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[(Graph)图论模板]

[Yucept]

Yucept 的图论模板

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Yucept 的图论模板

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网络流

算法模板

Dinic

```
struct Edge {
  int from, to, cap, flow;
  Edge(){}
  Edge(int _from, int _to, int _cap, int _flow)
  : from(_from), to(_to), cap(_cap), flow(_flow){}
};
struct Dinic {
  int n, m, s, t;
  Edge edges[maxm * 2];
  int head[maxn];
  int next[maxm * 2];
  bool inq[maxn];
  int d[maxn];
  int cur[maxn];
  void init(int n){
     this->n = n;
     m = 0;
     memset(head, -1, sizeof(head[0]) * (n + 1));
  void AddEdge(int from, int to, int cap) {
     next[m] = head[from];
```

```
edges[m] = Edge(from, to, cap, 0);
  head[from] = m ++; //
  next[m] = head[to];
  edges[m] = Edge(to, from, 0, 0);
  head[to] = m ++;
bool bfs() {
  memset(inq, false, sizeof(inq[0]) * (n + 1));
  queue<int> q;
  q.push(s);
  d[s] = 0; inq[s] = true;
  while(!q.empty()) {
     int u = q.front(); q.pop();
     for(int i = head[u]; i != -1; i = next[i]) {
       Edge& e = edges[i];
       int v = e.to;
       if(!inq[v] \&\& e.cap > e.flow) {
          inq[v] = true;
          d[v] = d[u] + 1;
          q.push(v);
          if(v == t) return true;
  return false;
```

```
int dfs(int u, int a) {
    if(u == t \parallel a == 0) return a;
    int flow = 0, f;
     for(int& i = cur[u]; i != -1; i = next[i]) {
       Edge& e = edges[i];
       int v = e.to;
       if(d[u] + 1 == d[v] && (f = dfs(v, min(a, e.cap - e.flow))) > 0) 
          e.flow += f:
          edges[i ^1].flow -= f;
          flow += f;
          a -= f;
          if(a == 0) break;
     return flow;
  int Maxflow(int s, int t) {
     this->s = s; this->t = t;
    int flow = 0;
     while(bfs()) {
       memcpy(cur, head, sizeof(head[0]) * (n + 1));
       flow += dfs(s, inf);
     return flow;
};
```

ISAP

```
struct ISAP {
  int n, m, s, t;
  Edge edges[maxm * 2];
  int head[maxn];
  int next[maxm * 2];
  bool inq[maxn];
  int d[maxn];
  int cur[maxn];
  int p[maxn];
  int num[maxn];
  void init(int n){
     this->n = n;
     m = 0:
     memset(head, -1, sizeof(head[0]) * (n + 1));
  void AddEdge(int from, int to, int cap) {
     next[m] = head[from];
     edges[m] = Edge(from, to, cap, 0);
     head[from] = m ++;
     next[m] = head[to];
     edges[m] = Edge(to, from, 0, 0);
     head[to] = m ++;
  bool bfs() {
     memset(inq, false, sizeof(inq[0]) * (n + 1));
```

```
queue<int> q;
  q.push(t);
  d[t] = 0; inq[t] = true;
  while(!q.empty()) {
    int u = q.front(); q.pop();
    for(int i = head[u]; i != -1; i = next[i]) {
       Edge& e = edges[i];
       int v = e.to;
       if(!inq[v] \&\& e.cap >= e.flow) {
          //printf("u = %d, v = %d\n", u, v);
          inq[v] = true;
          d[v] = d[u] + 1;
          q.push(v);
  return inq[s];
int Augment() {
  int x = t, a = \inf;
  while(x != s) {
    Edge& e = edges[p[x]];
    a = min(a, e.cap - e.flow);
    x = edges[p[x]].from;
  x = t;
```

```
while(x != s) {
     edges[p[x]].flow += a;
     edges[p[x] \land 1].flow -= a;
     x = edges[p[x]].from;
  return a;
int Maxflow(int s, int t) {
  this->s = s; this->t = t;
  int flow = 0;
  bfs();
  memset(num, 0, sizeof(num[0]) * (n + 1));
  for(int i = 0; i < n; i ++) num[d[i]] ++;
  int x = s;
  memcpy(cur, head, sizeof(head[0]) * (n + 1));
  while (d[x] < n) {
    if(x == t) {
       flow += Augment();
       x = s;
     bool ok = false;
     for(int i = cur[x]; i != -1; i = next[i]) {
       Edge& e = edges[i];
       if(e.cap > e.flow && d[x] == d[e.to] + 1) {
          ok = true;
          p[e.to] = i;
```

```
cur[x] = i;
            x = e.to;
            break;
       if(!ok) {
         int mt = n - 1;
         for(int i = head[x]; i != -1; i = next[i]) {
           Edge& e = edges[i];
           if(e.cap > e.flow) mt = min(mt, d[e.to]);
         if(-- num[d[x]] == 0) break; //gap 忧化
         num[d[x] = mt + 1] ++;
         cur[x] = head[x];
         if(x != s) x = edges[p[x]].from;
    return flow;
};
普通费用流
struct MCMF {
  int n, m, s, t;
  Edge edges[maxm * 2];
  int head[maxn];
  int next[maxm * 2];
```

```
bool inq[maxn];
int d[maxn];
int p[maxn];
int a[maxn];
void init(int n) {
  this->n = n;
  m = 0;
  memset(head, -1, sizeof(head[0]) * (n + 1));
void AddEdge(int from, int to, int cap, int cost) {
  next[m] = head[from];
  edges[m] = Edge(from, to, cap, 0, cost);
  head[from] = m ++;
  next[m] = head[to];
  edges[m] = Edge(to, from, 0, 0, -cost);
  head[to] = m ++;
bool BellmanFord(int s, int t, int& flow, int& cost) {
  for(int i = 0; i < n; i ++) d[i] = inf;
  memset(inq, false, sizeof(inq[0]) * (n + 1));
  d[s] = 0; inq[s] = true; p[s] = 0; a[s] = inf;
  queue<int> q;
  q.push(s);
  while(!q.empty()) {
     int u = q.front(); q.pop();
     inq[u] = false;
```

```
for(int i = head[u]; i != -1; i = next[i]) {
       Edge& e = edges[i]; int v = e.to;
       if(e.cap > e.flow && d[v] > d[u] + e.cost) {
          d[v] = d[u] + e.cost;
          p[v] = i;
          a[v] = min(a[u], e.cap - e.flow);
          if(!inq[v]) \{q.push(v); inq[v] = true;\}
  if(d[t] == inf) return false;
  flow += a[t];
  cost += d[t] * a[t];
  int u = t;
  while(u != s) {
     edges[p[u]].flow += a[t];
    edges[p[u] ^1].flow -= a[t];
    u = edges[p[u]].from;
  return true;
int Mincost(int s, int t) {
  int flow = 0, cost = 0;
  while(BellmanFord(s, t, flow, cost));
  return cost;
```

```
ZKW(高效)费用流
struct ZKW_flow{
  int st, ed, n, m;
  int head[maxn];
  int cap[maxm], cost[maxm], to[maxm], next[maxm];
  void init(int n){
     this->n = n:
     memset(head, -1, sizeof(head[0])* (n + 1));
    m = 0;
  void AddEdge(int u, int v, int c, int w) {
     cap[m] = c; cost[m] = w; to[m] = v;
     next[m] = head[u]; head[u] = m ++;
    cap[m] = 0; cost[m] = -w; to[m] = u;
     next[m] = head[v]; head[v] = m++;
  int dis[maxn];
  void Dijkstra() {
     for(int i = 0; i \le n; ++i) dis[i] = inf;
     priority_queue<pii, vector<pii>, greater<pii>> Q;
     dis[st] = 0;
     Q.push(make_pair(0, st));
     while(!Q.empty()){
       int u = Q.top().second, d = Q.top().first;
       Q.pop();
```

```
if(dis[u]!=d) continue;
    for(int p = head[u]; p != -1; p = next[p]){
       int &v = to[p];
       if(cap[p] \&\& dis[v] > d + cost[p]){
         dis[v] = d + cost[p];
         Q.push(make_pair(dis[v], v));
  for(int i = 0; i \le n; ++ i) dis[i] = dis[ed] - dis[i];
int MinCost, MaxFlow;
bool used[maxm];
int Add_Flow(int u, int flow){
  if(u == ed)
    MaxFlow += flow;
    MinCost += dis[st] * flow;
    return flow;
  used[u] = true;
  int now = flow;
  for(int p = head[u]; p != -1; p = next[p])
    int &v = to[p];
    if(cap[p] \&\& !used[v] \&\& dis[u] == dis[v] + cost[p]){
       int tmp = Add_Flow(v, min(now, cap[p]));
       cap[p] = tmp;
```

```
cap[p^1] += tmp;
        now = tmp;
       if(!now) break;
  return flow - now;
bool modify_label(){
  int d = inf;
  for(int u = 0; u \le n; ++ u) if(used[u])
    for(int p = head[u]; p != -1; p = next[p]){
       int &v = to[p];
       if(cap[p] \&\& !used[v]) d = min(d, dis[v] + cost[p] - dis[u]);
  if(d == inf) return false;
  for(int i = 0; i \le n; ++ i) if(used[i]) dis[i] += d;
  return true;
int MCF(int s, int t){
  st = s, ed = t;
  MinCost = MaxFlow = 0;
  Dijkstra();
  while(true){
     while(true){
       for(int i = 0; i \le n; ++ i) used[i] = false;
       if(!Add_Flow(st, inf)) break;
```

```
if(!modify_label()) break;
    FLOW = MaxFlow; //定义的 FLOW 为最终求得的流量
    return MinCost;
};
建模
矩阵取值问题 (费用流)
1.矩阵的每个格子里有一个数字,从左上角走到右下角,走 k 次,
求最大的数字之和。(格子里的数字第一次走过后变成0)。
MCMF mc:
int N, K;
int src, sink;
int mat[55][55];
void Prepare() {
  int Nt = N * N;
  mc.init(2 * Nt + 2);
  src = 0; sink = 2 * Nt + 1;
  for(int i = 1; i \le N; i ++) {
   for(int j = 1; j \le N; j ++) {
     scanf("%d", &mat[i][j]); //每个格子的数值
  for(int i = 1; i \le N; i ++) {
```

```
for(int j = 1; j \le N; j ++) {
      int u = (i - 1) * N + j, v;
      mc.AddEdge(u, u + Nt, 1, -mat[i][j]); //拆点,每个点只有第一次有值
      mc.AddEdge(u, u + Nt, K - 1, 0); //剩下 K-1 次的值为 0
      if(j < N) {
        v = u + 1;
        mc.AddEdge(u + Nt, v, K, 0);
      if(i < N) {
        v = u + N;
        mc.AddEdge(u + Nt, v, K, 0);
  mc.AddEdge(src, 1, K, 0); //超级源点到左上角 1 号点,容量为 K,限制 k
  mc.AddEdge(2 * Nt, sink, K, 0);
int main()
  while(scanf("%d%d", &N, &K) != EOF) {
    Prepare();
    int ans = mc.Mincost(src, sink);
    printf("%d\n", -ans); //这是一个求最大的问题, 所以要将格子的值取相
反数之后求最小费用再将最后的结果求相反数。
```

```
return 0;
最大流+二分
二分一般是用来限制源点流出的流量。
对于点有流量限制(比如只能走一次)可以拆点变成边来限制
拆点要准确,注意入点和出点。
*/
ISAP sap;
int Tot;
int N, M, K, src, sink;
bool like[maxn/4][maxn/4];
bool check(int cap) {
  sap.init(4 * N + 2);
  for(int i = 1; i \le N; i ++) {
    for(int j = 1; j \le N; j ++) {
      if(like[i][j]) sap.AddEdge(i, j + 3 * N, 1);
      else sap.AddEdge(i + N, j + 2 * N, 1);
  for(int i = 1; i \le N; i ++) {
    sap.AddEdge(i, i + N, K);
    sap.AddEdge(i + 2 * N, i + 3 * N, K); //点的限制为 K
    sap.AddEdge(src, i, cap); //源点流出的限制
    sap.AddEdge(i + 3 * N, sink, cap);
```

```
int tflow = sap.Maxflow(src, sink);
  //printf("cap = %d, flow = %d\n", cap, tflow);
  if(tflow == N * cap) return true; //判断满流
  return false;
void Prepare() {
  scanf("%d%d%d", &N, &M, &K);
  memset(like, false, sizeof(like));
  int u, v;
  while(M --) {
     scanf("%d%d", &u, &v);
     like[u][v] = true;
  src = 0, sink = 4 * N + 1;
int main()
  //freopen("input.txt", "r", stdin);
  scanf("%d", &Tot);
  for(int ac = 1; ac \leq Tot; ac ++) {
     Prepare();
     int l = 0, r = N, ans;
     while(l \le r) {
       ans = (1 + r) >> 1;
       if(check(ans)) l = ans + 1;
       else r = ans - 1;
```

```
printf("%d\n", 1 - 1);
  return 0;
最大权闭合子图
N个点需要一定费用来建设, M条边可以获得一定收益
连边: N 个点为花费点, M 条边为获利点
src 向获利点连边,容量为 profit
获利点向与其相关的花费点连边,容量为 inf
花费点向 sink 连边,容量为 cost
void Prepare() {
 src = 0; sink = N + M + 1;
 Di.init(N + M + 2);
 for(int i = 1; i \le N; i ++) {
    scanf("%d", &cost);
