

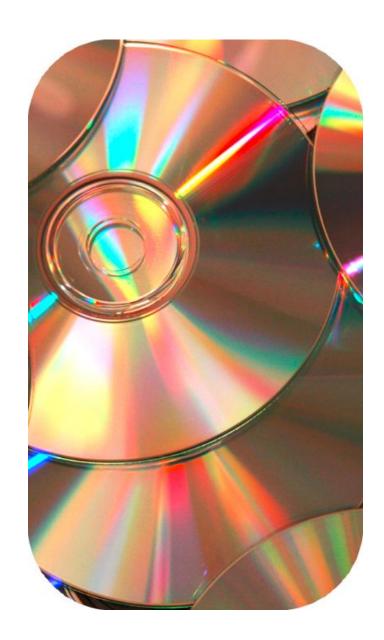
Hadoop Distributed File System (HDFS)





Agenda

- Overview
- Architecture
 - NameNode
 - DataNode
- Blocks and Replication
- File System Shell
- Web Console



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Two Key Aspects of Hadoop

HDFS ted

- Distributed
- Reliable
- Commodity gear





MapReduce

- Parallel Programming
- Fault Tolerant



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Hadoop Distributed File System (HDFS)

- Distributed, scalable, fault tolerant, high throughput
- Data access through MapReduce
- Files split into blocks
- 3 replicas for each piece of data by default
- Can create, delete, copy, but NOT update
- Designed for streaming reads, not random access
- Data locality: processing data on or near the physical storage to decrease transmission of data



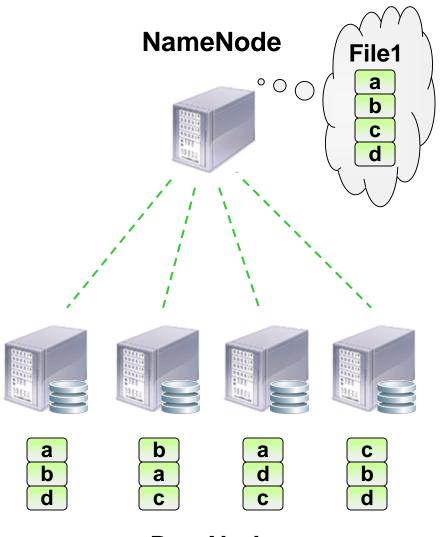




HDFS – Architecture

Master / Slave architecture

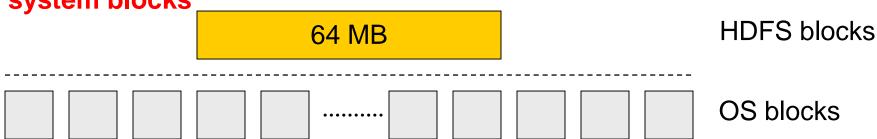
- Master: NameNode
 - manages the file system namespace and metadata
 - FsImage
 - EditLog
 - regulates client access to files
- Slave: DataNode
 - many per cluster
 - manages storage attached to the nodes
 - periodically reports status to NameNode



DataNodes

HDFS - Blocks

- HDFS is designed to support very large files
- Each file is split into blocks
 - Hadoop default: 64MB
 - BigInsights default: 128MB
- Blocks reside on different physical DataNode
- Behind the scenes, 1 HDFS block is supported by multiple operating system blocks



• If a file or a chunk of the file is smaller than the block size, only needed space is used. E.g.: a 210MB file is split as

64 MB 64 MB	64 MB	18 MB
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HDFS – Replication

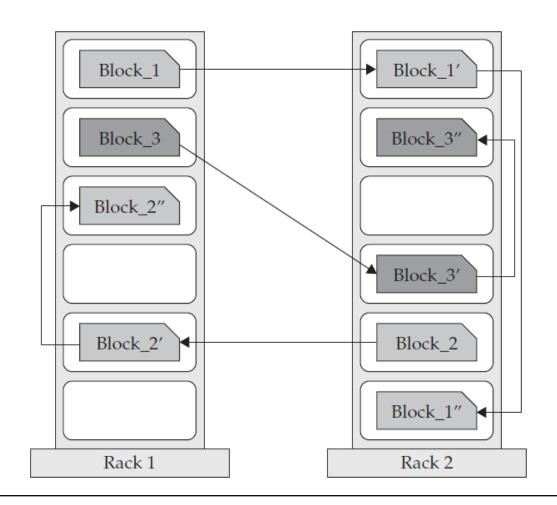
Blocks of data are replicated to multiple nodes

