Chapter 1 Introduction

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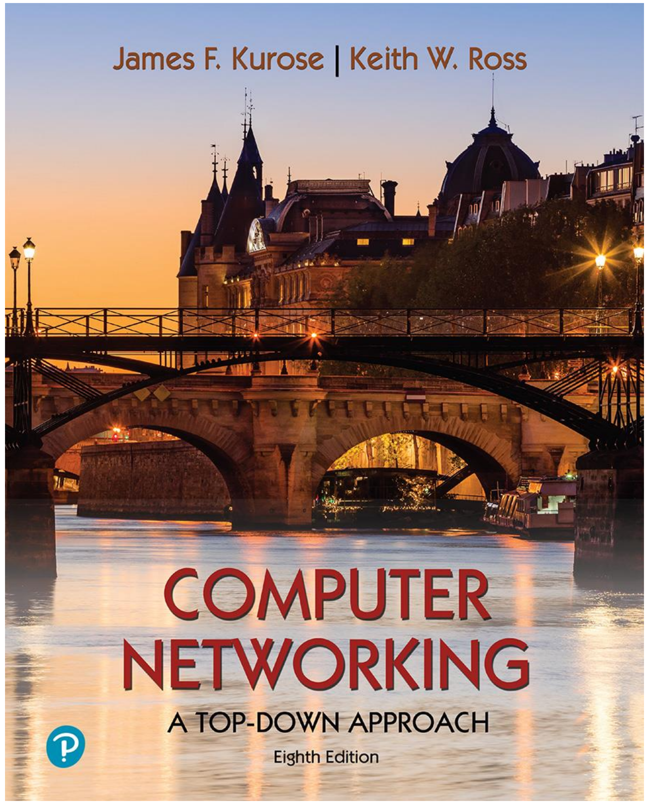
Thanks and enjoy! JFK/KWR



Computer Networking: A Top-Down Approach
8th edition
Jim Kurose, Keith Ross
Pearson, 2020

Chapter 1: introduction

Chapter goal:

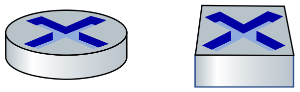
- Get "feel," "big picture," introduction to terminology
  - more depth, detail *later* in course

Overview/roadmap:

- What *is* the Internet? What *is* a protocol?
- Network edge: hosts, access network, physical media
- Network core: packet/circuit switching, internet structure
- Performance: loss, delay, throughput
- Protocol layers, service models
- Security
- History

The Internet: a "nuts and bolts" view

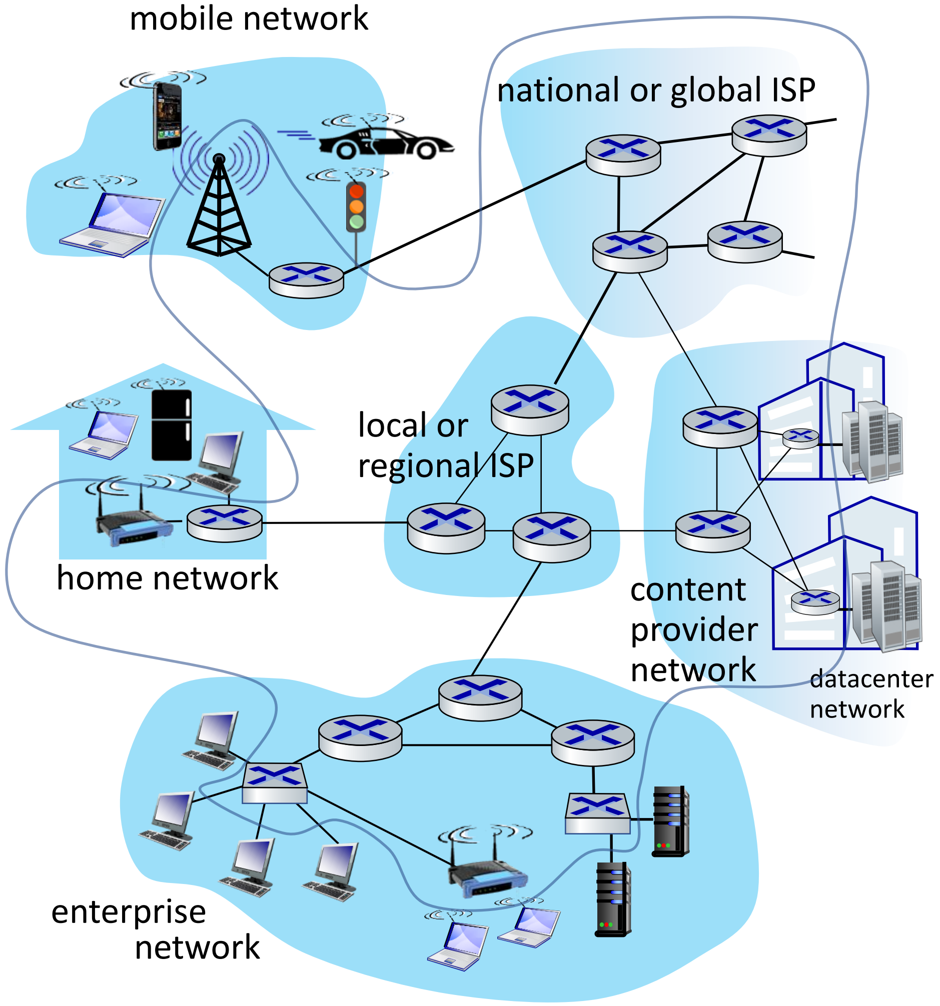
Billions of connected computing *devices*:
- *hosts = end systems*
- running *network apps* at Internet's "edge"



Packet switches: forward packets (chunks of data)
- *routers, switches*

Communication links
- fiber, copper, radio, satellite
- transmission rate: *bandwidth*

Networks
- collection of devices, routers, links: managed by an organization



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اینترنت چی هست؟
اولین تعریف براساس اجزای اینترنت است:
در لبه اینترنت ما دیوایس های کامپیوتری رو داریم که از طریق اینترنت بهم وصل شدند که بهشون host or end systems میگیم
لینک های ارتباطی به شکل های مختلف به اینترنت وصل می شوند - این لینک ارتباطی یک عرض باندی در اختیار ما می گذاره که با استفاده از اون اطلاعات رو رد و بدل میکنیم
اجزای داخلی شبکه که این اطلاعات را رد و بدل میکنند که بهشون سویچ یا روتر می گیم
این اطلاعات به صورت دیجیتال هستند ینی صفر و یک و بعد دسته بندی میشن و در قالب دسته هایی که بهشون packet switches میگیم ارسال میشن و سوییچ ها و روترها به عنوان نودهای شبکه این ها رو رد وبدل می کنند و سوییچ می کنند که به مقصد برسند

“Fun” Internet-connected devices

Amazon Echo

Internet refrigerator

IP picture frame

Pacemaker & Monitor

Tweet-a-watt: monitor energy use

bikes

cars

scooters

Web-enabled toaster + weather forecaster

Slingbox: remote control cable TV

Security Camera

AR devices

Internet phones

Gaming devices

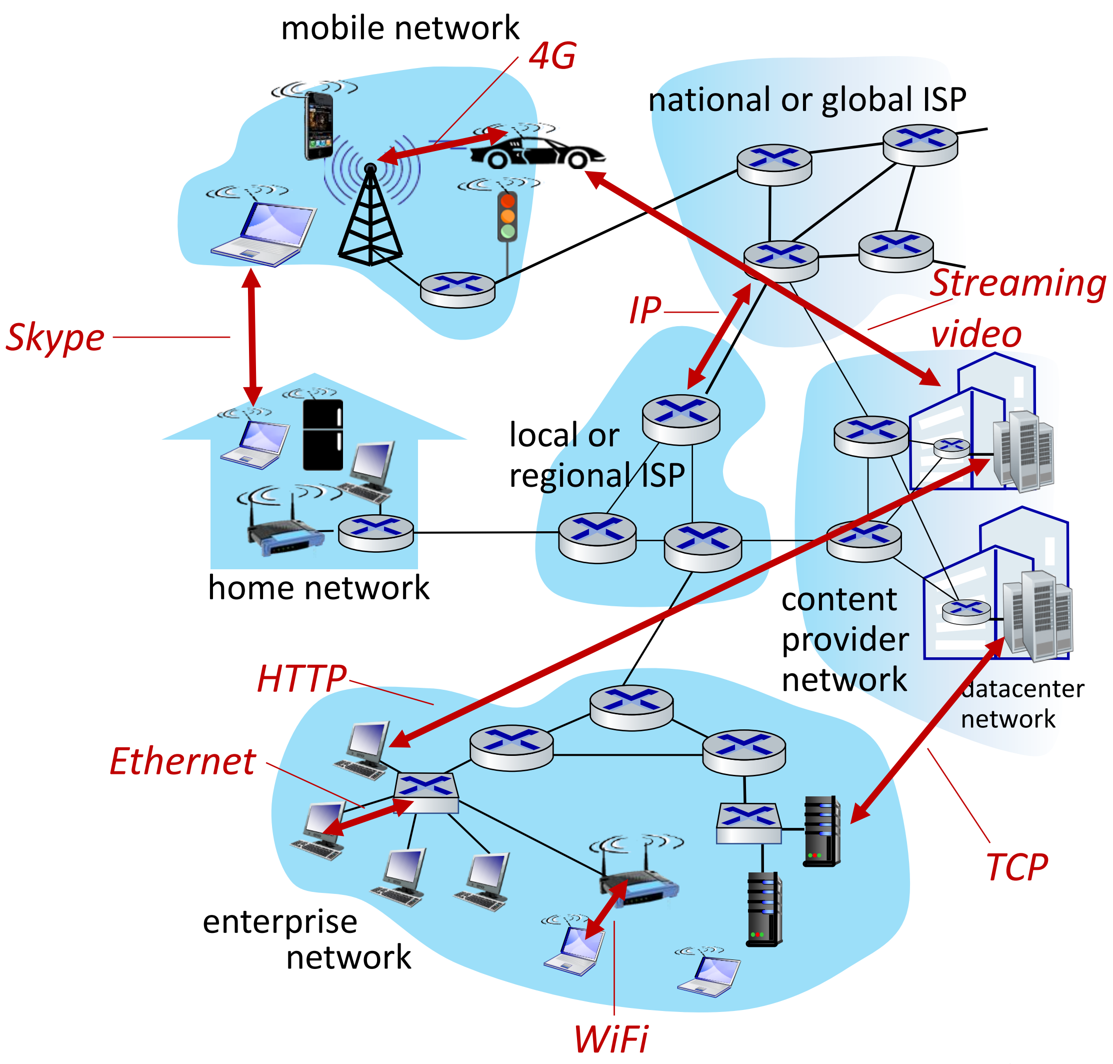
sensorized, bed mattress

Fitbit

Others?

The Internet: a "nuts and bolts" view

- *Internet: "network of networks"*
  - Interconnected ISPs
- *protocols* are *everywhere*
  - control sending, receiving of messages
  - e.g., HTTP (Web), streaming video, Skype, TCP, IP, WiFi, 4G, Ethernet
- *Internet standards*
  - RFC: Request for Comments
  - IETF: Internet Engineering Task Force



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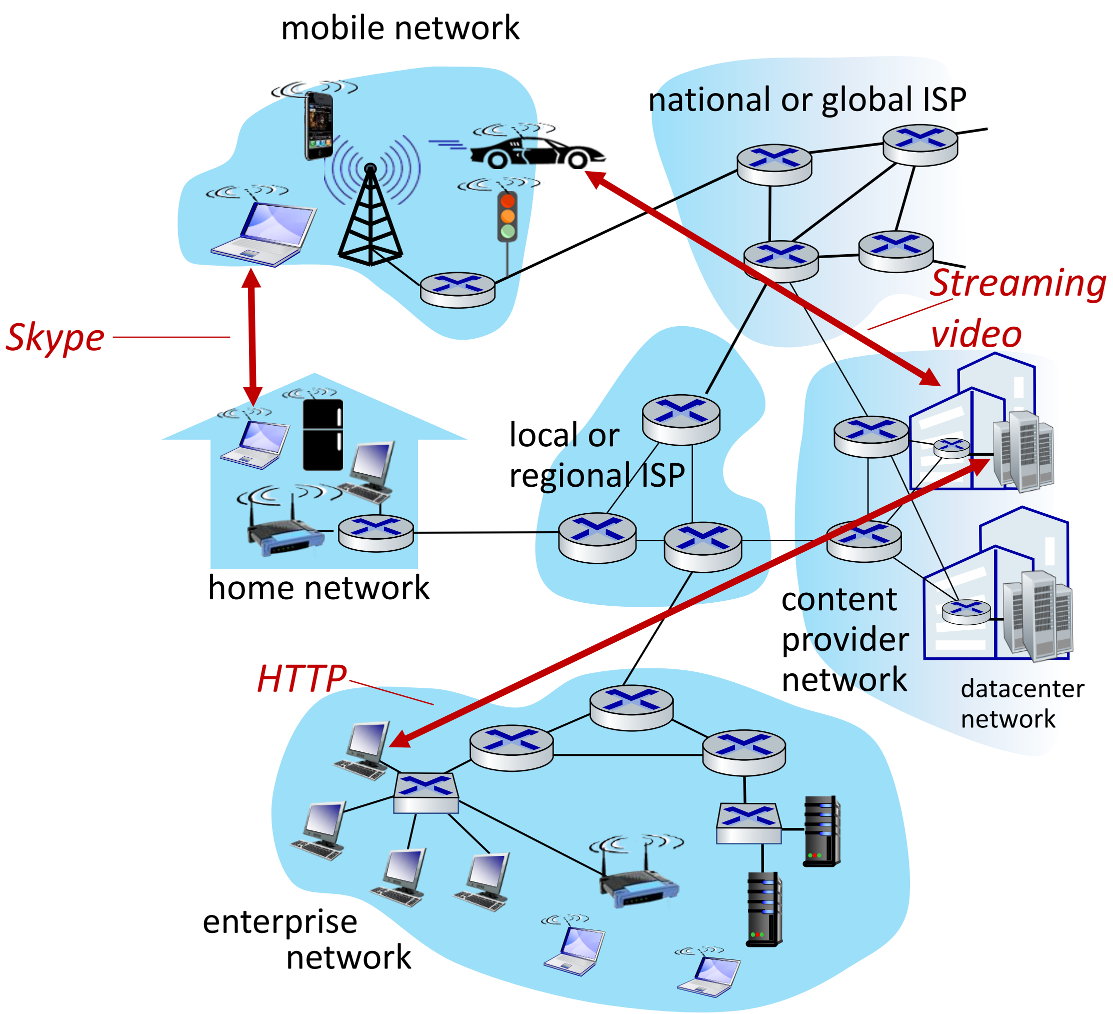
از نظر ساختاری تعریف دیگری از اینترنت: شبکه ای از شبکه ها است
برای این ارتباط هم ارتباط بین شبکه ها رو داریم و بعد پروتکل های ارتباطی داریم که براساس این پروتکل ها است که این عملکرد شبکه هم گسترش هم پیدا میکنه بین شبکه ها

The Internet: a "services" view

- *Infrastructure* that provides services to applications:
  - Web, streaming video, multimedia teleconferencing, email, games, e-commerce, social media, inter-connected appliances, …
- provides *programming interface* to distributed applications:
  - "hooks" allowing sending/receiving apps to "connect" to, use Internet transport service
  - provides service options, analogous to postal service



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تعریف سوم برای اینترنت:
نگاه سرویس و نگاه کاربرد
براساس این تعریف اینترنت یک زیرساختی است که یکسری سرویس یا کاربردهایی رو ارائه رو میده مثل وب سایت ها، ایمیل یا ....
اینترنت اون محیطی است که اینترفس های لازم رو در اختیار می ذاره برای اینکه این اپلیکیشن ها و این سرویس ها روی اینترنت بتونه برقرار بشه و این سرویس به کاربر ارائه میشه
این تعریف یک تعریف انقلابی است برای اینترنت ینی از زمانی که این نگاه شکل گرفت اینترنت گسترش پیدا کرد
پس این نگاه میگه که شبکه یک چیزی است که این سرویس ها رو به ما ارائه میده ینی سرویس وب، ایمیل، voip و ...

What's a protocol?

Human protocols:

- "what's the time?"
- "I have a question"
- introductions

Rules for:

… specific messages sent

… specific actions taken when message received, or other events

Network protocols:

- computers (devices) rather than humans
- all communication activity in Internet governed by protocols

Protocols define the format, order of messages sent and received among network entities, and actions taken on message transmission, receipt

-

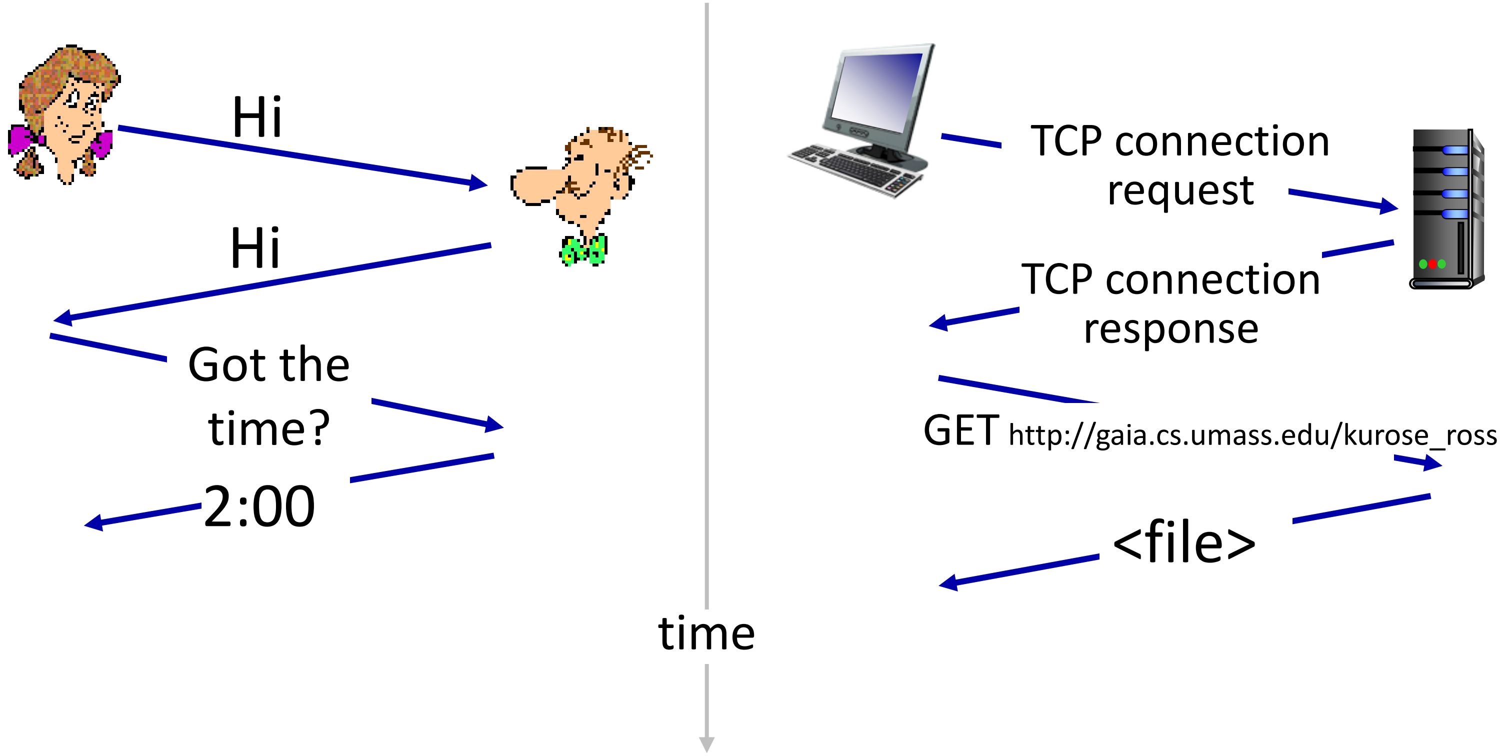
پروتکل ینی یک کاری رو به چه شکلی انجام بدیم

هر کاری در اینترنت بخواد انجام بگیره براساس یک پروتکل باید باشه ینی نحوه انجام کار باید تعریف بشه و بعد روش توافق بشه که استاندارد بشه پس پروتکل های استاندارد لازمه کار شبکه هستند

پس پروتکل ترتیب انجام کار، شکل انجام کار، فرمت بسته ها، ترتیب مسیج ها و خود مسیج ها که چه چیزی رو دارند منتقل می کنند توی پروتکل مشخص میشه

What’s a protocol?

A human protocol and a computer network protocol:



Q: other human protocols?

