

CSC 330 / Spring 2026

OCaml Matching, Patterns

Ibrahim Numanagić /nu-ma-nagg-ich/

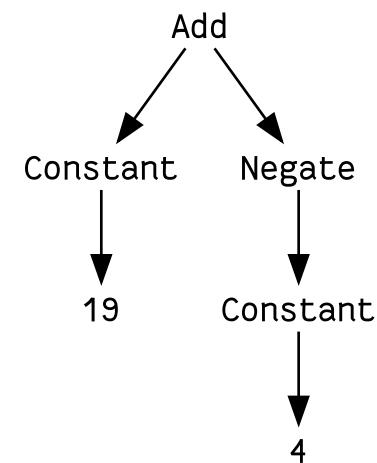
based on Dan Grossman's lecture materials

Recursive variant types

- ▶ Concise, elegant way to define different tree types
 - Nodes are tagged with a constructor with children for carried data

```
type expr =
| Constant of int
| Negate of expr
| Add of expr * expr
| Mul of expr * expr

Add (Constant 19, Negate (Constant 4))
```



- ▶ Functions over **expr** will typically use recursion over the tree structure

Defining our own option types

- ▶ Using what we already know:

```
type int_option = NoInt | OneInt of int
```

- ▶ Using new feature of defining our own *type constructor*:

```
type 'a myoption = MyNone | MySome of 'a
```

- ▶ Exactly how “built-in” options are defined (in standard library):

```
type 'a option = None | Some of 'a
```

So use pattern-matching!

- ▶ `None` and `Some` are (just) constructors for `'a option`
- ▶ *Build* with constructors and *use* with match expressions
 - Do not use `Option.get` and `Option.is_some` anymore
 - Get the advantages of variant types and pattern matching

Demonstration 1: Matching

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

Demonstration 1: Matching

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# (* Variants can be recursive, describing recursive data structures like trees *)
type expr =
| Constant of int
| Negate of expr
| Add of expr * expr
| Mul of expr * expr

let rec eval e =
  match e with
  | Constant i -> i
  | Negate e1 -> -eval e1
  | Add (e1, e2) -> eval e1 + eval e2
  | Mul (e1, e2) -> eval e1 * eval e2

let rec max_const (e : expr) : int =
  let max (x, y) = if x > y then x else y in
  match e with
  | Constant i -> i
  | Negate e1 -> max_const e1
  | Add (e1, e2) -> max (max_const e1, max_const e2)
  | Mul (e1, e2) -> max (max_const e1, max_const e2)

let rec has_const_not_under_add e =
  match e with
  | Constant i -> true
  | Negate e1 -> has_const_not_under_add e1
  | Add (e1, e2) -> false
  | Mul (e1, e2) -> has_const_not_under_add e1 || has_const_not_under_add e2
```

Demonstration 1: Matching

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# (* Variants can be recursive, describing recursive data structures like trees *)
type expr =
| Constant of int
| Negate of expr
| Add of expr * expr
| Mul of expr * expr

let rec eval e =
  match e with
  | Constant i -> i
  | Negate e1 -> -eval e1
  | Add (e1, e2) -> eval e1 + eval e2
  | Mul (e1, e2) -> eval e1 * eval e2

let rec max_const (e : expr) : int =
  let max (x, y) = if x > y then x else y in
  match e with
  | Constant i -> i
  | Negate e1 -> max_const e1
  | Add (e1, e2) -> max (max_const e1, max_const e2)
  | Mul (e1, e2) -> max (max_const e1, max_const e2)

let rec has_const_not_under_add e =
  match e with
  | Constant i -> true
  | Negate e1 -> has_const_not_under_add e1
  | Add (e1, e2) -> false
  | Mul (e1, e2) -> has_const_not_under_add e1 || has_const_not_under_add e2

type expr =
  Constant of int
| Negate of expr
| Add of expr * expr
| Mul of expr * expr
val eval : expr -> int = <fun>
val max_const : expr -> int = <fun>
val has_const_not_under_add : expr -> bool = <fun>
val number_of_adds : expr -> int = <fun>
val example_exp : expr = Add (Constant 19, Negate (Constant 4))
val example_ans : int = 15
val example_addcount : int = 2
```

