

# Homework1 Solution

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## Problem1

a. Prove  $2n + \Theta(n^2) = \Theta(n^2)$

证明:

先证  $2n + \Theta(n^2) \subseteq \Theta(n^2)$

$\forall f(n) \in 2n + \Theta(n^2), \exists c_1, c_2, n_0 \in \mathbb{R}^+, s.t. \forall n \geq n_0$

$$0 \leq c_1(2n + n^2) \leq f(n) \leq c_2(2n + n^2) \quad (1)$$

令  $n_1 = \max(n_0, 2), \exists n_1 \in \mathbb{R}^+, s.t. \forall n \geq n_1$

$$0 \leq c_1 n^2 \leq f(n) \leq c_2(2n + n^2) \leq 2c_2 n^2 \quad (2)$$

则  $f(n) \in \Theta(n^2)$ , 也即  $2n + \Theta(n^2) \subseteq \Theta(n^2)$

再证  $\Theta(n^2) \subseteq 2n + \Theta(n^2)$

$\forall f(n) \in \Theta(n^2), \exists c_1, c_2, n_0 \in \mathbb{R}^+, s.t. \forall n \geq n_0$

$$0 \leq c_1 n^2 \leq f(n) \leq c_2 n^2 \quad (3)$$

令  $n_1 = \max(n_0, 2), \exists n_1 \in \mathbb{R}^+, s.t. \forall n \geq n_1$

$$0 \leq \frac{c_1}{2}(2n + n^2) \leq c_1 n^2 \leq f(n) \leq c_2 n^2 \leq c_2(2n + n^2) \quad (4)$$

则  $f(n) \in 2n + \Theta(n^2)$ , 也即  $\Theta(n^2) \subseteq 2n + \Theta(n^2)$

综上所述,  $2n + \Theta(n^2) = \Theta(n^2)$

证毕!

**b. Prove**  $\Theta(g(n)) \cap o(g(n)) = \emptyset$

**证明:**

$\forall f(n) \in \Theta(g(n)), \exists c_1, c_2, n_0 \in \mathbb{R}^+, s.t. \forall n \geq n_0$

$$0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n) \quad (5)$$

假设  $f(n) \in o(g(n)), \forall c_3 > 0, \exists n_1 \in \mathbb{R}^+, s.t. \forall n \geq n_1$

$$0 \leq f(n) < c_3 g(n) \quad (6)$$

取  $c_3 = c_1, n_2 = \max(n_0, n_1), s.t. \forall n \geq n_2$

$$c_3 g(n) = c_1 g(n) \leq f(n) \quad (7)$$

这与(6)式矛盾, 故假设不成立, 也即  $\forall f(n) \in \Theta(g(n)), f(n) \notin o(g(n))$

综上所述,  $\Theta(g(n)) \cap o(g(n)) = \emptyset$

证毕!

**c. Prove**  $\Theta(g(n)) \cup o(g(n)) \neq O(g(n))$

**证明:**

取  $g(n) = 1, f(n) = \left| \sin\left(\frac{\pi n}{2}\right) \right|$

下面说明  $f(n) \in O(g(n))$  且  $f(n) \notin \Theta(g(n)) \cup o(g(n))$

取  $c_1 = 1, n_0 = 1, \forall n \leq n_0$

$$0 \leq f(n) \leq c_1 g(n) = c_1 = 1 \quad (8)$$

则  $f(n) \in O(g(n))$

因为  $\forall n = 2k, k = 1, 2, 3, \dots, f(n) = 0$ , 所以  $\forall c_2, n_0 \in \mathbb{R}^+, \exists n > n_0, n = 2k, s.t.$

$$f(n) = 0 < c_2 g(n) = c_2 \quad (9)$$

也即  $f(n) \notin \Theta(g(n))$

取  $c_3 = \frac{1}{2}, \forall n_1 \in \mathbb{R}^+, \exists n > n_1, n = 2k + 1, s.t.$

$$\frac{1}{2} = c_3 g(n) < f(n) = 1 \quad (10)$$

也即  $f(n) \notin o(g(n))$

综上所述,  $f(n) \in O(g(n))$  且  $f(n) \notin \Theta(g(n)) \cup o(g(n))$

则  $\Theta(g(n)) \cup o(g(n)) \neq O(g(n))$

证毕!

**d. Prove**  $\max(f(n), g(n)) = \Theta(f(n) + g(n))$

**证明:**

令  $h(n) = \max(f(n), g(n)), \exists n_0 \in \mathbb{R}^+, s.t. \forall n > n_0$

$$h(n) \geq f(n) \quad (11)$$

$$h(n) \geq g(n) \quad (12)$$

$$h(n) \leq f(n) + f(n) + 2g(n) \quad (13)$$

整理可得

$$0 \leq \frac{1}{2}(f(n) + g(n)) \leq h(n) = \max(f(n), g(n)) \leq 2(f(n) + g(n)) \quad (14)$$

则可得  $\max(f(n), g(n)) = \Theta(f(n) + g(n))$

证毕!

