

$$\llbracket \lambda x. e \rrbracket \eta = \phi(\lambda a \in D_1. \llbracket e \rrbracket [\eta | x:a])$$

Prop 10.8. Well-defined. That β is continuous.

Prop 10.12., Prop 10.13 and slightly more.

If $e_0 \rightarrow e_1$, then
 $\llbracket e_0 \rrbracket \eta = \llbracket e_1 \rrbracket \eta$ for all $\eta \in D_1$.

The contraction relation preserves the semantics.
 Note that this implies that the reduction relation \rightarrow^* also preserves

② Here are the denotational semantics for the
 normal-order evaluation and then for the
 eager evaluation.

normal-order evaluation.

$$D_2 = (V_2)_1, \quad V_2 \xrightarrow[\psi]{\phi} [D_2 \rightarrow_c D_2]$$

$$\llbracket - \rrbracket_n \in [\langle \text{exp} \rangle \rightarrow [D_2 \xrightarrow[\psi]{\phi} D_2]]$$

$$\llbracket x \rrbracket_n \eta = \eta(x)$$

$$\llbracket e_0 e_1 \rrbracket_n \eta = \begin{cases} \perp & \text{if } \llbracket e_0 \rrbracket_n \eta = \perp \\ \phi(\llbracket e_0 \rrbracket_n \eta)(\llbracket e_1 \rrbracket_n \eta) & \text{otherwise} \end{cases}$$

$$\llbracket \lambda x. e \rrbracket_n \eta = \phi(\lambda a \in D_2. \llbracket e \rrbracket_n [\eta | x:a]) \quad (\text{i.e. } \llbracket e \rrbracket_n \eta \in V_2)$$

eager evaluation.

$$D_3 = (V_3)_1, \quad V_3 \xrightarrow[\psi]{\phi} [V_3 \rightarrow_c D_3]$$

$$\llbracket - \rrbracket_e \in [\langle \text{exp} \rangle \rightarrow [V_3 \xrightarrow[\psi]{\phi} D_3]]$$

$$\llbracket v \rrbracket_e \eta = \eta(v)$$

$$\llbracket e_0 e_1 \rrbracket_e \eta = \begin{cases} \perp & \text{if } \llbracket e_0 \rrbracket_e \eta = \perp \text{ or } \llbracket e_1 \rrbracket_e \eta = \perp \\ \phi(\llbracket e_0 \rrbracket_e \eta)(\llbracket e_1 \rrbracket_e \eta) & \text{otherwise} \end{cases}$$

$$\llbracket \lambda v. e \rrbracket_e \eta = \phi(\lambda a \in V_3. \llbracket e \rrbracket_e [\eta | v:a])$$

Both semantics are well-defined. They validate
 the normal-order evaluation relation and the
 eager evaluation relation. That is,

$$(e \Rightarrow e') \Rightarrow \llbracket e \rrbracket_n \eta = \llbracket e' \rrbracket_n \eta \text{ for all } \eta \in D_2$$

$$(e \Rightarrow_e e') \Rightarrow \llbracket e \rrbracket_e \eta = \llbracket e' \rrbracket_e \eta \text{ for all } \eta \in V_3$$