

Searching Arrays

LAST

Arrays
Safety

TODAY

Linear search
Correctness
Contract failure
Contract exploits

NEXT

Cost
Sorting arrays

Array search

```
int search(int x, int[] A, int n)
{
    for (i = 0; i < n; i++)
    {
        if (A[i] == x) return i;
    }
}
```

Is $A[i]$ safe?

What is a reasonable precondition?

safe but ...

```
int search(int x, int[] A, int n)
//@requires n == \length(A);
{
    for (int i = 0; i < n; i++)
        //@loop_invariant 0 <= i;
        {
            if (A[i] == x) return i;
        }
}
```

```
int search(int x, int[] A, int n)
//@requires n == \length(A);
{
    for (int i = 0; i < n; i++)
        //@loop_invariant 0 <= i;
        {
            if (A[i] == x) return i;
        }
}
```

What do we return if x is not found?

```

int search(int x, int[] A, int n)
//@requires n == \length(A);
{
    for (int i = 0; i < n; i++)
        //@loop_invariant 0 <= i;
        {
            if (A[i] == x) return i;
        }
    return -1;
}

```

Client using search:

```

... //allocate B
int i = search(12, B, 5);
if (i != -1) {
    B[i] = 13; //change 12 to 13
}
...
}

```

How do we know $B[i]$ is safe in the client's code?

The client sees only the contracts

```
int search(int x, int[] A, int n)
//@requires n == \length(A);
/*@ensures \result == -1 ||
           A[\result] == x);
  @*/
```

```
{
  for (int i = 0; i < n; i++)
    //@loop_invariant 0 <= i;
    {
      if (A[i] == x) return i;
    }
  return -1;
}
```

← What if A[\result] is unsafe?

```

int search(int x, int[] A, int n)
//@requires n == \length(A);
/*@ensures \result == -1 ||
           (0 <= \result && \result < n && A[\result] == x);
  @*/
{
  for (int i = 0; i < n; i++)
  //@loop_invariant 0 <= i;
  {
    if (A[i] == x) return i;
  }
  return -1;
}

```



We rely on short-short-circuiting
behavior of && to ensure the safety of A[result]

What is wrong with this code?

```
int search(int x, int[] A, int n)
//@requires n == \length(A);
/*@ensures \result == -1 ||
           (0 <= \result && \result < n && A[\result] == x);
  @*/
{
  return -1;
}
```

Contract exploit!

**A better contract for
correctness**

Return -1 **only if** x is not in **A[0,n)**

Example:

A

0	1	2	3	4	5
5	6	9	3	2	1

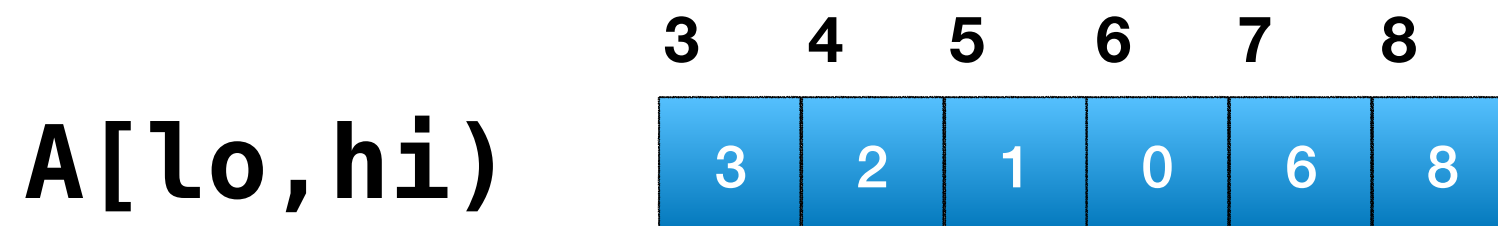
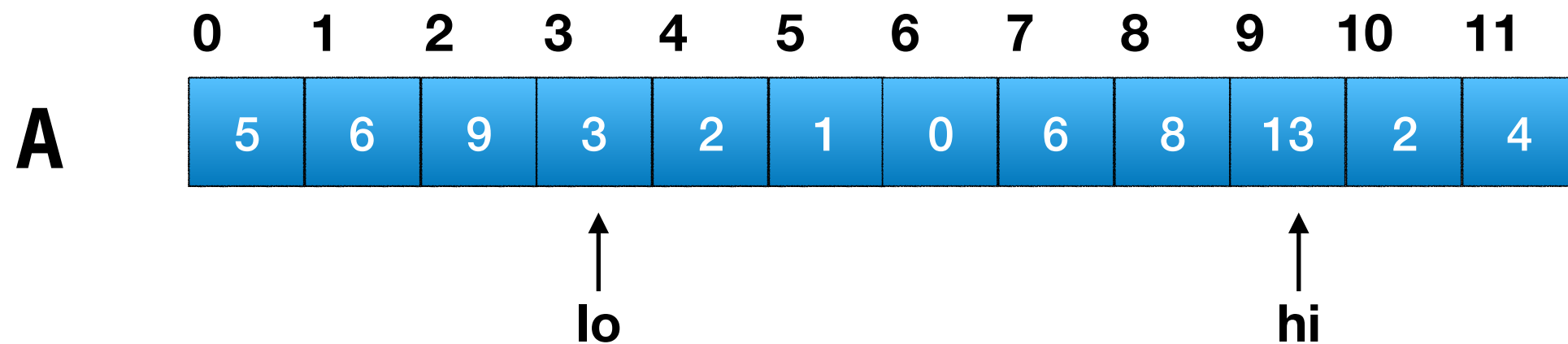
A[0,6)

