

MAKE UP CLASS

Tomorrow @ 10am

ER3123



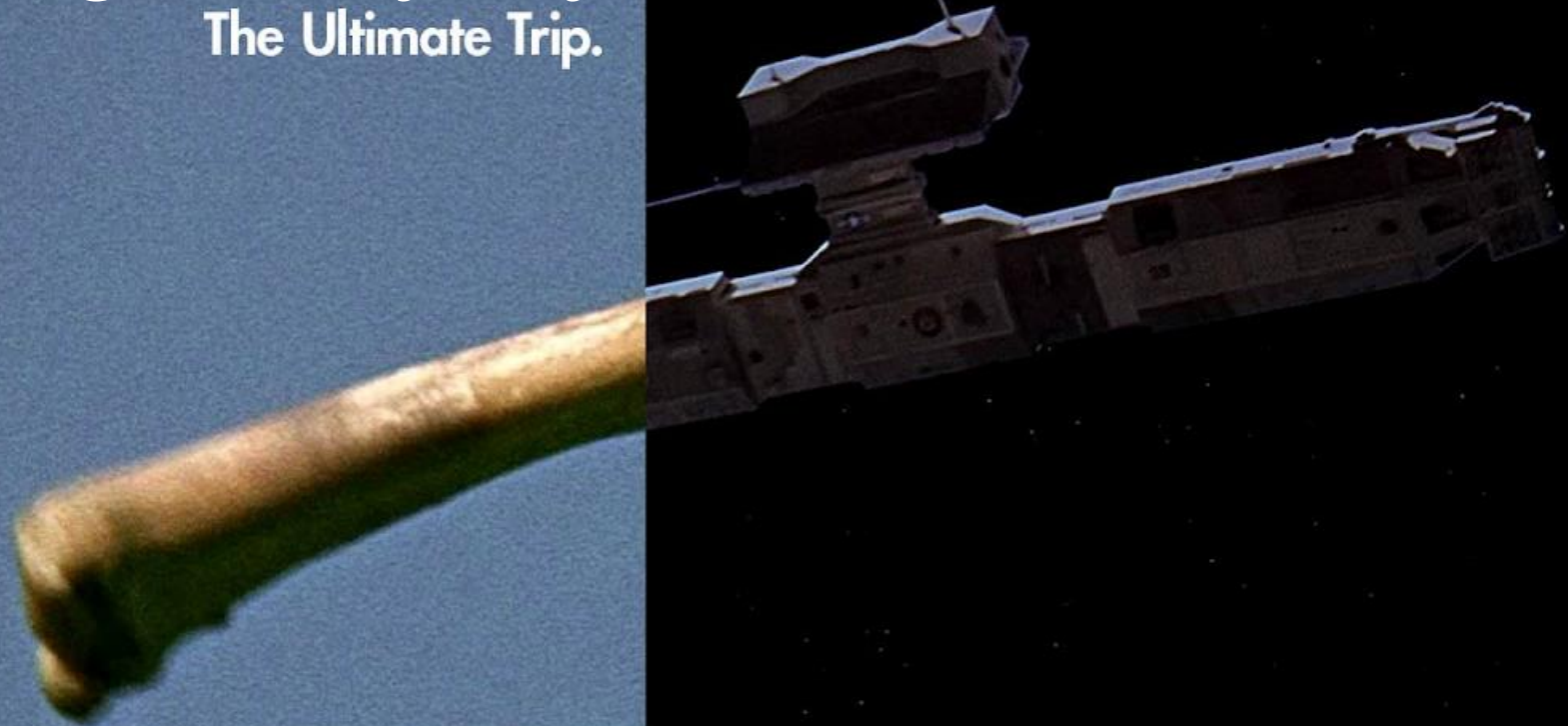
LEC05 & LAB05

Lec03 & Lab03 Marks

Lec02 Keys

W2022: A Digital Odyssey

The Ultimate Trip.

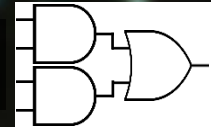


Number Systems | $(12)_{10} \rightarrow (1100)_2$

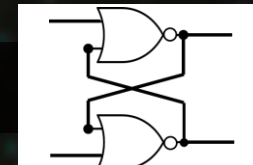
Logic Gates |



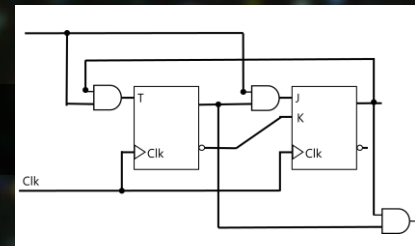
Combinational Logic |

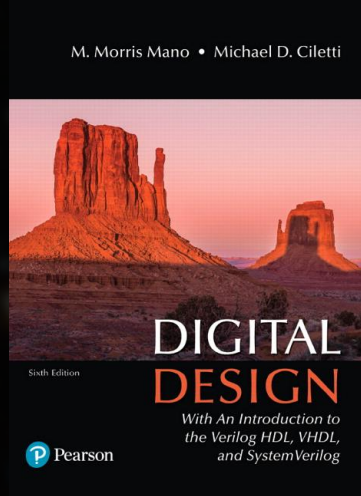


Flip-Flop |



Sequential Logic |





Chapter 2

Boolean Algebra and Logic Gates

BINARY COMPUTER

More Reliability in Engineering
Deep Logic Foundation



2001: A Space Odyssey (1968), Stanley Kubrick, Arthur C. Clarke

Ape Learning Scene: https://www.youtube.com/watch?v=VABNA_an2A0

Bone/Satellite Match Cut: <https://www.youtube.com/watch?v=avjdKTqiVvQ>

DESIGN COMPUTER

Positive Logic
Button-Up Approach

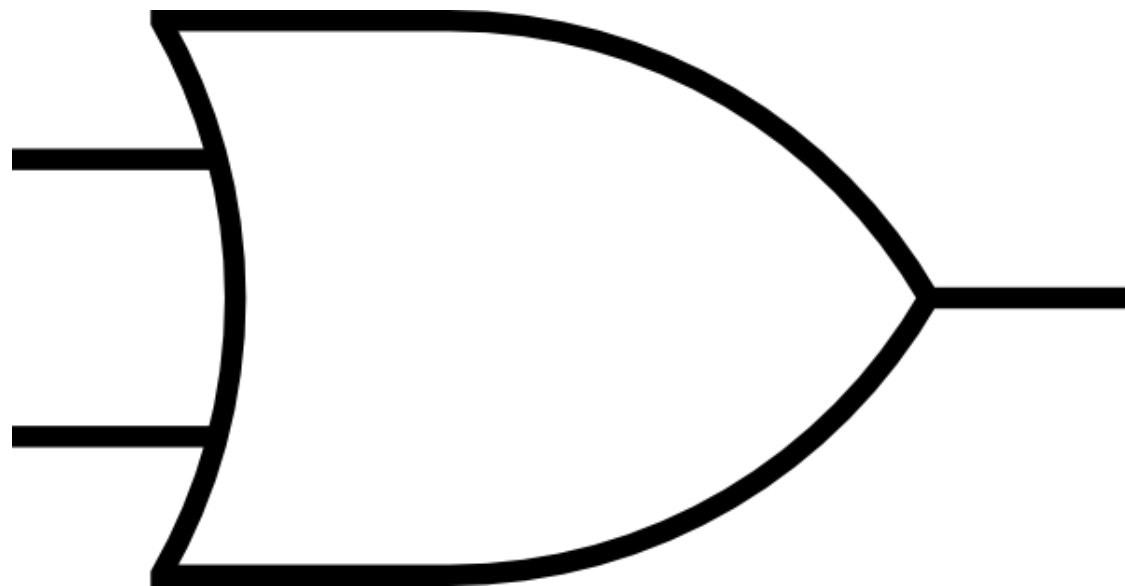
DESIGN COMPUTER

Positive Logic
Button-Up Approach

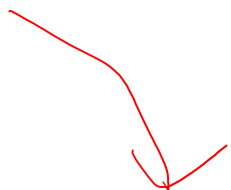
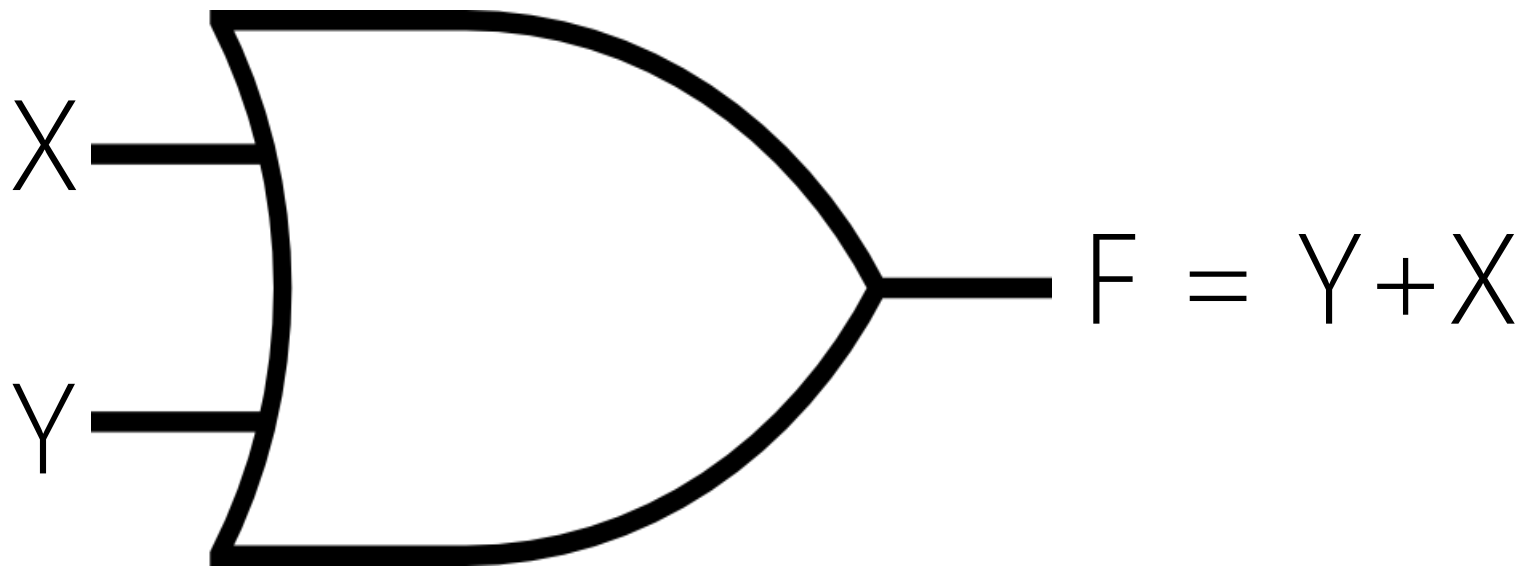
Finding simpler, but equivalent, computers reduces the overall cost!
Rely primarily on mathematical methods in Boolean algebra!

BUILD COMPUTER

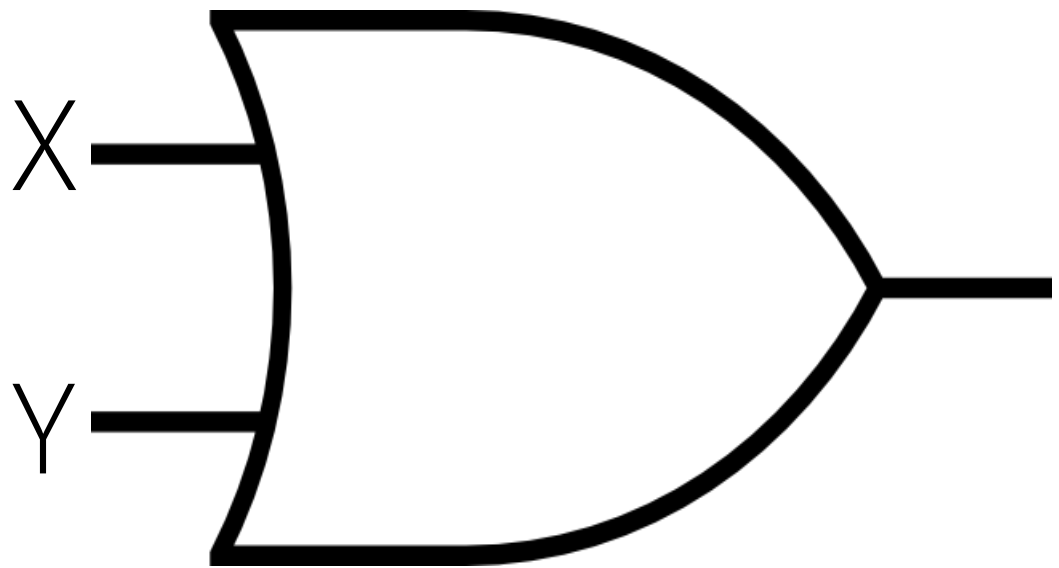
Electrical and Computer Engineering



OR



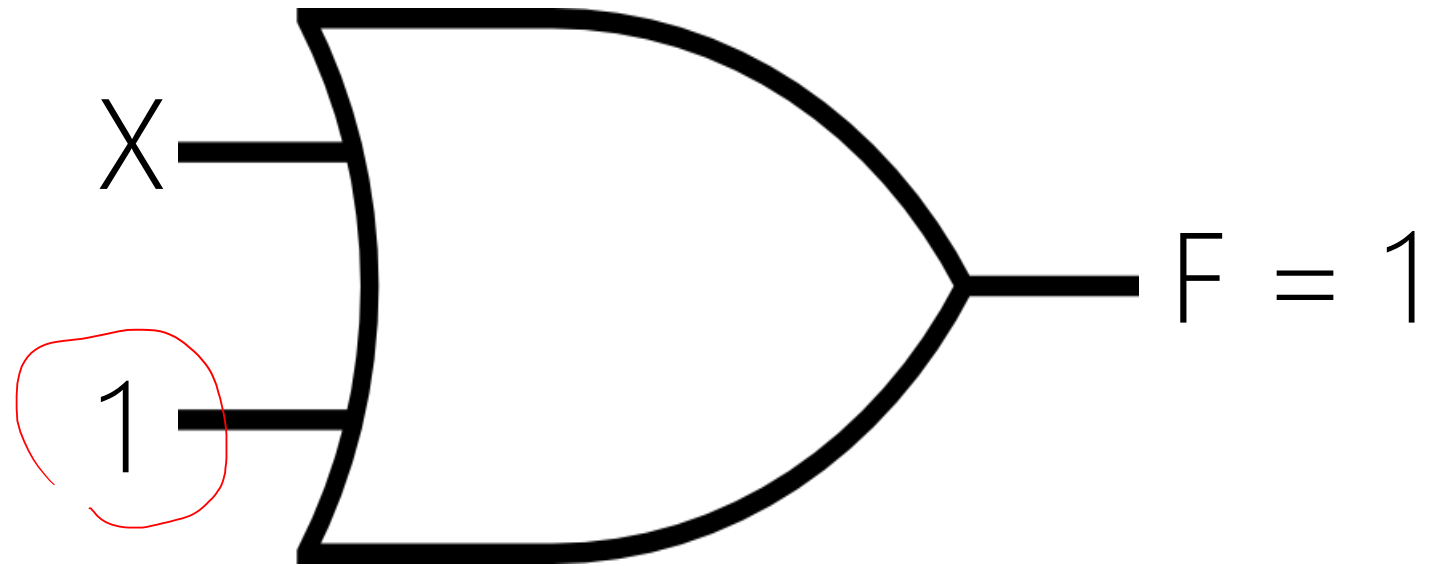
Y	X	Y OR X	$Y + X$
0	0		<u>0</u>
0	1		1
1	0		1
1	1		1



$$F = Y + X = X + Y$$

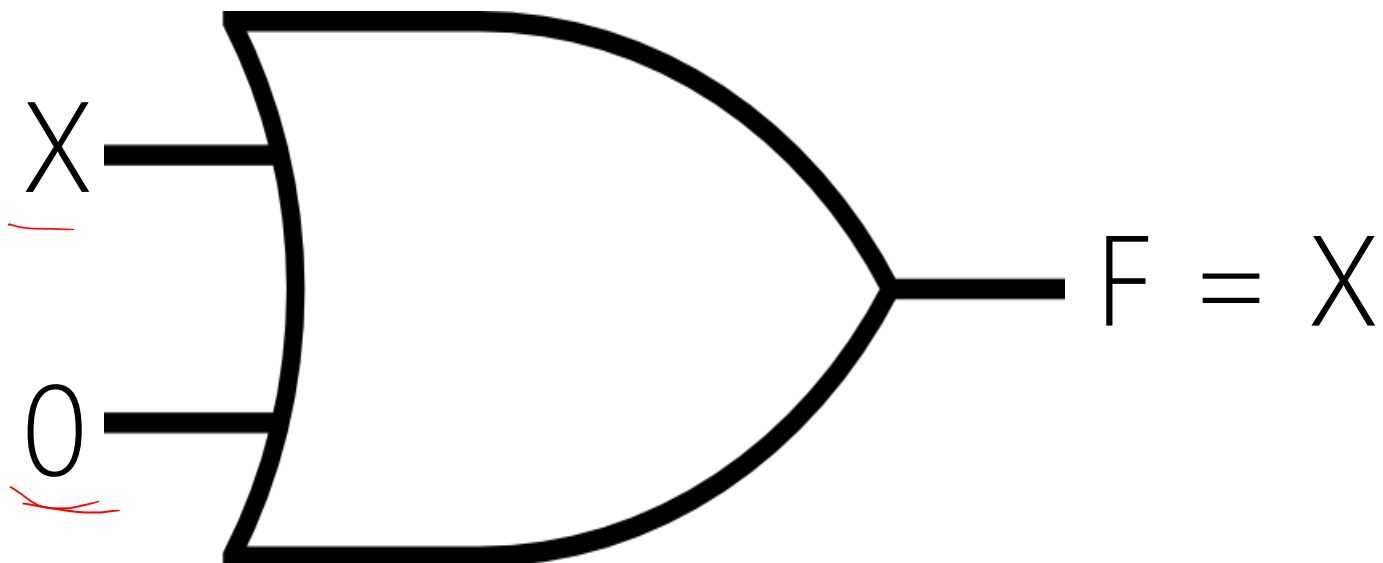
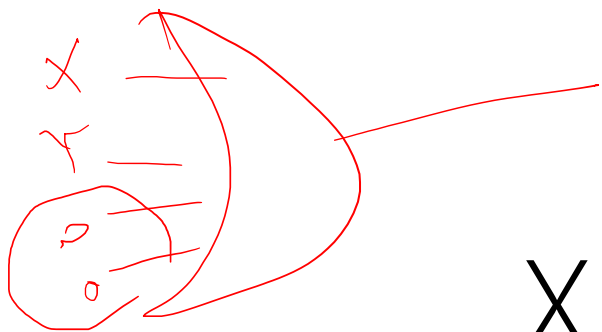
Commutative

