

Big Oh Considerations, Reading

Did you notice a bit of a disconnect between the previous two modules? Module 7 was about algorithmic efficiency and how we prefer $O(1)$ to $O(n)$. Then in Module 8's associative arrays we applied $O(n)$ loops to just about everything in the template! It's like we just threw out all the work we did in the previous module.

Well, *this* module takes a look at the $O(n)$ issue and looks for a solution.

The Square Bracket Operator

With the `StaticArray` and `DynamicArray` templates, the square bracket operator is efficient (in a big oh sense) because the numeric index can be used to directly calculate the memory location of the indexed value and find it with $O(1)$. It works like this, even though we don't actually do the math: for a **`DynamicArray<int> a;`** the 9th value is at **`a[9]`**. In the function, using array notation as we do, the computer calculates the *offset* in the "values" array as **`9*sizeof(int)`** and goes directly there to retrieve the value. No traversing, no counting. Just a simple calc resulting in $O(1)$.

But with the `AssociativeArray` template, the square bracket operator uses a for-loop to traverse the whole array of nodes, searching them one-by-one for a match to the **`typename K`** parameter, using equals-equals. That's $O(n)$. What would be really nice is if we could get that to be $O(1)$ the same way that `StaticArray::operator[]` does. What *it* does is *calculate* the memory location. `AssociativeArray::operator[]` needs a way to *calculate* where the value associated with its **`typename K`** parameter is, instead of *searching* for it.

The solution is to **convert the parameter into an offset in a simple math calculation**, like we do with numeric whole number parameters.