- Behavior is controlled by replication factor, configurable per file
- Default is 3 replicas

Common case:

- one replica on one node in the local rack
- another replica on a different node in the local rack
- and the last on a different node in a different rack

This cuts inter-rack network bandwidth, which improves write performance



Setting Rack Topology (Rack Awareness)

- Can be defined by script which specifies which node is on which rack.
- Script is referenced in topology.script.property.file in core-site.xml.
 - Example of property:

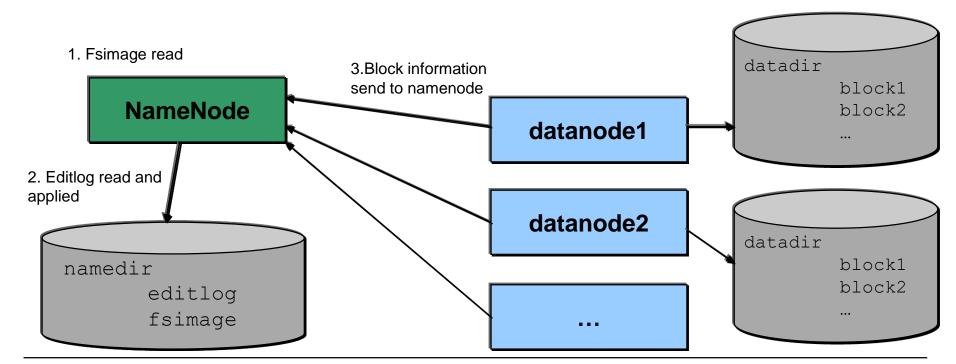
```
<name>topology.script.file.name</name>
    <value>/opt/ibm/biginsights/hadoop-conf/rack-aware.sh</value>
```

■ The network topology script (topology.script.file.name in the above example) receives as arguments one or more IP addresses of nodes in the cluster. It returns on stdout a list of rack names, one for each input. The input and output order must be consistent.



Namenode Startup

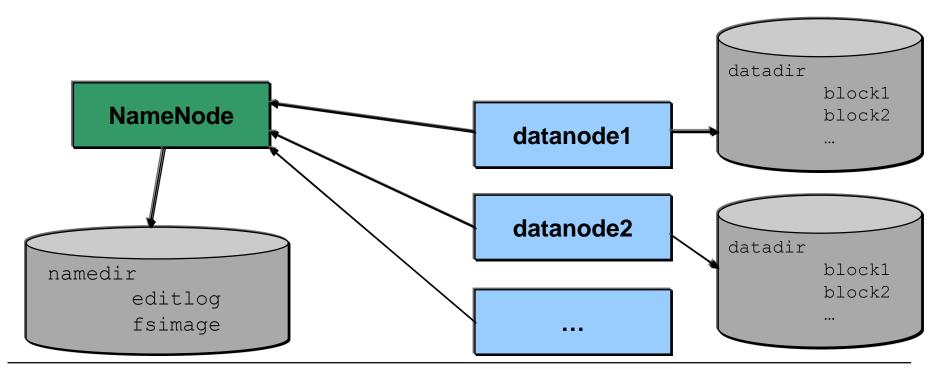
- 1. NameNode reads fsimage in memory
- 2. NameNode applies editlog changes
- 3. NameNode waits for block data from data nodes
 - Namenode doesn't store block information
 - Namenode exits safemode when 99.9% of blocks have at least one copy accounted for





