a87
3dd4
d8a8
147f
3c79
4643
95cf
95cf
427e
c464
427e
95cf
427e
2daf
700f
8482
8fd0
be8e
4e19
b5d5
f024
95cf
95cf
329b
```

3.2 Manacher algorithm

```

81d4 struct Manacher {
cd09     int Len;
9255     vector<int> lc;
b301     string s;
427e
ec07     void work() {
c033         lc[1] = 1;
6bef         int k = 1;
427e
491f         for (int i = 2; i <= Len; i++) {
7957             int p = k + lc[k] - 1;
5e04             if (i <= p) {
24a1                 lc[i] = min(lc[2 * k - i], p - i + 1);
8e2e             } else {
e0e5                 lc[i] = 1;
95cf             }
74ff             while (s[i + lc[i]] == s[i - lc[i]]) lc[i]++;
2b9a             if (i + lc[i] > k + lc[k]) k = i;
95cf         }
95cf     }
427e
bfd5     void init(const char *tt) {
aaaf         int len = strlen(tt);
f701         s.resize(len * 2 + 10);
7045         lc.resize(len * 2 + 10);
8e13         s[0] = '*';
ae54         s[1] = '#';
1321         for (int i = 0; i < len; i++) {
e995             s[i * 2 + 2] = tt[i];
69fd             s[i * 2 + 1] = '#';
95cf         }
43fd         s[len * 2 + 1] = '#';
75d1         s[len * 2 + 2] = '\0';
61f7         Len = len * 2 + 2;
3e7a         work();
95cf     }
427e
b194     pair<int, int> maxpal(int l, int r) {
901a         int center = l + r + 1;
ffb2         int rad = lc[center] / 2;
ab54         int rmid = (l + r + 1) / 2;

```

```

        int r1 = rmid - rad, rr = rmid + rad - 1;
        if ((r ^ 1) & 1) {
        } else rr++;
        return {max(l, r1), min(r, rr)};
    }
};

```

```

17e4
3908
69f3
69dc
95cf
329b

```

3.3 Aho-corasick automaton

```

struct AC : Trie {
    int fail[MAXN];
    int last[MAXN];

    void construct() {
        queue<int> q;
        fail[0] = 0;
        rep(c, CHARN) {
            if (int u = tr[0][c]) {
                fail[u] = 0;
                q.push(u);
                last[u] = 0;
            }
        }
        while (!q.empty()) {
            int r = q.front();
            q.pop();
            rep(c, CHARN) {
                int u = tr[r][c];
                if (!u) {
                    tr[r][c] = tr[fail[r]][c];
                    continue;
                }
                q.push(u);
                int v = fail[r];
                while (v && !tr[v][c]) v = fail[v];
                fail[u] = tr[v][c];
                last[u] = tag[fail[u]] ? fail[u] : last[fail[u]];
            }
        }
    }

    void found(int pos, int j) {

```

```

a1ad
9143
daca
427e
8690
93d2
a7a6
ce3c
b1c6
a506
3e14
f689
95cf
95cf
cc78
31f0
15dd
ce3c
ab59
0ef5
9d58
b333
95cf
3e14
b3ff
d2ea
c275
654c
95cf
95cf
95cf
427e
7752

```

```

043e     if (j) {
427e         // ! add codes for having found word with tag[j]
4a96         found(pos, last[j]);
95cf     }
95cf }
427e
9785 void find(const char* text) { // must be called after construct()
80a4     int p = 0, c, len = strlen(text);
9c94     rep(i, len) {
b3db         c = id(text[i]);
f119         p = tr[p][c];
f08e         if (tag[p])
389b             found(i, p);
1e67         else if (last[p])
299e             found(i, last[p]);
95cf     }
95cf }
329b };

```

3.4 Suffix array

The character immediately after the end of the string **MUST** be set to the **UNIQUE SMALLEST** element.

Usage:

s[]	the source string
sa[i]	the index of starting position of i -th suffix
rk[i]	the number of suffixes less than the suffix starting from i
h[i]	the longest common prefix between the i -th and $(i-1)$ -th lexicographically smallest suffixes
n	size of source string
m	size of character set

```

de09 void radix_sort(int x[], int y[], int sa[], int n, int m) {
ec00     static int cnt[1000005]; // size > max(n, m)
6066     fill(cnt, cnt + m, 0);
93b7     rep(i, n) cnt[x[y[i]]]++;
9154     partial_sum(cnt, cnt + m, cnt);
acac     for (int i = n - 1; i >= 0; i--) sa[--cnt[x[y[i]]]] = y[i];
95cf }
427e
c939 void suffix_array(int s[], int sa[], int rk[], int n, int m) {
a69a     static int y[1000005]; // size > n
7306     copy(s, s + n, rk);

```

```

iota(y, y + n, 0);
radix_sort(rk, y, sa, n, m);
for (int j = 1, p = 0; j <= n; j <= 1, m = p, p = 0) {
    for (int i = n - j; i < n; i++) y[p++] = i;
    rep(i, n) if (sa[i] >= j) y[p++] = sa[i] - j;
    radix_sort(rk, y, sa, n, m + 1);
    swap_ranges(rk, rk + n, y);
    rk[sa[0]] = p = 1;
    for (int i = 1; i < n; i++)
        rk[sa[i]] = ((y[sa[i]] == y[sa[i-1]] and y[sa[i]+j] == y[sa[i-1]+j])
            ? p : ++p);
    if (p == n) break;
}
rep(i, n) rk[sa[i]] = i;
}

void calc_height(int s[], int sa[], int rk[], int h[], int n) {
    int k = 0;
    h[0] = 0;
    rep(i, n) {
        k = max(k - 1, 0);
        if (rk[i]) while (s[i+k] == s[sa[rk[i]-1]+k]) ++k;
        h[rk[i]] = k;
    }
}

```

3.5 Trie

```

const int MAXN = 12000;
const int CHARN = 26;

inline int id(char c) { return c - 'a'; }

struct Trie {
    int n;
    int tr[MAXN][CHARN]; // Trie tree, 0 denotes fail
    int tag[MAXN];

    Trie() {
        memset(tr[0], 0, sizeof(tr[0]));
        tag[0] = 0;
        n = 1;
    }
}

```

```

95cf }
427e
427e // tag should not be 0
30b0 void add(const char* s, int t) {
d50a     int p = 0, c, len = strlen(s);
9c94     rep(i, len) {
3140         c = id(s[i]);
d6c8         if (!tr[p][c]) {
26dd             memset(tr[n], 0, sizeof(tr[n]));
2e5c             tag[n] = 0;
73bb             tr[p][c] = n++;
95cf         }
f119         p = tr[p][c];
95cf     }
35ef     tag[p] = t;
95cf }
427e
427e // returns 0 if not found
427e // AC automaton does not need this function
216c int search(const char* s) {
d50a     int p = 0, c, len = strlen(s);
9c94     rep(i, len) {
3140         c = id(s[i]);
f339         if (!tr[p][c]) return 0;
f119         p = tr[p][c];
95cf     }
840e     return tag[p];
95cf }
329b };

```

3.6 Rolling hash

PLEASE call `init_hash()` in `int main()`!

Usage:

`build(str)` Construct the hasher with given string.
`operator()(l, r)` Get hash value of substring $[l, r)$.

```

1e42 const LL mod = 1006658951440146419, g = 967;
9f60 const int MAXN = 200005;
0291 LL pg[MAXN];
427e
dfe7 inline LL mul(LL x, LL y) { return __int128_t(x) * y % mod; }
427e

```

```

void init_hash() { // must be called in `int main()`
    pg[0] = 1;
    for (int i = 1; i < MAXN; i++) pg[i] = mul(pg[i-1], g);
}

struct hasher {
    LL val[MAXN];

    void build(const char *str) { // assume lower-case letter only
        for (int i = 0; str[i]; i++)
            val[i+1] = (mul(val[i], g) + str[i]) % mod;
    }

    LL operator() (int l, int r) { // [l, r)
        return (val[r] - mul(val[l], pg[r-l]) + mod) % mod;
    }
};

```

599a
286f
4af8
95cf
427e
7e62
534a
427e
4554
f937
9645
95cf
427e
19f8
9986
95cf
329b

4 Math

4.1 Extended Euclidean algorithm and Chinese remainder theorem

```

void exgcd(LL a, LL b, LL &g, LL &x, LL &y) {
    if (!b) g = a, x = 1, y = 0;
    else {
        exgcd(b, a % b, g, y, x);
        y -= x * (a / b);
    }
}

LL crt(LL r[], LL p[], int n) {
    LL q = 1, ret = 0;
    rep (i, n) q *= p[i];
    rep (i, n) {
        LL m = q / p[i];
        LL d, x, y;
        exgcd(p[i], m, d, x, y);
        ret = (ret + y * m * r[i]) % q;
    }
    return (q + ret) % q;
}

```

4fba
7db6
037f
ffca
d798
95cf
95cf
427e
e491
84e6
00d9
be8e
98b4
9f4f
b082
3cd3
95cf
2e47
95cf

4.2 Matrix powermod

```

44b4 const int MAXN = 105;
92df const LL modular = 1000000007;
5c83 int n; // order of matrices
427e
8864 struct matrix{
3180     LL m[MAXN][MAXN];
427e
43c5     void operator *=(matrix& a){
e735         static LL t[MAXN][MAXN];
34d7         Rep (i, n){
4c11             Rep (j, n){
ee1e                 t[i][j] = 0;
c4a7                 Rep (k, n){
fcac                     t[i][j] += (m[i][k] * a.m[k][j]) % modular;
199e                     t[i][j] %= modular;
95cf                 }
95cf             }
95cf         }
dad4         memcpy(m, t, sizeof(t));
95cf     }
329b };
427e
63d8 matrix r;
3ec2 void m_powmod(matrix& b, LL e){
83f0     memset(r.m, 0, sizeof(r.m));
a7c3     Rep(i, n)
de64         r.m[i][i] = 1;
3e90     while (e){
5a0e         if (e & 1) r *= b;
35c5         b *= b;
16fc         e >>= 1;
95cf     }
95cf }

```

4.3 Linear basis