X	Y	X OR Y	X+Y
0	0	0	
0	1	1	
1	0	1	
1	1	1	



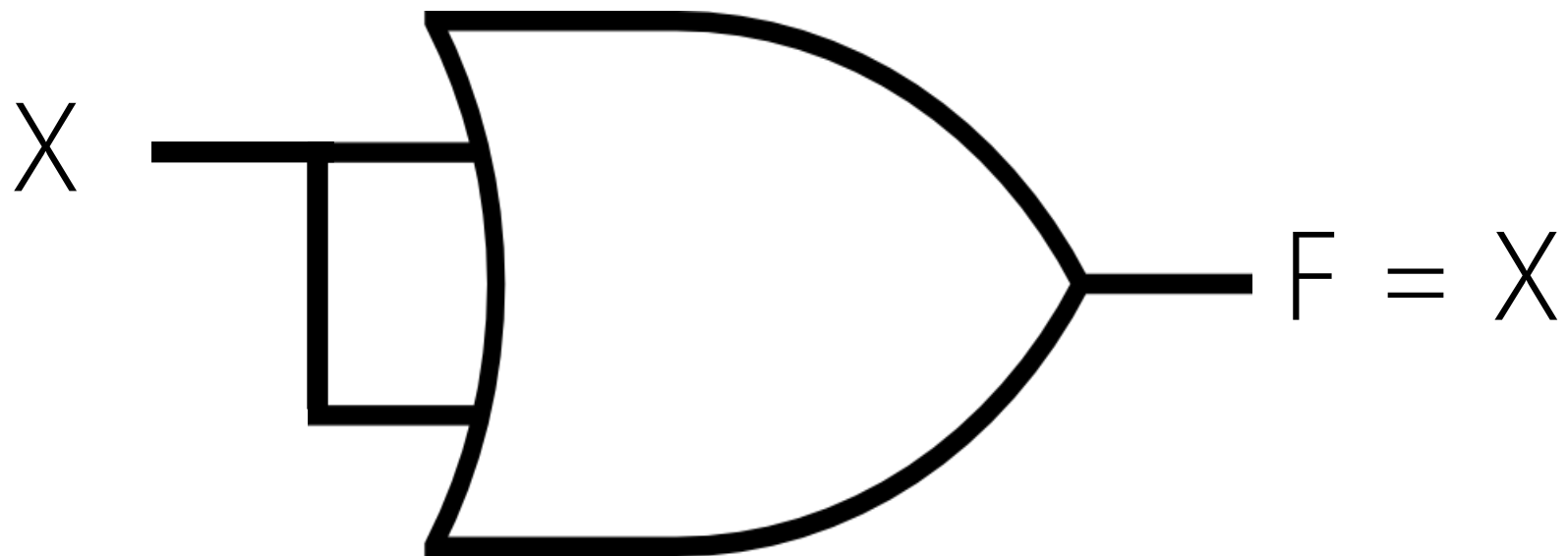
$$F = X + 1 = 1$$

Y	X	Y+X
1	0	1
1	1	1



Y	X	Y+X
0	0	0
0	1	1

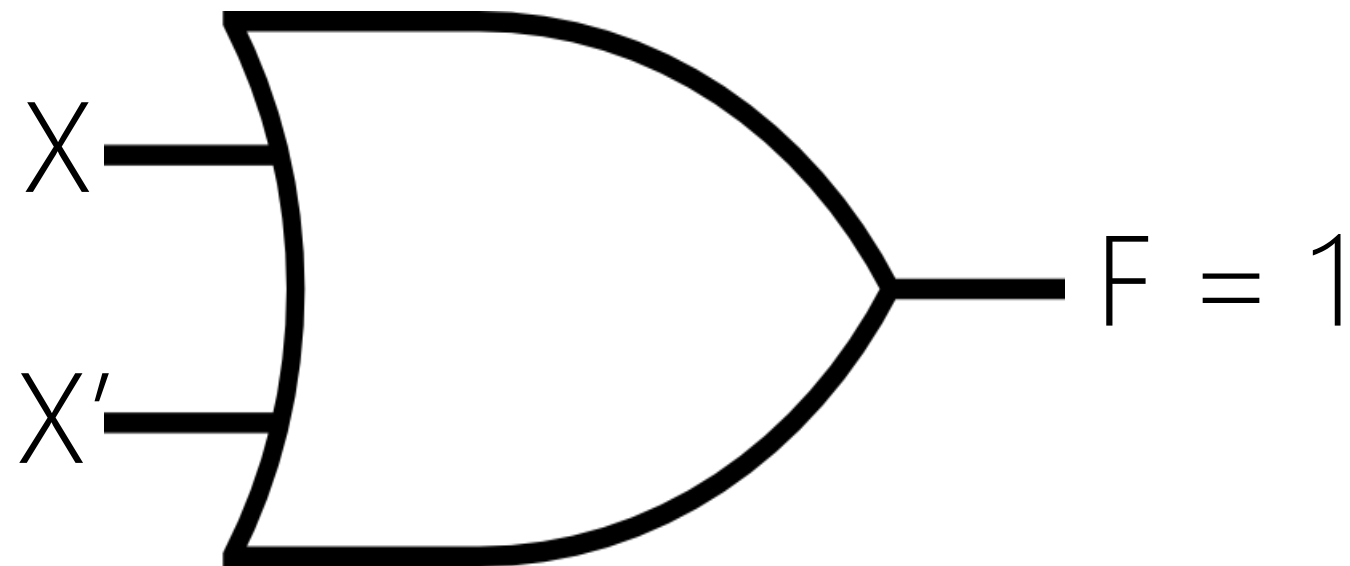
$$F = X + 0 = X$$



X	X	$X+X$
<u>0</u>	<u>0</u>	<u>0</u>
1	1	1

$$F = X + X = X$$

<u>X'</u>	<u>X</u>	X' + X
1	0	<u>1</u>
0	1	1



$$F = X + X' = 1$$

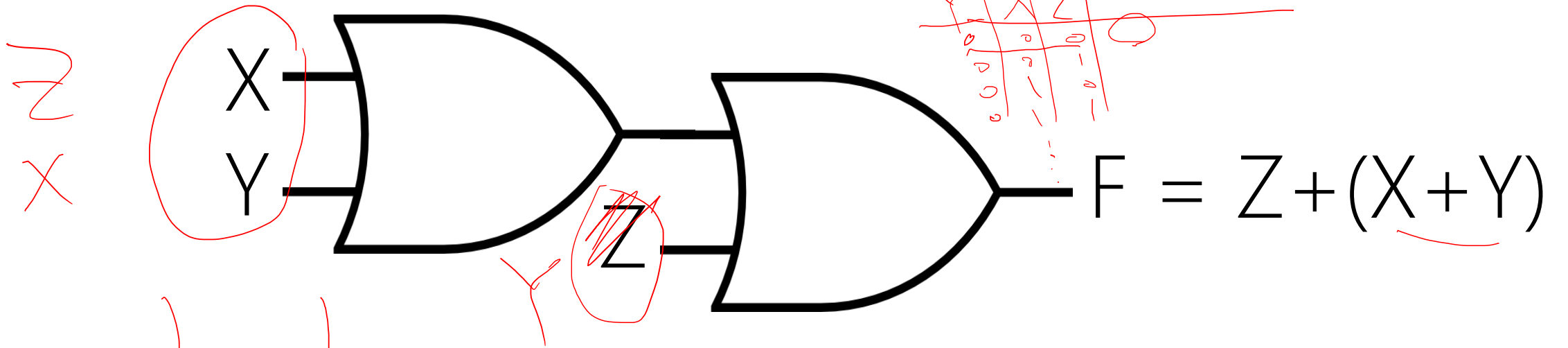
DESIGN

given the functionality, design the structure of a system

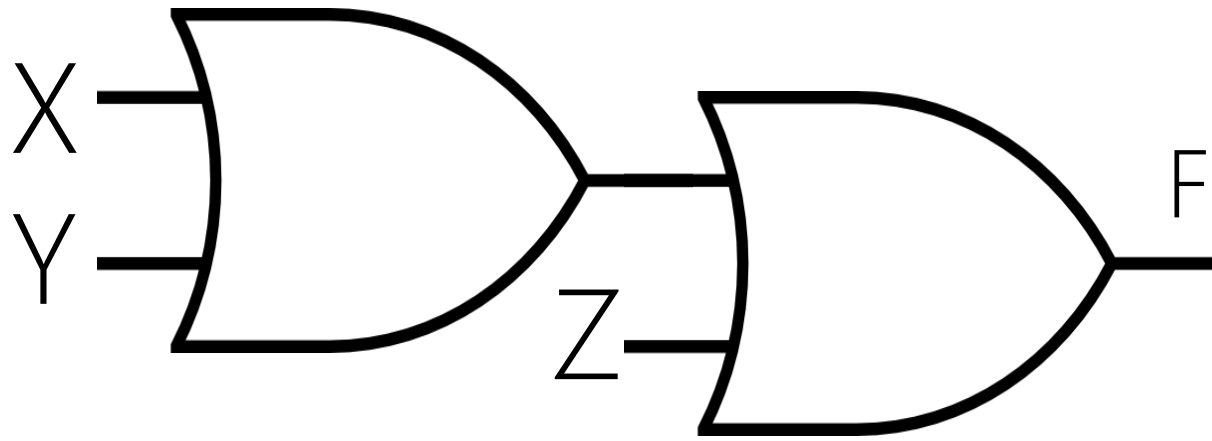
3-INPUT OR

DESIGN PATTERNS

Using Same or Similar Previous Designs for New Designs



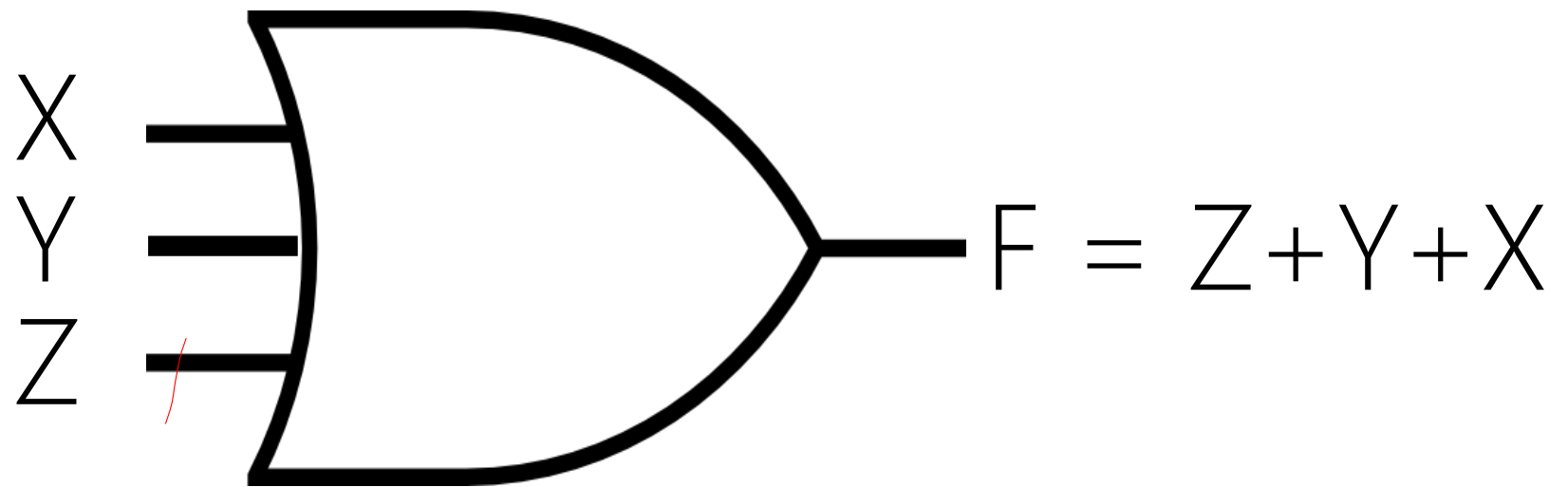
Z	Y	X	Z+(X+Y)	
0	0	0	0	0
0	0	1	1	1
0	1	0	1	1
0	1	1	1	1
1	0	0	1	1
1	0	1	1	1
1	1	0	1	1
1	1	1	1	1



$$F = Z + (Y + X) = Z + (X + Y) = (Z + X) + Y \\ = Z + Y + X$$

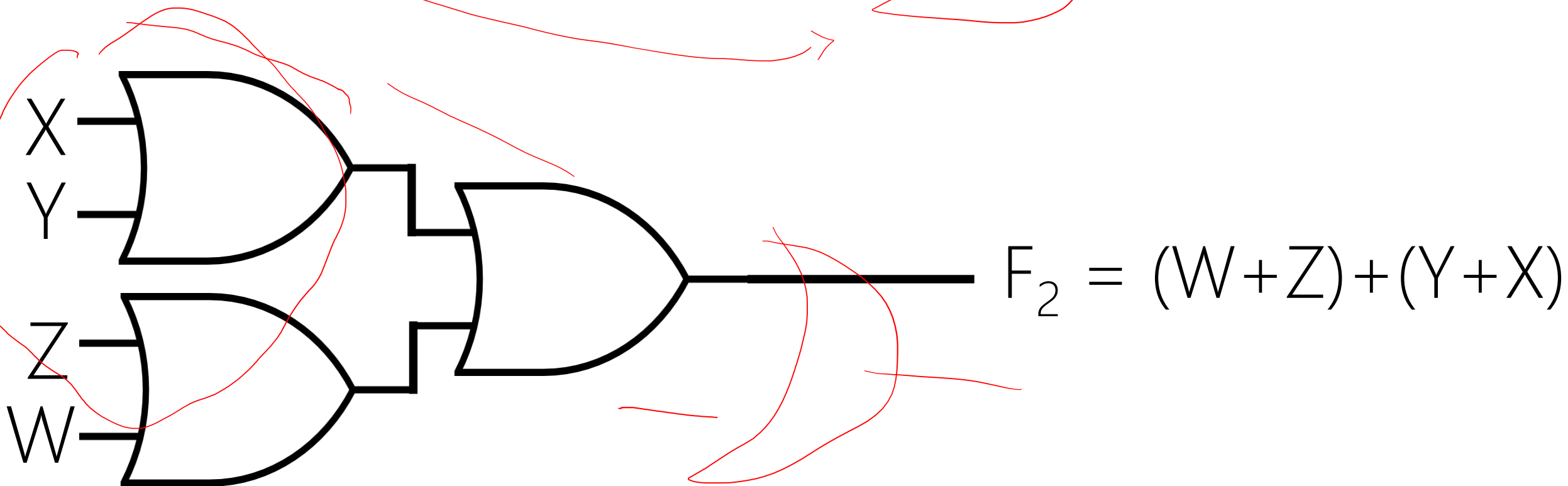
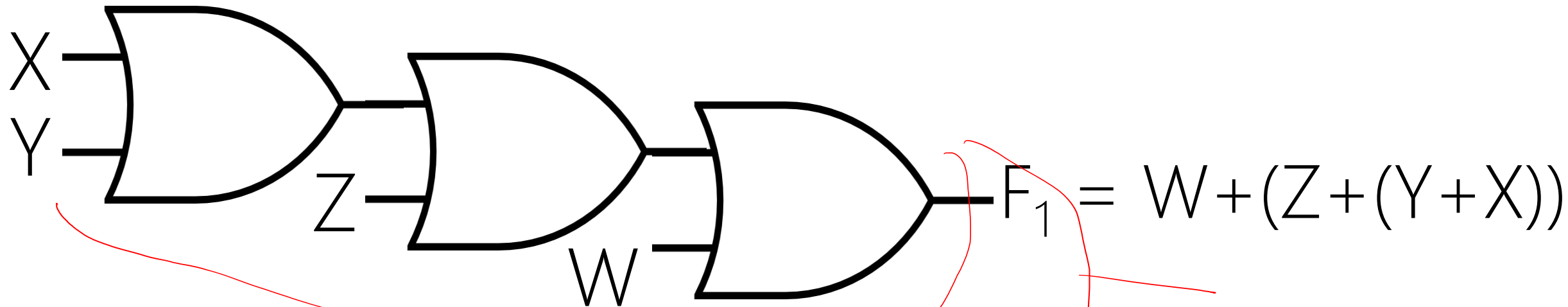
Associative

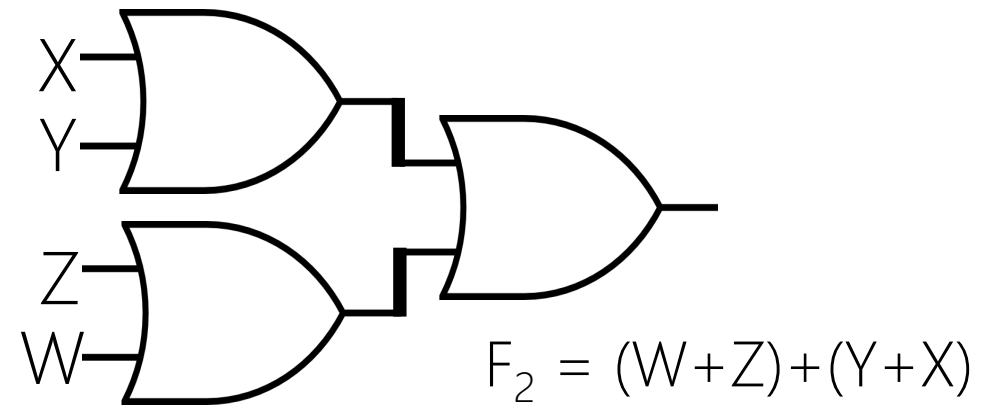
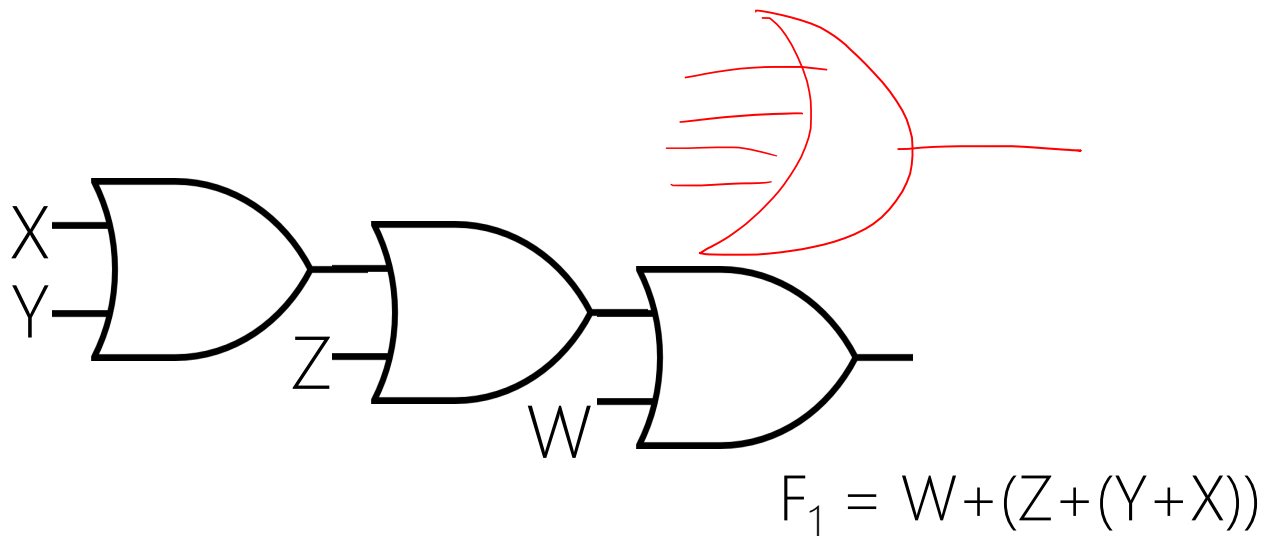
Z	Y	X	$Z + (Y + X)$	$Z + (X + Y)$	$(Z + X) + Y$	ZXY
0	0	0	0	0	0	0
0	0	1	1	1	1	0
0	1	0	1	1	1	0
0	1	1	1	1	1	0
1	0	0	1	1	1	0
1	0	1	1	1	1	0
1	1	0	1	1	1	0
1	1	1	1	1	1	0