Demonstration 1: Matching

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

Demonstration 1: Matching

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# (* Same features already used can almost define option *)
type int_option = NoInt | OneInt of int

let rec sum_int_options1 xs =
  if xs = [] then 0
  else
    match List.hd xs with
    | NoInt -> sum_int_options1 (List.tl xs)
    | OneInt i -> i + sum_int_options1 (List.tl xs)

let test1 = sum_int_options1 [ NoInt; OneInt 7; NoInt; OneInt 2; OneInt 1 ] ;;
```

Demonstration 1: Matching

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!  
  
# (* Same features already used can almost define option *)  
type int_option = NoInt | OneInt of int  
  
let rec sum_int_options1 xs =  
  if xs = [] then 0  
  else  
    match List.hd xs with  
    | NoInt -> sum_int_options1 (List.tl xs)  
    | OneInt i -> i + sum_int_options1 (List.tl xs)  
  
let test1 = sum_int_options1 [ NoInt; OneInt 7; NoInt; OneInt 2; OneInt 1 ] ;;  
  
type int_option = NoInt | OneInt of int  
val sum_int_options1 : int_option list -> int = <fun>  
val test1 : int = 10
```

Demonstration 1: Matching

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!  
  
# (* Same features already used can almost define option *)  
type int_option = NoInt | OneInt of int  
  
let rec sum_int_options1 xs =  
  if xs = [] then 0  
  else  
    match List.hd xs with  
    | NoInt -> sum_int_options1 (List.tl xs)  
    | OneInt i -> i + sum_int_options1 (List.tl xs)  
  
let test1 = sum_int_options1 [ NoInt; OneInt 7; NoInt; OneInt 2; OneInt 1 ] ;;  
  
type int_option = NoInt | OneInt of int  
val sum_int_options1 : int_option list -> int = <fun>  
val test1 : int = 10  
  
# (* In fact, we /can/ define our own polymorphic variant types *)  
type 'a my_option = MyNone | MySome of 'a  
  
let rec sum_int_options2 xs =  
  if xs = [] then 0  
  else  
    match List.hd xs with  
    | MyNone -> sum_int_options2 (List.tl xs)  
    | MySome i -> i + sum_int_options2 (List.tl xs)  
  
let test2 = sum_int_options2 [ MyNone; MySome 7; MyNone; MySome 2; MySome 1 ] ;;
```

Demonstration 1: Matching

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!

# (* Same features already used can almost define option *)
type int_option = NoInt | OneInt of int

let rec sum_int_options1 xs =
  if xs = [] then 0
  else
    match List.hd xs with
    | NoInt -> sum_int_options1 (List.tl xs)
    | OneInt i -> i + sum_int_options1 (List.tl xs)

let test1 = sum_int_options1 [ NoInt; OneInt 7; NoInt; OneInt 2; OneInt 1 ] ;;

type int_option = NoInt | OneInt of int
val sum_int_options1 : int_option list -> int = <fun>
val test1 : int = 10

# (* In fact, we /can/ define our own polymorphic variant types *)
type 'a my_option = MyNone | MySome of 'a

let rec sum_int_options2 xs =
  if xs = [] then 0
  else
    match List.hd xs with
    | MyNone -> sum_int_options2 (List.tl xs)
    | MySome i -> i + sum_int_options2 (List.tl xs)

let test2 = sum_int_options2 [ MyNone; MySome 7; MyNone; MySome 2; MySome 1 ] ;;

type 'a my_option = MyNone | MySome of 'a
val sum_int_options2 : int my_option list -> int = <fun>
val test2 : int = 10
```

Demonstration 1: Matching

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

Demonstration 1: Matching

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# (* Indeed, the option type constructor is not "built in" at all; just in standard library *)
(* type 'a option = None | Some of 'a *)
(* from now on, use pattern-matching for options *not* the previous way we showed to use them *)
let rec sum_int_options3 xs =
  if xs = [] then 0
  else
    match List.hd xs with
    | None -> sum_int_options3 (List.tl xs)
    | Some i -> i + sum_int_options3 (List.tl xs)

let test3 = sum_int_options3 [ None; Some 7; None; Some 2; Some 1 ] ;;
```

Demonstration 1: Matching

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# (* Indeed, the option type constructor is not "built in" at all; just in standard library *)
(* type 'a option = None | Some of 'a *)
(* from now on, use pattern-matching for options *not* the previous way we showed to use them *)
let rec sum_int_options3 xs =
  if xs = [] then 0
  else
    match List.hd xs with
    | None -> sum_int_options3 (List.tl xs)
    | Some i -> i + sum_int_options3 (List.tl xs)

let test3 = sum_int_options3 [ None; Some 7; None; Some 2; Some 1 ] ;;

val sum_int_options3 : int option list -> int = <fun>
val test3 : int = 10
```

How about lists?