```

8b44 const int MAXD = 30;
03a6 struct linearbasis {
3558     ULL b[MAXD] = {};
427e

```

```

bool insert(LL v) {
    for (int j = MAXD - 1; j >= 0; j--) {
        if (!(v & (1ll << j))) continue;
        if (b[j]) v ^= b[j]
        else {
            for (int k = 0; k < j; k++)
                if (v & (1ll << k)) v ^= b[k];
            for (int k = j + 1; k < MAXD; k++)
                if (b[k] & (1ll << j)) b[k] ^= v;
            b[j] = v;
            return true;
        }
    }
    return false;
}

```

4.4 Gauss elimination over finite field

```

const LL p = 1000000007;

LL powmod(LL b, LL e) {
    LL r = 1;
    while (e) {
        if (e & 1) r = r * b % p;
        b = b * b % p;
        e >>= 1;
    }
    return r;
}

typedef vector<LL> VLL;
typedef vector<VLL> WLL;

```

```

LL gauss(WLL &a, WLL &b) {
    const int n = a.size(), m = b[0].size();
    vector<int> irow(n), icol(n), ipiv(n);
    LL det = 1;

    rep (i, n) {
        int pj = -1, pk = -1;
        rep (j, n) if (!ipiv[j])

```

1566
9b2b
de36
ee78
037f
7836
f0b4
b0aa
46c9
8295
3361
95cf
95cf
438e
95cf
329b

b784
427e
2a2c
95a2
3e90
1783
5549
16fc
95cf
547e
95cf
427e
c130
42ac
427e
2c62
561b
a25e
2976
427e
be8e
d2b5
6b4a

```

e582     rep (k, n) if (!ipiv[k])
6112         if (pj == -1 || a[j][k] > a[pj][pk]) {
a905             pj = j;
657b             pk = k;
95cf         }
d480     if (a[pj][pk] == 0) return 0;
0305     ipiv[pk]++;
8dad     swap(a[pj], a[pk]);
aad8     swap(b[pj], b[pk]);
be4d     if (pj != pk) det = (p - det) % p;
d080     irow[i] = pj;
f156     icol[i] = pk;
427e
4ecd     LL c = powmod(a[pk][pk], p - 2);
865b     det = det * a[pk][pk] % p;
c36a     a[pk][pk] = 1;
dd36     rep (j, n) a[pk][j] = a[pk][j] * c % p;
1b23     rep (j, m) b[pk][j] = b[pk][j] * c % p;
f8f3     rep (j, n) if (j != pk) {
e97f         c = a[j][pk];
c449         a[j][pk] = 0;
820b         rep (k, n) a[j][k] = (a[j][k] + p - a[pk][k] * c % p) % p;
f039         rep (k, m) b[j][k] = (b[j][k] + p - b[pk][k] * c % p) % p;
95cf     }
95cf }
427e
37e1     for (int j = n - 1; j >= 0; j--) if (irow[j] != icol[j]) {
50dc         for (int k = 0; k < n; k++) swap(a[k][irow[j]], a[k][icol[j]]);
95cf     }
f27f     return det;
95cf }

```

4.5 Berlekamp-Massey algorithm

```

d790 vector<int> berlekamp(const vector<int>& a) {
4166     vector<int> p = {1}, r = {1};
baed     int dif = 1;
8bc9     rep (i, a.size()) {
3e58         int u = 0;
ac8e         rep (j, p.size())
a488             u = (u + 111 * p[j] * a[i-j]) % mod;
eae9         if (u == 0) {

```

```

        r.insert(r.begin(), 0);
    } else {
        auto op = p;
        p.resize(max(p.size(), r.size() + 1));
        int idif = inv(dif);
        rep (j, r.size())
            p[j+1] =
                (p[j+1] - 111 * r[j] * idif % mod * u % mod + mod) % mod;
        dif = u;
        r = op;
    }
}
return p;
}

```

```

b14c
8e2e
0c78
02f6
786b
9b57
793c
1836
644c
bc58
95cf
95cf
e149
95cf

```

4.6 Fast Walsh-Hadamard transform

```

void fwt(int* a, int n){
    for (int d = 1; d < n; d <= 1)
        for (int i = 0; i < n; i += d < 1)
            rep (j, d){
                int x = a[i+j], y = a[i+j+d];
                // a[i+j] = x+y, a[i+j+d] = x-y;    // xor
                // a[i+j] = x+y;                    // and
                // a[i+j+d] = x+y;                    // or
            }
}

```

```

061e
5595
05f2
b833
7796
427e
427e
427e
95cf
95cf
427e
4db1

```

```

void ifwt(int* a, int n){
    for (int d = 1; d < n; d <= 1)
        for (int i = 0; i < n; i += d < 1)
            rep (j, d){
                int x = a[i+j], y = a[i+j+d];
                // a[i+j] = (x+y)/2, a[i+j+d] = (x-y)/2;    // xor
                // a[i+j] = x-y;                                // and
                // a[i+j+d] = y-x;                                // or
            }
}

```

```

5595
05f2
b833
7796
427e
427e
427e
95cf
95cf
427e
2ab6

```

```

void conv(int* a, int* b, int n){
    fwt(a, n);
    fwt(b, n);

```

```

950a
e427

```

```

8a42     rep(i, n) a[i] *= b[i];
430f     ifwt(a, n);
95cf }

```

4.7 Fast fourier transform

```

4e09 const int NMAX = 1<<20;
427e
3fbf typedef complex<double> cplx;
427e
abd1 const double PI = 2*acos(0.0);
12af struct FFT{
c47c     int rev[NMAX];
27d7     cplx omega[NMAX], oinv[NMAX];
9827     int K, N;
427e
1442     FFT(int k){
e209         K = k; N = 1 << k;
b393         rep (i, N){
7ba3             rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
1908             omega[i] = polar(1.0, 2.0 * PI / N * i);
a166             oinv[i] = conj(omega[i]);
95cf         }
95cf     }
427e
b941     void dft(cplx* a, cplx* w){
a215         rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
ac6e         for (int l = 2; l <= N; l *= 2){
2969             int m = l/2;
b3cf             for (cplx* p = a; p != a + N; p += l)
c24f                 rep (k, m){
fe06                     cplx t = w[N/l*k] * p[k+m];
ecbf                     p[k+m] = p[k] - t; p[k] += t;
95cf                 }
95cf             }
95cf         }
427e
617b     void fft(cplx* a){dft(a, omega);}
a123     void ifft(cplx* a){
3b2f         dft(a, oinv);
57fc         rep (i, N) a[i] /= N;
95cf     }

```

```

void conv(cplx* a, cplx* b){
    fft(a); fft(b);
    rep (i, N) a[i] *= b[i];
    ifft(a);
}
};

```

```

427e
bdc0
6497
12a5
f84e
95cf
329b

```

4.8 Number theoretic transform

```

const int NMAX = 1<<21;

// 998244353 = 7*17*2^23+1, G = 3
const int P = 1004535809, G = 3; // = 479*2^21+1

struct NTT{
    int rev[NMAX];
    LL omega[NMAX], oinv[NMAX];
    int g, g_inv; // g:  $g_n = G^{(P-1)/n}$ 
    int K, N;

    LL powmod(LL b, LL e){
        LL r = 1;
        while (e){
            if (e&1) r = r * b % P;
            b = b * b % P;
            e >>= 1;
        }
        return r;
    }

    NTT(int k){
        K = k; N = 1 << k;
        g = powmod(G, (P-1)/N);
        g_inv = powmod(g, N-1);
        omega[0] = oinv[0] = 1;
        rep (i, N){
            rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
            if (i){
                omega[i] = omega[i-1] * g % P;
                oinv[i] = oinv[i-1] * g_inv % P;
            }
        }
    }
}

```

```

4ab9
427e
427e
fb9a
427e
87ab
c47c
0eda
81af
9827
427e
2a2c
95a2
3e90
6624
489e
16fc
95cf
547e
95cf
427e
f420
e209
7652
4b3a
e04f
b393
7ba3
ad4f
8d8b
9e14
95cf

```

```

95cf    }
95cf    }
427e
9668    void _ntt(LL* a, LL* w){
a215        rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
ac6e        for (int l = 2; l <= N; l *= 2){
2969            int m = l/2;
7a1d            for (LL* p = a; p != a + N; p += l)
c24f                rep (k, m){
0ad3                    LL t = w[N/l*k] * p[k+m] % P;
6209                    p[k+m] = (p[k] - t + P) % P;
fa1b                    p[k] = (p[k] + t) % P;
95cf                }
95cf            }
95cf        }
427e
92ea    void ntt(LL* a){_ntt(a, omega);}
5daf    void intt(LL* a){
1f2a        LL inv = powmod(N, P-2);
9910        _ntt(a, oinv);
a873        rep (i, N) a[i] = a[i] * inv % P;
95cf    }
427e
3a5b    void conv(LL* a, LL* b){
ad16        ntt(a); ntt(b);
e49e        rep (i, N) a[i] = a[i] * b[i] % P;
5748        intt(a);
95cf    }
329b    };

```

4.9 Sieve of Euler

```

cfc3    const int MAXX = 1e7+5;
5861    bool p[MAXX];
73ae    int prime[MAXX], sz;
427e
9bc6    void sieve(){
9628        p[0] = p[1] = 1;
1ec8        for (int i = 2; i < MAXX; i++){
bf28            if (!p[i]) prime[sz++] = i;
e82c            for (int j = 0; j < sz && i*prime[j] < MAXX; j++){
b6a9                p[i*prime[j]] = 1;

```

```

        if (i % prime[j] == 0) break;
    }
}

```

4.10 Sieve of Euler (General)