Adding file

- 1. File is added to NameNode memory and persisted in editlog
- 2. Data is written in blocks to datanodes
 - Datanode starts chained copy to two other datanodes
 - If at least one write for each block succeeds, write is successful





Managing Cluster

Adding Data Node

- Start new datanode (pointing to namenode)
- If required run balancer (hadoop balancer) to rebalance blocks

Remove Node

- Simply remove datanode
- Better: Add node to exclude file and wait till all blocks have been moved
- Can be checked in server admin console server:50070

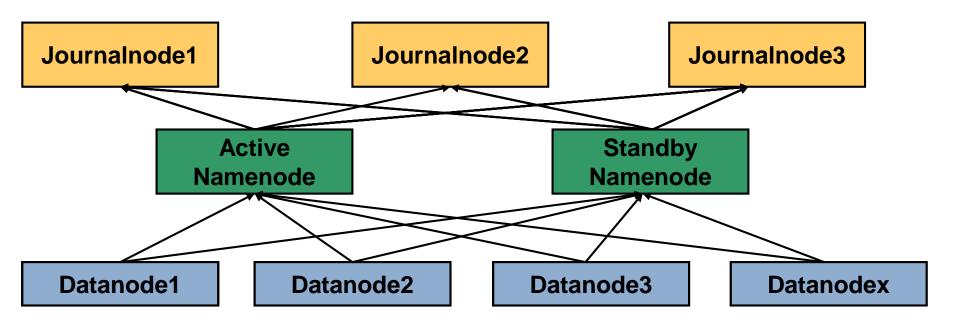
Checking filesystem health

- **Use** hadoop fsck



HDFS-2 Namenode HA

- HDFS-2 adds Namenode High Availability
- Standby Namenode needs filesystem transactions and block locations for fast failover
- Every filesystem modification is logged to at least 3 quorum journal nodes by active Namenode
 - Standby Node applies changes from journal nodes as they occur
 - Majority of journal nodes define reality
 - Split Brain is avoided by Journalnodes (They will only allow one Namenode to write to them)
- Datanodes send block locations and heartbeats to both Namenodes
- Memory state of Standby Namenode is very close to Active Namenode
 - → Much faster failover than cold start





Federated Namenode (HDFS2)

New in Hadoop2 Namenodes can be federated

- Historically Namenodes would become a bottleneck on huge clusters
- One million blocks or ~100TB of data require roughly one GB of RAM in Namenode

Blockpools

- Administrator can create separate blockpools/namespaces with different namenodes
- Datanodes register on all Namenodes
- Datanodes store data of all blockpools (otherwise you could setup separate clusters)
- New ClusterID identifies all namenodes in a cluster.
- A Namespace and its block pool together are called Namespace Volume
- You define which blockpool to use by connecting to a specific Namenode
- Each Namenode still has its own separate backup/secondary/checkpoint node

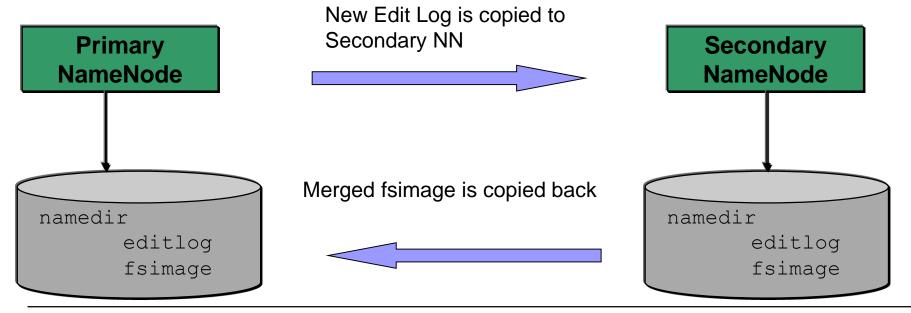
Benefits

- One Namenode failure will not impact other Blockpools
- Better scalability for large numbers of file operations



