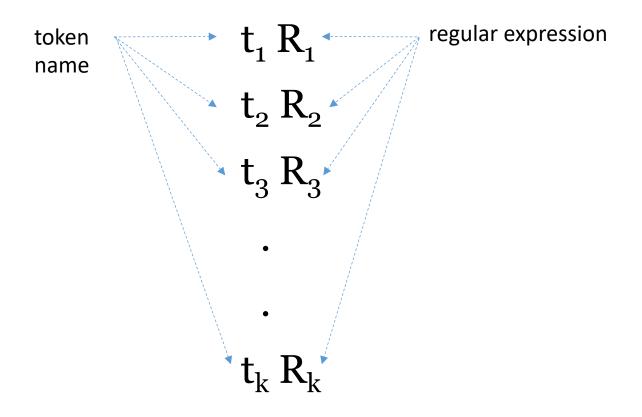
# Implementing my\_GetToken() Automatically

Rida Bazzi

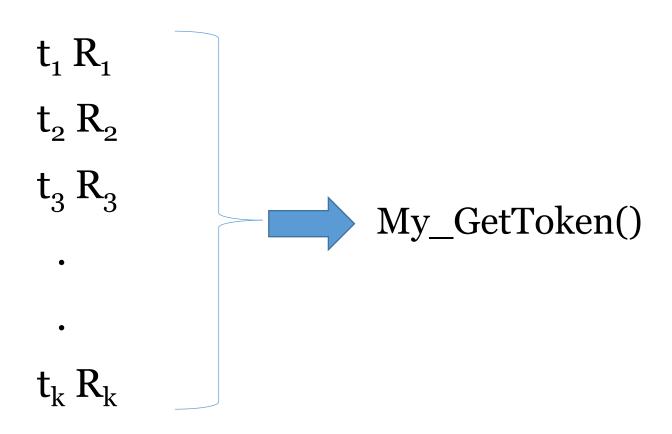
CSE 340 FALL 2021

# This is not meant to replace the project description!

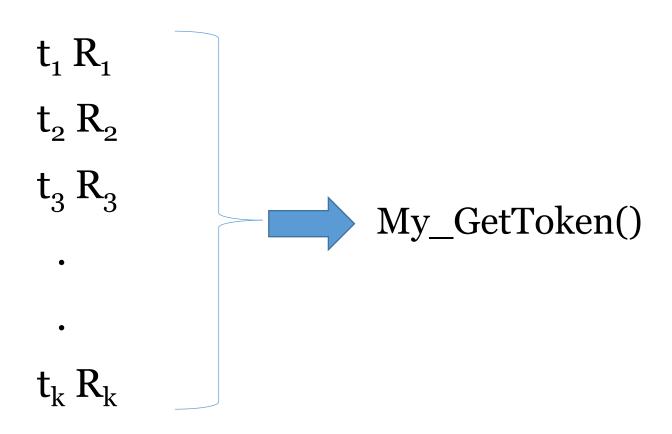
#### Generating My\_GetToken() automatically



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### Implementing my\_GetToken() automatically

- Given a list of token names and regular expressions, one for each token, implement my\_GetToken() function
- The function should correctly implement
  - longest matching prefix rule (we have seen this in the first week)
  - breaking ties according to list order (we have seen this also)
- We can assume that none of the regular expressions given for the tokens has epsilon in its language, which is to be expected, because epsilon is not a token!

**Note**: the language of a regular expression is the set of strings that the expression represents. See project document for more details.

# function match()

#### **Function**

match(REG r, string s, int p)

#### Input is

- r REG
- s is a string
- *p* is a position in string s

#### **Behavior**

- 1. determine the longest possible substring of s, starting at position p, that matches the regular expression represented by r
- 2. If there is no match, that will be indicated.
- 3. Note that match need not return an actual substring, it just needs to return a position p' corresponding to the end of the substring. The substring would be between p and p'

# function my\_getToken()

#### **Function**

my\_GetToken(Token\_List *L*, string *s*, int *p*)

#### Input

- *L* is a list of tokens, where each entry in the list consists of a token name and a REG
- s is a string
- *p* is a position in string s

#### **Behavior**

- 1. call match(r,s,p) for each REG in the list L
- 2. For each REG r, records the longest matching prefix obtained from the call match(r,s,p)
- 3. Returns the token for which match(r,s,p) returns the longest amongst all the prefixes obtained in step 2 and advance the position to reflect that the input is consumed
- 4. If there is a tie, return the token listed first in the list

#### Plan

- I will explain how to construct REGs
- I will then explain how to implement the function match
- Then I will explain how to implement my\_getToken()

# Constructing REGs

- REGs will be constructed recursively as shown on the following pages.
- Each REG is a directed graph
- Each REG has two special nodes
  - a starting node
  - an accepting node
- Labels on edges are characters of the alphabet or epsilon (for which we will use \_ (underscore) as the label

# Approach

1. Transform each expression into a graph that I will call REG (Regular Expression Graph)

- 2. Write a function match (REG r, string s, int p) that, given
  - 1. a REG *r*
  - 2. a string s
  - 3. a position p in the string s,

returns the longest matching substring in *s*, starting at position *p*, that is in the language of the expression of REG r

# Approach (cont'd)

3. Write a function GetToken(Token\_List *L*, string *s*, int *p*)

# Approach (cont'd)

3. Write a function my\_GetToken(Token\_List *L*, string *s*, int *p*) that,

#### given

- *L* is a list of tokens, where each entry in the list consists of a token name and a REG
- s is a string
- *p* is a position in string s

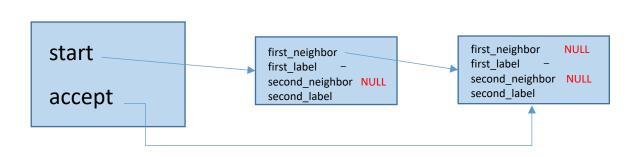
#### will

- 1. call match(r,s,p) for each REG in the list L
- 2. For each REG r, records the longest matching prefix obtained from the call match(r,s,p)
- 3. Returns the token for which match(r,s,p) returns the longest amongst all the prefixes obtained in step 2 and advance the position to reflect that the input is consumed
- 4. If there is a tie, return the token listed first in the list

#### **REG** for \_



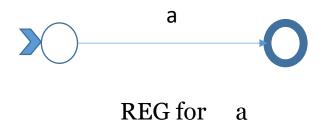
If the regular expression is \_ , the REG can be constructed immediately. The resulting REG is illustrated above as a graph and below as a data structure



#### Notation

- accept node
- > start node

#### REG for a

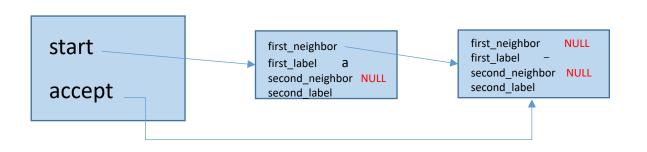


Notation

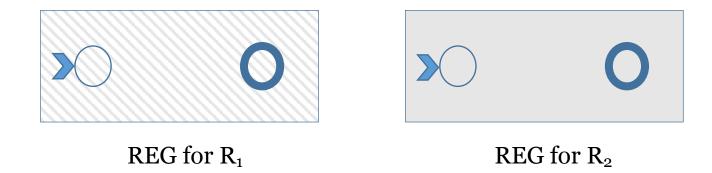
If the regular expression is a, where a is a character of the alphabet or a digit, the REG can also be constructed immediately. The resulting REG is illustrated above as a graph and below as a data structure

accept node

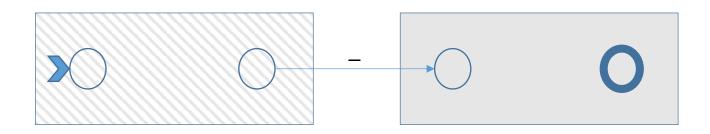
> start node



#### REG for $(R_1).(R_2)$



If we have two expressions  $R_1$  and  $R_2$ , we can construct the REG for  $(R_1).(R_2)$  from the REGs of  $R_1$  and  $R_2$  as shown below.



REG for  $(R_1).(R_2)$ 

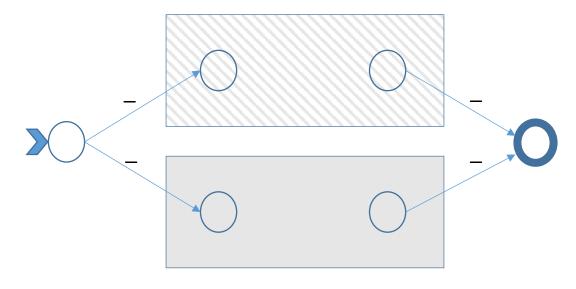
# REG for $(R_1)|(R_2)$



REG for R<sub>1</sub>



REG for R<sub>2</sub>

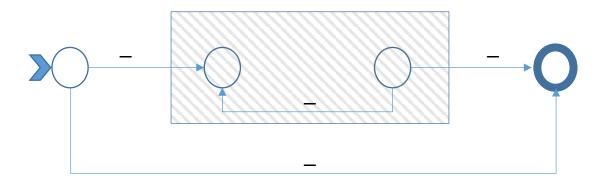


REG for  $(R_1)|(R_2)$ 

## REG for (R)\*



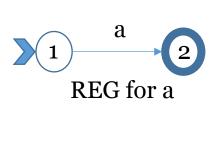
REG for R



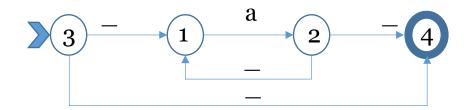
REG for  $(R)^*$ 

# Examples

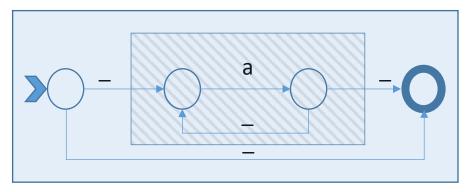
- In what follows we show how the construction works for a couple examples
- I will first show the graph illustration then I will show how the data structures look