```

namespace sieve {
constexpr int MAXN = 10000007;
bool p[MAXN]; // true if not prime
int prime[MAXN], sz;
int pval[MAXN], pcnt[MAXN];
int f[MAXN];

void exec(int N = MAXN) {
    p[0] = p[1] = 1;

    pval[1] = 1;
    pcnt[1] = 0;
    f[1] = 1;

    for (int i = 2; i < N; i++) {
        if (!p[i]) {
            prime[sz++] = i;
            for (LL j = i; j < N; j *= i) {
                int b = j / i;
                pval[j] = i * pval[b];
                pcnt[j] = pcnt[b] + 1;
                f[j] = _____; // f[j] = f(i^pcnt[j])
            }
        }
        for (int j = 0; i * prime[j] < N; j++) {
            int x = i * prime[j]; p[x] = 1;
            if (i % prime[j] == 0) {
                pval[x] = pval[i] * prime[j];
                pcnt[x] = pcnt[i] + 1;
            } else {
                pval[x] = prime[j];
                pcnt[x] = 1;
            }
            if (x != pval[x]) {
                f[x] = f[x / pval[x]] * f[pval[x]]

```

5f51
95cf
95cf
95cf

b62e
6589
e982
6ae8
cbf7
6030
427e
76f6
9628
427e
8a8a
bdda
c6b9
427e
a643
01d6
b2b2
37d9
758c
81fd
e0f3
a96c
95cf
95cf
34c0
f87a
20cc
9985
3f93
8e2e
cc91
6322
95cf
6191
d614

```

95cf      }
5f51      if (i % prime[j] == 0) break;
95cf    }
95cf  }
95cf  }
95cf  }

```

4.11 Miller-Rabin primality test

The array `a[]` (excluding `senitel`, i.e. `LLONG_MAX`) should be

```

{2}                when  $n < 2,047$ .
{2, 7, 61}          when  $n < 4,759,123,141 (2^{32})$ .
{2, 3, 5, 7, 11}    when  $n < 2.1 \times 10^{12}$ .
{2, 325, 9375, 28178, 450775, 9780504, 1795265022} when  $n < 2^{64}$ .

```

```

f16f bool test(LL n){
59f2   if (n < 3) return n==2;
427e   // ! The array a[] should be modified if the range of x changes.
3f11   const LL a[] = {2LL, 7LL, 61LL, LLONG_MAX};
c320   LL r = 0, d = n-1, x;
f410   while (~d & 1) d >>= 1, r++;
2975   for (int i=0; a[i] < n; i++){
ece1     x = powmod(a[i], d, n); // ! powmod must use for 64bit mulmod
7f99     if (x == 1 || x == n-1) goto next;
e257     rep (i, r) {
d7ff       x = mulmod(x, x, n);
8d2e       if (x == n-1) goto next;
95cf     }
438e     return false;
d490 next:;
95cf   }
3361   return true;
95cf }

```

4.12 Pollard's rho algorithm

```

2e6b ULL gcd(ULL a, ULL b) {return b ? gcd(b, a % b) : a;}
427e
54a5 ULL PollardRho(ULL n){
45eb   ULL c, x, y, d = n;

```

```

if (~n&1) return 2;
while (d == n){
  x = y = 2;
  d = 1;
  c = rand() % (n - 1) + 1;
  while (d == 1){
    x = (mulmod(x, x, n) + c) % n;
    y = (mulmod(y, y, n) + c) % n;
    y = (mulmod(y, y, n) + c) % n;
    d = gcd(x>y ? x-y : y-x, n);
  }
}
return d;
}

```

```

d3e5
3c69
0964
4753
5952
9e5b
33d5
e1bf
e1bf
a313
95cf
95cf
5d89
95cf

```

5 Graph Theory

5.1 Strongly connected component

```

const int MAXV = 100005;

struct graph{
  vector<int> adj[MAXV];
  stack<int> s;
  int V; // number of vertices
  int pre[MAXV], lnk[MAXV], scc[MAXV];
  int time, sccn;

  void add_edge(int u, int v){
    adj[u].push_back(v);
  }

  void dfs(int u){
    pre[u] = lnk[u] = ++time;
    s.push(u);
    for (int v : adj[u]){
      if (!pre[v]){
        dfs(v);
        lnk[u] = min(lnk[u], lnk[v]);
      } else if (!scc[v]){
        lnk[u] = min(lnk[u], pre[v]);
      }
    }
  }
}

```

```

837c
427e
2ea0
88e3
9cad
3d02
8b6c
27ee
427e
bfab
c71a
95cf
427e
d714
7e41
80f6
18f6
173e
5f3c
002c
6068
d5df

```

```

95cf    }
95cf    }
8de2    if (lnk[u] == pre[u]){
660f        sccn++;
3c9e        int x;
a69f        do {
3834            x = s.top(); s.pop();
b0e9            scc[x] = sccn;
6757        } while (x != u);
95cf    }
95cf    }
427e
4c88    void find_scc(){
f4a2        time = sccn = 0;
8de7        memset(scc, 0, sizeof scc);
8c2f        memset(pre, 0, sizeof pre);
6901        Rep (i, V){
56d1            if (!pre[i]) dfs(i);
95cf        }
95cf    }
427e
27ce    vector<int> adjc[MAXV];
364d    void contract(){
1a1e        Rep (i, V)
21a2            rep (j, adj[i].size()){
b730                if (scc[i] != scc[adj[i][j]])
b46e                    adjc[scc[i]].push_back(scc[adj[i][j]]);
95cf            }
95cf        }
329b    };

```

5.2 Vertex biconnected component

```

0f42    const int MAXN = 100005;
2ea0    struct graph {
33ae        int pre[MAXN], iscut[MAXN], bccno[MAXN], dfs_clock, bcc_cnt;
848f        vector<int> adj[MAXN], bcc[MAXN];
6b06        set<pair<int, int>> bcce[MAXN];
427e
76f7        stack<pair<int, int>> s;
427e
bfab        void add_edge(int u, int v) {

```

```

        adj[u].push_back(v);
        adj[v].push_back(u);
    }

int dfs(int u, int fa) {
    int lowu = pre[u] = ++dfs_clock;
    int child = 0;
    for (int v : adj[u]) {
        if (!pre[v]) {
            s.push({u, v});
            child++;
            int lowv = dfs(v, u);
            lowu = min(lowu, lowv);
            if (lowv >= pre[u]) {
                iscut[u] = 1;
                bcc[bcc_cnt].clear();
                bcce[bcc_cnt].clear();
                while (1) {
                    int xu, xv;
                    tie(xu, xv) = s.top(); s.pop();
                    bcce[bcc_cnt].insert({min(xu, xv), max(xu, xv)});
                    if (bccno[xu] != bcc_cnt) {
                        bcc[bcc_cnt].push_back(xu);
                        bccno[xu] = bcc_cnt;
                    }
                    if (bccno[xv] != bcc_cnt) {
                        bcc[bcc_cnt].push_back(xv);
                        bccno[xv] = bcc_cnt;
                    }
                    if (xu == u && xv == v) break;
                }
                bcc_cnt++;
            }
        } else if (pre[v] < pre[u] && v != fa) {
            s.push({u, v});
            lowu = min(lowu, pre[v]);
        }
    }
    if (fa < 0 && child == 1) iscut[u] = 0;
    return lowu;
}

void find_bcc(int n) {
    memset(pre, 0, sizeof pre);

```

```

c71a
a717
95cf
427e
7d3c
9fe6
ec14
18f6
173e
e7f8
fdcf
f851
189c
b687
6323
57eb
90b8
a147
a6a3
a0c3
0ef5
3db2
e0db
d27f
95cf
f357
752b
57c9
95cf
7096
95cf
03f5
95cf
7470
e7f8
f115
95cf
95cf
e104
1160
95cf
427e
17be
8c2f

```

```
e2d2     memset(iscut, 0, sizeof iscut);
40d3     memset(bccno, -1, sizeof bccno);
fae2     dfs_clock = bcc_cnt = 0;
5c63     rep (i, n) if (!pre[i]) dfs(i, -1);
95cf     }
329b     };
```

5.3 Cut vertices

If the graph is unconnected, the algorithm should be run on each component.

Usage:

tarjan(u, fa) Run Tarjan's algorithm on tree rooted at fa. Please call with identical u and fa.

```
9f60     const int MAXN = 200005;
0b32     vector<int> adj[MAXN];
18e4     int dfn[MAXN], low[MAXN], idx;
d39d     bool cut[MAXN];
427e
bfa8     void add_edge(int u, int v) {
c71a         adj[u].push_back(v);
a717         adj[v].push_back(u);
95cf     }
427e
50aa     void tarjan(int u, int fa) {
9891         dfn[u] = low[u] = ++idx;
ec14         int child = 0;
18f6         for (int v : adj[u]) {
3c64             if (!dfn[v]) {
9636                 tarjan(v, fa); low[u] = min(low[u], low[v]);
f368                 if (low[v] >= dfn[u] && u != fa) cut[u] = true;
7923                 child += u == fa;
95cf             }
769a             low[u] = min(low[u], dfn[v]);
95cf         }
7927         if (u == fa && child > 1) cut[u] = true;
95cf     }
```

5.4 Minimum spanning arborescence (Chu-Liu)

All vertices are 1-based.

Usage:

getans(n, root, edges) Compute the total size of MSA rooted at root.

Time Complexity: $O(|V||E|)$

```
struct edge {
    int u, v;
    LL w;
};

const int MAXN = 10005;
LL in[MAXN];
int pre[MAXN], vis[MAXN], id[MAXN];

LL getans(int n, int rt, vector<edge>& edges) {
    LL ans = 0;
    int cnt = 0;
    while (1) {
        Rep (i, n) in[i] = LLONG_MAX, id[i] = vis[i] = 0;
        for (auto e : edges) {
            if (e.u != e.v and e.w < in[e.v]) {
                pre[e.v] = e.u;
                in[e.v] = e.w;
            }
        }
        in[rt] = 0;
        Rep (i, n) {
            if (in[i] == LLONG_MAX) return -1;
            ans += in[i];
            int u;
            for (u = i; u != rt && vis[u] != i && !id[u]; u = pre[u])
                vis[u] = i;
            if (u != rt && !id[u]) {
                id[u] = ++cnt;
                for (int v = pre[u]; v != u; v = pre[v])
                    id[v] = cnt;
            }
        }
        if (!cnt) return ans;
        Rep (i, n) if (!id[i]) id[i] = ++cnt;
        for (auto& e : edges) {
            LL laz = in[e.v];
            e.u = id[e.u];
            e.v = id[e.v];
            if (e.u != e.v) e.w -= laz;
        }
    }
}
```

```

95cf      }
6cc4      n = cnt; rt = id[rt]; cnt = 0;
95cf      }
95cf  }
```

5.5 Maximum flow (Dinic)

Usage:

add_edge(u, v, c) Add an edge from u to v with capacity c .

max_flow(s, t) Compute maximum flow from s to t .

Time Complexity: For general graph, $O(V^2E)$; for network with unit capacity, $O(\min\{V^{2/3}, \sqrt{E}\}E)$; for bipartite network, $O(\sqrt{VE})$.

