

Topics in Number Theory

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Introduction

Sug Woo Shin taught a course (18.787) on Topics in Number Theory at MIT in Fall 2012. These are my “live- \TeX ed” notes from the course. The template is borrowed from Akhil Mathew.

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Lecture 1

Thu. 9/6/12

Course website: <http://math.mit.edu/~swshin/Fall12-18787>

§1 Overview

In this course, we will cover abelian varieties and p -divisible groups, also known as Barsotti-Tate groups. We first build some basic knowledge and apply it to some interesting problems in number theory. Our main reference is Abelian Varieties, by Mumford. We will

1. *classify* abelian varieties over finite fields \mathbb{F}_p and algebraic closures of finite fields $\overline{\mathbb{F}_p}$ (Honda-Tate Theory). We will also classify p -divisible groups up to isogeny (Dieudonné, Manin).

With some more work, we can get classification up to isomorphism.

Studying a variety over finite fields helps us understand abelian varieties over global fields, because when we study a global problem, one way to get information is to reduce modulo a prime and study over the variety over the special fiber.

2. go from characteristic p (\mathbb{F}_p) to characteristic 0 (e.g. \mathbb{Q}_p) using *deformations*.
 - The Serre-Tate Theorem will tell us that deformations of abelian varieties are basically deformations of p -divisible groups.
 - The theory of Grothendieck-Messing will reduce deformations of p -divisible groups to some linear algebra.

To understand abelian varieties and p -divisible groups, we first need to understand group schemes. An abelian variety is a special type of group scheme, while a p -divisible group is an inductive limit of group schemes.

§2 Review: Yoneda Lemma and T -valued points

This is not part of the lecture. I include this section as a reference.

2.1 The Yoneda Lemma

Lemma 1 (Yoneda Lemma): lem:yoneda Let \mathcal{C} be a locally small category. Let h_A denote the functor $\text{Hom}(\bullet, A) : \mathcal{C} \rightarrow (\text{Sets})$ and h^A denote the contravariant functor $\text{Hom}(A, \bullet)$ (i.e. it is a functor $\mathcal{C}^{\text{op}} \rightarrow (\text{Sets})$).

1. (Covariant version) Let F be functor from \mathcal{C} to (Sets) . As functors $(\text{Set})^{\mathcal{C}} \times \mathcal{C} \rightarrow (\text{Set})$, we have $\text{Nat}(h^A, F) \cong F(A)$. (F is in $(\text{Set})^{\mathcal{C}}$, A is in \mathcal{C} , and $\text{Nat}(h^A, F) \cong F(A)$ is a set.)

2. (Contravariant version) Let F be a contravariant functor from \mathcal{C} to (Sets). As functors $(\text{Set})^{\mathcal{C}^{\text{op}}} \times \mathcal{C} \rightarrow (\text{Set})$, we have $\text{Nat}(h_A, F) \cong F(A)$.

Corollary 2 (Yoneda Embedding): cor:yoneda

1. The embedding $h^\bullet : \mathcal{C}^{\text{op}} \rightarrow (\text{Set})^{\mathcal{C}}$ given by sending $A \mapsto h^A = \text{Hom}_{\mathcal{C}}(A, \bullet)$ is fully faithful. (The morphism $f : A \rightarrow B$ gets sent to $f \circ \bullet$.)
2. The embedding $h_\bullet : \mathcal{C} \rightarrow (\text{Set})^{\mathcal{C}^{\text{op}}}$ given by sending $A \mapsto h_A = \text{Hom}_{\mathcal{C}}(\bullet, A)$ is fully faithful. (The morphism $f : A \rightarrow B$ gets sent to $\bullet \circ f$.)

Remark: • A category is **locally small** if homomorphisms between any two objects form a set.

- $(\text{Set})^{\mathcal{C}^{\text{op}}}$ is the category of *contravariant functors* $\mathcal{C} \rightarrow (\text{Set})$.
- $\text{Hom}(A, B)$ has just the structure of a set.
- $\text{Nat}(G, F)$ denotes the set of natural transformations between G and F .
- A functor Φ is **fully faithful** if $\Phi_{A,B} : \text{Hom}(A, B) \rightarrow \text{Hom}(\Phi(A), \Phi(B))$ is bijective for any objects A and B . This basically means that Φ embeds the first category into the second, and there aren't any "extra" maps between embedded objects that are present in B but not A .
- We say a functor $F : \mathcal{C} \rightarrow (\text{Set})$ is **representable** if $F \cong h^A$ for some A (and ditto for the contravariant case).