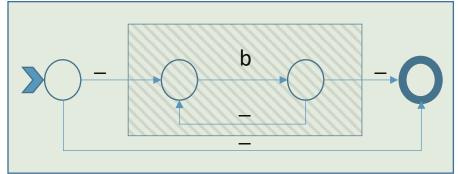






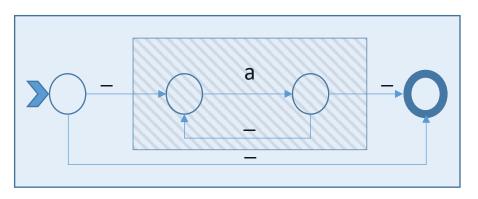
REG for (a)\*

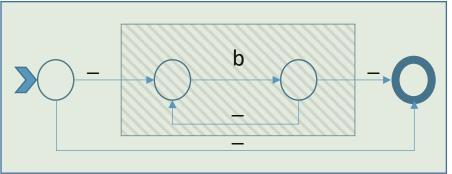




REG for (a)\*

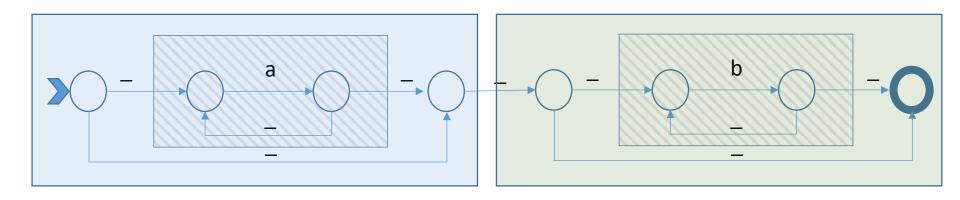
REG for (b)\*



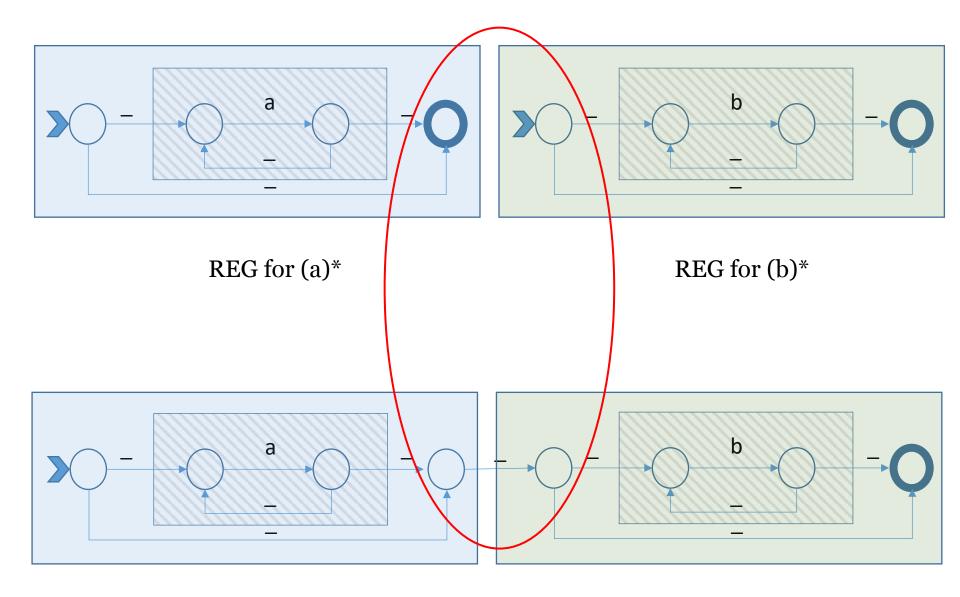


REG for  $(a)^*$ 

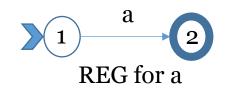
REG for (b)\*

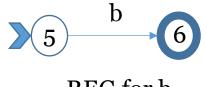


REG for  $((a)^*).((b)^*)$ 

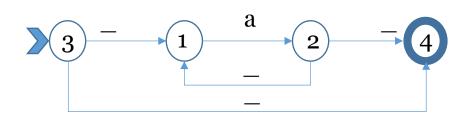


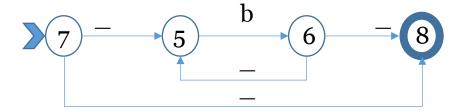
REG for  $((a)^*).((b)^*)$ 





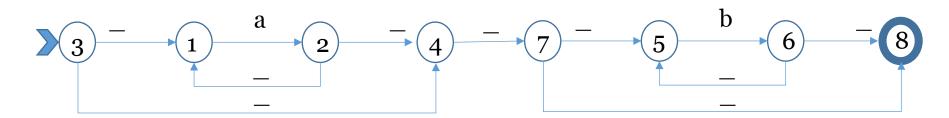
REG for b



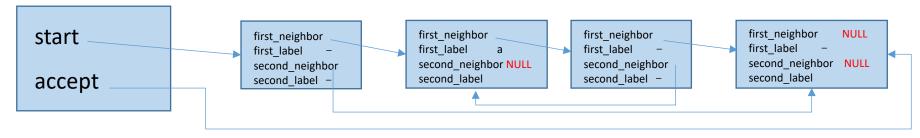


REG for  $(a)^*$ 

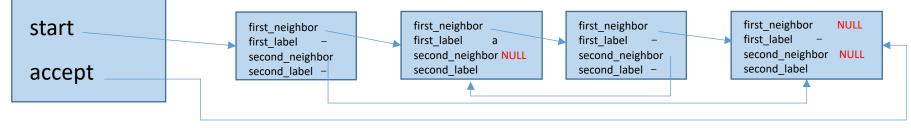
REG for (b)\*



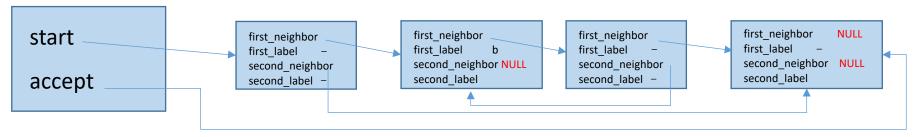
REG for  $((a)^*).((b)^*)$ 



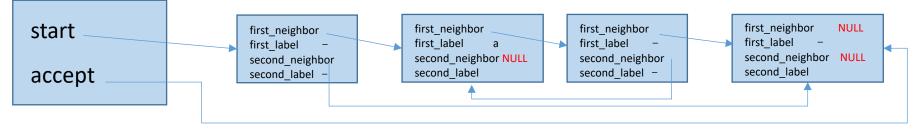
REG Data Structure for (a)\*



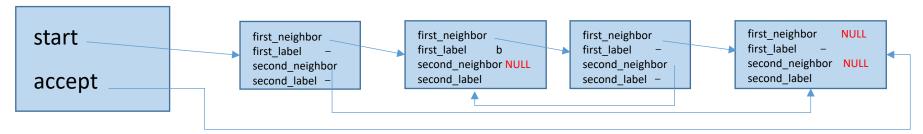
REG Data Structure for (a)\*



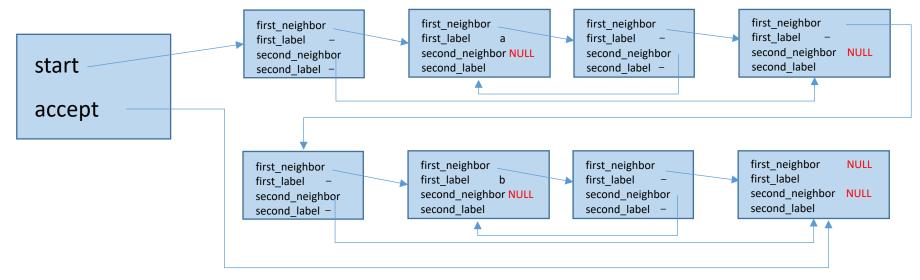
REG Data Structure for (b)\*



REG Data Structure for (a)\*



REG Data Structure for (b)\*



REG Data Structure for  $((a)^*).((b)^*)$ 

# Another example

- In what follows, I will assume that parse\_expr() returns REGs as discussed above
- I will show a step by step execution
   parse\_expr() on the expression ((a)\*).((b)\*)

parse\_expr() Input( ( a ) \* ).( ( b ) \* )

consume LPAREN

( ( a ) \* ).( ( b ) \* )

consume LPAREN
R1 = parse\_expr()

( ( a ) \* ).( ( b ) \* )

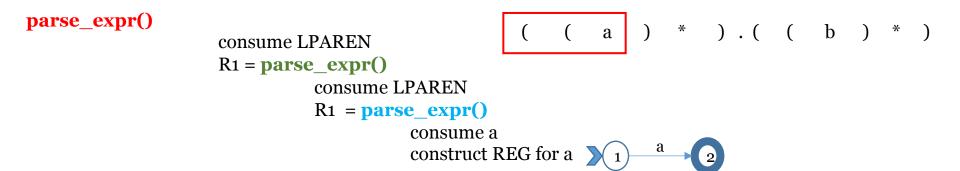
consume LPAREN
R1 = parse\_expr()
consume LPAREN

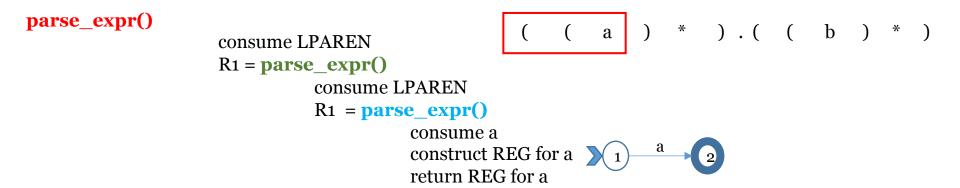
( ( a ) \* ).( ( b ) \* )

