# Written Assignment 7

Daniel Detore CS 135-B/LF

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1.a

$$T(1) = 2T(\lfloor \frac{1}{4} \rfloor) + \sqrt{1}$$
$$= 2T(0) + 1$$
$$= 2(1) + 1$$
$$T(1) = 3$$

$$T(2) = 2T(\lfloor \frac{2}{4} \rfloor) + \sqrt{2}$$
$$= 2T(0) + \sqrt{2}$$
$$T(2) = 2 + \sqrt{2}$$
$$\approx 3.41421$$

## 1.b

We need to prove that  $T(n) \leq 3\sqrt{n}\log_2 n$  for all  $n \geq 2$ . Let's check this premise for basis n=2:

$$T(2) \leq 3\sqrt{2}\log_2(2)$$
 
$$2 + \sqrt{2} \leq 3\sqrt{2}$$
 
$$2 \leq 2\sqrt{2}$$
 
$$1 \leq \sqrt{2}$$
 
$$\top$$

Inductive hypothesis: Let k be an integer  $\geq 2$ . We'll assume that  $T(i) \leq 3\sqrt{i} \log_2 i$  for all  $2 \leq i \leq k, i \in \mathbb{Z}$ . Now we must prove that

$$T(k+1) \le 3\sqrt{k+1}\log_2(k+1).$$

$$T(k+1) = 2T(\left\lfloor \frac{k+1}{4} \right\rfloor) + \sqrt{k+1}$$
 By definition of  $T$  
$$\leq 2(3\sqrt{\left\lfloor \frac{k+1}{4} \right\rfloor} \log_2 \left\lfloor \frac{k+1}{4} \right\rfloor) + \sqrt{k+1}$$
 By the inductive hypothesis 
$$\leq 2(3\sqrt{\frac{k+1}{4}} \log_2 (\frac{k+1}{4})) + \sqrt{k+1}$$
 
$$= 2(\frac{3}{2}\sqrt{k+1} \log_2 (\frac{k+1}{4})) + \sqrt{k+1}$$
 By algebra 
$$= 3\sqrt{k+1} \log_2 (\frac{k+1}{4}) + \sqrt{k+1}$$
 By algebra 
$$= 3\sqrt{k+1} (\log_2 (k+1) - \log_2 (4)) + \sqrt{k+1}$$
 
$$= 3\sqrt{k+1} (\log_2 (k+1) - 2) + \sqrt{k+1}$$
 
$$= 3\sqrt{k+1} \log_2 (k+1) - 6\sqrt{k+1} + \sqrt{k+1}$$
 
$$= 3\sqrt{k+1} \log_2 (k+1) - 5\sqrt{k+1}$$
 
$$< 3\sqrt{k+1} \log_2 (k+1)$$

 $\mathbf{2}$ 

Let  $a, b, c, d, m \in \mathbb{Z}$  with  $m \geq 2$ . We'll prove that  $a \equiv b \pmod{m} \land c \equiv d \pmod{m} \implies a - c \equiv b - d \pmod{m}$ . Given

$$\begin{cases} a \equiv b \pmod{m} \\ c \equiv d \pmod{m} \end{cases}$$

we can break this down into

$$\begin{cases} a \mod m = b \mod m \\ c \mod m = d \mod m \end{cases}$$

by the definition of congruence. We'll scale the second equation by -1,

$$\begin{cases} a \mod m &= b \mod m \\ (-c) \mod m &= (-d) \mod m \end{cases}$$

add the equations together as one can do with a system,

$$(a-c) \mod m = (b-d) \mod m$$

and rewrite this as the congruence

$$a - c \equiv b - d \pmod{m}$$
.

3

No. For k = 6, 2 \* 3 \* 5 \* 7 \* 11 \* 13 + 1 = 30031. 30031 has factors  $\{59, 509\}$  which means 30031 is not prime. Therefore  $p_1 p_2 \cdots p_k + 1$  is not always prime.

## 4

No. Since 30 is not divisible by 29, let m = 30:

$$2m^2 + 29$$
  
=  $2(30)^2 + 29$   
=  $1829$ 

1829 has factors  $\{31, 59\}$  which means 1829 is not prime. Therefore  $2m^2 + 29$  does not always produce a prime number when m is not divisible by 29.

### 5

Let's prove a basis where k = 1: and since  $1^4 - 1 = 0$  and 0 is divisible by 16, this basis holds. Let's assume that k is divisible by 16. With this we can say that there is some integer s for which  $k^4 - 1 = 16s$ .

Let's prove that k + 2 is divisible by 16 (k + 2 because k and this equation's result must be odd). We can do this by proving that there is some integer which, when multiplied by 16, gives some expression of k. For some integer t:

$$16t = (k+2)^4 - 1$$

$$= (k^2 + 4k + 4)^2 - 1$$

$$= k^4 + 8k^3 + 24k^2 + 32k + 16 - 1$$

$$= k^4 - 1 + 8k^3 + 24k^2 + 32k + 16$$

By the inductive hypothesis: there is some integer s for which  $k^4 - 1 = 16s$ 

$$= 16s + 8k^3 + 24k^2 + 32k + 16$$
$$= 16s + 8k^2(k+3) + 16(2k+1)$$

Since k is an integer, we can reasonably expect 2k + 1 to be an integer as well.  $8k^2(k+3)$  will be a multiple of 8 and some odd number  $k^2(k+3)$ , which is guaranteed to be divisible by 16. We can think of this equation as

(Some multiple of 
$$16$$
) = (The sum of multiples of  $16$ )

which proves that k + 2 produces a multiple of 16, which proves that for all odd integers  $k \ge 1$ ,  $k^4 - 1$  is divisible by 16.

### 6

Now we have a lot of easily-calculable modulus operations.

 $= 100 \cdot ((3^7 \bmod 13)^3 \bmod 13)^2 \bmod 13$ 

 $= 100 \cdot 1^2 \bmod 13$ 

 $= 100 \cdot 1 \bmod 13$ 

 $= 100 \bmod 13$ 

= 9