Proof of Corollary 1. We show (2) of the lemma implies (2) of the corollary; (1) is entirely analogous. Set $F = h_B$ to get

$$\text{Nat}(h_A, h_B) \cong h_B(A).$$

Now a natural transformation is just a morphism in the functor category, so $\text{Nat}(h_A, h_B) = \text{Hom}_{(\text{Set})^{\mathcal{C}^{\text{op}}}}(h_A, h_B)$, and by definition $h_B(A) = \text{Hom}(A, B)$, so we get

$$\text{Hom}_{(\text{Set})^{\mathcal{C}^{\text{op}}}}(h_A, h_B) \cong \text{Hom}(A, B).$$

This is exactly the condition to be fully faithful. □

One way to think of this is that an object is determined by how other objects map into it.¹

¹As mentioned here <http://mathoverflow.net/questions/3184/philosophical-meaning-of-the-yoneda-lemma/3223#3223>, if one thinks of objects of a category as particles and morphisms as ways to smash one particle into another particle, then the Yoneda lemma says that it is possible to determine the identity of a particle by smashing other particles into it.

2.2 T -valued points

Definition 3: Let X and T be objects in a locally small category. Define the set of T -valued points of X to be

$$T(X) := \text{Hom}(T, X).$$

In many cases we can think of “ T -valued points” as a generalization of “points” of X . For example, suppose T is a singleton set $\{\cdot\}$ and X is a set, then a T -valued point is just a point of X .

The main application to algebraic geometry can be seen through the following example.

Example 4: ex:T-points Let R be an integral domain and V a variety over R . Let $T = \text{Spec}(R)$ and X be the scheme corresponding to V . Then the T -points of X are exactly the points of V .

To see this, it’s sufficient just to consider the affine case. Suppose $V \in R^n$ is defined by f_1, \dots, f_m . By Lemma 5, to give a morphism

$$T = \text{Spec}(R) \rightarrow X = \text{Spec} \left(\frac{R[x_1, \dots, x_n]}{(f_1, \dots, f_m)} \right)$$

is the same as giving a R -algebra homomorphism

$$\frac{R[x_1, \dots, x_n]}{(f_1, \dots, f_m)} \rightarrow R,$$

which is just an assignment

$$(x_1, \dots, x_n) \mapsto (a_1, \dots, a_n) \text{ such that } f_i(x_1, \dots, x_n) = 0 \text{ for some } i,$$

i.e. a point of V .

Lemma 5 (cf. Hartshorne, II, Exercise 2.4): **lem:spec-fff** Let R be a ring. Then Spec is a fully faithful contravariant functor from the category of R -algebras to schemes over $\text{Spec}(R)$.

Example 4 is the most intuitive example. However, the power of the viewpoint of $X(T)$ is that we can consider more generalized points. For instance, letting R be a field k ,

- a $\text{Spec}(k[t])$ point is a one-parameter family of k -points, and
- a $\text{Spec} \left(\frac{k[t]}{(t^2)} \right)$ point is a k -point with a Zariski tangent vector.



The Yoneda Embedding tells us that we can identify a scheme X with the **functor of points** $h_X(\bullet) = X(\bullet)$ —i.e. with $X(T)$, the T -points of X , as T ranges over all schemes—without losing any information. A functor $X \rightarrow Y$ becomes a natural transformation $h_X = X(\bullet) \rightarrow h_Y = Y(\bullet)$, i.e. maps of sets $X(T) \rightarrow Y(T)$ for each T , that are functorial over T .

§3 Group schemes

3.1 Definition of group schemes

We will define group schemes over a fixed scheme S .

Definition 6: Let S be a scheme. Define (Sch/S) , the category of S -schemes, as follows.

- The objects are schemes T with a structure map to S , $T \rightarrow S$.

$$\begin{array}{c} T \\ \downarrow \\ S \end{array}$$

- The morphisms are

$$\text{Hom} \left(\begin{array}{c} T \quad T' \\ \downarrow f \quad \downarrow f' \\ S \quad S \end{array} \right) = \left\{ g : \begin{array}{ccc} T & \xrightarrow{g} & T' \\ & f \searrow & \swarrow f' \\ & S & \end{array} \text{ commutes} \right\}$$

called S -morphisms.