- ▶ We can use the same features and self-reference for our own list type constructor:

```
type 'a my_list = Empty | Cons of 'a * 'a my_list
```

- ▶ *Build* with constructors:

```
let xs = Cons (3, Cons (4, Cons (5, Empty)))
let ss = Cons ("hi", Cons ("bye", Empty))
```

- ▶ *Use* with match expressions:

```
let rec length xs =
  match xs with
  | Empty -> 0
  | Cons (x, xs') -> 1 + length xs'
```

OCaml's built-in lists

- ▶ Could have done *exactly* what we just did...
- ▶ but instead it is using special *syntax*:
 - Constructors are `[]` and `::` with `::` written **infix**
 - and/but use match expressions to use lists
 - ▶ **do not** use `List.hd`, `List.tl` and `e = []` anymore!

```
let rec length xs =
  match xs with
  | [] -> 0
  | x :: xs' -> 1 + length xs'
```

Why then `List.hd` and `List.tl`?

- ▶ Occasionally convenient
 - e.g., for passing functions to other functions (next week)
- ▶ A stepping stone for us to the more elegant and less error-prone approach (match expressions warn on missing cases!)
- ▶ Again, do not use these functions on remaining assignments
- ▶ And note how easy they are to define:

```
let hd xs =
  match xs with
  | [] -> failwith "list empty"
  | x :: _ -> x
  (* / x :: xs' -> x
     is also OK, but rather unstylish;
     compiler might even bark *)
```

```
let isEmpty xs =
  match xs with
  | [] -> true
  | _ -> false
  (* / x :: xs' -> false
     is also OK, but again not co
```

Demonstration 2: Polymorphic lists

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

Demonstration 2: Polymorphic lists

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# type 'a my_list = Empty | Cons of 'a * 'a my_list

let rec sum_int_options4 xs =
  match xs with
  | Empty -> 0
  | Cons (x, xs') ->
    (match x with
     | None -> sum_int_options4 xs'
     | Some i -> i + sum_int_options4 xs')
let test4 = sum_int_options4
  (Cons (None, Cons (Some 7, Cons (None, Cons (Some 2, Cons (Some 1, Empty))))))) ;;
```

Demonstration 2: Polymorphic lists

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# type 'a my_list = Empty | Cons of 'a * 'a my_list

let rec sum_int_options4 xs =
  match xs with
  | Empty -> 0
  | Cons (x, xs') ->
    (match x with
     | None -> sum_int_options4 xs'
     | Some i -> i + sum_int_options4 xs')
let test4 = sum_int_options4
  (Cons (None, Cons (Some 7, Cons (None, Cons (Some 2, Cons (Some 1, Empty)))))) ;;

type 'a my_list = Empty | Cons of 'a * 'a my_list
val sum_int_options4 : int option my_list -> int = <fun>
val test4 : int = 10
```

Demonstration 2: Polymorphic lists

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# type 'a my_list = Empty | Cons of 'a * 'a my_list

let rec sum_int_options4 xs =
  match xs with
  | Empty -> 0
  | Cons (x, xs') ->
    (match x with
     | None -> sum_int_options4 xs'
     | Some i -> i + sum_int_options4 xs')
let test4 = sum_int_options4
  (Cons (None, Cons (Some 7, Cons (None, Cons (Some 2, Cons (Some 1, Empty))))))) ;;

type 'a my_list = Empty | Cons of 'a * 'a my_list
val sum_int_options4 : int option my_list -> int = <fun>
val test4 : int = 10

# (* This is exactly how built-in lists are defined /except/ special syntax [] and :: !
   So yes, we can pattern-match with those constructors and should no longer use = [], 
   List.hd, or List.tl (!! *) )
let rec sum_int_options5 xs =
  match xs with
  | [] -> 0
  | x :: xs' -> ( (* parentheses are optional *)
    match x with
    | None -> sum_int_options5 xs'
    | Some i -> i + sum_int_options5 xs')
let test5 = sum_int_options5 [ None; Some 7; None; Some 2; Some 1 ] ;;
```

Demonstration 2: Polymorphic lists

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# type 'a my_list = Empty | Cons of 'a * 'a my_list

let rec sum_int_options4 xs =
  match xs with
  | Empty -> 0
  | Cons (x, xs') ->
    (match x with
     | None -> sum_int_options4 xs'
     | Some i -> i + sum_int_options4 xs')
let test4 = sum_int_options4
  (Cons (None, Cons (Some 7, Cons (None, Cons (Some 2, Cons (Some 1, Empty))))))) ;;

type 'a my_list = Empty | Cons of 'a * 'a my_list
val sum_int_options4 : int option my_list -> int = <fun>
val test4 : int = 10

# (* This is exactly how built-in lists are defined /except/ special syntax [] and :: !
   So yes, we can pattern-match with those constructors and should no longer use = [], 
   List.hd, or List.tl (!! *) )
let rec sum_int_options5 xs =
  match xs with
  | [] -> 0
  | x :: xs' -> ( (* parentheses are optional *)
    match x with
    | None -> sum_int_options5 xs'
    | Some i -> i + sum_int_options5 xs')
let test5 = sum_int_options5 [ None; Some 7; None; Some 2; Some 1 ] ;;

type 'a my_list = Empty | Cons of 'a * 'a my_list
val sum_int_options4 : int option my_list -> int = <fun>
val test4 : int = 10
val sum_int_options5 : int option list -> int = <fun>
val test5 : int = 10
```

