

Theorem 11.5

The class H_{pm} of hash functions defined by equations (11.3) and (11.4) is universal.

Proof

Consider two distinct keys k and l from Z_p , so that $k \neq l$. Without loss of generality, assume $k > l$.

For a given hash function h_{ab} we let

$$\begin{aligned} r &= (ak + b) \bmod p, \\ s &= (al + b) \bmod p. \end{aligned}$$

We first prove that $r \neq s$. Assume

$$r = s,$$

We have

$$r - s \equiv (ak + b) - (al + b) \equiv a(k - l) \equiv 0 \pmod{p}.$$

Then

$$p \mid a(k - l)$$

On the other hand, because $a \in Z_p^*$, we get $p \nmid a$. Because $(k - l) \in Z_p$, we get $p \nmid (k - l)$.

So

$$p \nmid a(k - l).$$

This is a contradiction.

So we have

$$r \neq s$$

Since $r \equiv s \pmod{m}$

$$\begin{aligned} r - s &\equiv [(ak + b) \bmod p] - [(al + b) \bmod p] \equiv 0 \pmod{m} \\ [a(k - l) \bmod p] &\equiv 0 \pmod{m} \quad (*) \end{aligned}$$

The possible value of $[a(k - l) \bmod p]$ is $m, 2 \cdot m, \dots, [p/m] \cdot m$. Its value cannot be zero because $r \neq s$.

So the number of different a satisfying formulae (*) is at most $[p/m]$. For each value of a , there are p different values of b . So the number of different h_{ab} satisfying $h_{ab}(k) = h_{ab}(l)$ is at most

$$\left\lfloor \frac{p}{m} \right\rfloor \cdot p \leq \frac{(p-1)p}{m} = \frac{|H_{pm}|}{m}$$

So the class H_{pm} of hash functions is universal.