   Di.AddEdge(i, sink, cost);
 int u, v, pf;
 ans = 0;
  for(int i = 1; i \le M; i ++) {
    scanf("%d%d%d", &u, &v, &pf);
    ans += pf;
    Di.AddEdge(i + N, u, inf);
```

```
Di.AddEdge(i + N, v, inf);
    Di.AddEdge(src, i + N, pf);
int main()
  while(scanf("%d%d", &N, &M)!= EOF) {
    Prepare();
    printf("%d\n", ans - Di.Maxflow(src, sink));
  return 0;
动态流+分层拆点
分层拆点要注意连边不出现问题
出点向入点连边
*/
Dinic Di;
int Tot;
int N, M, src, sink;
vector<int> G[60];
int D, W, C[60];
struct Pt{
  int x, y;
Pt pt[60];
```

```
int sqr(int x) {
  return x * x;
bool Near(Pt a, Pt b) {
  return sqr(a.x - b.x) + sqr(a.y - b.y) \le D * D;
void Prepare() {
  for(int i = 1; i \le N; i ++) {
     scanf("%d%d%d", &pt[i].x, &pt[i].y, &C[i]);
  for(int i = 1; i \le N; i ++) {
    for(int j = i + 1; j \le N; j ++) {
       if(Near(pt[i], pt[j])) {
         G[i].push_back(j); G[j].push_back(i);
  src = 0; sink = 2*(N + M) * N + 1;
  Di.init(sink + 1);
void conect(int cur) {
  int u, v;
  for(int i = 1; i \le N; i ++) {
    u = 2 * cur * N + i, v = u + N;
    Di.AddEdge(u, v, C[i]); //对每层的点拆点,表示点的容量限制
     if(pt[i].y <= D) { //超级源能到的点连边
```

```
Di.AddEdge(src, u, M);
     if(pt[i].y + D >= W) { //能到超级汇的连边
       Di.AddEdge(v, sink, M);
  if(cur) {
    for(int i = 1; i \le N; i ++) {
       u = 2 * cur * N + i;
       int sz = G[i].size();
       for(int j = 0; j < sz; j ++) {
         v = (2 * cur - 1) * N + G[i][j];
         Di.AddEdge(v, u, inf); //上一层的出点到下一层的入点连边
  }}
int main()
  //freopen("input.txt", "r", stdin);
  while(scanf("%d%d%d%d", &N, &M, &D, &W) != EOF) {
    Prepare();
     if(D \ge W) \{ printf("1\n"); continue; \}
     if(N == 0) {printf("IMPOSSIBLE\n"); continue;}
     int cur, ans = 0;
     for(cur = 0; cur < N + M; cur ++) {
       conect(cur);
```

```
ans += Di.Maxflow(src, sink);
       if(ans \geq M) break;
    if(cur == N + M) printf("IMPOSSIBLE\n");
    else printf("%d\n", cur + 2);
  return 0;
混合图欧拉回路
ISAP sap;
int Tot;
int N, M, src, sink;
int din[maxn], dout[maxn];
void Prepare() {
  scanf("%d%d", &N, &M);
  memset(din, 0, sizeof(din));
  memset(dout, 0, sizeof(dout));
  int u, v, kind;
  sap.init(N + 2);
  src = 0, sink = N + 1;
  while(M --) {
    scanf("%d%d%d", &u, &v, &kind);
    din[v] ++, dout[u] ++;
    if(!kind) sap.AddEdge(u, v, 1);
```

```
bool check() {
  int tflow = 0;
  for(int i = 1; i \le N; i ++) {
     if((din[i] + dout[i]) & 1) {
        return false;
  for(int i = 1; i \le N; i ++) {
     int x = abs(din[i] - dout[i]) / 2;
     if(din[i] > dout[i])
       sap.AddEdge(i, sink, x);
        tflow += x;
     if(dout[i] > din[i])
        sap.AddEdge(src, i, x);
  if(tflow == sap.Maxflow(src, sink)) return true;
  return false;
int main()
  //freopen("input.txt", "r", stdin);
  scanf("%d", &Tot);
  for(int ac = 1; ac <= Tot; ac ++) {
     Prepare();
     bool ok = check();
```

```
if(!ok) printf("impossible\n");
  else printf("possible\n");
}
return 0;
}
```

最短路

算法模板

Dijkstra

```
struct HeapNode {
  int d, u;
  HeapNode(){}
  HeapNode(int _d, int _u)
  : d(_d), u(_u) {}
  bool operator < (const HeapNode& rhs) const {
    return d > rhs.d;
  }
};
struct Dijkstra {
  int n, m;
  Edge edges[maxm];
  int next[maxm];
  int head[maxn];
  int d[maxn];
  int layer[maxn];
```

```
bool done[maxn];
void init(int n) {
  this->n = n;
  m = 0;
  memset(head, -1, sizeof(head[0]) * (n + 1));
void AddEdge(int from, int to, int dist) {
  next[m] = head[from];
  edges[m] = Edge(from, to, dist);
  head[from] = m ++;
void dijkstra(int src, int des) {
  priority_queue<HeapNode> q;
  for(int i = 1; i \le n; i ++) d[i] = inf;
  d[src] = 0;
  memset(done, false, sizeof(done[0]) * (n + 1));
  q.push(HeapNode(0, src));
  while(!q.empty()) {
     HeapNode x = q.top(); q.pop();
     int u = x.u;
     //printf("%d %d\n", x.u, x.d);
     if(done[u]) continue;
     done[u] = true;
     for(int i = head[u]; i != -1; i = next[i]) {
       Edge& e = edges[i];
       int v = e.to;
```

```
if(d[v] > d[u] + e.dist) {
             d[v] = d[u] + e.dist;
            // printf("u = %d, v = %d, d[v] = %d\n", u, v, d[v]);
             q.push(HeapNode(d[v], v));
};
Floyd
typedef long long LL;
const LL inf = -1;
LL d[maxn][maxn];
void floyd(int n) {
     for(int k = 1; k \le n; k ++) {
          for(int i = 1; i \le n; i ++) {
               for(int j = 1; j \le n; j ++) {
                     if(d[i][k] == inf || d[k][j] == inf) continue;
                     if(d[i][j] > d[i][k] + d[k][j] \parallel d[i][j] == inf)
                          d[i][j] = d[i][k] + d[k][j];
SPFA
/*
```

有向边,求单源点到所有点的最短路,直接 SPFA 求多源点到单终点的最短路,反向建图,就变成 单源点到多终点的最短路问题

```
*/
struct Edge{
  int from, to, dist;
  Edge(){}
  Edge(int _from, int _to, int _dist)
  : from(_from), to(_to), dist(_dist){}
};
struct SPFA {
  int n, m;
  Edge edges[maxm];
  int next[maxm];
  int height[maxm]; //某些题目的限制条件
  int head[maxn];
  int inq[maxn];
  int d[maxn];
  int cnt[maxn];
  void init(int n) {
     this->n = n;
     m = 0;
     memset(head, -1, sizeof(head[0]) * (n + 1));
  void AddEdge(int from, int to, int dist) {
     next[m] = head[from];
```

```
edges[m] = Edge(from, to, dist);
  head[from] = m ++;
int spfa(int src, int des, int limit) {
  queue<int>q;
  for(int i = 0; i \le n; i ++){
     d[i] = \inf; inq[i] = false;
  inq[src] = true; d[src] = 0;
  q.push(src);
  while(!q.empty()){
     int u = q.front(); q.pop();
     inq[u] = false;
     for(int i = head[u]; i != -1; i = next[i]){
       Edge& e = edges[i];
       int v = e.to;
       if(height[i] < limit) continue;</pre>
       if(d[v] > d[u] + e.dist){
          d[v] = d[u] + e.dist;
          if(!inq[v]){
             inq[v] = true; q.push(v);
  if(d[des] == inf) return -1;
```

```
return d[des];
一些题型
求最短路次短路条数
/***
求 s 到 t 的最短路与次短路(这里要求只比最短路多 1)的条数之和
 联想到最小,次小的一种更新关系:
if(x<最小)更新最小,次小
else if(==最小)更新方法数
else if(x<次小)更新次小
else if(x==次小)更新方法数
 同时记录 s 到 u 最短,次短路及方法数
 用一个堆每次取最小的, 更新完后再入堆
还是那个原理,第一次遇到的就是最优的,然后 done 标记为真
方法数注意是加法原理, 不是乘法
***/
struct HeapNode{
 int d, u;
 int mark; //代表种类, 次短为 1, 最短为 0
 HeapNode(){}
 HeapNode(int _d, int _u, int _mark)
 : d(_d), u(_u), mark(_mark){}
 bool operator < (const HeapNode& rhs) const{
   return rhs.d < d;
```

```
};
struct Edge{
  int from, to, dist;
  Edge(){}
  Edge(int _from, int _to, int _dist)
  : from(_from), to(_to), dist(_dist){}
};
int n, m, ans, s, t;
vector<int> G[maxn];
vector<Edge> edges;
bool done[maxn][2];
//d[t][0]是最短路的值, d[t][1]是次短路的值
int d[maxn][2], c[maxn][2];
//c[t][0] 最短路条数, c[t][1] 次短路条数
void init(){
  edges.clear();
  for(int i = 0; i \le n; i + +)
    G[i].clear();
void AddEdge(int from, int to, int dist){
  edges.push_back(Edge(from, to, dist));
  int m = edges.size();
  G[from].push_back(m - 1);
```

```
void Dijkstra(int s, int t){
  priority_queue<HeapNode> q;
  for(int j = 0; j < 2; j ++)
     for(int i = 0; i \le n; i + +)
       d[i][j] = inf; done[i][j] = false; c[i][j] = 0;
  d[s][0] = 0; c[s][0] = 1;
  q.push(HeapNode(d[s][0], s, 0));
  while(!q.empty()){
     HeapNode x = q.top(); q.pop();
     int u = x.u, mark = x.mark;
     if(done[u][mark]) continue;
     done[u][mark] = true;
     int sz = G[u].size();
     for(int i = 0; i < sz; i ++){
       Edge& e = edges[G[u][i]];
       int v = e.to;
       int dis = d[u][mark] + e.dist;
       if(d[v][0] > dis){
          if(d[v][0] != inf){
            d[v][1] = d[v][0];
            c[v][1] = c[v][0];
            q.push(HeapNode(d[v][1], v, 1));
          d[v][0] = dis;
          c[v][0] = c[u][mark];
```

```
q.push(HeapNode(d[v][0], v, 0));
       else if(d[v][0] == dis)
         c[v][0] += c[u][mark];
       else if(d[v][1] > dis)
         d[v][1] = dis;
         c[v][1] = c[u][mark];
         q.push(HeapNode(d[v][1], v, 1));
       else if(d[v][1] == dis)
         c[v][1] += c[u][mark];
int main()
  int Tot;
  scanf("%d", &Tot);
  while(Tot --){
    scanf("%d%d", &n, &m);
    init();
    int u, v, w;
    while(m --){
       scanf("%d%d%d", &u, &v, &w);
       AddEdge(u, v, w);
    scanf("%d%d", &s, &t);
```

```
Dijkstra(s, t);
     if(d[t][0] + 1 == d[t][1])
       ans = c[t][0] + c[t][1];
     else ans = c[t][0];
     printf("%d\n", ans);
  return 0;
A star 求 k 短路
起点和终点一样时要将 K+1
struct HeapNode{
  int d, u;
  HeapNode(){}
  HeapNode(int _d, int _u)
  : d(_d), u(_u)\{\}
  bool operator < (const HeapNode& rhs) const{</pre>
     return rhs.d < d;
};
struct Edge{
  int from, to, dist;
  Edge(){}
  Edge(int _from, int _to, int _dist)
  : from(_from), to(_to), dist(_dist){}
```

```
};
struct KSP {
  int n, m;
  int Ghead[maxn];
  int Fhead[maxn];
  int next[maxm * 2]; //边的数组可以尽量开大些
  Edge edges[maxm * 2];
  int cnt[maxn], d[maxn];
  void init(int n){
    this->n = n;
    m = 0;
    memset(Ghead, -1, sizeof(Ghead[0]) * (n + 1));
    memset(Fhead, -1, sizeof(Fhead[0]) * (n + 1));
  void AddEdge(int from, int to, int dist){
    next[m] = Ghead[from];
    edges[m] = Edge(from, to, dist);
    Ghead[from] = m ++;
    next[m] = Fhead[from];
    edges[m] = Edge(to, from, dist);
    Fhead[to] = m ++;
  void Dijkstra(int src){
    priority_queue<HeapNode> q;
    for(int i = 0; i \le n; i + +)
       d[i] = \inf; cnt[i] = 0;
```

```
d[src] = 0;
  q.push(HeapNode(d[src], src));
  while(!q.empty()){
    HeapNode x = q.top(); q.pop();
     int u = x.u;
     if(cnt[u]) continue;
     cnt[u] = 1;
     for(int i = \text{Fhead}[u]; i != -1; i = \text{next}[i]){
       Edge& e = edges[i];
       int v = e.to;
       if(d[v] > d[u] + e.dist)
          d[v] = d[u] + e.dist;
          q.push(HeapNode(d[v], v));
int Kth_Astar(int src, int des, int k){
  priority_queue<HeapNode> q;
  if(d[src] == inf) return -1;
  q.push(HeapNode(d[src], src));
  memset(cnt, 0, sizeof(cnt[0]) * (n + 1));
  while(!q.empty()){
     HeapNode x = q.top(); q.pop();
     int u = x.u;
```

```
cnt[u] ++;
       if(cnt[u] > k) continue;
       if(cnt[des] == k) return x.d;
       for(int i = Ghead[u]; i != -1; i = next[i]){
         Edge& e = edges[i];
         int v = e.to;
