

Netzwerkkodierung in Theorie und Praxis

Praktische Anwendungen der Netzwerkkodierung

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Theoretische Nachrichtentechnik











16:40-18:10

06.Apr.2016 L2

11.Apr.2016 L3

14.Apr.2016 E1

20.Apr.2016 L5

27.Apr.2016 L6

28.Apr.2016 E2

16:40-18:10 13.Apr.2016 L4 VMB/0E02/U

GÖR/0127/U

VMB/0E02/U

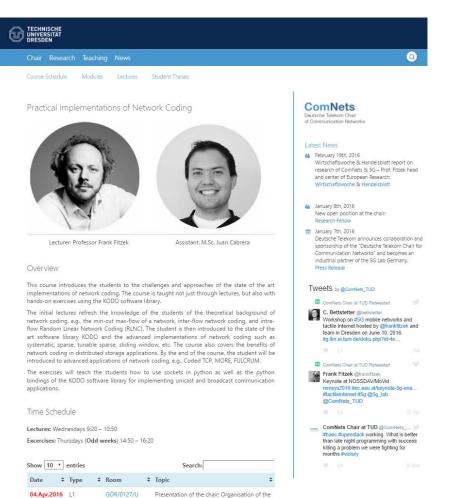
GÖR/0229/U

VMB/0E02/U

VMB/0F02/U

GÖR/0229/U

Lecture / Exercise Dates - tinyurl.com/zooafld



course: 5G Intro: Butterfly: min cut max flow.

Inter Flow NC; Index Coding; Zick Zack

Random Linear Network Coding (Basics)

UDP transmissions with python sockets.

RLNC advanced (sparse, tunable)

Analog Inter Flow Network Coding

Codina: CATWOMAN

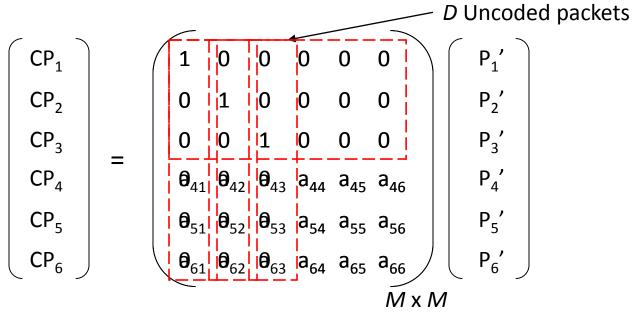
- Here all information for the lecture and the exercise can be found.
- Slides
- Links
 - Steinwurf
 - Python
 - KODOMARK (google play)

Please check every week!



A Practical Guide to RLNC Libraries

Systematic RLNC



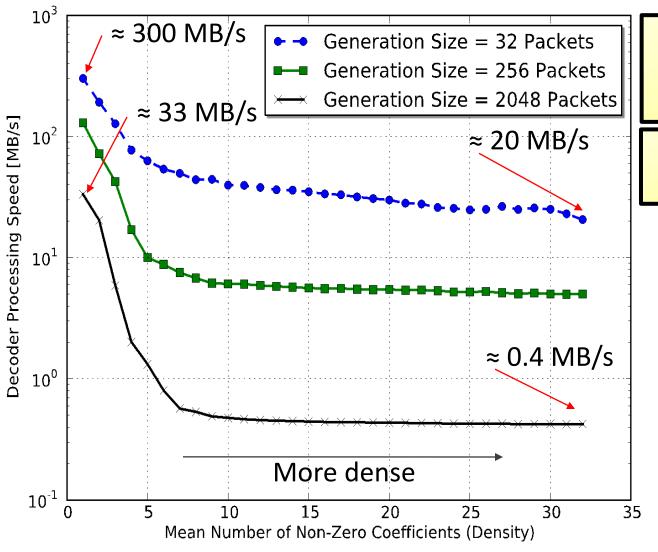
• Operations first elimination (Product):

- \mathbf{D} (M-D)
- Gaussian elimination n x n matrix, n = M D requires An3 + Bn2 + Cn operations
- Distribution of D determines average # of operations
 - Linked to channel model
- Erasures IID Be(Pe):

$$Q(M^{2}Ped)^{3})+B'(MPe)^{2}+C'(MPe)^{3}$$



Sparse Network Codes



One or two orders of magnitude in the coding speed by sparsity.