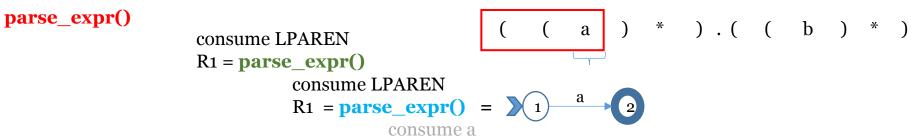
consume LPAREN
R1 = parse\_expr()
consume LPAREN
R1 = parse\_expr()

( a ) \* ).( ( b ) \* )

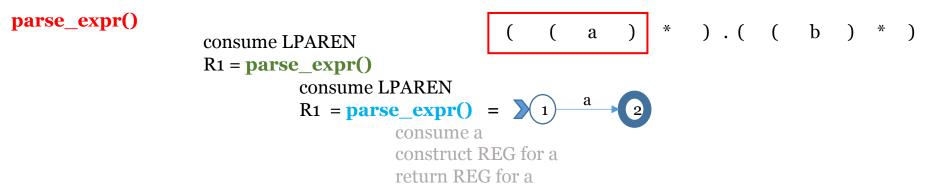
consume a



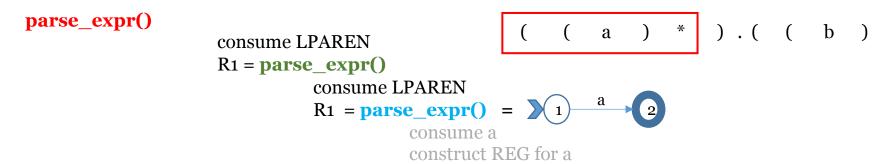




consume a construct REG for a return REG for a

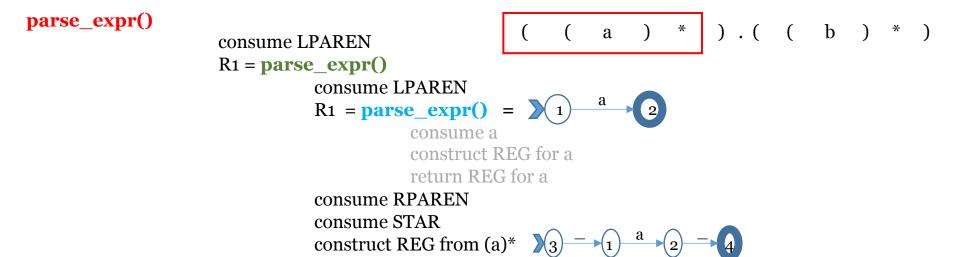


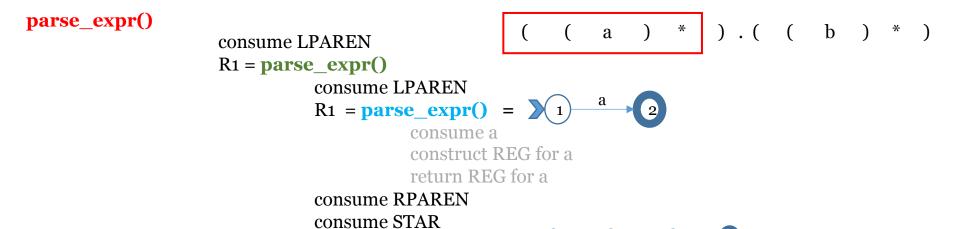
consume RPAREN



return REG for a

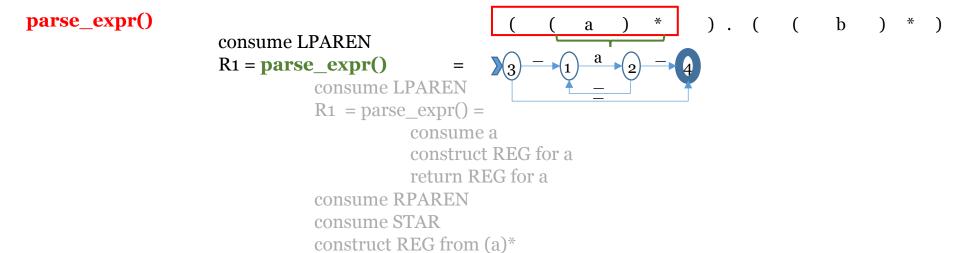
consume RPAREN consume STAR



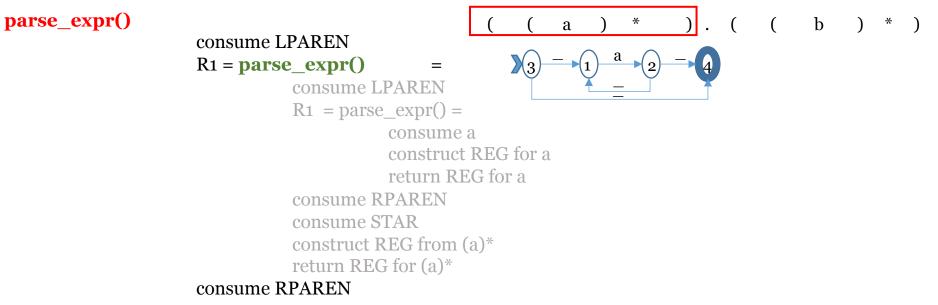


construct REG from (a)\*

return REG for (a)\*



return REG for (a)\*



consume LPAREN
R1 = parse\_expr() =
consume LPAREN
R1 = parse\_expr() =
consume a
construct REG for a
return REG for a
consume STAR
construct REG from (a)\*
return REG for (a)\*

consume RPAREN consume DOT

consume LPAREN
R1 = parse\_expr() =
consume LPAREN
R1 = parse\_expr() =
consume a
construct REG for a
return REG for a
consume STAR
construct REG from (a)\*
return REG for (a)\*
consume RPAREN
consume RPAREN
consume RPAREN
consume RPAREN
consume RPAREN
consume DOT

consume LPAREN

consume LPAREN R1 = parse\_expr() consume LPAREN R1 = parse\_expr() = consume a construct REG for a return REG for a consume RPAREN consume STAR construct REG from (a)\* return REG for (a)\* consume RPAREN consume DOT consume LPAREN R2 = parse\_expr()

consume LPAREN R1 = parse\_expr() consume LPAREN R1 = parse\_expr() = consume a construct REG for a return REG for a consume RPAREN consume STAR construct REG from (a)\* return REG for (a)\* consume RPAREN consume DOT consume LPAREN R2 = parse\_expr() consume LPAREN

```
consume LPAREN
R1 = parse_expr()
         consume LPAREN
         R1 = parse_expr() =
                   consume a
                   construct REG for a
                   return REG for a
         consume RPAREN
          consume STAR
         construct REG from (a)*
         return REG for (a)*
consume RPAREN
consume DOT
consume LPAREN
R2 = parse_expr()
         consume LPAREN
         R1 = parse_expr()
```

```
a
consume LPAREN
R1 = parse_expr()
          consume LPAREN
          R1 = parse_expr() =
                    consume a
                    construct REG for a
                    return REG for a
          consume RPAREN
          consume STAR
          construct REG from (a)*
          return REG for (a)*
consume RPAREN
consume DOT
consume LPAREN
R2 = parse_expr()
          consume LPAREN
          R1 = parse_expr()
```

consume b

consume LPAREN R1 = parse\_expr() consume LPAREN  $R1 = parse_expr() =$ consume a construct REG for a return REG for a consume RPAREN consume STAR construct REG from (a)\* return REG for (a)\* consume RPAREN consume DOT consume LPAREN  $R2 = parse_expr()$ consume LPAREN R1 = parse\_expr() consume b construct REG for b  $\searrow$  (5)

consume LPAREN R1 = parse\_expr() consume LPAREN  $R1 = parse_expr() =$ consume a construct REG for a return REG for a consume RPAREN consume STAR construct REG from (a)\* return REG for (a)\* consume RPAREN consume DOT consume LPAREN  $R2 = parse_expr()$ consume LPAREN  $R1 = parse_expr()$ consume b construct REG for b > (5)

return REG for b

consume LPAREN R1 = parse\_expr() consume LPAREN  $R1 = parse_expr() =$ consume a construct REG for a return REG for a consume RPAREN consume STAR construct REG from (a)\* return REG for (a)\* consume RPAREN consume DOT consume LPAREN  $R2 = parse_expr()$ consume LPAREN  $R1 = parse_expr()$ consume b construct REG for b

return REG for b

b consume LPAREN R1 = parse\_expr() consume LPAREN  $R1 = parse_expr() =$ consume a construct REG for a return REG for a consume RPAREN consume STAR construct REG from (a)\* return REG for (a)\* consume RPAREN consume DOT consume LPAREN  $R2 = parse_expr()$ consume LPAREN  $R1 = parse_expr()$ consume b construct REG for b return REG for b

consume RPAREN

b a consume LPAREN  $R1 = parse_expr()$ consume LPAREN  $R1 = parse_expr() =$ consume a construct REG for a return REG for a consume RPAREN consume STAR construct REG from (a)\* return REG for (a)\* consume RPAREN consume DOT consume LPAREN  $R2 = parse_expr()$ consume LPAREN  $R1 = parse_expr()$ consume b construct REG for b return REG for b consume RPAREN consume STAR

b a consume LPAREN R1 = parse\_expr() consume LPAREN  $R1 = parse\_expr() =$ consume a construct REG for a return REG for a consume RPAREN consume STAR construct REG from (a)\* return REG for (a)\* consume RPAREN consume DOT consume LPAREN R2 = parse\_expr() consume LPAREN  $R1 = parse_expr()$ consume b construct REG for b return REG for b consume RPAREN consume STAR construct REG for (b)\*