Z	Y	X	$Z+(Y+X)$	$Z+(X+Y)$	$(Z+X)+Y$	ZXY
0	0	0			0	
0	0	1			1	
0	1	0			1	
0	1	1			1	
1	0	0			1	
1	0	1			1	
1	1	0			1	
1	1	1			1	

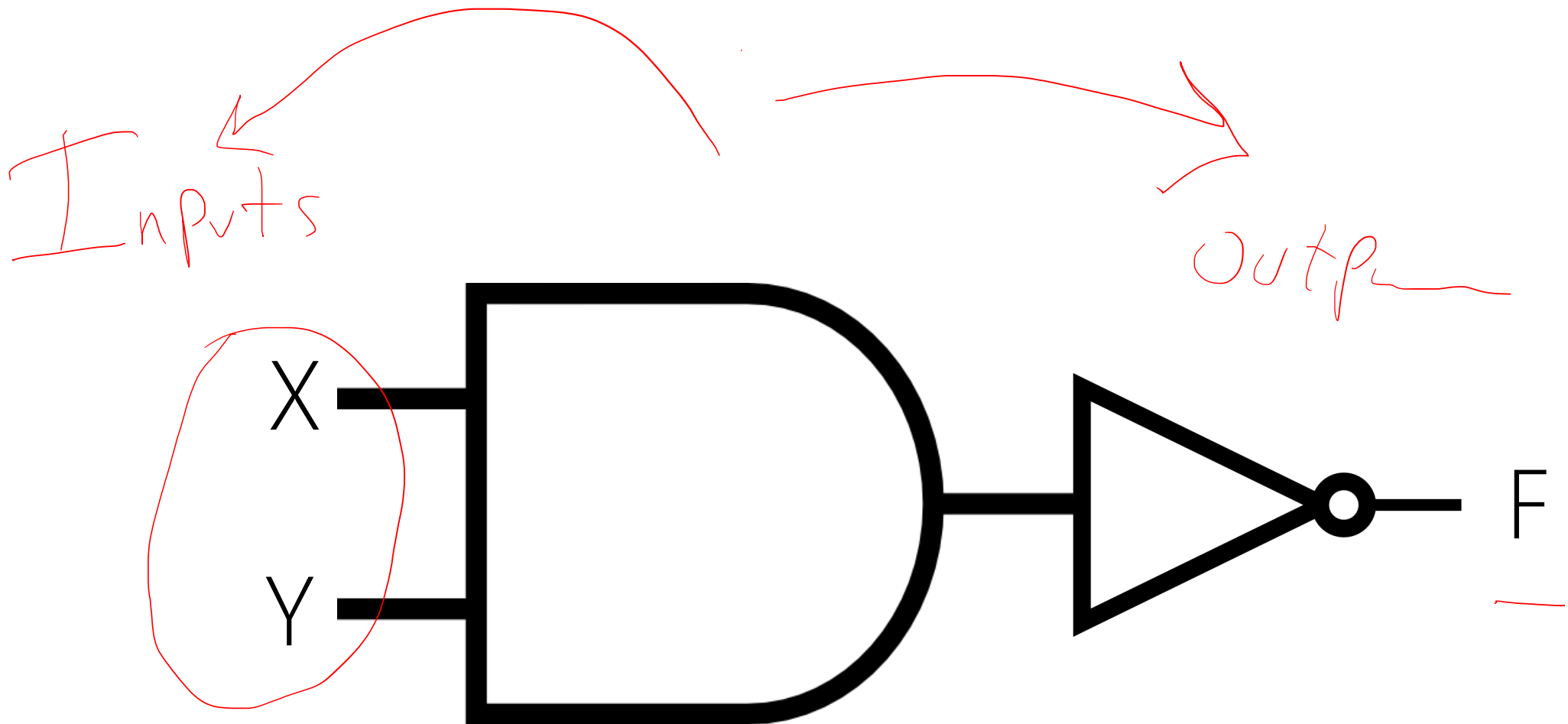
4-INPUT OR

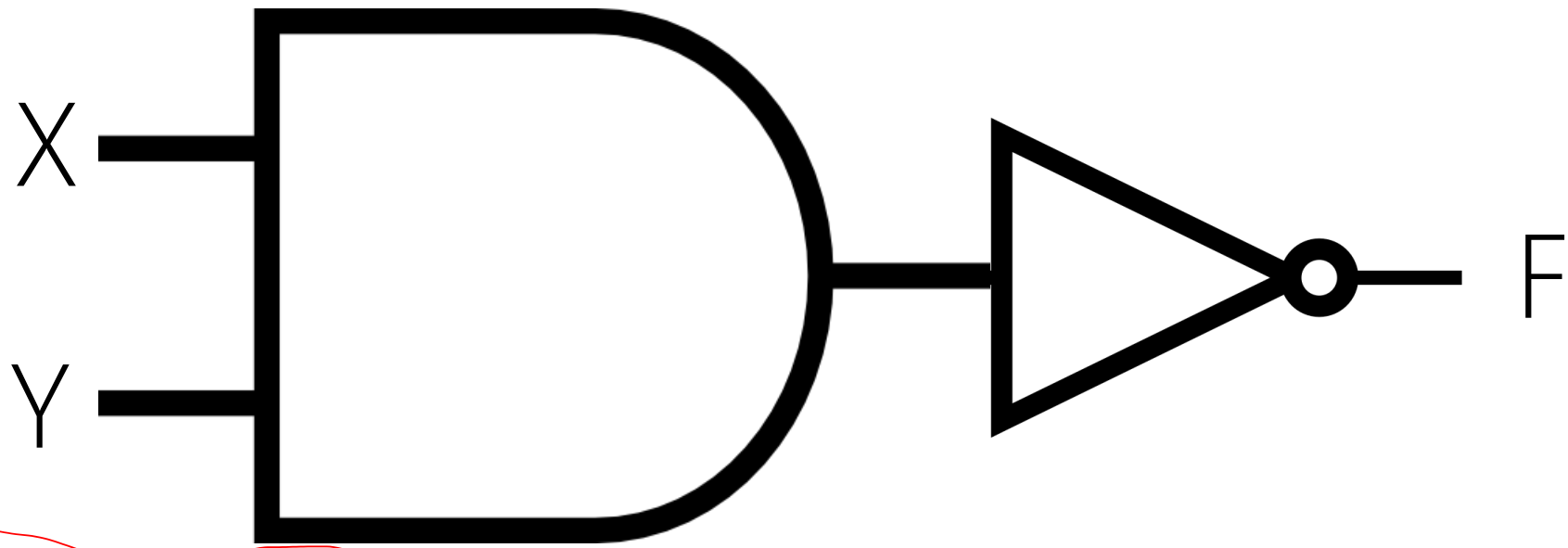




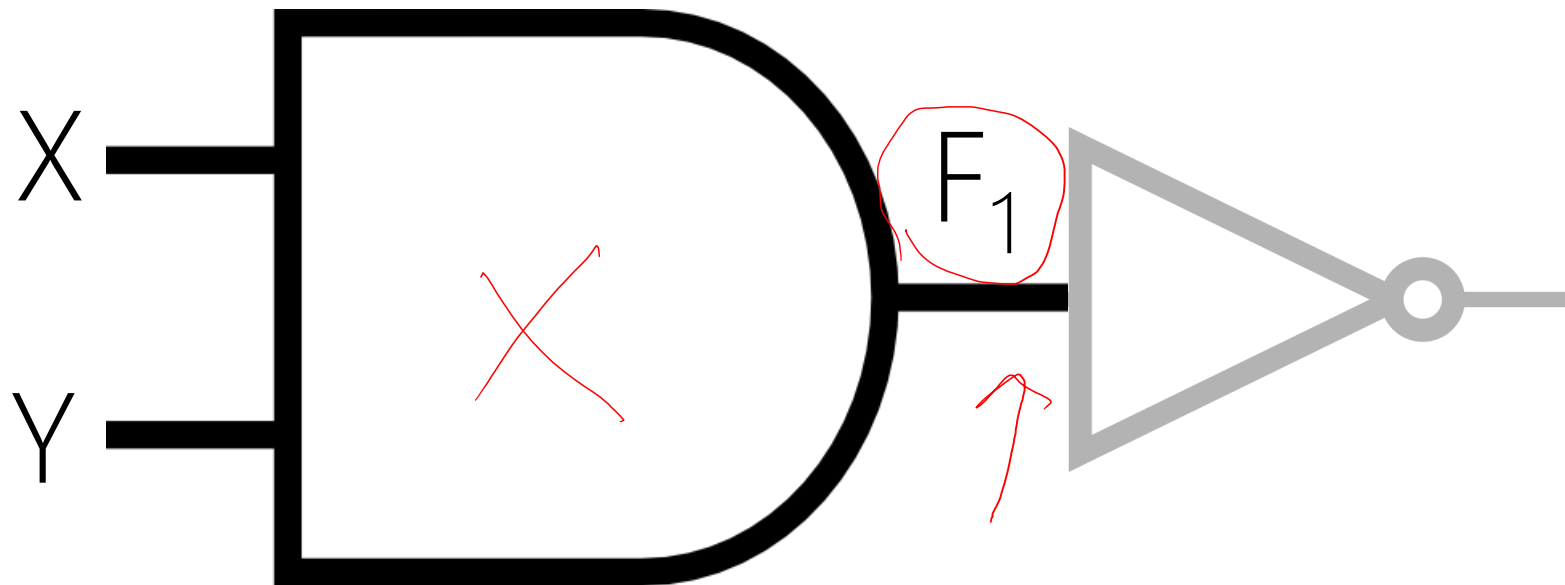
$F = W + Z + Y + X$	F_1	F_2
Effective (True)	Yes	Yes
Efficient (Fast)	Hmm, 3 levels, No!	Yes! 2 levels
Min. Cost	3 gates, Yes	3 gates, Yes

DESIGN vs. ANALYSIS

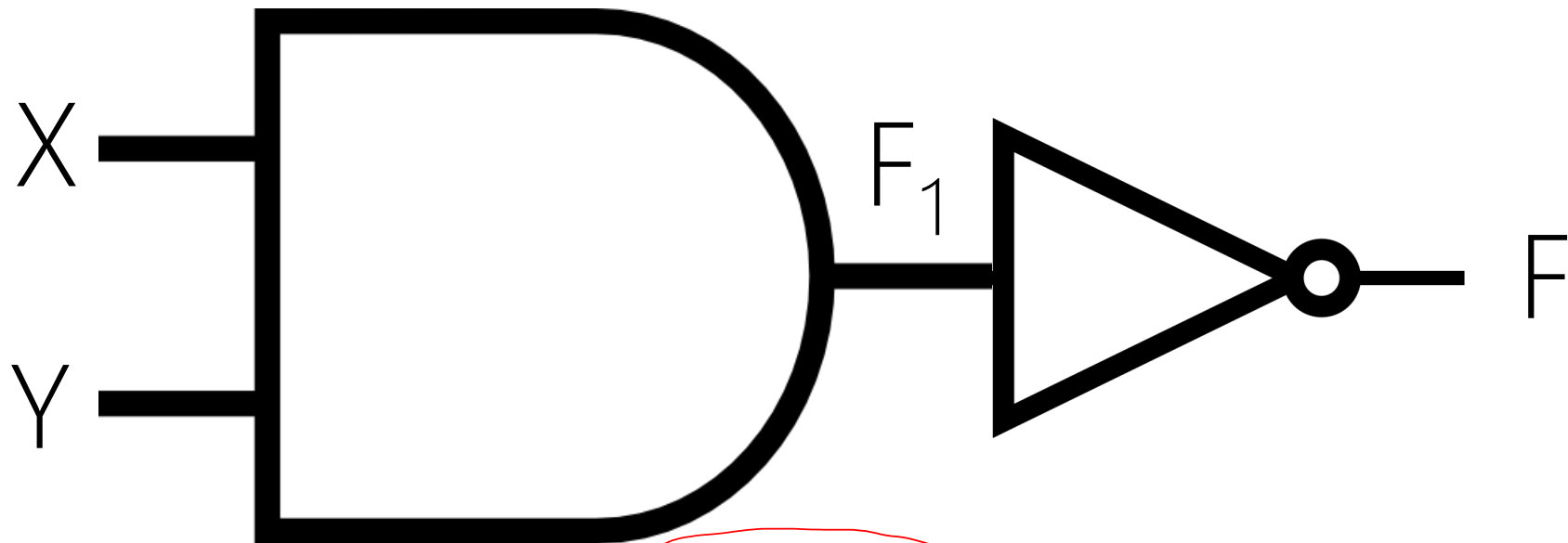




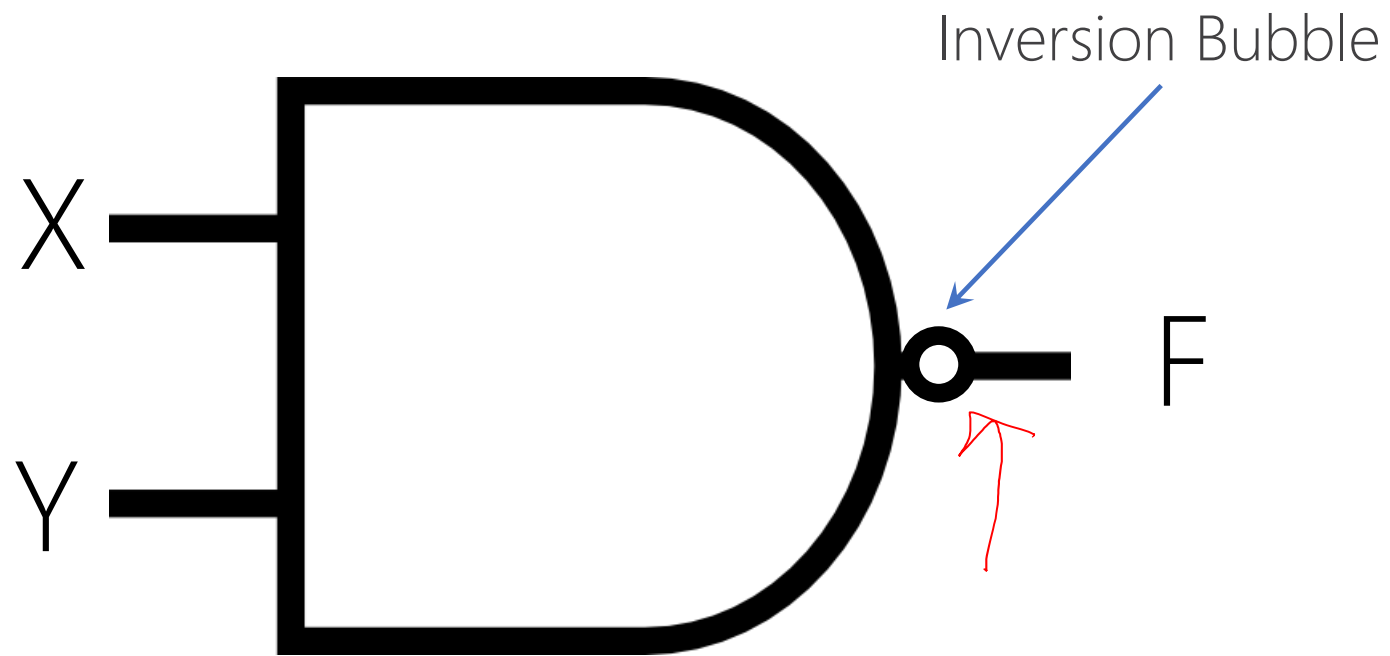
Y	X	F = ?
0	0	?
0	1	?
1	0	?
1	1	?



