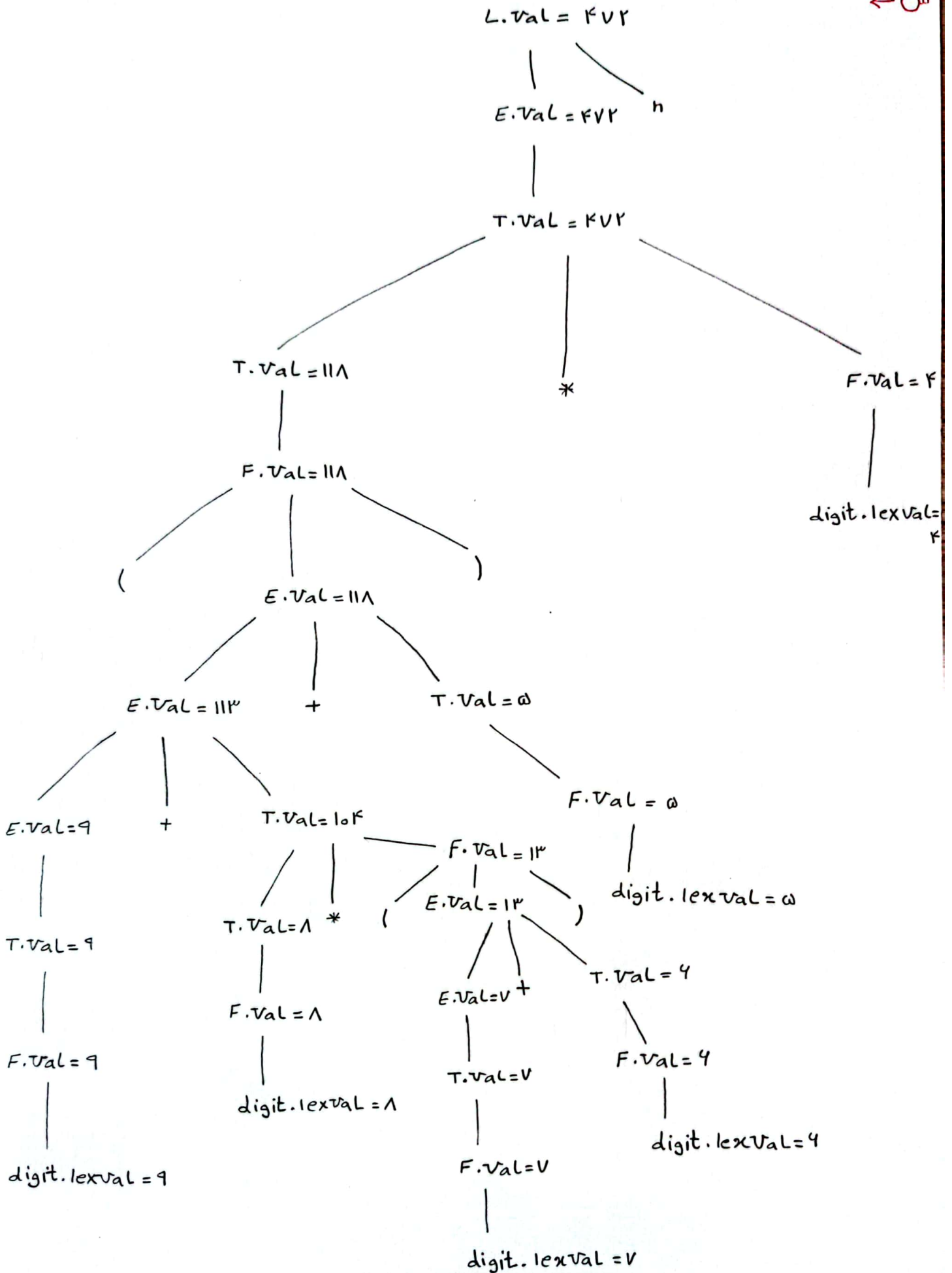
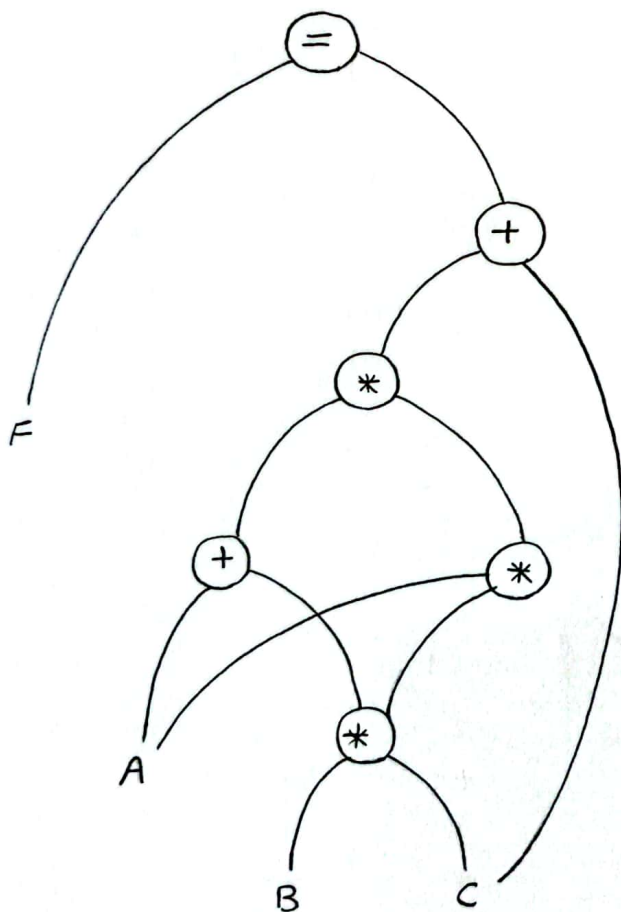


