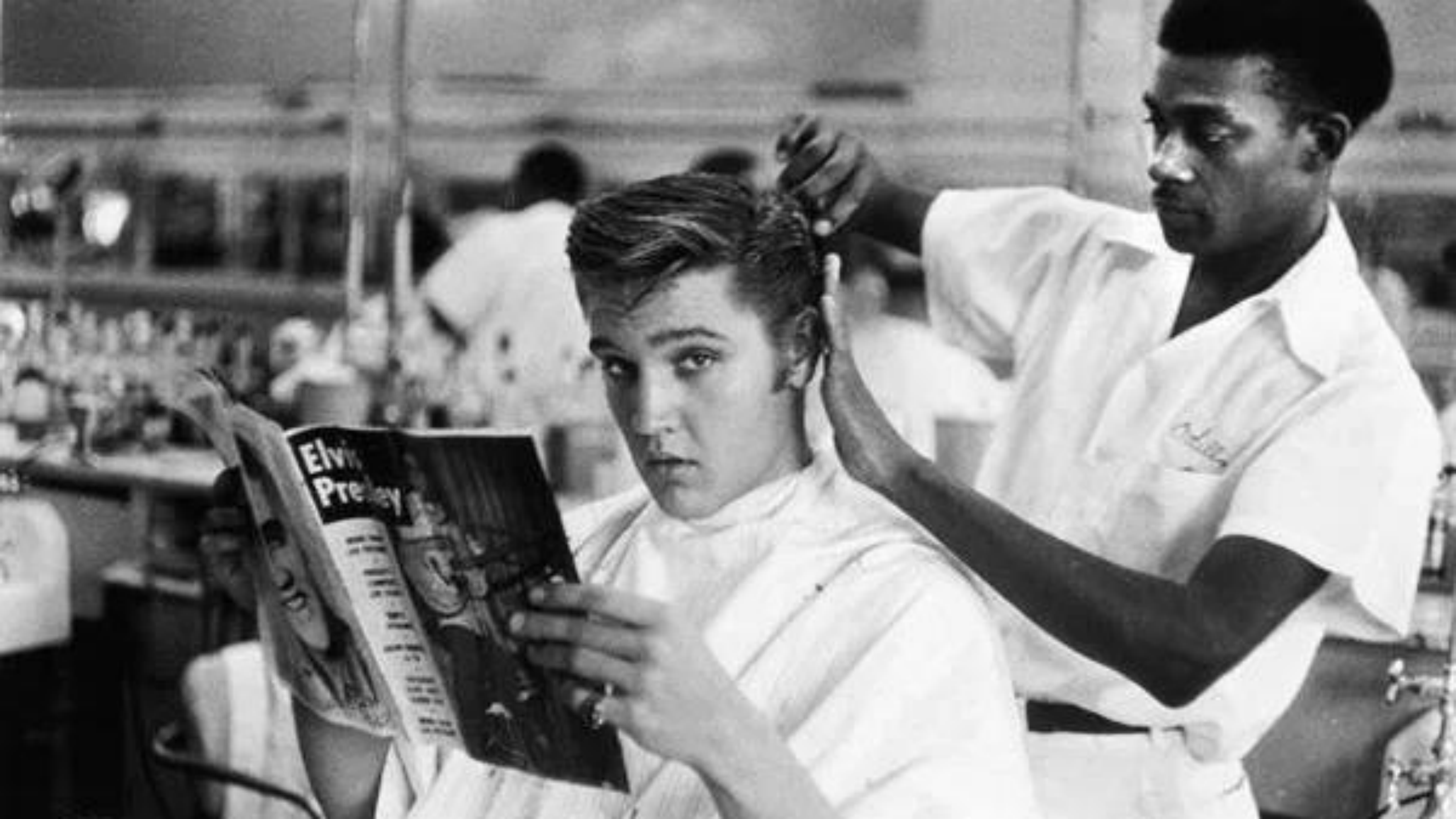

Elvi

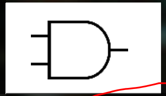
Lab07 & Lec07

Informal Survey on Course Delivery

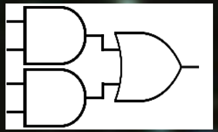
https://forms.gle/ra95dP8sMigfGP18A


W2022: A Digital Odyssey
The Ultimate Trip.