Secondary NameNode

- During operation primary Namenode cannot merge fsImage and editlog
- This is done on the secondary namenode
 - Every couple minutes, secondary namenode copies new edit log from primary NN
 - Merges editLog into fsimage
 - Copies the new merged fsImage back to primary namenode
- Not HA but faster startup time
 - Secondary NN does not have complete image. In-flight transactions would be lost
 - Primary Namenode needs to merge less during startup
- Was temporarily deprecated because of Namenode HA but has some advantages
 - (no need for Quorum nodes, less network traffic, less moving parts)





Possible FileSystem Setup

GPFS

- No single point of failure
- Posix compliance
- Advanced features like cold storage, backup and restore

Hadoop 2 with HA

- No single point of failure
- Wide community support

Hadoop 2 without HA (or Hadoop 1.x in older versions)

- Copy namedir to NFS (RAID)
- Have virtual IP for backup namenode
- Still some failover time to read blocks, no instant failover but less overhead

fs – file system shell

- File System Shell (fs)
 - Invoked as follows:

hadoop fs <args>

- Example:
 - Listing the current directory in hdfs

hadoop fs -ls.

fs – file system shell

- FS shell commands take URIs as argument
 - URI format: scheme://authority/path
- Scheme:
 - For the local filesystem, the scheme is file
 - For HDFS, the scheme is hdfs
- Authority is the hostname and port of the NameNode

```
hadoop fs -copyFromLocal

file:///myfile.txt

hdfs://localhost:9000/user/keith/myfile.txt
```

- Scheme and authority are optional
 - Defaults are taken from configuration file core-site.xml



fs – file system shell

- Many POSIX-like commands
 - cat, chgrp, chmod, chown, cp, du, ls, mkdir, mv, rm, stat, tail
- Some HDFS-specific commands
 - copyFromLocal, put, copyToLocal, get, getmerge, setrep

HDFS – FS shell commands

- copyFromLocal / put
 - Copy files from the local file system into fs

hadoop fs -copyFromLocal <localsrc> .. <dst>

Or

hadoop fs -put <localsrc> .. <dst>

HDFS – FS shell commands

- copyToLocal / get
 - Copy files from fs into the local file system

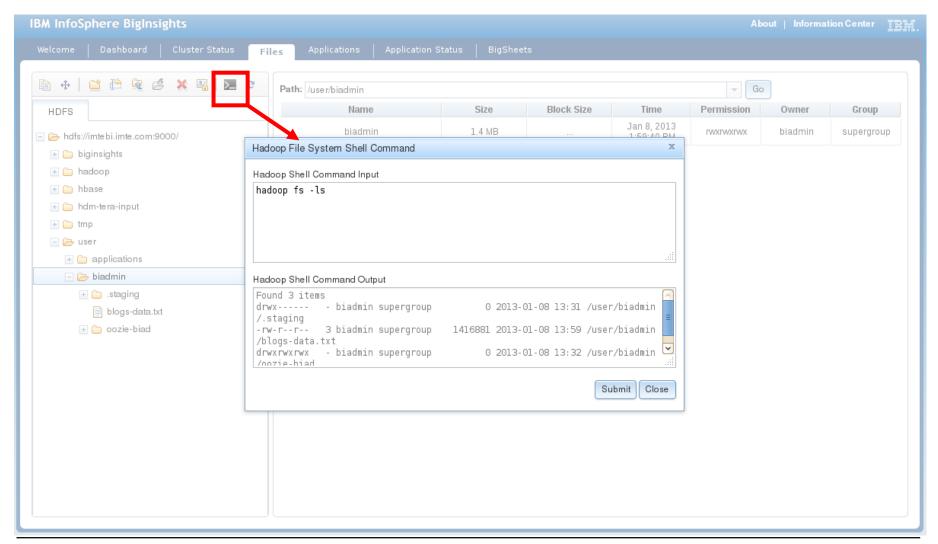
```
hadoop fs -copyToLocal [-ignorecrc] [-crc] <src> <localdst>
```

Or

```
hadoop fs -get [-ignorecrc] [-crc] <src> <localdst>
```



Files Tab – hadoop shell command





Questions?



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