

# OPERATOR PRECEDENCE PARSING

CSE 340 FALL 2021

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Notes based on the  
“Dragon Book”

# Parsing Operator Grammars

The grammar we have seen for expressions does not include the operator minus ('-').

This is not an oversight!

We can write the following grammar

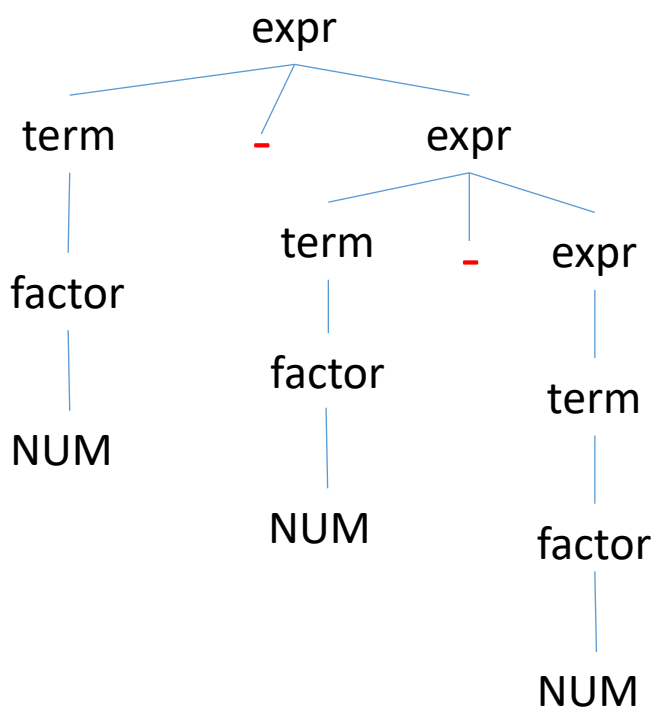
Expr  $\rightarrow$  term - Expr  
Expr  $\rightarrow$  term + Expr  
Expr  $\rightarrow$  term

but that would not work!

How do we parse the following?

1 - 2 - 3

According to the grammar above, we get



According to this tree,

1 - 2 - 3 = 2 !!!

# Parsing expressions with minus

The issue is that minus is left associative and the grammar treats minus as right-associative

Left associative grouping (**correct**)

$$1 - 2 - 3$$

$$(1 - 2) - 3$$

$$((1-2) - 3)$$

Right associative grouping (**wrong**)

