

* PAC for inconsistent-hypo sets.

① Supp. we use agnostic PAC to analyse algorithms in the "inconsistent-hyp set" case. Also, suppose that the algo. works by solving the following optimization pb:

ERM (empirical risk minimization).

$$h_S^{\text{ERM}} = \underset{h \in \mathcal{H}^{\text{eff}}}{\text{argmin}} \widehat{R}_S(h) \quad \text{for a given seq. } S \text{ of samples.}$$

Denote this algo. by A_{ERM}.

② Can we prove something like Thm 2.5 that says that if $|\mathcal{H}^{\text{eff}}|$ is at most exponential in n , A_{ERM} is an agnostic PAC algo.?

③ The answer is yes. Here is the thm.

Thm $\forall D \in \mathcal{P}_{\mathcal{X} \times \mathcal{Y}} \quad \forall \varepsilon, \delta > 0 \quad \forall m \in \mathbb{N},$

if $m \geq \frac{2}{\varepsilon^2} \left(\log |\mathcal{H}^{\text{eff}}| + \log \frac{2}{\delta} \right)$, then

$$\Pr_{\mathcal{D}^m} [R_{\mathcal{D}^m}(h_S^{\text{ERM}}) - \min_{h \in \mathcal{H}^{\text{eff}}} R(h) \leq \varepsilon] \geq 1 - \delta.$$

We first note the following proposition from Ch 1.

$$\begin{aligned} \text{Prop 1.1} \quad & \Pr_{\mathcal{D}^m} [R_{\mathcal{D}^m}(h_S^{\text{ERM}}) - \min_{h \in \mathcal{H}^{\text{eff}}} R(h) > \varepsilon] \\ & \leq \Pr_{\mathcal{D}^m} [\max_{h \in \mathcal{H}^{\text{eff}}} |R(h) - \widehat{R}_S(h)| > \frac{\varepsilon}{2}]. \end{aligned}$$

for all $\varepsilon > 0$ and $m \in \mathbb{N}$.

Proof. Let $h_0 = \underset{h \in \mathcal{H}}{\operatorname{arg\min}} R(h)$ (which is defined since \mathcal{H} is finite).

Also, let $S = ((x_1, y_1), \dots, (x_m, y_m))$ be the seq. of samples drawn independently from D . Then,

$$\begin{aligned} & |R(h_S^{\text{ERM}}) - \min_{h \in \mathcal{H}} R(h)| = R(h_S^{\text{ERM}}) - R(h_0) \\ &= R(h_S^{\text{ERM}}) - \hat{R}_S(h_S^{\text{ERM}}) + \hat{R}_S(h_S^{\text{ERM}}) - R(h_0) \\ &\leq R(h_S^{\text{ERM}}) - \hat{R}_S(h_S^{\text{ERM}}) + \hat{R}_S(h_0) - R(h_0) \\ &\leq 2 \max_{h \in \mathcal{H}} |R(h) - \hat{R}_S(h)| \end{aligned}$$

So, $\Pr[\min_{h \in \mathcal{H}} R(h) > \varepsilon] \leq \Pr[\max_{h \in \mathcal{H}} |R(h) - \hat{R}_S(h)| > \frac{\varepsilon}{2}]$. \square

Before proving the theorem, we note that the following variant of Thm 2.13 also holds essentially by the same proof.

Thm 2.13 b. $\forall D \in \mathcal{P}(\exists x, y) \quad \forall \varepsilon, S > 0 \quad \exists m \in \mathbb{N}$.

If $m \geq \frac{1}{2\varepsilon^2} (\log |\mathcal{H}| + \log \frac{2}{S})$, then new to this thm 2.13 b.

(missed in Thm 2.13 by mistake)

$\Pr[\forall h \in \mathcal{H}, |R(h) - \hat{R}_S(h)| \leq \varepsilon] \geq 1 - S$.

Using what we have shown so far, we prove our thm. J

Proof of the thm

We will show $\Pr_{S \sim D^m} [\min_{h \in \mathcal{H}} R(h) > \varepsilon] \leq S$.

By Prop. 4.1,

$$P \left[\min_{h \in D^m} R(h) - \min_{h \in \mathcal{H}} R(h) > \varepsilon \right]$$

$$\leq P \left[\max_{h \in D^m} |R(h) - \hat{R}_S(h)| > \frac{\varepsilon}{2} \right]$$

By assumption of the thm.,

$$m \geq \frac{1}{2\epsilon^2} \left(\log |\mathcal{H}| + \log \frac{2}{\delta} \right).$$

Thus, we can use Thm 2.13b with $\frac{\varepsilon}{2}$, and derive:

$$\therefore \leq S.$$

D.