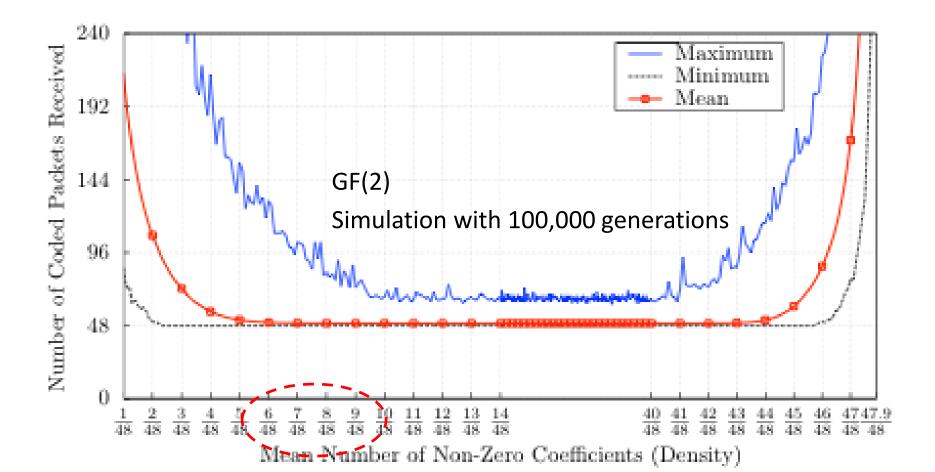
Dualism Theory and Implementation!





Sparse Network Codes

- Can we operate in the "speed up" region while maintaining good performance?
- Yes, but the key is not to use a fixed density
- (Tunable sparse network coding, 2012)

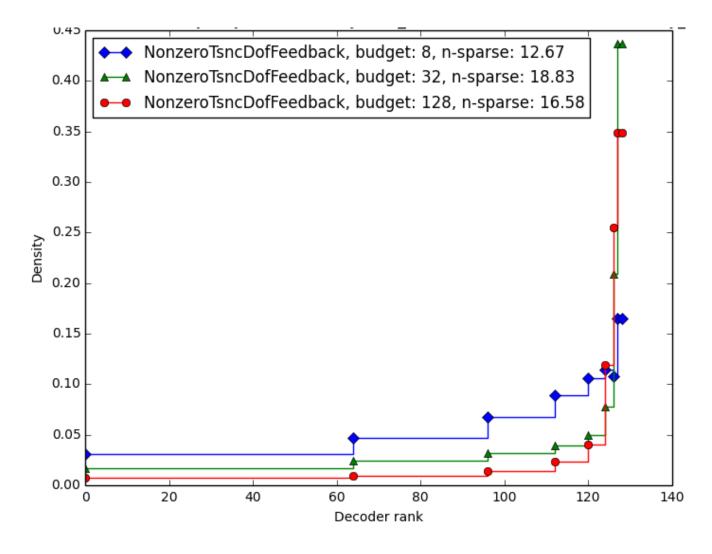




Tunable Sparse Network Codes

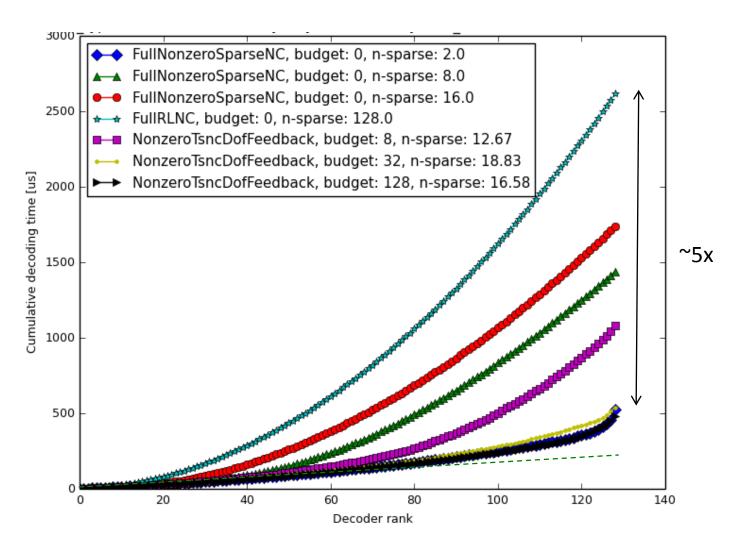
Scheme with feedback: targets specific performance