$$1 - 2 - 3$$

$$1 - (2 - 3)$$

$$(1 - (2 - 3))$$

# Parsing expressions with minus

We can attempt to fix the problem by using the following grammar

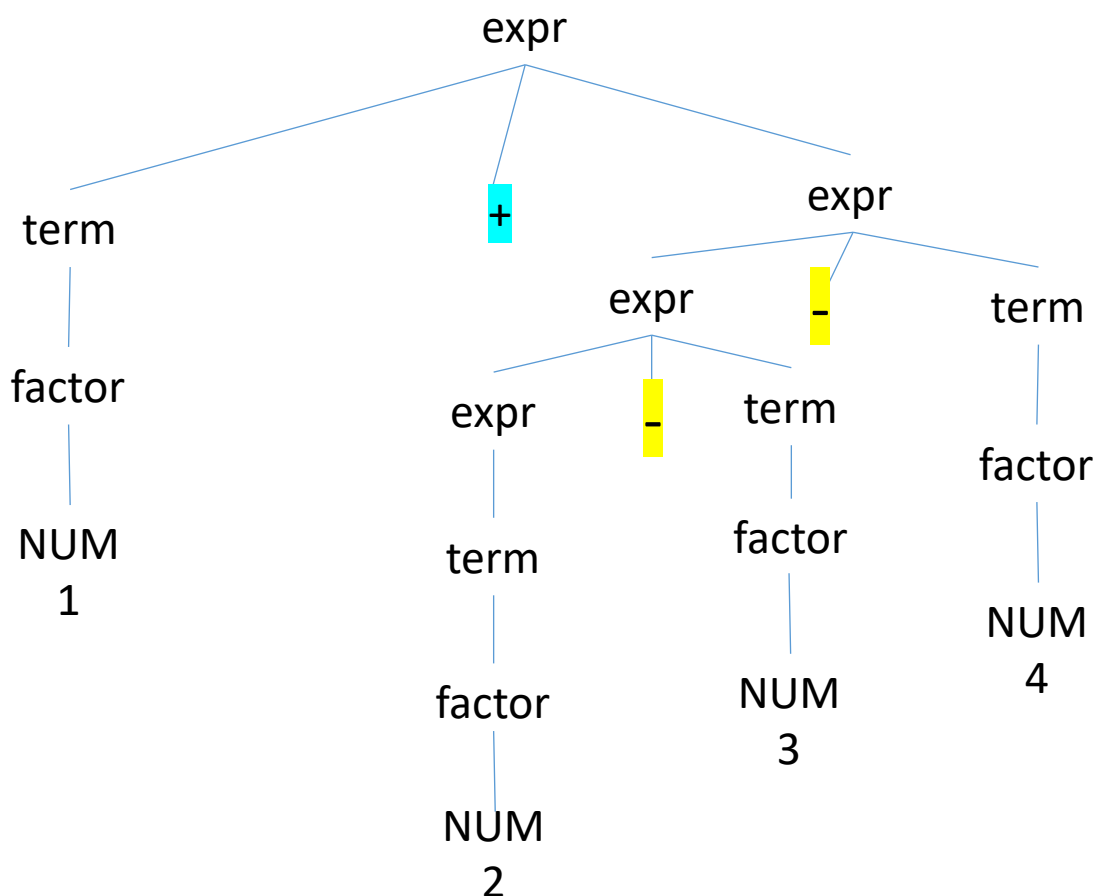
$\text{expr} \rightarrow \text{expr} - \text{term}$

$\text{expr} \rightarrow \text{term} + \text{expr}$

$\text{expr} \rightarrow \text{term}$

This grammar would give the following parsing for

1 + 2 - 3 - 4



# Parsing expressions with minus

We can attempt to fix the problem by using the following grammar

`expr -> expr - term`

`expr -> term + expr`

`expr -> term`

We cannot parse this grammar with a recursive descent parser!

```
parse_expr()
{
    // expr -> expr - term

    ....

    parse_expr() // infinite loop !!
}
```

We need another way to parse such expressions!

# A NEAT TRICK FROM FORTRAN COMPILER!

$a + b * c - d$

add ( ( ( at the beginning

replace every + with ) ) ) + ( ( (

replace every - with ) ) ) - ( ( (

replace every \* with ) ) \* ( (

replace every ^ with ) ^ (

add ) ) ) at the end

We get

(( ( a ) ) ) + ( ( ( b ) ) ) \* ( ( c ) ) ) - ( ( ( d ) ) )

(( ( a ) ) ) + ( ( ( b ) ) ) \* ( ( c ) ) ) - ( ( ( d ) ) )

Always works !

We can then parse with a simple parser that only has to worry about matching parentheses

# Operator Grammar

A grammar is called an operator grammar if

1. there is no righthand side of a rule which has two adjacent non-terminal
2. there is no rule of the form  $A \rightarrow \varepsilon$

**Example 1**      $E \rightarrow E A E \mid ( E ) \mid -E \mid ID$   
                   $A \rightarrow + \mid - \mid * \mid / \mid ^$

is not an operator grammar because of  $E A E$  has three adjacent non-terminals

**Example 2**      $E \rightarrow E + E \mid E - E \mid E * E \mid E / E$   
                   $\mid E ^ E \mid ( E ) \mid - E \mid ID$

is an operator grammar

# OPERATOR PRECEDENCE RELATIONSHIPS

To parse operator grammar, we first define parsing precedence relationships between the terminals of the grammar

We also introduce a new symbol \$

parsing precedence relationships

- $<\cdot$  yields precedence to
- $\cdot>$  takes precedence over
- $\doteq$  has the same precedence as

These are not the same as the operator precedence levels.

These are used in guiding the parsing

You can think of  $<\cdot$   $\cdot>$  as matching parentheses that group what appears between them (this should become clearer with the examples)

There is a theory to determine these relationship for an unambiguous operator grammar. We will only look at heuristics for common expressions.

