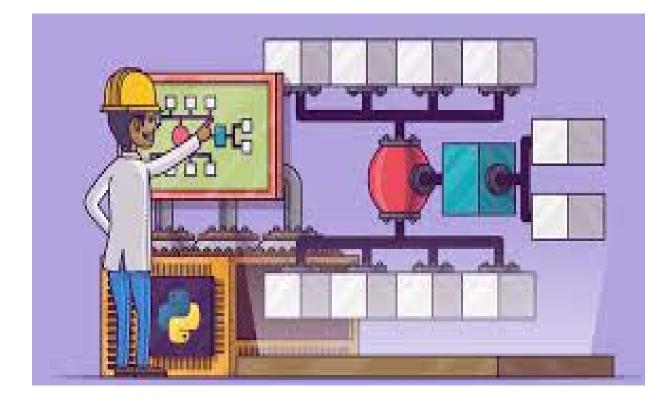


# ساختمان داده ها

مدرس: سمانه حسینی سمنانی

دانشگاه صنعتی اصفهان- دانشکده برق و کامپیوتر





- مفاهيم اوليه
- پیمایش درخت
- درخت دودویی معادل



مفاهيم اوليه



## مفاهيم اوليه

- درخت آزاد
- درخت ریشه دار
  - اولاد یک گره
  - اجداد یک گره
    - جنگل
- رابطه پدر و فرزندی
- زیر درخت ریشه دار
  - درجه هر نود



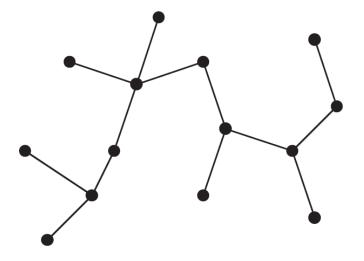
## مفاهيم اوليه

- عمق یک گره
- ارتفاع نود V
- ارتفاع درخت
- درخت متوازن
- درخت کاملا متوازن
  - درخت **k** تایی
    - درخت مرتب



• درخت یا درخت آزاد

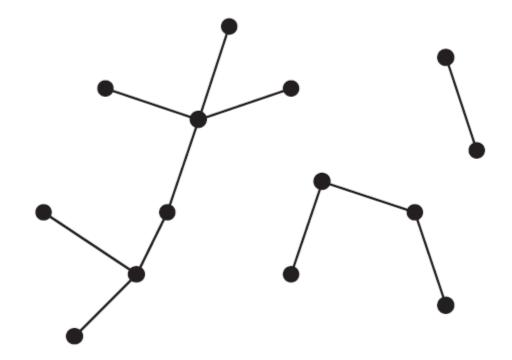
• A free tree is a connected, acyclic, undirected graph.





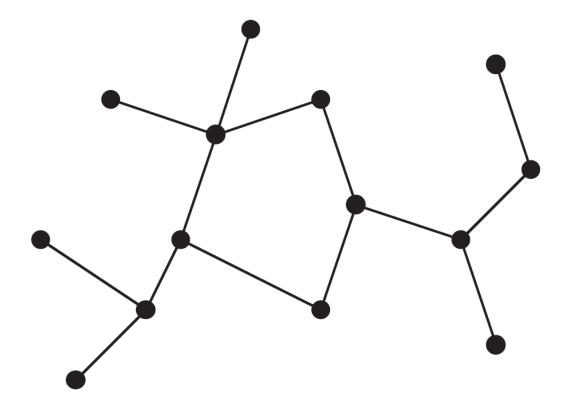


• If an undirected graph is acyclic but possibly disconnected, it is a forest.





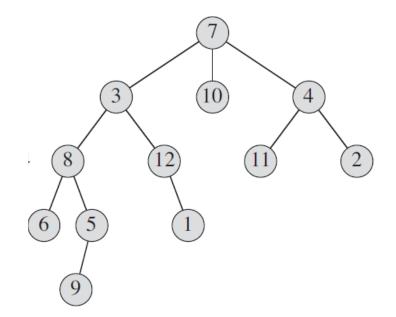
Tree or forest?