Demonstration 2: Polymorphic lists

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# type 'a my_list = Empty | Cons of 'a * 'a my_list

let rec sum_int_options4 xs =
  match xs with
  | Empty -> 0
  | Cons (x, xs') ->
    (match x with
     | None -> sum_int_options4 xs'
     | Some i -> i + sum_int_options4 xs')
let test4 = sum_int_options4
  (Cons (None, Cons (Some 7, Cons (None, Cons (Some 2, Cons (Some 1, Empty))))))) ;;

type 'a my_list = Empty | Cons of 'a * 'a my_list
val sum_int_options4 : int option my_list -> int = <fun>
val test4 : int = 10

# (* This is exactly how built-in lists are defined /except/ special syntax [] and :: !
   So yes, we can pattern-match with those constructors and should no longer use = [], List.hd, or List.tl (!! *) )
let rec sum_int_options5 xs =
  match xs with
  | [] -> 0
  | x :: xs' -> ( (* parentheses are optional *)
    match x with
    | None -> sum_int_options5 xs'
    | Some i -> i + sum_int_options5 xs')
let test5 = sum_int_options5 [ None; Some 7; None; Some 2; Some 1 ] ;;

type 'a my_list = Empty | Cons of 'a * 'a my_list
val sum_int_options4 : int option my_list -> int = <fun>
val test4 : int = 10
val sum_int_options5 : int option list -> int = <fun>
val test5 : int = 10

# (* Spoiler alert: nested patterns can make this even more concise
   we aren't /quite/ there yet, but this is the style we expect on hw2 *)
let rec sum_int_options6 xs =
  match xs with
  | [] -> 0
  | None :: xs' -> sum_int_options6 xs'
  | Some i :: xs' -> i + sum_int_options6 xs'
let test6 = sum_int_options6 [ None; Some 7; None; Some 2; Some 1 ] ;;
```

Demonstration 2: Polymorphic lists

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# type 'a my_list = Empty | Cons of 'a * 'a my_list

let rec sum_int_options4 xs =
  match xs with
  | Empty -> 0
  | Cons (x, xs') ->
    (match x with
     | None -> sum_int_options4 xs'
     | Some i -> i + sum_int_options4 xs')
let test4 = sum_int_options4
  (Cons (None, Cons (Some 7, Cons (None, Cons (Some 2, Cons (Some 1, Empty))))))) ;;

type 'a my_list = Empty | Cons of 'a * 'a my_list
val sum_int_options4 : int option my_list -> int = <fun>
val test4 : int = 10

# (* This is exactly how built-in lists are defined /except/ special syntax [] and :: !
   So yes, we can pattern-match with those constructors and should no longer use = [], List.hd, or List.tl (!! *) )
let rec sum_int_options5 xs =
  match xs with
  | [] -> 0
  | x :: xs' -> ( (* parentheses are optional *)
    match x with
    | None -> sum_int_options5 xs'
    | Some i -> i + sum_int_options5 xs')
let test5 = sum_int_options5 [ None; Some 7; None; Some 2; Some 1 ] ;;

type 'a my_list = Empty | Cons of 'a * 'a my_list
val sum_int_options4 : int option my_list -> int = <fun>
val test4 : int = 10
val sum_int_options5 : int option list -> int = <fun>
val test5 : int = 10

# (* Spoiler alert: nested patterns can make this even more concise
   we aren't /quite/ there yet, but this is the style we expect on hw2 *)
let rec sum_int_options6 xs =
  match xs with
  | [] -> 0
  | None :: xs' -> sum_int_options6 xs'
  | Some i :: xs' -> i + sum_int_options6 xs'
let test6 = sum_int_options6 [ None; Some 7; None; Some 2; Some 1 ] ;;

val sum_int_options6 : int option list -> int = <fun>
val test6 : int = 10
```

Patterns

How languages define variant types

- ▶ OCaml built powerful pattern-matching into the language and used it as *the fundamental primitive* for one-of types
- ▶ Other languages make other choices. For example OCaml *could* have:
 - kept variant type definitions as-is
 - had them automatically introduce *functions* like `isEmpty`, `hd` etc.