The parsing algorithm assumes that we already have a table that defines these relationships



## EXAMPLE

[illegible]

# Parsing Algorithm

**Input**         $w \$$   
**Output**     parse tree with E in all internal nodes  
**Initially**   stack contains \$, scanning starts at the start of w

**repeat**

```
if $ is on top of the stack and lexer.peek() = $ // EOF
    return;
else
{
    t = lexer.peek(); b = t.type;           // next token from w
    a = stack.terminalpeek().type;         // terminal at the top of stack
                                           // or just below if top is non-terminal

    if (table[a][b] == '<.') | ( table[a][b] = '=') // shift
        t = lexer.getToken();
        stack.push(t)
    else if (table[a][b] == '>') // reduce
    {
        RHS = an empty stack
        repeat
            s = stack.pop() // pop terminals and
                           // non-terminals
            if s is a terminal
                last_popped_term = s
            RHS.push(s)
        until ( ( is_a_terminal(stack.peek()) ) and
                ( table[stack.terminalpeek()][last_popped_term] == '<.' ) )

        if E → RHS rule exists // RHS calculated above
        {
            reduce E → RHS
            push E // we can think of E as the
                  // root of subtree for E → RHS
        }
        else
            syntax_error()
    }
    else
        syntax_error();
}
```

**Note:**

- `stack.peek()` peeks at the symbol at the top of the stack, which could be a terminal or a non-terminal.
- `stack.terminalpeek()` peeks at the terminal closest to the top of the stack.

## EXAMPLE

[illegible]

## EXAMPLE

## Stack

\$

Input

$$a + b * (c + d) - a$$

## Action

 $\prec$ [illegible]

## EXAMPLE

# Stack

\$

 $\prec$ 

## Input

$$a + b * (c + d) - a \text{ shift}$$

## Action

shift

[illegible]

## EXAMPLE

# Stack

\$

\$a

## Input

$a + b * (c + d) - a$  shift

$$+ b * (c + d) - a$$

## Action

shift

[illegible]

## EXAMPLE

## Stack

\$

\$a

 $\cdot \succ$ 

## Input

$$a + b * (c + d) - a \text{ shift}$$
$$+ b * (c + d) - a$$

## Action

shift

[illegible]

## EXAMPLE

## Stack

\$

\$a

 $\cdot \succ$ 

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$

## Action

shift

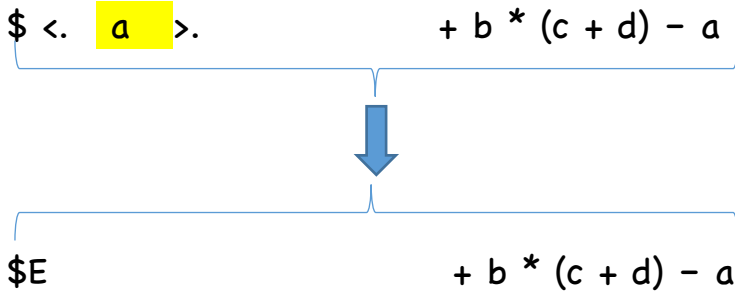
reduce  $E \rightarrow ID$

[illegible]



## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce E -> ID

[illegible]

## EXAMPLE

## Stack

\$

\$a

$\$E$

id

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

[illegible]

## EXAMPLE

## Stack

\$

\$a

$\$E$

id

 $\dot{\gamma}$ 

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

id

Input

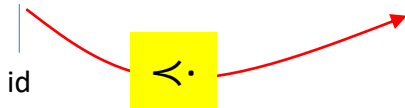
$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

shift

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

$\$E_+$

id

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

shift

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

$\$E_+$

id

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$

## Action

shift

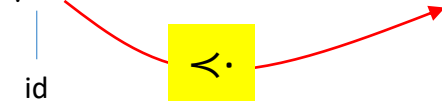
reduce  $E \rightarrow ID$

shift

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
<del>\$E+</del>	<del><math>b * (c + d) - a</math></del>	shift

[illegible]

## EXAMPLE

## Stack

\$

\$a

$\$E_+$

$\$E+b$

id

Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$
$$* (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