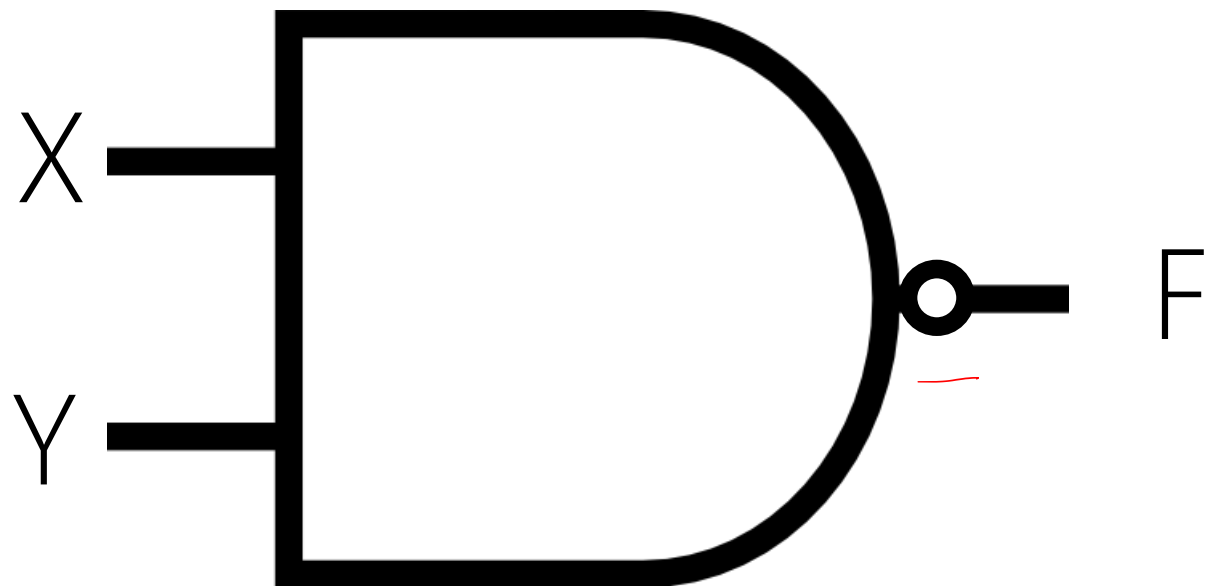
Y	X	$F_1 = YX$
0	0	<u>0</u>
0	1	<u>0</u>
1	0	<u>0</u>
1	1	1



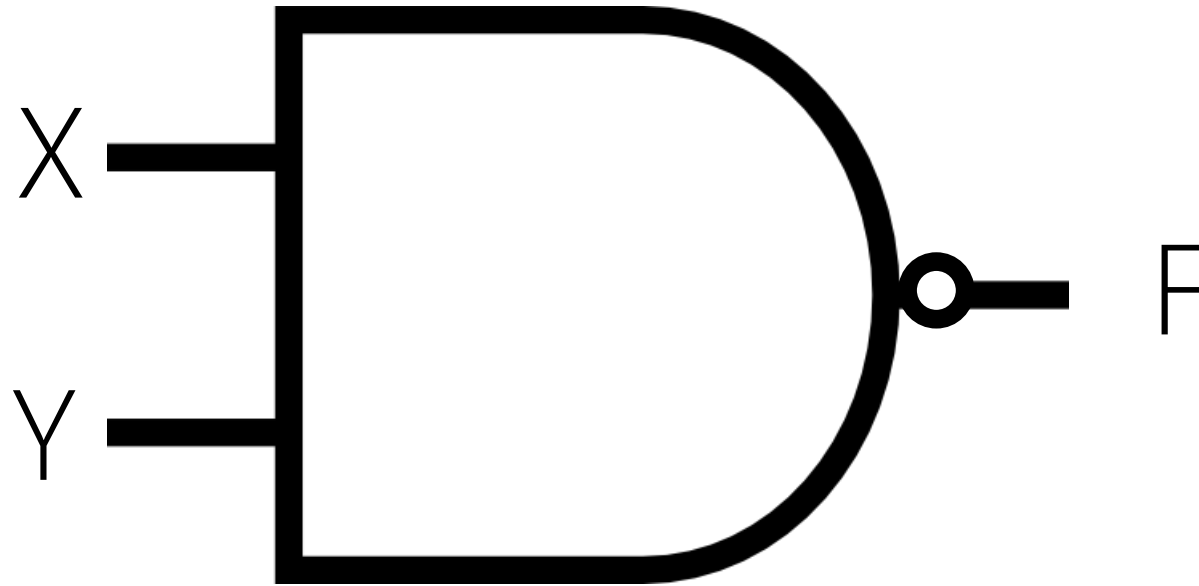
Y	X	$F_1 = YX$	$F = (YX)'$
0	0	0	1
0	1	0	1
1	0	0	1
1	1	1	0



NAND (Not – AND)



Y	X	$F = (YX)'$	$F = Y \uparrow X$
0	0	1	1
0	1	1	1
1	0	1	1
1	1	0	0



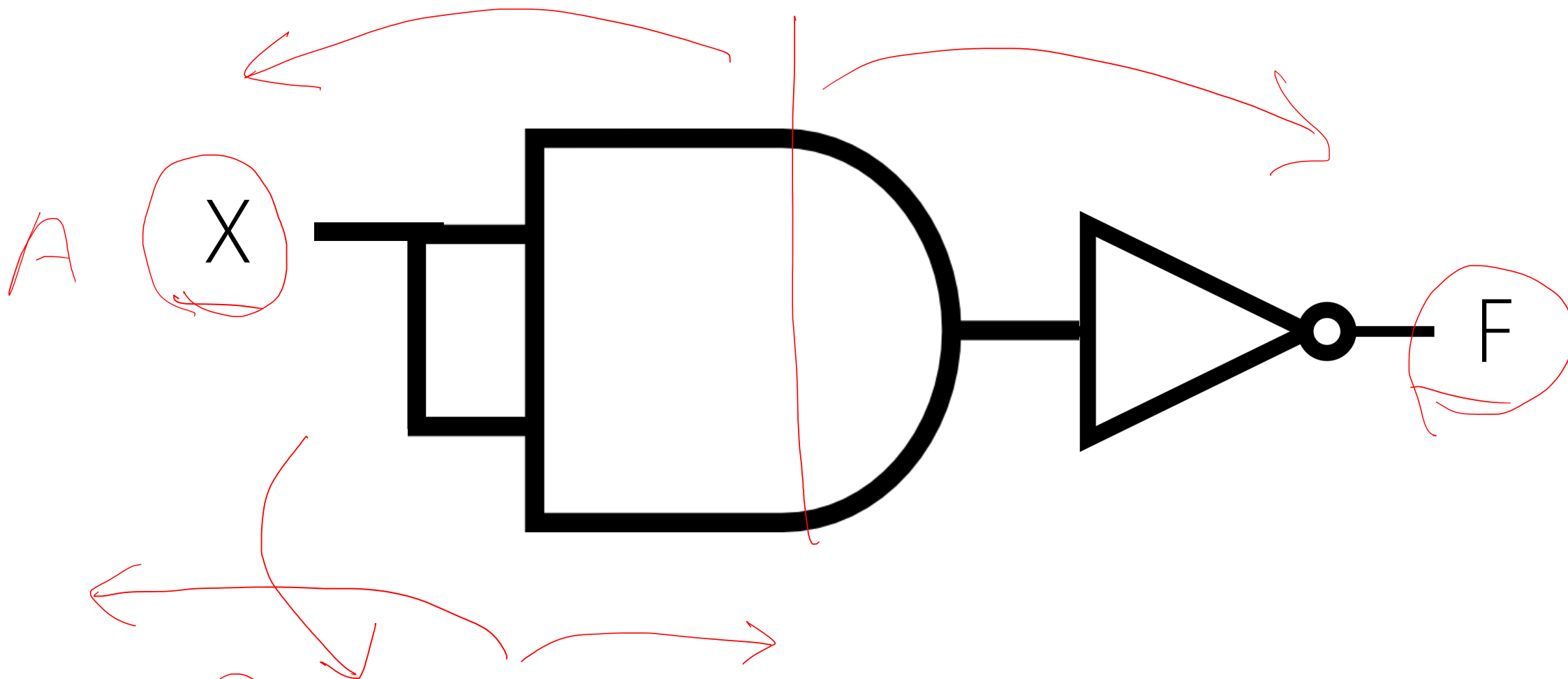
$$F = (YX)' = (XY)' = Y \uparrow X = X \uparrow Y$$

Commutative

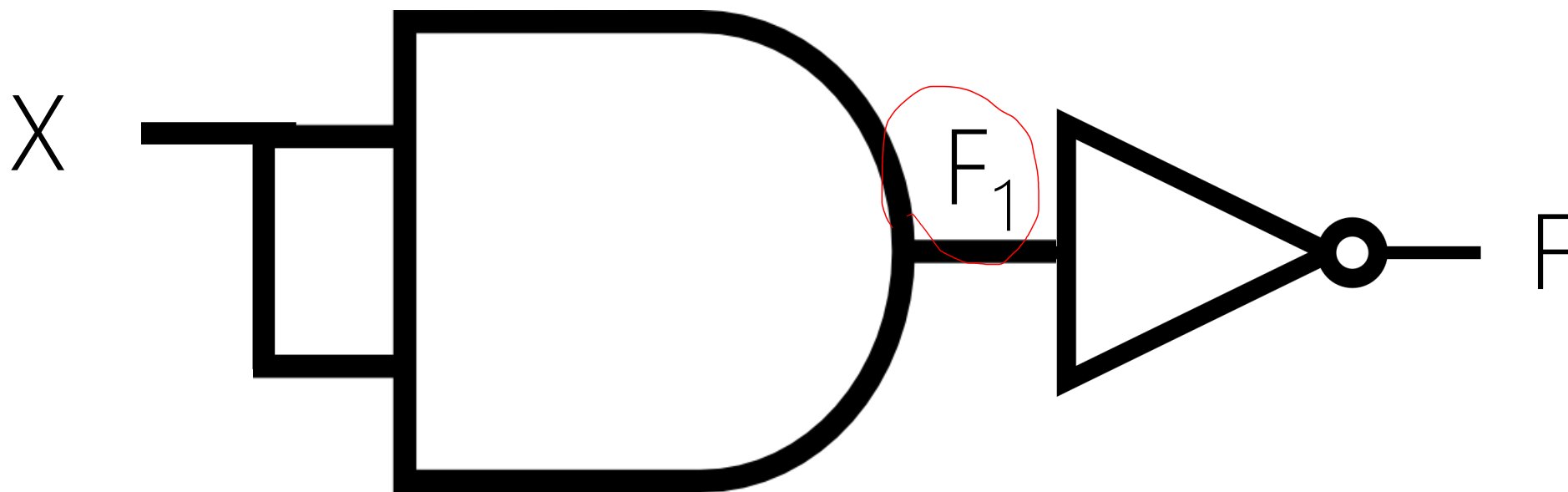
ANALYSIS

given the structure of a system, find its functionality.

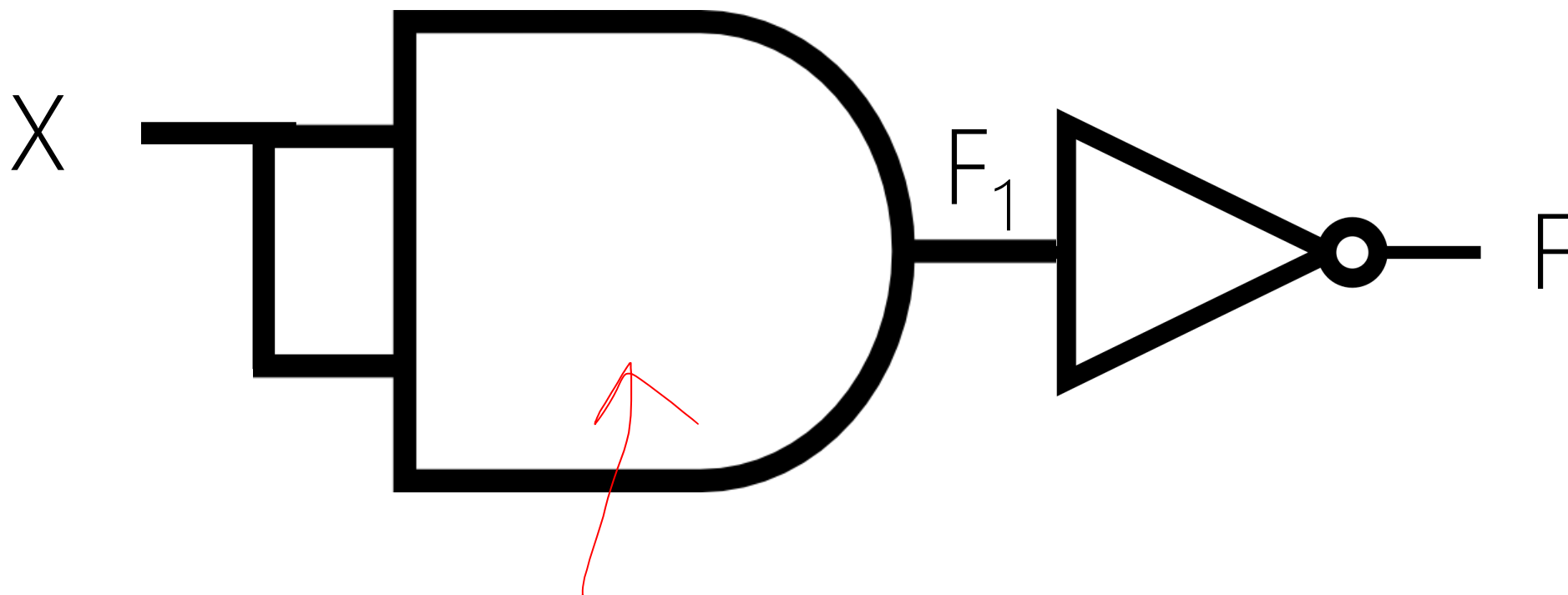
determine the functionality exhibited by a structure.



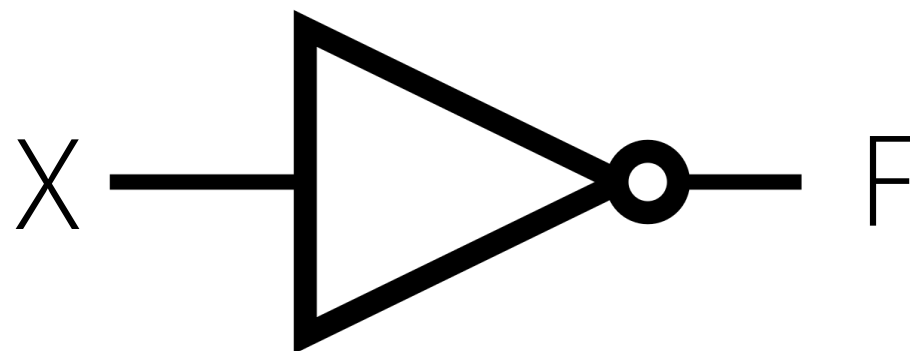
X		F = ?	
0	→		
1	→		



X	$F_1 = \underline{XX}$	F = ?
0	0	
1	1	



X	$F_1 = XX$	$F = (XX)'$
0	0	1
1	1	0

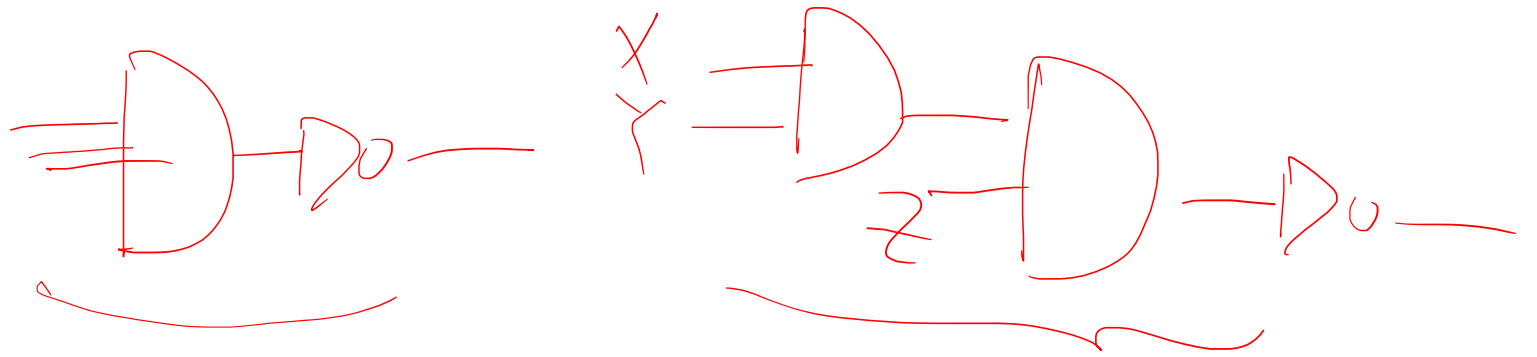


X	$F = (XX)' = X \uparrow X = X'$
0	1
1	0

DESIGN

given the **functionality**, design the structure of a system

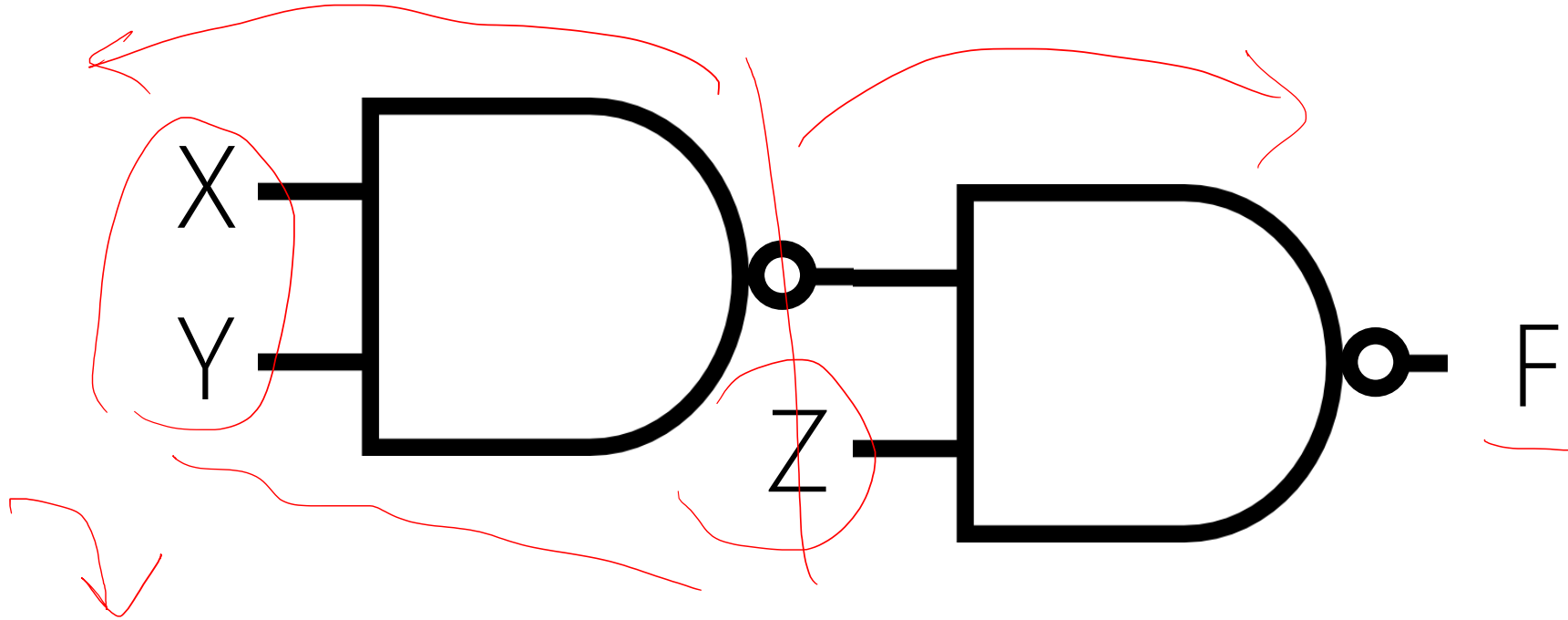
3-INPUT NAND



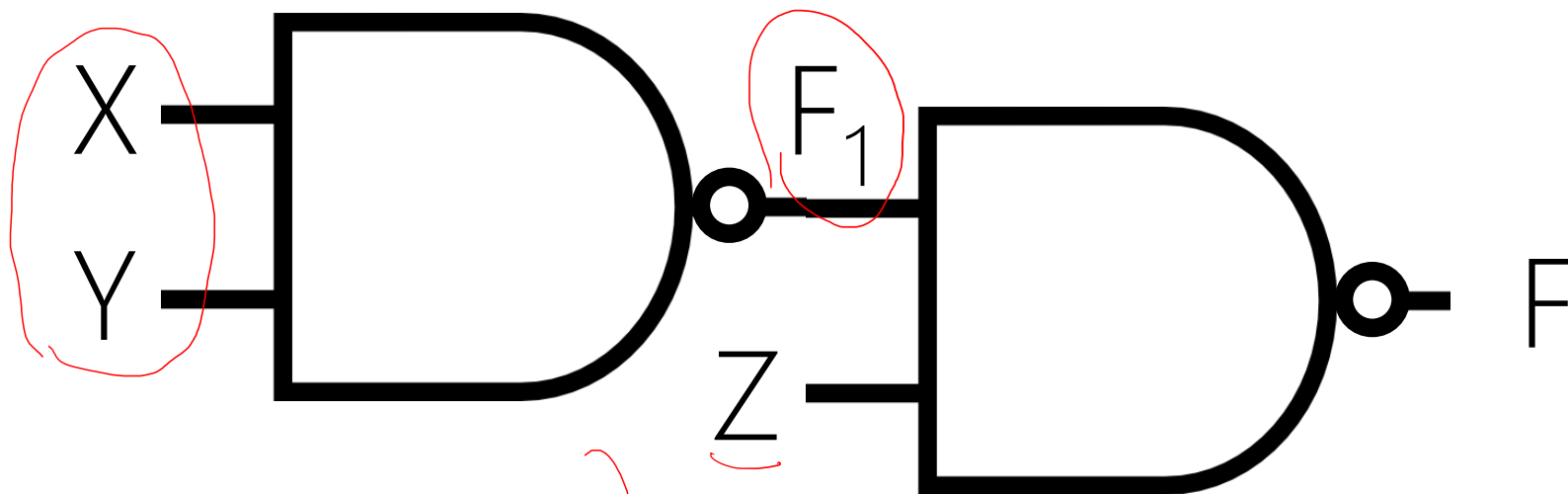
Z	Y	X	$F = (ZYX)'$
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	0