More pattern-matching

- ▶ OCaml also has pattern-matching for each-of types
- ▶ Pattern `(x, y, z)` matches triples `(v1, v2, v3)`
 - Similar for any width tuple, naturally
 - Similar patterns `{ f1 = x1; ...; fN = xN }` for records
- ▶ Useful for **nested patterns** (next lecture)
- ▶ But also useful to make let expressions even *more powerful*

Bad style

- ▶ One-branch match expressions make semantic sense but are ugly!
- ▶ **Don't do this:**

```
let sum_triple tr =
  match tr with
  | (x, y, z) -> x + y + z
```

Let expressions and patterns

- ▶ Let expression syntax is actually

```
let p = e1 in e2
```

where **p** is a *pattern*!

- Top-level let bindings can also do this
- Great for extracting multiple pieces of data at once
- No more need for **fst**, **snd**, **fst3**, **snd3**, **thd3**, and so on

```
let sum_triple tr =
  let (x, y, z) = tr in
    x + y + z
```

Another new feature

- ▶ Function arguments can also use patterns!
- ▶ **Function call semantics:** pattern-match the argument, extracting data for body

```
let sum_triple tr =
  let (x, y, z) = tr in
  x + y + z
let sum_triple (x, y, z) =
  x + y + z
let ten = sum_triple (5, 2, 3)
```

Behold! 😊

Compare:

- ▶ A function that takes a triple and sums its int components:

```
let sum_triple (x, y, z) =  
  x + y + z
```

- ▶ A function that takes three ints and sums them:

```
let sum_triple (x, y, z) =  
  x + y + z
```

- ▶ See the difference? (me neither...) 😊 🎉 😊 🎉 😊

The truth about functions

- ▶ In OCaml, every function takes exactly one argument
- ▶ We can *simulate* multiple arguments with tuples and tuple-patterns
 - Elegant, compositional design
 - Enables sometimes-convenient things
 - ▶ Return a tuple from one function and pass it like multiple arguments to another
- ▶ Soon we will see a different and more common way to simulate multiple arguments that *also* has built-in syntactic support

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# (* Pattern-matching is the normal ML way to use lists; let's revisit prior functions *)
let rec length xs =
  match xs with
  | [] -> 0
  | x :: xs' -> 1 + length xs'

let rec append (xs, ys) =
  match xs with [] -> ys | x :: xs' -> x :: append (xs', ys)

let rec concat ss =
  match ss with
  | [] -> ""
  | s :: ss' -> s ^ concat ss' ;;
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# (* Pattern-matching is the normal ML way to use lists; let's revisit prior functions *)
let rec length xs =
  match xs with
  | [] -> 0
  | x :: xs' -> 1 + length xs'

let rec append (xs, ys) =
  match xs with [] -> ys | x :: xs' -> x :: append (xs', ys)

let rec concat ss =
  match ss with
  | [] -> ""
  | s :: ss' -> s ^ concat ss' ;;

val length : 'a list -> int = <fun>
val append : 'a list * 'a list -> 'a list = <fun>
val concat : string list -> string = <fun>
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# (* Pattern-matching is the normal ML way to use lists; let's revisit prior functions *)
let rec length xs =
  match xs with
  | [] -> 0
  | x :: xs' -> 1 + length xs'

let rec append (xs, ys) =
  match xs with [] -> ys | x :: xs' -> x :: append (xs', ys)

let rec concat ss =
  match ss with
  | [] -> ""
  | s :: ss' -> s ^ concat ss' ;;

val length : 'a list -> int = <fun>
val append : 'a list * 'a list -> 'a list = <fun>
val concat : string list -> string = <fun>

# (* Pattern-matching for each-of types (tuples shown; records can also be pattern-matched)*)

(* 🚫 Terrible style never used: one-arm match expressions *)
let sum_triple1 tr = match tr with x, y, z -> x + y + z

(* 😊 Appropriate style: let expression syntax is /actually/: let p = e1 in e2 *)
let sum_triple2 tr =
  let x, y, z = tr in
  x + y + z

(* 😎 Even better when useful: can put a pattern right in the function binding: let rec f p = e *)
let sum_triple3 (x, y, z) = x + y + z
(* in fact, thanks to a convenient fib, that's what we have done since Lecture 2! *) ;;
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# (* Pattern-matching is the normal ML way to use lists; let's revisit prior functions *)
let rec length xs =
  match xs with
  | [] -> 0
  | x :: xs' -> 1 + length xs'

let rec append (xs, ys) =
  match xs with [] -> ys | x :: xs' -> x :: append (xs', ys)

let rec concat ss =
  match ss with
  | [] -> ""
  | s :: ss' -> s ^ concat ss' ;;

val length : 'a list -> int = <fun>
val append : 'a list * 'a list -> 'a list = <fun>
val concat : string list -> string = <fun>