```

bcf8 struct edge{
60e2     int from, to;
5e6d     LL cap, flow;
329b };
427e
e2cd const int MAXN = 1005;
9062 struct Dinic {
4dbf     int n, m, s, t;
9f0c     vector<edge> edges;
b891     vector<int> G[MAXN];
bbb6     bool vis[MAXN];
b40a     int d[MAXN];
ddec     int cur[MAXN];
427e
5973     void add_edge(int from, int to, LL cap) {
7b55         edges.push_back(edge{from, to, cap, 0});
1db7         edges.push_back(edge{to, from, 0, 0});
fe77         m = edges.size();
dff5         G[from].push_back(m-2);
8f2d         G[to].push_back(m-1);
95cf     }
427e
1836     bool bfs() {
3b73         memset(vis, 0, sizeof(vis));
93d2         queue<int> q;
5d13         q.push(s);
2cd2         vis[s] = 1;
721d         d[s] = 0;
cc78         while (!q.empty()) {
66ba             int x = q.front(); q.pop();
```

```

        for (int i = 0; i < G[x].size(); i++) {
            edge& e = edges[G[x][i]];
            if (!vis[e.to] && e.cap > e.flow) {
                vis[e.to] = 1;
                d[e.to] = d[x] + 1;
                q.push(e.to);
            }
        }
    }
    return vis[t];
}

LL dfs(int x, LL a) {
    if (x == t || a == 0) return a;
    LL flow = 0, f;
    for (int& i = cur[x]; i < G[x].size(); i++) {
        edge& e = edges[G[x][i]];
        if (d[x] + 1 == d[e.to] && (f = dfs(e.to, min(a, e.cap-e.flow))) > 0)
        {
            e.flow += f;
            edges[G[x][i]^1].flow -= f;
            flow += f;
            a -= f;
            if(a == 0) break;
        }
    }
    return flow;
}

LL max_flow(int s, int t) {
    this->s = s; this->t = t;
    LL flow = 0;
    while (bfs()) {
        memset(cur, 0, sizeof(cur));
        flow += dfs(s, LLONG_MAX);
    }
    return flow;
}

vector<int> min_cut() { // call this after maxflow
    vector<int> ans;
    for (int i = 0; i < edges.size(); i++) {
        edge& e = edges[i];
        if(vis[e.from] && !vis[e.to] && e.cap > 0) ans.push_back(i);
```



```

95cf    }
4206    return ans;
95cf    }
329b   };

```

5.6 Maximum cardinality bipartite matching (Hungarian)

```

302f   #include <bits/stdc++.h>
421c   using namespace std;
427e
0d6c   #define rep(i, n) for (int i = 0; i < (n); i++)
cfe3   #define Rep(i, n) for (int i = 1; i <= (n); i++)
8843   #define range(x) (x).begin(), (x).end()
5cad   typedef long long LL;
427e
84ee   struct Hungarian{
fbf6       int nx, ny;
9ec6       vector<int> mx, my;
9d4c       vector<vector<int>> > e;
edec       vector<bool> mark;
427e
8324       void init(int nx, int ny){
c1d1           this->nx = nx;
f9c1           this->ny = ny;
ac92           mx.resize(nx); my.resize(ny);
3f11           e.clear(); e.resize(nx);
1023           mark.resize(nx);
95cf       }
427e
4589       inline void add(int a, int b){
486c           e[a].push_back(b);
95cf       }
427e
0c2b       bool augment(int i){
207c           if (!mark[i]) {
dae4               mark[i] = true;
6a1e               for (int j : e[i]){
0892                   if (my[j] == -1 || augment(my[j])){
9ca3                       mx[i] = j; my[j] = i;
3361                       return true;
95cf                   }
95cf           }

```

```

    }
    return false;
}

int match(){
    int ret = 0;
    fill(range(mx), -1);
    fill(range(my), -1);
    rep (i, nx){
        fill(range(mark), false);
        if (augment(i)) ret++;
    }
    return ret;
}
};

```

```

95cf
438e
95cf
427e
3fac
5b57
b0f1
b957
4ed1
13a5
cc89
95cf
ee0f
95cf
329b

```

5.7 Maximum matching of general graph (Edmond's blossom)

Usage:

init(n)	Initialize the template with n vertices, numbered from 1.
add_edge(u, v)	Add an undirected edge uv .
solve()	Find the maximum matching. Return the number of matched edges.
mate[]	The mate of a matched vertex. If it is not matched, then the value is 0.

Time Complexity: $O(|V|^3)$, but extremely fast in practice.

```

const int MAXN = 1024;
struct Blossom {

    vector<int> adj[MAXN];
    queue<int> q;
    int n;
    int label[MAXN], mate[MAXN], save[MAXN], used[MAXN];

    void init(int nv) {
        n = nv; for (auto& v : adj) v.clear();
        fill(range(label), 0); fill(range(mate), 0);
        fill(range(save), 0); fill(range(used), 0);
    }

    void add_edge(int u, int v) { adj[u].push_back(v); adj[v].push_back(u); }

```

```

c041
6ab1
427e
0b32
93d2
5c83
0de2
427e
2186
3728
477d
bb35
95cf
427e
c2dd

```

```

427e void rematch(int x, int y) {
2a48     int m = mate[x]; mate[x] = y;
8af8     if (mate[m] == x) {
1aa4         if (label[x] <= n) {
f4ba             mate[m] = label[x]; rematch(label[x], m);
740a         } else {
8e2e             int a = 1 + (label[x] - n - 1) / n;
3341             int b = 1 + (label[x] - n - 1) % n;
2885             rematch(a, b); rematch(b, a);
ef33         }
95cf     }
95cf }
427e
8a50 void traverse(int x) {
43c0     Rep (i, n) save[i] = mate[i];
2ef7     rematch(x, x);
34d7     Rep (i, n) {
62c5         if (mate[i] != save[i]) used[i] ++;
97ef         mate[i] = save[i];
95cf     }
95cf }
427e
8bf8 void relabel(int x, int y) {
d101     Rep (i, n) used[i] = 0;
c4ea     traverse(x); traverse(y);
34d7     Rep (i, n) {
dee9         if (used[i] == 1 and label[i] < 0) {
1c22             label[i] = n + x + (y - 1) * n;
eb31             q.push(i);
95cf         }
95cf     }
95cf }
427e
a0ce int solve() {
34d7     Rep (i, n) {
a073         if (mate[i]) continue;
1fc0         Rep (j, n) label[j] = -1;
7676         label[i] = 0; q = queue<int>(); q.push(i);
1c7d         while (q.size()) {
66ba             int x = q.front(); q.pop();
b98c             for (int y : adj[x]) {
c07f                 if (mate[y] == 0 and i != y) {
7f36                     mate[y] = x; rematch(x, y); q = queue<int>(); break;

```

```

        }
        if (label[y] >= 0) { relabel(x, y); continue; }
        if (label[mate[y]] < 0) {
            label[mate[y]] = x; q.push(mate[y]);
        }
    }
}
int cnt = 0;
Rep (i, n) cnt += (mate[i] > i);
return cnt;
}
};

```

5.8 Minimum cost maximum flow

```

struct edge{
    int from, to;
    int cap, flow;
    LL cost;
};

const LL INF = LLONG_MAX / 2;
const int MAXN = 5005;
struct MCMF {
    int s, t, n, m;
    vector<edge> edges;
    vector<int> G[MAXN];
    bool inq[MAXN]; // queue
    LL d[MAXN]; // distance
    int p[MAXN]; // previous
    int a[MAXN]; // improvement
};

void add_edge(int from, int to, int cap, LL cost) {
    edges.push_back(edge{from, to, cap, 0, cost});
    edges.push_back(edge{to, from, 0, 0, -cost});
    m = edges.size();
    G[from].push_back(m-2);
    G[to].push_back(m-1);
}

bool spfa(){

```

```

93d2     queue<int> q;
8494     fill(d, d + MAXN, INF); d[s] = 0;
fd48     memset(inq, 0, sizeof(inq));
5e7c     q.push(s); inq[s] = true;
2dae     p[s] = 0; a[s] = INT_MAX;
cc78     while (!q.empty()){
b0aa         int u = q.front(); q.pop(); inq[u] = false;
3bba         for (int i : G[u]) {
56d8             edge& e = edges[i];
3601             if (e.cap > e.flow && d[e.to] > d[u] + e.cost){
55bc                 d[e.to] = d[u] + e.cost;
0bea                 p[e.to] = G[u][i];
8249                 a[e.to] = min(a[u], e.cap - e.flow);
e5d3                 if (!inq[e.to]) q.push(e.to), inq[e.to] = true;
95cf             }
95cf         }
95cf     }
6d7c     return d[t] != INF;
95cf }
427e
71a4 void augment(){
06f1     int u = t;
b19d     while (u != s){
db09         edges[p[u]].flow += a[t];
25a9         edges[p[u]^1].flow -= a[t];
e6c9         u = edges[p[u]].from;
95cf     }
95cf }
427e
6e20 #ifndef GIVEN_FLOW
5972     bool min_cost(int s, int t, int f, LL& cost) {
590d         this->s = s; this->t = t;
21d4         int flow = 0;
23cb         cost = 0;
22dc         while (spfa()) {
bcd8             augment();
a671             if (flow + a[t] >= f){
b14d                 cost += (f - flow) * d[t]; flow = f;
3361                 return true;
8e2e             } else {
2a83                 flow += a[t]; cost += a[t] * d[t];
95cf             }
95cf         }
438e     return false;

```