degradation



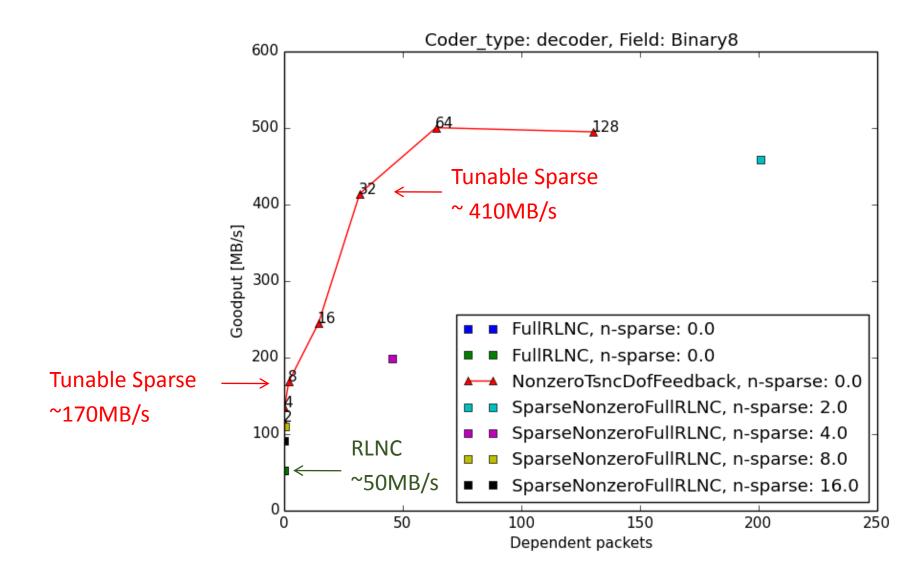


Tunable Sparse Network Codes



Scheme with feedback: targets specific performance degradation Example: 4 additional transmissions, i.e., 4% overhead

Tunable Sparse Network Codes

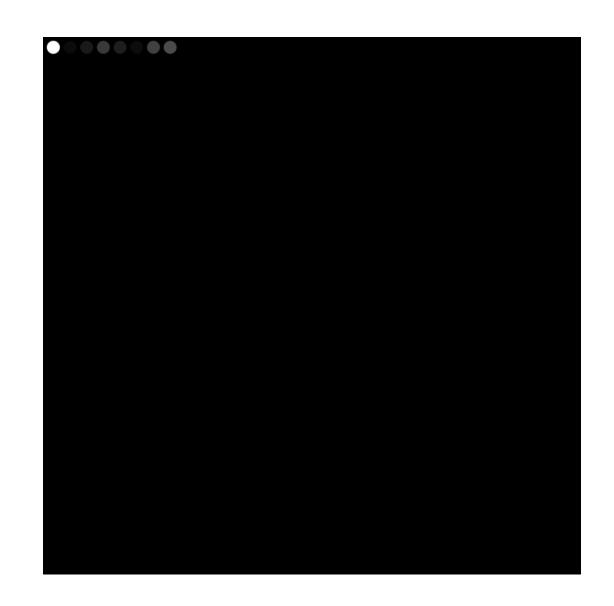




Sliding Window

Terminology

- On the fly encoding/decoding: Newly incoming packets are added to the block.
- Sliding window: Newly incoming packets are added to the block and acknowledged packets are removed.
- Sliding window can work without blocks