         q.push(HeapNode((x.d - d[u] + e.dist + d[v]), v));
    return -1;
};
int N, M, K, S, T;
KSP ks;
int main()
  while(scanf("%d%d", &N, &M) != EOF){
    ks.init(N);
    int u, v, w;
    while(M --)
       scanf("%d%d%d", &u, &v, &w);
       ks.AddEdge(u, v, w);
    scanf("%d%d%d", &S, &T, &K);
    if(S == T) K ++;
    ks.Dijkstra(T);
```

```
int ans = ks.Kth_Astar(S, T, K);
    printf("%d\n", ans);
  return 0;
二维最短路
用 D[U][FUEL] 表示走到 U 点剩 FUEL 油量时的最小费用
*/
struct HeapNode{
  int d, u, fuel;
  bool operator < (const HeapNode& rhs) const{
    return rhs.d < d;
struct Edge{
  int from, to, dist;
};
int n, m;
vector<int> G[maxn];
vector<Edge> edges;
bool done[maxn];
int d[maxn][maxc], p[maxn];
int f[maxn];
void init(){
  edges.clear();
```

```
for(int i = 0; i \le n; i ++){
     G[i].clear();
void addedge(int from, int to, int dist){
  edges.push_back((Edge){from, to, dist});
  edges.push_back((Edge){to, from, dist});
  int m = edges.size();
  G[from].push_back(m - 2);
  G[to].push_back(m - 1);
int Dijkstra(int s, int t, int cap){
  priority_queue<HeapNode> q;
  for(int i = 0; i \le n; i + +)
    for(int j = 0; j \le cap; j ++){
       d[i][j] = inf;
  d[s][0] = 0;
  q.push((HeapNode)\{d[s][0], s, 0\});
  while(!q.empty()){
    HeapNode x = q.top(); q.pop();
     int u = x.u;
     if(u == t) return x.d;
    //每次加一单位的油扩展结点
    int fuel = x.fuel + 1;
```

```
int cost = x.d + p[u];
     if(fuel \le cap \&\& d[u][fuel] > cost)
        q.push((HeapNode){cost, u, fuel});
        d[u][fuel] = cost;
       //printf("u = %d, fuel = %d\n", u, fuel);
     int SIZE = G[u].size();
     for(int i = 0; i < SIZE; i ++){
       Edge& e = edges[G[u][i]];
       int v = e.to;
       if(x.fuel \ge e.dist \&\& d[v][x.fuel - e.dist] > x.d){
          d[v][x.fuel - e.dist] = x.d;
          q.push((HeapNode){d[v][x.fuel - e.dist], v, x.fuel - e.dist});
          //printf("v = %d, fuel = %d\n", v, x.fuel - e.dist);
  return -1;
int main()
  while(scanf("%d%d", &n, &m) != EOF){
     for(int i = 0; i < n; i ++) scanf("%d", &p[i]);
     while(m --){
        int u, v, d;
```

```
vector<int> G[maxn];
      scanf("%d%d%d", &u, &v, &d);
                                                                             vector<Edge> edges;
      addedge(u, v, d);
                                                                             void init(){
    int q;
                                                                                 edges.clear();
    scanf("%d", &q);
                                                                                 for(int i = 0; i <= n; i ++)
    while(q --){
                                                                                      G[i].clear(); pre[i] = -1;
      int c, s, e;
      scanf("%d%d%d", &c, &s, &e);
                                                                             void addedge(int from, int to, int weight){
      int ans = Dijkstra(s, e, c);
      if(ans != -1) printf("%d\n", ans);
                                                                                 edges.push_back((Edge){from, to, weight});
      else printf("impossible\n");
                                                                                 edges.push_back((Edge){to, from, weight});
                                                                                 int m = edges.size();
                                                                                 G[from].push_back(m - 2);
  return 0:
                                                                                 G[to].push_back(m - 1);
删除最短路上的一条边再求最短路
                                                                             int spfa(int s, int t, int pt, int num){ //pt 为 2 代表删边后求最短路
//要删除一条最短路上的边使从 1 到 n 的时间最长甚至不可达。
                                                                                 queue<int> q;
//枚举每条边,然后找最坏情况。
                                                                                 memset(ing, false, sizeof ing);
                                                                                 for(int i = 0; i \le n; i ++) d[i] = inf;
struct Edge{
                                                                                 inq[s] = true; d[s] = 0;
    int from, to, weight;
                                                                                 q.push(s);
};
                                                                                 while(!q.empty()){
                                                                                      int u = q.front(); q.pop();
int n, m;
int inq[maxn], d[maxn];
                                                                                      inq[u] = false;
int pre[maxn], pe[maxn]; //pre[]是当前点在最短路中的前置点, pe[]存的是边
                                                                                      for(int i = 0; i < G[u].size(); i ++){
的编号
                                                                                          Edge& e = edges[G[u][i]];
```

```
int v = e.to;
              if(num == G[u][i]) continue;
              if(d[u] + e.weight < d[v])
                   d[v] = d[u] + e.weight;
                   if(pt == 1) {
                         pre[v] = u;
                         pe[v] = G[u][i];
                   if(!inq[v]){
                        q.push(v);
                        inq[v] = true;
    if(d[t] == inf) return -1;
    return d[t];
int main()
    int T;
     scanf("%d", &T);
     while(T --){
         scanf("%d%d", &n, &m);
         init();
         while(m --){
```

```
int u, v, w;
              scanf("%d%d%d", &u, &v, &w);
              addedge(u, v, w);
         int s = 1, t = n;
         int ans = spfa(s, t, 1, 2 * m + 1);
         if(ans == -1) {
              printf("-1\n");
              continue;
         int u = t;
         while(u != s){
              int res = spfa(s, t, 2, pe[u]);
              if(res == -1){ //如果删除某条边不可达,则 ans = -1
                  ans = -1;
                   break;
              ans = max(ans, res);
              u = pre[u];
         printf("%d\n", ans);
    return 0;
二分枚举+最短路
SPFA sp;
```

```
int N, M;
int Tot;
int maxh, ans;
int main()
  int Tot;
  scanf("%d", &Tot);
  for(int ac = 1; ac \leq Tot; ac ++){
    scanf("%d%d", &N, &M);
    sp.init(N);
    if(ac > 1) printf("\n");
    maxh = 0;
    int u, v, h, d;
    while(M --)
       scanf("%d%d%d", &u, &v, &h);
       maxh = max(h, maxh);
       d = 1;
       sp.height[sp.m] = h;
       sp.AddEdge(u, v, d);
       sp.height[sp.m] = h;
       sp.AddEdge(v, u, d);
    int L = 0, R = maxh, mid;
    while(L \le R)
       mid = (L + R) / 2;
       ans = sp.spfa(1, N, mid);
```

```
if(ans == -1) R = mid - 1;
      else L = mid + 1;
    printf("Scenario #%d:\n", ac);
    printf("%d\n", R);
  return 0;
分层最短路
N 个点, 然后有 N 层, 要假如 2*N 个点。
总共是 3*N 个点。
点 1~N 就是对应的实际的点 1~N. 要求的就是 1 到 N 的最短路。
然后点 N+1~3*N 是 N 层拆出出来的点。
然后用优先队列优化的 Dijkstra 就可以搞出最短路了
*/
Dijkstra sp;
void Prepare() {
  scanf("%d%d%d", &N, &M, &C);
  for(int i = 1; i \le N; i ++) {G[i].clear();}
  for(int i = 1; i \le N; i ++) {
    scanf("%d", &sp.layer[i]);
    G[sp.layer[i]].push_back(i);
  sp.init(3*N);
  //printf("cnt = %d\n", cnt);
```

```
for(int i = 1; i \le N; i + +) {
         int sz = G[i].size();
         for(int j = 0; j < sz; j ++) {
             int v = G[i][j];
             //printf("i = %d, v = %d\n", i, v);
              sp.AddEdge(2*N + i, v, 0); //v 属于第 i 层
             sp.AddEdge(v, i + N, 0);
    /***注意上面层的入点是下面层的出点******/
    for(int i = 2; i \le N; i ++) {
         sp.AddEdge(N + i, 2*N + (i-1), C); //第 i 层到第 i-1 层
         sp.AddEdge(N + (i-1), 2*N + i, C); //第 i-1 层到第 i 层
  int u, v, d;
  while(M --)
    scanf("%d%d%d", &u, &v, &d);
    sp.AddEdge(u, v, d);
    sp.AddEdge(v, u, d);
int main()
  //freopen("input.txt", "r", stdin);
  int Tot:
  scanf("%d", &Tot);
```

```
for(int ac = 1; ac \leftarrow Tot; ac ++){
     Prepare();
     printf("Case #%d: ", ac);
     sp.dijkstra(1, N);
     ans = sp.d[N];
     if(ans < inf) printf("%d\n", ans);
         else printf("-1\n");
  return 0;
生成树&并查集
种类并查集 (POJ 1182)
int pa[maxn], ra[maxn];
const int same = 0; //同类
const int enemy = 1; //被父节点吃
const int food = 2; //吃父节点
int n, m;
int findset(int x){
     if(pa[x] == x) return x;
     int tmp = pa[x];
     pa[x] = findset(pa[x]);
    ra[x] = (ra[tmp] + ra[x]) \% 3;
    return pa[x];
void init()
```

```
for(int i = 0; i \le n; i ++){
         pa[i] = i; ra[i] = same;
void merge(int x, int y, int u, int v, int d){
    pa[v] = u;
    ra[v] = (3 - ra[y] + (d - 1) + ra[x]) \% 3;
    // 3 - ra[y]是 v 和 y 的关系 d - 1 是 y 和 x 的关系
    // ra[x]是 x 和 u 的关系, 得出 v 和 u 的关系
int main()
    int d, x, y;
    int ans;
    scanf("%d%d", &n, &m);
         init();
         ans = 0;
         while(m --){
              scanf("%d%d%d", &d, &x, &y);
              if(x > n || y > n)
                   ans ++; continue;
              if(d == 2 \&\& x == y) {
                   ans ++; continue;
```

```
int u = findset(x), v = findset(y);
              if(u != v)
                   merge(x, y, u, v, d);
              else {
                   if(d == 1){
                        if(ra[x] != ra[y]) ans ++;
                   if(d == 2){
                        if((ra[y] + 3 - ra[x]) \% 3! = 1) ans ++;
         printf("%d\n", ans);
     return 0;
并查集求集合个数 (倒过来用)
int pa[maxn], ans[maxn], e[maxn];
bool mark[maxn];
struct Edge{
     int u, v;
}edges[maxn];
int findset(int x) {
     return pa[x] == x ? x : (pa[x] = findset(pa[x]));
```

```
int N, M, Q;
int main()
     while(scanf("%d%d", &N, &M) != EOF) {
          for(int i = 1; i \le N; i ++) pa[i] = i;
          for(int i = 1; i \le M; i ++) {
               scanf("%d%d", &edges[i].u, &edges[i].v);
               mark[i] = false;
          scanf("%d", &Q);
          for(int i = 1; i \le Q; i ++) {
               int id;
               scanf("%d", &id);
               mark[id] = true;