# درخت ریشه دار

- A **rooted tree** is a free tree in which one of the vertices is distinguished from the others (root).
- We often refer to a vertex of a rooted tree as a node of the tree.

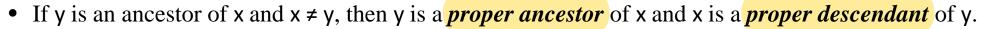


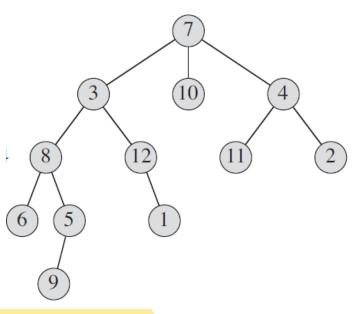
- يال ها جهت دهي مي شوند.
- از بالا به پایین در شکل روبرو.
  - درجه ریشه = صفر
- درجه ورودی هر نود غیر از ریشه = یک
- از اینجا به بعد وقتی میگیم درخت، منظور درخت ریشه دار است.



# رابطه پدر و فرزندی

- Consider a node x in a rooted tree T with root r.
- We call any node y on the unique simple path from r to x an *ancestor* of x.
- If y is an ancestor of x, then x is a *descendant* of y.
- Every node is both an ancestor and a descendant of itself.

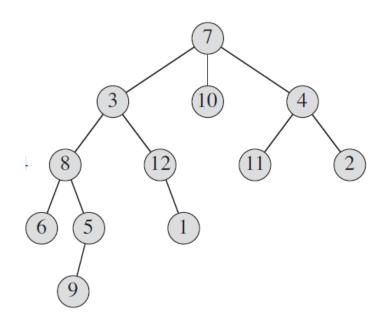






#### رابطه پدر و فرزندی

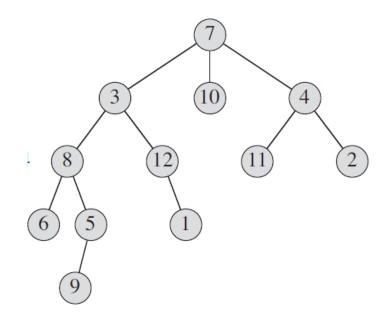
- If the last edge on the simple path from the root r of a tree T to a node x is (y,x), then y is the *parent* of x, and x is a *child* of y.
- The root is the only node in T with no parent.
- If two nodes have the same parent, they are *siblings*.
- A node with no children is a *leaf* or *external node*
- A non leaf node is an *internal node*.





#### زیر درخت ریشه دار

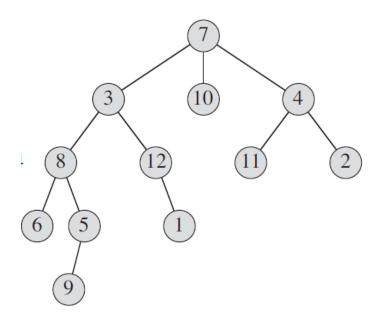
- The *subtree rooted at* x is the tree induced by descendants of x, rooted at x.
- For example, the subtree rooted at node 8 contains nodes 8, 6, 5, and 9.





# درجه هر نود

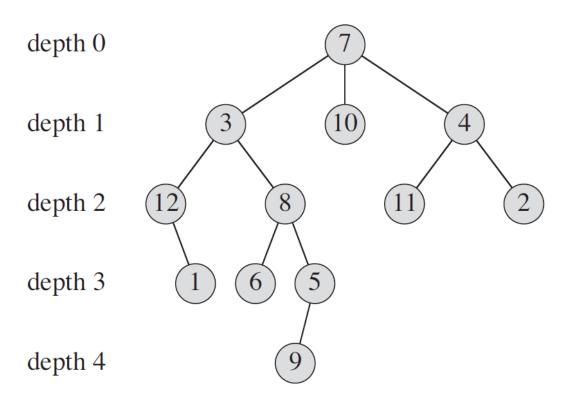
• The number of children of a node x in a rooted tree T equals the *degree* of x.