b a consume LPAREN R1 = parse\_expr() consume LPAREN  $R1 = parse\_expr() =$ consume a construct REG for a return REG for a consume RPAREN consume STAR construct REG from (a)\* return REG for (a)\* consume RPAREN consume DOT consume LPAREN  $R2 = parse_expr()$ consume LPAREN  $R1 = parse_expr()$ consume b construct REG for b return REG for b consume RPAREN consume STAR construct REG for (b)\* return REG for (b)\*

```
b
                                       a
consume LPAREN
R1 = parse_expr()
          consume LPAREN
          R1 = parse\_expr() =
                    consume a
                    construct REG for a
                    return REG for a
          consume RPAREN
          consume STAR
          construct REG from (a)*
          return REG for (a)*
consume RPAREN
consume DOT
consume LPAREN
R2 = parse_expr()
          consume LPAREN
          R1 = parse_expr()
                    consume b
                    construct REG for b
                    return REG for b
          consume RPAREN
          consume STAR
          construct REG for (b)*
          return REG for (b)*
```

```
b
                                       a
consume LPAREN
R1 = parse_expr()
          consume LPAREN
          R1 = parse\_expr() =
                    consume a
                    construct REG for a
                    return REG for a
          consume RPAREN
          consume STAR
          construct REG from (a)*
          return REG for (a)*
consume RPAREN
consume DOT
consume LPAREN
R2 = parse expr()
          consume LPAREN
          R1 = parse_expr()
                    consume b
                    construct REG for b
                    return REG for b
          consume RPAREN
          consume STAR
          construct REG for (b)*
          return REG for (b)*
```

parse\_expr()

consume RPAREN

parse\_expr() b a consume LPAREN  $R1 = parse_expr()$ consume LPAREN  $R1 = parse\_expr() =$ consume a construct REG for a return REG for a consume RPAREN consume STAR construct REG from (a)\* return REG for (a)\* consume RPAREN consume DOT consume LPAREN R2 = parse expr()consume LPAREN  $R1 = parse_expr()$ consume b construct REG for b return REG for b consume RPAREN consume STAR construct REG for (b)\* return REG for (b)\* consume RPAREN

construct REG for ((a\*).((b)\*)

parse\_expr() b a consume LPAREN R1 = parse\_expr() consume LPAREN  $R1 = parse\_expr() =$ consume a construct REG for a return REG for a consume RPAREN consume STAR construct REG from (a)\* return REG for (a)\* consume RPAREN consume DOT consume LPAREN R2 = parse expr()consume LPAREN R1 = parse\_expr() consume b construct REG for b return REG for b consume RPAREN consume STAR construct REG for (b)\* return REG for (b)\* consume RPAREN

> construct REG for  $((a^*).((b)^*)$ return REG for  $((a)^*).((b)^*)$

### How to match strings?

So far, we have seen how to construct REGs. Next, I will show how to determine if a string belongs to the language of a regular expression

- 1. I will start by defining paths a
- 2. Then I will define accepting path
- 3. Then, I will state the MAIN theorem for REGs which allows us to determine if a string is in the language of a regular expression by looking for an accepting path in the REG of the expression.
- 4. Then I will give examples of paths for a specific regular expression

### How to match strings?

After I am done illustrating path, I will show an equivalent formulation in terms of reachable nodes

- 5. I will start by introducing reachable nodes state the equivalence between finding accepting paths and reachable nodes
- 6. Then I will give the condition for determining if a string in in the language of a regular expression in terms of reachable nodes
- 7. Then I will give examples of how reachable nodes are calculated (the pseudocode is given in the project description)
- 8. I will conclude by summarizing how longest matching prefix is found using reachable nodes

#### 1. Definition of a path

- A *path* in a REG r is a sequence of nodes  $n_1 n_2 ...$   $n_k$  starting from the starting node of r such that consecutive nodes in the sequence are adjacent in the REG:
  - $n_1$  is the starting node of r
  - for each consecutive nodes  $n_i$  and  $n_{i+1}$ , there is an edge from  $n_i$  to  $n_{i+1}$  in r

• The *string of a path* is the string obtained by concatenating all labels on the edges of the path

#### 2. Definition of accepting path

- An *accepting path* in a REG r is a sequence of nodes  $n_1 n_2 ... n_k$  starting from the starting node of r and ending in the accepting node of r such that consecutive nodes in the sequence are adjacent in the REG:
  - $n_1$  is the starting node of r
  - $n_k$  is the accepting node of r
  - for each consecutive nodes  $n_i$  and  $n_{i+1}$ , there is an edge from  $n_i$  to  $n_{i+1}$  in r

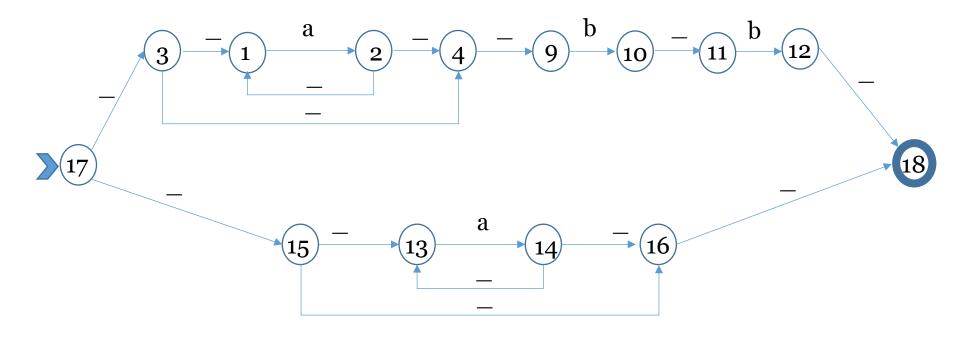
• The *string of an accepting path* is the string obtained by concatenating all labels on the edges of the accepting path

### 3. MAIN Theorem for REGs

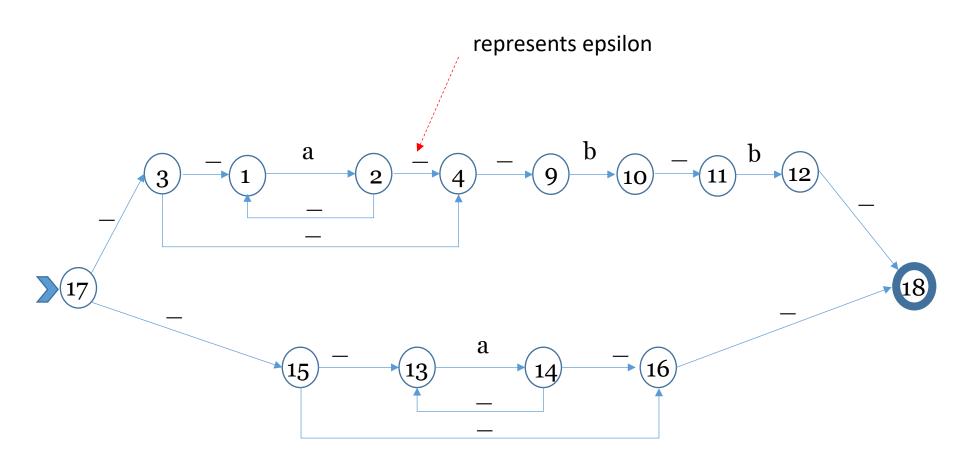
• Let *R* be a regular expression and *r* be its REG

• **Theorem**: A string w is in L(R) if and only if there is an accepting path whose string is w