0	1	2	3	4	5
5	6	9	3	2	1

A[0,3)

0	1	2
5	6	9

Array segments



lo == hi means an empty segment


$x \notin A[\text{lo}, \text{hi})$

```
bool is_in(int x, int[] A, int lo, int hi)
//@requires 0 <= lo && lo <= hi && hi == \length(A);
```

```
int search(int x, int[] A, int n)
//@requires n == \length(A);
/*@ensures (\result == -1 _____) ||
           (0 <= \result && \result < n && A[\result] == x);
  @*/
{

}
```

$x \notin A[0, n)$



```

bool is_in(int x, int[] A, int lo, int hi)
//@requires 0 <= lo && lo <= hi && hi <= \length(A);
{
    if (lo == hi) return false;
    return A[lo] == x || is_in(x, A, lo+1, hi);
}

int search(int x, int[] A, int n)
//@requires n == \length(A);
/*@ensures (\result == -1 && !is_in(x, A, 0, n)) ||
           (0 <= \result && \result < n && A[\result] == x);
@*/
{
    for (int i = 0; i < n; i++)
        //@loop_invariant 0 <= i;
        {
            if (A[i] == x) {
                return i;
            }
        }
    return -1;
}

```

Is the postcondition satisfied whenever the function returns?

Proving correctness


```

9int search(int x, int[] A, int n)
10//@requires n == \length(A);
11/*@ensures (\result == -1 && !is_in(x, A, 0, n)) ||
12
13           (0 <= \result && \result < n && A[\result] == x);
14  @*/
15{
16  for (int i = 0; i < n; i++)
17    //@loop_invariant 0 <= i;
18    {
19      if (A[i] == x) {
20        return i; ←
21      }
22    }
23  return -1; ←

```

return i

$0 \leq i$ (line 16, loop invariant)
 $i < n$ (line 15, loop guard)
 $A[i] == x$ (line 18, if condition)

return -1

How do we say logically that x is not in $A[0, n)$?
 Introduce a loop invariant?

```

9 int search(int x, int[] A, int n)
10 // @requires n == \length(A);
11 /* @ensures (\result == -1 && !is_in(x, A, 0, n)) ||
12          (0 <= \result && \result < n && A[\result] == x);
13  */
14 {
15     for (int i = 0; i < n; i++)
16         // @loop_invariant 0 <= i;
17         // @loop_invariant !is_in(x, A, 0, i);       $x \notin A[0,i]$ 
18         {
19             if (A[i] == x) {
20                 return i;
21             }
22         }
23     return -1;
24 }

```

INIT:

$!is_in(x, A, 0, 0)$ ($A[0,0]$ is empty)

PRES:

Assume $!is_in(x, A, 0, i)$ $x \notin A[0,i]$

Show $!is_in(x, A, 0, i')$ $x \notin A[0,i']$

$i' = i+1$ (line 15, loop increment)

So, we need to show $x \notin A[0,i+1]$, which means
 $!is_in(x, A, 0, i) \ \&\& \ A[i] \neq x$

If $A[i] == x$ (no proof obligation because we
 exit the loop)

Otherwise, immediate

Back to proving the postcondition

```
9 int search(int x, int[] A, int n)
10 //@requires n == \length(A);
11 /*@ensures (\result == -1 && !is_in(x, A, 0, n)) ||
12         (0 <= \result && \result < n && A[\result] == x);
13  */
14 {
15     for (int i = 0; i < n; i++)
16         //@loop_invariant 0 <= i;
17         //@loop_invariant !is_in(x, A, 0, i);
18     {
19         if (A[i] == x) {
20             return i;
21         }
22     }
23     return -1;
24 }
```

EXIT:

$!is_in(x, A, 0, i)$ (line 17, LI)

$i \geq n$ (negation of the loop guard)

But we need $i == n$.

Add $i \leq n$ as a loop invariant?

```

9 int search(int x, int[] A, int n)
10 //@requires n == \length(A);
11 /*@ensures (\result == -1 && !is_in(x, A, 0, n)) ||
12         (0 <= \result && \result < n && A[\result] == x);
13  */
14 {
15     for (int i = 0; i < n; i++)
16         //@loop_invariant 0 <= i && i <= n;
17         //@loop_invariant !is_in(x, A, 0, i);
18     {
19         if (A[i] == x) {
20             return i;
21         }
22     }
23     //@assert !is_in(x, A, 0, n);
24     return -1;
25 }

```

EXIT:

$\text{!is_in}(x, A, 0, i)$ (line 17, LI)
 $i \geq n$ (negation of the loop guard)
 $i \leq n$ (line 16, loop invariant)
 $i == n$ and $\text{!is_in}(x, A, 0, n)$

More contract exploits

```
int search(int x, int[] A, int n)
//@requires n == \length(A);
/*@ensures (\result == -1 && !is_in(x, A, 0, n)) ||
           (0 <= \result && \result < n && A[\result] == x);
@*/
{
    for (int i = 0; i < n; i++)
    {
        A[i] == x; // modify A[i]
        return i;  // return i
    }
    return -1;
}
```

```
int search(int x, int[] A, int n)
//@requires n == \length(A);
/*@ensures (\result == -1 && !is_in(x, A, 0, n)) ||
           (0 <= \result && \result < n && A[\result] == x);
  @*/
{
  for (int i = 0; i < n; i++)
  {
    A[i] == x+1;
  }
  return -1;
}
```