

Supplementary Material for “Battery Dispatch Optimization for Electric Vehicle Aggregators: A Decomposition-Coordination-based Least Squares Approach with Disjunctive Cuts ”

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I. PROOF OF TIGHTENED FEASIBILITY SPACE

In the B&B procedure by DCHR, the upper and lower bounds of e_i^{t+1} are denoted as $\mathcal{L}_1^{upper} = e_{0,i} + \eta_{c,i} \bar{P}(t - 1 + \lambda_{c,i}^t) \Delta T$ and $\mathcal{L}_1^{lower} = e_{0,i} - \frac{1}{\eta_{d,i}} \bar{P}(t - 1 + \lambda_{d,i}^t) \Delta T$, respectively. However, according to the BDO model (1a)-(1d), we express the upper and lower bounds of e_i^{t+1} by $\mathcal{L}_2^{upper} = e_{0,i} + \eta_{c,i} \bar{P}(t) \Delta T$ and $\mathcal{L}_2^{lower} = e_{0,i} - \frac{1}{\eta_{d,i}} \bar{P}(t) \Delta T$, respectively. Hence, we achieve the difference in algebra $\mathcal{L}_1^{upper} - \mathcal{L}_2^{upper} = (\lambda_{c,i}^t - 1) \eta_{c,i} \bar{P} \Delta T$. Since $\lambda_{c,i}^t \in \{0, 1\}$, this clearly renders $\mathcal{L}_1^{upper} \leq \mathcal{L}_2^{upper}$. Similarly, we can also achieve the lower bound satisfying $\mathcal{L}_1^{lower} \geq \mathcal{L}_2^{lower}$ due to $\lambda_{d,i}^t \in \{0, 1\}$. We have proved that DCHR has tighter relaxation lower and upper bounds than those by the BDO model (1a)-(1d) for $\Omega_{op,i}$. ■

II. PROOF OF THEOREM

Initially, since $\sum_{i=1}^N p_{sig,i}^{t*} = p_{ref}^t$, we can rewrite F^* as

$$F^* = \sum_{t=1}^T \left\| \sum_{i=1}^N (c^T x_i^{t*} - p_{sig,i}^{t*}) \right\|_2 = \sum_{t=1}^T \left\| \sum_{i=1}^N \alpha_i^* \right\|_2 \quad (A-1)$$

where the optimal auxiliary variable $\alpha_i^* = c^T x_i^{t*} - p_{sig,i}^{t*}$.

Thus, our goal is to prove that the following two objective functions can get converged over $(x_i^t, z_i^t) \in \mathcal{X}_i, \forall i \in \mathcal{E}$, with the same optimal solution (x_i^{t*}, z_i^{t*}) for the MIQP-based model (1a)-(1d), namely

$$\operatorname{argmin}_{(x_i^t, z_i^t) \in \mathcal{X}_i} \sum_{t=1}^T \left\| \sum_{i=1}^N \alpha_i^* \right\|_2 \iff \operatorname{argmin}_{(x_i^t, z_i^t) \in \mathcal{X}_i} \sum_{t=1}^T \sum_{i=1}^N \|\alpha_i^*\|_2 \quad (A-2)$$

At time t , when $\left\| \sum_{i=1}^N \alpha_i^* \right\|_2$ can be minimized as a non-zero number, we suppose if there only $\exists \alpha^*$ as the non-zero number for the optimal solution $c^T x_i^{t*}$. Then, we have

$$\left\| \sum_{i=1}^N \alpha_i^* \right\|_2 = \|\alpha_1^* + \alpha_2^* \dots + \alpha_N^*\|_2 = \|\alpha^* + \alpha^* + \dots + \alpha^*\|_2 = N \|\alpha^*\|_2 \quad (A-3)$$

It is clear that equality (A-3) can be also equivalent to

$$N \|\alpha^*\|_2 = \|\alpha^*\|_2 + \|\alpha^*\|_2 + \dots + \|\alpha^*\|_2 = \sum_{i=1}^N \|\alpha_i^*\|_2 \quad (A-4)$$

Besides, if $\alpha^* = 0$, then the $\left\| \sum_{i=1}^N \alpha_i^* \right\|_2$ can reach zero. This indicates that $\sum_{i=1}^N \|\alpha_i^*\|_2 = \left\| \sum_{i=1}^N \alpha_i^* \right\|_2 = 0$ naturally holds. Based on the above, we can derive (A-5) for $t \in \mathcal{T}$ by

$$\sum_{t=1}^T \left\| \sum_{i=1}^N \alpha_i^* \right\|_2 = \sum_{t=1}^T \sum_{i=1}^N \|\alpha_i^*\|_2 \quad (A-5)$$

Thus, this equality (A-5) indicates that (A-2) can stand at the same optimal solution F^* over the same feasibility space \mathcal{X}_i , which proves this Theorem. ■