DESIGN PATTERNS

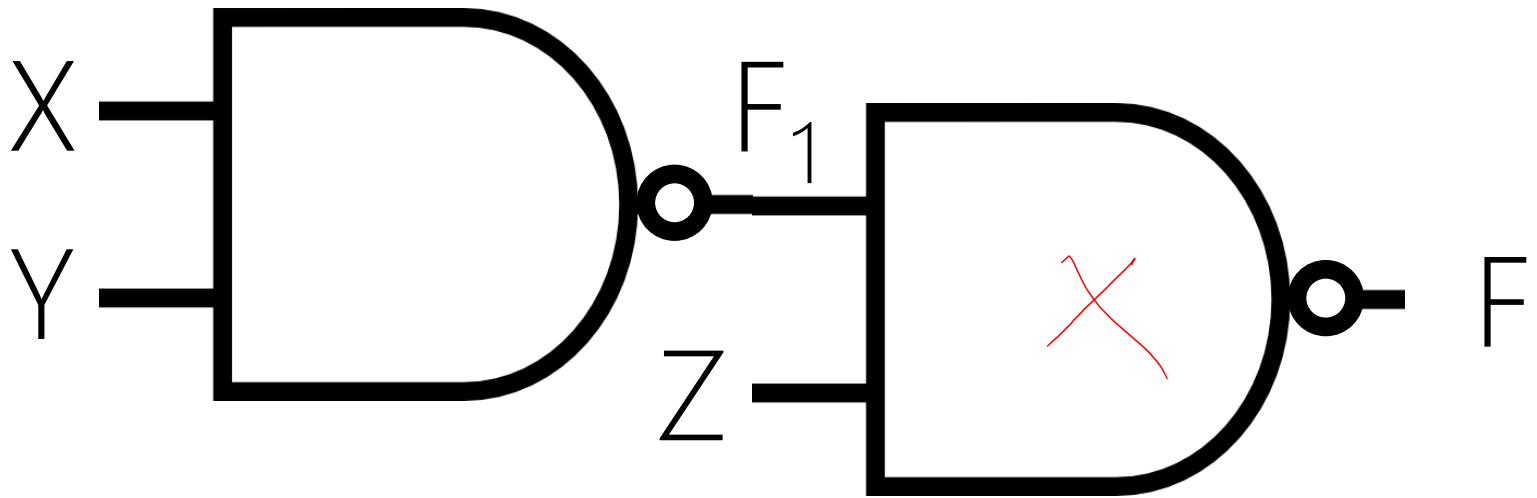
Using Same or Similar Previous Designs for New Designs



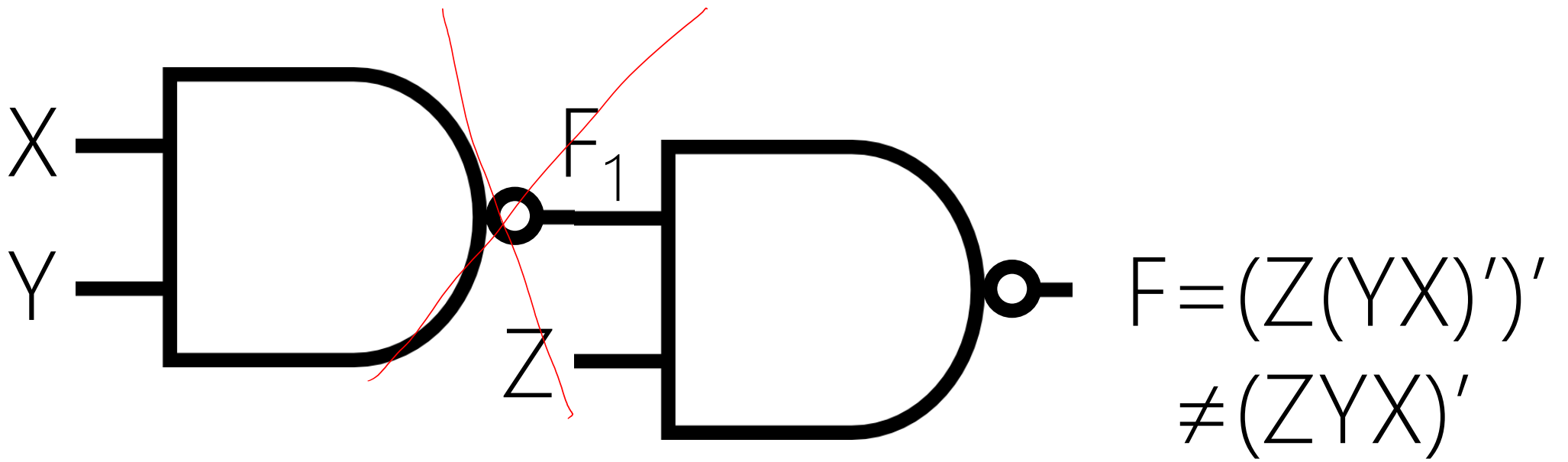
Z	Y	X	F = ?
0	0	0	
0	0	1	
0	1	0	
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	



Z	Y	X	$F_1 = (YX)'$	F = ?
0	0	0	1	
0	0	1	1	
0	1	0	1	
0	1	1	0	
1	0	0	1	
1	0	1	1	
1	1	0	1	
1	1	1	0	



Z	Y	X	$F_1 = (YX)'$	$F = (ZF_1)' = (Z(YX))'$
0	0	0	1	1
0	0	1	1	1
0	1	0	1	1
0	1	1	0	1
1	0	0	1	0
1	0	1	1	0
1	1	0	1	0
1	1	1	0	1



Z	Y	X	$F_1 = (YX)'$	$F = (ZF_1)' = (Z(YX)')'$	$F = (ZYX)'$
0	0	0	1	1	1
0	0	1	1	1	1
0	1	0	1	1	1
0	1	1	0	1	1
1	0	0	1	0	1
1	0	1	1	0	1
1	1	0	1	0	1
1	1	1	0	1	0

DESIGN PATTERNS

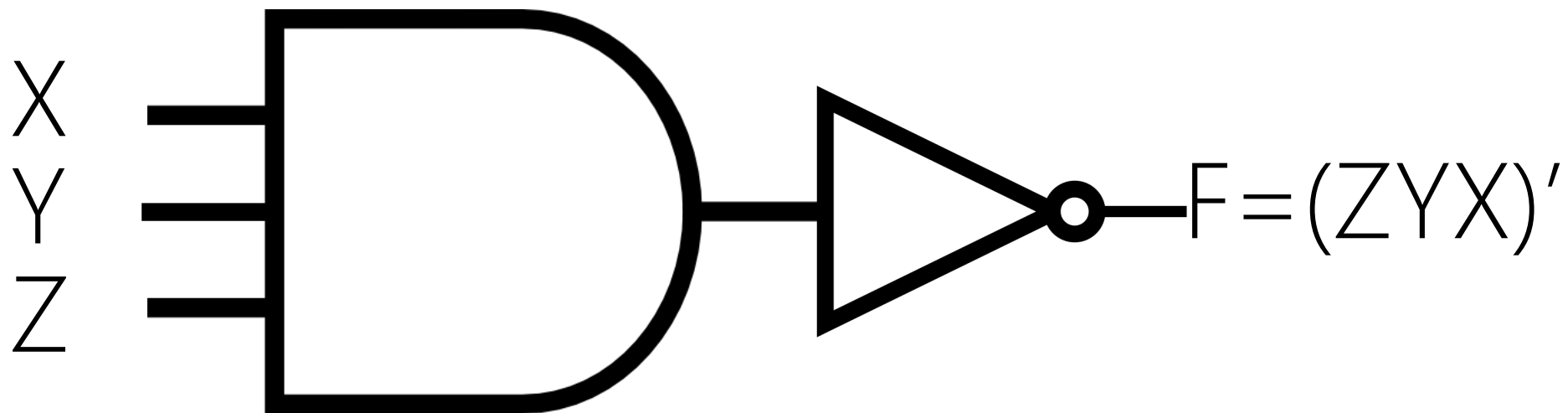


Using Same or Similar Previous Designs for New Designs

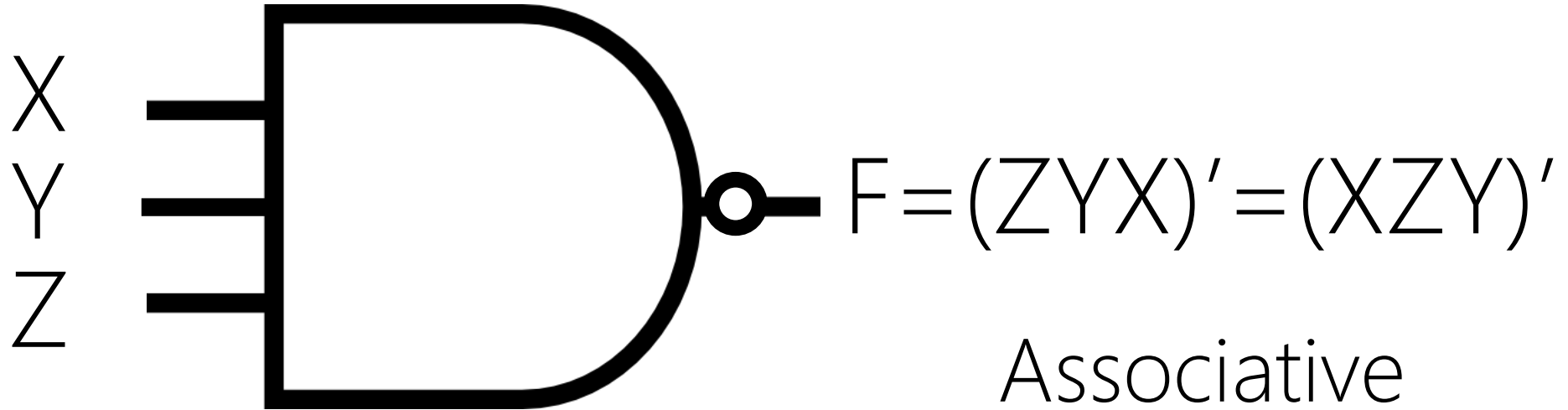
DESIGN → ANALYSIS → EVALUATION

Always Check

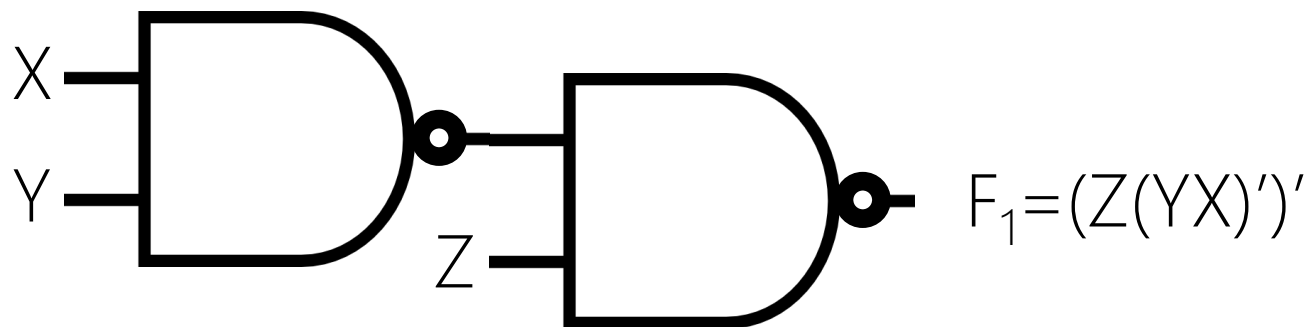
NOT (3-INPUT AND)



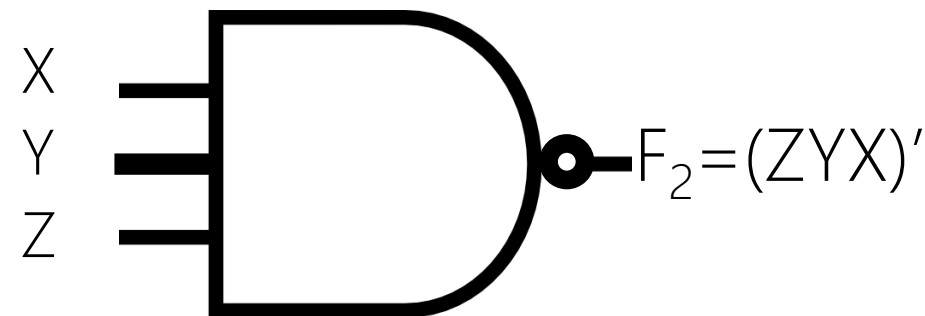
Z	Y	X	$F = (ZYX)'$	$F = (ZYX)'$
0	0	0	1	1
0	0	1	1	1
0	1	0	1	1
0	1	1	1	1
1	0	0	1	1
1	0	1	1	1
1	1	0	1	1
1	1	1	0	0