```

    }
#else
    int min_cost(int s, int t, LL& cost) {
        this->s = s; this->t = t;
        int flow = 0;
        cost = 0;
        while (spfa()) {
            augment();
            flow += a[t]; cost += a[t] * d[t];
        }
        return flow;
    }
#endif
};

```

95cf
a8cb
f9a9
590d
21d4
23cb
22dc
bcd8
2a83
95cf
84fb
95cf
1937
329b

5.9 Global minimum cut (Stoer-Wagner)

Usage:

stoer(w)

Compute the global minimum cut of the graph specified by the **symmetric** adjacent matrix w (0-based). Return the capacity of the cut and the indices of one part of the cut.

Time Complexity: $O(|V|^3)$

```

typedef vector<LL> VI;
typedef vector<VI> VVI;

pair<LL, VI> stoer(VVI &w) {
    int n = w.size();
    VI used(n), c, bestc;
    LL bestw = -1;

    for (int ph = n - 1; ph >= 0; ph--) {
        VI wt = w[0], added = used;
        int prev, last = 0;
        rep (i, ph) {
            prev = last;
            last = -1;
            for (int j = 1; j < n; j++)
                if (!added[j] && (last == -1 || wt[j] > wt[last]))
                    last = j;
            if (i == ph - 1) {
                rep (j, n) w[prev][j] += w[last][j];
            }
        }
    }
}

```

f9d7
045e
427e
f012
66f7
4d98
329d
427e
cd21
ec6e
f20e
4b32
8bfc
0706
4942
c4b9
887d
71bc
9cfa

```

1f25         rep (j, n) w[j][prev] = w[prev][j];
5613         used[last] = true;
8e11         c.push_back(last);
bb8e         if (bestw == -1 || wt[last] < bestw) {
bab6             bestc = c;
372e             bestw = wt[last];
95cf         }
8e2e     } else {
caeb         rep (j, n) wt[j] += w[last][j];
8b92         added[last] = true;
95cf     }
95cf }
95cf }
038c     return {bestw, bestc};
95cf }

```

5.10 Fast LCA

All indices of the tree are 1-based.

Usage:

```

preprocess(root)    Initialize with tree rooted at root.
lca(u, v)           Query the lowest common ancestor of  $u$  and  $v$ .

```

```

0e34 const int MAXN = 500005;
0b32 vector<int> adj[MAXN];
fccb int id[MAXN], nid;
1356 pair<int, int> st[MAXN << 1][33 - __builtin_clz(MAXN)];
427e
e16d void dfs(int u, int p, int d) {
0df2     st[id[u] = nid++][0] = {d, u};
18f6     for (int v : adj[u]) {
bd87         if (v == p) continue;
f58c         dfs(v, u, d + 1);
08ad         st[nid++][0] = {d, u};
95cf     }
95cf }
427e
3d1b void preprocess(int root) {
3269     nid = 0;
91e1     dfs(root, 0, 1);
5e98     int l = 31 - __builtin_clz(nid);
213b     rep (j, l) rep (i, 1+nid-(1<<j))
1131         st[i][j+1] = min(st[i][j], st[i+(1<<j)][j]);

```

```

}

int lca(int u, int v) {
    tie(u, v) = minmax(id[u], id[v]);
    int k = 31 - __builtin_clz(v-u+1);
    return min(st[u][k], st[v-(1<<k)+1][k]).second;
}

```

5.11 Heavy-light decomposition

Time Complexity: The decomposition itself takes linear time. Each query takes $O(\log n)$ operations.

```

const int MAXN = 100005;
vector<int> adj[MAXN];
int sz[MAXN], top[MAXN], fa[MAXN], son[MAXN], depth[MAXN], id[MAXN];

void dfs1(int x, int dep, int par){
    depth[x] = dep;
    sz[x] = 1;
    fa[x] = par;
    int maxn = 0, s = 0;
    for (int c: adj[x]){
        if (c == par) continue;
        dfs1(c, dep + 1, x);
        sz[x] += sz[c];
        if (sz[c] > maxn){
            maxn = sz[c];
            s = c;
        }
    }
    son[x] = s;
}

int cid = 0;
void dfs2(int x, int t){
    top[x] = t;
    id[x] = ++cid;
    if (son[x]) dfs2(son[x], t);
    for (int c: adj[x]){
        if (c == fa[x]) continue;
        if (c == son[x]) continue;
        else dfs2(c, c);
    }
}

```

```

95cf    }
95cf }
427e
0f04 void decomp(int root){
9fa4     dfs1(root, 1, 0);
1c88     dfs2(root, root);
95cf }
427e
2c98 void query(int u, int v){
03a1     while (top[u] != top[v]){
45ec         if (depth[top[u]] < depth[top[v]]) swap(u, v);
427e         // id[top[u]] to id[u]
005b         u = fa[top[u]];
95cf     }
6083     if (depth[u] > depth[v]) swap(u, v);
427e     // id[u] to id[v]
95cf }

```

5.12 Centroid decomposition

Note that the centroid here is not the exact centroid of the graph. It only guarantees that the size of each subtree does not exceed half of that of the original tree. This is enough to guarantee the correct time complexity. All vertices are numbered from 1. Call `decomp(root)` to use.

Usage:

`decomp(u, p)` Decompose the tree rooted at u with parent p .

Time Complexity: The decomposition itself takes $O(n \log n)$ time.

```

1fb6 vector<int> adj[100005];
88e0 int sz[100005], sum;
427e
f93d void getsz(int u, int p) {
5b36     sz[u] = 1; sum++;
18f6     for (int v : adj[u]) {
bd87         if (v == p) continue;
e3cb         getsz(v, u);
8449         sz[u] += sz[v];
95cf     }
95cf }
427e
67f9 int getcent(int u, int p) {
d51f     for (int v : adj[u])
76e4         if (v != p and sz[v] > sum / 2)

```

```

        return getcent(v, u);
    return u;
}

void decompose(int u) {
    sum = 0; getsz(u, 0);
    u = getcent(u, 0); // update u to the centroid

    for (int v : adj[u]) {
        // get answer for subtree v
    }
    // get answer for the whole tree
    // don't forget to count the centroid itself

    for (int v : adj[u]) { // divide and conquer
        adj[v].erase(find(range(adj[v]), u));
        decompose(v);
        adj[v].push_back(u); // restore deleted edge
    }
}

```

5.13 DSU on tree

This implementation avoids parallel existence of multiple data structures but requires that the data structure is invertible. To use this template, implement `merge`, `enter`, `leave` as needed; first call `decomp(root, 0)`, then call `work(root, 0, false)`. Labels of vertices start from 1.

Usage:

`decomp(u, p)` Decompose the tree u .
`work(u, p, keep)` Work for subtree u . When `keep` is set, information is not cleared.

Time Complexity: $O(n \log n)$ times the complexity for `merge`, `enter`, `leave`.

```

vector<int> adj[100005];
int sz[100005], son[100005];

void decomp(int u, int p) {
    sz[u] = 1;
    for (int v : adj[u]) {
        if (v == p) continue;
        decomp(v, u);
        sz[u] += sz[v];
    }
}

```

```

d28c         if (sz[v] > sz[son[u]]) son[u] = v;
95cf     }
95cf }
427e
b7ec template <typename T>
62f5 void trav(T fn, int u, int p) {
4412     fn(u);
30b3     for (int v : adj[u]) if (v != p) trav(fn, v, u);
95cf }
427e
7467 #define for_light(v) for (int v : adj[u]) if (v != p and v != son[u])
33ff void work(int u, int p, bool keep) {
72a2     for_light(v) work(v, u, 0); // process light children
427e
427e     // process heavy child
427e     // current data structure contains info of heavy child
9866     if (son[u]) work(son[u], u, 1);
427e
18a9     auto merge = [u] (int c) { /* count contribution of c */ };
1ab0     auto enter = [] (int c) { /* add vertex c */ };
f241     auto leave = [] (int c) { /* remove vertex c */ };
427e
3d3b     for_light(v) {
74c6         trav(merge, v, u);
c13d         trav(enter, v, u);
95cf     }
427e
427e     // count answer for root and add it
427e     // Warning: special check may apply to root!
c54f     merge(u);
9dec     enter(u);
427e
427e     // Leave current tree
4e3e     if (!keep) trav(leave, u, p);
95cf }

```

6 Data Structures

6.1 Fenwick tree (point update range query)

```

9976 struct bit_purq { // point update, range query

```

```

int N;
vector<LL> tr;

void init(int n) { // fill the array with 0
    tr.resize(N = n + 5);
}

LL sum(int n) {
    LL ans = 0;
    while (n) {
        ans += tr[n];
        n &= n - 1;
    }
    return ans;
}

void add(int n, LL x){
    while (n < N) {
        tr[n] += x;
        n += n & -n;
    }
}
};

```

6.2 Fenwick tree (range update point query)

```

struct bit_rupq{ // range update, point query
    int N;
    vector<LL> tr;

    void init(int n) { // fill the array with 0
        tr.resize(N = n + 5);
    }

    LL query(int n) {
        LL ans = 0;
        while (n < N) {
            ans += tr[n];
            n += n & -n;
        }
        return ans;
    }
}

```

```

d7af
99ff
427e
d34f
1010
95cf
427e
63d0
f7ff
e290
0715
c0d4
95cf
4206
95cf
427e
f4bd
ad20
6c81
0af5
95cf
95cf
329b

```

```

3d03
d7af
99ff
427e
d34f
1010
95cf
427e
38d4
f7ff
ad20
0715
0af5
95cf
4206
95cf

```

```

427e void add(int n, LL x) {
f4bd     while (n){
e290         tr[n] += x;
6c81         n &= n - 1;
c0d4     }
95cf }
95cf }
329b };

```

6.3 Segment tree

