# Linear Optimization - Homework 2

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## Problem 1. $(1^{st} \text{ Bullet})$

Formally, following the argument on pg. 83, if **d** is a feasible direction at **x**, that means  $\exists \theta > 0$  s.t.  $\mathbf{x} + \theta \mathbf{d} \in \text{the feasible polyhedron}$ , P. Since we're only looking for feasible solutions,  $\mathbf{A}(\mathbf{x} + \theta \mathbf{d}) = \mathbf{b}$ . Also, since x is feasible,  $\mathbf{A}\mathbf{x} = \mathbf{b}$ . Since,  $\theta > 0$  and subtracting b from both sides, it is apparent Ad = 0.

Now for the second part of the statement, we're moving in a feasible direction in the polyhedron, so  $\mathbf{x} + \theta \mathbf{d} \geq 0$ . That means for a given index, i,  $x_i + \theta d_i \geq 0$ . Since  $\theta > 0$ ,  $d_i \geq 0$  for any index where  $x_i = 0$ .

### (2<sup>nd</sup> Bullet)

We should follow the same analysis as above.

If we can move in the d direction then  $\mathbf{A}(x+\theta\mathbf{d}) = \mathbf{b}$ . Since  $\mathbf{A}\mathbf{x} = \mathbf{b}$ , those terms cancel leaving  $\theta \mathbf{A}\mathbf{d} = 0$ . We divide out  $\theta$  leaving  $\mathbf{A}\mathbf{d} = 0$ . This proves the first statement of the set.

For,  $\mathbf{Dd} \leq 0$ , we again start with the definition of moving in a feasible direction. This gives us  $\mathbf{D}(\mathbf{x} + \theta \mathbf{d}) \leq \mathbf{f}$ . From the problem statement,  $\mathbf{Dx} \leq \mathbf{f}$ , so those terms cancel. That leaves us with  $\theta \mathbf{Dd} \leq 0$ . We divide out  $\theta$ , leaving us with the second statement defining the set,  $\mathbf{Dd} \leq 0$ .

**Problem 2.** (1<sup>st</sup> Bullet) The extreme points will be on the axes of the polyhedron. That means the extreme points will be  $\{(0,0,\frac{1}{3}),(0,\frac{1}{2},0),(1,0,0),(0,0,0)\}$ .

 $(2^{nd} \text{ Bullet})$  Assuming cost can not be negative,

$$\mathbf{x} = \begin{pmatrix} 0 \\ 0 \\ 1/3 \end{pmatrix} \implies \mathbf{c} = \begin{pmatrix} c_1 \\ c_2 \\ 0 \end{pmatrix} \forall c_1, c_2$$

$$\mathbf{x} = \begin{pmatrix} 0 \\ 1/2 \\ 0 \end{pmatrix} \implies \mathbf{c} = \begin{pmatrix} c_1 \\ 0 \\ c_3 \end{pmatrix} \forall c_1, c_3$$

$$\mathbf{x} = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix} \implies \mathbf{c} = \begin{pmatrix} 0 \\ c_2 \\ c_3 \end{pmatrix} \forall c_2, c_3$$

$$\mathbf{x} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \implies \mathbf{c} = \begin{pmatrix} c_1 \\ c_2 \\ c_3 \end{pmatrix} \forall c_2, c_3$$

Problem 3.12. (a) In standard form, our problem is described as

minimize 
$$-2x_1 - x_2$$
st 
$$x_1 - x_2 + x_3 = 2$$

$$x_1 + x_2 + x_4 = 6$$

$$x_1, x_2, x_3, x_4 \ge 0$$

BFS of  $(x_1, x_2) = (0, 0)$  yields a cost of 0 and  $x_3 = 2$  and  $x_4 = 6$ . (b)

That means we have an optimal solution at  $x_1 = 4, x_2 = 2$ . (c) Graph goes here.

#### Problem 4.

$$y_{1} \begin{bmatrix} & x_{1} & x_{2} & x_{3} & x_{4} & x_{5} & y_{1} & y_{2} & y_{3} \\ \hline -5 & -1 & -1 & -3 & -1 & -2 & 0 & 0 & 0 \\ \hline 2 & 1 & 3 & 0 & 4 & 1^{*} & 1 & 0 & 0 \\ 2 & 1 & 2 & 0 & -3 & 1 & 0 & 1 & 0 \\ y_{3} & 1 & -1 & -4 & 3 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

$$(4)$$

We still have  $y_2$  in the basis which we must drive out. We apply a change of basis to drive it out.

This means our Phase I BFS is  $(x_5, x_2, x_3) = (2, 0, \frac{1}{3})$ . We now re-introduce the original objective function.

#### Problem 5. (a)

(b) The price of gasoline can go arbitrarily high because that does not change the geometry of the solution.

**Problem 6.** We formulate this as a transportation problem. The supplier nodes S will be the teams with "supply" or number of wins  $\mathbf{x} = (x_1, x_2, ..., x_n)$ . The consumer nodes, C, will be games between  $S_i$  and  $S_j$ . That means there will be a total number of  $\frac{n(n-1)k}{2}$  which is equal to  $\sum_{i=1}^{n} x_i$ . The arcs are between nodes in S and C and they represent the number of games won by i against team j. We then solve the transportation problem for a our test vector  $\mathbf{x}$ . If the problem for  $\mathbf{x}$  is feasible then  $\mathbf{x}$  is a valid solution.