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Assignment-1

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Que 1:-

$$t_{u=10} = 5 \text{ sec}$$
$$t_{u=50} = ?$$

$$Rt^2 = t$$

$$R \times 100 = 5$$

$$R = \frac{1}{20}$$

$$t_{u=50} = \frac{1}{20} \times 50 \times 50$$
$$= 125 \text{ sec}$$

Que 2:-

$$T_A = u^3, T_B = 2u^2$$

At Breaking point:-

$$u^3 = 2u^2$$

$$u = 2$$

Que 3:-

$$f(u) = u \cdot 2^u, g(u) = 4^u$$

Applying limit rule:-

$$\lim_{u \rightarrow \infty} \frac{f(u)}{g(u)} \Rightarrow \lim_{u \rightarrow \infty} \frac{u \cdot 2^u}{4^u} \Rightarrow \lim_{u \rightarrow \infty} \frac{u \cdot 2^u}{2^{2u}}$$

Applying L' Hospital:-

$$\lim_{u \rightarrow \infty} \frac{1}{2^u \log 2}$$

Hence  $f(u)$  is in  $o(g(u))$ .



Ques 4: -

Let any polynomial function,  $P(x)$  and  $\log(x)$  function

Using limit rule,  $\lim_{x \rightarrow \infty} \frac{\log(x)}{P(x)}$

Applying L'Hospital rule: -

$$\lim_{x \rightarrow \infty} \frac{1}{x \cdot P'(x)}$$

$\neq 0$

Thus,  $\log x$  grows slower than all functions.

Let two log fun.  $\log_a x$  and  $\log_b(x+1)$

Applying limit rule: -

$$\lim_{x \rightarrow \infty} \frac{\log_a x}{\log_b(x+1)} \Rightarrow \lim_{x \rightarrow \infty} \frac{\log_a x}{\log_b x + \log(1+1/x)}$$

$$\lim_{x \rightarrow \infty} \frac{\log_a x}{\log_b x}$$

$$\Rightarrow \lim_{x \rightarrow \infty} \frac{\ln b}{\ln a}$$

$\Rightarrow$  constant

Hence all log functions grow at same rate.



$$f(x) = x^4 + \log x + 17, g(x) = x^4$$

$$\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = \lim_{x \rightarrow \infty} \frac{x^4 + \log x + 17}{x^4}$$

$$\lim_{x \rightarrow \infty} \frac{4x^3 + 1/x}{4x^3}$$

$$\lim_{x \rightarrow \infty} \frac{4x^3 + 1}{4x^3} \quad \frac{1}{4x^4}$$

$\Rightarrow 1$  (constant)

Hence  $f(x)$  is  $O(\log(x))$ .

Que 7:-

(a)

$$R=1$$

$$R=2$$

$$R=3$$

⋮

$$R=n-1$$

$$O(n-1) \Rightarrow O(n)$$

(b)

$$\text{for } i=1 \quad j=2, 3, \dots, n \quad n-1$$

$$j=2 \quad j=3, 4, \dots, n \quad n-2$$

$$\vdots$$

$$i=n-1 \quad n \quad (n-1)$$

$$(n-1)(n-2) + (n-3) + (n-4) + \dots + 1$$

$$\Rightarrow \frac{n(n-1)}{2}$$

$$\Rightarrow O(n^2)$$

Que 9:-

$$T_A = 100^{2n}$$

$$T_B = 2^{n^4}$$

Using limit rule:-

$$\lim_{n \rightarrow \infty} \frac{2^{n^4}}{100^{2n}} \left( \frac{\infty}{\infty} \right)$$

$$\lim_{n \rightarrow \infty} \frac{4n^3}{100^{2n} \log 100}$$

$$\lim_{n \rightarrow \infty} \frac{2n}{100^{2n} (\log(100))^4} = 0$$

Hence,  $T_A$  grows faster when  $n \rightarrow \infty$

Que 10:-

$$f(n) = n \log n, g(n) = \log(n!)$$

$$\Rightarrow \log(1 \cdot 2 \cdot 3 \cdot 4 \cdots (n-1) \cdot n)$$

$$\Rightarrow \log\left(\frac{n^{n/2}}{2}\right)$$

$$\approx \frac{n}{2} \log \frac{n}{2}$$

$$\Rightarrow \frac{n}{2} \log n - \frac{n}{2} \log 2$$

$$\approx \frac{n}{2} (\log n - \log 2) \approx O(n \log n)$$

Hence,  $f(n) \in O(g(n))$



Que 11.  
(a)

$$2^{n-1} + 4^{n+1}$$

for 0

$$c_1 g(n) \leq f(n) \leq c_2 g(n)$$

$$2^{n-1} + 2^{2n+2}$$

$$\frac{2^n + 4 \cdot 2^{2n}}{2}$$

$$\Rightarrow 2^{2n} \left( \frac{1}{2} + 1 \right)$$

Highest order term  $2^{2n}$

$$\Rightarrow O(4^n)$$

(b)

$$(2x^2 + 6)^8$$

$$\text{Highest order term} = (2x^2)^8$$

$$= 2x^{16}$$

$$O(2x^{16})$$

Que 12:-

$$T_A = x^2$$

$$T_B = x + 2$$

Breaking point:-

$$x^2 = x + 2$$

$$x^2 - x - 2 = 0$$

$$(x-2)(x+1) = 0$$

$$x = 2$$



Ques:-

a)

Avg. Case (O)  
 Performs avg. no.  
 of steps per input  
 data of  $n$  elements.  
 Averaged over all  
 possible inputs.

Ex:-  $O(n \log n)$

Worst case (O)

performs max. no.  
 of steps on input  
 data of  $n$  elements.  
 Input is arbitrary  
 and output is upper  
 bound.

Ex:-  $O(n^2)$

(b)

Worst case (O)  
 It is arbitrary  
 and O/P is upper bound.

Asymptotically Bounded  
 Growth of a  
 running time to  
 within certain  
 factors below.

Ques:-

$$f(n) = n^2$$

$$1 + 2 + 3 + \dots + (n-1) + n = \frac{n(n+1)}{2}$$

whenever  $n$  doubles, the running time inc(1) fold.  
 Quad. algo. are for relatively smaller problems.