```

3942 LL p;
1ebb const int MAXN = 4 * 100006;
451a struct segtree {
27be     int l[MAXN], m[MAXN], r[MAXN];
4510     LL val[MAXN], tadd[MAXN], tmul[MAXN];
427e
ac35 #define lson (o<<1)
1294 #define rson (o<<1|1)
427e
1344 void pull(int o) {
bbe9     val[o] = (val[lson] + val[rson]) % p;
95cf }
427e
e4bc void push_add(int o, LL x) {
5dd6     val[o] = (val[o] + x * (r[o] - l[o])) % p;
6eff     tadd[o] = (tadd[o] + x) % p;
95cf }
427e
d658 void push_mul(int o, LL x) {
b82c     val[o] = val[o] * x % p;
aa86     tadd[o] = tadd[o] * x % p;
649f     tmul[o] = tmul[o] * x % p;
95cf }
427e
b149 void push(int o) {
3159     if (l[o] == m[o]) return;
0a90     if (tmul[o] != 1) {
0f4a         push_mul(lson, tmul[o]);
045e         push_mul(rson, tmul[o]);
ac0a         tmul[o] = 1;
95cf     }

```

```

if (tadd[o]) {
    push_add(lson, tadd[o]);
    push_add(rson, tadd[o]);
    tadd[o] = 0;
}
}

void build(int o, int ll, int rr) {
    int mm = (ll + rr) / 2;
    l[o] = ll; r[o] = rr; m[o] = mm;
    tmul[o] = 1;
    if (ll == mm) {
        scanf("%lld", val + o);
        val[o] %= p;
    } else {
        build(lson, ll, mm);
        build(rson, mm, rr);
        pull(o);
    }
}

void add(int o, int ll, int rr, LL x) {
    if (ll <= l[o] && r[o] <= rr) {
        push_add(o, x);
    } else {
        push(o);
        if (m[o] > ll) add(lson, ll, rr, x);
        if (m[o] < rr) add(rson, ll, rr, x);
        pull(o);
    }
}

void mul(int o, int ll, int rr, LL x) {
    if (ll <= l[o] && r[o] <= rr) {
        push_mul(o, x);
    } else {
        push(o);
        if (ll < m[o]) mul(lson, ll, rr, x);
        if (m[o] < rr) mul(rson, ll, rr, x);
        pull(o);
    }
}

LL query(int o, int ll, int rr) {

```

```

1b82
9547
0e73
6234
95cf
95cf
427e
471c
0e87
9d27
ac0a
5c92
001f
e5b6
8e2e
7293
5e67
ba26
95cf
95cf
427e
4406
3c16
db32
8e2e
c4b0
4305
d5a6
ba26
95cf
95cf
427e
48cd
3c16
e7d0
8e2e
c4b0
d1ba
67f3
ba26
95cf
95cf
427e
0f62

```

```

3c16     if (ll <= l[o] && r[o] <= rr) {
6dfe         return val[o];
8e2e     } else {
c4b0         push(o);
462a         if (rr <= m[o]) return query(lson, ll, rr);
5cca         if (ll >= m[o]) return query(rson, ll, rr);
bbf9         return query(lson, ll, rr) + query(rson, ll, rr);
95cf     }
95cf     }
4d99 } seg;

```

6.4 Treap

Self-balanced binary search tree which supports split and merge.

Usage:

push(x)	Push lazy tags to children.
pull(x)	Update statistics of node x .
Init(x, v)	Initialize node x with value v .
Add(x, v)	Apply addition to subtree x .
Reverse(x)	Apply reversion to subtree x .
Merge(x, y)	Merge trees rooted at x and y . Return the root of new tree.
Split(t, k, x, y)	Split out the left k elements of tree t . The roots of left part and right part are stored in x and y , respectively.
init(n)	Initialize the treap with array of size n .
work(op, l, r)	Range operation over $[l, r)$.

Time Complexity: Expected $O(\log n)$ per operation.

```

9f60 const int MAXN = 200005;
a7c5 mt19937 gen(time(NULL));
9542 struct Treap {
6d61     int ch[MAXN][2];
3948     int sz[MAXN], key[MAXN], val[MAXN];
5d9a     int add[MAXN], rev[MAXN];
2b1b     LL sum[MAXN] = {0};
a773     int maxv[MAXN] = {INT_MIN}, minv[MAXN] = {INT_MAX};
427e
a629     void Init(int x, int v) {
5a00         ch[x][0] = ch[x][1] = 0;
d8cd         key[x] = gen(); val[x] = v; pull(x);
95cf     }
427e
3bf9     void pull(int x) {

```

```

        sz[x] = 1 + sz[ch[x][0]] + sz[ch[x][1]];
        sum[x] = val[x] + sum[ch[x][0]] + sum[ch[x][1]];
        maxv[x] = max({val[x], maxv[ch[x][0]], maxv[ch[x][1]]});
        minv[x] = min({val[x], minv[ch[x][0]], minv[ch[x][1]]});
    }

```

```

void Add(int x, int a) {
    val[x] += a; add[x] += a;
    sum[x] += LL(sz[x]) * a; maxv[x] += a; minv[x] += a;
}

```

```

void Reverse(int x) {
    rev[x] ^= 1;
    swap(ch[x][0], ch[x][1]);
}

```

```

void push(int x) {
    for (int c : ch[x]) if (c) {
        Add(c, add[x]);
        if (rev[x]) Reverse(c);
    }
    add[x] = 0; rev[x] = 0;
}

```

```

int Merge(int x, int y) {
    if (!x || !y) return x | y;
    push(x); push(y);
    if (key[x] > key[y]) {
        ch[x][1] = Merge(ch[x][1], y); pull(x); return x;
    } else {
        ch[y][0] = Merge(x, ch[y][0]); pull(y); return y;
    }
}

```

```

void Split(int t, int k, int &x, int &y) {
    if (t == 0) { x = y = 0; return; }
    push(t);
    if (sz[ch[t][0]] < k) {
        x = t; Split(ch[t][1], k - sz[ch[t][0]] - 1, ch[t][1], y);
    } else {
        y = t; Split(ch[t][0], k, x, ch[t][0]);
    }
    if (x) pull(x); if (y) pull(y);
}

```

e1c3
99f8
94e9
6bb9
95cf
427e
8c8e
a7b1
832a
95cf
427e
aaf6
52c6
7850
95cf
427e
1a53
5fe5
fd76
7a53
95cf
49ee
95cf
427e
9d2c
1b09
cd7e
bffa
a3df
8e2e
bf9e
95cf
95cf
427e
dc7e
6303
f26b
3465
ffd8
8e2e
8a23
95cf
89e3
95cf


```

b1f4 } treap;
427e
24b6 int root;
427e
d34f void init(int n) {
34d7     Rep (i, n) {
7681         int x; scanf("%d", &x);
0ed8         treap.Init(i, x);
bcc8         root = (i == 1) ? 1 : treap.Merge(root, i);
95cf     }
95cf }
427e
d030 void work(int op, int l, int r) {
6639     int tl, tm, tr;
b6c4     treap.Split(root, l, tl, tm);
8de3     treap.Split(tm, r - 1, tm, tr);
3658     if (op == 1) {
c039         int x; scanf("%d", &x); treap.Add(tm, x);
1dcb     } else if (op == 2) {
ae78         treap.Reverse(tm);
581d     } else if (op == 3) {
e092         printf("%lld,%d,%d\n",
867f             treap.sum[tm], treap.minv[tm], treap.maxv[tm]);
95cf     }
6188     root = treap.Merge(treap.Merge(tl, tm), tr);
95cf }

```

6.5 Link/cut tree

Dynamic connectivity of undirected acyclic graph. Support single-vertex update, path aggregation and relative LCA query. Vertices are numbered from 1. Zero initialization is enough except for the statistic information.

Usage:

pull(x)	Update statistics of node x .
Root(u)	Get the root of tree where vertex u is in.
Link(u, v)	Link two unconnected trees.
Cut(u, v)	Cut an existent edge.
Query(u, v)	Path aggregation.
Update(u, x)	Single point modification.
LCA(u, v, root)	Get the lowest common ancestor of u and v in tree rooted at root.

Time Complexity: $O(\log n)$ per operation

```

const int MAXN = 1000005;
struct LCT {
    int fa[MAXN], ch[MAXN][2], val[MAXN], sum[MAXN];
    bool rev[MAXN];

    bool isroot(int x) { return ch[fa[x]][0] == x || ch[fa[x]][1] == x; }
    void pull(int x) { sum[x] = val[x] ^ sum[ch[x][0]] ^ sum[ch[x][1]]; }
    void reverse(int x) { swap(ch[x][0], ch[x][1]); rev[x] ^= 1; }
    void push(int x) {
        if (rev[x]) rep (i, 2) if (ch[x][i]) reverse(ch[x][i]); rev[x] = 0;
    }
    void rotate(int x) {
        int y = fa[x], z = fa[y], k = ch[y][1] == x, w = ch[x][!k];
        if (isroot(y)) ch[z][ch[z][1] == y] = x;
        ch[x][!k] = y; ch[y][k] = w; if (w) fa[w] = y;
        fa[y] = x; fa[x] = z; pull(y);
    }
    void pushall(int x) { if (isroot(x)) pushall(fa[x]); push(x); }
    void splay(int x) {
        int y = x, z = 0;
        for (pushall(y); isroot(x); rotate(x)) {
            y = fa[x]; z = fa[y];
            if (isroot(y)) rotate((ch[y][0] == x) ^ (ch[z][0] == y) ? x : y);
        }
        pull(x);
    }
    void access(int x) {
        int z = x;
        for (int y = 0; x; x = fa[y = x]) { splay(x); ch[x][1] = y; pull(x); }
        splay(z);
    }
    void chroot(int x) { access(x); reverse(x); }
    void split(int x, int y) { chroot(x); access(y); }

    int Root(int x) {
        for (access(x); ch[x][0]; x = ch[x][0]) push(x);
        splay(x); return x;
    }
    void Link(int u, int v) { chroot(u); fa[u] = v; }
    void Cut(int u, int v) { split(u, v); fa[u] = ch[v][0] = 0; pull(v); }
    int Query(int u, int v) { split(u, v); return sum[v]; }
    void Update(int u, int x) { splay(u); val[u] = x; }
    int LCA(int x, int y, int root) {

```

2e73
ca06
6a6d
c6e1
427e
eba3
f19f
1c4d
1a53
89a0
95cf
425f
51af
e1fe
1e6f
6d09
95cf
52c6
f69c
d095
c494
ceef
4449
95cf
78a0
95cf
6229
1548
8854
7afd
95cf
a067
126d
427e
d87a
f4f1
0d77
95cf
9e46
7c10
0691
a999
1f42