# (* Pattern-matching for each-of types (tuples shown; records can also be pattern-matched)*)

(* 🚫 Terrible style never used: one-arm match expressions *)
let sum_triple1 tr = match tr with x, y, z -> x + y + z

(* 😊 Appropriate style: let expression syntax is /actually/: let p = e1 in e2 *)
let sum_triple2 tr =
  let x, y, z = tr in
  x + y + z

(* 😎 Even better when useful: can put a pattern right in the function binding: let rec f p = e *)
let sum_triple3 (x, y, z) = x + y + z
(* in fact, thanks to a convenient fib, that's what we have done since Lecture 2! *) ;;

val length : 'a list -> int = <fun>
val append : 'a list * 'a list -> 'a list = <fun>
val concat : string list -> string = <fun>
val sum_triple1 : int * int * int -> int = <fun>
val sum_triple2 : int * int * int -> int = <fun>
val sum_triple3 : int * int * int -> int = <fun>
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# (* And one more nested-patterns spoiler: *)
let rec sum_pairs xs =
  match xs with
  | [] -> 0
  | (x, y) :: xs' -> x + y + sum_pairs xs'

(* Cute example of expressiveness of functions actually taking one tuple *)
let rotate_left (x, y, z) = (y, z, x)
let rotate_right tr = rotate_left (rotate_left tr)

(* Just as never use one-branch match expressions with each-of patterns,
   it is also usually bad style to use let expressions with one-of patterns:
   you will get a warning at compile-time plus a run-time exception if match fails *)
let get_risky1 opt =
  match opt with
  | None -> failwith "nopes"
  | Some v -> v
let get_risky2 opt = let (Some v) = opt in v ;;
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)

# (* And one more nested-patterns spoiler: *)
let rec sum_pairs xs =
  match xs with
  | [] -> 0
  | (x, y) :: xs' -> x + y + sum_pairs xs'

(* Cute example of expressiveness of functions actually taking one tuple *)
let rotate_left (x, y, z) = (y, z, x)
let rotate_right tr = rotate_left (rotate_left tr)

(* Just as never use one-branch match expressions with each-of patterns,
   it is also usually bad style to use let expressions with one-of patterns:
   you will get a warning at compile-time plus a run-time exception if match fails *)
let get_risky1 opt =
  match opt with
  | None -> failwith "nopes"
  | Some v -> v
let get_risky2 opt = let (Some v) = opt in v ;;

Lines 20-21, characters 2-3:
Warning 8 [partial-match]: this pattern-matching is not exhaustive.
  Here is an example of a case that is not matched: None

val sum_pairs : (int * int) list -> int = <fun>
val rotate_left : 'a * 'b * 'c -> 'b * 'c * 'a = <fun>
val rotate_right : 'a * 'b * 'c -> 'c * 'a * 'b = <fun>
val get_risky1 : 'a option -> 'a = <fun>
val get_risky2 : 'a option -> 'a = <fun>
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# (* And one more nested-patterns spoiler: *)
let rec sum_pairs xs =
  match xs with
  | [] -> 0
  | (x, y) :: xs' -> x + y + sum_pairs xs'

(* Cute example of expressiveness of functions actually taking one tuple *)
let rotate_left (x, y, z) = (y, z, x)
let rotate_right tr = rotate_left (rotate_left tr)

(* Just as never use one-branch match expressions with each-of patterns,
   it is also usually bad style to use let expressions with one-of patterns:
   you will get a warning at compile-time plus a run-time exception if match fails *)
let get_risky1 opt =
  match opt with
  | None -> failwith "nopes"
  | Some v -> v
let get_risky2 opt = let (Some v) = opt in v ;;

Lines 20-21, characters 2-3:
Warning 8 [partial-match]: this pattern-matching is not exhaustive.
  Here is an example of a case that is not matched: None

val sum_pairs : (int * int) list -> int = <fun>
val rotate_left : 'a * 'b * 'c -> 'b * 'c * 'a = <fun>
val rotate_right : 'a * 'b * 'c -> 'c * 'a * 'b = <fun>
val get_risky1 : 'a option -> 'a = <fun>
val get_risky2 : 'a option -> 'a = <fun>

# let get_risky3 (Some v) = v
get_risky1 None ;;
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)

# (* And one more nested-patterns spoiler: *)
let rec sum_pairs xs =
  match xs with
  | [] -> 0
  | (x, y) :: xs' -> x + y + sum_pairs xs'

