

8-30-2021

CSE 340

Last time

- Syntax analysis
- getToken() function and its behavior
- difference between syntax and semantics

Today

- getToken()
- peek()
- recursive descent parsing

getToken() and peek()

Assume that the input is already broken down into a sequence of tokens that are stored in an array



We also have an index that initially points to the first entry (index = 1 initially)

If there are no tokens, the array is empty and $\#tokens = 0$ but index = 1 initially.

- 1 2 3 ... : if index > $\#tokens$

gettoken() : if index > # tokens
return EOF

else tok = token-array[index]
index = index + 1 // consume the token
return tok;

peek(int howfar): if index + howfar - 1 > # tokens
return EOF
else return

how far to
peek. > 0

token-array[index + howfar - 1]

// does not consume

// tokens and does

// not change index

Example

token list

IF = { "if" }

ID

NUM

input : cf 1if 1 if 1 iff 123

space
which is a delimiter

1	2	3	4
ID, "cf 1if 1"	IF, ""	ID "iff"	NUM "123"

Initially index = 1

$\text{gettoken}() \rightarrow \text{ID}, "iff" \quad \text{index} = 2$
 $\text{peek}(1) \rightarrow \text{IF}, ""$
 $\text{peek}(2) \rightarrow \text{ID}, "iff"$
 $\text{peek}(4) \rightarrow \text{EOF}$
 $\text{peek}(1) \rightarrow \text{IF}, ""$
 $\text{gettoken}() \rightarrow \text{IF}, ""$
 $\text{peek}(1) \rightarrow \text{ID}, "iff"$

Parsing by example

Grammar for expressions

$\text{expr} \rightarrow \text{term PLUS expr}$

$\text{expr} \rightarrow \text{term}$

$\text{term} \rightarrow \text{factor MULT term}$

$\text{term} \rightarrow \text{factor}$

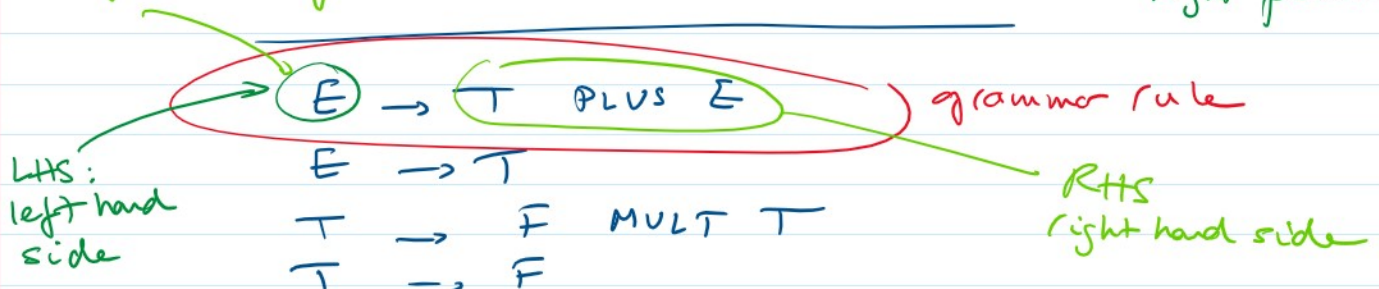
$\text{factor} \rightarrow \text{NUM}$

$\text{factor} \rightarrow \text{ID}$

$\text{factor} \rightarrow \text{LPAREN } E \text{ RPAREN}$

top level symbol

left and right parentheses



$F \rightarrow \text{NUM} \mid \text{ID} \mid \text{LPAREN } E \text{ RPAREN}$
 OR \nearrow Example (1+2)

The LHS's of rules are non-terminals
 The RHS's of rules are sequence of non-terminals and terminals (tokens)

$E \rightarrow T \text{ PLUS } E$
 $\uparrow \quad \uparrow \quad \uparrow \quad \uparrow$
 NT NT Terminal NT
 (token)
 Rule

Example 1 + 2 * 3

$E \rightarrow T + E$
 $E \rightarrow T$
 $T \rightarrow F * T$
 $T \rightarrow F$
 $F \rightarrow \text{ID} \mid \text{NUM} \mid (E)$

NUM PLUS NUM MULT NUM
 F | F MULT T
 T PLUS E
 E

Recursive descent parsing:

- One parsing function for each non-terminal
- When `parse-x()` is called at a given point in the input:
 - an `x` will be "consumed"
 - no more no less OR
- a syntax error is raised

- to "parse" terminals, we call `gettoken()`

parse_input() // input consists of one expr

{ $1 + 2 \times 3$ before

parse- $\in()$; 1 + 2 * 3 \uparrow after
consumed

```
t = gettoken();
```

```
if (t.token-type != EOF)
    syntax_error();
```

$$\equiv \text{expect}(\text{EOF})$$
$$\left. \begin{array}{l} \text{ } \\ \text{ } \\ \text{ } \end{array} \right\}$$

Parse- $F()$

§ // $F \rightarrow NVM$

$$\parallel f \rightarrow \pm D$$
$$\parallel f \rightarrow (\in)$$
$$t = \text{peek}(L);$$

if (t.token-type == NVM)

expect (num); // $E \rightarrow \text{Num}$

else if (t.token-type == ID)

$\text{expect}(\perp D); \quad // \quad f \rightarrow \perp D$

else if (t.token-type == LPAREN)

expect (LPAREN): // (

parse - E(); // E

```

    expect (LBRACE); // (
    parse_E(); // E
    expect (RPAREN); // )
}
else
    syntax_error();
}

```

```

parse_T()
{
    //  $T \rightarrow F \text{ MULT } T$ 
    //  $T \rightarrow F$ 

    parse_F(); // F

    t = peek(t);
    if (t.token_type == MULT)
    {
        expect (MULT); // MULT
        parse_T(); // T
    } else if ((t.token_type == EOF) |
               (t.token_type == RPAREN) |
               (t.token_type == PLUS))
        return;
    else
        syntax_error();
}

```