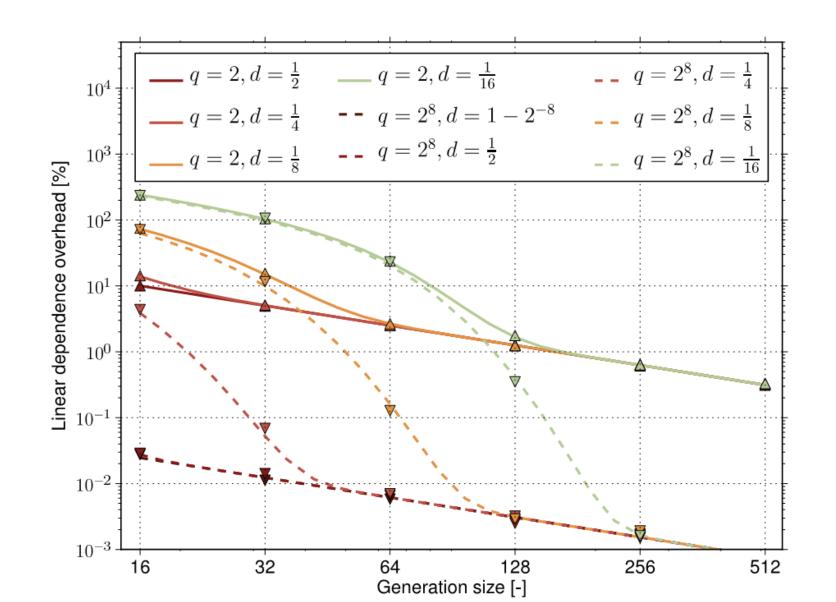


KODO: Overhead

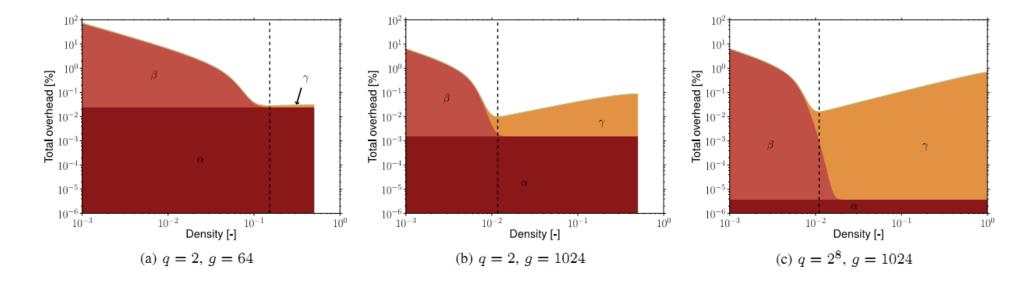
TECHNISCHE What is overhead?

- Overhead by encoding vector
 - Generation size G
 - Field size F
 - Payload P
 - Number of bits for each packet A = G * log2(F)
- Linear dependent retransmissions
 - In case a packet is linear dependent it has to be retransmitted (together with the header)
 - A = P + G * log2(F)

TECHNISCHE F/G and Linear Dependency



Designing rules ...

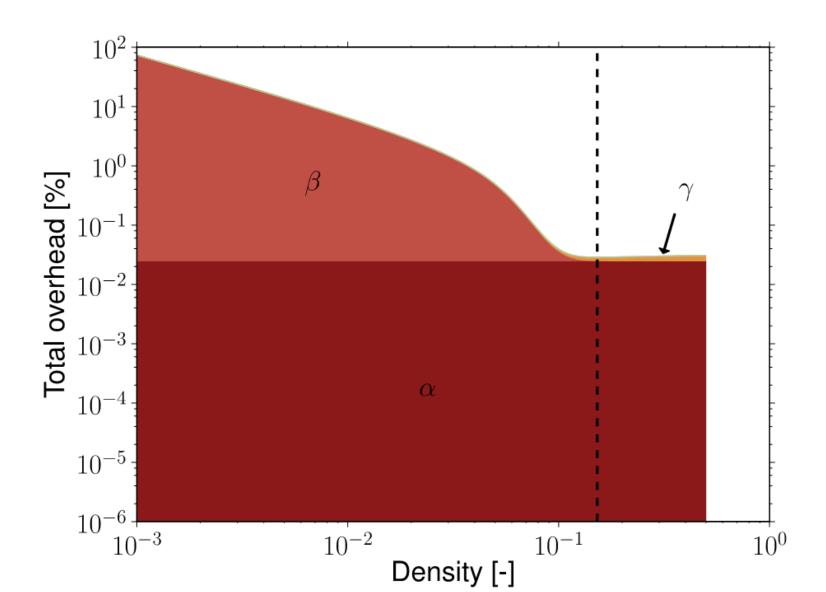


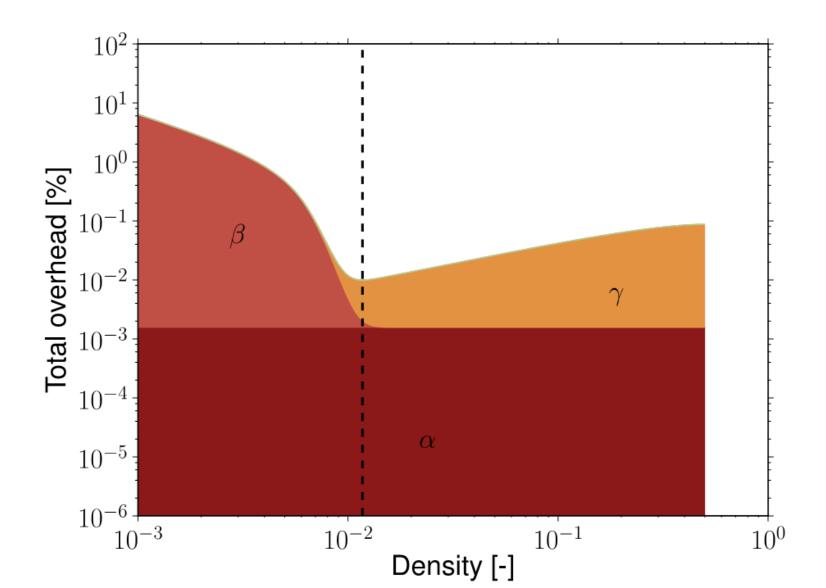
a = additional transmissions due to field size