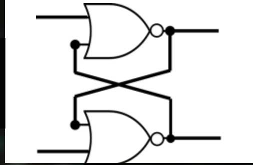
Number Systems | $(12)_{10} \rightarrow (1100)_2$

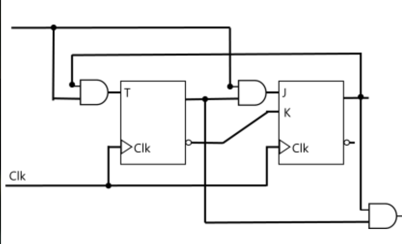
Logic Gates | 

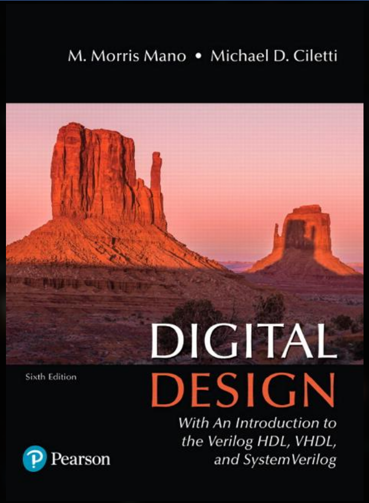
SoP

Combinational Logic | 

POS

Flip-Flop | 

Sequential Logic | 



Chapter 2: Boolean Algebra and Logic Gates
Chapter 3: Gate-Level Minimization

Review SoP & PoS

http://etc.ch/mJFi

MINIMIZATION

aka. Simplification

Same Effective Design but More Efficient

MINIMIZATION

Number of Gates
Number of Inputs (2-input vs 4-input)
Number of Interconnections
Propagation Time
Cost of Gates
Circuit Area
. . .
A circuit may not satisfy all due to conflicting constraints!

3-level

SOP
POS

2-level

MINIMIZATION

aka. Simplification

SoP (ANDs-OR) → Simplify → NAND

PoS (ORs-AND) → Simplify → NOR

MINIMIZATION

I) Boolean Algebra (algebraically)

MINIMIZATION

II) Map (Karnaugh map, K-map)