# عمق هر نود

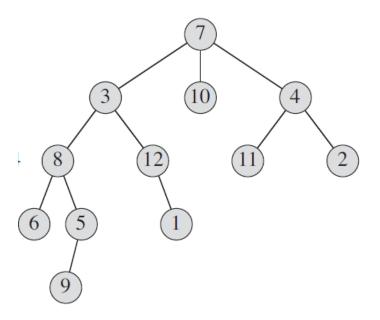
• The length of the simple path from the root r to a node x is the *depth* of x in T.





# ارتفاع هر نود

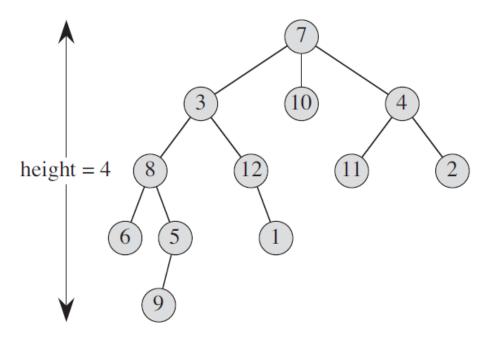
• The *height* of a node in a tree is the number of edges on the longest simple downward path from the node to a leaf, and the height of a tree is the height of its root.





#### ارتفاع درخت

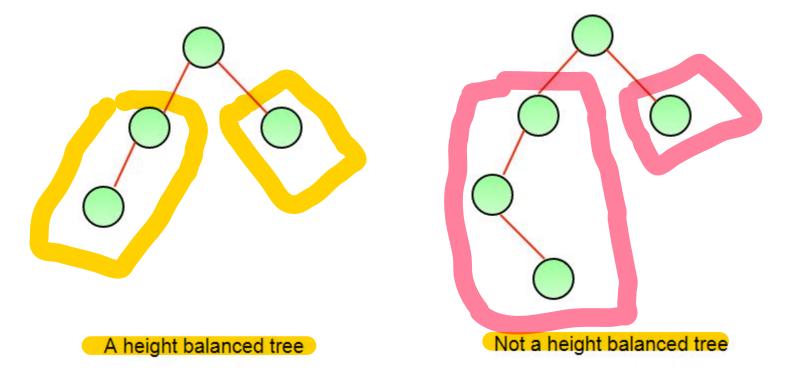
- *Height* of a tree is the height of its root.
- The height of a tree is also equal to the largest depth of any node in the tree.





# درخت متوازن

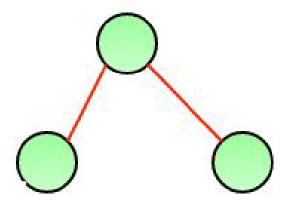
• A balanced tree is a tree where The left and right subtrees' heights differ by at most one.





# درخت كاملا متوازن

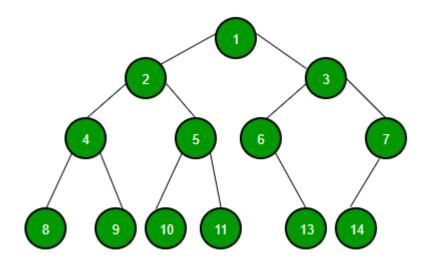
• A complete balanced tree is a tree where The left and right subtrees' heights differ by at most zero.





# درخت k تایی

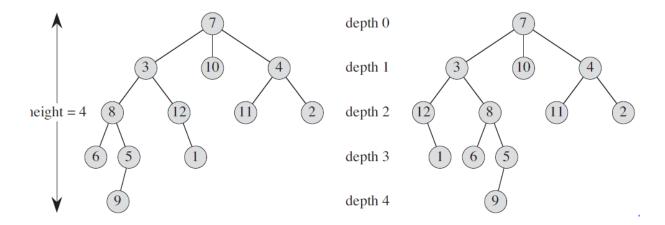
• درخت k تایی درختی است که در آن هر گره حداکثر k گره فرزند دارد.