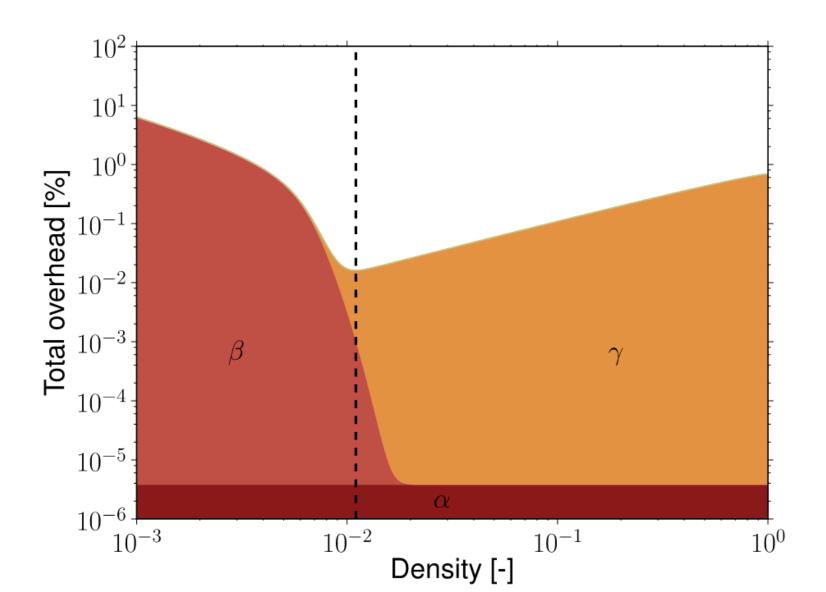
b = additional packets due to linear dependency

c = bits in the encoding vector

Janus Heide and Morten V. Pedersen and Frank H.P. Fitzek and Muriel Medard. **On Code Parameters and Coding Vector Representation for Practical RLNC.** 2011. in *IEEE International Conference on Communications (ICC) - Communication Theory Symposium*. Kyoto, Japan.





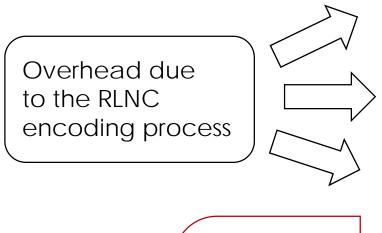




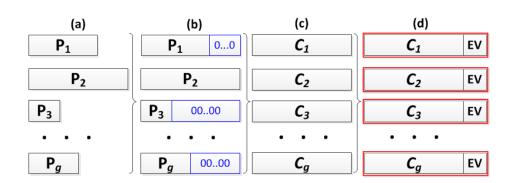
Effects of Heterogeneous Packet lengths on Network Coding



Effects of Heterogeneous Packet lengths on Network Coding



- (a) Coding overhead: defined by the Encoding Vector size, containing the coding coefficients.
- (b) Linear dependency overhead: need of extra transmissions when linear dependent combinations are received
- (c) Padding overhead: created by coding with the biggest data packet of a generation with different sizes



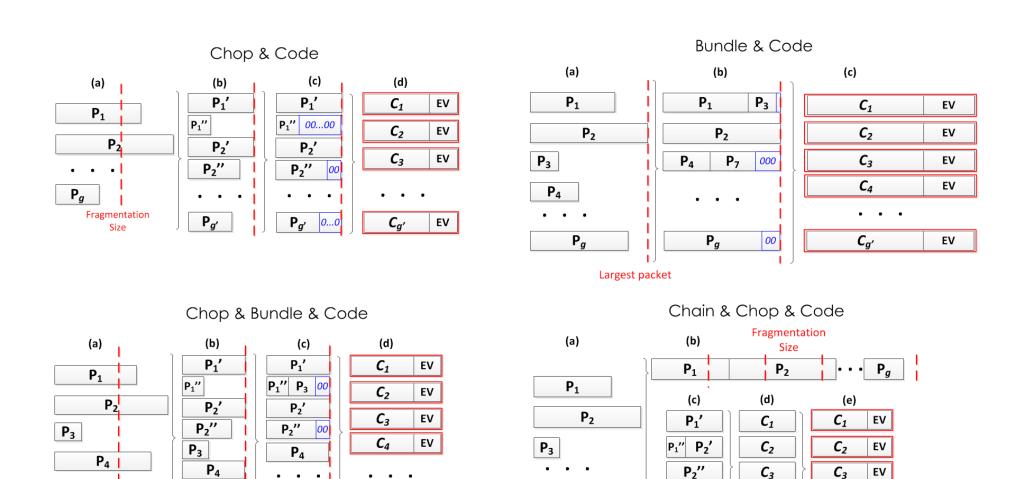
Naive assumption: all packets in a generation have the same size