ALGEBRA

A set of elements

A set of operators

A set of axioms | postulates | assumptions | definitions

ALGEBRA

A set of elements: e.g.,

{0,1}

{-1,0,1}

{1,2,3,4,...}

{..., -2,-1,0,1,2,...}

ALGEBRA

A set of operators

$\{+, \times\}$

$\{+, -, \times, \div\}$

$\{\sim, \leq, \geq\}$

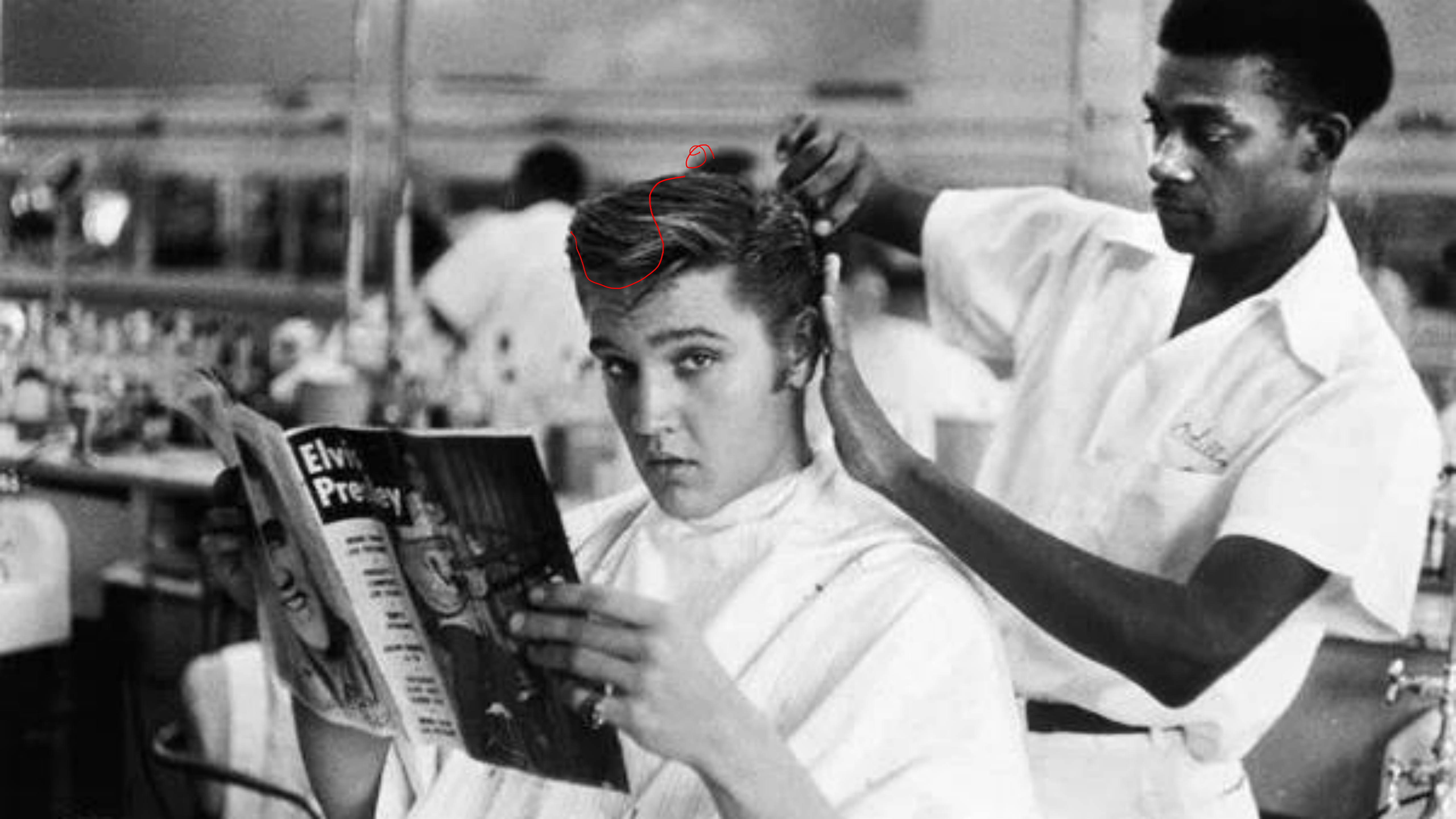
$x \; ? \; y : z$; $y = (x > 2) \; ? \; \text{'Hello'} : \text{'Bye'}$