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پروتکل ینی یک ارتباط آگاهانه می خواد برقرار بشه

Chapter 1: roadmap

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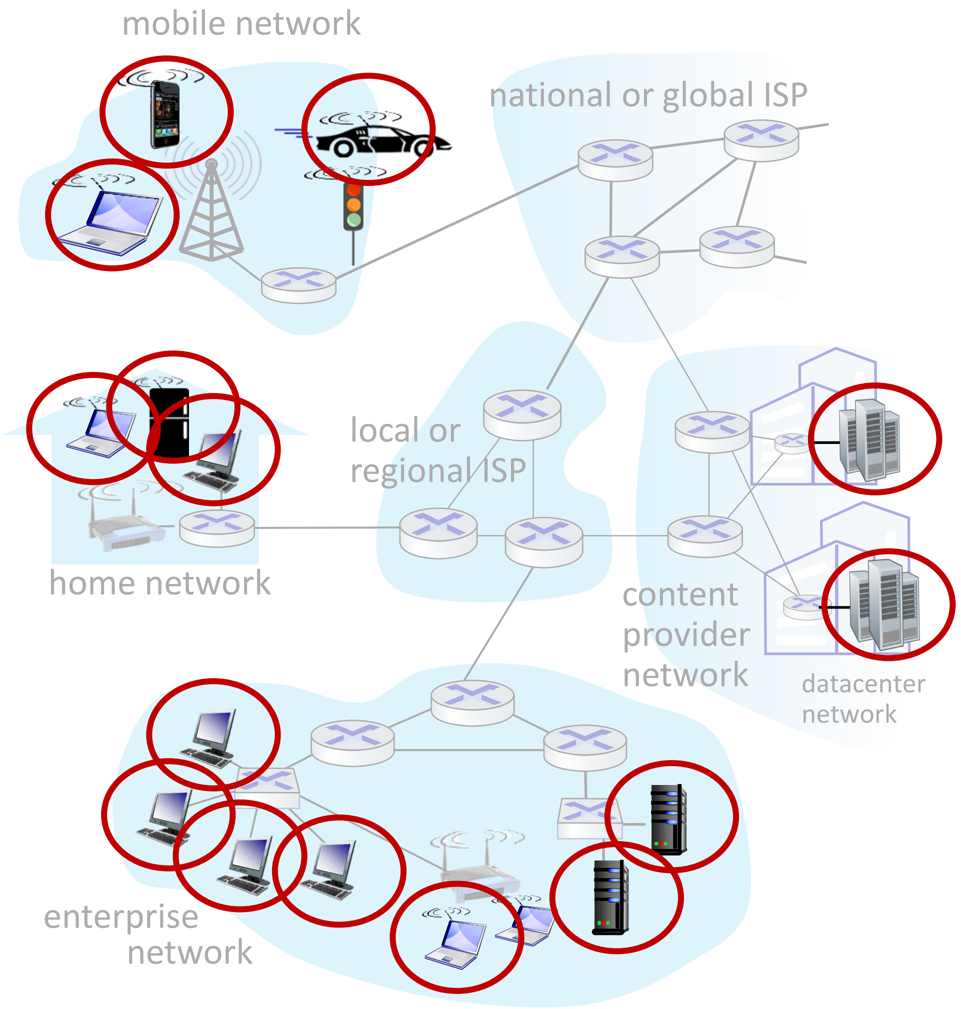
A closer look at Internet structure

Network edge:

- hosts: clients and servers
- servers often in data centers



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از نظر ساختاری و سازمانی گفتیم اینترنت شبکه ای از شبکه ها است ولی این شبکه ها همشون عین هم نیستند مثلا شبکه های edge داریم یا شبکه های access رو داریم یا شبکه های core داریم و این ها کارشون متفاوته و بر این اساس معماری هاشون هم متفاوت است

در کل شبکه ی شبکه ها از شبکه های مختلفی تشکیل شده که نوعشون و جایگاهشون متفاوت است

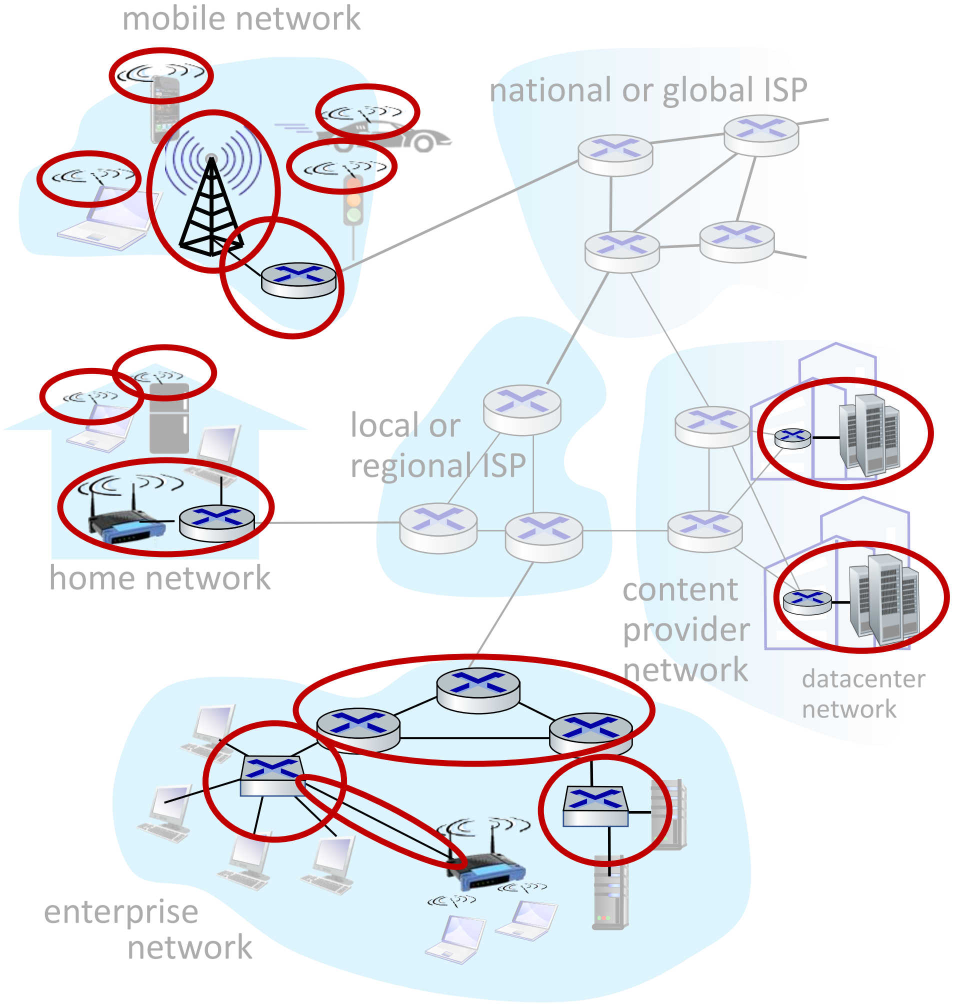
A closer look at Internet structure

Network edge:

- hosts: clients and servers
- servers often in data centers

Access networks, physical media:

- wired, wireless communication links



A closer look at Internet structure

Network edge:

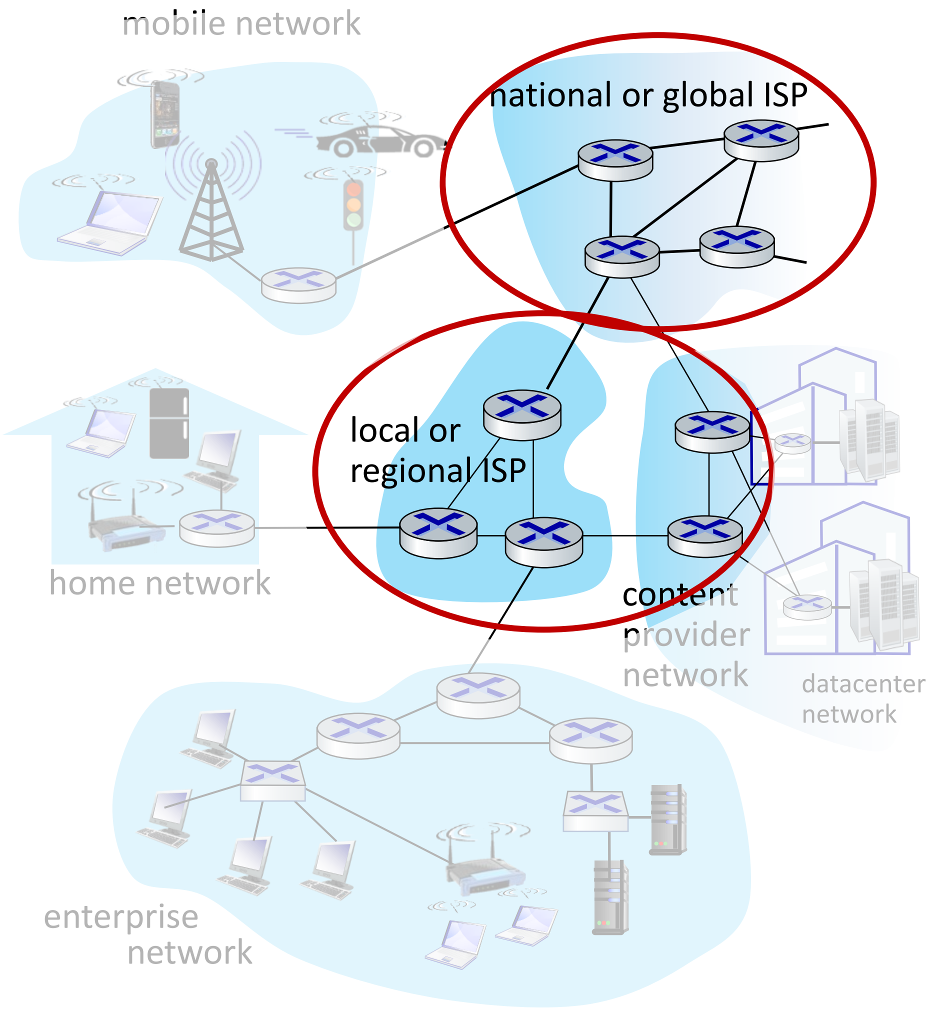
- hosts: clients and servers
- servers often in data centers

Access networks, physical media:

- wired, wireless communication links

Network core:

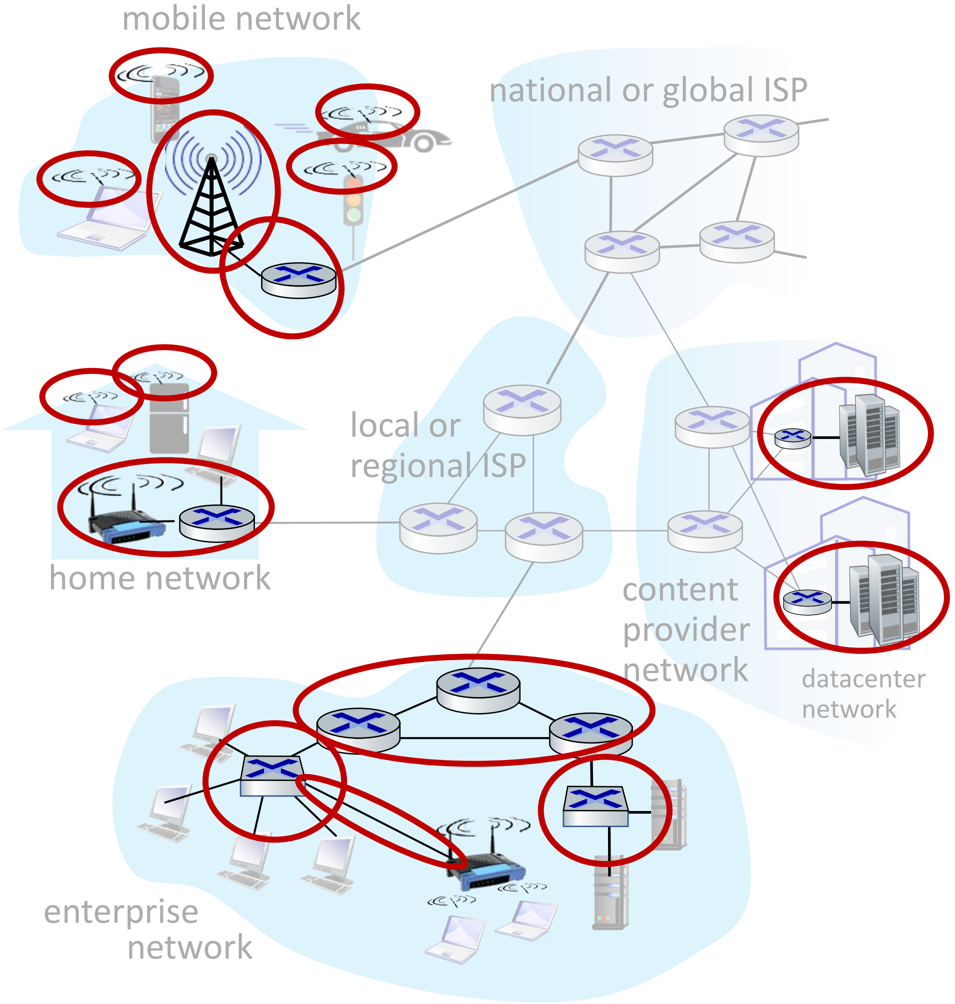
- interconnected routers
- network of networks



Access networks and physical media

Q: How to connect end systems to edge router?

- residential access nets
- institutional access networks (school, company)
- mobile access networks (WiFi, 4G/5G)



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access network ها شبکه های ارتباط مشتری با اینترنت است

ینی از طریق این ماشین انتهایی چجوری به شبکه وصل میشه؟

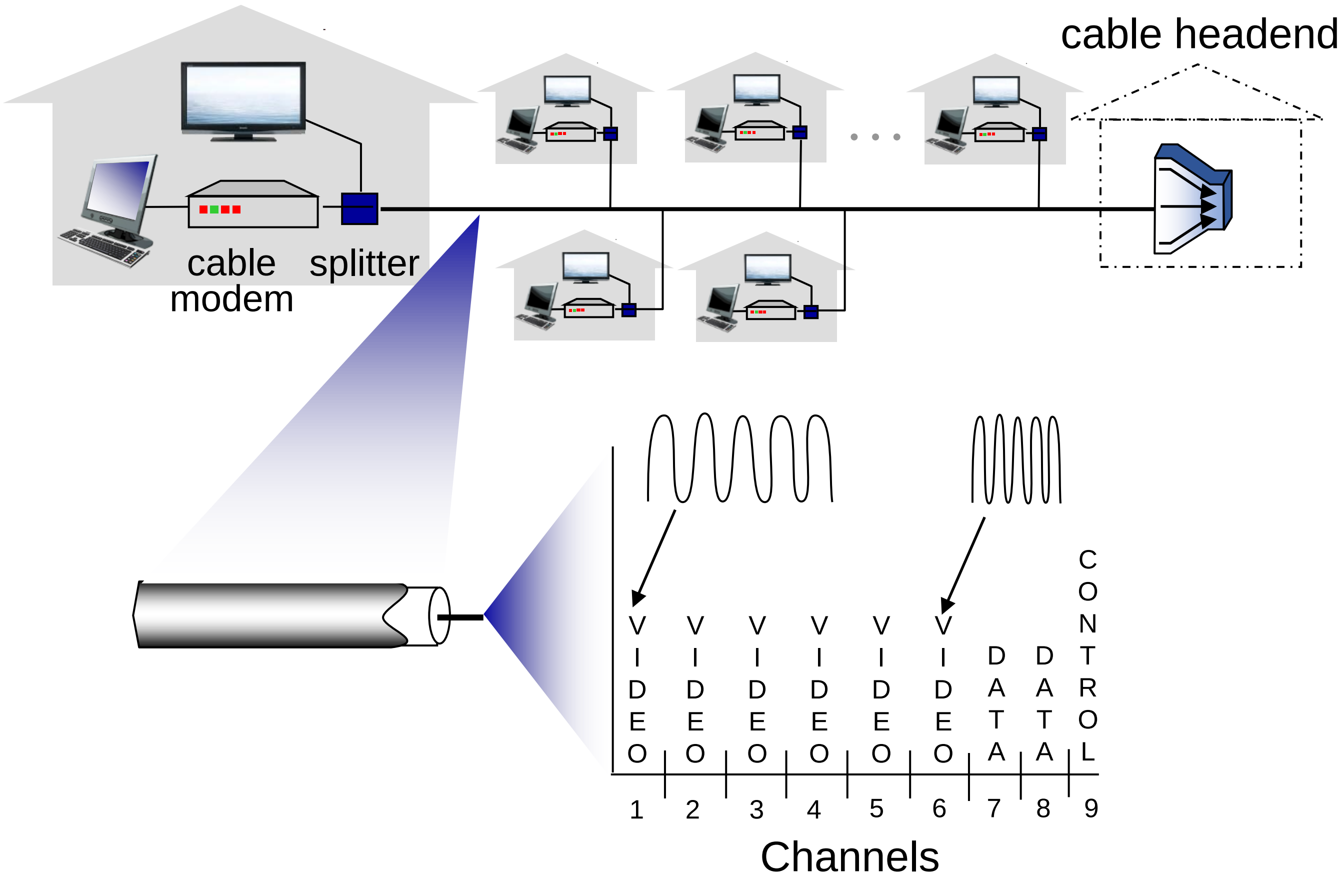
اون قسمت از شبکه که اتصال دیوایس انتهایی به شبکه رو تامین می کنه بهش میگیم access network

شاید متنوع ترین قسمت شبکه توی Access networks اش است

در خونه ها بیشتر ADSL است

یک مسئله مهم در access بحث پهنای باند است یا bit per second ینی چند بیت در ثانیه می تونه Access networks ما متحمل بشه و این خیلی روی کیفیت اپلیکیشن ها تاثیر می ذاره

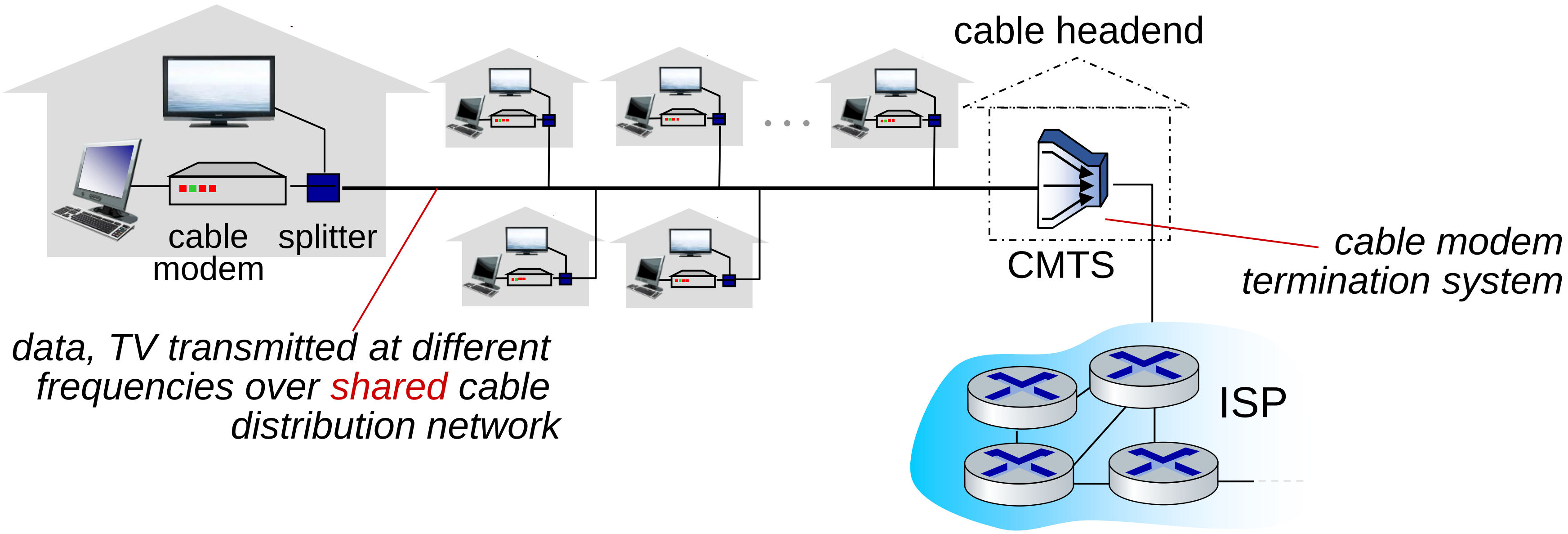
Access networks: cable-based access



frequency division multiplexing (FDM): different channels transmitted in different frequency bands

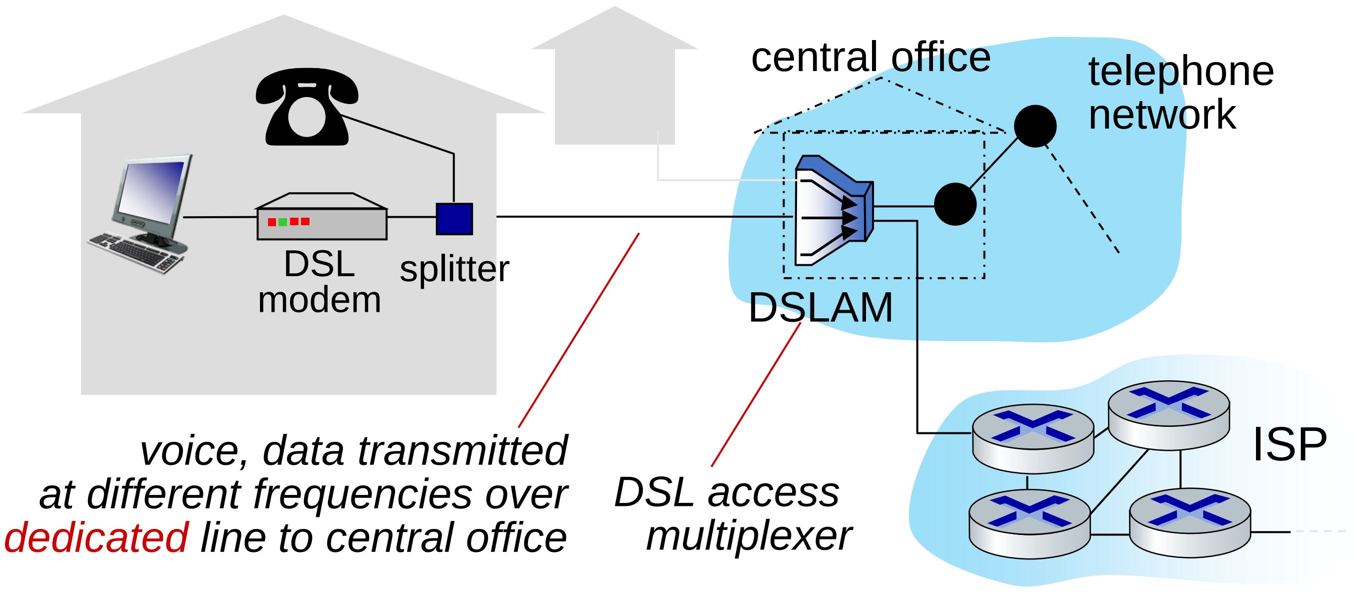
شبکه های دسترسی: دسترسی مبتنی بر کابل
مالتی پلکسی تقسیم فرکانس (FDM): کانال های مختلف در باندهای فرکانسی مختلف منتقل می شوند

Access networks: cable-based access



- HFC: hybrid fiber coax
  - asymmetric: up to 40 Mbps – 1.2 Gbps downstream transmission rate, 30-100 Mbps upstream transmission rate
- network of cable, fiber attaches homes to ISP router
  - homes *share access network* to cable headend

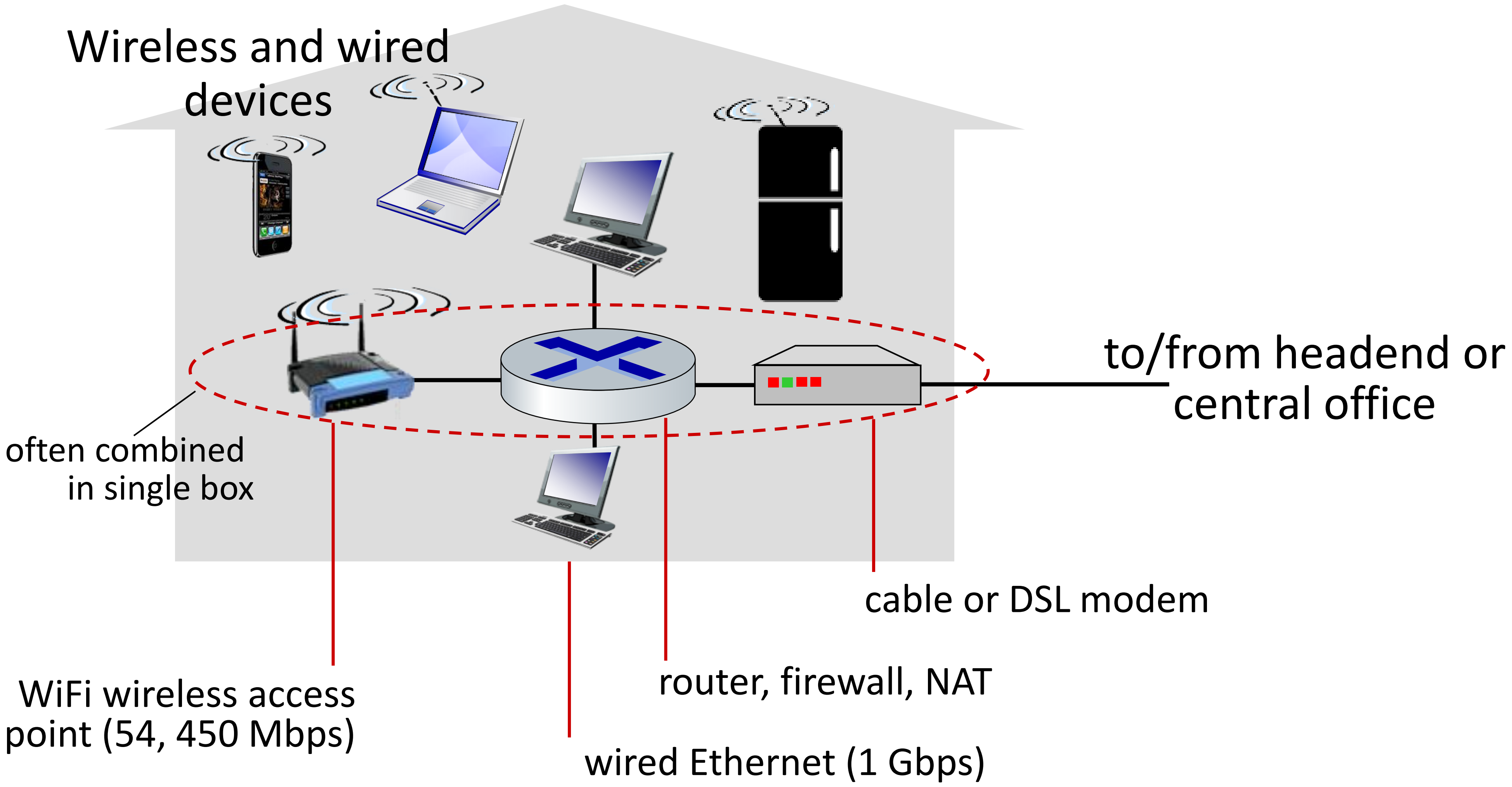
Access networks: digital subscriber line (DSL)