For short we'll write $\text{Hom}_S(T, T')$, the maps f, f' being implicit.

We now apply the philosophy of the previous section: to study X we study $h_X = X(\bullet)$. If $X \in (\text{Sch}/S)$ we get canonically

$$\begin{aligned} h_X : (\text{Sch}/S) &\rightarrow (\text{Sets}) \\ T &\mapsto \text{Hom}_S(T, X). \end{aligned}$$

Since T and X are S -schemes, we define the T -points of X to be $X(T) := \text{Hom}_S(T, X)$. The functor h_X sends

$$(T \xrightarrow{f} T') \mapsto (\text{Hom}_S(T, X) \xleftarrow{h_X(f) = \bullet \circ f} \text{Hom}_S(T', X)).$$

The Yoneda Embedding 2 tells us that h_\bullet is a fully faithful contravariant functor

$$\begin{aligned} h_\bullet : (\text{Sch}/S) &\rightarrow \text{Fun}^{\text{op}}((\text{Sch}/S), (\text{Sets})) = (\text{Sets})^{(\text{Sch}/S)^{\text{op}}} \\ X &\mapsto h_X. \end{aligned}$$

We say $h \in \text{Fun}^{\text{op}}((\text{Sch}/S), (\text{Sets}))$ is **representable** (by the scheme X) if $h \cong h_X$.

We have several equivalent definitions for a group scheme. The Yoneda Embedding gives the equivalence of the 2nd and 3rd definitions.

Definition 7: A **group scheme** G over S' any of the following three equivalent objects.

1. a group object in (Sch/S) , i.e. it is (G, \widetilde{h}_G) where $G \in (\text{Sch}/S)$ and the following commutes:

$$\begin{array}{ccc} (\text{Sch}/S) & \xrightarrow{\widetilde{h}_G} & (\text{Gps}) \\ & \searrow h_G \quad \swarrow \text{forgetful} & \\ & (\text{Set}) & \end{array}$$

2. (G, h_G) equipped with the following maps of sets

- e_T (identity): $\{\cdot\} \rightarrow G(T)$
- i_T (inverse): $G(T) \rightarrow G(T)$
- m_T (multiplication): $G(T) \times G(T) \rightarrow G(T)$.

such that $G(T)$ is a *group* under these operations, namely,

- (a) (Associativity) The following commutes:

$$\begin{array}{ccc} G(T) \times G(T) \times G(T) & \xrightarrow{(m_T, \text{id})} & G(T) \times G(T) \\ \downarrow (\text{id}, m_T) & & \downarrow m_T \\ G(T) \times G(T) & \xrightarrow{m_T} & G(T). \end{array}$$

Note: This represents associativity because going clockwise we get $(xy)z$ and going counterclockwise we get $x(yz)$.

- (b) (Inverse)

$$\begin{array}{ccccc} & & G(T) \times G(T) & & \\ & \nearrow (\text{id}, i_T) & & \searrow m_T & \\ G(T) & \xrightarrow{\text{structure}} & G(S) & \xrightarrow{e_T} & G(T) \\ & \searrow (i_T, \text{id}) & & \nearrow m_T & \\ & & G(T) \times G(T) & & \end{array}$$

Note: The top, middle, and bottom give xx^{-1} , e , and $x^{-1}x$, respectively, so commutativity gives $xx^{-1} = e = x^{-1}x$.

- (c) (Identity) Let $e'_T : G(T) \rightarrow G(T)$ be the composition of the structure map with $e : G(S) \rightarrow G(T)$.

$$\begin{array}{ccccc} & & G(T) \times G(T) & & \\ & \nearrow (\text{id}, e'_T) & & \searrow m_T & \\ G(T) & \xrightarrow{\text{id}} & G(T) & & G(T) \\ & \searrow (e'_T, \text{id}) & & \nearrow m_T & \\ & & G(T) \times G(T) & & \end{array}$$

Note: This gives $x \cdot e = x = e \cdot x$.

and these group operations are *functorial*, namely for all $T \xrightarrow{f} T'$ in (Sch/S) ,

•

$$\begin{array}{ccc} \{\cdot\} & \xrightarrow{e_T} & G(T) \\ & \searrow e_{T'} & \uparrow h_G(f) \\ & & G(T') \end{array}$$