س ←



production	semantic rules	← س ۲
1) $S \rightarrow L_1.L_2$	$S.val = L_1.val + L_2.val / L_2.\#$	
2) $S \rightarrow L$	$S.val = L.val$	
3) $L \rightarrow L_1 B$	$L.val = L_1.val * \gamma + B.val$	
	$L.\# = L_1.\# * \gamma$	
4) $L \rightarrow B$	$L.val = B.val$	
	$L.\# = \gamma$	
5) $B \rightarrow 1$	$B.val = 1$	
6) $B \rightarrow 0$	$B.val = 0$	



Triple:

address	operation	arg1	arg2	← س ک
0	+	z	2	
1	param	(0)		
2	param	x		
3	call	f	2	
4	*	c	d	
5	[ ] =	b	i	
6	-	(5)	(4)	
7	+	(6)	(3)	
8	[ ] =	a	i	
9	=	(8)	(7)	

Quadruple:

address	operation	arg1	arg2	result
0	+	z	2	t1
1	param	t1		
2	param	x		
3	call	f	2	t2
4	*	c	d	t3
5	[ ] =	b	i	t4
6	-	t4	t3	t5
7	+	t5	t2	t6
8	[ ] =	a	i	t7
9	=	t6		t7

الف)  $L_r: \text{if } x > 0 \text{ goto } L_r$   
 $\text{goto } L_1$

$L_r: \text{if } x < 100 \text{ goto } L_r$   
 $\text{goto } L_1$

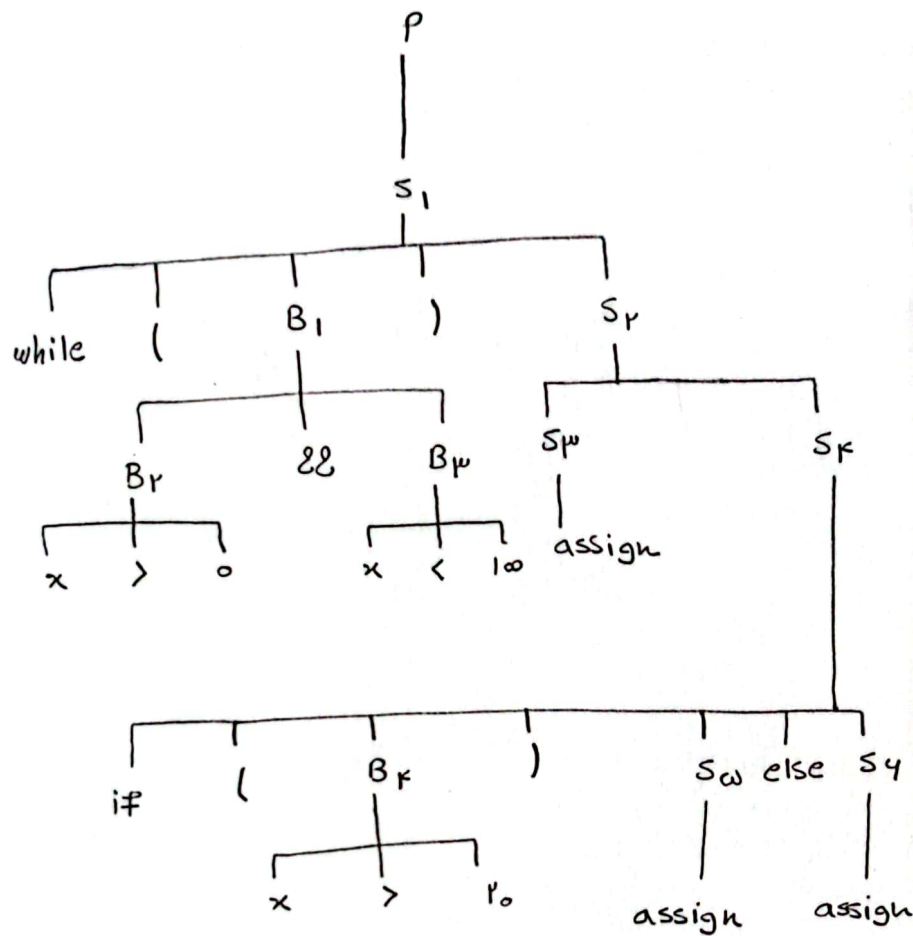
$L_p: x = x + 1$

$L_w: \text{if } x > 10 \text{ goto } L_y$   
 $\text{goto } L_v$

$L_y: x = x + 2$   
 ~~$\text{goto } L_p$~~

$L_v: x = x + 1$   
 $\text{goto } L_p$

$L_1:$



الف

$B_1: i = 1$   
goto  $B_r$

$B_r: j = 1$   
goto  $B_p$

$B_p: t_1 = (i \ll 3) + r * i$

$t_r = t_1 + j - 11$

$a[t_r] = 0.0$

$j = j + 1$

if  $j \leq 1$  goto  $B_p$

else goto  $B_r$

$B_r: i = i + 1$

if  $i \leq 10$  goto  $B_r$

else goto L

ب

$dp = 0$

$i = 0$

L:  $t_1 = i * 1$

$T_r = A[t_1]$

$t_r = B[t_1]$

$t_w = T_r * t_r$

$dp = dp + t_w$

$i = i + 1$

if  $i < n$  goto L