               e[i] = id;
          int cnt = N;
          for(int i = 1; i \le M; i ++) {
               if(!mark[i]) {
                    int x = findset(edges[i].u), y = findset(edges[i].v);
                    if(x != y) 
                         pa[x] = y;
                         cnt --;
```

```
for(int i = Q; i >= 1; i --) {
               ans[i] = cnt;
               int x = findset(edges[e[i]].u), y = findset(edges[e[i]].v);
               if(x != y) {
                    pa[x] = y;
                    cnt --;
          for(int i = 1; i < Q; i ++)
               printf("%d ", ans[i]);
          printf("%d\n", ans[Q]);
     return 0;
种类并查集+计数 (POJ 1988)
int pa[maxn], cnt[maxn], pl[maxn];
int n, p;
int findset(int x){
     if(pa[x] == x) return x;
     int tmp = pa[x];
     pa[x] = findset(pa[x]);
     pl[x] += pl[tmp];
     return pa[x];
```

```
void init(){
     for(int i = 0; i < maxn; i + + ){
          pa[i] = i; cnt[i] = 1; pl[i] = 0;
void merge(int u, int v){
     pa[v] = u;
     pl[v] += cnt[u];
     cnt[u] += cnt[v];
int main()
     while(scanf("%d", &p) != EOF){
          init();
          while(p --){
               char op[5];
               scanf("%s", op);
               int x, y;
               if(op[0] == 'M')\{
                    scanf("%d%d", &x, &y);
                    int u = findset(x), v = findset(y);
                    if(u != v)
                         merge(u, v);
               else {
```

```
scanf("%d", &x);
                   int u = findset(x);
                   printf("%d\n", cnt[u] - pl[x] - 1);
    return 0;
Kruskal
位运算枚举树中的点重新构图
int pa[maxn];
int findset(int x) {
  return pa[x] != x ? (pa[x] = findset(pa[x])) : x;
struct Edge {
  int from, to, dist;
  Edge(){}
  Edge(int _u, int _v, int _c)
  :from(_u), to(_v), dist(_c){}
  bool operator < (const Edge& rhs) const {
    return dist < rhs.dist;
Edge e[maxm];
int MST(int n, int m) {
```

```
sort(e, e + m);
  for(int i = 0; i < n; i ++) pa[i] = i;
  int cnt = 0, mst = 0;
  for(int i = 0; i < m; i ++) {
    int u = e[i].from, v = e[i].to;
    int x = findset(u), y = findset(v);
    int dist = e[i].dist;
     if(x != y) {
       pa[y] = x;
       mst += dist;
       if(++cnt == n - 1) return mst;
  return INF;
int n, m, K;
Edge ed[maxm];
int pl[maxn];
int id[maxn];
int main() {
  int Tot;
  scanf("%d", &Tot);
  for(int ac = 1; ac \leq Tot; ac ++) {
     scanf("%d%d%d", &N, &M, &K);
     for(int i = 0; i < N; i ++) scanf("%d", &pl[i]);
     for(int i = 0; i < M; i ++) {
```

```
int u, v, c;
  scanf("%d%d%d", &u, &v, &c);
  ed[i] = Edge(u-1, v-1, c);
int mst = pl[0];
int bestn = 1, ans = pl[0];
for(int S = 0; S < (1 << (N-1)); S ++) {
  n = 1; m = 0;
  memset(id, -1, sizeof(id));
  id[0] = 0;
  for(int i = 0; i < N-1; i ++)
    if((S>>i) \& 1) id[i+1] = n ++;
  for(int i = 0; i < M; i++) {
    int u = id[ed[i].from];
    int v = id[ed[i].to];
    if(u >= 0 \&\& v >= 0)
       e[m ++] = Edge(u, v, ed[i].dist);
  if(MST(n, m) \le K) {
    int tot = pl[0];
    for(int i = 0; i < N-1; i ++)
       if((S>>i) \& 1) tot += pl[i+1];
    if(tot > ans) {
       ans = tot; bestn = n;
```

```
printf("%d\n", ans);
  return 0;
K 度最小生成树
const int maxn = 30;
const int inf = 0x3f3f3f3f;
struct HeapNode{
    int u, d;
    HeapNode(){}
    HeapNode(int \_u,int \_d):u(\_u), d(\_d){}
    bool operator < (const HeapNode& rhs) const{
        return rhs.d < d:
};
map<string, int> mp;
int g[maxn][maxn]; //邻接矩阵存图
int lowc[maxn];
int col[maxn]; //记录每个点所属生成树
int pre[maxn]; //记录生成树中的父亲节点
int fst[maxn]; //记录每个生成树到根的最小距离的点
int max_side[maxn];
int n, m, k;
int Prim(int src, int id) {
  priority_queue<HeapNode> Q;
```

```
while(!Q.empty()) Q.pop();
  lowc[src] = 0;
  Q.push(HeapNode(src, 0));
  int ans = 0;
  while(!Q.empty()) {
     HeapNode x = Q.top(); Q.pop();
     int u = x.u;
     if(!col[u]) {
       col[u] = id;
       ans += lowc[u];
       for(int v = 1; v < n; v ++) {
          if(!col[v] \&\& g[u][v] != 0 \&\& lowc[v] > g[u][v]) {
            pre[v] = u;
            lowc[v] = g[u][v];
            Q.push(HeapNode(v, lowc[v]));
  return ans;
void update(int u, int fa, int maxside) {
  max\_side[u] = max(maxside, g[u][fa]);
  for(int v = 1; v < n; v ++) {
     if(v != fa && g[u][v] != 0 && (pre[u] == v \parallel pre[v] == u)) {
       update(v, u, max_side[u]);
```

```
void solve() {
  int v, ans, cnt;
  for(v = 0; v < n; v ++) {
    lowc[v] = inf;
     col[v] = pre[v] = fst[v] = 0;
  ans = 0, cnt = 1;
  for(v = 1; v < n; v ++) {
     if(!col[v]) {
       ans += Prim(v, cnt ++);
  for(v = 1; v < n; v ++) {
     int id = col[v];
     if(g[0][v] != 0 && (!fst[id] || g[0][v] < g[0][fst[id]])) 
       fst[id] = v;
  for(int i = 1; i < cnt; i ++) {
     ans += g[0][fst[i]];
     g[0][fst[i]] = g[fst[i]][0] = 0;
     update(fst[i], 0, 0);
```

```
k = k - cnt + 1;
  while(k --) {
     int tmp = 0;
     for(v = 1; v < n; v ++) {
       if(g[0][v] != 0 \&\& (tmp == 0 || max\_side[tmp] - g[0][tmp]
                                 < max\_side[v] - g[0][v]))
          tmp = v;
     if(max\_side[tmp] \le g[0][tmp]) break;
     ans = ans - max\_side[tmp] + g[0][tmp];
     g[0][tmp] = g[tmp][0] = 0;
     int p = 0;
     for(v = tmp; pre[v] != 0; v = pre[v]) {
       if(p == 0 || g[p][pre[p]] < g[v][pre[v]])
          p = v;
     pre[p] = 0;
     update(tmp, 0, 0);
  printf("Total miles driven: %d\n", ans);
int main(){
  //freopen("input.txt","r",stdin);
  char a[20], b[20];
  int c;
  while(scanf("%d", &m) != EOF) {
```

```
mp["Park"] = 0;
   n = 1;
   memset(g, 0, sizeof(g));
    while(m --) {
     scanf("%s %s %d",a,b, &c);
     if(!mp.count(a)) mp[a] = n++;
     if(!mp.count(b)) mp[b] = n++;
     int u = mp[a], v = mp[b];
     if(!g[u][v] || g[u][v] > c)
       g[u][v] = g[v][u] = c;
    scanf("%d", &k);
   solve();
  return 0;
斯坦纳树 (HDU 4085)
        斯坦纳树
给你 n,m,k , 分别表示有 n 个点, m 条边, 每条边有一个权值,
表示修复这条边需要的代价,从前 k 个点中任取一个使其和
后 k 个点中的某一个点, 通过边连接, 并且必须是一一对应,
问最小的代价是多少。
#include <stdio.h>
#include <string.h>
#include <stdlib.h>
```

```
#include <queue>
using namespace std;
const int inf = 0x3f3f3f3f3f;
const int maxn = 60;
struct Edge{
     int from, to, dist;
     Edge(){}
     Edge(int _u, int _v, int _d)
     :from(_u), to(_v), dist(_d){}
};
Edge ed[maxn*maxn];
int head[maxn], nxt[maxn*maxn], m;
int N, M, K, mask;
int st[maxn], vis[maxn][1<<11];
int d[maxn][1<<11], dp[1<<11];
void AddEdge(int from, int to, int dist) {
     nxt[m] = head[from]; ed[m] = Edge(from, to, dist);
     head[from] = m ++;
bool check(int S) {
     int res = 0;
     for(int i = 0; S; i++, S >>= 1){
         res += (S\&1)*(i< K?1:-1);
     return res == 0;
```

```
void init(int n){
     m = 0; mask = (1 << (2*K))-1;
     memset(st, 0, sizeof(st[0])*(n+1));
     memset(vis, 0, sizeof(vis));
     memset(head, -1, sizeof(head[0])*(n+1));
     for(int i = 1; i \le n; i ++) {
          for(int j = 0; j \le mask; j ++ d[i][j] = inf;
queue<int> Q;
void SPFA() {
     while(!Q.empty()) {
          int u = Q.front()/10000, S = Q.front()\%10000;
          vis[u][S] = 0;
          Q.pop();
          for(int i = head[u]; i != -1; i = nxt[i]) {
               int v = ed[i].to, w = ed[i].dist;
               if(d[v][S|st[v]] > d[u][S] + w)  {
                    d[v][S|st[v]] = d[u][S] + w;
                    if((S|st[v]) == S && !vis[v][S]) 
                         Q.push(v*10000 + S);
                         vis[v][S] = 1;
```

```
int main()
     int Tot;
     scanf("%d", &Tot);
     while(Tot --) {
          scanf("%d%d%d", &N, &M, &K);
          init(N);
          int u, v, w;
          while(M --) {
               scanf("%d%d%d", &u, &v, &w);
               AddEdge(u, v, w);
               AddEdge(v, u, w);
          for(int i = 1; i \le K; i + +) {
               st[i] = 1 << (i-1);
               d[i][st[i]] = 0;
               st[N-i+1] = (1 << (K + (i-1)));
               d[N-i+1][st[N-i+1]] = 0;
          for(int S = 0; S \le mask; S ++){
               for(int i = 1; i \le N; i ++) {
                    for(int p = (S-1)\&S; p; p = (p-1)\&S)
                         d[i][S] = min(d[i][S], d[i][p|st[i]] + d[i][(S-p)|st[i]]);
                    if(d[i][S]<inf && !vis[i][S])
                         Q.push(i*10000 + S), vis[i][S] = 1;
```

```
SPFA();
         for(int S = 0; S \le mask; S ++) {
              dp[S] = inf;
              for(int i = 1; i \le N; i ++)
                   dp[S] = min(dp[S], d[i][S]);
         for(int S = 1; S \le mask; S ++) {
              if(check(S)) {
                   for(int p = (S-1)\&S; p; p = (p-1)\&S) {
                        if(check(p))
                             dp[S] = min(dp[S], dp[p] + dp[S - p]);
         if(dp[mask] >= inf) puts("No solution");
         else printf("%d\n", dp[mask]);
     return 0;
有向树形图(UVA 11865)
typedef int type;
const int maxn = 100 + 10;
const int maxm = 10000 + 10;
const int inf = INT_MAX;
```

```
struct Edge{
  int from, to, b;
  type dist;
  Edge(){}
  Edge(int _from, int _to, int _b, type _dist)
  :from(_from), to(_to), b(_b), dist(_dist){}
  bool operator < (const Edge& rhs) const {
     return b > rhs.b;
}E[maxm];
struct DMST {
  int pre[maxn], ID[maxn], vis[maxn];
  type In[maxn];
  Edge edges[maxm];
  type Directed_MST(int root, int n, int m) {
     type ans = 0;
     while(true) {
       //找最小入边
       for(int i = 0; i < n; i ++) In[i] = inf;
       for(int i = 0; i < m; i ++) {
          int u = edges[i].from, v = edges[i].to;
          if(edges[i].dist < In[v] && u != v) 
            In[v] = edges[i].dist; pre[v] = u;
       for(int i = 0; i < n; i ++) {
```

```
if(i == root) continue;
  if(In[i] == inf) \{return -1;\}
//找环
int cntnode = 0;
memset(ID, -1, sizeof(ID));
memset(vis, -1, sizeof(vis));
In[root] = 0;
for(int i = 0; i < n; i ++) {
  ans += In[i];
  int v = i;
  while(vis[v] != i \&\& ID[v] == -1 \&\& v != root) {
     vis[v] = i; v = pre[v];
  if(v != root && ID[v] == -1) {
     for(int u = pre[v]; u != v; u = pre[u]) {
       ID[u] = cntnode;