Real life data: heterogenousity of packet lengths



Fragmentation

Packetization solutions



 P_g

P_g 0...0

 $C_{g'}$

 $C_{g'}$

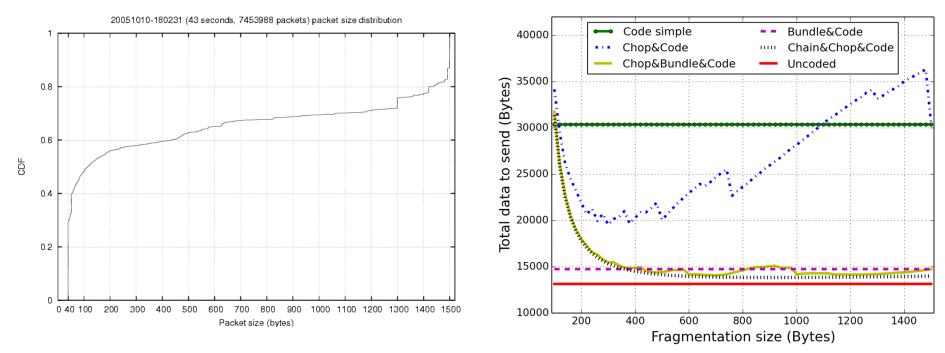
EV

EV

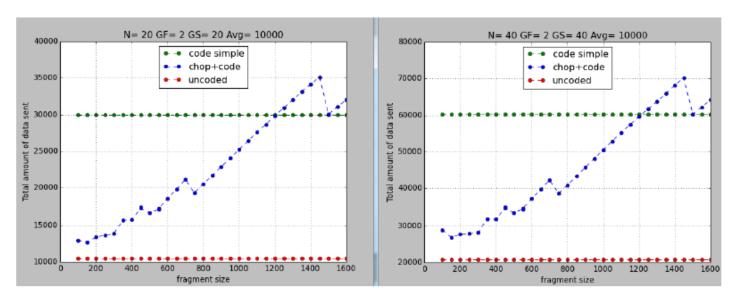


Internet data from CAIDA

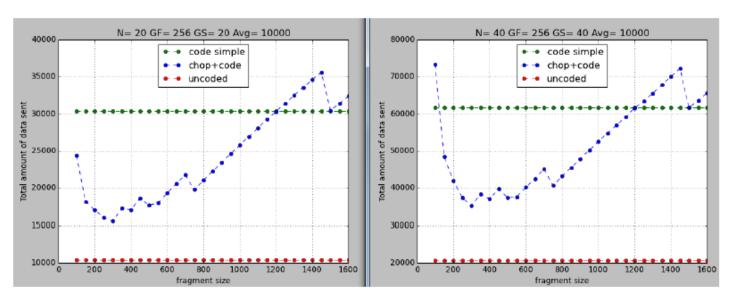
- Generations of 20 packets with field size q=2^8
- Packets are chosen randomly following the CDF, and run the simulations 10000 times and averaged the results
- Internet packet size distribution is mostly of 40 Bytes and 1500 Bytes (with approximate probabilities of 40% and 20% respectively). 1