(* Cute example of expressiveness of functions actually taking one tuple *)
let rotate_left (x, y, z) = (y, z, x)
let rotate_right tr = rotate_left (rotate_left tr)

(* Just as never use one-branch match expressions with each-of patterns,
   it is also usually bad style to use let expressions with one-of patterns:
   you will get a warning at compile-time plus a run-time exception if match fails *)
let get_risky1 opt =
  match opt with
  | None -> failwith "nopes"
  | Some v -> v
let get_risky2 opt = let (Some v) = opt in v ;;

Lines 20-21, characters 2-3:
Warning 8 [partial-match]: this pattern-matching is not exhaustive.
  Here is an example of a case that is not matched: None

val sum_pairs : (int * int) list -> int = <fun>
val rotate_left : 'a * 'b * 'c -> 'b * 'c * 'a = <fun>
val rotate_right : 'a * 'b * 'c -> 'c * 'a * 'b = <fun>
val get_risky1 : 'a option -> 'a = <fun>
val get_risky2 : 'a option -> 'a = <fun>

# let get_risky3 (Some v) = v
get_risky1 None ;;

Warning 8 [partial-match]: this pattern-matching is not exhaustive.
  Here is an example of a case that is not matched: None
val get_risky3 : (('a option -> 'a) -> 'b option -> 'c) option -> 'c = <fun>
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# type t = { a: int; b: string }
let v = { a = 5; b = "6" } ;;
let { a; b } = v in a + int_of_string b ;;
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# type t = { a: int; b: string }
let v = { a = 5; b = "6" } ;;
let { a; b } = v in a + int_of_string b ;;
type t = { a : int; b : string; }
val v : t = {a = 5; b = "6"}
- : int = 11
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# type t = { a: int; b: string }
let v = { a = 5; b = "6" } ;;

let { a; b } = v in a + int_of_string b ;;

type t = { a : int; b : string; }
val v : t = {a = 5; b = "6"}
- : int = 11

# match v with
| { a = a'; b = b' } -> a' + 1
| { a; _ } -> a - 1 ;;
```

Demonstration 3: Patterns

```
Welcome to utop version 2.16.0 (using OCaml version 5.4.0)!
```

```
# type t = { a: int; b: string }
let v = { a = 5; b = "6" } ;;

let { a; b } = v in a + int_of_string b ;;

type t = { a : int; b : string; }
val v : t = {a = 5; b = "6"}
- : int = 11

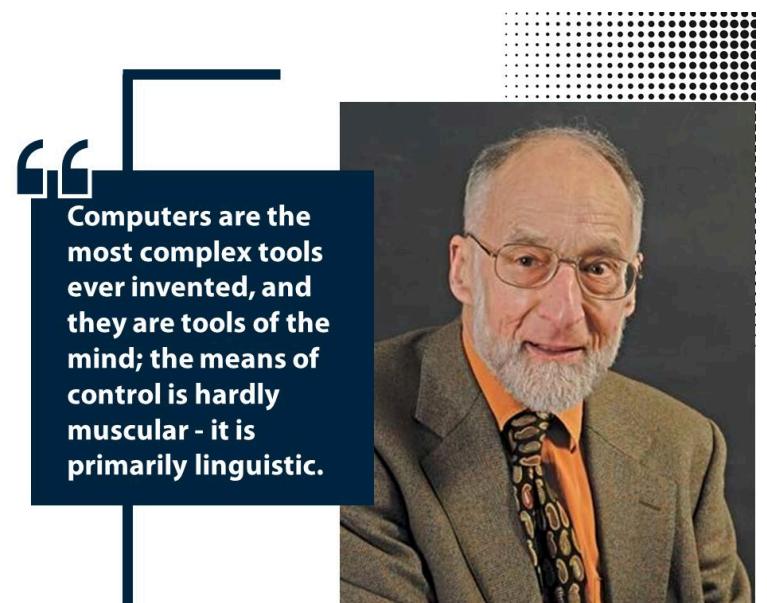
# match v with
| { a = a'; b = b' } -> a' + 1
| { a; _ } -> a - 1 ;;

Line 3, characters 2-10:
Warning 11 [redundant-case]: this match case is unused.
```

Some history ...

ML

- ▶ Typechecking: **Robin Milner** et al. at the Edinburgh Laboratory for Computer Science in Scotland
 - Based on theorem proving
 - Milner developed ML: *Meta Language* whose goal is to construct valid proofs
 - Famous typechecking (type inference) algorithm: Hindley–Milner–Damas
 - ▶ Similar to garbage collection
 - Won a Turing Award in 1991



Robin Milner
ACM A.M. Turing Laureate



ML: Example