# درخت مرتب

- An *ordered tree* is a rooted tree in which the **children** of each node are ordered.
- That is, if a node has k children, then there is a first child, a second child, . . . , and a  $k_{th}$  child.



As a rooted tree, it is identical to the tree in (a), but as an ordered tree it is different, since the children of node 3 appear in a different order.



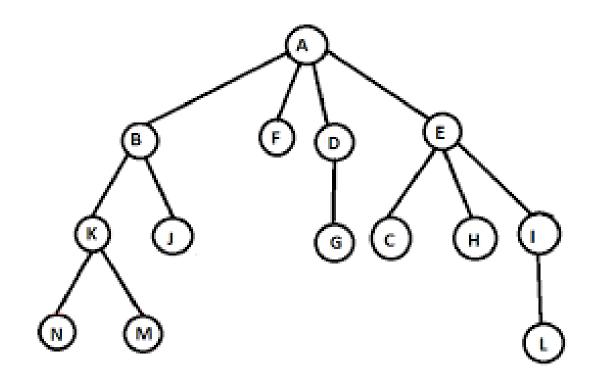
پیمایش درخت

Tree traversal



# روش های مختلف پیمایش درخت

- Random
- Preorder
- Inorder
- Postorder
- Mixed

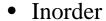




### روش های مختلف پیمایش درخت

#### • Preorder

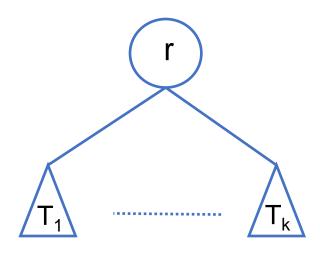
$$r$$
 ,  $T_1$   $T_2$   $\cdots$   $T_k$ 



$$T_1$$
,  $r$ ,  $T_2 \cdots T_k$ 



$$T_1 \cdots T_k r$$

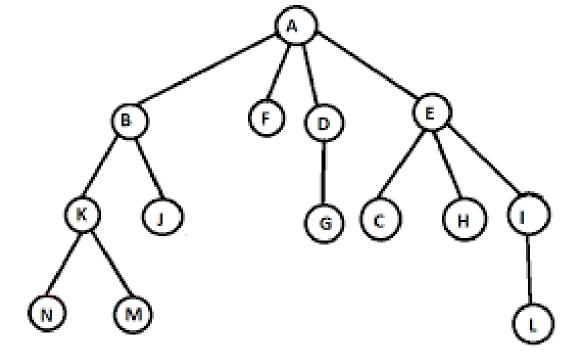


بازگشتی



#### مثال

- Preorder: A, B, K, N, M, J, F, D, G, E, C, H, I, L
- Inorder: N, K, M, B, J, A, F, G, D, C, E, H, L, I
- Postorder: N, M, K, J, B, F, G, D, C, H, L, I, E, A



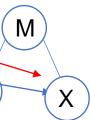


#### مثال

درخت دودویی

- Preorder: M, N, H, C, R, S, K, W, T, G, D, X, I, Y, A, J, P, O, E, Z, V, B, U, L, Q, F
- Postorder: C, W, T, K, S, G, R, H, D, N, A, O, E, P, J, Y, Z, I, B, Q, L, F, U, V, X, M

