

CMPT 280

Topic 13: k -D Trees

Mark G. Eramian

University of Saskatchewan

References

- Textbook, Chapter 13

2-D Trees

2-D Range Search

- A k -dimensional tree or k -D tree is a data structure that stores and organizes points in a d -dimensional space.
- A 2-D tree (kd -tree with $k = 2$) is a binary tree which stores an element and a two dimensional key in each node.
- It works like a combination of two ordered binary trees: one for the x dimension and one for the y dimension.
- At the even levels, the branch is based on the x -coordinate.
- At odd levels, the branch is based on the y -coordinate.

2-D Trees

2-D Range Search

Recall the 2D Range Searching Problem:

Given a rectangular region with corners (x_1, y_1) and (x_2, y_2) , determine the points in the region.

2D Range Searching:

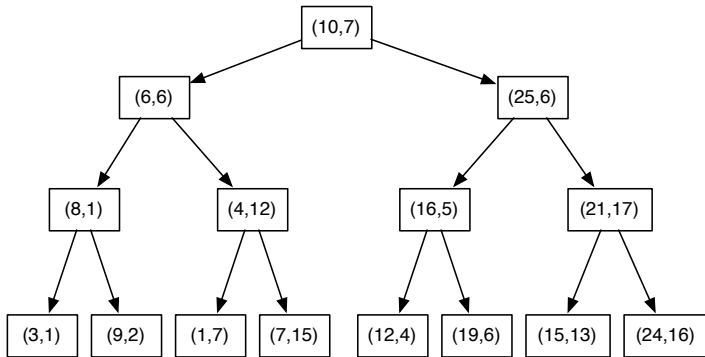
- Search the 2D kd-tree using the x range $[x_1, x_2]$ at the odd levels and the y range $[y_1, y_2]$ at even levels. When both subtrees of a node need to be searched, return the union of the found keys in the subtrees as was done in the 1D range search algorithm.
- Extra check to make sure that key in the current node has both coordinates in range, if it does, add it to the union.

2-D Trees

2-D Range Search

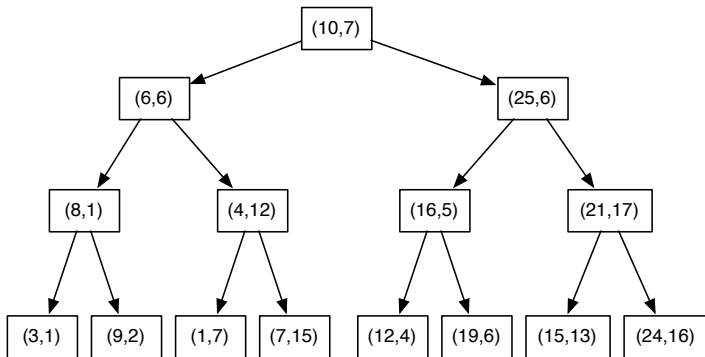
- Example: Find all people with age between 20 and 30 who are at least 6 feet tall.
- The basic technique can be applied to higher dimensional data. Each dimension of data takes its turn at being the one used for branching.
- This type of tree is great for database queries.

Exercise 1



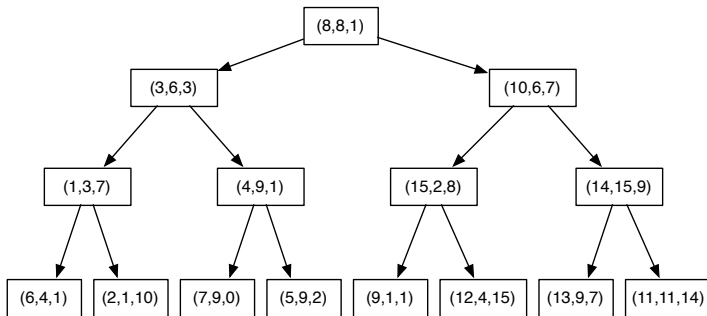
Which nodes will be visited in a 2D range search of this tree for the query rectangle with corners $(9, 13)$ and $(20, 16)$? Which keys are returned?

Exercise 2



Which nodes will be visited in a 2D range search of this tree for the query rectangle with corners $(9, 4)$ and $(20, 16)$? Which keys are returned?

Exercise 3



Which nodes will be visited in a 3D range search of this tree for the query cuboid with corners $(4, 6, 5)$ and $(9, 10, 12)$? Which keys are returned?

Sketch of Tree-Building Algorithm

Suppose we are given a set of key-element pairs. To build a 2-D tree, we do the following:

```
1 Order the elements by the x-coordinate of their keys.
2 Pick the key with the median x, make this element the root of the tree.
3
4 Recursively process the keys smaller than the median by
5 partitioning on the other coordinate. The median becomes the left
6 subtree of the root.
7
8 Recursively process the keys larger than the median by
9 partitioning on the other coordinate. The median becomes the right
10 subtree of the root.
```

Example: Building a 2-D tree

Given the following keys (we omit showing the associated elements):

$(1, 7), (9, 2), (21, 17), (3, 1), (10, 7), (24, 16), (4, 12), (12, 4),$

$(25, 6), (6, 6), (15, 13), (7, 15), (16, 5), (8, 1), (19, 6),$

let's build a 2-D tree. If we do this right, we should end up with the tree from Exercises 1 and 2. Let's go to the chalkboard...

The full example is written out in the Exercise Solutions for this topic.

Building k -D Trees

- Similar process for any k – just step through each dimension at each level.
- It's straightforward on paper, but in practice it's hard to implement efficiently. Sorts are slow: $O(n \log n)$.
- A lot of effort goes into doing this without sorting, and without using extra space beyond the input list of points and the k -D tree.
- Turns out, we don't need to sort. We just need to partition the keys into median, smaller than median and larger than median. There's an $O(n)$ algorithm for that!

Next Class

- Next class reading: Chapter 14: Other Trees