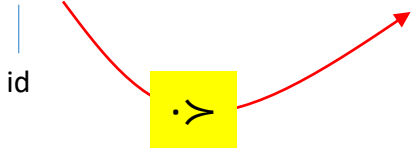
shift

	+	-	*	/	^	id	(	)	\$
+	·	·	·	·	·	·	·	·	·
-	·	·	·	·	·	·	·	·	·
*	·	·	·	·	·	·	·	·	·
/	·	·	·	·	·	·	·	·	·
^	·	·	·	·	·	·	·	·	·
id	·	·	·	·	·			·	·
(	·	·	·	·	·	·	·	=	
)	·	·	·	·	·			·	·
\$	·	·	·	·	·	·	·		



## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	

[illegible]

## EXAMPLE

## Stack

\$

\$a

$\$E_+$

$\$E+b$

id

 $\cdot \succ$ 

Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$
$$* (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

shift

reduce  $E \rightarrow ID$

[illegible]

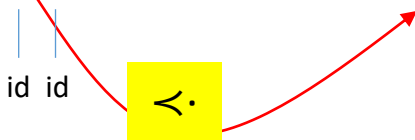
## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	
<div style="display: flex; justify-content: space-around; width: 100px;"> <div style="border-left: 1px solid blue; height: 20px; margin-bottom: 5px;"></div> <div style="border-left: 1px solid blue; height: 20px; margin-bottom: 5px;"></div> </div> <div style="display: flex; justify-content: space-around; width: 100px;"> <span>id</span> <span>id</span> </div>		

[illegible]

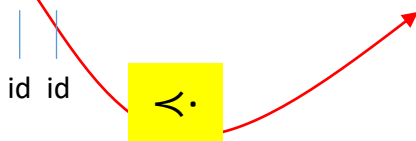
## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	

[illegible]

# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
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\$E+E	$* (c + d) - a$	shift

[illegible]

## EXAMPLE

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\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	

id id

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
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[illegible]

## EXAMPLE

Stack	Input	Action
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[illegible]



## EXAMPLE

Stack	Input	Action
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[illegible]

## EXAMPLE

Stack	Input	Action
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\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	
<div> <div></div> <div></div> </div> <div>id id</div>		

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
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\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	

Diagram illustrating the LR(0) item sets and transitions for the grammar:

```

graph TD
    S["id id"] --> L["<."]
    L --> S
    L --> R["c + d) - a"]
  
```

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift

id id

$< \cdot$

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
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\$E+E*	$(c + d) - a$	shift
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id id

[illegible]

## EXAMPLE

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\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	

id

id

$\cdot >$

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
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id

id

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
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id id id		

[illegible]



## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
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\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	

id

id

id

$< \cdot$

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
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\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift

[illegible]

## EXAMPLE

Stack	Input	Action
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\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	

id id id

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
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\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	

id

id

id

$<.$

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
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\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift

id

id

id

$< .$

[illegible]

## EXAMPLE

Stack	Input	Action
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\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
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\$E+E	$* (c + d) - a$	shift
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\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	

id id id

[illegible]

# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
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\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	

id

id

id

→

$\cdot >$

→

[illegible]

# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
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\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$

[illegible]



## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	

id

id

id

id

	+	-	*	/	^	id	(	)	\$
+	→	→	→	→	→	→	→	→	→
-	→	→	→	→	→	→	→	→	→
*	→	→	→	→	→	→	→	→	→
/	→	→	→	→	→	→	→	→	→
^	→	→	→	→	→	→	→	→	→
id	→	→	→	→	→			→	→
(	→	→	→	→	→	→	→	=	
)	→	→	→	→	→			→	→
\$	→	→	→	→	→	→	→		

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	

Diagram illustrating the LR(0) item set  $E+E*E+$  and its transitions. The item set is shown in a yellow box. Red arrows indicate transitions: one to the right on the non-terminal  $E$  and one down on the terminal  $+$ . Below the item set, the stack contents  $id\ id\ id\ id$  are shown with vertical lines connecting them to the corresponding positions in the item set string.