•

$$\begin{array}{ccc} G(T) & \xrightarrow{i_T} & G(T) \\ h_G(f) \uparrow & & \uparrow h_G(f) \\ G(T') & \xrightarrow{i_{T'}} & G(T') \end{array}$$

•

$$\begin{array}{ccc} G(T) \times G(T) & \xrightarrow{i_T} & G(T) \\ h_G(f) \uparrow & & \uparrow h_G(f) \\ G(T') \times G(T') & \xrightarrow{i_{T'}} & G(T') \end{array}$$

3. (G, e, i, m) where $G \in (\text{Sch}/S)$,

$$\begin{aligned} e &: S \rightarrow G \\ i &: G \rightarrow G \\ m &: G \times G \rightarrow G \end{aligned}$$

and we have the analogues of the group laws in the 2nd definition, but with fiber product instead of product and with e, i, m instead of e_T, i_T, m_T .

(a) (Associativity)

$$\begin{array}{ccc} G \times_S G \times_S G & \xrightarrow{(m, \text{id})} & G \times_S G \\ \downarrow (\text{id}, m) & & \downarrow m \\ G \times_S G & \xrightarrow{m} & G. \end{array}$$

(b) (Inverse)

$$\begin{array}{ccccc} & & G \times_S G & & \\ & \nearrow (\text{id}, i) & & \searrow m & \\ G & \xrightarrow{\text{structure}} & S & \xrightarrow{e} & G \\ & \searrow (i, \text{id}) & & \nearrow m & \\ & & G \times_S G & & \end{array}$$

(c) (Identity) Let $e' : G \rightarrow G$ be the composition of the structure map with $e : S \rightarrow G$.

$$\begin{array}{ccc}
 & G \times_S G & \\
 (\text{id}, e') \nearrow & & \searrow m \\
 G & \xrightarrow{\text{id}} & G(T) \\
 (e', \text{id}) \searrow & & \nearrow m \\
 & G \times_S G &
 \end{array}$$

Proof of equivalence. The 1st and 2nd definition are equivalent: In the 2nd definition, the first set of conditions simply say $G(T)$ is a group, and the second set of conditions say that \widetilde{h}_G is a functor; i.e. it sends the scheme morphism f to a group homomorphism $\widetilde{h}_G(f)$.

The 2nd and 3rd definitions are equivalent: We go between G to $G(T)$ by the Yoneda embedding. h_G sends fiber products of schemes to products of sets. \square



We can understand group schemes as *schemes* with group axioms on schemes, or as *functors of points* with group axioms on the set of T -points for each T .

3.2 Examples of group schemes

Let $G = \text{Spec } A$ and $S = \text{Spec } R$. Suppose A is an R -algebra, so there is a natural structure map $G \rightarrow S$. We have by Lemma 5 that Spec is a contravariant fully faithful functor from $(R\text{-algebras})$ to $(\text{Sch}/\text{Spec } R)$:

$$\begin{array}{ccc}
 (\text{rings}) & \xrightarrow{\text{Spec}} & (\text{Sch}) \\
 \uparrow & & \uparrow \\
 (R\text{-algebras}) & \xrightarrow[\text{Spec}]{\text{f.f.}} & (\text{Sch}/\text{Spec } R)
 \end{array}$$

(Note the categories on the bottom are not full subcategories of the top.) As Lemma 5 says, S -morphisms between schemes over $\text{Spec } R$ are nothing but R -algebra homomorphisms in the opposite direction, so we can be more concrete. So giving $G = \text{Spec } A$ a group scheme structure, i.e. giving e, i, m for G , amounts to giving R -algebra maps (note $\text{Spec}(A \otimes_R A) = \text{Spec } A \times_{\text{Spec } R} \text{Spec } A$)

$$\begin{aligned}
 e &: A \rightarrow R \\
 i &: A \rightarrow A \\
 m &: A \rightarrow A \otimes_R A.
 \end{aligned}$$

such that R -algebra version of (a), (b), and (c) hold. (Just invert all the arrows in (a), (b), and (c), and replace the rings with schemes. A satisfying these axioms is called a **Hopf algebra**.)

We can give some common, concrete examples of group varieties.