```

6cb2     chroot(root); access(x); splay(y);
02e5     while (fa[y]) splay(y = fa[y]);
c218     return y;
95cf     }
329b     };

```

6.6 Balanced binary search tree from pb_ds

```

0475 #include <ext/pb_ds/assoc_container.hpp>
332d using namespace __gnu_pbds;
427e
43a7 tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update>
      rkt;
427e // null_tree_node_update
427e
427e // SAMPLE USAGE
190e rkt.insert(x);           // insert element
05d4 rkt.erase(x);           // erase element
add5 rkt.order_of_key(x);    // obtain the number of elements less than x
b064 rkt.find_by_order(i);    // iterator to i-th (numbered from 0) smallest element
c103 rkt.lower_bound(x);
4ff4 rkt.upper_bound(x);
b19b rkt.join(rkt2);         // merge tree (only if their ranges do not intersect)
cb47 rkt.split(x, rkt2);      // split all elements greater than x to rkt2

```

6.7 Persistent segment tree, range k-th query

```

f1a7 struct node {
2ff6     static int n, pos;
427e
7cec     int value;
70e2     node *left, *right;
427e
20b0     void* operator new(size_t size);
427e
3dc0     static node* Build(int l, int r) {
b6c5         node* a = new node;
ce96         if (r > l + 1) {
181e             int mid = (l + r) / 2;
3ba2             a->left = Build(l, mid);
8aaf             a->right = Build(mid, r);

```

```

} else {
    a->value = 0;
}
return a;
}

```

```

static node* init(int size) {
    n = size;
    pos = 0;
    return Build(0, n);
}

```

```

static int Query(node* lt, node *rt, int l, int r, int k) {
    if (r == l + 1) return l;
    int mid = (l + r) / 2;
    if (rt->left->value - lt->left->value < k) {
        k -= rt->left->value - lt->left->value;
        return Query(lt->right, rt->right, mid, r, k);
    } else {
        return Query(lt->left, rt->left, l, mid, k);
    }
}

```

```

static int query(node* lt, node *rt, int k) {
    return Query(lt, rt, 0, n, k);
}

```

```

node *Inc(int l, int r, int pos) const {
    node* a = new node(*this);
    if (r > l + 1) {
        int mid = (l + r) / 2;
        if (pos < mid)
            a->left = left->Inc(l, mid, pos);
        else
            a->right = right->Inc(mid, r, pos);
    }
    a->value++;
    return a;
}

```

```

node *inc(int index) {
    return Inc(0, n, index);
}
} nodes[8000000];

```

```

8e2e
bfc4
95cf
5ffd
95cf
427e
5a45
2c46
7ee3
be52
95cf
427e
93c0
d30c
181e
cb5a
8edb
2412
8e2e
0119
95cf
95cf
427e
c9ad
9e27
95cf
427e
b19c
5794
ce96
181e
203d
f44a
649a
1024
95cf
2b3e
5ffd
95cf
427e
e80f
c246
95cf
865a

```

```

427e
99ce int node::n, node::pos;
1987 inline void* node::operator new(size_t size) {
bb3c     return nodes + (pos++);
95cf }

```

6.8 Block list

All indices are 0-based. All ranges are left-closed right-open.

Usage:

block::fix()	Apply tags to the current block.
Init(l, r)	Range initializer.
Reverse(l, r)	Reverse the range.
Add(l, r, x)	Add x to the range.
Query(l, r)	Range aggregation.

```

fd9e const int BLOCK = 800;
76b3 typedef vector<int> vi;
427e
a771 struct block {
8fbc     vi data;
e3b5     LL sum; int minv, maxv;
41db     int add; bool rev;
427e
d7eb     block(vi&& vec) : data(move(vec)),
1f0c         sum(accumulate(range(data), 0ll)),
8216         minv(*min_element(range(data))),
527d         maxv(*max_element(range(data))),
6437         add(0), rev(0) { }
427e
b919     void fix() {
0694         if (rev) reverse(range(data));         rev = 0;
0527         if (add) for (int& x : data) x += add;   add = 0;
95cf     }
427e
8bc4     void merge(block& another) {
b895         fix(); another.fix();
f516         vi temp(move(data));
d02c         temp.insert(temp.end(), range(another.data));
88ea         *this = block(move(temp));
95cf     }
427e
42e8     block split(int pos) {

```

```

        fix();
        block result(vi(data.begin() + pos, data.end()));
        data.resize(pos); *this = block(move(data));
        return result;
    }
};

typedef list<block>::iterator lit;

struct blocklist {
    list<block> blk;

    void maintain() {
        lit it = blk.begin();
        while (it != blk.end() && next(it) != blk.end()) {
            lit it2 = it;
            while (next(it2) != blk.end() &&
                    it2->data.size() + next(it2)->data.size() <= BLOCK) {
                it2->merge(*next(it2));
                blk.erase(next(it2));
            }
            ++it;
        }
    }

    lit split(int pos) {
        for (lit it = blk.begin(); ; it++) {
            if (pos == 0) return it;
            while (it->data.size() > pos)
                blk.insert(next(it), it->split(pos));
            pos -= it->data.size();
        }
    }

    void Init(int *l, int *r) {
        for (int *cur = l; cur < r; cur += BLOCK)
            blk.emplace_back(vi(cur, min(cur + BLOCK, r)));
    }

    void Reverse(int l, int r) {
        lit it = split(l), it2 = split(r);
        reverse(it, it2);
        while (it != it2) {

```

```

3e79
ccab
861a
56b0
95cf
329b
427e
2a18
427e
ce14
5540
427e
7b8e
3131
4628
852d
188c
3600
93e1
e1fa
95cf
5771
95cf
95cf
427e
b7b3
2273
5502
8e85
2099
a5a1
427e
95cf
95cf
427e
1c7b
9919
8950
95cf
427e
a22f
997b
dfd0
8f89

```

```

6a06         it->rev ^= 1;
5283         it++;
95cf     }
b204     maintain();
95cf }
427e
3cce void Add(int l, int r, int x) {
997b     lit it = split(l), it2 = split(r);
8f89     while (it != it2) {
e927         it->sum += LL(x) * it->data.size();
03d3         it->minv += x; it->maxv += x;
4511         it->add += x; it++;
95cf     }
b204     maintain();
95cf }
427e
3ad3 void Query(int l, int r) {
997b     lit it = split(l), it2 = split(r);
c33d     LL sum = 0; int minv = INT_MAX, maxv = INT_MIN;
8f89     while (it != it2) {
e472         sum += it->sum;
72c4         minv = min(minv, it->minv);
e1c4         maxv = max(maxv, it->maxv);
5283         it++;
95cf     }
b204     maintain();
8792     printf("%lld_%d_%d\n", sum, minv, maxv);
95cf }
958e } lst;

```

6.9 Persistent block list

Block list that supports persistence. All indices are 0-based. All ranges are left-closed right-open. `std::shared_ptr` is used to ease memory management. One should modify the constructor of `block` to maintain extra information. Here we use this policy that the size of each block does not exceed `BLOCK`, while the sum of sizes of two adjacent blocks does not less than `BLOCK`.

When some operation that breaks block list property, please call `maintain` in time to restore the property.

Usage:

`maintain()` Maintain the block list property.

`split(pos)` Split the block list at position `pos`. Returns an iterator to a block starting at `pos`.

`sum(l, r)` An example function of list traversal between $[l, r)$.

Time Complexity: When `BLOCK` is properly selected, the time complexity is $O(\sqrt{n})$ per operation.

```

constexpr int BLOCK = 800;
typedef vector<int> vi;
typedef shared_ptr<vi> pvi;
typedef shared_ptr<const vi> pcvi;

struct block {
    pcvi data;
    LL sum;

    // add information to maintain
    block(pcvi ptr) :
        data(ptr),
        sum(accumulate(ptr->begin(), ptr->end(), 0ll))
    { }

    void merge(const block& another) {
        pvi temp = make_shared<vi>(data->begin(), data->end());
        temp->insert(temp->end(), another.data->begin(), another.data->end());
        *this = block(temp);
    }

    block split(int pos) {
        block result(make_shared<vi>(data->begin() + pos, data->end()));
        *this = block(make_shared<vi>(data->begin(), data->begin() + pos));
        return result;
    }
};

typedef list<block>::iterator lit;

struct blocklist {
    list<block> blk;

    void maintain() {
        lit it = blk.begin();
        while (it != blk.end() and next(it) != blk.end()) {
            lit it2 = it;

```

```

0b03         while (next(it2) != blk.end() and
029f             it2->data->size() + next(it2)->data->size() <= BLOCK) {
93e1             it2->merge(*next(it2));
e1fa             blk.erase(next(it2));
95cf         }
5771         ++it;
95cf     }
95cf }
427e
b7b3 lit split(int pos) {
2273     for (lit it = blk.begin(); ; it++) {
5502         if (pos == 0) return it;
d480         while (it->data->size() > pos) {
2099             blk.insert(next(it), it->split(pos));
95cf         }
a1c8         pos -= it->data->size();
95cf     }
95cf }
427e
fd38 LL sum(int l, int r) { // traverse
48b4     lit it1 = split(l), it2 = split(r);
ac09     LL res = 0;
9f1d     while (it1 != it2) {
8284         res += it1->sum;
61fd         it1++;
95cf     }
b204     maintain();
244d     return res;
95cf }
329b };

```

6.10 Sparse table, range extremum query

The array is 0-based and the range is closed.