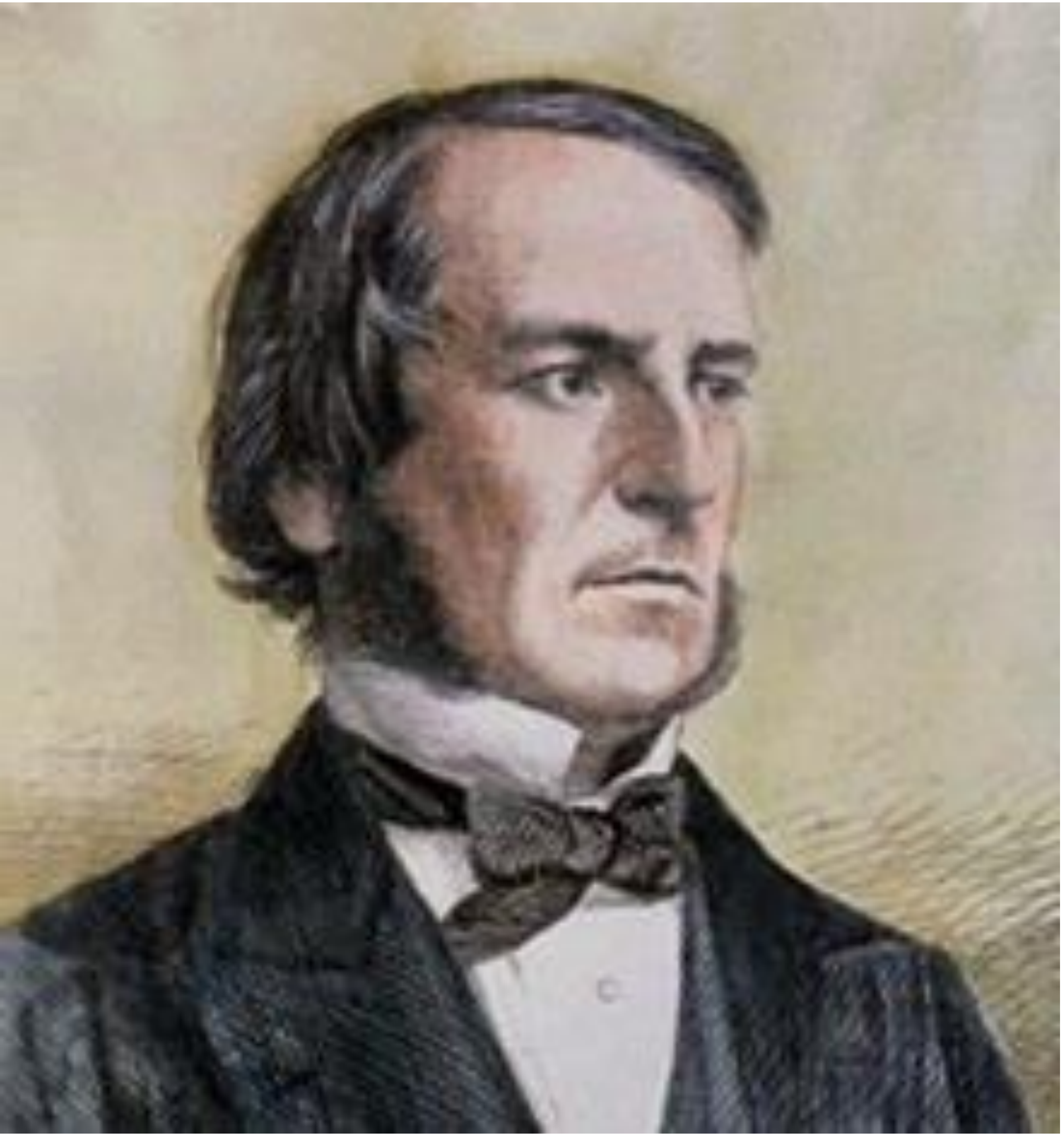
ALGEBRA

Unary: !x, -x, x′, …

Binary: x+y, x÷y, x^y, …

Ternary: x?y:z (Elvis), x BETWEEN y AND z (SQL)





George Boole (/buːl/)

Mathematician
Philosopher
Logician

The Laws of Thought (1854)

Boolean Algebra!


Edward Vermilye Huntington
1874 – 1952
In 1904, he put Boolean algebra
on a sound axiomatic foundation

POSTULATE

aka. Axiom, Assumption

https://en.wikipedia.org/wiki/Axiom

A statement that is taken to be TRUE

Serve as a premise or starting point for further reasoning

I. CLOSURE

A set is closed with respect to an operator § if the result of § ∈ S:

Unary: $x \in S$: $\S x \in S$

Binary: $x, y \in S$: $x \S y \in S$

I. CLOSURE

S = {..., -2,-1,0,1,2,...}

§ = + (addition), - (subtraction), × (multiplication), – (negation), ! (fact)

S is closed with respect to +, -, ×, –, !

I. CLOSURE

$S = \{\dots, -2,-1,0,1,2,\dots\}$

$\S =$ ^ (power)

?

I. CLOSURE

$S = \{\ldots, -2, -1, 0, 1, 2, \ldots\}$

§ = ^ (power)

S is NOT closed with respect to ^ for $2^{(-1)} \notin S$

I. CLOSURE

$S = \{0,1\}$

$§ = +$ (OR), $\times$ (AND), $'$ (NOT)

S is closed with respect to $+$, $\times$, $'$

I. CLOSURE

$S = \{0, 1\}$

$\S = -$ (negation)

S is NOT closed with respect to $-$ for $(-1) \notin S$

II. COMMUTATIVE

A binary operator § on a set S is commutative iff for all x,y ∈ S:

$$x \mathbin{\S} y = y \mathbin{\S} x$$

(x § y may or may not be in S)

II. COMMUTATIVE

$S = \{\ldots, -2,-1,0,1,2,\ldots\}$

§ = + (addition), × (multiplication)