Z	Y	X	$F = (ZYX)'$	$F = (ZYX)'$
0	0	0	1	1
0	0	1	1	1
0	1	0	1	1
0	1	1	1	1
1	0	0	1	1
1	0	1	1	1
1	1	0	1	1
1	1	1	0	0

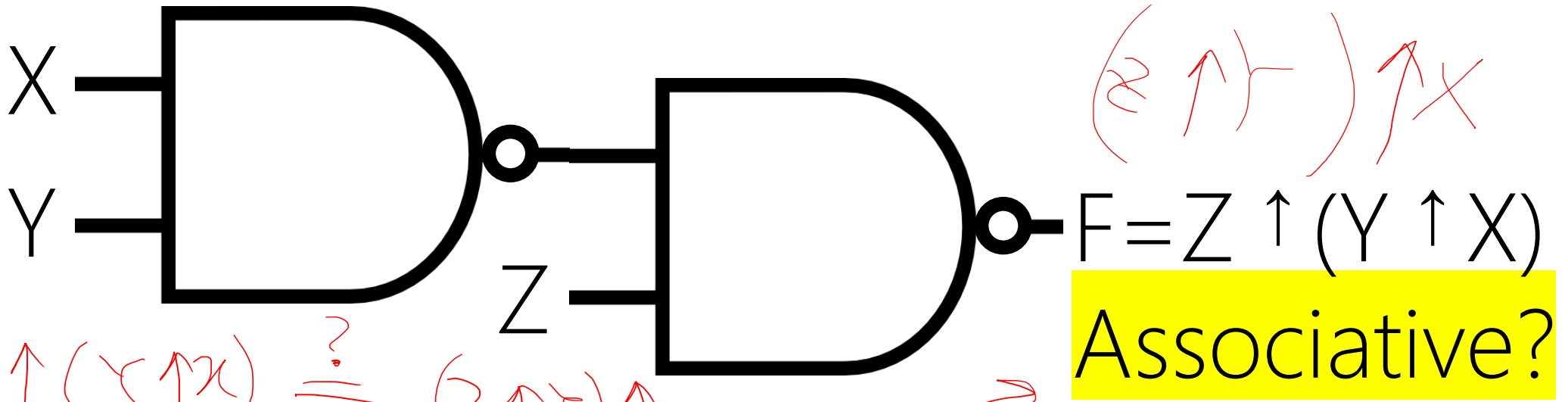


A



B

$F = (ZYX)'$	F_1	F_2
Effective (True)	No!	Yes
Efficient (Fast)	⊗	⊗
Min. Cost	⊗	⊗



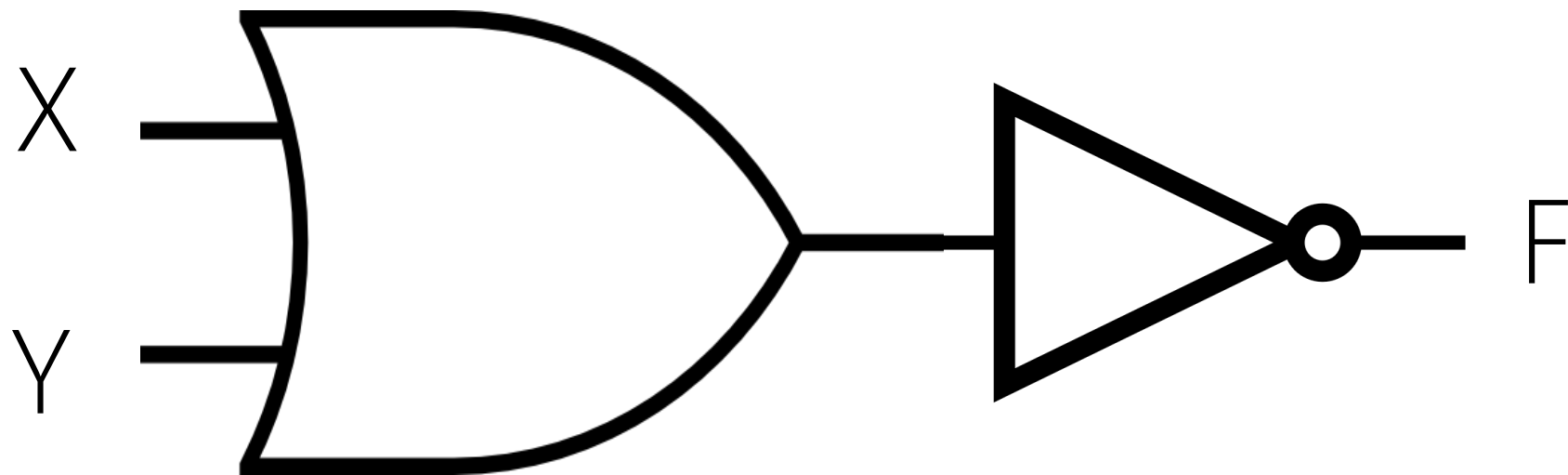
~~*~~

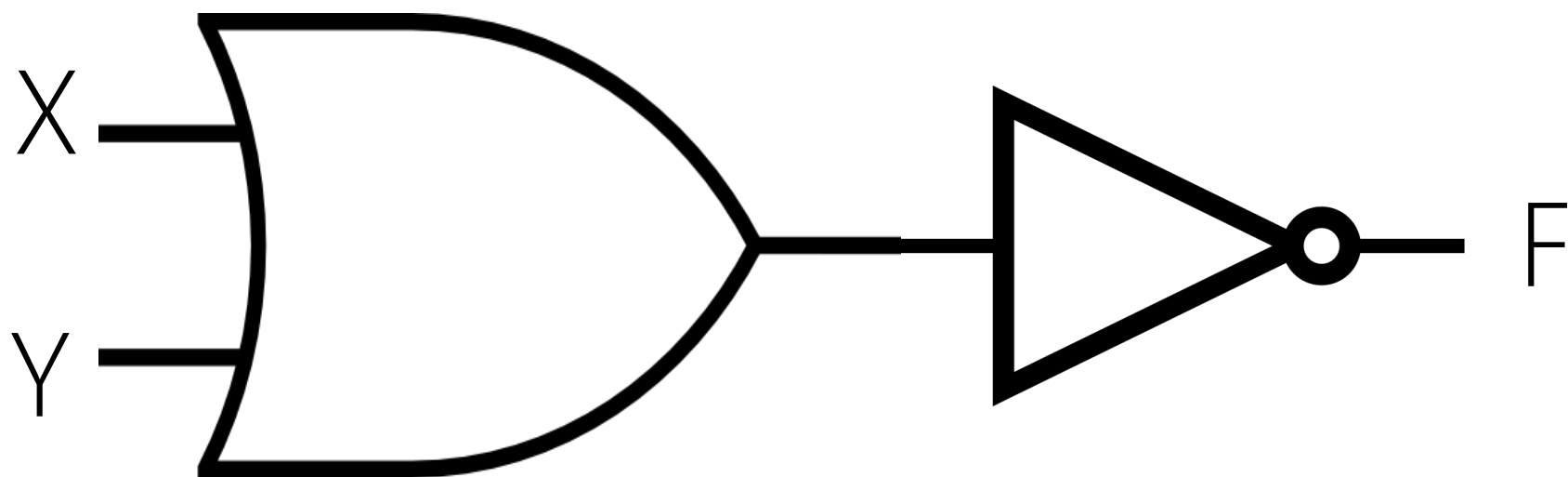
$$Z \uparrow (Y \uparrow X) \stackrel{?}{=} (Z \uparrow Y) \uparrow X$$

$$F = (Z(YX))' = Z \uparrow (Y \uparrow X)$$

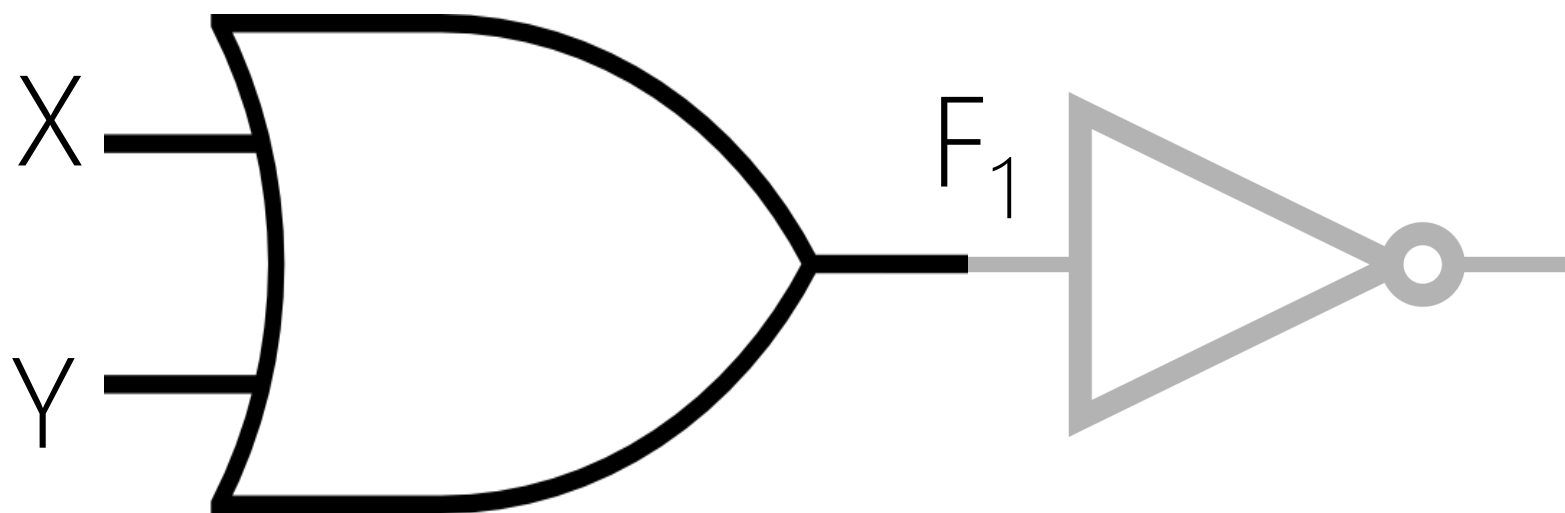
Z	Y	X	
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

ANALYSIS

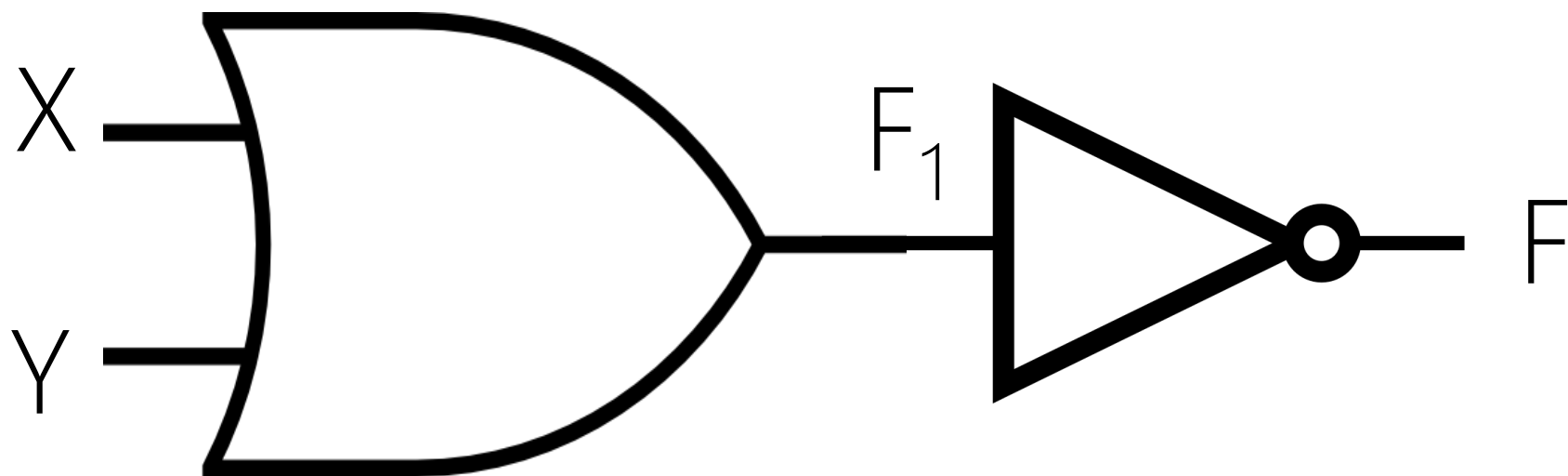




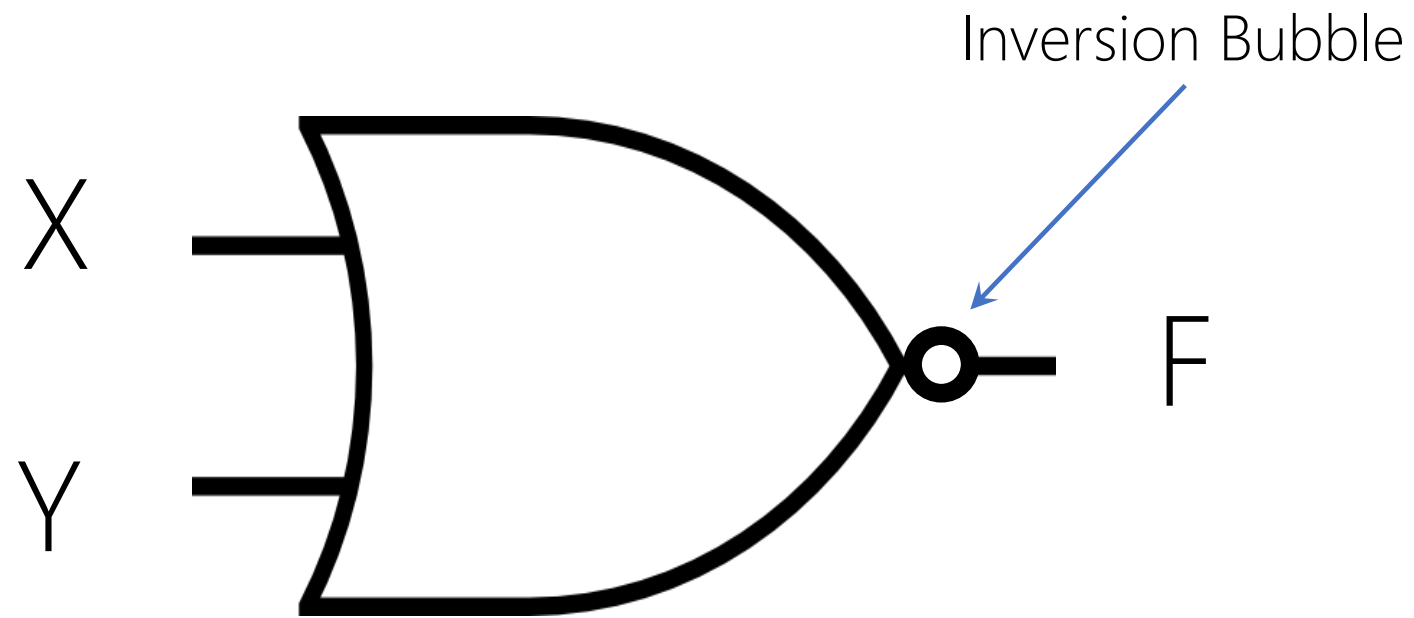
Y	X	F = ?
0	0	?
0	1	?
1	0	?
1	1	?



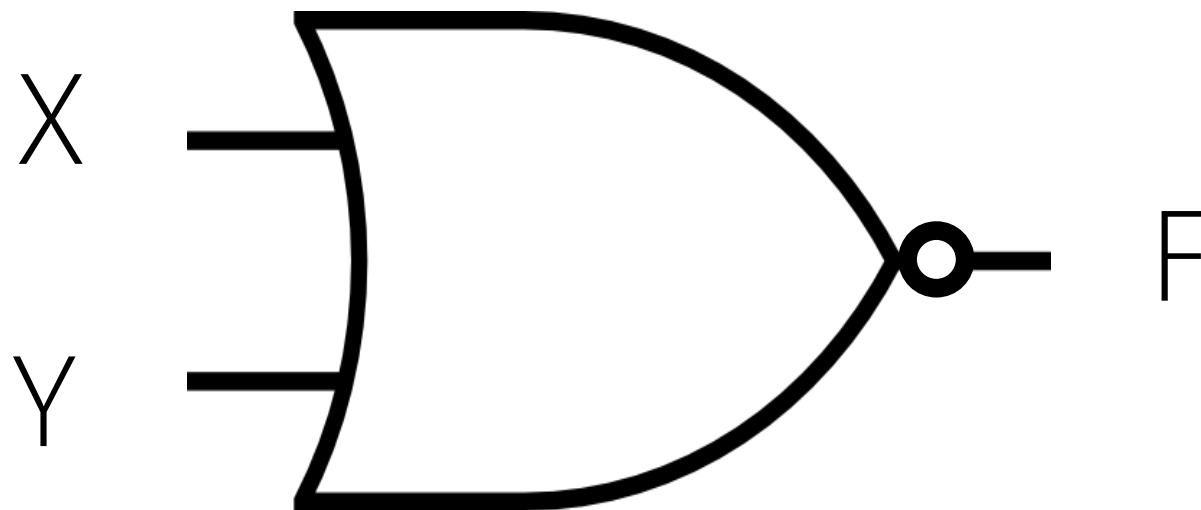
Y	X	$F_1 = Y + X$
0	0	0
0	1	1
1	0	1
1	1	1



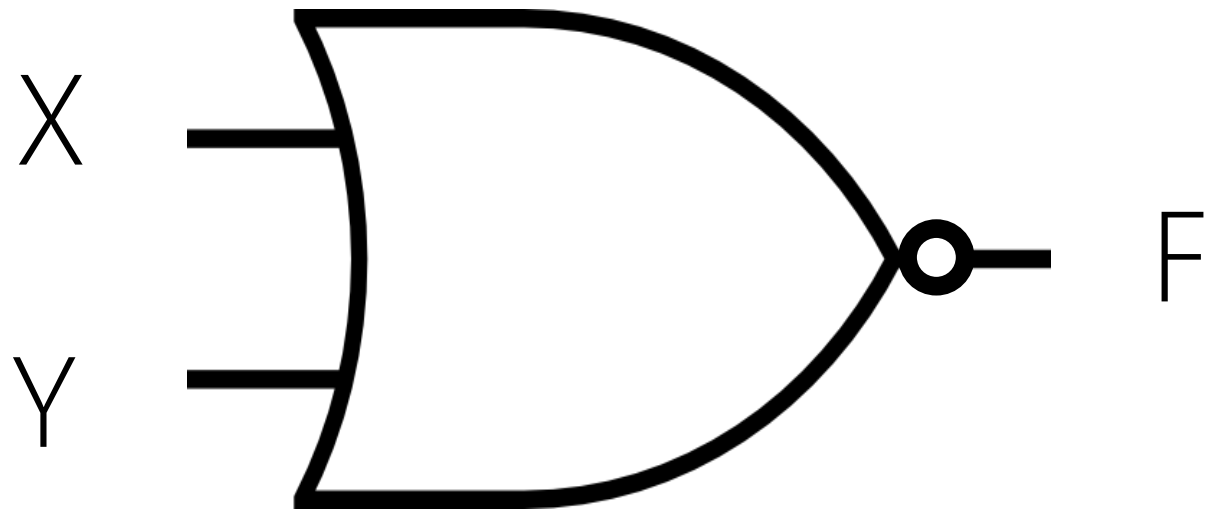
Y	X	$F_1 = Y + X$	$F = (Y + X)'$
0	0	0	1
0	1	1	0
1	0	1	0
1	1	1	0



NOR (Not – OR)

