

3. Calloc and Throw.

① Some programming languages allow continuations to be denotable values, and provide language constructs for manipulating continuation values.

② Semantically, this means that we change V as follows:

$$V \xrightarrow[\phi]{\phi} V_{int} + V_{bool} + V_{fun} + V_{tuple} + V_{alt} + \underline{V_{cont.}}$$

Syntactically, it involves adding the following two constructs:

$\langle \text{exp} \rangle ::= \text{calloc} \langle \text{exp} \rangle \mid \text{throw} \langle \text{exp} \rangle \langle \text{exp} \rangle$

calloc expects a function as its argument. ($\text{calloc}(\lambda f. e)$) reifies the current continuation, binds f to it, and executes e . The bound continuation f can be invoked by throw , as in $(\text{throw}. f \ 3)$. This calls the continuation f with the value 3.

③ Here are the semantic clauses for calloc and throw :

$$\llbracket \text{calloc } e \rrbracket \eta \kappa = \llbracket e \rrbracket \eta (\lambda f. f(\psi \langle \kappa \rangle)) \kappa \text{ fun.}$$

$$\llbracket \text{throw } e \ e' \rrbracket \eta \kappa = \llbracket e \rrbracket \eta (\lambda \kappa'. \llbracket e' \rrbracket \eta \kappa') \text{ cont.}$$

Intuitively, in

$(\text{calloc } \lambda k. \dots \text{throw } k \ 3 \dots)$

— can be viewed as putting a label denoted by k , and — can be understood as a goto to this label.

④ What are the results of the following expressions?

(i) $\text{calloc}(\lambda k. 2 + \text{throw } k \ (3 \times 4))$.

(ii) $\left(\text{calloc}(\lambda k. \lambda x. \text{throw } k \ (\lambda y. x + y)) \right) 6$

The first example can be understood as skipping some part of computation. The second shows how we can repeat the computation of some part of an expression using continuation.

⑤ The next example is likely very hard to understand because it uses features not explained so far and it is also quite tricky. The example is from p290 of the textbook.

that we would like to implement a

Imagine routine backtrack that takes a function and tries the function with a parameter amb representing a nondeterministic choice between true and false. It collects results of all those choices and returns.

For instance, a list of all those results.

$\text{backtrack}(\lambda \text{amb.}$

if amb then (if amb then 0 else 1)
(empty tuple.)
else (if amb then 2 else 3)