     ID[v] = cntnode ++;
if(cntnode == 0) break;
for(int i = 0; i < n; i + +) {
  if(ID[i] == -1) {
     ID[i] = cntnode ++;
```

```
//缩点,重新标记
       for(int i = 0; i < m; i ++) {
          int v = edges[i].to;
          edges[i].from = ID[edges[i].from];
          edges[i].to = ID[edges[i].to];
          if(edges[i].from != edges[i].to) {
            edges[i].dist -= In[v];
       n = cntnode;
       root = ID[root];
     return ans;
};
DMST dem;
int N, M, C;
int Tot;
bool check(int cnt) {
  //printf("cnt = %d\n", cnt);
  for(int i = 0; i < cnt; i ++) {
     dem.edges[i] = E[i];
  int ans = dem.Directed MST(0, N, cnt);
  //printf("b = %d, ans = %d\n", E[cnt].b, ans);
```

```
if(ans == -1 \parallel ans > C) return false;
  return true;
int main()
  //freopen("input.txt", "r", stdin);
  scanf("%d", &Tot);
  for(int ac = 1; ac \leq Tot; ac ++) {
     scanf("%d%d%d", &N, &M, &C);
    int u, v, b, c;
     for(int i = 0; i < M; i ++) {
       scanf("%d%d%d%d", &u, &v, &b, &c);
       E[i] = Edge(u, v, b, c);
     sort(E, E + M);
     int l = 1, r = M, mid;
     int ans = -1;
     while(l \le r) {
       mid = (1 + r) >> 1;
       if(check(mid)) {
          r = mid - 1; ans = E[mid - 1].b;
       else l = mid + 1;
     if(ans \geq 0) printf("%d kbps\n", ans);
     else printf("streaming not possible.\n");
```

```
return 0;
最小直径生成树(URAL 1569)
#include <stdio.h>
#include <string.h>
#include <stdlib.h>
#include <algorithm>
#include <vector>
#include <queue>
using namespace std;
typedef pair<int, int> pii;
const int inf = 0x3f3f3f3f;
const int maxn = 200 + 10;
int w[maxn][maxn];
int d[maxn][maxn];
int pre[maxn];
bool inq[maxn];
int dp[maxn], ans;
int cu, cv;
void init(int n) {
    for(int i = 1; i \le n; i ++) {
         for(int j = 1; j \le n; j ++) {
              w[i][j] = inf;
```

```
w[i][i] = 0;
void Floyd(int n) {
     for(int k = 1; k \le n; k + +)
     for(int i = 1; i \le n; i ++)
     for(int j = 1; j \le n; j ++) {
          if(d[i][j] > d[i][k] + d[k][j])
               d[i][j] = d[i][k] + d[k][j];
void update(int u, int v, int n) {
     vector<pii> vt, tk;
     for(int i = 1; i \le n; i ++) {
          vt.push_back(make_pair(d[u][i], d[v][i]));
     sort(vt.begin(), vt.end());
     /*按照 d[u][i] 从小到大排, 考虑 d[u][i] < d[u][i]
     && d[v][i] > d[u][j] (i < j)*/
     int sz = vt.size();
     for(int i = 0; i < sz; i ++) {
          while(!tk.empty() && tk.back().second <= vt[i].second) {</pre>
               tk.pop_back();
          tk.push_back(vt[i]);
```

```
int D = inf;
int tmp;
if(tk.size() == 1) {
     if(tk[0].first < tk[0].second) {
          tmp = 0;
          D = tk[0].first << 2;
     }else {
          tmp = w[u][v] << 1;
          D = tk[0].second << 2;
} else {
     sz = tk.size();
     for(int i = 1; i < sz; i ++) {
          if(D > ((w[u][v] + tk[i - 1].first + tk[i].second) << 1)){
               tmp = w[u][v] + tk[i].second - tk[i - 1].first;
               D = ((w[u][v] + tk[i - 1].first + tk[i].second) << 1);
if(D < ans) {
     cu = u;
     cv = v;
     ans = D;
     dp[u] = tmp;
     dp[v] = (w[u][v] << 1) - tmp;
```

```
vector<pii> edges;
void Spfa(int n) {
     queue<int>q;
     memset(pre, -1, sizeof(pre));
     memset(inq, false, sizeof(inq));
     for(int i = 1; i \le n; i ++) {
          if(i != cu && i != cv) {
               dp[i] = inf;
     inq[cu] = inq[cv] = true;
     q.push(cu); q.push(cv);
     while(!q.empty()) {
          int u = q.front(); q.pop();
          inq[u] = false;
          for(int v = 1; v \le n; v ++) {
               if(w[u][v] != inf) {
                    if(dp[v] > dp[u] + (w[u][v] << 1)) {
                         dp[v] = dp[u] + (w[u][v] << 1);
                         pre[v] = u;
                         if(!inq[v]) {
                              q.push(v);
                              inq[v] = true;
```

```
edges.clear();
          for(int u = 1; u \le n; u ++) {
               if(pre[u] != -1) {
                   edges.push_back(make_pair(min(u,
                                                           pre[u]),
                                                                        max(u,
pre[u])));
         if(cv != cu) {
               edges.push_back(make_pair(min(cv, cu), max(cv, cu)));
int N, M;
int main()
     //freopen("input.txt", "r", stdin);
     while(scanf("%d%d", &N, &M) != EOF) {
         init(N);
         int a, b;
         while(M --) {
               scanf("%d%d", &a, &b);
               w[a][b] = w[b][a] = min(1, w[a][b]);
         memcpy(d, w, sizeof(d));
```

```
Floyd(N);
         ans = inf;
         for(int i = 1; i \le N; i ++) {
                                                                                    const int inf = 0x3f3f3f3f;
              for(int j = 1; j \le N; j ++) {
                                                                                    const int maxn = 100 + 10;
                   if(w[i][j] != inf) {
                                                                                    const int logmaxn = 32;
                                                                                    const int maxm = 200 + 10;
                         update(i, j, N);
                                                                                    int n, m;
                                                                                    int pa[maxn];
                                                                                    int maxcost[maxn][logmaxn], fa[maxn];
         //printf("%d\n", ans / 2);
                                                                                    int anc[maxn][logmaxn];
         Spfa(N);
                                                                                    int L[maxn], cost[maxn];
                                                                                    bool vis[maxm];
         sort(edges.begin(), edges.end());
         for(vector<pii>::iterator it = edges.begin(); it != edges.end(); it ++) {
                                                                                    Edge e[maxm];
              printf("%d %d\n", it->first, it->second);
                                                                                    vector<int> G[maxn], C[maxn];
                                                                                    int findset(int x){
                                                                                       return pa[x] == x ? x : (pa[x] = findset(pa[x]));
     return 0;
                                                                                    void dfs(int u, int pa, int level) {
次小生成树(UVA 10462)
                                                                                       L[u] = level;
                                                                                       int sz = G[u].size();
struct Edge{
  int from, to, dist;
                                                                                       for(int i = 0; i < sz; i ++)
  Edge(){}
  Edge(int _from, int _to, int _dist)
                                                                                         int v = G[u][i];
  :from(_from), to(_to), dist(_dist){}
                                                                                         if(v != pa) {
  bool operator < (const Edge& rhs) const{
                                                                                            fa[v] = u, cost[v] = C[u][i];
     return dist < rhs.dist; //在用到优先队列时要换成 >
                                                                                            dfs(v, u, level + 1);
```

```
void preprocess(){
  for(int i = 0; i < n; i + +)
     anc[i][0] = fa[i]; maxcost[i][0] = cost[i];
  for(int j = 1; (1 << j) < n; j ++)
     anc[i][j] = -1;
  for(int j = 1; (1 << j) < n; j ++)
     for(int i = 0; i < n; i ++)
       if(anc[i][j-1]!=-1) {
          int a = anc[i][j-1];
          anc[i][j] = anc[a][j-1];
          \max(st[i][j] = \max(\max(st[i][j-1], \max(st[i][j-1]);
int query(int p, int q){
  int log;
  if(L[p] < L[q]) swap(p, q);
  for(log = 1; (1 << log) <= L[p]; log ++); log --;
  int ans = -inf;
  for(int i = log; i >= 0; i --) {
     if(L[p] - (1 << i) >= L[q]) {
       ans = max(ans, maxcost[p][i]); p = anc[p][i];
```

```
if(p == q) return ans;
  for(int i = log; i >= 0; i --) {
     if(anc[p][i] != -1 && anc[p][i] != anc[q][i]) {
       ans = max(ans, maxcost[p][i]); p = anc[p][i];
       ans = max(ans, maxcost[q][i]); q = anc[q][i];
  ans = max(ans, cost[p]);
  ans = max(ans, cost[q]);
  return ans;
int Tot;
int main()
  //freopen("input.txt", "r", stdin);
  scanf("%d", &Tot);
  for(int ac = 1; ac \leq Tot; ac ++) {
     scanf("%d%d", &n, &m);
     for(int i = 0; i < n; i ++){
        pa[i] = i;
       G[i].clear(); C[i].clear();
     for(int i = 0; i < m; i ++){
        int a, b, c;
        scanf("%d%d%d", &a, &b, &c);
```

```
a --; b --;
  e[i] = Edge(a, b, c);
sort(e, e + m);
int ans = 0, cnt = 1;
memset(vis, false, sizeof vis);
for(int i = 0; i < m; i + +)
  int x = e[i].from, y = e[i].to, d = e[i].dist;
  int u = findset(x), v = findset(y);
  if(u != v)
     vis[i] = true; pa[v] = u; ans += d;
     G[x].push_back(y); C[x].push_back(d);
     G[y].push_back(x); C[y].push_back(d);
     cnt ++;
printf("Case #%d: ", ac);
if(cnt < n) {
  printf("No way\n");
  continue;
dfs(0, -1, 0);
preprocess();
int smst = inf;
for(int i = 0; i < m; i ++) {
  if(!vis[i]) {
```

```
int u = e[i].from, v = e[i].to;
          smst = min(ans + e[i].dist - query(u, v), smst);
     if(smst == inf) printf("No second way\n");
     else printf("%d\n", smst);
  return 0;
最小生成树最大边最大
typedef pair<int, int> pii;
const double inf = 0x3f3f3f3f;
const int maxn = 2000 + 10;
const int maxm = 1000000 + 100;
struct Edge{
  int from, to, dist;
  Edge(){}
  Edge(int _from, int _to, int _dist)
  :from(_from), to(_to), dist(_dist){}
  bool operator < (const Edge& rhs) const {
