# **Functional Programming in Scala**

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#### **Outline**

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# **Functional Programming**

### Function are values, i.e., a function can be

		value	i unction
•	Anonymous	3	x => x + 1
•	Assigned to a variable	x = 3	f = X => X + 1
•	Passed as input/output parameter	f(3)	f(x => x + 1)
•	Created dynamically	3 + 4	f ∘ g

Value

Function

# **Fundamental Theory**

- Imperative languages ⇒ Von Neumann Architecture
  - Efficiency
- Functional languages ⇒ Lambda Calculus
  - A solid theoretical basis that is also closer to the user, but
  - relatively unconcerned with the architecture of the machines on which programs will run

#### **Mathematical Functions**

- A mathematical function is
  - a mapping of members of one set, called the domain set, to another set, called the range set
- A lambda expression specifies the parameter(s) and the mapping of a function in the following form λ(x) x \* x \* x for the function cube
   cube(x) = x \* x \* x
- Lambda expressions describe nameless functions
- Lambda expressions are applied to parameter(s) by placing the parameter(s) after the expression
   (λ(x) x \* x \* x)(2) which evaluates to 8

# **Higher-order Functions**

- A higher-order function is one that either takes functions as parameters or yields a function as its result, or both
- For example,
  - Function composition
  - Apply-to-all
  - Forall/Exists
  - Insert-left/Insert-right
  - Functions as parameters
  - Closures

# **Function Composition**

#### A function that

- takes two functions as parameters and
- yields a function whose value is the first actual parameter function applied to the application of the second

```
f \circ g = f : (g : x)
For f(x) = x + 2; g(x) = x * x; f \circ g(x) = x * x + 2
```

```
val f = (x:Double) => x + 2
val g = (x:Double) => x * x
val h = f compose g
h(3)
val k = f andThen g
k(3)
```

# **Apply-to-all**

#### A functional form that

- takes a single function as a parameter and
- yields a list of values obtained by applying the given function to each element of a list of parameters

$$\alpha f :< x_1, x_2, ..., x_n > = < f : x_1, f : x_2, ..., f : x_2 >$$

For 
$$h(x)=x^*x \Rightarrow \alpha h:(1,2,3)$$
 yields (1,4,9)

```
List(2,3,4).map((x:Int) => x * x)

def inc (x:Int) = x + 1

List(4,5,6).map(inc)
```

#### Forall/Exists

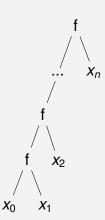
#### A functional form that

- takes a single **predicate** function as a parameter and
- yields a value obtained by applying the given function to each element of a list of parameters and take the and/or of the results

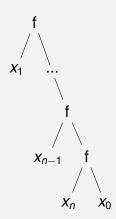
```
\forall f : \langle x_1, x_2, ..., x_n \rangle = \bigcap f : x_i
\exists f : \langle x_1, x_2, ..., x_n \rangle = \bigcup f : x_i
```

# Insert Left / Insert Right

$$/f :< X_0 >, < X_1, X_2, ..., X_n >$$



$$f: < x_0 >, < x_1, x_2, ..., x_n >$$



### Insert Left / Insert Right (cont'd)

#### **Functions as Parameters**

In user-defined functions, functions can be passed as parameters.

```
def apply(x:Int)(f:Int=>Int) = f(x)
val inc1 = (x:Int) => x + 1
val sq = (x:Int) => x * x
val fl = List(inc1, sq)
fl.map(apply(3)) //yield List(4,9)
```

# Closure [1]

"An object is data with functions. A closure is a function with data." - John D. Cook

# Closure = function + binding of its free variables

# **Currying functions [2]**

```
f: X_1 \times X_2 \times ... \times X_n \to Y

curry: f: X_1 \to X_2 \to ... \to X_n \to Y

Example in Scala
```

```
def add(x:Int, y:Int) = x + y
add(1,3)
add(1) add(1)(3)
def plus(x:Int)(y:Int) = x + y
plus(1)(3)
val inc1 = plus(1) _
inc1(3)
val addCurried = (add _).curried
val plusUncurried = Function.uncurried(plus _)
```

Read more on Partially Applied Functions [2]