	+	-	*	/	^	id	(	)	\$
+	→	→	→	→	→	→	→	→	→
-	→	→	→	→	→	→	→	→	→
*	→	→	→	→	→	→	→	→	→
/	→	→	→	→	→	→	→	→	→
^	→	→	→	→	→	→	→	→	→
id	→	→	→	→	→			→	→
(	→	→	→	→	→	→	→	=	
)	→	→	→	→	→			→	→
\$	→	→	→	→	→	→	→		

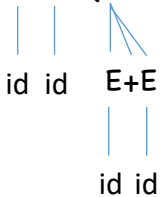
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$

	+	-	*	/	^	id	(	)	\$
+	→	→	→	→	→	→	→	→	→
-	→	→	→	→	→	→	→	→	→
*	→	→	→	→	→	→	→	→	→
/	→	→	→	→	→	→	→	→	→
^	→	→	→	→	→	→	→	→	→
id	→	→	→	→	→			→	→
(	→	→	→	→	→	→	→	=	
)	→	→	→	→	→			→	→
\$	→	→	→	→	→	→	→		

# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	



## When do we stop popping the stack?

$$\$ E + E * ( E + E )$$

\$ \lambda. + \lambda. \* \lambda. ( E + E )

when the  $E + E$  is popped, the following hold

1. the top of the stack is a terminal which is (
2. the last popped terminal is +
3. ( < . +

so we stop

what is popped is between a pair  $\langle \cdot \rangle$ :  $\langle \cdot$  RHS of reduction  $\cdot \rangle$

[illegible]

# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	

Diagram illustrating the stack state and input stream during a shift-reduce parse:

- Stack:  $\$E+E*(E$
- Input:  $id\ id\ E+E\ ) - a$
- The stack is partitioned into three sections:  $id\ id$  (under the first  $E$ ),  $E+E$  (under the second  $E$ ), and a yellow box containing  $\dot{=}$  (under the third  $E$ ).
- Blue lines show the derivation of the first  $E$  from  $id\ id$  and the second  $E$  from  $E+E$ .
- A red arrow points from the yellow box to the closing parenthesis  $)$  in the input, indicating the next shift action.

# EXAMPLE

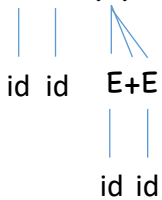
Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift

Diagram illustrating the stack state and input stream during a shift action. The stack contains  $E+E*(E$ . The input stream is  $a + b * (c + d) - a$ . The diagram shows the stack divided into three parts:  $id\ id$ ,  $E+E$ , and a yellow box containing a division symbol  $\div$ . Blue lines connect the  $id\ id$  part of the stack to the  $id\ id$  part of the input stream. Blue lines connect the  $E+E$  part of the stack to the  $E+E$  part of the input stream. A red arrow points from the yellow box to the  $\div$  symbol in the input stream.

# EXAMPLE

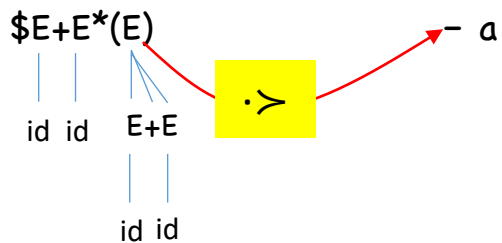
Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	





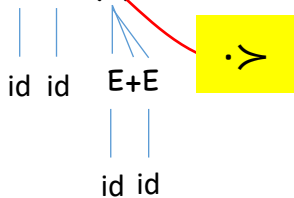
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift



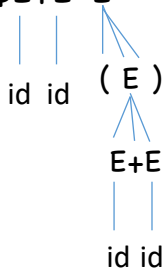
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$



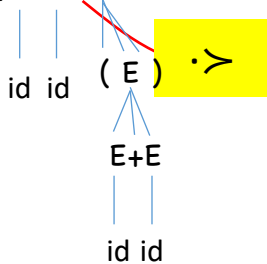
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	



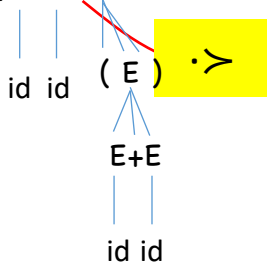
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	



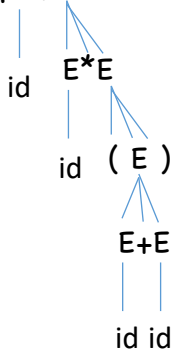
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	reduce $E \rightarrow E*E$