+ and × are commutative on S for $x+y=y+x$ and $x \times y = y \times x$

II. COMMUTATIVE

$S = \{\dots, -2, -1, 0, 1, 2, \dots\}$

$§ = -$ (subtraction)

- is NOT commutative on S for $x - y \neq y - x$

II. COMMUTATIVE

S = {0,1}

§ = + (OR), × (AND), ↑ (NAND), ↓ (NOR), ⊕ (XOR), ⊙ (XNOR)

All are commutative on S for x § y=y § x

III. ASSOCIATIVE

A binary operator § on a set S is associative iff for all x,y,z ∈ S:

x § (y § z)= (x § y) § z = x § y § z

III. ASSOCIATIVE

$S = \{\ldots, -2,-1,0,1,2,\ldots\}$

§ = + (addition), × (multiplication), - (subtraction)

+, -, × are associative on S for

$x\pm(y\pm z)=(x\pm y) \pm z=x\pm y\pm z$

$x\times(y\times z)=(x\times y)\times z=x\times y\times z$

$\forall x, y, z \quad (x^y)^z$

$\exists (x, y, z) \Longrightarrow \quad \stackrel{?}{=}$

$x^{(y^z)}$

III. ASSOCIATIVE

S = {…, -2,-1,0,1,2,…}

§ = ^ (power)

?

$x = 2 \quad (2^1)^3 = 8$

$y = 1 \Longrightarrow \quad \neq$

$z = 3 \quad 2^{(1^3)} = 2$

III. ASSOCIATIVE

$S = \{\dots, -2, -1, 0, 1, 2, \dots\}$

§ = ^ (power)

^ is NOT associative on S for $2\wedge(1\wedge 3) \neq (2\wedge 1)\wedge 3$

$(x+[y)+z]$

III. ASSOCIATIVE

$S = \{0,1\}$

§ = + (OR), × (AND)

+,× are associative on S for

$$x+(y+z)=(x+y)+z=x+y+z$$

$$x\times(y\times z)=(x\times y)\times z=x\times y\times z$$

IV. DISTRIBUTIVE

If § and † are two binary operators on a set S, § is distributive over † iff:

Left Distributivity: x § (y † z) = (x § y) † (x § z)
Right Distributivity: (y § x) † (z § x) = (y † z) § x

IV. DISTRIBUTIVE

If § and † are two binary operators on a set S, § is distributive over † iff:

If § Commutative: x § (y † z) = (x § y) † (x § z) = (y † z) § x

IV. DISTRIBUTIVE

$S=\{\ldots, -2, -1, 0, 1, 2, \ldots\}$

$+$ (Addition), $\times$ (Multiplication)

1 $x + (y \times z) \overset{?}{\Leftrightarrow} (x + y) \times (x + z) \overset{?}{\Leftrightarrow} (y \times z) + x$

2 $x \times (y + z) \overset{?}{\Leftrightarrow} (x \times y) + (x \times z) \overset{?}{\Leftrightarrow} (y + z) \times x$

IV. DISTRIBUTIVE

$S=\{\ldots, -2, -1, 0, 1, 2, \ldots\}$

$+$ (Addition), $\times$ (Multiplication)

$$x + (y \times z) \neq (x + y) \times (x + z) \neq (y \times z) + x$$

$$x \times (y + z) = (x \times y) + (x \times z) = (y + z) \times x$$

IV. DISTRIBUTIVE

$S=\{0,1\}$

$+$ (OR), $\times$ (AND)

$$x + (y \times z) \overset{?}{\Leftrightarrow} (x + y) \times (x + z) \overset{?}{\Leftrightarrow} (y \times z) + x$$

$$x \times (y + z) \overset{?}{\Leftrightarrow} (x \times y) + (x \times z) \overset{?}{\Leftrightarrow} (y + z) \times x$$

V. IDENTITY

$e \in S$ is an identity element w.r.t binary operator § iff for all $x \in S$:

$x \mathbin{\S} e = x = e \mathbin{\S} x$

V. IDENTITY

$S = \{\ldots, -2, -1, 0, 1, 2, \ldots\}$

$\S = +$ (addition), $\times$ (multiplication)