# **Immutability**

- Immutable: Cannot change
- In Java, strings are immutable "Hello".toUpper() doesn't change "Hello" but returns a new string "HELLO"
- In Scala, val is immutable val num = 12 num = 10 // wrong
- Pure functional programming: No mutations
- Don't mutate—always return the result as a new value
- Functions that don't mutate state are inherently parallelizable

# **Example on Immutability**

```
abstract class IntStack
   def push(x: Int): IntStack =
                    new IntNonEmptyStack(x, this)
   def isEmpty: Boolean
   def top: Int
   def pop: IntStack
class StackEmpty extends IntStack
   def isEmpty = true
   def top = error("EmptyStack.top")
   def pop = error("EmptyStack.pop")
class IntNonEmptyStack(elem: Int, rest: IntStack)
                           extends IntStack
   def isEmpty = false
   def top = elem
   def pop = rest
```

# Everything is an expression

- Body of a function is an expression, i.e. evaluating to a value
- If there are many expressions, the value of the last executed expression will be returned
- No return required

```
def fact(x:Int):Int =
   if (x == 0) 1 else x * fact(x - 1)
val s = for (x <- 1 to 25 if x*x > 50) yield 2*x
```

# Pattern Matching [3]

Like switch, but much more powerful

```
Syntax:
<exp > match {
case <pattern 1 > => <exp1>
case <pattern 2 > => <exp2>
Example in Scala,
def mathTest(x : Int): String = x match {
    case 1 => "one"
    case 2 => "two"
    case => "many"
```

#### Recursion

- With recursive functions, return type is obligated
   def fact(n:Int):Int = if (x == 0) 1 else n \* fact(n 1)
- Need def because the name is used on the right
- Iteration (while, for) can always be expressed as recursion

# Lazy evaluation [4]

- Expressions are eagerly evaluated, by default, where they appeared
- Lazy evaluation means the expression is just evaluated when the associated variable is firstly referred.
  - lazy val x = 1 + y

The expression 1 + y is evaluated just when x is firstly used

Pass-by-name parameter
 def foo(b:Boolean,x:=>Int,y:=>Int) = if (b) x else y foo(a==0,1,b/a)

# Type Inference

- Scala is strongly typed. Any value has a type guarantee
- Just like C++: char\* greeting = "Hello"; greeting = 42;//Type Error
- But without having to declare types:
   val greeting = "Hello";
- Contrast with scripting languages such as JavaScript var greeting = 'Hello';//This is JavaScript greeting = 42;//Ok alert(greeting.length);//Runtime Error

# Type Inference (cont'd)

 Can override inferred type (only to a supertype, of course).

```
var greeting : Any = "Hello"
greeting = 42//Ok
```

- Parameter type must be declared def mistery (x) = 42 \* x // Error def mistery (x : Int) = 42 \* x // Ok
- Return type can be inferred
   If x is Int, then 42 \* x is also Int
- Exception, recursive function requires return type declared

#### **Parameter Inference from Context**

- When a function parameter type is known, an anonymous function can be supplied without specifying its parameter types def twice(f: (Int)=>Int,x:Int) = f(f(x)) twice(x=>42\*x, 3) // Ok, x:Int is inferred from context
   Very useful when calling library functions
  - List(1,2,3).filter(x=> x%2==0)
     List[A].filter(p:(A)=>Boolean) : List[A]
    - A is Int since List(1,2,3) is a List[Int]
    - p: must be (Int) => Boolean
    - X must be Int

# **Parameter Simplification**

- Ok to omit () around a single inferred parameter List(1, 2, 3).filter(x => x % 2 == 0) List(1, 2, 3).sortWith((x,y) => x > y) // need () with 2 or more parameters
- Use \_ for a parameter that occurs only once in a body List(1, 2, 3).filter(\_ % 2 == 0) List(1, 2, 3).sortWith(\_ > \_)

### **Summary**

- Functional programming languages use function application, conditional expressions, recursion, and functional forms to control program execution instead of imperative features such as variables and assignments
- Purely functional languages have advantages over imperative alternatives, but their lower efficiency on existing machine architectures has prevented them from enjoying widespread use

#### References I

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- [4] Control Abstraction, http://www.artima.com/pins1ed/control-abstraction.html, 19 06 2014.