- use *existing* telephone line to central office DSLAM
  - data over DSL phone line goes to Internet
  - voice over DSL phone line goes to telephone net
- 24-52 Mbps dedicated downstream transmission rate
- 3.5-16 Mbps dedicated upstream transmission rate

-
توی خونه ها عمدتا از ADSL استفاده میشه و دلیلش این است که شبکه تلفنی قبل از شبکه دیتا گسترش پیدا کرده بود به همه جا رسوخ پیدا کرده بود
بعد که اینترنت مطرح شد به جای اینکه لینک جدید کشیده بشه و از اون استفاده بشه گفتن از لینک هایی که وجود داره استفاده بشه ینی از شبکه های تلفنی به عنوان واسطه استفاده بشه واسه اینترنت
در ADSL مودم DSL در خونه گذاشته میشه (بیشتر شکل رو ببین)
DSLAM میاد تلفن و دیتا رو ازهم جدا میکنه و تلفن داده میشه به سوییچ های تلفنی ولی دیتا داده میشه به نود ISP
شبکه ISP وصله به شبکه اینترنت
روی خط ارتباطی بین خونه و مرکز سوییچ تلفن ما زوج سیم تلفن داریم که ویس و دیتا با هم منتقل میشن و مالتی پلکس میشه تو فرکانس های رنج های مختلف

Access networks: home networks



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در داخل خود خونه دسترسی خود شبکه access در نقطه انتهایی است

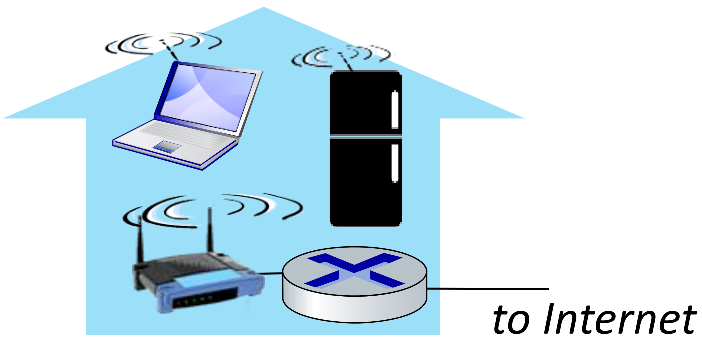
Wireless access networks

Shared *wireless* access network connects end system to router

- via base station aka "access point"

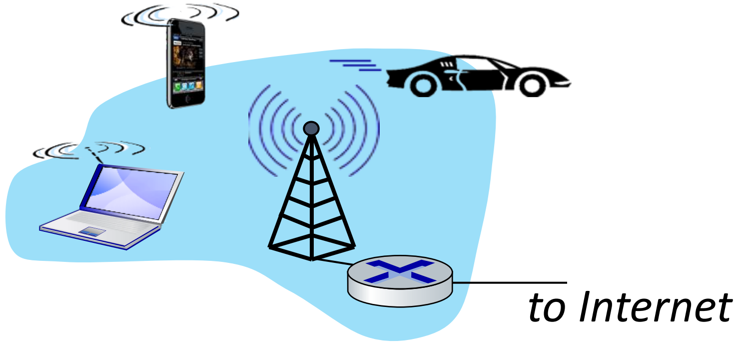
Wireless local area networks (WLANs)

- typically within or around building (~100 ft)
- 802.11b/g/n (WiFi): 11, 54, 450 Mbps transmission rate



Wide-area cellular access networks

- provided by mobile, cellular network operator (10's km)
- 10's Mbps
- 4G cellular networks (5G coming)



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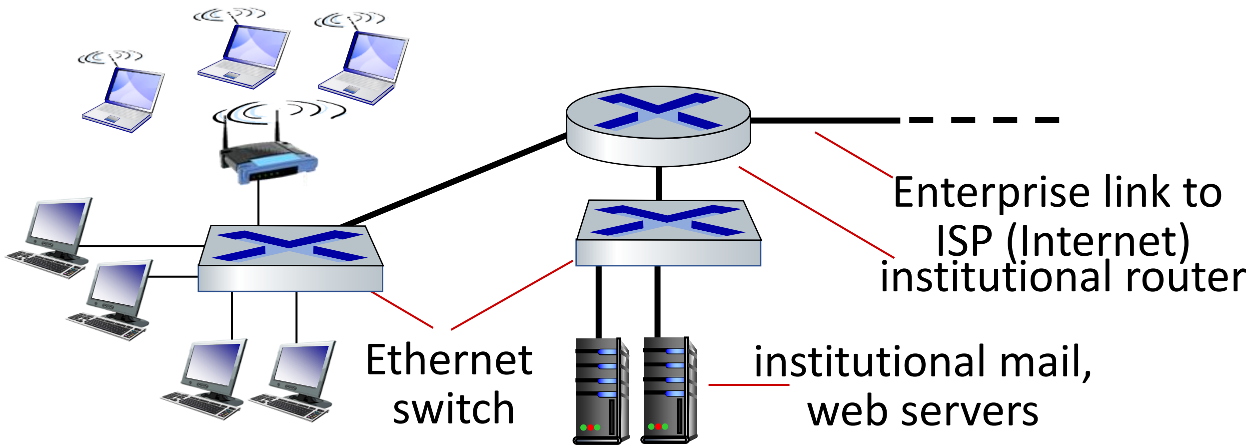
شبکه های وایرلس به صورت محلی دیوایس ها می تونن بهشون وصل بشن در فاصله مثلا حدود 30 متر حداکثر و سرعت ها نسبت به adsl بالاتر است
در سطح وسیع تر استفاده از شبکه های وایرلس wide-area مطرح است که یکی از اینا همون شبکه های cellular است ینی شبکه های موبایل

Access networks: enterprise networks

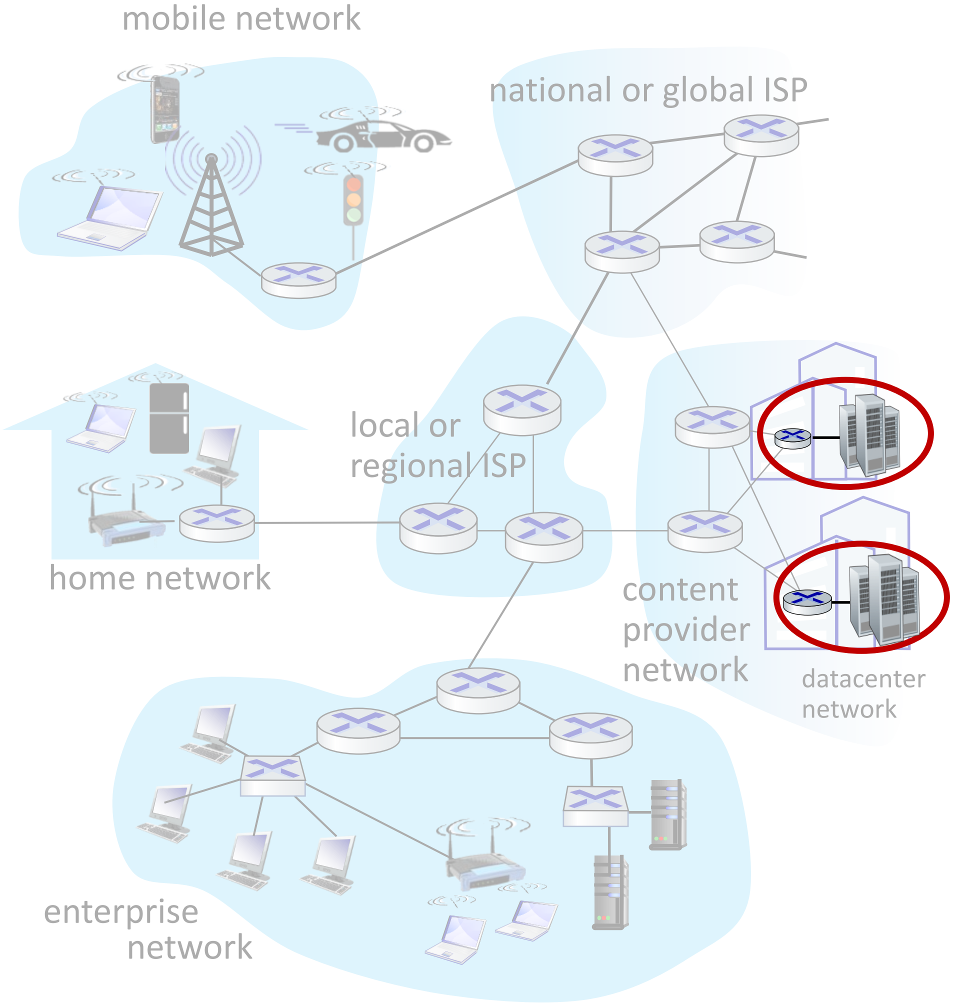


- companies, universities, etc.
- mix of wired, wireless link technologies, connecting a mix of switches and routers (we'll cover differences shortly)
  - Ethernet: wired access at 100Mbps, 1Gbps, 10Gbps
  - WiFi: wireless access points at 11, 54, 450 Mbps

Access networks: data center networks

- high-bandwidth links (10s to 100s Gbps) connect hundreds to thousands of servers together, and to Internet

Courtesy: Massachusetts Green High Performance Computing Center (mghpcc.org)

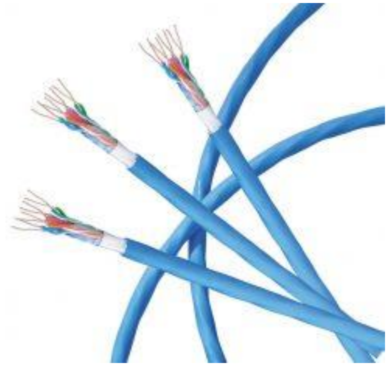


Links: physical media

- bit: propagates between transmitter/receiver pairs
- physical link: what lies between transmitter & receiver
- guided media:
  - signals propagate in solid media: copper, fiber, coax
- unguided media:
  - signals propagate freely, e.g., radio

Twisted pair (TP)

- two insulated copper wires
  - Category 5: 100 Mbps, 1 Gbps Ethernet
  - Category 6: 10Gbps Ethernet



Links: physical media

Coaxial cable:

- two concentric copper conductors
- bidirectional
- broadband:
  - multiple frequency channels on cable
  - 100’s Mbps per channel

Fiber optic cable:

- glass fiber carrying light pulses, each pulse a bit
- high-speed operation:
  - high-speed point-to-point transmission (10’s-100’s Gbps)
- low error rate:
  - repeaters spaced far apart
  - immune to electromagnetic noise

Links: physical media

Wireless radio

- signal carried in various "bands" in electromagnetic spectrum
- no physical "wire"
- broadcast, "half-duplex" (sender to receiver)
- propagation environment effects:
  - reflection
  - obstruction by objects
  - Interference/noise

Radio link types:

- Wireless LAN (WiFi)
  - 10-100's Mbps; 10's of meters
- wide-area (e.g., 4G cellular)
  - 10's Mbps over ~10 Km
- Bluetooth: cable replacement
  - short distances, limited rates
- terrestrial microwave
  - point-to-point; 45 Mbps channels
- satellite
  - up to 45 Mbps per channel
  - 270 msec end-end delay

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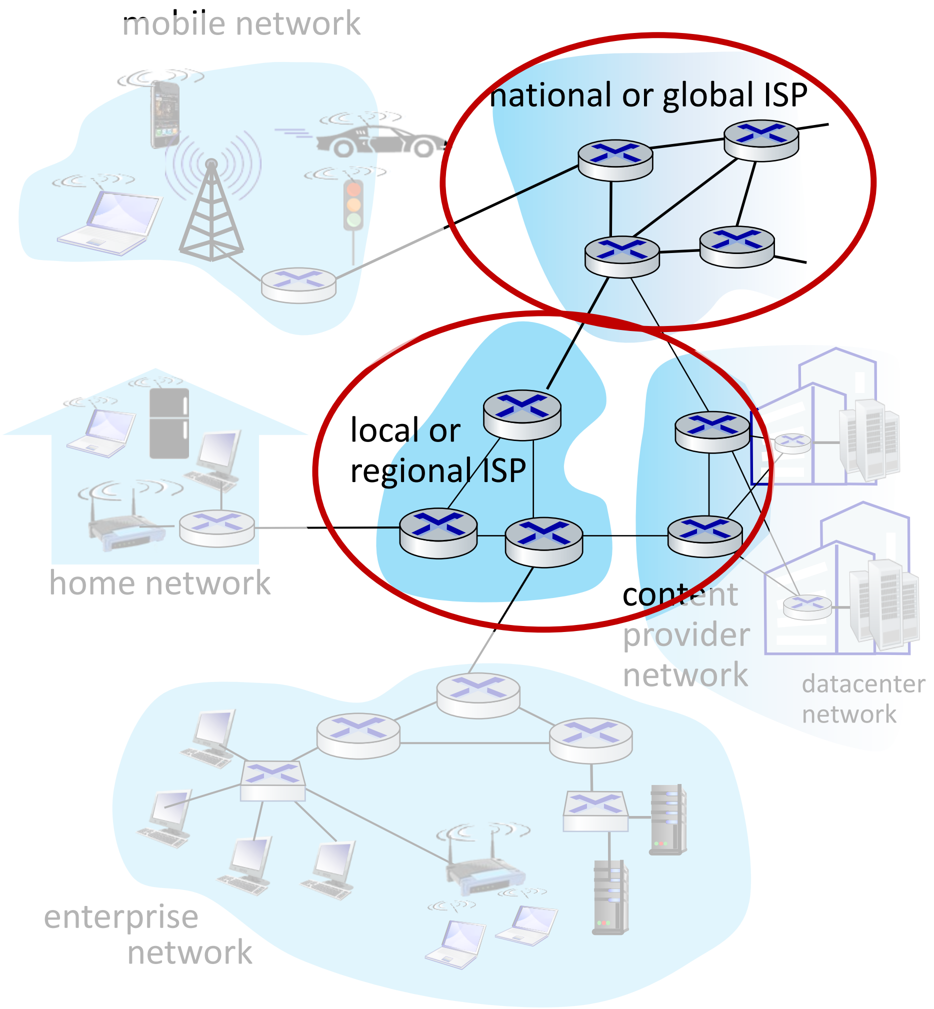
The network core

- mesh of interconnected routers
- packet-switching: hosts break application-layer messages into *packets*
  - network forwards packets from one router to the next, across links on path from source to destination



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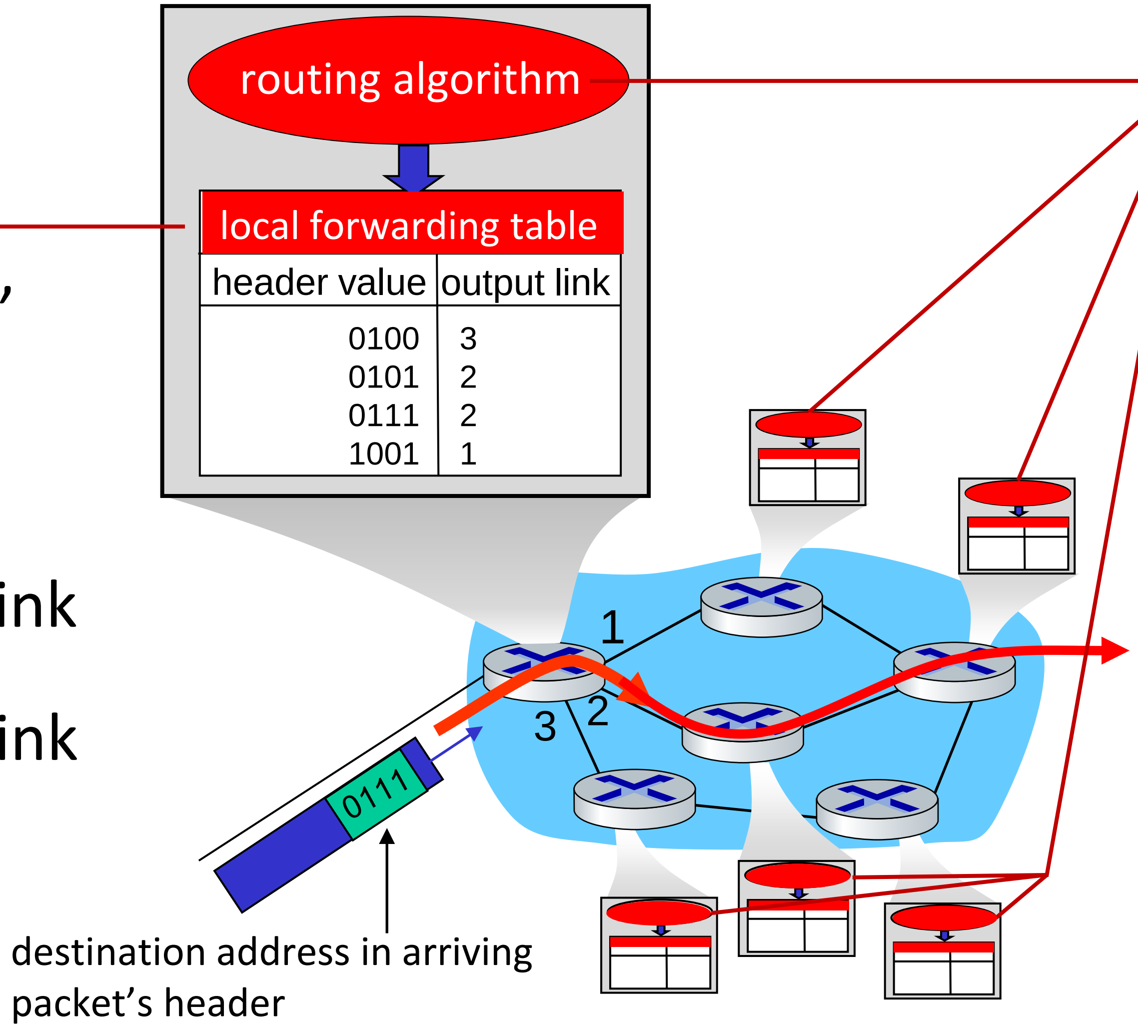
core ارتباط این شبکه های access رو با هم دیگه تامین میکنند و همینطور ارتباطشون با اینترنت هم هست
تو core دیوایس های انتهایی وجود نداره بلکه همین شبکه های access به core وصل میشوند و core وظیفش این است که به بهترین شکل و به صورت بهینه ارتباط بین همه این دیوایس های انتهایی رو و شبکه های access رو با هم فراهم بکنه برای چی؟ برای همین اپلیکیشن هایی که گفتیم توی این دیوایس ها استفاده میشه توسط کاربران
این اپلیکیشن ها پکت ها رو می فرستند و از طریق access منتقل میشه به core و core باید براساس اینکه این بسته به کدوم شبکه access انتهایی یا به کدوم یوزر رسونده بشه اون بسته رو باید هدایت بکنه از طریق اون شبکه مشی که توی core وجود داره

Two key network-core functions

Forwarding:

- aka “switching”
- *local* action: move arriving packets from router’s input link to appropriate router output link



Routing:

- *global* action: determine source-destination paths taken by packets
- routing algorithms

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این بسته ها یا پکت ها به یک نود می رسند و باید روش پردازش انجام بشه و به مسیر درست هدایت بشه و به نود بعدی هدایت بشه که به مقصد برسه
روتر دوتا کار اساسی اینجا انجام میده:
یکیش routing است: ینی تشخیص اینکه برای هر مقصد چه مسیری مناسب است
دومیش forwarding ینی بسته که میاد ادرس مقصد روش زده میشه و میاد این ادرس مقصد رو نگاه میکنه و بعد براساس اطلاعات routing خودش که قبلا مشخص کرده برای هر مقصد می بینه برای این مقصد کدوم مناسبه و مثلا میشه خروجی دو و می فرستتش به خروجی دو
routing algorithm که محاسبه میکنه برای هر مقصدی چه مسیری خوبه

routing
45 h
3,053 miles
48 h
3,309 miles
United States
Victoria
Seattle
WASHINGTON
MONTANA
NORTH DAKOTA
MINNESOTA
Quebec City
Montreal
Ottawa
MAINE
Portland
OREGON
IDAHO
WYOMING
SOUTH DAKOTA
Minneapolis
WISCONSIN
MICHIGAN
Toronto
VERMONT
NEW HAMPSHIRE
Detroit
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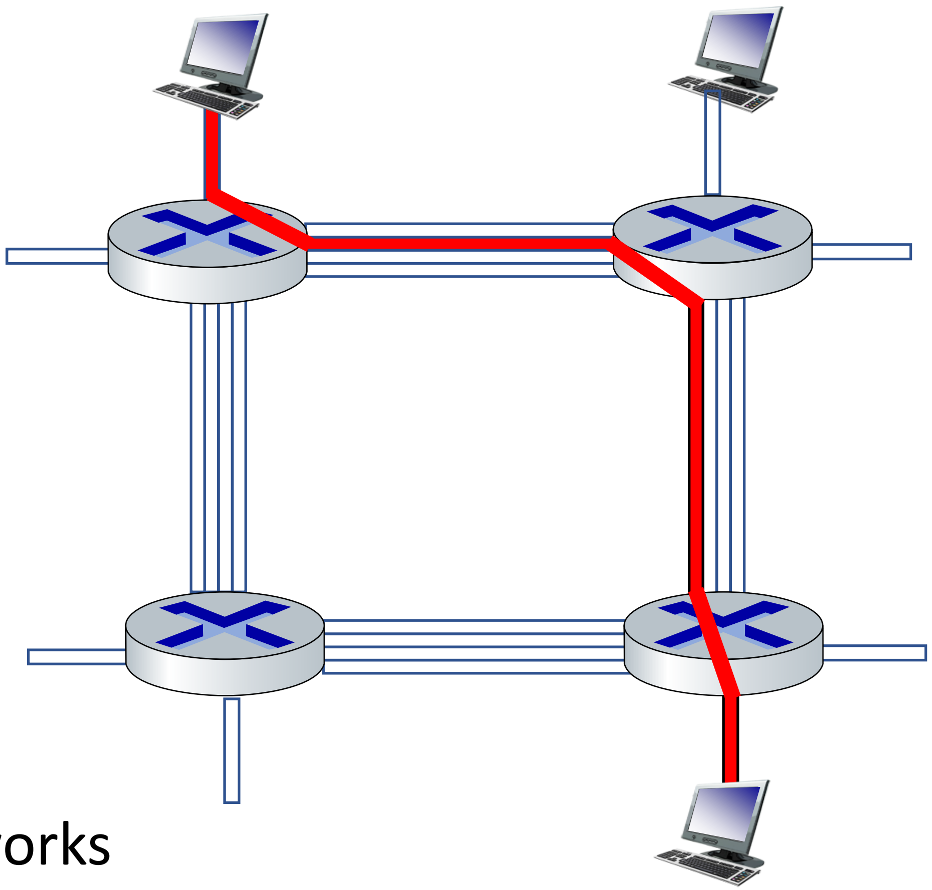
ILLINOIS
INDIANA
Indianapolis
OHIO
PENNSYLVANIA
New York
Philadelphia
NJ
MARYLAND
DE
WEST VIRGINIA
VIRGINIA
KENTUCKY
San Jose
CALIFORNIA
Las Vegas
Nashville
TENNESSEE
OKLAHOMA
ARKANSAS
NORTH CAROLINA
Charlotte
SOUTH CAROLINA
Los Angeles
ARIZONA
NEW MEXICO
MISSISSIPPI
ALABAMA
Atlanta
GEORGIA
San Diego
Phoenix
Dallas
Tucson
El Paso
TEXAS
BAJA CALIFORNIA
SONORA
CHIHUAHUA
Austin
Houston
LOUISIANA
New Orleans
Jacksonville
San Antonio
Orlando
Tampa
FLORIDA
COAHUILA


forwarding
forwarding
45 h
3,053 miles
48 h
3,309 miles
San Jose
Northampton
Sacramento
Cleveland
United States