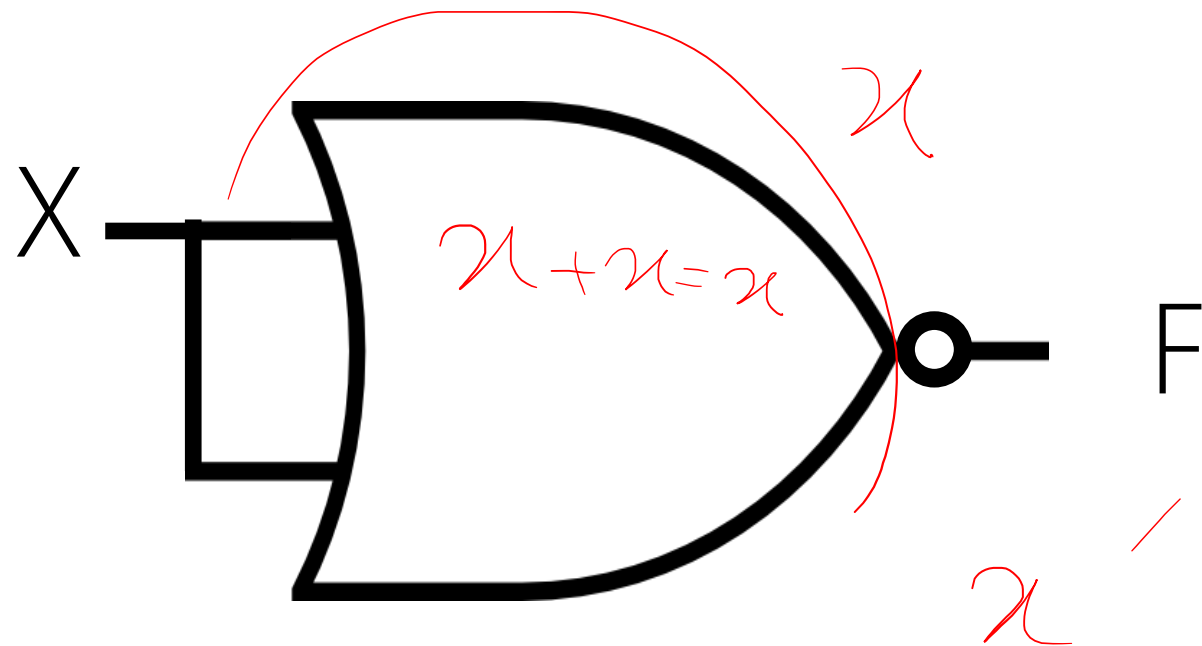


Y	X	$F = (Y + X)'$	$F = Y \downarrow X$
0	0	1	
0	1	0	
1	0	0	
1	1	0	

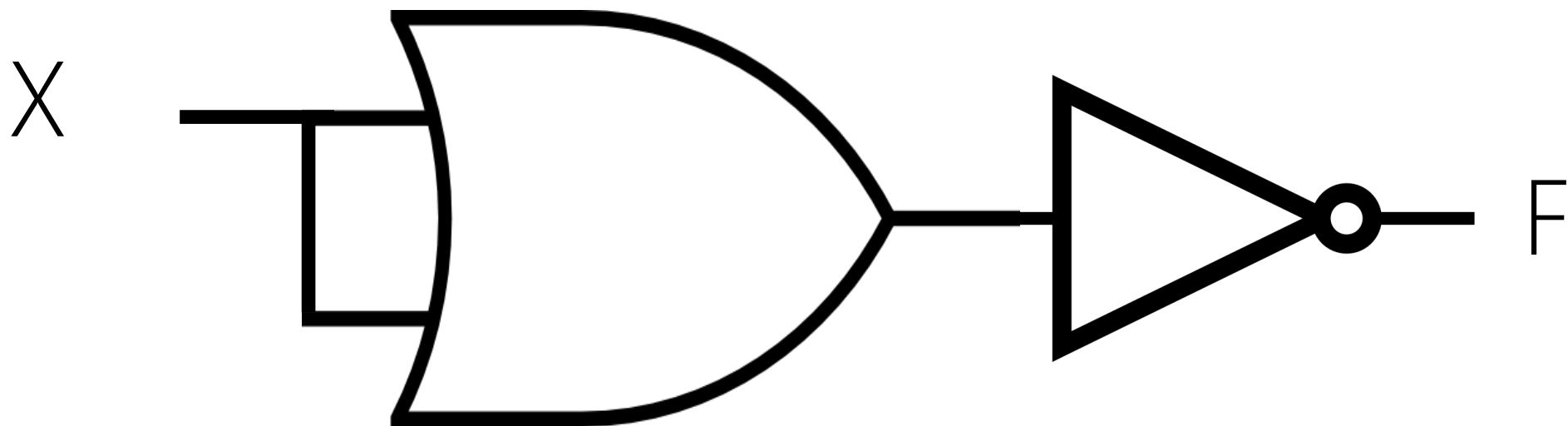


$$F = (Y + X)' = (X + Y)' = Y \downarrow X = X \downarrow Y$$

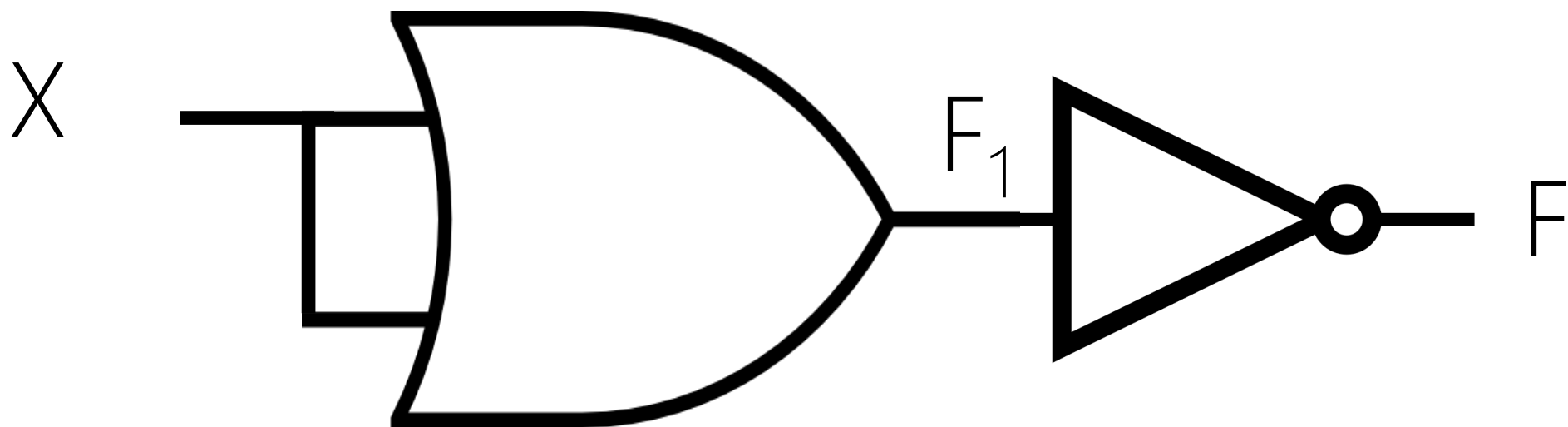
Commutative



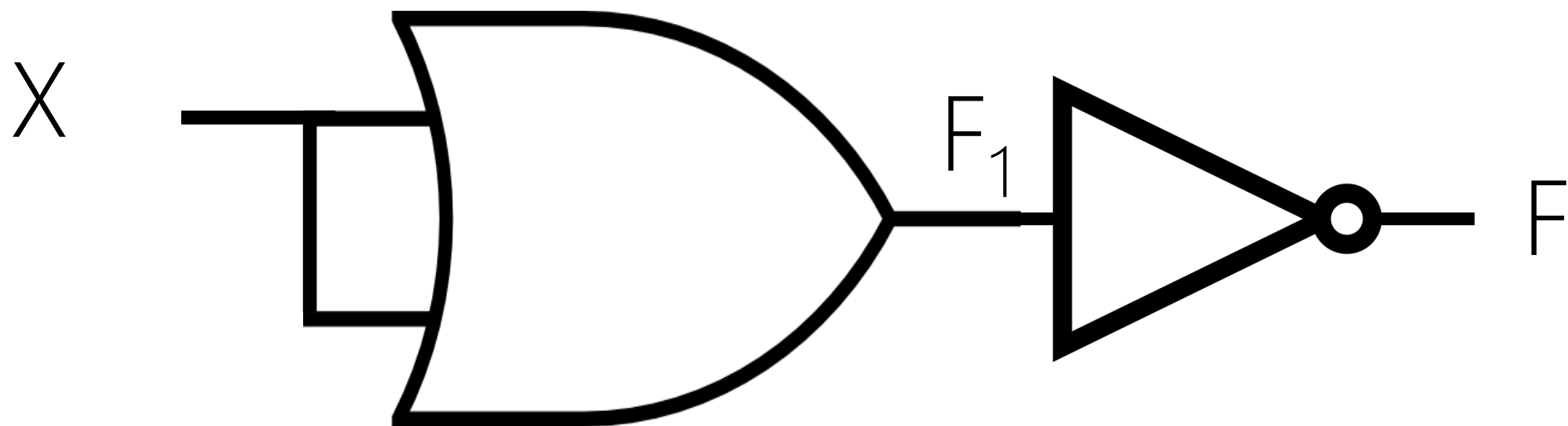
F = ?



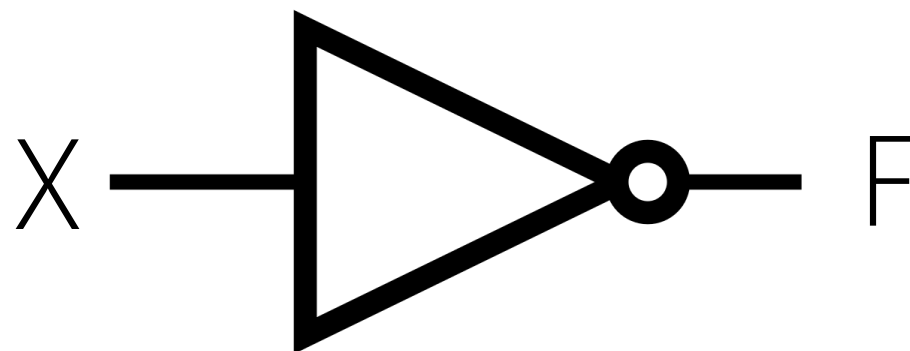
X	F = ?
0	
1	



X	$F_1 = X + X$	$F = ?$
0	0	
1	1	



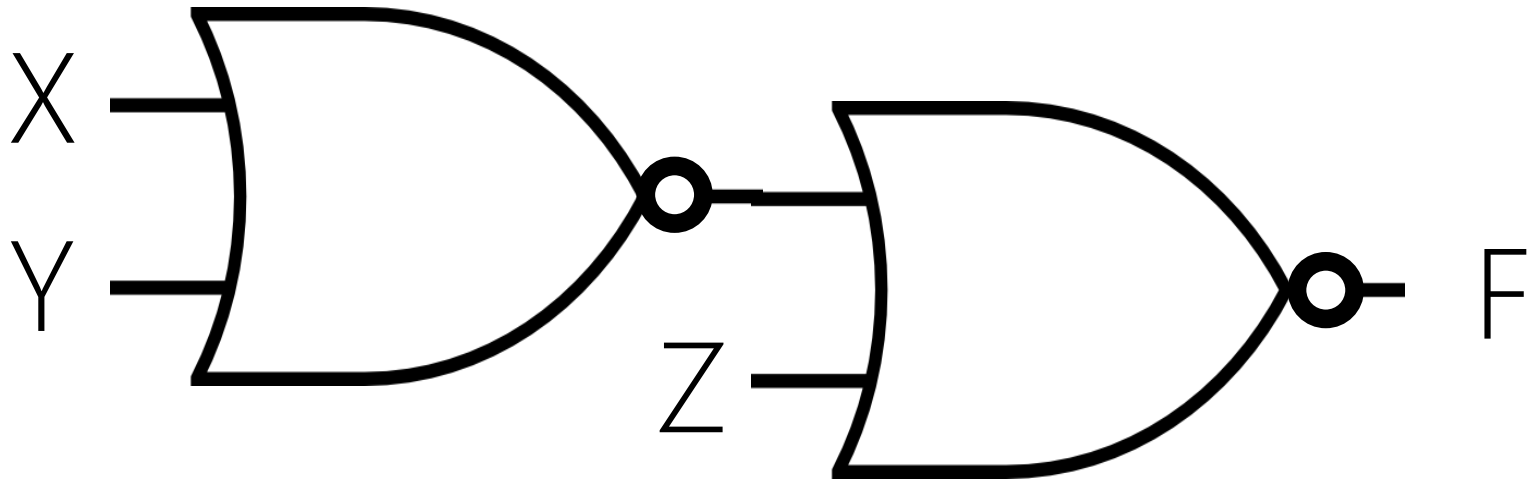
X	$F_1 = X + X$	$F = (X + X)'$
0	0	1
1	1	0



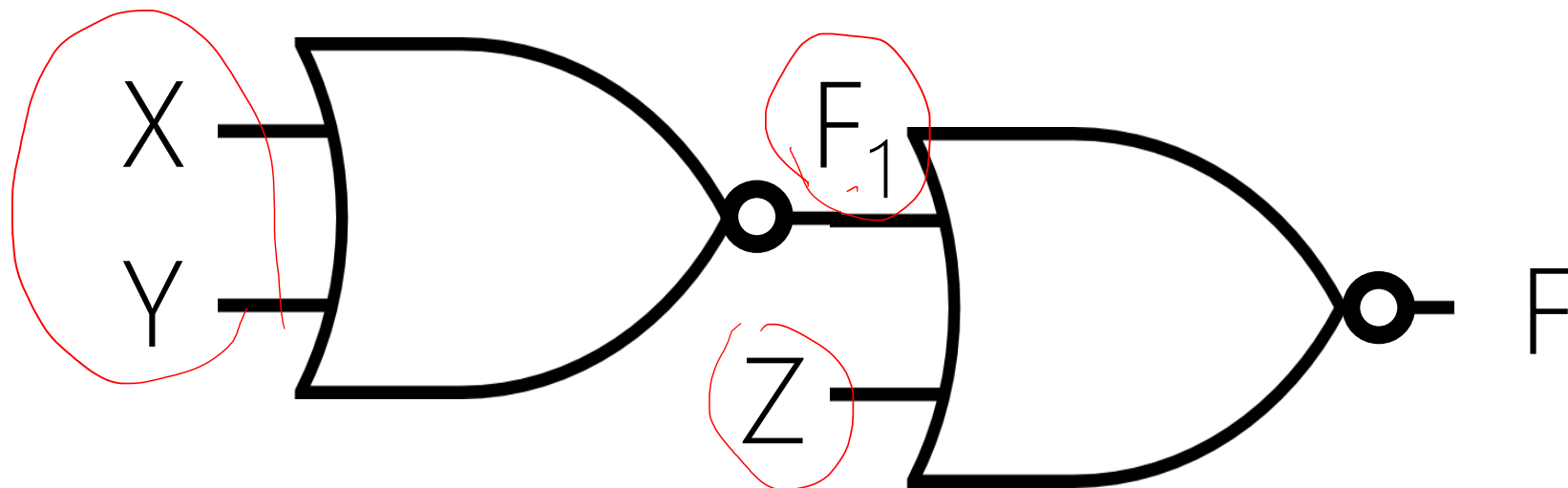
X	$F = (X+X)' = \underline{X \downarrow X} = X'$
0	1
1	0

3-INPUT NOR

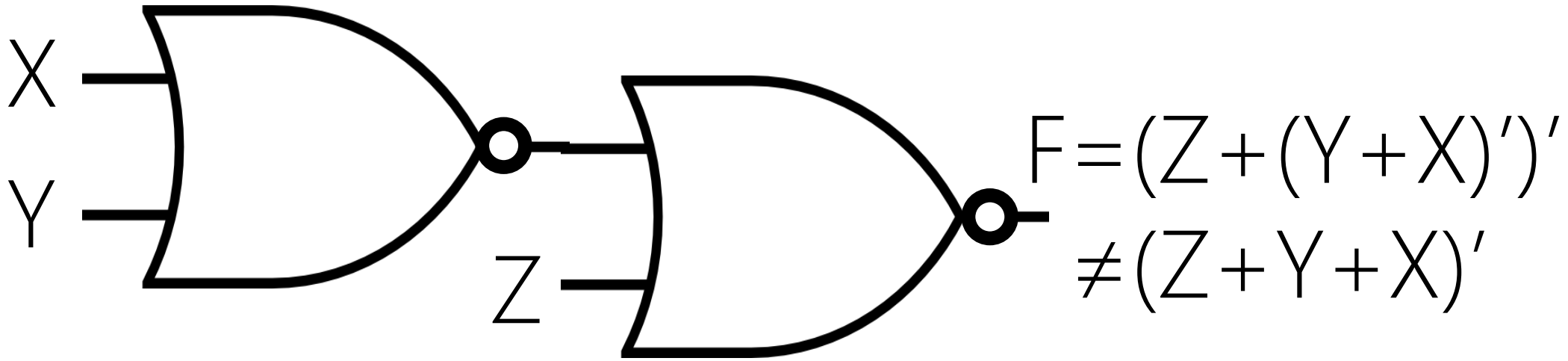
Z			Y			X			$F = (Z + Y + X)'$		
0	0	0	0	0	0	0	0	0	1		
0	0	1	0	0	1	0	0	1	0		
0	1	0	0	1	0	0	0	1	0		
0	1	1	0	1	1	0	0	1	0		
1	0	0	1	0	0	0	1	1	0		
1	0	1	1	0	1	0	1	1	0		
1	1	0	1	1	0	0	1	1	0		
1	1	1	1	1	1	0	1	1	0		



Z	Y	X	F = ?
0	0	0	
0	0	1	
0	1	0	
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	



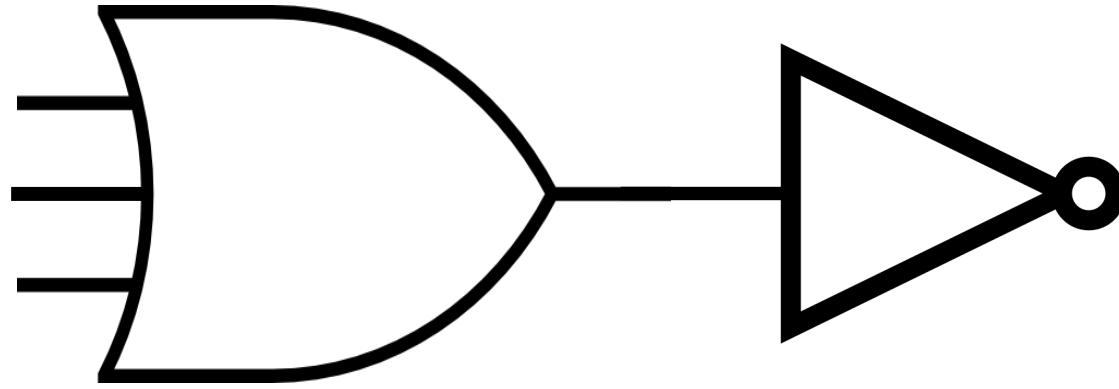
Z	Y	X	$F_1 = (Y+X)'$	$F = (Z+F_1)' = (Z+(Y+X)')'$
0	0	0	1	0
0	0	1	0	1
0	1	0	0	1
0	1	1	0	1
1	0	0	1	0
1	0	1	0	0
1	1	0	0	0
1	1	1	0	0



Z	Y	X	$F = (Z + F_1)' = (Z + (Y + X)')'$	$F = (Z + Y + X)'$
0	0	0	0	1
0	0	1	1	0
0	1	0	1	0
0	1	1	1	0
1	0	0	0	0
1	0	1	0	0
1	1	0	0	0
1	1	1	0	0

NOT (3-INPUT OR)