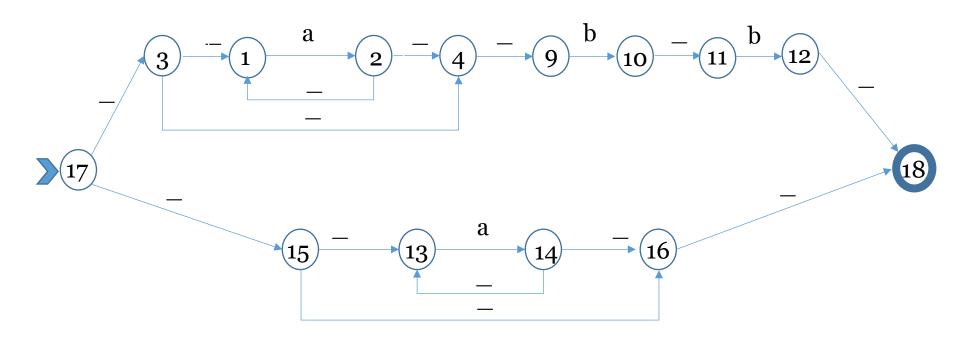
#### 4. Example of path



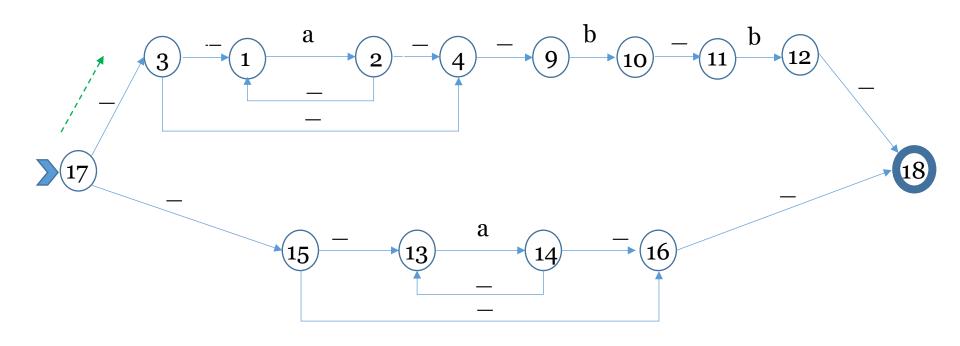
#### 4. Example of path



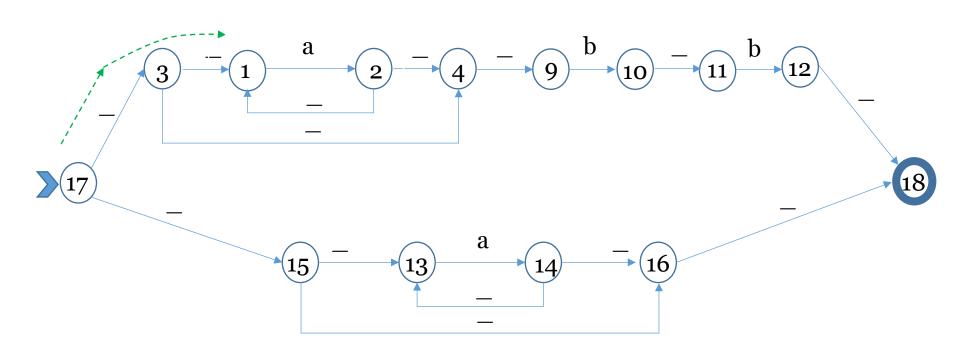
aabb =



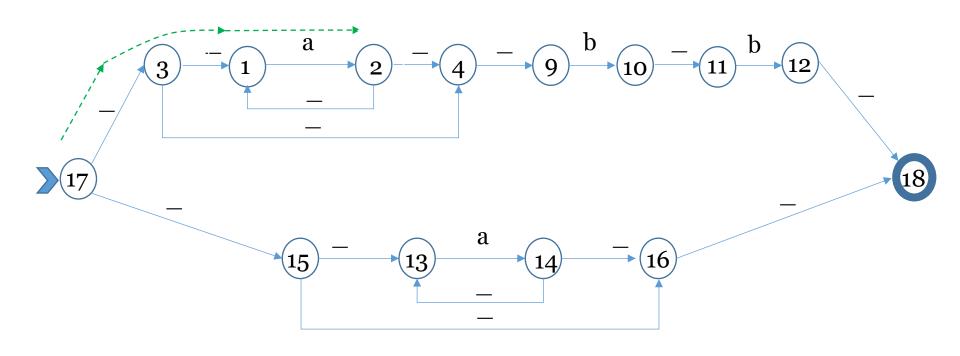
 $aabb = _{-}$ 



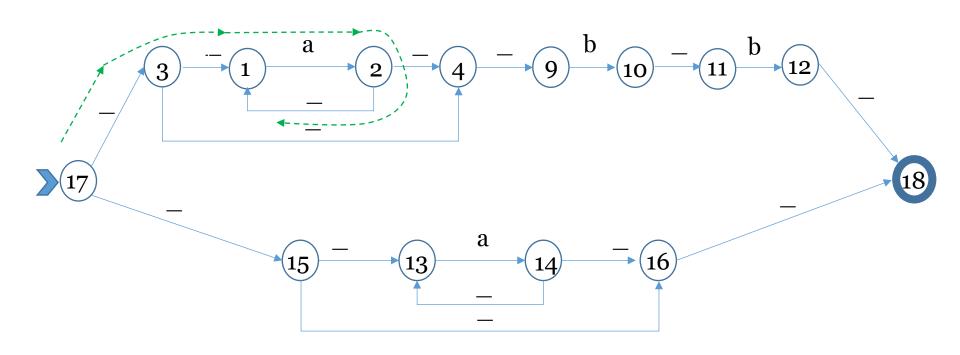
aabb = \_ \_



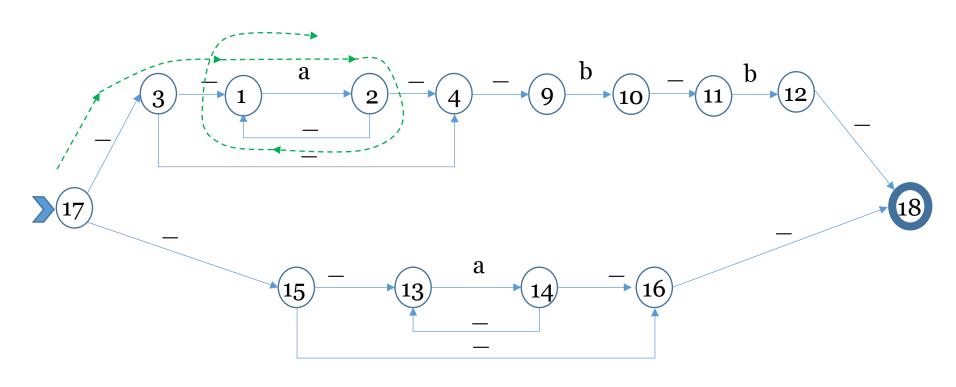
 $aabb = \underline{\phantom{a}} a$ 



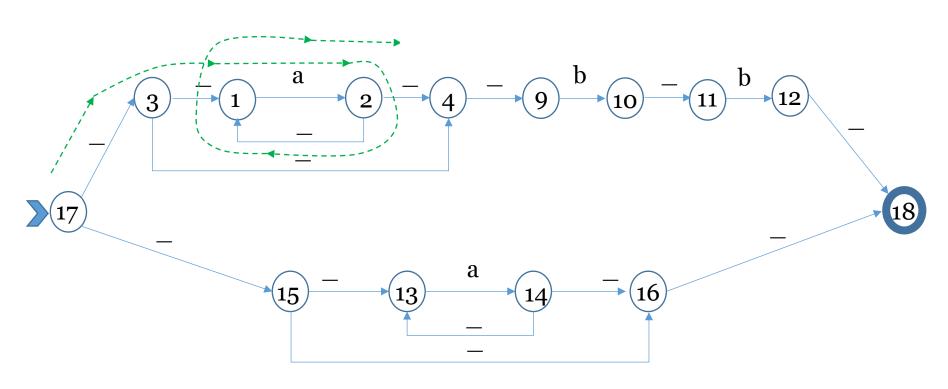
aabb = \_ \_ a \_



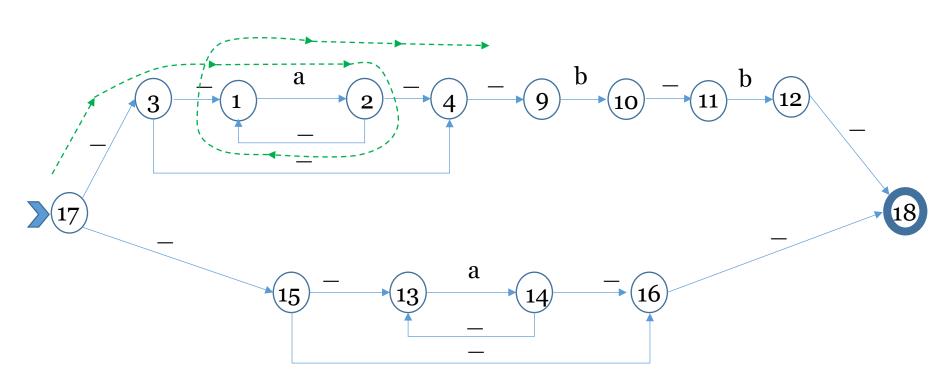
aabb = \_ \_ a \_ a



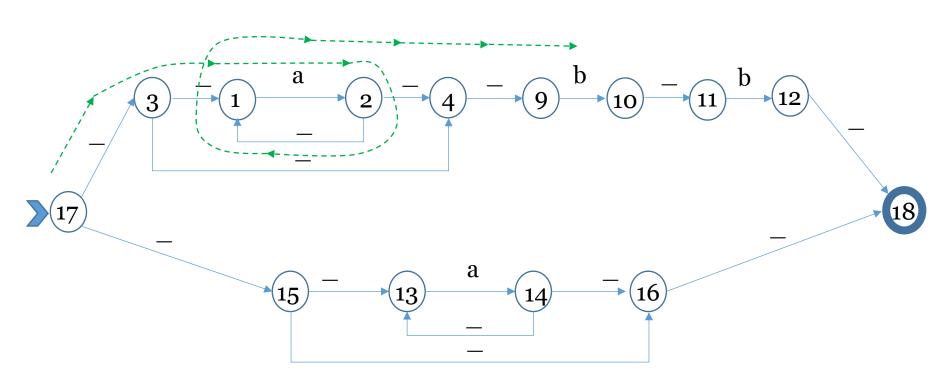
aabb = \_ \_ a \_ a \_



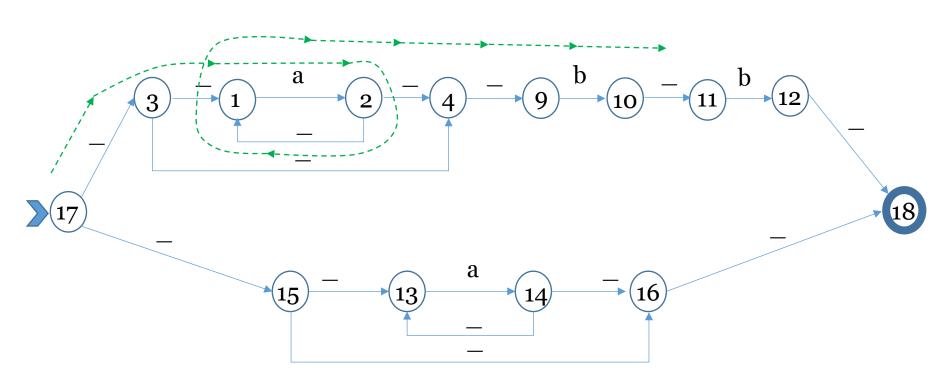
aabb = \_ \_ a \_ a \_ \_ \_



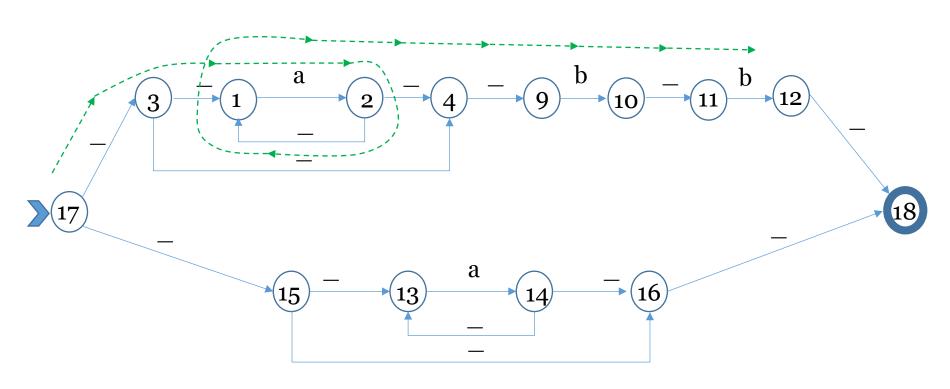
aabb = \_ \_ a \_ a \_ \_ b



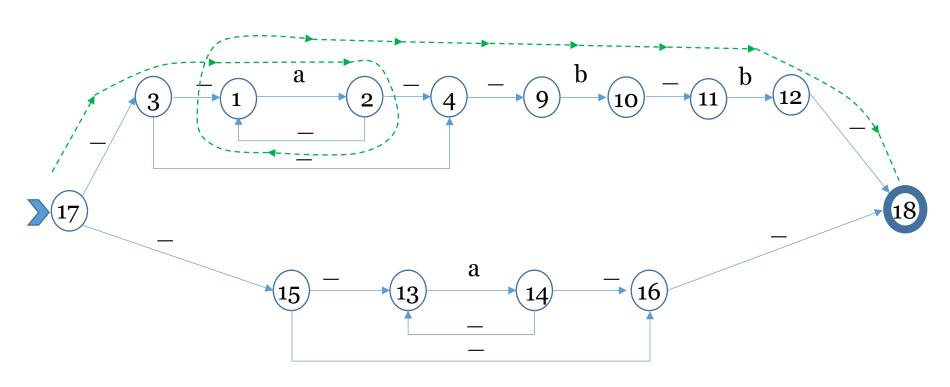
aabb = \_ \_ a \_ a \_ \_ b \_

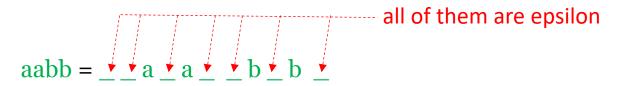