Alternative to packet switching: circuit switching

end-end resources allocated to, reserved for "call" between source and destination

- in diagram, each link has four circuits.
  - call gets 2nd circuit in top link and 1st circuit in right link.
- dedicated resources: no sharing
  - circuit-like (guaranteed) performance
- circuit segment idle if not used by call (no sharing)
- commonly used in traditional telephone networks

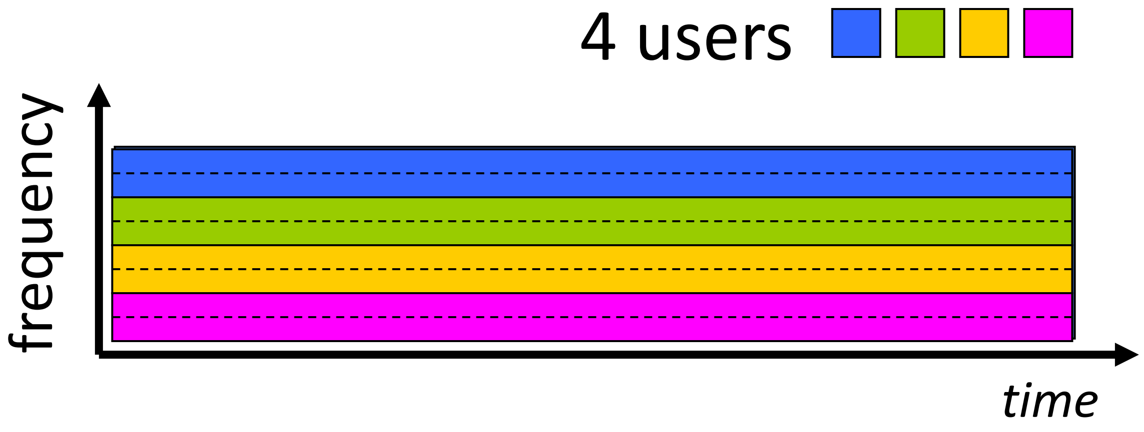


* Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/interactive

Circuit switching: FDM and TDM

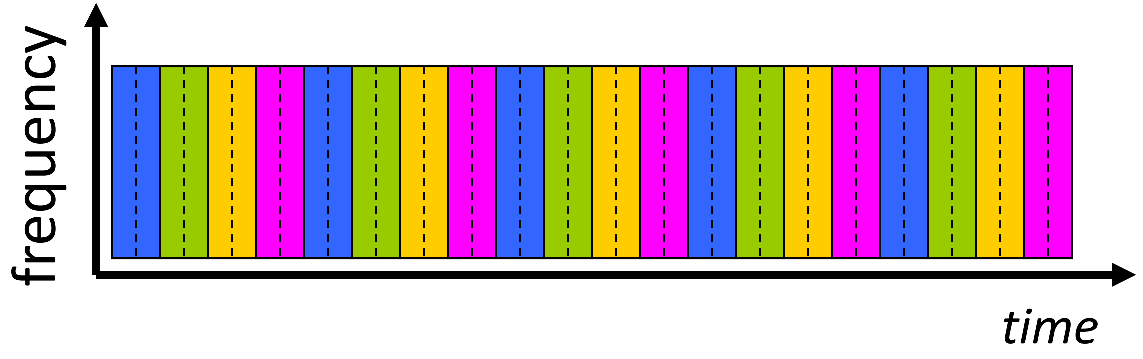
Frequency Division Multiplexing (FDM)

- optical, electromagnetic frequencies divided into (narrow) frequency bands
- each call allocated its own band, can transmit at max rate of that narrow band



Time Division Multiplexing (TDM)

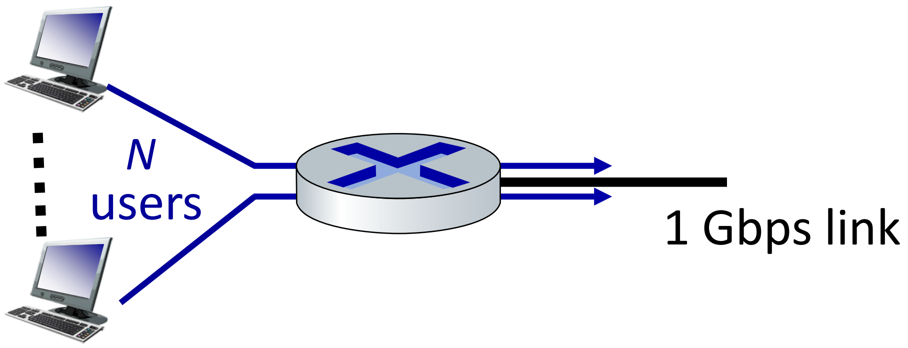
- time divided into slots
- each call allocated periodic slot(s), can transmit at maximum rate of (wider) frequency band (only) during its time slot(s)



Packet switching versus circuit switching

example:

- 1 Gb/s link
- each user:
  - 100 Mb/s when "active"
  - active 10% of time



Q: how many users can use this network under circuit-switching and packet switching?

- *circuit-switching:* 10 users
- *packet switching:* with 35 users, probability > 10 active at same time is less than .0004 *

Q: how did we get value 0.0004?

A: HW problem (for those with course in probability only)

* Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/interactive

Packet switching versus circuit switching

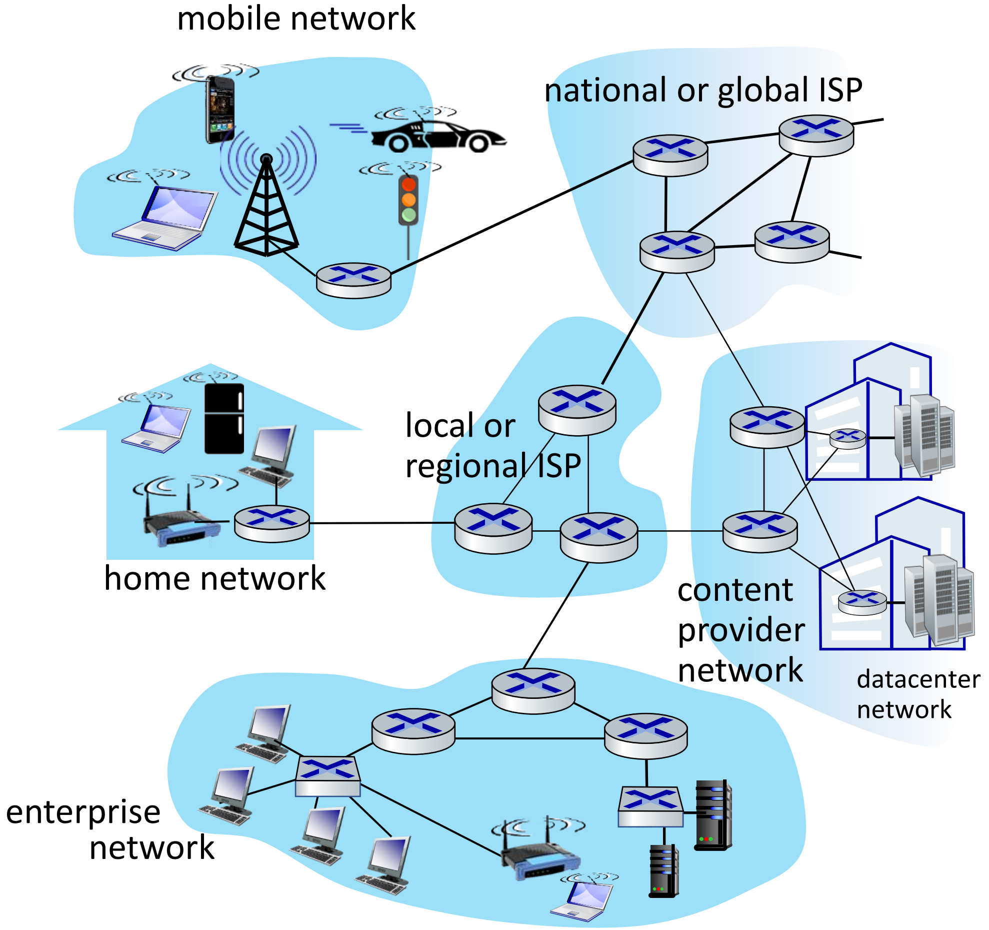
Is packet switching a “slam dunk winner”?

- great for “bursty” data – sometimes has data to send, but at other times not
  - resource sharing
  - simpler, no call setup
- excessive congestion possible: packet delay and loss due to buffer overflow
  - protocols needed for reliable data transfer, congestion control
- *Q:* How to provide circuit-like behavior with packet-switching?
  - “It’s complicated.” We’ll study various techniques that try to make packet switching as “circuit-like” as possible.

Q: human analogies of reserved resources (circuit switching) versus on-demand allocation (packet switching)?

Internet structure: a "network of networks"

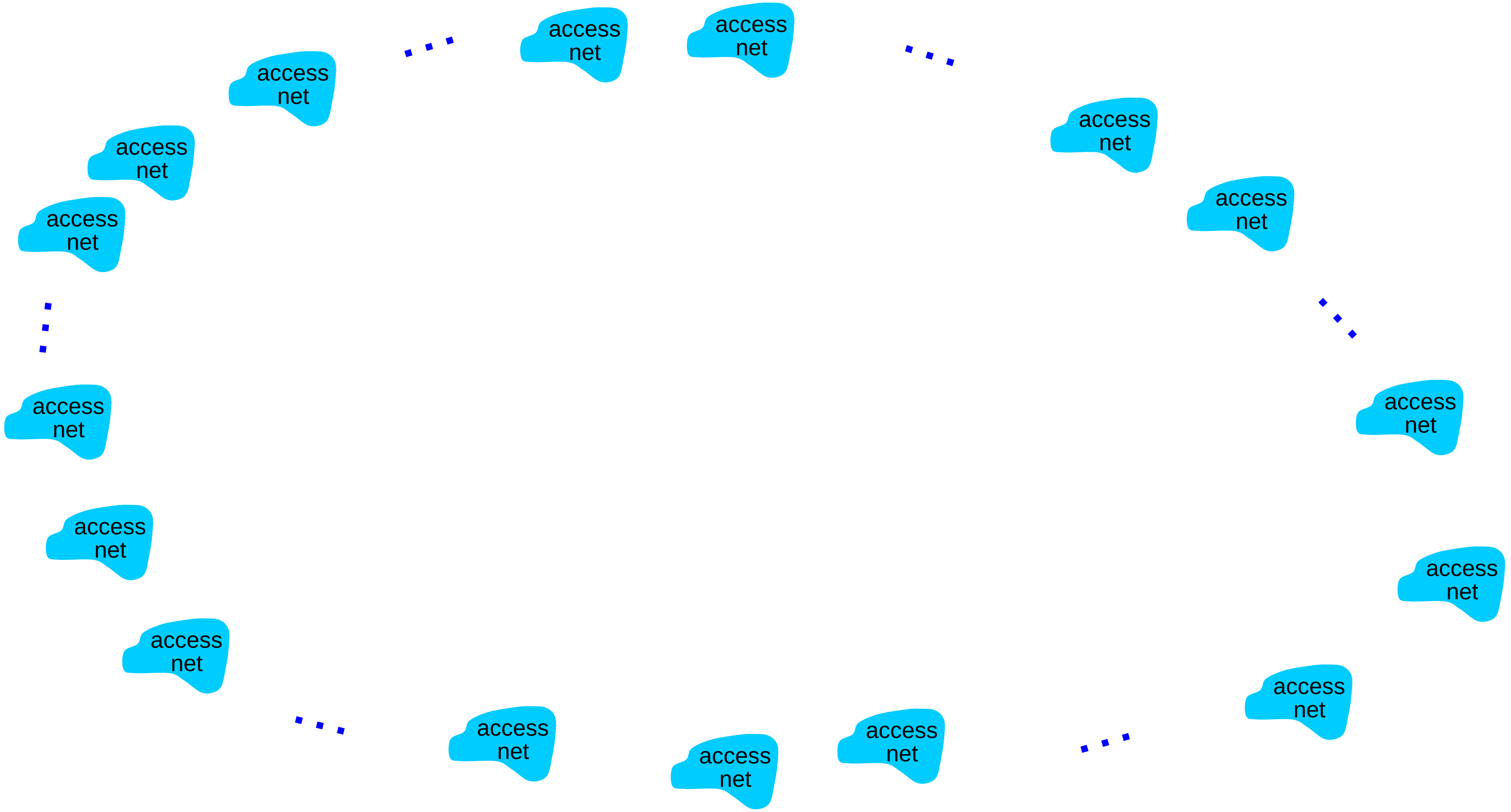
- hosts connect to Internet via access Internet Service Providers (ISPs)
- access ISPs in turn must be interconnected
  - so that *any* two hosts *(anywhere!)* can send packets to each other
- resulting network of networks is very complex
  - evolution driven by economics, national policies



Let's take a stepwise approach to describe current Internet structure

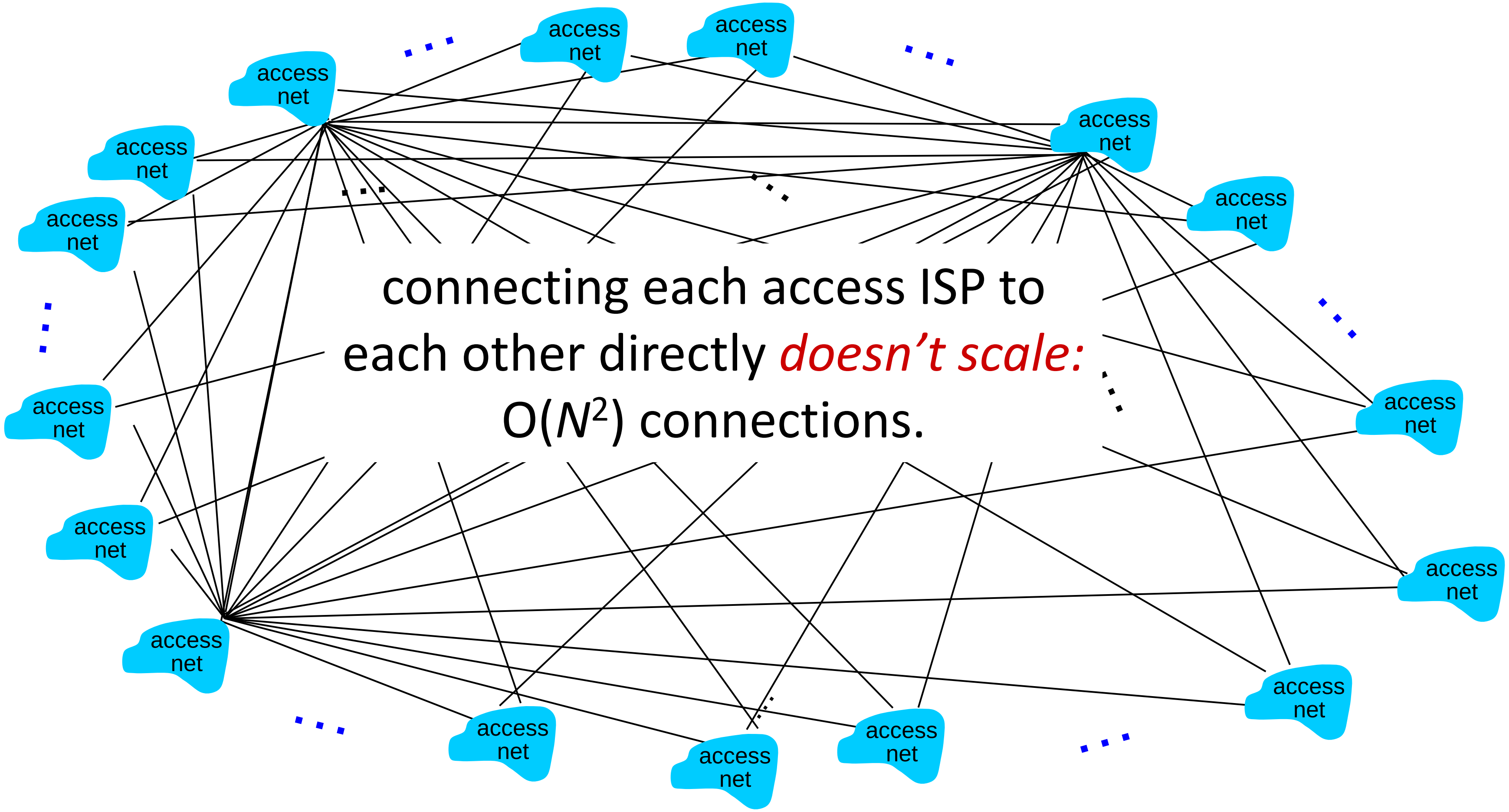
Internet structure: a "network of networks"

Question: given *millions* of access ISPs, how to connect them together?



Internet structure: a "network of networks"

Question: given *millions* of access ISPs, how to connect them together?

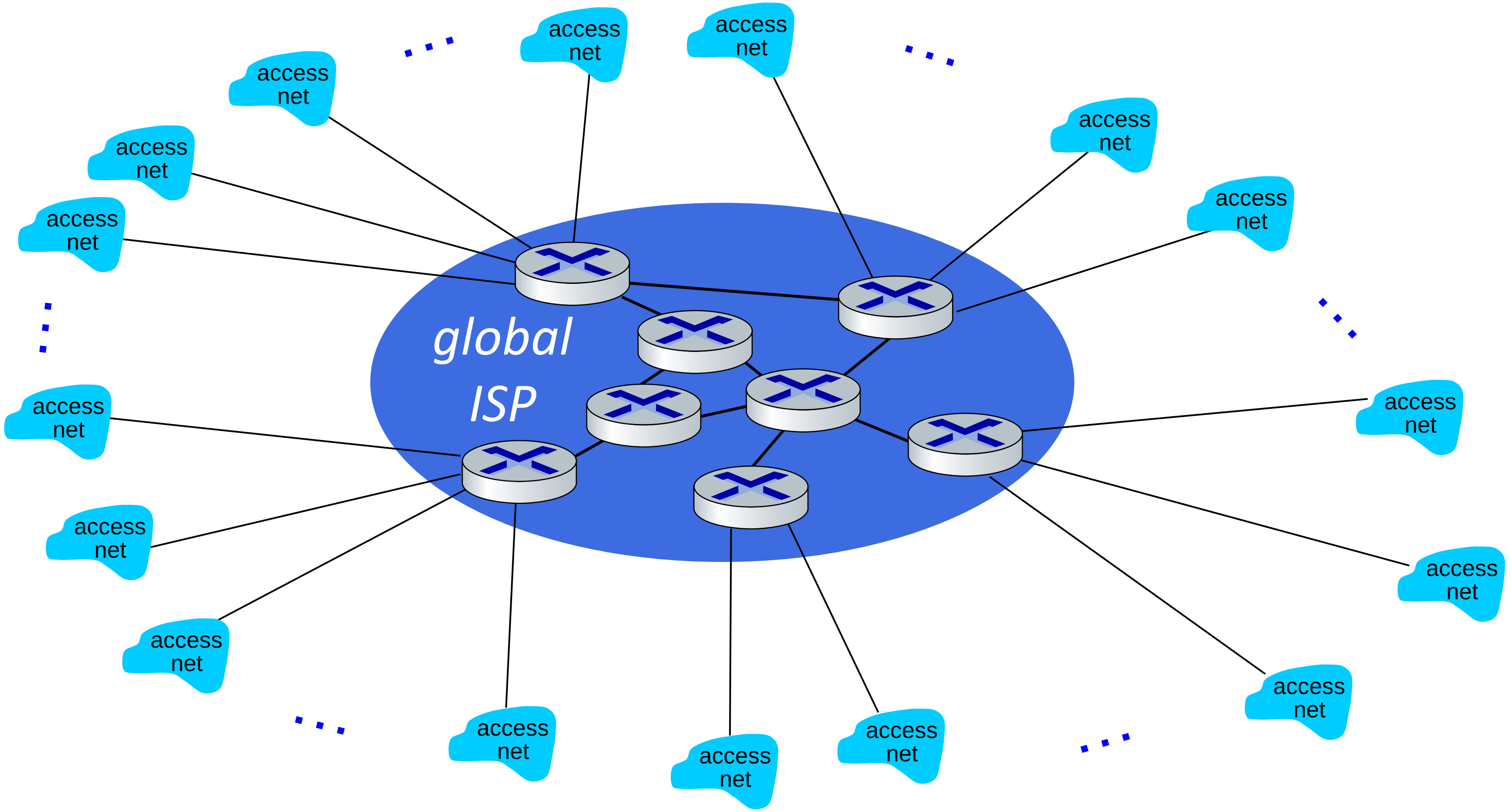


connecting each access ISP to each other directly *doesn't scale:* O(N^2) connections.

Internet structure: a “network of networks”

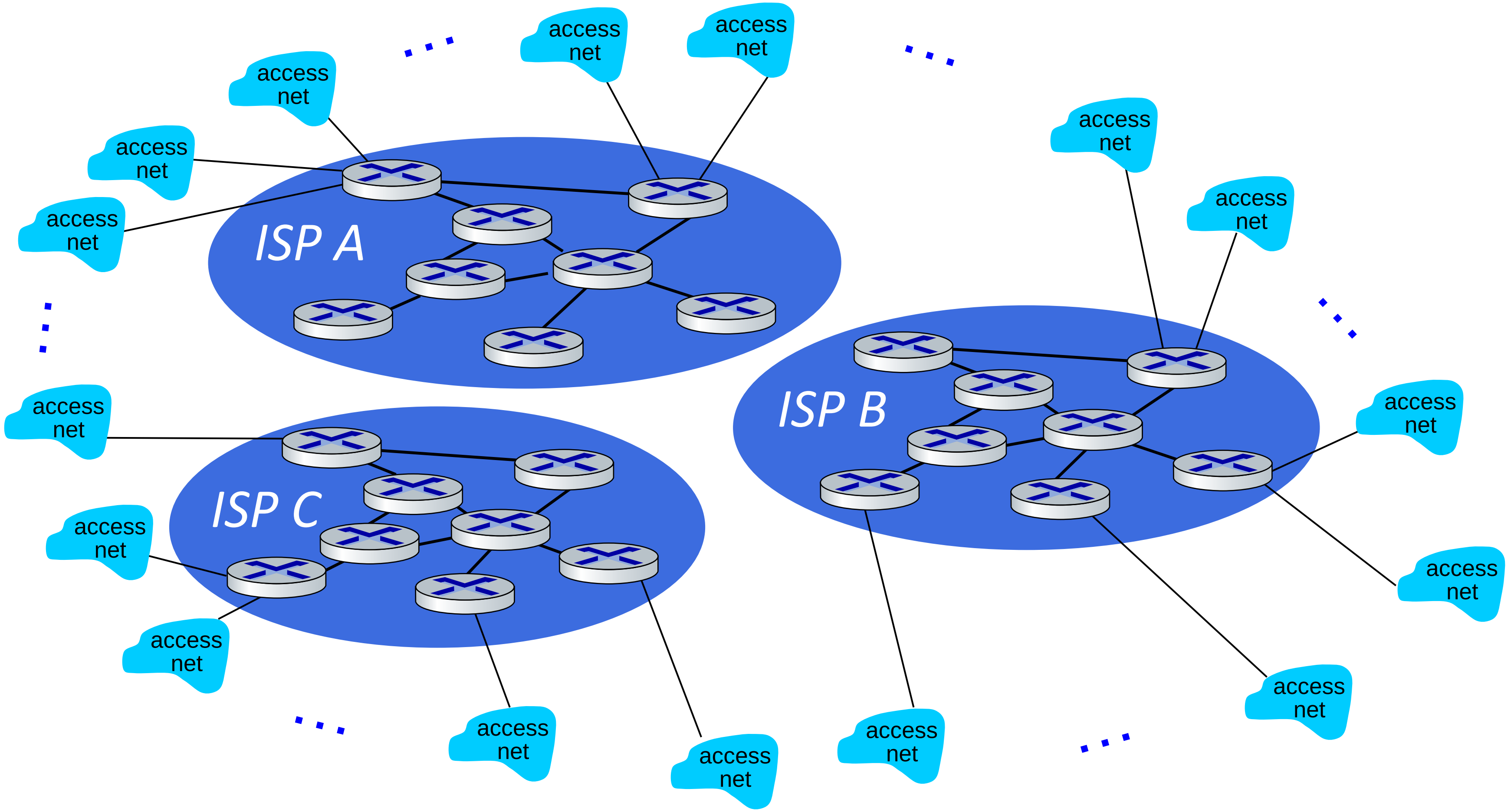
Option: connect each access ISP to one global transit ISP?