     return dist < rhs.dist;
};
int N, M;
int mst, trz;
bool intree[maxm];
```

```
Edge e[maxm];
int dp[maxn][maxn];
int ca[maxn], cb[maxn];
int cacnt, cbcnt;
int res;
int main()
  //freopen("input.txt", "r", stdin);
  int k, i, j;
  int u, v, w;
  while(scanf("%d%d", &N, &M) != EOF) {
    for(i = 0; i < M; i ++) {
       scanf("%d%d%d", &e[i].from, &e[i].to, &e[i].dist);
       intree[i] = 0;
     sort(e, e + M);
    for(i = 1; i \le N; i ++) {
       for(j = 1; j \le N; j ++) {
         dp[i][j] = -1;
       dp[i][i] = 0;
     mst = trz = 0;
    for(k = 0; k < M; k ++) {
       u = e[k].from; v = e[k].to; w = e[k].dist;
       if(dp[u][v]!=-1) continue; //已经在一个块中
```

```
intree[k] = 1;
  cacnt = cbcnt = 0;
  for(i = 1; i \le N; i ++) {
    if(dp[i][u]!=-1) { //找成环的边
       ca[cacnt ++] = i;
    if(dp[i][v] != -1) {
       cb[cbcnt ++] = i;
  int ut, vt;
  for(i = 0; i < cacnt; i ++) {
    for(j = 0; j < cbcnt; j ++) {
       ut = ca[i]; vt = cb[j];
       if(ut != vt) { //构成环的最大权值
          dp[ut][vt] = dp[vt][ut] = max(w, max(dp[ut][u], dp[vt][v]));
  mst += w;
  trz ++;
  if(trz == N - 1) break;
if(trz < N - 1) {
  puts("disconnected"); continue;
```

```
res = mst;
   for(i = 0; i < M; i ++) {
     if(!intree[i]) {
       w = mst - e[i].dist - dp[e[i].from][e[i].to];
       res = min(res, w);
    printf("%d\n", res);
 return 0;
2-sat
建图
用 2*u 表示选择 u, 2*u+1 表示不选择 u
对于有限制(如 u, v 不能同时出现)
如果选 u, 不选 v, 那么 2*u-->2*v+1 连一条边, 选 v 不选 u 同理
更广泛一点,对于 a 来表示 u 的一种情况, b 表示 v 的一种情况
则 a-->b^1 连一条边,由对称性 b-->a^1 连一条边
Dfs 版
struct TwoSAT {
 int n, m;
 int head[maxn * 2];
 int to[maxm * 2];
  int next[maxm * 2];
  bool mark[maxn * 2];
```

```
int S[maxn * 2], c;
void init(int n) {
  this->n = n:
  m = 0;
  memset(head, -1, sizeof(head[0]) * (2*n + 1));
  memset(mark, false, sizeof(mark[0]) * (2*n + 1));
bool dfs(int x) {
  if(mark[x ^ 1]) return false;
  if(mark[x]) return true;
  mark[x] = true;
  S[c ++] = x;
  for(int i = head[x]; i != -1; i = next[i]) {
     if(!dfs(to[i])) return false;
  return true;
/*连边直接上条件*/
void AddEdge(int x, int y) {
  next[m] = head[x]; to[m] = y; head[x] = m ++;
bool tsat() {
  for(int i = 0; i < n * 2; i += 2) {
     if(!mark[i] && !mark[i + 1]) {
       c = 0;
       if(!dfs(i)) {
```

```
while(c > 0) mark[S[-- c]] = false;
            if(!dfs(i + 1)) return false;
     return true;
};
Tarjan 版
struct Edge{
  int to, next;
  Edge(){}
  Edge(int _to, int _next) : to(_to), next(_next){}
};
struct Tarjan {
  int m;
  int head[maxn];
  Edge edges[maxm];
  int pre[maxn], sccno[maxn], lowlink[maxn], scc_cnt, dfs_clock;
  int S[maxn], top;
  int cnt[maxn]; //记录每个强连通分量中点的个数
  void init() {
    m = 0;
    memset(head, -1, sizeof(head));
    memset(cnt, 0, sizeof(cnt));
```

```
void AddEdge(int from, int to) {
  int next = head[from];
  head[from] = m;
  edges[m ++] = Edge(to, next);
void dfs(int u) {
  pre[u] = lowlink[u] = ++ dfs_clock;
  S[top ++] = u;
  for(int i = head[u]; i != -1; i = edges[i].next) {
    int v = edges[i].to;
    if(!pre[v]) {
       dfs(v);
       lowlink[u] = min(lowlink[u], lowlink[v]);
     }else if(!sccno[v]) {
       lowlink[u] = min(lowlink[u], pre[v]);
  if(pre[u] == lowlink[u]) {
     scc_cnt ++;
     while(true) {
       cnt[scc_cnt] ++;
       int x = S[--top];
       sccno[x] = scc\_cnt;
       if(x == u) break;
```

```
void find_scc(int n) {
     dfs\_clock = top = scc\_cnt = 0;
     memset(pre, 0, sizeof(pre));
     memset(sccno, 0, sizeof(sccno));
     for(int i = 0; i < n; i ++) {
       if(!pre[i]) dfs(i);
};
Tarjan sat;
int N, M;
void Prepare() {
  sat.init();
  int u, v;
  while(M --) {
     scanf("%d%d", &u, &v);
     if(u > 0 \&\& v > 0) {
       sat.AddEdge(2*(u - 1) + 1, 2*(v - 1));
       sat.AddEdge(2*(v-1) + 1, 2*(u-1));
     if(u < 0 && v < 0) {
       sat.AddEdge(2*(-u - 1), 2*(-v - 1) + 1);
       sat.AddEdge(2*(-v - 1), 2*(-u - 1) + 1);
```

```
if(u > 0 \&\& v < 0) {
       sat.AddEdge(2*(u - 1) + 1, 2*(-v - 1) + 1);
       sat.AddEdge(2*(-v - 1), 2*(u - 1));
     if(u < 0 && v > 0) {
       sat.AddEdge(2*(-u - 1), 2*(v - 1));
       sat.AddEdge(2*(v-1) + 1, 2*(-u-1) + 1);
int main()
  while(scanf("%d%d", &N, &M) != EOF) {
     Prepare();
     sat.find scc(N * 2);
     bool ok = true;
     for(int i = 0; i < N * 2; i += 2) {
       if(sat.sccno[i] == sat.sccno[i + 1])
         ok = false;
     if(ok) puts("1");
     else puts("0");
  return 0;
2-sat 输出解
```

```
TwoSAT sat;
int N;
bool ok;
int S[maxn], T[maxn], D[maxn];
bool check(int s1, int t1, int s2, int t2) {
  if((s1 < s2 && s2 < t1) ||
       (s1 < t2 \&\& t2 < t1) \parallel
       (s2 <= s1 && t1 <= t2)
  ) return true;
  return false;
void Prepare() {
  int u, v;
  sat.init(N);
  ok = true;
  int hh, mm;
  for(int i = 0; i < N; i ++) {
     scanf("%d:%d", &hh, &mm);
     S[i] = hh * 60 + mm;
     scanf("%d:%d %d", &hh, &mm, &D[i]);
     T[i] = hh * 60 + mm;
     if(T[i] < S[i] + D[i]) ok = false;
     for(int j = 0; j < i; j ++) {
       u = 2 * i; v = 2 * j;
       int s1 = S[i], t1 = S[i] + D[i];
```

```
int s2 = T[i] - D[i], t2 = T[i];
       int s3 = S[j], t3 = S[j] + D[j];
       int s4 = T[j] - D[j], t4 = T[j];
       if(check(s1, t1, s3, t3)) {
           sat.AddEdge(u, v ^ 1);
       if(check(s1, t1, s4, t4)) {
           sat.AddEdge(u, v);
       if(check(s2, t2, s3, t3)){
           sat.AddEdge(u ^ 1, v ^ 1);
       if(check(s2, t2, s4, t4)){
           sat.AddEdge(u ^ 1, v);
int main()
  while(scanf("%d", &N) != EOF) {
     Prepare();
     if(!ok) {
        puts("NO");
        continue;
```

Yucept 的图论模板

```
ok = sat.tsat();
                                                                                     if(e[i].u < e[i].u && e[i].v && e[i].v < e[i].v) return true;
     if(!ok) puts("NO");
                                                                                     return false;
     else{
       puts("YES");
                                                                                   void Prepare() {
       for(int i = 0; i < 2 * N; i += 2) {
                                                                                     int u, v;
         if(sat.mark[i]) {
                                                                                     sat.init(M);
            printf("%02d:%02d %02d:%02d\n", S[i/2]/60, S[i/2]%60,
                                                                                     for(int i = 0; i < M; i ++) {
                    (S[i/2] + D[i/2])/60, (S[i/2] + D[i/2])\%60);
                                                                                        scanf("%d%d", &u, &v);
                                                                                        if(u > v) swap(u, v);
         else printf("%02d:%02d %02d:%02d\n", (T[i/2] - D[i/2])/60,
                                                                                        e[i].u = u, e[i].v = v;
(T[i/2]-D[i/2])\%60,
                                                                                        for(int j = 0; j < i; j ++) {
                    T[i/2]/60, T[i/2]\%60);
                                                                                          if(check(i, j)) {
                                                                                             int a = 2 * i, b = 2 * j;
                                                                                             sat.AddEdge(a, b ^ 1);
                                                                                             sat.AddEdge(a ^ 1, b);
                                                                                             sat.AddEdge(b, a ^ 1);
  return 0;
                                                                                             sat.AddEdge(b ^ 1, a);
与边相关的 2-sat
TwoSAT sat;
int N, M;
struct Edge{
  int u, v;
                                                                                   int main()
};
Edge e[maxn];
                                                                                     while(scanf("%d%d", &N, &M) != EOF) {
bool check(int i, int j) { //判断是否可以共存
                                                                                        Prepare();
  if(e[i].u < e[i].u && e[i].u < e[i].v && e[i].v > e[i].v) return true;
                                                                                        bool ok = sat.tsat();
```

```
if(!ok) puts("the evil panda is lying again");
     else puts("panda is telling the truth...");
  return 0;
位枚举+2-sat
TwoSAT sat;
int N;
int mat[maxn][maxn];
bool ok;
void Prepare() {
  for(int i = 0; i < N; i ++) {
     for(int j = 0; j < N; j ++) {
       scanf("%d", &mat[i][j]);
  ok = true;
  for(int i = 0; i < N; i ++) {
     if(mat[i][i] != 0) \{ok = false; break;\}
     for(int j = 0; j < N; j ++) {
       if(mat[i][j] != mat[j][i]) {
          ok = false; break;
```

```
void conect(int x, int y, int c) {
  if(x % 2 == 0 && y % 2 == 0) { //与
    if(c == 1) \{
       sat.AddEdge(2*x, 2*y);
       sat.AddEdge(2*y, 2*x);
    else {
       sat.AddEdge(2*x, 2*y + 1);
       sat.AddEdge(2*y, 2*x + 1);
  else if(x % 2 == 1 && y % 2 == 1) { //或
    if(c == 1) \{
       sat.AddEdge(2*x + 1, 2*y);
       sat.AddEdge(2*y + 1, 2*x);
    else {
       sat.AddEdge(2*x + 1, 2*y + 1);
       sat.AddEdge(2*y + 1, 2*x + 1);
  else { //异或
    if(c == 1) \{
       sat.AddEdge(2*x + 1, 2*y);
       sat.AddEdge(2*y + 1, 2*x);
       sat.AddEdge(2*x, 2*y + 1);
```

```
sat.AddEdge(2*y, 2*x + 1);
     else {
       sat.AddEdge(2*x + 1, 2*y + 1);
       sat.AddEdge(2*x, 2*y);
       sat.AddEdge(2*y, 2*x);
       sat.AddEdge(2*y + 1, 2*x + 1);
int main()
  while(scanf("%d", &N) != EOF) {
    Prepare();
     if(!ok) {puts("NO"); continue;}
    for(int k = 0; k < 32; k ++) {
       sat.init(N);
       for(int i = 0; i < N; i ++) {
         for(int j = 0; j < N; j ++) {
            if(i == j) continue;
            int c = ((mat[i][j] >> k) \& 1);