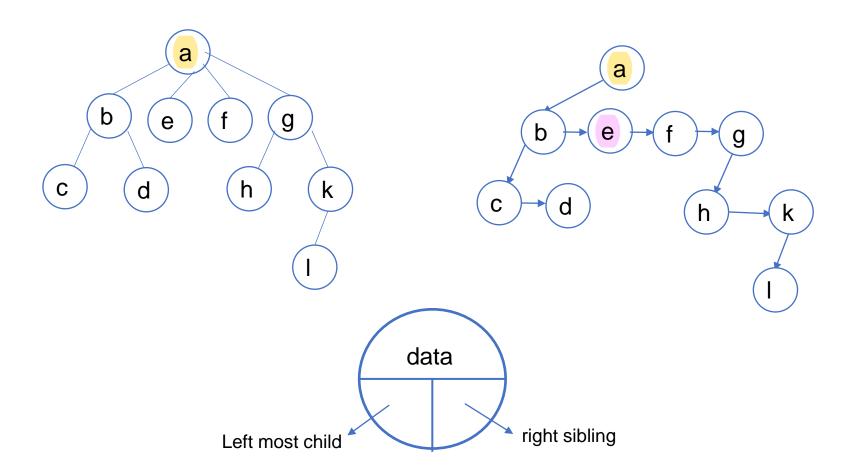
• Inorder?



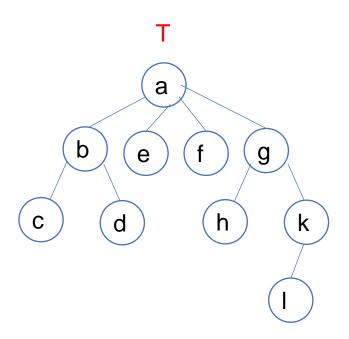


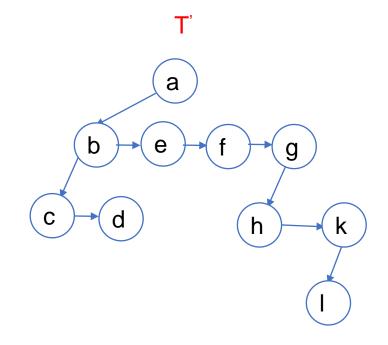
# درخت دودویی معادل









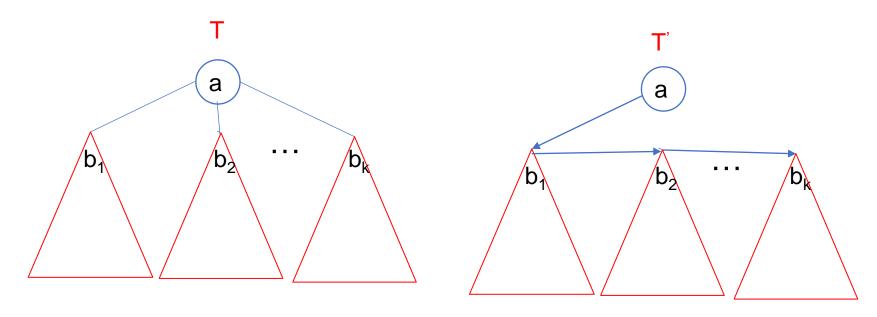


preorder(T) = a, b, c, d, e, f, g, h, k, l

preorder(T') = a, b, c, d, e, f, g, h, k, l

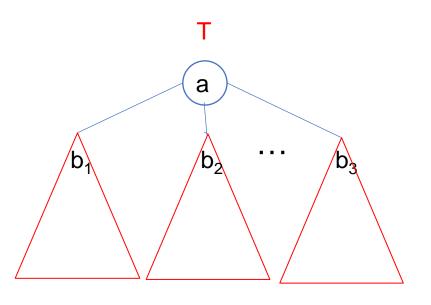
- چرا برابر شد؟
- آیا در حالت کلی این طور است؟

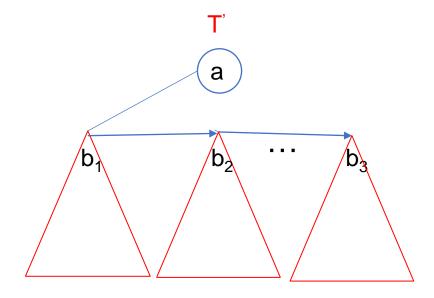




- چرا برابر شد؟
- آیا در حالت کلی این طور است؟







$$\begin{array}{c}
?\\
postorder(T') = postorder((T))\\
?\\
inorder(T) = inorder(T')
\end{array}$$