Customer and provider ISPs have economic agreement.



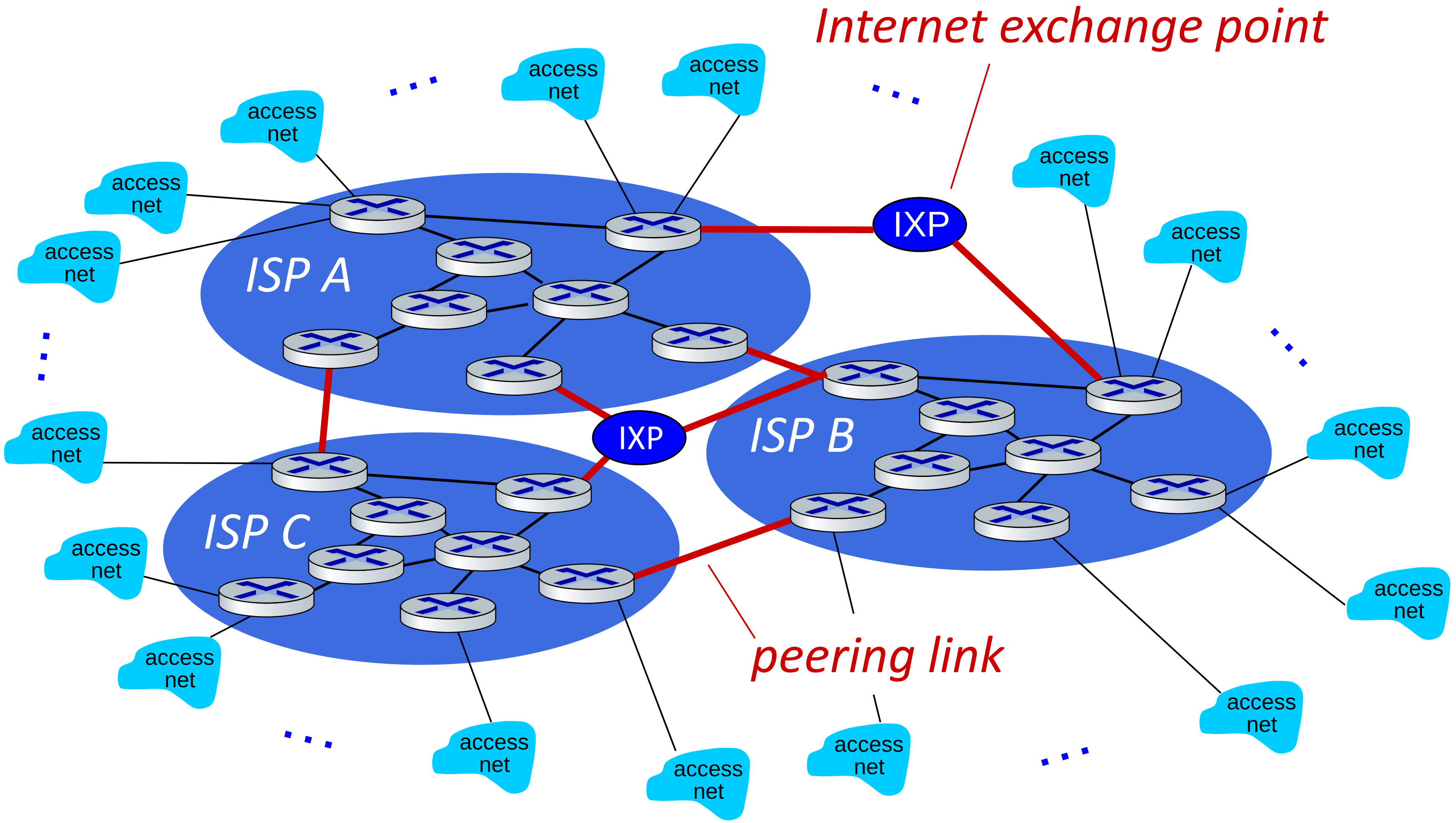
Internet structure: a "network of networks"

But if one global ISP is viable business, there will be competitors ….



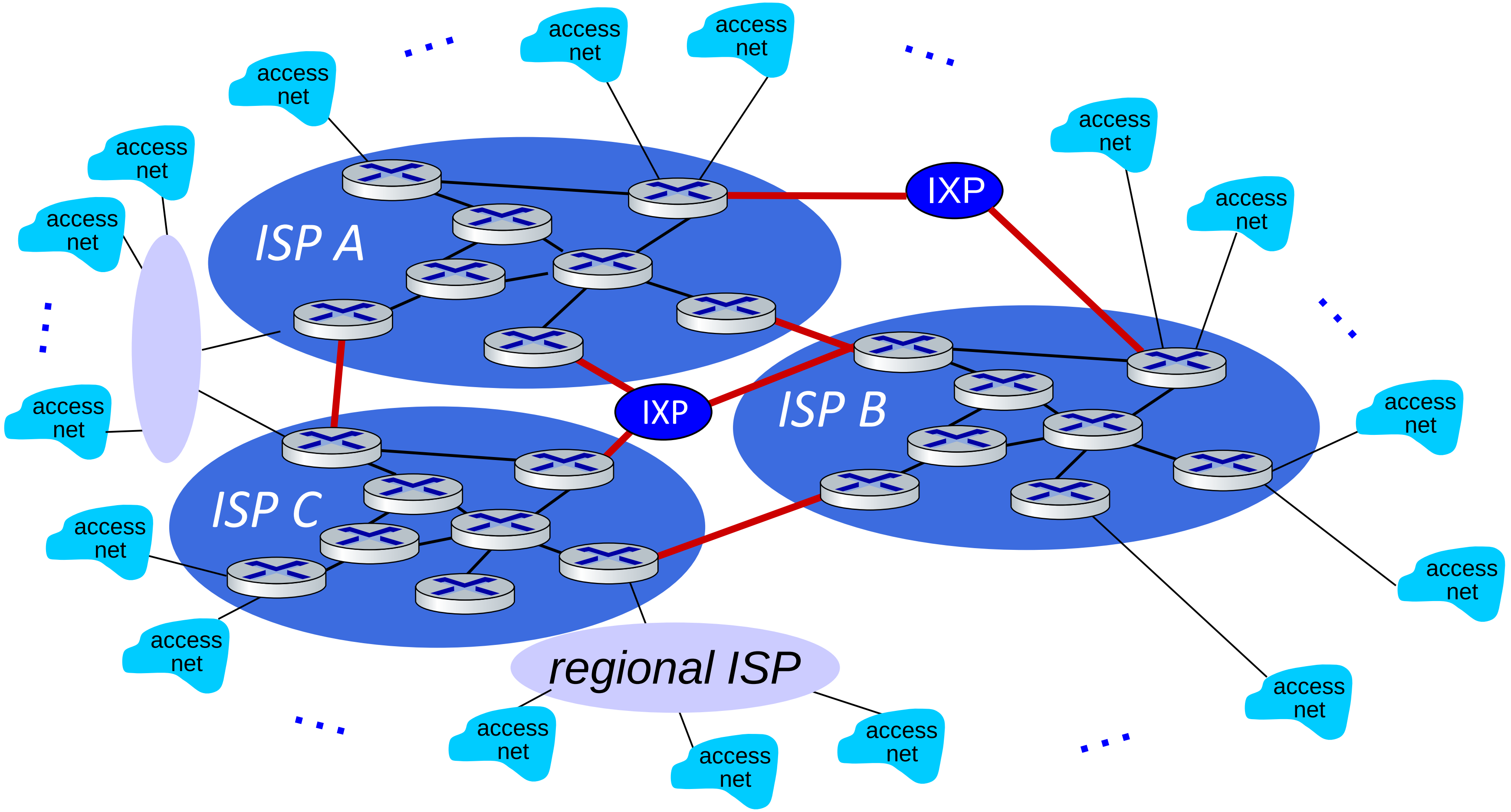
Internet structure: a “network of networks”

But if one global ISP is viable business, there will be competitors …. who will want to be connected



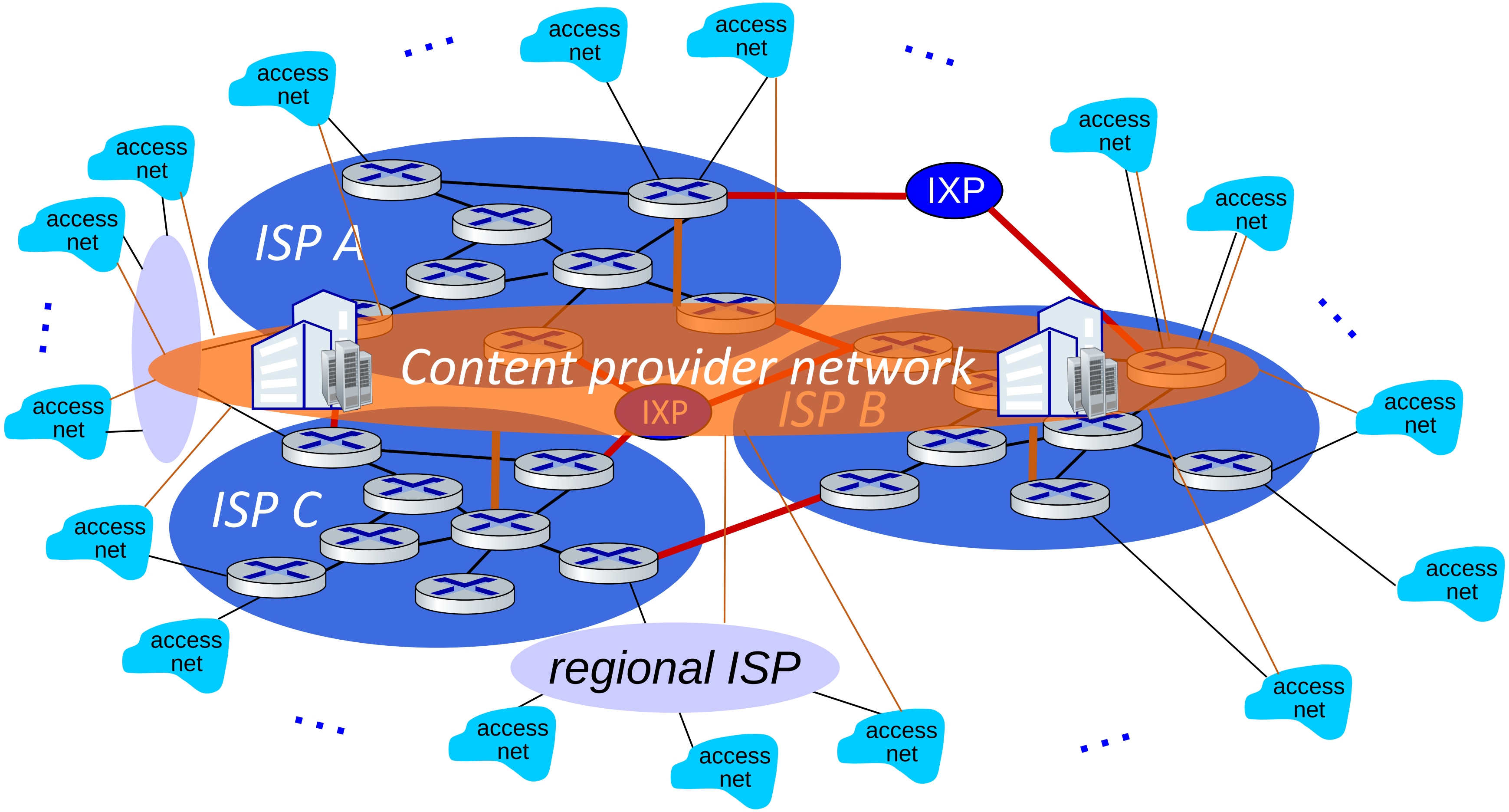
Internet structure: a "network of networks"

… and regional networks may arise to connect access nets to ISPs

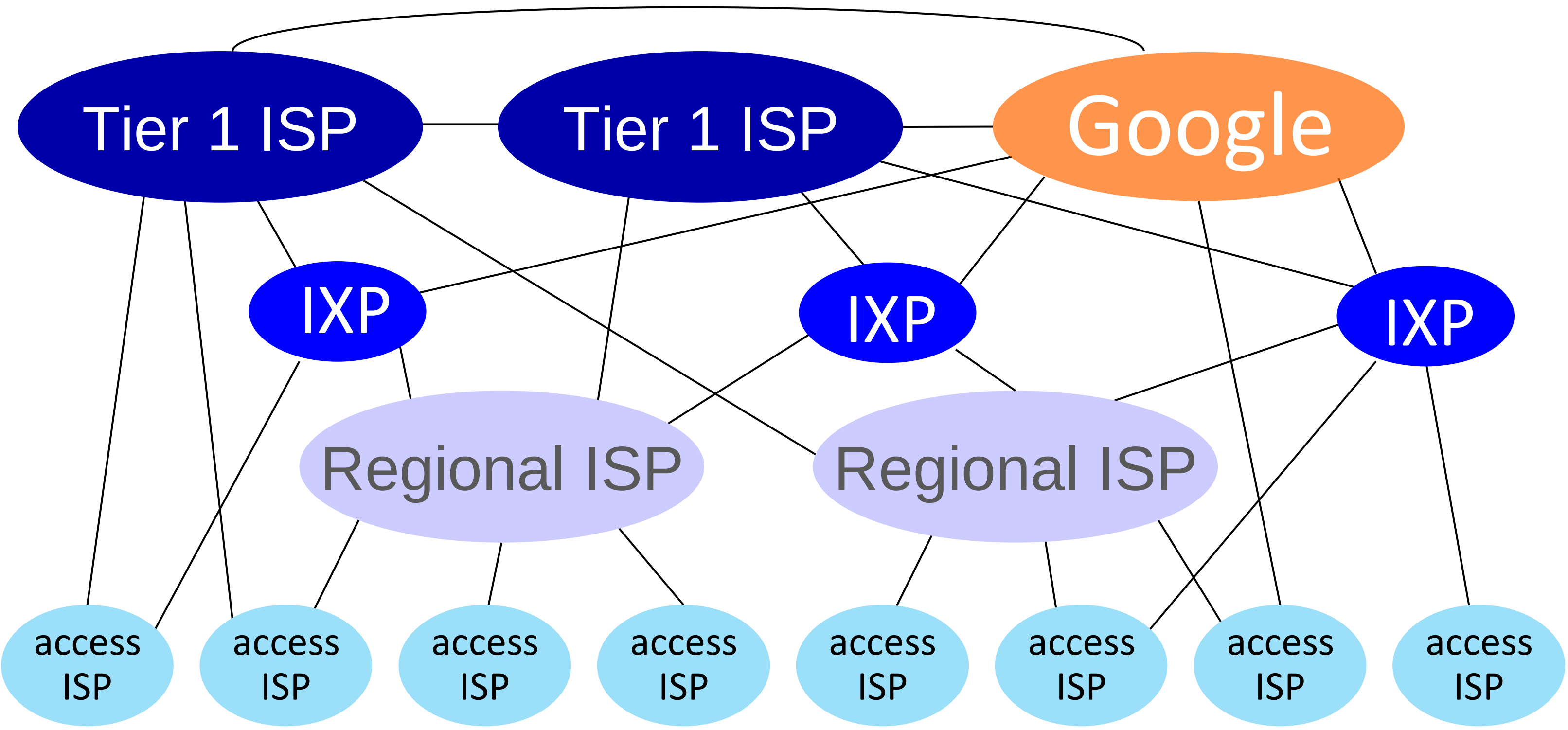


Internet structure: a "network of networks"

... and content provider networks (e.g., Google, Microsoft, Akamai) may run their own network, to bring services, content close to end users



Internet structure: a "network of networks"



At "center": small # of well-connected large networks

- "tier-1" commercial ISPs (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
- content provider networks (e.g., Google, Facebook): private network that connects its data centers to Internet, often bypassing tier-1, regional ISPs

Chapter 1: roadmap

- What *is* the Internet?
- What *is* a protocol?
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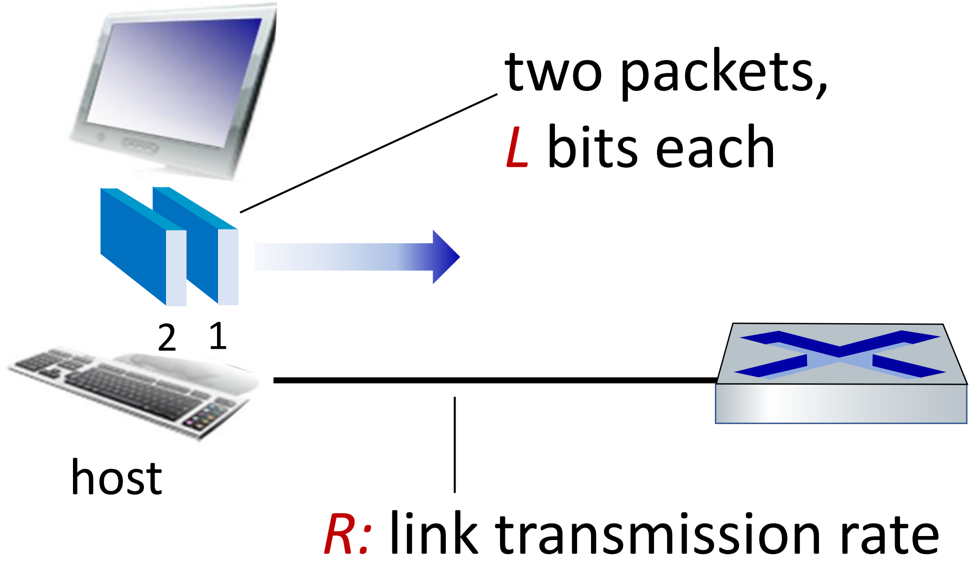
- Network core: packet/circuit switching, internet structure
- Performance: delay, loss, throughput
- Security
- Protocol layers, service models
- History

Host: sends *packets* of data

host sending function:

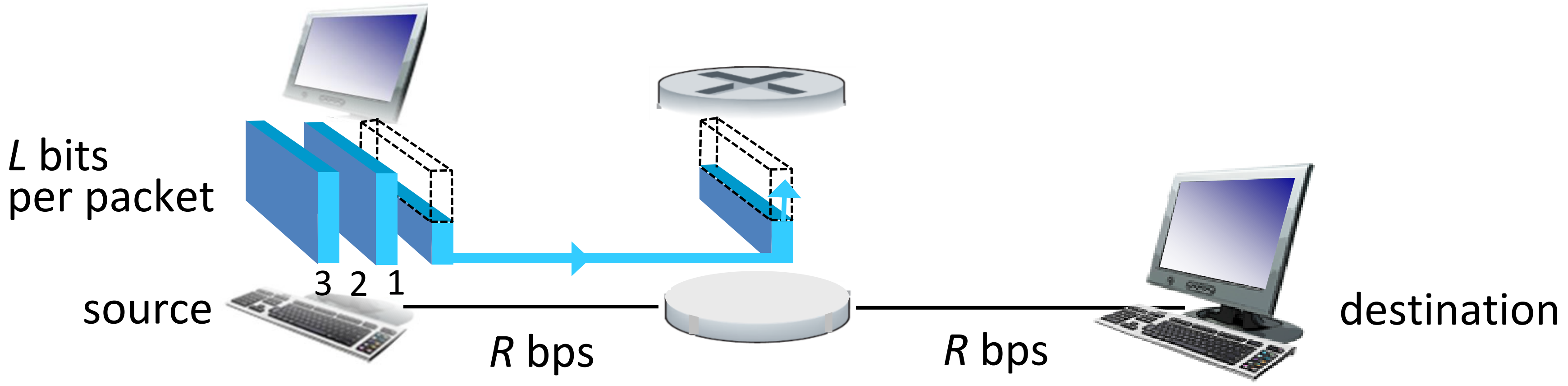
- takes application message
- breaks into smaller chunks, known as *packets*, of length *L* bits
- transmits packet into access network at *transmission rate R*
  - link transmission rate, aka *link capacity, aka link bandwidth*



$$\text{packet transmission delay} = \text{time needed to transmit } L\text{-bit packet into link} = \frac{L \text{ (bits)}}{R \text{ (bits/sec)}}$$

-
تاخیر انتقال بسته: L/R

Packet-switching: store-and-forward

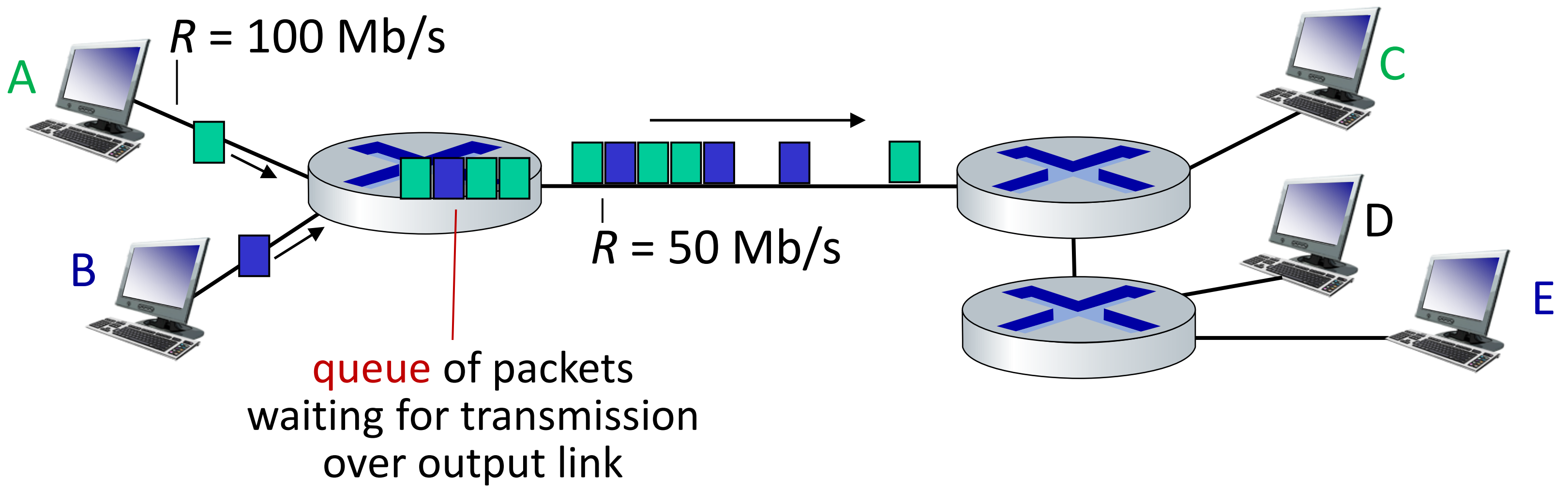


- packet transmission delay: takes L/R seconds to transmit (push out) L-bit packet into link at R bps
- *store and forward: entire* packet must arrive at router before it can be transmitted on next link