$e_+ = 0 : x+0=0+x=x$

$e_\times = 1 : x\times 1=1\times x=x$

V. IDENTITY

$S = \{\ldots, -2,-1,0,1,2,\ldots\}$

$§ = -$ (subtraction)

$e = ?$ 0

$x - 0 = x$

$\neq$

$0 - x = -x$

V. IDENTITY

S = {..., -2,-1,0,1,2,...}

§ = - (subtraction)

e = 0

$x-0 \neq 0-x \neq x$

NO identity element for – (subtraction) in S

V. IDENTITY

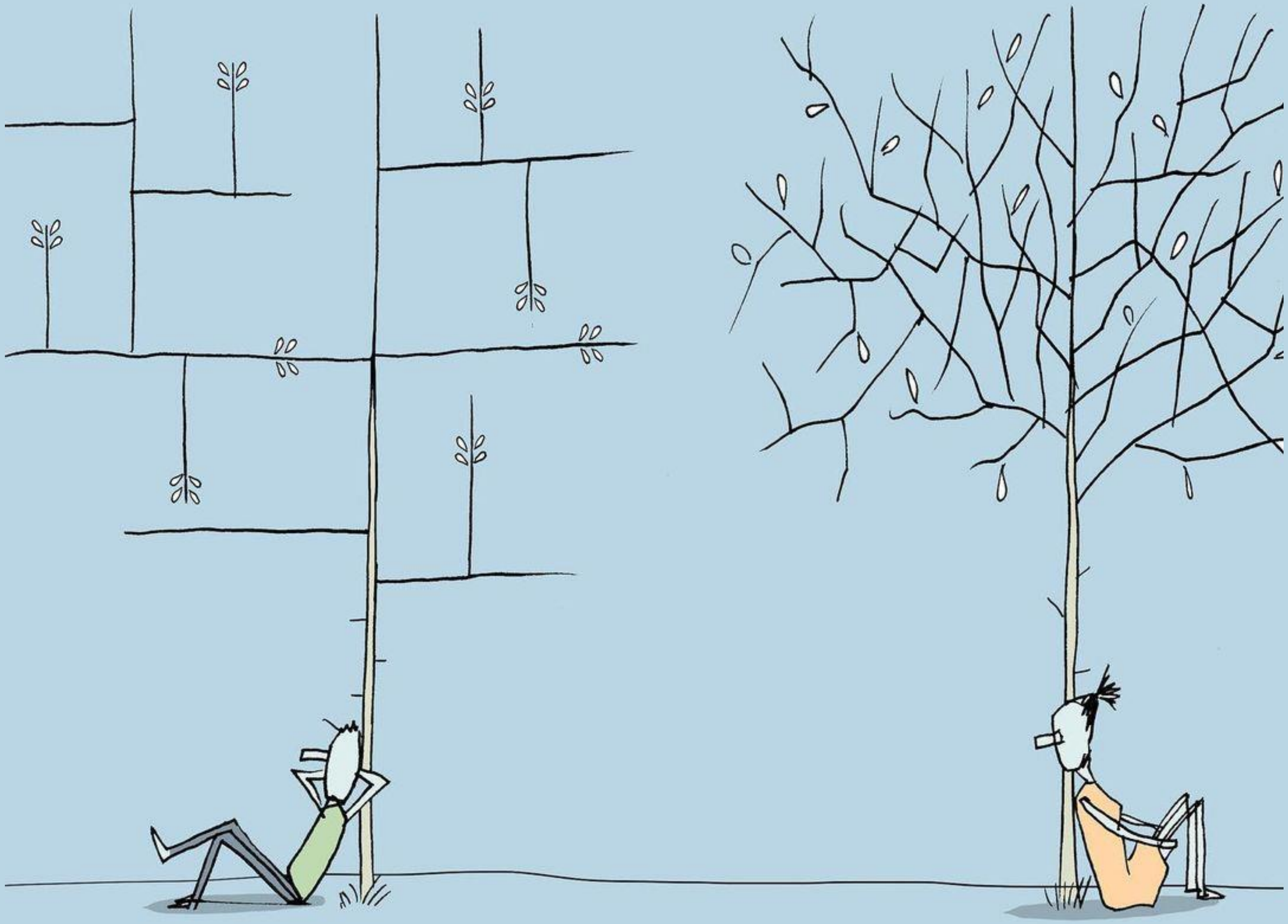
$S = \{0,1\}$

$\S = +\ (\text{OR}),\ \times\ (\text{AND})$

$e_{+} = 0 : x+0=0+x=x$

$e_{\times} = 1 : x\times 1=1\times x=x$

https://www.yuvalrob.com/



$x \rightarrow y$

$x^{-1} = y$

VI. INVERSE

$y^{-1} = x$

For all x∈S, there should be y∈S w.r.t binary operator § iff:

$x \S y = e_{\S} = y \S x$

We denote $y = x^{-1}$ and $x = y^{-1}$

VI. INVERSE

2 → −2

= 0

$S = \{\dots, -2, -1, 0, 1, 2, \dots\}$

$\S = +$ (addition)

$x + (-x) = (-x) + x = e_+ = 0$

$x^{-1} = -x$

$2 \times \frac{1}{2} = e_\times = 1$

$\frac{1}{2} \notin S$

VI. INVERSE

$S = \{\ldots, -2,-1,0,1,2,\ldots\}$

$\S = \times$ (multiplication)

S does not have the inverse property for $\times$ since

$2\times\frac{1}{2} = \frac{1}{2}\times 2 = e_\times = 1$ but $\frac{1}{2} \notin S$

$$0 + \overset{0}{?} = e_+ = 0$$

$$1 + \cancel{?} = e_+ = 0$$

VI. INVERSE

$$S = \{0,1\}$$

$$\S = + \text{ (OR)}, \times \text{ (AND)}$$

?

VI. INVERSE

$S = \{0,1\}$

$\S = +\ (\mathrm{OR}),\ \times\ (\mathrm{AND})$

$1+?=?+1 \neq e_{+} = 0$

$0\times ?=?\times 0 \neq e_{\times} = 1$

$x^{-1} = y \longrightarrow \S$

$x' = y$ → § †

VII. COMPLEMENT

For all x∈S, there should be y∈S w.r.t binary operators § and † iff:

$$x \,\S\, y = e_{\dagger} = y \,\S\, x$$

$$x \dagger y = e_{\S} = y \dagger x$$

We denote $y = x'$ and $x = y'$

$0' = 1$

VII. COMPLEMENT (NOT)

$S = \{0,1\}$

$\S = +\ (\text{OR}),\ \times\ (\text{AND})$

$1+0=0+1 = e_\times = 1$

$0\times1=1\times0 = e_+ = 0$

$0' = 1,\ 1' = 0$

BOOLEAN ALGEBRA

- A set S with at least two elements x, y and x ≠ y.
- Two binary operators § and †

- S is closed, commutative, distributive, and complement w.r.t § , †
- $e_§$ and $e_†$ exist

Claude Elwood Shannon

Mathematician
Electrical Engineer
Cryptographer

M.Sc. Thesis (1937)
A Symbolic Analysis of Relay and Switching Circuits

$x = 0$
$y = 1$

Switching Algebra!
2-valued Boolean algebra

SWITCHING ALGEBRA

- $S = \{0,1\}$
- $\S = \times$ (AND), $\dagger = +$ (OR)

- S is closed, commutative, distributive, complement w.r.t $\times$, $+$
- $e_\times = 1$ and $e_+ = 0$

$$0 \times 1 = e_+ = 0$$
$$0 + 1 = e_\times = 1$$
$$0' = 1;\ 1' = 0$$

SWITCHING ALGEBRA
IS-A
BOOLEAN ALGEBRA

It satisfies all conditions of Boolean algebra!

Prove → Book: 2.3 axiomatic definition of Boolean algebra

Another sample of algebra in CS: Relational Algebra (SQL)

Is relational algebra a Boolean algebra? Check this when you take COMP-3150: Database Management Systems!

[in]EQUALITY PROOF

Design A =?= Design B

If A≠B, Pick the Effective Design

If A==B, Pick the Efficient Design

[in]EQUALITY PROOF

prove by truth table

For equality proof, $F_1=F_2$, for all possibility in the input variables (all rows), both side of equation must have equal value for same input variables.

For inequality proof, $F_1 \neq F_2$, find at least one possibility (a row) that have different values.

[in]EQUALITY PROOF

Prove by postulates