```

db63 const int MAXN = 100007;
b330 int a[MAXN];
69ae int st[MAXN][32 - __builtin_clz(MAXN)];
427e
8041 inline int ext(int x, int y){return x>y?x:y;} // ! max
427e
d34f void init(int n){
ce01     int l = 31 - __builtin_clz(n);

```

```

rep (i, n) st[i][0] = a[i];
rep (j, l)
    rep (i, 1+n-(1<<j))
        st[i][j+1] = ext(st[i][j], st[i+(1<<j)][j]);
}

int rmq(int l, int r){
    int k = 31 - __builtin_clz(r-l+1);
    return ext(st[l][k], st[r-(1<<k)+1][k]);
}

```

7 Geometrics

7.1 2D geometric template

```

#include <bits/stdc++.h>
using namespace std;

typedef int T;
typedef struct pt {
    T x, y;
    T operator , (pt a) { return x*a.x + y*a.y; } // inner product
    T operator * (pt a) { return x*a.y - y*a.x; } // outer product
    pt operator + (pt a) { return {x+a.x, y+a.y}; }
    pt operator - (pt a) { return {x-a.x, y-a.y}; }

    pt operator * (T k) { return {x*k, y*k}; }
    pt operator - () { return {-x, -y}; }
} vec;

typedef pair<pt, pt> seg;

bool ptOnSeg(pt& p, seg& s){
    vec v1 = s.first - p, v2 = s.second - p;
    return (v1, v2) <= 0 && v1 * v2 == 0;
}

// 0 not on segment
// 1 on segment except vertices
// 2 on vertices
int ptOnSeg2(pt& p, seg& s){

```

cf75
b811
6937
082a
95cf
427e
c863
92f5
baa2
95cf

302f
421c
427e
4553
c0ae
7a9d
ffaa
3ec7
221a
8b34
427e
368b
90f4
ba8c
427e
0ea6
427e
8d6e
ce77
de97
95cf
427e
427e
427e
427e
8421

```

ce77     vec v1 = s.first - p, v2 = s.second - p;
70ca     T ip = (v1, v2);
8b14     if (v1 * v2 != 0 || ip > 0) return 0;
0847     return (v1, v2) ? 1 : 2;
95cf }
427e
427e // if two orthogonal rectangles do not touch, return true
72bb inline bool nIntRectRect(seg a, seg b){
f9ac     return min(a.first.x, a.second.x) > max(b.first.x, b.second.x) ||
f486         min(a.first.y, a.second.y) > max(b.first.y, b.second.y) ||
39ce         min(b.first.x, b.second.x) > max(a.first.x, a.second.x) ||
80c7         min(b.first.y, b.second.y) > max(a.first.y, a.second.y);
95cf }
427e
427e // >0 in order
427e // <0 out of order
427e // =0 not standard
7538 inline double rotOrder(vec a, vec b, vec c){return double(a*b)*(b*c);}
427e
31ed inline bool intersect(seg a, seg b){
427e     // ! if (nIntRectRect(a, b)) return false; // if commented, assume that a
        and b are non-collinear
cb52     return rotOrder(b.first-a.first, a.second-a.first, b.second-a.first) >= 0 &&
059e         rotOrder(a.first-b.first, b.second-b.first, a.second-b.first) >= 0;
95cf }
427e
427e // 0 not intersect
427e // 1 standard intersection
427e // 2 vertex-line intersection
427e // 3 vertex-vertex intersection
427e // 4 collinear and have common point(s)
4d19 int intersect2(seg& a, seg& b){
5dc4     if (nIntRectRect(a, b)) return 0;
42c0     vec va = a.second - a.first, vb = b.second - b.first;
2096     double j1 = rotOrder(b.first-a.first, va, b.second-a.first),
72fe         j2 = rotOrder(a.first-b.first, vb, a.second-b.first);
5ac6     if (j1 < 0 || j2 < 0) return 0;
9400     if (j1 != 0 && j2 != 0) return 1;
83db     if (j1 == 0 && j2 == 0){
6b0c         if (va * vb == 0) return 4; else return 3;
fb17     } else return 2;
95cf }
427e
2c68 template <typename Tp = T>

```

```

inline pt getIntersection(pt P, vec v, pt Q, vec w){
    static_assert(is_same<Tp, double>::value, "must_be_double!");
    return P + v * (w*(P-Q)/(v*w));
}

// -1 outside the polygon
// 0 on the border of the polygon
// 1 inside the polygon
int ptOnPoly(pt p, pt* poly, int n){
    int wn = 0;
    for (int i = 0; i < n; i++) {

        T k, d1 = poly[i].y - p.y, d2 = poly[(i+1)%n].y - p.y;
        if (k = (poly[(i+1)%n] - poly[i])*(p - poly[i])){
            if (k > 0 && d1 <= 0 && d2 > 0) wn++;
            if (k < 0 && d2 <= 0 && d1 > 0) wn--;
        } else return 0;
    }
    return wn ? 1 : -1;
}

istream& operator >> (istream& lhs, pt& rhs){
    lhs >> rhs.x >> rhs.y;
    return lhs;
}

istream& operator >> (istream& lhs, seg& rhs){
    lhs >> rhs.first >> rhs.second;
    return lhs;
}

```

```

5894
6850
7c9a
95cf
427e
427e
427e
427e
cbdd
5fb4
1294
427e
3cae
b957
8c40
3c4d
aad3
95cf
0a5f
95cf
427e
d4a3
fa86
331a
95cf
427e
07ae
5cab
331a
95cf

```

8 Appendices

8.1 Primes

8.1.1 First primes

p	$g(p)$	p	$g(p)$	p	$g(p)$	p	$g(p)$	p	$g(p)$
2	1	3	2	5	2	7	3	11	2
13	2	17	3	19	2	23	5	29	2
31	3	37	2	41	6	43	3	47	5
53	2	59	2	61	2	67	2	71	7
73	5	79	3	83	2	89	3	97	5
101	2	103	5	107	2	109	6	113	3
127	3	131	2	137	3	139	2	149	2
151	6	157	5	163	2	167	5	173	2
179	2	181	2	191	19	193	5	197	2
199	3	211	2	223	3	227	2	229	6

8.1.2 Arbitrary length primes

$\lg p$	p	$g(p)$	p	$g(p)$
3	967	5	1031	14
4	9859	2	10273	10
5	96331	10	102931	3
6	958543	6	1031137	5
7	9594539	2	10169651	2
8	96243449	3	103211039	7
9	980483981	2	1042484357	2
10	9858935453	2	10261276009	7
11	95748666809	3	101759940101	2
12	950781833849	3	1012797784423	5
13	9739822952371	7	10037217092377	7
14	96181051140397	5	104974966380359	11
15	981030138360889	13	1029038416465403	2
16	9655206098080843	3	10116299875820773	2
17	97687777921994419	3	101506415998163437	2

8.1.3 $\sim 1 \times 10^9$

p	$g(p)$	p	$g(p)$	p	$g(p)$
954854573	3	967607731	2	973215833	3
975831713	3	978949117	2	980766497	3
983879921	3	985918807	3	986608921	29
991136977	5	991752599	13	997137961	11
1003911991	3	1009775293	2	1012423549	6
1021000537	5	1023976897	7	1024153643	2
1037027287	3	1038812881	11	1044754639	3
1045125617	3	1047411427	3	1047753349	6

8.1.4 $\sim 1 \times 10^{18}$

p	$g(p)$	p	$g(p)$
951970612352230049	3	963284339889659609	3
967495386904694119	3	969751761517096213	2
983238274281901499	2	984647442475101409	23
989286107138674069	11	1002507954383424641	3
1006658951440146419	2	1020152326159075903	3
1034876265966119449	7	1042753851435034019	2
1043609016597371563	2	1045571042176595707	2
1048364250160580293	2	1049495624119026949	2

8.2 Pell's equation

$x^2 - ny^2 = 1$, where n is a positive nonsquare integer.

Let (x_0, y_0) be the smallest positive solution of the equation, then the k -th solution is:

$$\begin{pmatrix} x_k \\ y_k \end{pmatrix} = \begin{pmatrix} x_0 & ny_0 \\ y_0 & x_0 \end{pmatrix}^k \begin{pmatrix} x_0 \\ y_0 \end{pmatrix}$$

Some smallest solutions to Pell's equation:

n	2	3	5	6	7	8	10	11	12	13	14	15	17	18	19	20
x	3	2	9	5	8	3	19	10	7	649	15	4	33	17	170	9
y	2	1	4	2	3	1	6	3	2	180	4	1	8	4	39	2

8.3 Burnside's lemma and Polya's enumeration theorem

The Burnside's lemma says that

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

where G is a group acting on X , X^g is the set of elements in X that are fixed by g , i.e. $X^g = \{x \in X : gx = x\}$.

The unweighted version of Pólya enumeration theorem says that

$$|Y^X/G| = \frac{1}{|G|} \sum_{g \in G} m^{c_g}$$

where $m = |X|$ is the number of colors, c_g is the number of the cycles of permutation g .

8.4 Lagrange's interpolation

For sample points $(x_0, y_0), \dots, (x_k, y_k)$, define

$$l_j(x) = \prod_{0 \leq m \leq k, m \neq j} \frac{x - x_m}{x_j - x_m}$$

then the Lagrange polynomial is

$$L(x) = \sum_{j=0}^k y_j l_j(x).$$

To use the script below, type two lines

```
x0 x1 x2 ... xn
x0 x1 x2 ... xn
```

the script will print the fractional coefficient of the polynomial in ascending exponent order.

```
#!/usr/bin/python2
from fractions import *

def polyadd(a, b) : return map(lambda x, y : (x or 0) + (y or 0), a, b)

def polymul(a, b) :
    p = [0] * (len(a)+len(b)-1)
    for e1, c1 in enumerate(a) :
        for e2, c2 in enumerate(b) :
            p[e1 + e2] += c1 * c2
    return p

x, y = [map(Fraction, raw_input().split()) for _ in 0,0]
n = len(x)
lj = [reduce(polymul, [[-x[m]/(x[j]-x[m]), 1/(x[j]-x[m])]
    for m in range(n) if m != j]) for j in range(n)]
print ' '.join(map(str, reduce(polyadd,
    map(lambda a, b : [x * a for x in b], y, lj))))
```

6dc9
4b2b
427e
bbbe
427e
796b
83e4
f697
156c
dfce
5849
427e
f06d
e80a
a649
9dfa
46f9
d754