Example 8: Define the additive group scheme $\mathbb{G}_{a, \text{Spec } R}$ as follows. (First we consider the 3rd definition.) Let $A = R[t]$ and let $\mathbb{G}_{a, \text{Spec } R} = \text{Spec } A$ be the scheme with e, i , and m induced by the R -algebra homomorphisms

$$\begin{aligned} e : R[t] &\rightarrow R & f &\mapsto f(0) \\ i : R[t] &\rightarrow R[t] & f &\mapsto f(-t) \\ m : R[t] &\rightarrow R[t'] \otimes_R R[t''] \cong R[t', t''] & f &\mapsto f(t' + t'') \end{aligned}$$

For instance, for $R = k$, on points the group operation is just addition. Indeed, the map m gives $\text{Spec } R[t', t''] \rightarrow \text{Spec } R[t]$ that sends the ideal $(t' - a, t'' - b)$ to $m^{-1}((t' - a, t'' - b)) = (t - (a + b))$, i.e. sends the point (a, b) to the point $a + b$.

Now consider $\mathbb{G}_{a, \text{Spec } R}(\text{Spec } R')$ where R' is a R -algebra. Using the 2nd definition, $\mathbb{G}_{a, \text{Spec } R}(\text{Spec } R')$ consists of maps $\text{Spec } R' \rightarrow \text{Spec } R[t]$ —i.e. maps $R[t] \rightarrow R'$, which together with the group axioms, means

$$\mathbb{G}_{a, \text{Spec } R}(\text{Spec } R') = (R', +).$$

(Check this.)

Example 9: Define the multiplicative group scheme $\mathbb{G}_{m, \text{Spec } R}$ as follows. Let $A = R[t, t^{-1}]$; let $\mathbb{G}_{m, \text{Spec } R}$ be the scheme with e, i, m induced by the R -algebra homomorphisms

$$\begin{aligned} e : f &\mapsto f(1) \\ i : f &\mapsto f(t^{-1}) \\ m : f &\mapsto f(t't''). \end{aligned}$$

(Note 1 is the multiplicative identity so we look at $f(1)$ not $f(0)$.)

From a different angle, we get

$$\mathbb{G}_{m, \text{Spec } R}(\text{Spec } R') = (R'^{\times}, \cdot).$$

(When we consider $\text{Spec } R' \rightarrow \text{Spec } R[t, t^{-1}]$, i.e. maps $R[t, t^{-1}] \rightarrow R'$, the image of t must be an invertible element.)

Remark: For any ring R , $\mathbb{G}_{a, \text{Spec } R} \cong \mathbb{G}_{a, \text{Spec } \mathbb{Z}} \times_{\text{Spec } \mathbb{Z}} \text{Spec } R$, by defining an isomorphism on the level of points. The same is true for $\mathbb{G}_{m, \text{Spec } R}$.

Example 10: To define the additive and multiplicative group schemes for general S , we need to use relative Spec.

$$\begin{aligned} \mathbb{G}_{a, S} &:= \underline{\text{Spec}}(\mathcal{O}_S[t]) \\ \mathbb{G}_{m, S} &:= \underline{\text{Spec}}(\mathcal{O}_S[t, t^{-1}]). \end{aligned}$$

with e, i , and m defined similarly. (See Hartshorne II.5 for review on \mathcal{O}_S and http://en.wikipedia.org/wiki/Spectrum_of_a_ring#Global_Spec for review on relative spec. $\mathcal{O}_S[t]$

means replace the $\text{Spec } A(U)$ in the wikipedia definition by $\text{Spec } A(U)[t]$, and likewise for $\mathcal{O}_S[t, t^{-1}]$. We are basically cover \mathcal{O}_S by affine schemes, constructing a polynomial algebra over each affine scheme, and patching them together.)

Define

$$\begin{aligned}\text{GL}_{n,S}(T) &= \text{GL}_n(\mathcal{O}_T(T)) \\ &= M_n(\mathcal{O}_T(T))^\times\end{aligned}$$

Taking $n = 1$, we recover the multiplicative group scheme: $\text{GL}_{1,S} = \mathbb{G}_{m,S}$.

Example 11: Define

$$\mu_{n,S} = \underline{\text{Spec}} \mathcal{O}_S[t]/(t^n - 1).$$

Here e, i, m are the same as for $\mathbb{G}_{m,S}$. Alternatively,

$$\mu_{n,S}(T) = \{x \in \mathcal{O}_T(T) : x^n = 1\}.$$

(The image of t should satisfy $t^n = 1$.)

