

## HOMEWORK 5 SOLUTIONS

### 1. No change to the array

j	Action	Array Contents
0	$A[0] = 1 \leq 8$ $i = 0$ Exchange $A[0]$ with $A[0]$	{ 1, 5, 3, 2, 0, 4, 6, 8 }
1	$A[1] = 5 \leq 8$ $i = 1$ Exchange $A[1]$ with $A[1]$	{ 1, 5, 3, 2, 0, 4, 6, 8 }
2	$A[2] = 3 \leq 8$ $i = 2$ Exchange $A[2]$ with $A[2]$	{ 1, 5, 3, 2, 0, 4, 6, 8 }
3	$A[3] = 2 \leq 8$ $i = 3$ Exchange $A[3]$ with $A[3]$	{ 1, 5, 3, 2, 0, 4, 6, 8 }
4	$A[4] = 0 \leq 8$ $i = 4$ Exchange $A[4]$ with $A[4]$	{ 1, 5, 3, 2, 0, 4, 6, 8 }
5	$A[5] = 4 \leq 8$ $i = 5$ Exchange $A[5]$ with $A[5]$	{ 1, 5, 3, 2, 0, 4, 6, 8 }
6	$A[6] = 6 \leq 8$ $i = 6$ Exchange $A[6]$ with $A[6]$	{ 1, 5, 3, 2, 0, 4, 6, 8 }
7	(Exit the loop) Exchange $A[7]$ with $A[7]$ Return 7	{ 1, 5, 3, 2, 0, 4, 6, 8 }

3.

One possible  $O(n)$  solution: Have two pointers, one starts from the front and the other starts from the back.

while front < back

- if front is -ve and back is +ve, increment front and decrement back
- If front is +ve and back is -ve, swap them, increment front and decrement back
- If front is -ve but back is also -ve, increment front
- If front is +ve and back is +ve, decrement back

Another possible  $O(n)$  solution is to start both pointers at the front. One pointer keeps moving forward, and when it finds a negative number, it swaps it with the other pointer and both pointers are incremented.

Another possible option is to use one round of QuickSort Partition using 0 as the pivot value

4.

#	Solution/Points
1	Sorting is beneficial to find the mth smallest item. Total runtime would be $O(n \log n)$ (to sort) + $O(m) = O(n \log n)$ . Without sorting, it could take $O(mn)$ which is quadratic.
2	Can be done in $O(n \log n)$ with sorting. Sort the array, then loop through it and keep track of the max count. Total run time would be $O(n \log n) + O(n) = O(n \log n)$ . Without sorting, it could take upto $O(n^2)$ . (You may also do it without sorting, by using Hashing. In that case it would only take $O(n)$ to find the mode)
3	Yes sorting helps. Sort both arrays merge to find the median. Can be done in $O(n \log n)$ with sorting.
4	No sorting doesn't help in this case. Can be done in $O(n)$ without sorting, by just keeping track of a counter, looping through the array and incrementing the counter every time the particular item is seen. Sorting on the other hand would take $O(n \log n)$ runtime, which is less efficient.
5	Nope no sorting here. Can be done in just one iteration i.e. $O(n)$ , by looping through the array and keeping track of the minimum item seen so far. Sorting would take a minimum of $O(n \log n)$ which is less efficient