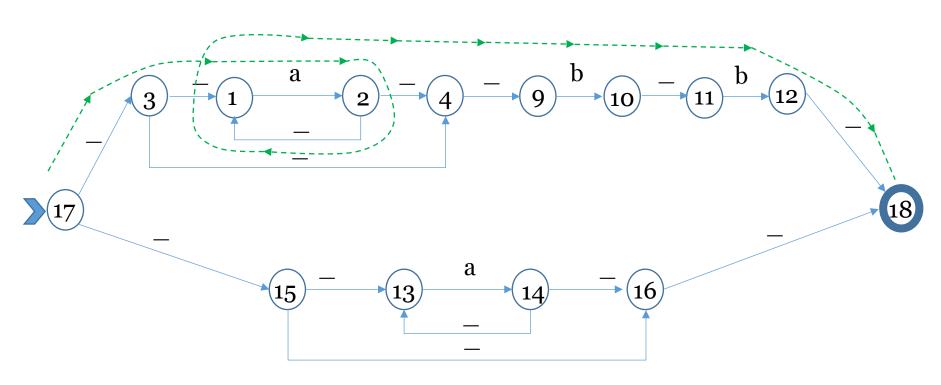
aabb = \_ \_ a \_ a \_ \_ b \_ b



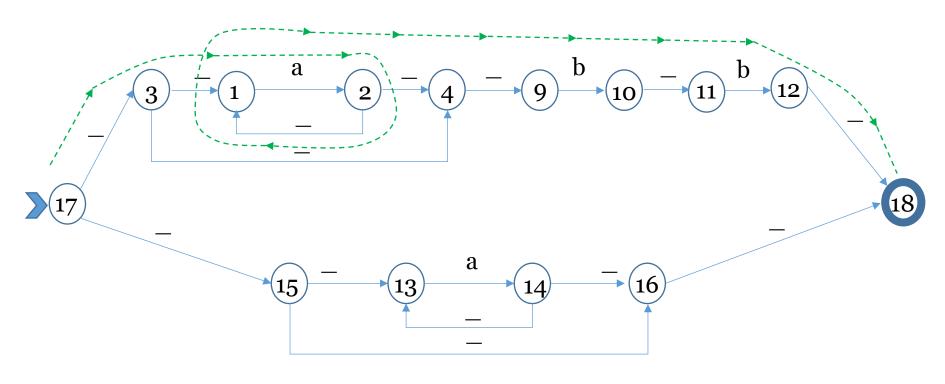
aabb = \_ \_ a \_ a \_ \_ b \_ b \_

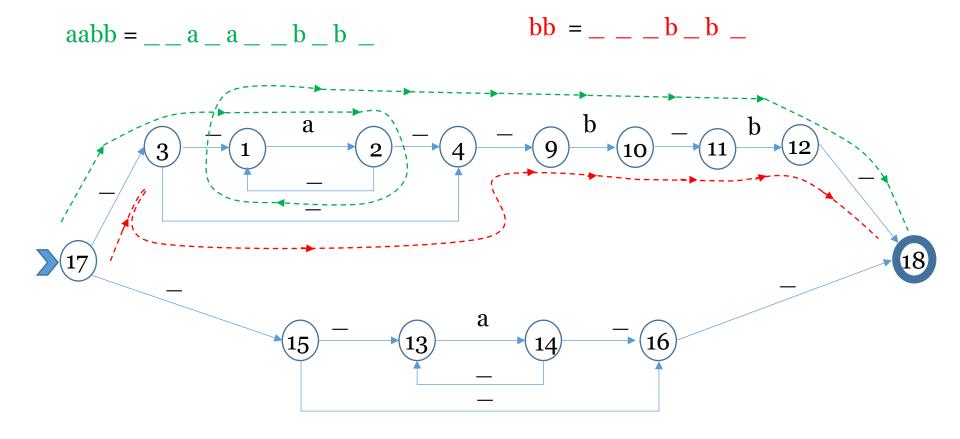












aaa = \_ \_ a \_ a \_ a \_ \_

REG for 
$$(((a)^*).((b).(b))) | ((a)^*)$$

#### 5. Definition of reachable nodes

• For a path  $n_1 n_2 ... n_k$  whose string is w, we say that  $n_k$  is **reachable by consuming** w

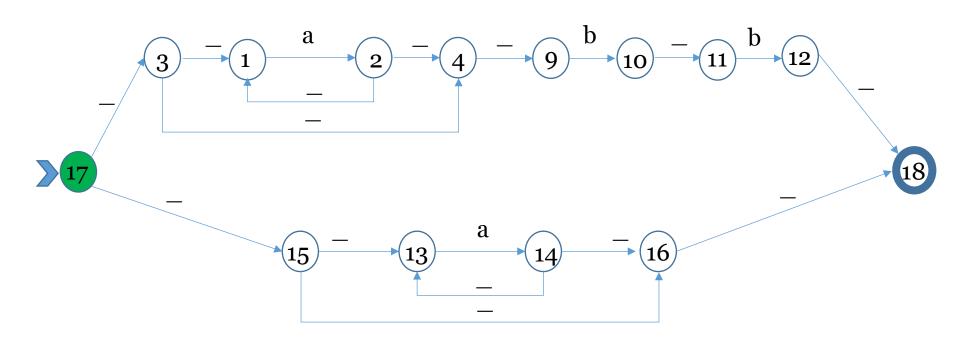
• If w is the string of an accepting path, then the accepting node is reachable by consuming w.

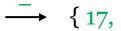
# 6. Restating the MAIN Theorem for REGs

• Let *R* be a regular expression and *r* be its REG

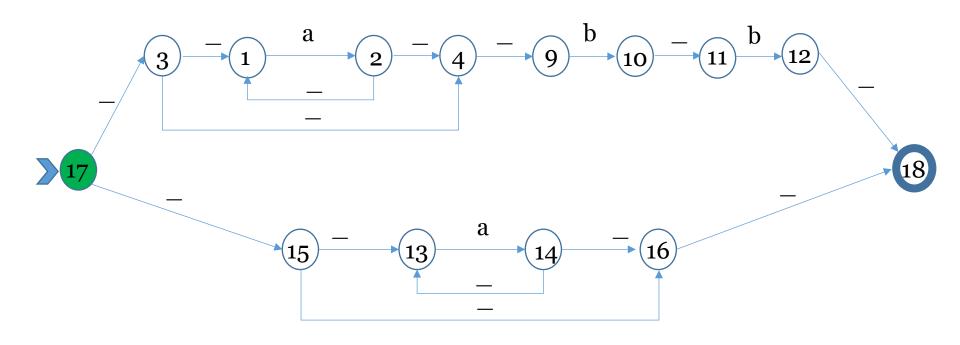
• **Theorem**: A string w is in L(R) if and only if the accepting node of r, the REG of R, is reachable by consuming w