TECHNISCHE UNIVERSITÄT Packet Size Comparison 1



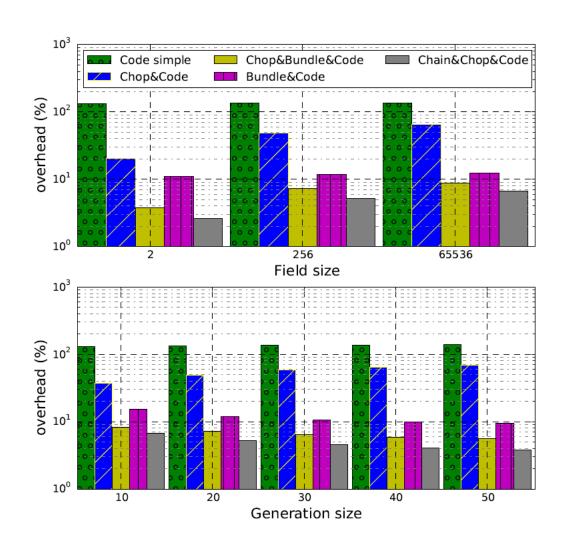
GF=2^8





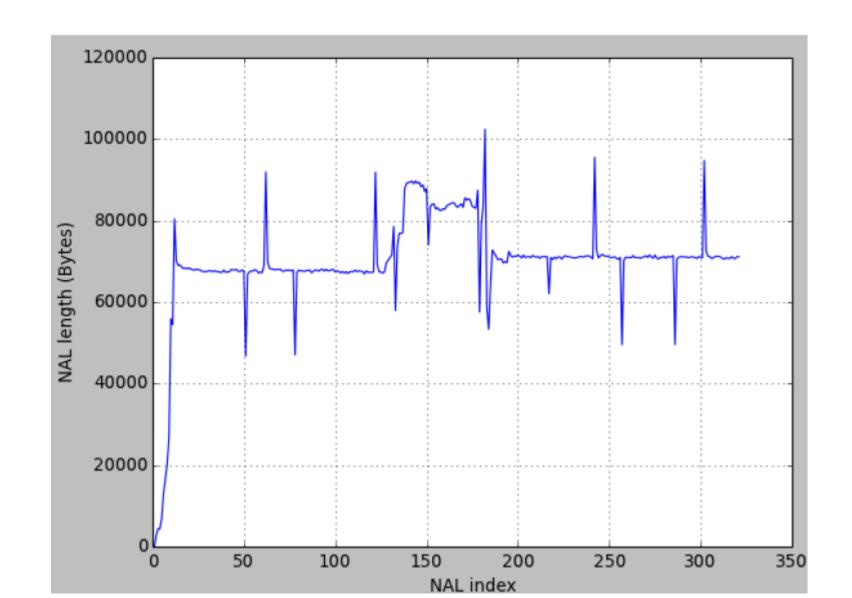
Internet data from CAIDA (2)

- Total coding overhead (%) added for each solution respect the uncoded data for:
 - different field sizes (and generation size constant of 20 packets)
 - different generation sizes (and field size constant q=28)
- Coding directly adds more than 100% of overhead.
- Chain& Chop& Code achieves to reduce to less than 5% for a large generation size.



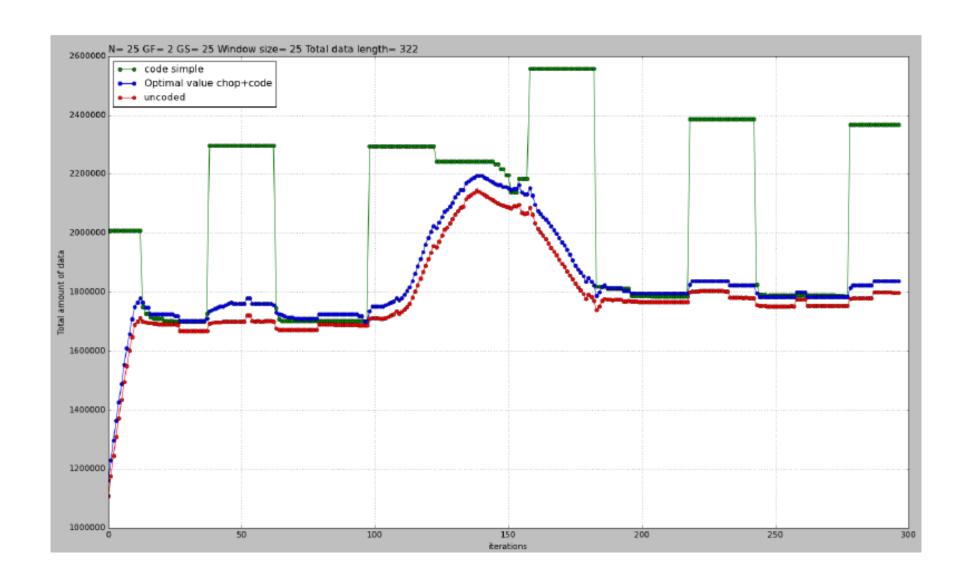


Packet Size Example 2





TECHNISCHE UNIVERSITÄT Packet Size Comparison 2





KODO: Energy consumption and complexity





TABLE I OVERVIEW OF THE FINITE FIELD IMPLEMENTATIONS USED.

GF(2)	The binary field $GF(2)$ was implemented using only the XOR operation, which typically yields very efficient implementations on most modern processors.
$GF(2^8)$	The binary extension field $GF(2^8)$ uses full multiplication and division lookup tables. Using this approach will replace the required polynomial multiplication or division of two field elements with a single table lookup. More details on these lookup tables are described in [17]. Addition and subtraction require only a bit-by-bit XOR.
$GF(2^{16})$	Computing a full multiplication and division table for $GF(2^{16})$ will typically require too much memory for most platforms (in the order of GBs). Instead, our implementation uses an extended logarithmic lookup, which reduces the memory requirements to the order of KBs. Addition and subtraction are performed as bit-by-bit XOR operation. Details are provided in [17].
$GF(2^{32}-5)$	This field is also known as an OPF (Optimal Prime Field). It has the advantage that it does not require any lookup tables and uses the standard arithmetic logic unit of the processor for all arithmetic operations. A detailed description of this field can be found in [18].