```
let fact n =
  letref count=n and result=1
  in    if count=0
        then result
        else loop count,result := count-1,count*result;;
# fact = - : (int -> int)

fact 4;;
# 24 : int
```

On to OCaml...

- ▶ In the early '80s, there was a schism in the ML community:
 -  vs 
 - ▶  at Inria (Gérard Huet, Xavier Leroy, Jérôme Vouillon et al.) developed **CAML** (*Categorical Abstract Machine Language*) and later **OCaml** (*Objective CAML*; 1996)
 - ▶  developed **Standard ML**
 - ▶ Both are very similar
 - Microsoft introduced a variant of OCaml suited for .NET called **F#** (2005)



Weekly edition

The world in brief

War in the Middle East

War in Ukraine

United States

The world economy

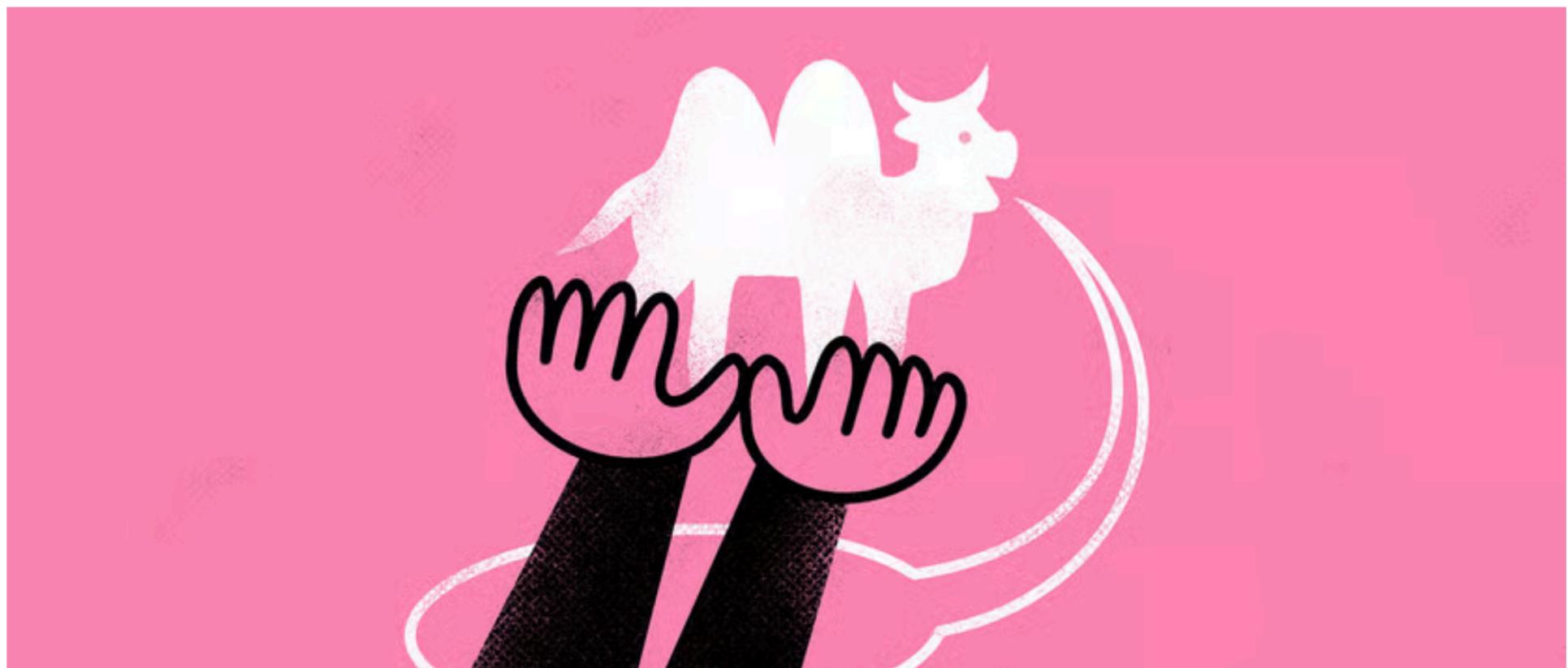
Business

Finance & economics | A large hump

Jane Street's sneaky retention tactic

It involves the use of an obscure, French programming language

 Share



ML

- ▶ Common features of ML languages:

- Similar syntax
- Static type system with parametric polymorphism (about that later!)
- Functional programming
- Modularity
- General purpose
- Emphasis on theorem proving and static analysis ([Coq/Rocq](#) theorem prover)