## Problem2

a.  $T(n) = 2T(\sqrt{n}) + 1$

解:

令  $m = \lg n$ , 则有  $n = 2^m$ , 带入原式得

$$T(2^m) = 2T(2^{\frac{m}{2}}) + 1 \quad (15)$$

令  $G(n) = T(2^n)$ , 可得

$$G(m) = 2G(\frac{m}{2}) + 1 \quad (16)$$

根据主定理,  $a = 2, b = 2, f(n) = 1$ , 则  $n^{\log_b a} = n$ , 取  $\epsilon = 1$ , 则有

$$1 = f(n) = \Theta(n^{\log_b a - \epsilon}) = \Theta(1) \quad (17)$$

则  $G(n) = \Theta(n)$ , 又因为  $G(n) = T(2^n)$ , 可得

$$T(n) = \Theta(\lg n) \quad (18)$$

b.  $nT(n) = (n - 2)T(n - 1) + 2$

解:

原式可化为

$$n(T(n) - 1) = (n - 2)(T(n - 1) - 1) \quad (19)$$

令  $G(n) = T(n) - 1$ , 则上式可化为

$$\frac{G(n)}{G(n - 1)} = \frac{n - 2}{n} \quad (20)$$

累乘可得

$$\frac{G(n)}{G(2)} = \frac{G(n)}{G(n - 1)} \frac{G(n - 1)}{G(n - 2)} \cdots \frac{G(3)}{G(2)} = \frac{n - 2}{n} \frac{n - 3}{n - 1} \cdots \frac{1}{3} = \frac{2(n - 2)!}{n!} = \frac{2}{n(n - 1)} \quad (21)$$

则可得

$$T(n) = \frac{2G(2)}{n(n - 1)} + 1 \quad (22)$$

也即

$$T(n) = \Theta\left(\frac{1}{n^2}\right) \quad (23)$$

## Problem3

a. 解:

下表给出排序结果, 从上到下渐进增长率增大, 同一行中从左向右渐进增长率增大

等价类	函数
$\Theta(1)$	$1, n^{\frac{1}{lgn}}$
$\Theta(lg(lg^*n))$	$lg(lg^*n)$
$\Theta(lg^*(lgn))$	$lg^*(lgn)$
$\Theta(lg^*n)$	$lg^*n$
$\Theta(2^{lg^*n})$	$2^{lg^*n}$
$\Theta(lg(lgn))$	$ln(lnn)$
$\Theta(\sqrt{lgn})$	$\sqrt{lgn}$
$\Theta(lgn)$	$lnn$
$\Theta(lg^2n)$	$lg^2n$
$\Theta(2^{\sqrt{2lgn}})$	$2^{\sqrt{2lgn}}$
$\Theta(n^{\frac{1}{2}})$	$\sqrt{2}^{lgn}$
$\Theta(n)$	$2^{lgn}, n$
$\Theta(nlgn)$	$nlgn, lg(n!)$
$\Theta(n^2)$	$n^2, 4^{lgn}$
$\Theta(n^3)$	$n^3$
$\Theta((lgn)!)$	$(lgn)!$
$\Theta(lgn^{lgn})$	$lgn^{lgn}, n^{lg(lgn)}$
$\Theta((\frac{3}{2})^n)$	$(\frac{3}{2})^n$
$\Theta(2^n)$	$2^n$
$\Theta(n2^n)$	$n2^n$
$\Theta(e^n)$	$e^n$
$\Theta(n!)$	$n!$
$\Theta((n+1)!)$	$(n+1)!$
$\Theta(2^{2^n})$	$2^{2^n}$
$\Theta(2^{2^{n+1}})$	$2^{2^{n+1}}$

b. 解:

令  $f(n) = 2^{2^{n+2}} + (-1)^n 2^{2^{n+2}}$ , 当  $n$  为偶数时,  $f(n) = \Theta(2^{2^{n+2}})$ , 当  $n$  为奇数时,  $f(n) = 0$   
 因此,  $\forall c_1, c_2, n_0 \in \mathbb{R}^+, \exists n > n_0, s.t.$

$$f(n) > c_1 g_i(n) \quad (n = 2k) \quad (24)$$

$$f(n) < c_2 g_i(n) \quad (n = 2k + 1) \quad (25)$$

则对于所有  $g_i(n)$ , 上述  $f(n)$  既不是  $O(g_i(n))$  也不是  $\Omega(g_i(n))$