Lecture 25: The floor function

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Lemma. For any real number x, there is a unique integer n such that n < x < n+1.

(In mathematical language, $\forall x \in \mathbb{R}, \exists ! n \in \mathbb{Z}, n \leq x < n+1.$)

Proof. Existence. Case 1. x≥0.

Consider the set $2m \in \mathbb{Z} \mid o \le m \le x$. This is a finite set. So it has a maximum (this can be proved for any finite subset of \mathbb{R} using induction on the cardinality of the finite set.) Let $n = \max 2m \in \mathbb{Z} \mid o \le m \le x$. So

- 1) $0 \le n \le x$ and 2) $n+1 \notin \{m \in \mathbb{Z} \mid 0 \le m \le x\}$.
- 2 implies that either n+1 < 0 or n+1 > x.
- ① implies $n+1 \ge 1$. So n+1 > x, Hence, by ①, we get $n \le x < n+1$.

Case 2. X ∈ Z.

In this case x <x <x+1 and n=x works.

Case 3. X < 0 and X & Z.

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 $x < \infty \implies -x > \infty \implies \text{ by case 1, there is } m \in \mathbb{Z} \text{ such that}$ $m \le -x < m+1.$

Since $x \notin \mathbb{Z}$, $m \neq -x$. Hence m < -x < m+1. Therefore -(m+1) < x < -m = -(m+1)+1.

 \Rightarrow $-(m+1) \leq \chi < -(m+1)+1$.

So n = -(m+1) works.

Suppose to the contrary that for some $n_1, n_2 \in \mathbb{Z}$.

we have $n_1 \neq n_2$ and $n_1 \leq x < n_1 + 1$ and $n_2 \leq x < n_2 + 1$.

So either $n_1 > n_2$ or $n_2 > n_1$. By symmetry, it is enough

to deal with the case n₁>n₂.

 $\Rightarrow \chi \geq n_1 \geq n_2 + 1 > \chi \Rightarrow \chi > \chi$ which is a contradiction.

Lecture 25: Division algorithm

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Theorem (Division algorithm) For any integers a and b,

 $b \neq 0$, there is a unique pair of integers (q, r) such that

(1)
$$a = bq + r$$
 (2) $0 \le r < |b|$.

Proof. Case 1. b>0.

Existence Claim $q = \lfloor a/b \rfloor$ and r = a - bq is such a pair.

- · q is an integer by the definition of the floor function.
- $a, b, q \in \mathbb{Z} \implies r = a bq \in \mathbb{Z}$
- Property (1) is clear: $r = a bq \implies a = bq + r$.

Now we show Property (2):

$$\lfloor \frac{\alpha}{b} \rfloor \leq \frac{\alpha}{b} < \lfloor \frac{\alpha}{b} \rfloor + 1 \implies q \leq \frac{\alpha}{b} < q + 1$$

Since b>0, we get bq < a < bq+b

$$\Rightarrow$$
 $0 \le a - bq < b$ (adding - bq to all sides.)

$$\Rightarrow$$
 $0 \le r < b = |b|$.

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Case 2. b<0.

Claim q = -[-9], r = a - bq satisfy the mentioned properties.

- . As before we can see q and r are integers which satisfy the 1^{st} property.
 - . Now we show Property (2):

$$\lfloor -\frac{\alpha}{b} \rfloor \leq -\frac{\alpha}{b} < \lfloor -\frac{\alpha}{b} \rfloor + 1 \implies -q \leq -\frac{\alpha}{b} < -q + 1$$

Since -b>o, we get

$$(-b)(-q) \le (-b)(-\frac{a}{b}) < (-b)(-q+1)$$

$$\Rightarrow$$
 bq $\leq a < bq - b - bq + |b|$

$$\Rightarrow$$
 o $\leq a - bq < |b|$

Uniqueness. We have to prove

$$\begin{pmatrix}
a = bq_1 + r_1 & and & a = bq_2 + r_2 \\
o \le r_1 < |b| & and & o \le r_2 < |b|
\end{pmatrix} \Rightarrow \begin{cases}
q_1 = q_2, \\
r_1 = r_2.
\end{cases}$$

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$$a = b q + r; \Rightarrow \begin{cases} \frac{\alpha}{|b|} = \frac{b q_i}{|b|} + \frac{r_i}{|b|} \\ 0 \le r_i < |b| \end{cases}$$

$$\Rightarrow \frac{b}{|b|} q_{i} \leq \frac{b}{|b|} q_{i} + \frac{r_{i}}{|b|} < \frac{b}{|b|} q_{i} + 1$$

$$\sim \frac{a}{|b|}$$

Notice that
$$\left(\frac{b}{|b|} = 1 \text{ if } b>0\right)$$
 and $\left(\frac{b}{|b|} = -1 \text{ if } b<0\right)$.

In particular, by q e Z. So

$$\frac{b}{|b|} q_i = \lfloor \frac{9}{|b|} \rfloor \Rightarrow q_i = \frac{|b|}{b} \lfloor \frac{9}{|b|} \rfloor$$

[These are true for i=1 and i=2.] So

$$q_1 = \frac{|b|}{b} \left[\frac{a}{|b|} \right] = q_2$$

Hence
$$r_1 = a - bq_1 = a - bq_2 = r_2$$
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