

INF2080

Oblig 1

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February 4, 2018

Problem 1: Regular languages

Let A and B be regular languages defined by DFAs \mathcal{A} and \mathcal{B} . Let $n_{\mathcal{A}}$ and $n_{\mathcal{B}}$ be the number of states in \mathcal{A} and \mathcal{B} , respectively.

Problem 1a

What are the worst-case (highest) number of states in **DFAs** for the languages $A \cap B$ and A^* ?

$A \cap B : n_{\mathcal{A}} * n_{\mathcal{B}}$

$A^* : n_{\mathcal{A}}$

Problem 1b

What are the worst-case (highest) number of states in **NFAs** for the languages $A \cap B$, AB and A^* ?

$A \cap B : n_{\mathcal{A}} * n_{\mathcal{B}}$

$AB : n_{\mathcal{A}} + n_{\mathcal{B}}$

$A^* : n_{\mathcal{A}} + 1$

Problem 1c

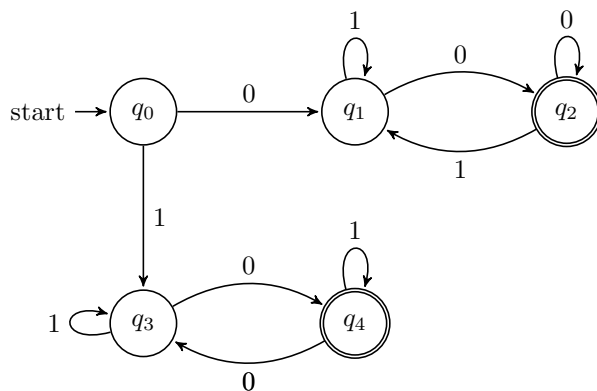
Create a regular expression defining the same language as the NFA

$(aa+)^*b(cc+)^*$

Problem 1d

Create a DFA for the language

$\{w \mid w \text{ contains equally many occurrences of the substrings } 01 \text{ and } 10\}$.



Problem 2: all-NFAs

An all-NFA is defined in Sipser, problem 1.43 as a 5-tuple $(Q, \Sigma, \delta, q_0, F)$ that accepts $x \in \Sigma^*$ if *every* possible state that M could reach after reading input x is in F (as opposed to *at least one*).

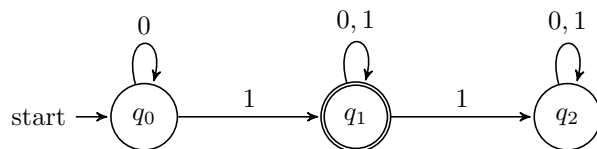
If any branch in an all-NFA computation reaches an implicit or explicit sink state, the input is not accepted.

Show how an all-NFA can be converted to an equivalent DFA.

Hint: Adjust the conversion from NFA to DFA shown in the lectures.

For this problem I will convert two NFA's, one all-NFA and one normal NFA. They both have the same states and transitions, but when converting to a DFA, the accepted states are different.

Here is the NFA:



The NFA:

$L = \{w \mid w \text{ contains at least one } 1 \text{ and an arbitrary number of } 0\text{s}\}$.

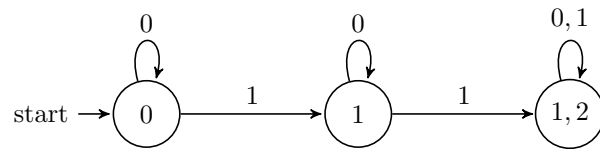
The all-NFA:

$$A = \{w \mid w \text{ contains exactly one 1 and an arbitrary number of 0s}\}.$$

The states in the DFA is the set of states you can go to from the NFA-state(s).
 The final states of a DFA converted from the NFA is the states that contain at least one final state from the NFA.

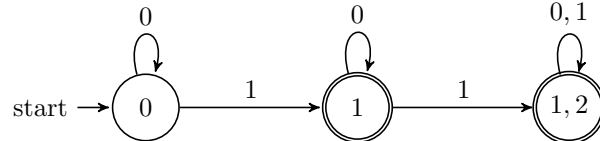
The final states of a DFA converted from the all-NFA is the states that contain only final-states from the all-NFA.

When converting the state diagram over to a DFA you get the following diagram:



The difference of converting from an all-NFA to a DFA and an NFA to a DFA is when you are choosing the accepting states. As stated above, when converting from an NFA, the DFA-state only has to have one of the accepting states in the state-set, whereas for it to be an all-NFA all of the states in the state-set in the DFA has to be an accepted state.

The DFA from the NFA will look like this:



Whereas the all-NFA conversion will look like this:

