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- 1) The first line, assignment, is  $\theta(1)$ . The second line, a for loop, is 100n + 1. The third line is 100n. Therefore the  $1+100n+100n = 200n+2 = \theta(n)$ .
- 2) The first line, assignment is theta(1). The second line, a for loop, is  $(n^3)/4 + 1$  because for each iteration, i decreases for 4. Likewise for the  $3^{rd}$  line is  $(n^3)/4$ . Therefore  $1+2(n^3/4)+1=\theta(1/2*n^3+2)=\theta(n^3)$ .
- 3) The first line is 1. The second line is  $\log(n)$  because for every iteration, it decreases by half. The  $3^{rd}$  line is also  $\log(n)$  times. Hence  $1 + \log(n) + \log(n) + 1 = \log(2n) + 2 = \theta(\log(2n))$ .
- 4) The first line, assignment is theta(1). The second line, is 1. The third line is 1 because even though j increases by a multiple of 2 on each iteration, the total number of iteration is constant. The third line is a constant 1.  $1+1+1+1=1=1=100*\log 100+2=0(\log 100)$ .
- 5) The first line, which is a assignment, is 1. The second line, which iterates n times is n. The third line, which iterates for every I, is n/2. Therefore, the third line, num++ would be n/2 also. Thus the runtime of this algorithm is  $1 + n + n/2 + n/2 = 2n+1 = \theta(n)$ .
- 6) the first line is 1. The second line is log n because of the multiple of i by 2 for each iteration. The third line is n/4 because for every iteration, j increases by 4. hence the  $3^{rd}$  line is also n/4. Thus runtime is  $1 + n/4*log(n) + 1 + n/4*1 = n/4*log(n)+3+n/4 = \theta(nlog(n))$ .
- 7) The first for loop is 2n+1 because it iterates through all of 2n+1 for checking.  $2^{Nd}$  line is 2n.  $3^{Rd}$  line is n/3+2 because the  $4^{th}$  line is n; the last line is n also. Thus the runtime is  $2n+1+n/3+2+n=n/3+2n+3=\theta(n/3)$ .
- 8) The first line is 1. The  $2^{nd}$  line is (n-1+1) \* (