Input  $aba = ab_a$ 



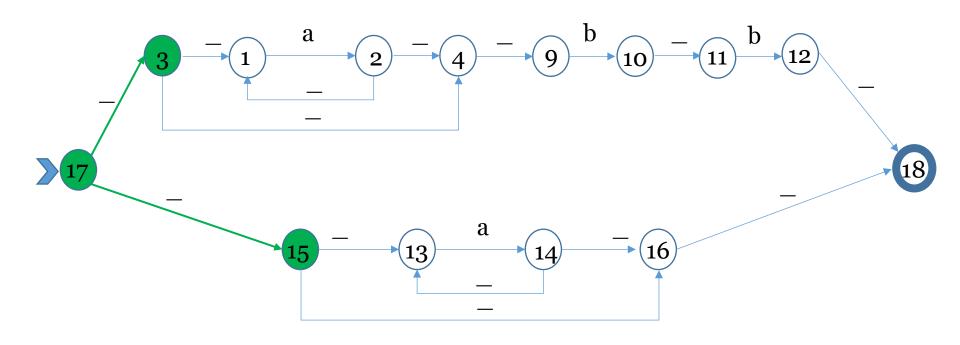


Initially state 17 is reachable by consuming nothing. It is the initial state



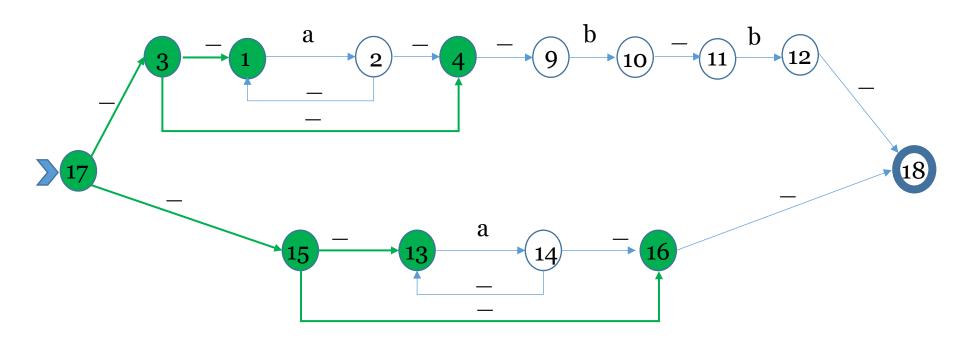
 $\stackrel{-}{\longrightarrow}$  { 17, 3, 15

From state 17 is we can go to states 3 and 15 by consuming nothing

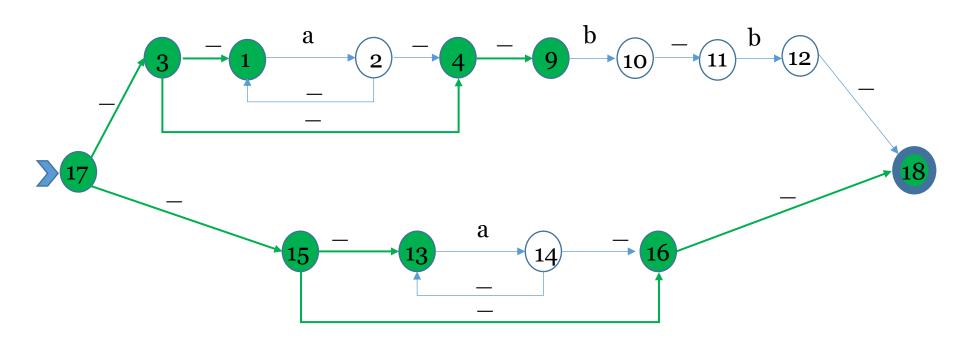


 $\xrightarrow{-}$  { 17, 3, 15, 1, 4, 13, 16

we can go from state 17 to states 3, 15 1, 4, 13 and 16 by consuming nothing

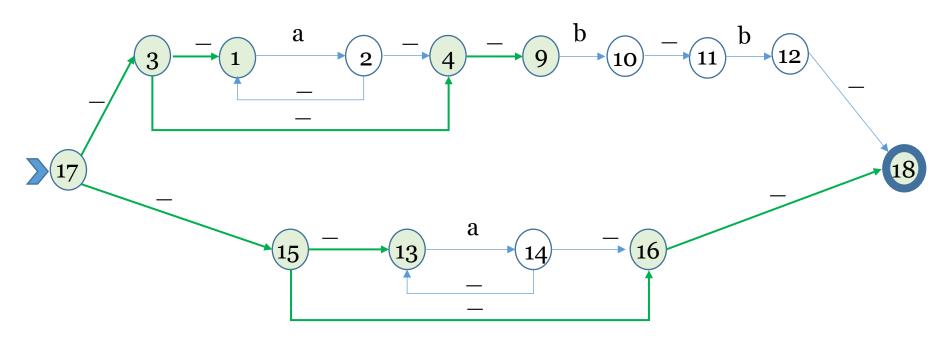


```
\stackrel{-}{\longrightarrow} { 17, 3, 15, 1, 4, 13, 16, 9, 18 }
```



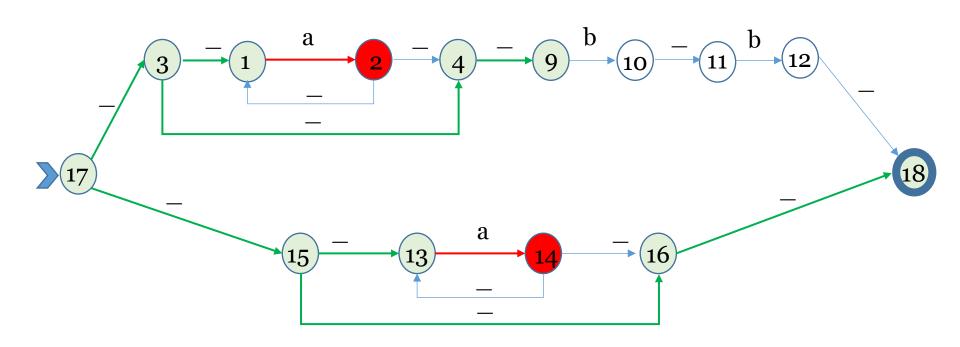
$$\stackrel{-}{\longrightarrow}$$
 { 17, 3, 15, 1, 4, 13, 16, 9, **18** }

The set represented by this regular expression contains epsilon because accept state 18 is reachable from state 17 without consuming any input

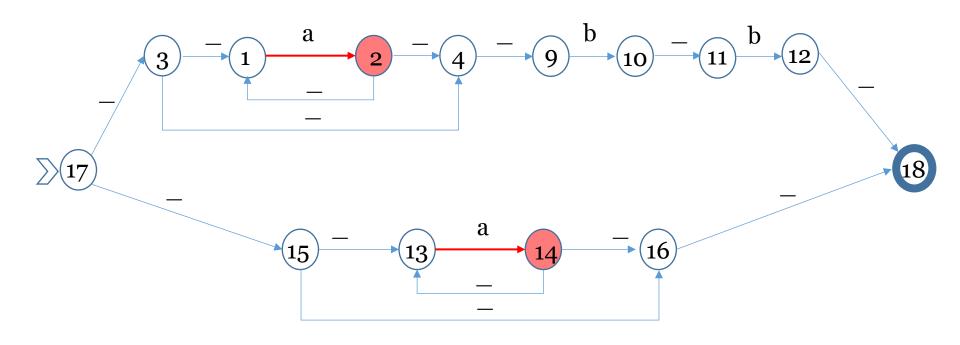


$$\xrightarrow{-}$$
 { 17, 3, 15, 1, 4, 13, 16, 9, **18** }  $\xrightarrow{a}$  { **2**, **14** }

from the set of nodes that we obtained, we can go to the set { 2, 14 } by consuming a

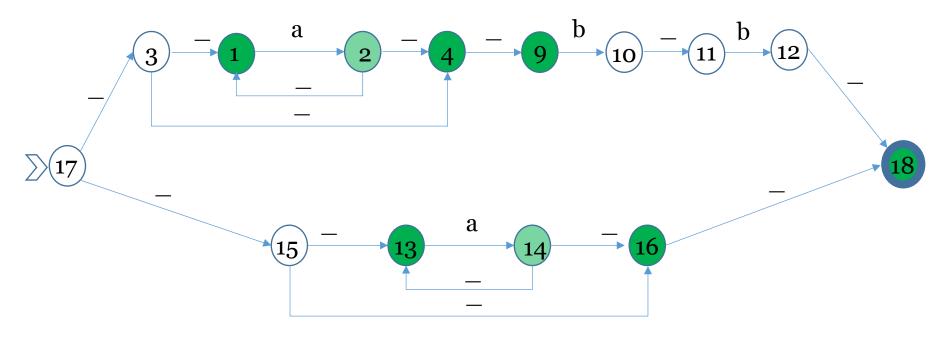


$$\xrightarrow{-}$$
 { 17, 3, 15, 1, 4, 13, 16, 9, **18** }  $\xrightarrow{a}$  { **2**, 14 }



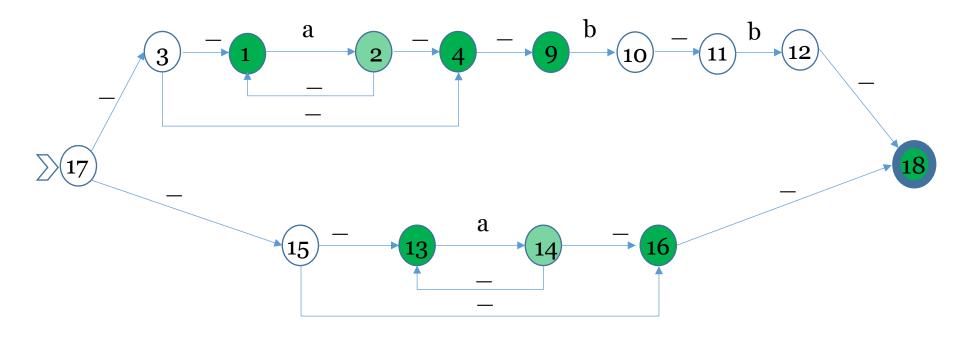
$$\xrightarrow{-} \{17, 3, 15, 1, 4, 13, 16, 9, 18\} \xrightarrow{a} \{2, 14\} \xrightarrow{-} \{1, 2, 4, 9, 13, 14, 16, 18\}$$

Then we can go to the set of nodes { 1, 2, 4, 9, 13, 14, 16, 18} by consuming nothing



