

Machete de interpretación y semántica denotacional



Intérprete con estrategia Call By Name (CBN)



$$\frac{\Gamma' \vdash M \hookrightarrow V}{\Gamma, x = \langle M, \Gamma' \rangle, \Delta \vdash x \hookrightarrow V} \quad x \notin D(\Delta)$$
$$\frac{\Gamma \vdash M \hookrightarrow \langle x, M', \Gamma' \rangle \quad \Gamma', x = \langle N, \Gamma \rangle \vdash M' \hookrightarrow V}{\Gamma \vdash MN \hookrightarrow V}$$
$$\frac{}{\Gamma \vdash \lambda x.M \hookrightarrow \langle x, M, \Gamma \rangle}$$
$$\frac{\Gamma \vdash \text{True} \hookrightarrow \text{True}}{\Gamma \vdash \text{if } M \text{ then } N_1 \text{ else } N_2 \hookrightarrow V} \quad \frac{\Gamma \vdash M \hookrightarrow \text{True} \quad \Gamma \vdash N_1 \hookrightarrow V}{\Gamma \vdash \text{if } M \text{ then } N_1 \text{ else } N_2 \hookrightarrow V}$$
$$\frac{\Gamma \vdash \text{False} \hookrightarrow \text{False}}{\Gamma \vdash \text{if } M \text{ then } N_1 \text{ else } N_2 \hookrightarrow V} \quad \frac{\Gamma \vdash M \hookrightarrow \text{False} \quad \Gamma \vdash N_2 \hookrightarrow V}{\Gamma \vdash \text{if } M \text{ then } N_1 \text{ else } N_2 \hookrightarrow V}$$
$$\frac{\Gamma, x = \langle \mu x.M, \Gamma \rangle \vdash M \hookrightarrow V}{\Gamma \vdash \mu x.M \hookrightarrow V}$$

Intérprete con estrategia Call By Value (CBV)



$$\begin{array}{c} \overline{\Gamma, x = V, \Delta \vdash x \hookrightarrow V} \quad x \notin D(\Delta) \qquad \overline{\Gamma' \vdash \mu y.M \rightarrow V} \quad x \notin D(\Delta) \\ \overline{\Gamma, x = \langle \mu y.M, \Gamma' \rangle, \Delta \vdash x \hookrightarrow V} \\ \\ \overline{\Gamma \vdash \lambda x.M \hookrightarrow \langle x, M, \Gamma \rangle} \qquad \frac{\Gamma \vdash N \hookrightarrow W \quad \Gamma \vdash M \hookrightarrow \langle x, M', \Gamma' \rangle \quad \Gamma', x = W \vdash M' \hookrightarrow V}{\Gamma \vdash MN \hookrightarrow V} \\ \\ \overline{\Gamma \vdash \text{True} \hookrightarrow \text{True}} \qquad \overline{\Gamma \vdash \text{False} \hookrightarrow \text{False}} \\ \frac{\Gamma \vdash M \hookrightarrow \text{True} \quad \Gamma \vdash N_1 \hookrightarrow V}{\Gamma \vdash \text{if } M \text{ then } N_1 \text{ else } N_2 \hookrightarrow V} \qquad \frac{\Gamma \vdash M \hookrightarrow \text{False} \quad \Gamma \vdash N_2 \hookrightarrow V}{\Gamma \vdash \text{if } M \text{ then } N_1 \text{ else } N_2 \hookrightarrow V} \\ \frac{\Gamma, x = \langle \mu x.M, \Gamma \rangle \vdash M \hookrightarrow V}{\Gamma \vdash \mu x.M \hookrightarrow V} \end{array}$$

Extensión de los intérpretes con números naturales



CBN y CBV

$$\frac{}{\Gamma \vdash \text{zero} \hookrightarrow \text{zero}}$$

$$\frac{\Gamma \vdash M \hookrightarrow \text{zero}}{\Gamma \vdash \text{pred}(M) \hookrightarrow \text{zero}}$$

$$\frac{\Gamma \vdash M \hookrightarrow \text{zero}}{\Gamma \vdash \text{isZero}(M) \hookrightarrow \text{True}}$$

$$\frac{\Gamma \vdash M \hookrightarrow V}{\Gamma \vdash \text{succ}(M) \hookrightarrow \text{succ}(V)}$$

$$\frac{\Gamma \vdash M \hookrightarrow \text{succ}(V)}{\Gamma \vdash \text{pred}(M) \hookrightarrow V}$$

$$\frac{\Gamma \vdash M \hookrightarrow \text{succ}(V)}{\Gamma \vdash \text{isZero}(M) \hookrightarrow \text{False}}$$

Semántica denotacional del Cálculo Lambda (sin error)



$$\begin{aligned} \llbracket x \rrbracket_v &= v(x) \\ \llbracket \lambda x:\tau. M \rrbracket_v &= V^{\llbracket \tau \rrbracket} \mapsto \llbracket M \rrbracket_{v, x=V} \\ \llbracket MN \rrbracket_v &= \llbracket M \rrbracket_v \llbracket N \rrbracket_v \\ \llbracket \text{True} \rrbracket_v &= \text{true} \\ \llbracket \text{False} \rrbracket_v &= \text{false} \\ \llbracket 0 \rrbracket_v &= 0 \\ \llbracket \text{succ}(M) \rrbracket_v &= \llbracket M \rrbracket_v + 1 \\ \llbracket \mu x:\tau. M \rrbracket_v &= \text{FIX}(V^{\llbracket \tau \rrbracket} \mapsto \llbracket M \rrbracket_{v, x=V}) \end{aligned}$$
$$\llbracket \text{if } M \text{ then } N \text{ else } O \rrbracket_v = \begin{cases} \llbracket N \rrbracket_v & \text{si } \llbracket M \rrbracket_v = \text{true} \\ \llbracket O \rrbracket_v & \text{si } \llbracket M \rrbracket_v = \text{false} \\ \perp & \text{si } \llbracket M \rrbracket_v = \perp \end{cases}$$
$$\llbracket \text{pred}(M) \rrbracket_v = \begin{cases} 0 & \text{si } \llbracket M \rrbracket_v = 0 \\ \llbracket M \rrbracket_v - 1 & \text{si no} \end{cases}$$
$$\llbracket \text{isZero}(M) \rrbracket_v = \begin{cases} \text{true} & \text{si } \llbracket M \rrbracket_v = 0 \\ \text{false} & \text{si } \llbracket M \rrbracket_v = n \neq 0 \\ \perp & \text{si } \llbracket M \rrbracket_v = \perp \end{cases}$$