One-hop numerical example:

- L = 10 Kbits
- R = 100 Mbps
- one-hop transmission delay = 0.1 msec

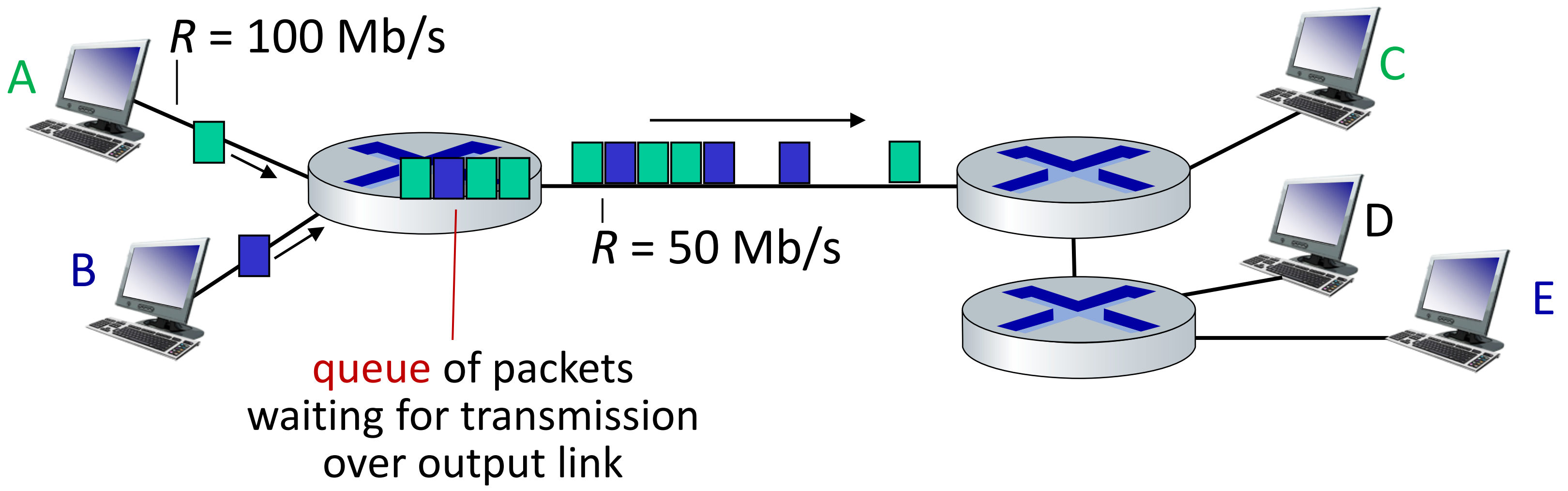
Packet-switching: queueing



Queueing occurs when work arrives faster than it can be serviced:

-
بحث صف:
صف چرا اتفاق می افته؟ چون در نود میانی در core شبکه پکت های مختلف از مبداهای مختلف وارد می شوند و به مقصد های مختلف می روند ولی در یک مسیری از مسیر ممکنه مسیرشون مشترک باشه
این صف تاخیر ایجاد میکنه و اگر این صف پر بشه پکت loss اتفاق می افته

Packet-switching: queueing

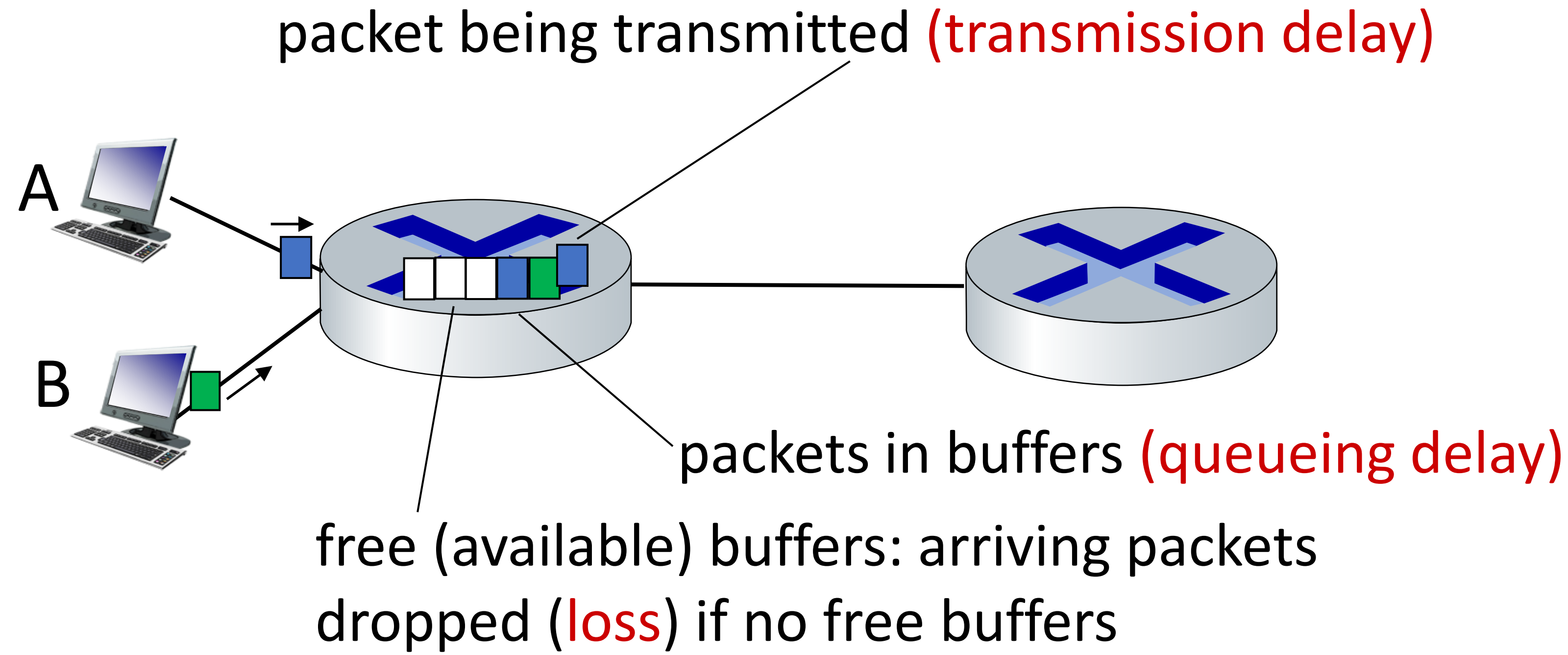


Packet queuing and loss: if arrival rate (in bps) to link exceeds transmission rate (bps) of link for some period of time:

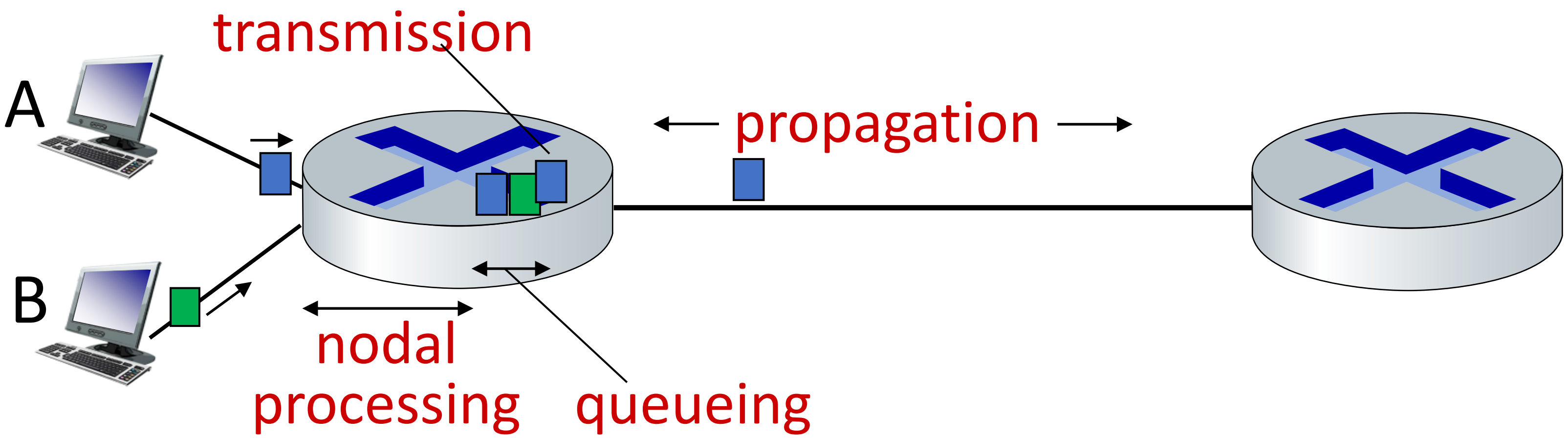
- packets will queue, waiting to be transmitted on output link
- packets can be dropped (lost) if memory (buffer) in router fills up

How do packet delay and loss occur?

- packets *queue* in router buffers, waiting for turn for transmission
  - queue length grows when arrival rate to link (temporarily) exceeds output link capacity
- packet *loss* occurs when memory to hold queued packets fills up



Packet delay: four sources



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

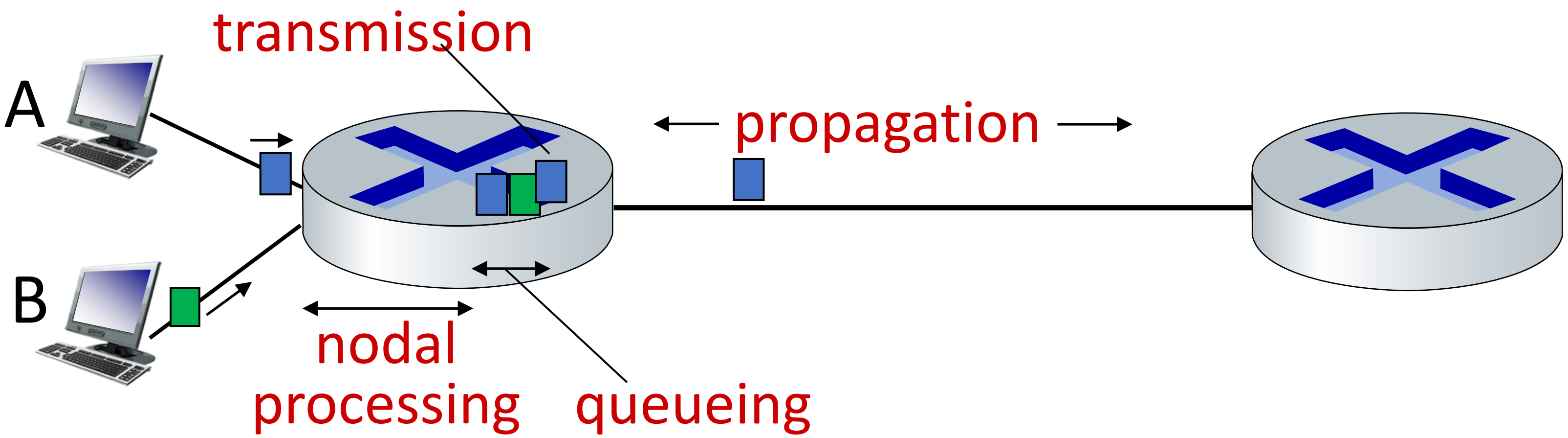
d_{proc}: nodal processing

- check bit errors
- determine output link
- typically < microsecs

d_{queue}: queueing delay

- time waiting at output link for transmission
- depends on congestion level of router

Packet delay: four sources



$$d_{nodal} = d_{proc} + d_{queue} + d_{trans} + d_{prop}$$

d_{trans}: transmission delay:
- *L*: packet length (bits)
- *R*: link *transmission rate (bps)*
- $d_{trans} = L/R$

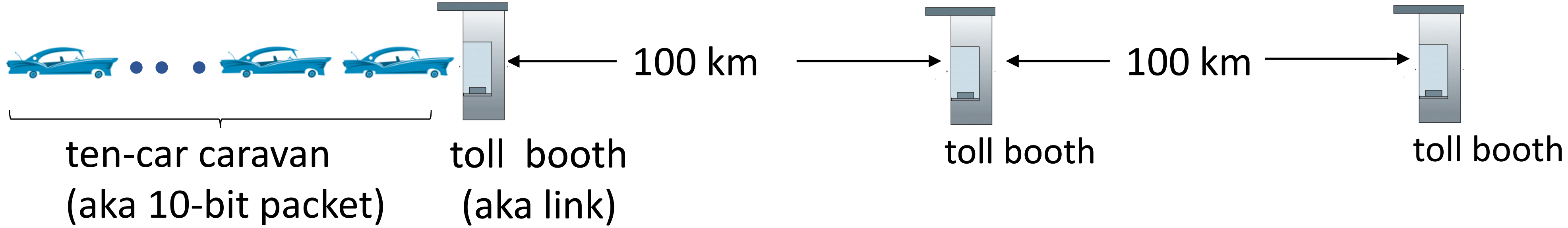
d_{prop}: propagation delay:
- *d*: length of physical link
- *s*: propagation speed (~2x10^8 m/sec)
- $d_{prop} = d/s$

d_{trans} and d_{prop}
very different

-

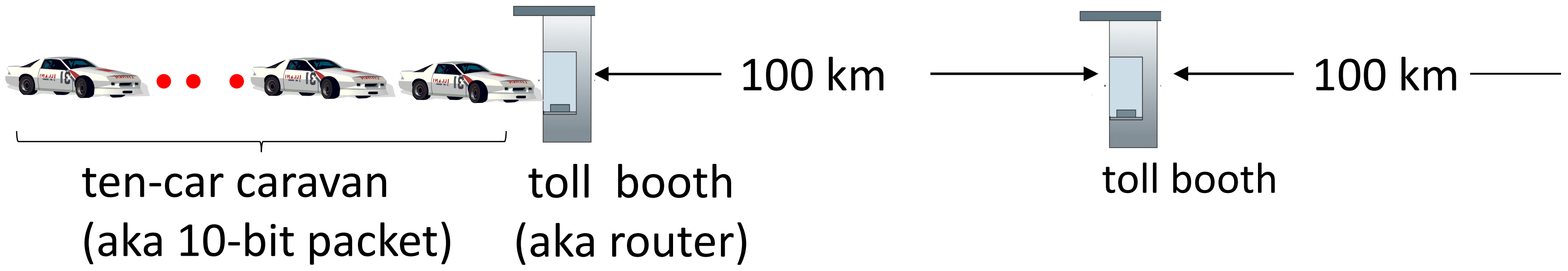
تاخیر ارسال: L/R ثانیه طول میکشه که این ارسال بشه
تاخیر processing ینی توی هر نود بسته باید دریافت بشه و پراسس بشه و بعد ارسال بشه
تاخیر propagation : زمانی که این بسته طول لینک رو طی میکنه این زمان با زمان transmission فرق داره توی transmission این است که چقدر طول می کشه بسته ارسال بشه و این بسته propagation میشه توی لینک ینی چقدر طول میکشه این لینک رو طی بکنه و به مقصد برسه

Caravan analogy



- car ~ bit; caravan ~ packet; toll service ~ link transmission
- toll booth takes 12 sec to service car (bit transmission time)
- "propagate" at 100 km/hr
- *Q:* How long until caravan is lined up before 2nd toll booth?
- time to "push" entire caravan through toll booth onto highway = 12*10 = 120 sec
- time for last car to propagate from 1st to 2nd toll both: 100km/(100km/hr) = 1 hr
- *A:* 62 minutes

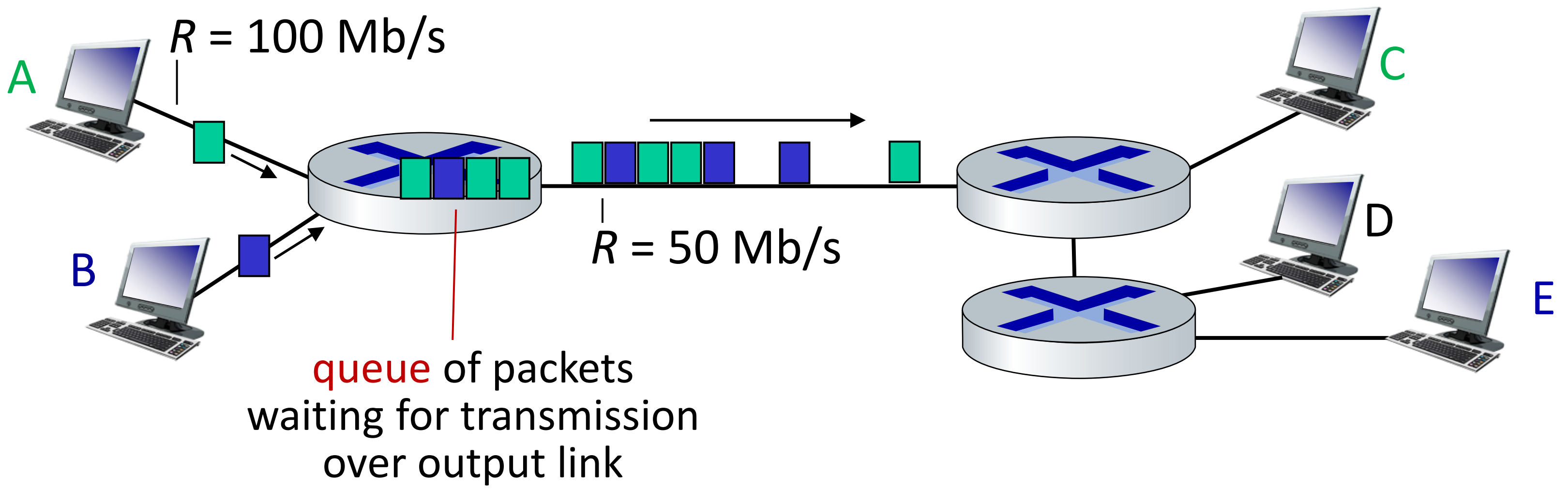
Caravan analogy



- suppose cars now "propagate" at 1000 km/hr
- and suppose toll booth now takes one min to service a car
- *Q: Will cars arrive to 2nd booth before all cars serviced at first booth?*

  A: Yes! after 7 min, first car arrives at second booth; three cars still at first booth

Packet-switching: queueing



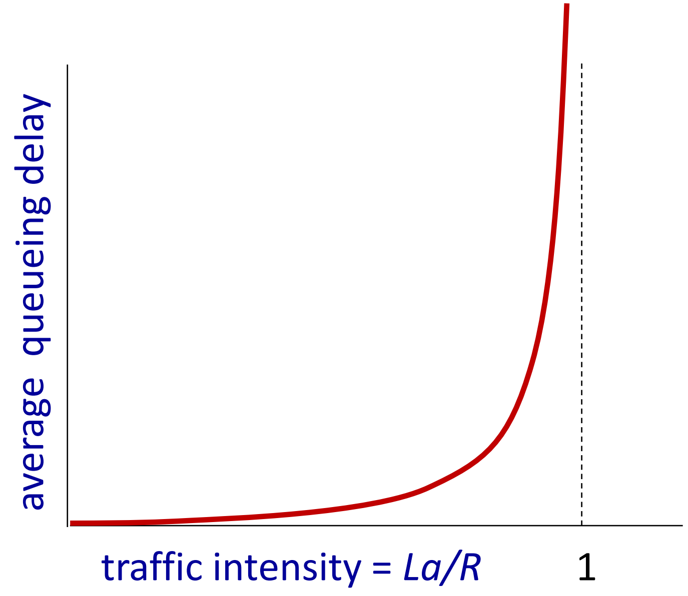
Queueing occurs when work arrives faster than it can be serviced:

Packet queueing delay (revisited)

- *a:* average packet arrival rate
- *L:* packet length (bits)
- *R:* link bandwidth (bit transmission rate)

$$\frac{L \cdot a}{R} : \quad \frac{\text{arrival rate of bits}}{\text{service rate of bits}}$$

"traffic intensity"



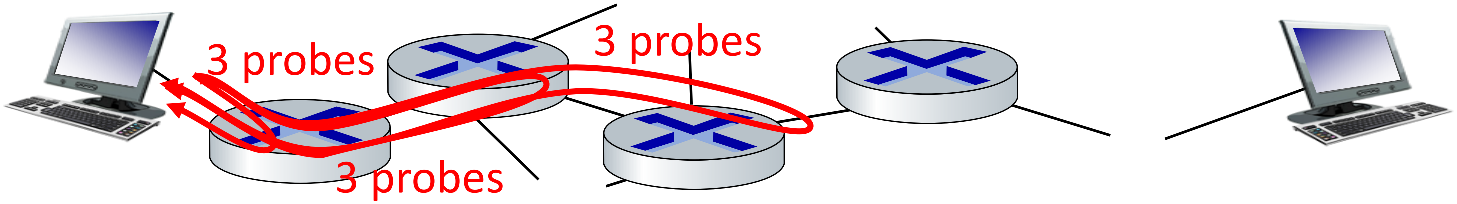
- *La/R* ~ 0: avg. queueing delay small
- *La/R* -> 1: avg. queueing delay large
- *La/R* > 1: more "work" arriving is more than can be serviced - average delay infinite!

La/R ~ 0

La/R -> 1

“Real” Internet delays and routes

- what do “real” Internet delay & loss look like?
- **traceroute** program: provides delay measurement from source to router along end-end Internet path towards destination. For all *i:*
  - sends three packets that will reach router *i* on path towards destination (with time-to-live field value of *i*)
  - router *i* will return packets to sender
  - sender measures time interval between transmission and reply



Real Internet delays and routes

traceroute: gaia.cs.umass.edu to www.eurecom.fr

3 delay measurements from
gaia.cs.umass.edu to cs-gw.cs.umass.edu

```
1  cs-gw (128.119.240.254)  1 ms  1 ms  2 ms
2  border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145)  1 ms  1 ms  2 ms
3  cht-vbns.gw.umass.edu (128.119.3.130)  6 ms 5 ms 5 ms
4  jn1-at1-0-0-19.wor.vbns.net (204.147.132.129)  16 ms 11 ms 13 ms
5  jn1-so7-0-0-0.wae.vbns.net (204.147.136.136)  21 ms 18 ms 18 ms
6  abilene-vbns.abilene.ucaid.edu (198.32.11.9)  22 ms  18 ms  22 ms
7  nycm-wash.abilene.ucaid.edu (198.32.8.46)  22 ms  22 ms  22 ms
8  62.40.103.253 (62.40.103.253)  104 ms 109 ms 106 ms
9  de2-1.de1.de.geant.net (62.40.96.129)  109 ms 102 ms 104 ms
10  de.fr1.fr.geant.net (62.40.96.50)  113 ms 121 ms 114 ms
11  renater-gw.fr1.fr.geant.net (62.40.103.54)  112 ms  114 ms  112 ms
12  nio-n2.cssi.renater.fr (193.51.206.13)  111 ms  114 ms  116 ms
13  nice.cssi.renater.fr (195.220.98.102)  123 ms  125 ms  124 ms
14  r3t2-nice.cssi.renater.fr (195.220.98.110)  126 ms  126 ms  124 ms
15  eurecom-valbonne.r3t2.ft.net (193.48.50.54)  135 ms  128 ms  133 ms
16  194.214.211.25 (194.214.211.25)  126 ms  128 ms  126 ms
17  * * *
18  * * *
19  fantasia.eurecom.fr (193.55.113.142)  132 ms  128 ms  136 ms
```

3 delay measurements
to border1-rt-fa5-1-0.gw.umass.edu

trans-oceanic link

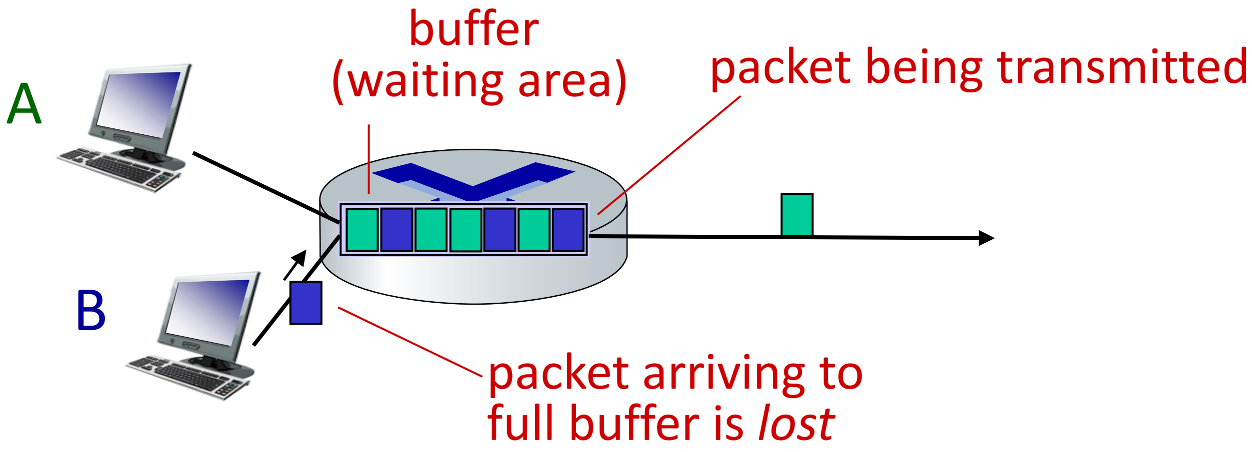
looks like delays *decrease*! Why?

