

# Practical Optimization Algorithms and Applications

## Chapter I: Introduction

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# General Information about the Course

- Prerequisites:
  - Some knowledge of linear algebra;
  - The standard sequence of calculus courses;
- Books:
  - J. Nocedal and S. Wright. *Numerical Optimization* SECOND EDITION. Springer, New York, USA., 2006.
  - W-Y. Sun and Y-X. Yuan. *Optimization Theory and Methods: Nonlinear Programming*. Springer, New York, USA., 2006.
- Additional Online Recourse:
  - <http://www.mcs.anl.gov/otc/Guide/>
  - <http://www.mcs.anl.gov/otc/Guide/SoftwareGuide>
  - <http://neos-server.org/neos/>

# What is Optimization?

**Optimization** is the minimization or maximization of a function subject to constraints on its variables.

- It traces the roots to the calculus of variations and the work of Euler and Lagrange.
- It is often called **mathematical programming**, a somewhat confusing term coined in the 1940s, before the word “programming” became inextricably linked with computer software.

We will use the following notation in this course:

- $x$  is the vector of **variables**, also called *unknowns* or *parameters*;
- $f$  is the **objective function**, a (scalar) function of  $x$  that we want to maximize or minimize;
- $c_i$  are **constraint** functions, which are scalar functions of  $x$  that define certain equations and inequations that the unknown vector  $x$  must satisfy.

# Mathematics Formulation

Using the notations, the optimization problem can be written as the following **Standardized Formulation**:

$$\min \quad f(x) \tag{1a}$$

$$\text{s.t.} \quad c_i(x) = 0, \quad i \in \mathcal{E} = \{1, \dots, m_e\} \tag{1b}$$

$$c_i(x) \geq 0, \quad i \in \mathcal{I} = \{m_e + 1, \dots, m\} \tag{1c}$$

Here “s.t.” means “subject to”,  $\mathcal{E}$  and  $\mathcal{I}$  are the sets of indices for equality and inequality constraints, respectively.

- If  $\mathcal{E} = \mathcal{I} = \emptyset$ , (1) is called **unconstrained optimization**;
- If  $\mathcal{E} \neq \emptyset$  or  $\mathcal{I} \neq \emptyset$ , (1) is called **constrained optimization**.

## Example: A Transportation Problem

A chemical company has 2 factories  $F_1$  and  $F_2$  and a dozen retail outlets  $R_1, R_2, \dots, R_{12}$ .

- Each factory  $F_i$  can produce  $a_i$  tons of a certain chemical product each week;  $a_i$  is called the *capacity* of the plant.
- Each retail outlet  $R_j$  has a known weekly *demand* of  $b_j$  tons of the product.
- The cost of shipping one tone of the product from factory  $F_i$  to retail outlet  $R_j$  is  $c_{ij}$

The problem is to determine how much of the product to ship from each factory to each outlet so as to satisfy all the requirements and minimize cost.

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$$\begin{aligned} \min \quad & \sum_{i,j=1}^{i=2,j=12} c_{ij}x_{ij} \\ \text{s.t.} \quad & \sum_{j=1}^{12} x_{ij} \leq a_i, \quad i = 1, 2, \\ & \sum_{i=1}^2 x_{ij} \geq b_j, \quad j = 1, \dots, 12, \\ & x_{ij} \geq 0, \quad i = 1, 2, j = 1, \dots, 12. \end{aligned}$$



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Because the objective and constraints are all linear function. The above problem is called **linear programming**.

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Suppose there are  $p$  factories instead of 2, i.e.  $F_1 \cdots F_p$ , and each factory has  $q$  outlets  $R_1, R_2, \dots, R_q$ ,  $p$  and  $q$  are positive integer. Then the model becomes

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$$\min \sum_{i,j=1}^{i=p,j=q} c_{ij} x_{ij} \quad (2a)$$

$$\text{s.t.} \quad \sum_{j=1}^q x_{ij} \leq a_i, \quad i = 1, \dots, p, \quad (2b)$$

$$\sum_{i=1}^p x_{ij} \geq b_j, \quad j = 1, \dots, q, \quad (2c)$$

$$x_{ij} \geq 0, \quad i = 1, \dots, p, j = 1, \dots, q. \quad (2d)$$

## Example: A Transportation Problem

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Because the objective is described as a nonlinear function, the above problem is called a **nonlinear programming**.

Thanks for your attention!