            conect(i, j, c);
       ok = sat.tsat(); if(!ok) break;
```

```
if(ok) puts("YES");
    else puts("NO");
  return 0;
连通分量
强连通
Tarjan
struct Edge{
  int to, next;
  Edge(){}
  Edge(int _to, int _next) : to(_to), next(_next){}
};
struct Tarjan {
  int m;
  int head[maxn];
  Edge edges[maxm];
  int pre[maxn], sccno[maxn], lowlink[maxn], scc_cnt, dfs_clock;
  int S[maxn], top;
  int cnt[maxn]; //记录每个强连通分量中点的个数
  void init() {
    m = 0;
```

memset(head, -1, sizeof(head));
memset(cnt, 0, sizeof(cnt));

```
void AddEdge(int from, int to) {
  int next = head[from];
  head[from] = m;
  edges[m ++] = Edge(to, next);
void dfs(int u) {
  pre[u] = lowlink[u] = ++ dfs_clock;
  S[top ++] = u;
  for(int i = head[u]; i != -1; i = edges[i].next) {
    int v = edges[i].to;
    if(!pre[v]) {
       dfs(v);
       lowlink[u] = min(lowlink[u], lowlink[v]);
     }else if(!sccno[v]) {
       lowlink[u] = min(lowlink[u], pre[v]);
  if(pre[u] == lowlink[u]) {
    scc_cnt ++;
    while(true) {
       cnt[scc_cnt] ++;
       int x = S[--top];
       sccno[x] = scc\_cnt;
       if(x == u) break;
```

```
void find_scc(int n) {
     dfs\_clock = top = scc\_cnt = 0;
     memset(pre, 0, sizeof(pre));
     memset(sccno, 0, sizeof(sccno));
     for(int i = 0; i < n; i ++) {
       if(!pre[i]) dfs(i);
最少加几条边让原图强连通
Tarjan tar;
int N, M;
bool out0[maxn], in0[maxn]; //
int Map[505][505];
void Pre() {
  tar.init();
  for(int i = 1; i \le N; i ++) {
     for(int j = 1; j \le M; j ++) {
       scanf("%d", &Map[i][j]);
  for(int i = 1; i \le N; i ++) {
     for(int j = 1; j \le M; j ++) {
       if(j < M) {
```

```
int u = (i - 1) * M + j - 1;
  int v = u + 1;
  if(Map[i][j+1] > Map[i][j]) {
     tar.AddEdge(v, u);
  else if(Map[i][j+1] == Map[i][j]) {
     tar.AddEdge(u, v); tar.AddEdge(v, u);
  else tar.AddEdge(u, v);
  //printf("u = %d, v = %d\n", u, v);
if(i < N) {
  int u = (i - 1) * M + j - 1;
  int v = u + M;
  if(Map[i][j] > Map[i+1][j]) {
     tar.AddEdge(u, v);
  else if(Map[i][j] == Map[i + 1][j]) {
     tar.AddEdge(u, v);
     tar.AddEdge(v, u);
  else tar.AddEdge(v, u);
  //printf("u = %d, v = %d \ n", u, v);
```

```
int main()
  while(scanf("%d%d", &M, &N) != EOF) {
    Pre();
    tar.find_scc(N * M);
    if(tar.scc_cnt == 1) {
       printf("0\n"); continue;
    for(int i = 1; i \le tar.scc\_cnt; i ++) {
      in0[i] = out0[i] = true; //true 表示入度或出度为 0
    for(int u = 0; u < N * M; u ++) {
       for(int i = tar.head[u]; i != -1; i = tar.edges[i].next) {
         int v = tar.edges[i].to;
         if(tar.sccno[v] != tar.sccno[u]) {
           in0[tar.sccno[v]] = out0[tar.sccno[u]] = false;
    int a, b;
    a = b = 0;
    for(int i = 1; i \le tar.scc\_cnt; i ++) {
       if(in0[i]) a ++; //入度为零的强连通分量个数
       if(out0[i]) b ++; //出度为零的强连通分量个数
```

```
int ans = max(a, b);
                                                                                   //freopen("input.txt", "r", stdin);
    printf("%d\n", ans);
                                                                                   scanf("%d", &Tot);
                                                                                   for(int ac = 1; ac \leq Tot; ac ++) {
  return 0;
                                                                                     Prepare();
                                                                                     tar.find_scc(N);
添加一些边使原图仍是简单图且非强连通图(要使添边数目最大)
                                                                                     printf("Case %d: ", ac);
将这个图变成两个强连通,而且是一个连通分量只能指向另一个连通分量,
                                                                                     if(tar.scc\_cnt == 1) \{printf("-1\n"); continue; \}
可以先缩点
                                                                                     for(int u = 0; u < N; u ++) {
Tarjan tar;
                                                                                       for(int i = tar.head[u]; i != -1; i = tar.edges[i].next) {
int N, M, Tot;
                                                                                          int v = tar.edges[i].to;
int in0[maxn], out0[maxn];
                                                                                          if(tar.sccno[u] != tar.sccno[v]) {
                                                                                            in0[tar.sccno[v]] ++; out0[tar.sccno[u]] ++;
int mint:
void Prepare() {
  scanf("%d%d", &N, &M);
  tar.init();
  int u, v;
                                                                                     mint = maxn;
  for(int i = 0; i < M; i ++) {
                                                                                     int cur;
    scanf("%d%d", &u, &v);
                                                                                     for(int i = 1; i \le tar.scc\_cnt; i ++) {
                                                                                       if(in0[i] == 0 || out0[i] == 0) {
    u --; v --;
    tar.AddEdge(u, v);
                                                                                          if(mint > tar.cnt[i]) {
                                                                                            mint = tar.cnt[i]; cur = i;
  memset(in0, 0, sizeof(in0[0]) * (N + 1));
  memset(out0, 0, sizeof(out0[0]) * (N + 1));
int main()
                                                                                     LL ans = (N - mint) * (N - mint - 1) + mint * (mint - 1);
```

```
ans += (mint * (N - mint));
    ans -= M;
    printf("%I64d\n", ans);
  return 0;
双连通
双连通找桥(有重边)
struct Edge{
  int from, to, id, mark;
  Edge(){}
  Edge(int _from, int _to, int _id, int _mark)
  :from(_from), to(_to), id(_id), mark(_mark){}
};
struct BCC {
  int n, m;
  int pre[maxn], dfs_clock;
  int head[maxn];
  int low[maxn];
  int next[maxm];
  Edge edges[maxm];
  vector<int> Bridges;
  void init(int n) {
    this->n = n;
    memset(head, -1, sizeof(head[0])*(n + 1));
```

```
memset(pre, 0, sizeof(pre[0])*(n + 1));
  memset(low, 0, sizeof(low[0])*(n + 1));
  dfs clock = 0;
  Bridges.clear();
  m = 0;
void AddEdge(int from, int to, int id){
  next[m] = head[from];
  edges[m] = Edge(from, to, id, 0);
  head[from] = m ++;
void dfs(int u, int fa) {
  low[u] = pre[u] = ++ dfs\_clock;
  for(int i = head[u]; i != -1; i = next[i]) {
     Edge& e = edges[i];
     int v = e.to;
     if(edges[i].mark) continue;
     edges[i].mark = edges[i ^1].mark = 1;
     if(!pre[v]) {
       dfs(v, u);
       low[u] = min(low[u], low[v]);
       if(low[v] > pre[u]) {
          Bridges.push_back(e.id);
          //printf("u = %d, v = %d\n", u, v);
     }else {
```

```
low[u] = min(low[u], pre[v]);
};
BCC bcc;
int N, M;
int weight[maxm];
void Pre() {
  bcc.init(N);
  int u, v;
  for(int i = 0; i < M; i ++) {
     scanf("%d%d%d", &u, &v, &weight[i]);
    //printf("u = %d, v = %d\n", u, v);
    bcc.AddEdge(u, v, i);
    bcc.AddEdge(v, u, i);
int main()
  //freopen("input.txt", "r", stdin);
  while(scanf("\%d\%d", &N, &M), N + M) {
    Pre();
     bcc.dfs(1, -1);
     bool flag = false;
    for(int i = 1; i \le N; i ++) {
```

```
if(!bcc.pre[i]) {flag = true; break;}
     if(flag) {
       printf("0\n"); continue;
    int sz = bcc.Bridges.size();
    if(sz == 0)  {
       printf("-1\n"); continue;
    //printf("sz = %d\n", sz);
     int ans = maxm;
     for(int i = 0; i < sz; i ++) {
       //printf("%d\n", bcc.edges[bcc.Bridges[i]].dist);
       ans = min(ans, weight[bcc.Bridges[i]]);
    if(ans == 0) ans ++;
    printf("%d\n", ans);
  return 0;
二分图
最少删除几条边使含奇圈的图变成二分图
int e[maxn][2];
int N, M, Tot;
int main()
```

```
scanf("%d", &Tot);
  while(Tot --) {
    scanf("%d%d", &N, &M);
   for(int i = 0; i < M; i ++) {
      scanf("%d%d", &e[i][0], &e[i][1]);
    int ans = M + 10, cnt:
   for(int i = 0; i < (1 << N); i ++) { //枚举所有染色的情况
     cnt = 0;
     for(int j = 0; j < M; j ++) {
        if(((i >> e[i][0]) \& 1) == ((i >> e[i][1]) \& 1)) {
          cnt ++;
      if(ans > cnt) ans = cnt;
    printf("%d\n", ans);
  return 0:
三正则图匹配
给一n个点的三正则图, 求最大匹配。
根据握手定理,n一定是偶数。
由于三正则图,而且题目提示是2边连通,所以图中不存在桥,
```

```
也就是一定可以找到一条回路经过每个顶点至少一次
(强连通的定义:强连通图一定存在一条回路记过每个顶点至
少一次)由于是三则图,每个顶点的度是3,如果这条回路经过
某个顶点 2 次, 那么这个顶点的度就是 4, 这个和条件矛盾。
这条经过每个顶点一次的交错路就可以作出 n/2 匹配。
*/
#include <stdio.h>
int K, N;
int main(){
 scanf("%d", &K);
 int u, v;
 for(int ac = 1; ac \leq K; ac ++) {
   scanf("%d", &N);
   for(int i = 0; i < N*3 / 2; i ++) scanf("%d%d", &u, &v);
   printf("%d\n", N / 2);
 return 0;
二分图最大匹配
最小点覆盖&最大独立集
最小点覆盖 = 最大匹配数
最大独立集建模:
将有联系的 X 点和 Y 点连边, 求最大匹配
然后点的总数 - 最大匹配数 = 独立点个数
最小路径覆盖=最大独立集
```

```
匈牙利
struct BPM {
  int n, m;
  int head[maxn];
  int next[12 * maxn];
  int to [12 * maxn];
  int left[maxn]; // left[i]为右边第 i 个点的匹配点编号, -1 表示不存在
  bool T[maxn]; // T[i]为右边第 i 个点是否已标记
  int right[maxn]; // 求最小覆盖用
  bool S[maxn]; // 求最小覆盖用
  void init(int n) {
    this->n = n;
    m = 0:
    memset(head, -1, sizeof(head[0])*(n + 1));
  void AddEdge(int u, int v) {
    next[m] = head[u]; to[m] = v; head[u] = m ++;
  bool dfs(int u) {
   S[u] = true;
   for(int i = head[u]; i != -1; i = next[i]) {
      int v = to[i];
      if(!T[v]) {
         T[v] = true;
         if(left[v] == -1 \parallel dfs(left[v])) {
           left[v] = u; right[u] = v;
```

```
return true;
    return false;
 int MaxMatch() {
    int ans = 0;
    memset(left, -1, sizeof(left[0])*(n + 1));
    memset(right, -1, sizeof(right[0])*(n + 1));
    for(int u = 0; u < n; u ++) {
      memset(T, false, sizeof(T[0])*(n + 1));
      memset(S, false, sizeof(S[0])*(n + 1));
      if(dfs(u)) ans ++;
         return ans;
};
HK 算法
struct HK {
  int left[maxn], right[maxn];
  int dx[maxn], dy[maxn];
  int n, m;
  int head[maxn];
  int to[maxm];
  int next[maxm];
```

```
void init(int n) {
  this->n = n;
  m = 0;
  memset(head, -1, sizeof(head[0]) * (n + 1));
void AddEdge(int u, int v) {
  next[m] = head[u];
  to[m] = v;
  head[u] = m ++;
bool bfs() {
  bool flag = false;