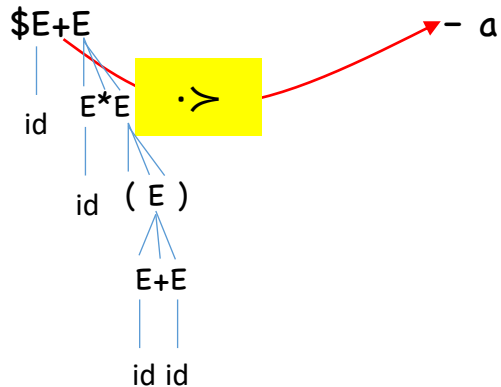
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	



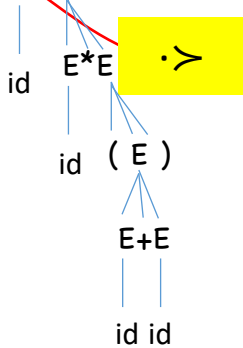
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	reduce $E \rightarrow E*E$



# EXAMPLE

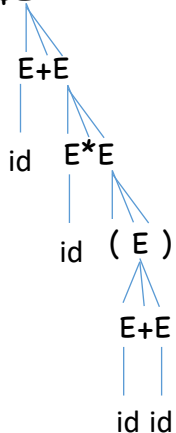
Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	reduce $E \rightarrow E+E$





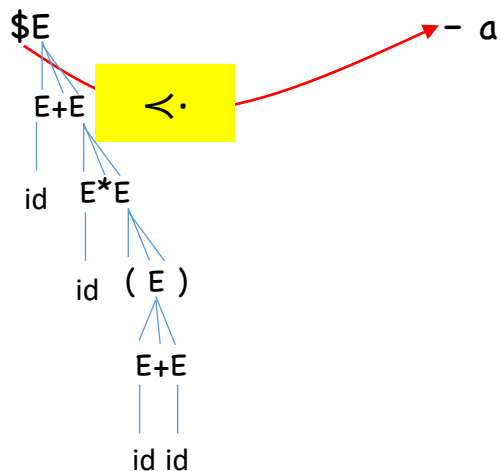
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	reduce $E \rightarrow E+E$
\$E	$- a$	



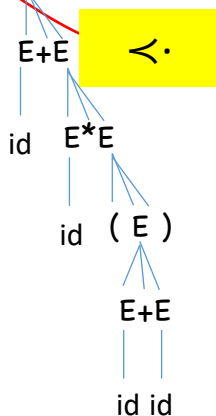
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	reduce $E \rightarrow E+E$



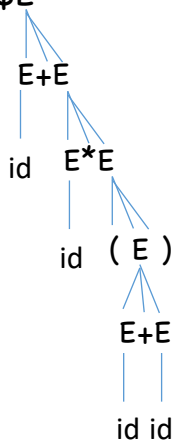
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	reduce $E \rightarrow E+E$
\$E	$- a$	shift



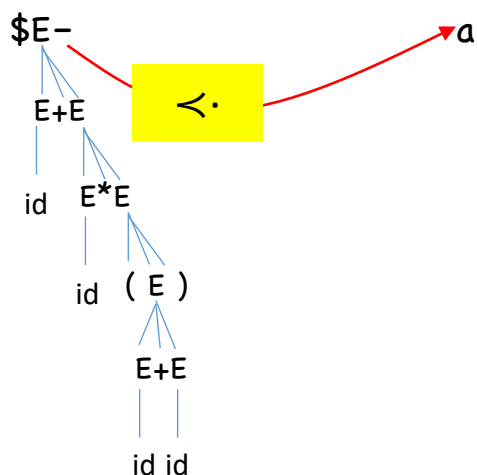
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	reduce $E \rightarrow E+E$
\$E	$- a$	shift
\$E-	$a$	



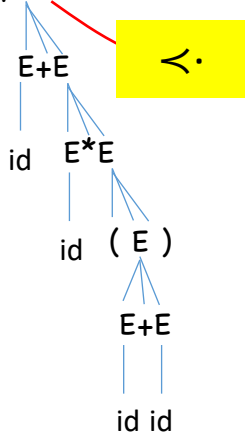
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
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\$E+E*(E	$+ d) - a$	shift
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\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	reduce $E \rightarrow E+E$
\$E	$- a$	shift