Example 12: ? Define the constant group scheme as follows: let H be an absolute group. Define

$$\underline{H}(T) = \text{Hom}(\pi_0(T), H) = \text{Hom}_{\text{cont}}(T, H)$$

where H is given the discrete topology in the last expression. ($\pi_0(T)$ means connected components of T .) Note $T \rightarrow T'$ gives $\underline{H}(T') \rightarrow \underline{H}(T)$. This gives us a group scheme.

Example 13: Let Γ be an abstract commutative group, and

$$G = \underline{\text{Spec}} \underbrace{\mathcal{O}_S[\Gamma]}_{\text{group algebra}}.$$

(If S is an affine scheme, we just get the group algebra.) Here

$$\mathcal{O}_S[\Gamma] = \bigoplus_{\gamma \in \Gamma} \mathcal{O}_S \cdot \gamma.$$

Define

$$\begin{array}{ll} e : \mathcal{O}_S[\Gamma] \rightarrow \mathcal{O}_S & \gamma \mapsto 1 \\ i : \mathcal{O}_S[\Gamma] \rightarrow \mathcal{O}_S[\Gamma] & \gamma \mapsto \gamma^{-1} \\ i : \mathcal{O}_S[\Gamma] \rightarrow \mathcal{O}_S[\Gamma] \otimes \mathcal{O}_S[\Gamma] & \gamma \mapsto \gamma \otimes \gamma. \end{array}$$

Problem 1: Check that this is a group scheme.

Check that if $\Gamma = \mathbb{Z}/n\mathbb{Z}$ we get μ_n , and if $\Gamma = \mathbb{Z}$ we get \mathbb{G}_m .

3.3 Morphisms between group scheme

The natural next step is to define a notion of morphisms between group schemes. As we've said, the objects of (Gp/S) to be the group schemes over S . The morphisms are

$$\begin{aligned} & \text{Hom}_{(\text{Gp}/S)}(G, G') \\ &:= \text{Hom}(\widetilde{h}_G, \widetilde{h}_{G'}) \text{ in } \text{Fun}((\text{Sch}/S), (\text{Gp})) \\ &= \left\{ \begin{array}{ccc} G & \xrightarrow{\quad} & G' \\ & \searrow \quad \swarrow & \\ & S & \end{array} : G(T) \rightarrow G'(T) \text{ is a group homomorphism, for every } T \in (\text{Sch}/S) \right\} \end{aligned}$$

Definition 14: A **subgroup scheme** is a subscheme $H \subseteq G$ such that $H(T) \subseteq G(T)$ (subgroup) for all $T \in (\text{Sch}/S)$. Equivalently, a subgroup scheme is (H, e_H, i_H, m_H) such that $H \subseteq G$, and the following commute:

$$\begin{array}{ccccc} S & \xrightarrow{e_H} & H & & H & \xrightarrow{i_H} & H & & H \times H & \xrightarrow{m_H} & H \\ & \searrow e_G & \downarrow & & \downarrow & & \downarrow & & \downarrow & & \downarrow \\ & & G & & G & \xrightarrow{i_G} & G & & G \times G & \xrightarrow{m_G} & G \end{array}$$

We want to define kernels and cokernels. Cokernels are more difficult; let's do kernels first.

Definition 15: Let $G \xrightarrow{f} H$ be in (Gp/S) . Define the kernel K/S to be the functor $K(\bullet)$ such that

$$K(T) = \ker(G(T) \xrightarrow{f(T)} H(T))$$

for all T/S .

Proof of well-definedness. It's not obvious that this functor is represented by a scheme! So let's call the functor $F(T) := \ker(G(T) \xrightarrow{f(T)} H(T))$ for now; we have to show there exists a scheme K such that $K(T) = F(T)$, i.e. we need to show F is represented by a scheme K . We do this by constructing K . Define $K := G \times_H S$, so we have the following diagram.

$$\begin{array}{ccc} K & \longrightarrow & G \\ \downarrow & & \downarrow f \\ S & \xrightarrow{e_H} & H \end{array}$$

Now take T -points and check that

$$K(T) = G(T) \times_{H(T)} S(T) = \{g \in G(T) : f(g) = 1_{H(T)}\} = \ker f(T).$$

The first equality is from definition of the fiber product (Here, $\times_{H(T)}$ denotes the set-theoretic pullback). \square

Quotients are hard; we'll get to them later.