  memset(dx, 0, sizeof(dx));
  memset(dy, 0, sizeof(dy));
  queue<int> Q;
  for(int i = 1; i \le n; i ++) {
     if(right[i] == -1) Q.push(i);
  while(!Q.empty()) {
     int u = Q.front();
     Q.pop();
     for(int i = head[u]; i != -1; i = next[i]) {
       int v = to[i];
       if(dy[v] == 0) {
          dy[v] = dx[u] + 1;
          if(left[v] == -1) flag = true;
```

```
else {
             dx[left[v]] = dy[v] + 1;
             Q.push(left[v]);
  return flag;
int dfs(int u) {
  for(int i = head[u]; i != -1; i = next[i]) {
     int v = to[i];
     if(dy[v] == dx[u] + 1) \{
        dy[v] = 0;
        if(left[v] == -1 \parallel dfs(left[v])) {
           left[v] = u;
           right[u] = v;
           return true;
  return false;
int MaxMatch() {
  int res = 0;
  memset(left, -1, sizeof(left));
```

```
memset(right, -1, sizeof(right));
    while(bfs()) {
       for(int i = 1; i \le n; i ++) {
         if(right[i] == -1) res += dfs(i);
     return res;
二分图多重最大匹配 (可加限制条件)
int N, D[maxn], W;
int an[8], Dsum;
struct BPM {
  int n, m;
  bool T[maxd];
  int left[maxn][maxd];
  int vcnt[maxd];
  bool G[maxd][maxn];
  void init(int n, int m) {
    this->n = n; this->m = m;
    memset(G, false, size of G);
  bool match(int u/*,int limit*/) { //可以加一个限制条件 limit
    for(int v = 0; v < m; v ++) {
       if(!T[v] && G[u][v]) {
         T[v] = true;
```

```
int limit = D[v];
       if(vcnt[v] < limit) {</pre>
          left[v][vcnt[v] ++] = u;
          return true;
       for(int i = 0; i < vcnt[v]; i ++) {
          if(match(left[v][i])) {
            left[v][i] = u;
            return true;
  return false;
int MaxMatch() { //求多重最大匹配
  int ans = 0;
  memset(vcnt, 0, sizeof(vcnt));
  memset(left, -1, sizeof(left));
  for(int u = 0; u < n; u ++) {
     memset(T, false, sizeof(T));
     if(match(u)) ans ++;
  return ans;
bool Perfect(){ //求多重完美匹配
```

```
memset(vcnt, 0, sizeof(vcnt));
     memset(left, -1, sizeof(left));
    for(int u = 0; u < n; u ++) {
       memset(T, false, sizeof(T));
       if(!match(u)) return false;
     return true;
};
BPM Mt;
char str[105];
void Pre() {
  scanf("%d", &N);
  Mt.init(7 * 50 + 1, N);
  Dsum = 0;
  for(int i = 0; i < N; i ++) {
    for(int j = 0; j < 7; j ++){
       scanf("%d", &an[j]);
     scanf("%d%d", &D[i], &W);
    Dsum += D[i];
     for(int j = 0; j < W; j ++) {
       for(int k = 0; k < 7; k + +)
         if(an[k] == 1) {
            Mt.G[j * 7 + k][i] = true;
```

```
int main()
  int Tot;
  scanf("%d", &Tot);
  while(Tot --){
     Pre();
     int ans = Mt.MaxMatch();
     if(ans == Dsum) printf("Yes\n");
     else printf("No\n");
  return 0;
奇偶匹配&位运算
const int \max n = 60 * 60 + 10;
const int dy[] = \{-2, -1, 1, 2, 2, 1, -1, -2, 0, 1, 0, -1\};
const int dx[] = \{-1, -2, -2, -1, 1, 2, 2, 1, -1, 0, 1, 0\};
BPM mt;
int N, M;
int mat[60][60];
bool check(int x, int y) {
  if(x < 0 \parallel y < 0 \parallel y >= M \parallel x >= N) return false;
   return true;
```

```
void conect(int x, int y, int val) {
  int u = x * M + y;
  for(int k = 0; k < 12; k ++) {
     if(!((val \gg k) & 1)) continue;
    int nx = x + dx[k], ny = y + dy[k];
     if(check(nx, ny) && mat[nx][ny] != -1) {
       int v = nx * M + ny;
       if((x + y) & 1) //x+y 为奇数的向偶数连边
          mt.AddEdge(u, v);
       if((nx + ny) \& 1)
          mt.AddEdge(v, u);
void Prepare() {
  mt.init(N*M);
  for(int i = 0; i < N; i ++) {
     for(int j = 0; j < M; j ++) {
       scanf("%d", &mat[i][j]);
  for(int i = 0; i < N; i ++) {
     for(int j = 0; j < M; j ++) {
       if(mat[i][j] != -1) {
          conect(i, j, mat[i][j]);
```

```
int main(){
  int ac = 1;
 while(scanf("%d%d", &N, &M), N + M) {
   Prepare();
   int ans = mt.MaxMatch();
   printf("%d. %d\n", ac ++, ans);
 return 0;
输入 N*N 矩形的数据,通过行交换或者列交换来使正对角线的数值都为
1.
      题解:按照正常逻辑也可以进行判断。
      转换成二分图匹配问题更简单一些 (行号和列号相等)
   • 左边为矩形的行数,右边为矩形中'1'所在的列。
struct BPM {
  int n, m;
 int left[maxn];
 bool T[maxn];
 int G[maxn][maxn];
 void init(int n, int m) {
   this->n = n;
   this->m = m;
```

```
bool dfs(int u) {
     for(int v = 1; v \le m; v ++) {
       if(!T[v] && G[u][v]) {
          T[v] = true;
          if(left[v] == -1 \parallel dfs(left[v])) {
             left[v] = u;
             return true;
     return false;
  bool Perfect() {
     memset(left, -1, sizeof(left));
     for(int u = 1; u \le n; u ++) {
       memset(T, false, sizeof(T));
       if(!dfs(u)) return false;
     return true;
};
BPM mt;
int N;
int x[maxn], y[maxn];
int cnt;
```

```
void Prepare() {
  mt.init(N, N);
  for(int i = 1; i \le N; i ++) {
     for(int j = 1; j \le N; j ++) {
        scanf("%d", &mt.G[i][j]);
int main()
  while(scanf("%d", &N) != EOF) {
     Prepare();
     bool ok = mt.Perfect();
     if(!ok) {printf("-1\n"); continue;}
     cnt = 0;
     for(int i = 1; i \le N; i ++) {
       int cur = i, j;
       for(j = i; j \le N; j ++) {
          if(mt.left[cur] >= mt.left[j]) cur = j;
        if(cur != i) {
          x[cnt] = cur, y[cnt ++] = i;
          swap(mt.left[cur], mt.left[i]);
     printf("%d\n", cnt);
```

```
for(int i = 0; i < cnt; i ++) {
       printf("C %d %d\n", x[i], y[i]);
  return 0;
一行变多行,一列变多列
X[i][j] 和 Y[i][j]分别表示第 i 行 j 列的新的 X,Y 坐标值
memset(X, -1, size of X);
  int x = -1;
  for(int i = 1; i \le R; i ++) {
    for(int j = 1; j \le C; j ++) {
       if(g[i][j] == '.');
       else {
         if(g[i][j-1]!= '*') X[i][j] = ++ x;
         else X[i][j] = x;
  memset(Y, -1, size of Y);
  int y = -1;
  for(int j = 1; j \le C; j ++) {
    for(int i = 1; i \le R; i ++) {
       if(g[i][j] == '.');
       else {
         if(g[i-1][j] != '*') Y[i][j] = ++ y;
```

```
else Y[i][j] = y;
 求最小覆盖。X和Y为最小覆盖中的点集
 int mincover(vector<int>& X, vector<int>& Y) {
    int ans = MaxMatch();
    memset(S, false, sizeof(S));
    memset(T, false, sizeof(T));
    for(int u = 0; u < n; u ++)
      if(right[u] == -1) dfs(u); // 从所有 X 未盖点出发增广
    for(int u = 0; u < n; u ++)
      if(!S[u]) X.push_back(u); // X 中的未标记点
    for(int v = 0; v < m; v ++)
      if(T[v]) Y.push back(v); // Y 中的已标记点
   return ans;
二分图最佳匹配 (最大权)
KM 模版
int W[maxn][maxn];
int lx[maxn], ly[maxn];
int left[maxn];
int slack[maxn]; //
bool S[maxn], T[maxn];
int ny, nx; //x 结点和 y 结点的个数
```

```
int N, M, Tot;
int Z;
bool dfs(int u) {
  S[u] = true;
  for(int v = 1; v \le ny; v + +) {
     if(T[v]) continue;
     int tmp = lx[u] + ly[v] - W[u][v];
     if(tmp == 0) 
       T[v] = true;
       if(left[v] == -1 \parallel dfs(left[v])) {
          left[v] = u;
          return true;
     } else if(slack[v] > tmp) {
       slack[v] = tmp;
  return false;
void update() {
  int a = inf;
  for(int j = 1; j \le ny; j ++) {
     if(!T[j]) {
       a = min(a, slack[j]);
```

```
for(int i = 1; i \le nx; i ++) {
     if(S[i]) lx[i] = a;
  for(int i = 1; i \le ny; i + +){
     if(T[i]) ly[i] += a;
     else{
        slack[i] = a;
void KM() {
  memset(left, -1, sizeof left);
  memset(ly, 0, size of ly);
  for(int i = 1; i \le nx; i ++) {
     lx[i] = -inf; //=0 || -inf 视具体情况
     for(int j = 1; j \le ny; j ++) {
       lx[i] = max(lx[i], W[i][j]);
  for(int i = 1; i \le nx; i ++) {
     for(int j = 1; j \le ny; j ++)
        slack[j] = inf;
     while(true) {
        memset(T, false, sizeof(T));
        memset(S, false, sizeof(S));
        if(dfs(i)) break;
```

```
else update();
void Prepare(){
  scanf("%d%d", &N, &M);
  ny = N; nx = N;
  for(int i = 1; i \le N; i ++) {
    for(int j = 1; j \le N; j ++) {
       W[i][j] = -inf;
  int u, v, w;
  while(M --) {
    scanf("%d%d%d", &u, &v, &w);
    W[u][v] = max(-w, W[u][v]);
int main()
  scanf("%d", &Tot);
  while(Tot --) {
    Prepare();
     KM();
    int ans = 0;
    for(int i = 1; i \le ny; i + +){
```

```
if(left[i] == -1) continue;
       ans += W[left[i]][i];
     printf("%d\n", -ans);
  return 0;
回溯求所有解
void Pre(){
  scanf("%d", &N);
  ny = N;
  nx = N;
  memset(W, 0, size of W);
  for(int i = 1; i \le N; i ++) {
     for(int j = 0; j < N; j ++) {
       scanf("%d", &Z);
       W[Z][i] = j;
  for(int i = 1; i \le N; i ++) {
     for(int j = 0; j < N; j ++) {
       scanf("%d", &Z);
       W[i][Z] = j;
```

```
void find(int dep, int sum) {
  if(sum > ans) return;
  if(dep > nx) {
     if(sum != ans) return;
     printf("Best Pairing %d\n", cnt ++);
     for(int i = 1; i \le ny; i ++) {
       printf("Supervisor %d with Employee %d\n", i, left[i]);
  else {
     for(int i = 1; i \le ny; i ++) {
       if(!vis[i]) {
          vis[i] = true;
          left[dep] = i;
          find(dep + 1, sum - W[dep][i]);
          vis[i] = false;
int main()
  scanf("%d", &Tot);
  for(int ac = 1; ac \leftarrow Tot; ac ++) {
     Pre();
     KM();
```

```
ans = 0;
     for(int i = 1; i \le ny; i + +){
       if(left[i] == -1) continue;
       ans -= W[left[i]][i];
     printf("Data Set %d, Best average difference: ", ac);
     printf("%.6f\n", (double)ans / 2 / nx);
     cnt = 1;
     memset(vis, false, sizeof(vis));
     find(1, 0);
     printf("\n");
  return 0;
偏好度
struct Brute {
  int H, A;
Brute St[maxn], Xt[maxn];
int V[maxn];
int check(Brute St, Brute Xt) {
  return (St.H - 1) / Xt.A >= (Xt.H - 1) / St.A ? 1 : -1;
void Prepare() {
  nx = ny = N;
  for(int i = 1; i \le N; i ++) scanf("%d", &V[i]);
```

```
for(int i = 1; i \le N; i ++) scanf("%d", &St[i].H);
  for(int i = 1; i \le N; i ++) scanf("%d", &Xt[i].H);
  for(int i = 1; i \le N; i ++) scanf("%d", &St[i].A);
  for(int i = 1; i \le N; i ++) scanf("%d", &Xt[i].A);
  for(int i = 1; i \le N; i ++) {
     for(int j = 1; j \le N; j ++) {
       W[i][j] = 10 * check(St[i], Xt[j]) * V[i];
       if(i == j) W[i][j] ++;
int main()
  while(scanf("%d", &N), N) {
     Prepare();
     KM();
     int cnt = 0, ans = 0;
     for(int i = 1; i \le N; i ++) {
       if(W[left[i]][i] % 10 != 0) {cnt ++; W[left[i]][i] --;}
       ans += W[left[i]][i];
     ans = 10;
     if(ans < 0) printf("Oh, I lose my dear seaco!\n");
     else printf("%d %.3f%%\n", ans, cnt * 100.0 / N);
  return 0;
```