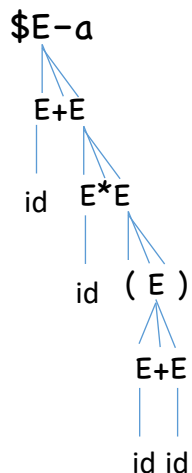
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
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\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
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\$E+E	$- a$	reduce $E \rightarrow E+E$
\$E	$- a$	shift
\$E-	$- a$	shift



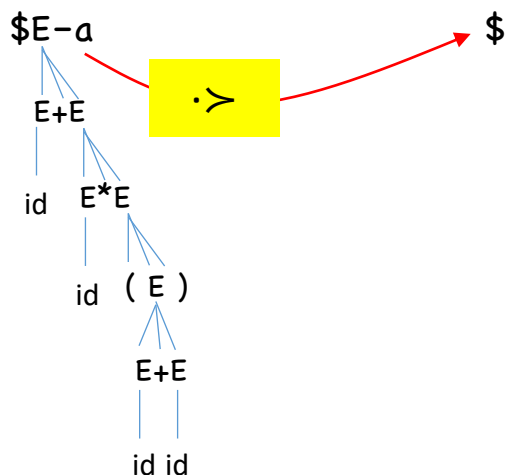
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
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\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
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\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	reduce $E \rightarrow E+E$
\$E	$- a$	shift
\$E-	$a$	shift



# EXAMPLE

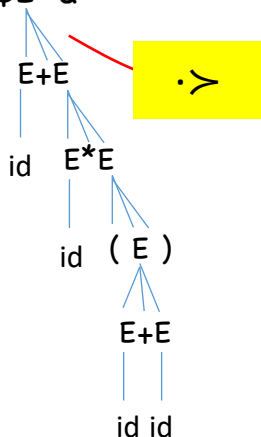
Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
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\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
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\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	reduce $E \rightarrow E+E$
\$E	$- a$	shift
\$E-	$a$	shift





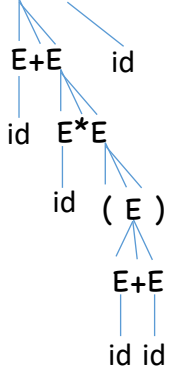
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
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\$E+E	$* (c + d) - a$	shift
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\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	reduce $E \rightarrow E+E$
\$E	$- a$	shift
\$E-	$a$	shift
\$E-a	\$	reduce $E \rightarrow ID$



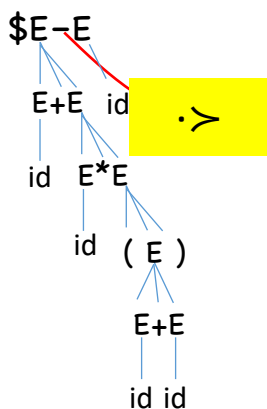
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
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\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	reduce $E \rightarrow E+E$
\$E	$- a$	shift
\$E-	$a$	shift
\$E-a	$\$$	reduce $E \rightarrow ID$
\$E-E	$\$$	



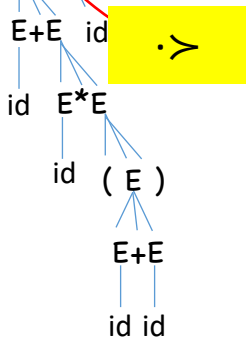
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
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\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
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\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	reduce $E \rightarrow E+E$
\$E	$- a$	shift
\$E-	$a$	shift
\$E-a	\$	reduce $E \rightarrow ID$



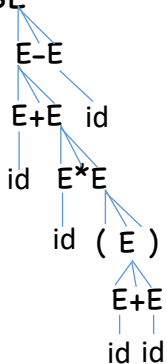
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
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\$E+E*(E	$+ d) - a$	shift
\$E+E*(E+	$d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$
\$E+E*(E	$) - a$	shift
\$E+E*(E)	$- a$	reduce $E \rightarrow (E)$
\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	reduce $E \rightarrow E+E$
\$E	$- a$	shift
\$E-	$a$	shift
\$E-a	$\$$	reduce $E \rightarrow ID$
\$E-E	$\$$	reduce $E \rightarrow E-E$