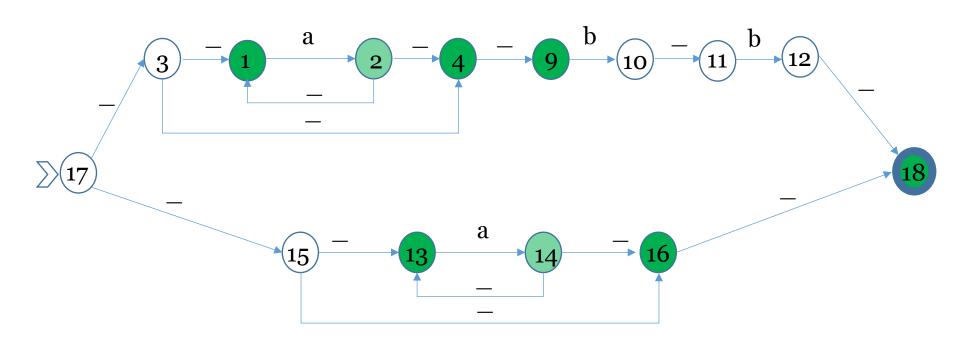
$$\xrightarrow{-} \{17, 3, 15, 1, 4, 13, 16, 9, 18\} \xrightarrow{a} \{2, 14\} \xrightarrow{-} \{1, 2, 4, 9, 13, 14, 16, 18\}$$

So, effectively, from the initial node 17, we can go to the set  $\{1, 2, 4, 9, 13, 14, 16, 18\}$  by consuming a. For the rest of the example, I will not add commentary



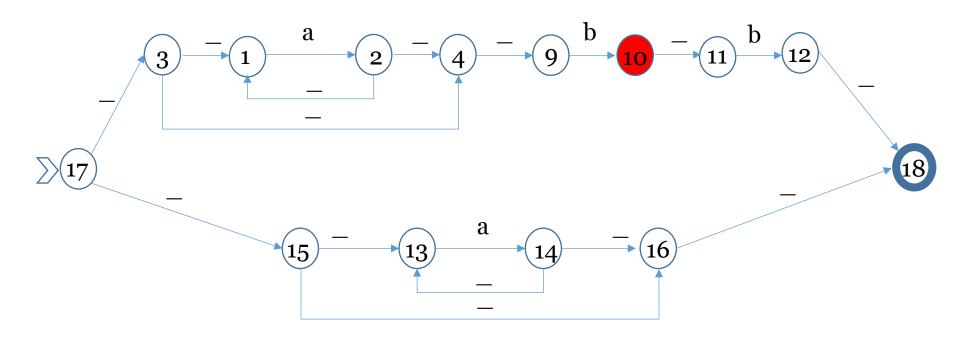
$$\xrightarrow{-} \{ 17, 3, 15, 1, 4, 13, 16, 9, 18 \} \xrightarrow{a} \{ 2, 14 \} \xrightarrow{-} \{ 1, 2, 4, 9, 13, 14, 16, 18 \}$$

$$\xrightarrow{b} \{$$



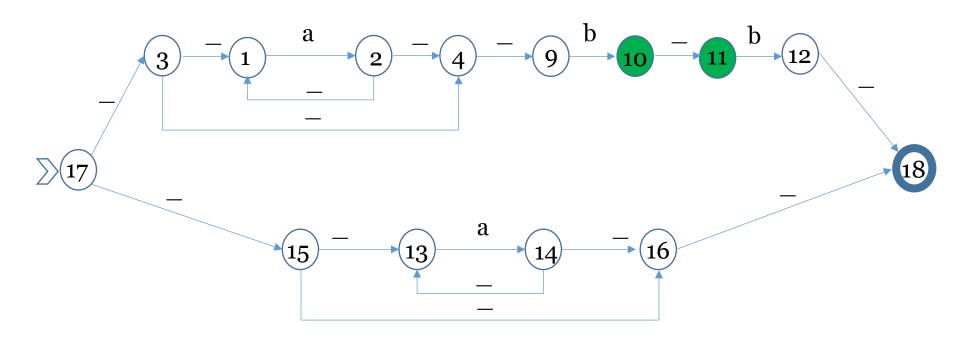
$$\xrightarrow{-} \{17, 3, 15, 1, 4, 13, 16, 9, 18\} \xrightarrow{a} \{2, 14\} \xrightarrow{-} \{1, 2, 4, 9, 13, 14, 16, 18\}$$

$$\xrightarrow{b} \{10\}$$



$$\xrightarrow{-} \{ 17, 3, 15, 1, 4, 13, 16, 9, 18 \} \xrightarrow{a} \{ 2, 14 \} \xrightarrow{-} \{ 1, 2, 4, 9, 13, 14, 16, 18 \}$$

$$\xrightarrow{b} \{ 10 \} \xrightarrow{-} \{ 10, 11 \}$$
 So far, ab is viable but not matching



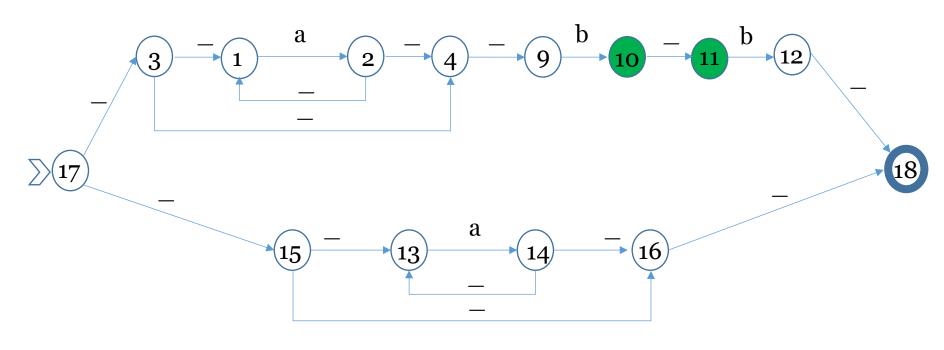
$$\xrightarrow{-} \{17, 3, 15, 1, 4, 13, 16, 9, 18\} \xrightarrow{a} \{2, 14\} \xrightarrow{-} \{1, 2, 4, 9, 13, 14, 16, 18\}$$

$$\stackrel{b}{\longrightarrow} \{10\} \stackrel{-}{\longrightarrow} \{10,11\}$$

So far, ab is viable but not matching.

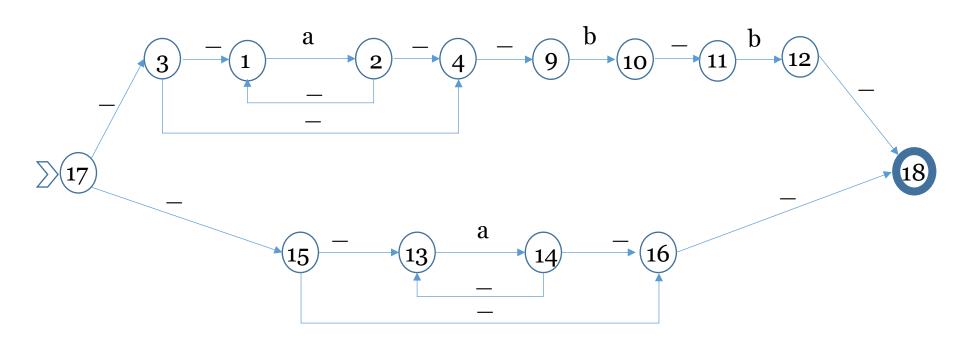
**Not matching means** that the set of nodes that can be reached by consuming ab does not contain the accept node

**viable** means that the set that can be reached by consuming ab is not empty, so there is still hope to reach the accepting state by consuming characters yet to come.



$$\xrightarrow{-} \{17, 3, 15, 1, 4, 13, 16, 9, \mathbf{18}\} \xrightarrow{a} \{2, \mathbf{14}\} \xrightarrow{-} \{1, 2, 4, 9, 13, 14, 16, \mathbf{18}\}$$

b {10} - {10,11} a { aba is not viable nor matching because no nodes can be reached by consuming aba

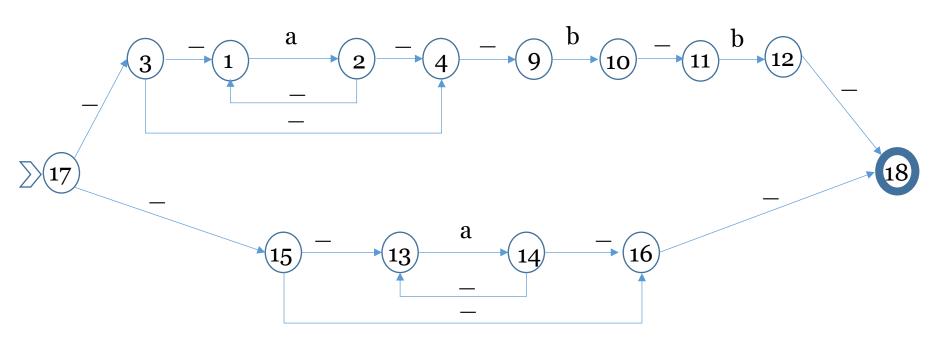


REG for 
$$(((a)^*).((b).(b))) | ((a)^*)$$

$$\xrightarrow{-} \{17, 3, 15, 1, 4, 13, 16, 9, 18\} \xrightarrow{a} \{2, 14\} \xrightarrow{-} \{1, 2, 4, 9, 13, 14, 16, 18\}$$

$$b \in \{10\}$$
  $\rightarrow \{10, 11\}$   $a \in \{10, 11\}$ 

after the "a" and returns ERROR



REG for 
$$(((a)^*).((b).(b))) | ((a)^*)$$

# 8. How is restated main theorem relevant?

#### Recall our task

- Given:
  - a REG r for regular expression R,
  - a string s, and
  - a position p in the string,

- Find:
  - longest possible substring *w* starting at *p* such that *w* is in L(*R*)

#### How is main theorem relevant?

#### • Given:

- a REG r for regular expression R,
- a string s, and
- a position p in the string,

#### • Find:

longest possible substring w starting at p such that the accepting state of r is reachable by consuming w