X
Y
Z

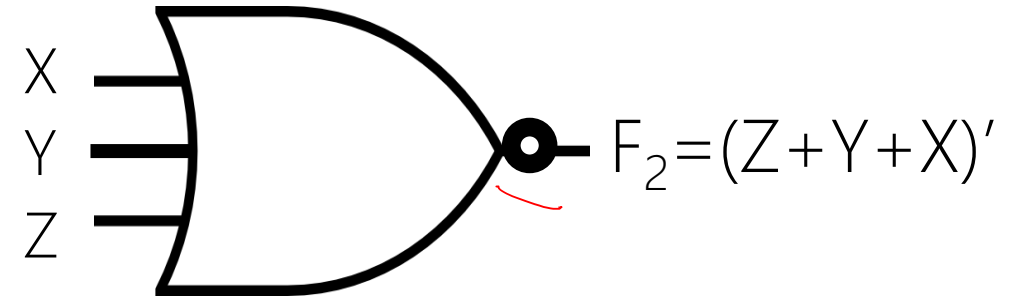
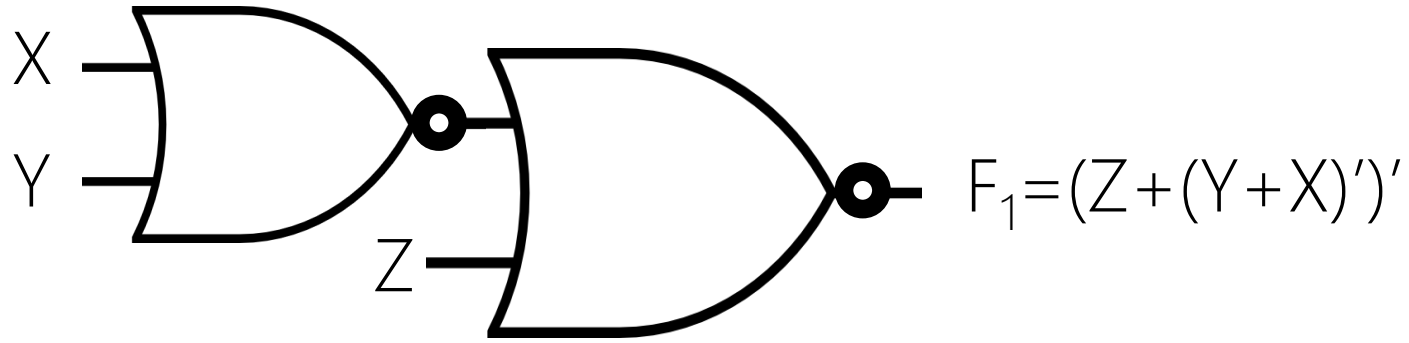


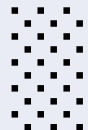
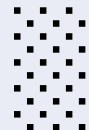
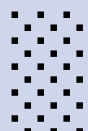
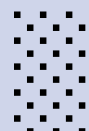
$$F = (Z + Y + X)'$$

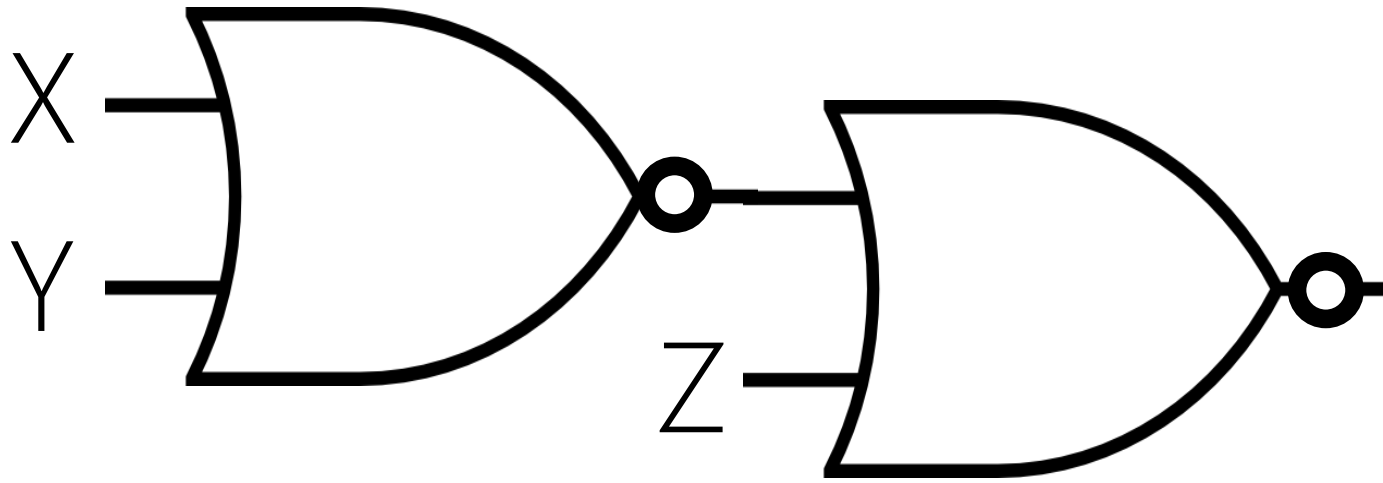
$$F = (X + Y + Z)'$$

Associative

Z	Y	X	$F = (Z + Y + X)'$	$F = (X + Y + Z)'$
0	0	0	1	1
0	0	1	0	0
0	1	0	0	0
0	1	1	0	0
1	0	0	0	0
1	0	1	0	0
1	1	0	0	0
1	1	1	0	0



$F = (Z + Y + X)'$	F_1	F_2
Effective (True)	No!	Yes
Efficient (Fast)		
Min. Cost		



$$(\overline{Z} \downarrow \overline{Y}) \downarrow X$$

$$F = Z \downarrow (Y \downarrow X)$$

Associative?

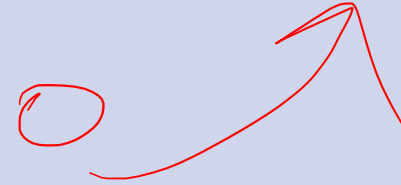
Z	Y	X	$F = (Z(YX)')' = Z \downarrow (Y \downarrow X)$
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

RECAP

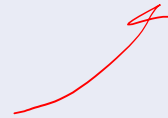
GATE

WHEN $F=1$

NOT



AND



OR

NAND



NOR



GATE	WHEN F=1
NOT	The input is 0
AND	All the inputs are 1
OR	At least one input is 1
NAND	At least one input is 0
NOR	All the inputs are 0



DESIGN

a design algorithm for any digital units (logic circuits), given truth table

1. minterm

aka. Standard Product

1 → 0 X' vs. X

1

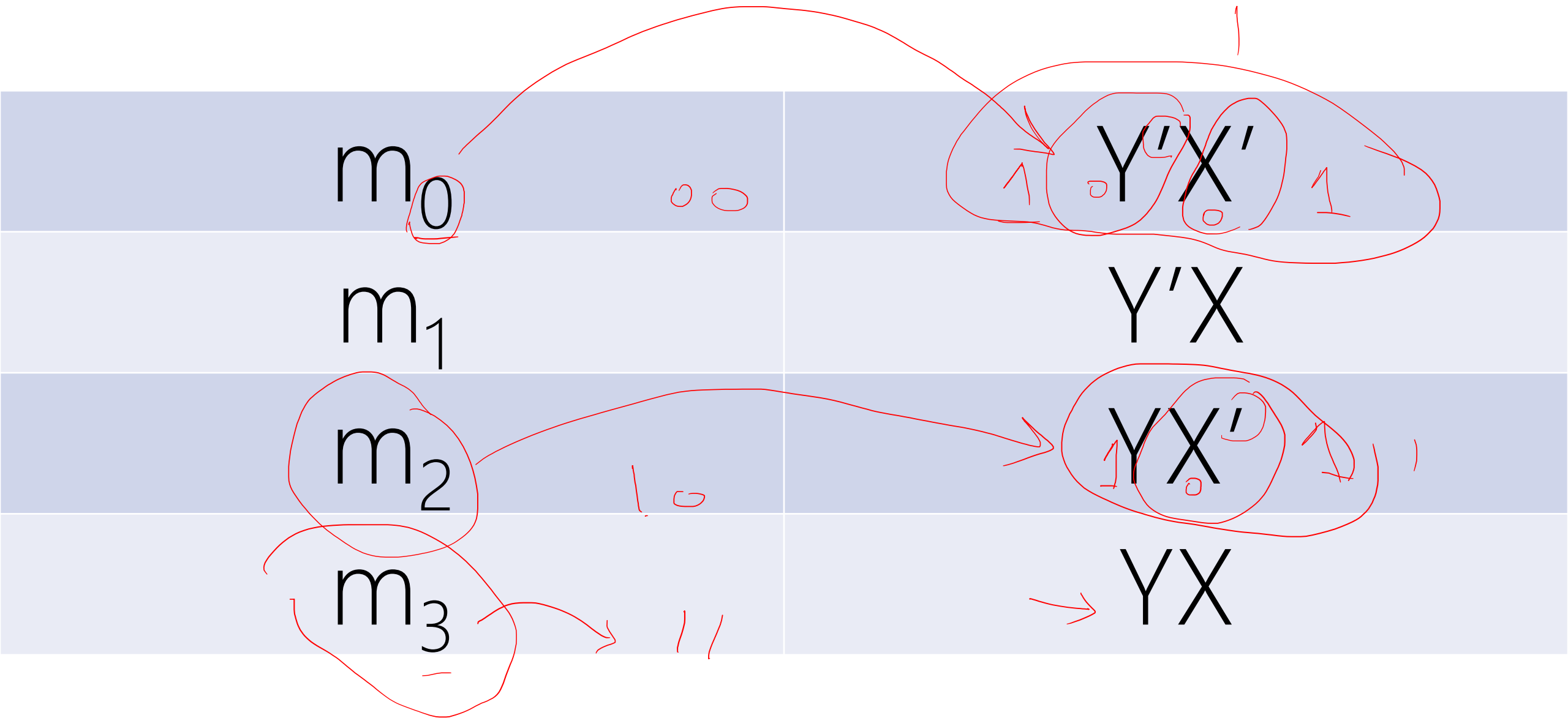
1 binary variable appear either:

- in its normal form X , or
- in its complement form X'



YX vs. YX' vs. $Y'X$ vs. $Y'X'$

2 binary variables appear either in one of these forms:



ZYX vs. ZYX' vs. ...

3 binary variables appear either in one of these forms: how many?

ZYX vs. ZYX' vs. ...

3 binary variables appear either in one of these forms: how many?

Each variable can take 2 forms (normal and complement)

We have 3 variables, $2 \times 2 \times 2 = 2^3 = 8$

m_0

$Z'Y'X'$

m_1

$Z'Y'X$

m_2

$Z'YX'$

m_3

$Z'YX$

m_4

$ZY'X'$

m_5

$ZY'X$

m_6

ZYX'

m_7

ZYX

$$A_{n-1} \cdots A_2 A_1 A_0 \text{ vs. } A_{n-1} \cdots A_2 A_1 A'_0 \dots$$

n binary variables appear either in one of these forms: how many?

Each variable can take 2 forms (normal and complement)

We have n variables, $2 \times 2 \times 2 \times \cdots \times 2 = 2^n$

m_0	$A'_{n-1} \cdots A'_2 A'_1 A'_0$
m_1	$A'_{n-1} \cdots A'_2 A'_1 A_0$
m_2	$A'_{n-1} \cdots A'_2 A_1 A'_0$
\vdots	\vdots
\vdots	\vdots
m_{2^n-3}	$A_{n-1} \cdots A_2 A'_1 A_0$
m_{2^n-2}	$A_{n-1} \cdots A_2 A_1 A'_0$
m_{2^n-1}	$A_{n-1} \cdots A_2 A_1 A_0$

2. TRUTH TABLE

en.wikipedia.org/wiki/Truth_table



<u>X</u>	<u>F</u> = F(X) = ?
→ 0	} ?
→ 1	} ?

X	$F = F(X) = 0$
0 }	{ 0
1	{ 0

X	$F = F(X) = X'$
0	1
1	0

m_0

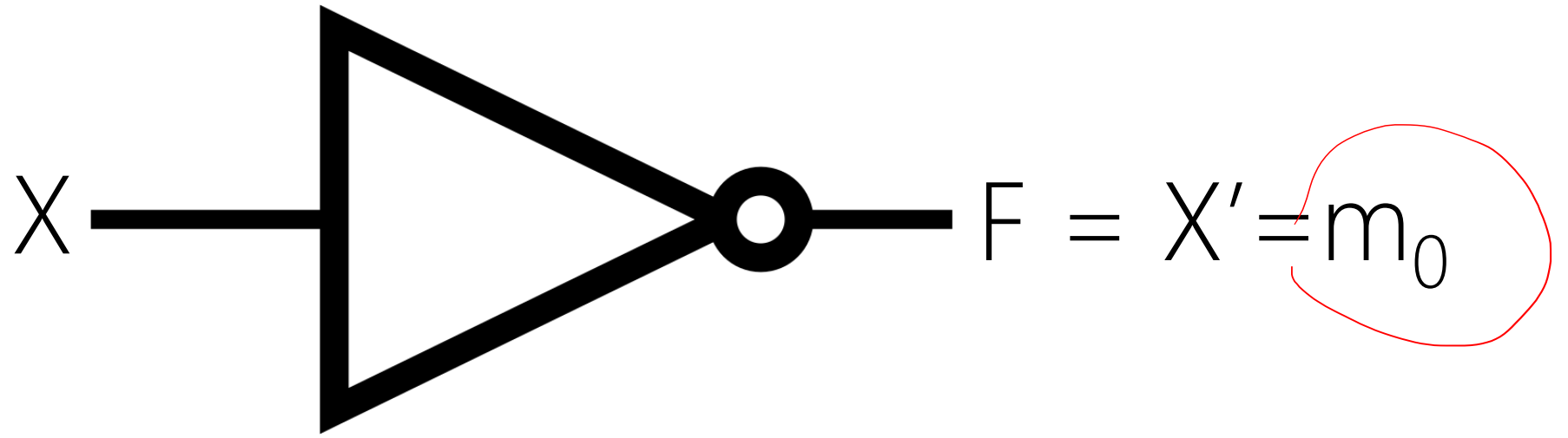
X'

m_1

X

X	$F = F(X) = X' = m_0$
0	1
1	0

X	$F = F(X) = X' = m_0$
0	1
1	0

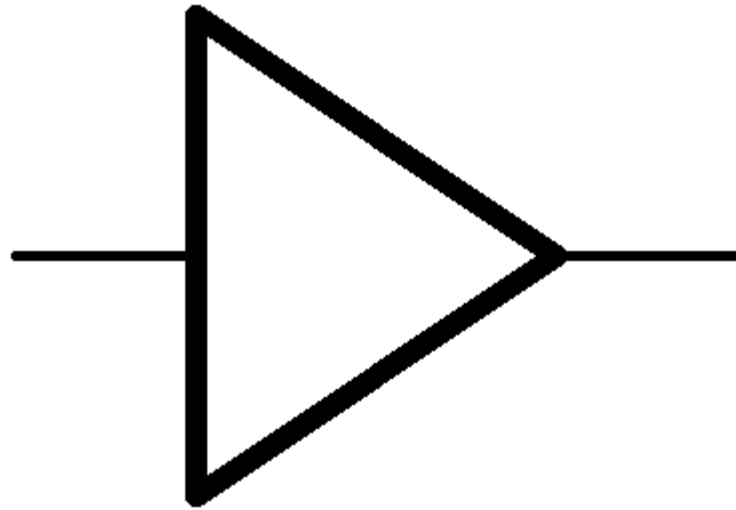


X	F = F(X) = X
0	0
1	1

X	$F = F(X) = X = m_1$
0	0
<u>1</u>	1

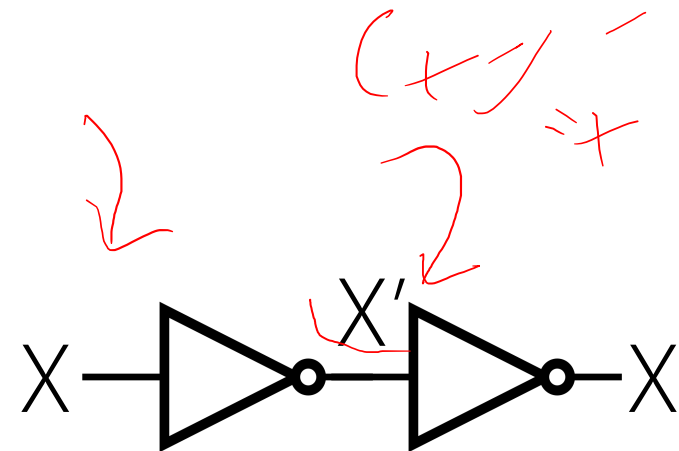
X	$F = F(X) = X = m$
0	0
1	1

X



$$F = X = m_1$$

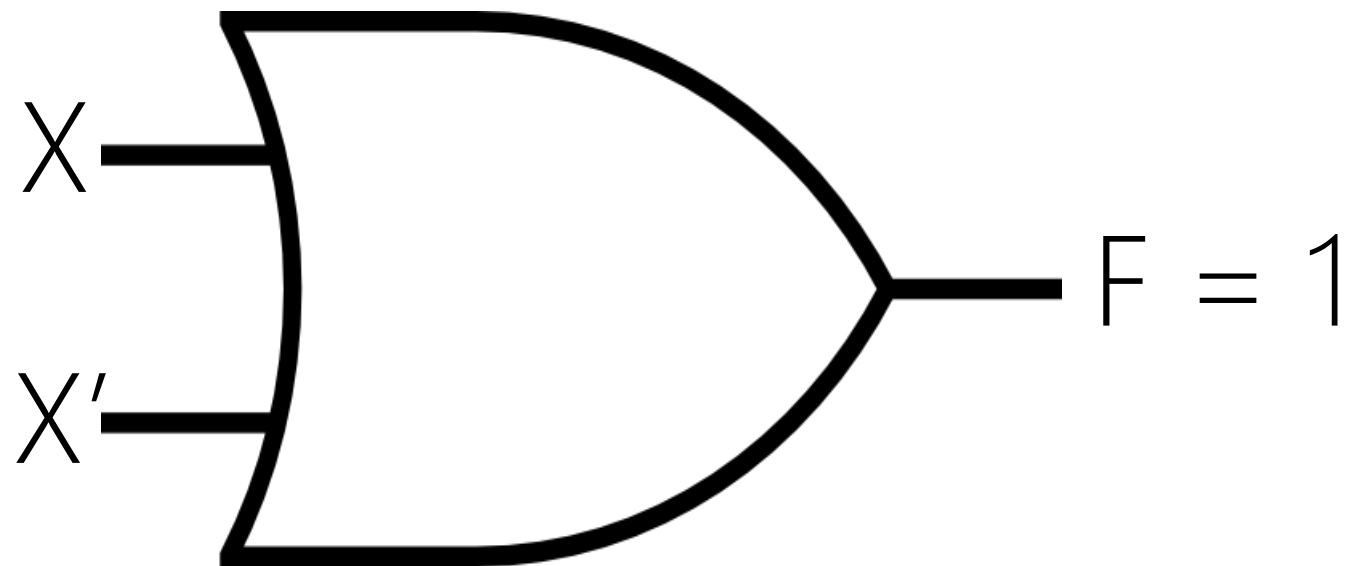
Digital Buffer



X	$F = F(X)$
0	1
1	1

X	$F = F(X) = X' = m_0$
0	1
1	1

X	$F = F(X) = X' + X = m_0 + m_1$
0	1
1	1



X'	X	X'+X
1	0	1
0	1	1

$$F = X + X' = 1$$

TRUTH TABLE \leftrightarrow minterm