* means no response (probe lost, router not replying)

* Do some traceroutes from exotic countries at www.traceroute.org

Packet loss

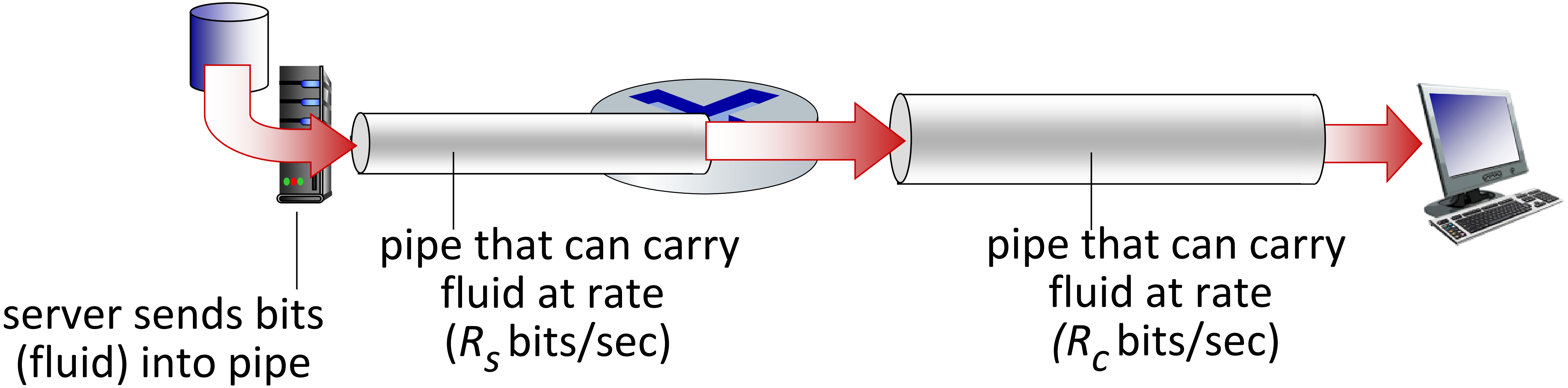
- queue (aka buffer) preceding link in buffer has finite capacity
- packet arriving to full queue dropped (aka lost)
- lost packet may be retransmitted by previous node, by source end system, or not at all



* Check out the Java applet for an interactive animation (on publisher's website) of queuing and loss

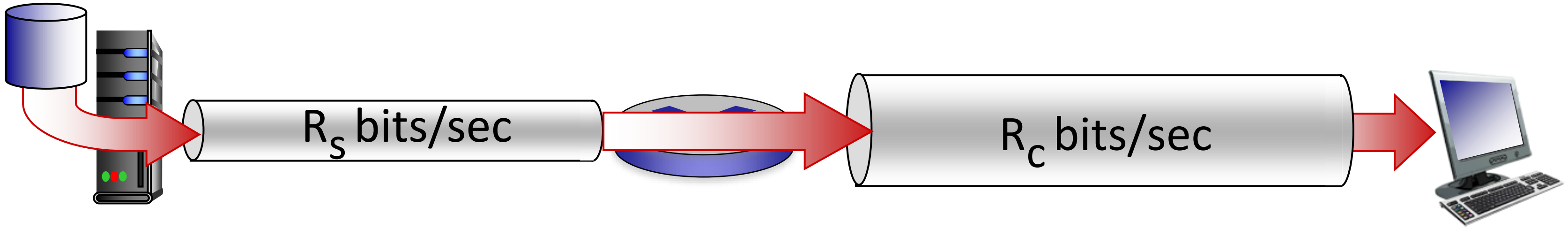
Throughput

- *throughput:* rate (bits/time unit) at which bits are being sent from sender to receiver
  - *instantaneous:* rate at given point in time
  - *average:* rate over longer period of time

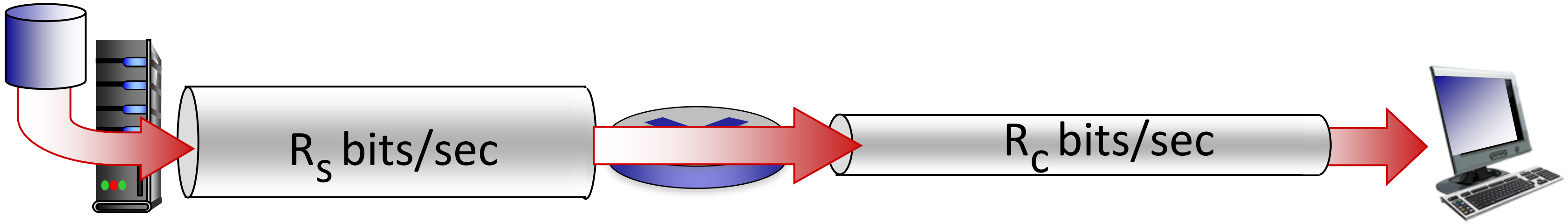


Throughput

$R_s < R_c$ What is average end-end throughput?



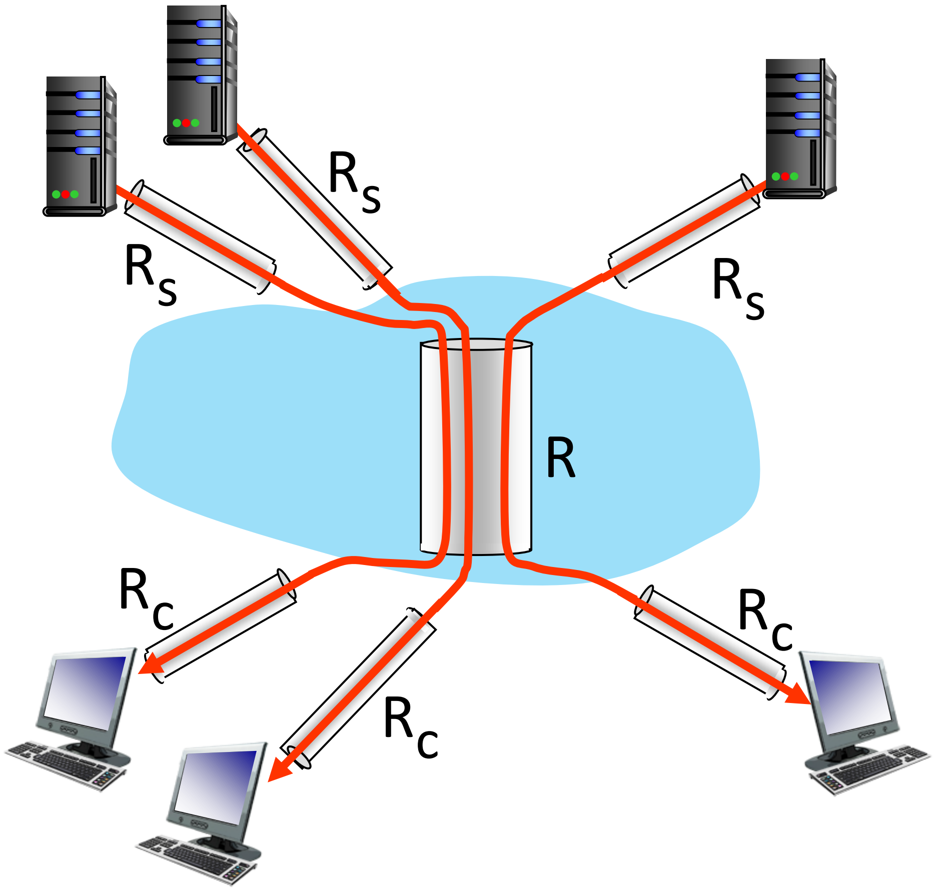
$R_s > R_c$ What is average end-end throughput?



bottleneck link

link on end-end path that constrains end-end throughput

Throughput: network scenario



10 connections (fairly) share backbone bottleneck link *R* bits/sec

- per-connection end-end throughput: $min(R_c,R_s,R/10)$
- in practice: R_c or R_s is often bottleneck

* Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/

Chapter 1: roadmap

- What *is* the Internet?
- What *is* a protocol?
- Network edge: hosts, access network, physical media
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- Protocol layers, service models
- History

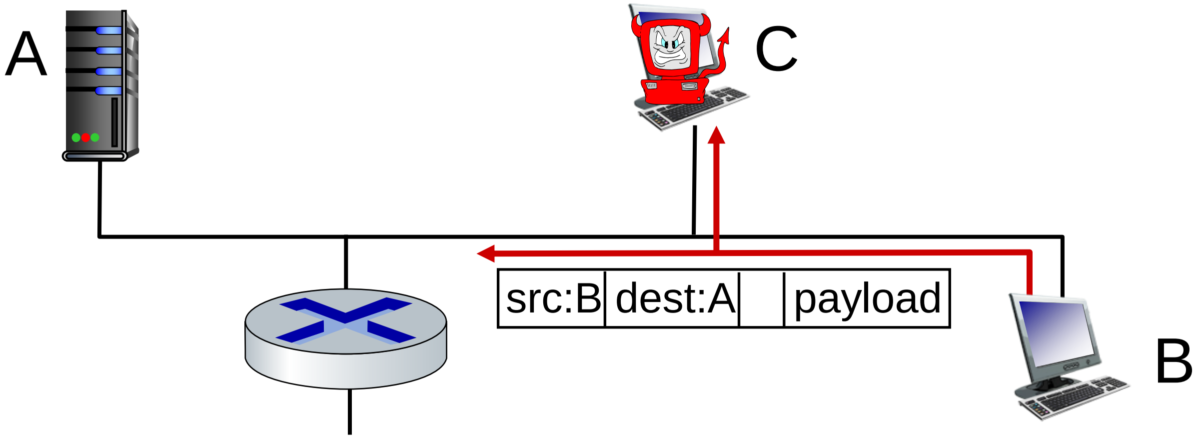
Network security

- Internet not originally designed with (much) security in mind
  - *original vision:* "a group of mutually trusting users attached to a transparent network" ☺
  - Internet protocol designers playing "catch-up"
  - security considerations in all layers!
- We now need to think about:
  - how bad guys can attack computer networks
  - how we can defend networks against attacks
  - how to design architectures that are immune to attacks

Bad guys: packet interception

packet "sniffing":

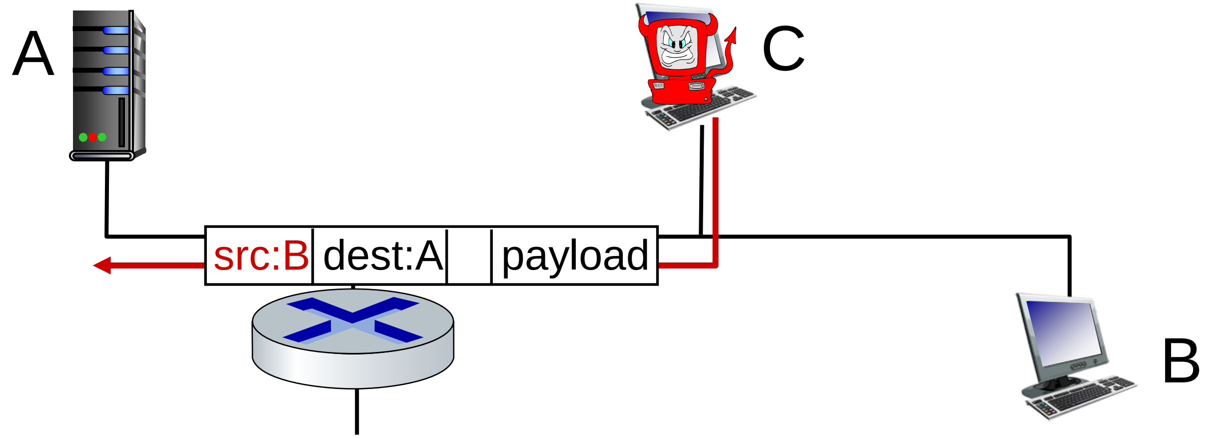
- broadcast media (shared Ethernet, wireless)
- promiscuous network interface reads/records all packets (e.g., including passwords!) passing by



Wireshark software used for our end-of-chapter labs is a (free) packet-sniffer

Bad guys: fake identity

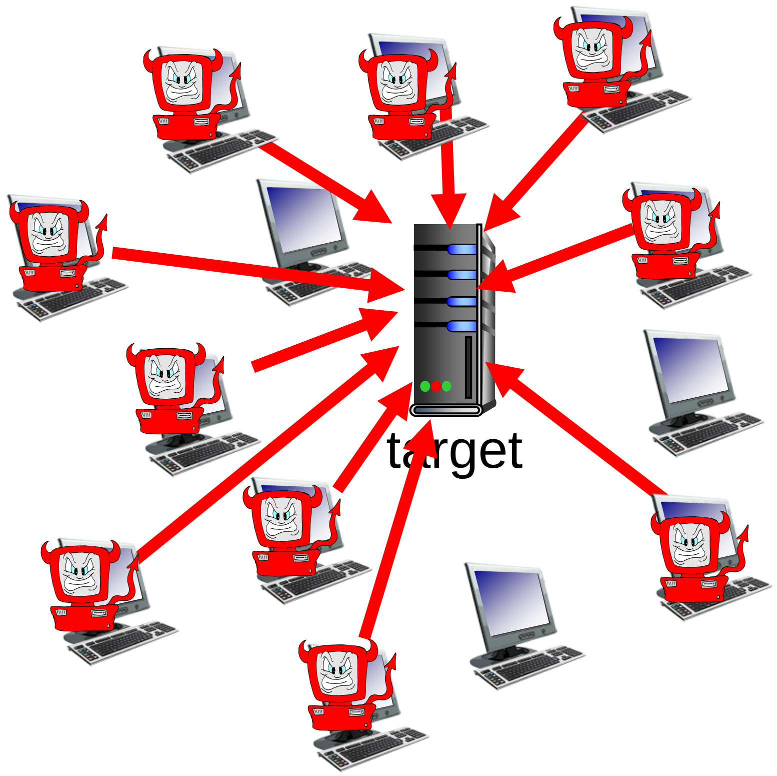
IP spoofing: injection of packet with false source address



Bad guys: denial of service

Denial of Service (DoS): attackers make resources (server, bandwidth) unavailable to legitimate traffic by overwhelming resource with bogus traffic

1. select target
2. break into hosts around the network (see botnet)
3. send packets to target from compromised hosts



Lines of defense:

- authentication: proving you are who you say you are
  - cellular networks provides hardware identity via SIM card; no such hardware assist in traditional Internet
- confidentiality: via encryption
- integrity checks: digital signatures prevent/detect tampering
- access restrictions: password-protected VPNs
- firewalls: specialized “middleboxes” in access and core networks:
  - off-by-default: filter incoming packets to restrict senders, receivers, applications
  - detecting/reacting to DOS attacks

… lots more on security (throughout, Chapter 8)

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Protocol "layers" and reference models

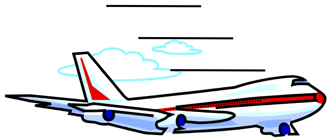
Networks are complex, with many "pieces":

- hosts
- routers
- links of various media
- applications
- protocols
- hardware, software

Question: is there any hope of *organizing* structure of network?

- and/or our *discussion* of networks?

Example: organization of air travel



end-to-end transfer of person plus baggage

ticket (purchase)	ticket (complain)
baggage (check)	baggage (claim)
gates (load)	gates (unload)
runway takeoff	runway landing
airplane routing	airplane routing

airplane routing

How would you *define/discuss* the *system* of airline travel?

- a series of steps, involving many services

Example: organization of air travel

layers: each layer implements a service

- via its own internal-layer actions
- relying on services provided by layer below

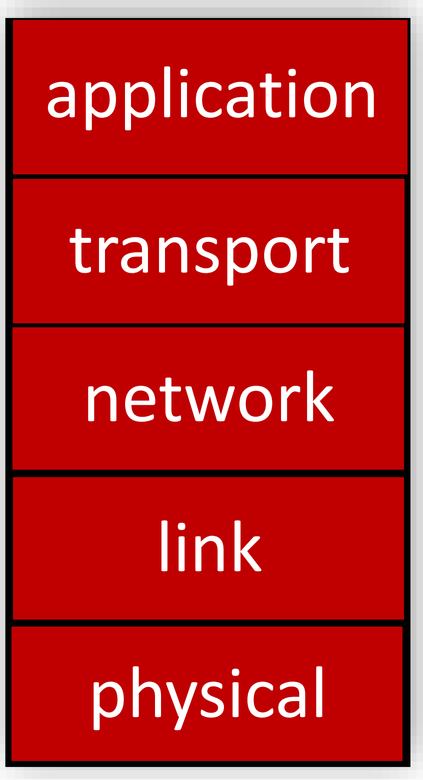
Why layering?

Approach to designing/discussing complex systems:

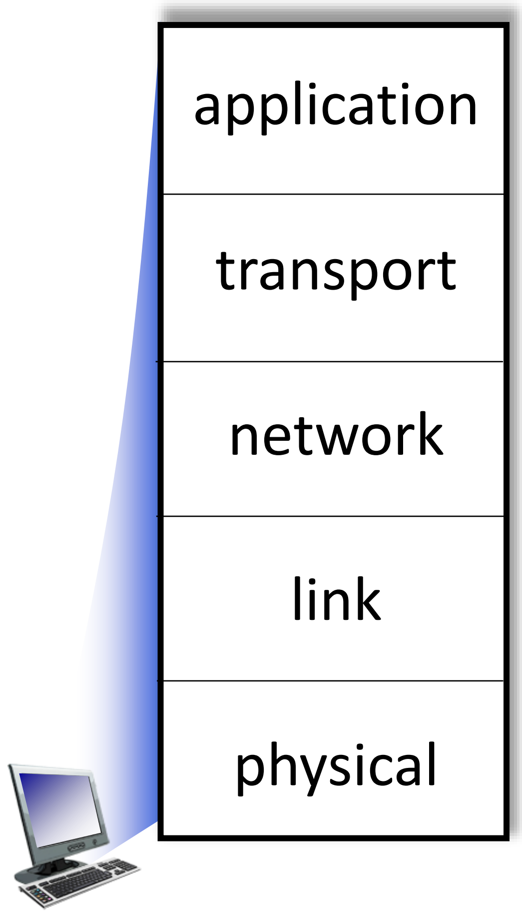
- explicit structure allows identification, relationship of system's pieces
  - layered *reference model* for discussion
- modularization eases maintenance, updating of system
  - change in layer's service *implementation*: transparent to rest of system
  - e.g., change in gate procedure doesn't affect rest of system

Layered Internet protocol stack

- *application:* supporting network applications
  - HTTP, IMAP, SMTP, DNS
- *transport:* process-process data transfer
  - TCP, UDP
- *network:* routing of datagrams from source to destination
  - IP, routing protocols
- *link:* data transfer between neighboring network elements
  - Ethernet, 802.11 (WiFi), PPP
- *physical:* bits “on the wire”



Services, Layering and Encapsulation



source

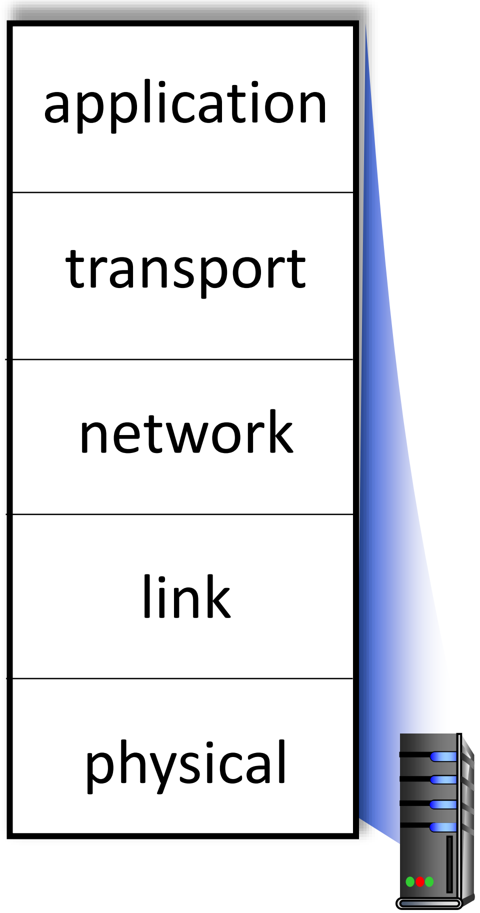
M

Application exchanges messages to implement some application service using *services* of transport layer

H_t M

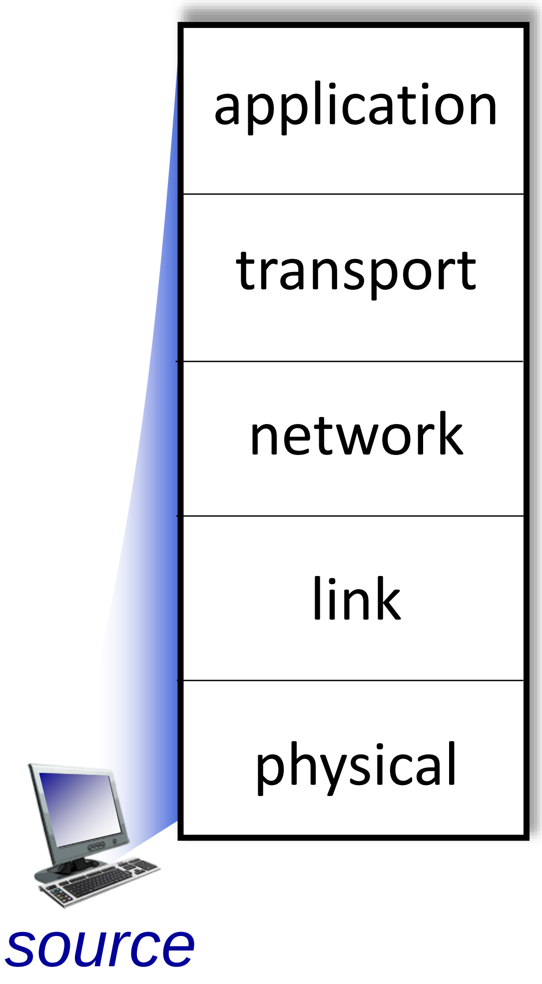
Transport-layer protocol transfers M (e.g., reliably) from one *process* to another, using services of network layer