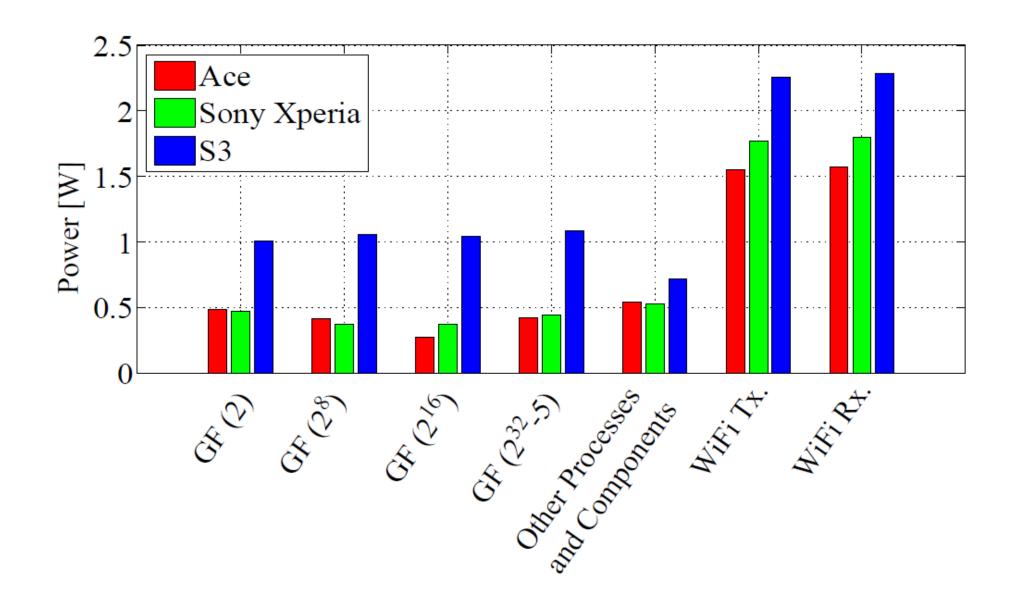




TABLE I
ENCODE AND DECODE PROCESSING SPEED (MBPS) AND AVERAGE PROCESSING POWER (W) FOR THREE MOBILE DEVICES.

Encoder		GF(2)			$GF(2^8)$			$GF(2^{16})$			$GF(2^{32}-5)$	
Pkts/Gen.	Ace	Xperia	S3	Ace	Xperia	S3	Ace	Xperia	S3	Ace	Xperia	S3
2	1675.8	3328.1	3916.22	238.4	463.2	669.71	45.5	100.5	216.67	175.6	293.1	396.52
4	921.1	1615.3	2413.61	120.1	233.0	342.34	18.9	46.2	99.14	89.1	148.0	209.18
8	527.9	938.2	1394.60	59.4	117.4	172.91	9.1	24.8	36.04	44.9	74.0	80.76
16	253.1	347.1	640.2	30.5	58.3	85.6	4.7	14.1	27.1	22.7	36.7	52.5
32	132.6	206.9	338.6	15.3	28.5	41.6	2.4	6.2	13.9	11.3	18.0	25.8
64	66.5	99.8	152.3	7.5	14.3	20.7	1.2	3.1	7.0	5.6	9.0	12.8
128	30.3	49.5	77.6	3.6	7.1	10.3	0.6	1.6	3.5	2.8	4.5	6.4
256	11.7	23.3	38.7	1.6	3.4	5.1	0.3	0.8	1.6	1.3	2.2	3.2
512	4.8	10.8	18.5	0.8	1.7	2.5	0.1	0.4	0.8	0.6	1.1	1.6
Decoder	Ace	Xperia	S3	Ace	Xperia	S3	Ace	Xperia	S3	Ace	Xperia	S3
2	2079.0	4005.0	5545.00	148.6	345.3	477.57	53.5	120.7	203.53	187.6	325.5	440.20
4	944.6	1774.0	2499.21	92.2	200.7	294.45	17.9	50.8	102.71	96.8	166.1	231.63
8	561.2	998.8	1384.16	51.9	108.9	158.78	8.7	24.4	45.01	51.1	82.2	103.40
16	287.0	527.3	783.7	27.7	56.2	82.7	5.1	13.1	25.7	25.0	39.9	56.6
32	150.6	256.2	403.2	14.4	27.8	40.3	2.5	7.3	13.3	12.0	18.7	26.7
64	67.6	107.8	169.6	7.3	13.8	20.0	1.2	3.5	6.9	5.6	8.6	12.2
128	31.9	49.8	77.3	3.4	6.6	9.6	0.6	1.6	3.4	2.4	3.7	5.4
256	12.7	24.4	38.4	1.4	3.0	4.5	0.3	0.7	1.4	0.9	1.5	2.1
512	4.8	10.7	18.1	0.6	1.3	1.9	0.1	0.3	0.6	0.3	0.5	0.8
Power (W)	0.486	0.474	1.012	0.421	0.377	1.057	0.279	0.378	1.041	0.425	0.448	1.090