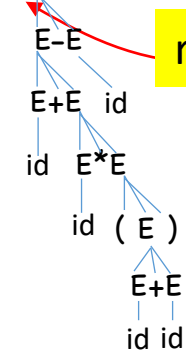
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E+	$b * (c + d) - a$	shift
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\$E+E	$* (c + d) - a$	shift
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$
\$E+E*(E	$+ d) - a$	shift
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\$E	$- a$	shift
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\$E-E	$\$$	reduce $E \rightarrow E-E$
\$E	$\$$	



# EXAMPLE

Stack	Input	Action
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\$E+E*E	$- a$	reduce $E \rightarrow E*E$
\$E+E	$- a$	reduce $E \rightarrow E+E$
\$E	$- a$	shift
\$E-	$a$	shift
\$E-a	$\$$	reduce $E \rightarrow ID$
\$E-E	$\$$	reduce $E \rightarrow E-E$
\$E	$\$$	



return

# Dealing with non-terminals

We have the non-terminals on the stack as they are pushed when a reduction occurs

Given that the grammar is an operator grammar, we can have at most one non-terminal on the top of the stack. There is always a terminal on the top of the stack or just below

In the algorithm we assume that `stack.terminalpeek()` ignores non-terminals and returns the terminal symbol at the top of the stack or just below

# HEURISTIC FOR DETERMINING PRECEDENCE RELATIONSHIPS

We assume we have a set of operators with

- precedence levels
- associativity (left or right)
- operators at the same level have the same associativity

We assume the input is of the form  $w \$$

We have the following heuristics for determining  $\prec$ ,  $\succ$ , and  $\doteq$  relationships between operators

1. if  $op1$  has higher precedence level than  $op2$ , then

- $op1 \succ op2$
- $op2 \prec op1$

**Example:**  $* \succ +$

$+ \prec *$

$\wedge \succ +$

$+ \prec \wedge$



# HEURISTIC FOR DETERMINING PRECEDENCE RELATIONSHIPS

2. if op1 and op2 are operators of the same operator precedence, possibly the same operator, then

- If they are left associative:

- $op1 \cdot > op2$

- $op2 \cdot > op1$

**Example.** + and - are left associative, so we have

+  $\cdot >$  +

+  $\cdot >$  -

-  $\cdot >$  +

-  $\cdot >$  -

- If they are right associative:

- $op1 < \cdot op2$

- $op2 < \cdot op1$

**Example.** ^ is right associative, so we have

^  $\cdot <$  ^

^  $< \cdot$  ^

# HEURISTIC FOR DETERMINING PRECEDENCE RELATIONSHIPS

3. Also, we have the following

1.	op	$\prec\cdot$	ID
2.	ID	$\cdot\succ$	op
3.	op	$\prec\cdot$	(
4.	(	$\prec\cdot$	op
5.	op	$\cdot\succ$	)
6.	)	$\cdot\succ$	op
7.	\$	$\prec\cdot$	op
8.	op	$\cdot\succ$	\$
9.	(	$\doteq$	)
10.	\$	$\prec\cdot$	id
11.	id	$\cdot\succ$	\$
12.	\$	$\prec\cdot$	(
13.	(	$\prec\cdot$	(
14.	(	$\prec\cdot$	id
15.	)	$\cdot\succ$	\$
16.	)	$\cdot\succ$	)
17.	id	$\cdot\succ$	)

# Unary Operators (one operand)

If we have a unary operator `uop` **that is not a binary operator**, we can support it as follows

- `op <· uop` for every other operator `op`. `op` can be unary or binary
- `uop <· op` if `uop` has lower operator precedence level than `op`
- `uop ·> op` if `uop` has higher operator precedence level than `op`

If we have a unary operator **that is also a binary operator**, like MINUS, we cannot incorporate it in the scheme!

**Example** `id*-id` is not easily parsed

One solution is to have the `getToken()` function make the distinction by looking at the context in which the operator appears.

**Example** In FORTRAN a minus sign is unary if the previous token is an operator, LPAREN, COMMA, or EQUAL

It is better to handle this in the lexer than it is to make the parser more complicated

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