- transport-layer protocol encapsulates application-layer message, M, with *transport* layer-layer header H_t to create a transport-layer segment
  - H_t used by transport layer protocol to implement its service



destination

Services, Layering and Encapsulation



M

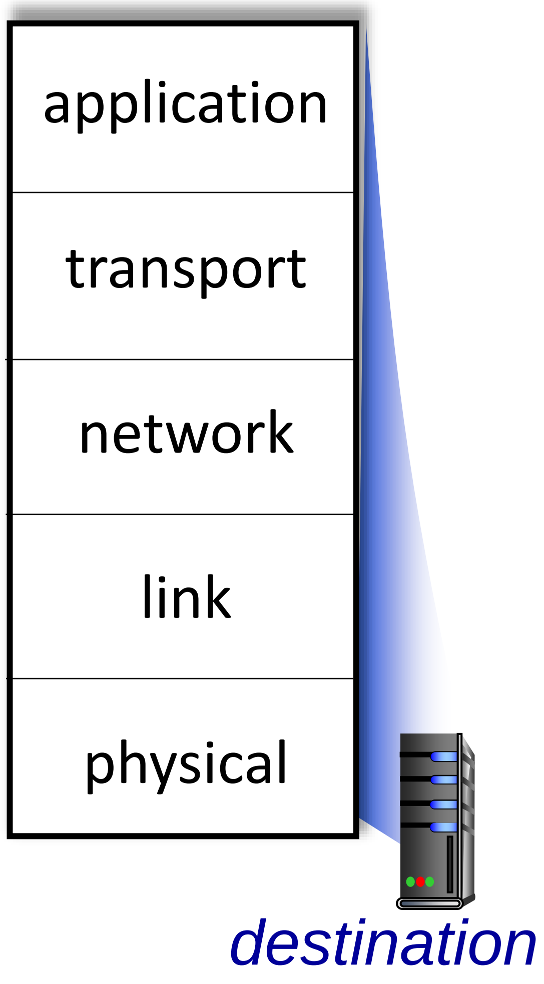
H_t M

Transport-layer protocol transfers M (e.g., reliably) from one *process* to another, using services of network layer

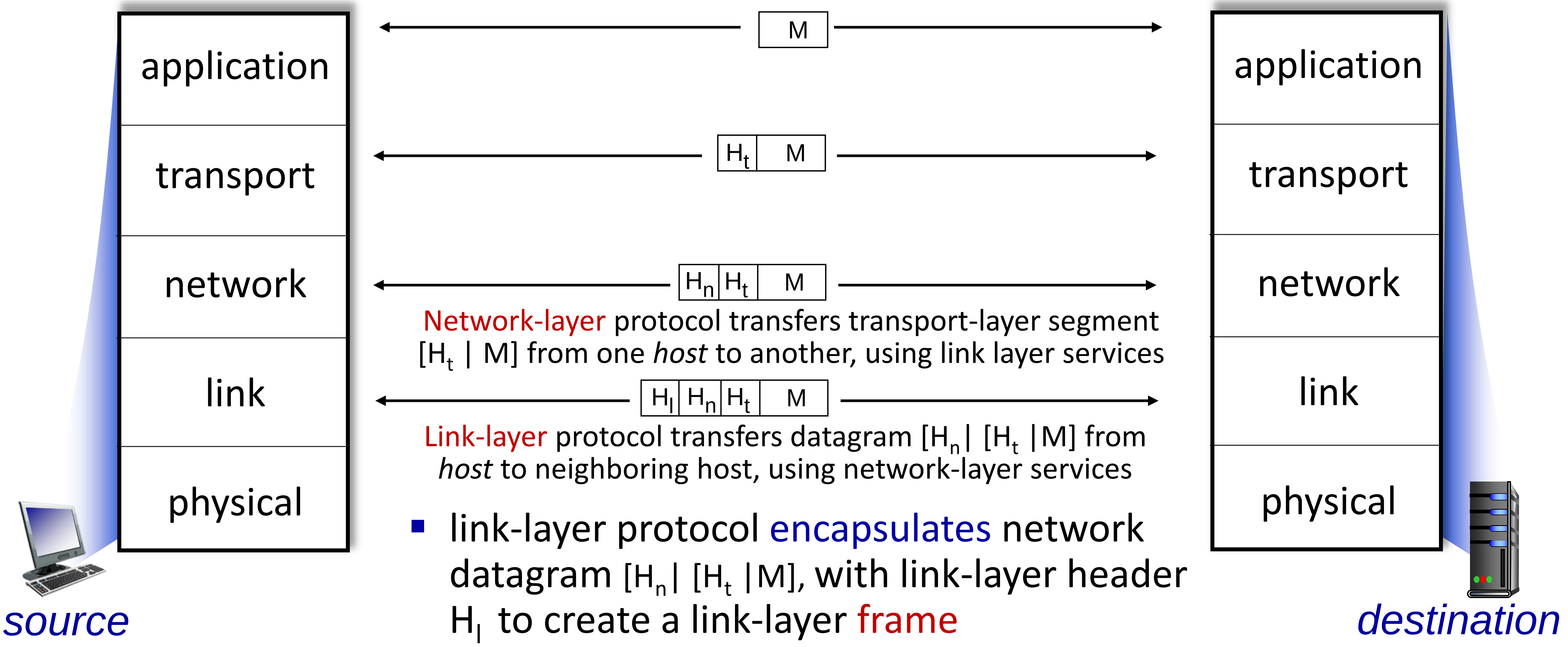
H_n H_t M

Network-layer protocol transfers transport-layer segment [H_t | M] from one *host* to another, using link layer services

- network-layer protocol encapsulates transport-layer segment [H_t | M] with network layer-layer header H_n to create a network-layer datagram
  - H_n used by network layer protocol to implement its service

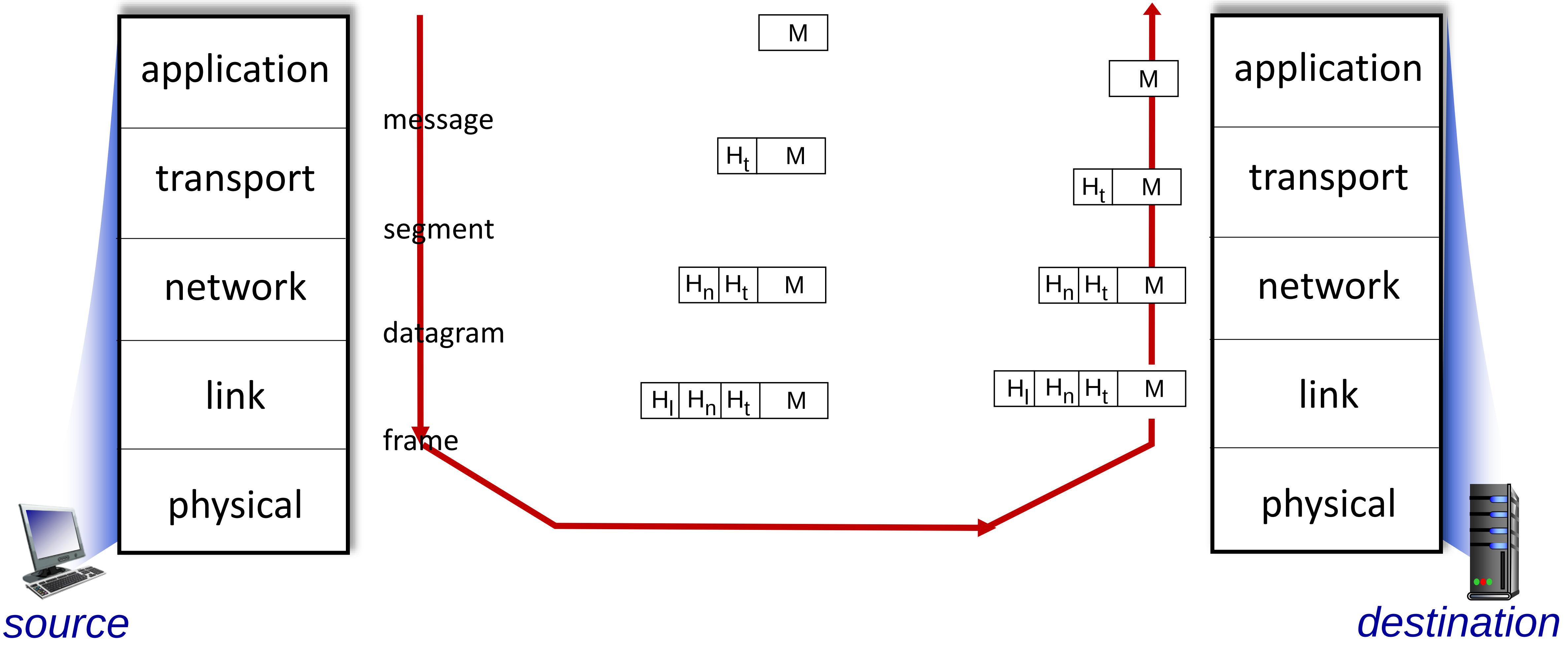


Services, Layering and Encapsulation

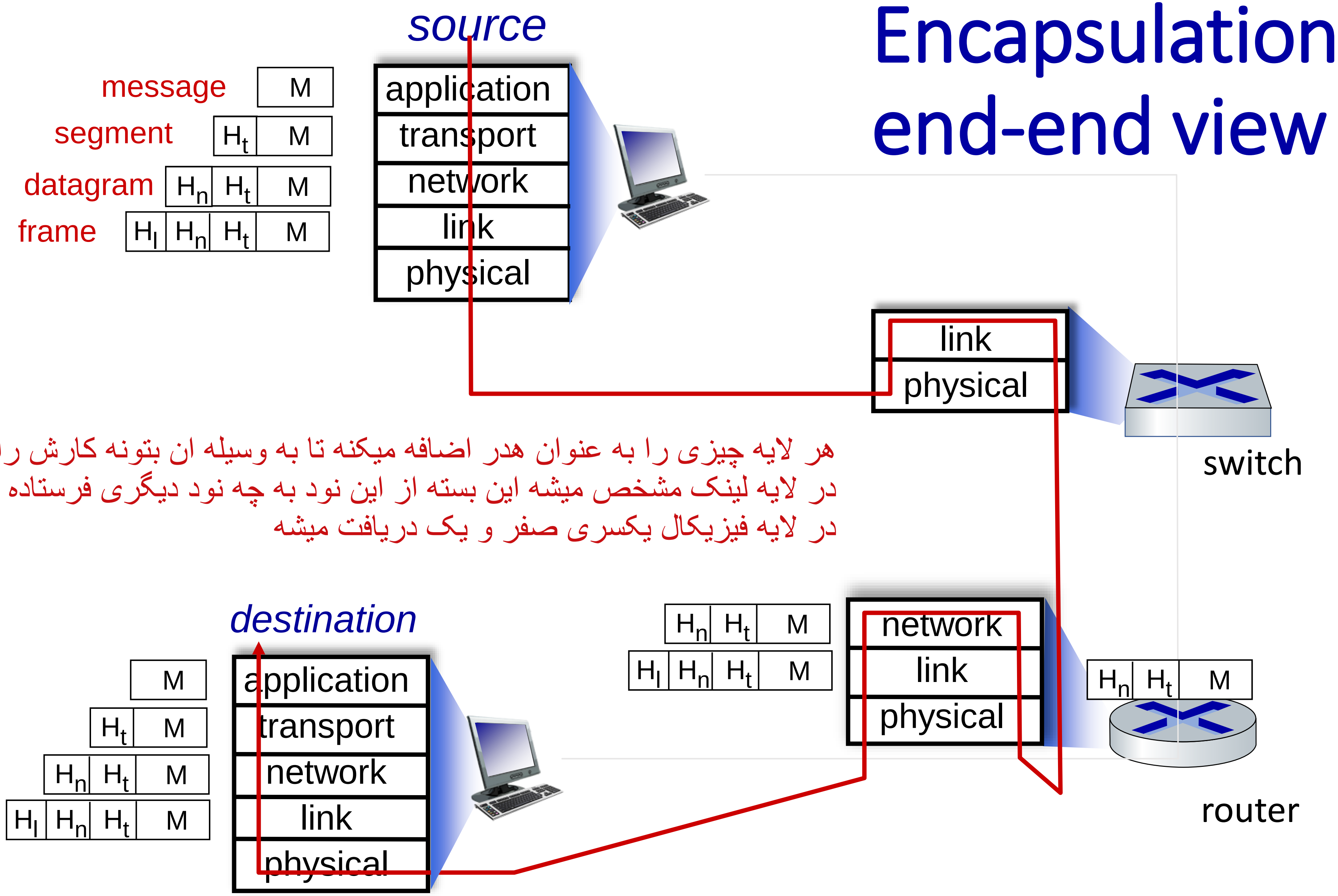


- link-layer protocol encapsulates network datagram $[H_n$| $[H_t$ |M], with link-layer header H_l to create a link-layer frame

Services, Layering and Encapsulation



Encapsulation: an end-end view



هر لایه چیزی را به عنوان هدر اضافه میکنه تا به وسیله ان بتونه کارش را انجام بده
در لایه لینک مشخص میشه این بسته از این نود به چه نود دیگری فرستاده بشه
در لایه فیزیکال یکسری صفر و یک دریافت میشه

Chapter 1: roadmap

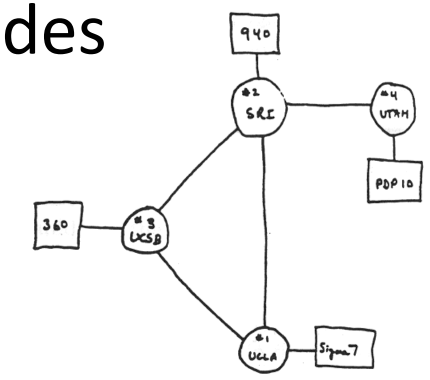
- What *is* the Internet?
- What *is* a protocol?
- Network edge: hosts, access network, physical media
- Network core: packet/circuit switching, internet structure
- Performance: loss, delay, throughput
- Security
- Protocol layers, service models
- History

Internet history

1961-1972: Early packet-switching principles

- 1961: Kleinrock - queueing theory shows effectiveness of packet-switching
- 1964: Baran - packet-switching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: first ARPAnet node operational

- 1972:
  - ARPAnet public demo
  - NCP (Network Control Protocol) first host-host protocol
  - first e-mail program
  - ARPAnet has 15 nodes



THE ARPA NETWORK

Internet history

1972-1980: Internetworking, new and proprietary networks

- 1970: ALOHAnet satellite network in Hawaii
- 1974: Cerf and Kahn - architecture for interconnecting networks
- 1976: Ethernet at Xerox PARC
- late70's: proprietary architectures: DECnet, SNA, XNA
- 1979: ARPAnet has 200 nodes

Cerf and Kahn's internetworking principles:

- minimalism, autonomy - no internal changes required to interconnect networks
- best-effort service model
- stateless routing
- decentralized control

define today's Internet architecture

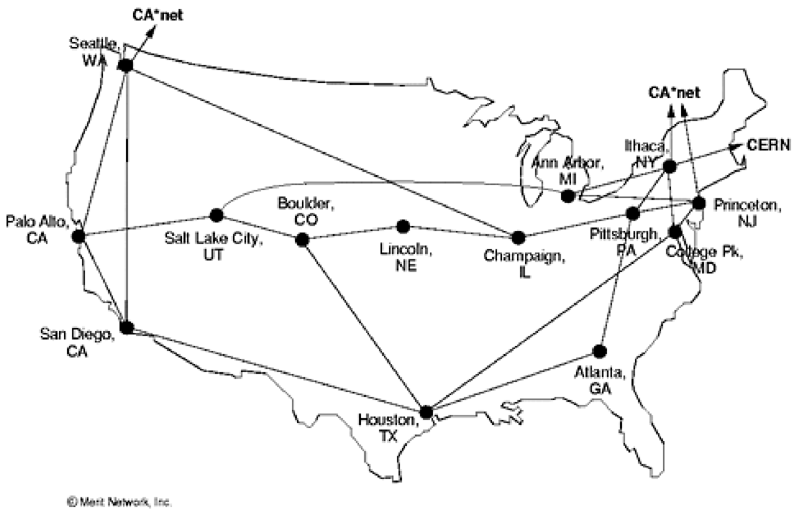
Internet history

1980-1990: new protocols, a proliferation of networks

- 1983: deployment of TCP/IP
- 1982: smtp e-mail protocol defined
- 1983: DNS defined for name-to-IP-address translation
- 1985: ftp protocol defined
- 1988: TCP congestion control

- new national networks: CSnet, BITnet, NSFnet, Minitel
- 100,000 hosts connected to confederation of networks

NSFNET T1 Network 1991



Internet history

1990, 2000s: commercialization, the Web, new applications

- early 1990s: ARPAnet decommissioned
- 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- early 1990s: Web
  - hypertext [Bush 1945, Nelson 1960's]
  - HTML, HTTP: Berners-Lee
  - 1994: Mosaic, later Netscape
  - late 1990s: commercialization of the Web

late 1990s – 2000s:

- more killer apps: instant messaging, P2P file sharing
- network security to forefront
- est. 50 million host, 100 million+ users
- backbone links running at Gbps

Internet history

2005-present: scale, SDN, mobility, cloud

- aggressive deployment of broadband home access (10-100's Mbps)
- 2008: software-defined networking (SDN)
- increasing ubiquity of high-speed wireless access: 4G/5G, WiFi
- service providers (Google, FB, Microsoft) create their own networks
  - bypass commercial Internet to connect "close" to end user, providing "instantaneous" access to social media, search, video content, …
- enterprises run their services in "cloud" (e.g., Amazon Web Services, Microsoft Azure)
- rise of smartphones: more mobile than fixed devices on Internet (2017)
- ~18B devices attached to Internet (2017)

Chapter 1: summary

We’ve covered a “ton” of material!

- Internet overview
- what’s a protocol?
- network edge, access network, core
  - packet-switching versus circuit-switching
  - Internet structure
- performance: loss, delay, throughput
- layering, service models
- security
- history

You now have:

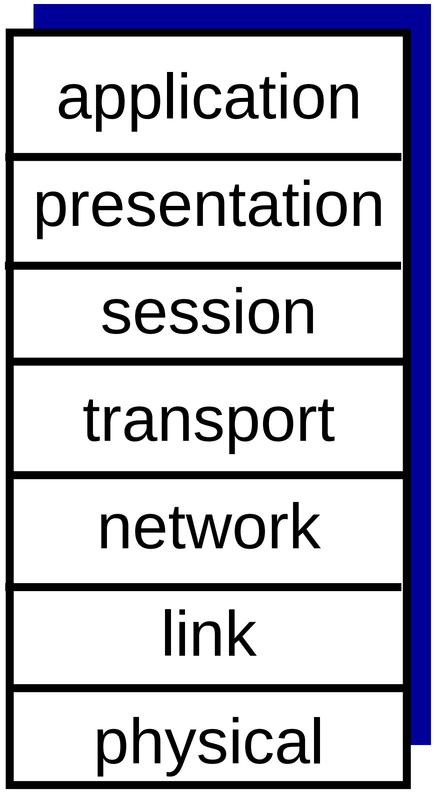
- context, overview, vocabulary, “feel” of networking
- more depth, detail, *and fun* to follow*!*

Additional Chapter 1 slides

ISO/OSI reference model

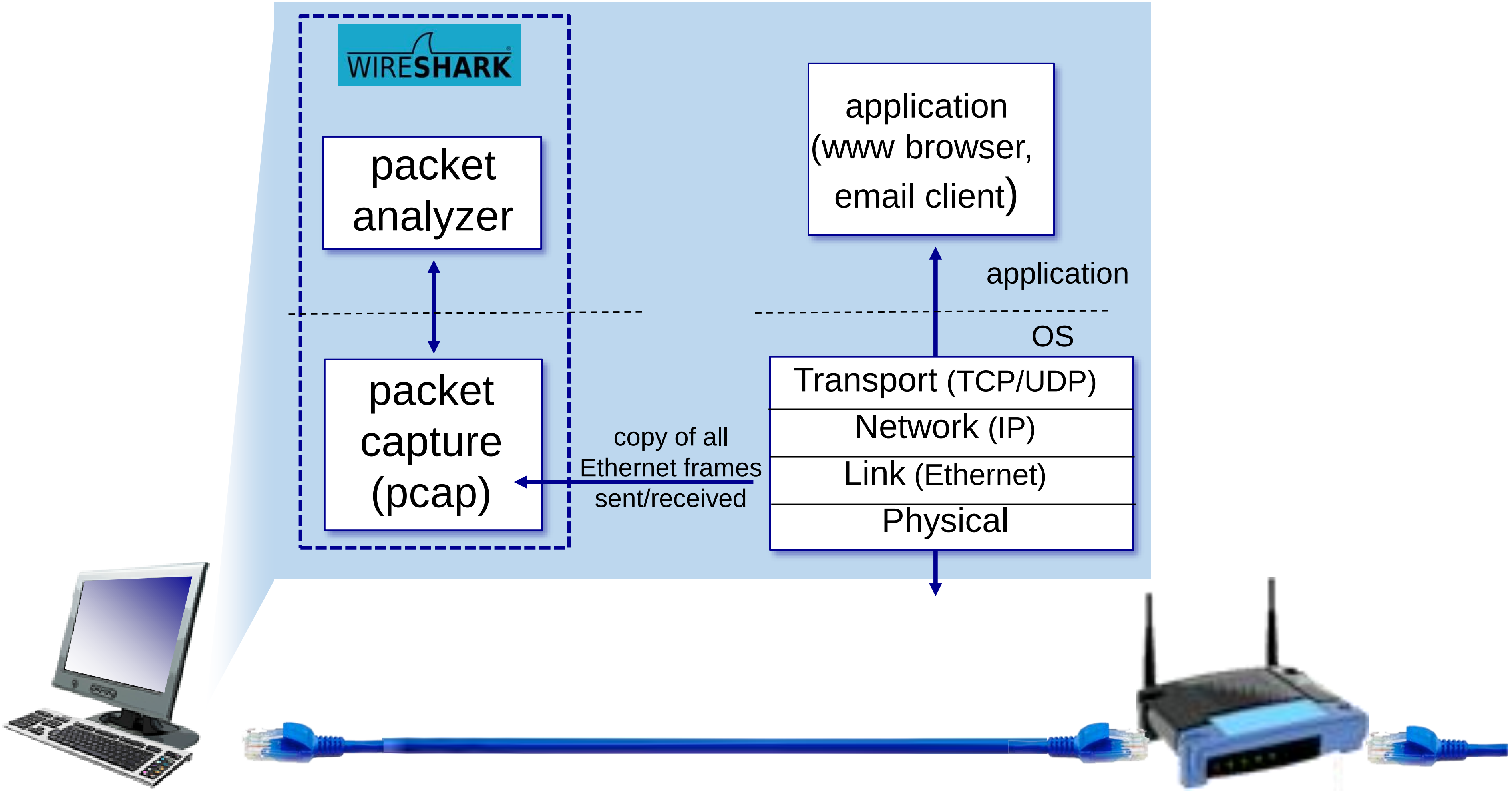
Two layers not found in Internet protocol stack!

- *presentation:* allow applications to interpret meaning of data, e.g., encryption, compression, machine-specific conventions
- *session:* synchronization, checkpointing, recovery of data exchange
- Internet stack "missing" these layers!
  - these services, *if needed,* must be implemented in application
  - needed?



The seven layer OSI/ISO reference model

Wireshark



WIRESHARK
packet analyzer
packet capture (pcap)
copy of all Ethernet frames sent/received
application (www browser, email client)
application
OS
Transport (TCP/UDP)
Network (IP)
Link (Ethernet)
Physical